

Counting complexity and oracles in natural computing

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Outline

- Complexity classes for parallel computing models
- Cellular automata in various geometric spaces
- Membrane systems
- Expanding cellular automata
- Conclusions and future work

**The first and second
machine classes**

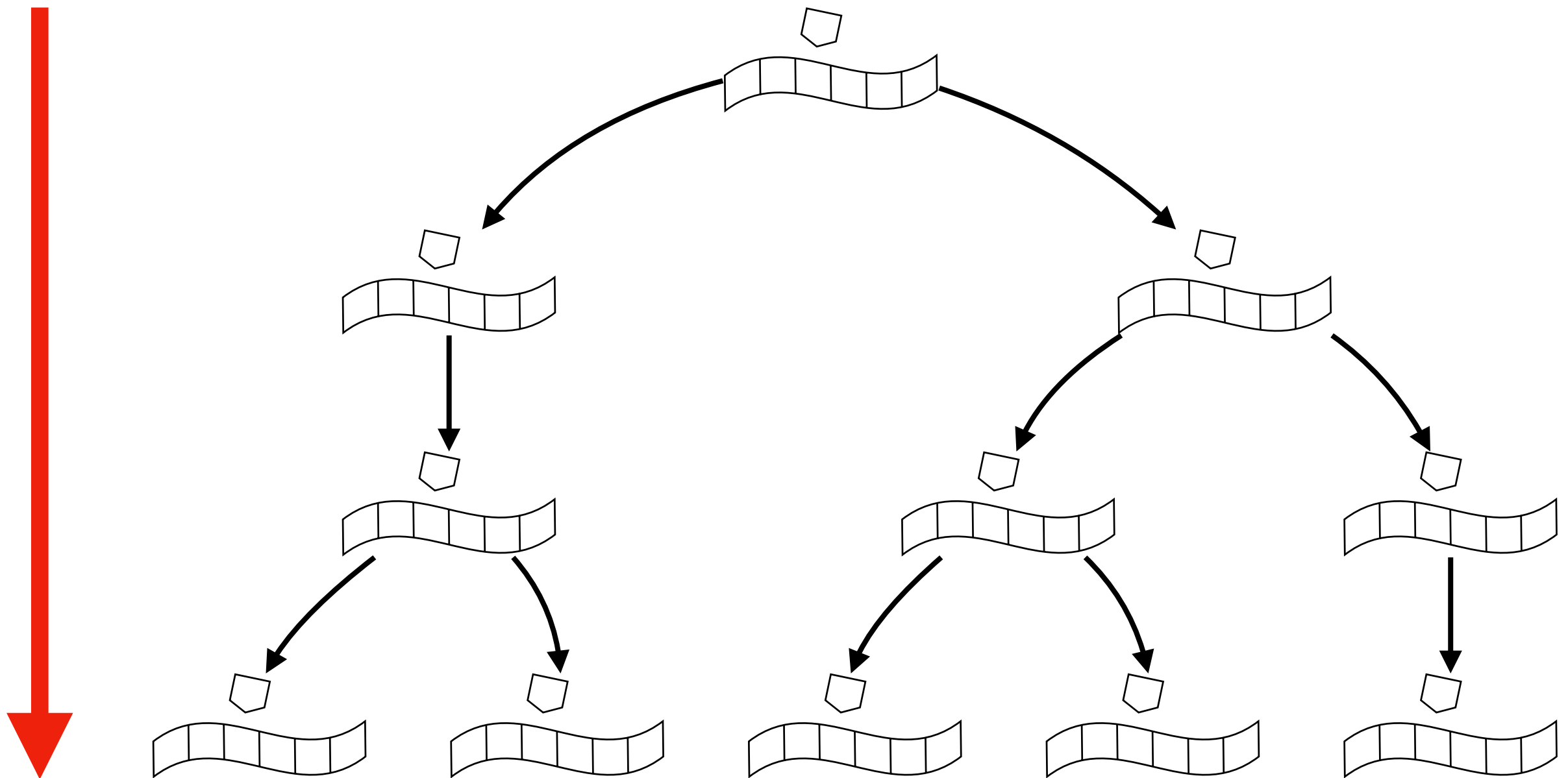
The first machine class and P

- Includes the **deterministic Turing machine** and all models that simulate and are simulated by it efficiently:
 - Random access machines (**RAM**) with constant-time **addition** and **subtraction**
 - **Cellular automata** with a finite initial configuration

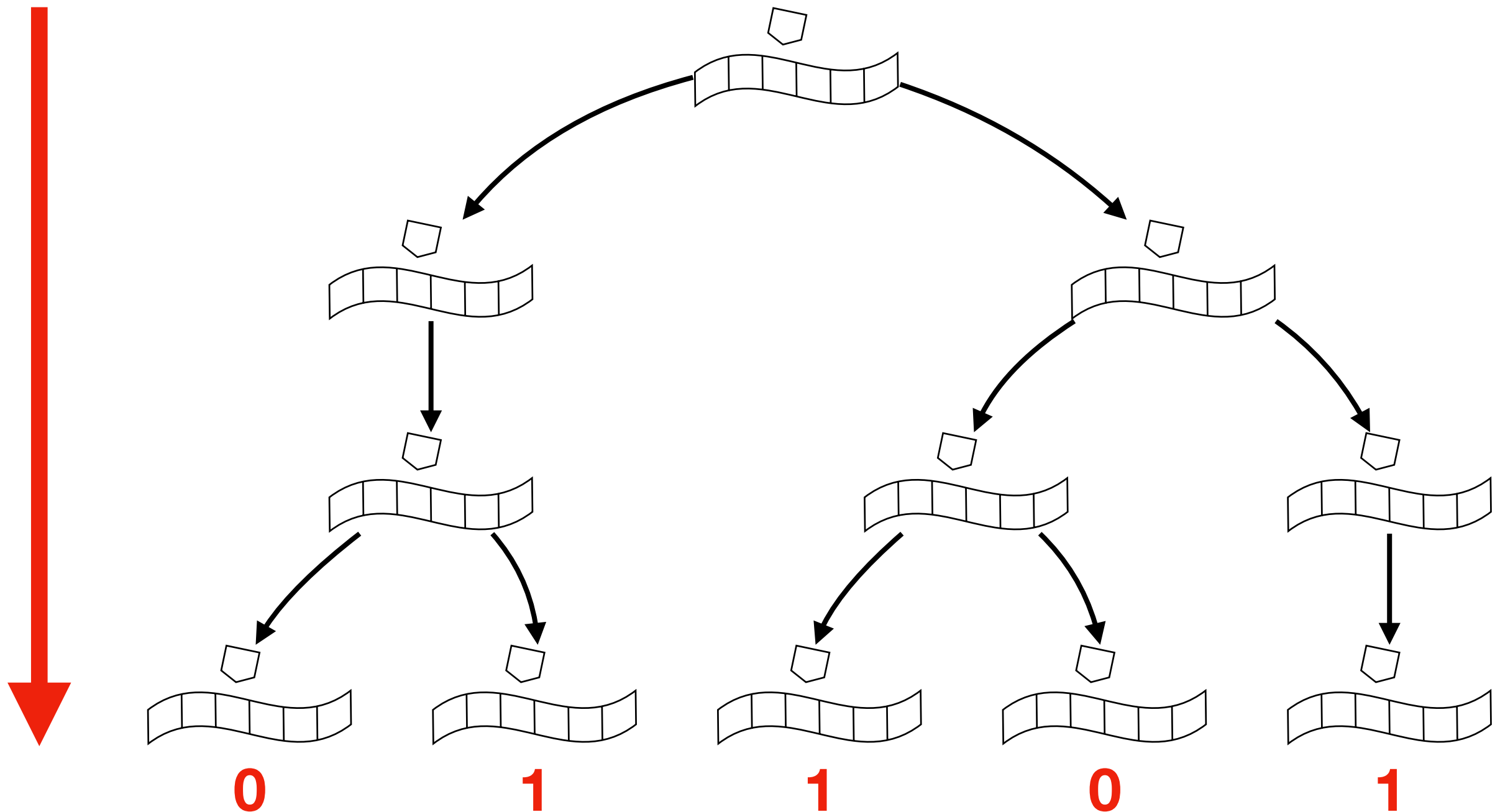
The second machine class and PSPACE

- Includes models of computation that solve in **polynomial time** what a Turing machine solves in **polynomial space**:
 - **Alternating** Turing machines
 - **Random access machines** including constant time **multiplication** and **division**
 - Parallel processes generated by **fork(2)** running on an **unbounded** number of processors
 - **Cellular automata** over **hyperbolic** grids

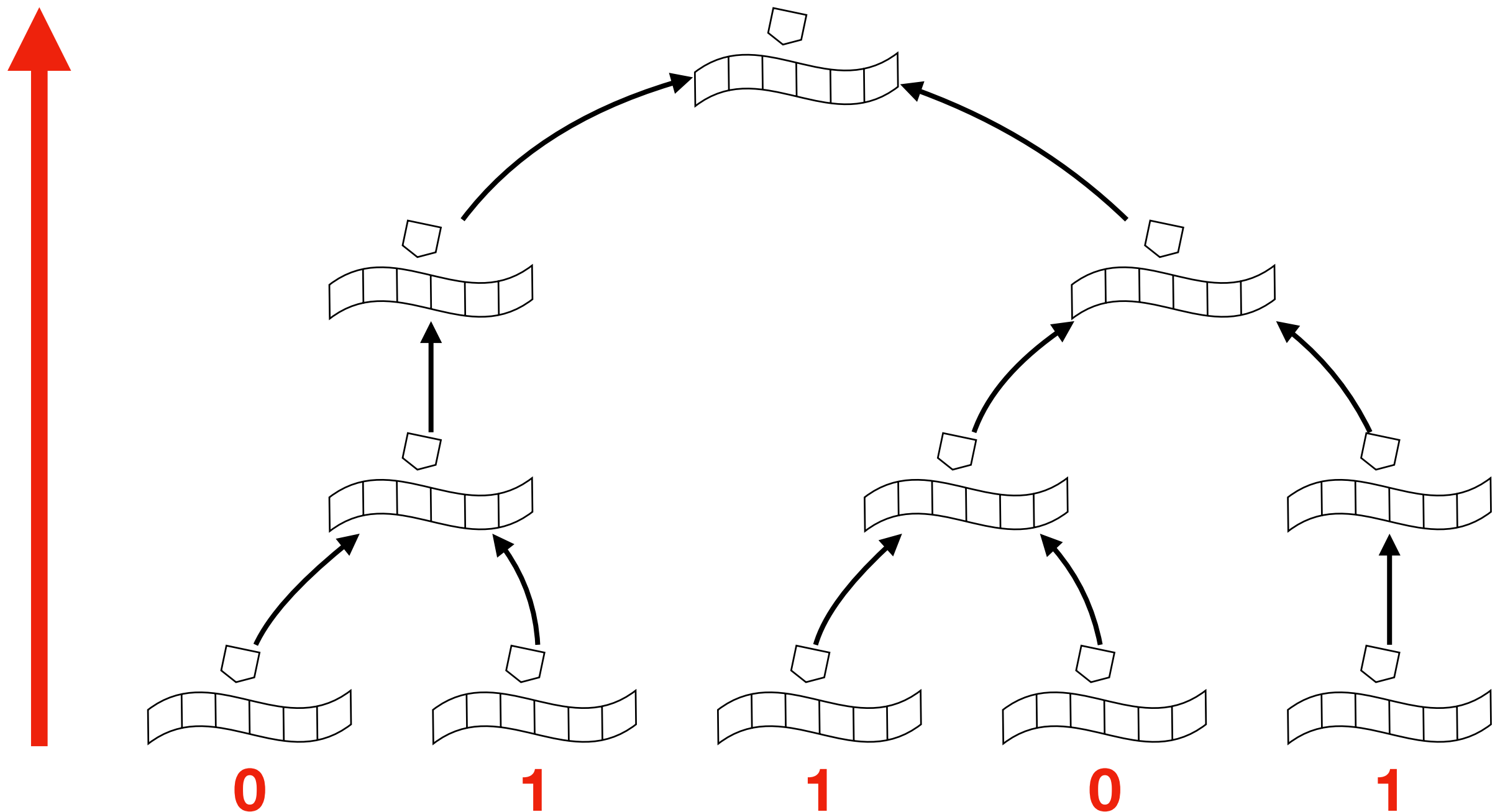
Nondeterministic Turing machines: NP



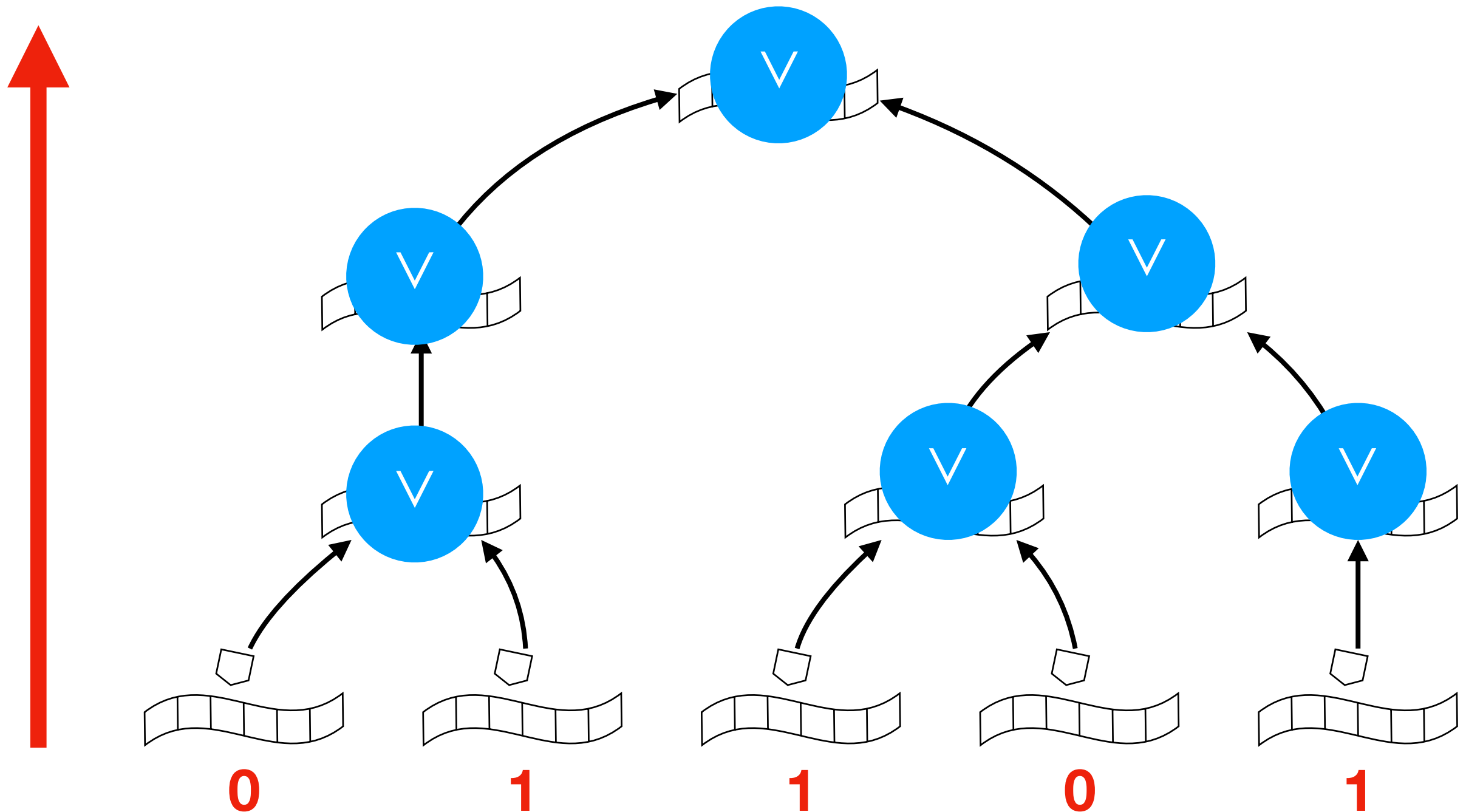
Nondeterministic Turing machines: NP



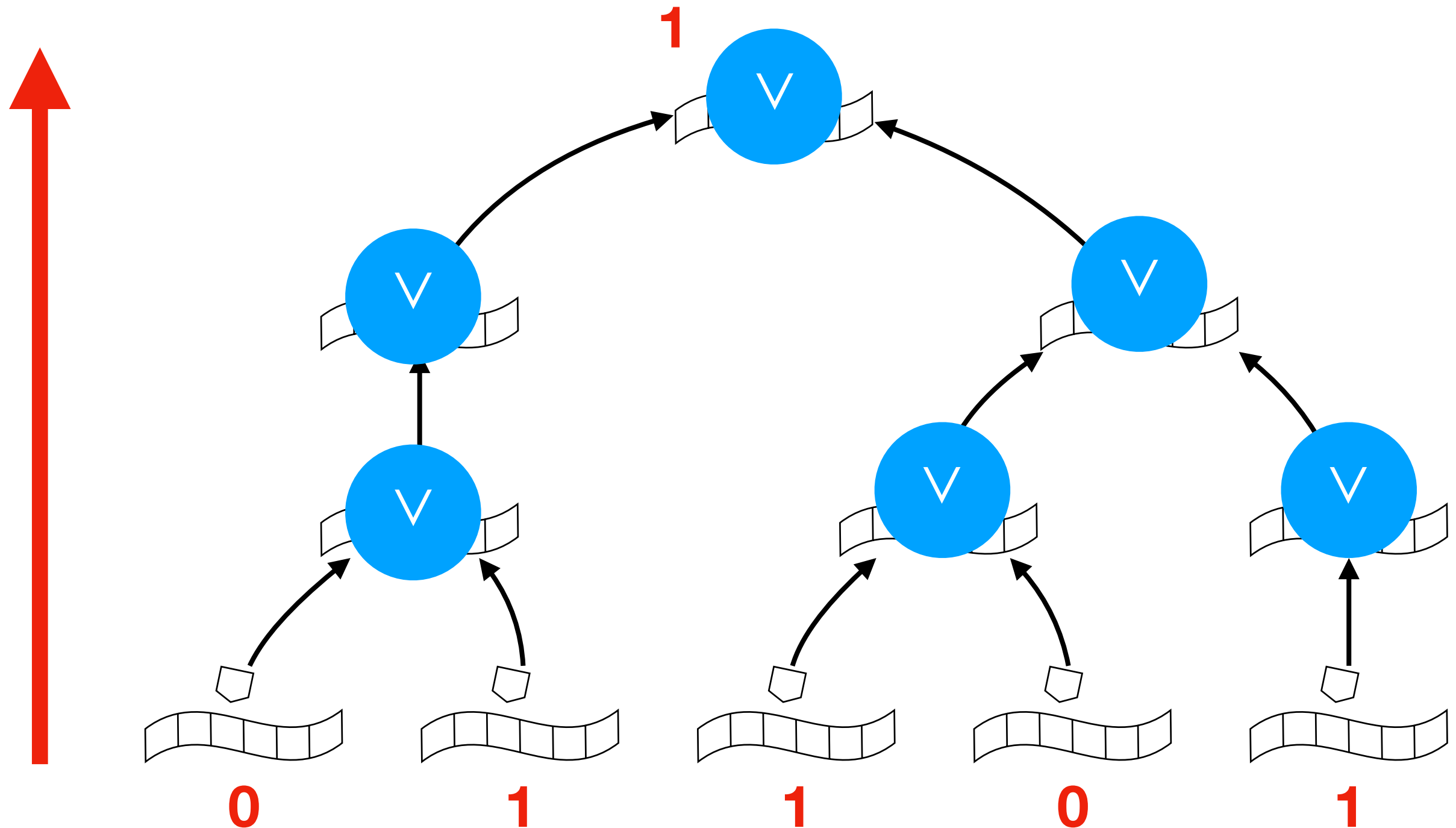
Nondeterministic Turing machines: NP



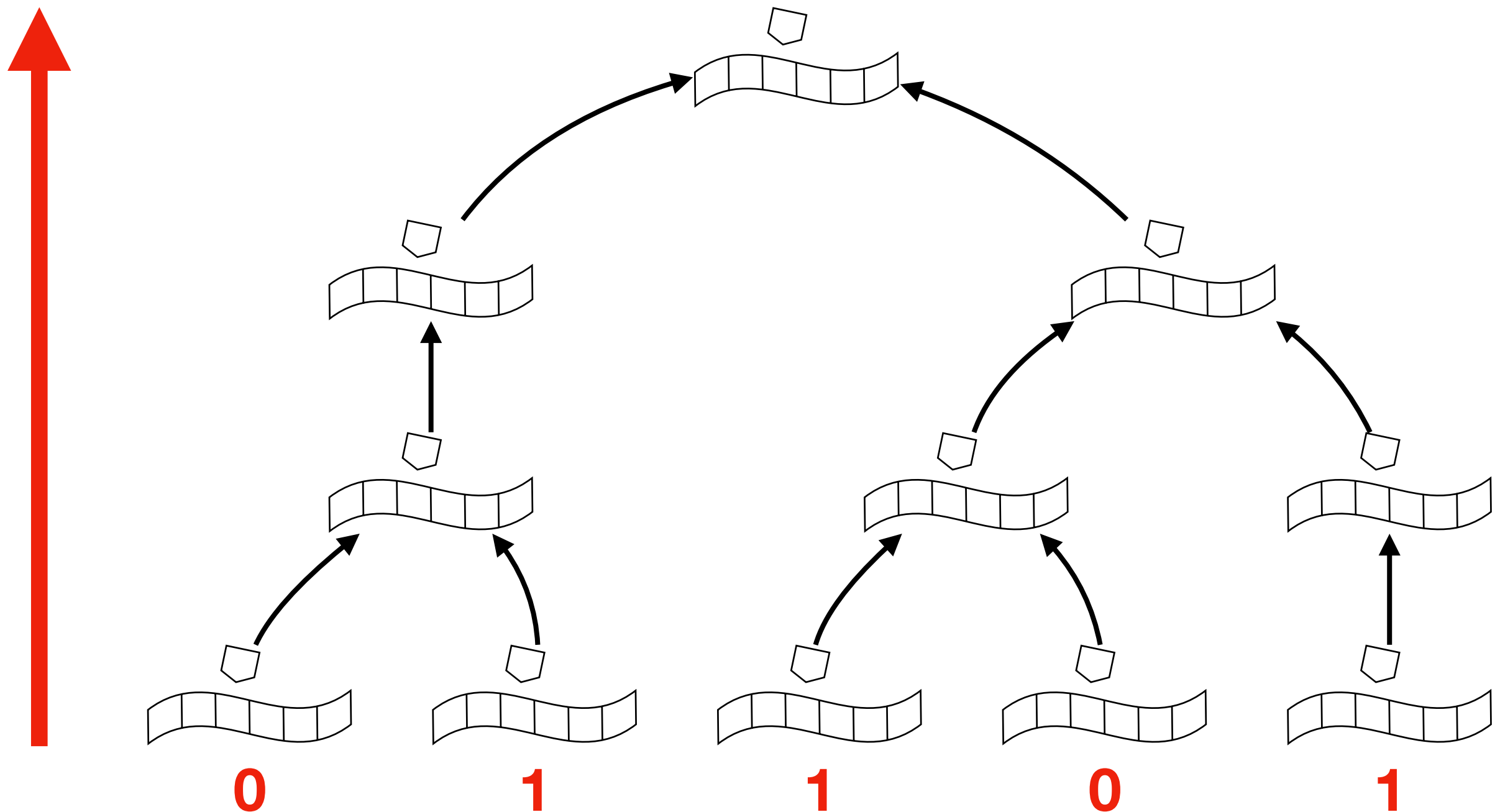
Nondeterministic Turing machines: NP



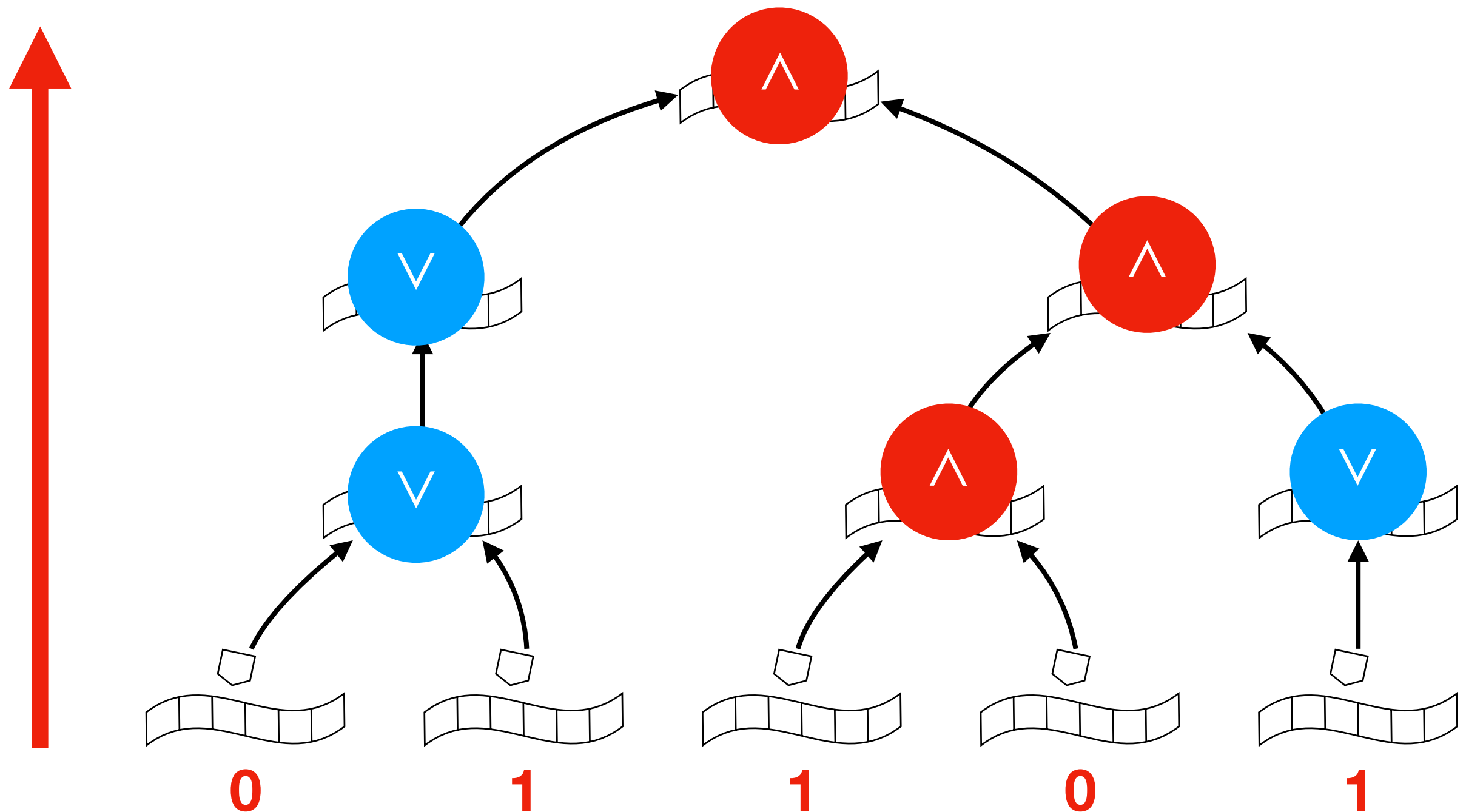
Nondeterministic Turing machines: NP



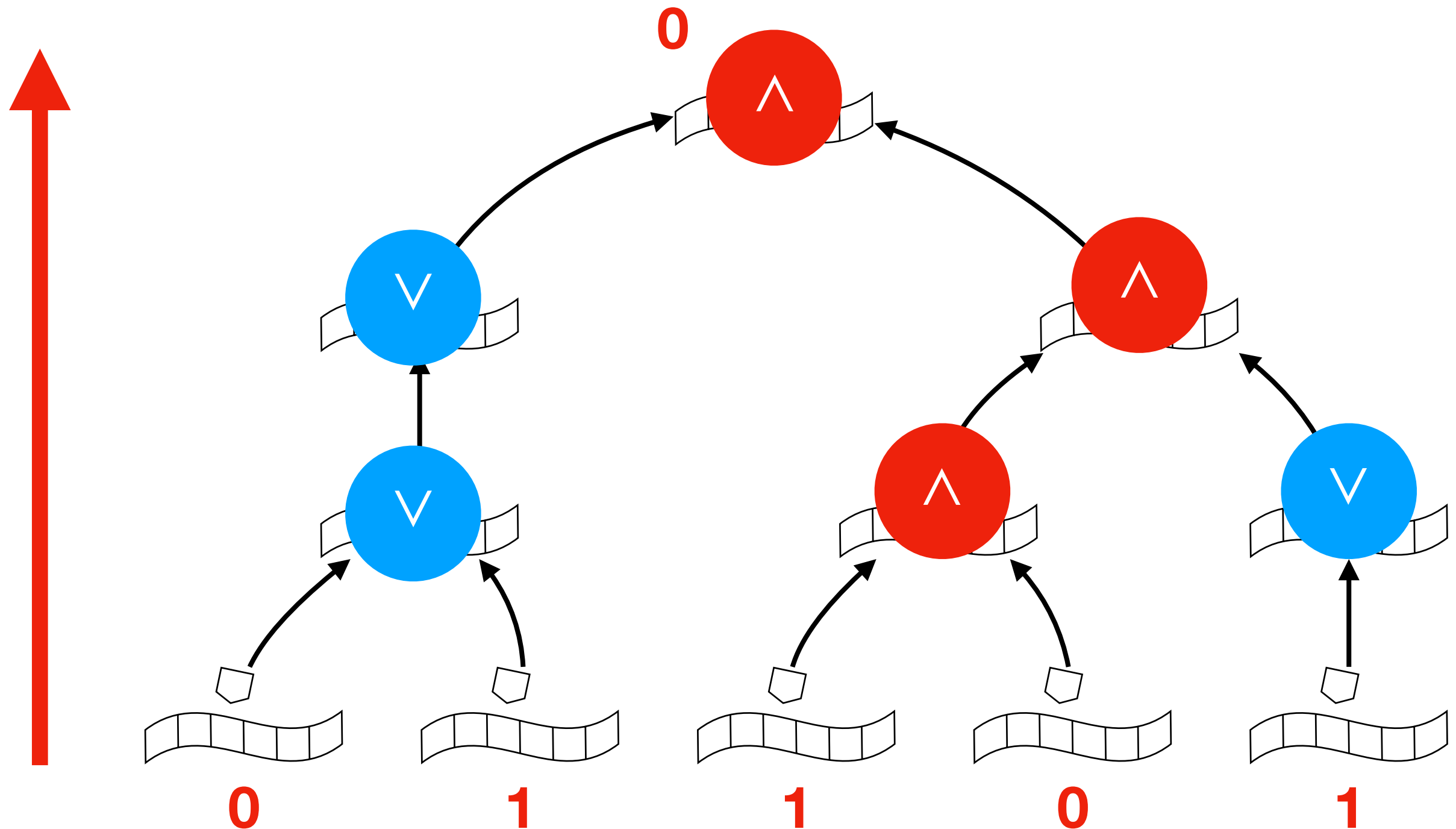
Alternating Turing machines: PSPACE



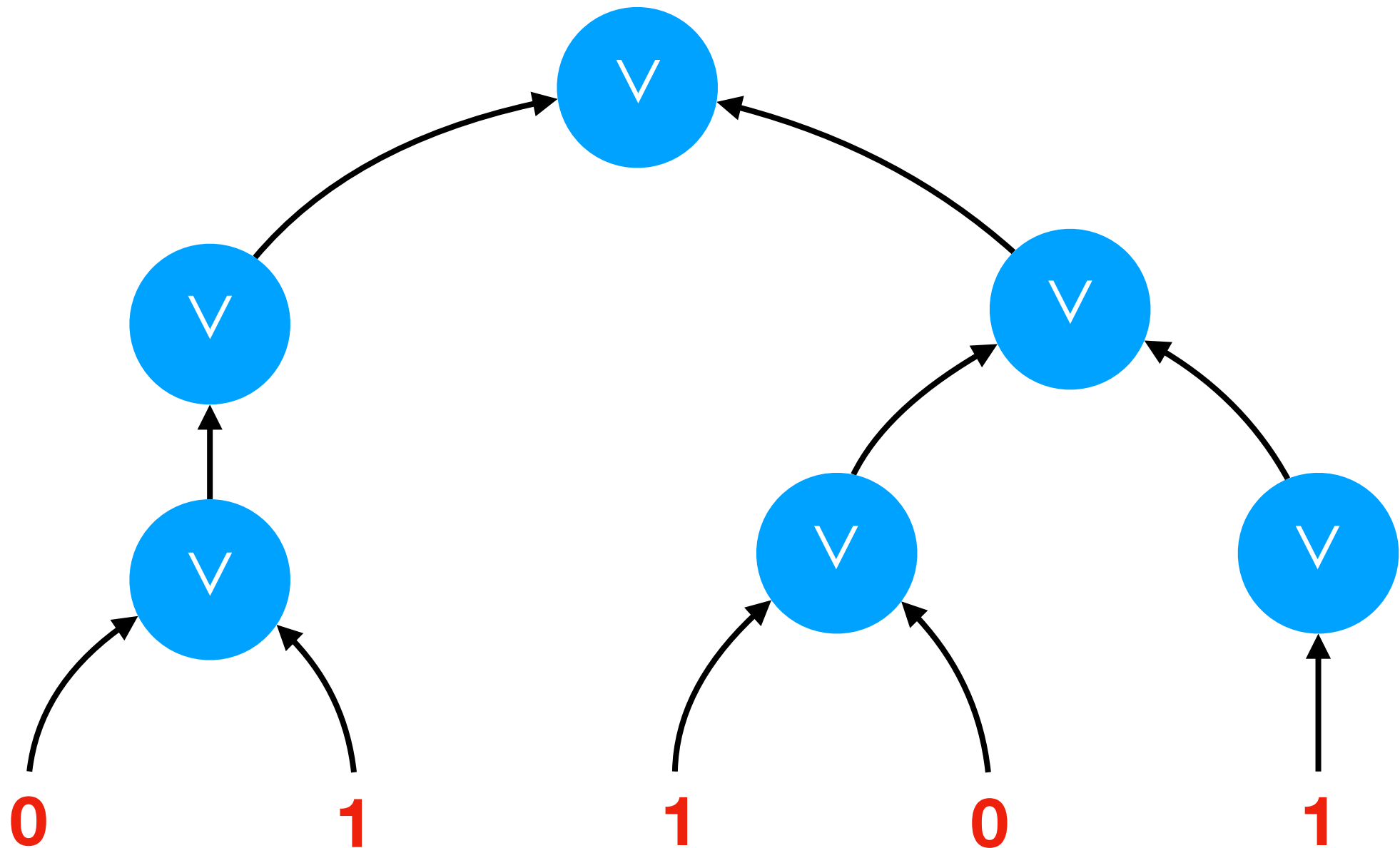
Alternating Turing machines: PSPACE



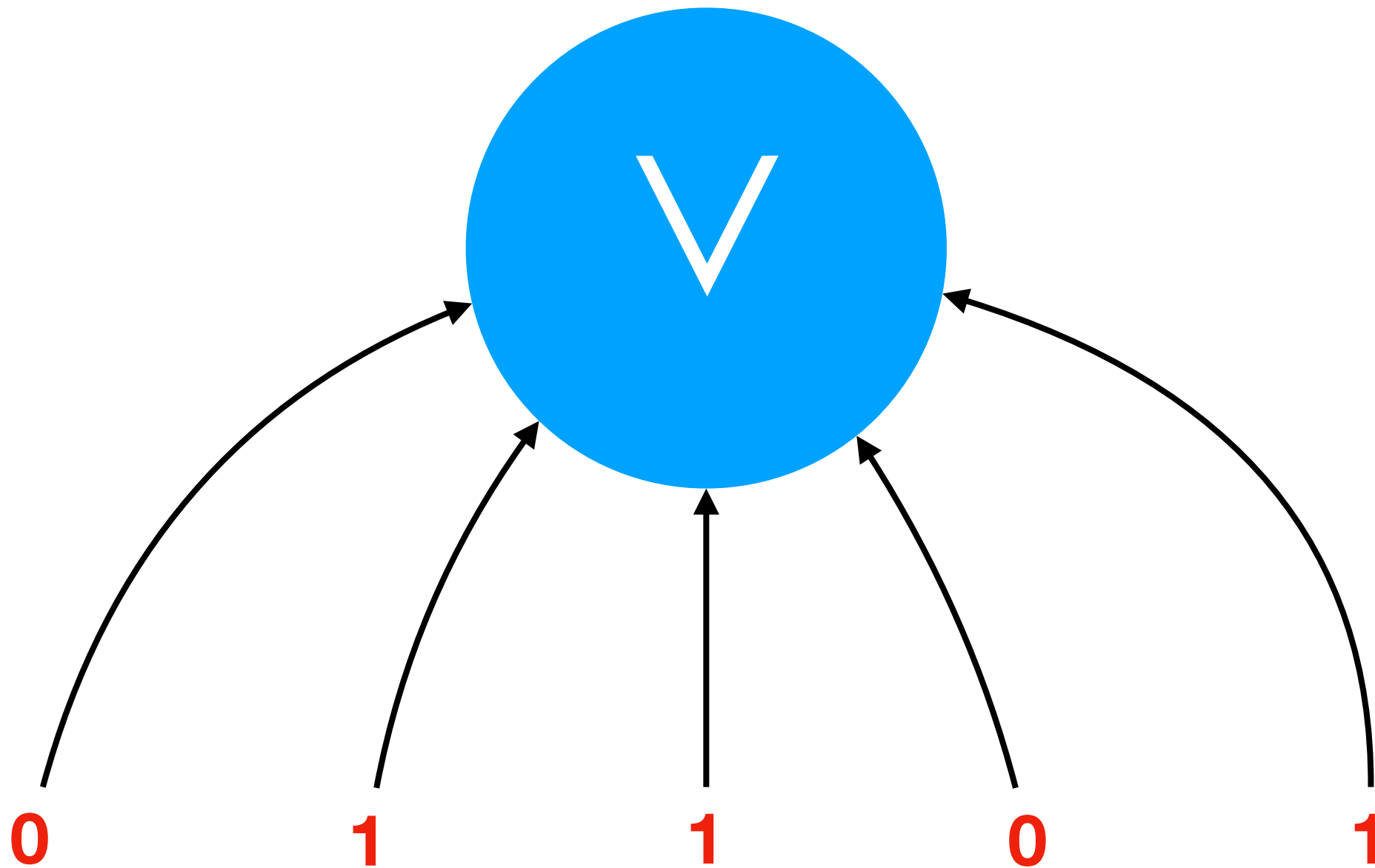
Alternating Turing machines: PSPACE



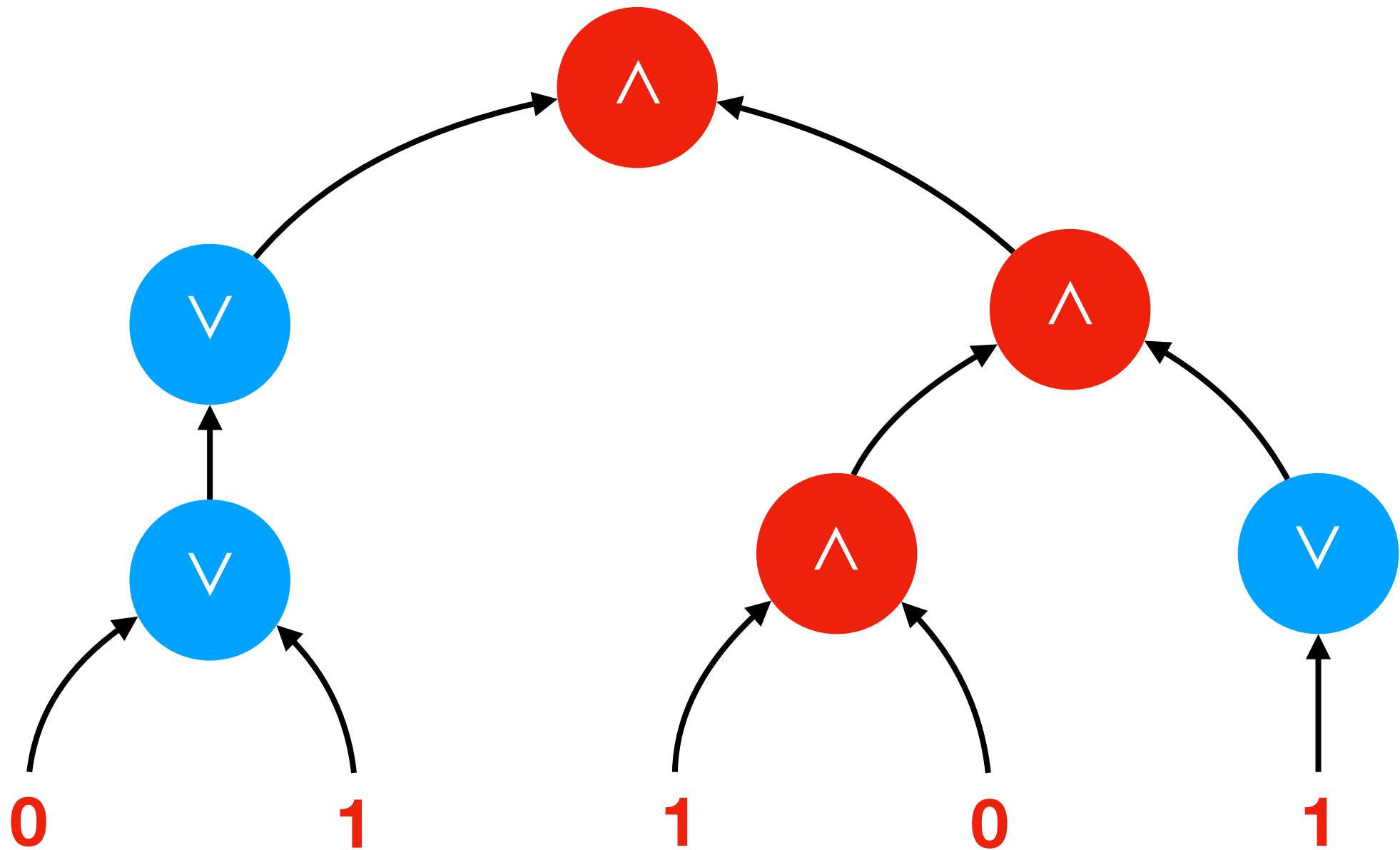
Flattening v-circuits



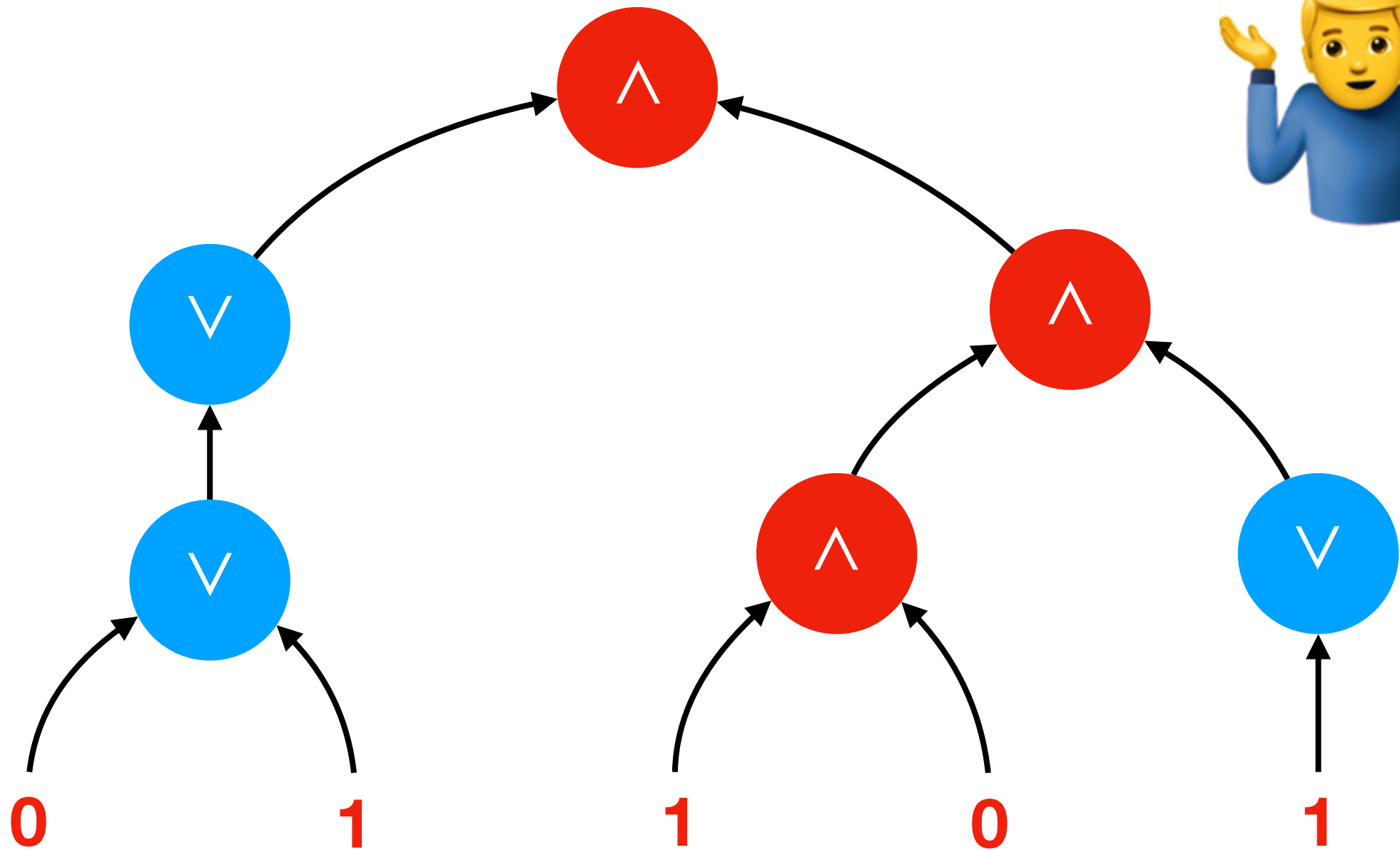
Flattening v-circuits



Flattening $\wedge \vee$ -circuits?

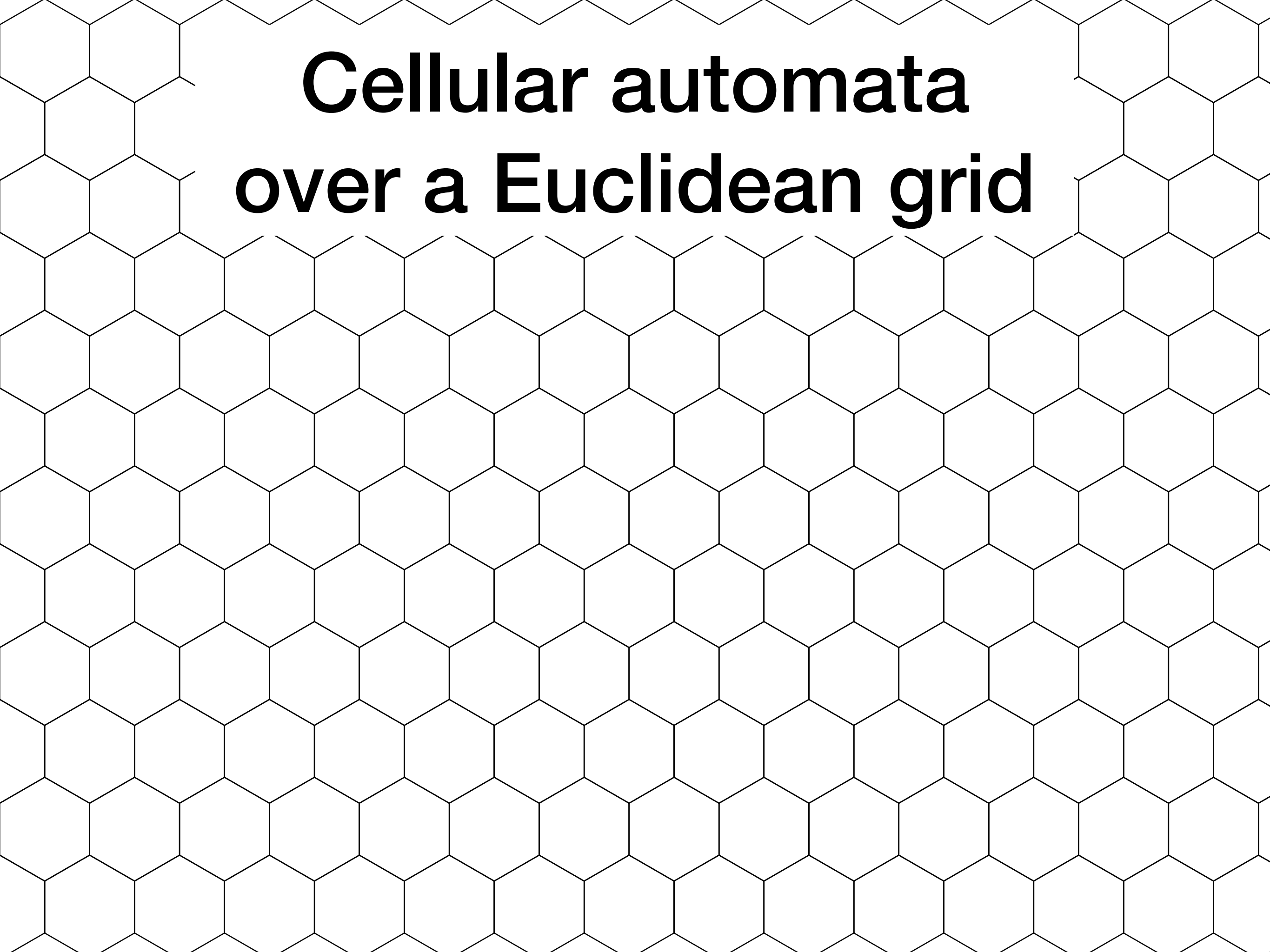


Flattening $\wedge \vee$ -circuits?

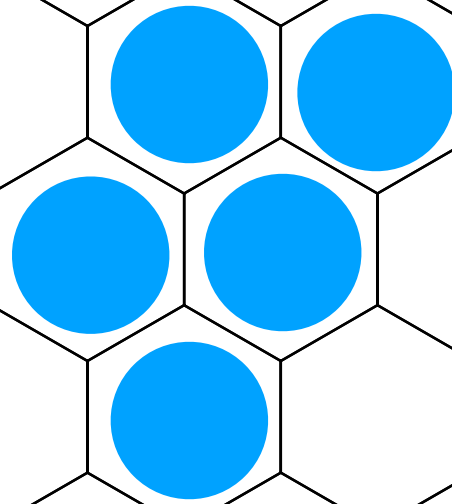


Computation space vs computation efficiency

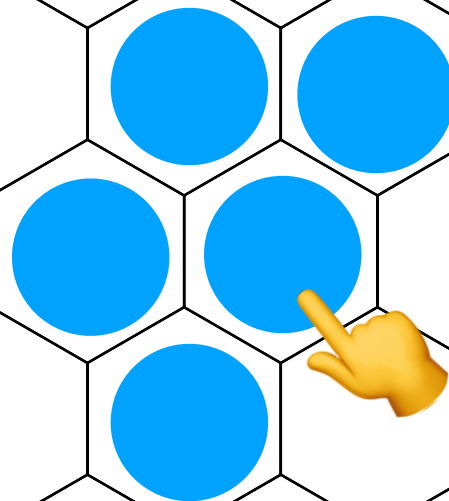
Cellular automata over a Euclidean grid



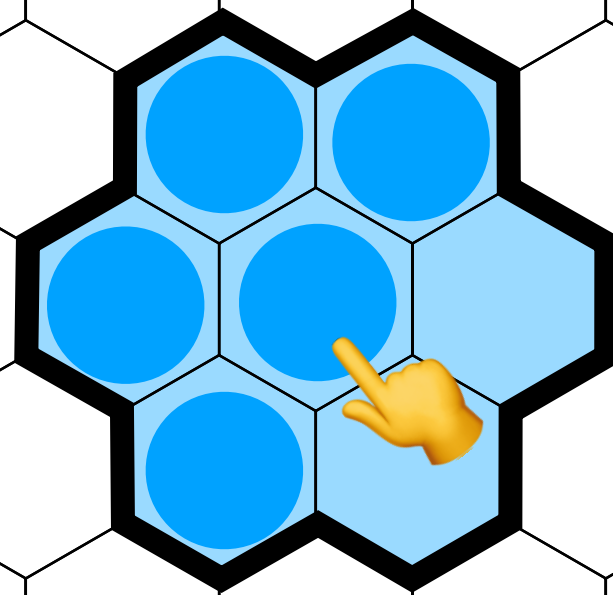
Cellular automata over a Euclidean grid



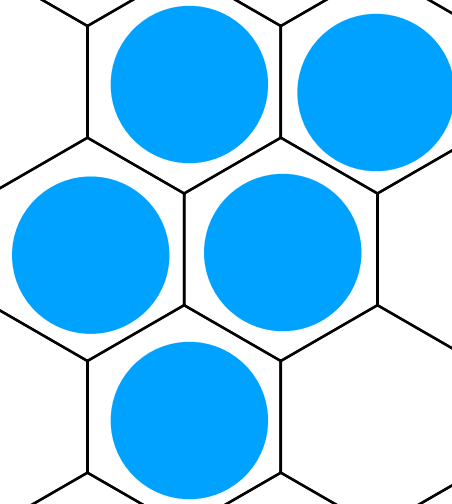
Cellular automata over a Euclidean grid



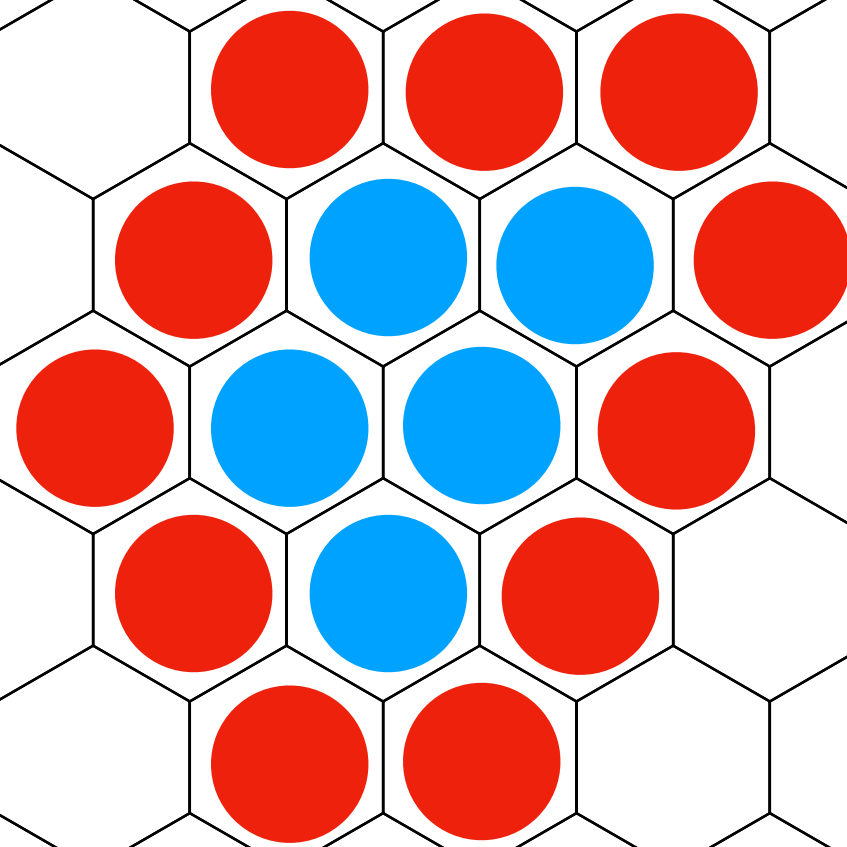
Cellular automata over a Euclidean grid



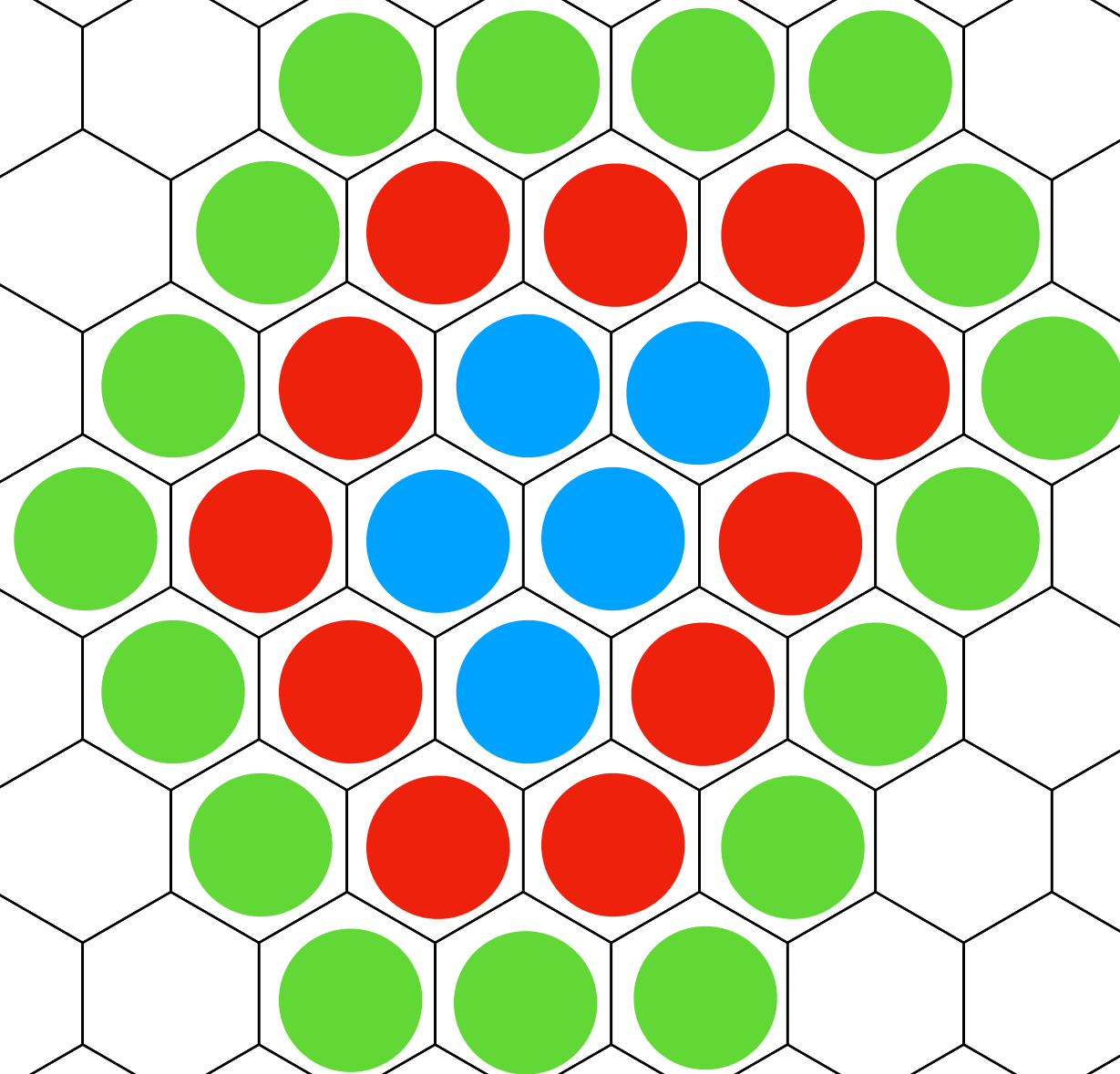
Cellular automata over a Euclidean grid



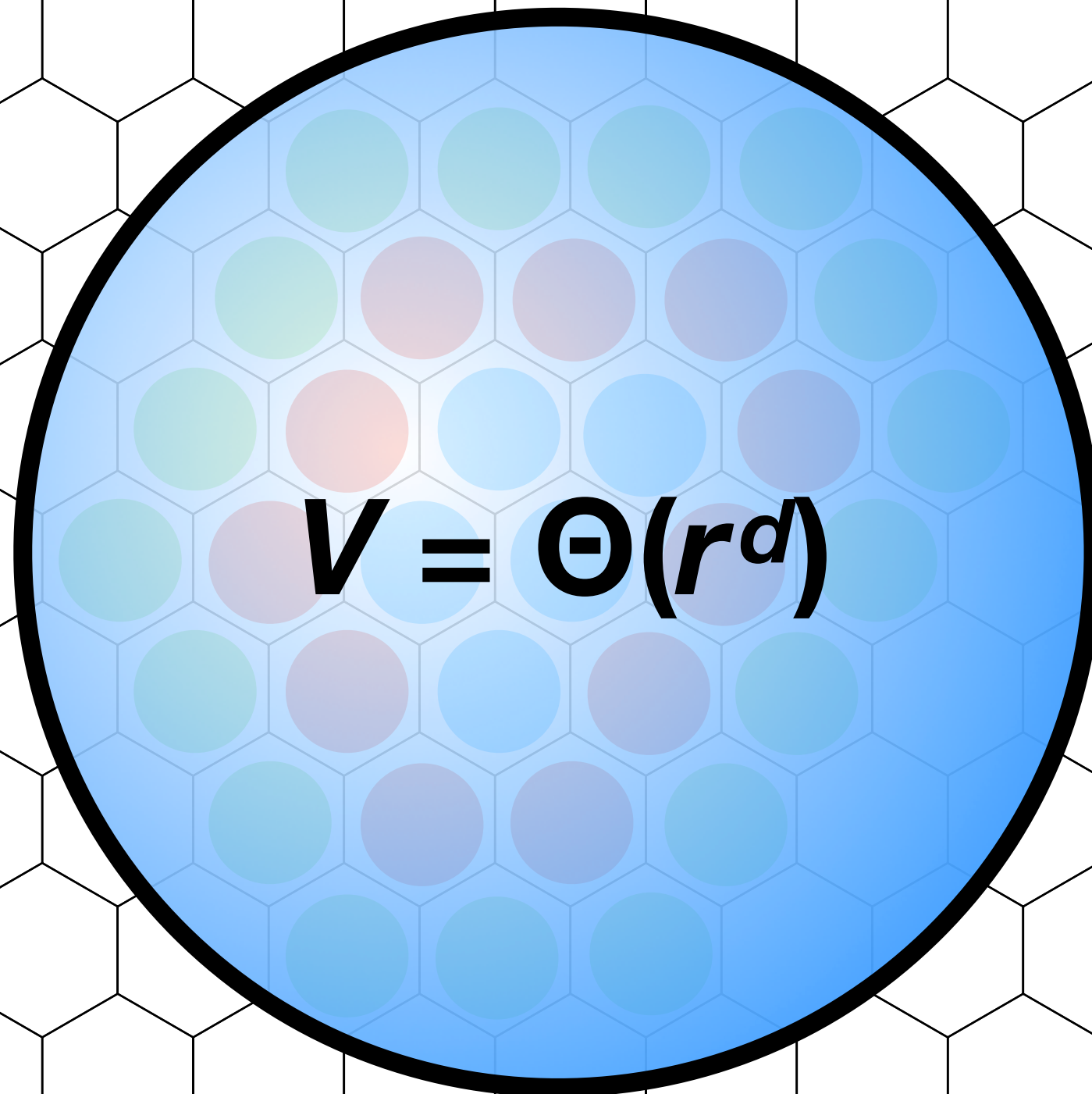
Cellular automata over a Euclidean grid



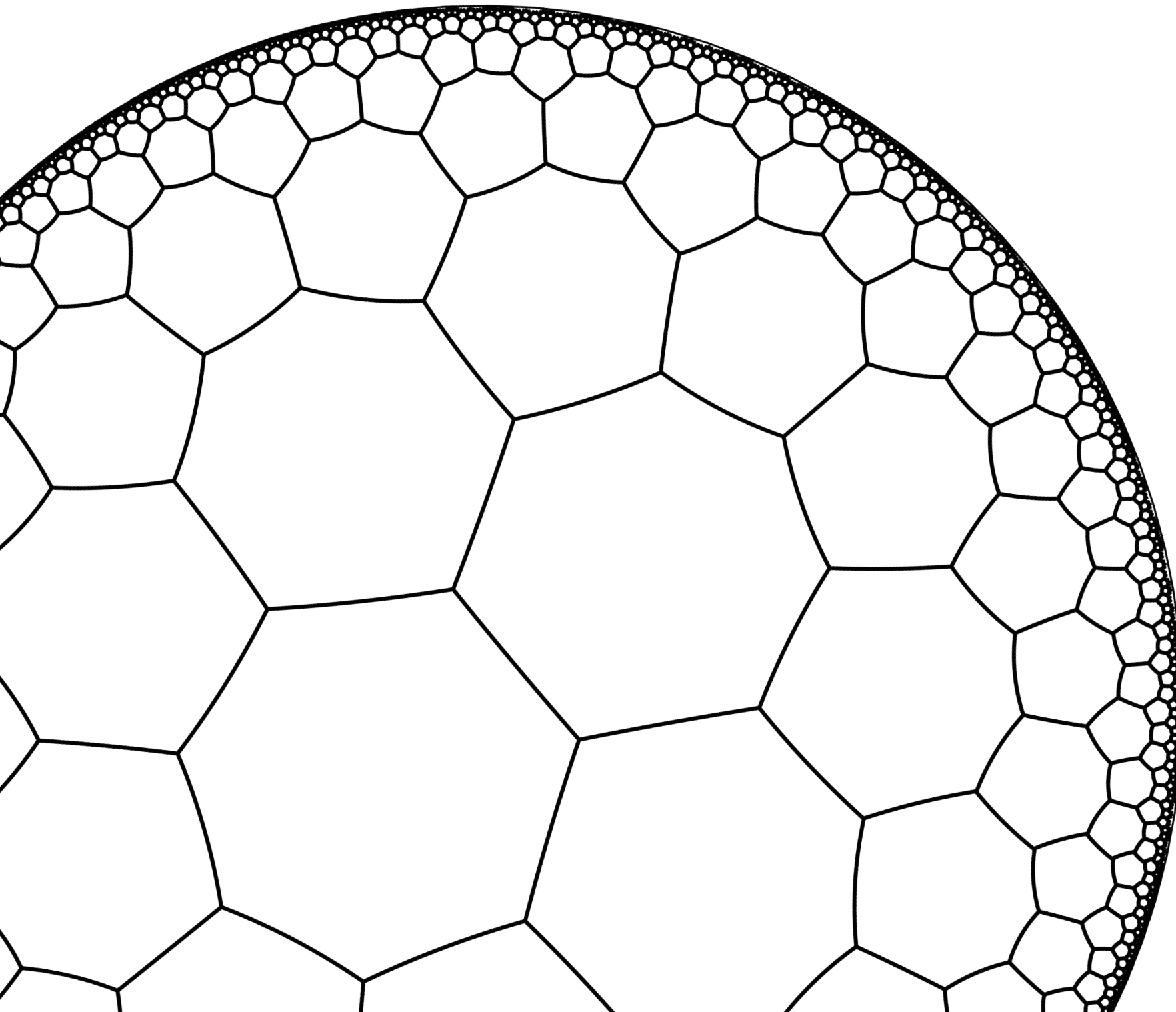
Cellular automata over a Euclidean grid



Cellular automata over a Euclidean grid

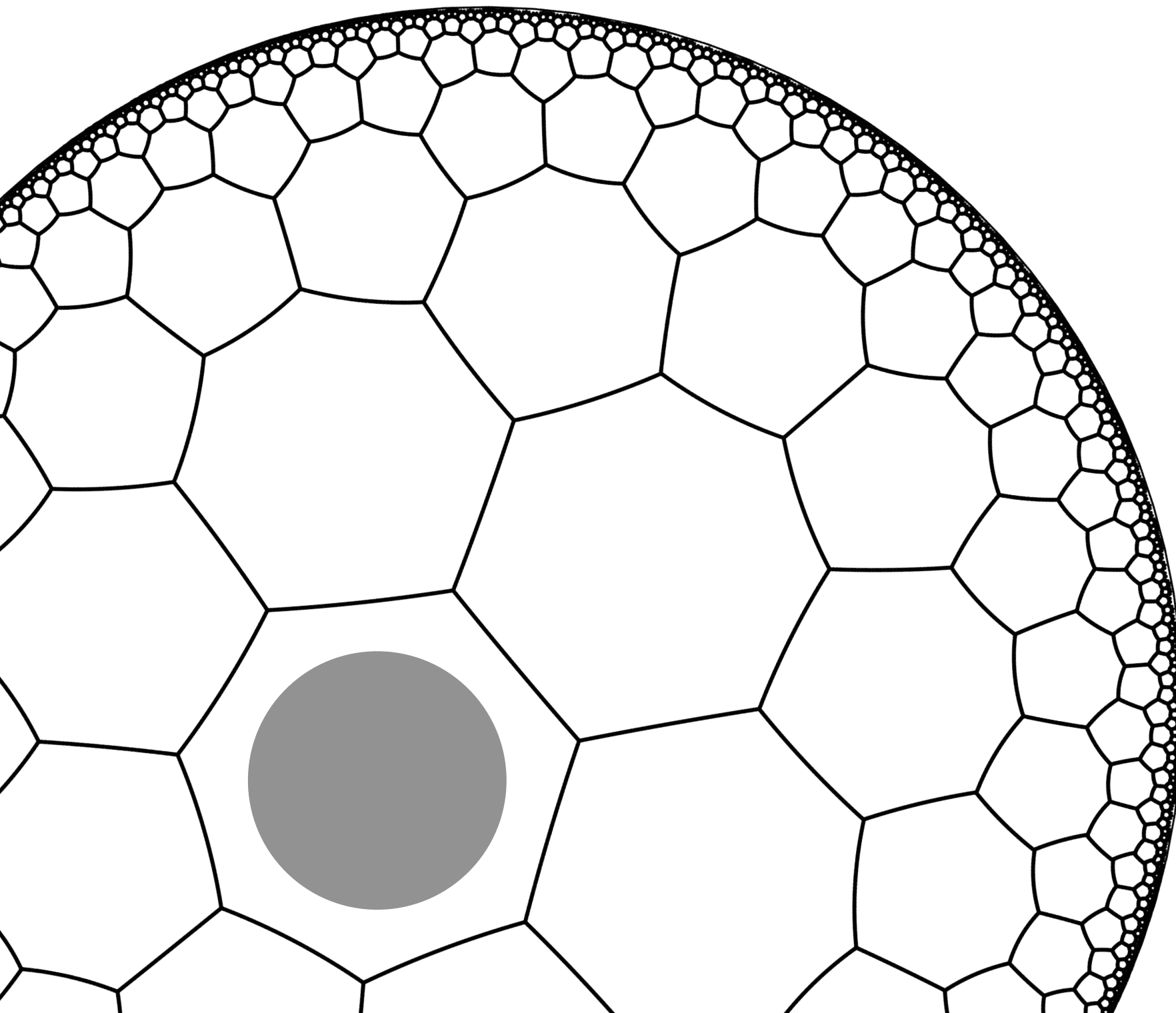


Cellular automata over a hyperbolic grid



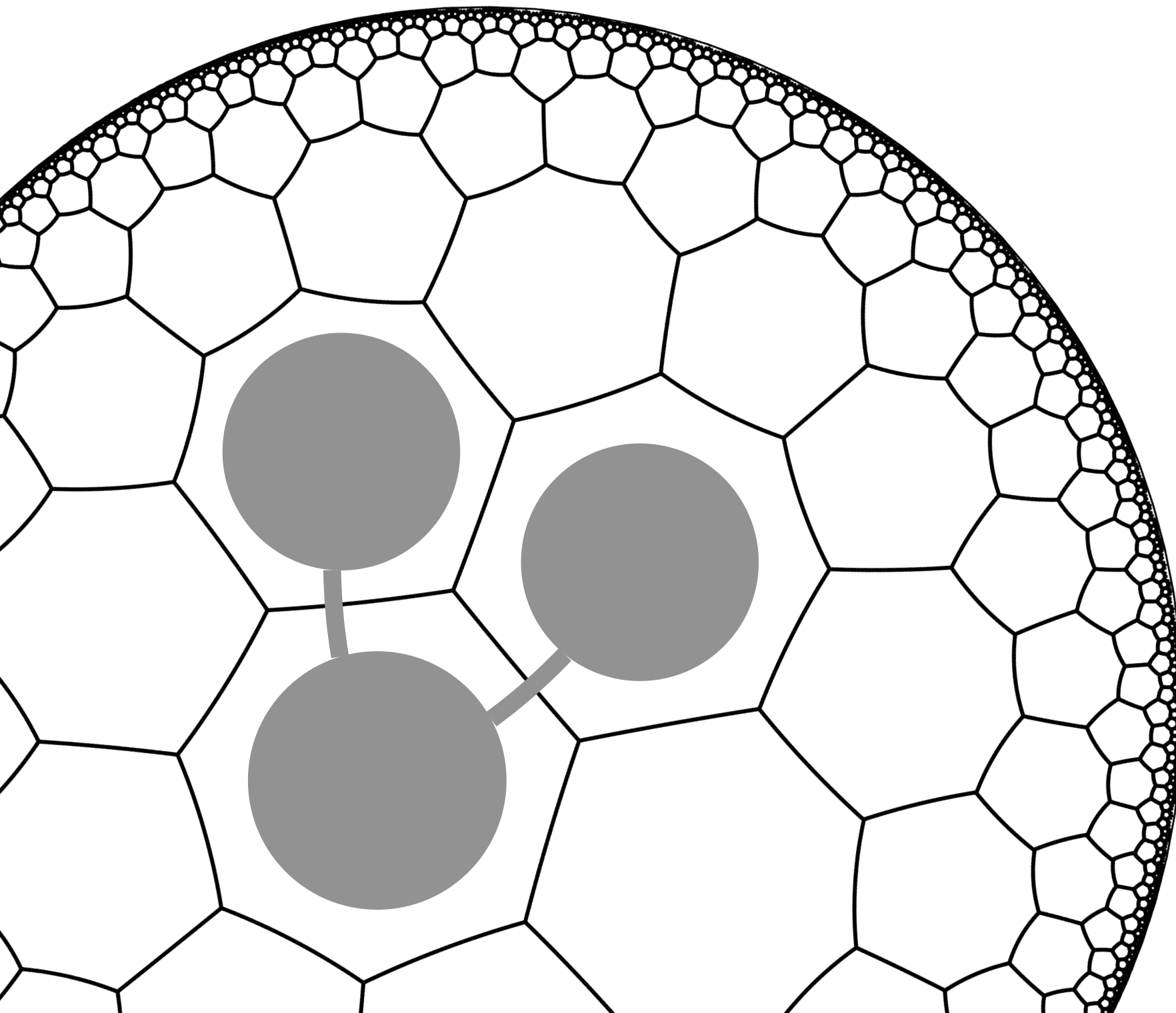
Pavage du plan hyperbolique par des heptagones, dans le modèle du disque de Poincaré. By Theon, used under CC BY-SA 3.0 <https://en.wikipedia.org/wiki/File:PavageHypPoincare2.svg>

Cellular automata over a hyperbolic grid



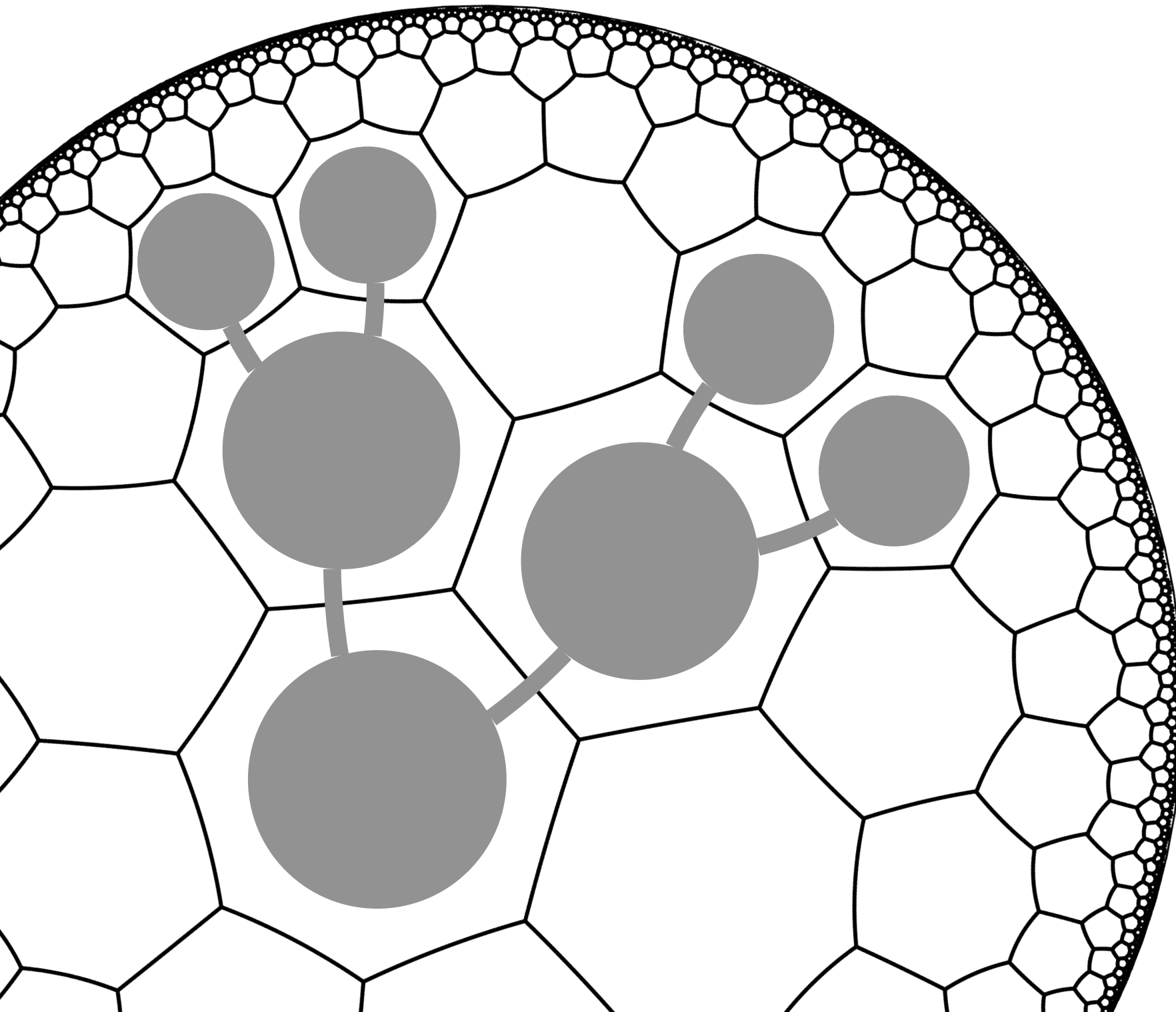
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Cellular automata over a hyperbolic grid



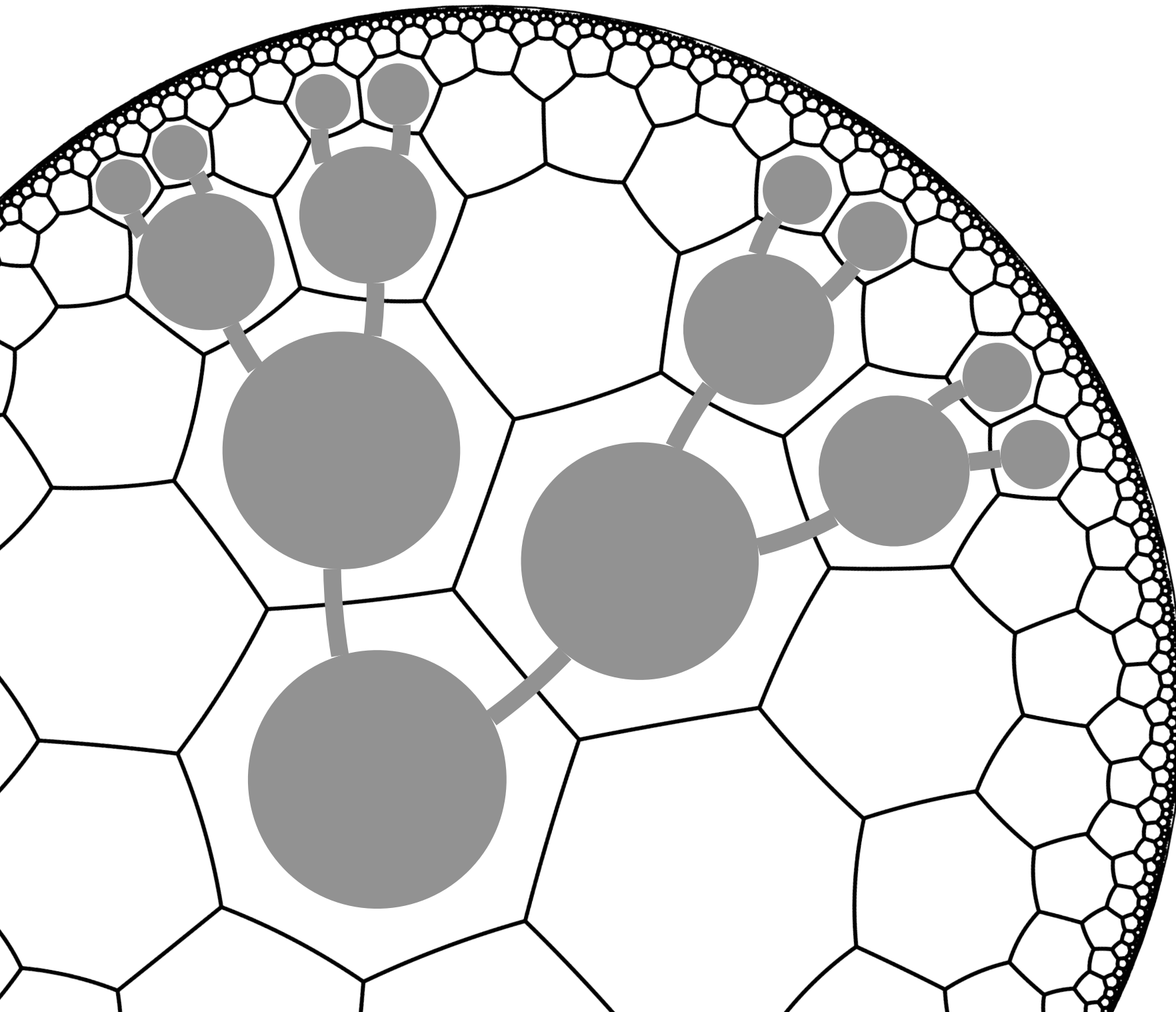
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Cellular automata over a hyperbolic grid



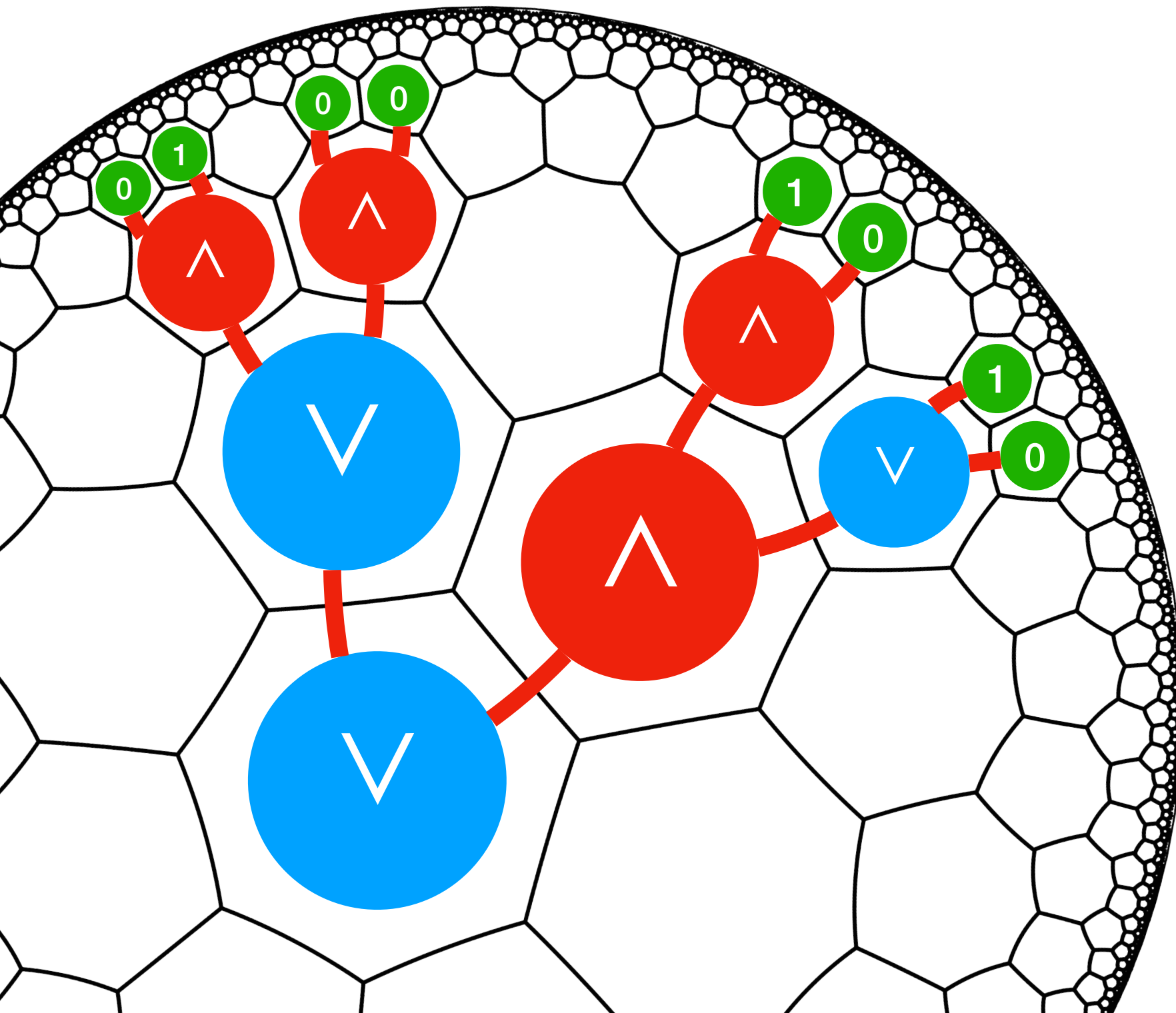
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Cellular automata over a hyperbolic grid



Pavage du plan hyperbolique par des heptagones, dans le modèle du disque de Poincaré. By Theon, used under CC BY-SA 3.0 <https://en.wikipedia.org/wiki/File:PavageHypPoincare2.svg>

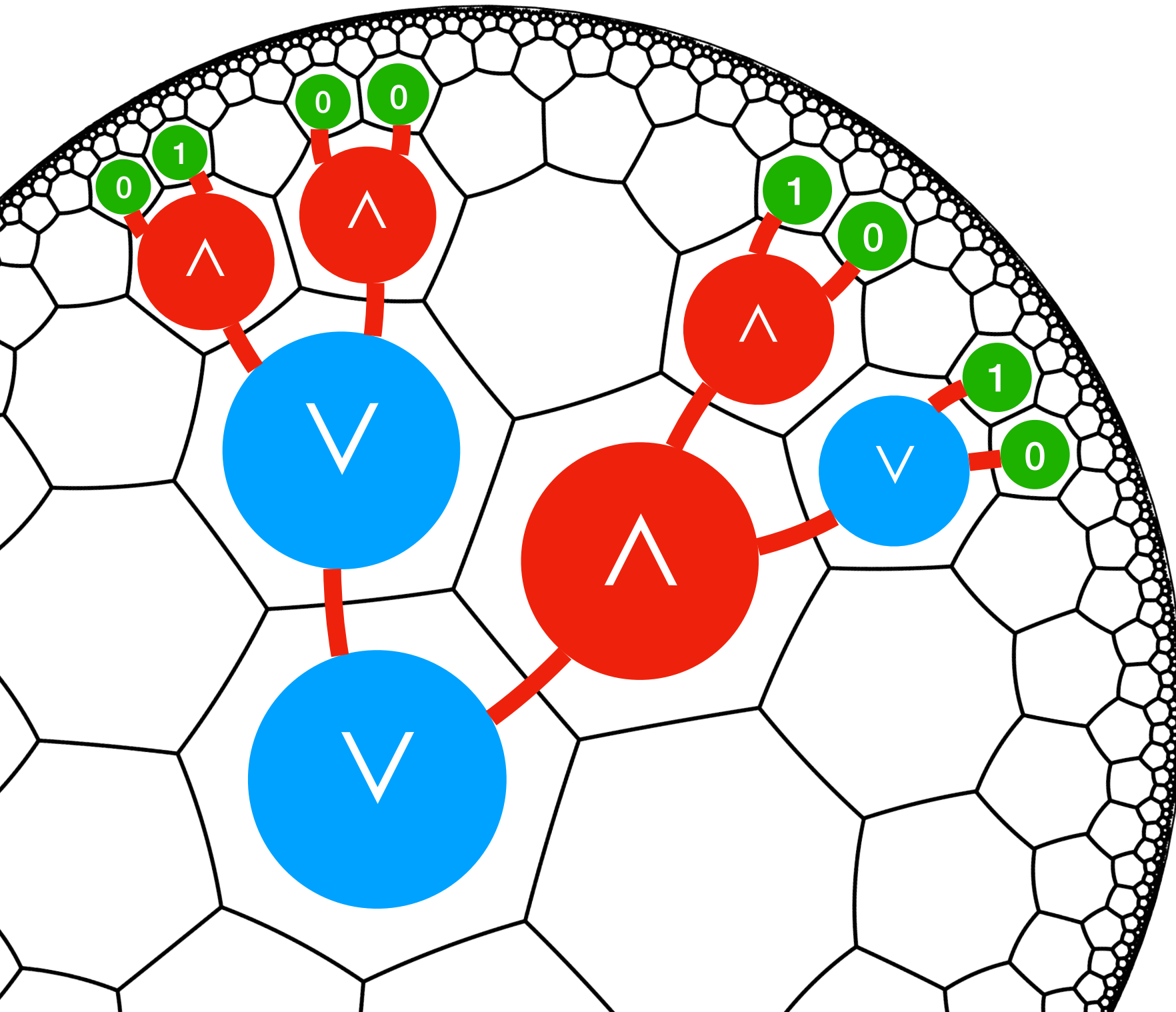
Cellular automata over a hyperbolic grid



Pavage du plan hyperbolique par des heptagones, dans le modèle du disque de Poincaré. By Theon, used under CC BY-SA 3.0 <https://en.wikipedia.org/wiki/File:PavageHypPoincare2.svg>

Cellular automata over a hyperbolic grid

$$V = \Omega(2^r)$$



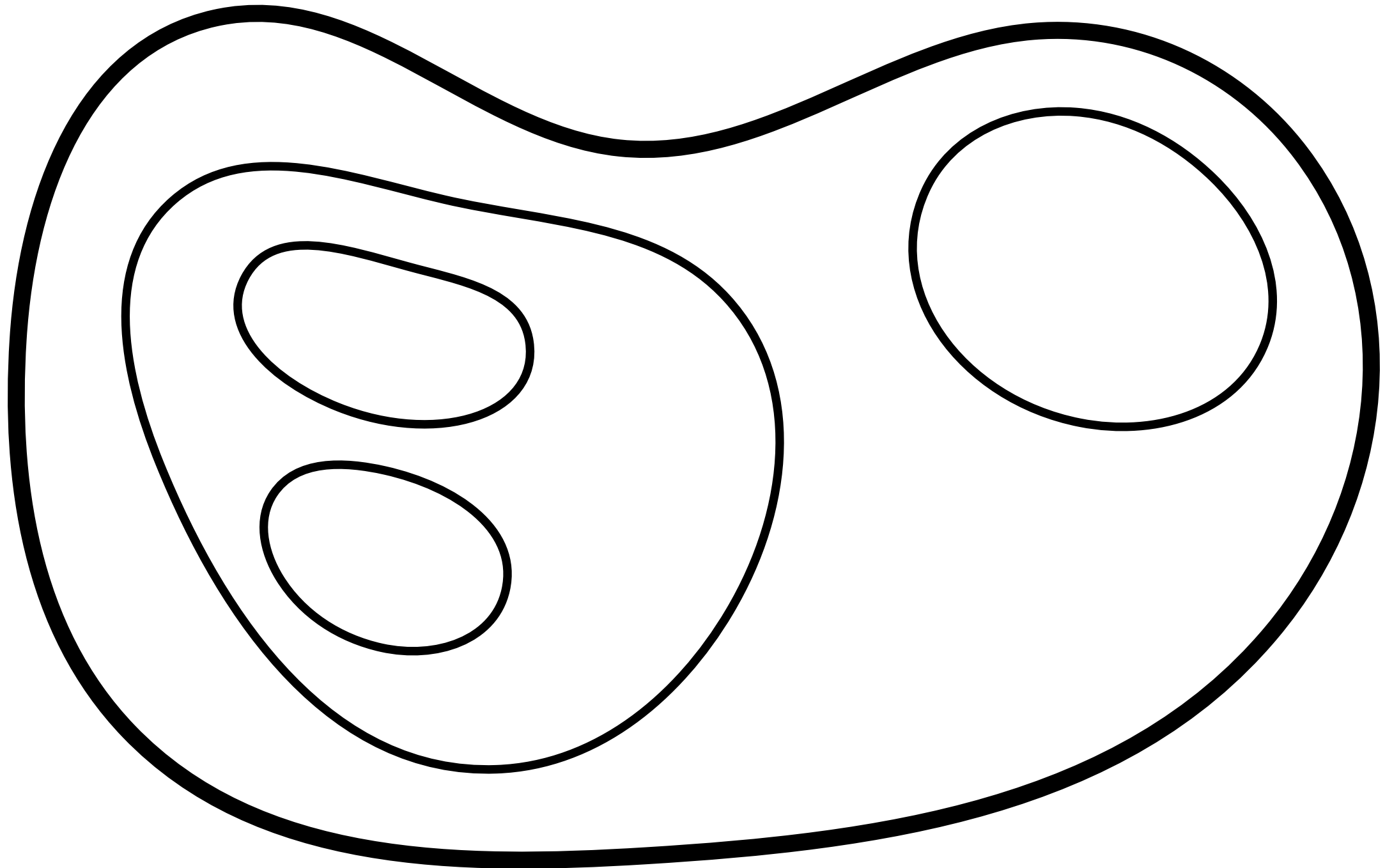
Pavage du plan hyperbolique par des heptagones, dans le modèle du disque de Poincaré. By Theon, used under CC BY-SA 3.0 <https://en.wikipedia.org/wiki/File:PavageHypPoincare2.svg>

Rule of thumb 👍

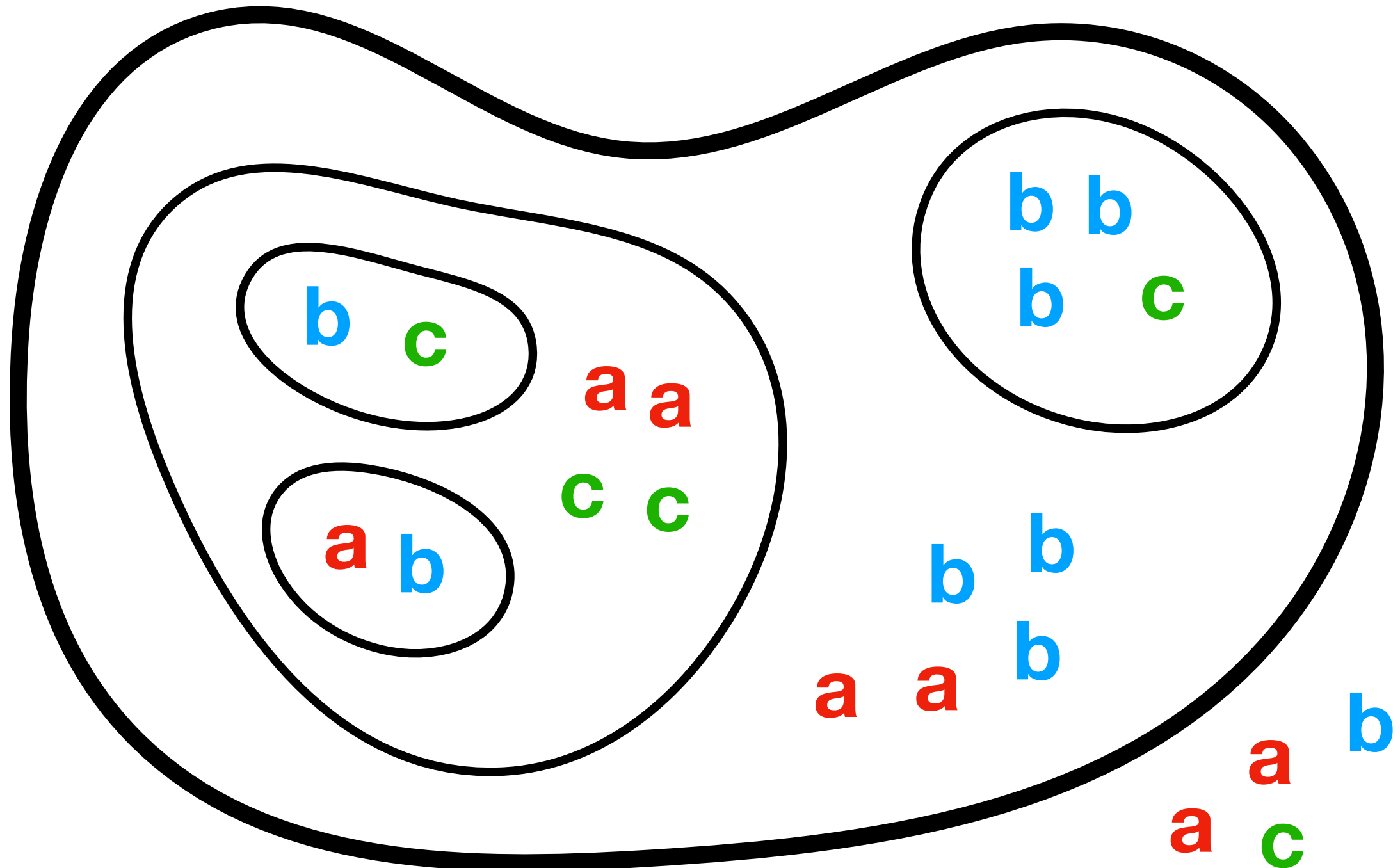
- Sequential machines are first class
- (Constant or polynomial) bounded parallel machines are also first class
- Unbounded (or exponential bounded) parallel machines are second class
- Apparently, this holds even for unconventional computing models 😞

**A “more unconventional”
model of computation:
membrane systems**

Membrane systems

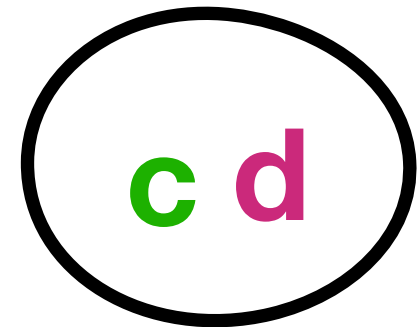
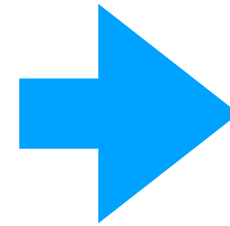
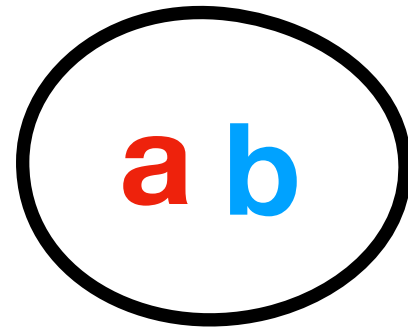


Membrane systems

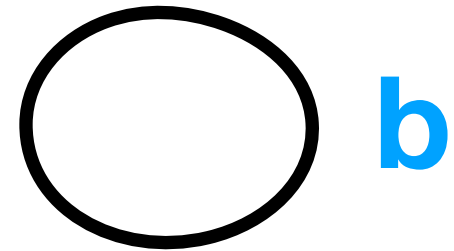
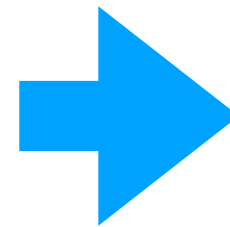
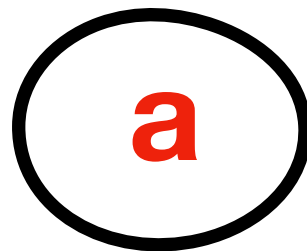


Evolution rules

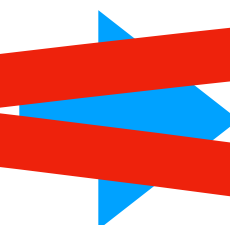
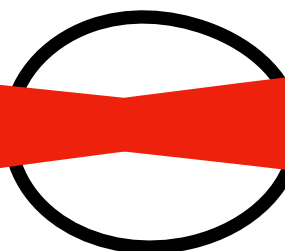
$[ab \rightarrow cd]$



$[a \rightarrow [] b]$

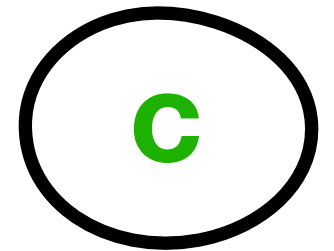
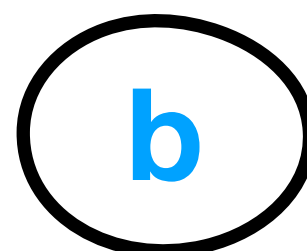
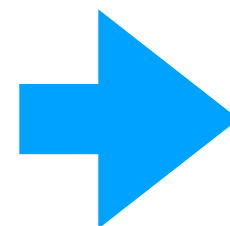
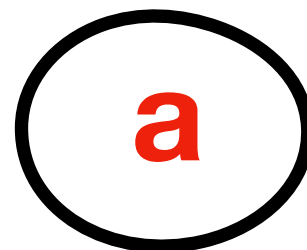


~~$a [] \rightarrow [b]$~~



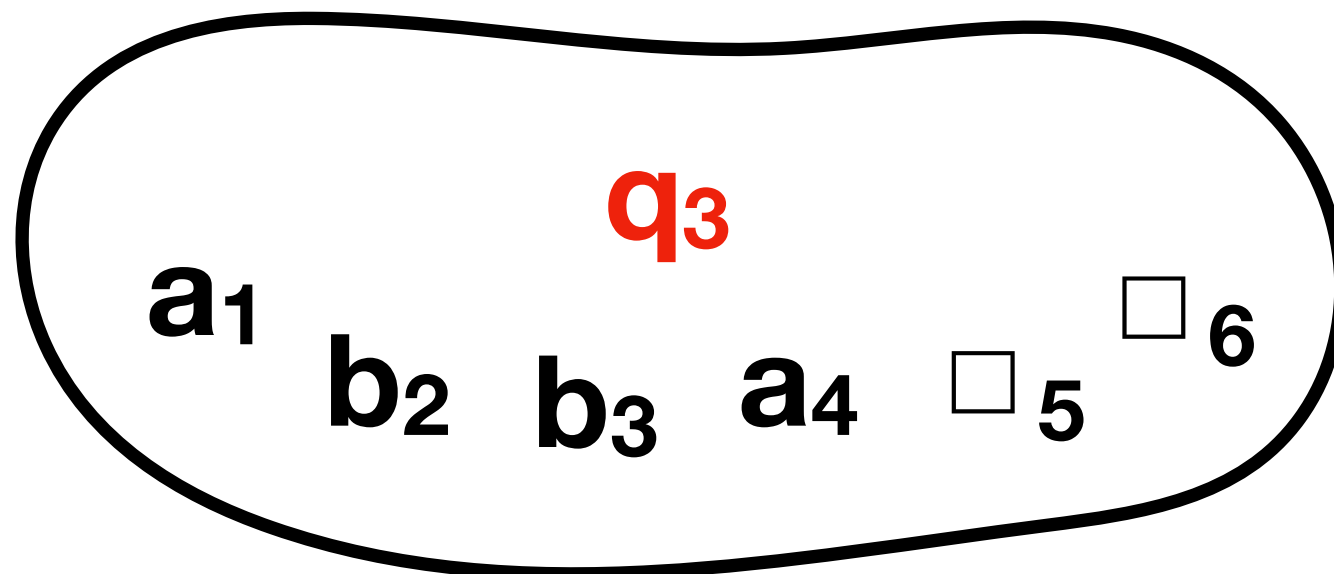
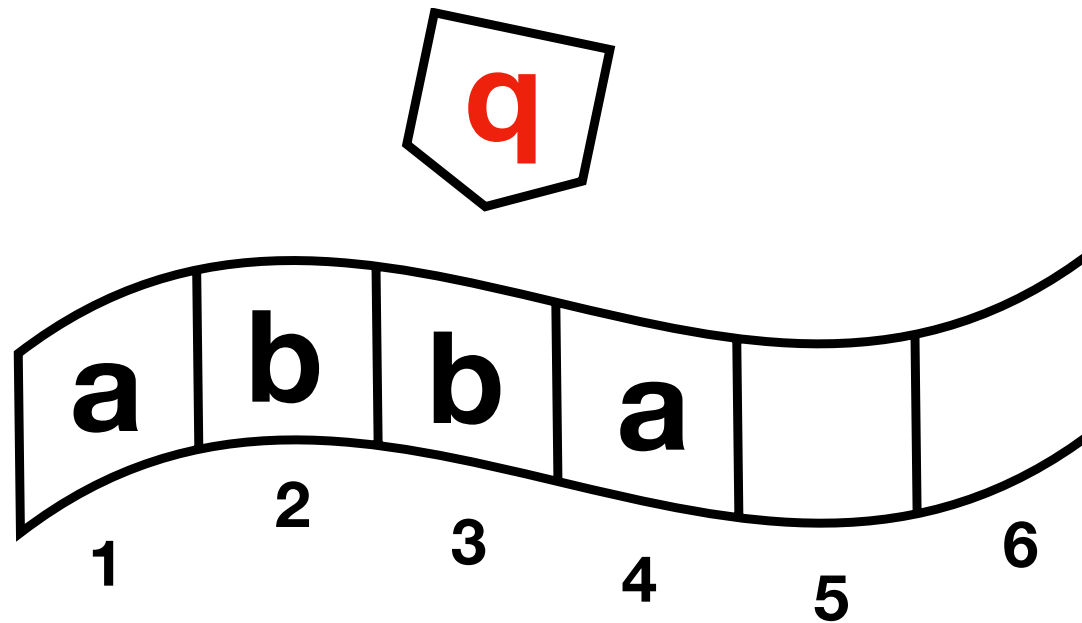
“monodirectional”

$[a \rightarrow [b] [c]]$

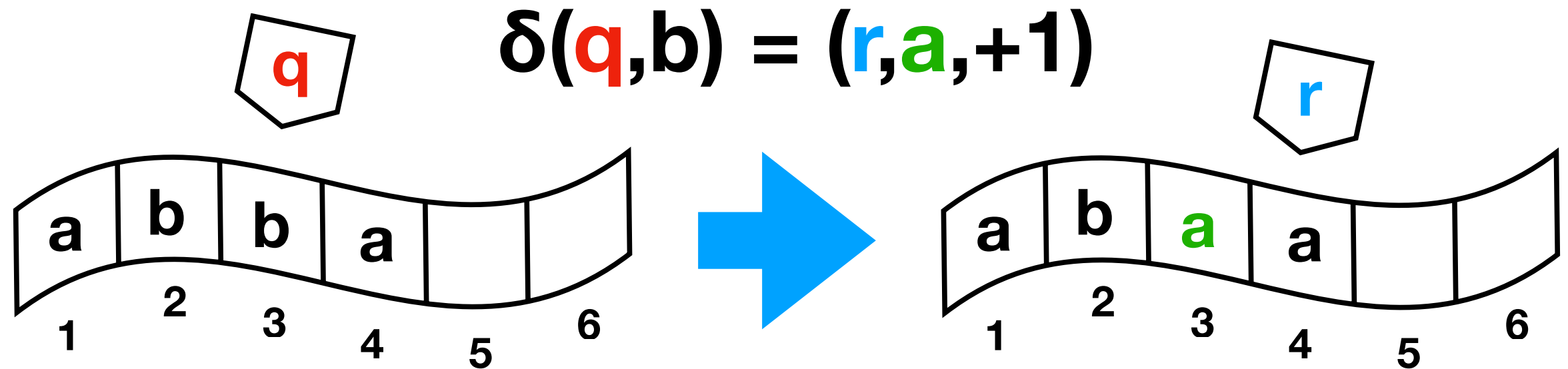


Simulating Turing machines with membrane systems

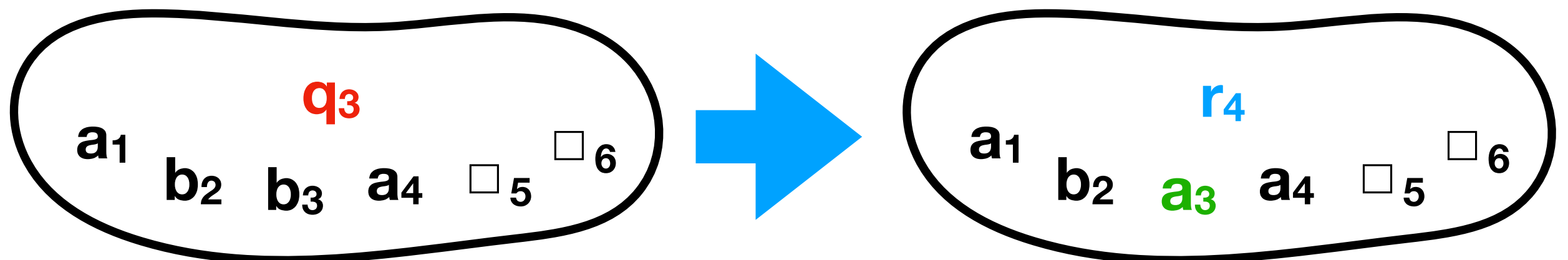
Encoding the configuration



Simulating transitions

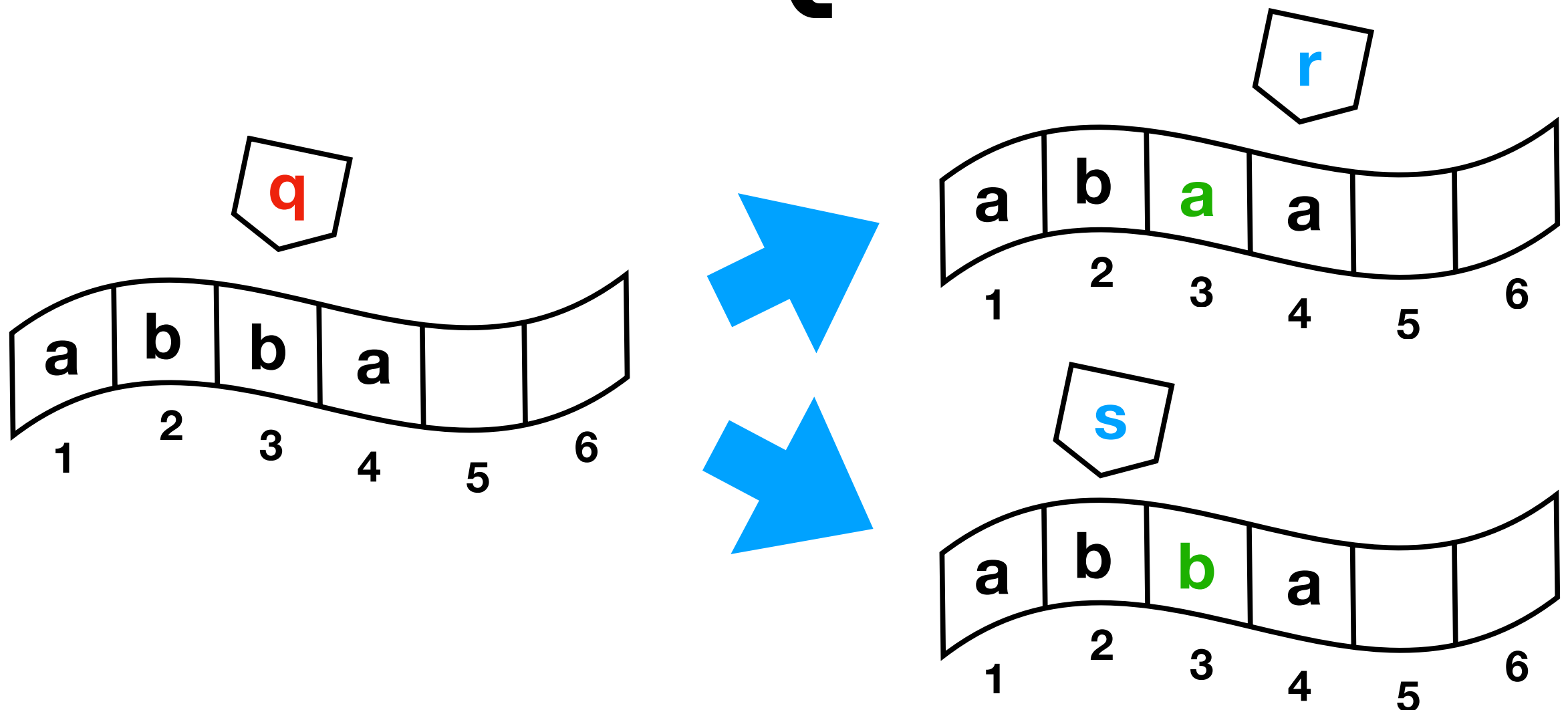


$$[\mathbf{q}_3 \mathbf{b}_3 \rightarrow \mathbf{r}_4 \mathbf{a}_3]$$



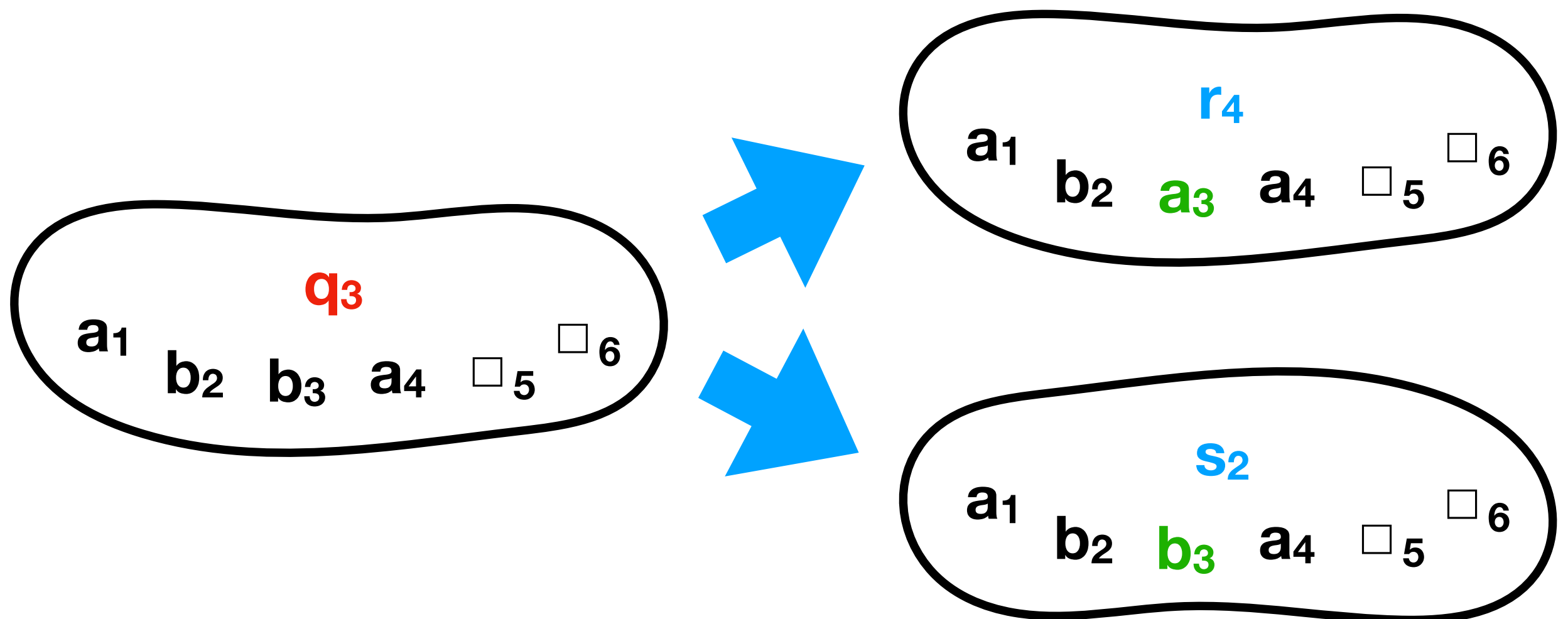
Simulating nondeterminism

$$\delta(\mathbf{q},b) = \begin{cases} (\mathbf{r},\mathbf{a},+1) \\ (\mathbf{s},\mathbf{b},-1) \end{cases}$$

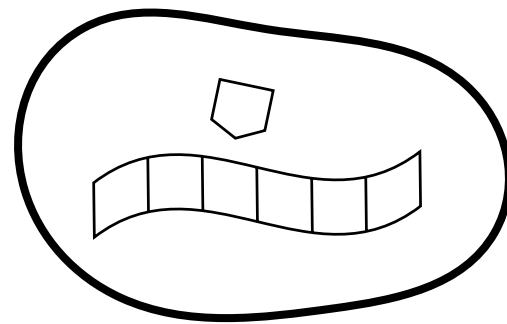


Simulating nondeterminism

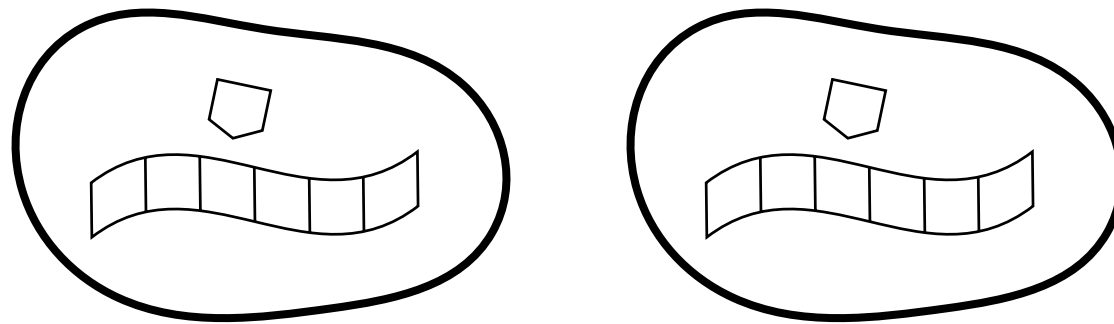
$$\delta(\mathbf{q}, \mathbf{b}) = \begin{cases} (\mathbf{r}, \mathbf{a}, +1) \\ (\mathbf{s}, \mathbf{b}, -1) \end{cases} \quad [\mathbf{q}_3 \mathbf{b}_3] \rightarrow [\mathbf{r}_4 \mathbf{a}_3] [\mathbf{s}_2 \mathbf{b}_3]$$



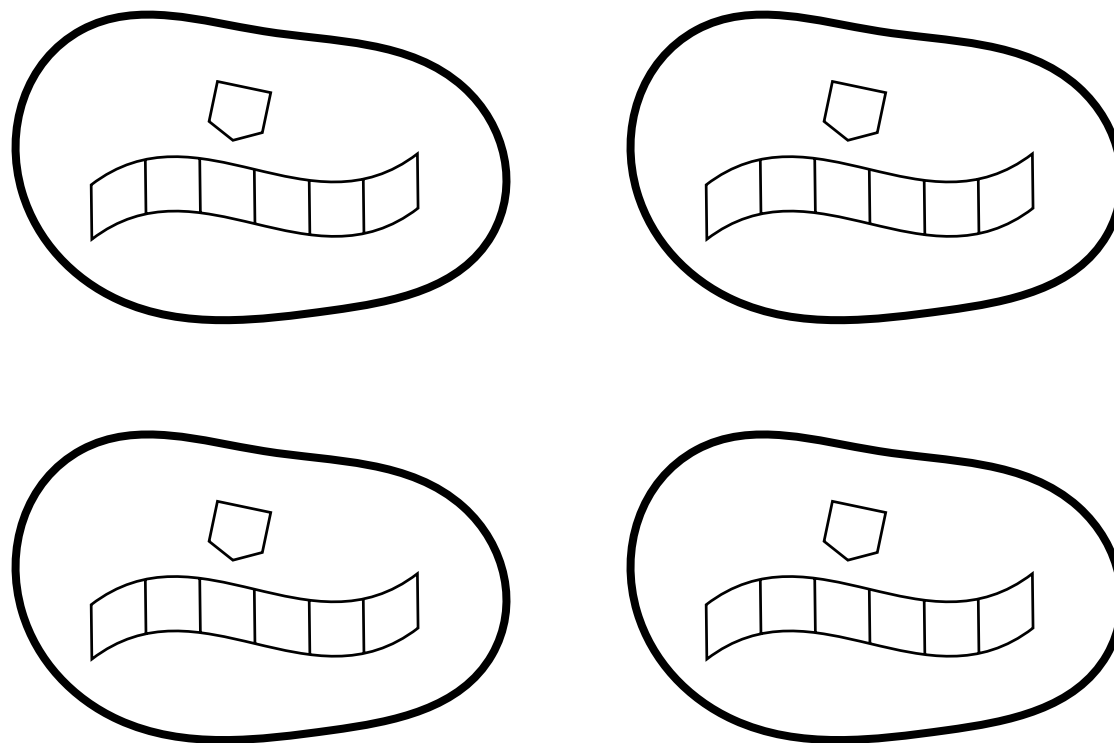
Simulating NDTM \Rightarrow NP



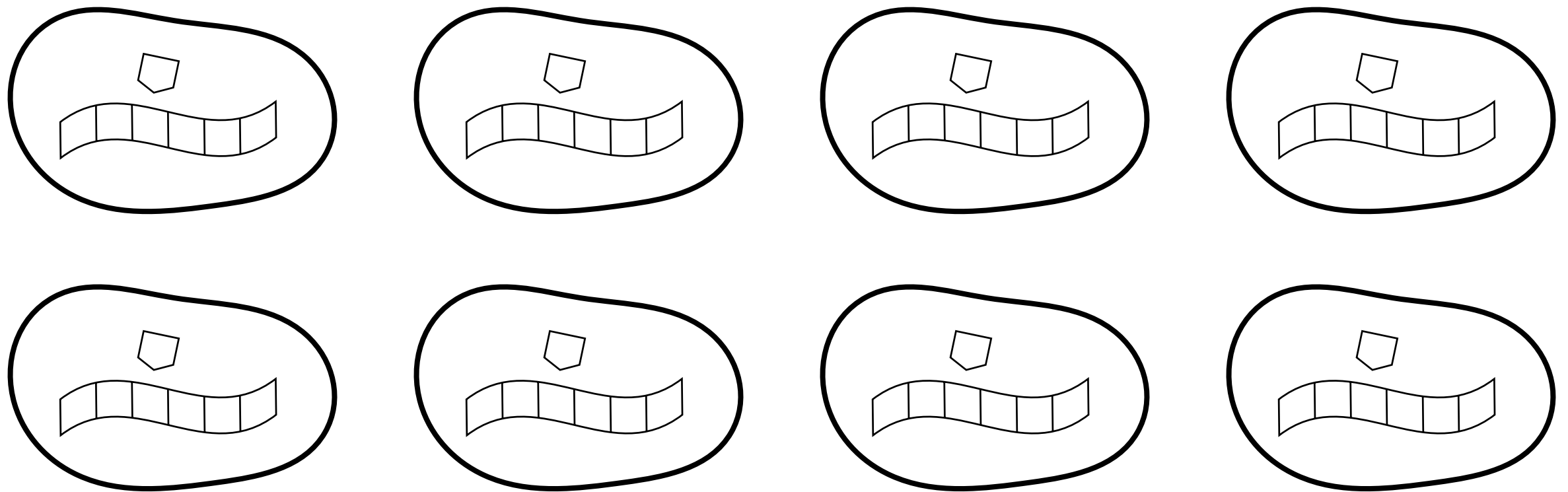
Simulating NDTM \Rightarrow NP



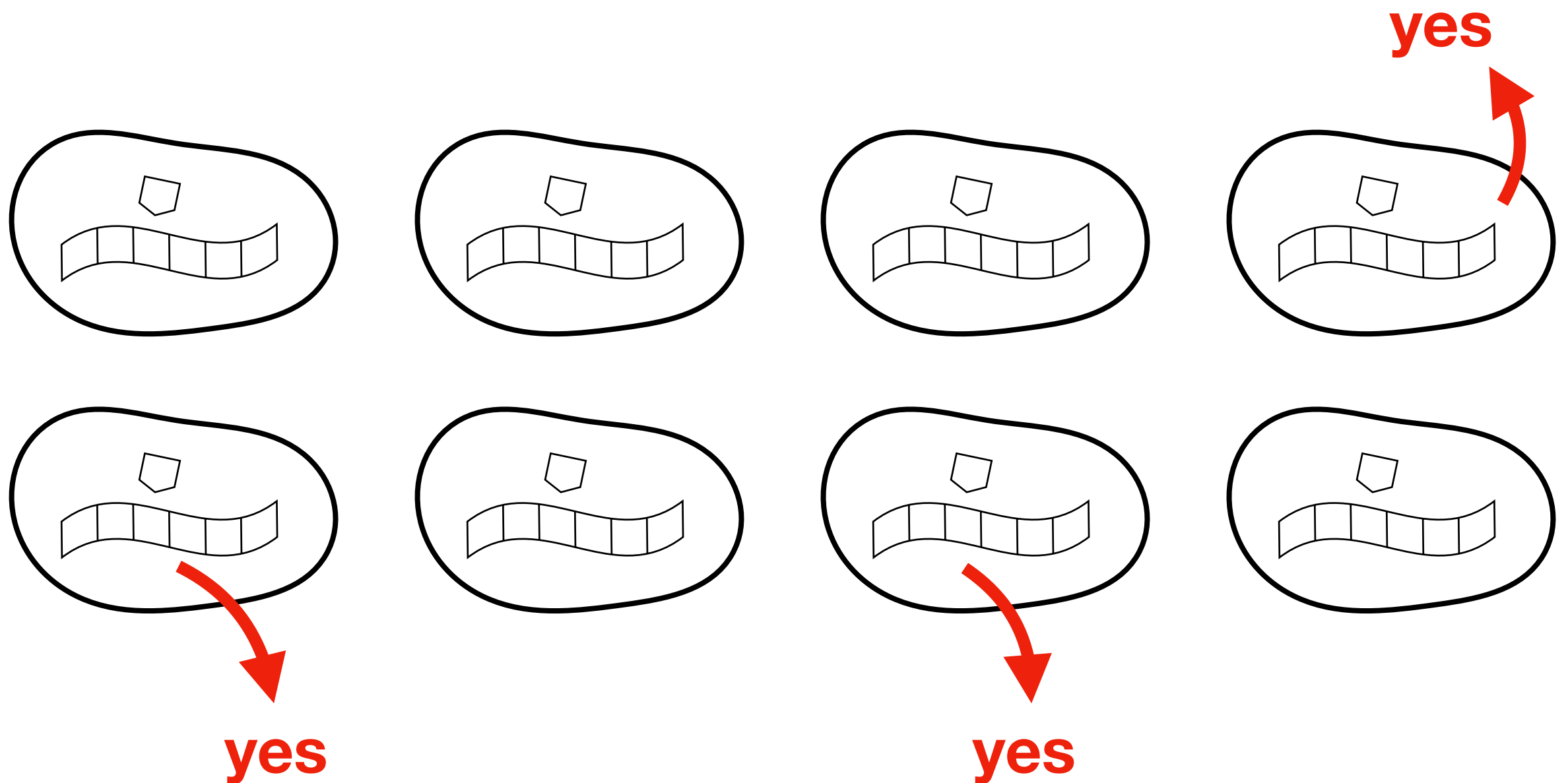
Simulating NDTM \Rightarrow NP



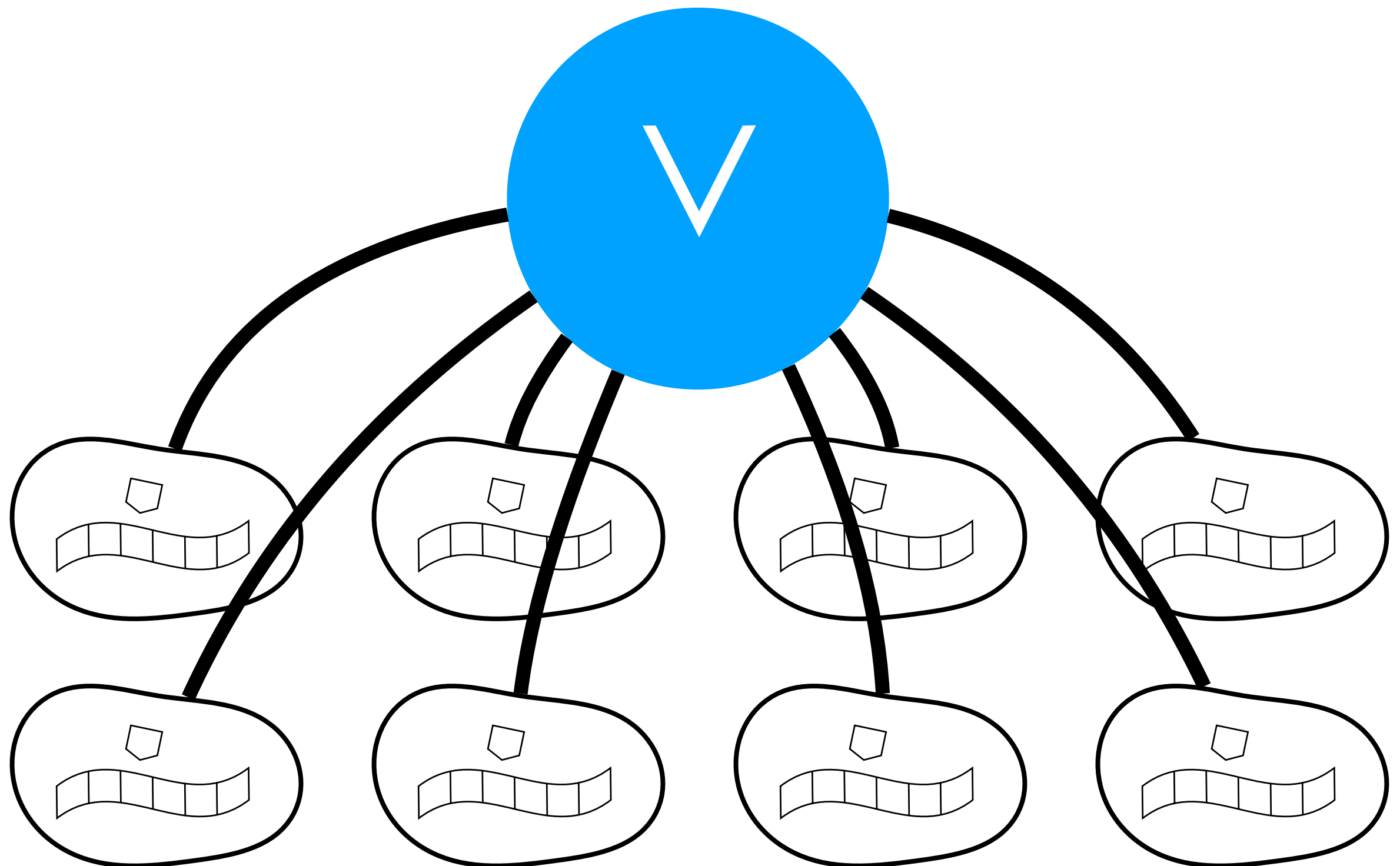
Simulating NDTM \Rightarrow NP



Simulating NDTM \Rightarrow NP



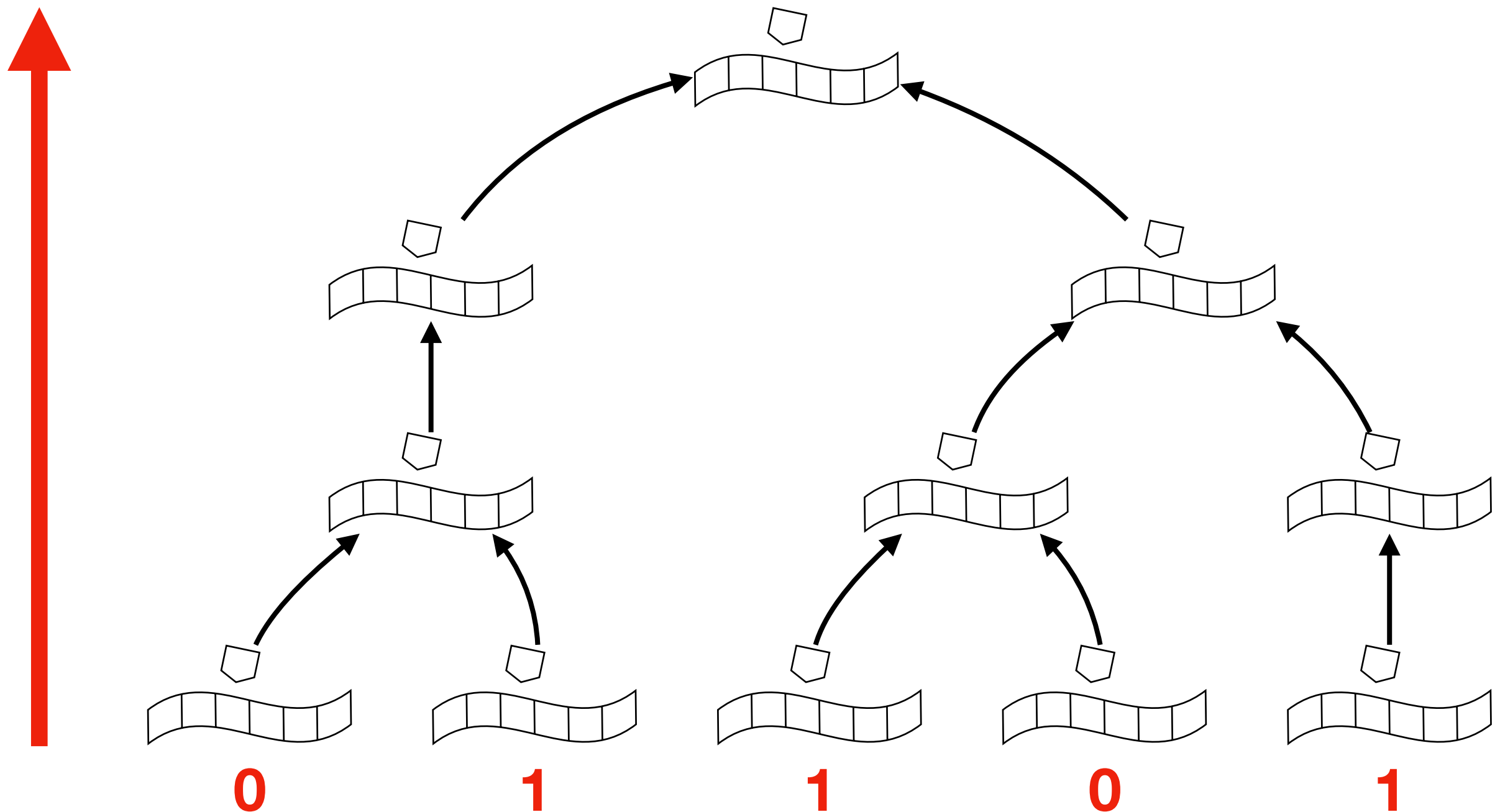
Simulating NDTM \Rightarrow NP



Counting complexity

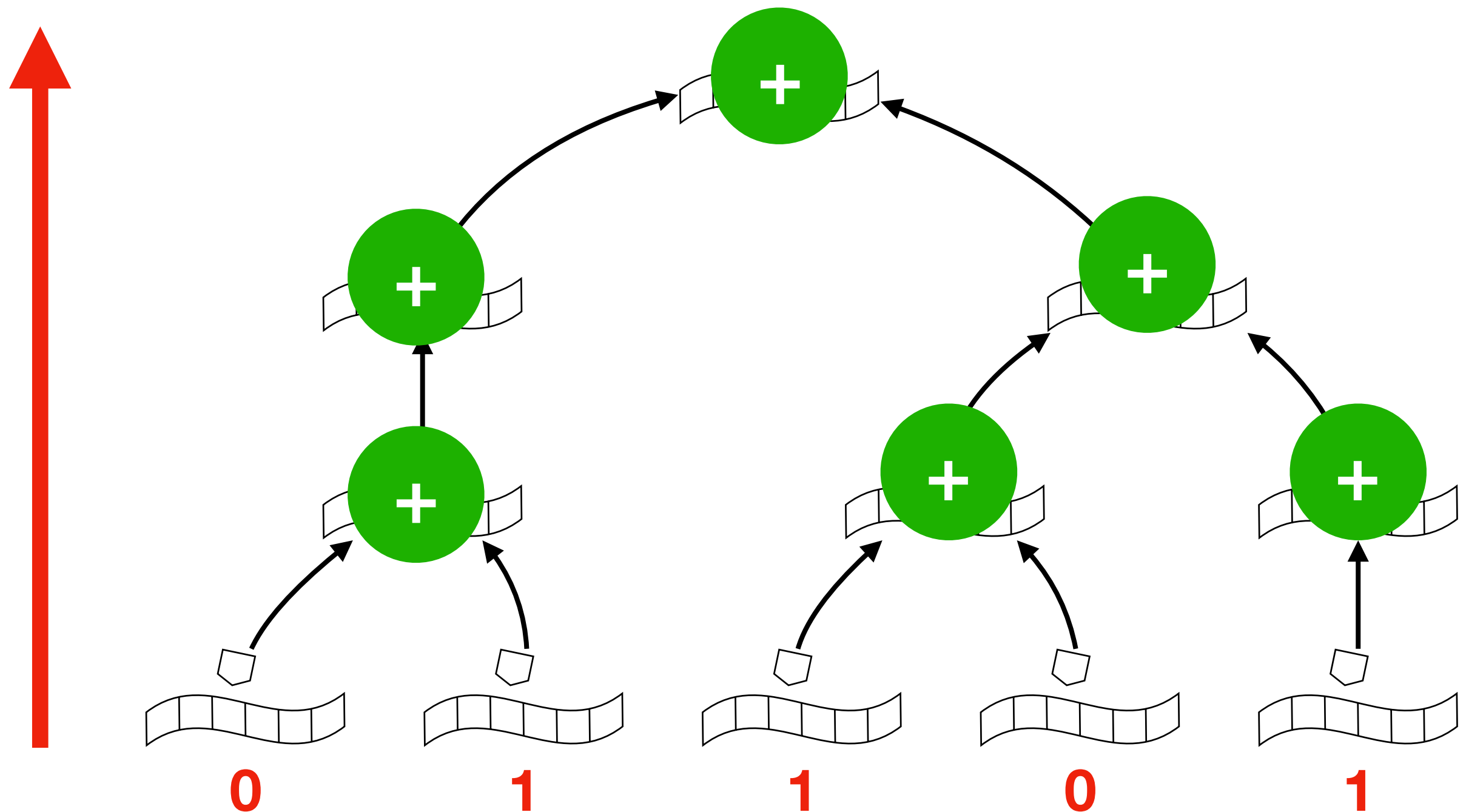
Counting

Turing machines: #P



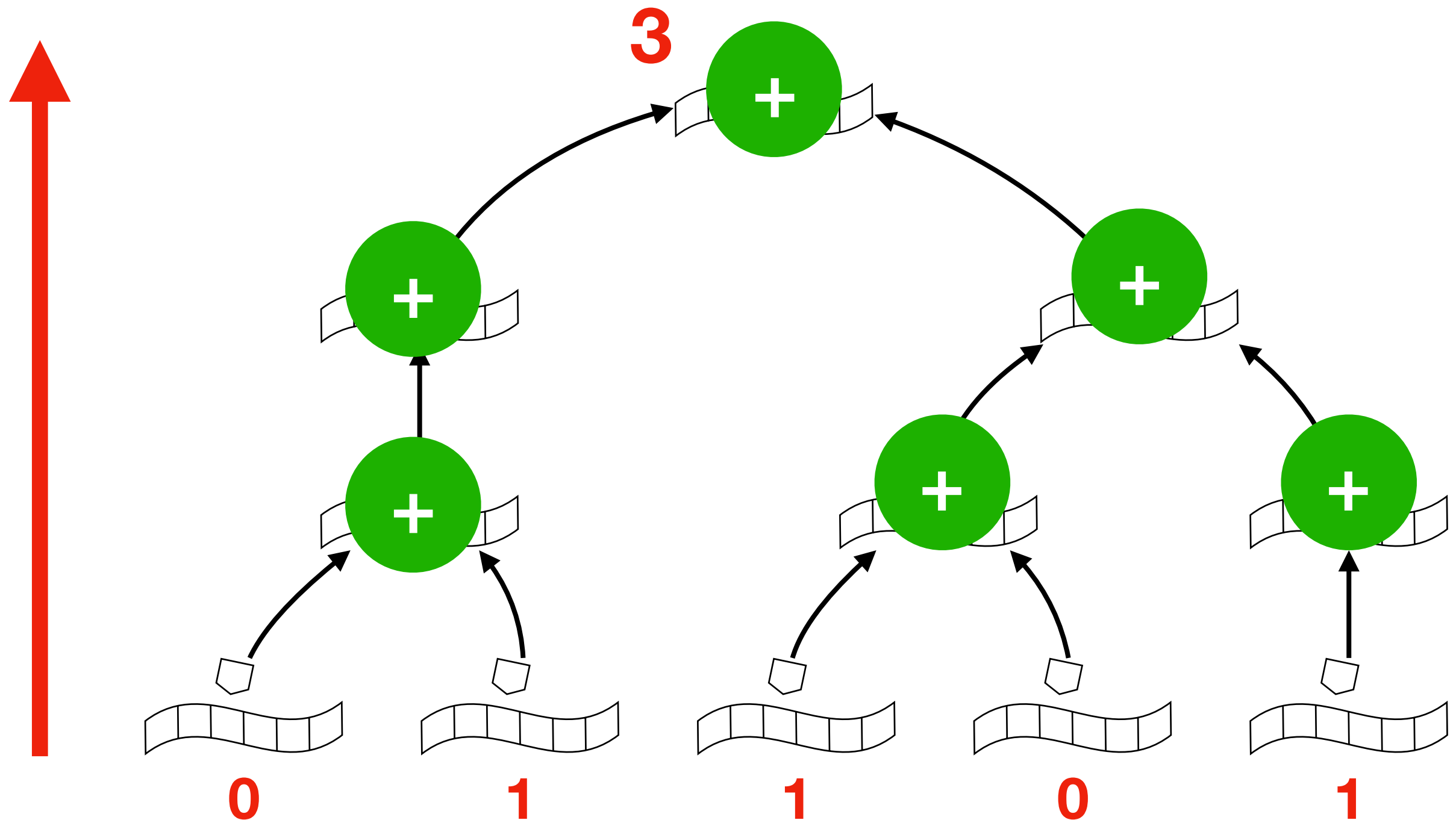
Counting

Turing machines: #P

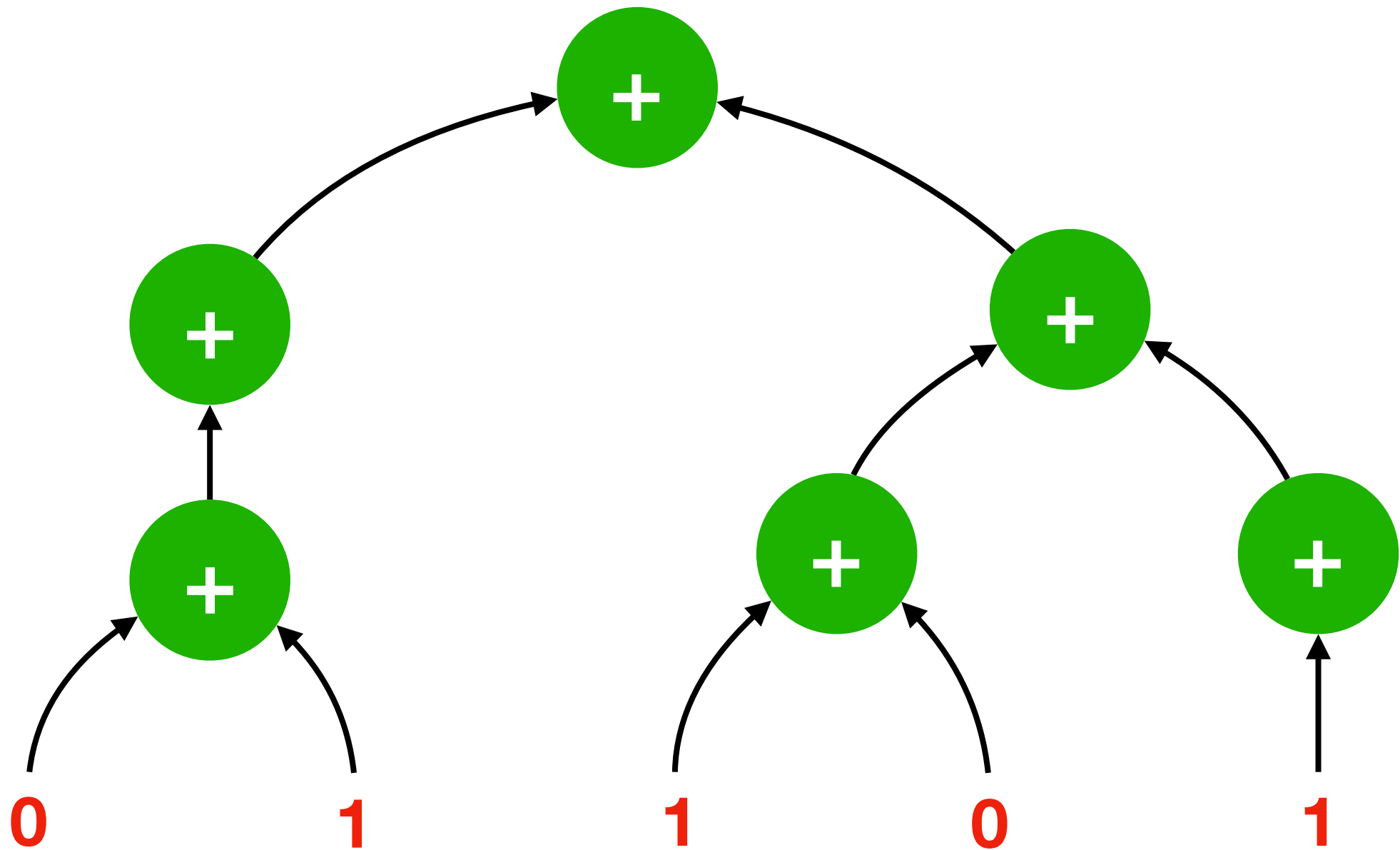


Counting

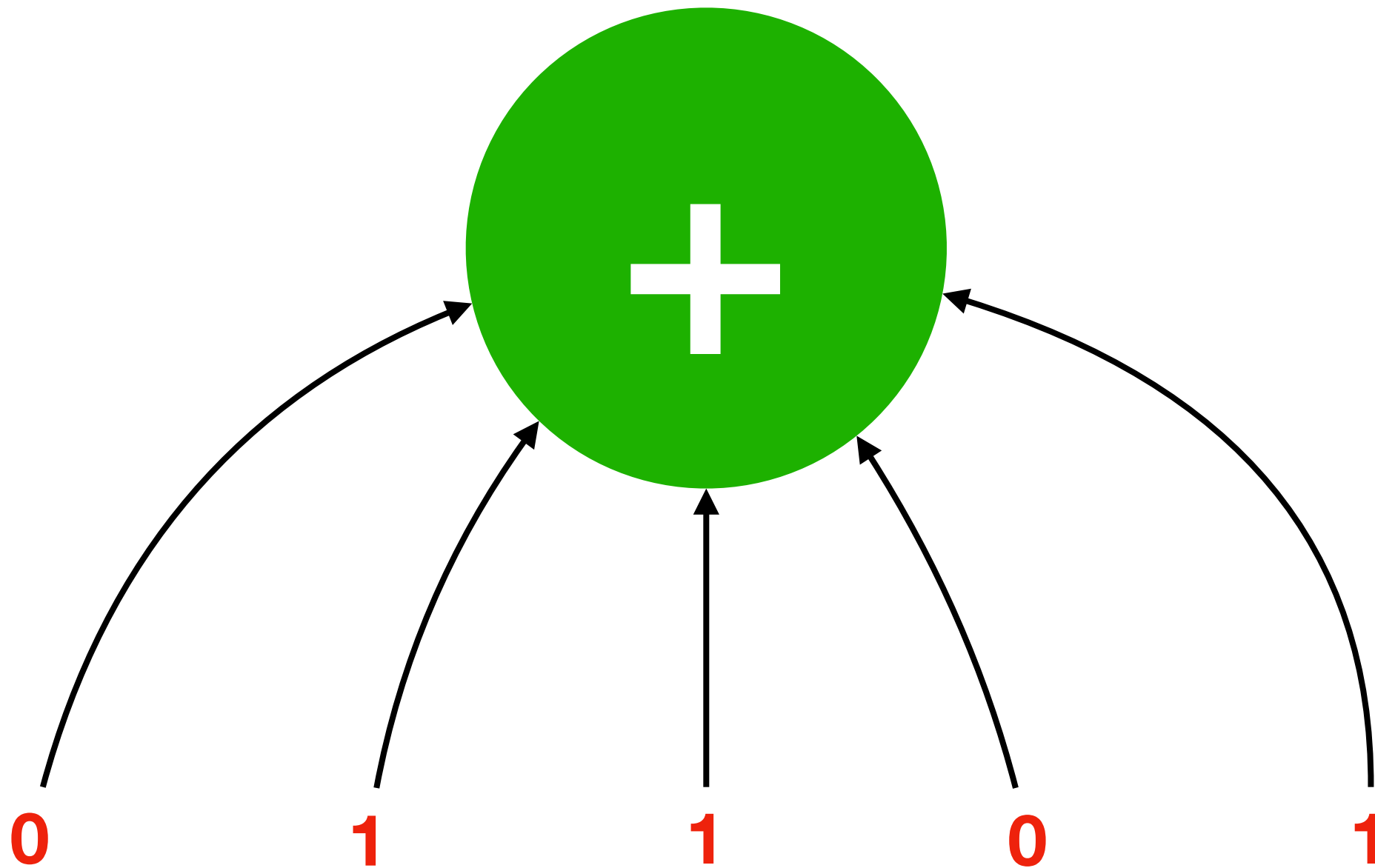
Turing machines: #P



Flattening \oplus -circuits

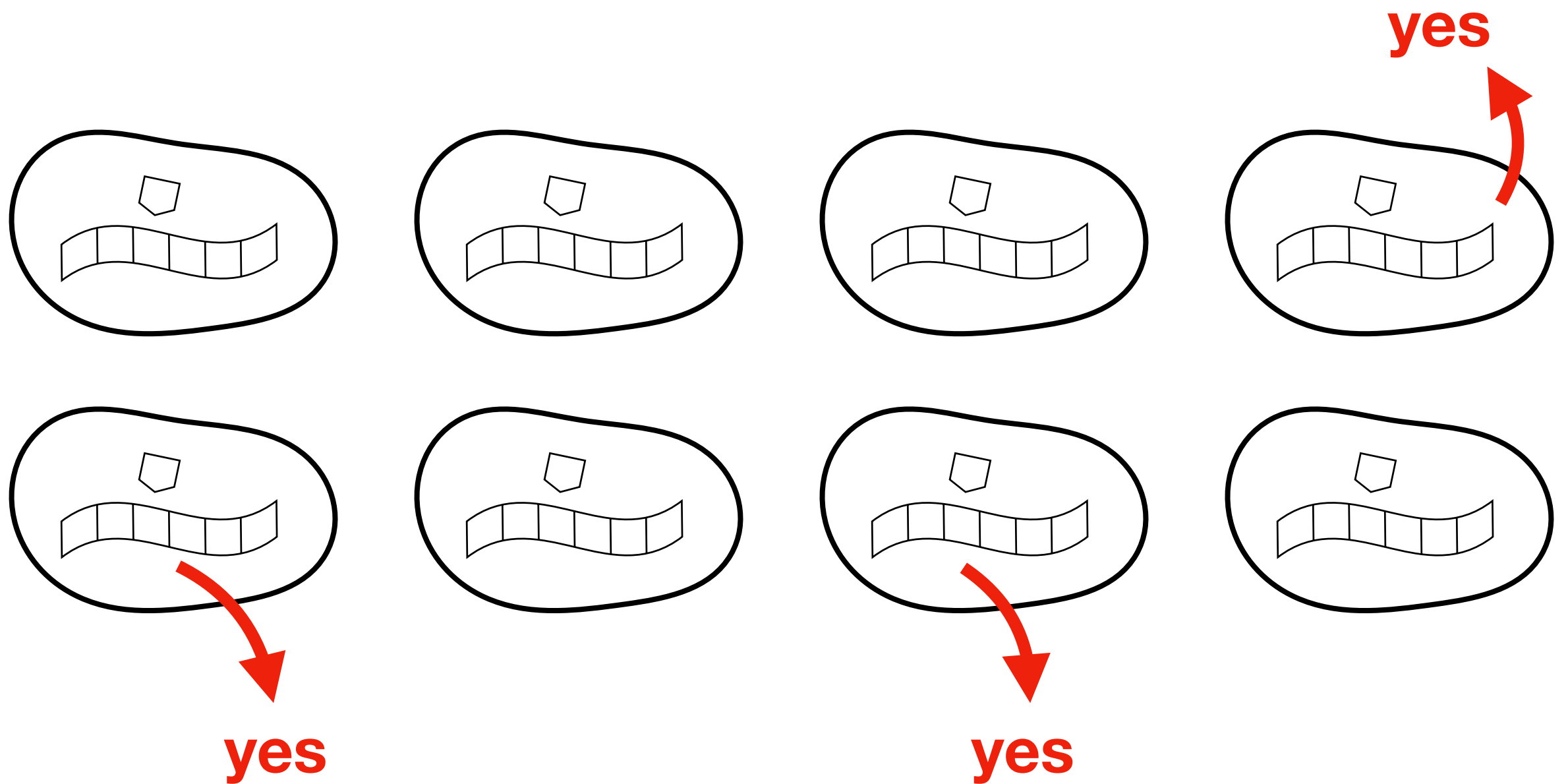


Flattening \oplus -circuits

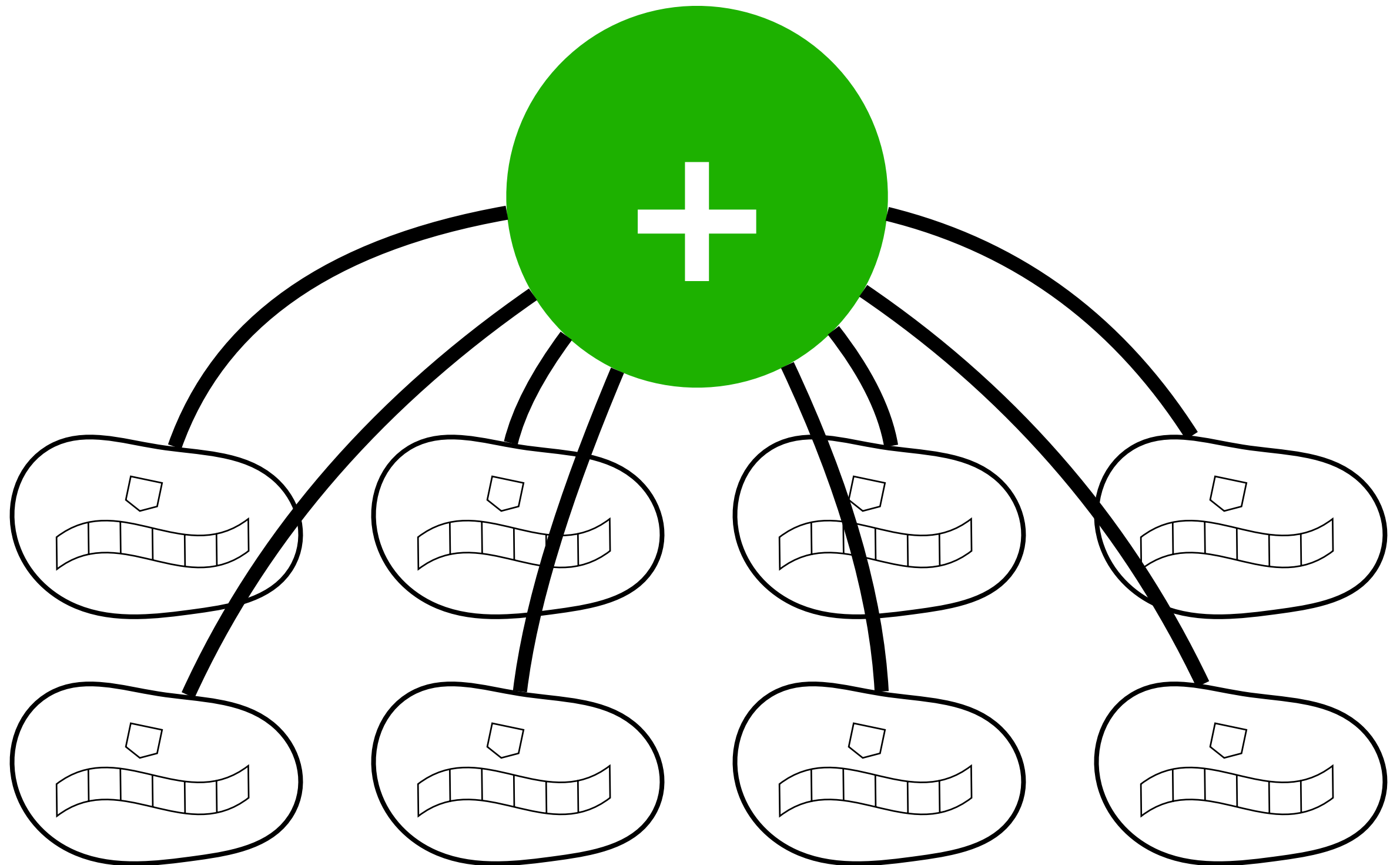


Counting with membrane systems

Simulating CTM \Rightarrow #P



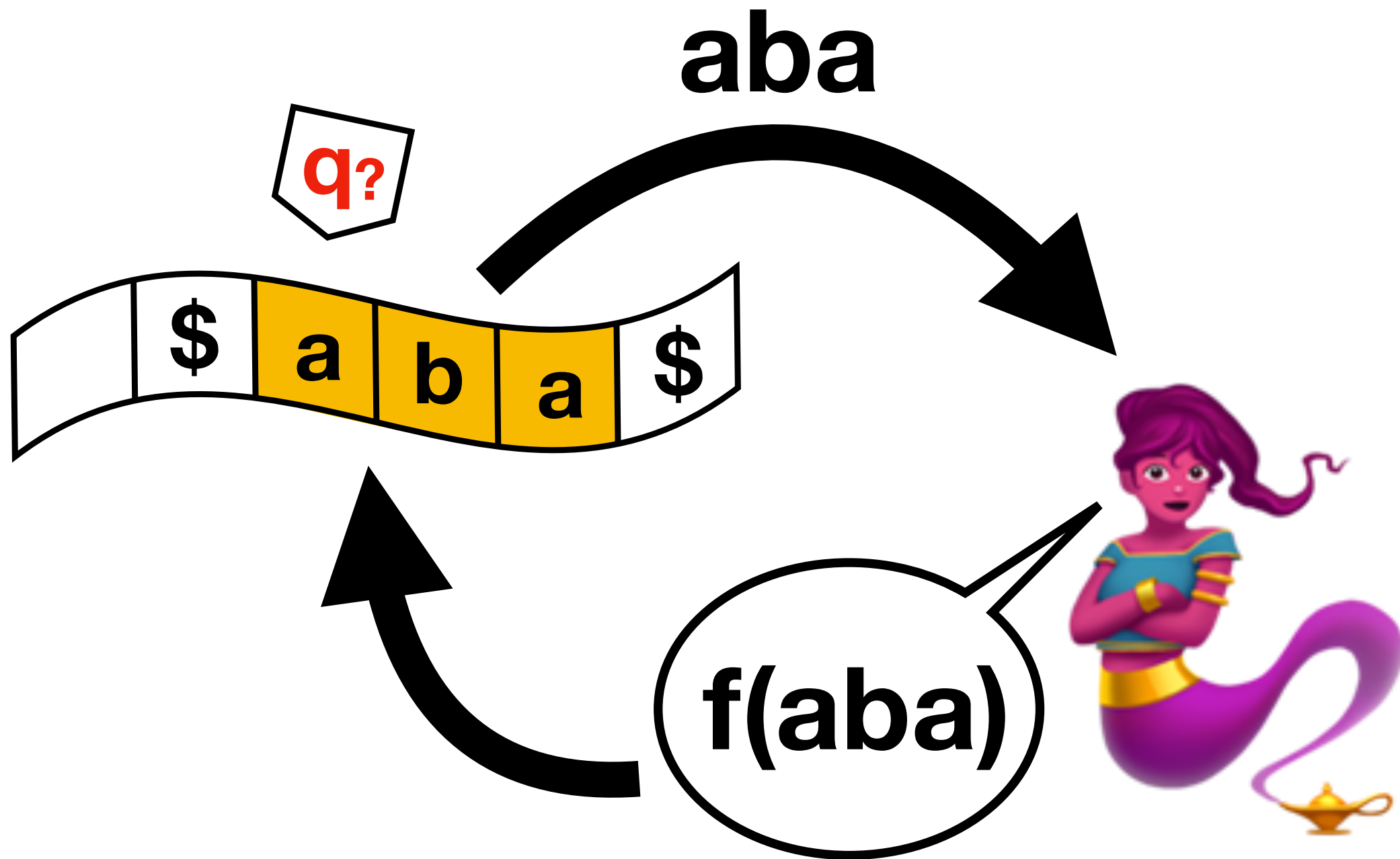
Simulating CTM \Rightarrow #P



Oracles

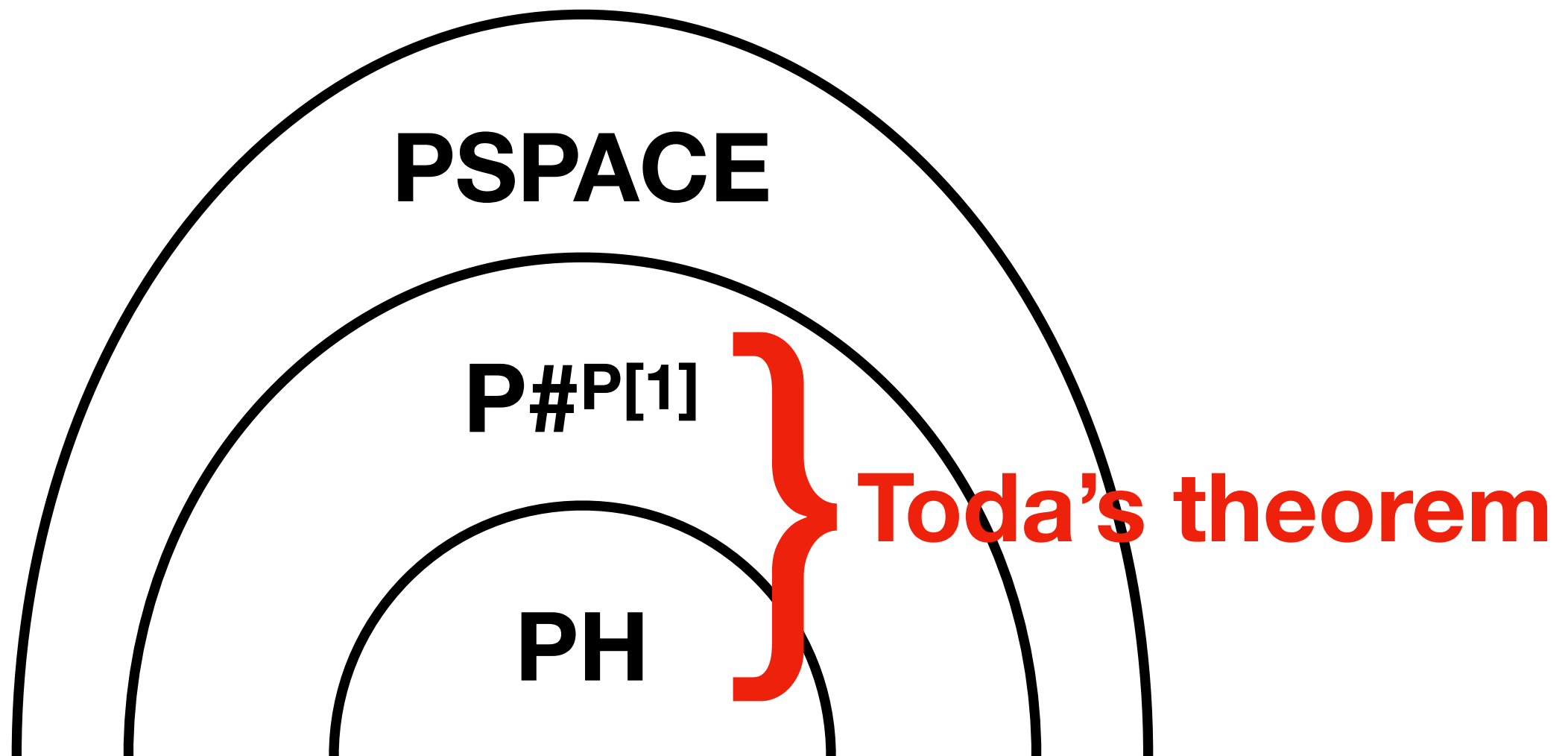


TM with oracle



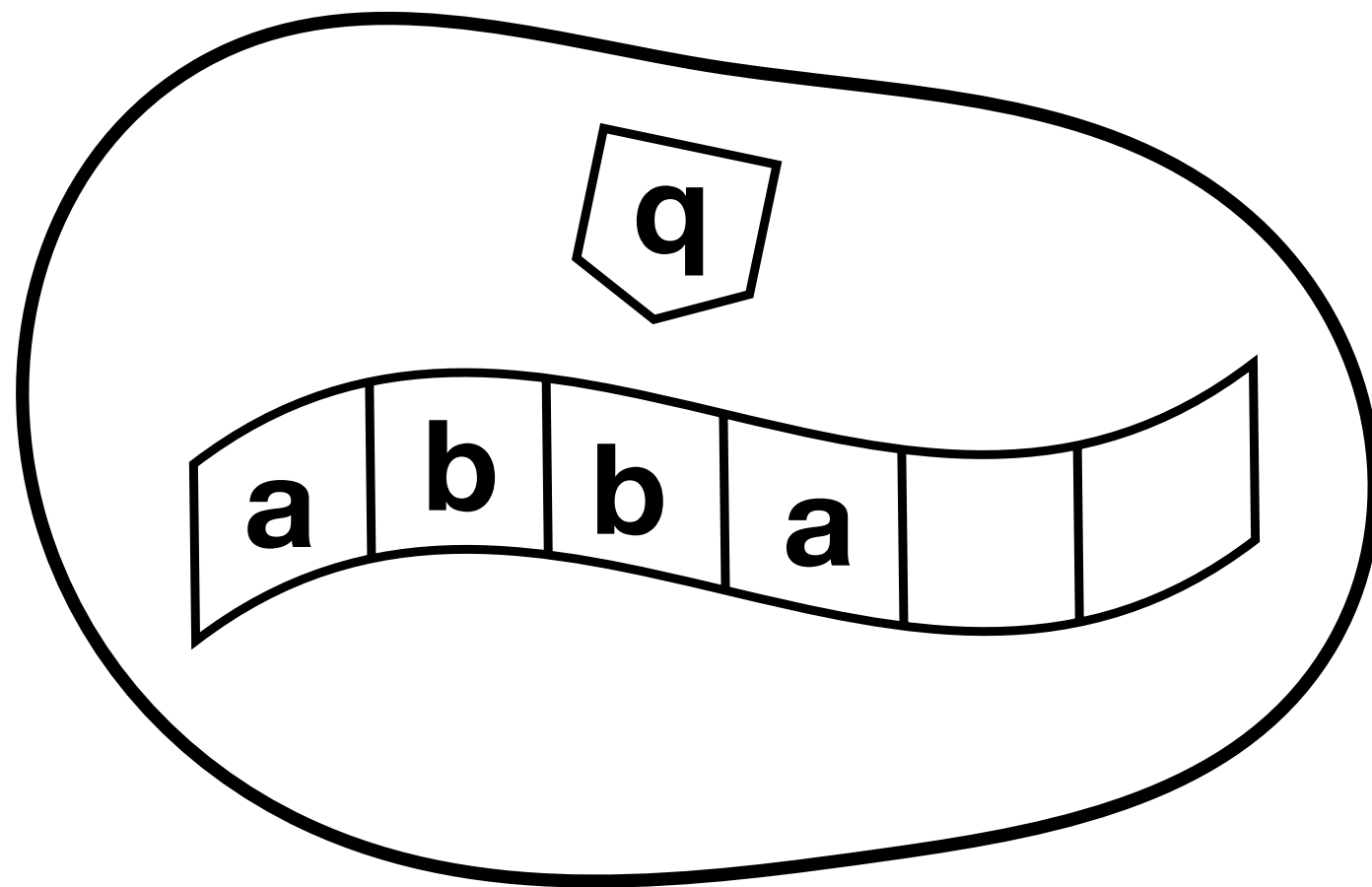
The complexity class $P^{\#P}[1]$

The class of problems solved by polynomial-time Turing machines with a **single query** to an oracle for a **#P**-complete problem

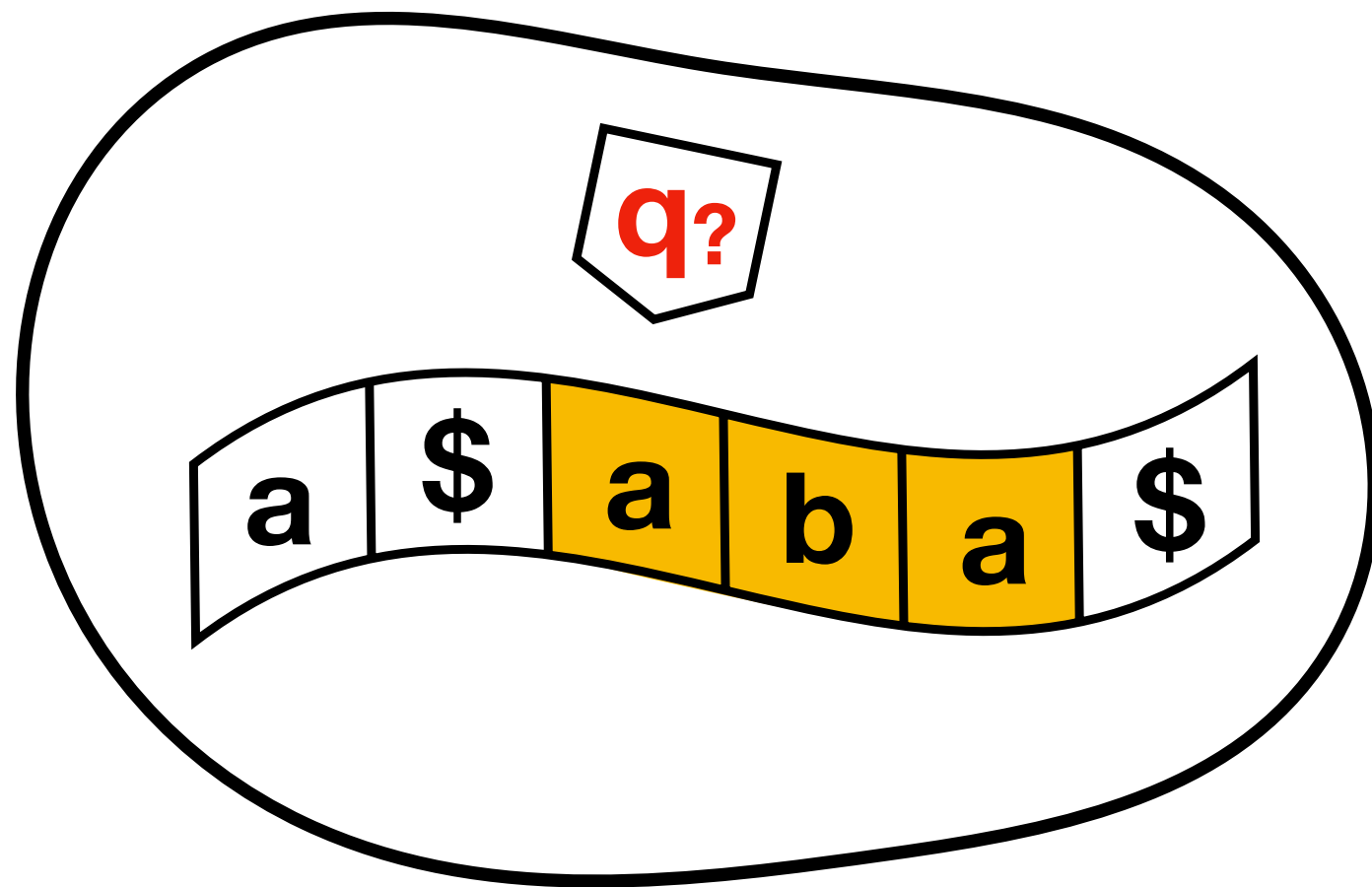


**Solving $P^{\#}P[1]$ with
(monodirectional, shallow)
membrane systems**

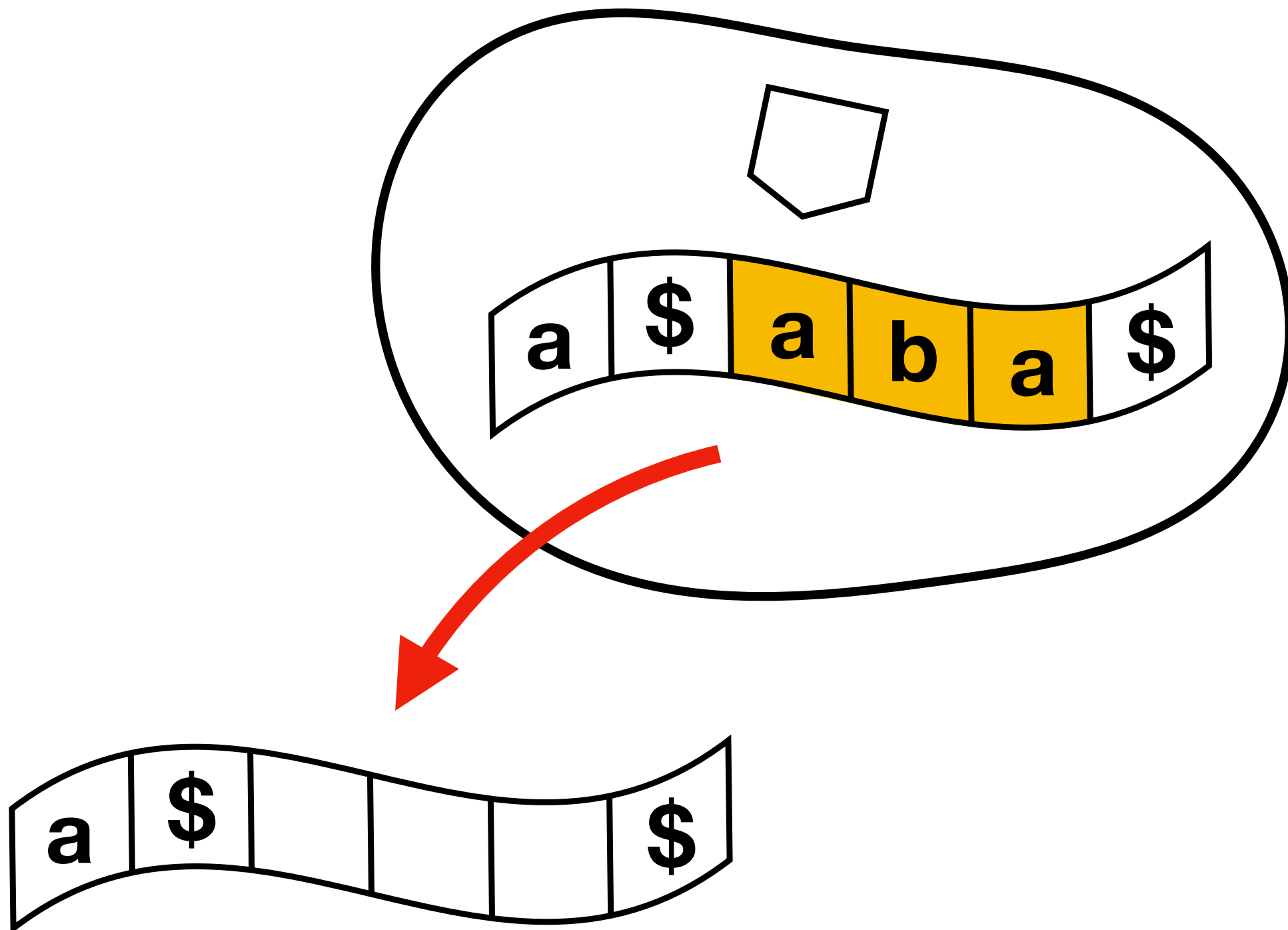
Pre-query computation



Entering the query state

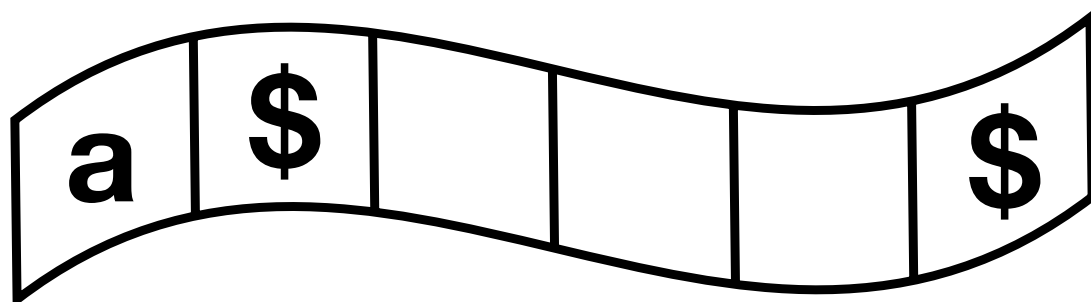
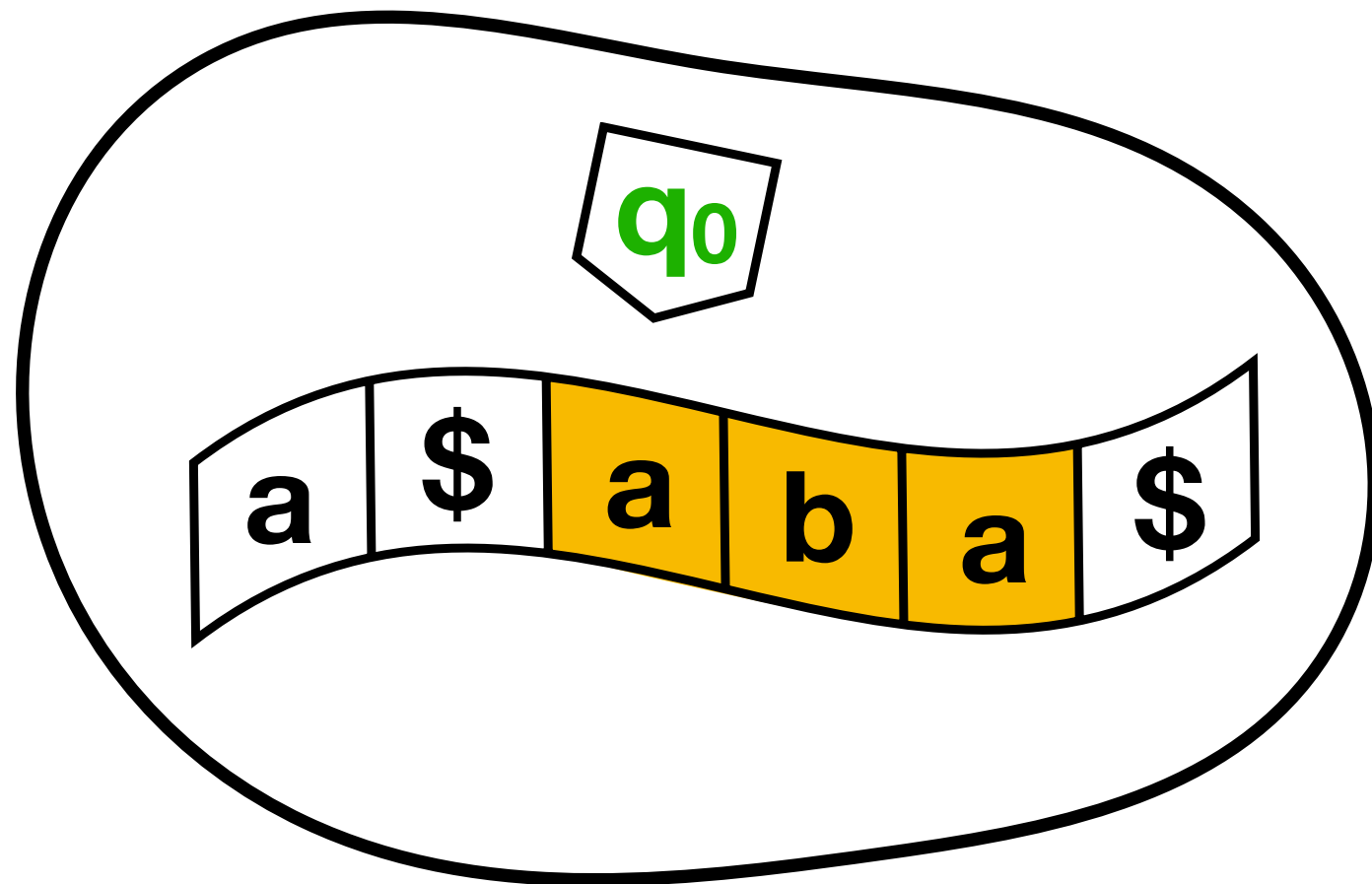


Entering the query state

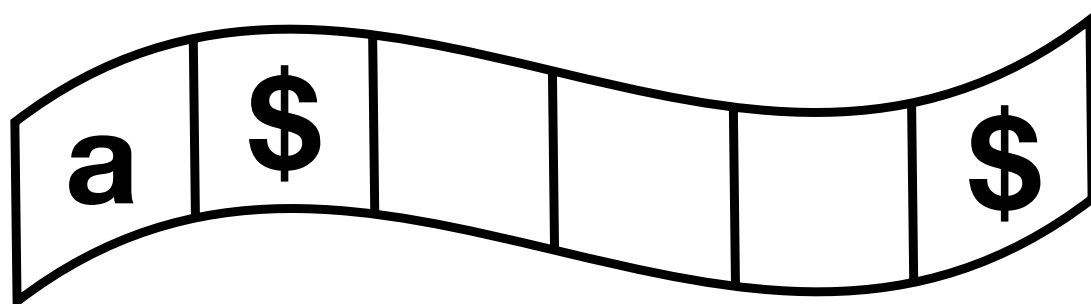
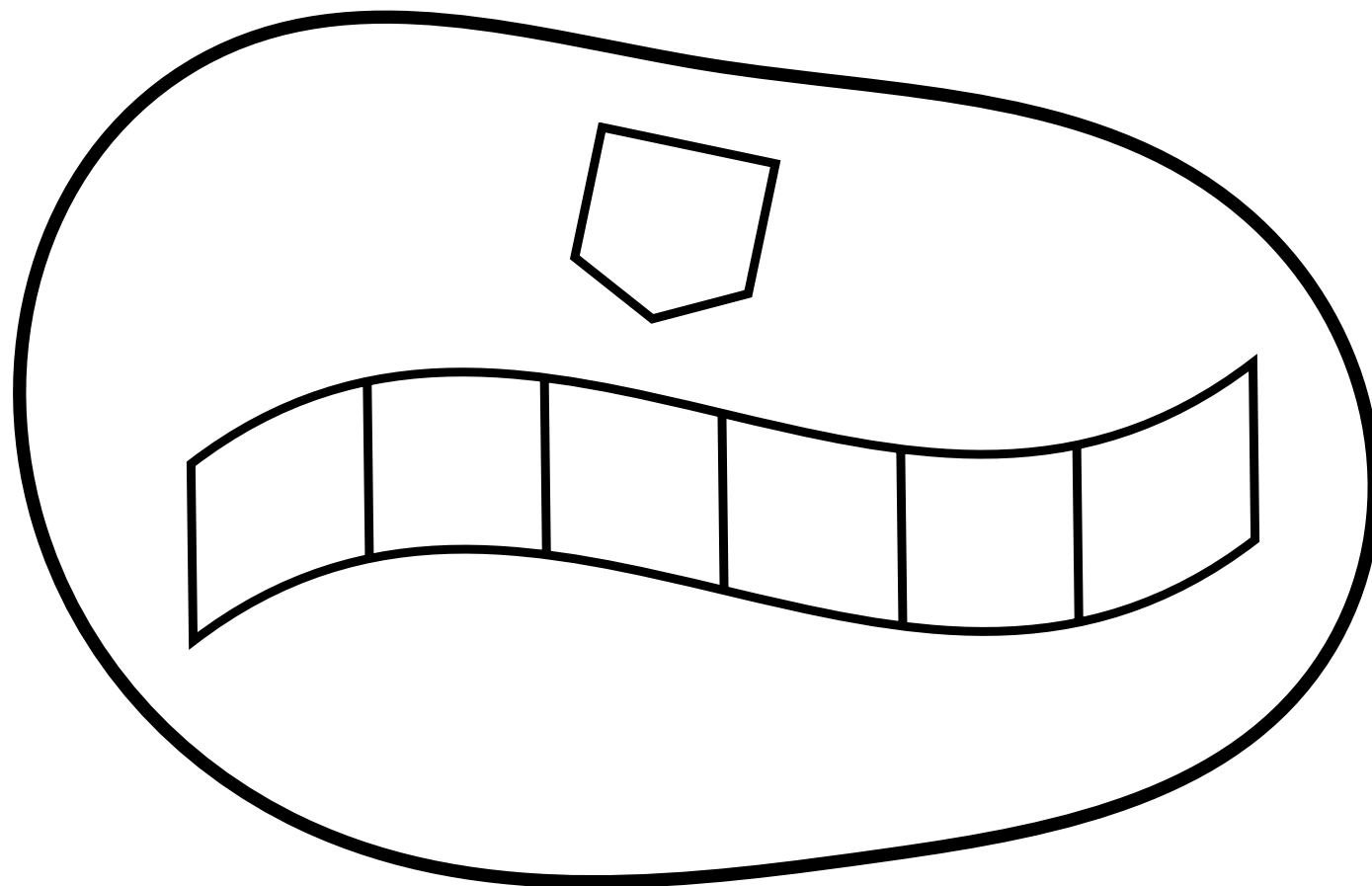


Entering the query state

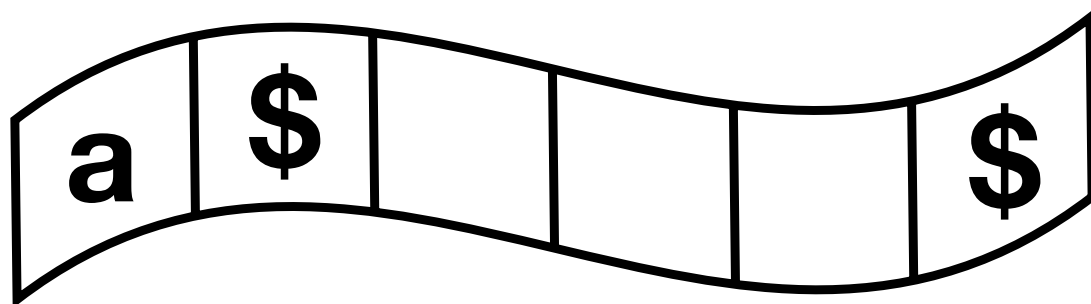
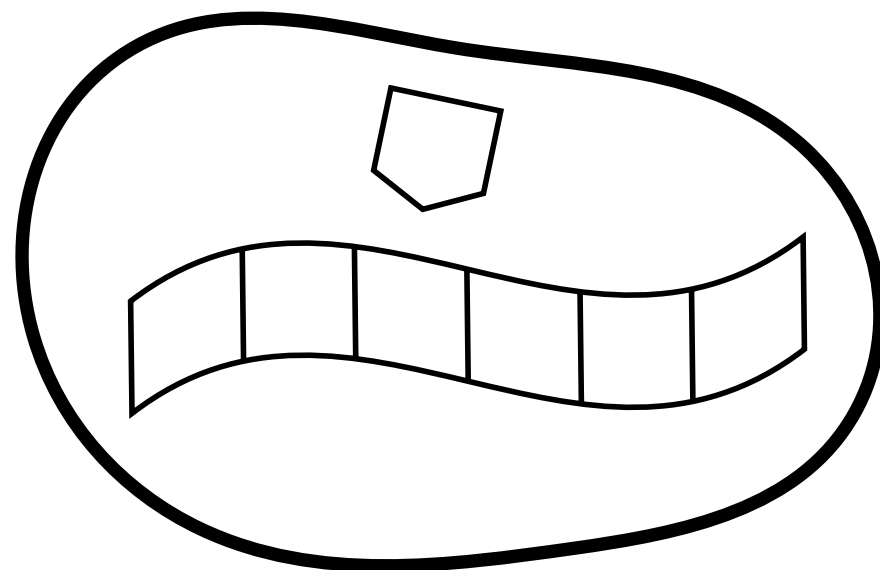
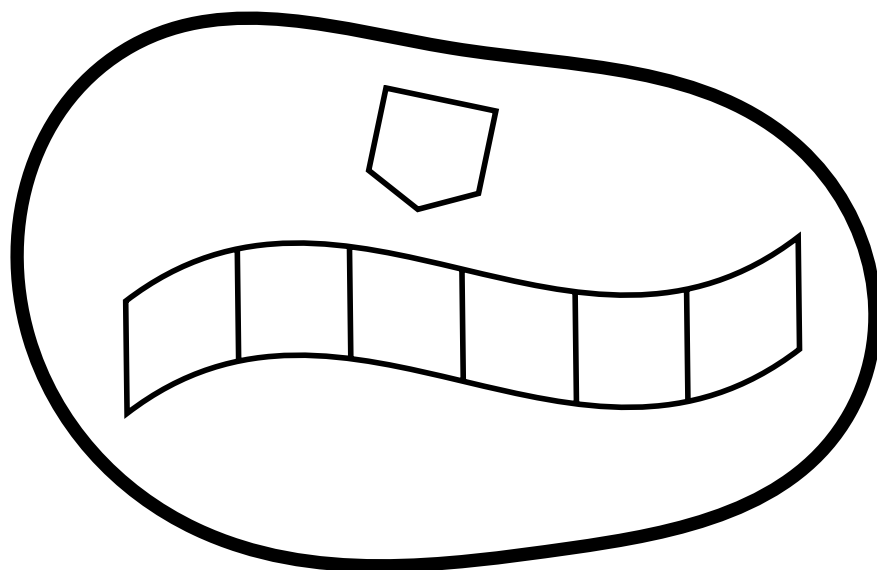
Initial state of a
TM solving
the #P-complete
oracle problem



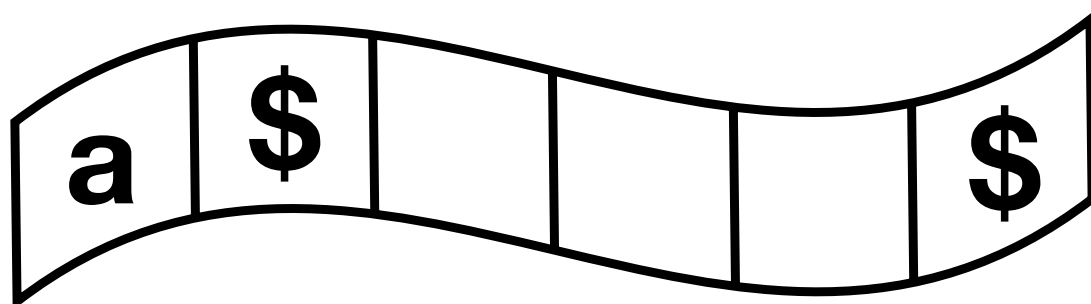
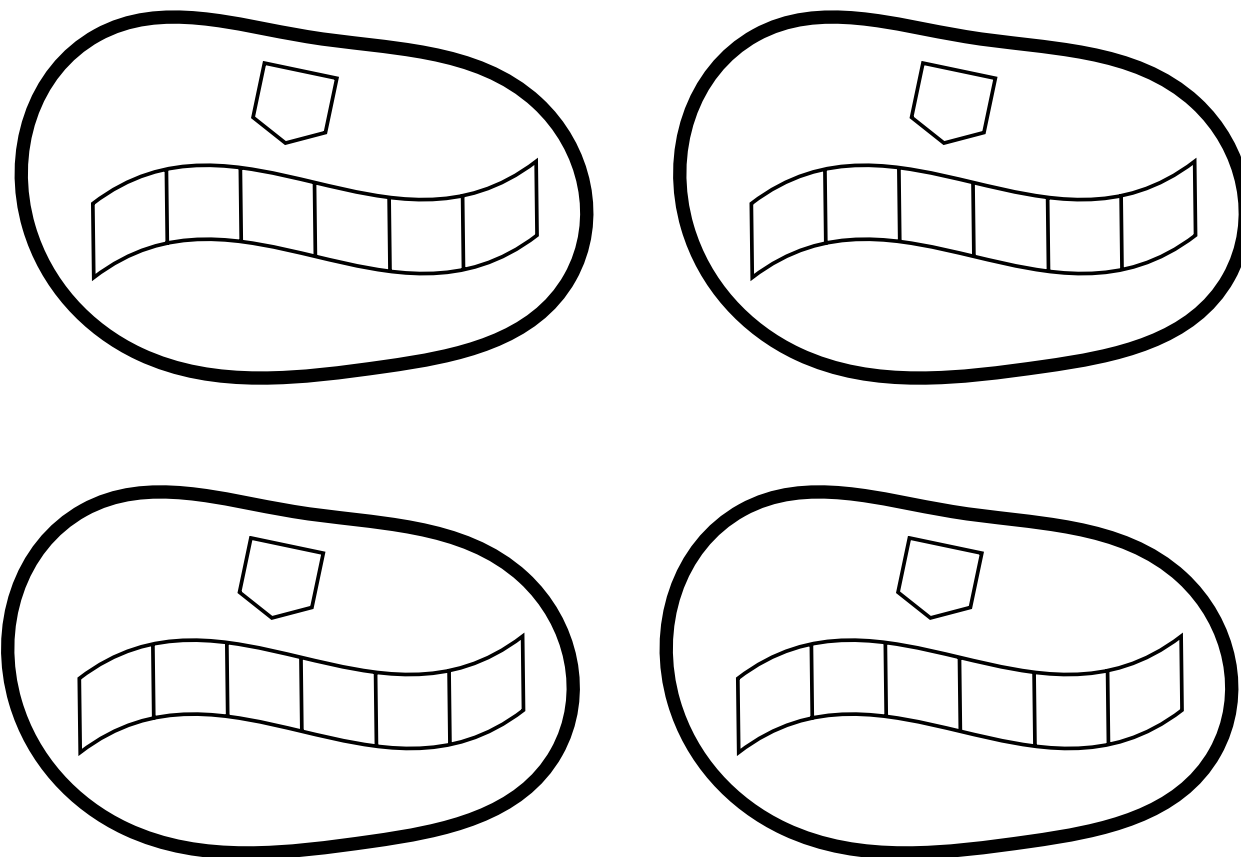
Simulating the auxiliary TM



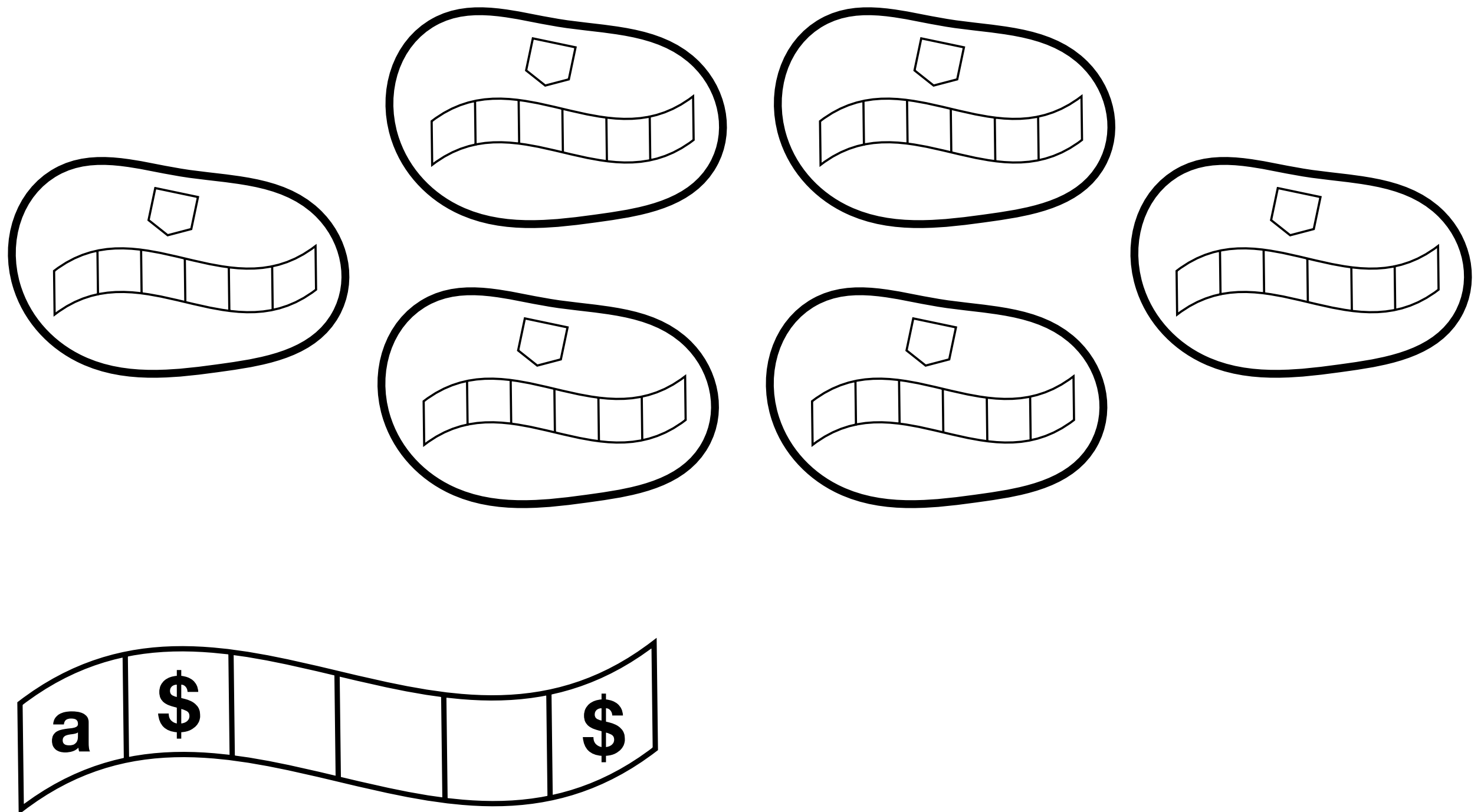
Simulating the auxiliary TM



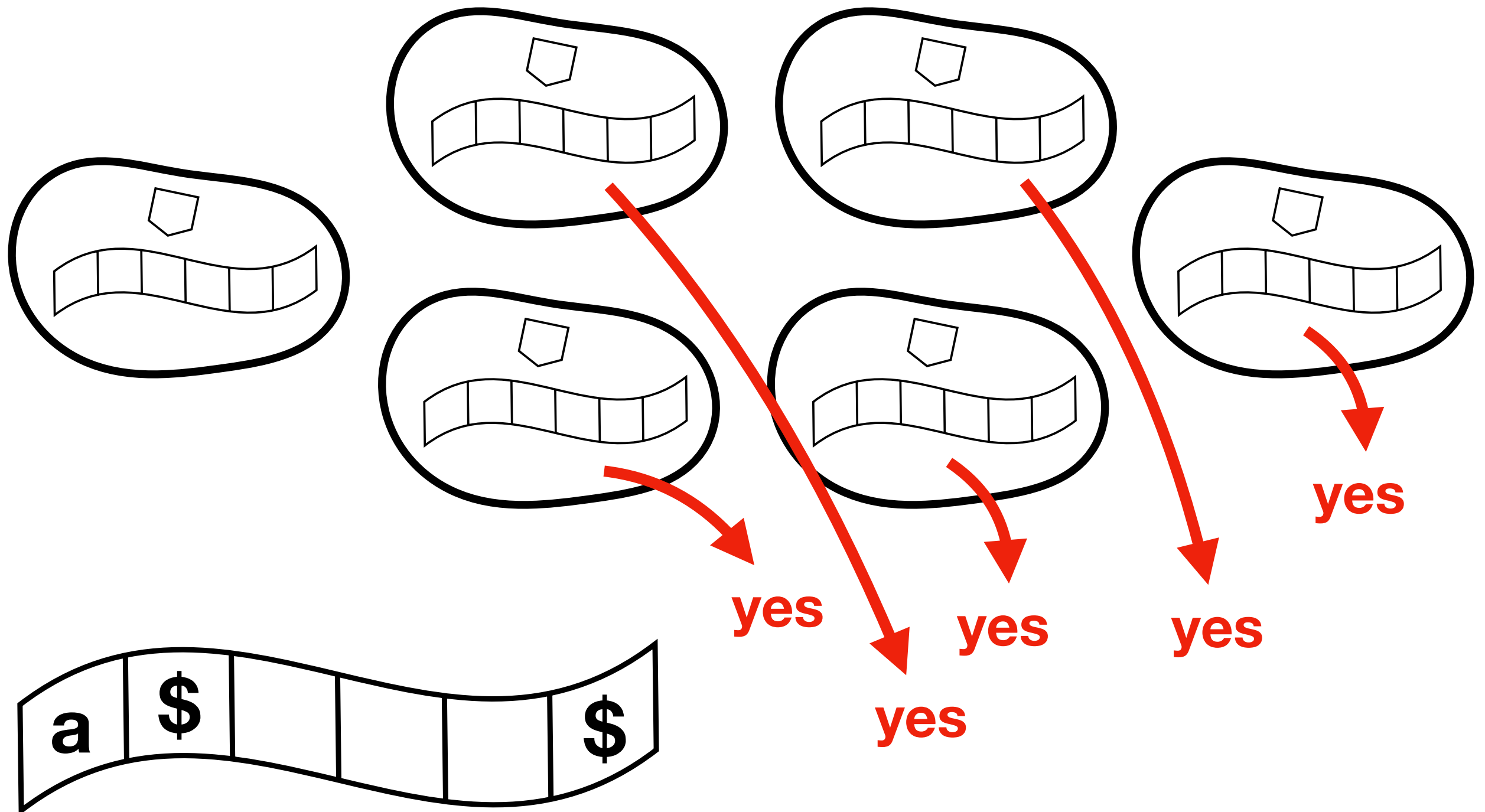
Simulating the auxiliary TM



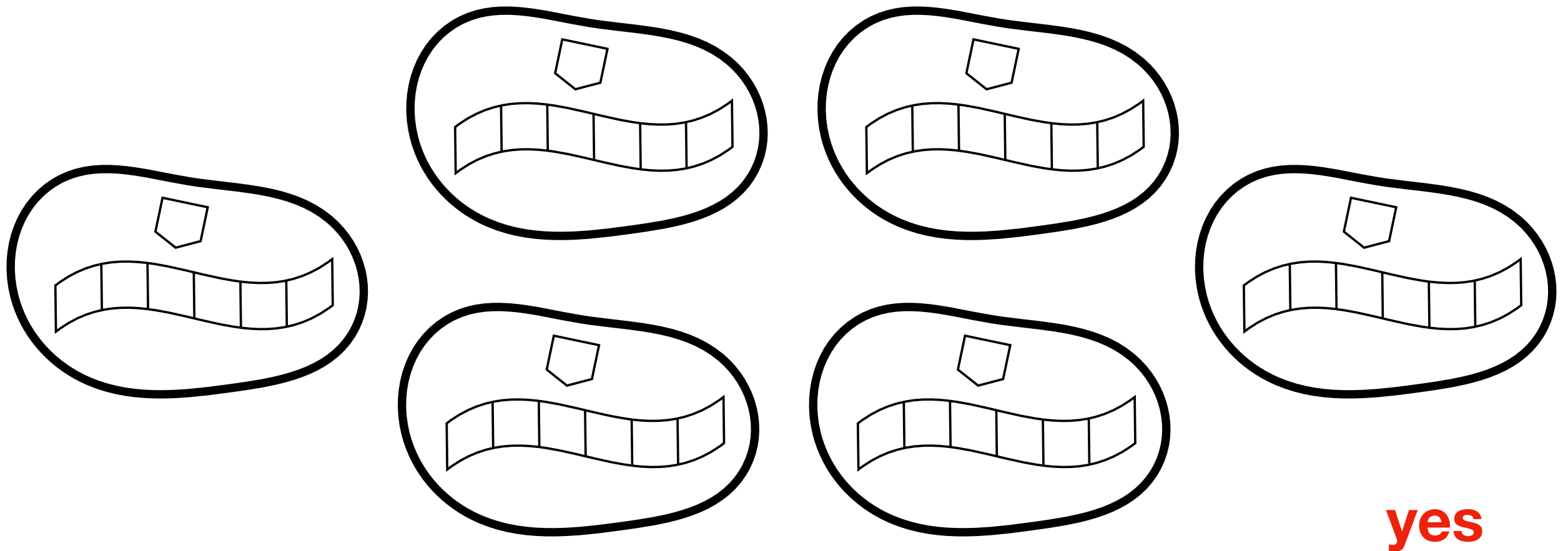
Simulating the auxiliary TM



Collecting the output



Converting unary to binary



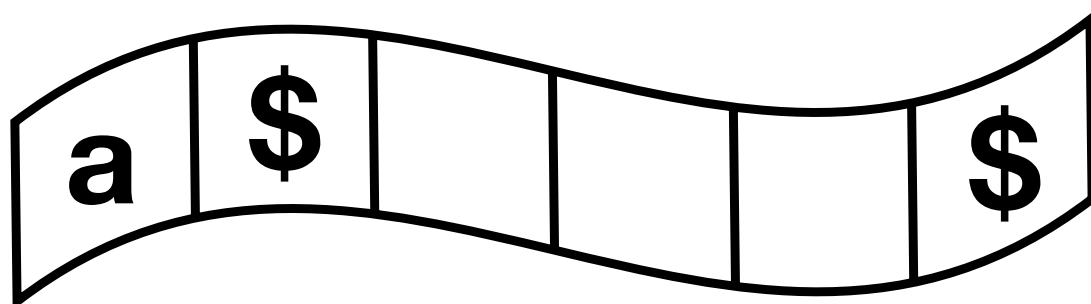
yes

yes

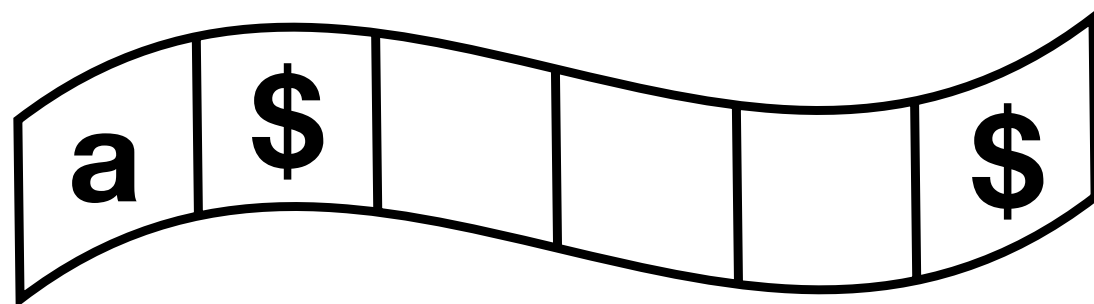
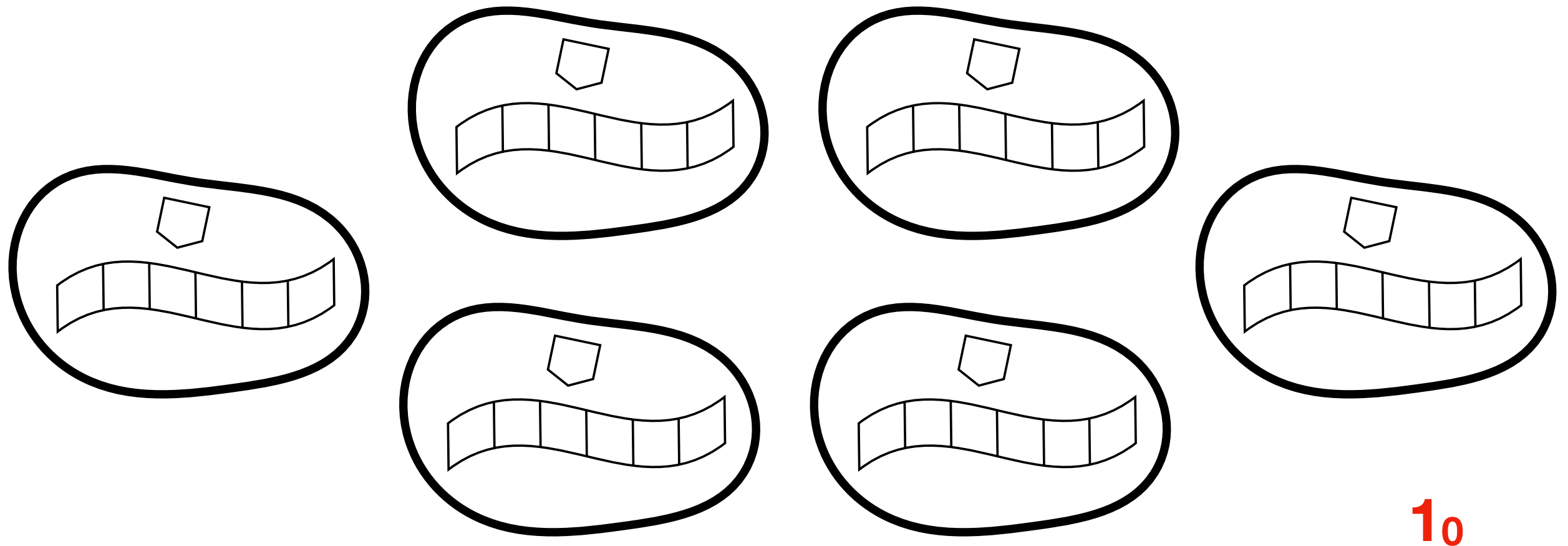
yes

yes

yes



Converting unary to binary



10

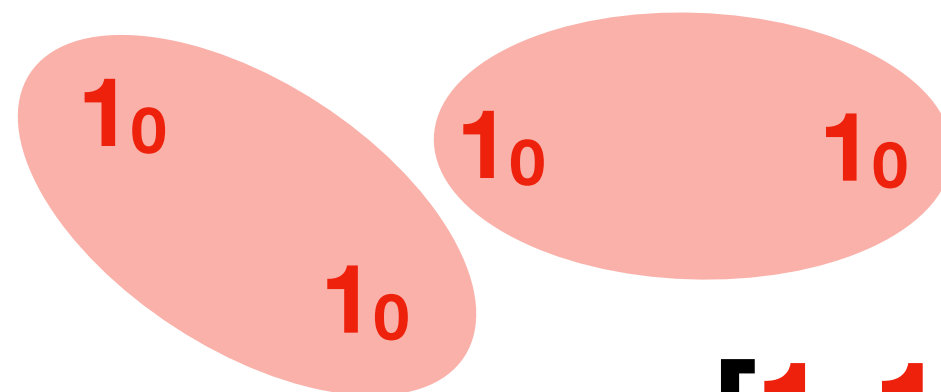
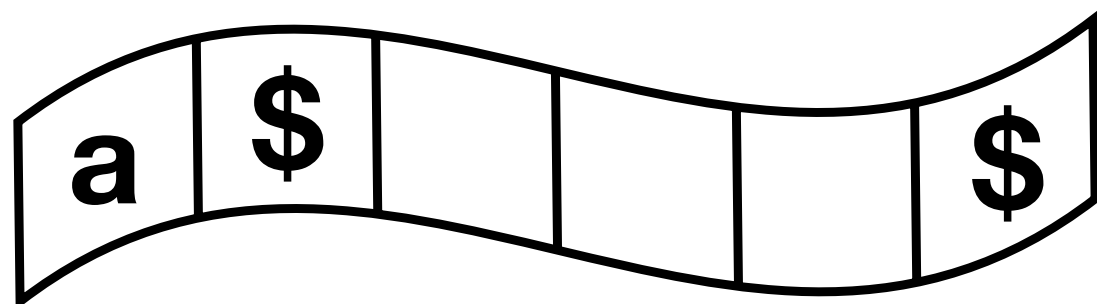
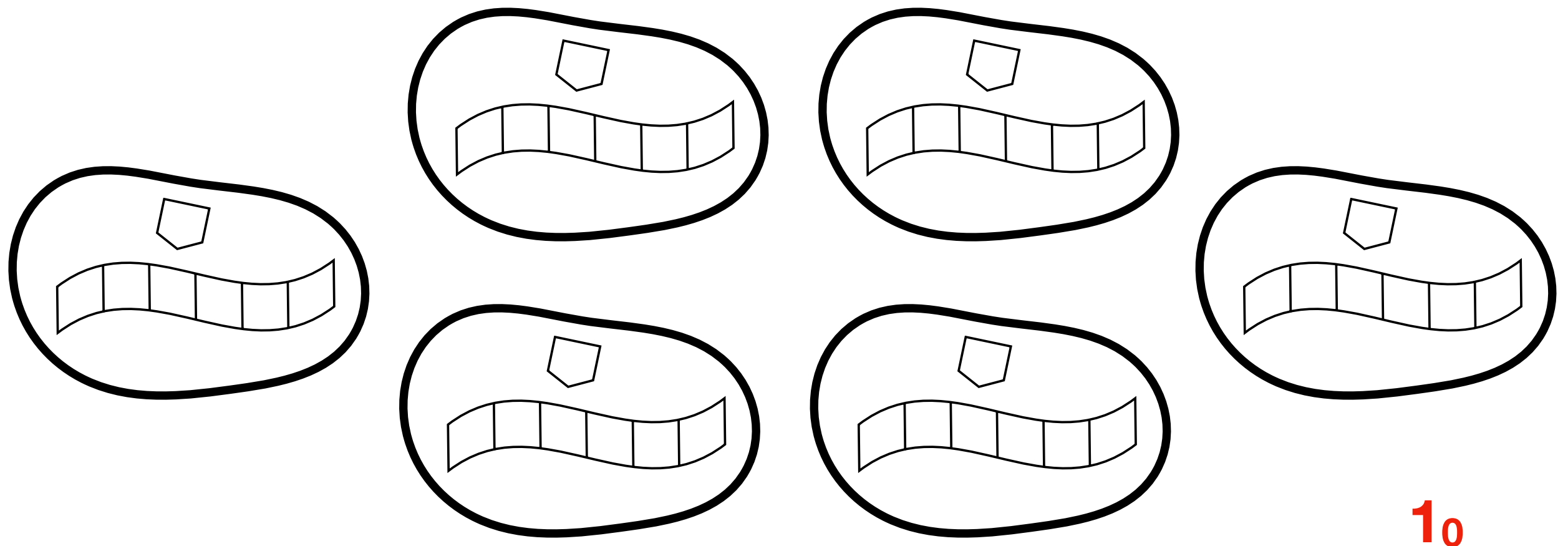
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10

10

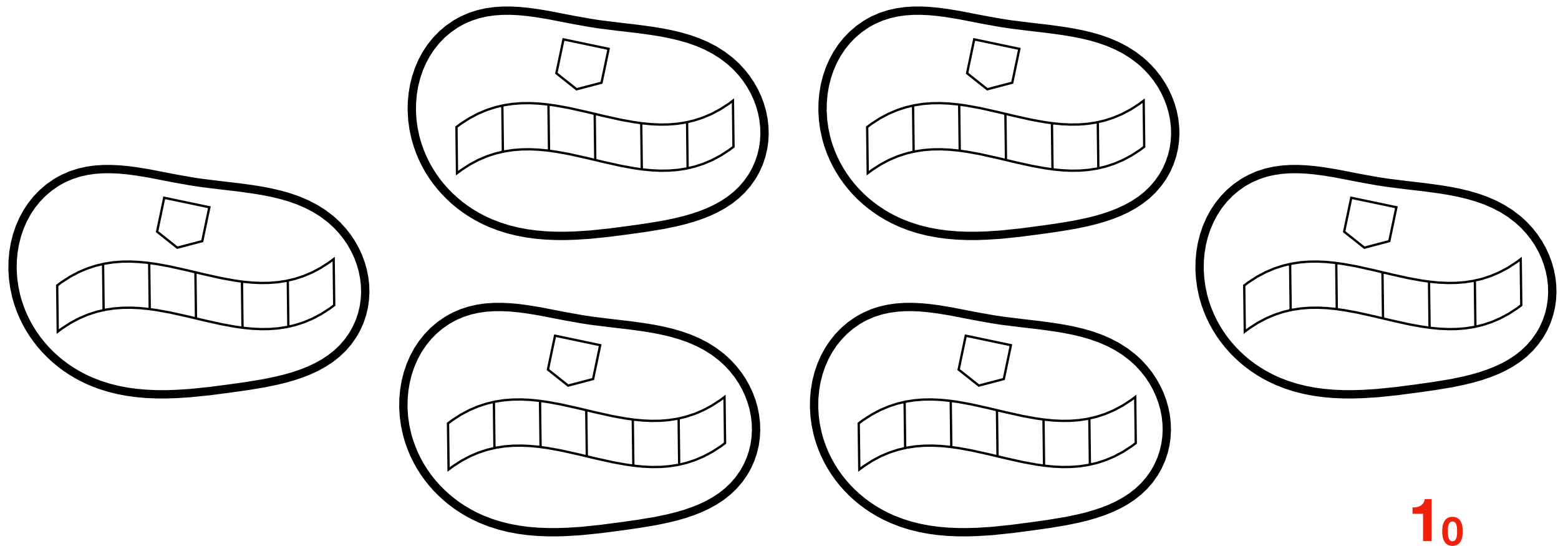
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Converting unary to binary



[1₀1₀ → 1₁]

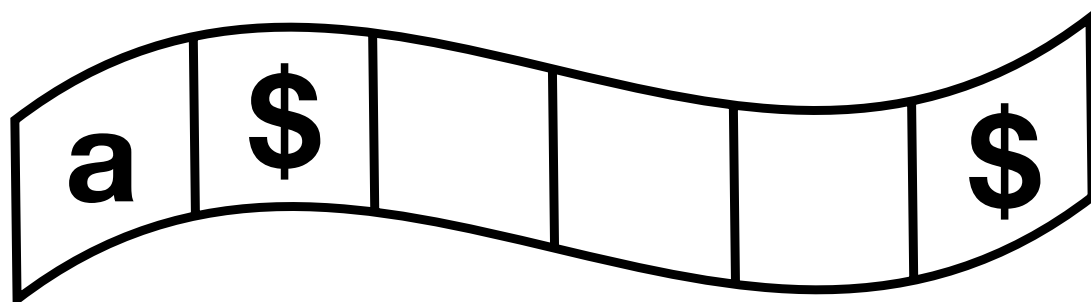
Converting unary to binary



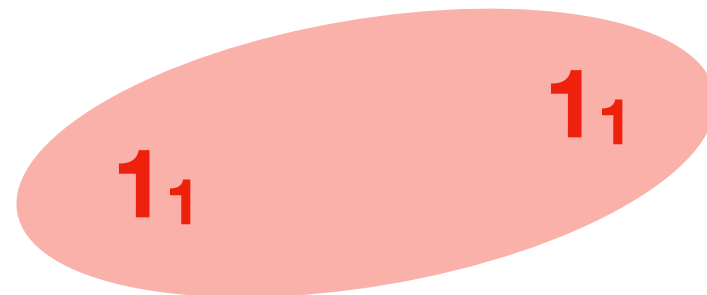
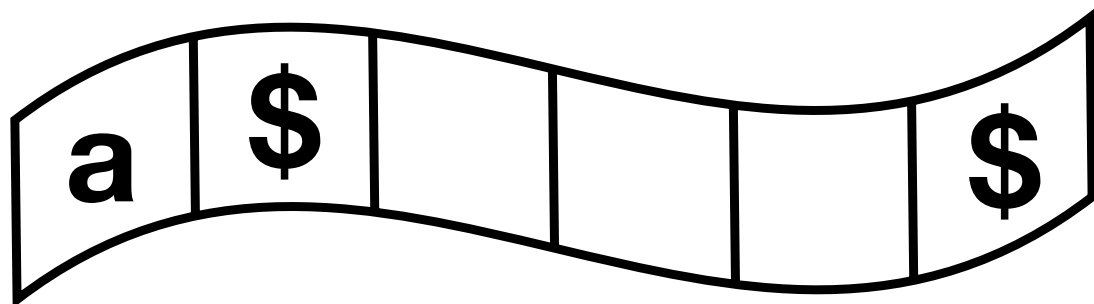
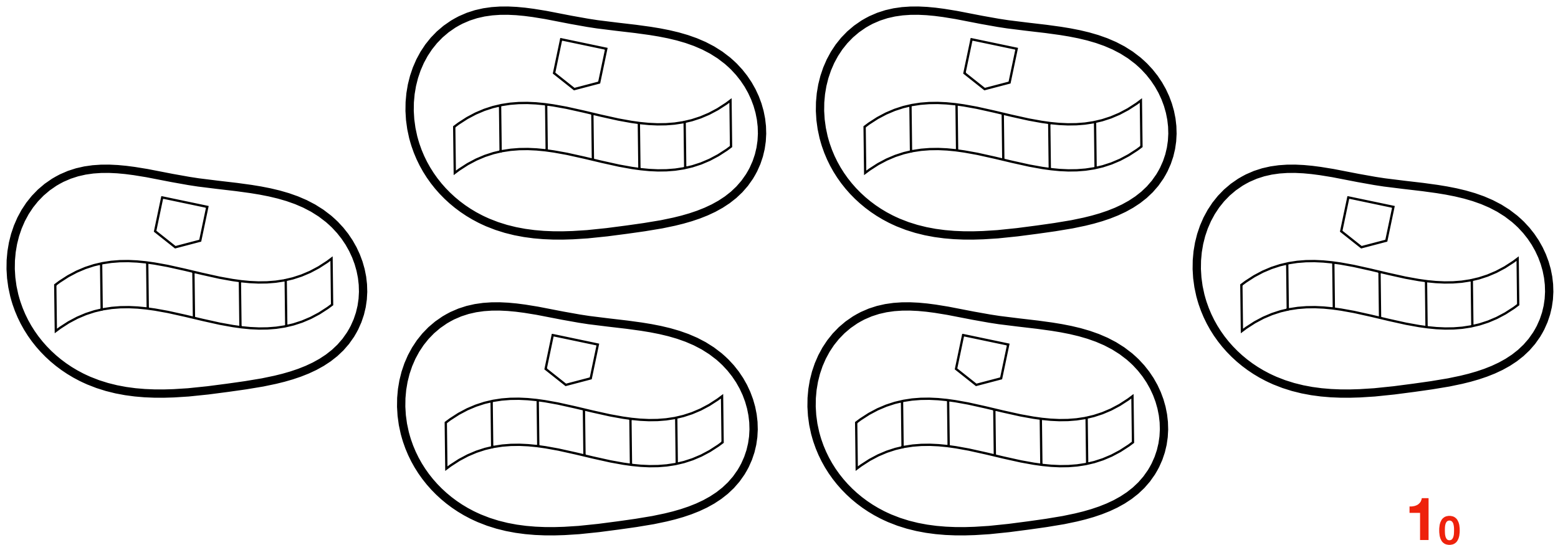
1₀

1₁

1₁

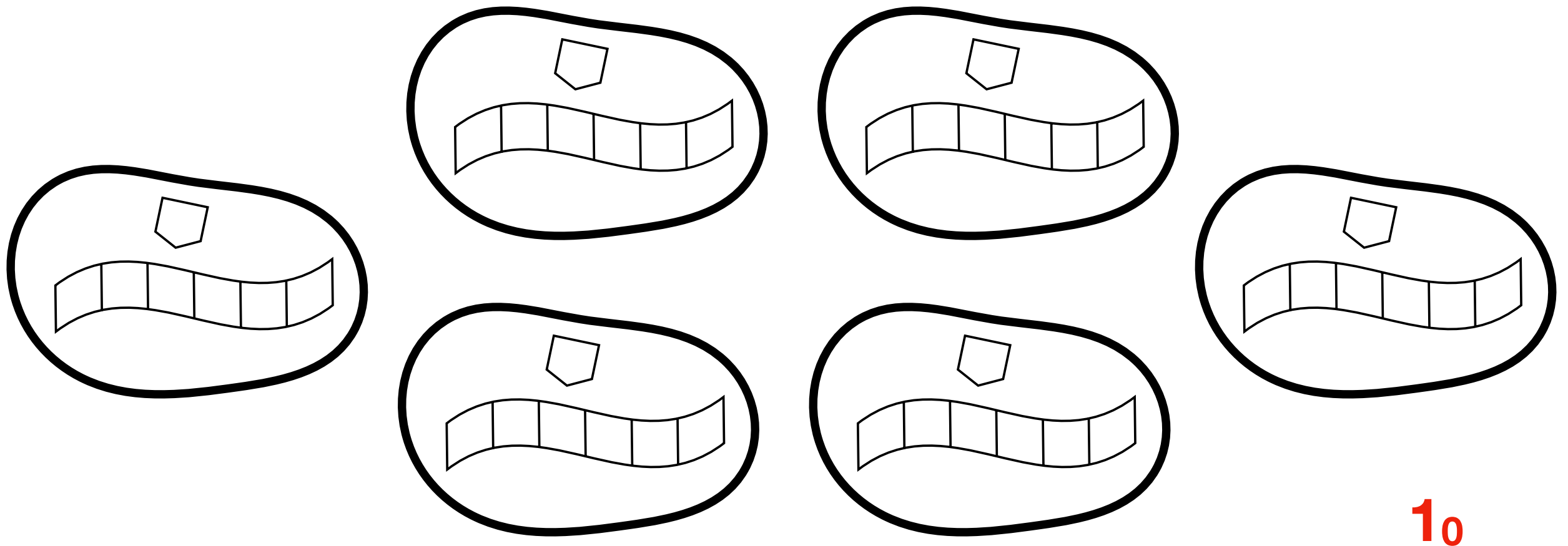


Converting unary to binary

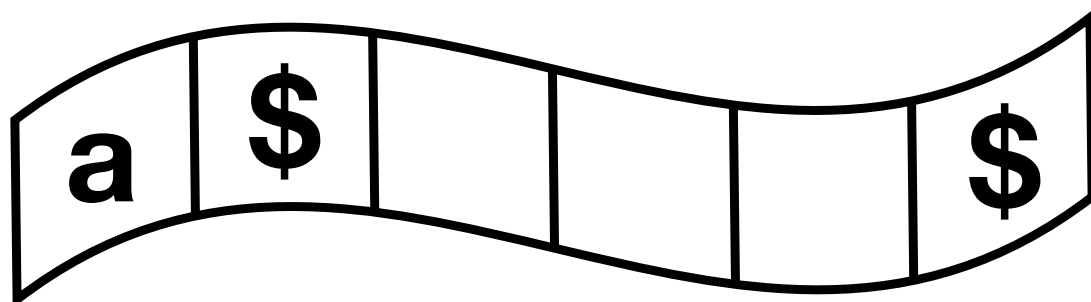


$[1_1 1_1 \rightarrow 1_2]$

Converting unary to binary

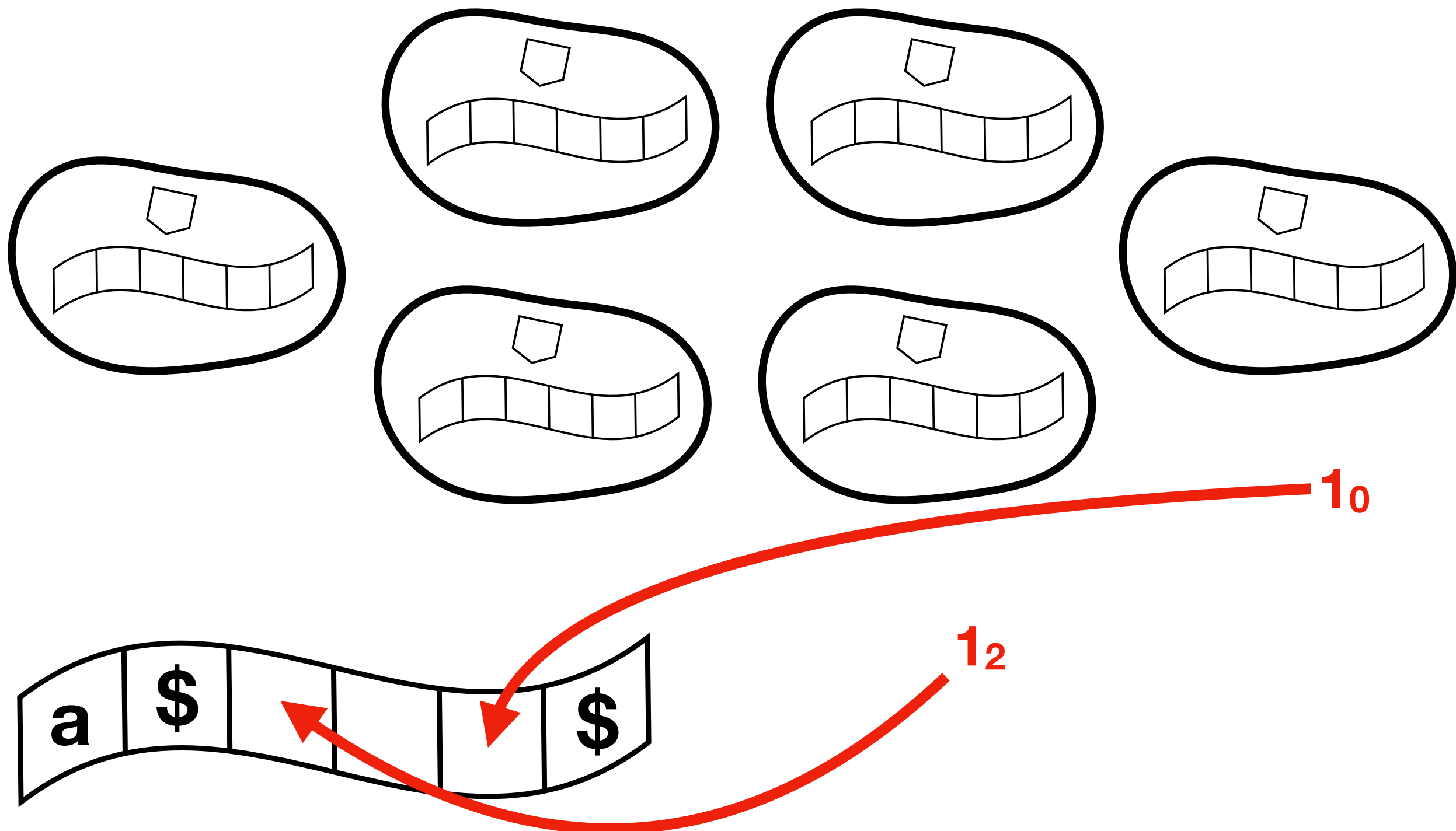


1₀

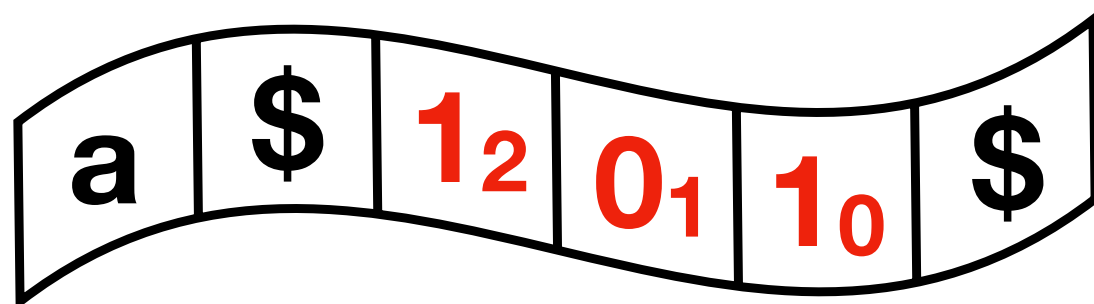
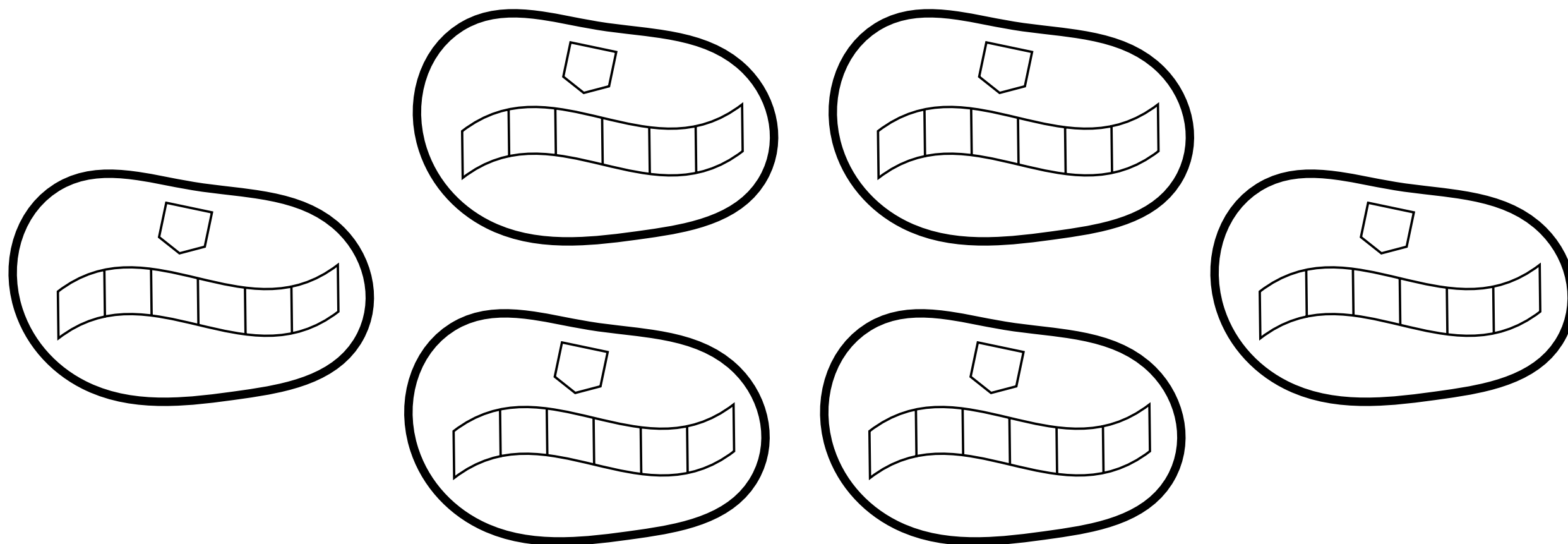


1₂

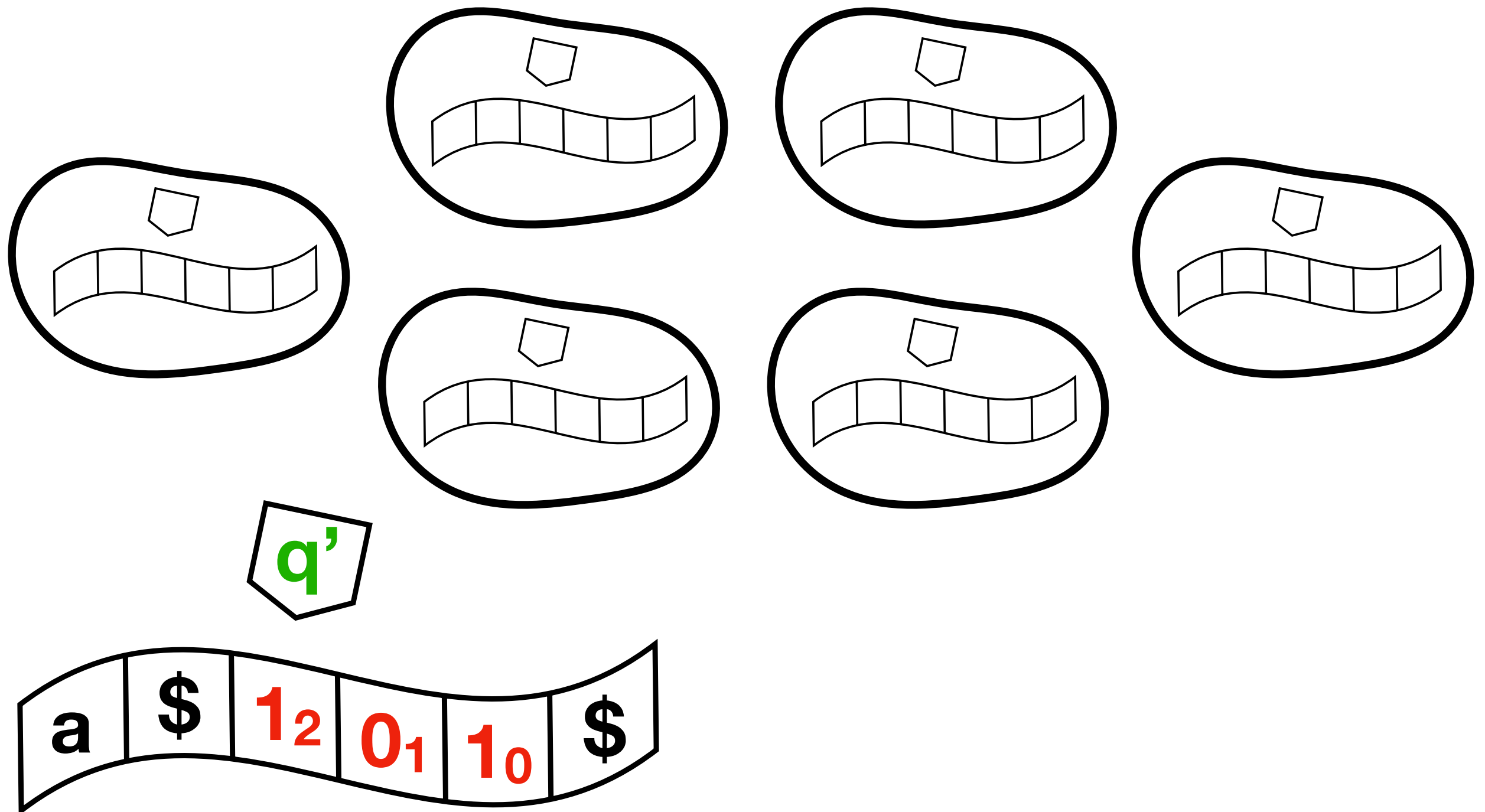
Answer on the tape



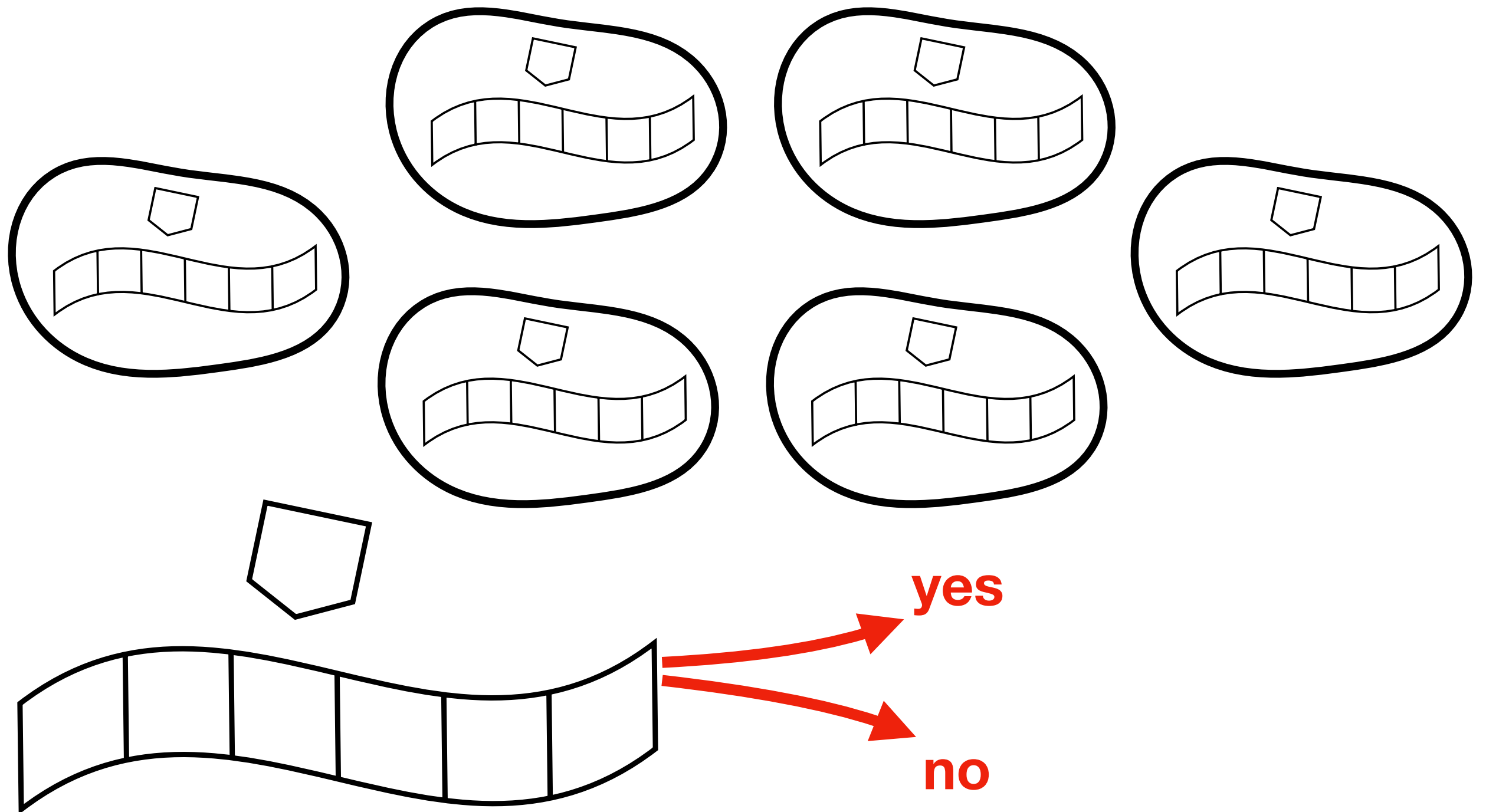
Answer on the tape



Resuming the simulation of the main TM



Final answer



**Simulating
(monodirectional, shallow)
membrane systems is in $P^{\#P}[1]$**

Counting the number of objects **a** sent out by a membrane at time **t** is in #P

- **for** $i := 0$ **to** t **do**
 - deterministically choose a maximal multisite of rules to apply inside the simulated membrane
 - apply all the rules except membrane division in deterministic polynomial time (“Milano Theorem”)
 - if applying membrane division, nondeterministically choose whether to simulate the left or the right resulting membrane
- if an object **a** was sent out in the last step then **accept**, otherwise **reject**

Lemma: $\mathbf{P\#P[1]} = \text{parallel } \mathbf{P\#P}$

- Trivially $\mathbf{P\#P[1]} \subseteq \text{parallel } \mathbf{P\#P}$
- A polynomial number of queries $f(x_1), \dots, f(x_m)$ can be replaced by a single query to

$$g(x_1 \$ x_2 \$ \dots \$ x_m) = \sum_{i=1}^n B^i \times f(x_i) \quad \Rightarrow \quad f(x_i) = \left\lfloor \frac{g(x_1 \$ x_2 \$ \dots \$ x_m)}{B^{i-1}} \right\rfloor$$

- The function **g** is also in $\mathbf{\#P}$ because this class is closed under sums and products

Simulating (shallow, omnidirectional) membrane systems in $P^{\#P}[1]$

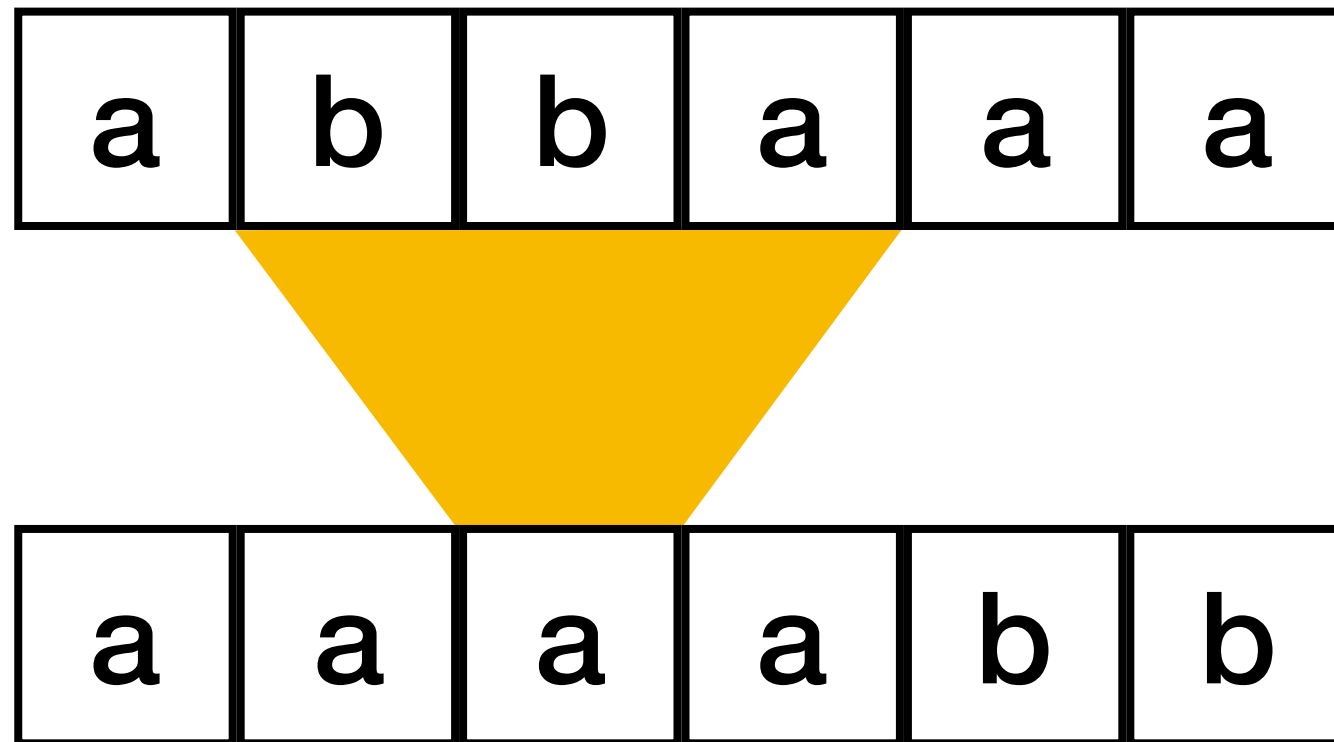
- for each membrane in the initial configuration, for each object type **a** and for each time step **t**, ask the oracle how many objects of type **a** are sent out by the membrane at time **t**
(note: **polynomial number of parallel queries!**)
- **while** the system has not produced the answer object **do**
 - simulate one step of the external environment deterministically (Milano Theorem)
 - **add the objects sent out from the membranes** (according to the queries asked) to the environment
- **accept** or **reject** according to the answer of the system simulated

Computational complexity of membrane systems

- No membranes (only environment) $\rightarrow \mathbf{P}$
- Shallow, monodirectional $\rightarrow \mathbf{P\#P[1]} = \text{parallel } \mathbf{P\#P}$
- Shallow, bidirectional $\rightarrow \mathbf{P\#P}$
- Constant depth k , bidirectional $\rightarrow \mathbf{PC_kP}$
where $\mathbf{C_0P} = \mathbf{P}$, $\mathbf{C_1P} = \mathbf{PP}$, $\mathbf{C_2P} = \mathbf{PP^{PP}}$, $\mathbf{C_kP} = \mathbf{PP^{C_{k-1}P}}$
is the **counting hierarchy** [work in progress]
- Unbounded depth, bidirectional $\rightarrow \mathbf{PSPACE}$

Expanding cellular automata (XCA)

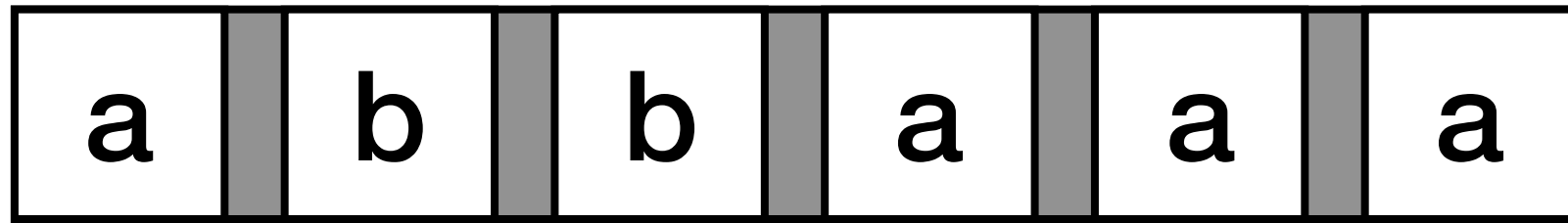
Expanding CA



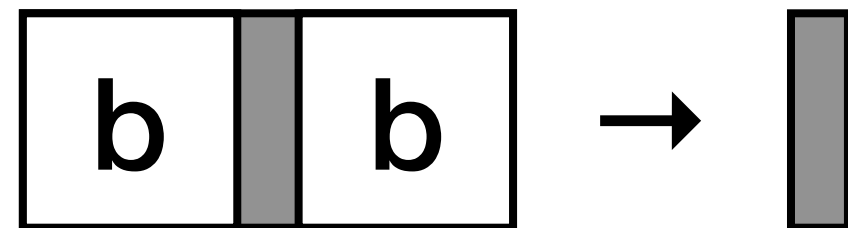
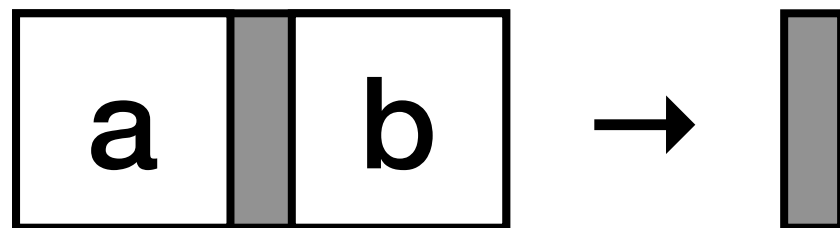
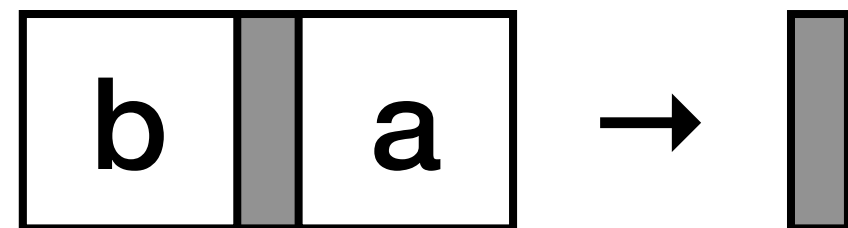
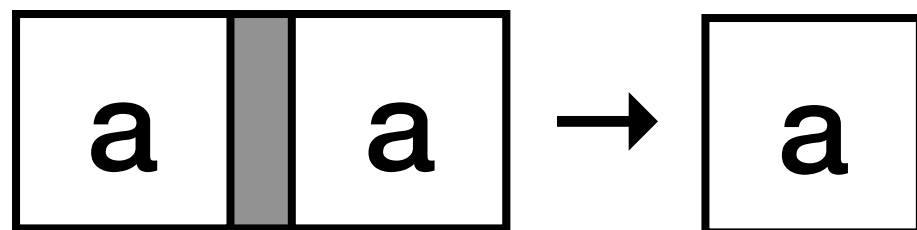
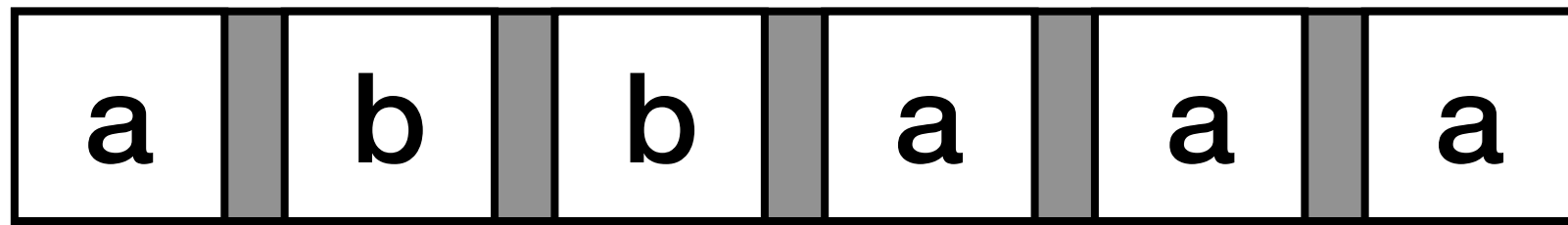
Expanding CA

a	b	b	a	a	a
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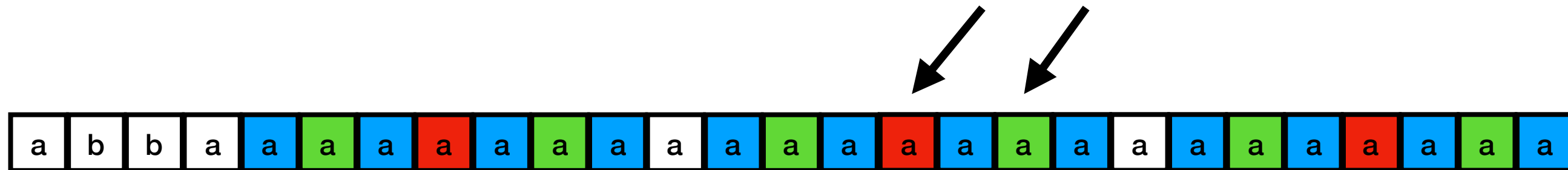
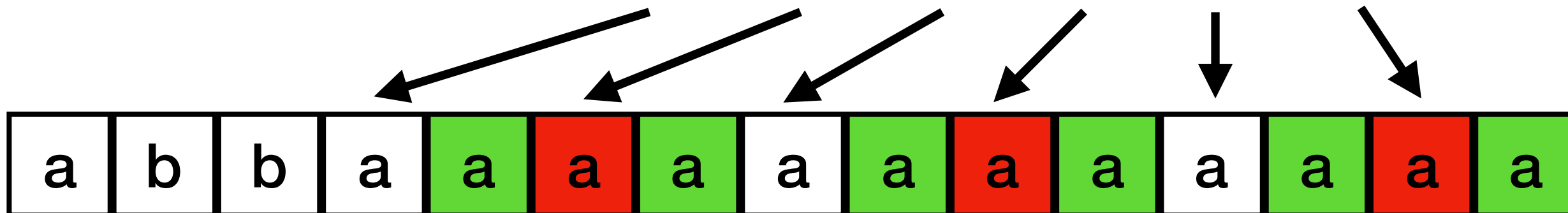
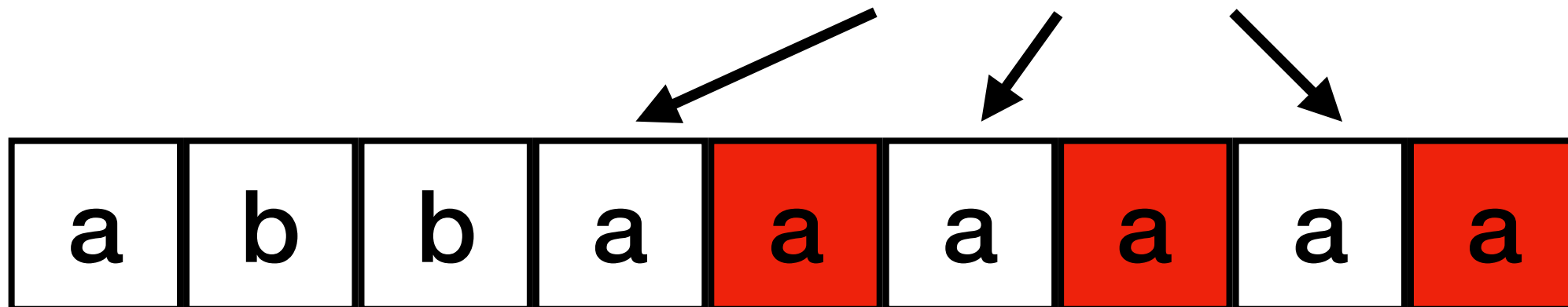
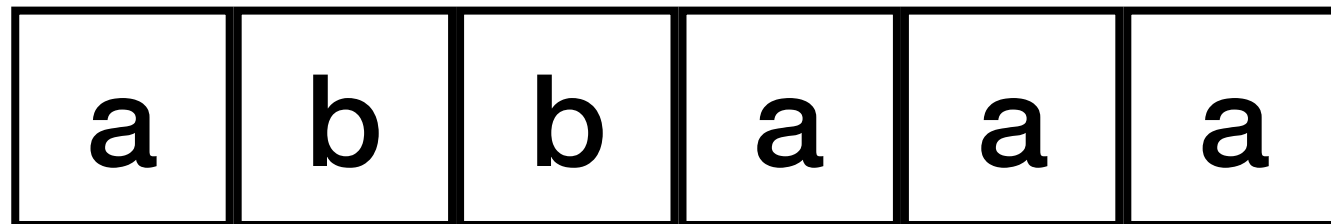
Expanding CA



Expanding CA



Expanding CA



Complexity results on XCA

- The class of problems solved in polynomial time by XCA is exactly the class of **problems truth-table reducible to NP**
- If **shrinking** (deleting cells) is also allowed, then the class becomes **PSPACE**

Conclusions and future work

Summary of results

- A lot of parallel computing models characterise **either P or PSPACE** when working in polynomial time
- Some variants of **membrane systems** characterise **more “exotic” complexity classes** with oracles, like **$P^{\#P[1]}$, $P^{\#P}$, P^{NP}**
- Expanding CA characterise the class of problems truth-table reducible to **NP**, which is somehow similar to oracle complexity classes

Conjectures and future work

- Find out **why** these models happen to characterise these exotic complexity classes
- Find out how the topology of the parallel computing units influences the efficiency:
 - Trees or stars for membrane systems
 - Linear or Euclidean grid for CA
 - Linear but expanding for XCA

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Thanks for your attention!

Merci de votre attention !