

A Brief Introduction to LLVM

Nick Sumner

(with a few modifications by Matt Dwyer)

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 - Machine code generation libraries
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 - A simple, typed IR (*bitcode*)
 - Program analysis / optimization libraries
 - Machine code generation libraries
 - Tools that compose the libraries to perform tasks
- Easy to add / remove / change functionality

How can you use it?

.Compiling programs to bitcode:

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clang -g -c -emit-llvm <sourcefile> -o <bitcode>.bc
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•Writing your own tools:

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./callcounter -static test.bc
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- Writing your own tools:

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- Reporting properties of the program:

Function Counts

=====

b : 2

a : 1

printf : 3

What is LLVM Bitcode?

- A (Relatively) Simple IR

```
#include<stdio.h>

void
foo(unsigned e) {
    for (unsigned i = 0; i < e; ++i) {
        printf("Hello\n");
    }
}

int
main(int argc, char **argv) {
    foo(argc);
    return 0;
}
```

Code

`clang -c -S -emit-llvm -O1 -g0`

```
@str = private constant [6 x i8] c"Hello\00"

define void @foo(i32) {
    %2 = icmp eq i32 %0, 0
    br i1 %2, label %3, label %4

; <label>:3:                                ; preds = %4, %1
    ret void

; <label>:4:                                ; preds = %1, %4
    %5 = phi i32 [ %7, %4 ], [ 0, %1 ]
    %6 = tail call i32 @printf@plt(8* getelementptr
        ([6 x i8], [6 x i8]* @str, i64 0, i64 0))
    %7 = add nuw i32 %5, 1
    %8 = icmp eq i32 %7, %0
    br i1 %8, label %3, label %4
}

define i32 @main(i32, i8** nocapture readnone) {
    tail call void @foo(i32 %0)
    ret i32 0
}
```

IR

What is LLVM Bitcode?

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clang -c -emit-llvm
(and llvm-dis)

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int  
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    return 0;  
}
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Functions

```
@str = private constant [6 x i8] c"Hello\00"
```

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define void @foo(i32) {  
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    br i1 %2, label %3, label %4  
  
; <label>:3:                                ; preds = %4, %1  
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; <label>:4:                                ; preds = %1, %4  
    %5 = phi i32 [ %7, %4 ], [ 0, %1 ]  
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Basic Blocks

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void  
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```
    for (unsigned i = 0; i < e; ++i) {  
        printf("Hello\n");  
    }
```

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int  
main(int argc, char **arg  
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    return 0;  
}
```

labels & predecessors

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```

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; <label>:3:                ; preds = %4, %1  
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```
; <label>:4:                ; preds = %1, %4
```

```
    %5 = phi i32 [ %7, %4 ], [ 0, %1 ]  
    call i32 @puts(i8* getelementptr  
        @str, i64 0, i64 0))
```

```
    %7 = add nuw i32 %5, 1  
    %8 = icmp eq i32 %7, %0  
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```

```
}
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foo(unsigned e) {
```

```
    for (unsigned i = 0; i < e; ++i) {  
        printf("Hello\n");  
    }
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```
}
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```
int  
main(int argc, char **argv) {  
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branches & successors

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; <label>:3:                                ; preds = %4, %1  
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```
; preds = %1, %4  
    %5 = phi i32 [ %7, %4 ], [ 0, %1 ]  
    %6 = tail call i32 @puts(i8* getelementptr  
        ([6 x i8], [6 x i8]* @str, i64 0, i64 0))  
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Instructions

```
@str = private constant [6 x i8] c"Hello\00"

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define i32 @main(i32, i8** nocapture readnone) {
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```

Inspecting Bitcode

- LLVM libraries help examine the bitcode
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```
Module& module = ...;  
for (Function& fun : module) {  
    for (BasicBlock& bb : fun) {  
        for (Instruction& i : bb) {
```

Iterate over the:

- Functions in a Module
- BasicBlocks in a Function
- Instructions in a BasicBlock

Inspecting Bitcode

• LLVM libraries help examine the bitcode

- Easy to examine and/or manipulate
- Many helpers (e.g. CallSite,)

```
Module& module = ...;
for (Function& fun : module) {
  for (BasicBlock& bb : fun) {
    for (Instruction& i : bb) {
      CallSite cs{&i};
      if (!cs.getInstruction()) {
        continue;
      }
    }
  }
}
```

CallSite helps you extract information from Call and Invoke instructions.

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Module &module = ...;
for (Function& fun : module) {
    for (BasicBlock& bb : fun) {
        for (Instruction& i : bb) {
            CallSite cs{&i};
            if (!cs.getInstruction()) {
                continue;
            }
            outs() << "Found a function call: " << i << "\n";
        }
    }
}
```

Inspecting Bitcode

- LLVM libraries help examine the bitcode
 - Easy to examine and/or manipulate
 - Many helpers (e.g. CallSite, outs(), dyn_cast)

```
Module &module = ...;
for (Function& fun : module) {
    for (BasicBlock& bb : fun) {
        for (Instruction& i : bb) {
            CallSite cs{&i};
            if (!cs.getInstruction()) {
                continue;
            }
            outs() << "Found a function call: " << i << "\n";
            Value* called = cs.getCalledValue()->stripPointerCasts();
            if (Function* f = dyn_cast<Function>(called)) {
                outs() << "Direct call to function: " << f->getName() << "\n";
            }
        }
    }
}
```

`dyn_cast()` efficiently checks the runtime types of LLVM IR components.

Dealing with SSA

- You may ask where certain values came from
 - Useful for tracking dependencies
 - “Where was this variable defined?”

Dealing with SSA

- You may ask where certain values came from
- LLVM IR is in SSA form
 - How many acronyms can I fit into one line?
 - What does this mean?
 - Why does it matter?

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```
void foo()  
  unsigned i = 0;  
  while (i < 10) {  
    i = i + 1;  
  }  
}
```

Dealing with SSA

- You may ask where certain values came from
- LLVM IR is in SSA form
 - How many acronyms can I fit into one line?
 - What does this mean?
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```
void foo()  
  unsigned i = 0;  
  while (i < 10) {  
    i = i + 1;  
  }  
}
```

What is the single definition of i at this point?

Dealing with SSA

- The phi instruction

- It selects which of the definitions to use
- Always at the start of a basic block

Dealing with SSA

• Thus the phi instruction

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- Always at the start of a basic block

```
void foo() {  
    unsigned i = 0;  
    while (i < 10) {  
        i = i + 1;  
    }  
}
```

```
define void @foo() {  
    br label %1  
  
; <label>:1 ; preds = %1, %0  
%i.phi = phi i32 [ 0, %0 ], [ %2, %1 ]  
%2 = add i32 %i.phi, 1  
%exitcond = icmp eq i32 %2, 10  
br i1 %exitcond, label %3, label %1  
  
; <label>:3 ; preds = %1  
ret void  
}
```

Dealing with SSA

• Thus the phi instruction

- It selects which of the definitions to use
- Always at the start of a basic block

```
void foo()  
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  }  
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```

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; <label>:1 ; preds = %1, %0  
%i.phi = phi i32 [ 0, %0 ], [ %2, %1 ]  
%2 = add i32 %i.phi, 1  
%exitcond = icmp eq i32 %2, 10  
br i1 %exitcond, label %3, label %1  
  
; <label>:3 ; preds = %1  
ret void  
}
```

Dependencies in General

• You can loop over the values an instruction uses

```
for (Use& u : inst->operands()) {  
    // inst uses the Value* u  
}
```


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```
for (Use& u : inst->operands()) {  
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}
```

```
for %a = %b + %c:  
    [%b, %c]
```

Dependencies in General

• You can loop over the values an instruction uses

```
for (Use& u : inst->operands()) {  
    // inst uses the Value* u  
}
```

• You can loop over the instructions that use a particular value

```
Instruction* inst = ...;  
for (User* user : inst->users())  
    if (auto* i = dyn_cast<Instruction>(user)) {  
        // inst is used by Instruction i  
    }
```

Dealing with Types

- LLVM IR is *strongly typed*
 - Every value has a type → `getType()`

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• A value must be explicitly cast to a new type

```
define i64 @trunc(i16 zeroext %a) {  
    %1 = zext i16 %a to i64  
    ret i64 %1  
}
```

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```
define i64 @trunc(i16 zeroext %a) {  
    %1 = zext i16 %a to i64  
    ret i64 %1  
}
```

- Also types for pointers, arrays, structs, etc.

- Strong typing means they take a bit more work

The `tipc` TIP compiler

- Compiles a TIP program into LLVM bitcode
- Consists of several phases
 - Parsing TIP (uses ANTLR4)
 - Builds an AST
 - Walks the AST to generate bitcode
 - Applies building LLVM optimizations
 - Emits bitcode to a .bc file
- To build an executable we link with libraries

The tipc TIP compiler

```
$ ./tipc --help
```

```
OVERVIEW: tipc - a TIP to llvm compiler
```

```
USAGE: tipc [options] <tip source file>
```

OPTIONS:

Generic Options:

```
-help          - Display available options (-help-hidden for more)
-help-list     - Display list of available options (-help-list-hidden
for more)
-version       - Display the version of this program
```

tipc Options:

Options for controlling the TIP compilation process.

```
-d             - disable bitcode optimization
-l            - pretty print with line numbers
-p            - pretty print
```