ECE 432/532 Programming for Parallel Processors

Sorting

- n keys and p = comm sz processes.
- n/p keys assigned to each process.
- No restrictions on which keys are assigned to which processes.
- When the algorithm terminates:
 - The keys assigned to each process should be sorted in (say) increasing order.
 - If 0 ≤ q < r < p, then each key assigned to process q should be less than or equal to every key assigned to process r.

Serial bubble sort

Serial bubble sort

- Is it good for parallel implementation?
- Example, you have the numbers 9, 5, 7
 - How would you sort them in serial way?
 - How would you sort them in parallel way?
- The order in which the "compare-swaps" take place is essential to the correctness of the algorithm

Odd-even transposition sort

- A sequence of phases.
- Even phases, compare swaps:

$$(a[0], a[1]), (a[2], a[3]), (a[4], a[5]), \dots$$

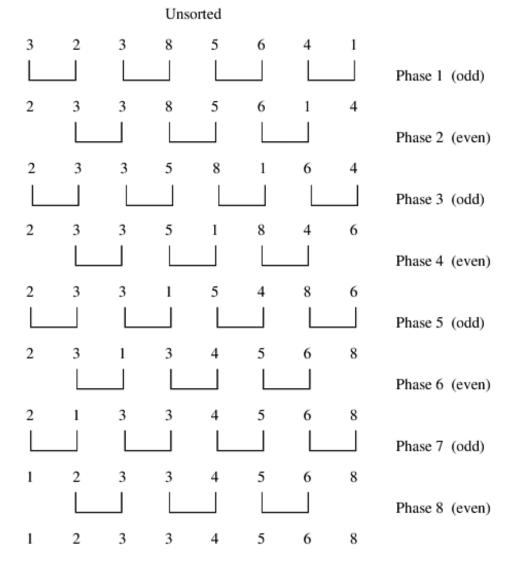
Odd phases, compare swaps:

$$(a[1], a[2]), (a[3], a[4]), (a[5], a[6]), \dots$$

Example

```
Start: 5, 9, 4, 3
Even phase: compare-swap (5,9) and (4,3) getting the list 5, 9, 3, 4
Odd phase: compare-swap (9,3) getting the list 5, 3, 9, 4
Even phase: compare-swap (5,3) and (9,4) getting the list 3, 5, 4, 9
Odd phase: compare-swap (5,4) getting the list 3, 4, 5, 9
```

Example



Sorted

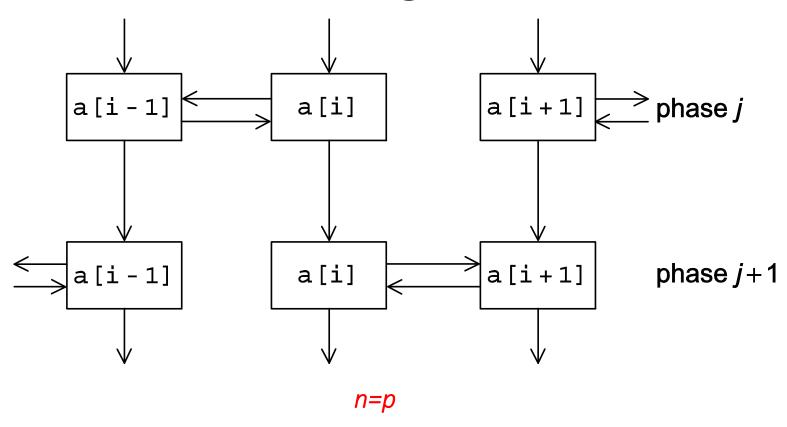
Serial odd-even transposition sort

```
void Odd_even_sort(
      int a[] /* in/out */,
      int n /* in */) {
   int phase, i, temp;
  for (phase = 0; phase < n; phase++)
      if (phase % 2 == 0) { /* Even phase */
        for (i = 1; i < n; i += 2)
            if (a[i-1] > a[i]) {
              temp = a[i];
              a[i] = a[i-1];
              a[i-1] = temp;
      } else { /* Odd phase */
        for (i = 1; i < n-1; i += 2)
           if (a[i] > a[i+1]) {
              temp = a[i];
              a[i] = a[i+1];
              a[i+1] = temp;
  /* Odd_even_sort */
```

Parallel odd-even transposition sort

- The odd-even transposition sort has considerably more opportunities for parallelism than bubble sort
- All of the compare-swaps in a single phase can happen simultaneously.
- Foster's methodology:
 - Tasks: Determine the value of a[i] at the end of phase j.
 - Communications: The task that's determining the value of a[i] needs to communicate with either the task determining the value of a[i-1] or a[i+1]. Also the value of a[i] at the end of phase j needs to be available for determining the value of a[i] at the end of phase j+1.

Communications among tasks in odd-even sort



Parallel odd-even transposition sort

	Process			
Time	0	1	2	3
Start	15, 11, 9, 16	3, 14, 8, 7	4, 6, 12, 10	5, 2, 13, 1
After Local Sort	9, 11, 15, 16	3, 7, 8, 14	4, 6, 10, 12	1, 2, 5, 13
After Phase 0	3, 7, 8, 9	11, 14, 15, 16	1, 2, 4, 5	6, 10, 12, 13
After Phase 1	3, 7, 8, 9	1, 2, 4, 5	11, 14, 15, 16	6, 10, 12, 13
After Phase 2	1, 2, 3, 4	5, 7, 8, 9	6, 10, 11, 12	13, 14, 15, 16
After Phase 3	1, 2, 3, 4	5, 6, 7, 8	9, 10, 11, 12	13, 14, 15, 16

When will the list be sorted? After how many phases?

Pseudo-code

```
Sort local keys;
for (phase = 0; phase < comm_sz; phase++) {
   partner = Compute_partner(phase, my_rank);
   if (I'm not idle) {
      Send my keys to partner;
      Receive keys from partner;
      if (my_rank < partner)
            Keep smaller keys;
      else
            Keep larger keys;
}
</pre>
```

Compute_partner

MPI PROC NULL is a constant defined by MPI. When it's used as the source or destination rank in a point-to-point communication, no communication will take place and the call to the communication will simply return