ECE 432/532 Programming for Parallel Processors

- In virtually all distributed-memory systems, communication can be *much* more expensive than local computation.
- For example, sending a double from one node to another will take far longer than adding two doubles stored in the local memory of a node.
- The cost of sending a fixed amount of data in multiple messages is much greater than the cost of sending a single message with the same amount of data.

VS

```
if (my_rank == 0)
    MPI_Send(x, 1000, MPI_DOUBLE, 1, 0, comm);
else /* my_rank == 1 */
    MPI_Recv(x, 1000, MPI_DOUBLE, 0, 0, comm, &status);
```

```
if (my_rank == 0)
    MPI_Send(x, 1000, MPI_DOUBLE, 1, 0, comm);
else /* my_rank == 1 */
    MPI_Recv(x, 1000, MPI_DOUBLE, 0, 0, comm, &status);
```

- In MPI, a derived datatype can be used to represent any collection of data items in memory by storing both the types of the items and their relative locations in memory.
- The idea is that if a function that sends data knows the types and the relative locations in memory of a collection of data items, it can collect the items from memory before they are sent.
- Similarly, a function that receives data can distribute the items into their correct destinations in memory when they're received.

- Example: in the trapezoidal rule program we needed to call MPI_Bcast three times:
 - once for the left endpoint a
 - once for the right endpoint b
 - and once for the number of trapezoids *n*.
- As an alternative, we could build a single derived datatype that consists of two doubles and one int.
 - need one call to MPI_Bcast
- On process 0, a,b, and n will be sent with the one call, while on the other processes, the values will be received with the call.

- Formally, a derived datatype consists of a sequence of basic MPI datatypes together with a displacement for each of the datatypes.
- In the trapezoidal rule example, suppose that on process 0 the variables a, b, and n are stored in memory locations with the following addresses

Variable	Address		
a	24		
b	40		
n	48		

 $\{(MPI_DOUBLE, 0), (MPI_DOUBLE, 16), (MPI_INT, 24)\}$

MPI_Type_create_struct

 Builds a derived datatype that consists of individual elements that have different basic types.

MPI_Type_create_struct

- The argument count is the number of elements in the datatype → three (for trapezoidal rule)
- The array_of_block_lengths, allows for the possibility that the individual data items might be arrays or subarrays
 - first element is an array containing five elements → array_of_blocklengths[0] = 5;
 - In our case \rightarrow int array_of_blocklengths[3] = {1, 1, 1};

MPI_Type_create_struct

 The third argument, array_of_displacements specifies the displacements, in bytes, from the start of the message:

```
array_of_displacements[] = {0, 16, 24};
```

The array_of_datatypes stores the MPI datatypes of the elements

```
MPI_Datatype array_of_types[3] = \{MPI_DOUBLE, MPI_DOUBLE, MPI_INT\};
```

MPI_Get_address

- Returns the address of the memory location referenced by location_p.
- The special type MPI_Aint is an integer type that is big enough to store an address on the system.

```
int MPI_Get_address(
    void* location_p /* in */,
    MPI_Aint* address_p /* out */);
```

```
MPI_Datatype input_mpi_t;
int array_of_blocklengths[3] = \{1, 1, 1\};
MPI_Aint a_addr, b_addr, n_addr;
MPI_Get_address(&a. &a_addr):
array_of_displacements[0] = 0;
MPI_Get_address(&b, &b_addr);
array_of_displacements[1] = b_addr - a_addr;
MPI_Get_address(&n. &n_addr):
array_of_displacements[2] = n_addr - a_addr;
array_of_displacements[] = {0, 16, 24};
MPI_Datatype array_of_types[3] = \{MPI_DOUBLE, MPI_DOUBLE, MPI_INT\};
MPI_Type_create_struct(3, array_of_blocklengths,
      array_of_displacements, array_of_types,
      &input_mpi_t);
```

MPI_Type_commit

- Before we can use input_mpi_t in a communication function, we must first commit it with a call to mpi_Type_commit
- Allows the MPI implementation to optimize its internal representation of the datatype for use in communication functions.

```
int MPI_Type_commit(MPI_Datatype* new_mpi_t_p /* in/out */);
```

MPI_Type_free

• When we're finished with our new type, this frees any additional storage used.

```
int MPI_Type_free(MPI_Datatype* old_mpi_t_p /* in/out */);
```

Get input function with a derived datatype (1)

Get input function with a derived datatype (2)

Get input function with a derived datatype (3)

```
void Get_input(int my_rank, int comm_sz, double* a_p, double* b_p,
     int * n p) {
  MPI Datatype input mpi t;
  Build_mpi_type(a_p, b_p, n_p, &input_mpi_t);
  if (my_rank == 0) {
     printf("Enter a, b, and n\n");
     scanf("%lf %lf %d", ap, bp, np);
  MPI Bcast(a p, 1, input mpi t, 0, MPI COMM WORLD);
  MPI_Type_free(&input_mpi_t);
  /* Get_input */
```



Performance evaluation

Performance evaluation

- We write parallel programs because we expect that they'll be faster than a serial program that solves the same problem.
 - How can we verify this?
- We're not interested in the time taken from the start of program execution to the end of program execution.
 - E.g. the time it takes to type in the matrix or print out the product
- We're only interested in the time it takes to do the actual operation

Performance evaluation

 MPI provides a function, MPI_Wtime, that returns the number of seconds that have elapsed since some time in the past:

```
double MPI_Wtime(void);
```

Elapsed parallel time

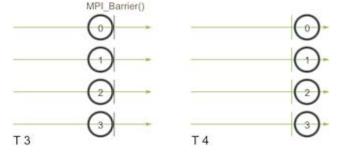
```
double MPI_Wtime(void);
```

 Returns the number of seconds that have elapsed since some time in the past.

MPI_Barrier

• Ensures that no process will return from calling it until every process

in the communicator has started calling it.



MPI Barrier()

Elapsed parallel time

```
int myrank, numprocs;
double mytime, maxtime, mintime, avgtime; /*variables used for gathering timing statistics*/
MPI Comm rank(MPI COMM WORLD, &myrank); MPI Comm size(MPI COMM WORLD, &numprocs);
MPI Barrier(MPI COMM WORLD); /*synchronize all processes*/
mytime = MPI_Wtime(); /*get time just before work section */
work();
mytime = MPI Wtime() - mytime; /*get time just after work section*/ /*compute max, min, and
average timing statistics*/
MPI Reduce(&mytime, &maxtime, 1, MPI DOUBLE, MPI MAX, 0, MPI COMM WORLD);
MPI Reduce(&mytime, &mintime, 1, MPI DOUBLE, MPI MIN, 0, MPI COMM WORLD);
MPI Reduce(&mytime, &avgtime, 1, MPI DOUBLE, MPI SUM, 0, MPI COMM WORLD);
if (myrank == 0) {
       avgtime /= numprocs;
       printf("Min: %lf Max: %lf Avg: %lf\n", mintime, maxtime,avgtime);
```

Elapsed serial time

- In this case, you don't need to link in the MPI libraries.
- Returns time in microseconds elapsed from some point in the past.

```
#include "timer.h"
. . .
double now;
. . .
GET_TIME(now);
```



Elapsed serial time

```
#include "timer.h"
. . .
double start, finish;
. . .
GET_TIME(start);
/* Code to be timed */
. . .
GET_TIME(finish);
printf("Elapsed time = %e seconds\n", finish-start);
```

Run-times of serial and parallel matrix-vector multiplication

	Order of Matrix				
comm_sz	1024	2048	4096	8192	16,384
1	4.1	16.0	64.0	270	1100
2	2.3	8.5	33.0	140	560
4	2.0	5.1	18.0	70	280
8	1.7	3.3	9.8	36	140
16	1.7	2.6	5.9	19	71

- If we fix *n* and increase comm_sz, the run-times usually decrease
- For small *n*, there is very little benefit in increasing comm_sz
- As we increase the problem size, the run-times increase
- As we increase the number of processes, the run-times typically decrease for a while. However, at some point, the run-times can actually start to get worse → Why?

Speedup and efficiency

$$S(n,p) = \frac{T_{\text{serial}}(n)}{T_{\text{parallel}}(n,p)}$$

	Order of Matrix				
comm_sz	1024	2048	4096	8192	16,384
1	1.0	1.0	1.0	1.0	1.0
2	1.8	1.9	1.9	1.9	2.0
4	2.1	3.1	3.6	3.9	3.9
8	2.4	4.8	6.5	7.5	7.9
16	2.4	6.2	10.8	14.2	15.5

E(n,p) =	S(n,p)		$T_{\text{serial}}(n)$		
	p	_	$p \times T_{\text{parallel}}(n, p)$		

	Order of Matrix				
comm_sz	1024	2048	4096	8192	16,384
1	1.00	1.00	1.00	1.00	1.00
2	0.89	0.94	0.97	0.96	0.98
4	0.51	0.78	0.89	0.96	0.98
8	0.30	0.61	0.82	0.94	0.98
16	0.15	0.39	0.68	0.89	0.97

Scalability

 A program is scalable if the problem size can be increased at a rate so that the efficiency doesn't decrease as the number of processes increase.



Scalability

- Programs that can maintain a constant efficiency without increasing the problem size are sometimes said to be strongly scalable.
- Programs that can maintain a constant efficiency if the problem size increases at the same rate as the number of processes are sometimes said to be weakly scalable.