

Static Single Assignment (SSA)

*The content of this lecture is adapted from the lectures of
Gennady Pekhimenko, Todd Mowry and Phillip Gibbons*

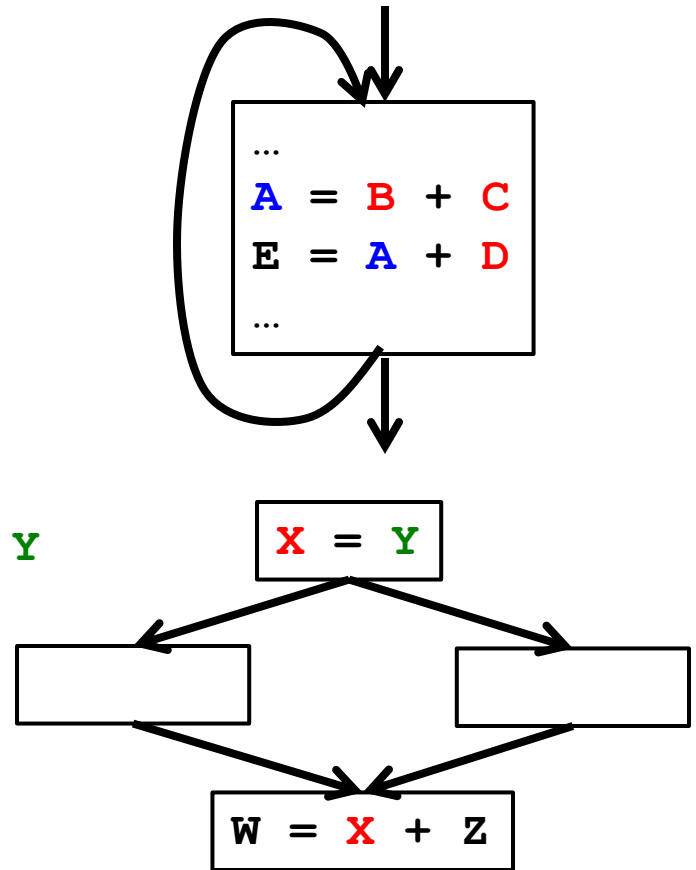
Where Is a Variable Defined or Used?

- Example: Loop-Invariant Code Motion

- Are **B**, **C**, and **D** only defined outside the loop?
- Other definitions of **A** inside the loop?
- Uses of **A** inside the loop?

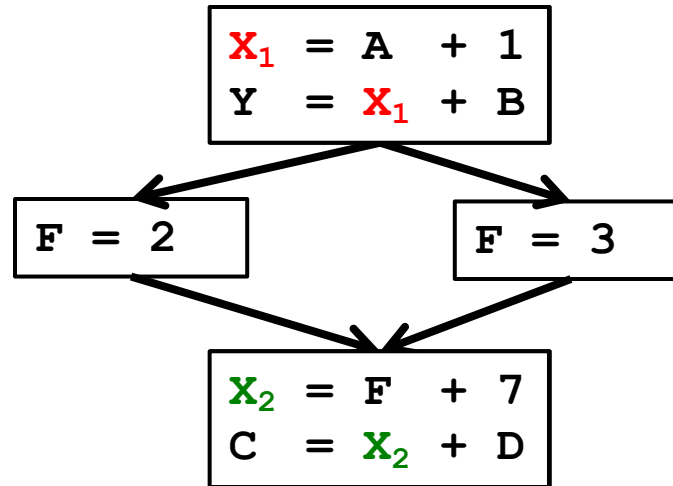
- Example: Copy Propagation

- For a given use of **X**:
 - Are all reaching definitions of **X**:
 - copies from same variable: e.g., **X** = **Y**
 - Where **Y** is not redefined since that copy?
- If so, substitute use of **X** with use of **Y**



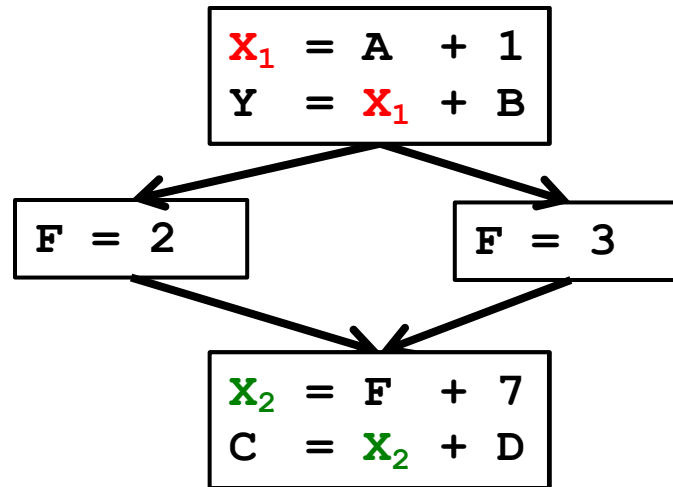
- It would be nice if we could *traverse directly* between related uses and def's
 - this would enable a form of *sparse* code analysis (skip over “don’t care” cases)

Appearances of Same Variable Name May Be Unrelated



- The values in reused storage locations may be provably independent
 - in which case the compiler can optimize them as separate values
- Compiler could use renaming to make these different versions more explicit

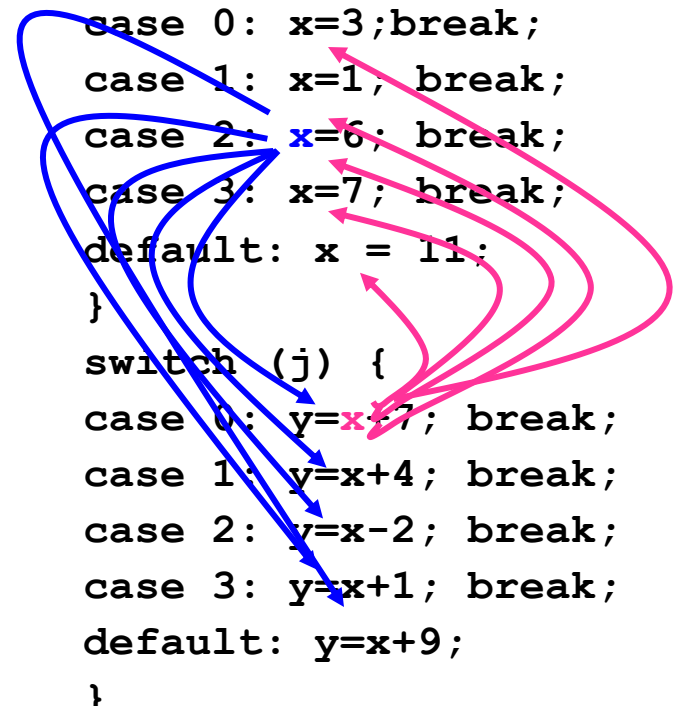
Definition-Use and Use-Definition Chains



- Use-Definition (UD) Chains:
 - for a given definition of a variable X , what are all of its uses?
- Definition-Use (DU) Chains:
 - for a given use of a variable X , what are all of the reaching definitions of X ?

DU and UD Chains Can Be Expensive

```
foo(int i, int j) {  
    ...  
    switch (i) {  
    case 0: x=3; break;  
    case 1: x=1, break;  
    case 2: x=6; break;  
    case 3: x=7; break;  
    default: x = 11;  
    }  
    switch (j) {  
    case 0: y=x+7; break;  
    case 1: y=x+4; break;  
    case 2: y=x-2; break;  
    case 3: y=x+1; break;  
    default: y=x+9;  
    }  
    ...  
}
```



In general,

N defs

M uses

$\Rightarrow O(NM)$ space and time

One solution: limit each variable to ONE definition site

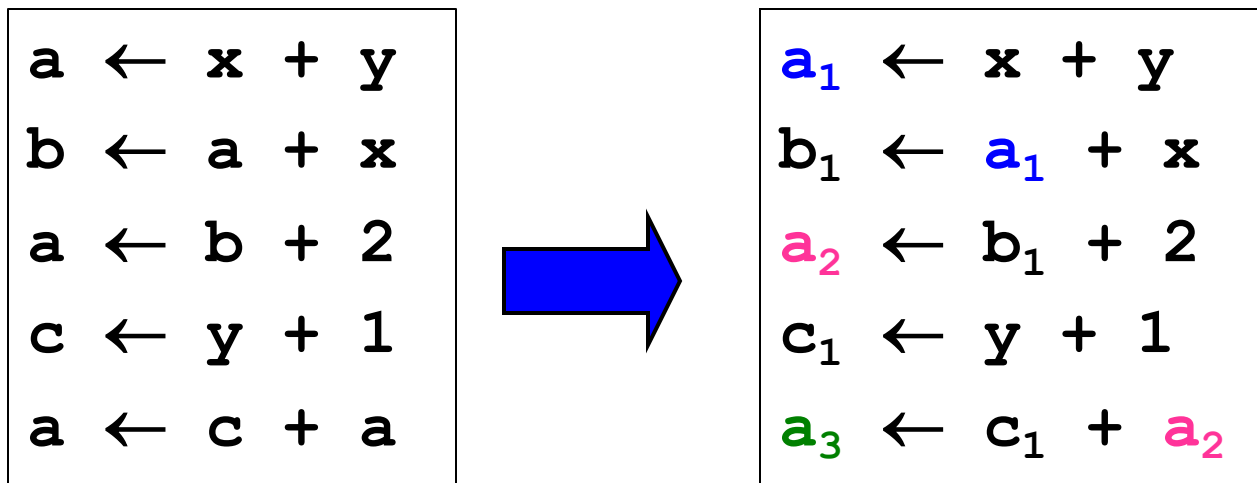
DU and UD Chains Can Be Expensive (2)

```
foo(int i, int j) {  
    ...  
    switch (i) {  
    case 0: x=3; break;  
    case 1: x=1; break;  
    case 2: x=6;  
    case 3: x=7;  
    default: x = 11;  
    }  
    x1 = x; // x1 is one of the above x's  
    switch (j) {  
    case 0: y=x1+7;  
    case 1: y=x1+4;  
    case 2: y=x1-2;  
    case 3: y=x1+1;  
    default: y=x1+9;  
    }  
    ...  
}
```

One solution: limit each variable to ONE definition site

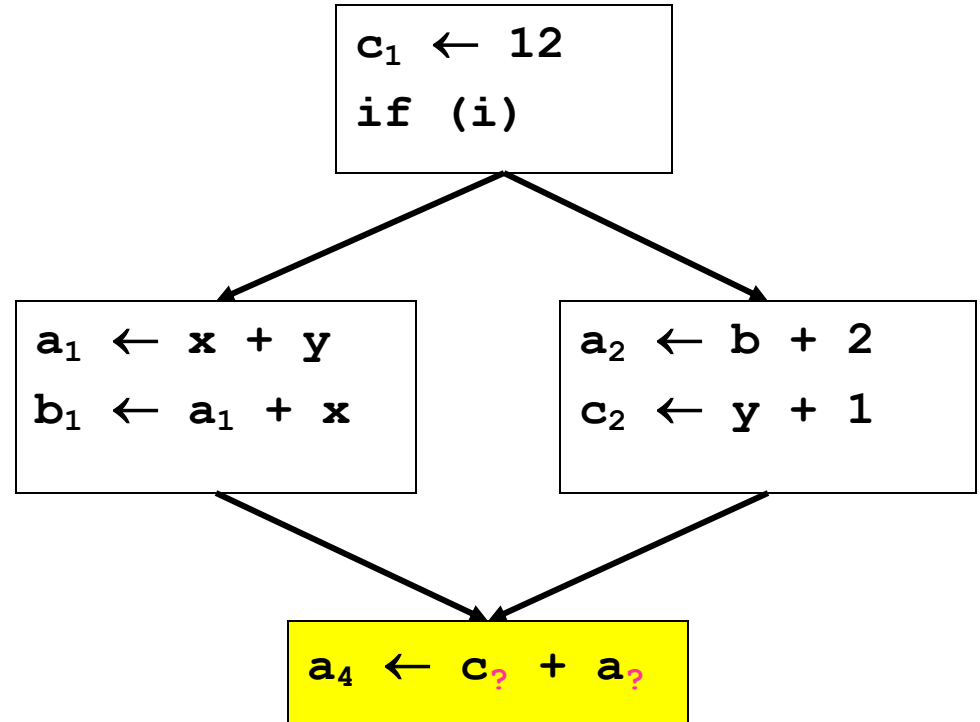
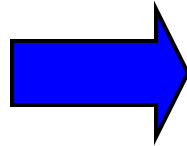
Static Single Assignment (SSA)

- Static single assignment is an IR where **every variable is assigned a value at most once** in the program text
- Easy for a basic block (reminiscent of Value Numbering):
 - Visit each instruction in program order:
 - LHS: **assign** to a **fresh version** of the variable
 - RHS: **use** the **most recent version** of each variable



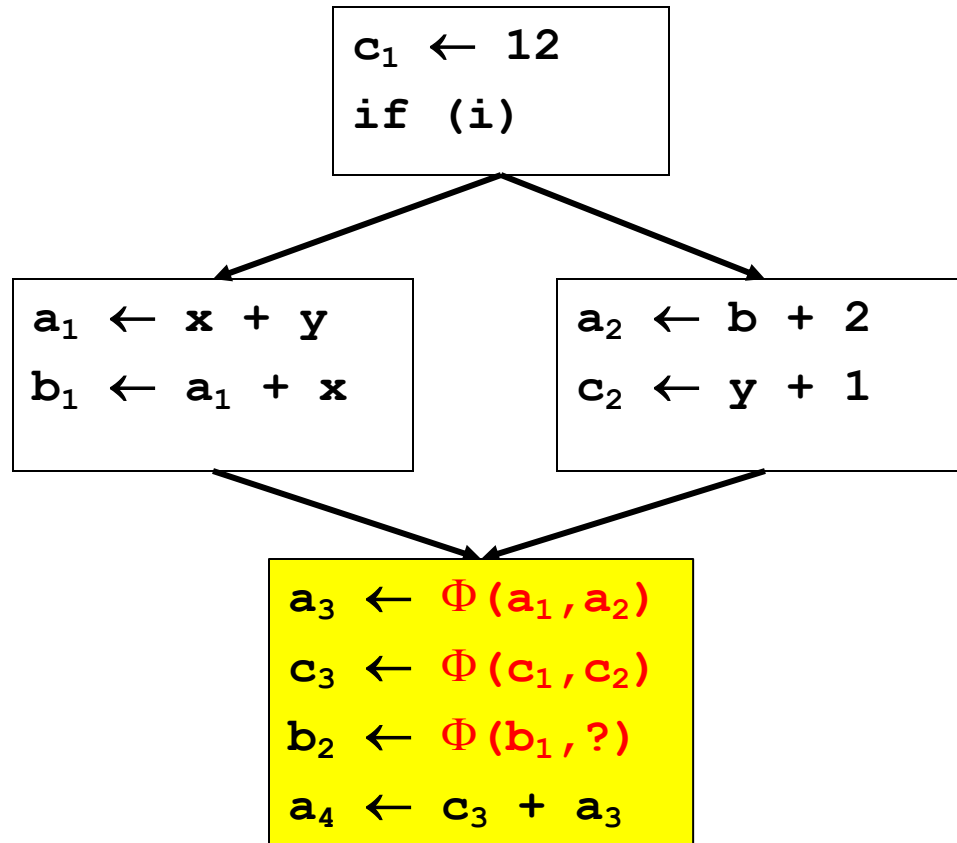
What about Joins in the CFG?

```
c ← 12
if (i) {
  a ← x + y
  b ← a + x
} else {
  a ← b + 2
  c ← y + 1
}
a ← c + a
```



→ Use a notational fiction: a Φ function

Merging at Joins: the Φ function



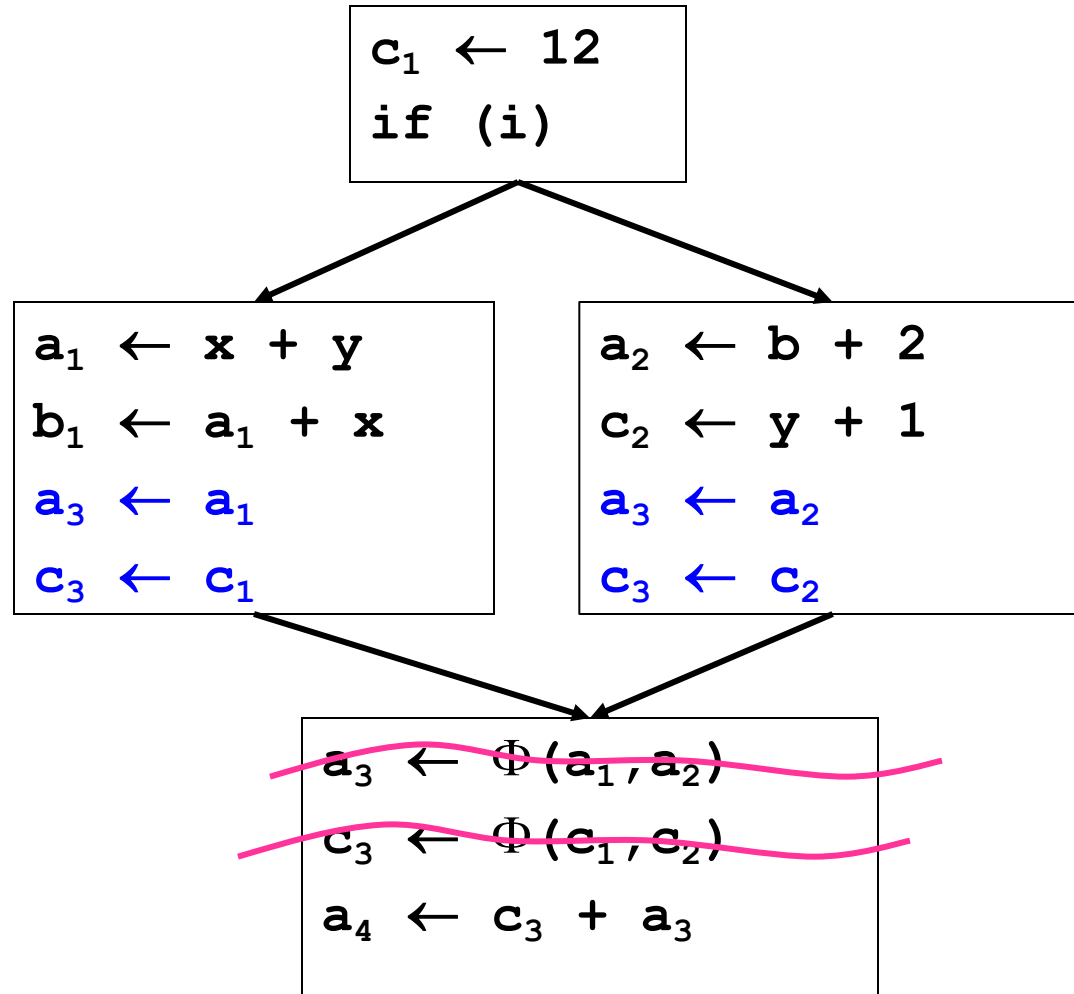
The Φ function

- Φ merges multiple definitions along multiple control paths into a single definition.
- At a basic block with p predecessors, there are p arguments to the Φ function.

$$x_{\text{new}} \leftarrow \Phi(x_1, x_2, \dots, x_p)$$

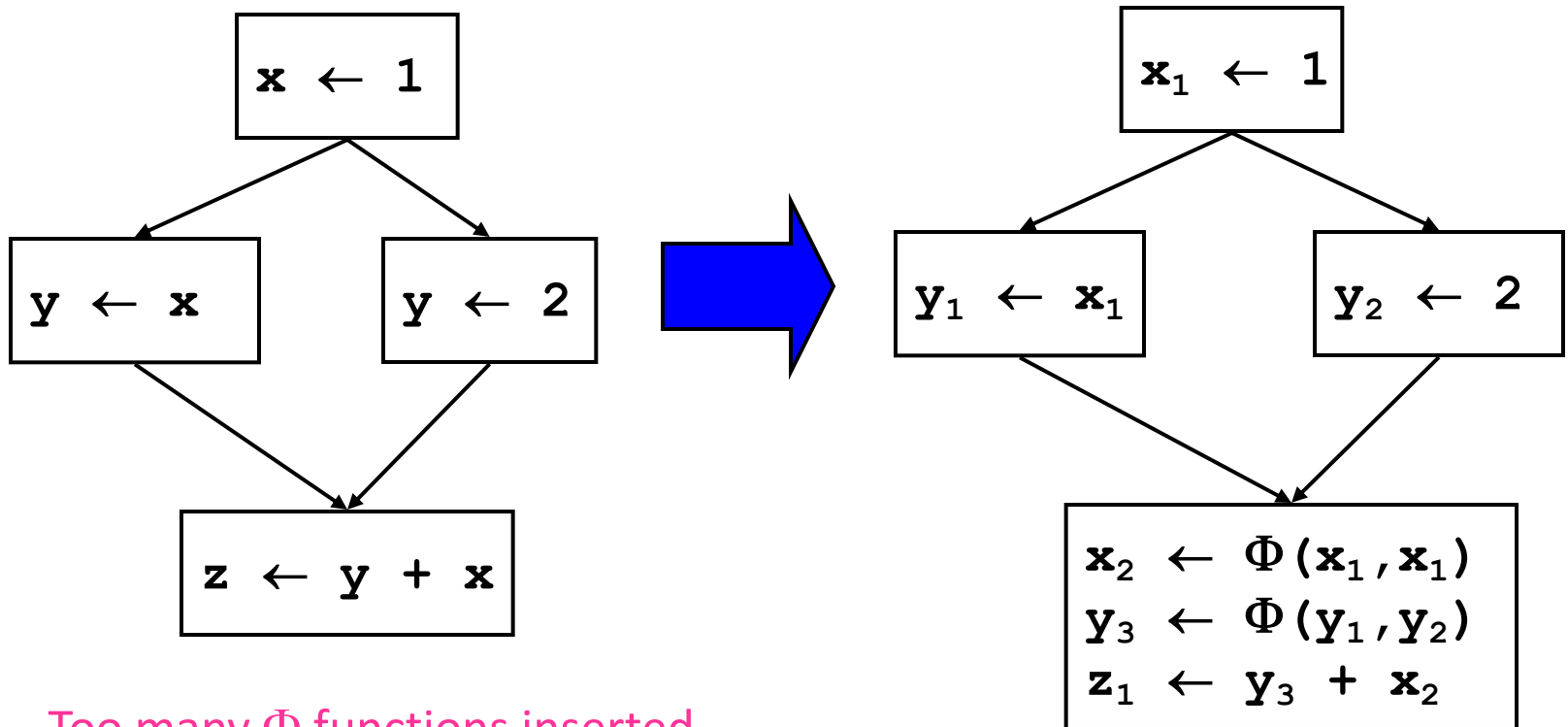
- How do we choose which x_i to use as value of Φ ?
 - For static analysis we just use \perp
 - For code gen, add moves on each incoming edge

“Implementing” Φ



Trivial SSA

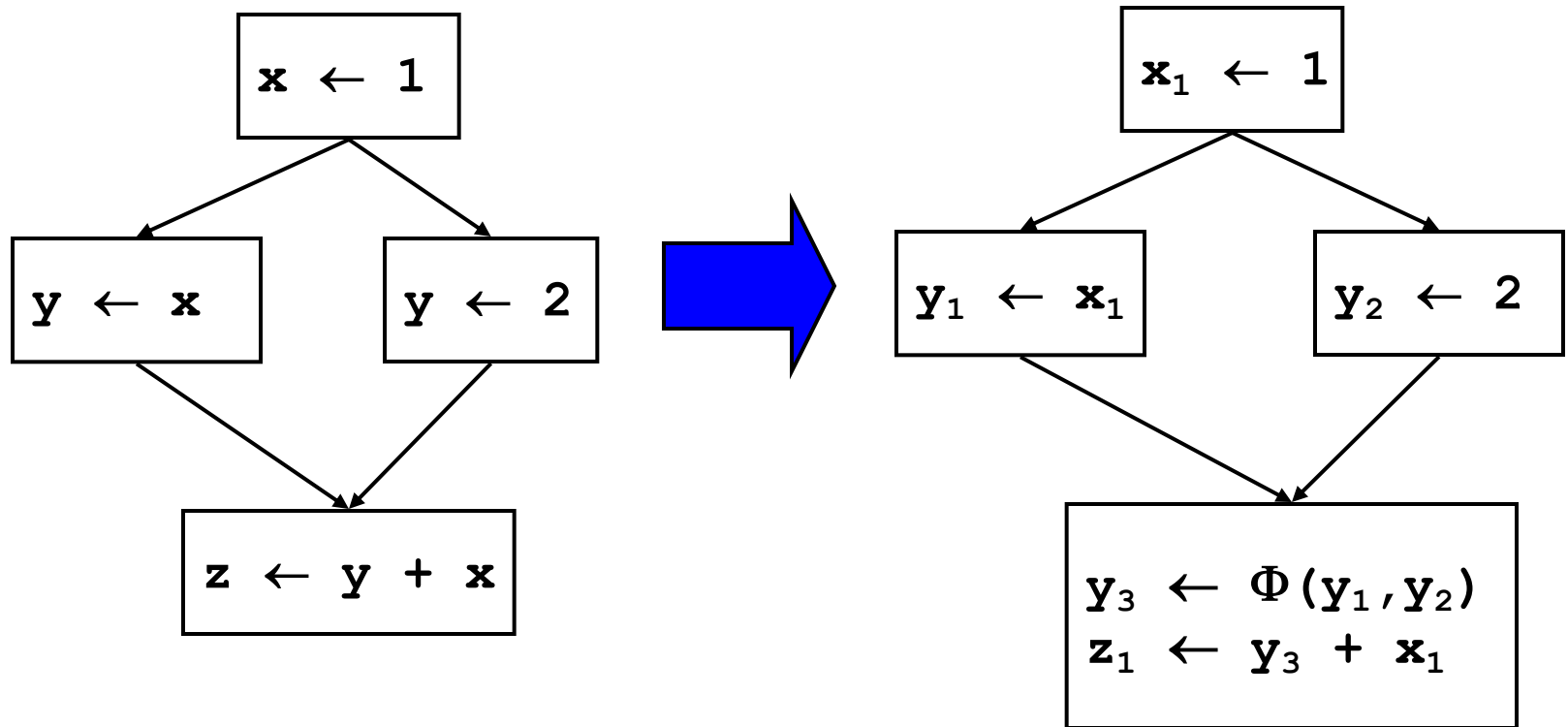
- Each assignment generates a fresh variable.
- At each join point insert Φ functions for **all live variables**.



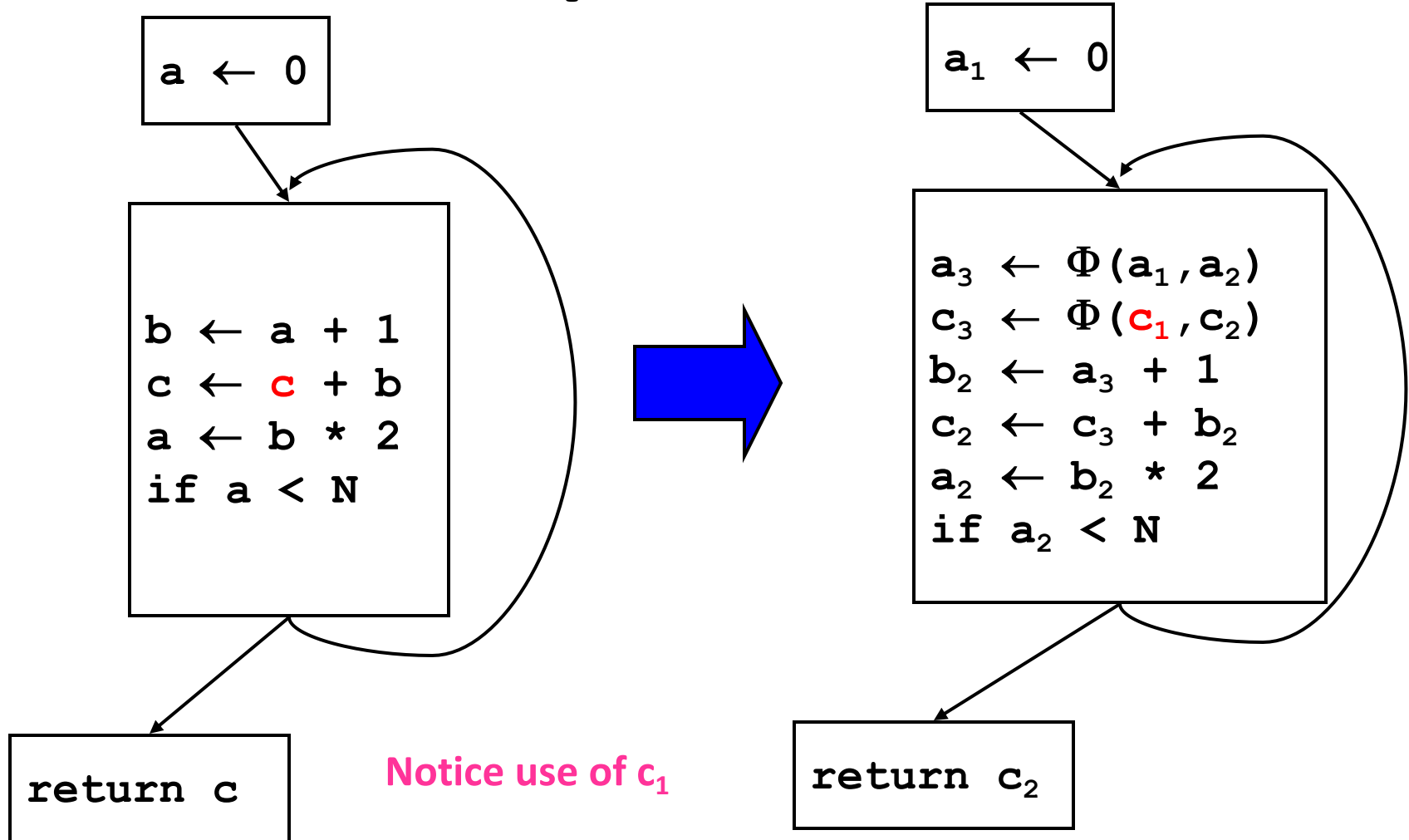
Too many Φ functions inserted.

Minimal SSA

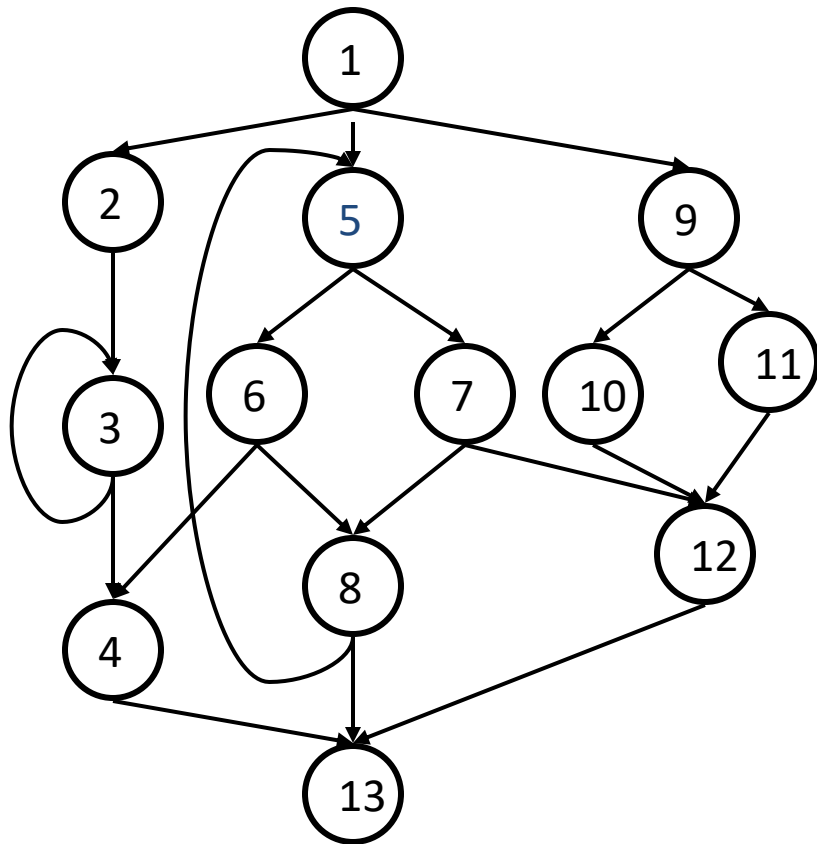
- Each assignment generates a fresh variable.
- At each join point insert Φ functions for **all live variables** with **multiple outstanding defs.**



Another Example



When Do We Insert Φ ?

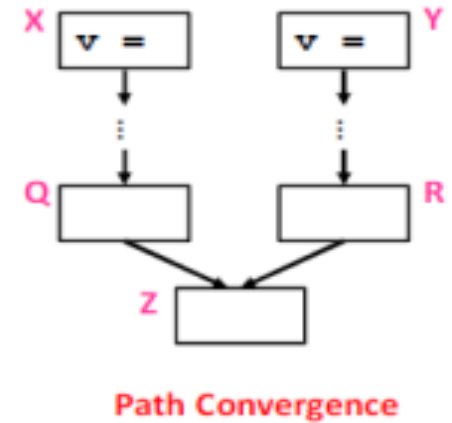


CFG

If there is a def of **a** in block **5**,
which nodes need a $\Phi()$?

When do we insert Φ ?

- We insert a Φ function for variable **A** in block **Z** iff:
 - **A** was defined more than once before
 - (i.e., **A** defined in **X** and **Y** AND $X \neq Y$)
 - There exists a non-empty path from **x** to **z**, P_{xz} ,
 - and a non-empty path from **y** to **z**, P_{yz} , s.t.
 - $P_{xz} \cap P_{yz} = \{z\}$
(Z is only common block along paths)
 - $z \notin P_{xq}$ or $z \notin P_{yr}$ where $P_{xz} = P_{xq} \rightarrow z$ and $P_{yz} = P_{yr} \rightarrow z$
(at least one path reaches Z for first time)
- Entry block contains an implicit def of all vars
- Note: $v = \Phi(\dots)$ is a def of **v**



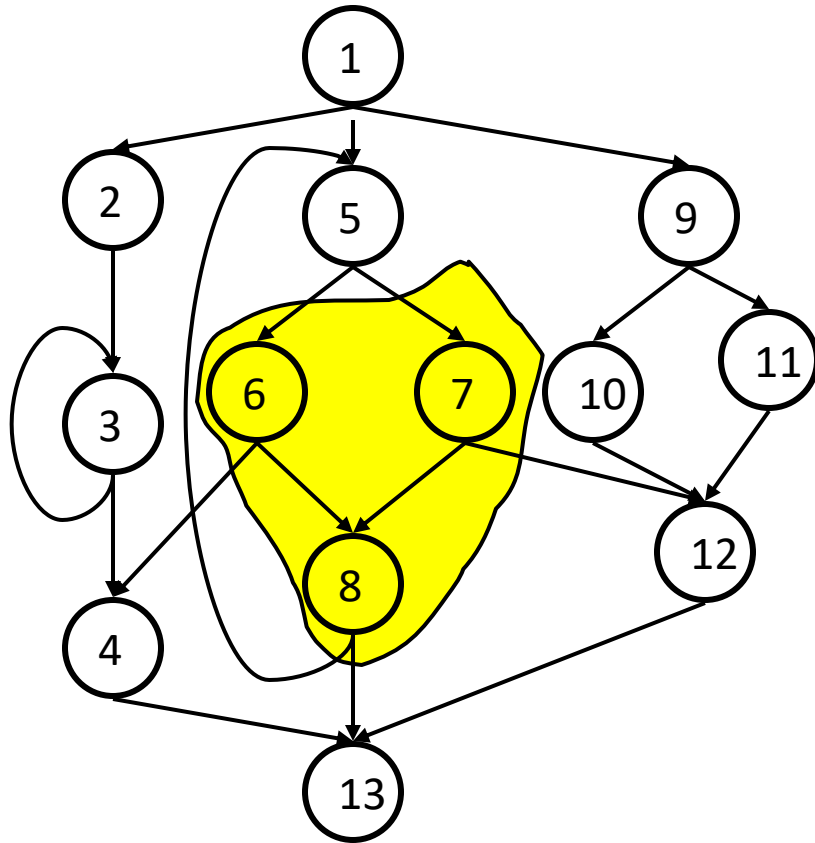
Control flow dominators

- Given two nodes, n and m , in a CFG
- n dominates m (aka n dom m)
 - if every path from the entry to m passes through n
- Every node trivially dominates itself
 - n strictly dominates m (aka n sdom m)
 - If n dom m and $n \neq m$
- Dominance relation induces a tree-shaped ordering on CFG nodes

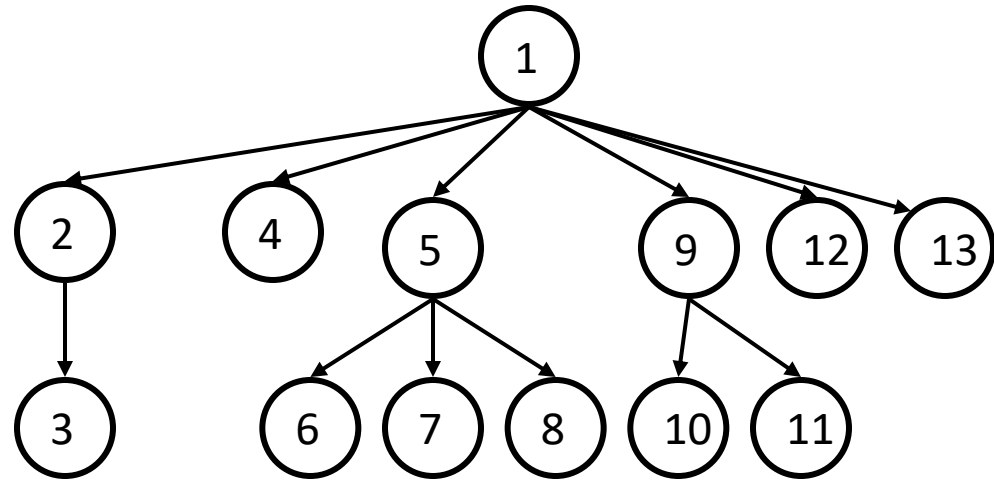
Dominance Property of SSA

- In SSA, definitions dominate uses.
 - If x_i is used in $x \leftarrow \Phi(\dots, x_i, \dots)$, then $\text{BB}(x_i)$ dominates i^{th} predecessor of $\text{BB}(\text{PHI})$
 - If x is used in $y \leftarrow \dots x \dots$, then $\text{BB}(x)$ dominates $\text{BB}(y)$
- We can use this for an efficient algorithm to convert to SSA

Dominance



CFG

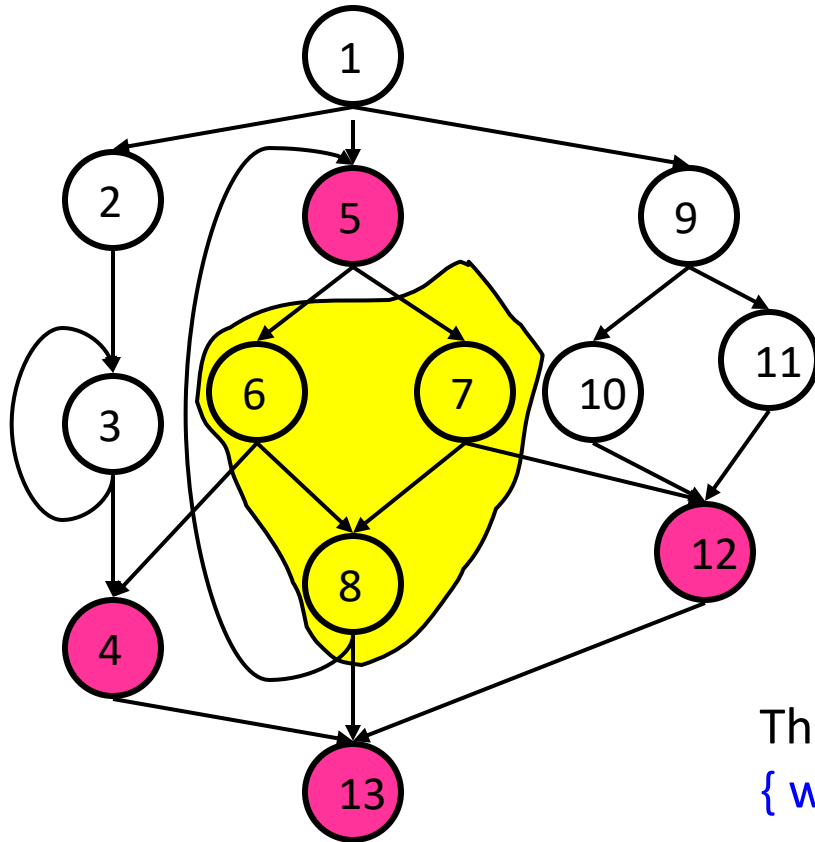


D-Tree

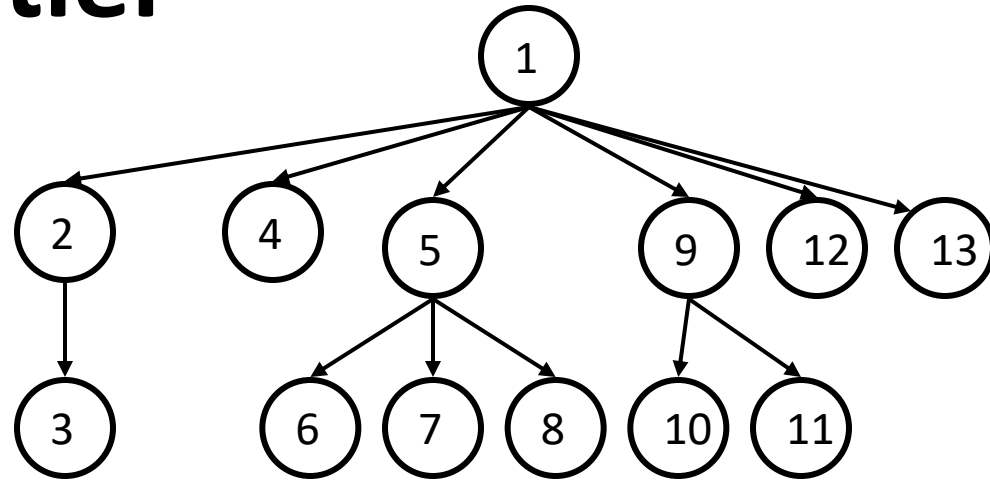
If there is a def of **a** in block **5**,
which nodes need a $\Phi()$?

x strictly dominates **w** (**x** **sdom** **w**) iff **x** **dom** **w** AND **x** \neq **w**

Dominance Frontier



CFG

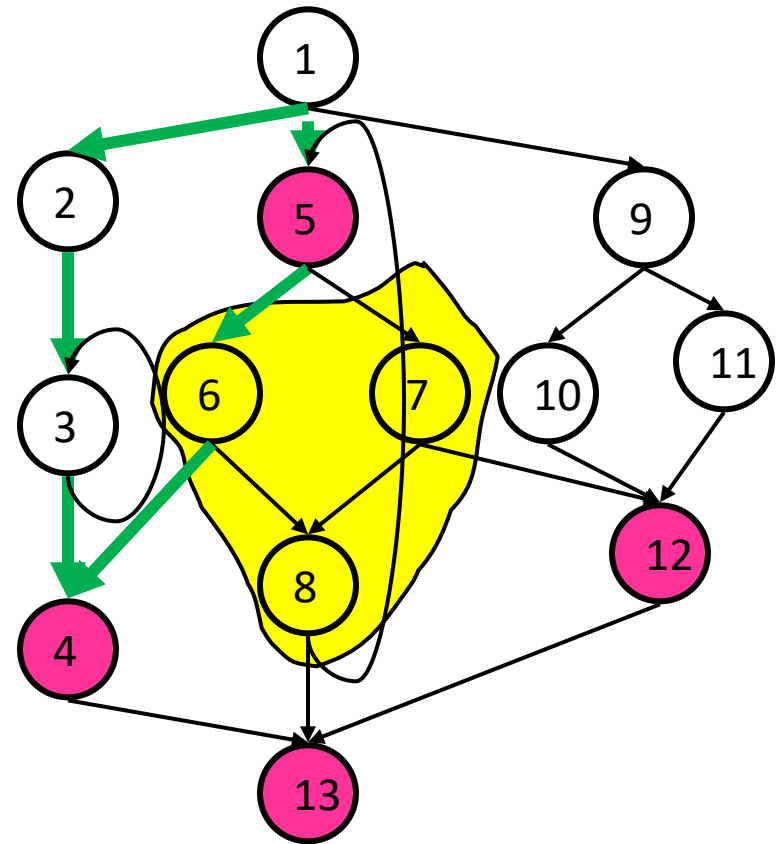
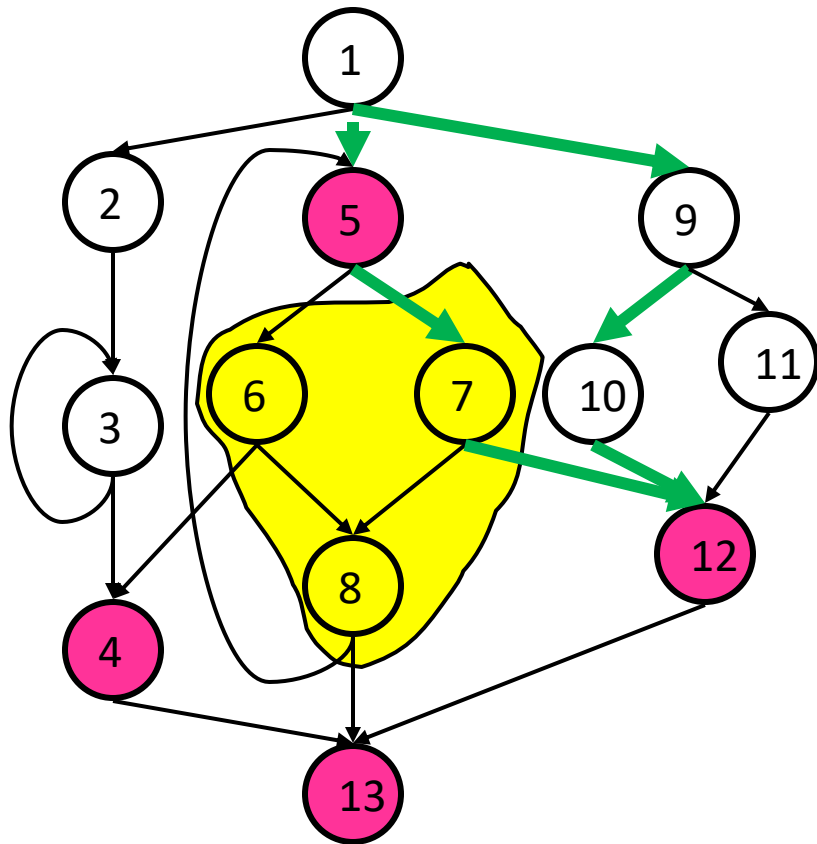


D-Tree

The **Dominance Frontier** of a node $x = \{ w \mid x \text{ dom pred}(w) \text{ AND } \neg(x \text{ sdom } w) \}$

x strictly dominates w ($x \text{ sdom } w$) iff $x \text{ dom } w$ AND $x \neq w$

Dominance Frontier and Path Convergence



If there is a def of **a** in block **5**,
nodes in $DF(5)$ need a $\Phi()$ for **a**

Using Dominance Frontier to Compute SSA

- place all $\Phi()$
- Rename all variables

Using Dominance Frontier to Place $\Phi()$

- Gather all the **defsites**, i.e., assignments, of every variable
- Then, for **every variable**
 - foreach **defsite**
 - foreach **node in DominanceFrontier(defsite)**
 - if we haven't put $\Phi()$ in node, then **put one in**
 - if this node didn't define the variable before, then **add this node to the defsites**
- This essentially computes the **Iterated Dominance Frontier** on the fly, **inserting the minimal number of $\Phi()$ necessary**

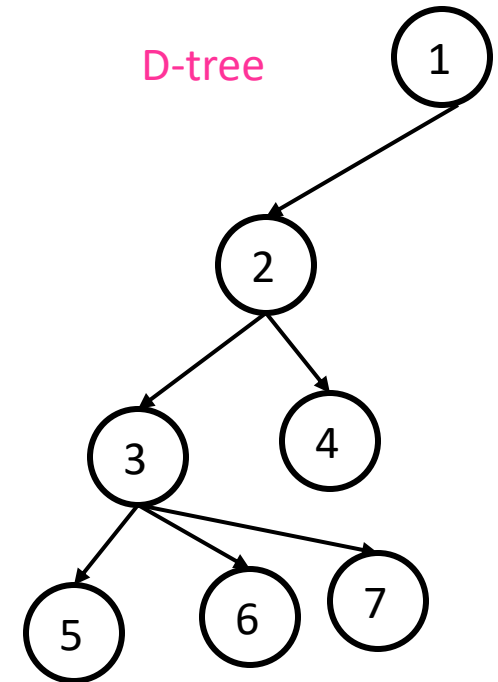
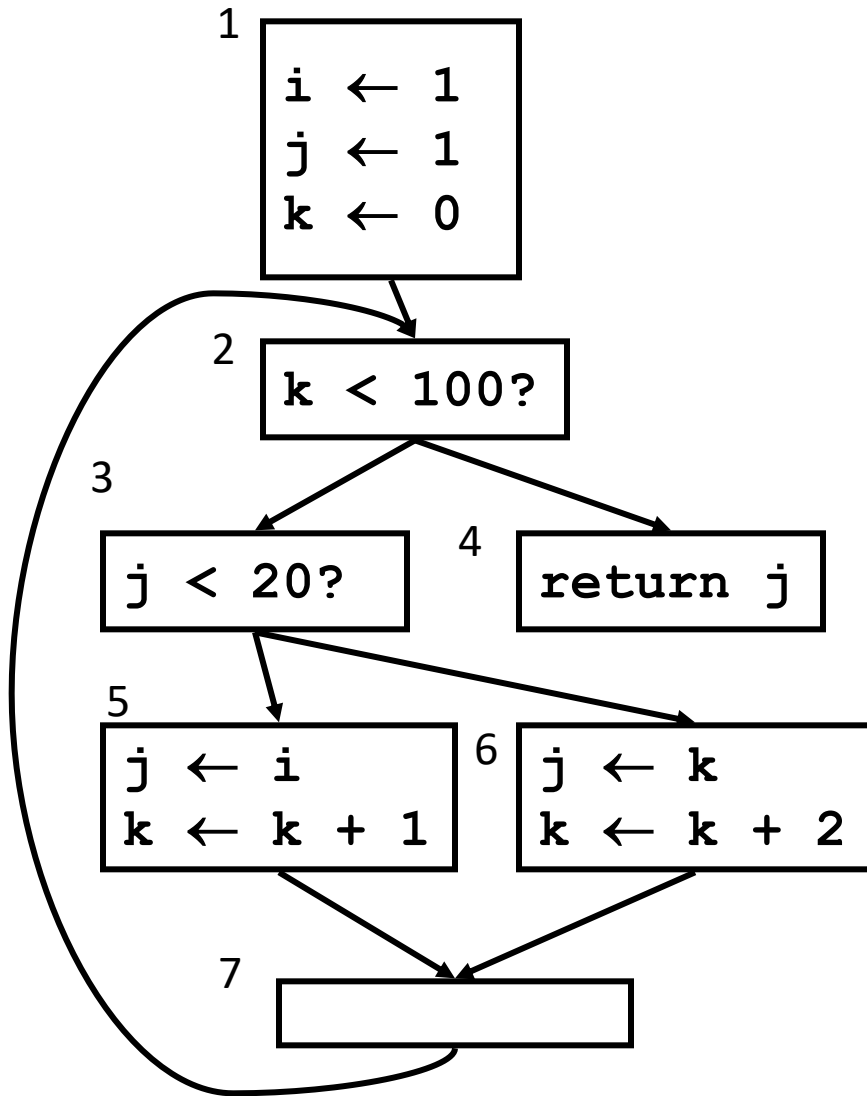
Using Dominance Frontier to Place $\Phi()$

```
foreach node n {
    foreach variable v defined in n {
        varsdef[n]  $\cup$ = {v}
        defsites[v]  $\cup$ = {n}
    }
}
foreach variable v {
    W = defsites[v]
    while W not empty {
        n = remove node from W
        foreach y in DF[n] {
            if y  $\notin$  PHI[v] {
                insert " $v \leftarrow \Phi(v,v,...)$ " at top of y
                PHI[v] = PHI[v]  $\cup$  {y}
                if v  $\notin$  varsdef[y]: W = W  $\cup$  {y}
            }
        }
    }
}
```

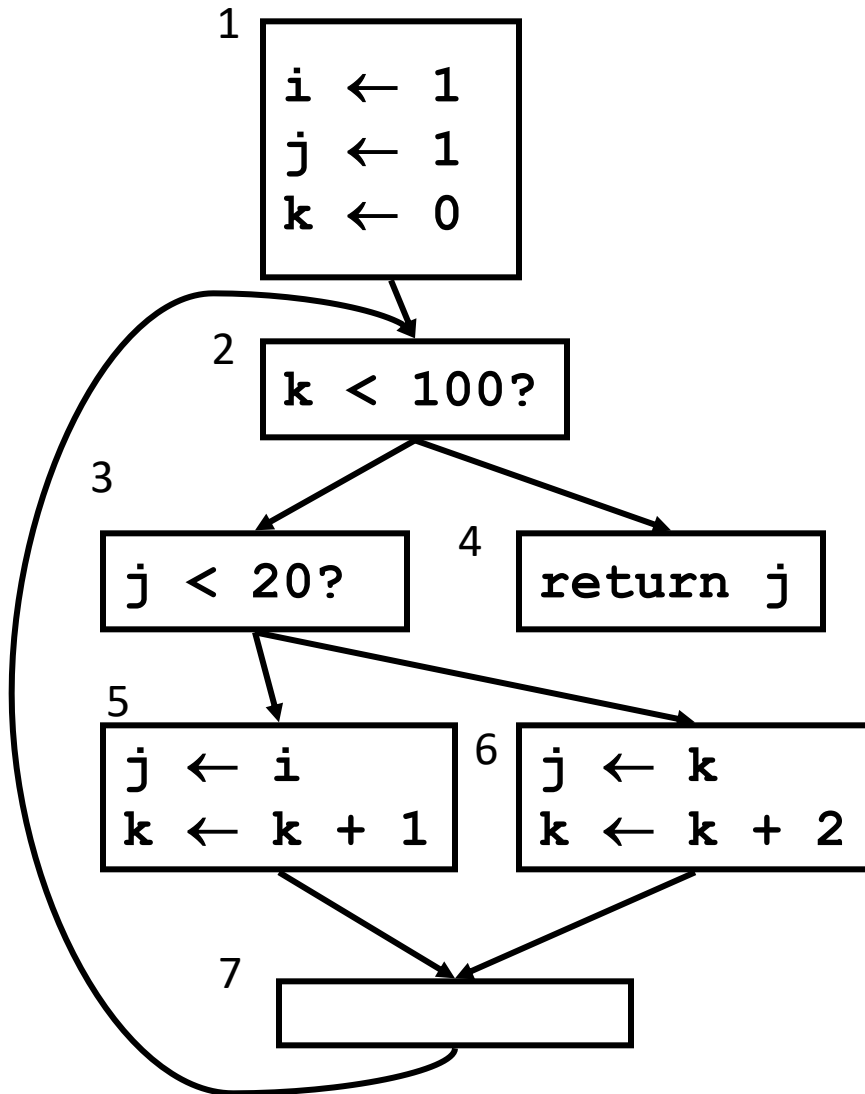

Renaming Variables

- Algorithm:
 - Walk the D-tree, renaming variables as you go
 - Replace uses with most recent renamed def
- For straight-line code this is easy
- What if there are branches and joins?
 - use the **closest def such that the def is above the use in the D-tree**
- Implementation:
 - for each var: **rename** (v)
 - **rename**(v):
 - replace uses with top of stack
 - at def: push onto stack
 - call **rename**(v) on all children in D-tree
 - if D-tree leaf, replace uses of CFG successors
 - for each def in this block pop from stack

Compute Dominance Tree

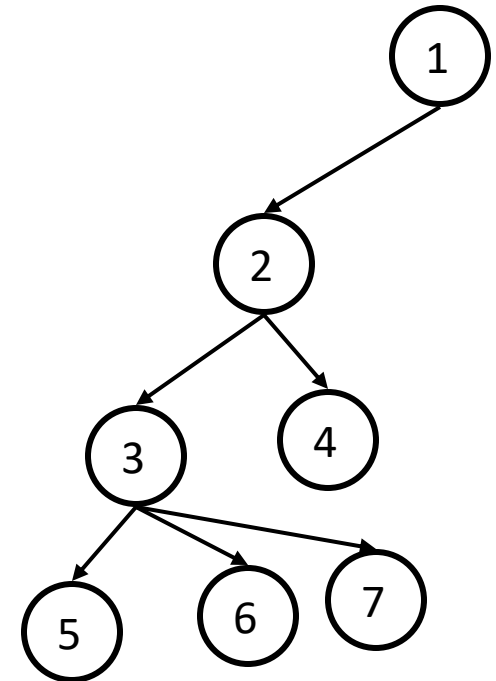


Compute Dominance Frontiers

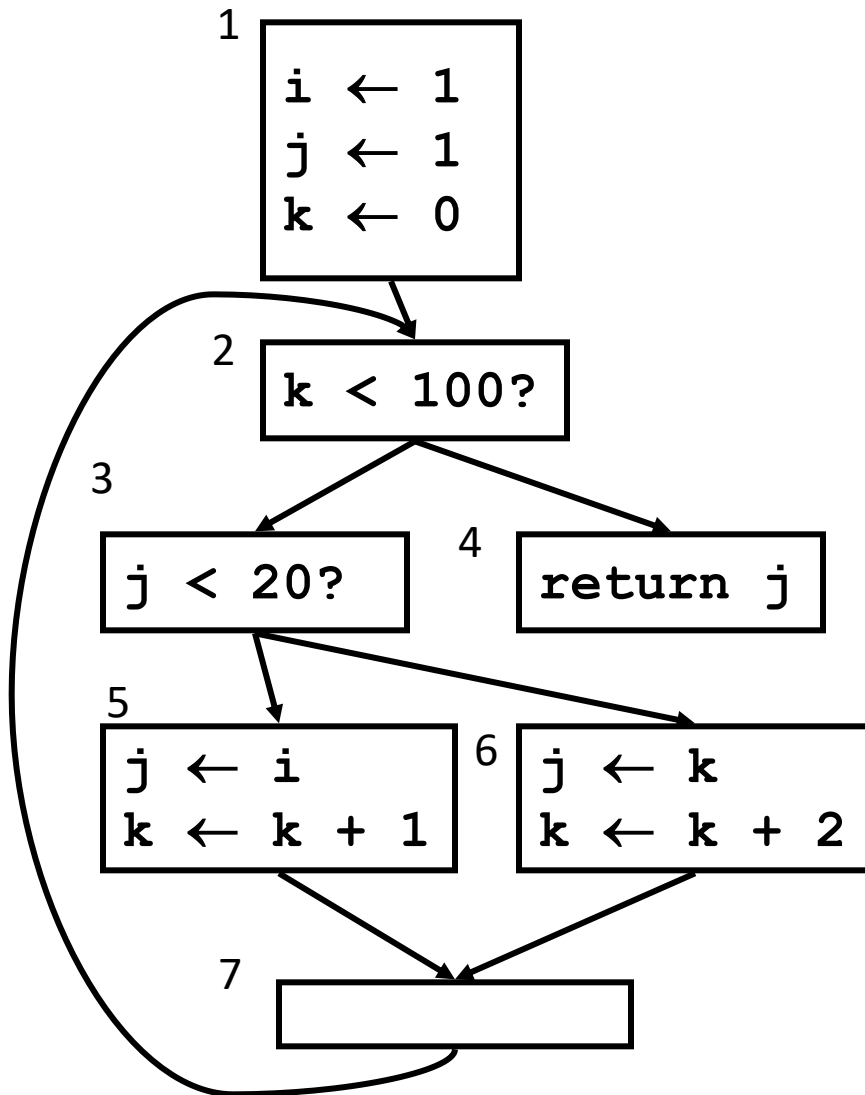


DFs

1	{}
2	{2}
3	{2}
4	{}
5	{7}
6	{7}
7	{2}



Insert $\Phi()$



DFs

1	{}
2	{2}
3	{2}
4	{}
5	{7}
6	{7}
7	{2}

varsdef[n]

1	{i,j,k}
2	{}
3	{}
4	{}
5	{j,k}
6	{j,k}
7	{}

defsites[v]

i	{1}
j	{1,5,6}
k	{1,5,6}

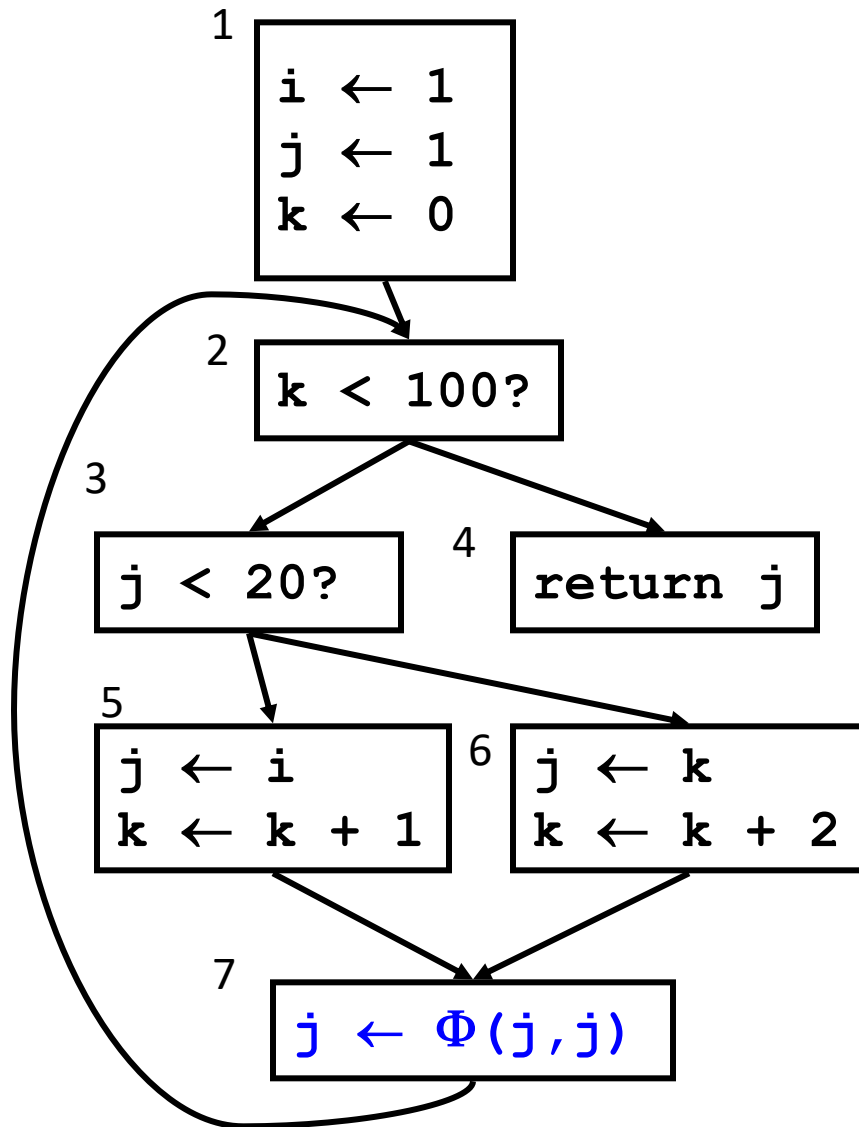
DFs

var i: $W=\{1\}$

var j: $W=\{1,5,6\}$

DF{1} DF{5}

Insert $\Phi()$



DFs

1	{}
2	{2}
3	{2}
4	{}
5	{7}
6	{7}
7	{2}

varsdef[n]

1	{i,j,k}
2	{}
3	{}
4	{}
5	{j,k}
6	{j,k}
7	{}

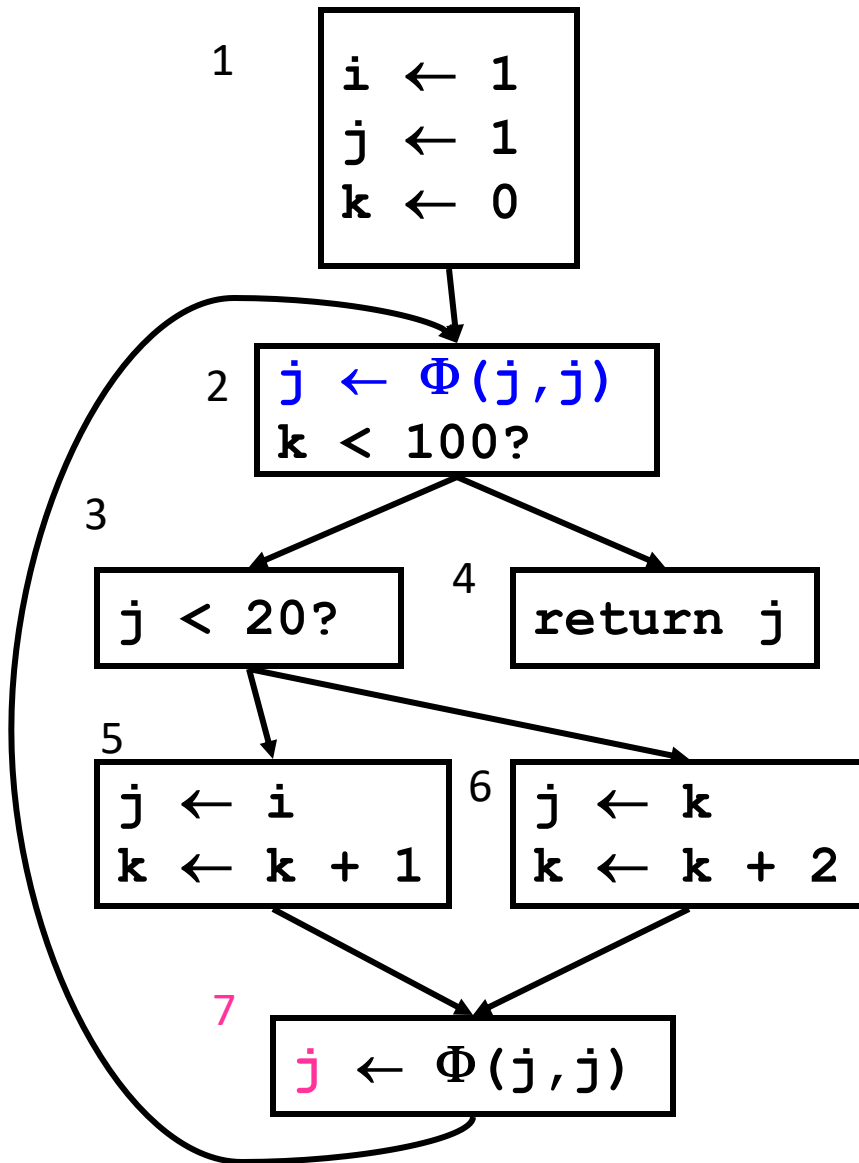
defsites[v]

i	{1}
j	{1,5,6}
k	{1,5,6}

DFs

var j: W={1,5,6}

DF{1} DF{5}



DFs

1	{}
2	{2}
3	{2}
4	{}
5	{7}
6	{7}
7	{2}

varsdef[n]

1	{i,j,k}
2	{}
3	{}
4	{}
5	{j,k}
6	{j,k}
7	{}

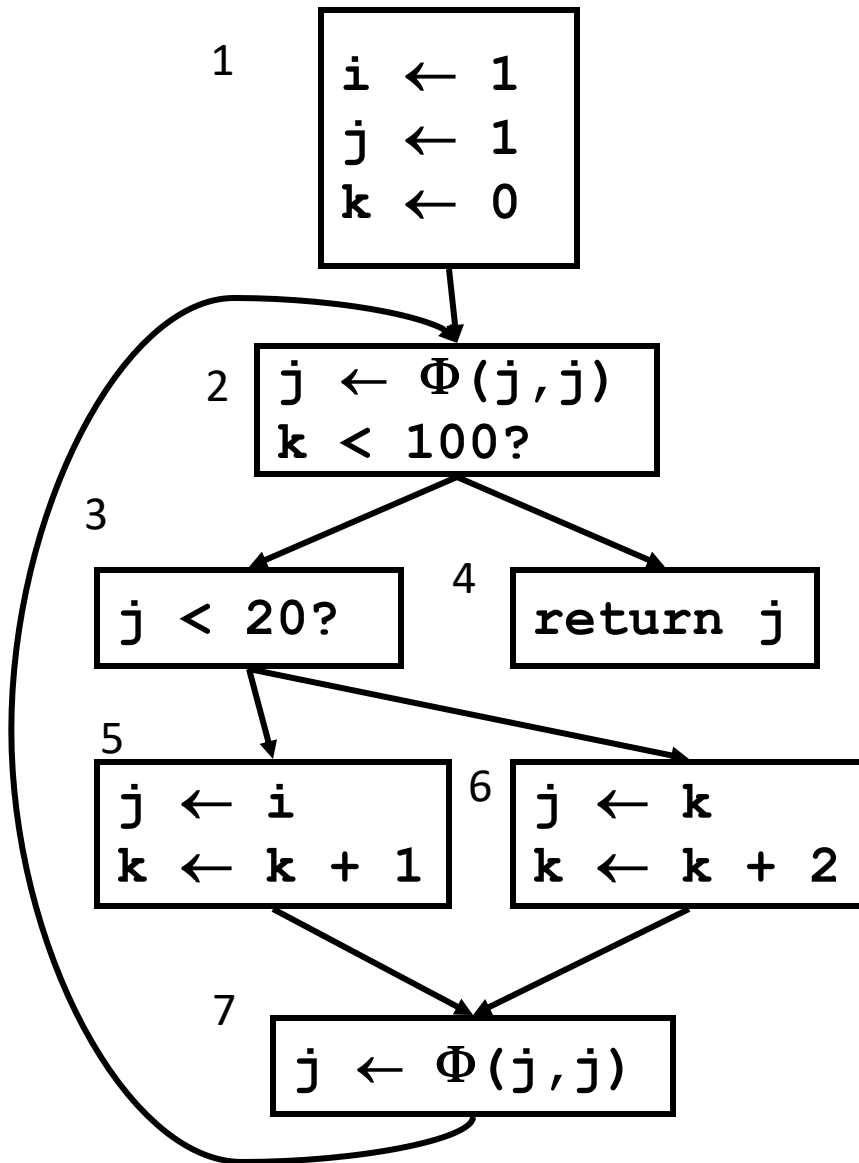
defsites[v]

i	{1}
j	{1,5,6,7}
k	{1,5,6}

DFs

var j: W={1,5,6,7}

DF{1} DF{5} DF{7}



DFs

1	{}
2	{2}
3	{2}
4	{}
5	{7}
6	{7}
7	{2}

varsdef[n]

1	{i,j,k}
2	{}
3	{}
4	{}
5	{j,k}
6	{j,k}
7	{}

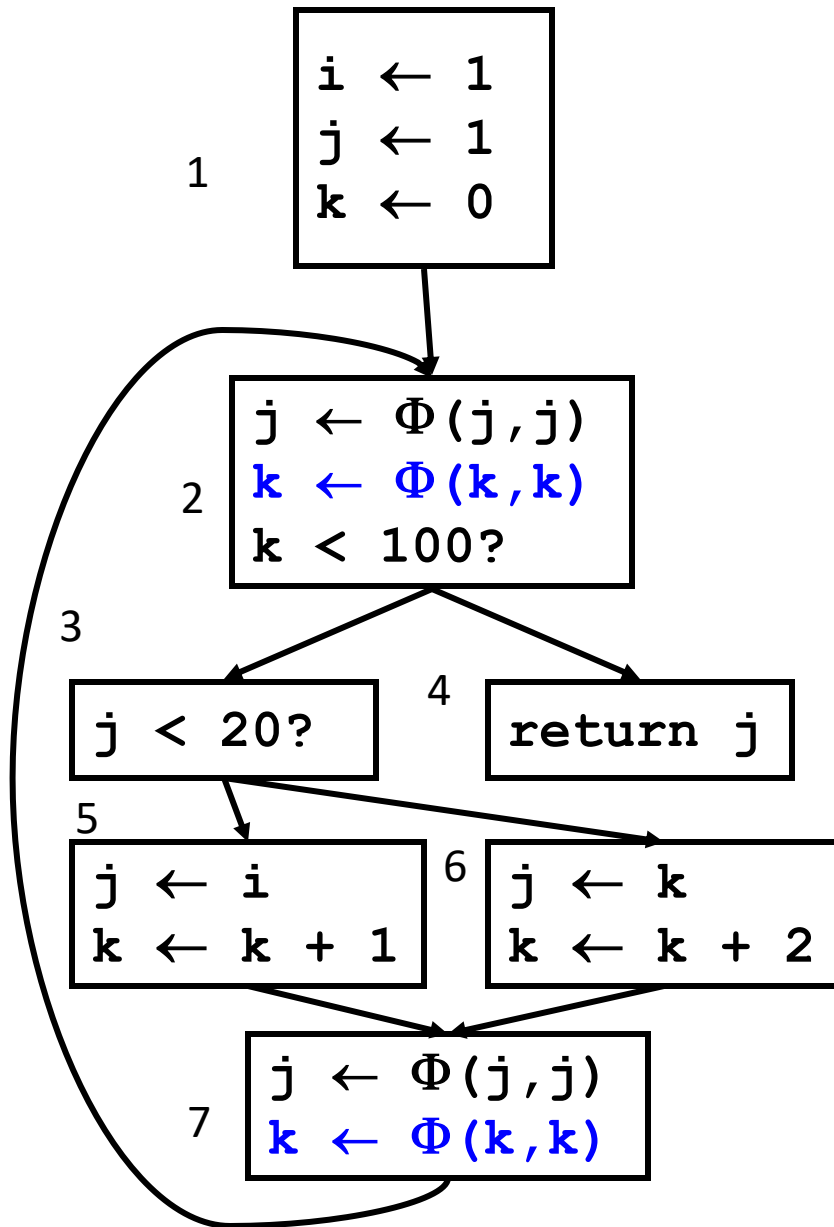
defsites[v]

i	{1}
j	{1,5,6}
k	{1,5,6}

DFs

var j: W={1,5,6,7}

DF{1} DF{5} DF{7} **DF{6}**



DFs

1	{}
2	{2}
3	{2}
4	{}
5	{7}
6	{7}
7	{2}

varsdef[n]

1	{i,j,k}
2	{}
3	{}
4	{}
5	{j,k}
6	{j,k}
7	{}

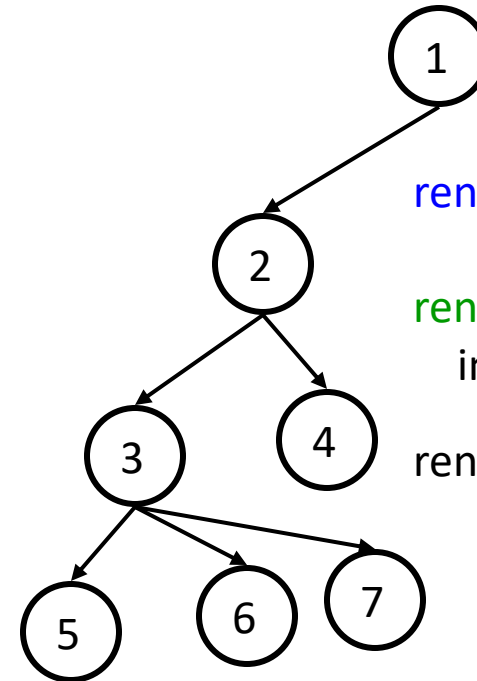
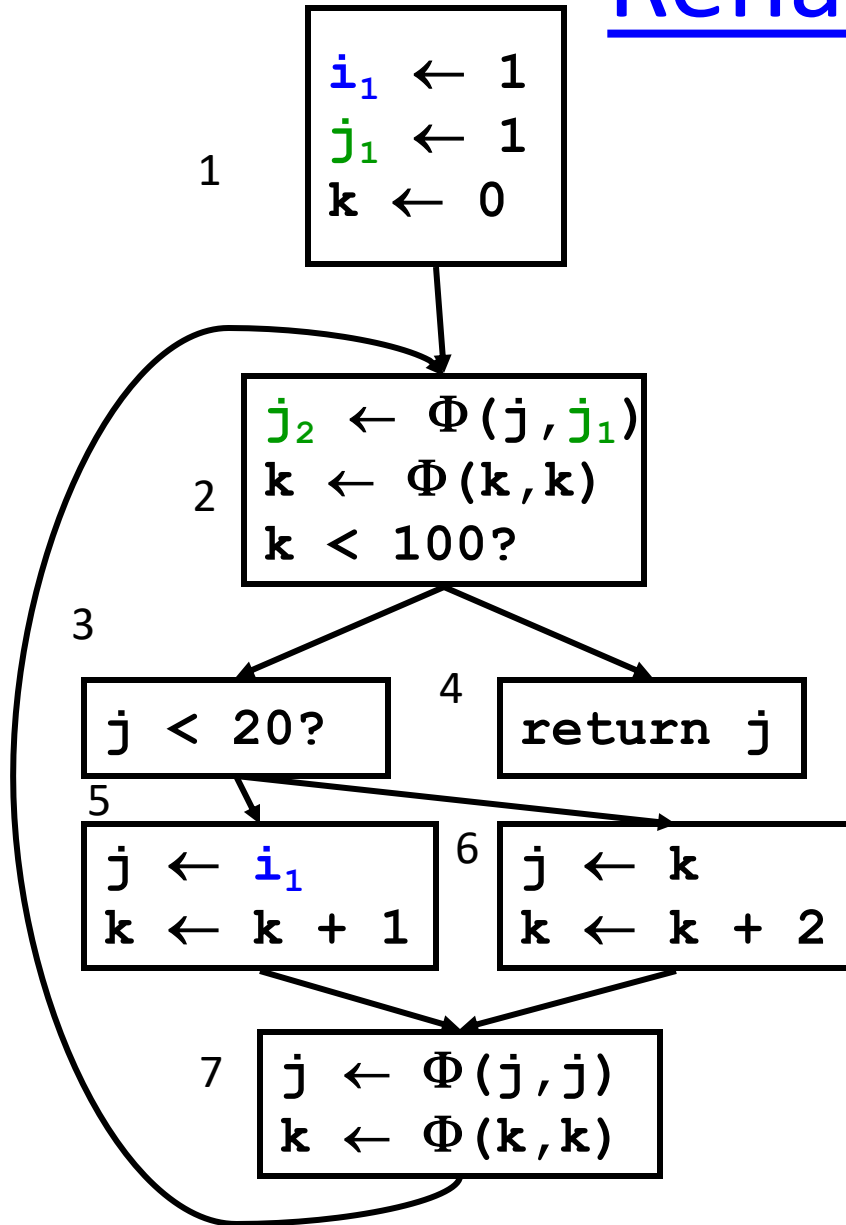
Def sites[v]

i	{1}
j	{1,5,6}
k	{1,5,6}

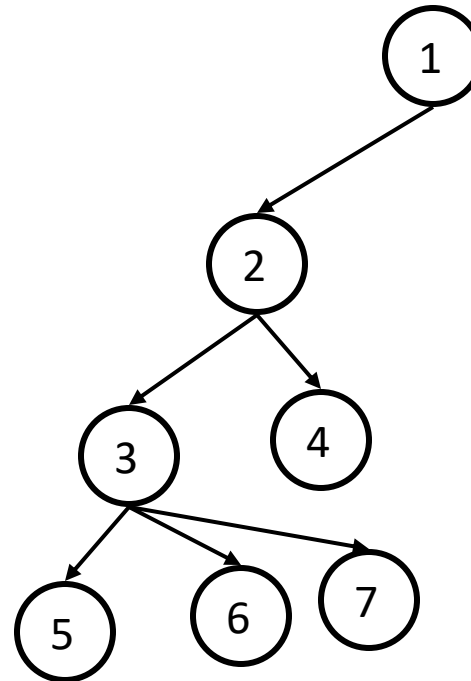
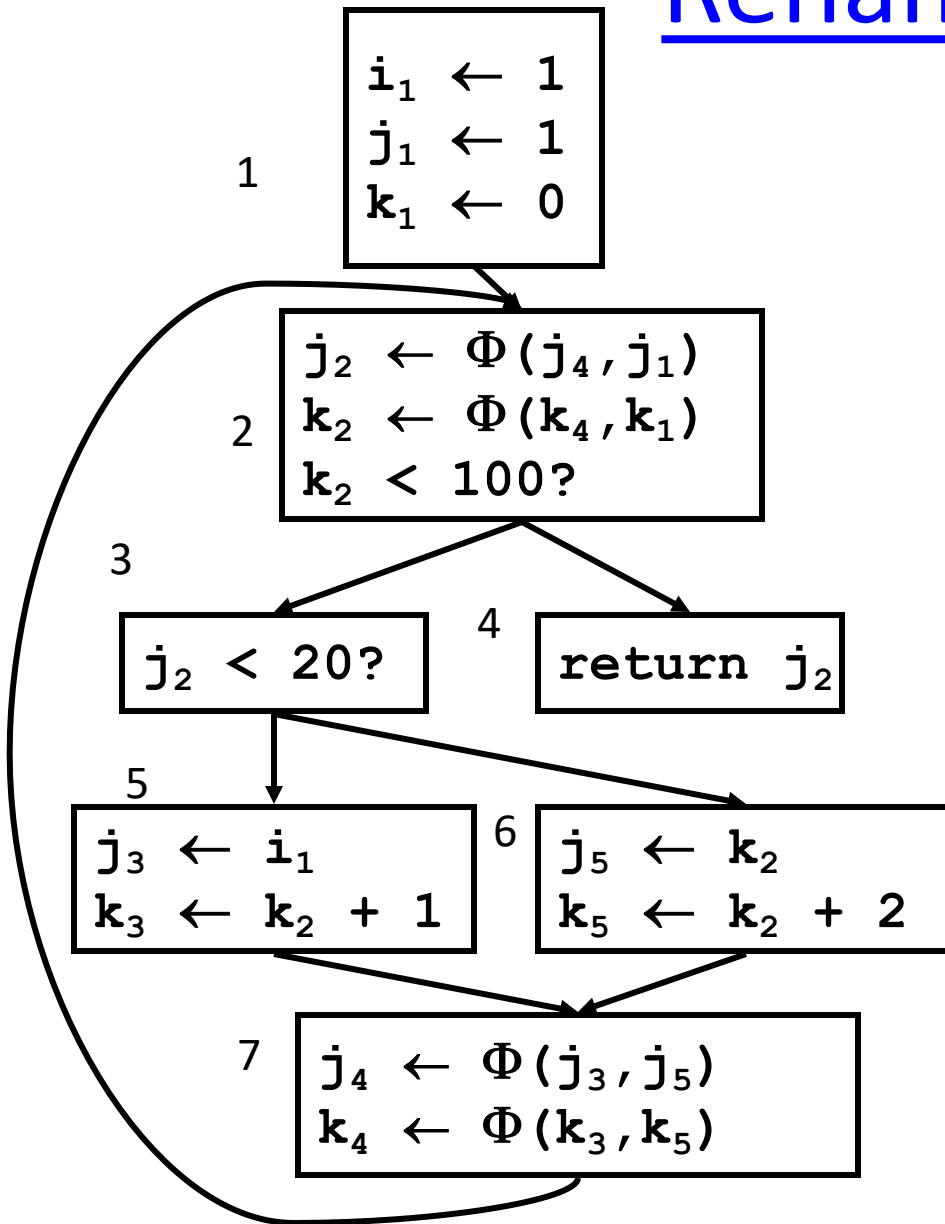
DFs

var **k**: W={1,5,6}

Rename Vars



Rename Vars



SSA Properties

- Only 1 assignment per variable
- Definitions dominate uses

Constant Propagation with SSA

- If “ $v \leftarrow c$ ”, replace all uses of v with c
- If “ $v \leftarrow \Phi(c, c, c)$ ” (each input is the same constant), replace all uses of v with c

```
W ← list of all defs
while !W.isEmpty {
    S ← remove from W
    if ((S has form " $v \leftarrow c$ ") ||
        (S has form " $v \leftarrow \Phi(c, \dots, c)$ ")) then {
        delete S
        foreach stmt U that uses  $v$  {
            replace  $v$  with  $c$  in U
            W.add(U)
        }
    }
}
```

Other Optimizations with SSA

- Copy propagation
 - delete “ $x \leftarrow \Phi(y,y,y)$ ” and replace all x with y
 - delete “ $x \leftarrow y$ ” and replace all x with y
- Constant Folding
 - (Also, constant conditions too!)

Constant Propagation

