

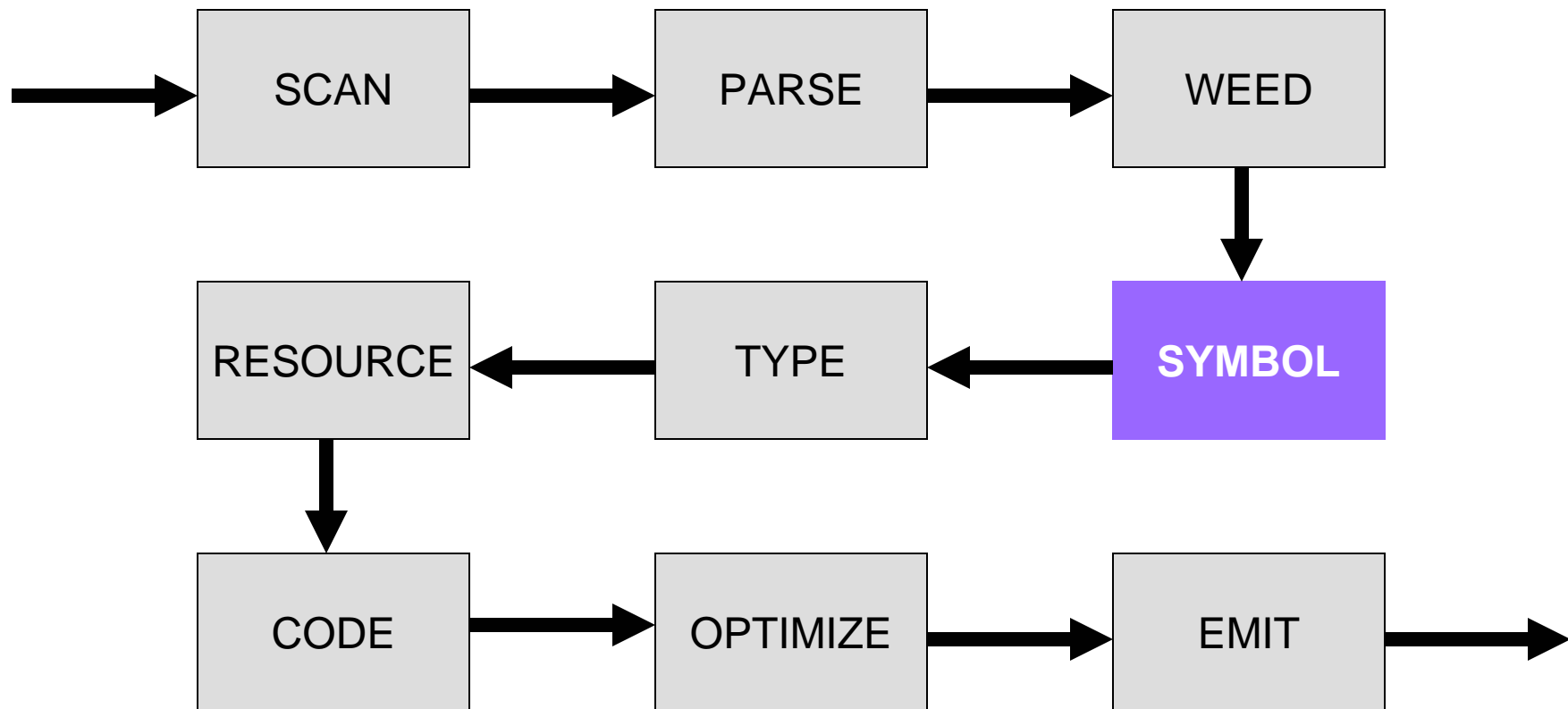


Compiler

Symbol Processing

© Copyright 2005, Matthew B. Dwyer and Robby. The syllabus and lectures for this course are copyrighted materials and may not be used in other course settings outside University of Nebraska-Lincoln and Kansas State University in their current form or modified form without express written permission of one of the copyright holders. During this course, students are prohibited from selling notes to or being paid for taking notes by any person or commercial firm without the express written permission of one of the copyright holders.

Compiler Architecture





Analyzing Identifiers

- Lexical analysis defines *form* of identifiers
- Syntactic analysis defines *where* identifiers can appear
- Symbol analysis defines *correlation* of definition and uses of identifiers
- Grammars are too weak for this

$$\{w\alpha w \mid w \in \Sigma^*\}$$

is not context-free



Symbol Table

- Maps identifiers to their meaning (i.e., definition)

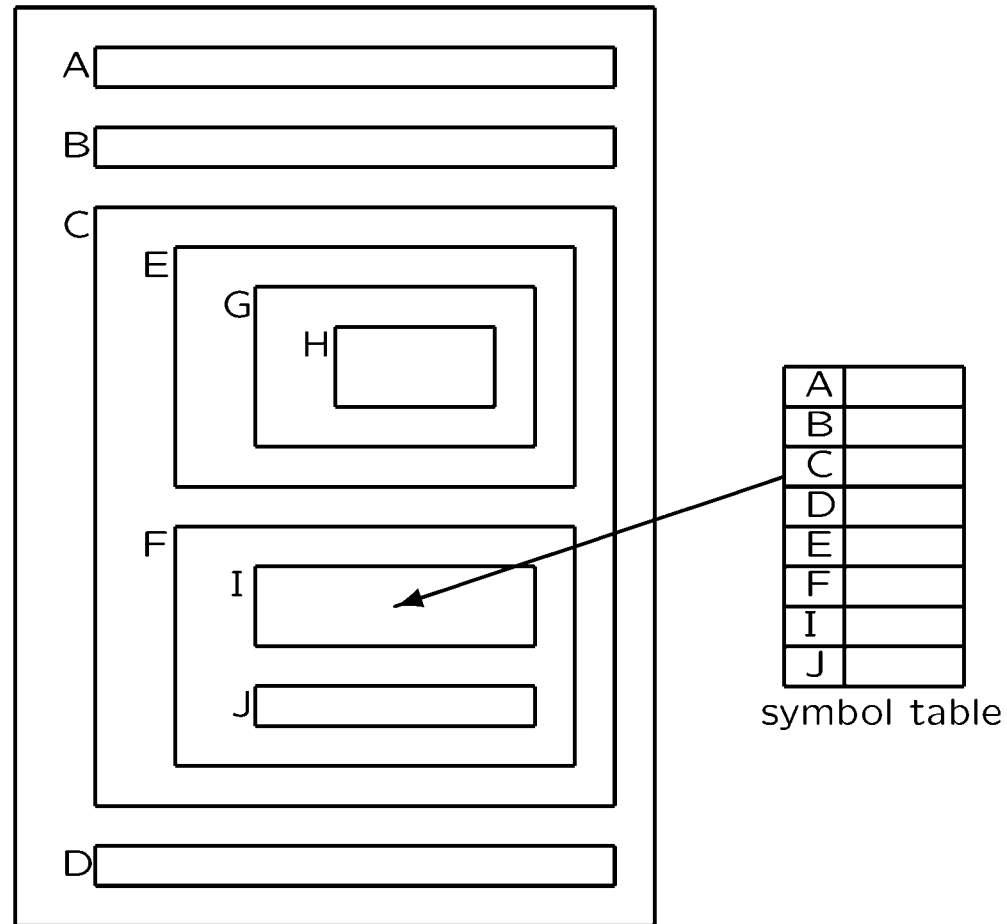
i	local	int
done	local	boolean
insert	method	ldots
x	formal	List
List	class	...
⋮	⋮	⋮



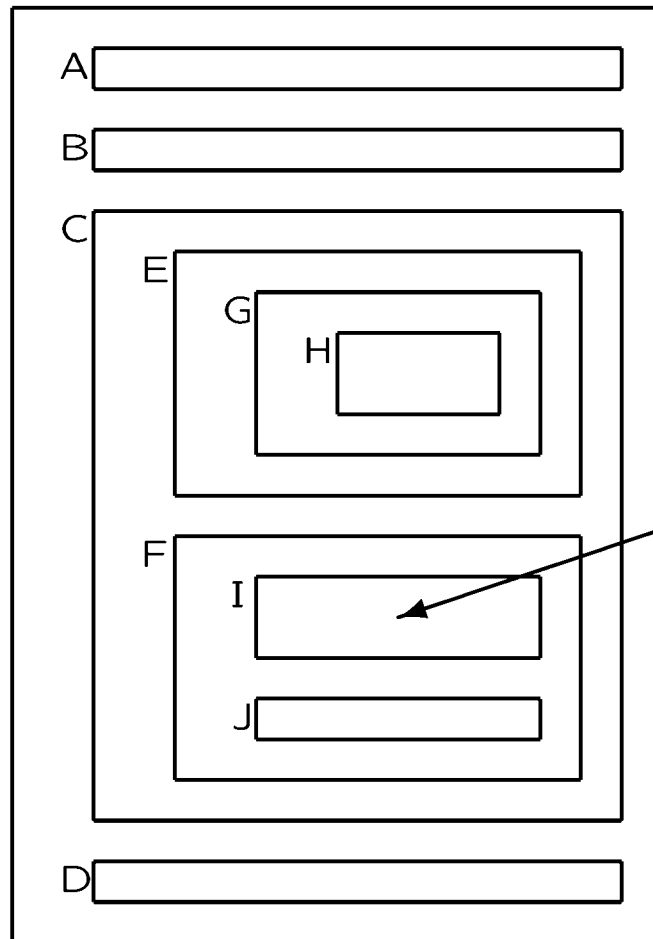
SJ Symbol Table Uses

- which classes are defined;
- which fields are defined;
- which methods are defined;
- what are the signatures of methods;
- are identifiers defined twice;
- are identifiers defined when used; and
- are identifiers used properly?

Static Nested Scope Rules



Nesting vs. Ordering

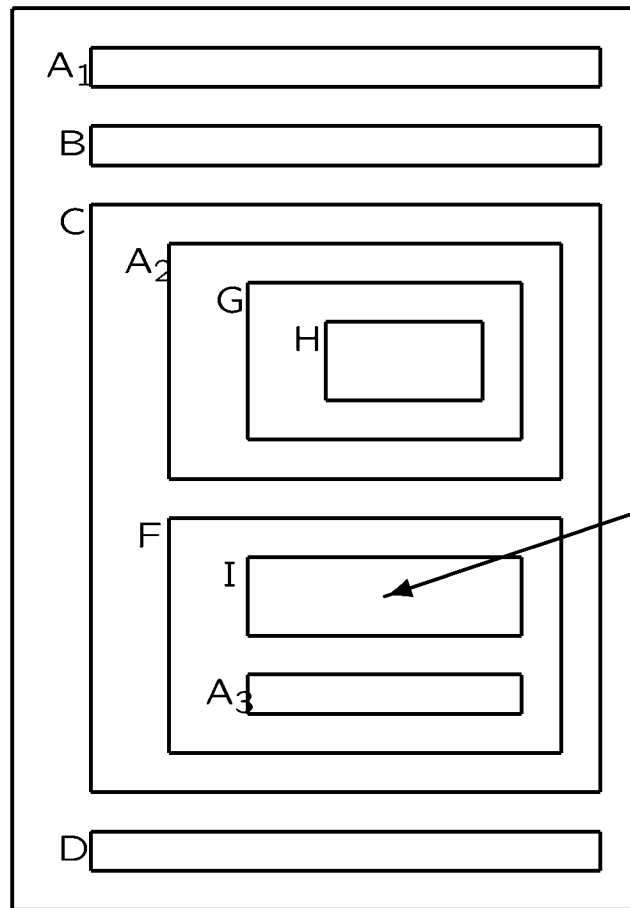


A	
B	
C	
D	
E	
F	
I	
J	

symbol table

Multiple passes are required to eliminate the need for *forward declarations*

Most Closely Nested Definition



A ₃	
B	
C	
D	
F	
I	

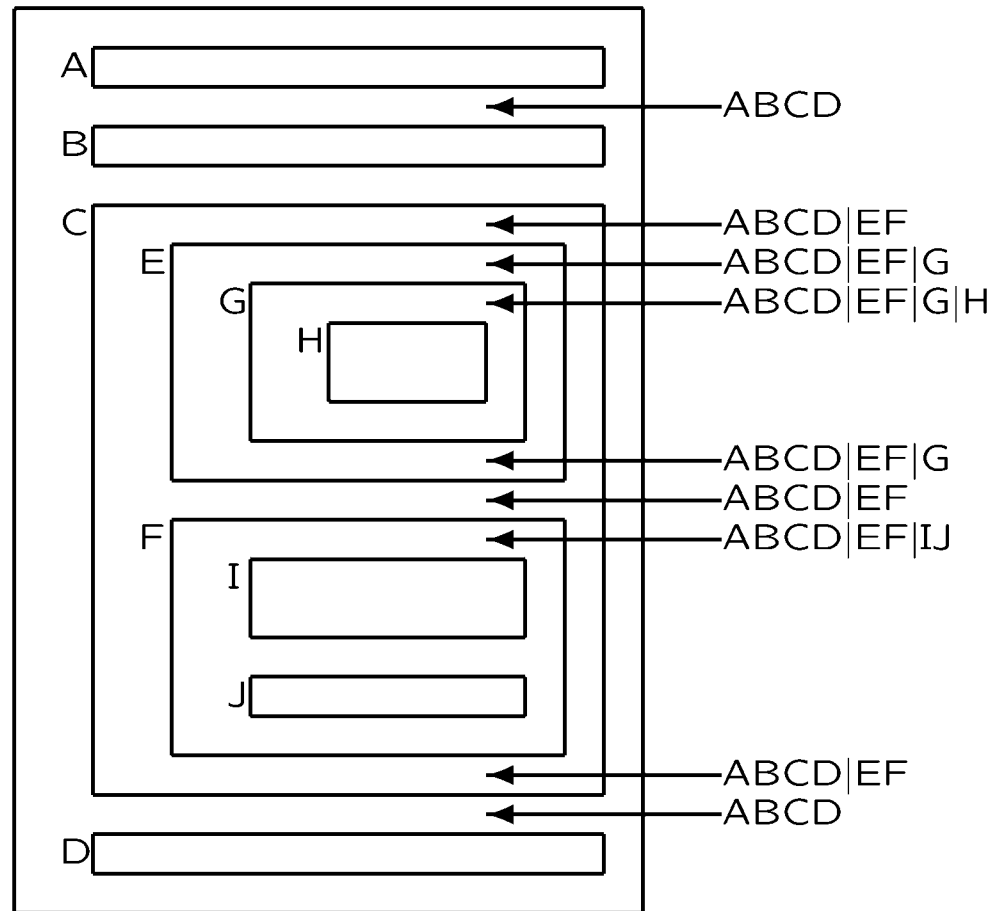
symbol table

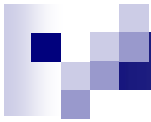
A₂ shadows A₁

A₃ shadows A₂, A₁

Identifiers at same level must be unique

Symbol Table acts like a Stack





Implementation

Symbol table is *stack of hash tables*

- ☐ each hash table contains the identifiers in a level;
- ☐ push a new hash table when a level is entered;
- ☐ each identifier is entered in the top hash table;
- ☐ it is an error if it is already there;
- ☐ a use of an identifier is looked up in the hash tables from top to bottom;
- ☐ it is an error if it is not found;
- ☐ pop a hash table when a level is left.



SJ Symbol Table

- Strict rules on variable naming and variable declarations simplify symbol resolutions
 - variables can only be declared in the beginning of method bodies
 - parameter and local variable names are not allowed to shadow field names
- Other rules
 - method overloading is disallowed
 - allows only one class, etc.

SJ SymbolTable API

[Overview](#) [Package](#) **[Class](#)** [Tree](#) [Deprecated](#) [Index](#) [Help](#)

[PREV CLASS](#) [NEXT CLASS](#)

[FRAMES](#) [NO FRAMES](#)

SUMMARY: [NESTED](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)

DETAIL: [FIELD](#) | [CONSTR](#) | [METHOD](#)

sjc.symboltable

Class SymbolTable

java.lang.Object
└─ sjc.symboltable.SymbolTable

```
public class SymbolTable
extends java.lang.Object
```

This class represents a symbol table for a StaticJava CompilationUnit.

Author:

[Robby](#)

Field Summary

java.util.Map<org.eclipse.jdt.core.dom.ASTNode, java.lang.Object>

[symbolMap](#)

Holds the map of: a SimpleName expression, i.e., a reference to a field, a method parameter, or a local variable, to its corresponding FieldDeclaration, SingleVariableDeclaration, or VariableDeclarationStatement, respectively, and a MethodInvocation expression to its corresponding MethodDeclaration or Method.

Method Summary

java.lang.String [toString\(\)](#)

Returns the String representation of this symbol table.



Mutually Recursive Definitions

- A single traversal of the parse tree is not enough.
- Make two passes:
 - collect definitions of field and method identifiers; and
 - analyze uses of field and method identifiers.
- In cases like recursive types, the definition is not completed before the second traversal.

SymbolTableBuilder API

[Overview](#) [Package](#) **[Class](#)** [Tree](#) [Deprecated](#) [Index](#) [Help](#)

[PREV CLASS](#) [NEXT CLASS](#)

SUMMARY: [NESTED](#) | [FIELD](#) | [CONSTR](#) | [METHOD](#)

[FRAMES](#) [NO FRAMES](#)

DETAIL: [FIELD](#) | [CONSTR](#) | [METHOD](#)

sjc.symboltable

Class SymbolTableBuilder

java.lang.Object

└─ sjc.symboltable.SymbolTableBuilder

```
public class SymbolTableBuilder
    extends java.lang.Object
```

This class is used to build symbol table for a StaticJava CompilationUnit. Note that the algorithm assumes that the JDT AST tree was built using the [ASTParser](#). That is, it assumes certain structures on the AST, e.g., a class does not have an instance method.

Author:

[Robby](#)

Nested Class Summary

static class	SymbolTableBuilder.Error
This class is used to signal an error in the process of building a symbol table.	

Method Summary

static SymbolTable	build (org.eclipse.jdt.core.dom.CompilationUnit cu)
Builds a SymbolTable for the given StaticJava CompilationUnit.	



AST Visitor Code for SymbolTableBuilder

```
protected static class Visitor extends ASTVisitor {  
    public @NonNull Map<ASTNode, Object> result =  
                                                new HashMap<ASTNode, Object>();  
  
    protected @NonNull Map<String, ASTNode> nameMap =  
                                                new HashMap<String, ASTNode>();  
  
    protected @NonNull Map<String, MethodDeclaration> methodMap =  
                                                new HashMap<String, MethodDeclaration>();  
  
    protected @NonNull Set<String> localNames = new HashSet<String>();  
  
    protected String className;  
  
    ...  
  
    protected void dispose() { ... }  
}
```



SymbolTableBuilder Code — TypeDeclaration (1)

```
@Override public boolean visit(TypeDeclaration node) {  
    // remembers the class name  
    className = node.getName().getIdentifier();  
  
    if ("java.lang.String".equals(className)) {  
        throw new Error(node, "Cannot redeclare the String class");  
    }  
  
    // visit field declarations and process method names first  
    // because we want to be able to resolve field and method names  
    // in the method bodies later on  
    for (Object o : node.bodyDeclarations()) {  
        if (o instanceof FieldDeclaration) {  
            ((ASTNode) o).accept(this);  
        } else if (o instanceof MethodDeclaration) {  
            MethodDeclaration md = (MethodDeclaration) o;  
            String methodName = md.getName().getIdentifier();  
            if (methodMap.containsKey(methodName)) { // throw error }  
            methodMap.put(methodName, md);  
        } else { // throw error }  
    }  
}
```




SymbolTableBuilder Code — TypeDeclaration (2)

```
// visit method declarations
for (Object o : node.bodyDeclarations()) {
    if (o instanceof MethodDeclaration) {
        ((ASTNode) o).accept(this);
    }
}
return false;
}
```



SymbolTableBuilder Code — FieldDeclaration

```
@Override public boolean visit(FieldDeclaration node) {  
    VariableDeclarationFragment vdf = (VariableDeclarationFragment) node  
                                     .fragments().get(0);  
  
    String name = vdf.getName().getIdentifier();  
    if (nameMap.containsKey(name)) {  
        throw new Error(node,  
            "Error in field declaration '" + name  
            + "' : the field name has been used in:\n"  
            + nameMap.get(name));  
    }  
    nameMap.put(name, node);  
    return false;  
}
```



SymbolTableBuilder Code — MethodDeclaration

```
@Override public boolean visit(MethodDeclaration node) {  
    for (Object o : node.parameters()) {  
        SingleVariableDeclaration svd = (SingleVariableDeclaration) o;  
        String name = svd.getName().getIdentifier();  
        if (nameMap.containsKey(name)) { // throw error }  
  
        localNames.add(name);  
        nameMap.put(name, svd);  
    }  
    node.getBody().accept(this);  
  
    for (String name: localNames) {  
        nameMap.remove(name);  
    }  
    localNames.clear();  
    return false;  
}
```



SymbolTableBuilder Code — VariableDeclarationStatement

```
@Override public boolean visit(VariableDeclarationStatement node) {  
    VariableDeclarationFragment vdf = (VariableDeclarationFragment) node  
                                     .fragments().get(0);  
  
    String name = vdf.getName().getIdentifier();  
    if (nameMap.containsKey(name)) {  
        throw new Error(node,  
            "Error in local variable declaration '" + name  
            + "' : the variable name has been used in:\n"  
            + nameMap.get(name));  
    }  
    localNames.add(name);  
    nameMap.put(name, node);  
    return false;  
}
```



SymbolTableBuilder Code — MethodInvocation

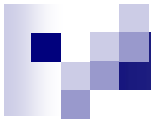
```
@Override public boolean visit(MethodInvocation node) {  
    // Note that we don't visit the MethodInvocation's simple name  
    // because we want visit(SimpleName) to resolve variable references  
    // instead of method names.  
    String methodName = node.getName().getIdentifier();  
  
    if (node.getExpression() == null  
        || className.equals(node.getExpression().toString())) {  
        if (methodMap.containsKey(methodName)) {  
            result.put(node, methodMap.get(methodName));  
        } else { // throw error }  
    } else {  
        // lib call, delay until type checking phase  
    }  
  
    for (Object e : node.arguments()) {  
        ((Expression) e).accept(this);  
    }  
    return false;  
}
```



SymbolTableBuilder Code — SimpleName

```
@Override public boolean visit(SimpleName node) {
    String varName = node.getIdentifier();
    ASTNode parent = node.getParent();

    // Note that we have to make sure that at this point, the node
    // only corresponds to a variable reference
    if (parent instanceof Expression || parent instanceof Statement) {
        if (nameMap.containsKey(varName)) {
            result.put(node, nameMap.get(varName));
        } else {
            throw new Error(node, "Cannot resolve symbol '" + varName
                             + "' in:\n" + parent);
        }
    }
    return false;
}
```



Testing Strategy

- The testing strategy for the symbol tables involves an extension of the pretty printer.
- A textual representation of the symbol table is printed once for every scope area.
- These tables are then compared to a corresponding manual construction for a sufficient collection of programs.
- Furthermore, every error message should be provoked by some test program.