

# An introduction to complexity theory

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# Complexity theory in a nutshell

“How hard can it be?”, *Jeremy Clarkson*

# What is complexity theory?

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What do we mean when we talk about a computer?

# Structure of part one

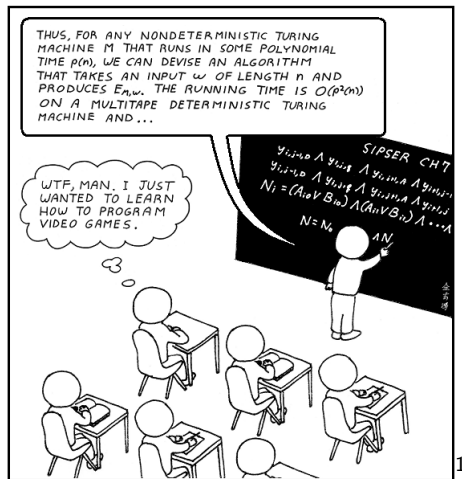
- What is a computer?
- What is a problem?
- How do we measure difficulty?

# Summary of part one

- What is a computer? *Deterministic Turing Machine, Non-Deterministic Turing Machine*
- What is a problem? *Deciding if a word is in a language, verifying that a word is in a language*
- How do we measure difficulty? *Upper bound of time for an input of length  $n$*



## End of part one



<sup>1</sup><http://abstrusegoose.com/206>

# Structure of part two

- Putting it all together!
- ...only to get another (very difficult) problem.
- How might we try to solve this new problem?

# Exercise Left for the Student

Does  $P = NP$ ?

(NB: You should probably refer to the literature before trying to solve this.)

# The $P$ versus $NP$ problem

Arguably first proposed by Gödel in a letter to von Neumann in 1956.<sup>2</sup>

First stated formally by Cook in 1971.<sup>3</sup>

Solving the problem will earn you a million dollars, courtesy of the Clay Institute.<sup>4</sup>

Aaronson has called it the most important Millenium Problem, as the answer could make the other problems significantly easier or harder to solve.<sup>5</sup>

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<sup>2</sup><https://ecommons.cornell.edu/bitstream/handle/1813/6910/89-994.pdf>

<sup>3</sup><http://dl.acm.org/citation.cfm?coll=GUIDE&dl=GUIDE&id=805047>

<sup>4</sup><http://www.claymath.org/millennium-problems>

<sup>5</sup><http://www.thenakedscientists.com/HTML/interviews/interview/1001376/>

# The easy side: $P \subseteq NP$

Recall that any  $TM$  is by definition non-deterministic.

Likewise, any polynomial-time  $TM$  is also non-deterministic.

So  $\forall L \in P, L \in NP$ . Hence  $P \subseteq NP$ .

# The easy side: $P \subseteq NP$

Alternative proof (using verification):

Let  $TM M$  decide  $L$  in polynomial time. Define  $V$  as follows:

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 $V(w, c) :$ 
if  $M(w)$  accepts then
  | accept
end
else
  | reject
end

```

$V$  verifies  $L$  in polynomial time. Hence  $P \subseteq NP$ .

# The harder side: Is $NP \subseteq P$

Another way to think of this problem is: *If a problem can be easily verified, can it be easily solved?*

## How might we answer this question?

Why not look at the hardest problems in  $NP$ ?

If  $P = NP$ , then even the hardest problems in  $NP$  will be solvable in polynomial time.

And if  $P \subset NP$ , then these are the problems that won't have a polynomial time solution, as could be checked by lower-bound analysis.

But how can we determine the hardest problems in  $NP$ ?



# Summary of part two

- Putting it all together!  $P$ ,  $NP$
- ...only to get another (very difficult) problem. *Are easy to verify problems easy to solve?*
- How might we try to solve this new problem?  *$NP$  – Complete problems*

# What else is there?

Recall our three questions from part one:

- What is a computer?
- What is a problem?
- How do we measure difficulty?

What if we answered these differently?

# What is a computer?

Probabilistic Turing Machines:  $BPP$ ,  $RP$

Quantum computers:  $EQP$ ,  $BQP$

Talking to another, more powerful computer:  $MA$ ,  $IP$

Time travel:  $P_{CTC}$

# What is a problem?

Computational problems:  $NP - Hard$

Counting problems:  $\#P$

Complementary problems:  $co - NP$

# How do we measure difficulty?

Exponential time:  $EXP$

Space complexity:  $PSPACE, EXPSPACE$

Linear time:  $LIN$

Sublinear working space:  $L$

# This is only the beginning

There are many more complexity classes out there, and very quickly relating them in a simple equation like this:

$$P \subseteq NP$$

Becomes this:

$$L \subseteq NL \subseteq P \subseteq NP \subseteq PSPACE = NPSPACE = IP = P_{CTC} \subseteq EXP \subseteq EXPSPACE$$

## The end

PROOF:

$$e^{i \cdot P_i} = -1$$

And,

$$P_i = P \cdot i$$

So,

$$e^{i \cdot P_i} = e^{P \cdot i \cdot i} = e^{-P}$$

So,

$$e^{-P} = -1$$

Squaring both sides,

$$e^{-2P} = 1$$

Which leaves

$$P = 0$$

Thus,

$$P = NP$$

QED

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<sup>6</sup><http://www.smbc-comics.com/?id=3919>

## The end



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