

# A friendly user guide to ELF micropatching with Shiva (alpha v1)

# **Shiva Manual**

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# 1 Introduction to Shiva

# **ELF program interpreters**

Shiva is an ELF microcode-patching technology that works similarly to the well known (but little understood) ELF dynamic linker "/lib/ld-linux.so". Every ELF program that is dynamically linked has a program header segment type called PT\_INTERP that describes the path to the program interpreter. The program interpreter is a separate ELF program loaded by the kernel that is responsible for helping to build the process runtime image by loading and linking extra modules containing code and data. Traditionally these modules are called shared libraries, and the program interpreter is "/lib/ld-linux.so". On aarch64 it is "/lib/ld-linux-aarch64.so.1"

Image 1.0: Requesting program interpreter

```
Elf file type is DYN (Shared object file)
Entry point 0x7e0
There are 8 program headers, starting at offset 64
rogram Headers:
            Offset
                                           PhysAddr
                            VirtAddr
 Type
                            MemSiz
                                            Flags Align
 PHDR
             0x00000000000001c0 0x00000000000001c0 R
 INTERP
             0x0000000000000001b 0x000000000000001b R
    [Requesting program interpreter: /lib/ld-linux-aarch64.so.1]
             LOAD
             0x0000000000000a98 0x0000000000000a98 R E
                                                 0x10000
 LOAD
             0x000000000000d40 0x00000000010d40 0x000000000010d40
             0x00000000000002d0 0x00000000000042e0 RW
                                                 0x10000
 DYNAMIC
             0x000000000000d50 0x000000000010d50 0x000000000010d50
             0x0000000000000200 0x0000000000000200
                                           RW
                                                 0x8
 NOTE
             0x000000000000021c 0x00000000000021c 0x00000000000021c
             0x0000000000000044 0x0000000000000044
                                           R
                                                 0 \times 4
 GNU STACK
             0x0000000000000000 0x000000000000000 RW
                                                 0x10
 GNU RELRO
             0x000000000000d40 0x00000000010d40 0x000000000010d40
             0x00000000000002c0 0x00000000000002c0 R
                                                 0x1
```

Just like any other interpreter, such as Python, or Java, the ELF program interpreter is responsible for performing various types of loading, patching and transformation.

# Shiva is an ELF program interpreter

Shiva is a newly designed program interpreter who's sole purpose is to load ELF microcode patches that are written in C and compiled into Shiva modules (ELF relocatable objects). Shiva modules are relocatable objects and by nature already contain the ELF meta-data (sections, symbols relocations) necessary to re-build much of the process image. Additionally Shiva has introduced some new concepts such as *Transformations* which are an extension to traditional ELF relocations, allowing for complex patching operations such as function splicing.

The core purpose of Shiva is to create a microcode patching system that seamlessly fits into the existing ELF ABI toolchain of compilers, linkers, and loaders. Shiva aims to load and link patches that are written 100% in C, in a way that is natural for developers who may not be reverse engineers but are experts in C. Shiva is in-itself a linker and a loader that works alongside and in conjunction with the existing Dynamic linker "ld-linux.so". To understand more technical details see "shiva\_final\_desing.pdf" within the documentation directory.

# Systems and arch's that Shiva supports

Currently Shiva's AMP tailored-microcode patching system runs on Linux AArch64, and supports AArch64 Linux ELF PIE binaries. Future support for ARM32, x86\_64/x86\_32 and other requested architectures. At the present moment Shiva doesn't support ET\_EXEC ELF binaries (non-pie) due to time constraints, but it's on the road map.

# 2 Installing Shiva-AArch64

## **Download Source code**

git clone git@github.com:advanced-microcode-patching/shiva git clone github.com/elfmaster/libelfmaster

# **Build source code**

#### Build and install libelfmaster

Shiva requires a custom libelfmaster branch that has not been merged into main yet. Specifically make sure to use the *aarch64\_support* branch.

cd libelfmaster git checkout aarch64\_support cd ./src sudo ./make.sh

#### **Build and install Shiva**

#### Commands:

cd shiva make make shiva-ld make patches sudo make install

#### **Build results:**

Shiva is a static ELF executable and installed to "/lib/shiva". The Shiva modules are typically stored in "/opt/shiva/modules". The example patches all live within "/git/shiva/modules/aarch64\_patches/".

# Build and install example patches

Please see the "shiva/modules/aarch64\_patches/" directory. All of the example patches live here. Let us illustrate an example of building and testing one such patch on a program called 'test\_rodata'.

#### Commands:

cd shiva/modules/aarch64\_patches/rodata\_interposing make sudo make install

Test the patch by setting the patch path and running "/lib/shiva" manually.

\$ SHIVA\_MODULE\_PATH=./ro\_patch.o /lib/shiva ./test\_rodata

Notice that the software prints the new string "The great arcanum", per the patch. Yet when the program is executed without Shiva it executes with its old behavior.

\$./test\_rodata

#### Install the patch permanently

You may also install the patch permanently with the Shiva prelinking tool. This tool updates the PT\_INTERP segment from "/lib/Id-linux.so" to "/lib/shiva", and updates the dynamic segment with several new tags describing the path to the patch that should be loaded and linked at runtime: "/opt/shiva/modules/ro\_patch.o".

\$ /bin/shiva-ld -e test\_rodata -p ro\_patch.o -i /lib/shiva -s /opt/shiva/modules -o test\_rodata.patched

Shiva now installs the patch in memory everytime it's executed

#### \$./test\_rodata.patched

The only discernible difference between test\_rodata.patched and test\_rodata are the PT\_INTERP segment modification and the PT\_DYNAMIC addition of two new entries describing the patch path. No actual code or data has changed within the ELF binary. Shiva performs all of the patching at runtime.

# 3 Writing patches with Shiva

#### The Shiva workflow

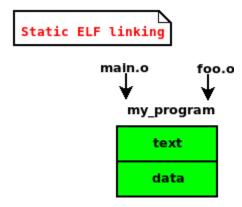
Before delving in to the use of Shiva, let us first explain the workflow of what the ELF compilation and linking process looks like before and after Shiva.

# **ELF Program compilation and linking example**

Standard program compilation begins with a compiler like gcc; the source code files (l.e. main.o and foo.o) are compiled and then get linked (By "/bin/ld") into a final output ELF executable my\_program.

# Static linking of objects

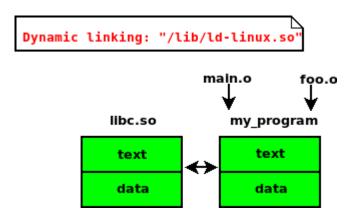
Image 3.0: static linking of two ELF relocatable objects with "/bin/ld"



Additionally these ELF relocatable object files contain calls to functions that live within ELF shared libraries. Once the object files (foo.o and main.o) are linked into an ELF executable (my\_program), additional Dynamic linking meta-data is created and stored into what is known as the PT\_DYNAMIC segment of the program header table. The PT\_DYNAMIC segment contains additional meta-data that is to be passed along to the Program interpreter at runtime. The program interpreter (aka: dynamic linker) is specified as a file path in the PT\_INTERP segment: "/lib/Id-linux.so". When the executable is run, the kernel loads and executes the interpreter program first: "/lib/Id-linux.so". The interpreter loads and links the needed shared libraries (i.e. libc.so). The program has now been compiled, and linked twice: once statically, and again dynamically.

#### Dynamically linked executable

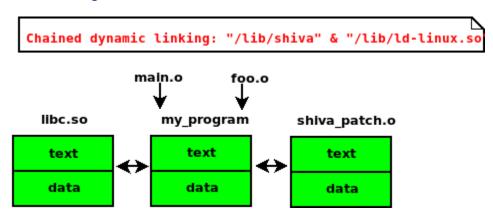
Image 3.1: Dynamically link libc.so in at runtime



A Shiva patch can be written in C and compiled as an ELF relocatable object file. In the same sense that "Id-linux.so" loads and links shared libraries (ELF ET\_DYN files), Shiva will load and link patch objects (ELF ET\_REL files) to re-write the program in memory just before execution.

#### Dynamically patched executable

Image 3.2: Chained linking with Shiva



# **ELF Symbol interposition**

The ELF format uses symbols to denote code and data. Program functions generally live in the .text section, and global variables live within the .data, .rodata, and .bss sections. Shiva makes use of the ELF relocation, section, and symbol data to allow patch developers to *interpose* existing symbols; in other words re-link functions and data that have symbols associated with them. As patch developers we will find this as a natural and beautifully convenient way to patch code and data.

## Citation 3.0: From "Oracle linkers and libraries guide"

"Interposition can occur when multiple instances of a symbol, having the same name, exist in different dynamic objects that have been loaded into a process. Under the default search model, symbol references are bound to the first definition that is found in the series of dependencies that have been loaded. This first symbol is said to interpose on the other symbols of the same name."

# Patching an .rodata string (Simple first patch)

This patching exercise lives in "modules/aarch64\_patches/rodata\_interposing". In the following example we are going to look at a simple program that prints a constant string "Arcana Technologies". Constant data is read-only and is usually stored in the .rodata section of a program. The .rodata section is generally put into text segment, or

a read-only segment that is adjacent to the text segment. A shiva module has it's own .rodata section within it's mapped text segment at runtime. Our goal with the following patch is to replace a read-only string with another by simply re-defining it in C.

#### lmage 2.0: Source code of program rodata\_test.c

```
#include <stdio.h>
const char rodata_string[] = "Arcana Technologies";
int main(void)
{
        printf("rodata_string: %s\n", rodata_string);
        exit(0);
}
```

NOTE: The source code to rodata\_test.c is only being shown for the sake of illustrating what the program does that we are patching. In may workflows the source code wouldn't even exist, and Shiva doesn't depend on or use the source in any way.

Image 3.3: Source code of patch object: ro\_patch.c

```
elfmaster@esoteric-aarch64:~/amp/shiva/modules/aarch64_patches/rodata_interposing$ cat ro_patch.c
const char rodata string[] = "The Great Arcanum";
```

# Building the patch

The patch source code "modules/aarch64\_patches/rodata\_interposing/ro\_patch.c" is a single line patch, redefining the symbol "rodata\_string[]". To build this patch we can simply run the following gcc command to build the patch:

\$ gcc -mcmodel=large -fno-pic -fno-stack-protector -c ro\_patch.c -o ro\_patch.o

The ELF ET\_REL object: ro\_patch.o describes a symbol with the same exact name as the symbol **rodata\_string** in the target executable. Shiva will see this when linking and it will link in the version of the symbol from our patch. This perfectly demonstrates the concept of symbol interposition in Shiva.

#### Testing our patch: ro\_patch.o

In this example we are patching the AArch64 ELF PIE executable: "modules/aarch64\_patches/rodata\_interposing/test\_rodata" Let's first run the program un-patched.

\$ ./test\_rodata rodata\_string: Arcana Technologies \$

The program simply prints the string "rodata\_string: Arcana technologies". To see if our patch worked we can invoke Shiva directly to load and link the executable with the patch at runtime. We must specify the path to our patch object with the SHIVA\_MODULE\_PATH environment variable.

\$ SHIVA\_MODULE\_PATH=./ro\_patch.o /lib/shiva ./test\_rodata rodata\_string: The Great Arcanum \$

When executing ./test\_rodata with "/lib/shiva" as the primary loader it installs the specified patch, and the new rodata\_string[] value "The Great Arcanum" is printed to stdout. When we run ./test\_rodata without "/lib/shiva" it does not install the patch and the original string "Arcana Technologies" is printed.

## Using shiva-ld to prelink the patch

The reality is that a user doesn't want to invoke "/lib/shiva" every single time they want to run a program that needs a patch. This is why "/lib/shiva" is designed to be an ELF interpreter like "Id-linux.so". We can use the Shiva prelinker "/bin/shiva-Id" to install the appropriate meta-data into the program we are patching:

- Update PT\_INTERP with the path to "/lib/shiva" (Replacing "/lib/ld-linux.so").
- Add several custom entries to PT\_DYNAMIC describing the patch basename: "ro\_patch.o" and the module search path "/opt/shiva/modules".

The following command will install the Shiva interpreter path and patch information into the ./test\_rodata binary.

\$ shiva-ld -e test\_rodata -p ro\_patch.o -i /lib/shiva -s /opt/shiva/modules

This next command copies the patch file into the correct search path: "/opt/shiva/modules"

\$ sudo cp ro\_patch.o /opt/shiva/modules

Before we run ./test\_rodata, let's take a quick look to see what modifications shiva-ld made to the executable:

Image 3.4: Requesting program interpreter is "/lib/shiva"

Program Headers:			
Type	0ffset	VirtAddr	PhysAddr
	FileSiz	MemSiz	Flags Align
PHDR	0x00000000000000040	0x000000000000000040	0x00000000000000040
	0x000000000000001c0	0x000000000000001c0	R 0x8
INTERP	0x00000000000000200	0x000000000000000200	0x00000000000000200
	0x0000000000000001b	0x0000000000000001b	R θx1
[Requesting	g program interprete	er: /lib/shiva]	
LOAD	0x00000000000000000	0x00000000000000000	0x00000000000000000
	0x00000000000000868	0x00000000000000868	R E 0x10000
LOAD	0x00000000000000d78	0x0000000000010d78	0x0000000000010d78
	0x00000000000000298	0x000000000000002a0	RW 0x10000
DYNAMIC	0x0000000000003000	0x0000000000012000	0x0000000000012000
	0x00000000000001d0	0x000000000000001d0	RW 0x8
LOAD	0x0000000000003000	0x0000000000012000	0x0000000000012000
	0x00000000000000259	0x00000000000000259	RWE 0x1000
GNU_STACK	0x00000000000000000	0x00000000000000000	0x00000000000000000
	0x00000000000000000	0x00000000000000000	RW 0x10
GNU_RELRO	0x00000000000000d78	0x0000000000010d78	0x0000000000010d78
	0x00000000000000288	0x00000000000000288	R θx1

In image 2.2 we can see that "/lib/shiva" is set as the new Interpreter path. The astute reader will also notice that there is a third PT\_LOAD segment marked RWE. The shiva-Id tool creates a new dynamic segment (As shown below in Image 2.3). In order to make room for this new and larger dynamic segment shiva-Id creates a new PT\_LOAD segment and updates the PT\_DYNAMIC segment program header so that it points into the new PT\_LOAD segment at 0x3000. The updated PT\_DYNAMIC segment is moved and stored in this new segment

Image 3.5: New dynamic segment contains 3 custom entries

```
ynamic section at offset 0x3000 contains 29 entries:
           Туре
                                          Name/Value
0x00000000000000001 (NEEDED)
                                         Shared library: [libc.so.6]
0x0000000000000000 (INIT)
0x00000000000000000 (FINI)
                                         0x81c
0x000000000000000019 (INIT ARRAY)
                                        0x10d78
0x0000000000000001b (INIT ARRAYSZ)
                                        8 (bytes)
0x00000000000000001a (FINI ARRAY)
                                         0x10d80
0x00000000000000001c (FINI ARRAYSZ)
                                         8 (bytes)
0x000000006ffffef5 (GNU HASH)
                                         0x388
0x00000000000000005 (STRTAB)
0x00000000000000006 (SYMTAB)
                                         0x280
0x00000000000000000 (STRSZ)
                                         142 (bytes)
0x0000000000000000 (SYMENT)
                                         24 (bytes)
0x00000000000000015 (DEBUG)
                                         \theta x \theta
0x00000000000000003 (PLTGOT)
                                         0x10f78
0x00000000000000000 (PLTRELSZ)
                                         144 (bytes)
0x00000000000000014 (PLTREL)
                                         RELA
0x00000000000000017 (JMPREL)
                                         0x540
0x00000000000000007 (RELA)
                                         0x450
                                        240 (bytes)
0x00000000000000000 (RELASZ)
0x00000000000000000 (RELAENT)
                                         24 (bytes)
0x0000000000000001e (FLAGS)
                                         BIND NOW
0x000000006ffffffb (FLAGS 1)
                                         Flags: NOW PIE
0x000000006ffffffe (VERNEED)
                                         0x430
0x000000006fffffff (VERNEEDNUM)
0x000000006ffffff0 (VERSYM)
                                         0x416
0x0000000060000018 (Operating System specific: 60000018)
                                                                          0x121d0
0x0000000000000017 (Operating System specific: 60000017)
                                                                           0x121e3
0x00000000000000019 (Operating System specific: 60000019)
0x0000000000000000 (NULL)
```

The above image 2.3 shows the dynamic segment of the ELF binary: test\_rodata, after we ran the shiva-Id tool on it. Notice the three *Operating System specific* entries at the tail end. The readelf utility doesn't know how to parse these custom entries. Here is what the actual entry types and values are:

(SHIVA\_NEEDED) Patch object: [ro\_patch.o]
(SHIVA\_SEARCH) Search path: [/opt/shiva/modules]
(SHIVA\_ORIG\_INTERP) Original RTLD: [/lib/ld-linux-aarch64.so.1]

This data is needed by "/lib/shiva" at runtime so that it can locate and load the patch object, and the original dynamic linker (See Chained Linking in section x.x).

## Final outcome of patching test\_rodata

Now every time that the test\_rodata executable is ran it will invoke the Shiva interpreter which will in turn load the patch object "/opt/shiva/modules/ro\_patch.o".

```
$ ./test_rodata
rodata_string: The Great Arcanum
$
```

We can observe that the program test\_rodata is now printing the patched version of rodata\_string[] when it is ran. This patch may appear permanent, but in effect the patch is actually being installed at runtime by "/lib/shiva" everytime the program runs.

#### Current limitations of patching const data

There are patching scenarios with *const* data that fail with symbol interposition due to code optimizations with read-only data. In the previous example we illustrated how to patch a read-only global variable defined in C as:

```
const char rodata_string[] = "Arcana Technologies";
```

The target program gets re-linked to use the patch version of *rodata\_string*. Using symbol interposition to overwrite *.rodata* variables works perfectly unless the read-only value is optimized out of memory and stored only in a register. Let's look at the following example:

Image 3.6: C code illustrating a read-only global variable

```
const int rodata_val = 5;
int main(void)
{
         printf(".rodata string: %d\n", rodata_val);
         exit(0);
}
```

After compiling the code above with gcc, even after all optimizations are disabled, you will get code similar to the following AArch64 assembly.

Image 3.7: Code optimization of read-only variable in aarch64 assembly

```
Dump of assembler code for function main:
  0x0000aaaaaaaaa774 <+0>:
                                stp
                                        x29, x30, [sp, #-16]!
  0x0000aaaaaaaaa778 <+4>:
                                        x29, sp
                                mov
                                                                         // #5
=> 0x0000aaaaaaaaa77c <+8>:
                                        w1, #0x5
                                mov
  0x0000aaaaaaaaa780 <+12>:
                                        x0, 0xaaaaaaaaa000
                                adrp
                                        x0, x0, #0x840
  0x0000aaaaaaaaa784 <+16>:
                                add
  0x0000aaaaaaaaa788 <+20>:
                                bl
                                        0xaaaaaaaaa660 <printf@plt>
                                                                         // #0
  0x0000aaaaaaaaa78c <+24>:
                                        w0, #0x0
                                mov
  0x0000aaaaaaaaa790 <+28>:
                                        0xaaaaaaaaa610 <exit@plt>
                                bl
```

In the code above we can see that the value 5 is being copied as an immediate value into a register. This is an optimization that removes the need for an adrp/add/ldr instruction trio, thus eliminating two instructions and the need for memory access. This optimization is inauspicious for symbol interposition based patching since the actual *.rodata* variables

memory is not being accessed and therefore there is no linking code to update. Instead there is only an instruction: "mov w1, #0x5". Clearly this instruction can be re-encoded rather simply, but what would a patch look like for this? For the curious reader, move ahead to the section titled "ELF Transformations & Function splicing".

## Symbol interposition on functions

Shiva gives patch writers the ability to replace a function trivially using symbol interposition. We've all seen userland rootkits that re-write libc.so functions by introducing a preloaded library with a new version of the function. In the same spirit Shiva allows patch writers to replace any function in the target ELF executable with a symbol of type STT\_FUNC and symbol bindings of STB\_GLOBAL or STB\_WEAK. The only exception is that the function main() cannot be replaced, although future support is on the way, as it depends on the glibc initialization code that transfers control to main from \_start.

Let's get started with a simple example of replacing a function called foo().

## Patching binary test\_func by replacing function foo()

\$ cd shiva\_examples/function\_interposing

Image 3.8: Original source code of the binary *test\_foo* 

```
elfmaster@esoteric-aarch64:~/amp/shiva_examples/function_interposing$ cat test_foo.c
#include <stdio.h>

int foo(void)
{
         printf("I am the original foo() function\n");
         return 1;
}

int main(void)
{
         int ret;
         ret = foo();
         printf("Return value: %#x\n", ret);
}
elfmaster@esoteric-aarch64:~/amp/shiva_examples/function_interposing$
```

We can clearly see that foo() is a global function called by main(). In order to replace foo() with our own version we can simply re-write the function by name in our patch source code.

## lmage 3.9: foo\_patch.c source code

```
elfmaster@esoteric-aarch64:~/amp/shiva_examples/function_interposing$ cat foo_patch.c
#include <stdio.h>
int foo(void)
{
    printf("I am the new function foo()\n");
    return 0x31337;
}
elfmaster@esoteric-aarch64:~/amp/shiva_examples/function_interposing$
```

# Symbol interposition on global data: .rodata, .bss, .data

Initialized global variables (that are not const) are stored within the .data section.

\$ cd modules/aarch64\_patches/dataonly\_interposing

Original source code of binary

Redefining global data types

**Patching STB\_LOCAL variables** 

Using the extern keyword in your patch code