

# Communication topologies in natural computing

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# Outline

- The first and second machine classes
- Communication topologies and their role
- Membrane computing
- Complexity theory of membrane systems
- A research project

# The first machine class and P

The deterministic **Turing machine** and all models that simulate and are simulated by it efficiently

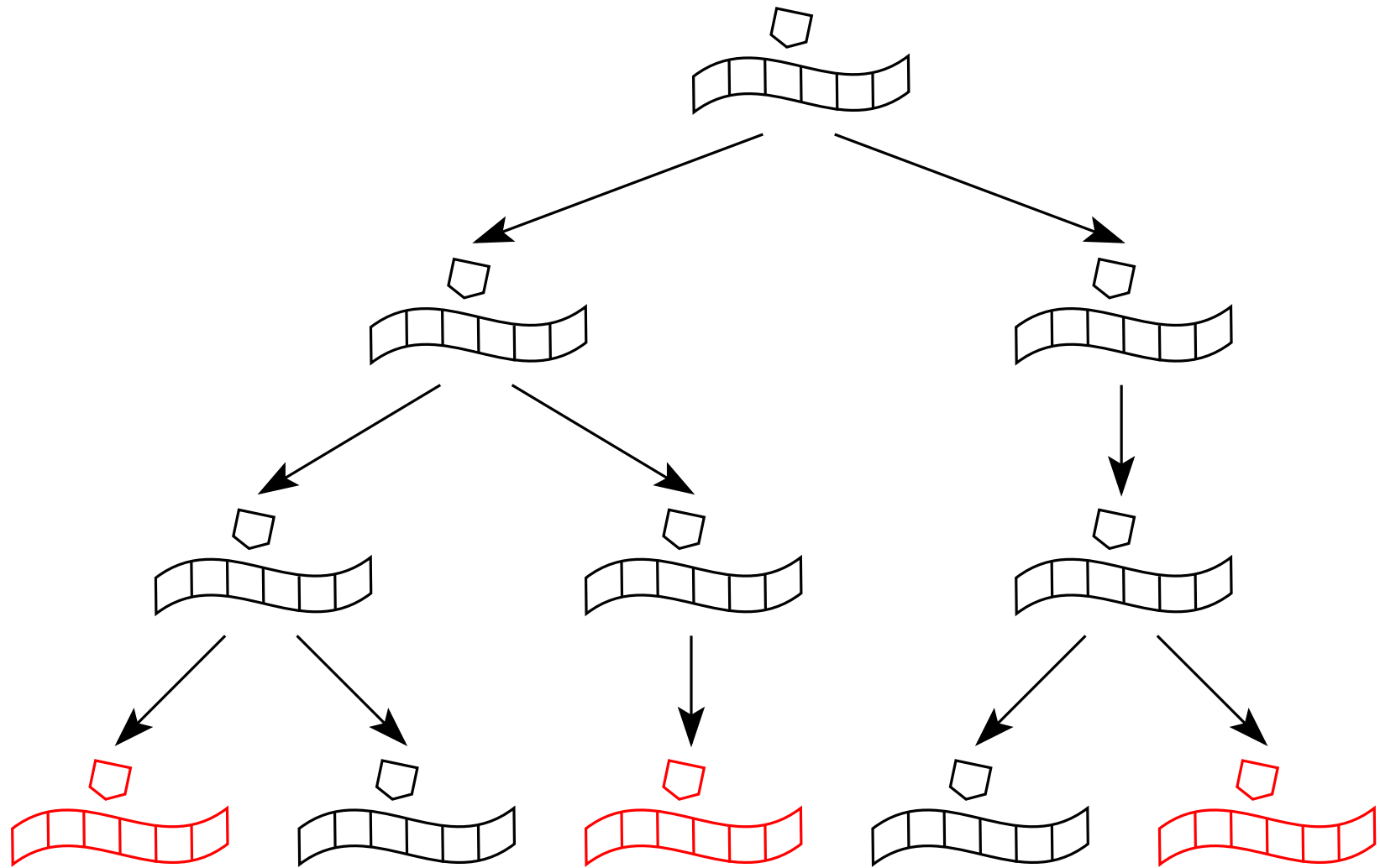
- Random access machines with arithmetic operations + and –
- **Cellular automata** with finite initial configuration

# The second machine class and PSPACE

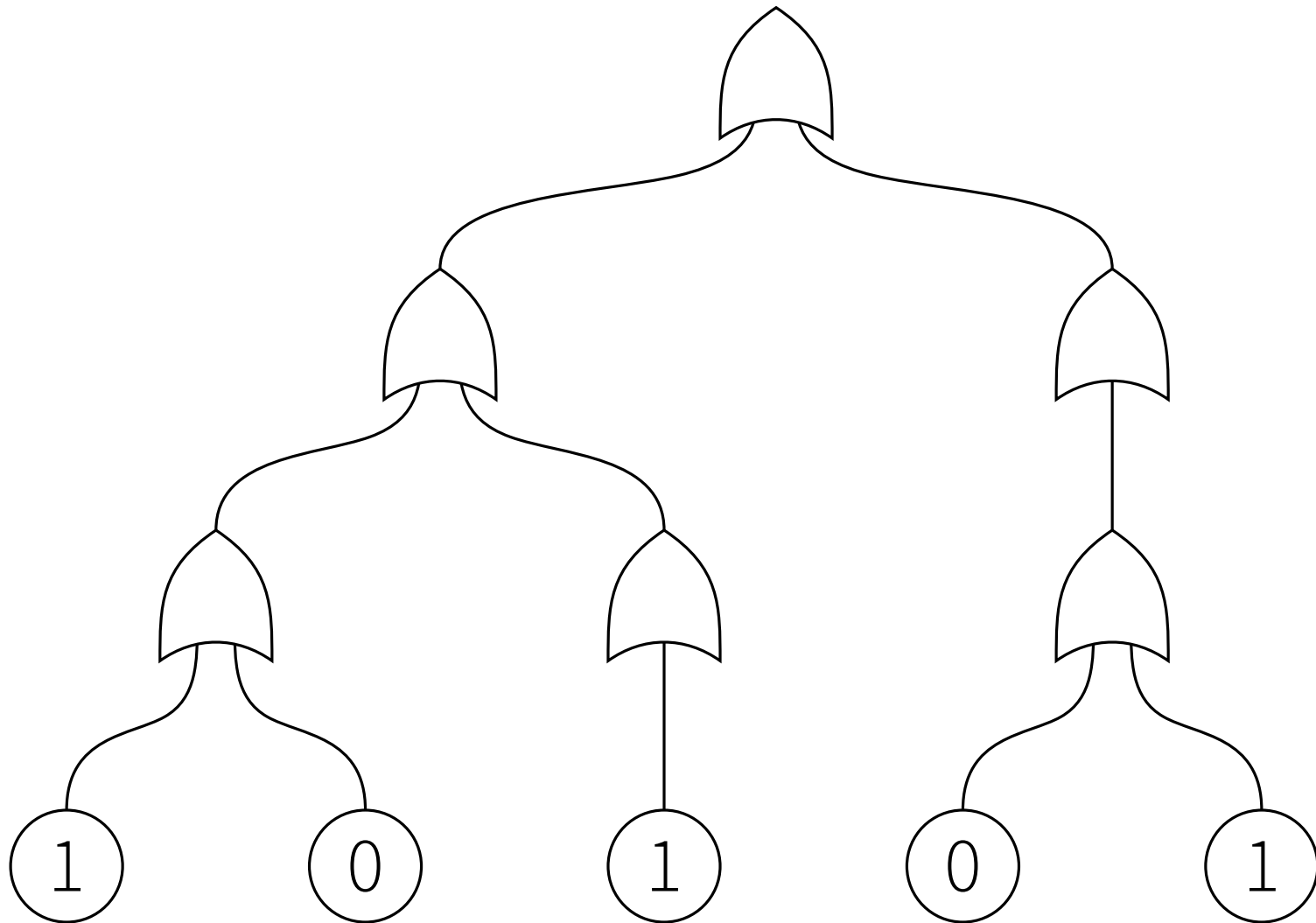
Computing models that solve in polynomial **time**  
what a Turing machine solves in polynomial **space**

- Alternating Turing machines
- **Random access machines** with arithmetic operations  $+$   $-$   $\times$   $\div$
- Parallel processes generated by **fork()** running on an unbounded number of processors
- Cellular automata over hyperbolic grids

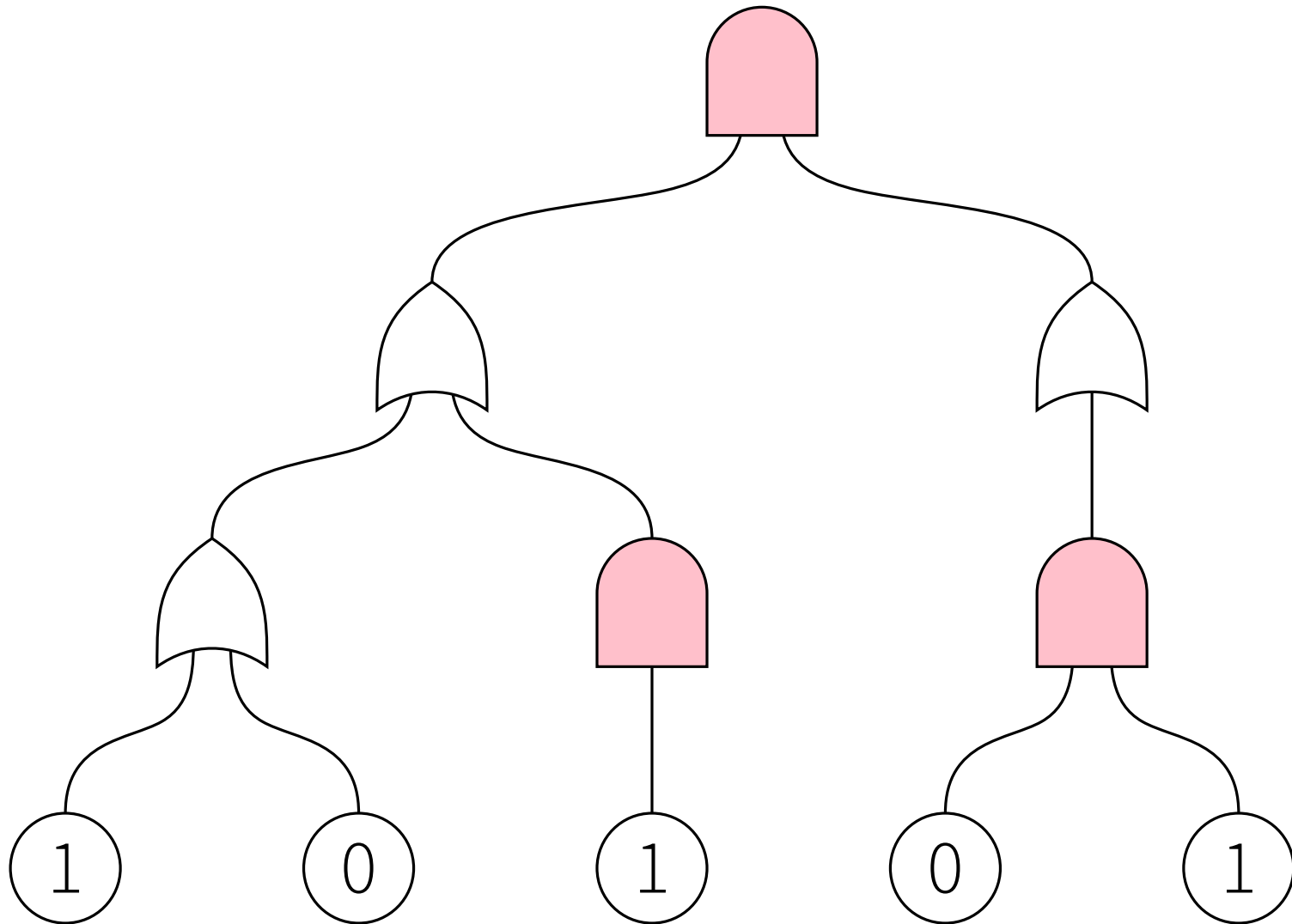
# Nondeterministic Turing machines: NP



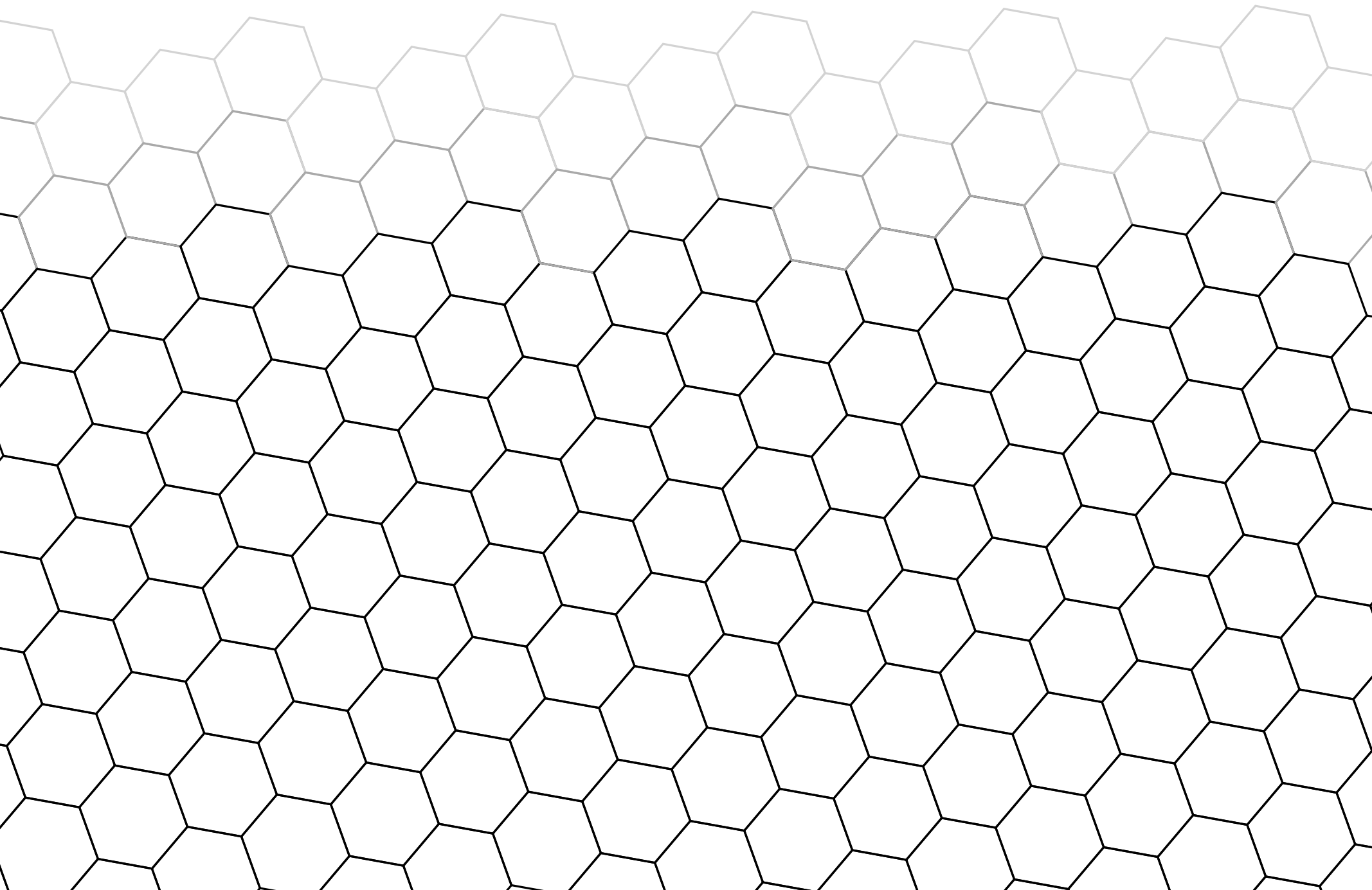
# Nondeterministic Turing machines: NP



# Alternating Turing machines: PSPACE

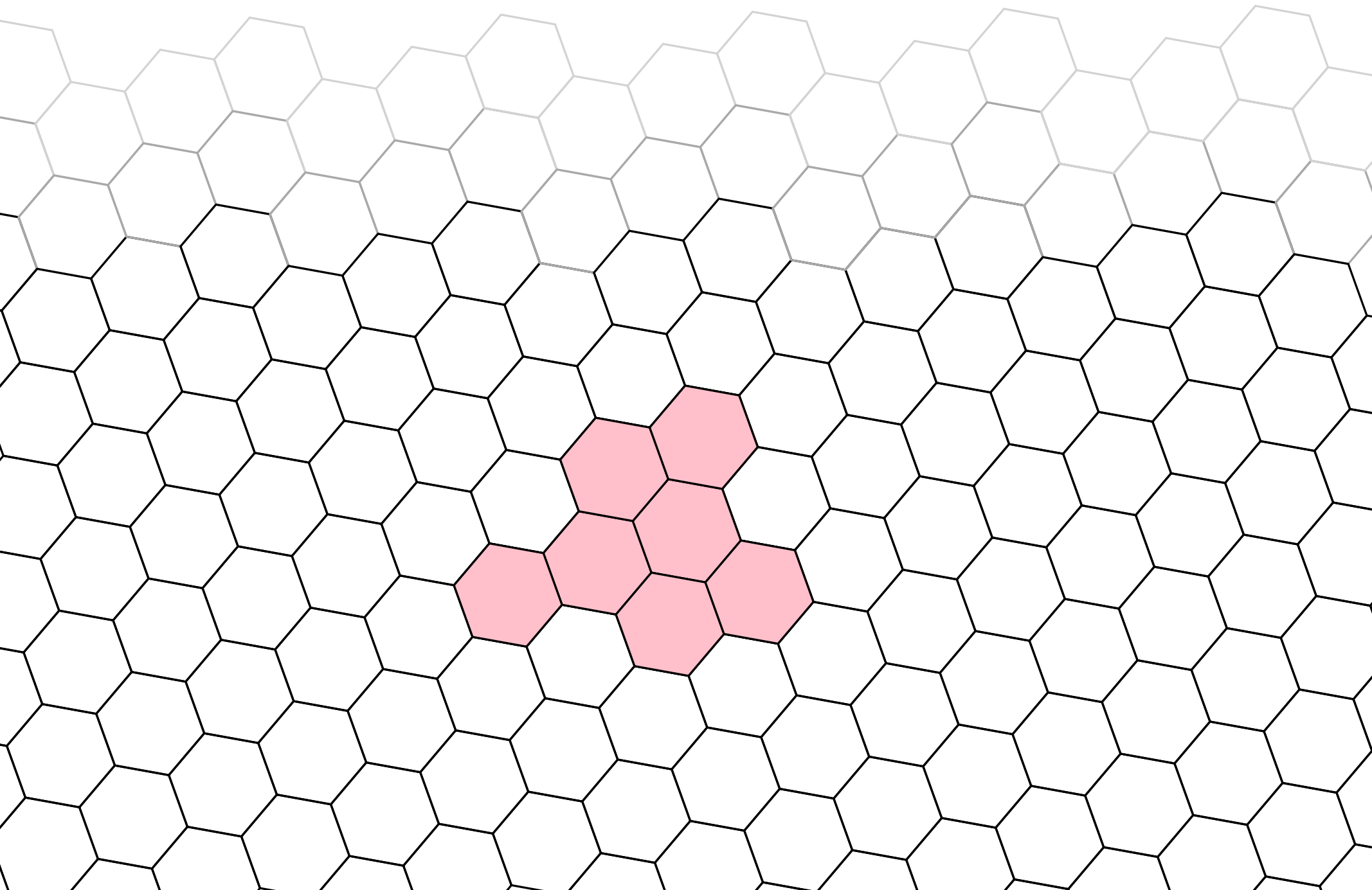


# Cellular automata over a Euclidean grid

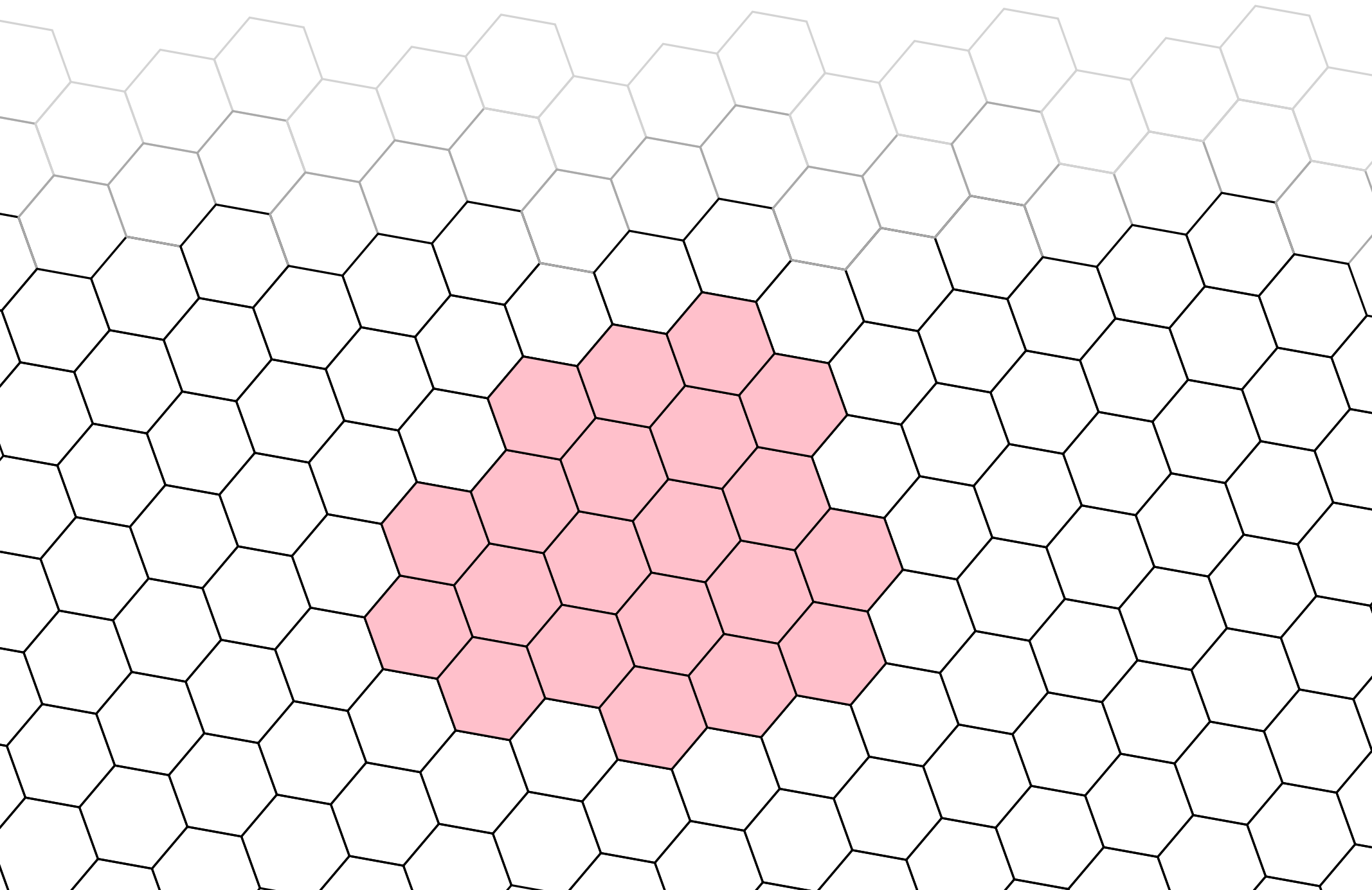




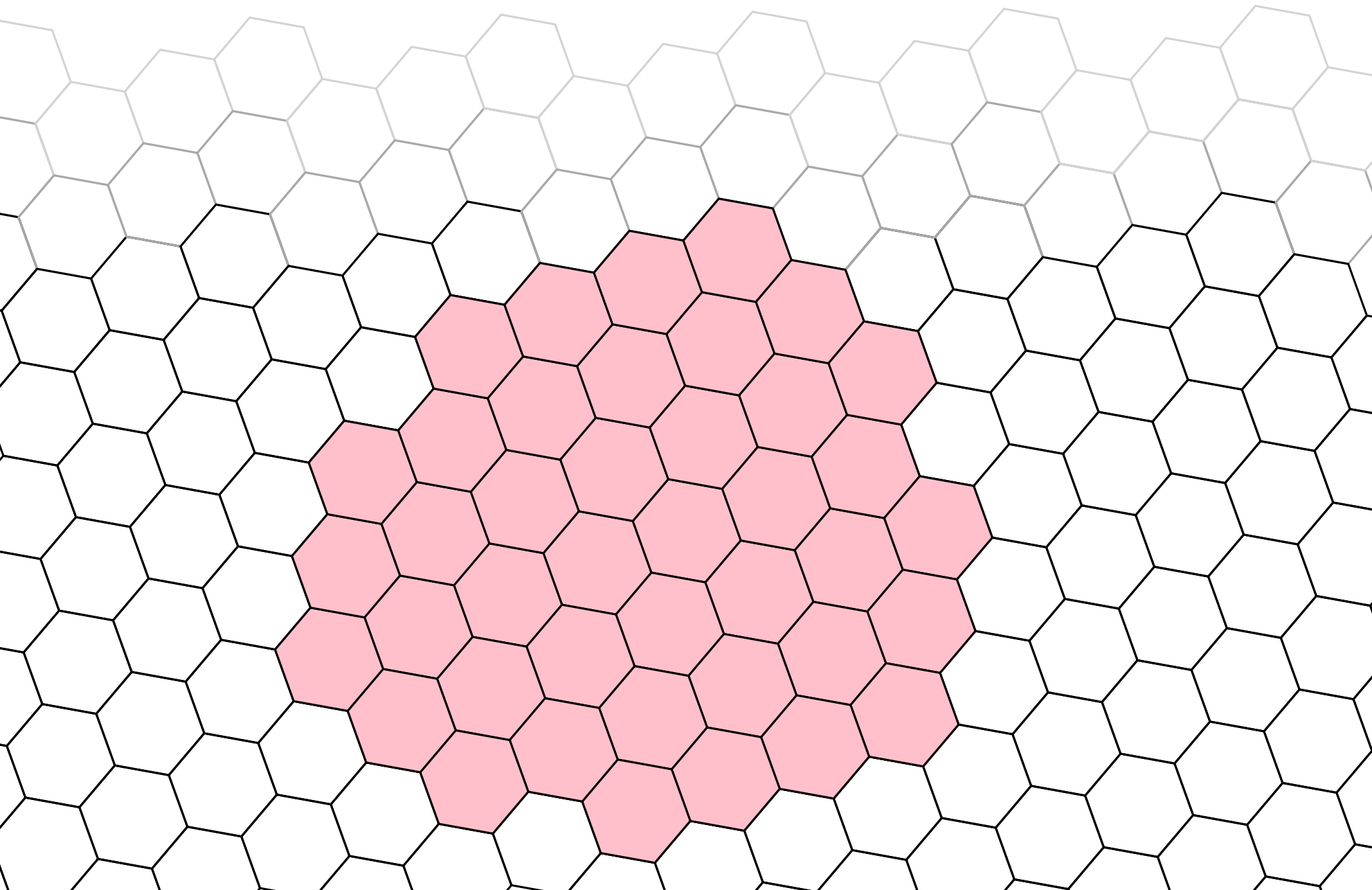
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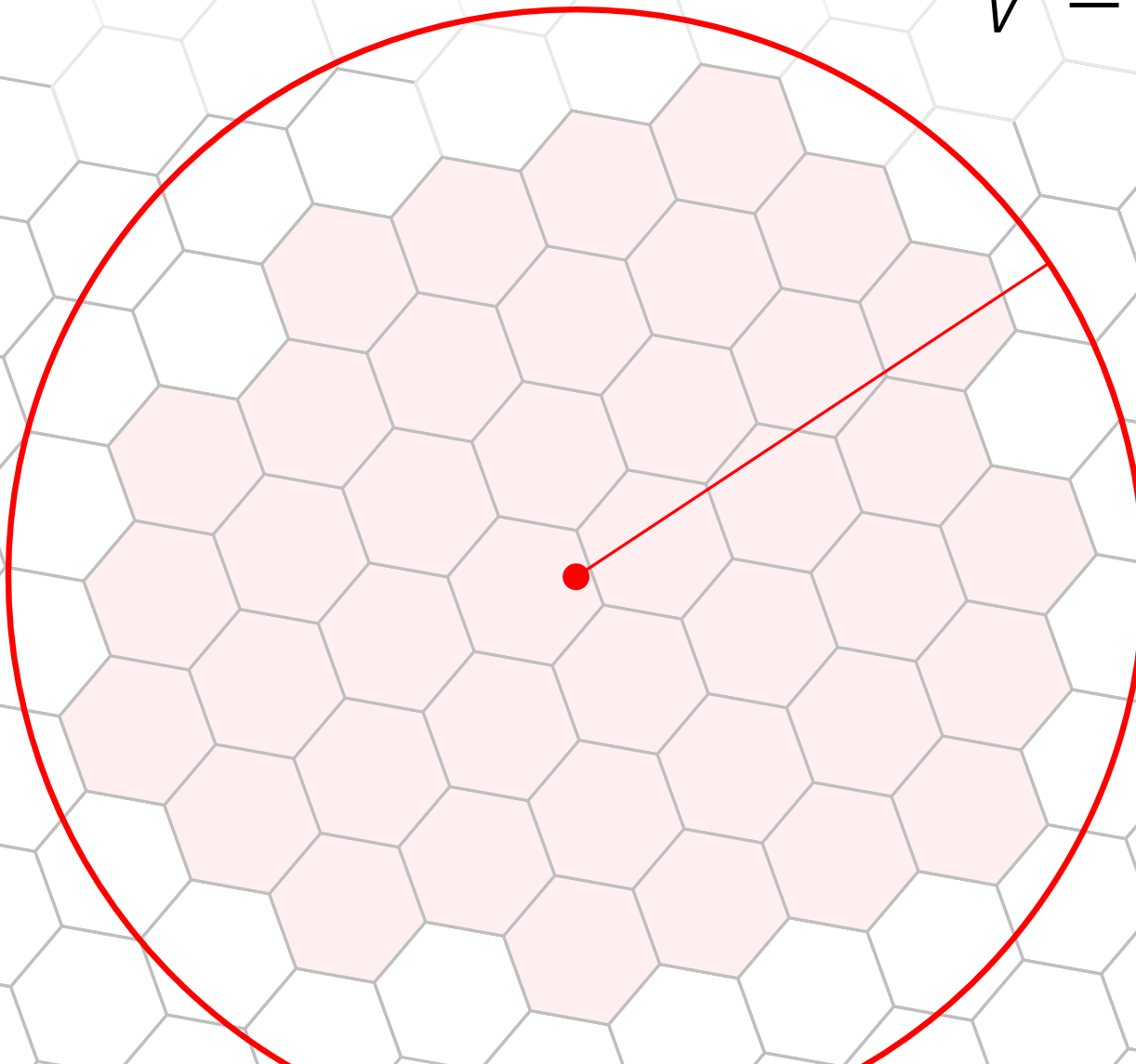


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# Cellular automata over a Euclidean grid

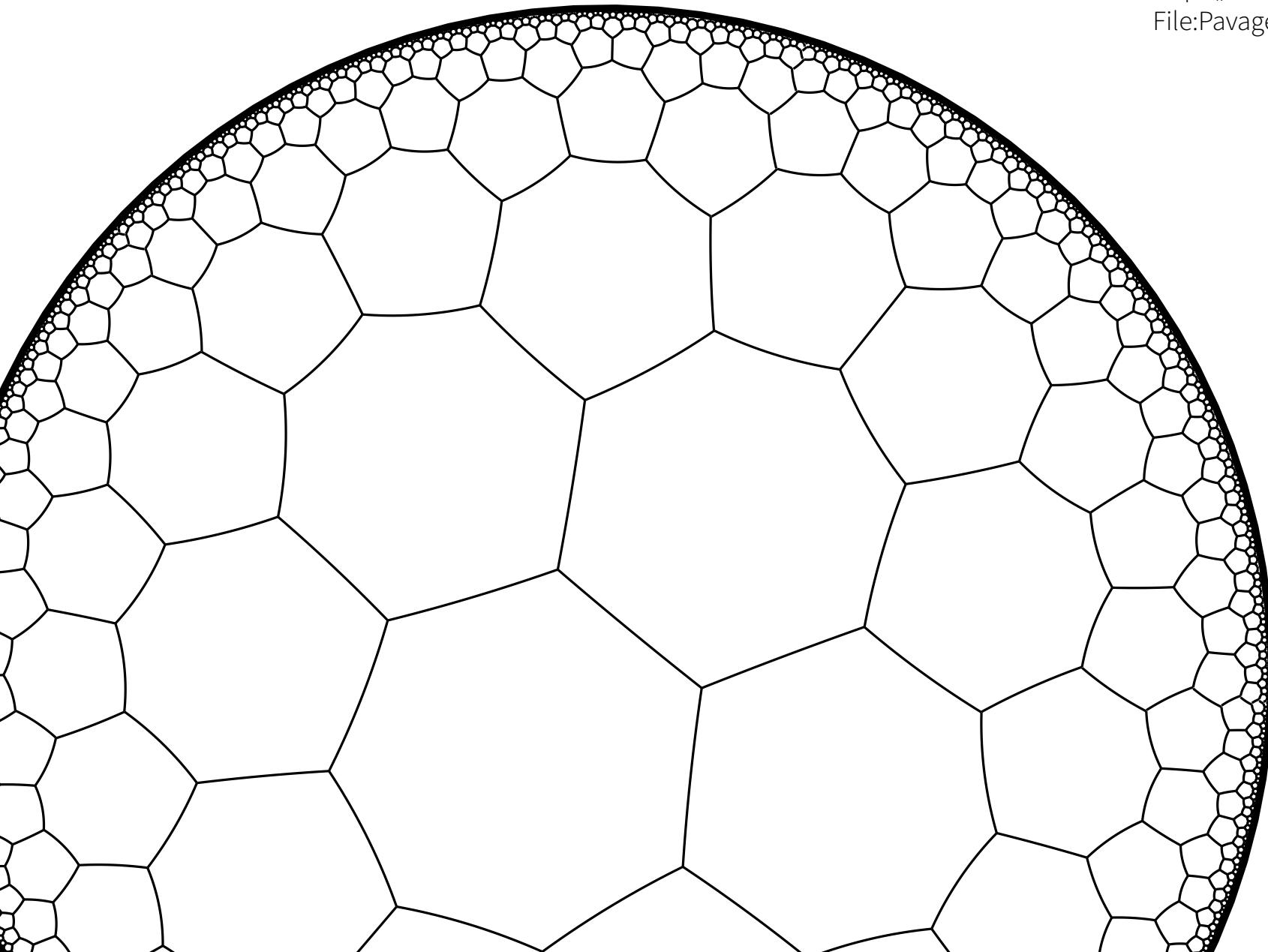
$$V = \Theta(r^d)$$



# Hyperbolic cellular automata

Pavage du plan hyperbolique par des heptagones, dans le modèle du disque de Poincaré. By Theon, used under CC BY-SA 3.0

<https://en.wikipedia.org/wiki/File:PavageHypPoincare2.svg>

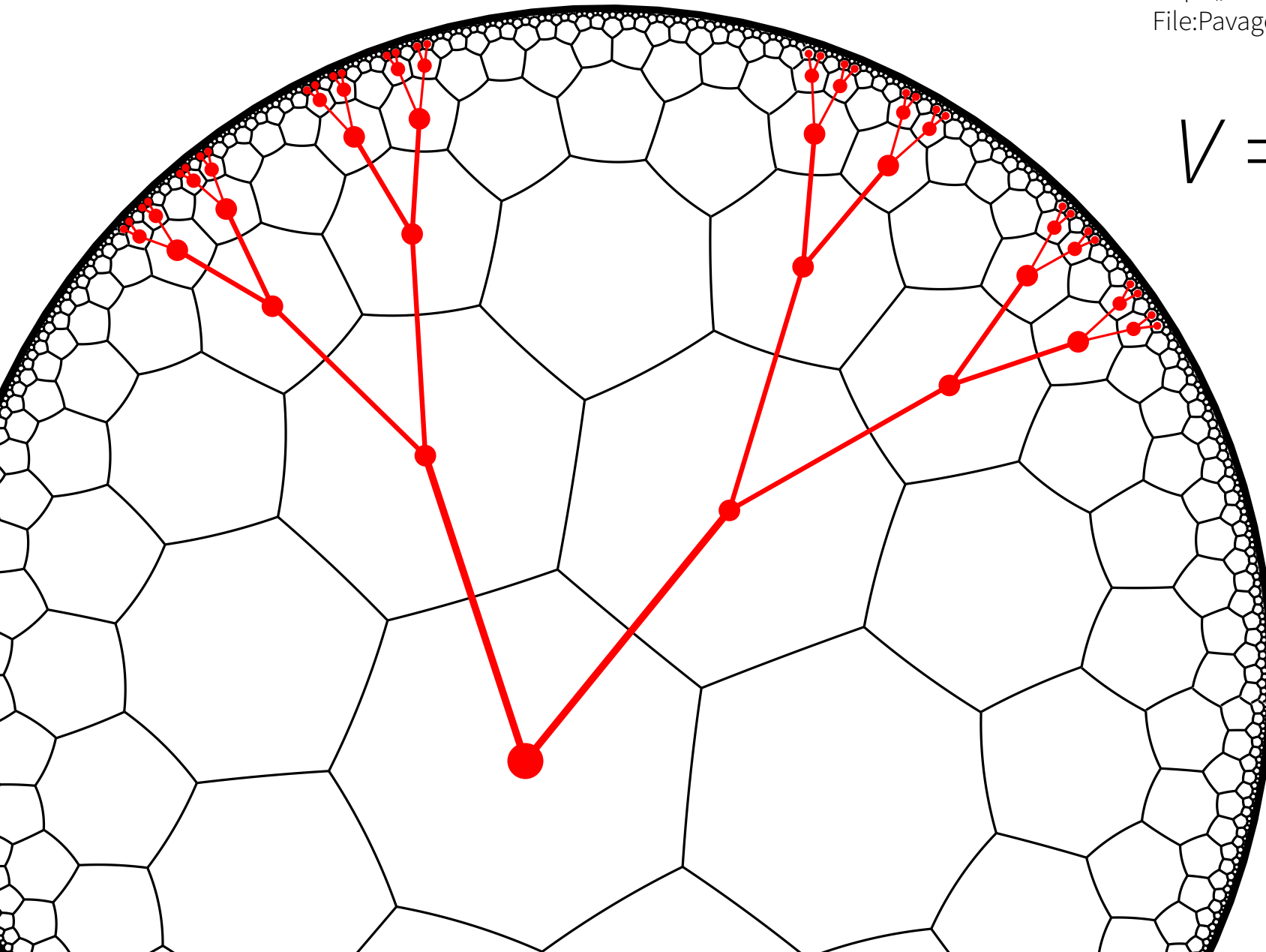


# Hyperbolic cellular automata

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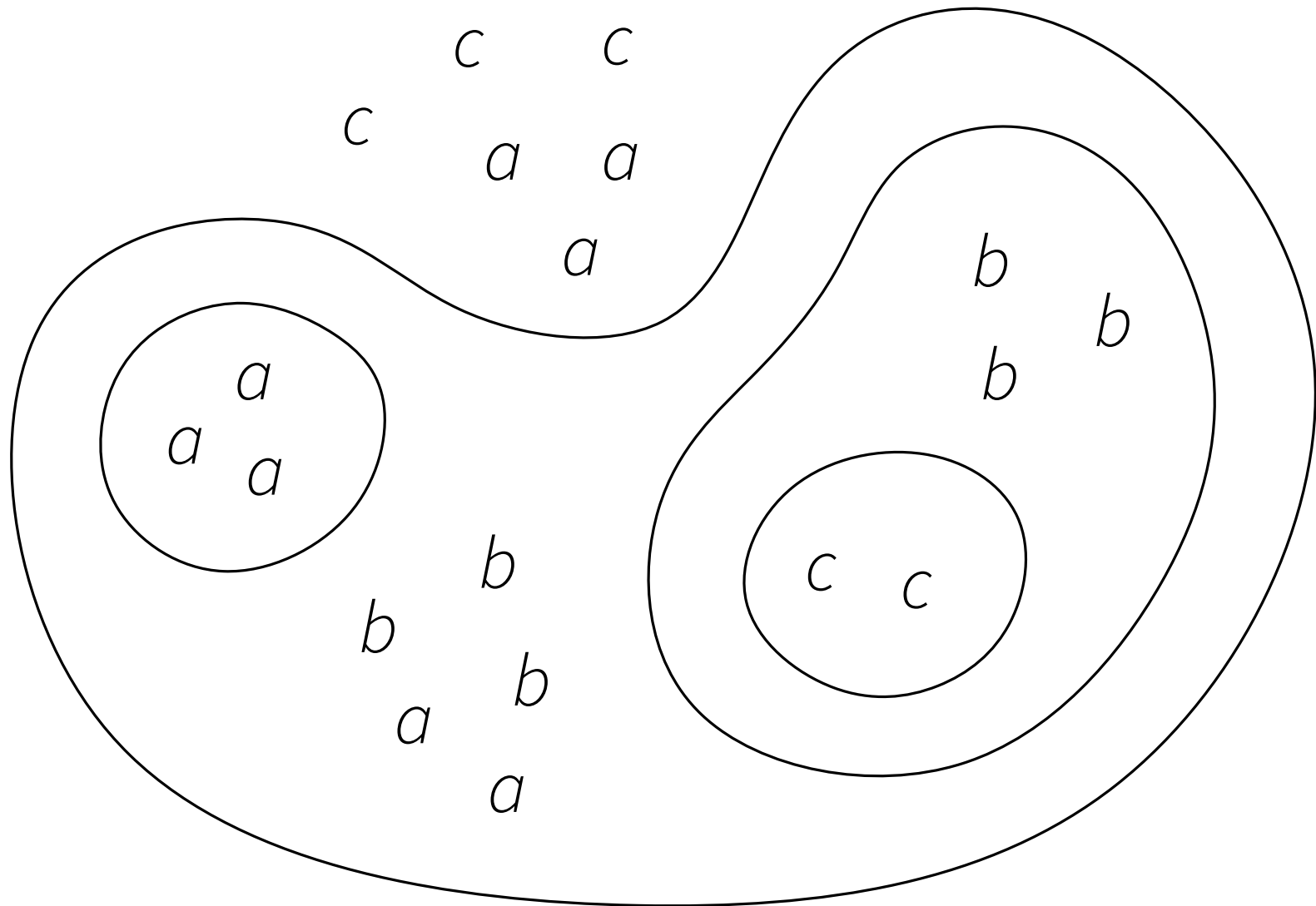
$$V = \Omega(2^r)$$



# Rule of thumb

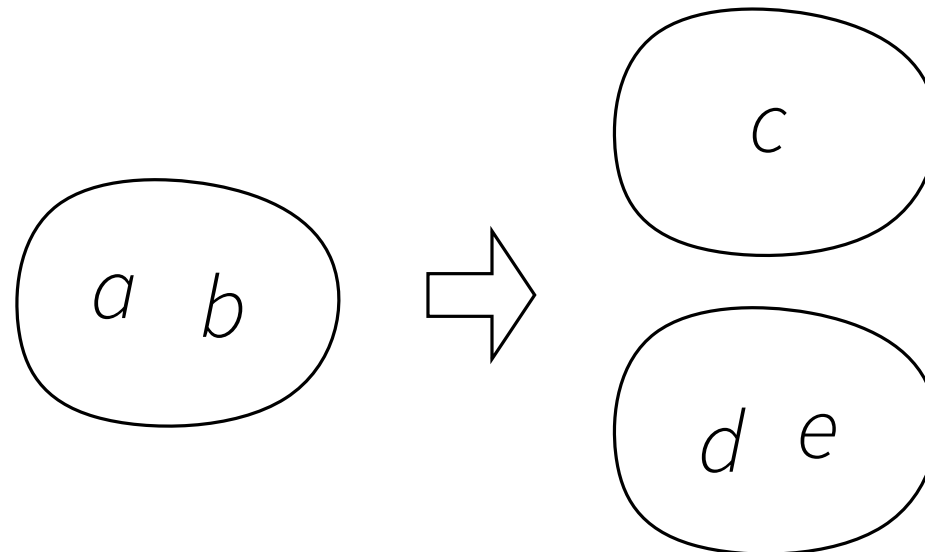
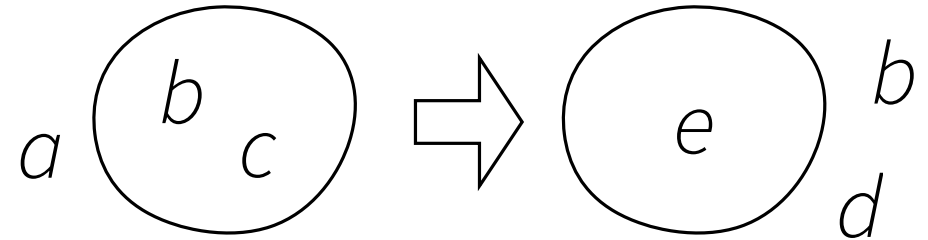
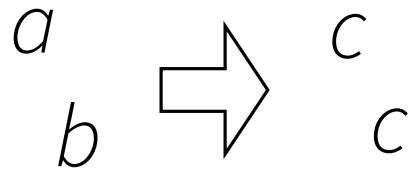
- Sequential machines are first class
- Bounded parallel machines are first class
- Unbounded parallel machines are second class

# Membrane systems (P systems)

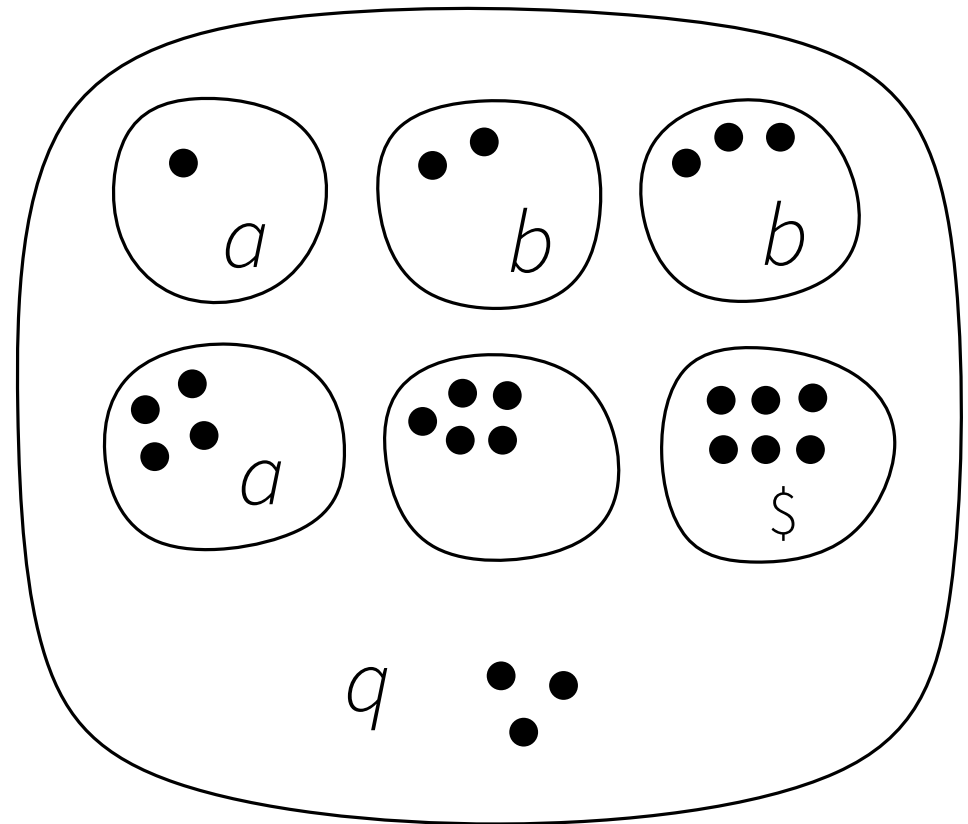
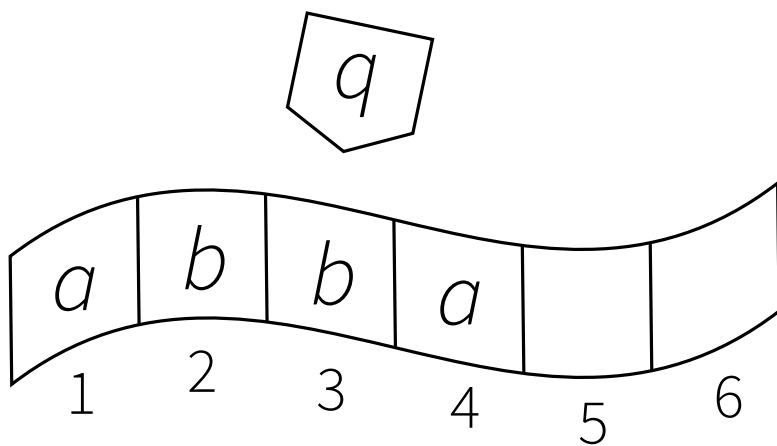




# Evolution rules and parallel semantics

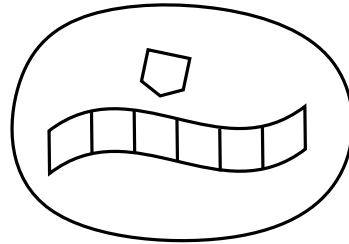


# Efficient universality of membrane systems



Alhazov, Leporati, Mauri, Porreca, Zandron: Space complexity equivalence of P systems with active membranes and Turing machines. Theoretical Computer Science 529, 69–81 (2014)

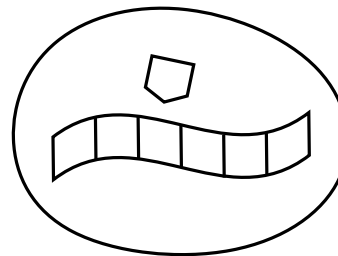
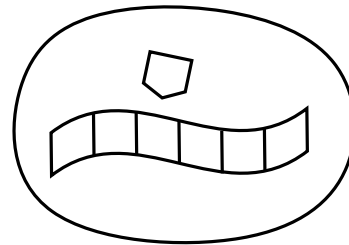
# Membrane systems solving NP problems



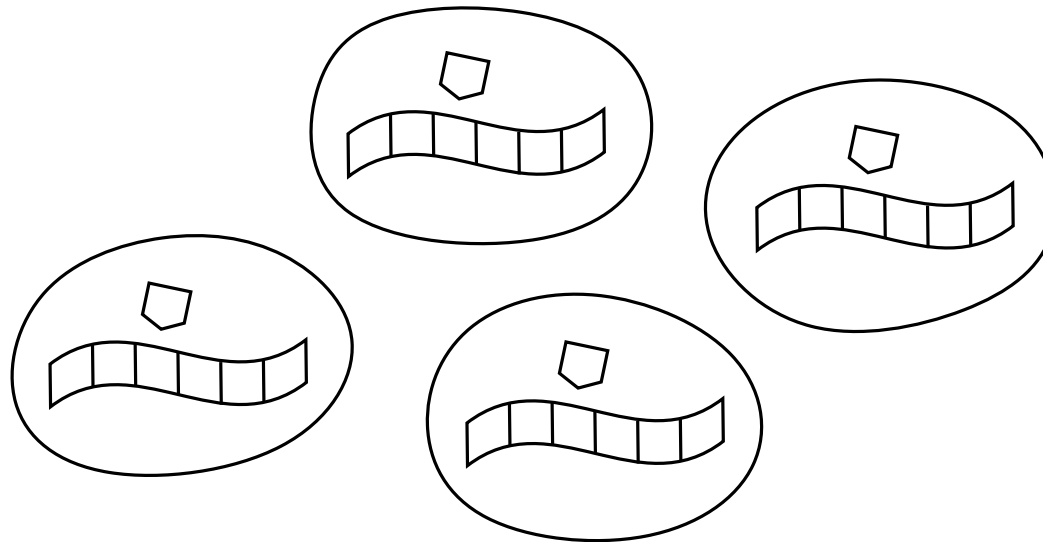
# Membrane systems solving NP problems

$$\delta(q, a) = \begin{cases} (q', b, +1) \\ (q', c, -1) \end{cases}$$

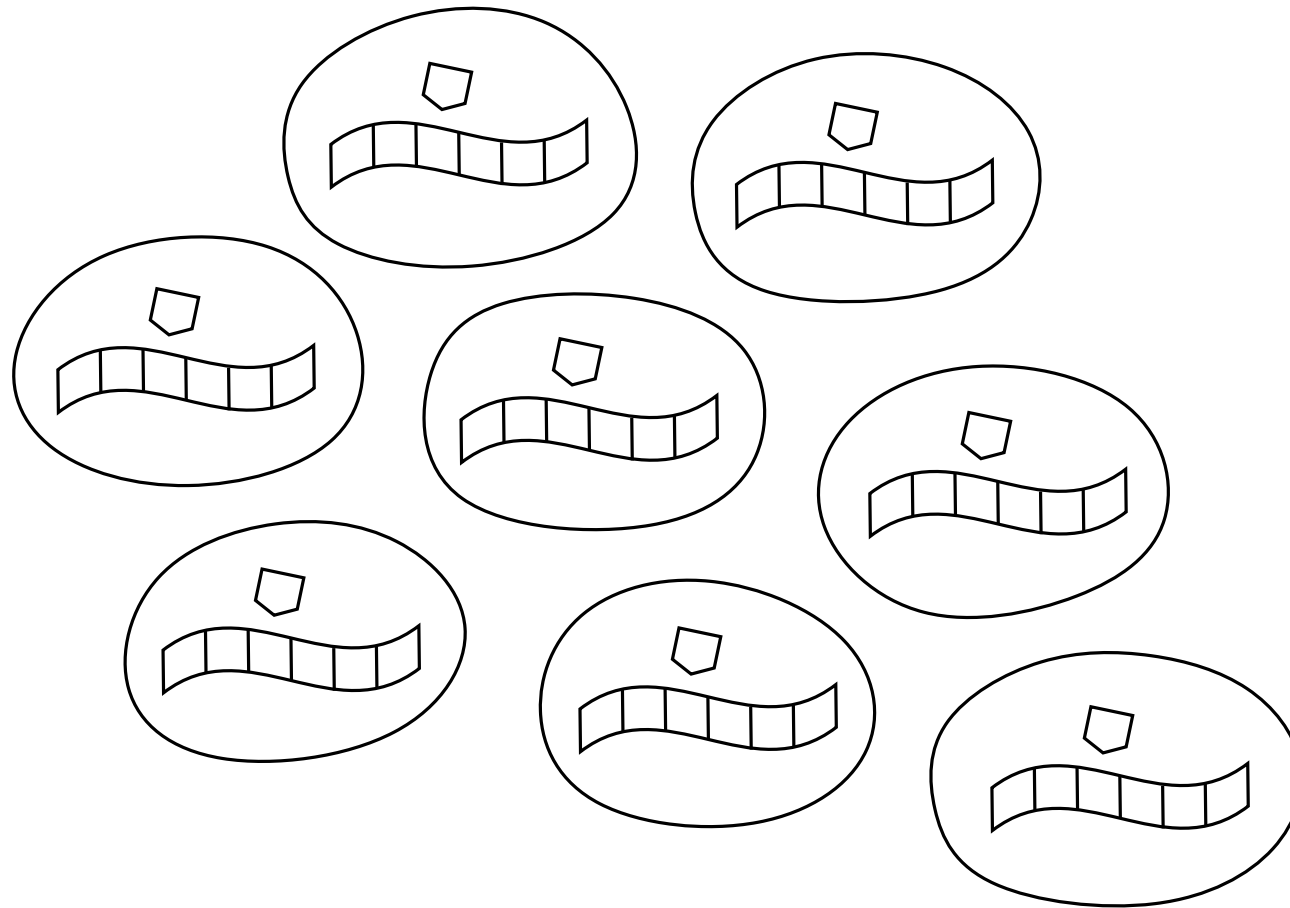
$$[q_3 \ a_3] \rightarrow [q'_4 \ b_3] [q''_2 \ c_3]$$



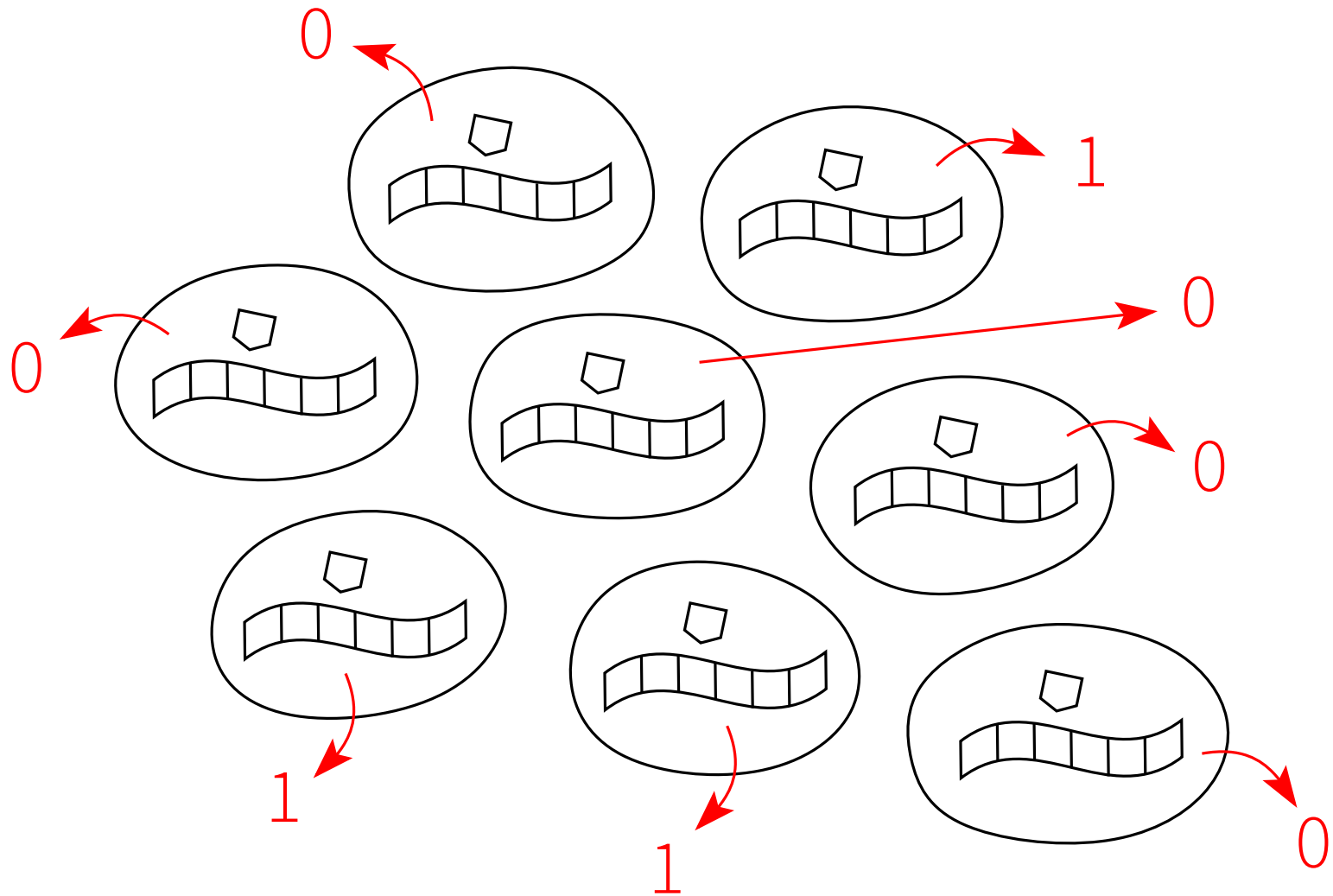
# Membrane systems solving NP problems



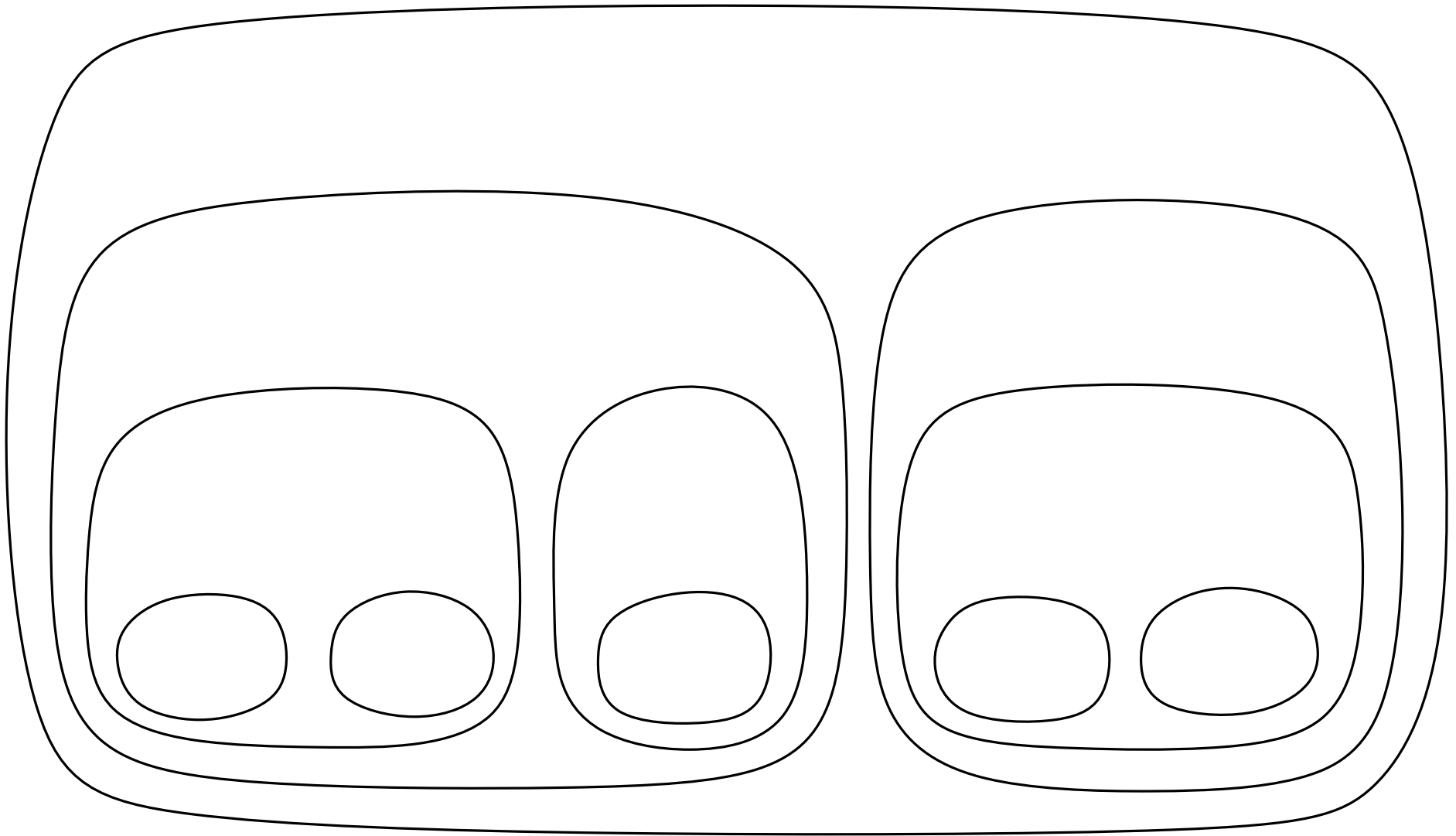
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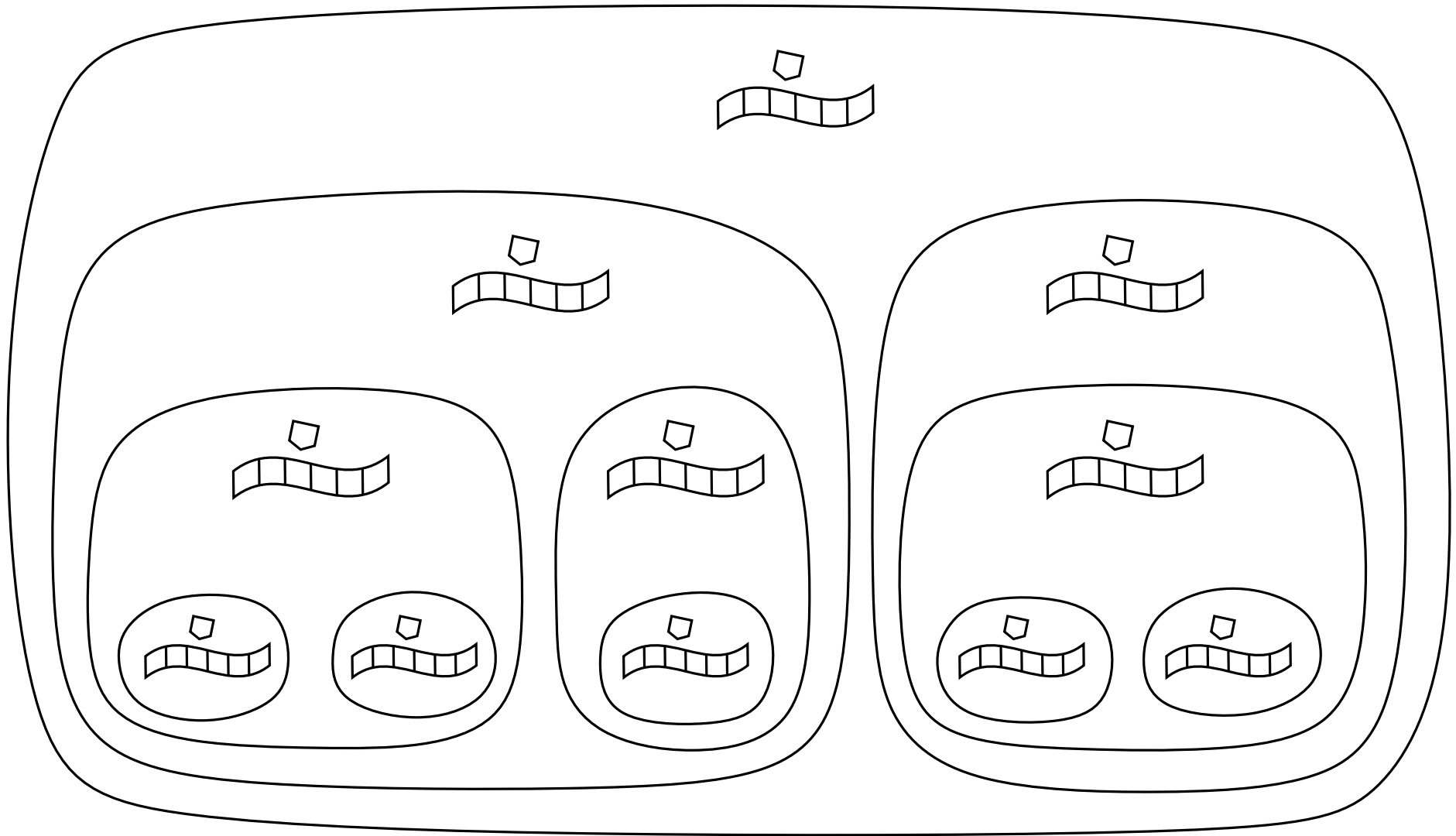
# Membrane systems are second class



Sosík, Rodríguez-Patón: Membrane computing and complexity theory: A characterization of PSPACE. Journal of Computer and System Sciences 73(1), 137–152 (2007)

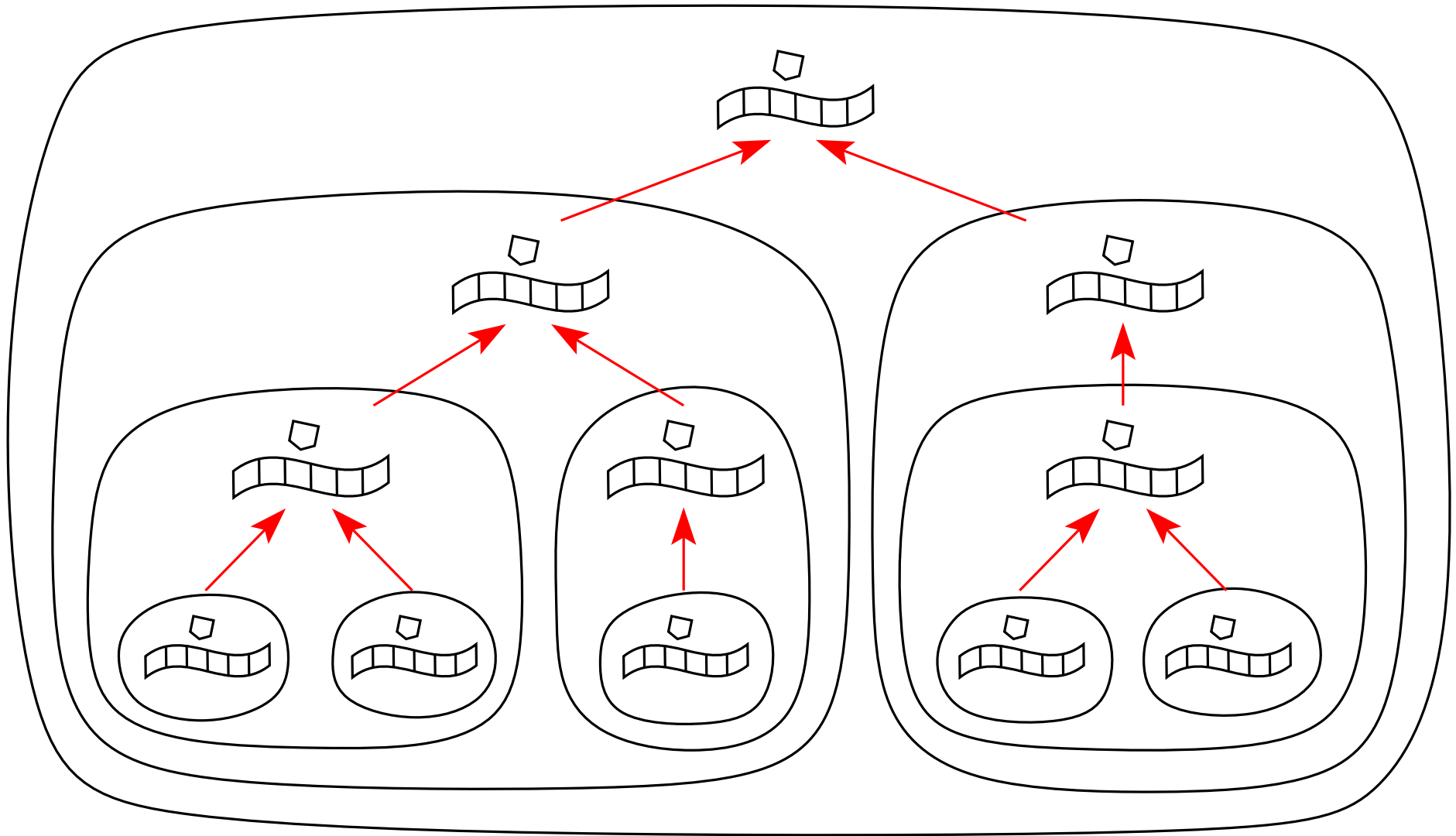


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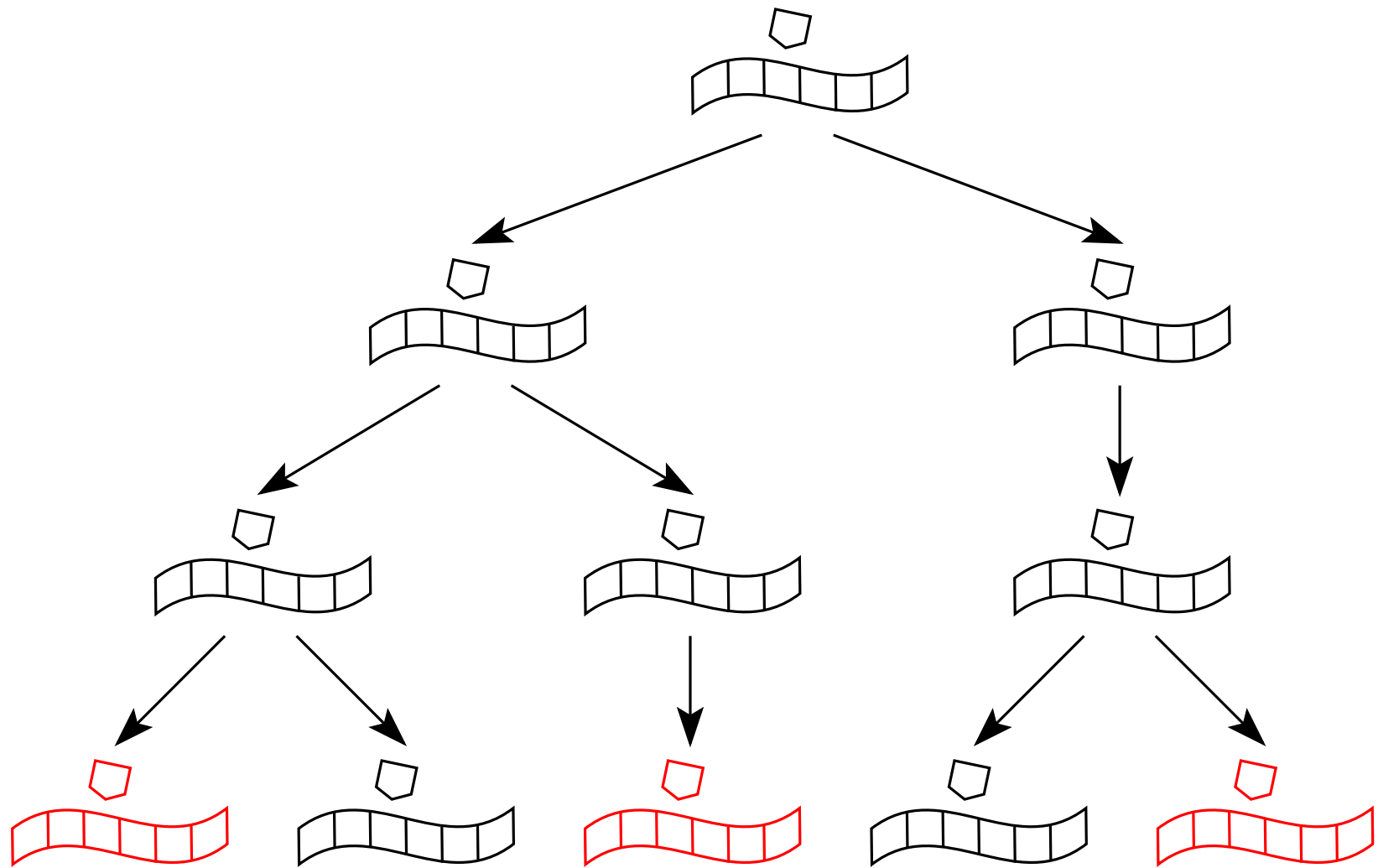
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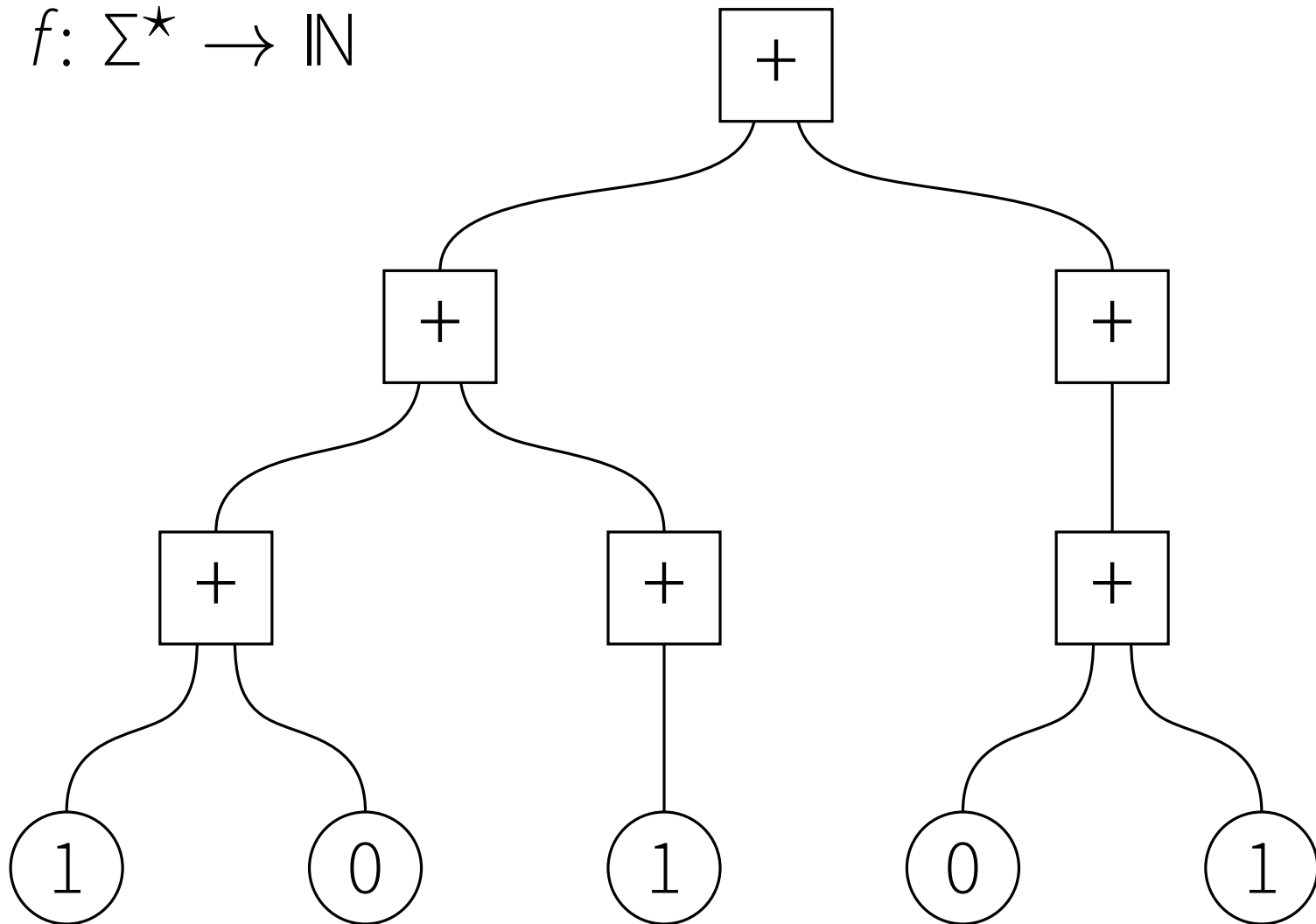
Sosík, Rodríguez-Patón: Membrane computing and complexity theory: A characterization of PSPACE. Journal of Computer and System Sciences 73(1), 137–152 (2007)

# Counting Turing machines: #P

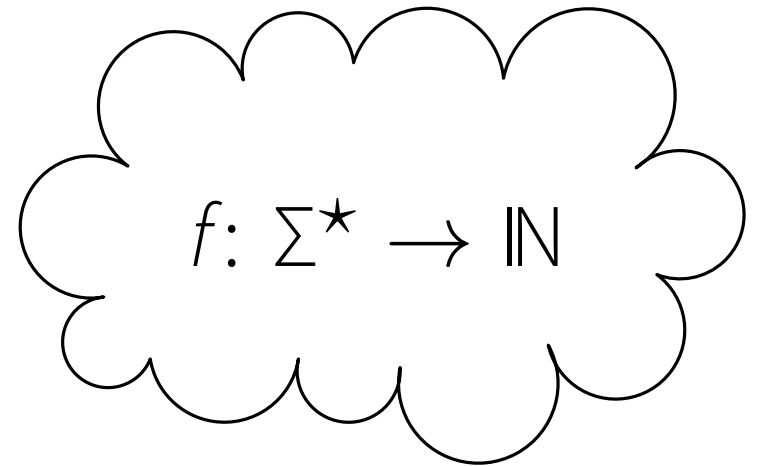
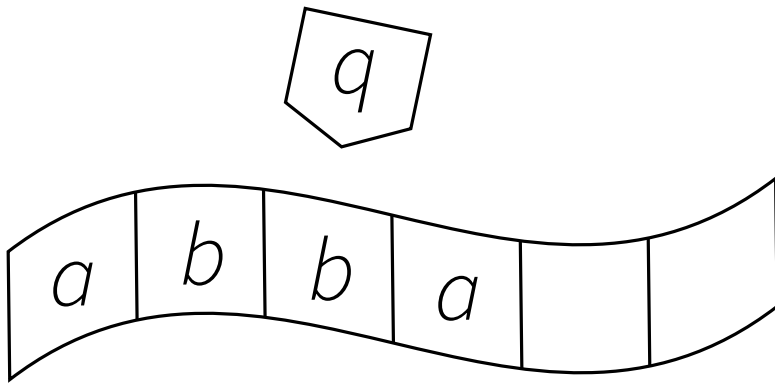


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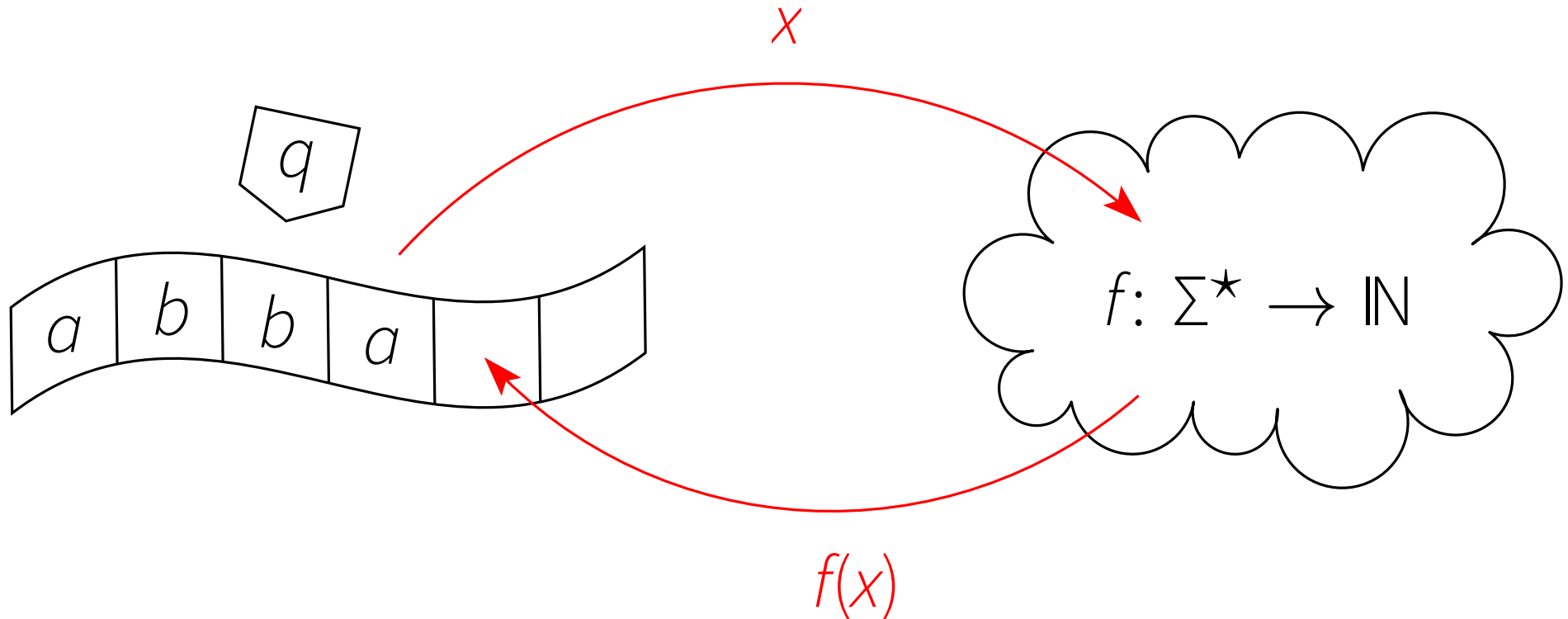
$$f: \Sigma^* \rightarrow \mathbb{N}$$



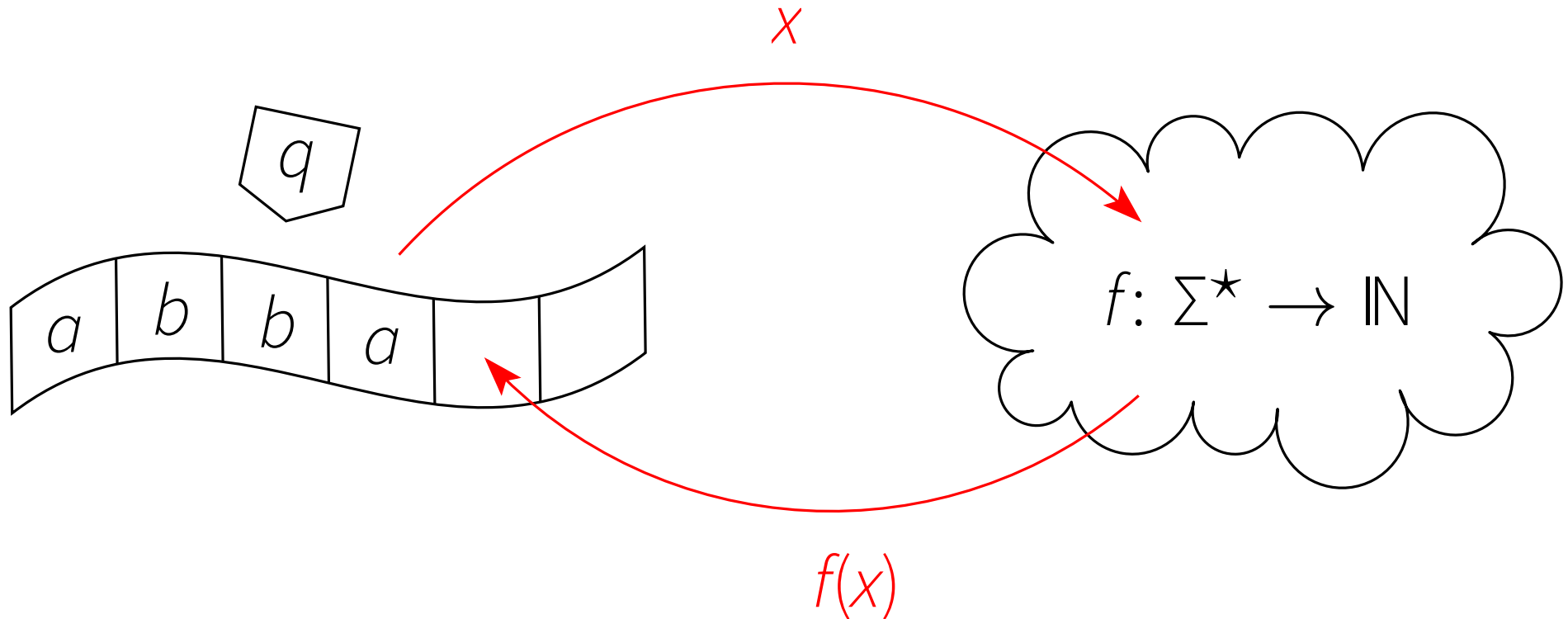
# Oracle Turing machines and the class $P^{\#P}$



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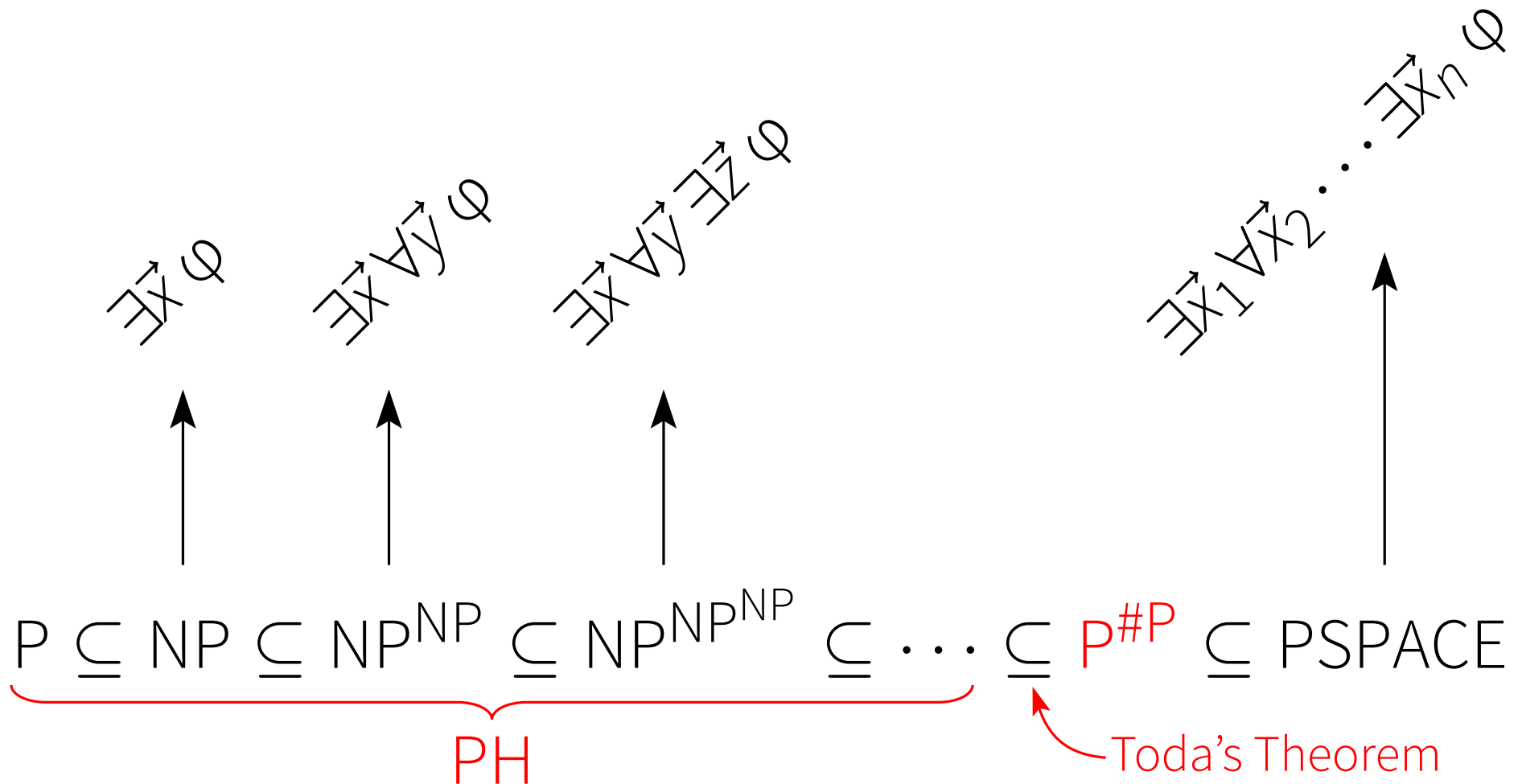
# Oracle Turing machines and the class $P^{\#P}$



$$\underbrace{P \subseteq NP \subseteq NP^{NP} \subseteq NP^{NP^{NP}} \subseteq \dots}_{PH} \subseteq P^{\#P} \subseteq PSPACE$$

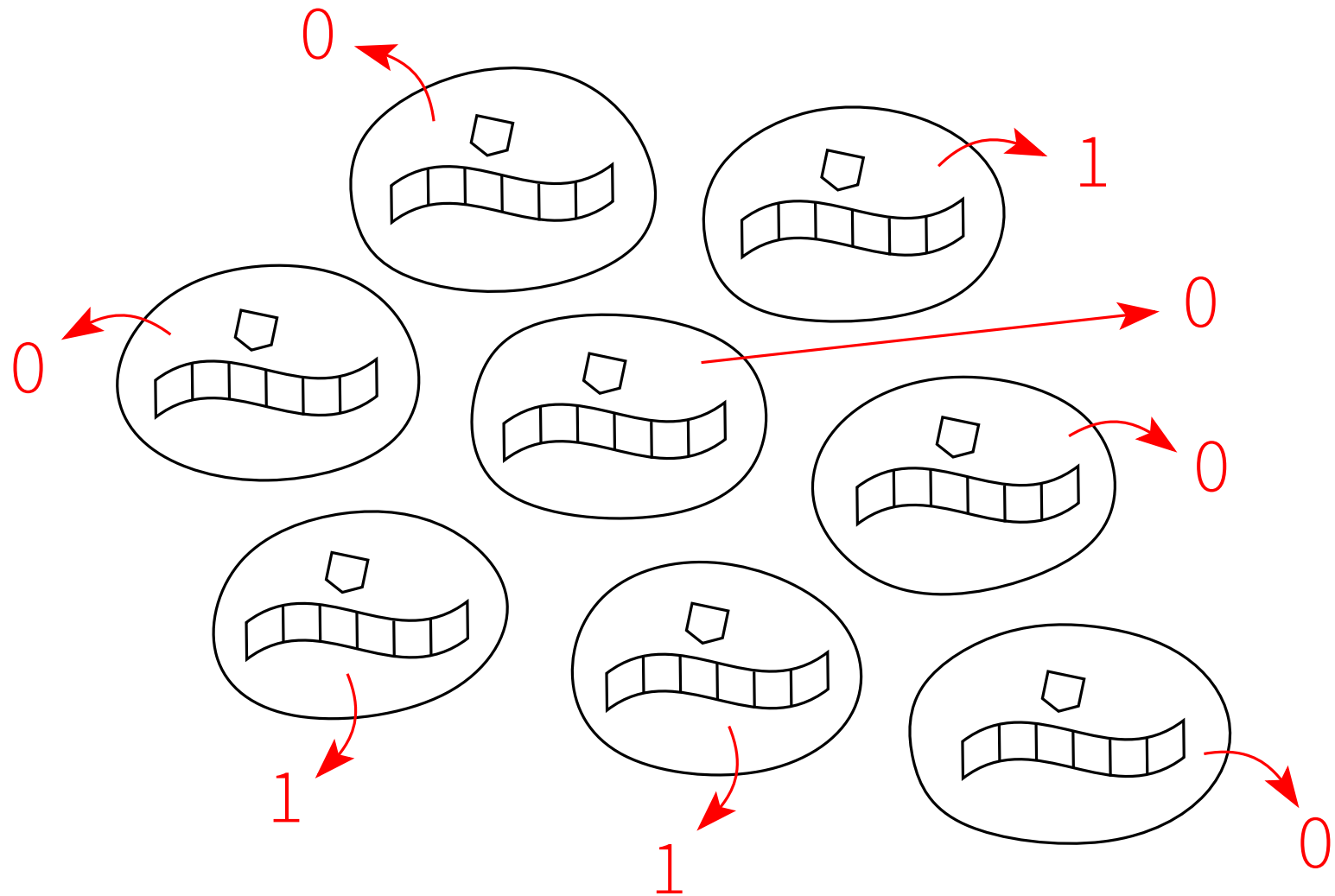
Toda's Theorem

# Oracle Turing machines and the class $P^{\#P}$



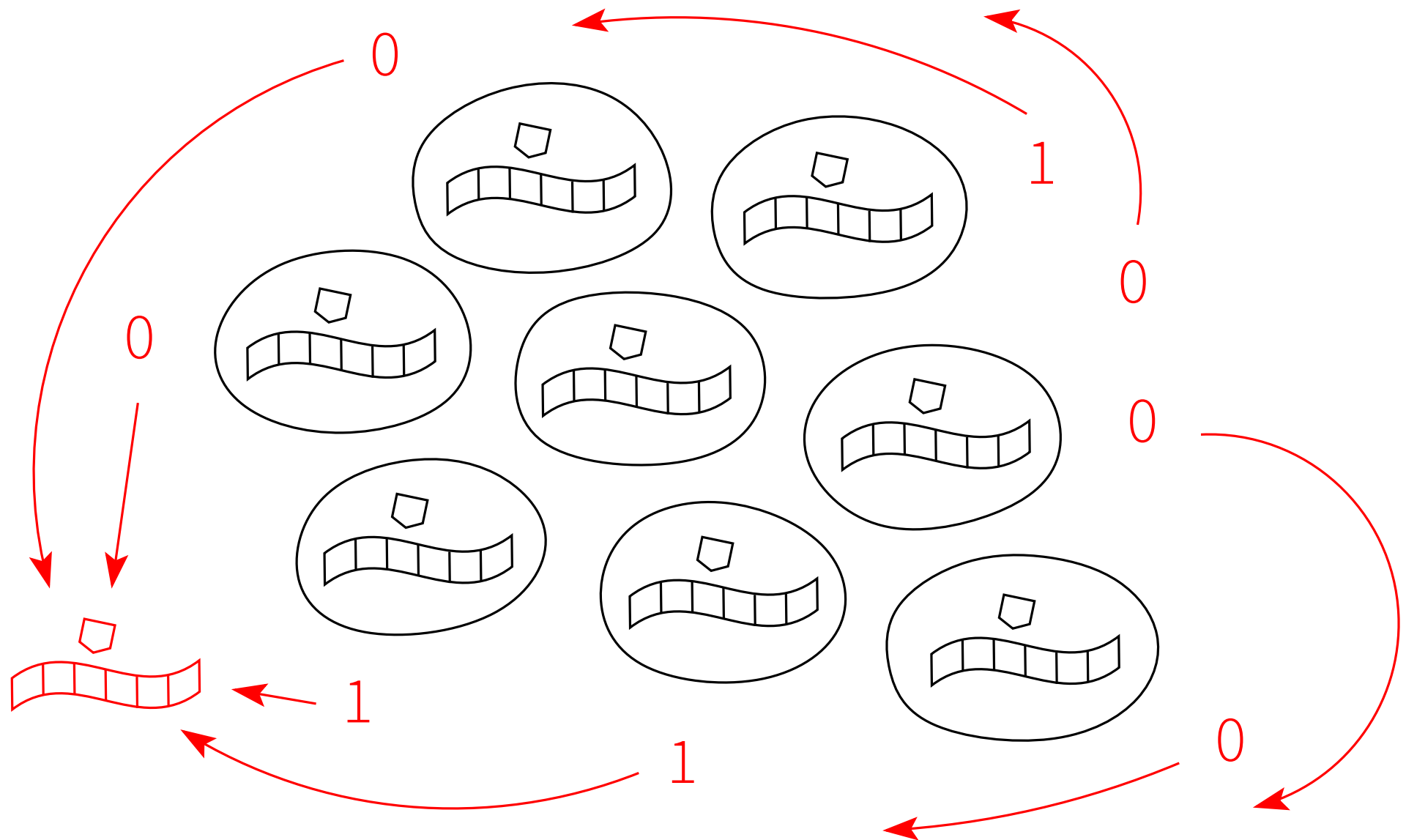


# Shallow membrane systems solve $P^{\#P}$



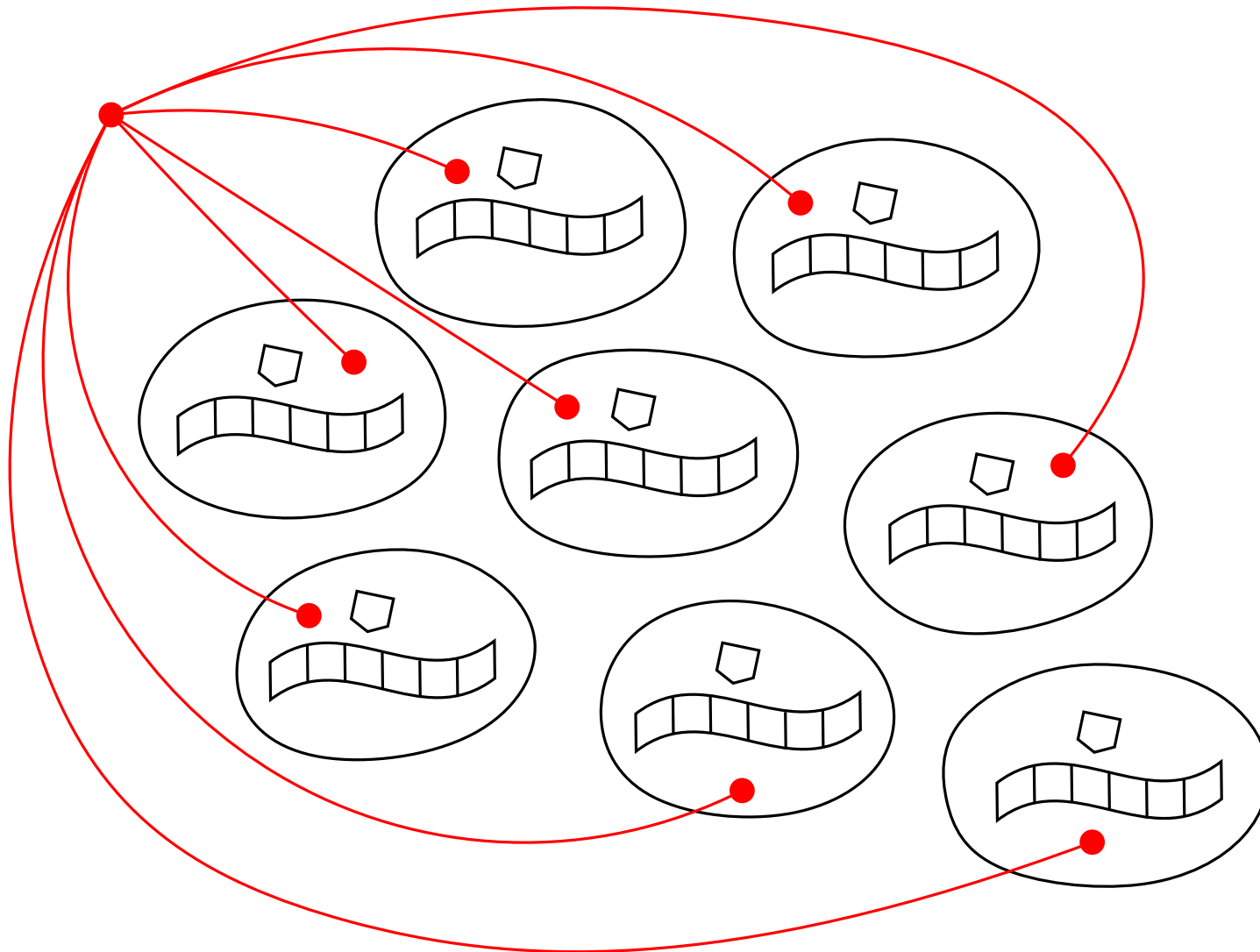
Leporati, Manzoni, Mauri, Porreca, Zandron: Characterising the complexity of tissue P systems with fission rules. Journal of Computer and System Sciences 90, 115–128 (2017)

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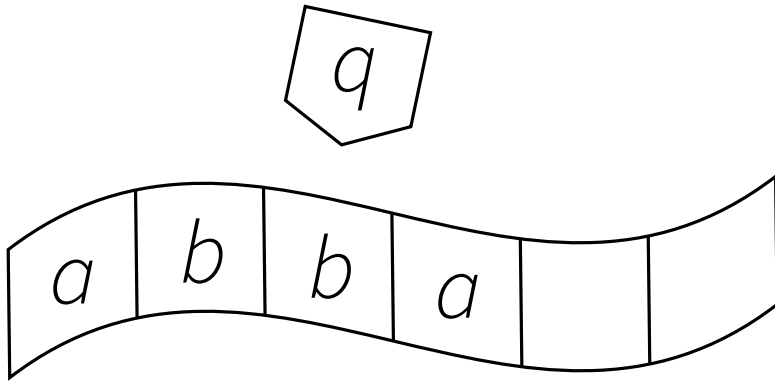
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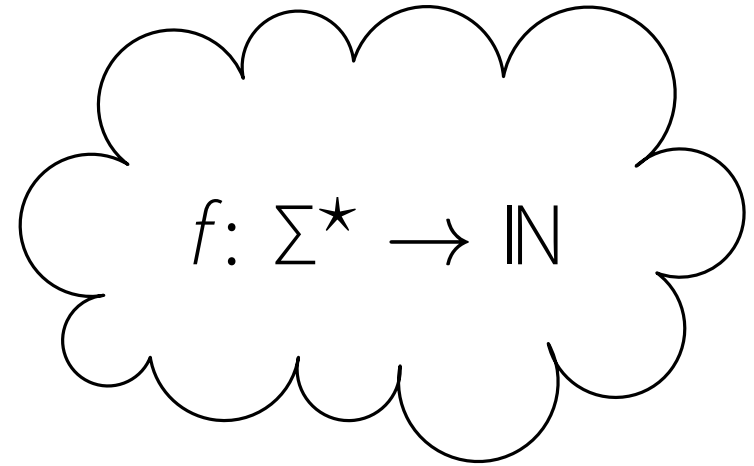


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# Simulating shallow membrane systems

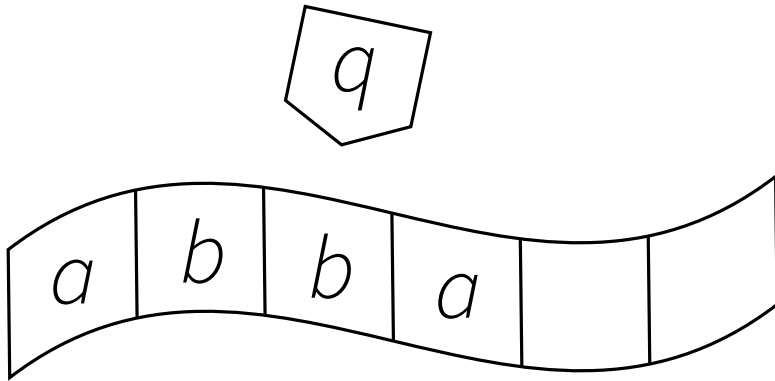


external environment

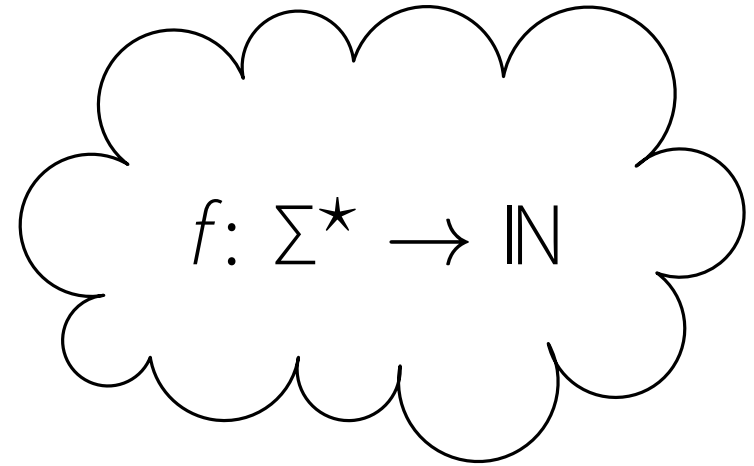


dividing membranes

# Simulating shallow membrane systems



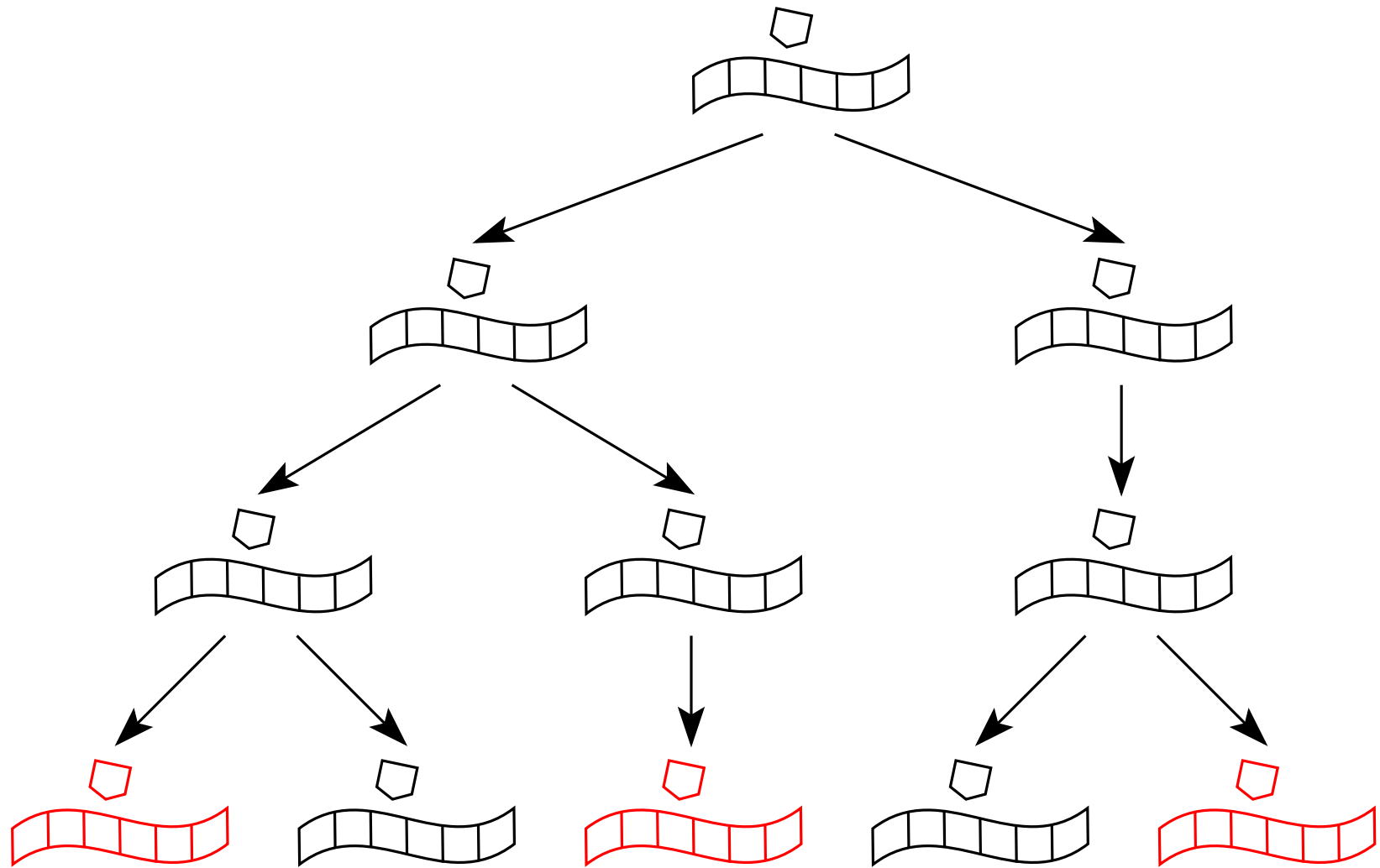
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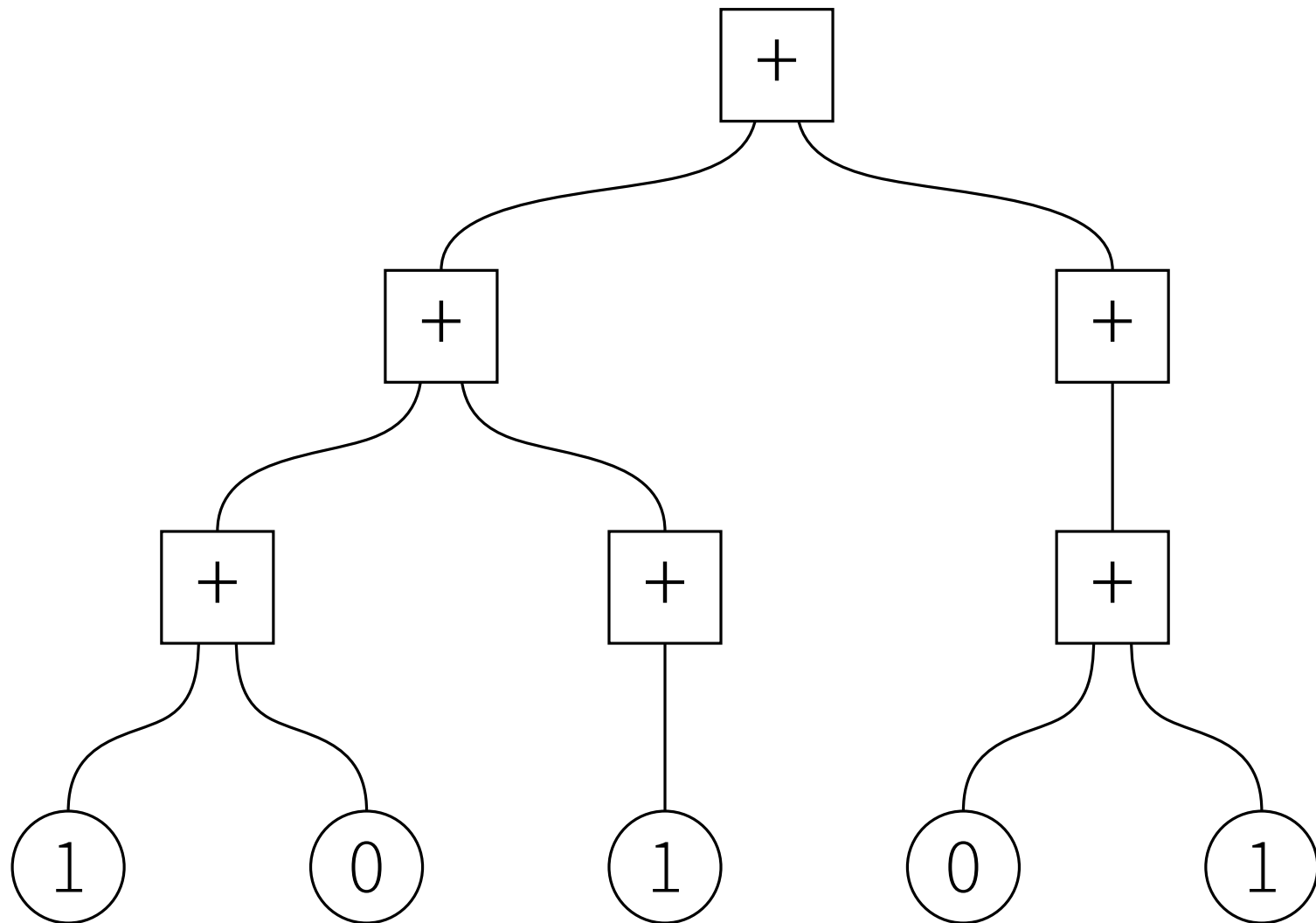
dividing membranes

$f(x) = f(x_1, \dots, x_t, a)$  = number of instances of  $a$  sent out  
at time  $t+1$  by dividing membranes  
which receive input  $x_i$  at time  $1 \leq i \leq t$

# Simulating parallelism nondeterministically



# Simulating parallelism nondeterministically



# Shallow membrane systems and $P^{\#P}$

The complexity class  $P^{\#P}$  is exactly the class of problems solved by membrane systems with **shallow (depth 1)** division in polynomial time



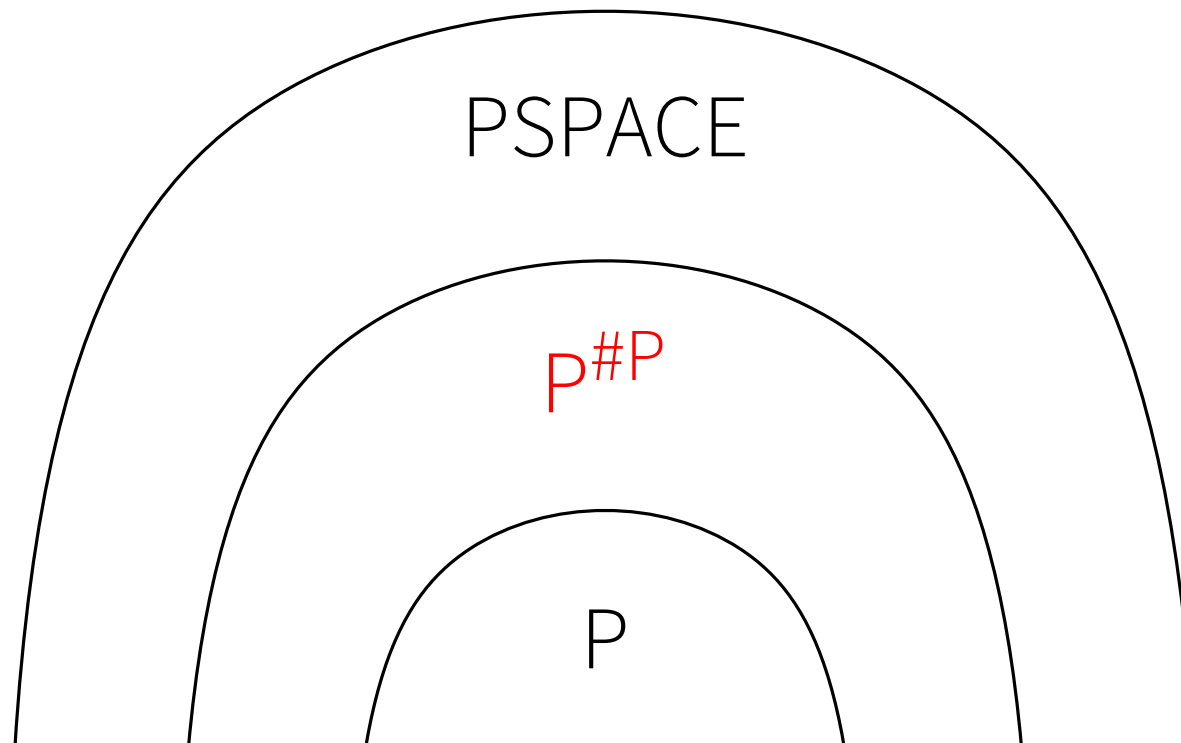
# Shallow membrane systems and $P^{\#P}$

The complexity class  $P$  is exactly the class of problems solved by membrane systems **without** division in polynomial time

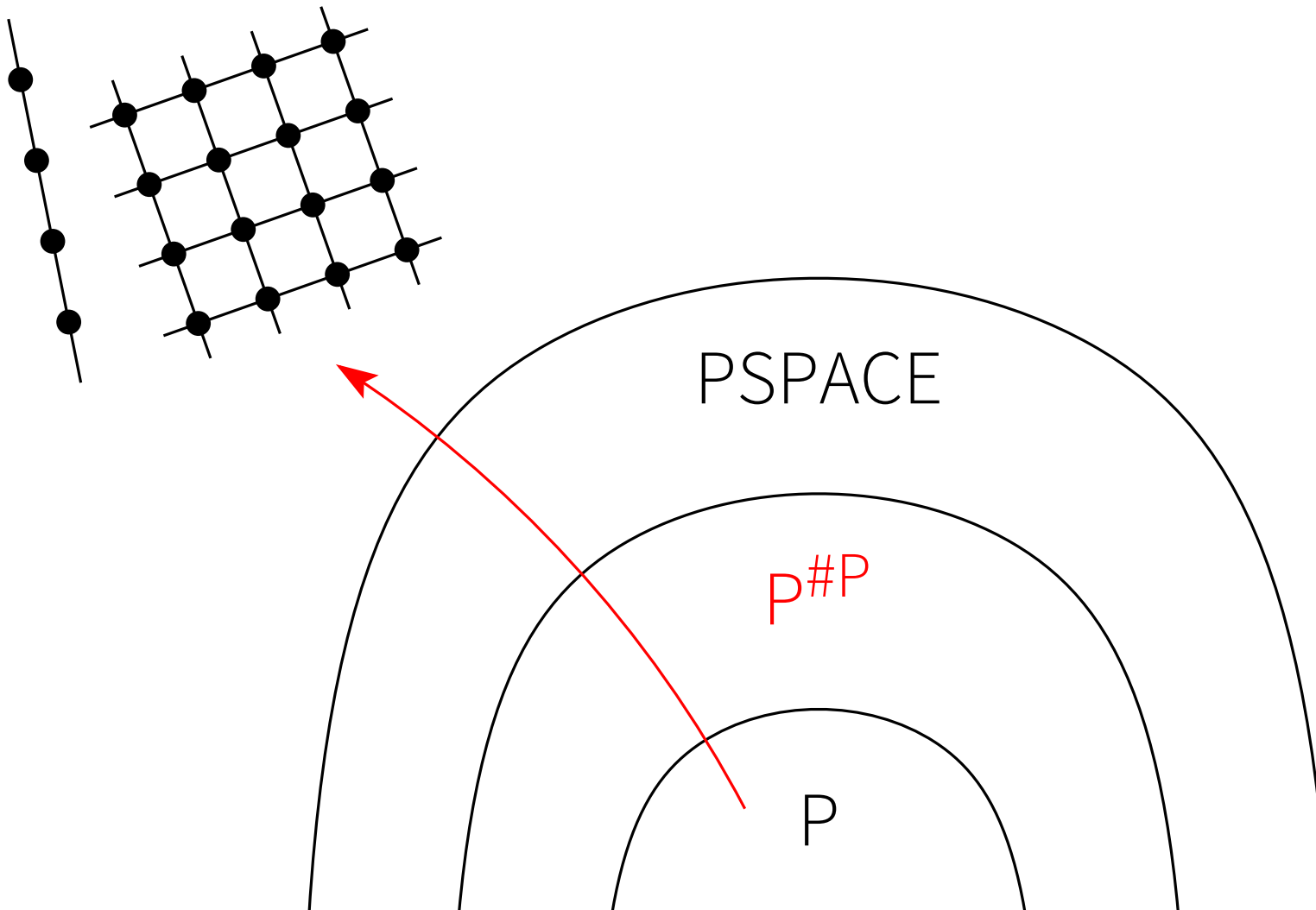
The complexity class  $P^{\#P}$  is exactly the class of problems solved by membrane systems with **shallow (depth 1)** division in polynomial time

The complexity class  $PSPACE$  is exactly the class of problems solved by membrane systems with **deep** division in polynomial time

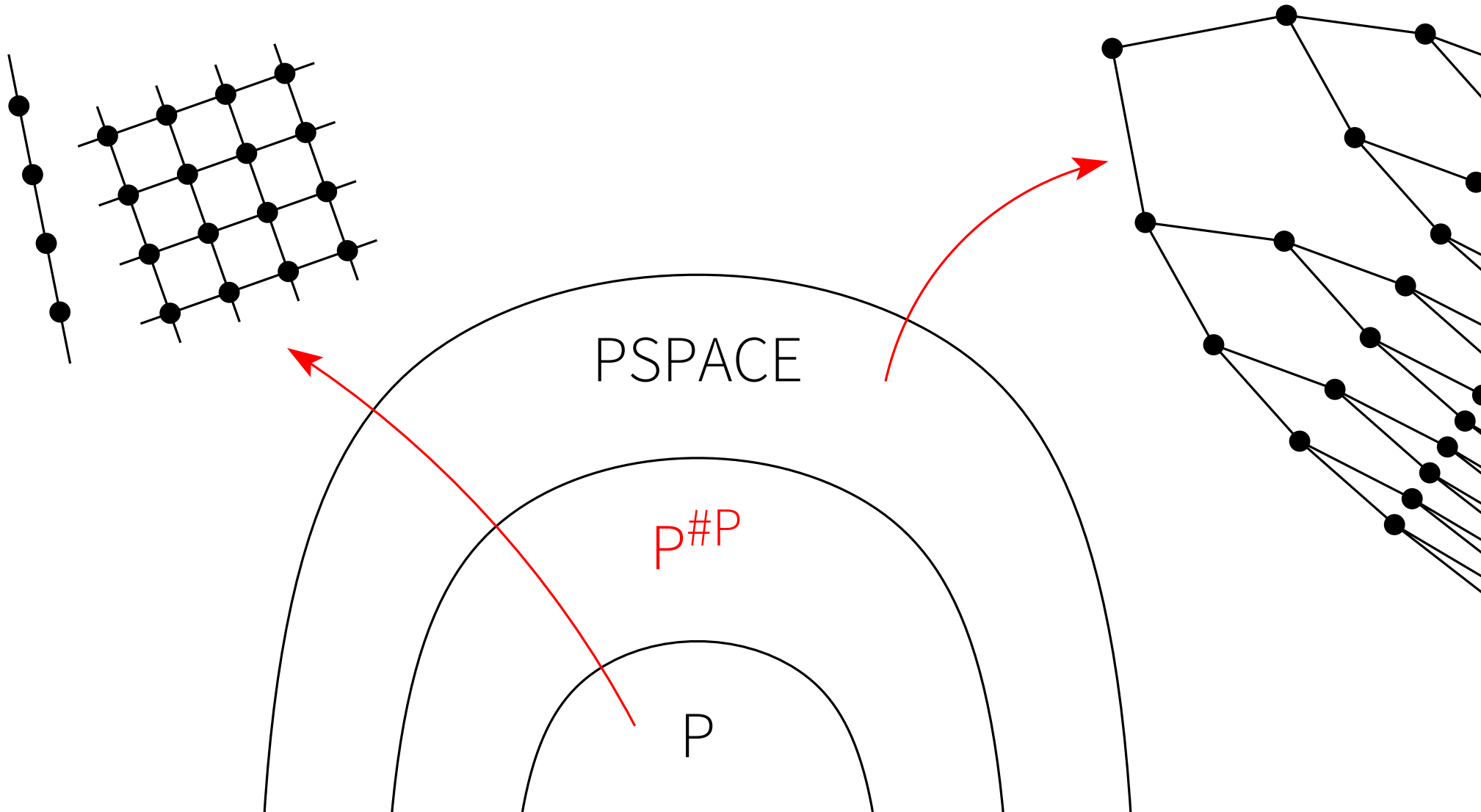
# Communication topologies and complexity



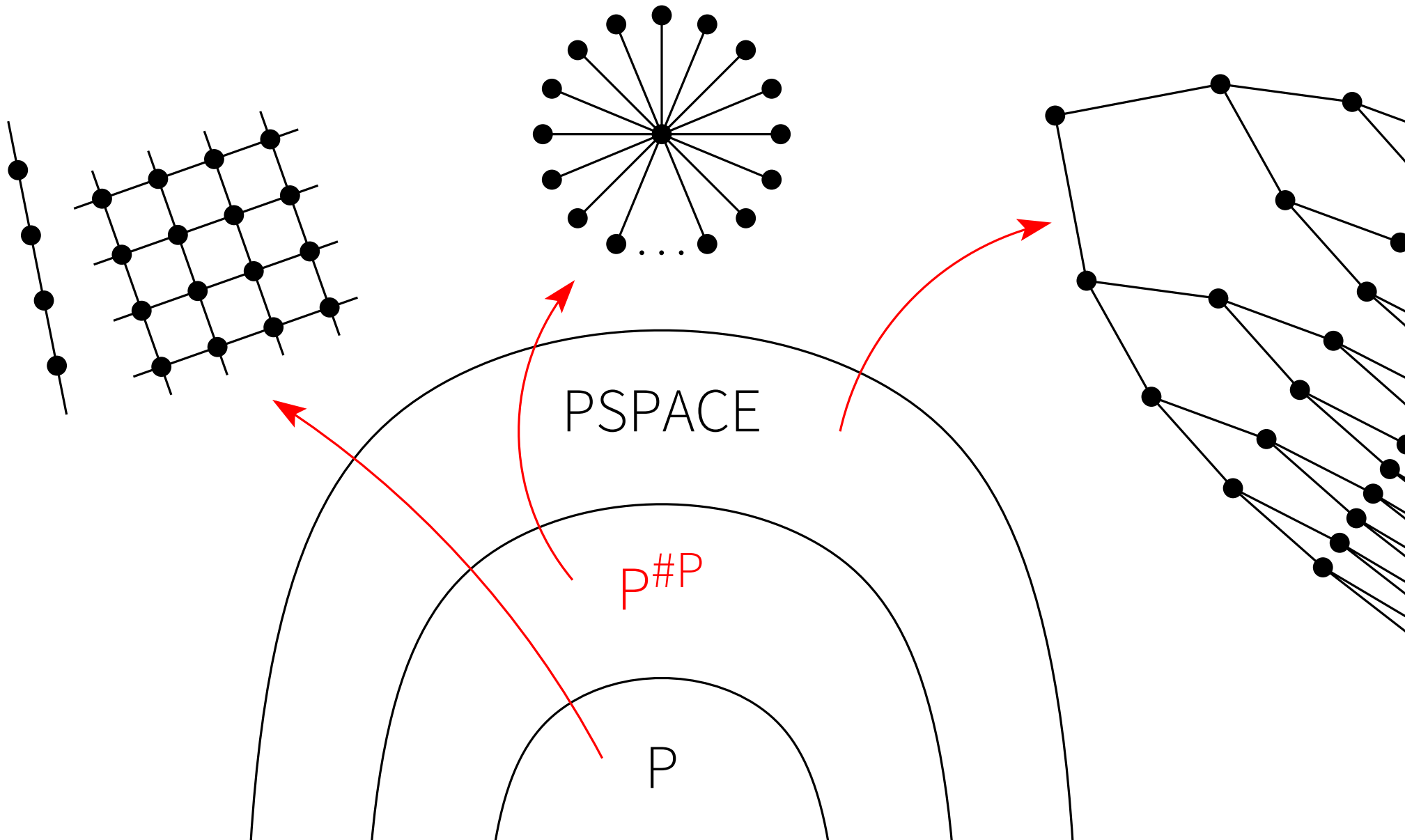
# Communication topologies and complexity



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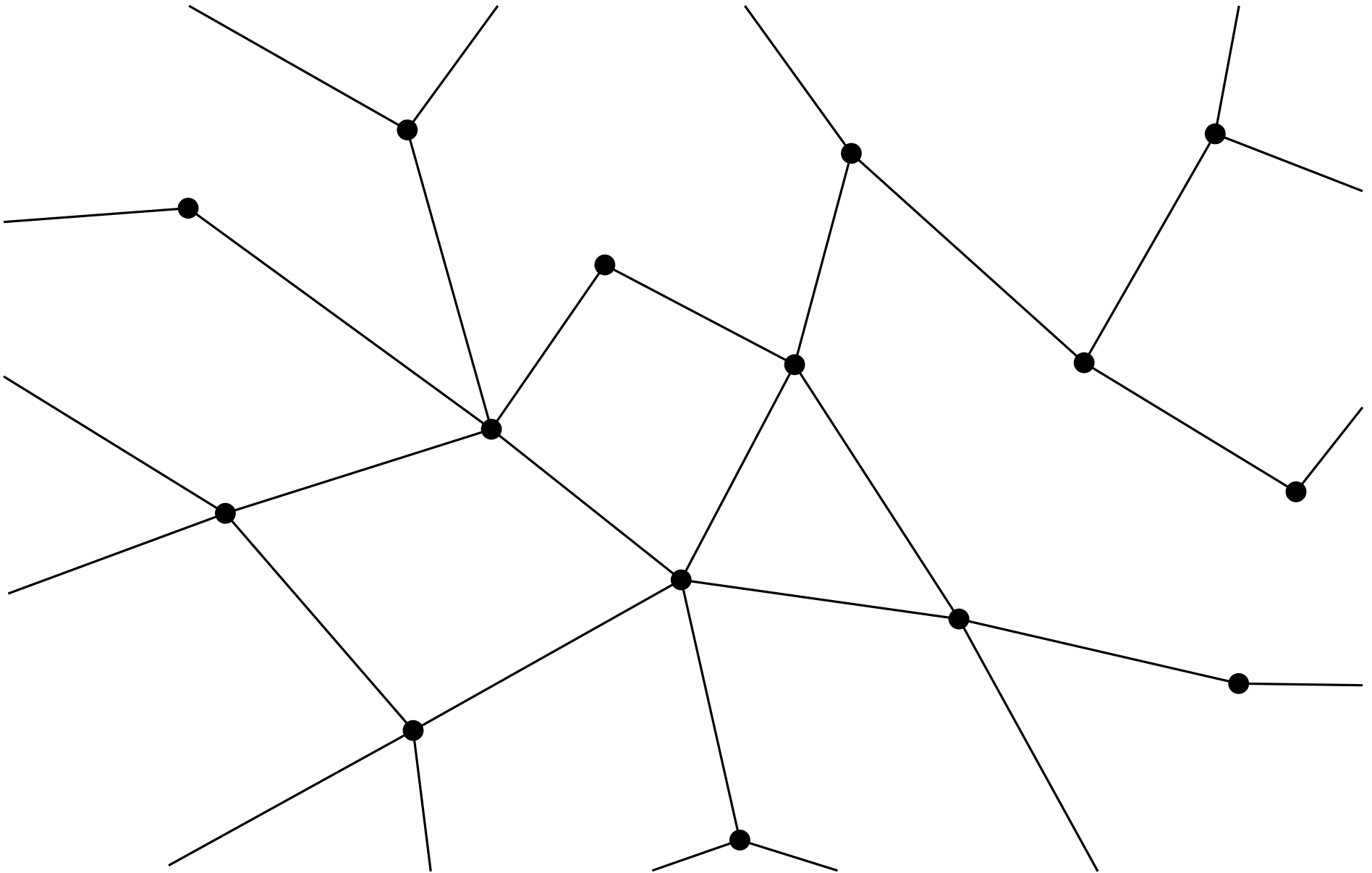


# A research project

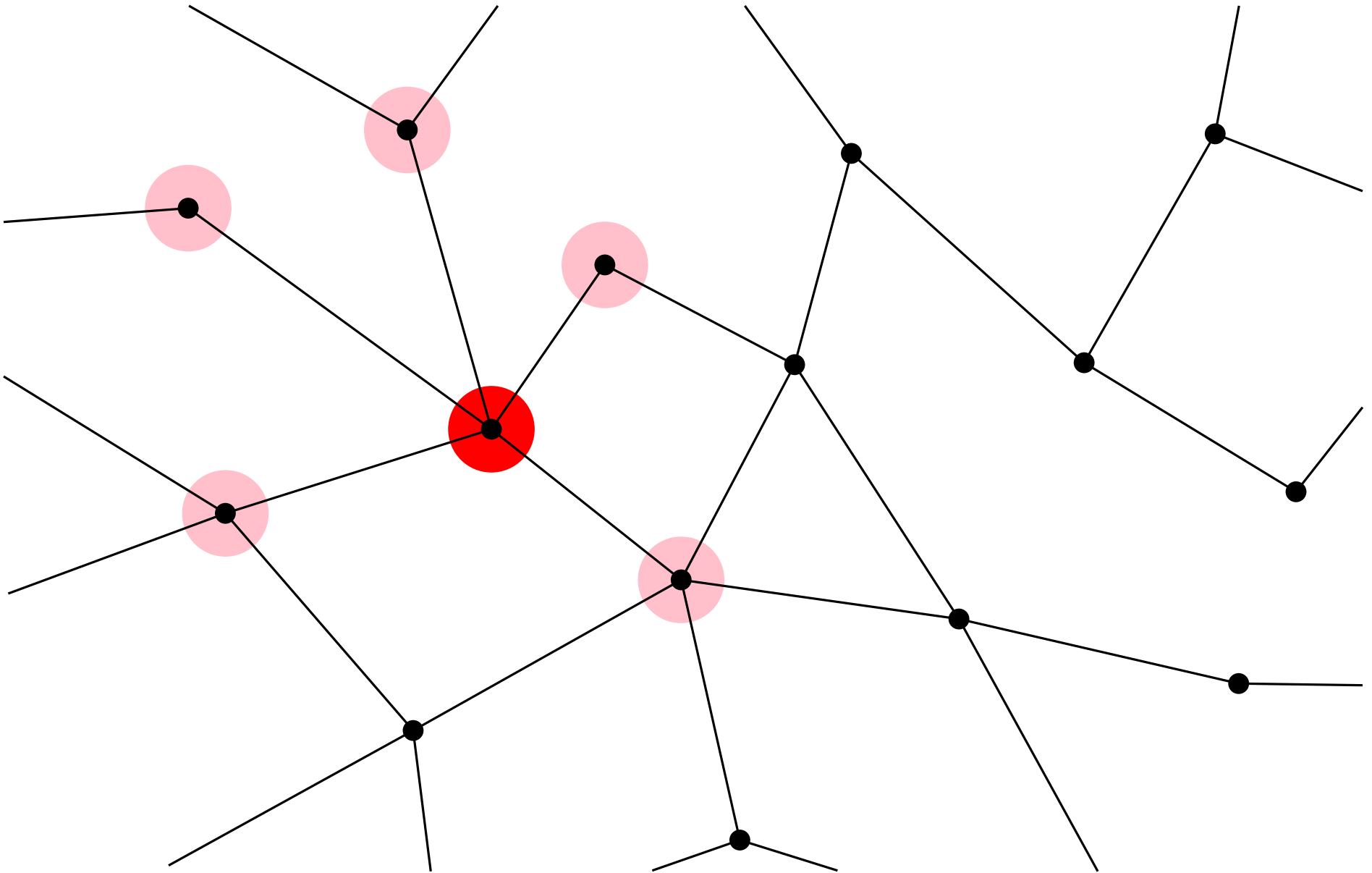
Investigate the role of communication topologies in parallel computing models, with a focus on natural computing

Which graph-theoretic, geometric, descriptive complexity properties of the communication topology influence the computational complexity?

# Automata networks over infinite graphs

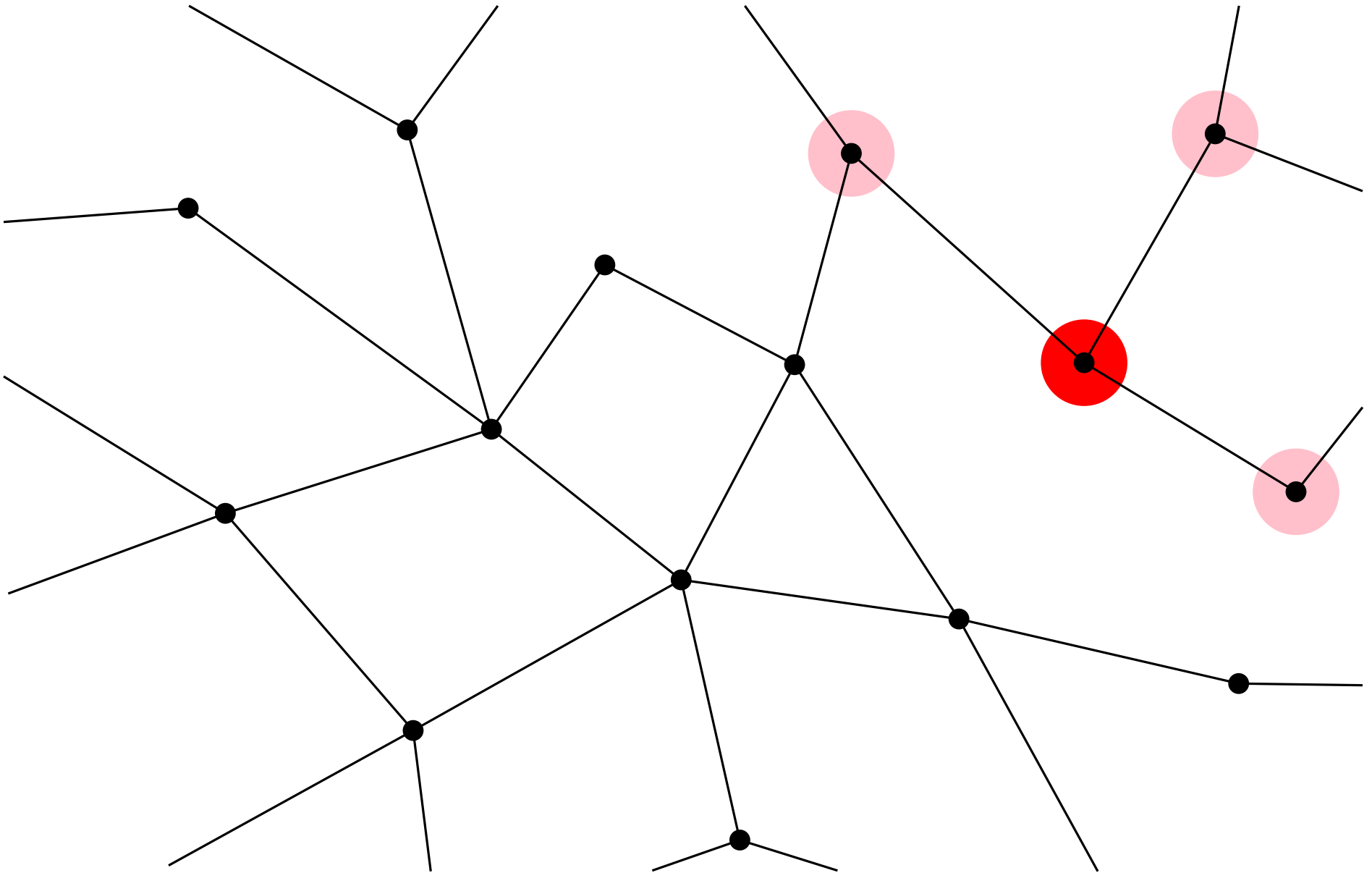


# Automata networks over infinite graphs

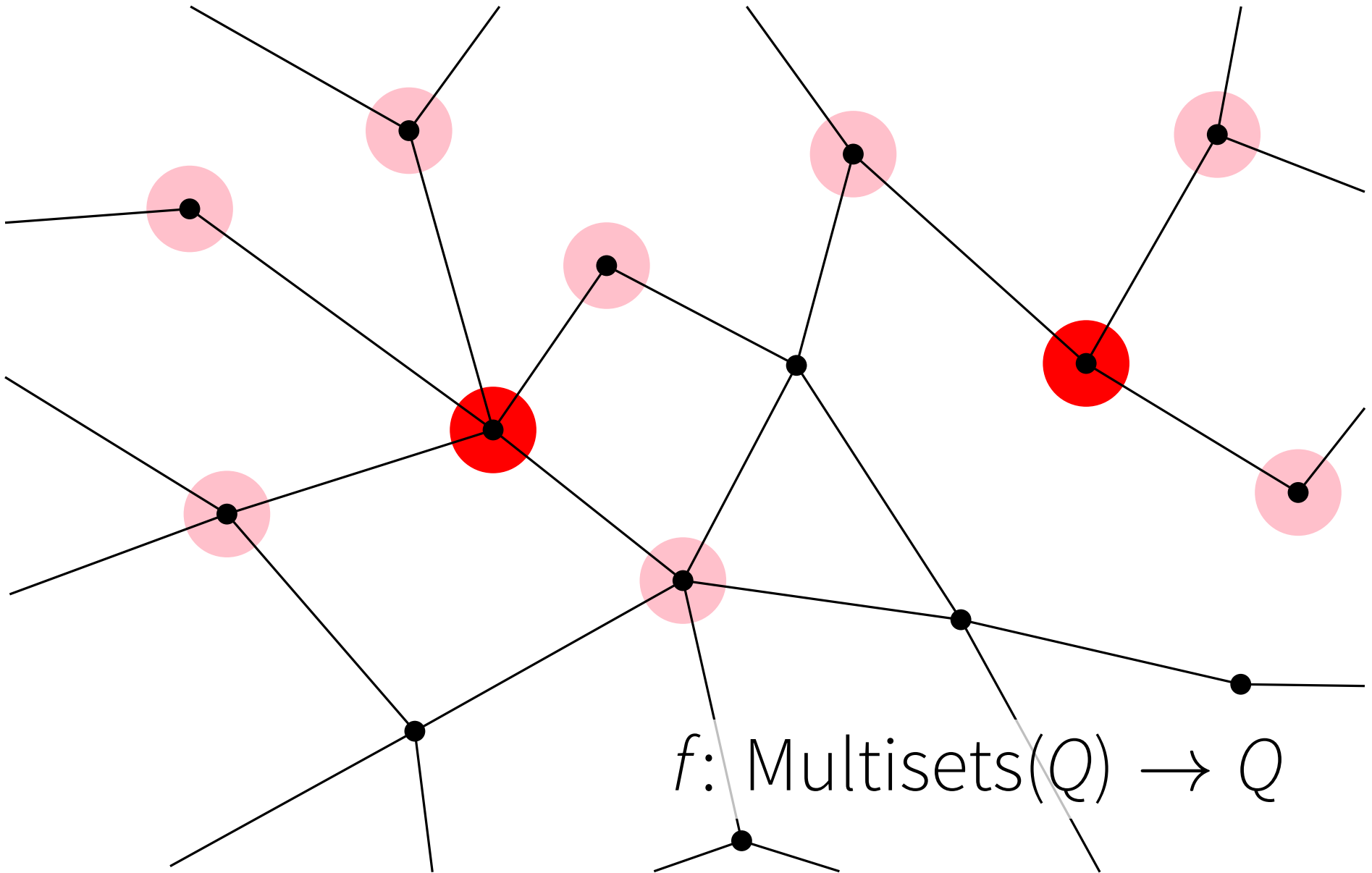




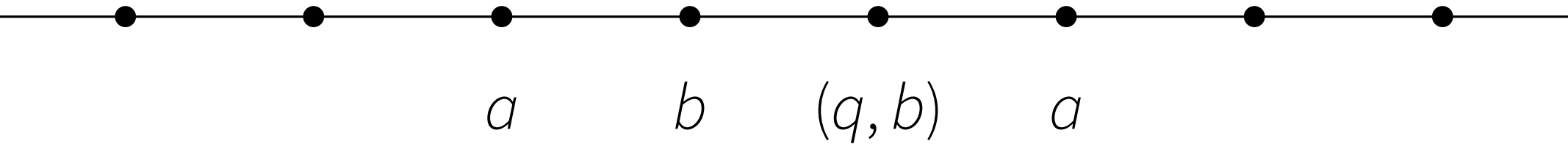
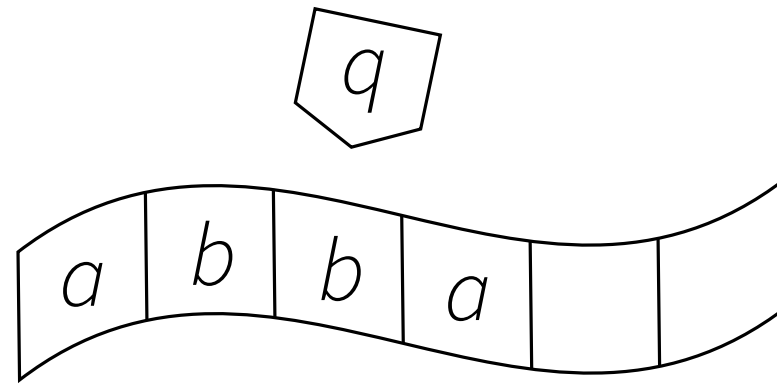
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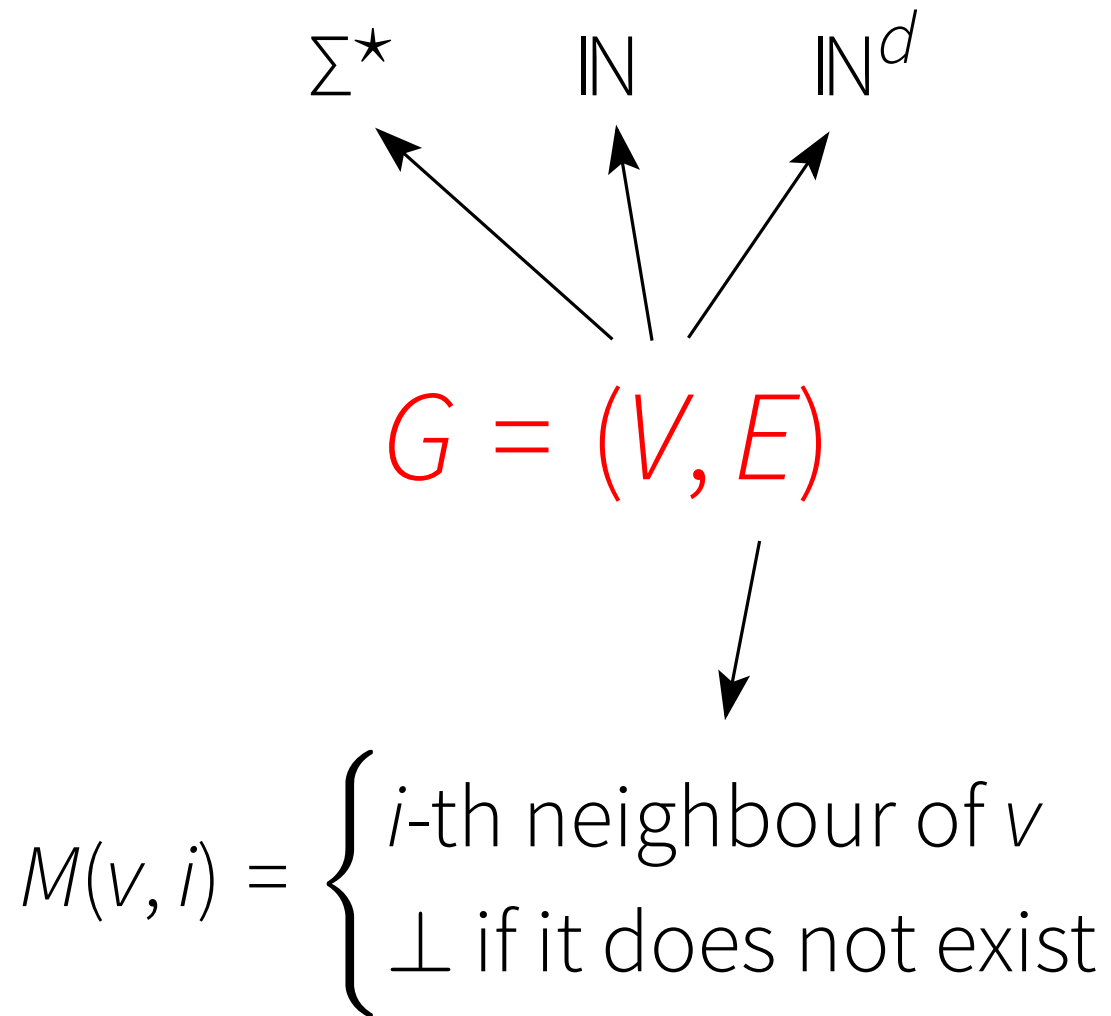
# 1D cellular automata as generalised TMs



# Descriptive complexity of infinite graphs

$$G = (V, E)$$

# Descriptive complexity of infinite graphs



# Things to do

- Choose restrictions on the local functions  
 $f: \text{Multisets}(Q) \rightarrow Q$ , e.g., threshold functions
- Choose a bound for the description complexity of the underlying graph
- Choose a way to encode the input in the initial configuration

# Generalised complexity classes

- $P(G)$  = problems solved in polynomial time over  $G$
- $PSPACE(G)$  = polynomial space over  $G$
- $EXPTIME(G)$  = exponential time over  $G$
- ...
- $LOGTIME(G)$  = logarithmic time over  $G$
- $NP(G)$  = nondeterministic polynomial time over  $G$

# Preliminary results

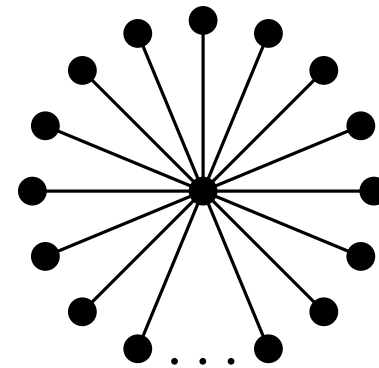
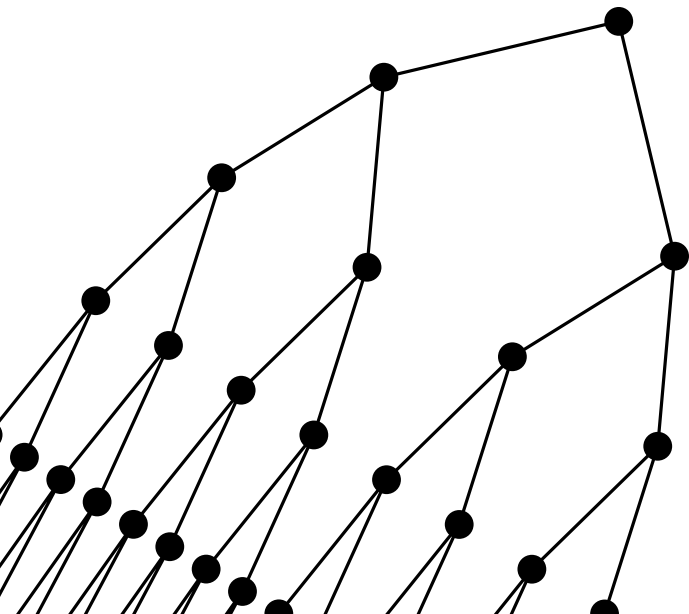
- $P(G) = P$
- $PSPACE(G) = PSPACE$
- $EXPTIME(G) = EXPTIME$

where  $G$  is the **linear** graph or, more generally,  
an **efficiently navigable** graph embeddable  
in the **Euclidean space**



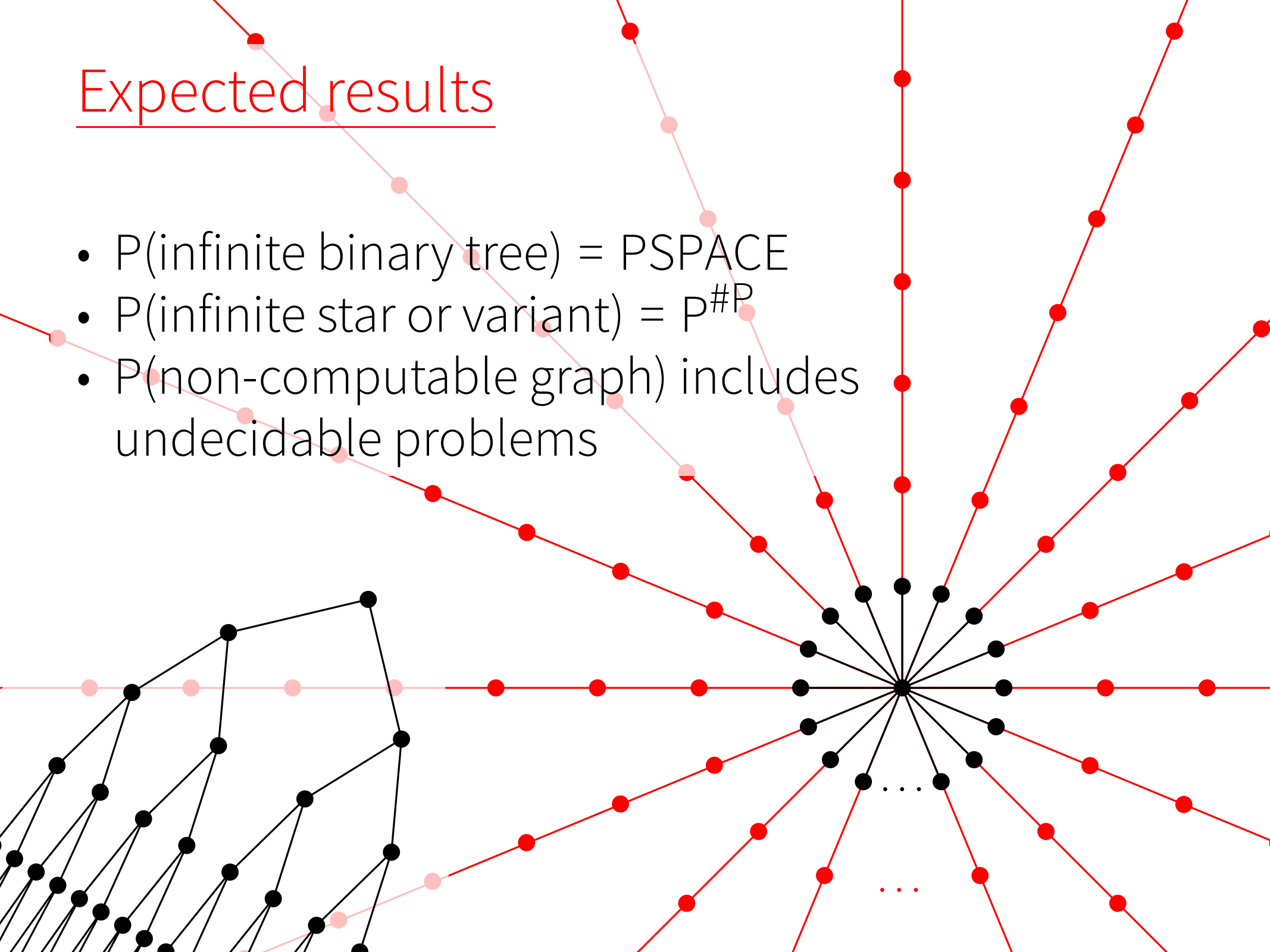
# Expected results

- $P(\text{infinite binary tree}) = PSPACE$
- $P(\text{infinite star or variant}) = P^{\#P}$
- $P(\text{non-computable graph})$  includes undecidable problems



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- $P(\text{infinite binary tree}) = PSPACE$
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- $P(\text{non-computable graph})$  includes undecidable problems



# Long-term goals: theory of computation

- Find graphs characterising the standard complexity classes
- ...or prove that it is impossible
- Define new complexity classes in terms of “natural” graphs
- Examine the complexity of simulating automata networks over certain graphs

# Long-term goals: distributed algorithms

- Find low-level algorithms (local rules) working with all graphs or certain families of graphs
- ...or prove impossibility results
- Investigate how the graph-theoretic and geometric properties can speed up or slow down the algorithms
- Example: algorithms running in time  $\Theta(n^{f(d)})$  in  $\mathbb{R}^d$

# Connections with other areas

- Exploit results on the dynamics of automata networks
- Provide bounds for applications of automata networks (e.g., biological modelling)
- Theory (and practice) of distributed algorithms
- Machine learning (e.g., “non-Euclidean learning”)
- Attack open problems in complexity theory

Thanks for your attention!

Merci de votre attention !

Any questions?