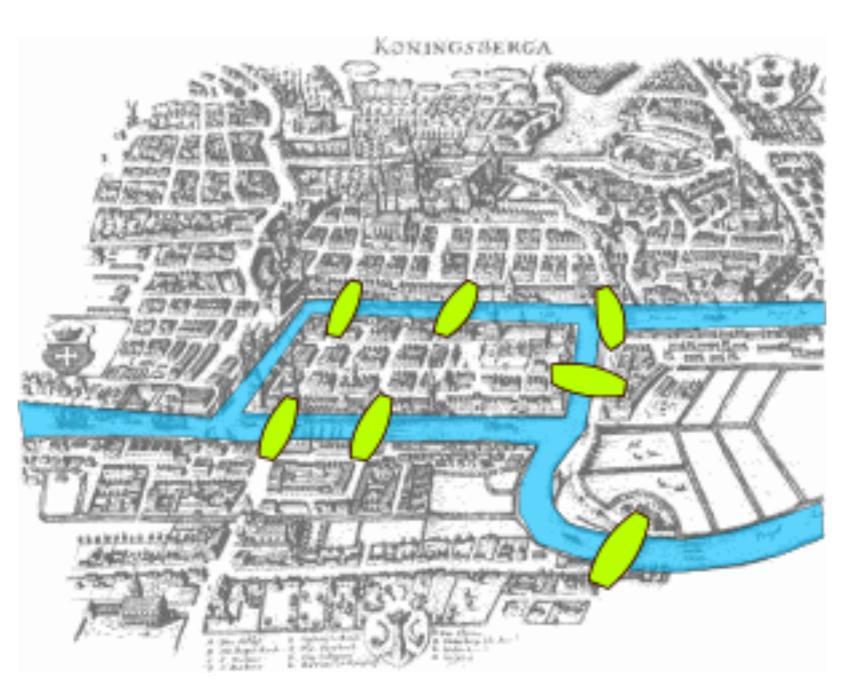
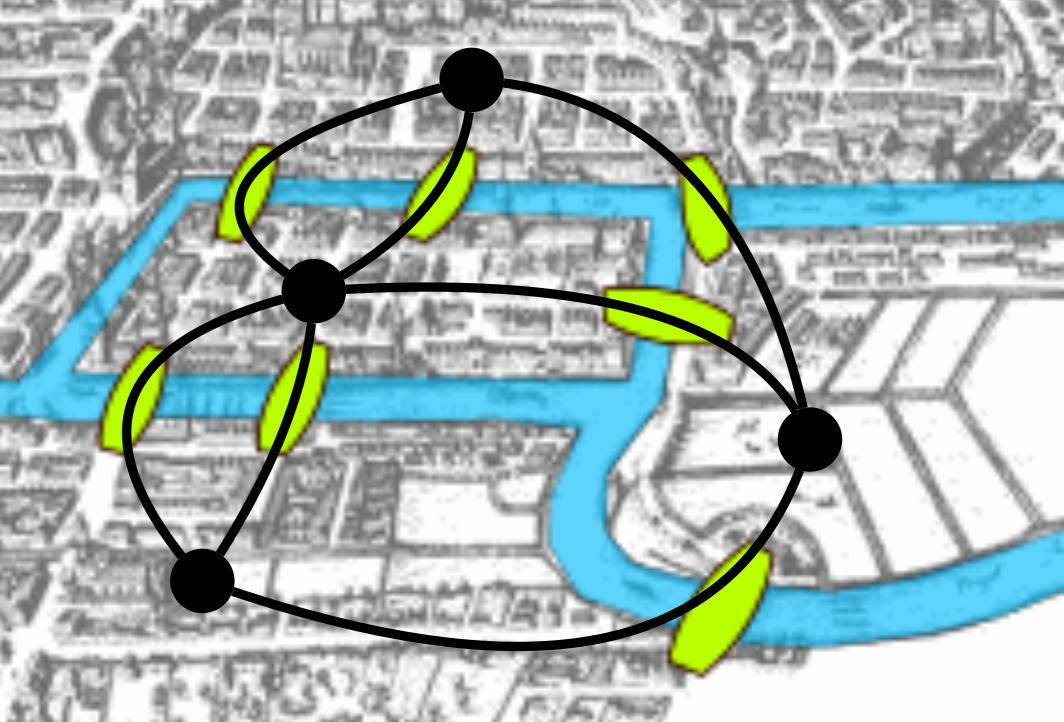
Introduction à l'informatique CM7

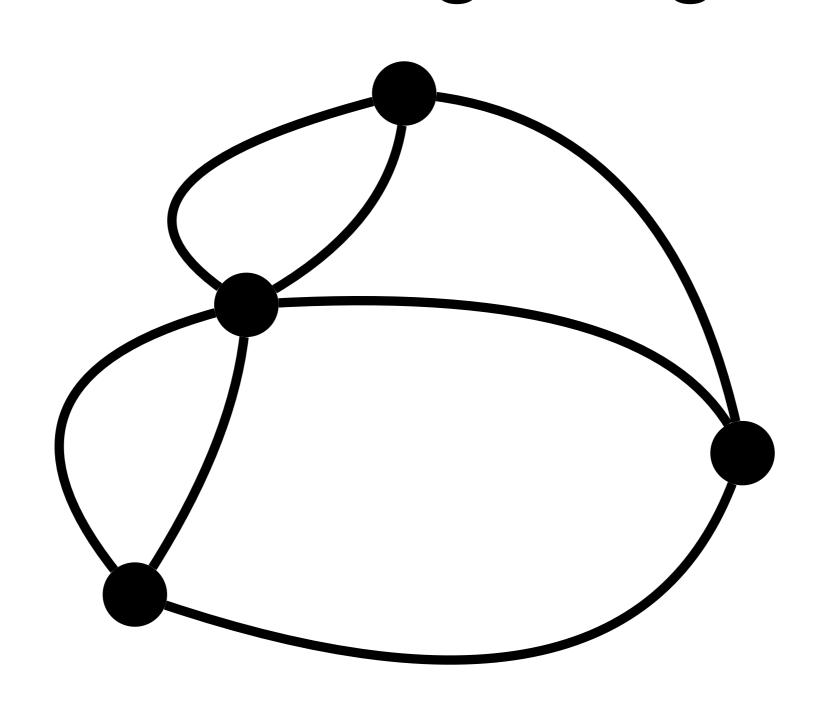
Antonio E. Porreca https://aeporreca.org/teaching

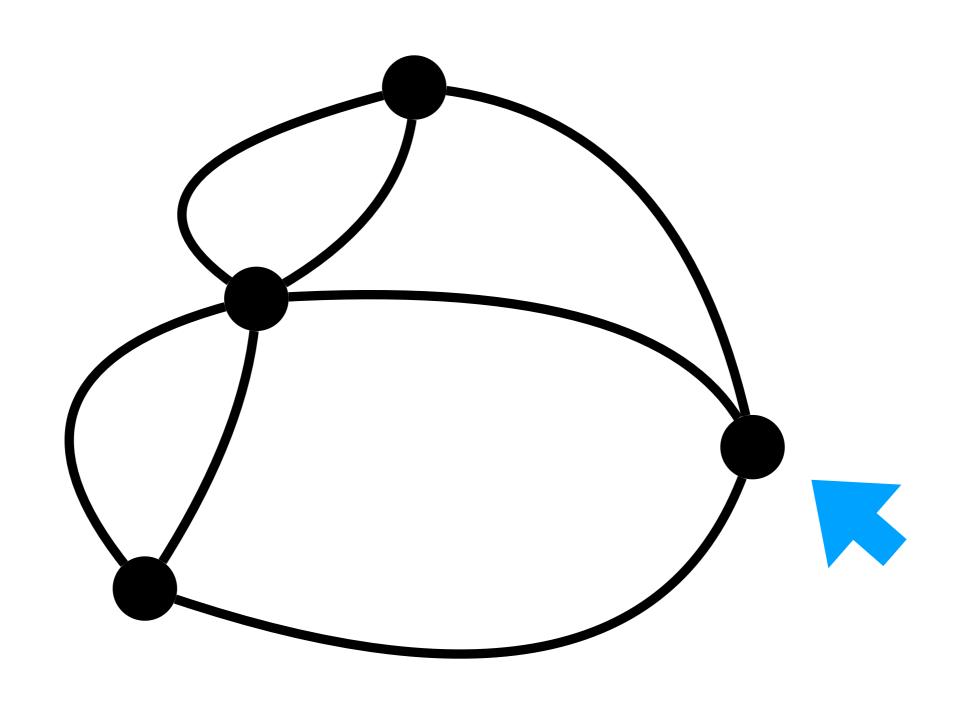


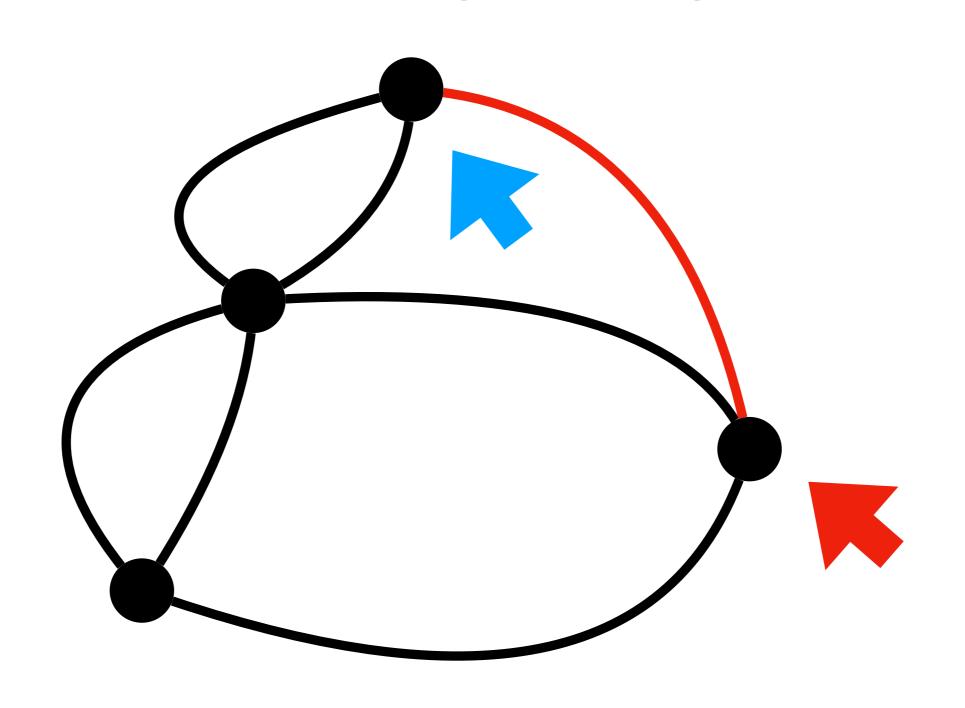
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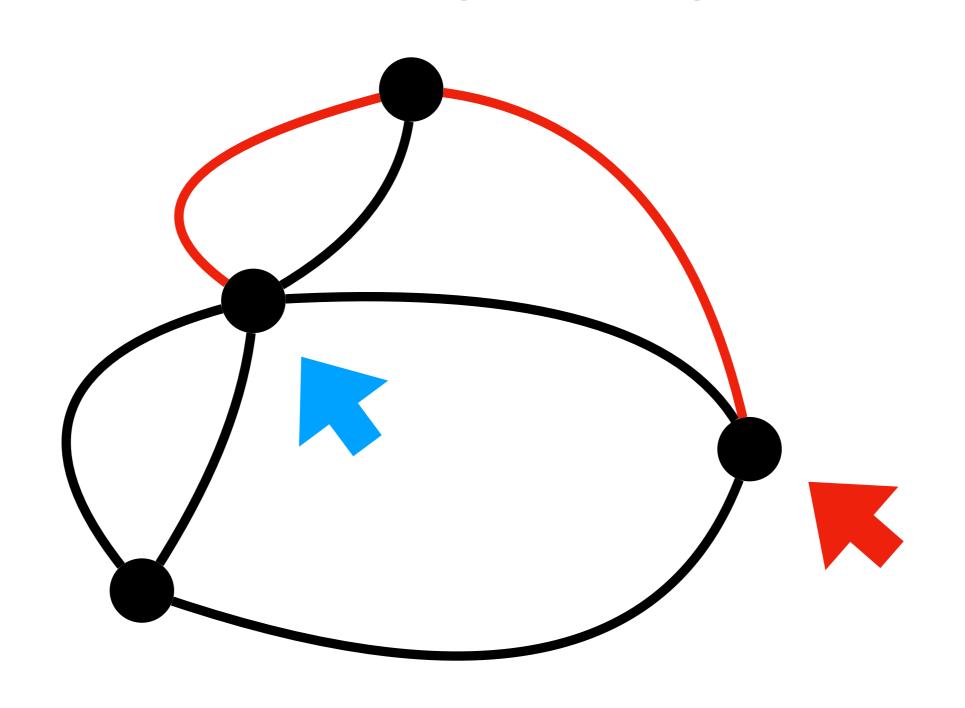
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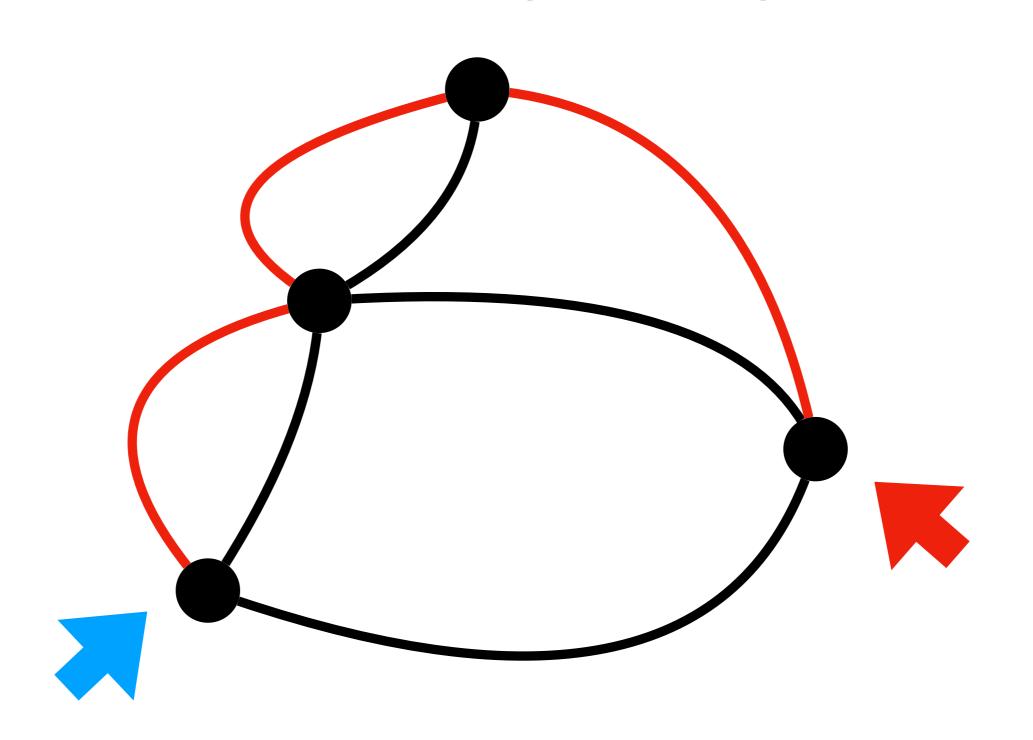


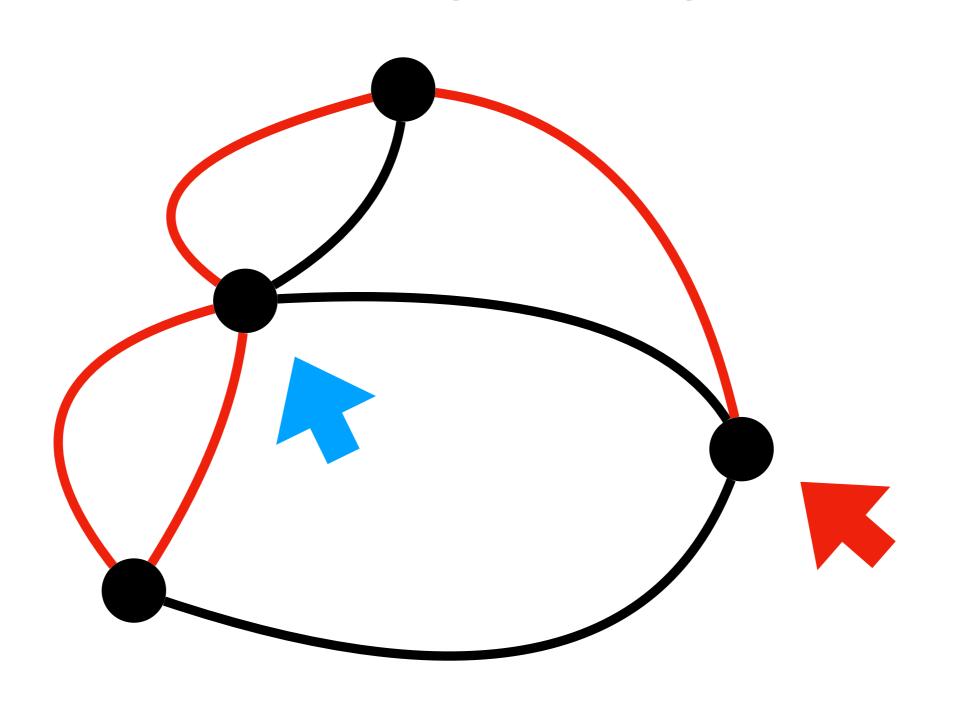


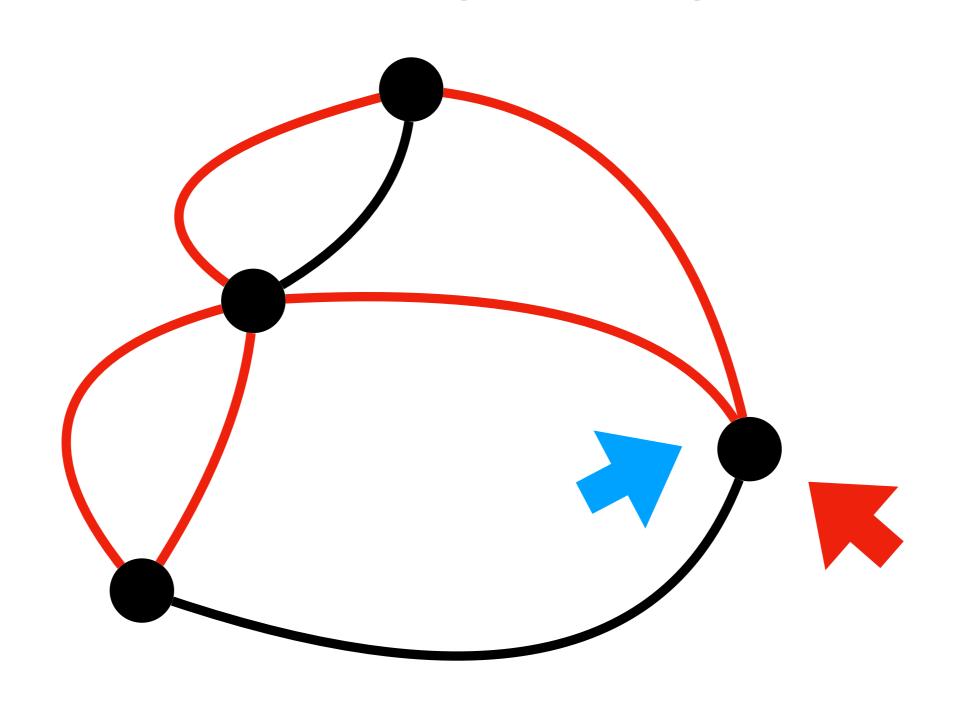


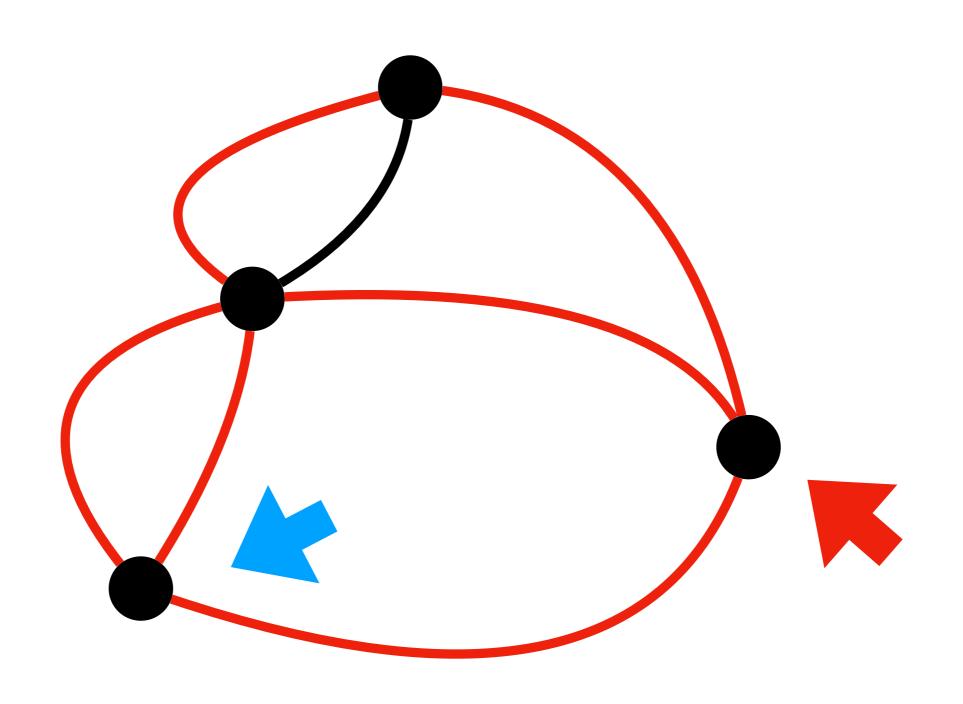


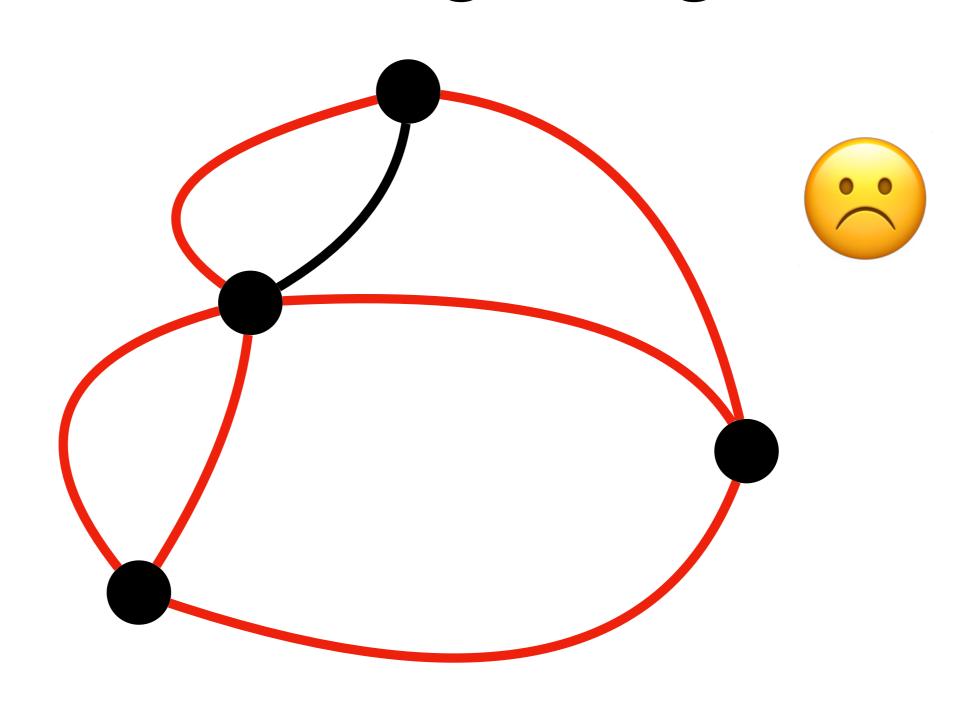












Il n'ya pas de solution!



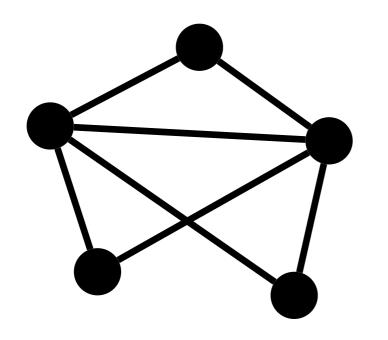
← Leonhard Euler

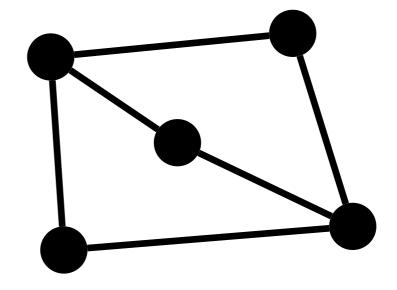
Théorème d'Euler (1736)

Un graphe non orienté connexe admet un cycle qui traverse chaque arête une et une seule fois (un « cycle eulérien ») si et seulement si chaque sommet a un nombre pair de voisins connectés par un arête.

Théorème d'Euler (1736)

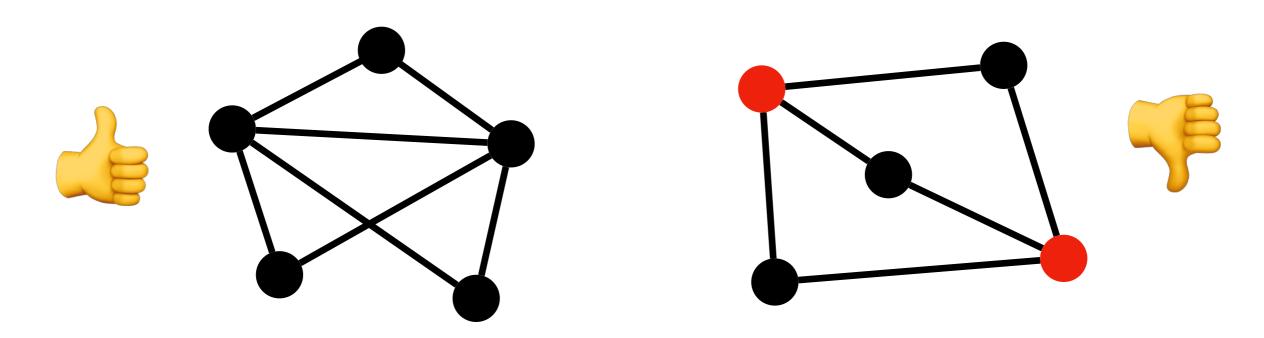
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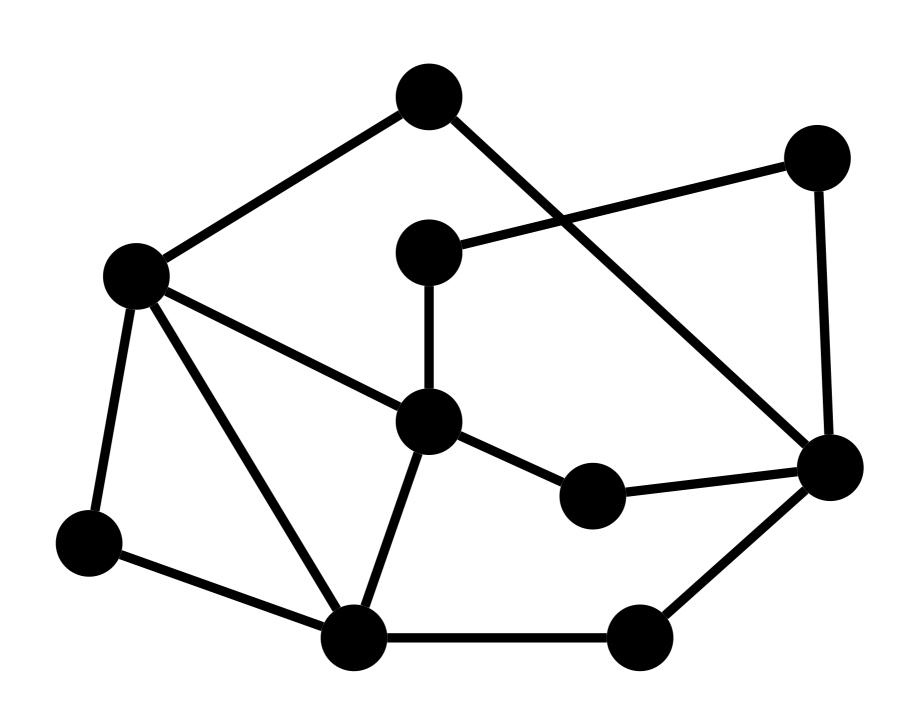


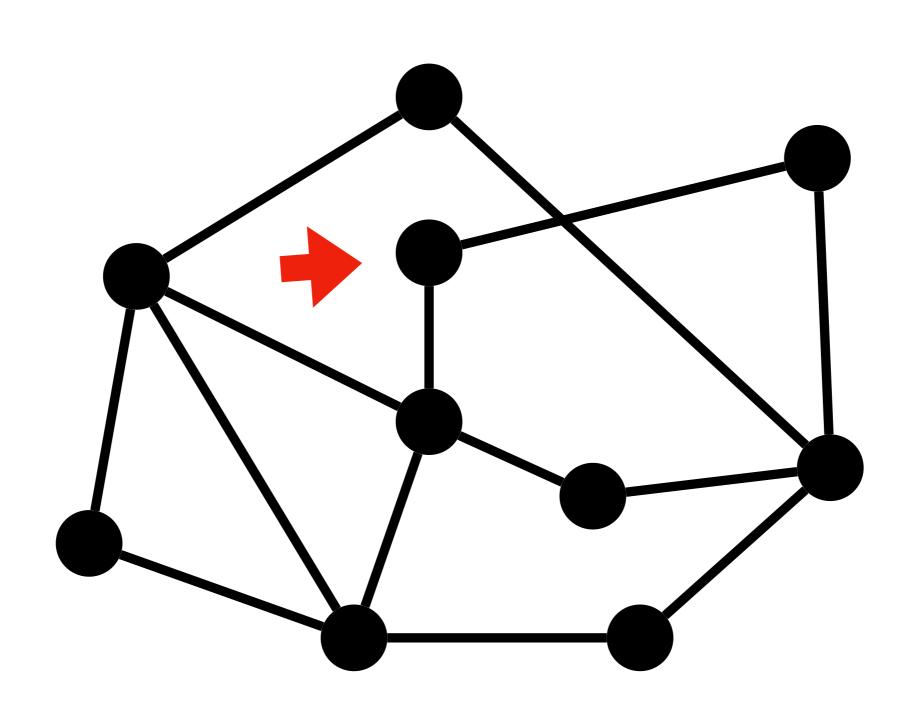


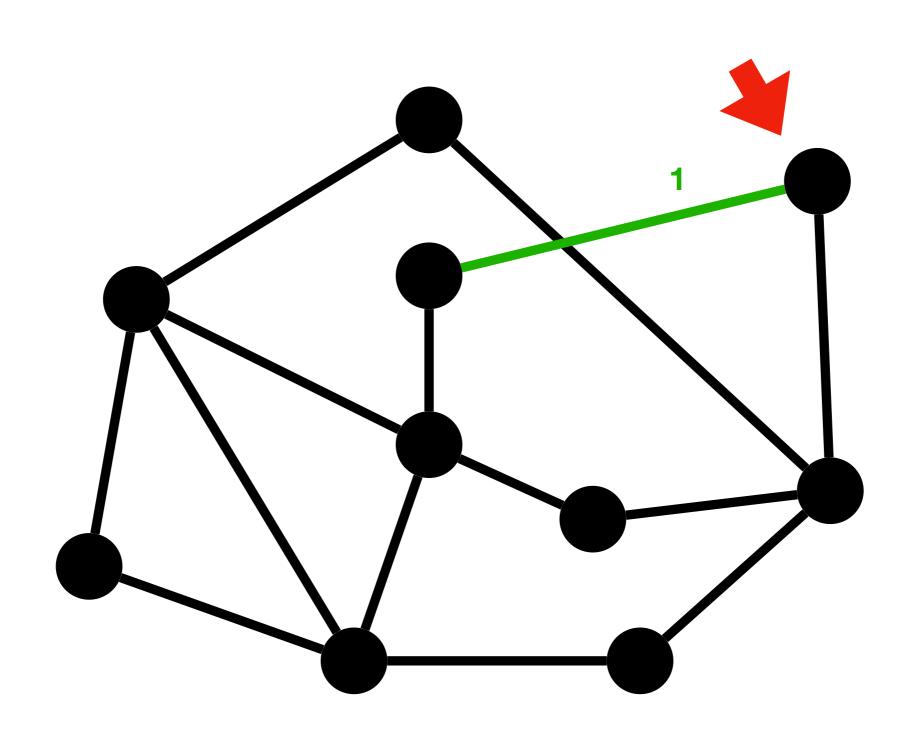
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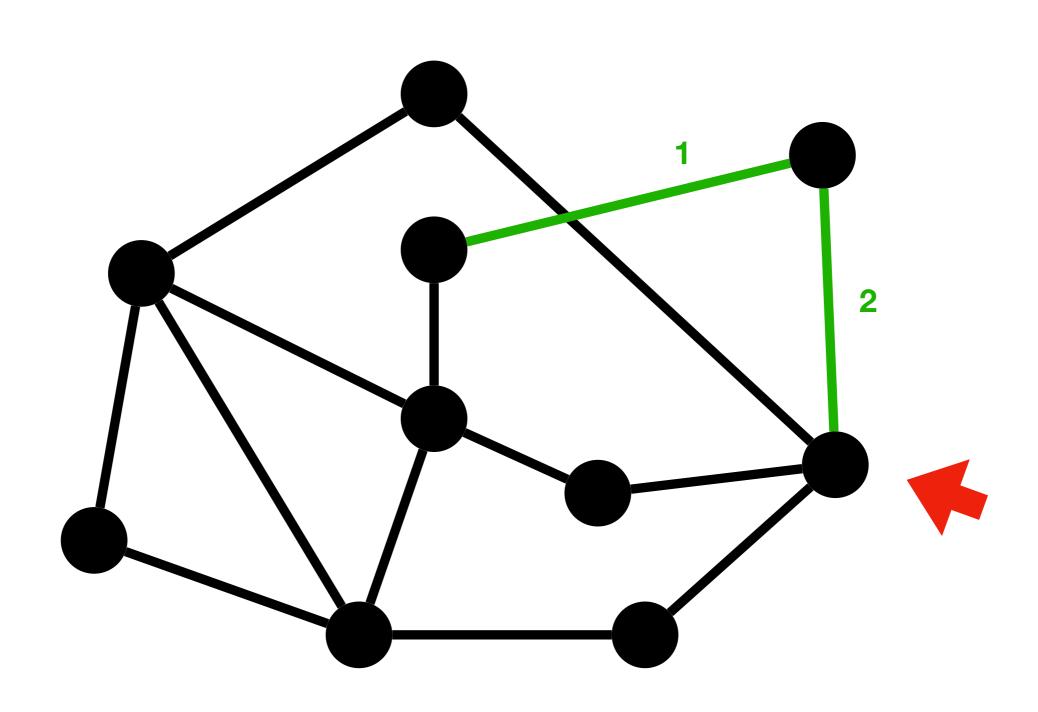
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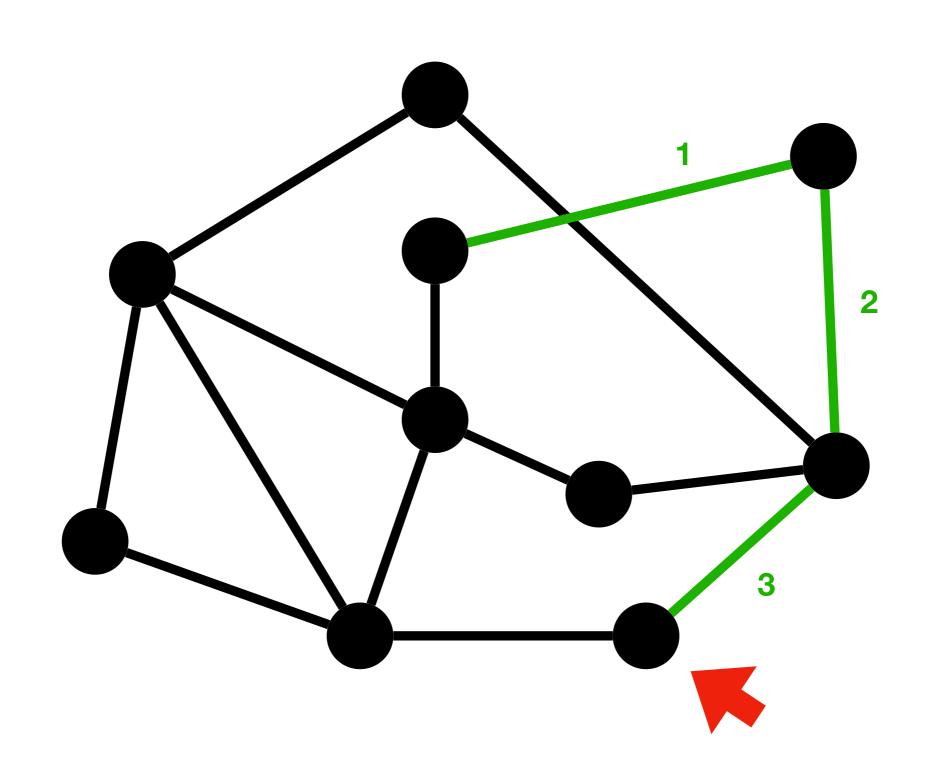


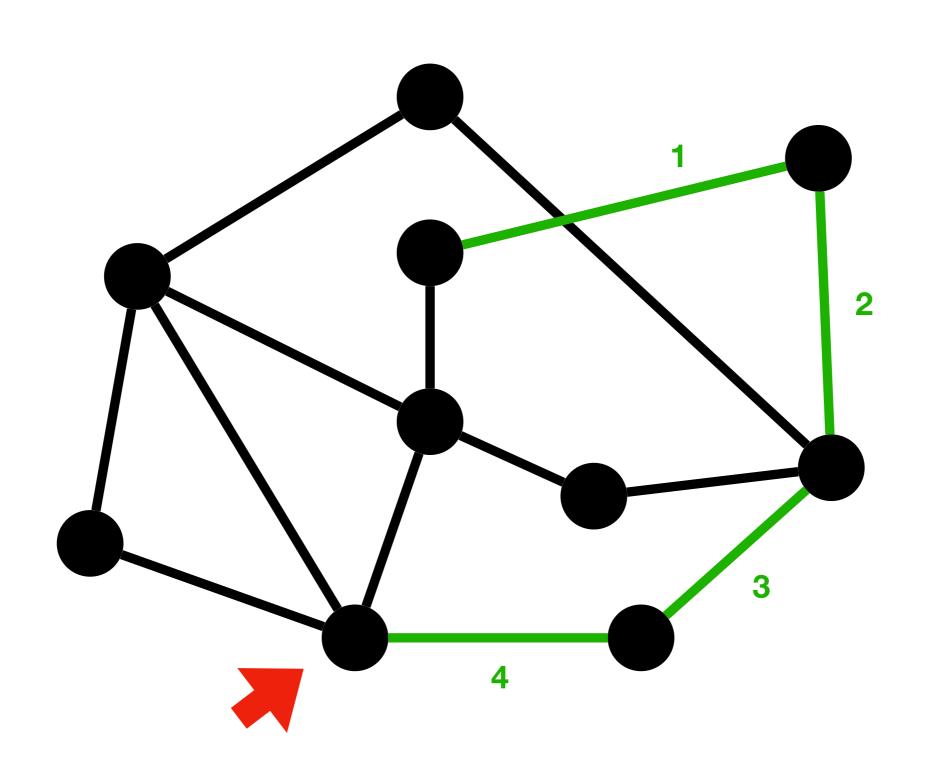


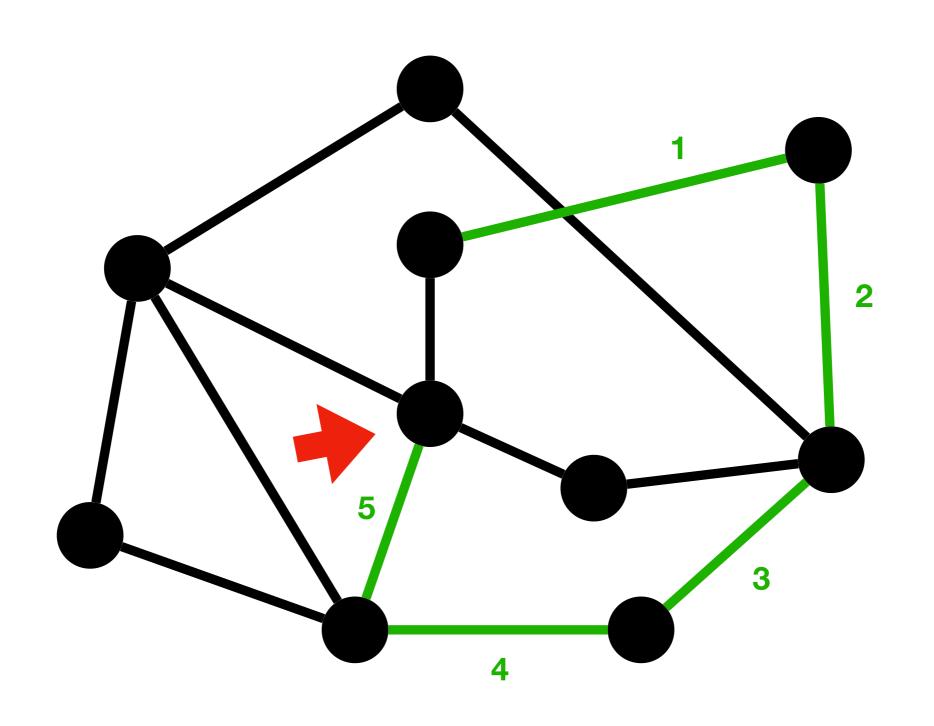


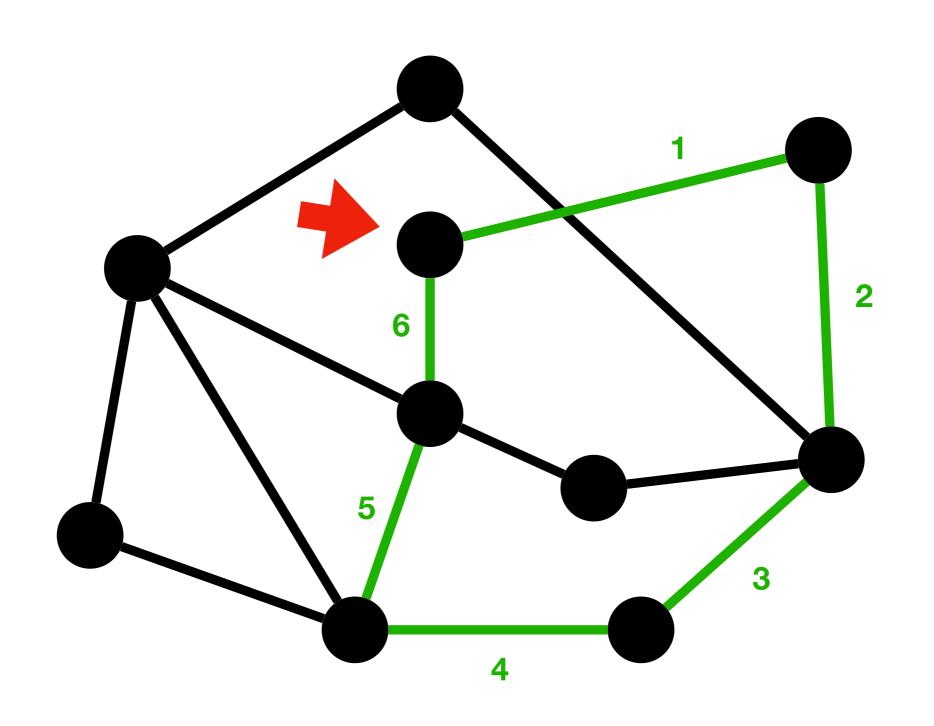


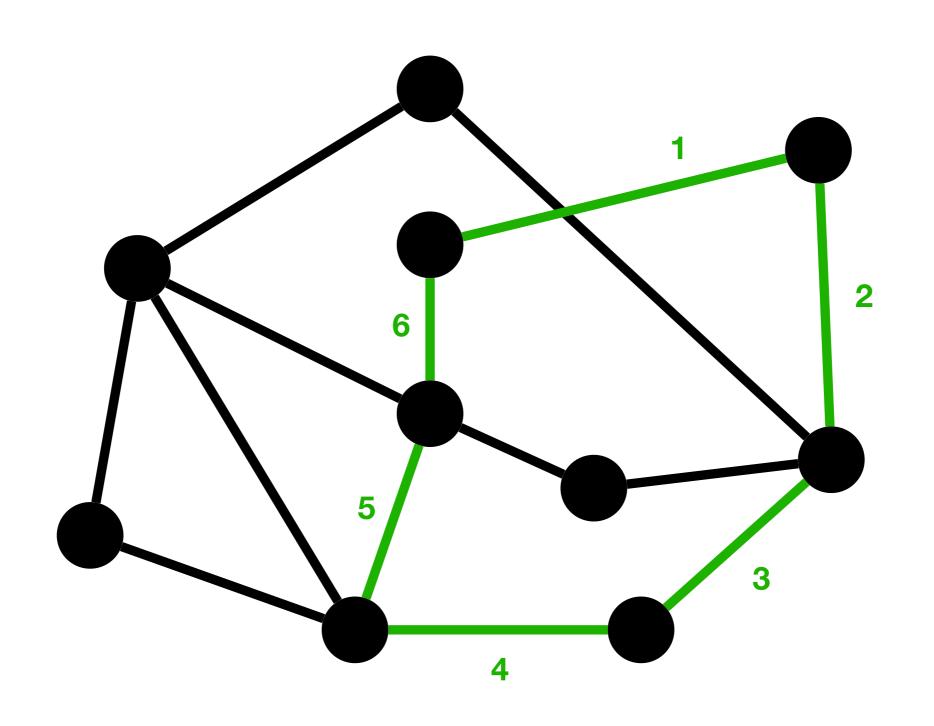


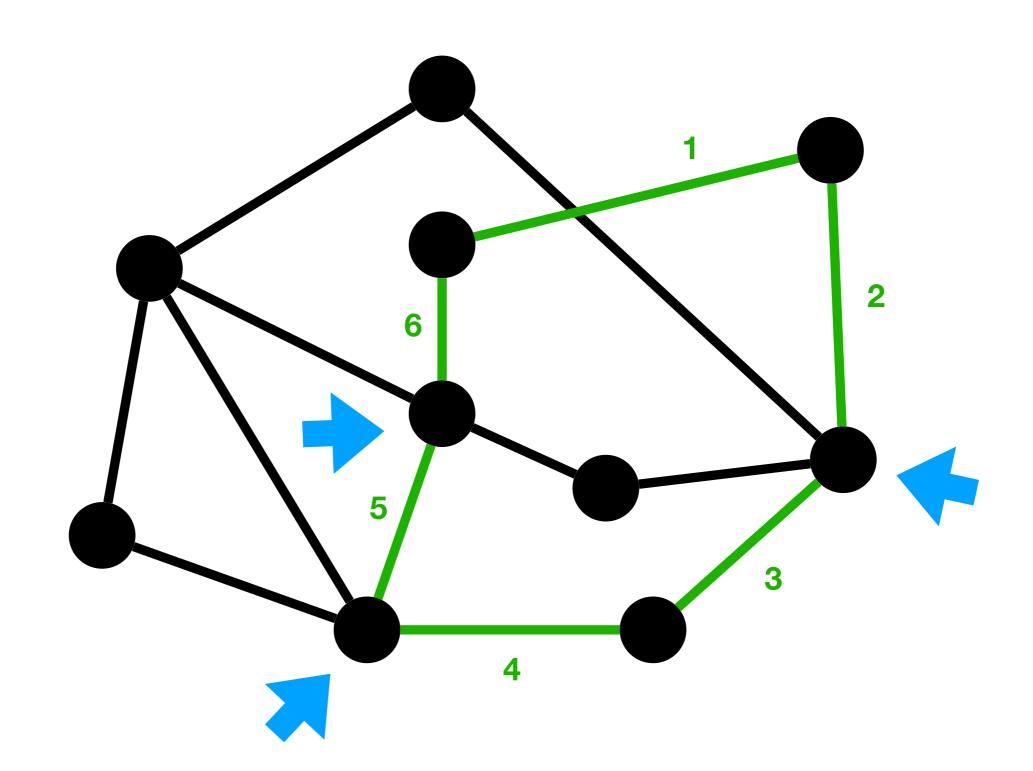


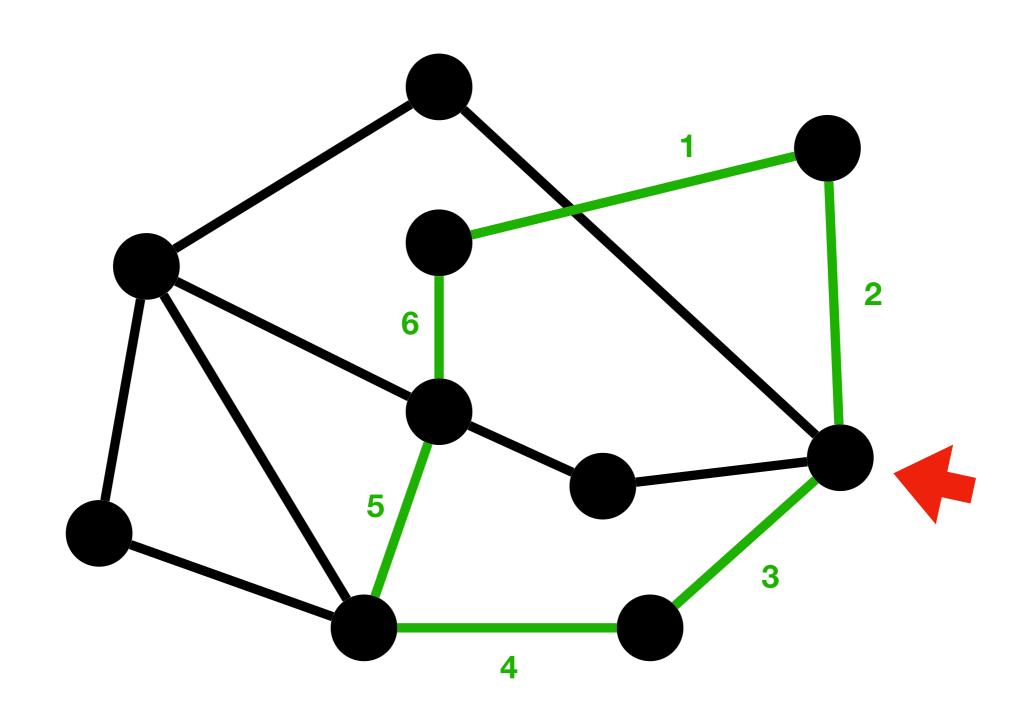


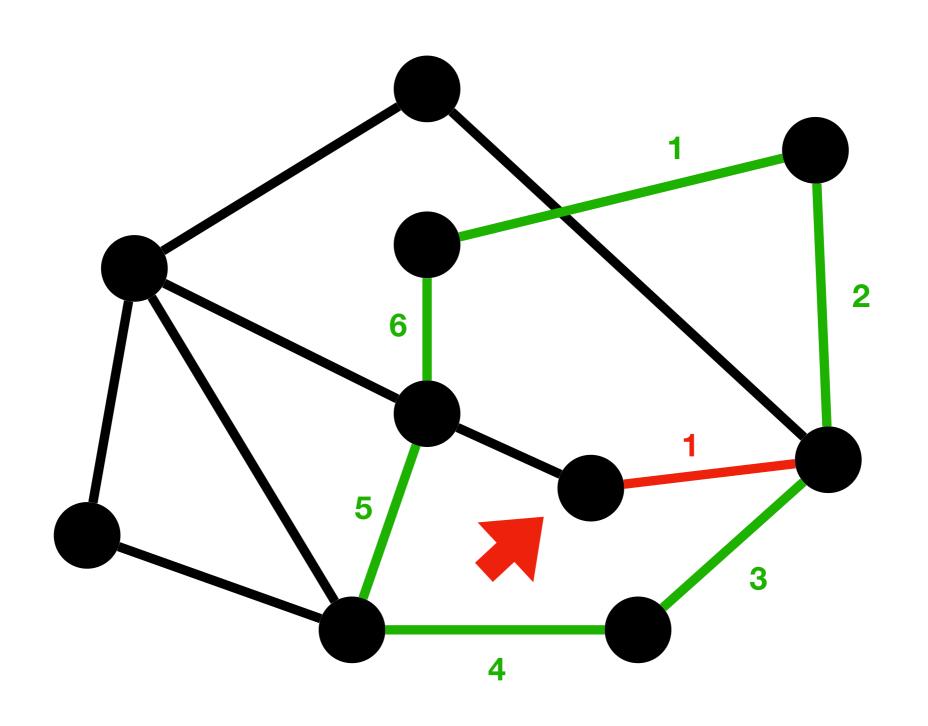


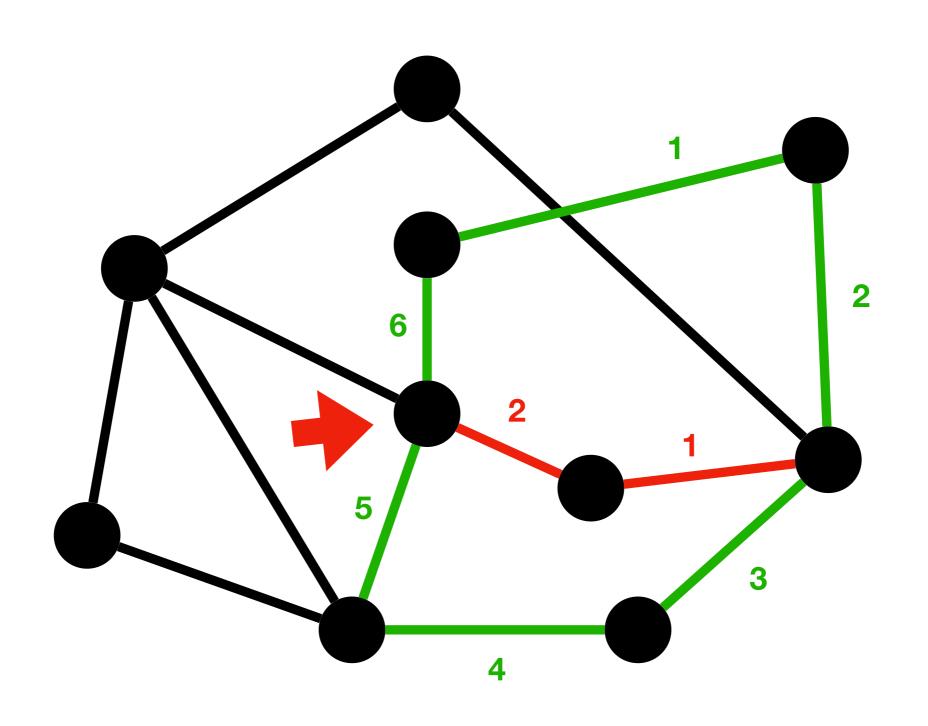


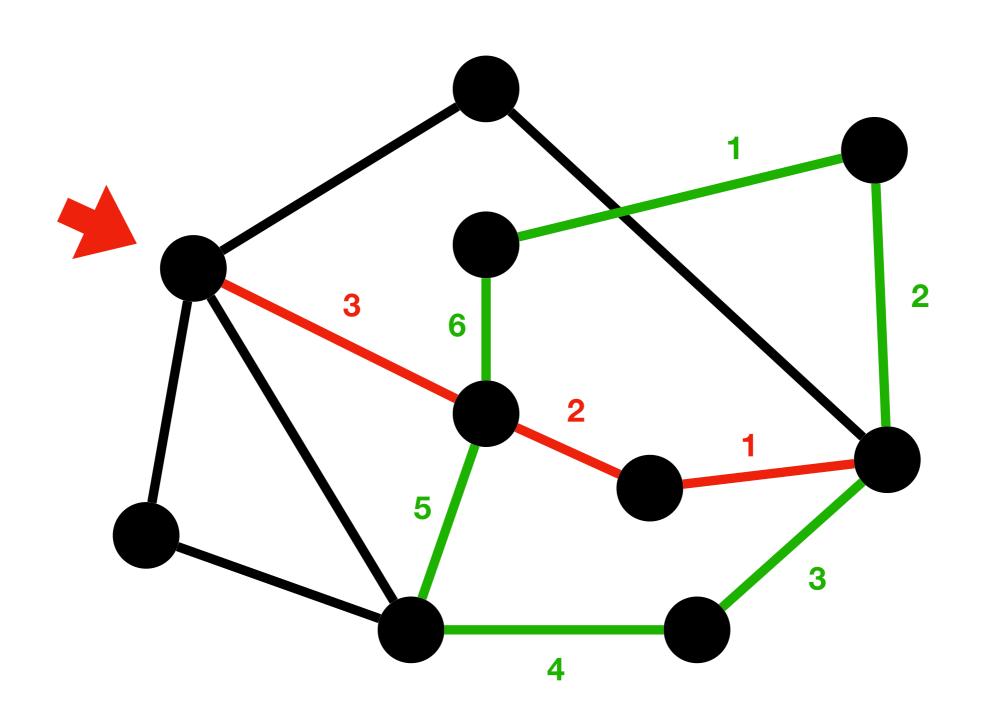


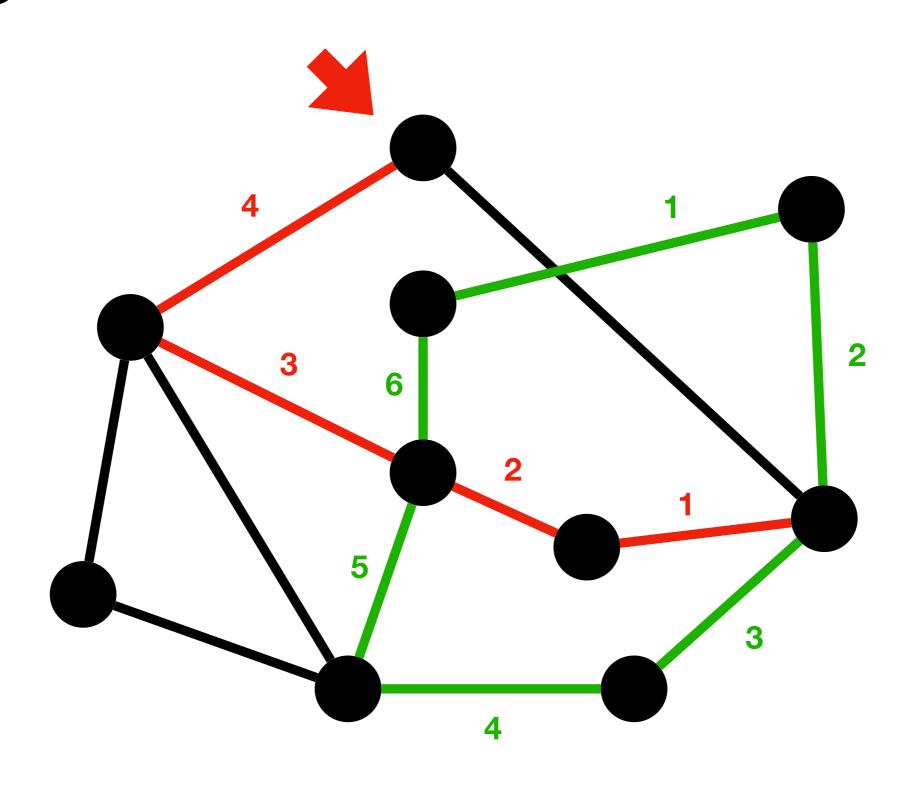


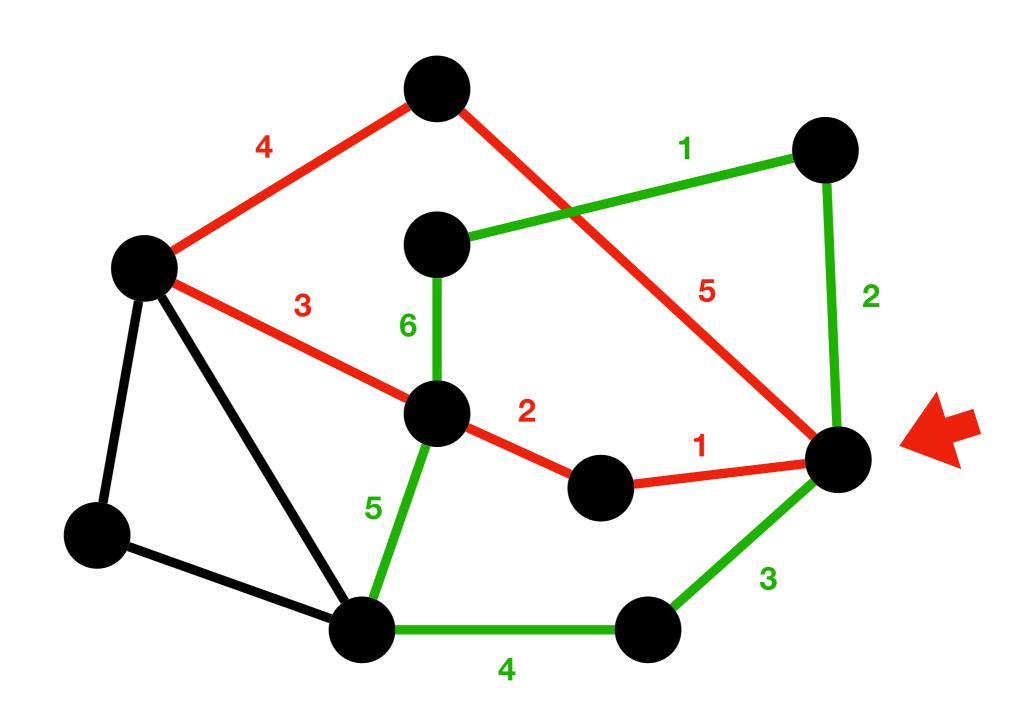


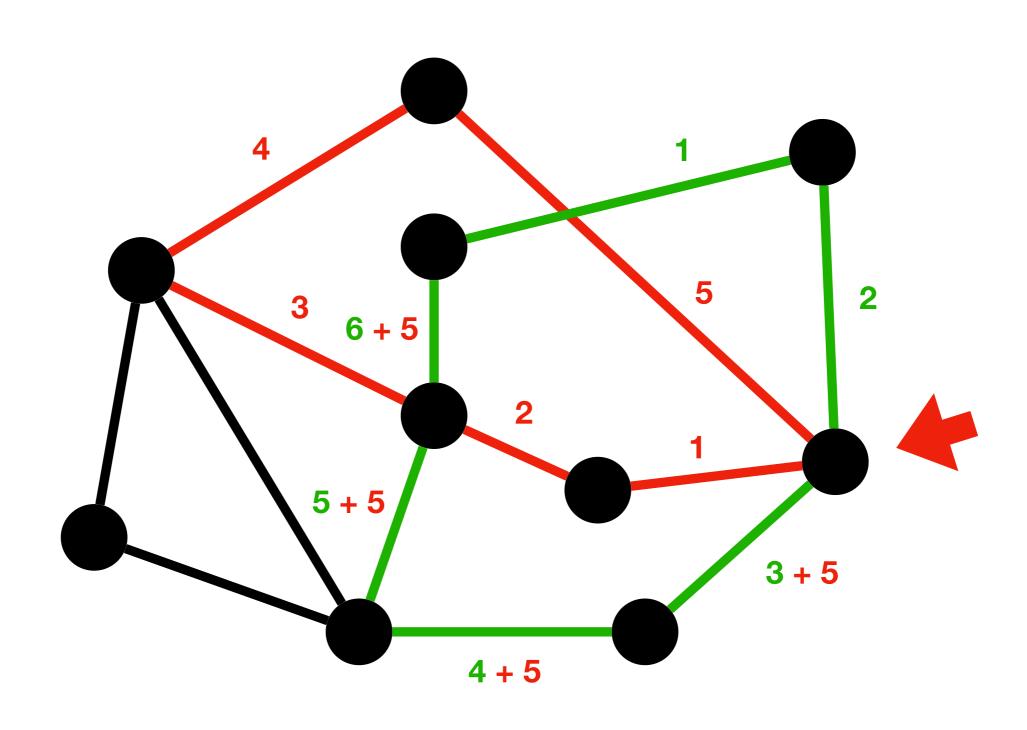


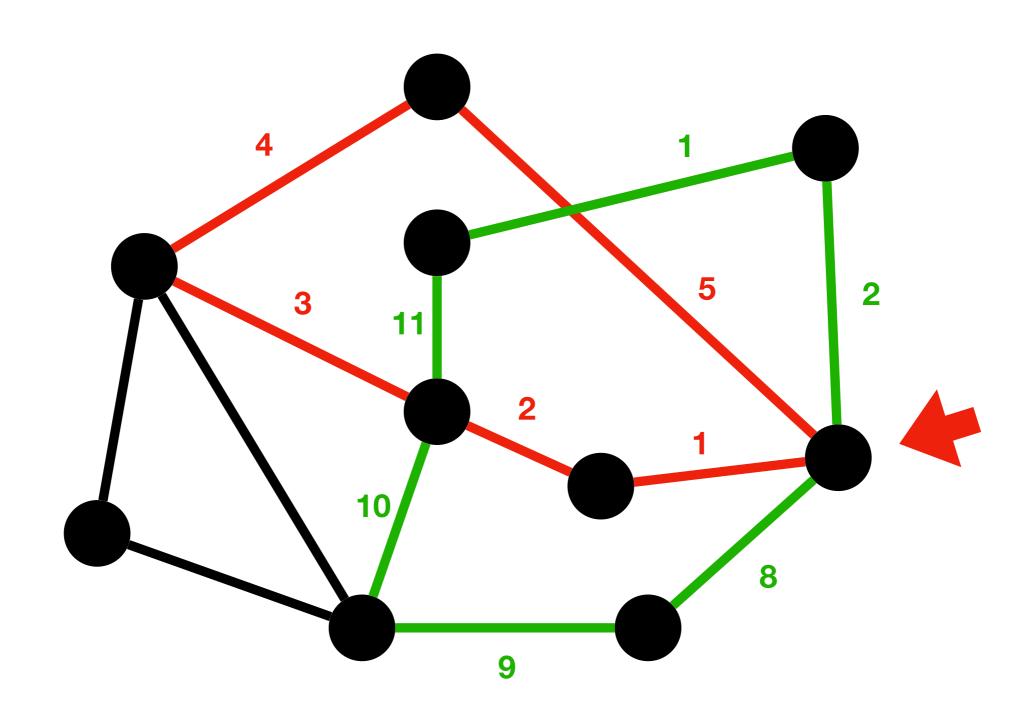


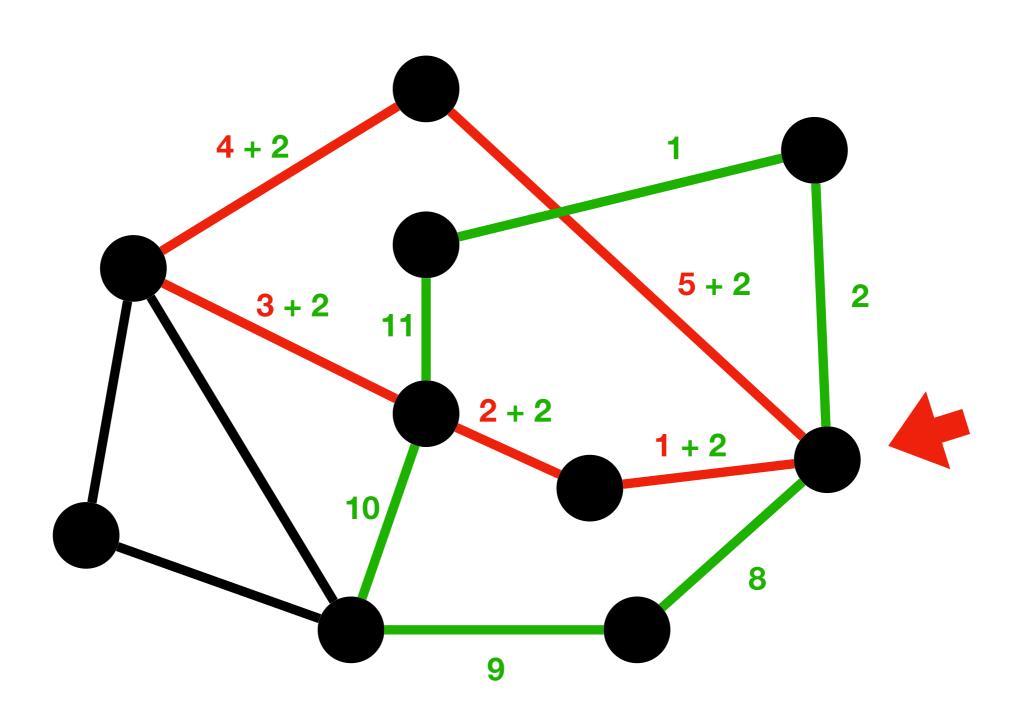


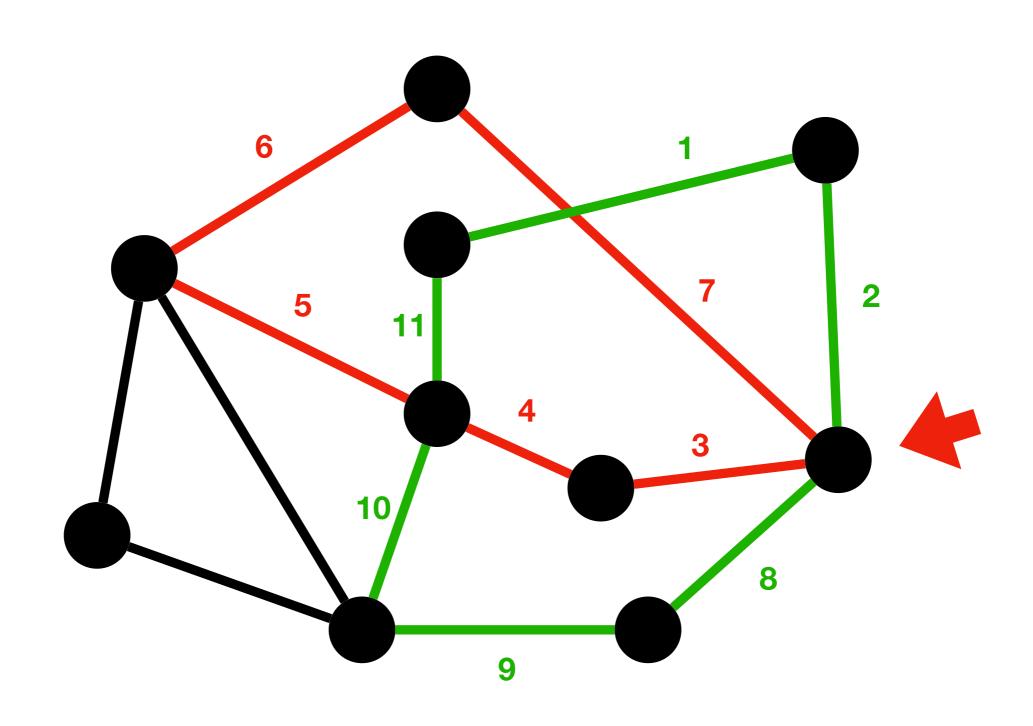


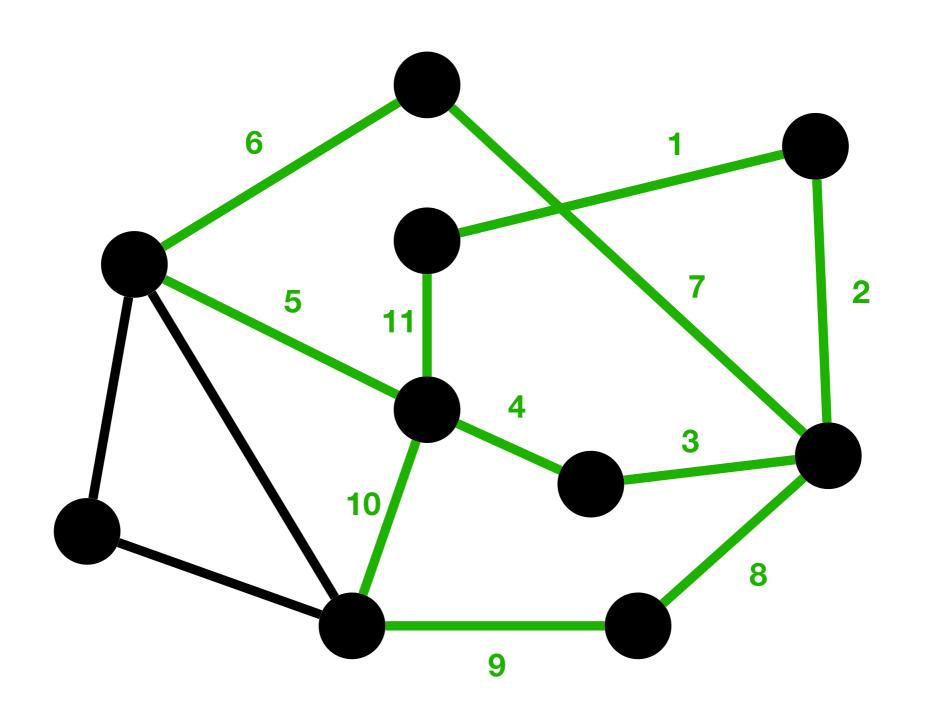


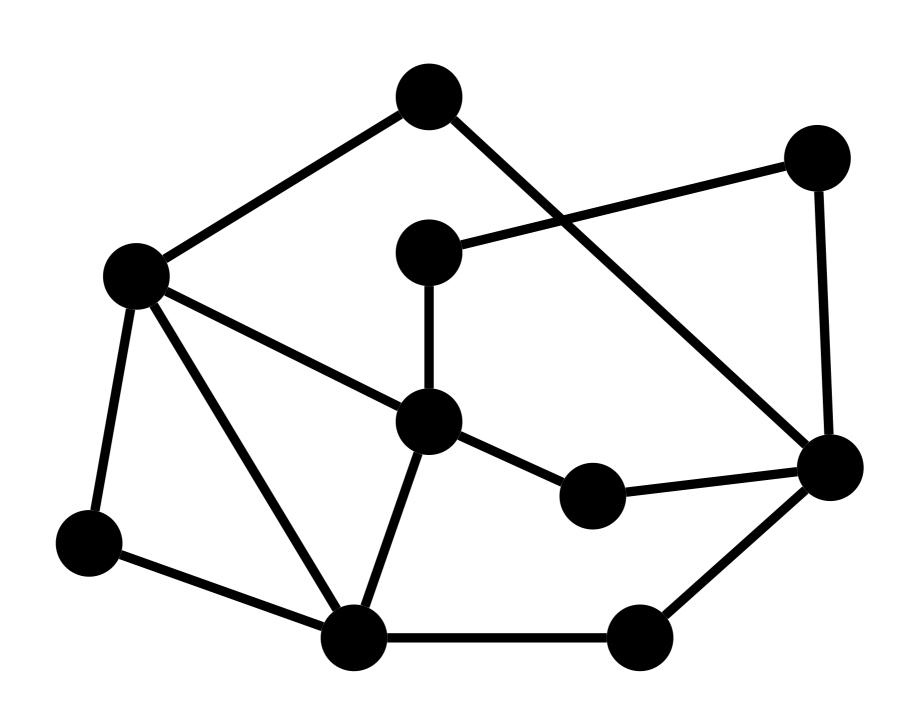


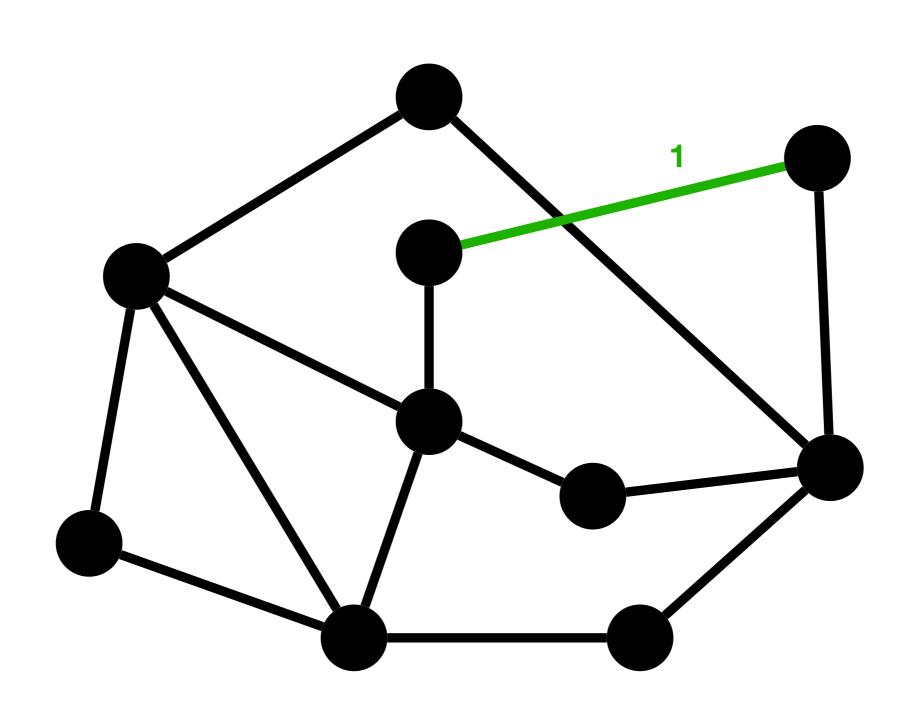


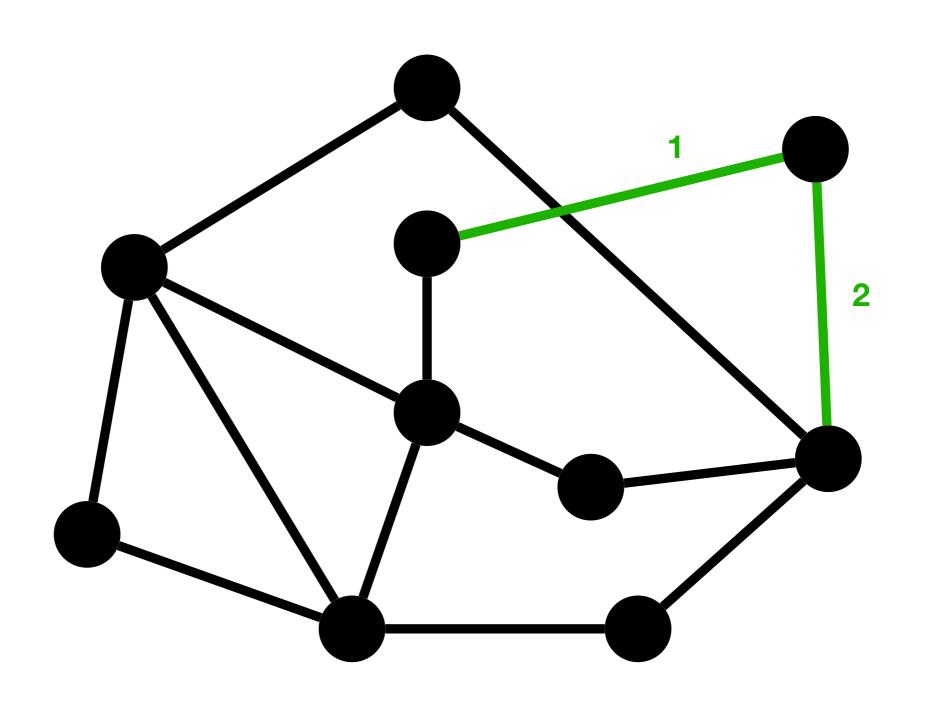


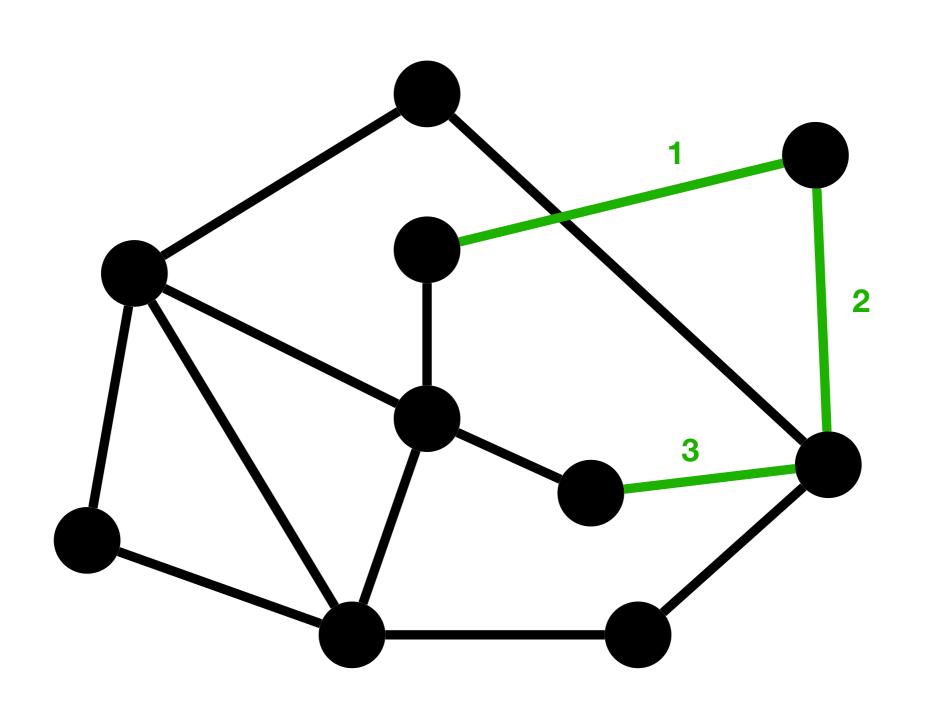


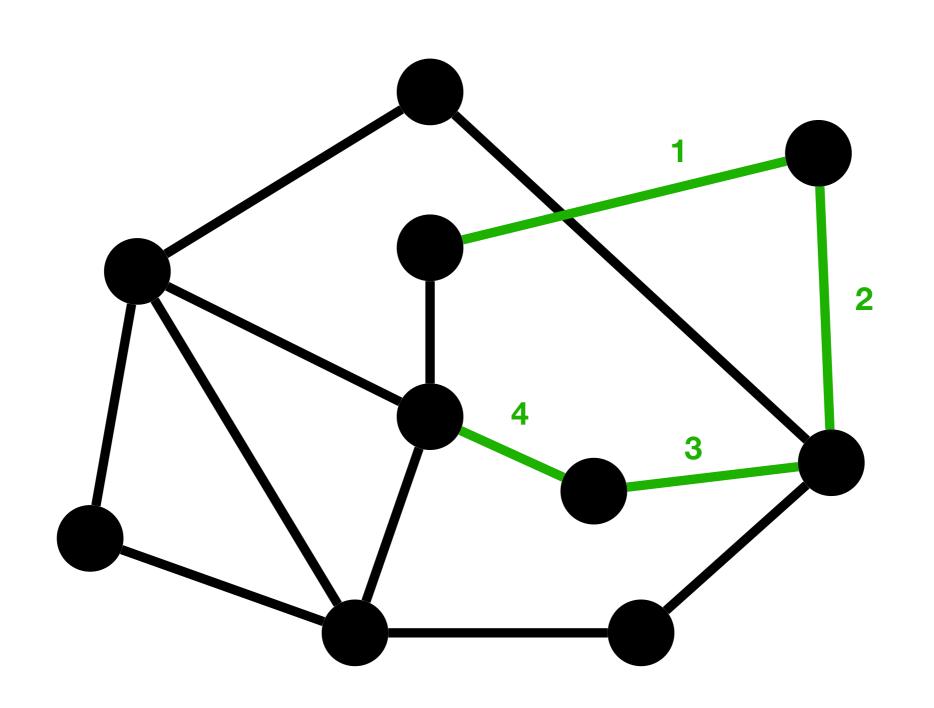


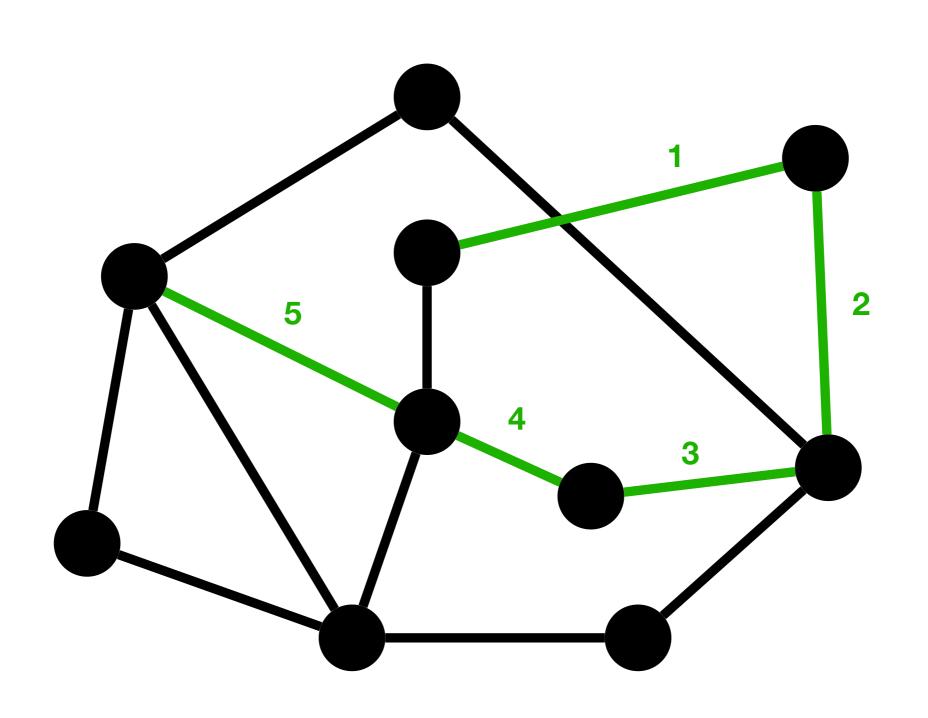


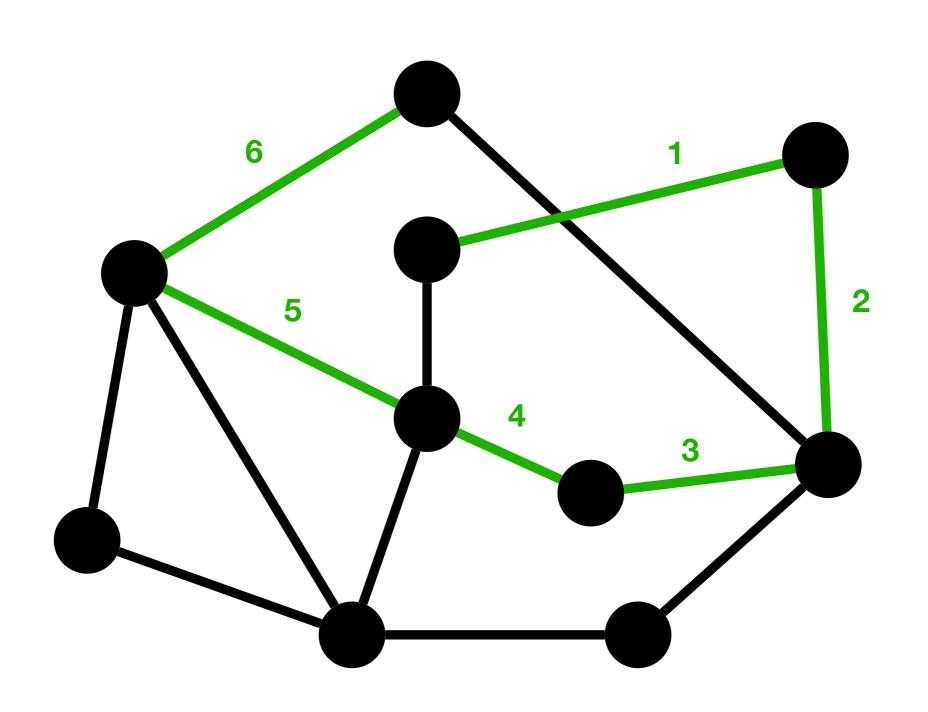


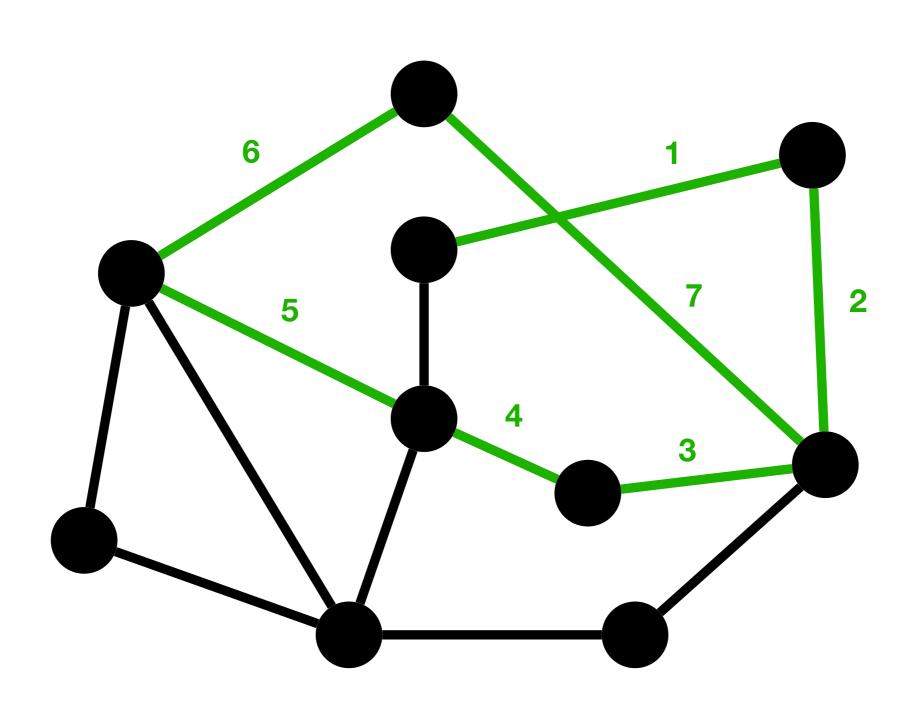


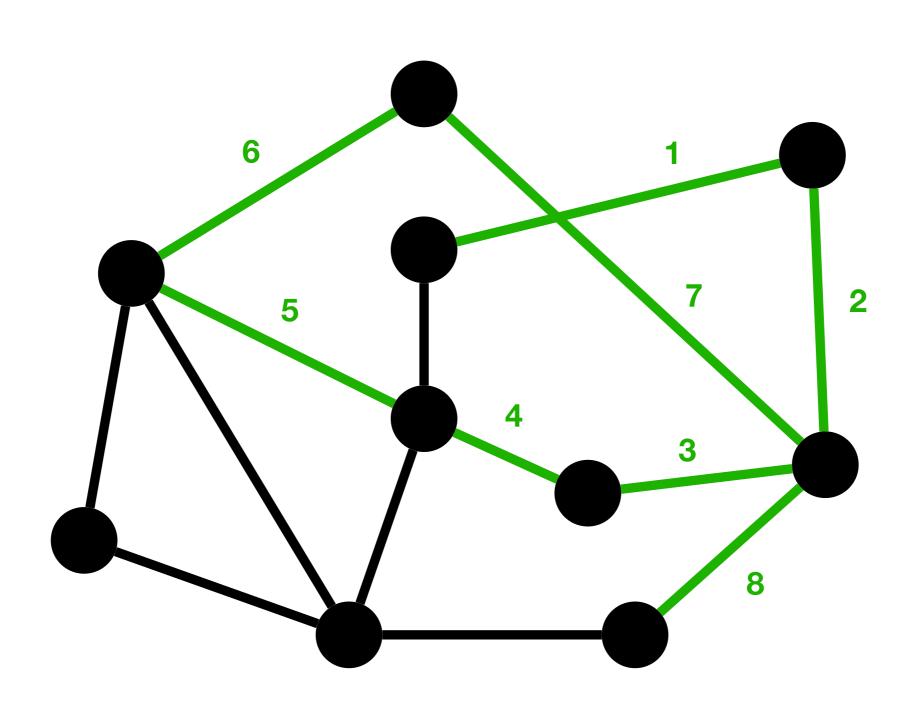


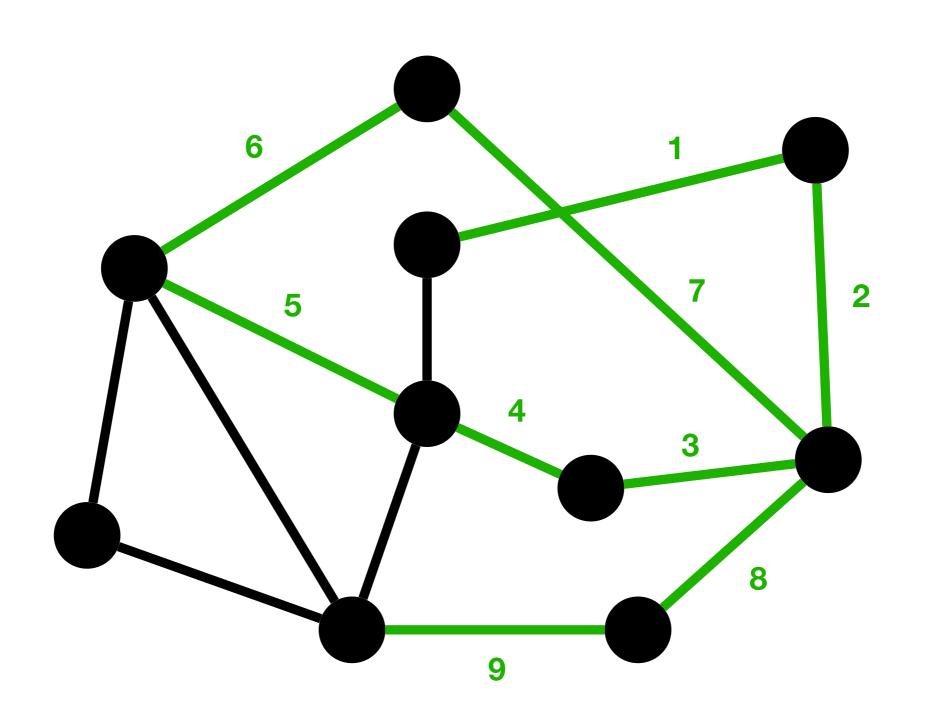


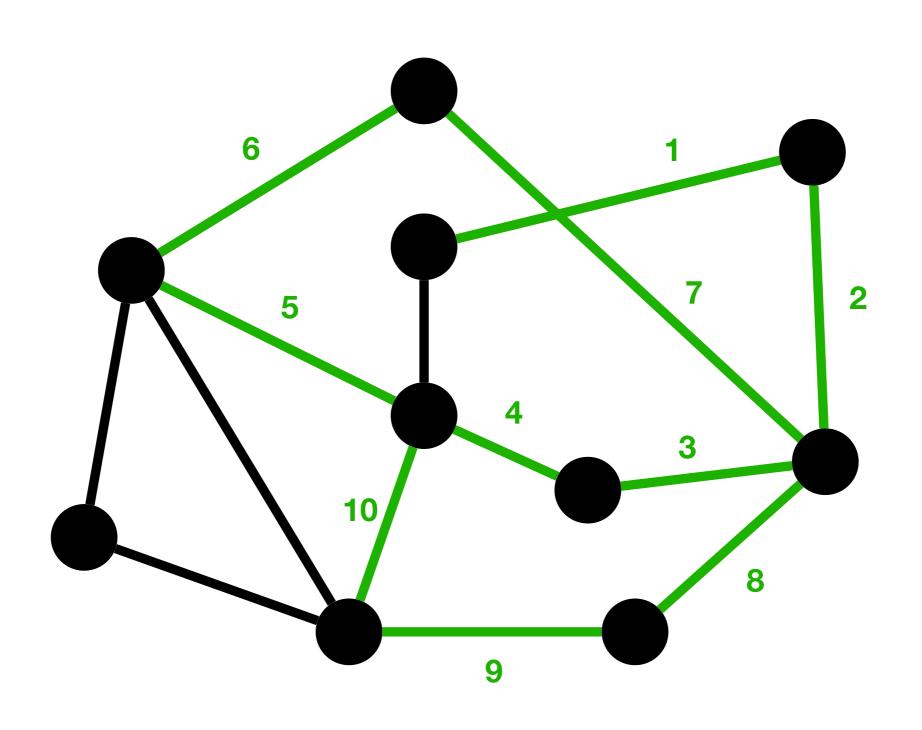


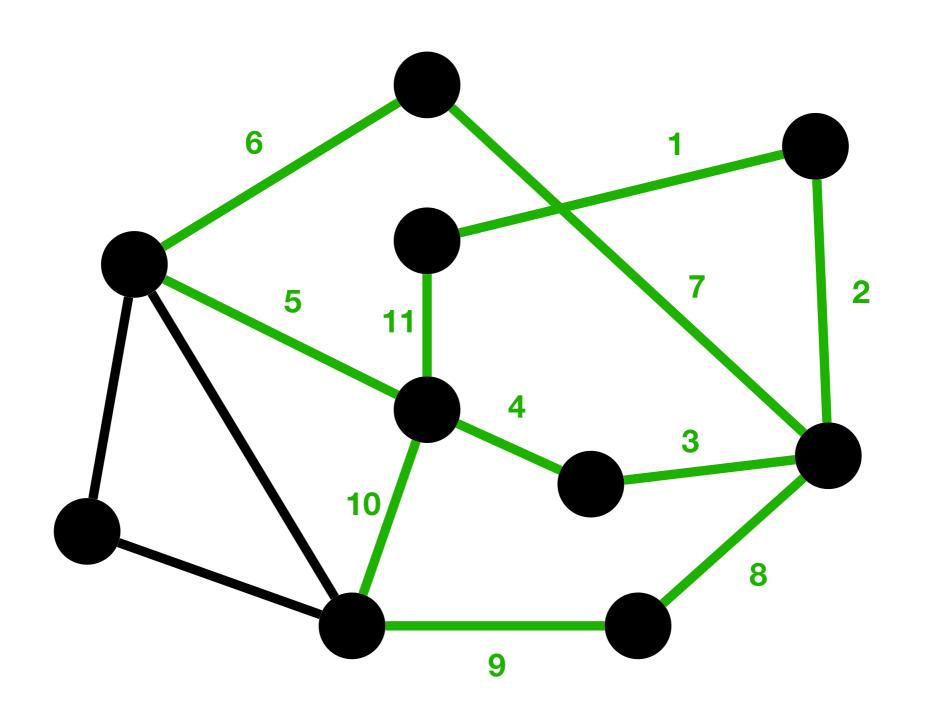


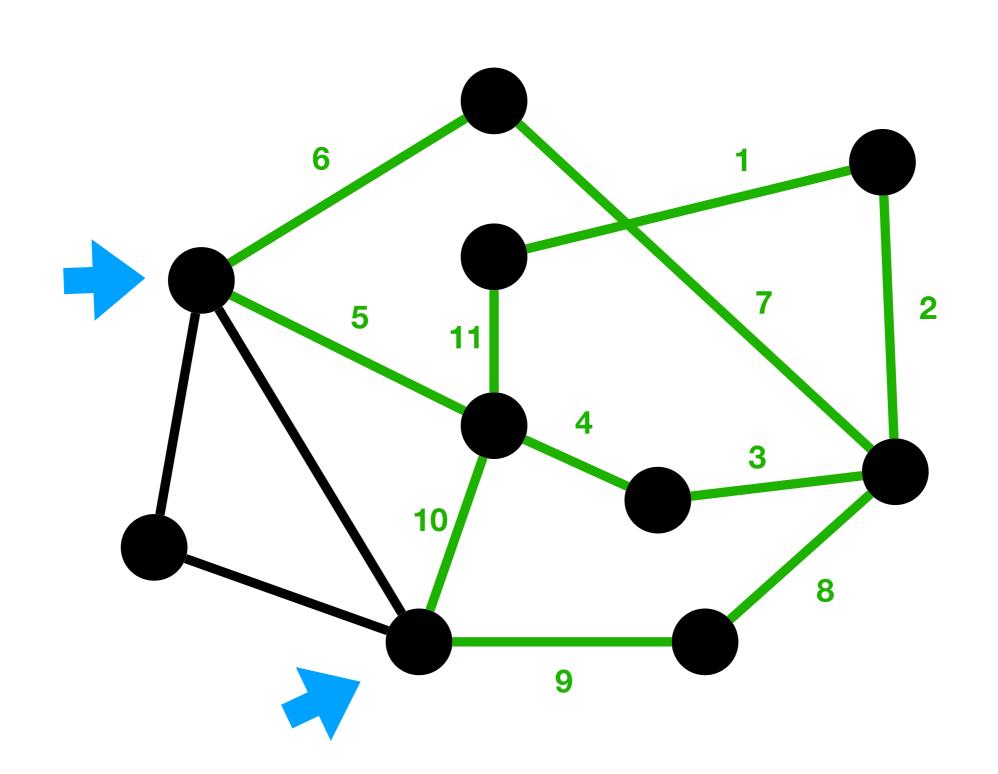


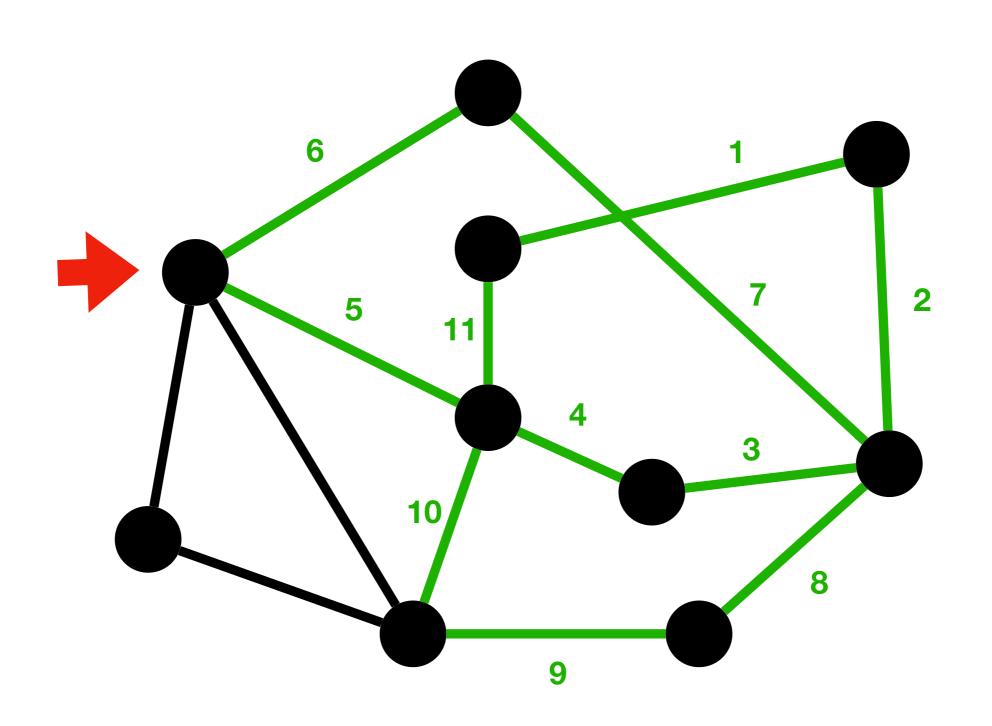


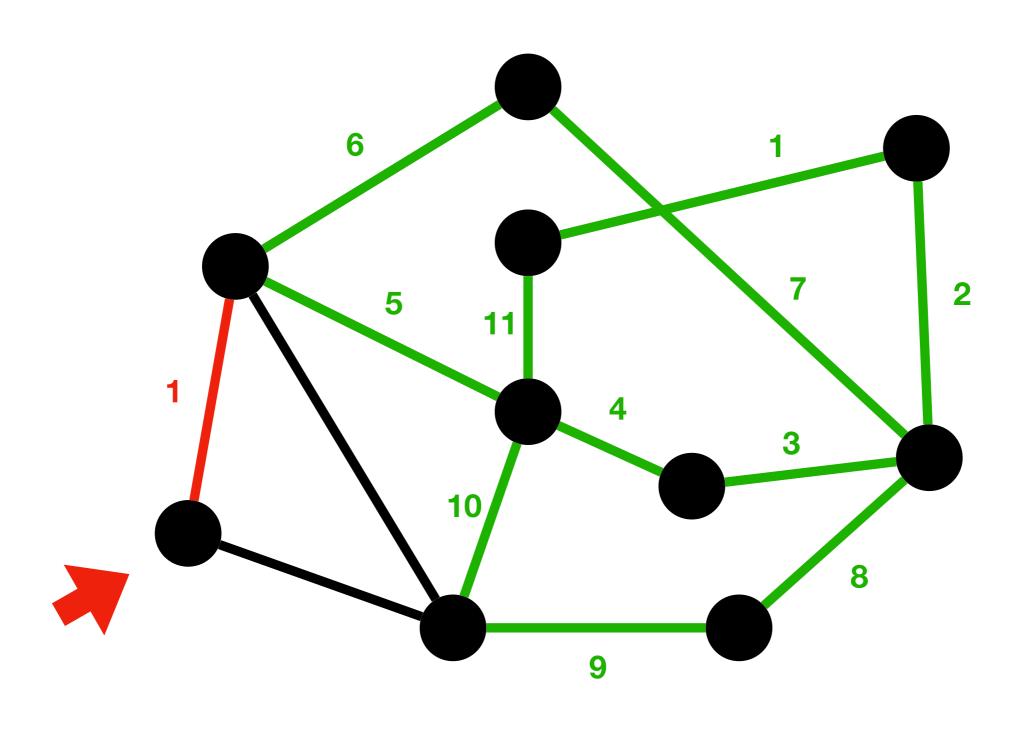


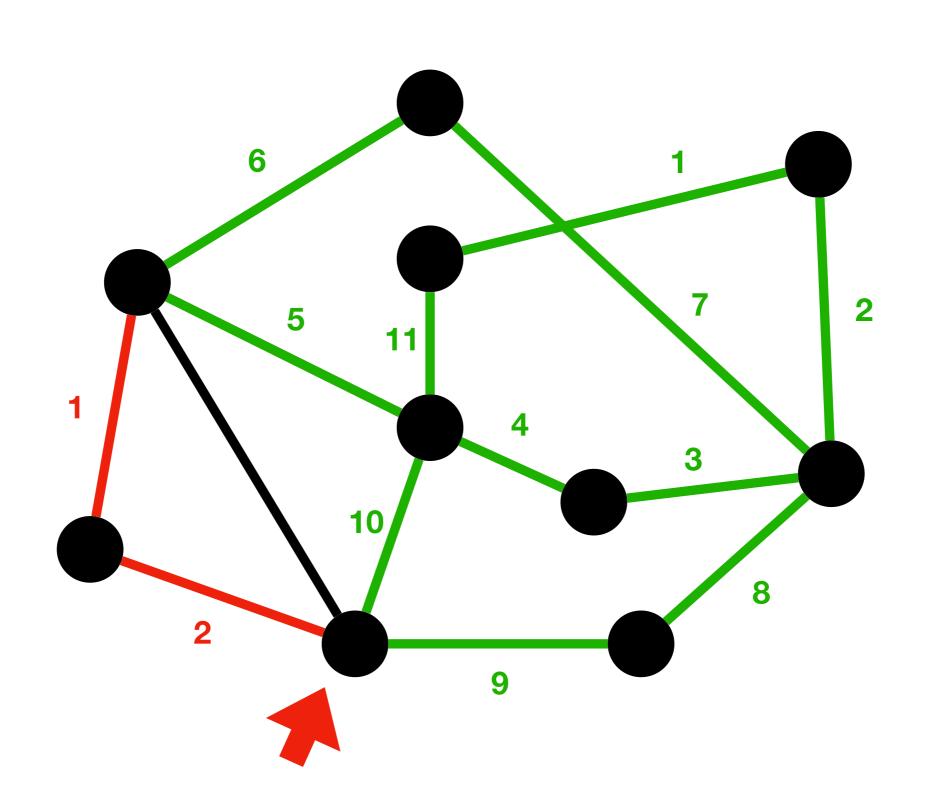


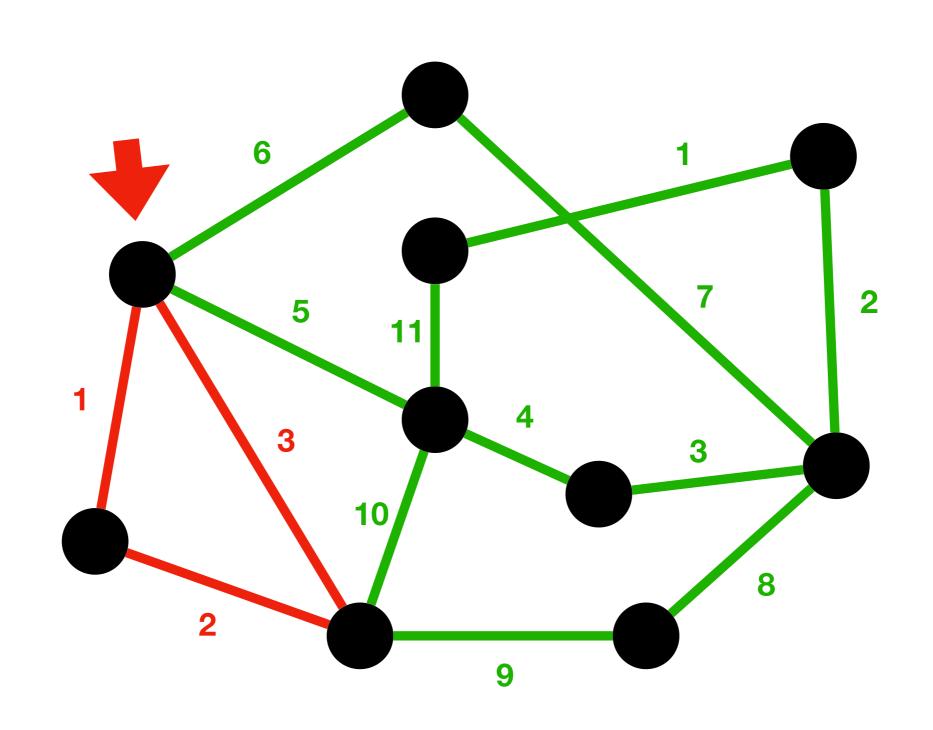


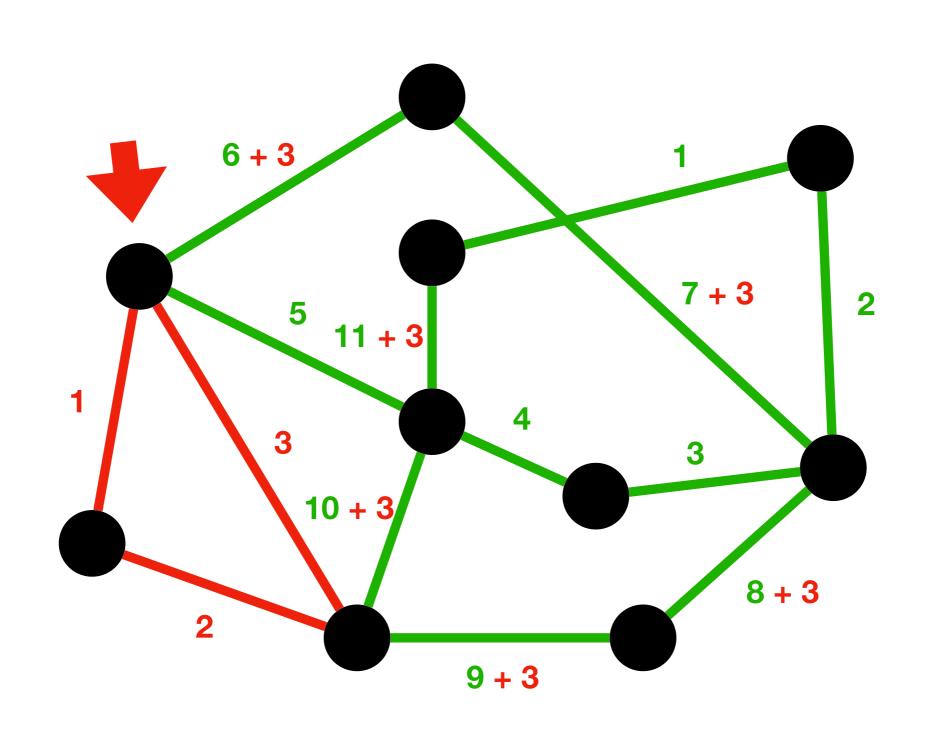


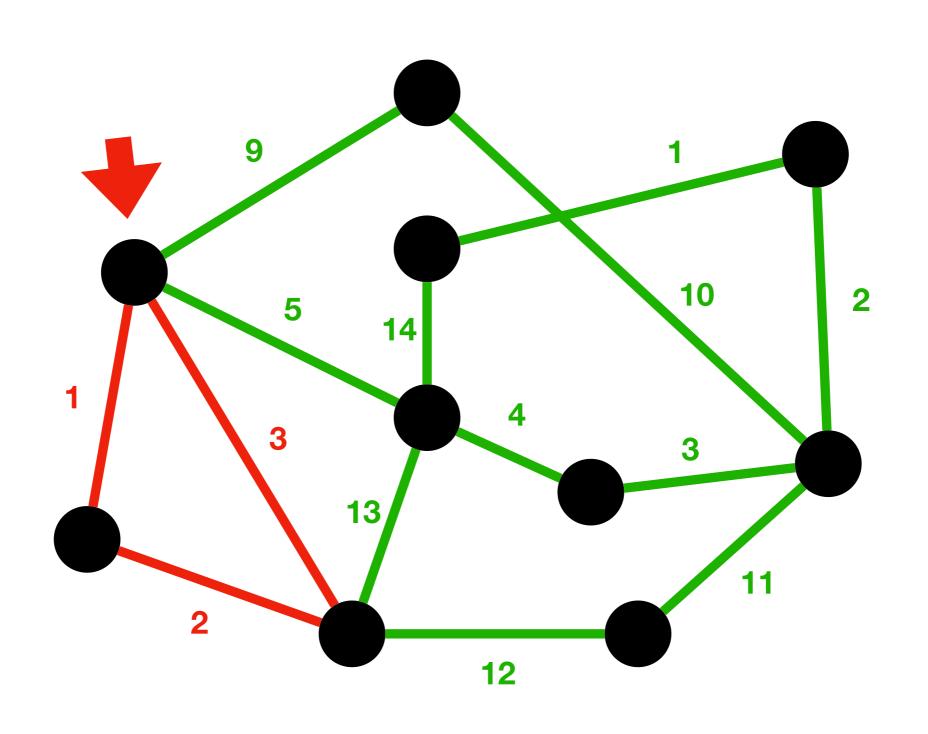


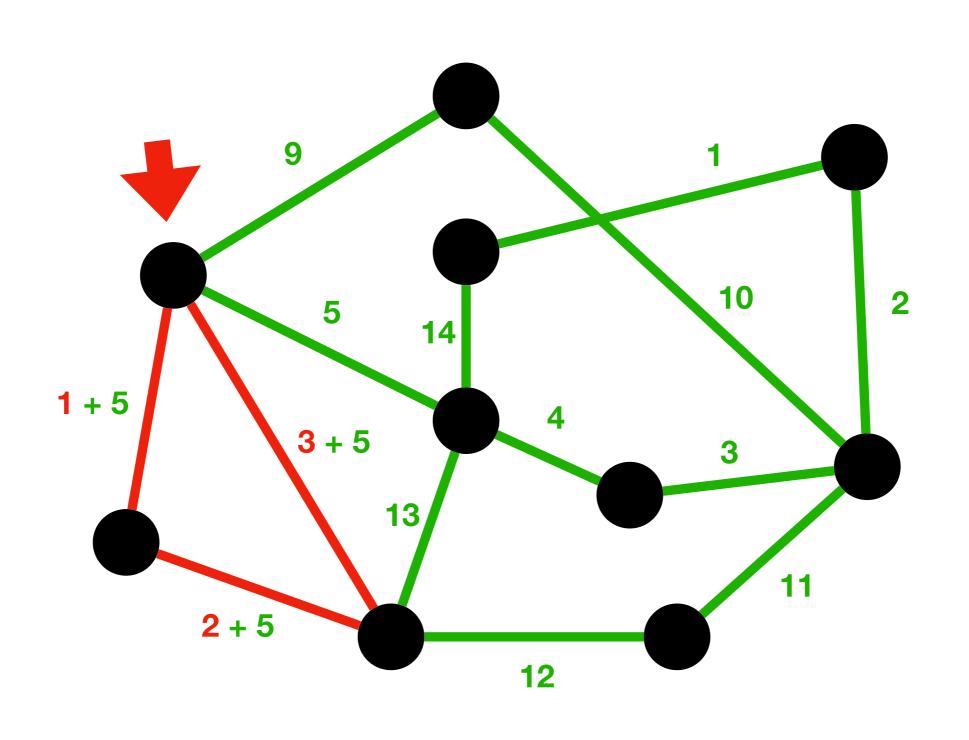


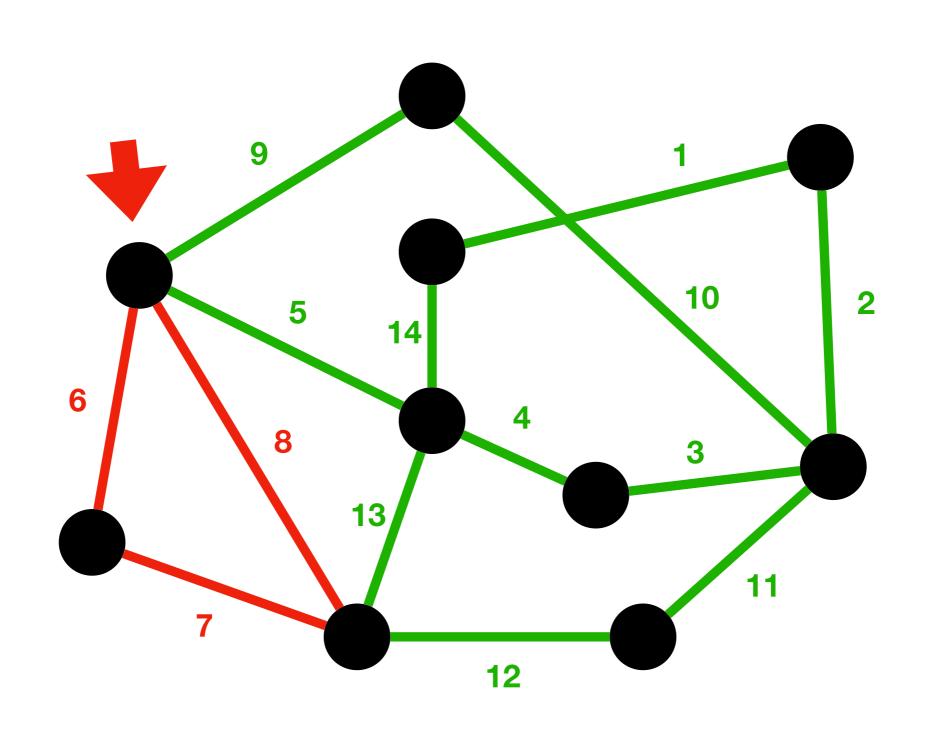


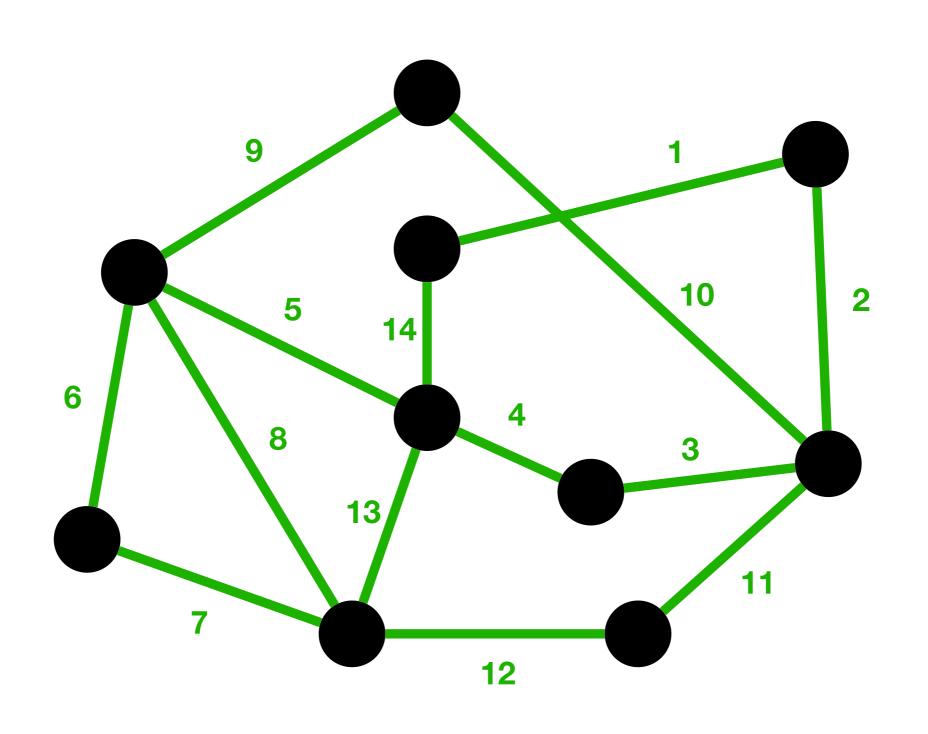


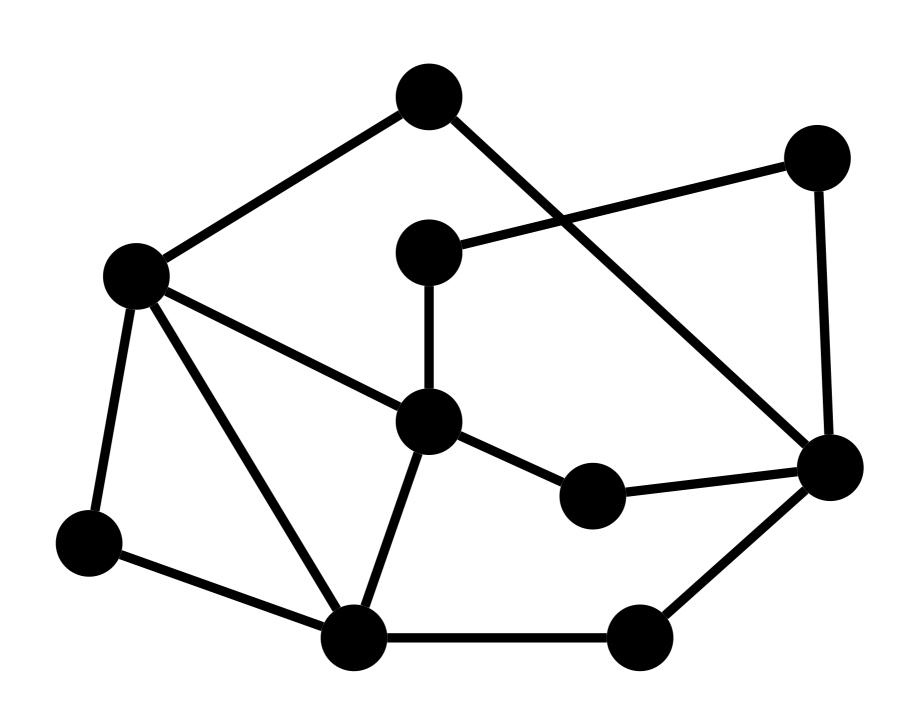


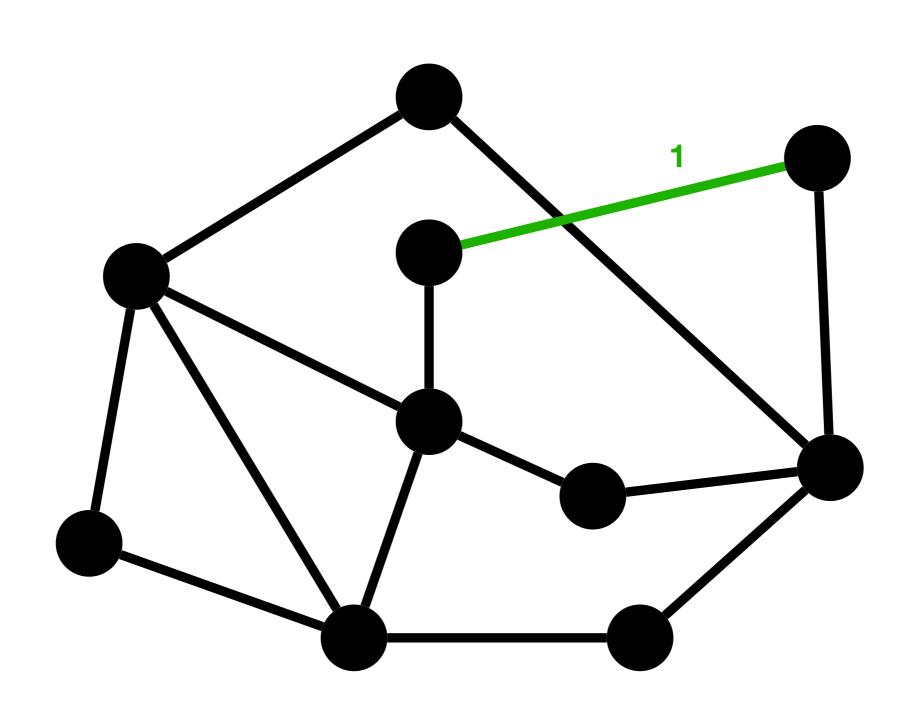


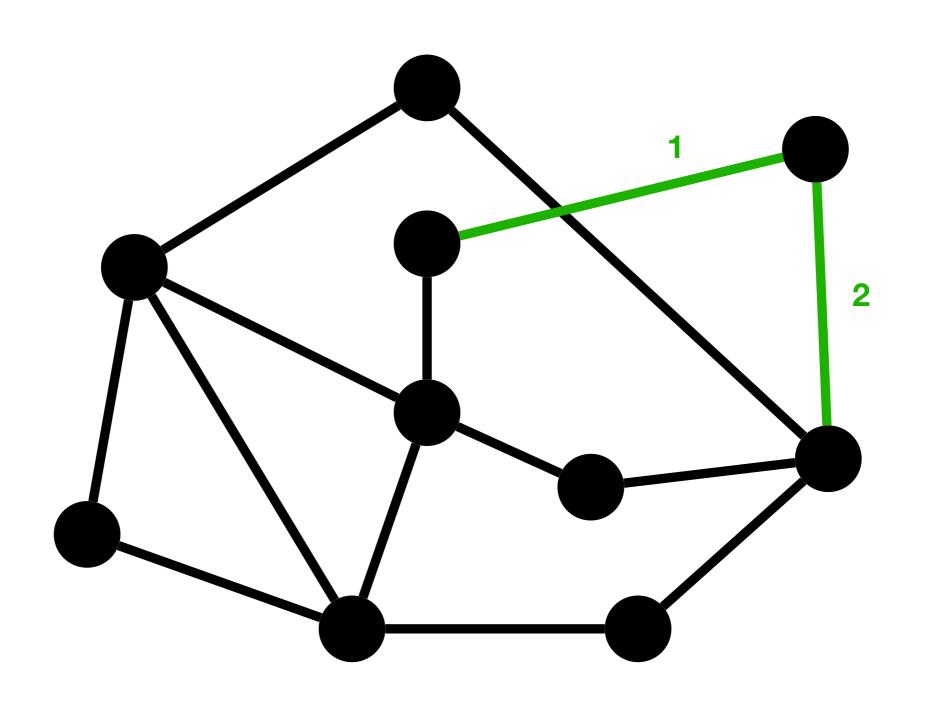


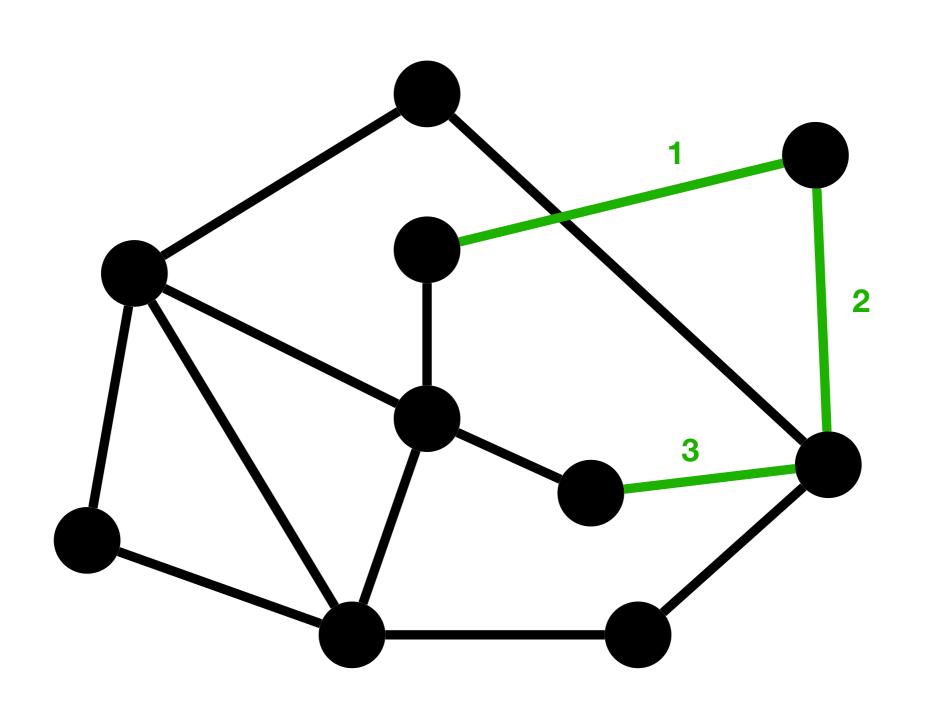


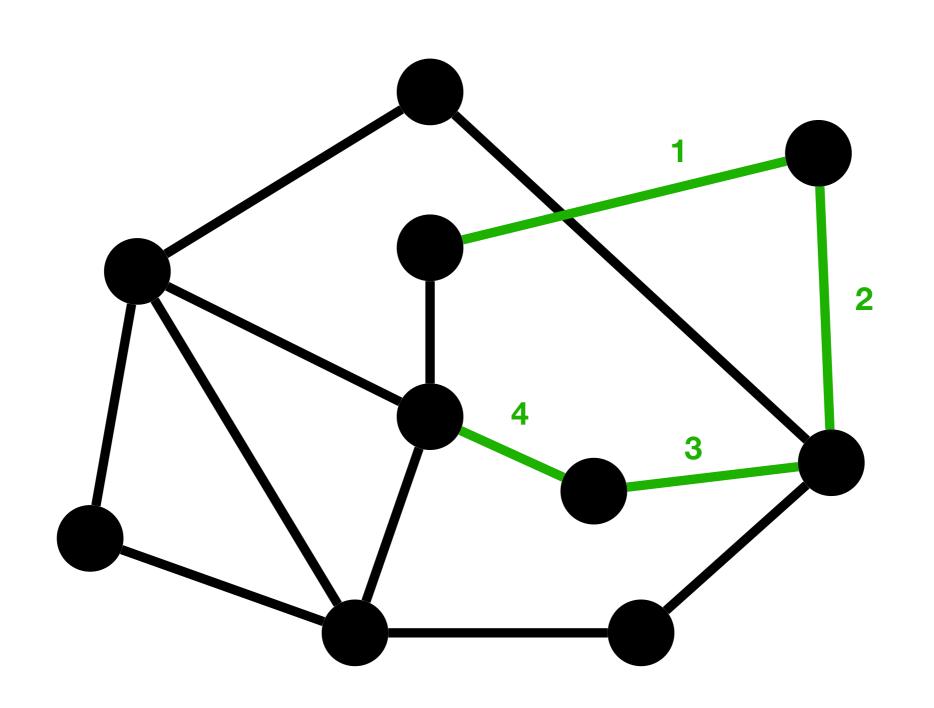


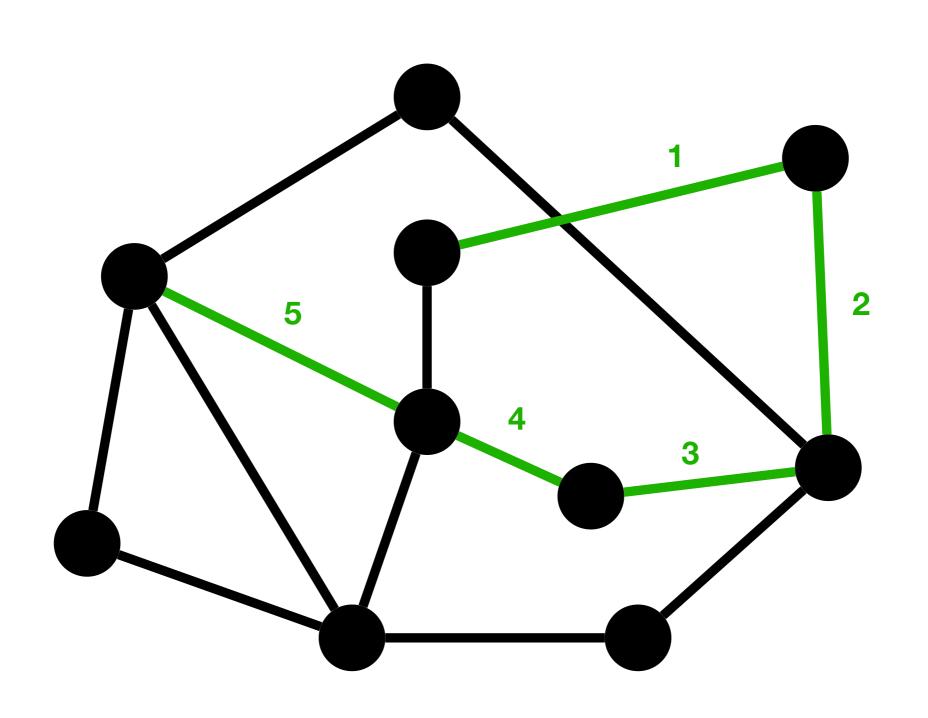


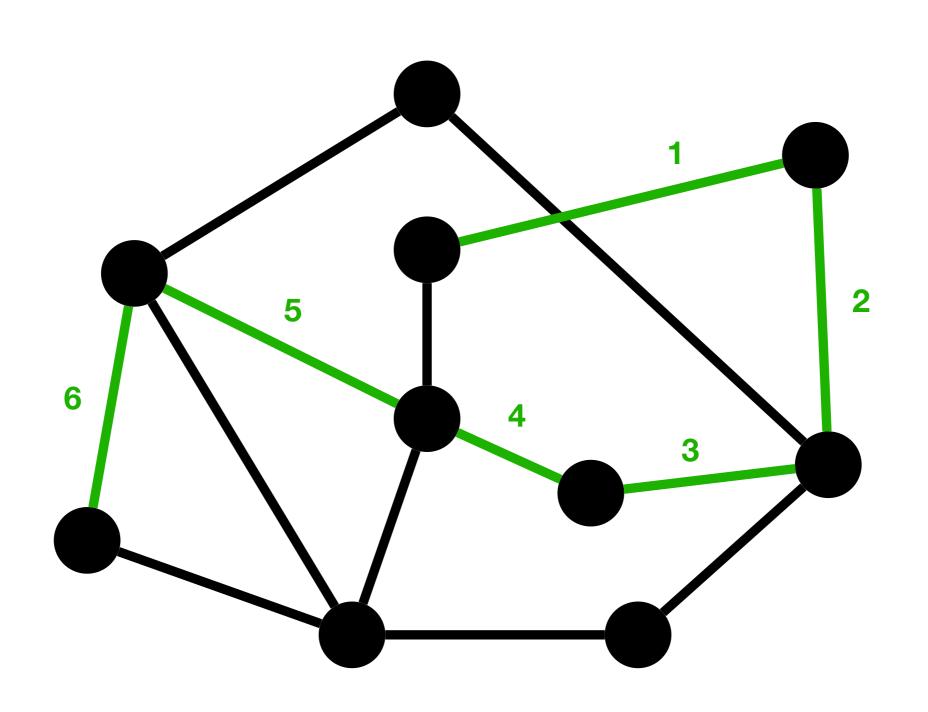


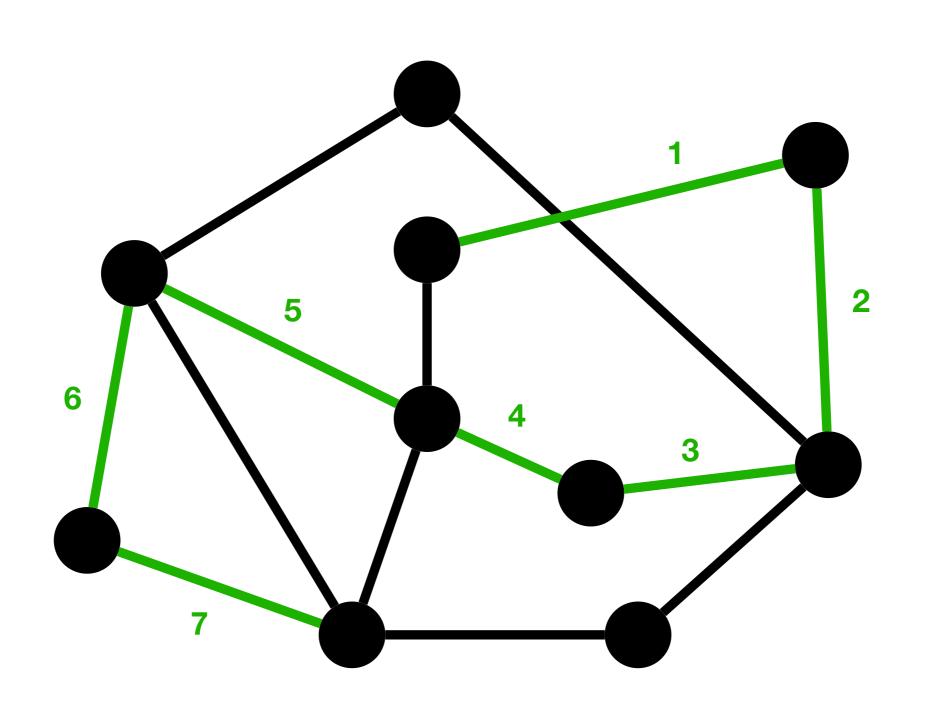


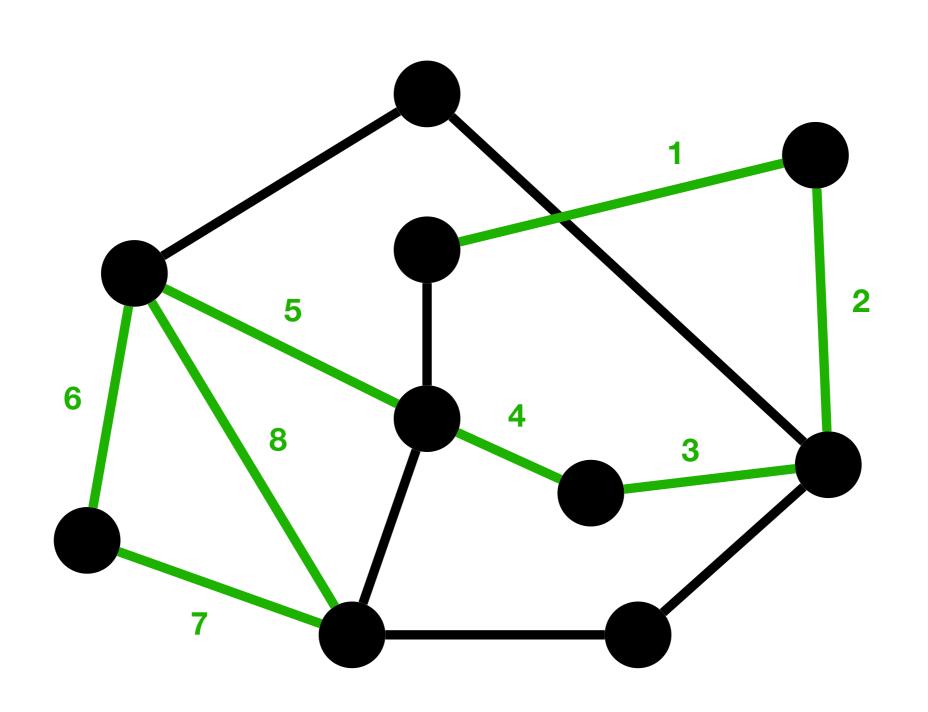


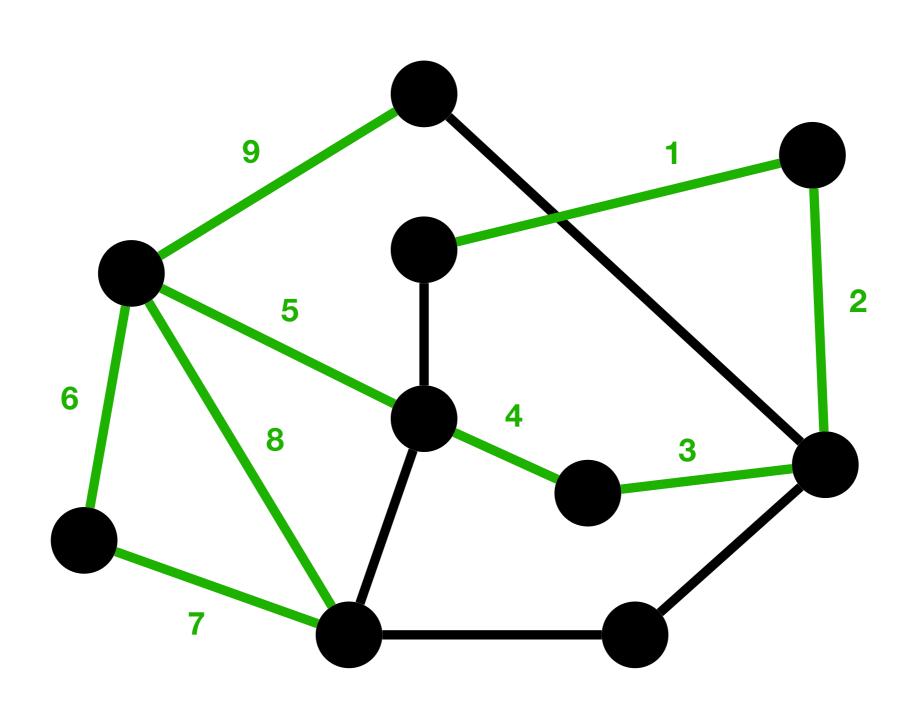


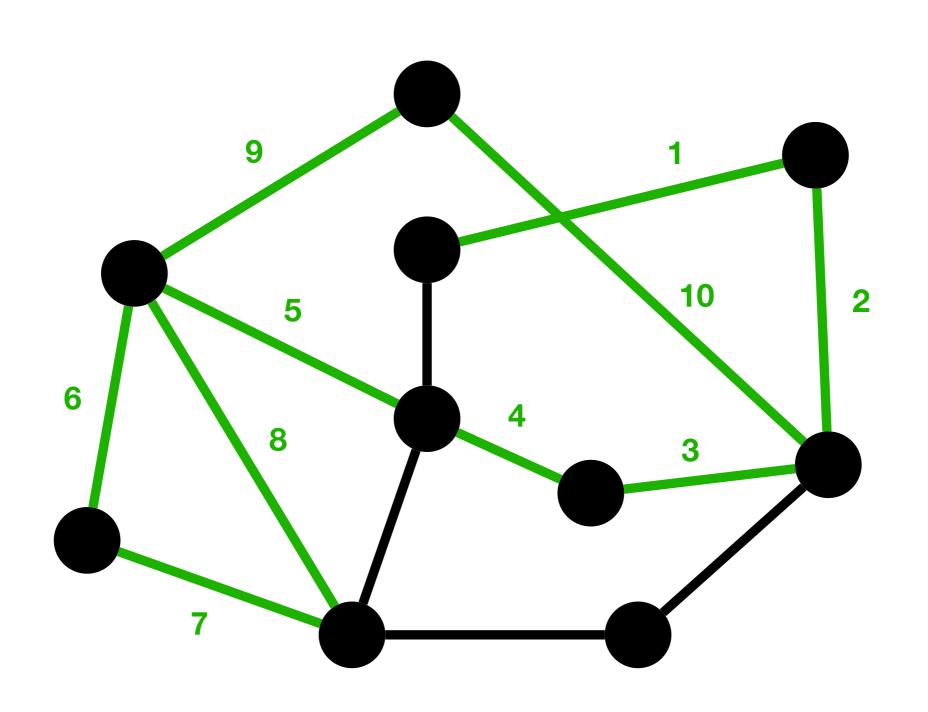


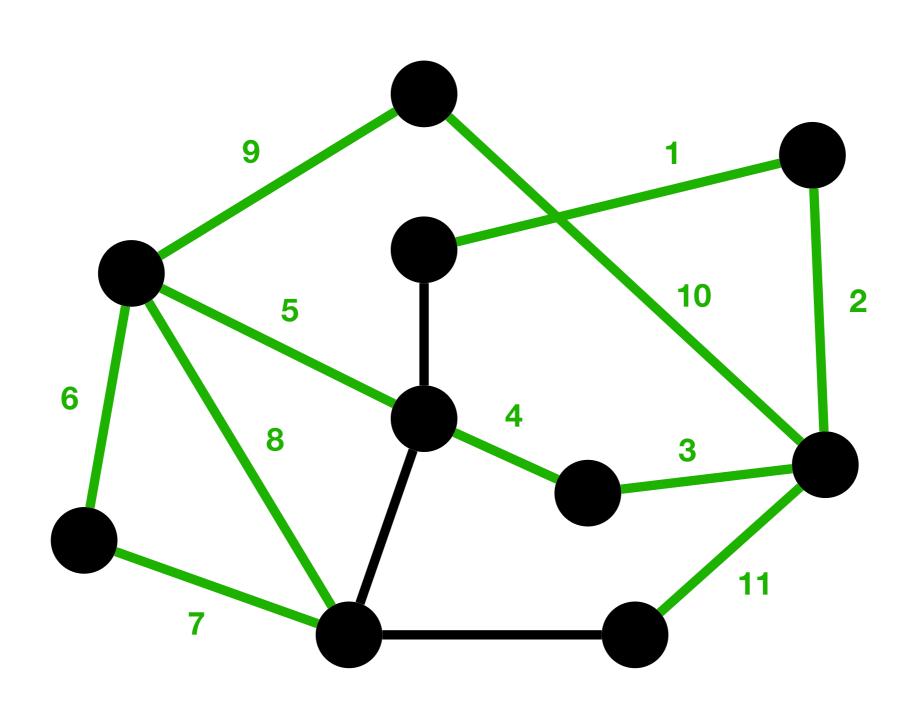


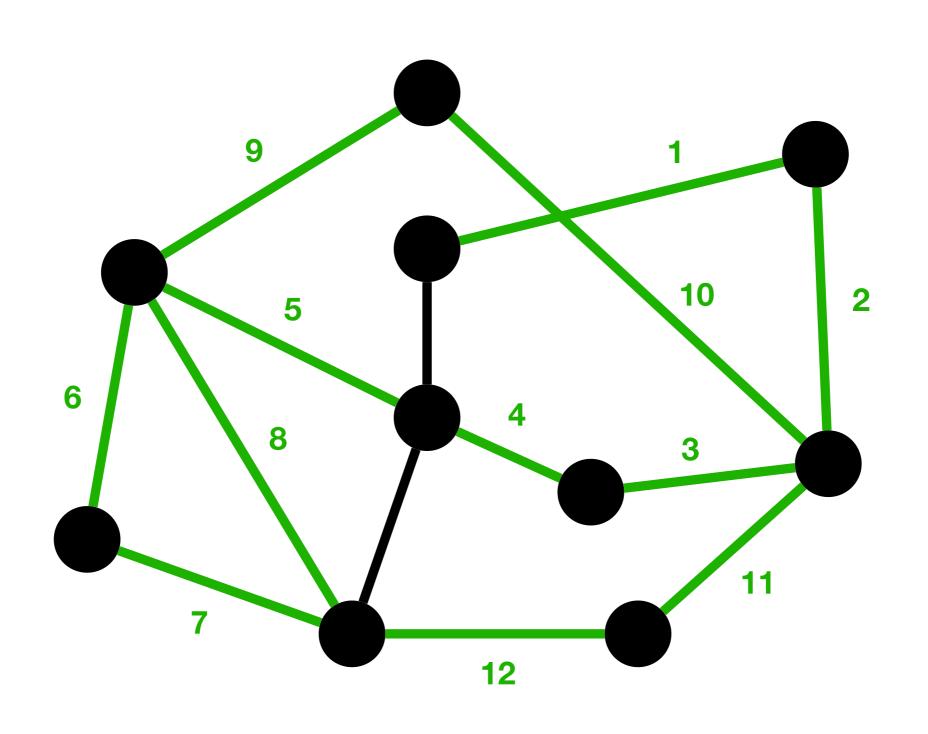


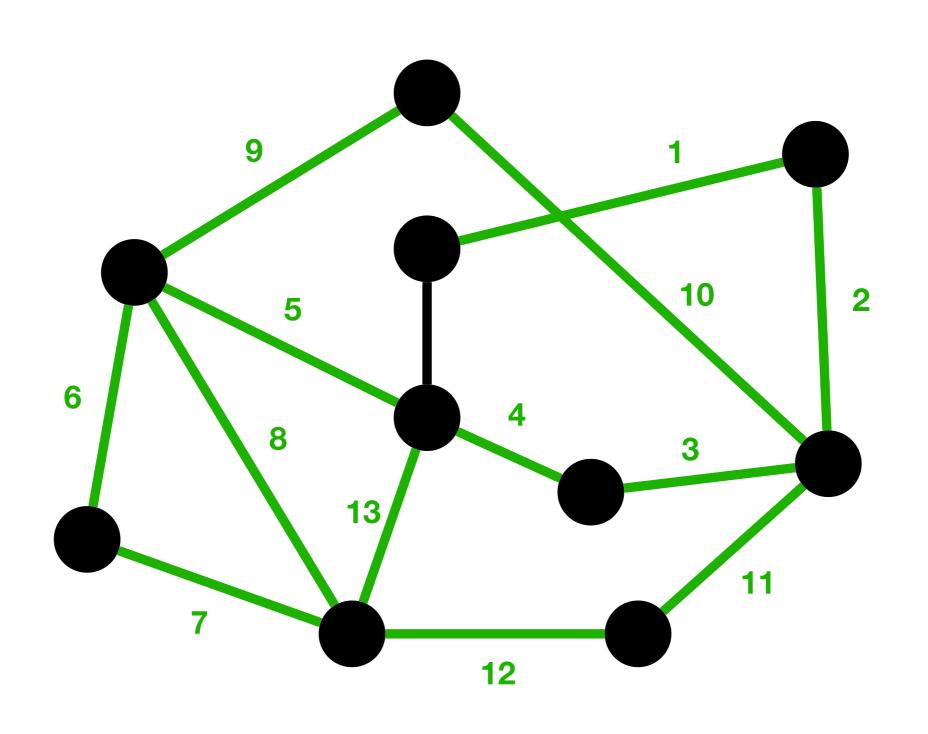


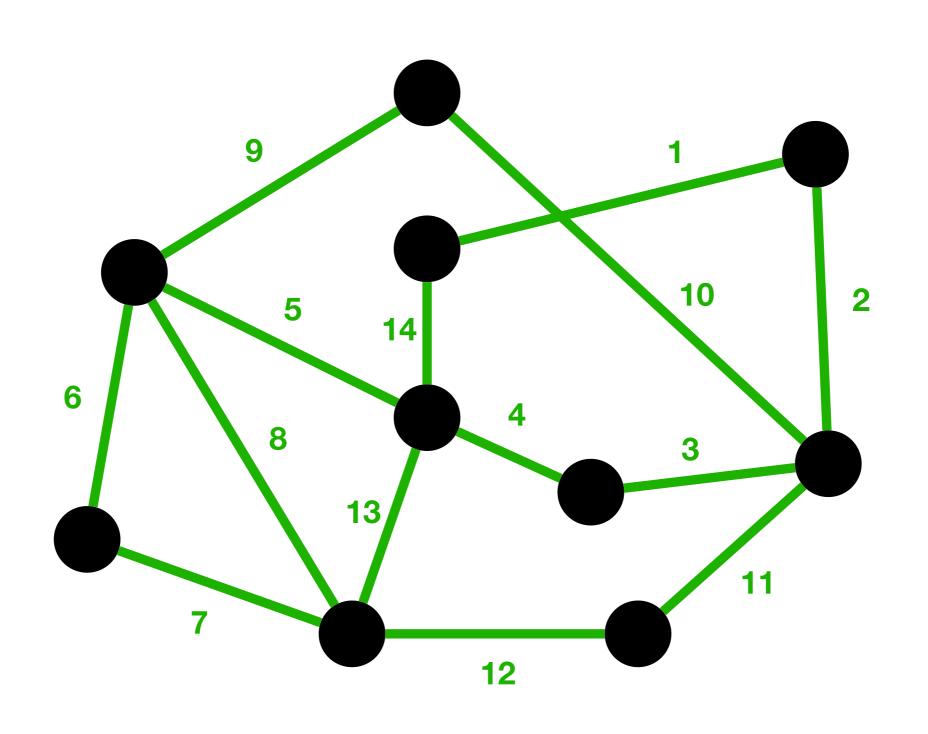












- Choisir n'importe quel sommet initial v
- Suivre un chemin arbitraire d'arêtes jusqu'à retourner à v, obtenant ainsi un cycle partiel c
- Tant qu'il y a des sommets u dans le cycle c avec des arêtes qu'on n'a pas encore choisit faire
 - Suivre un chemin de sommets à partir de u jusqu'à retourner à u, obtenant un cycle c'
 - Prolonger le cycle c par c'

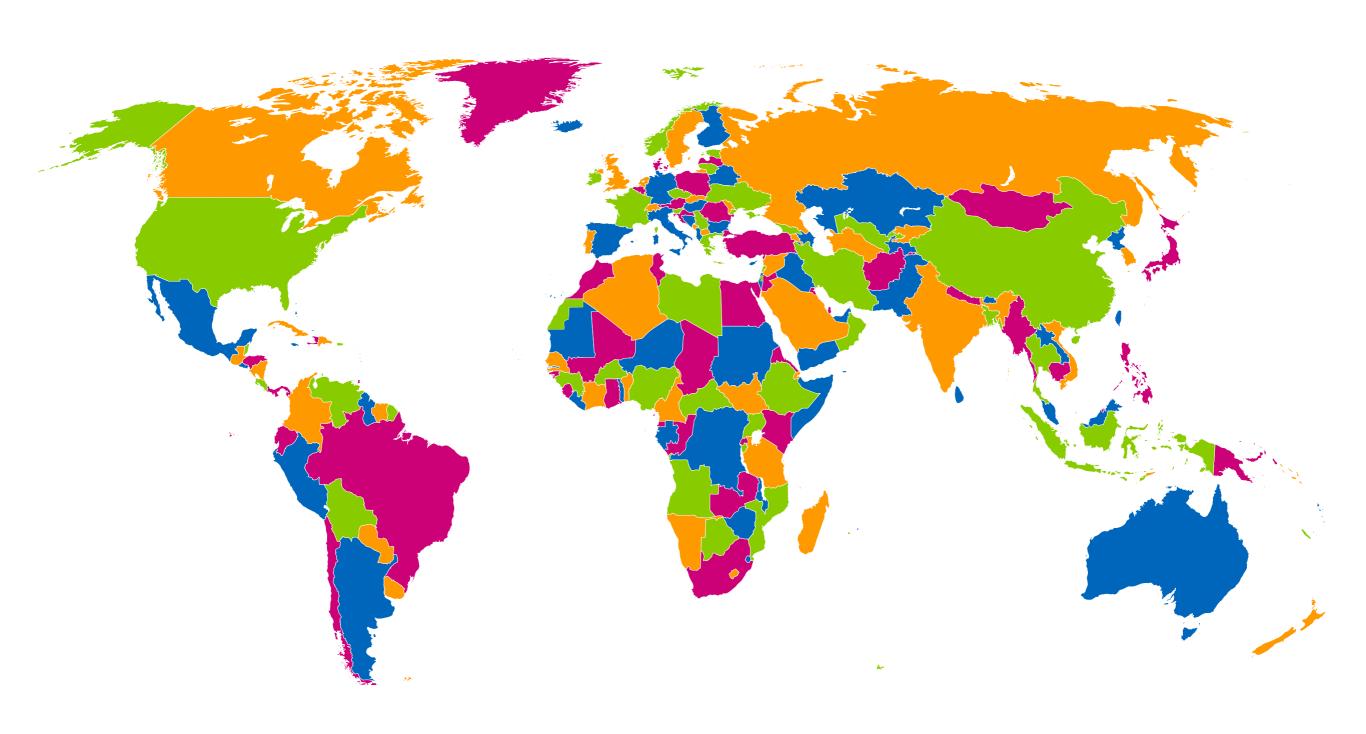
 « Choisir n'importe quel sommet initial v. Suivre un chemin arbitraire d'arêtes jusqu'à retourner à v, obtenant ainsi un cycle partiel c »

- « Choisir n'importe quel sommet initial v. Suivre un chemin arbitraire d'arêtes jusqu'à retourner à v, obtenant ainsi un cycle partiel c »
- On ne peut pas rester bloqué dans aucun sommet different de v, puisque ils ont toujours un nombre pair d'arêtes pas encore explorées (« si on entre dans un sommet, on peut toujours en sortir »)

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- Et, tôt ou tard, on retourne à *v*, puisque dans le pire des cas on épuise l'ensemble des autres arêtes, et *v* a un nombre **impair** d'arêtes inexplorées (donc il reste une arête pour y entrer)

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- Ça est également vrais pour chaque sommet u choisi après v

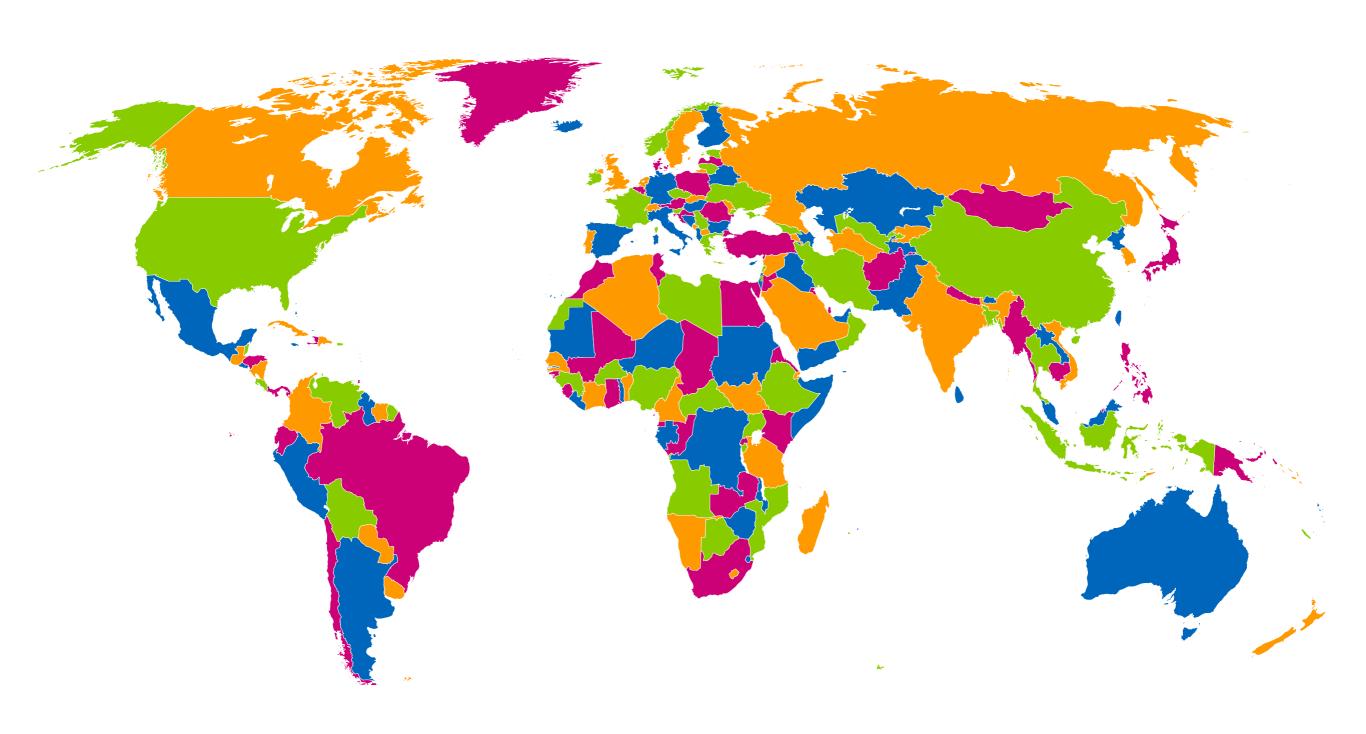
Coloration de cartes

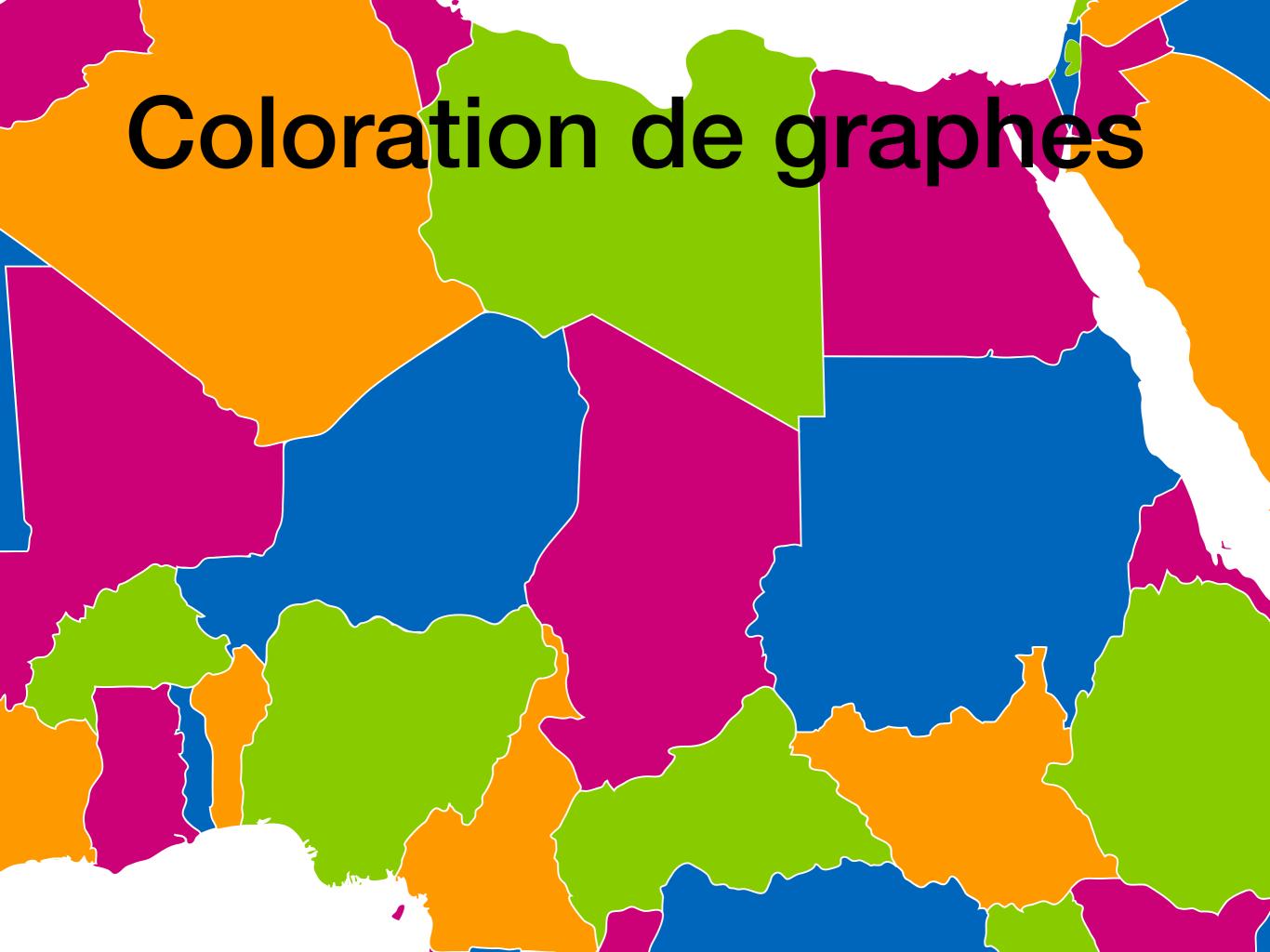


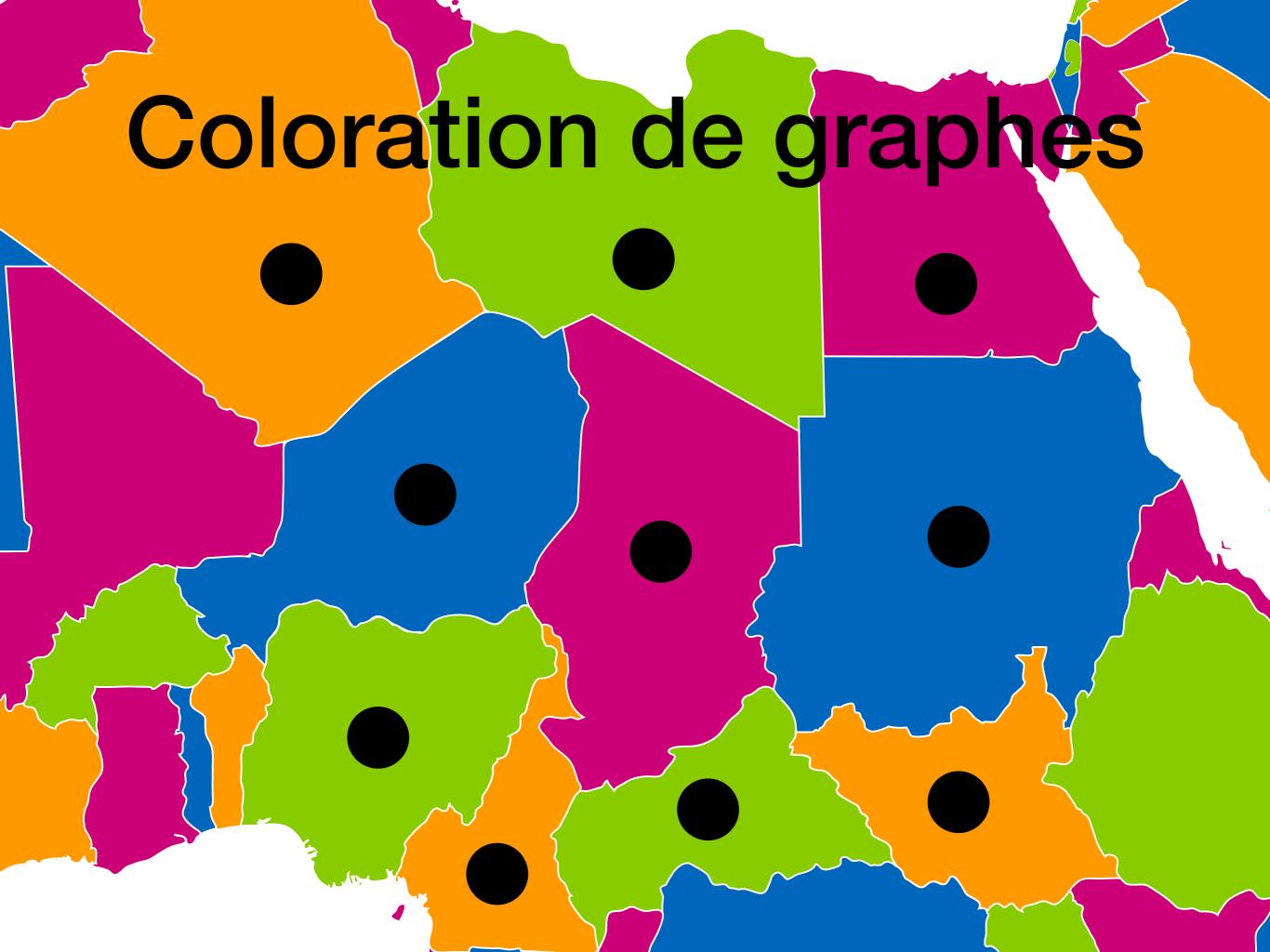
Théorème des quatre couleurs

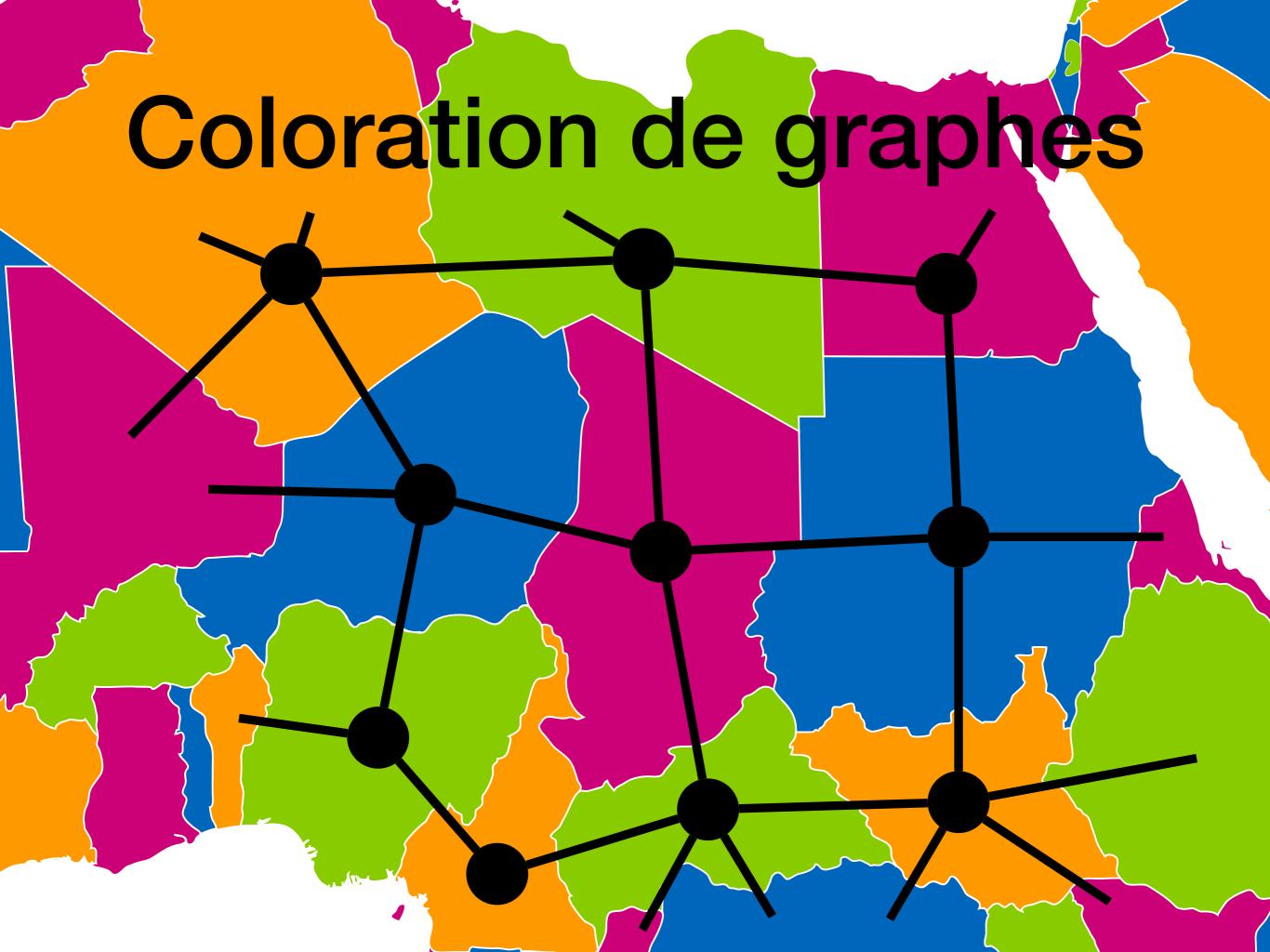
On peut colorer n'importe quelle carte avec un maximum de 4 couleurs de sorte que les pays (régions, etc.) adjacents reçoivent toujours deux couleurs distinctes

Coloration de graphes

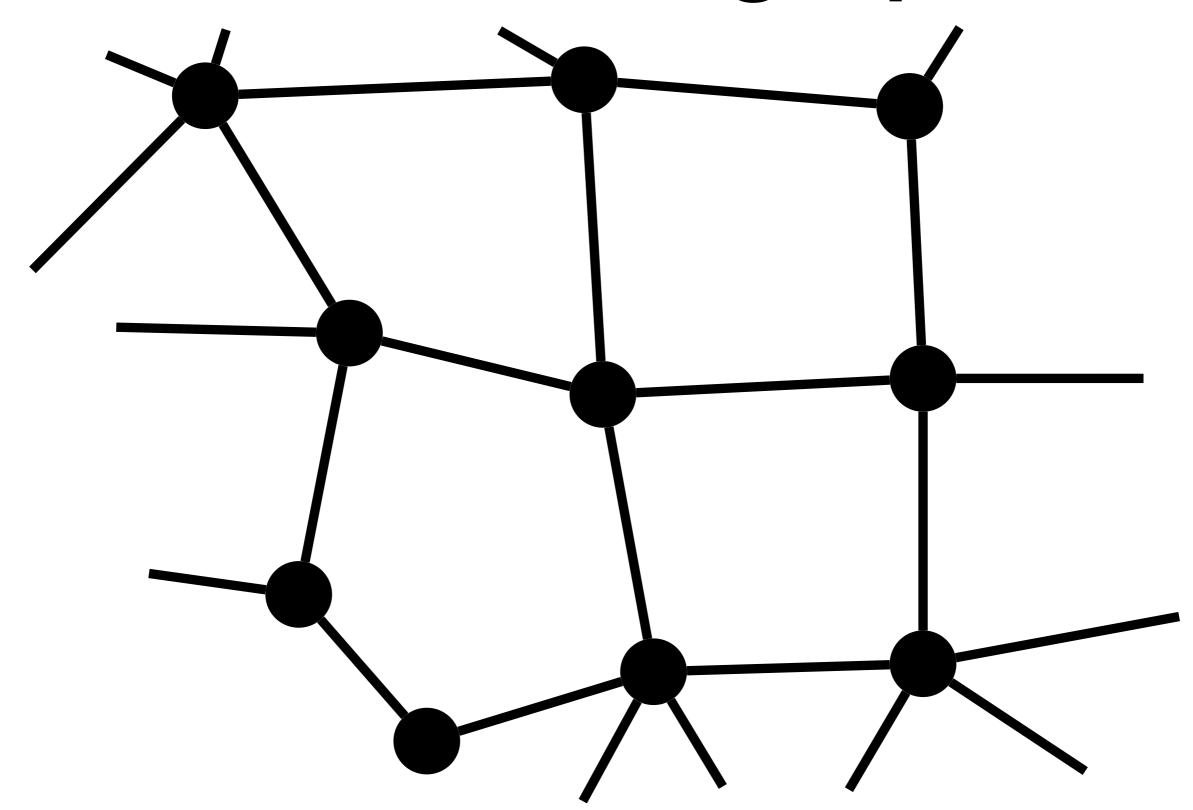




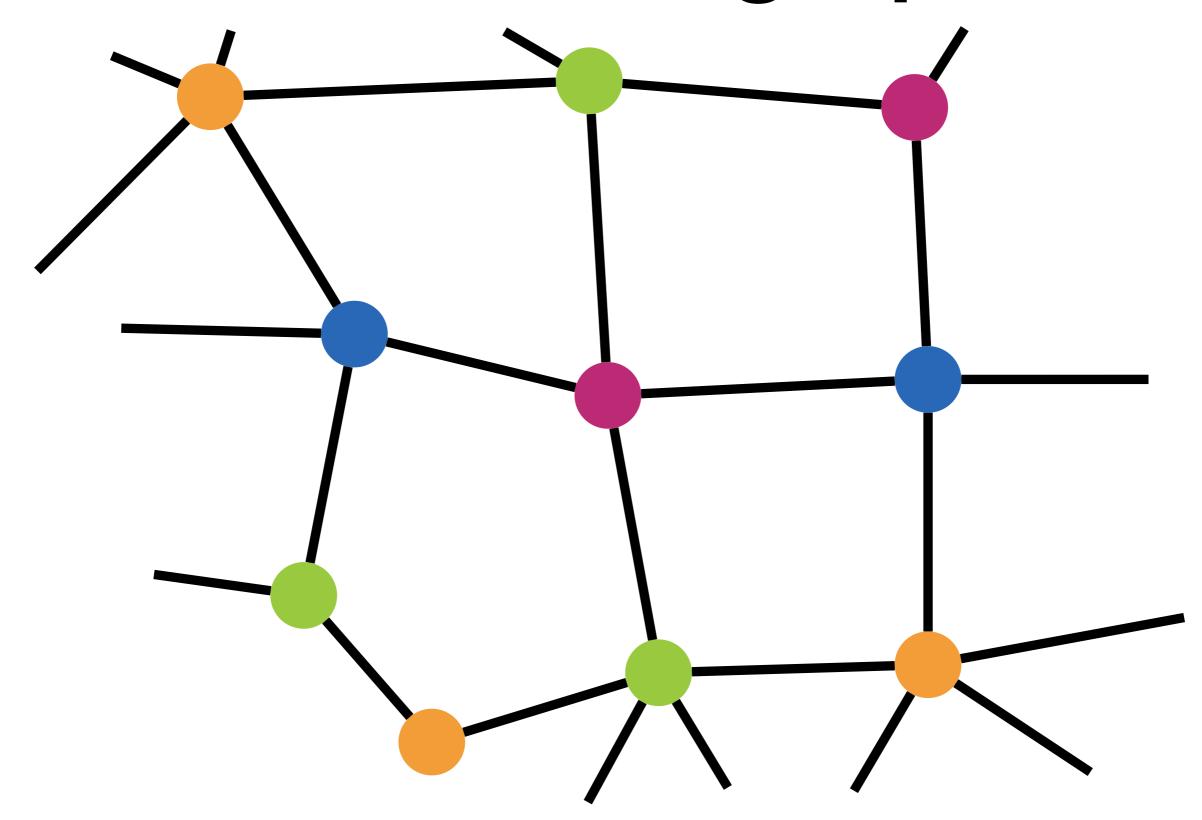




Coloration de graphes



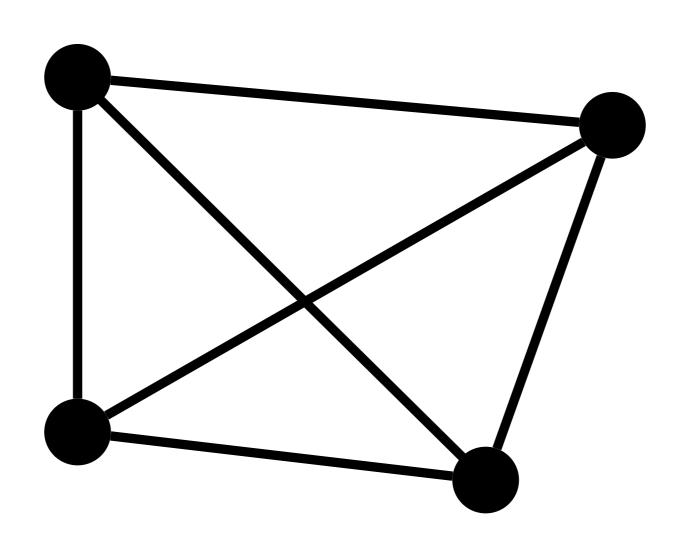
Coloration de graphes



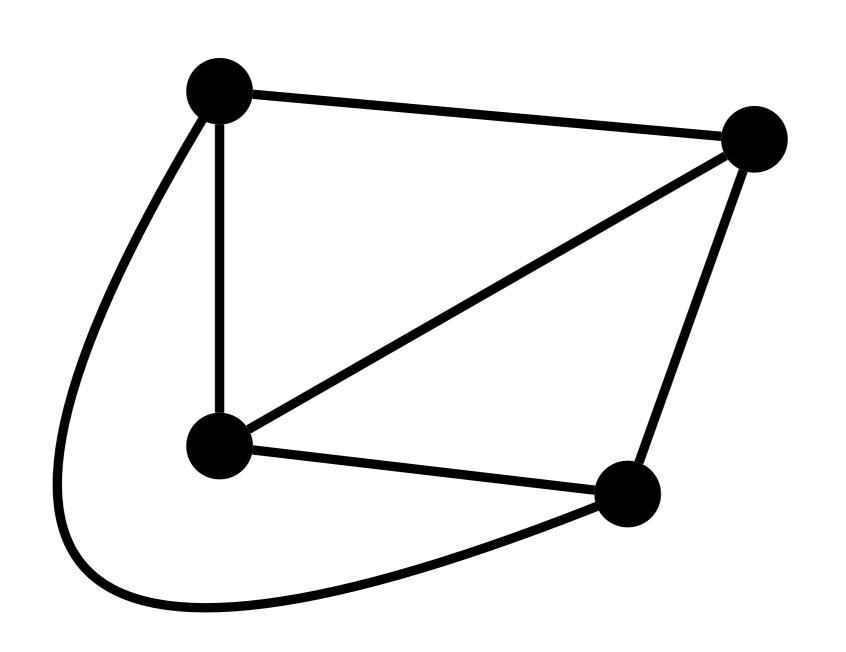
Graphes planaires

- Un graphe (non orienté) planaire est un graphe qu'on peut dessiner sur un plan (une feuille de papier) sans qu'aucune arête n'en croise une autre
- Un graphe non orienté G est planaire si et seulement si il existe une carte ayant G comme graphe associé

Graphes qui ne semblent pas planaires

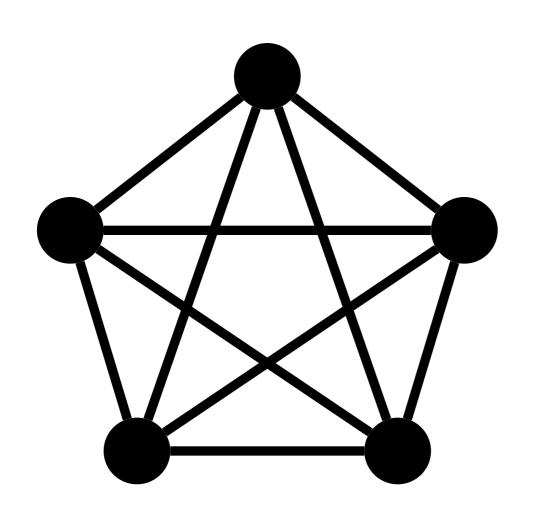


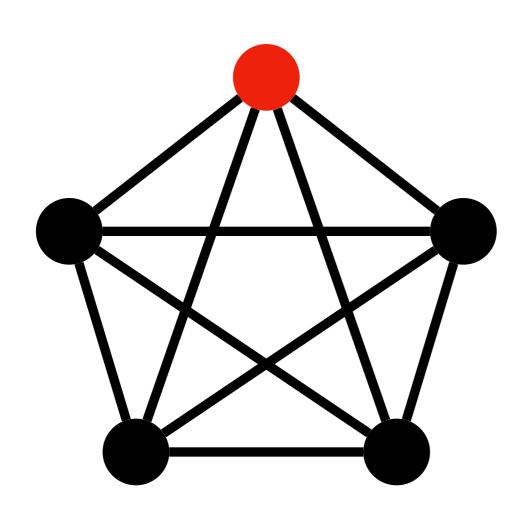
Graphes qui ne semblent pas planaires

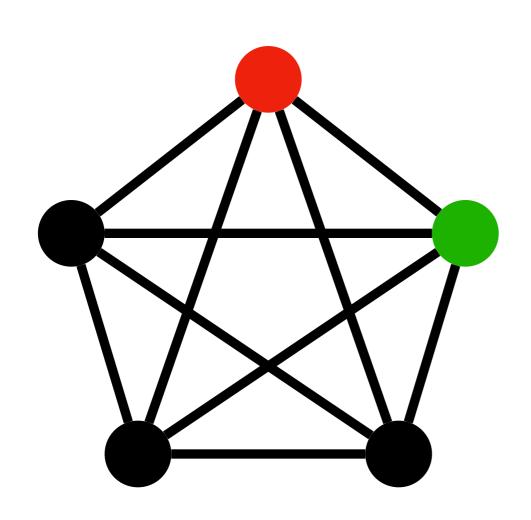


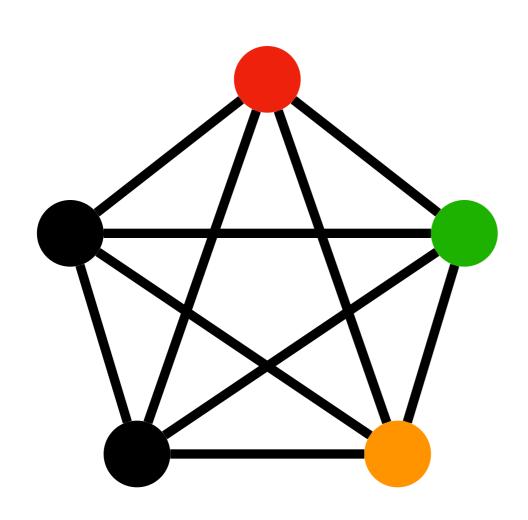
Théorème des quatre couleurs

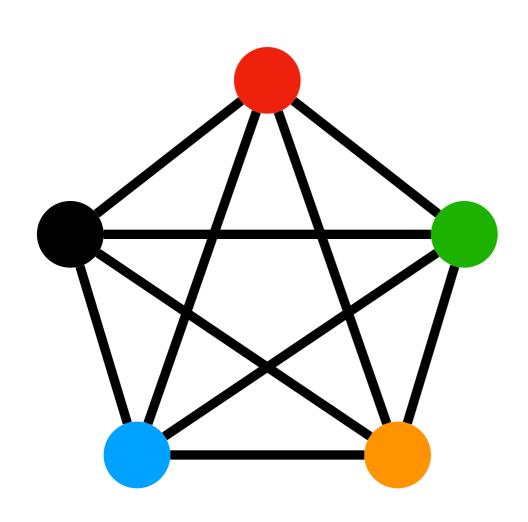
On peut colorer les sommets de n'importe quel graphe non orienté planaire avec un maximum de 4 couleurs de sorte que les sommets adjacents reçoivent toujours deux couleurs distinctes

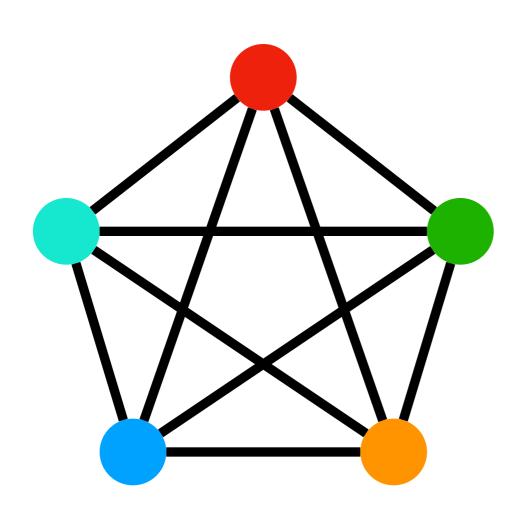


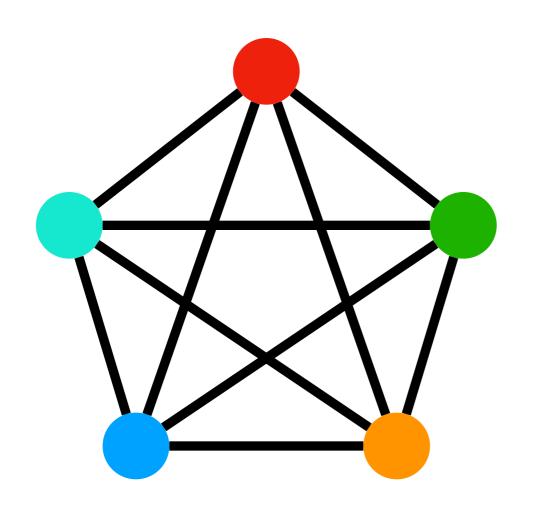


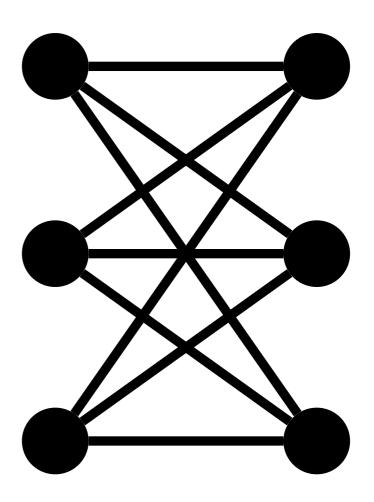


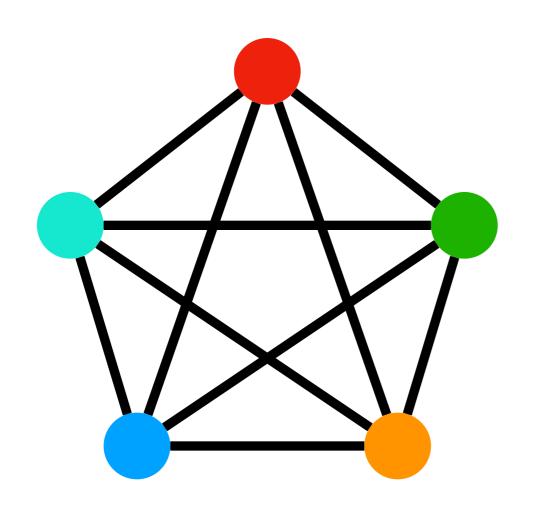


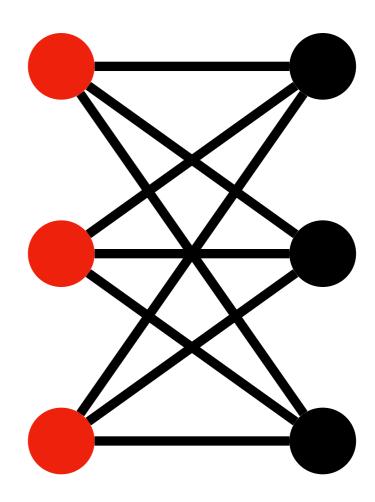


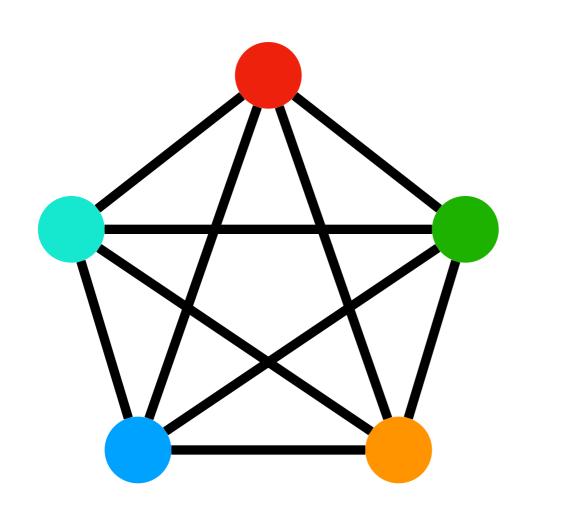


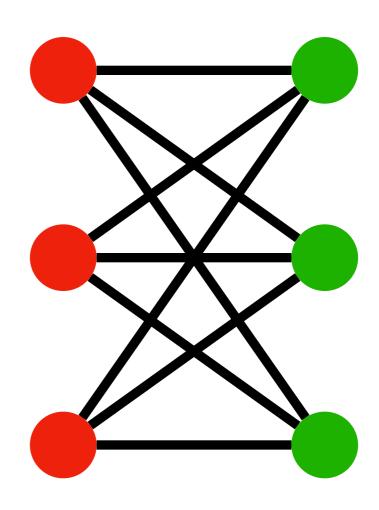


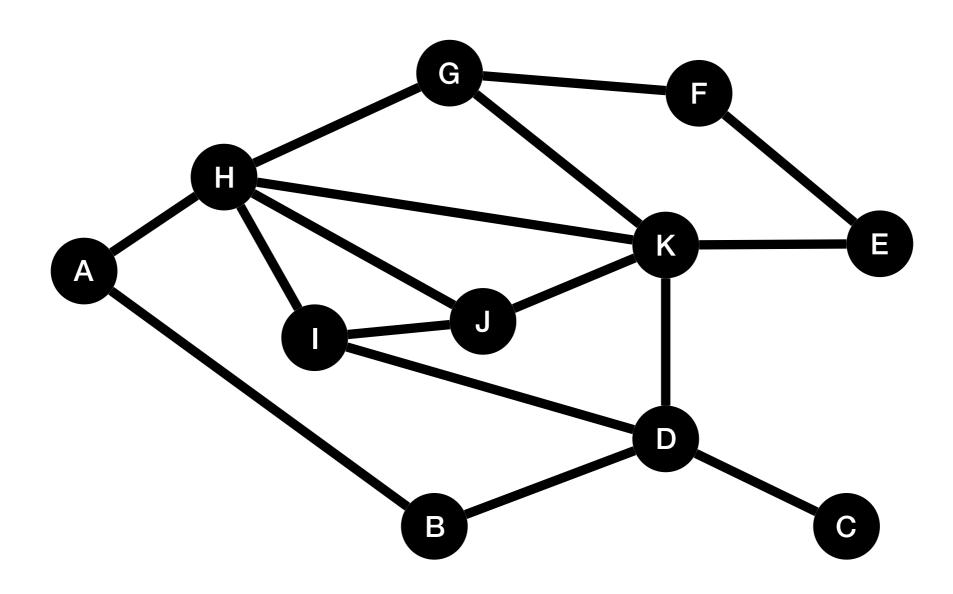


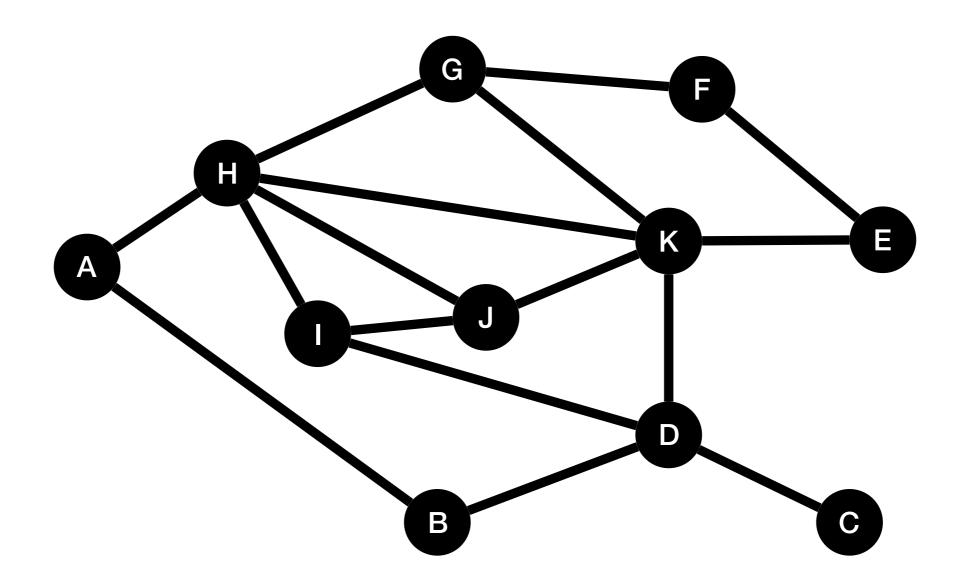




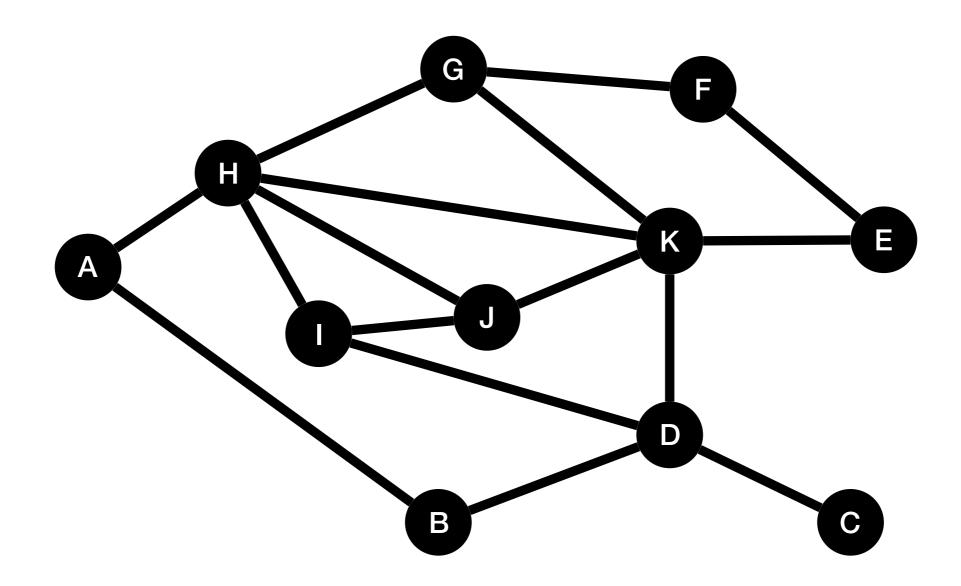






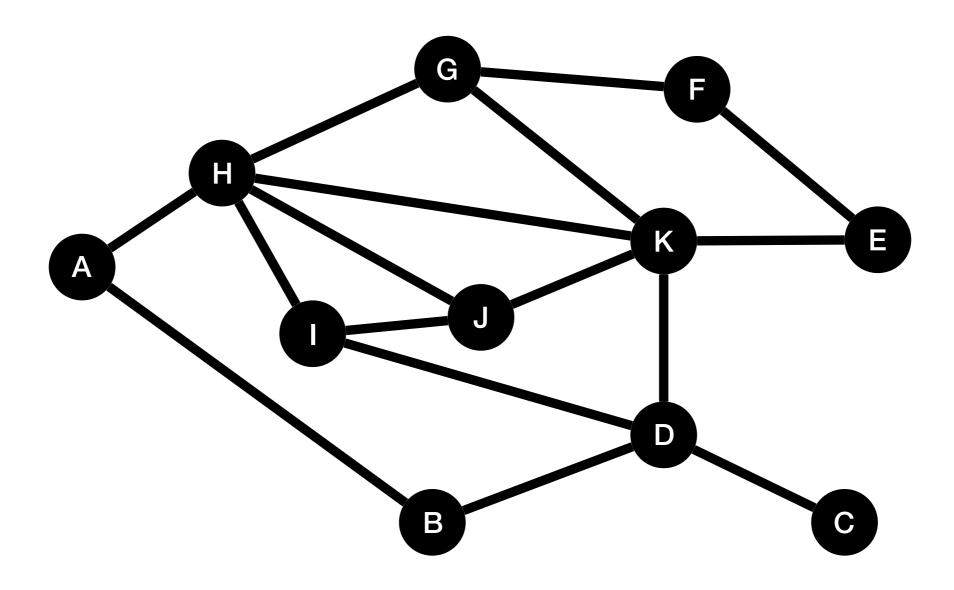


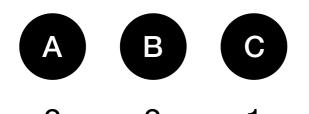


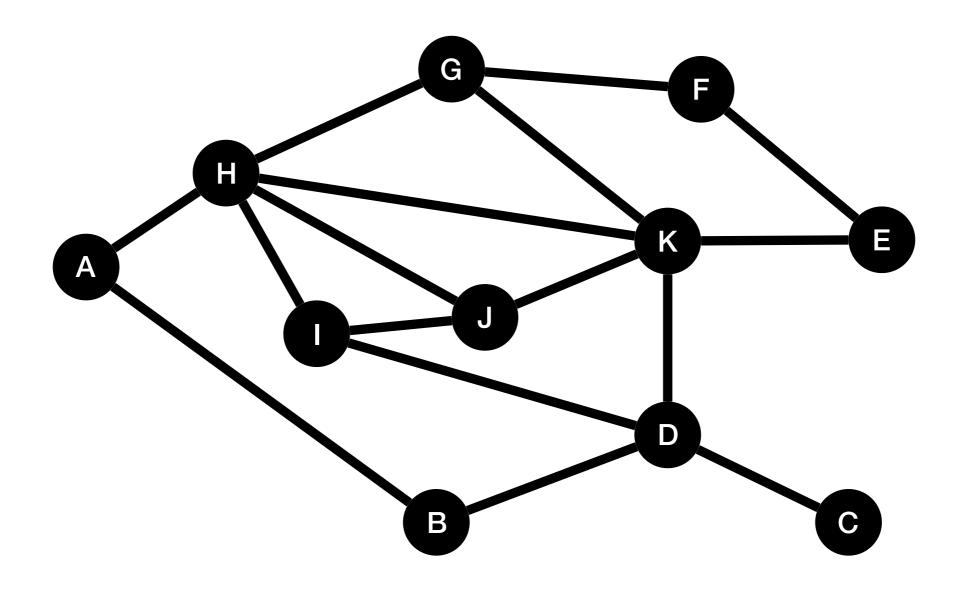


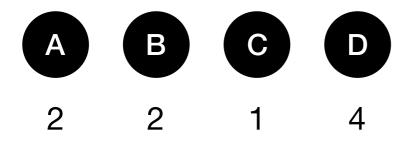


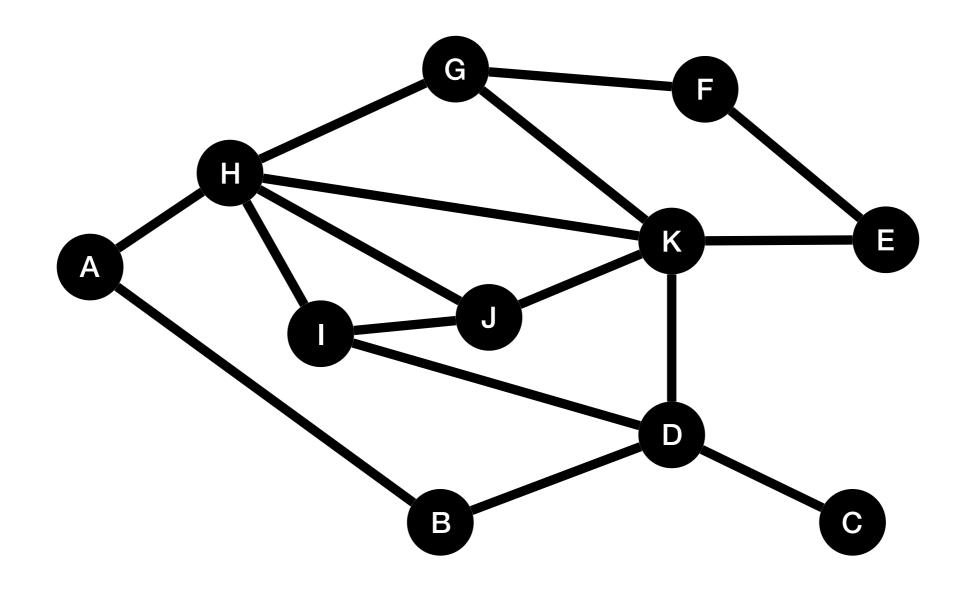
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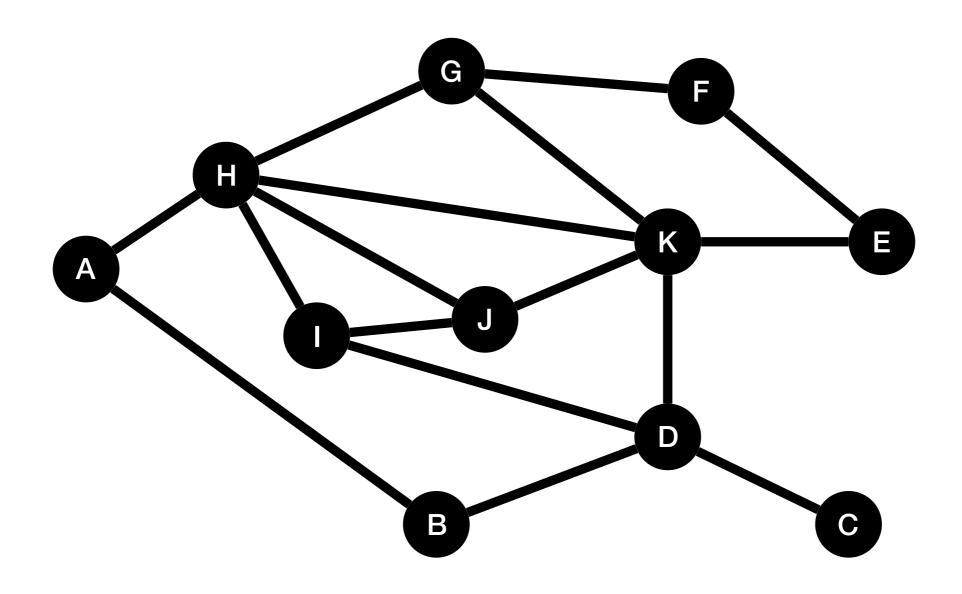






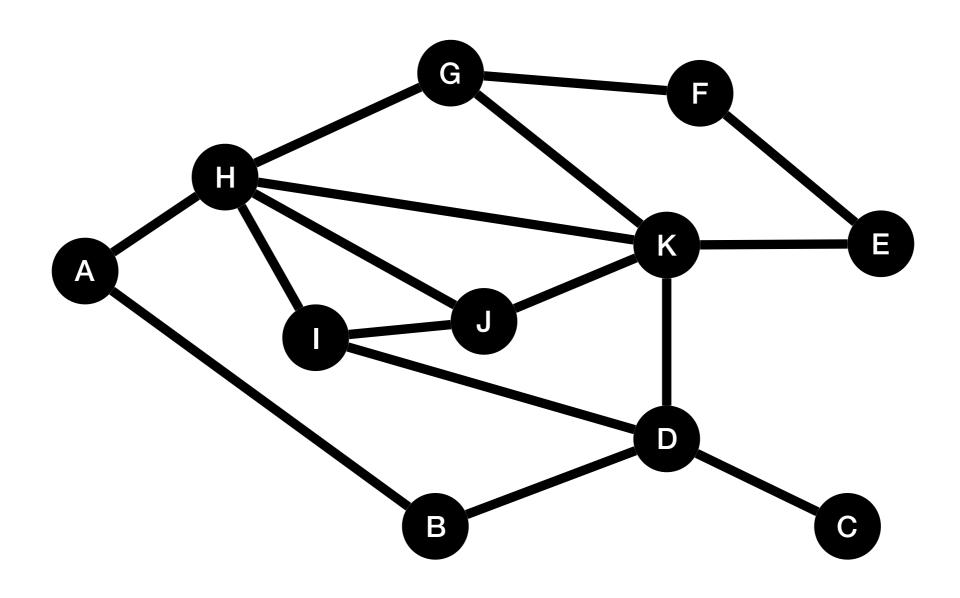


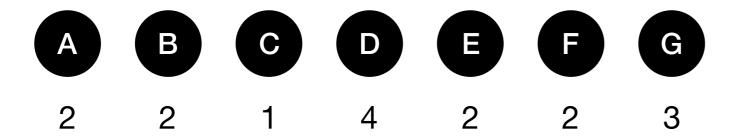
A B C D E
2 2 1 4 2

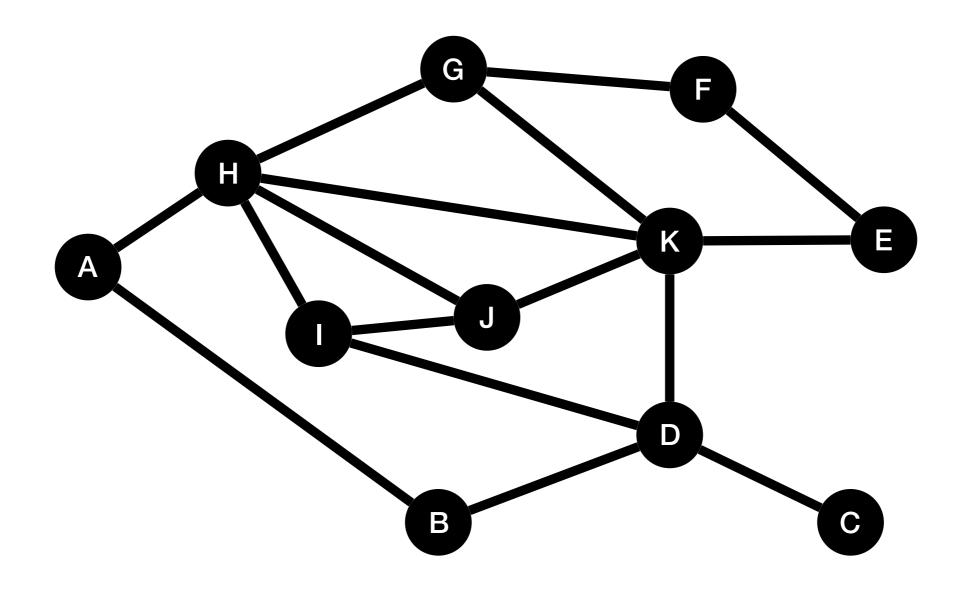


A B C D E F

2 2 1 4 2 2

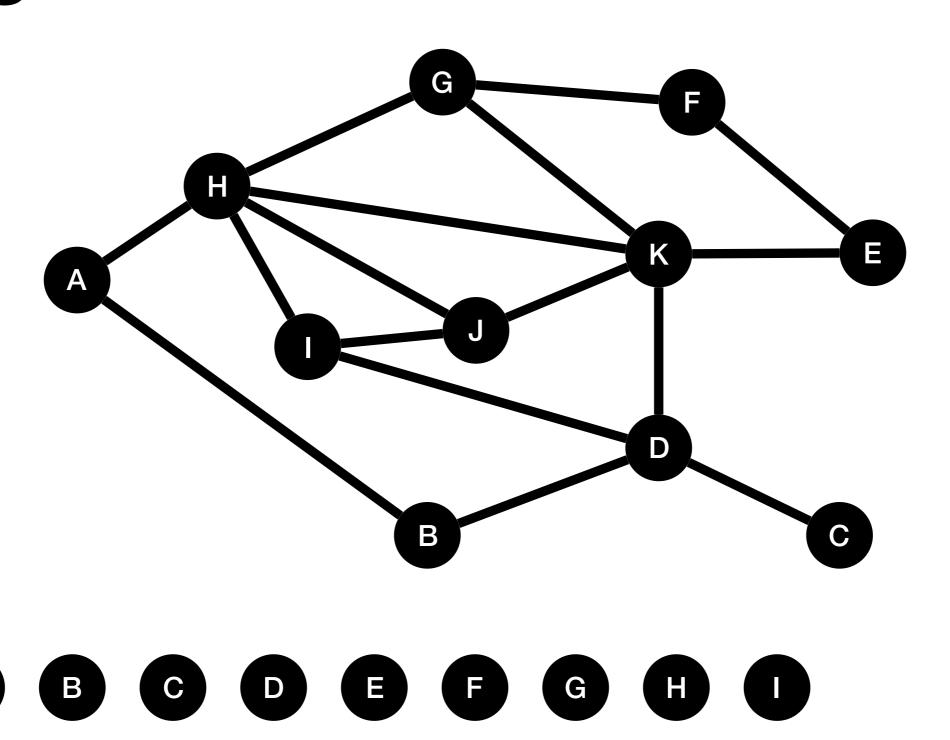


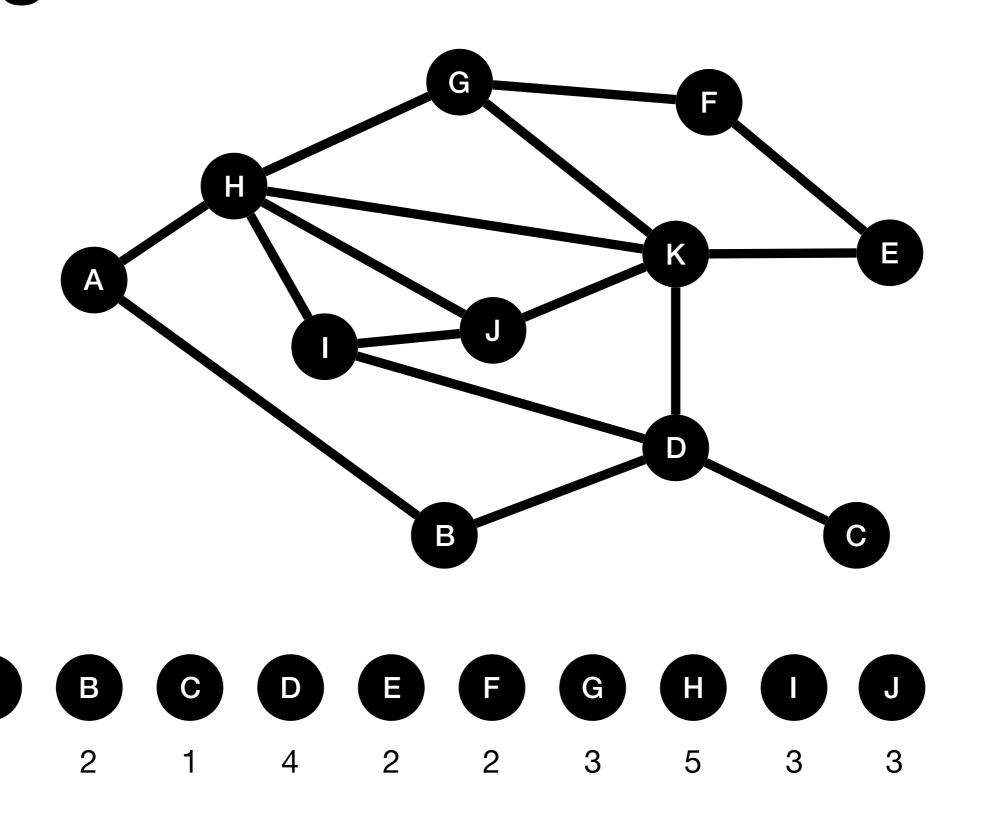


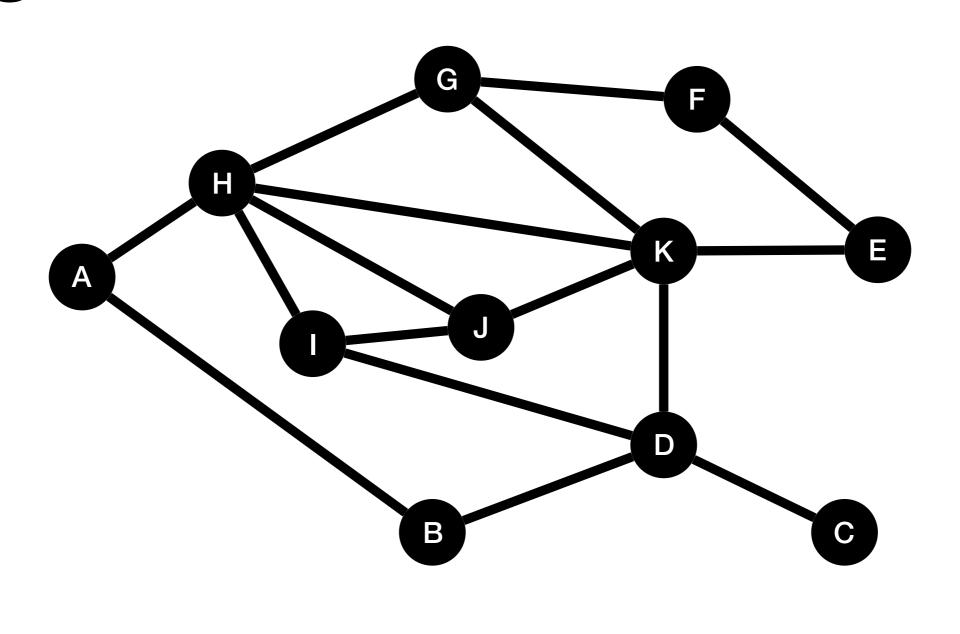


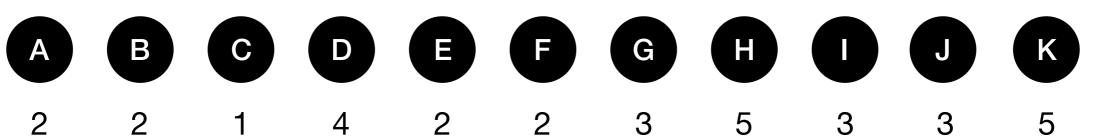
A B C D E F G H

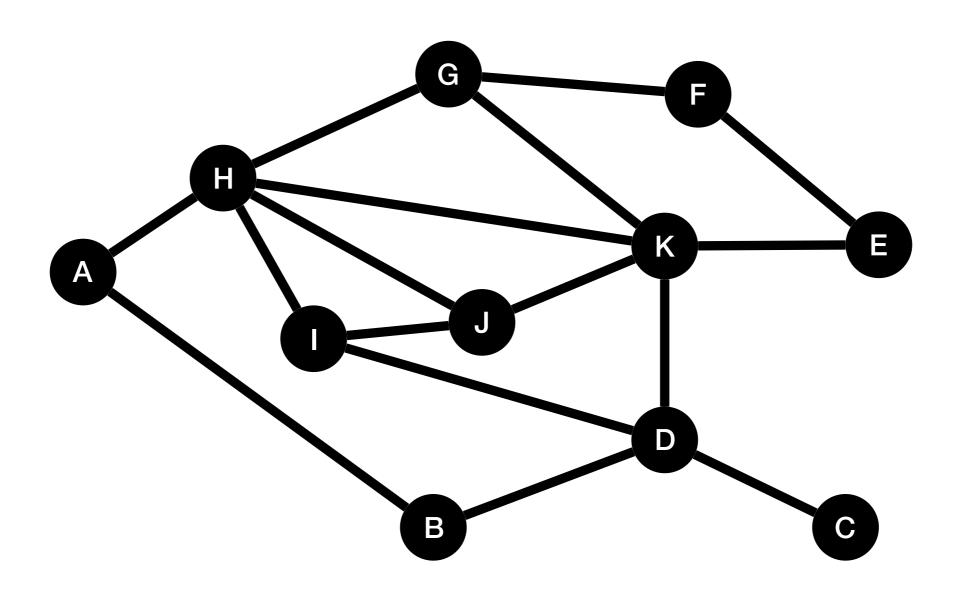
2 2 1 4 2 2 3 5

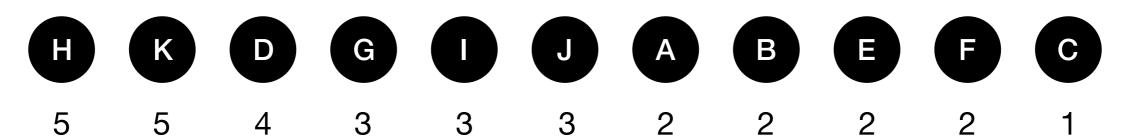


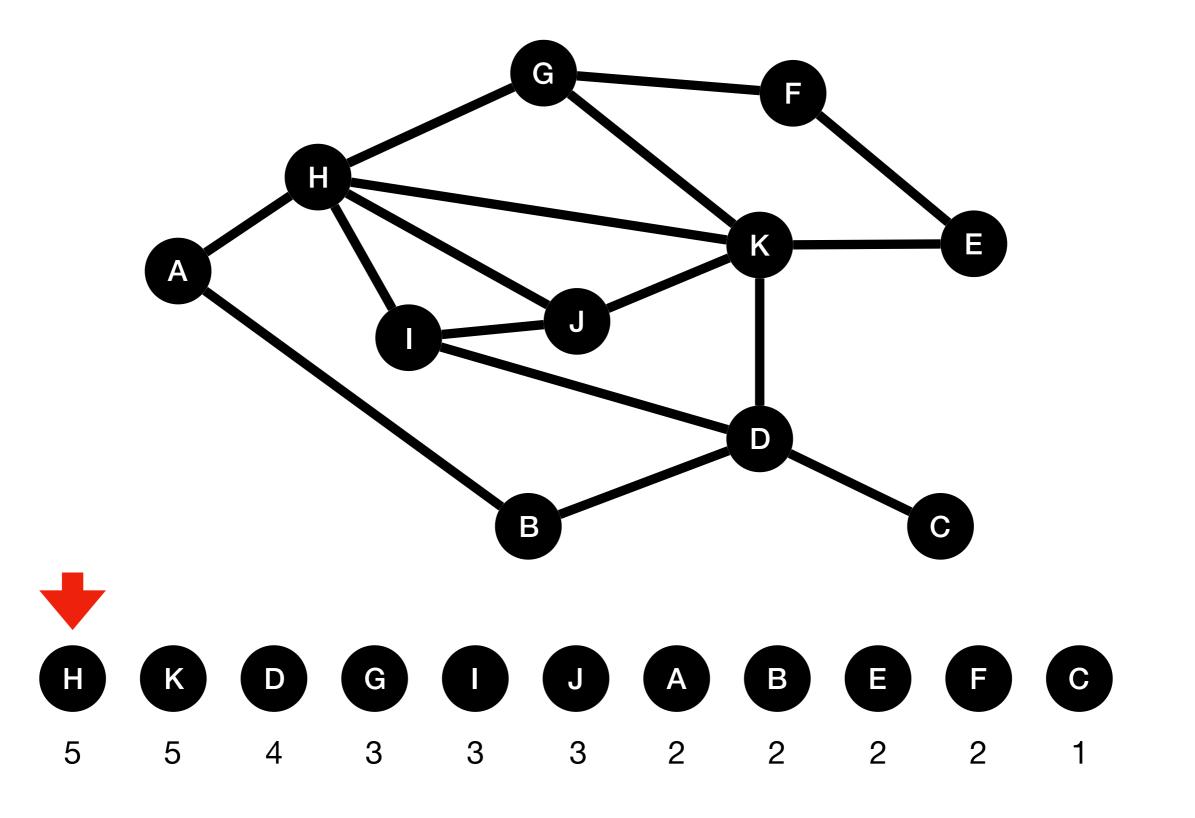


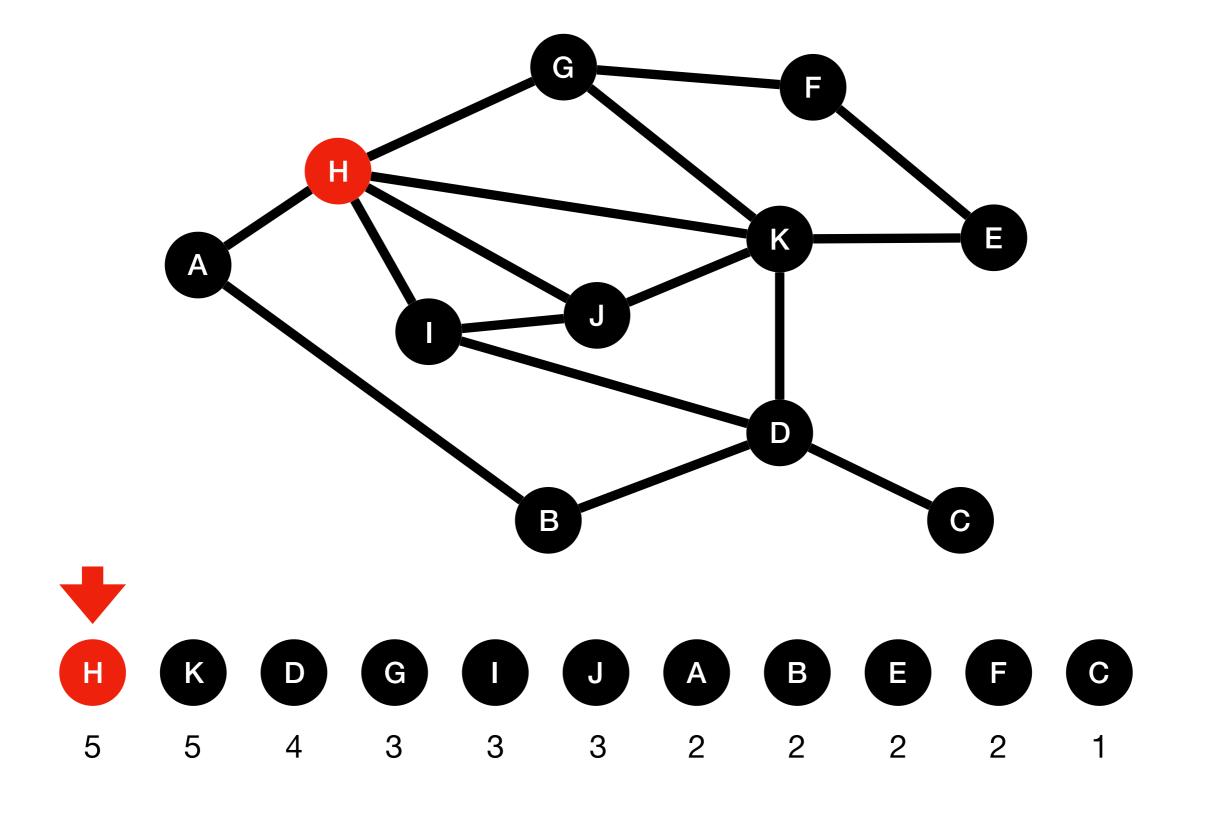


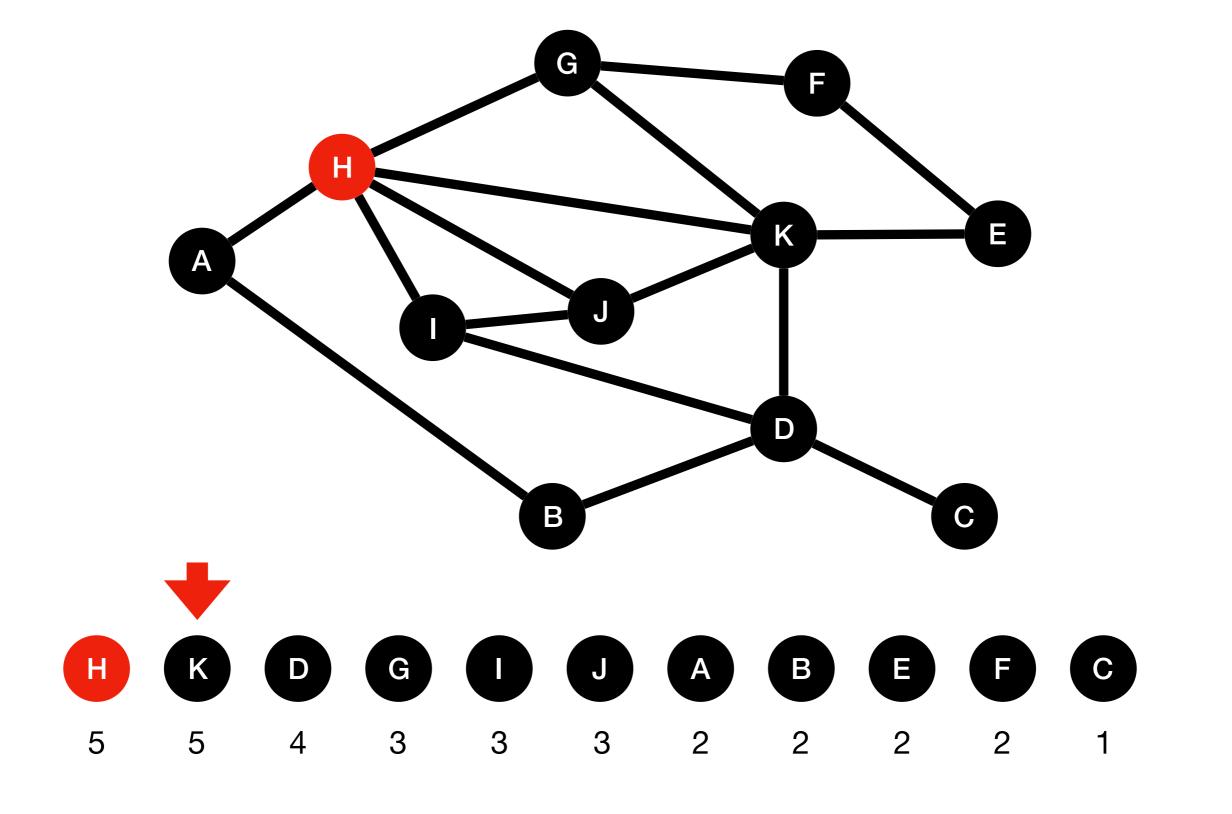


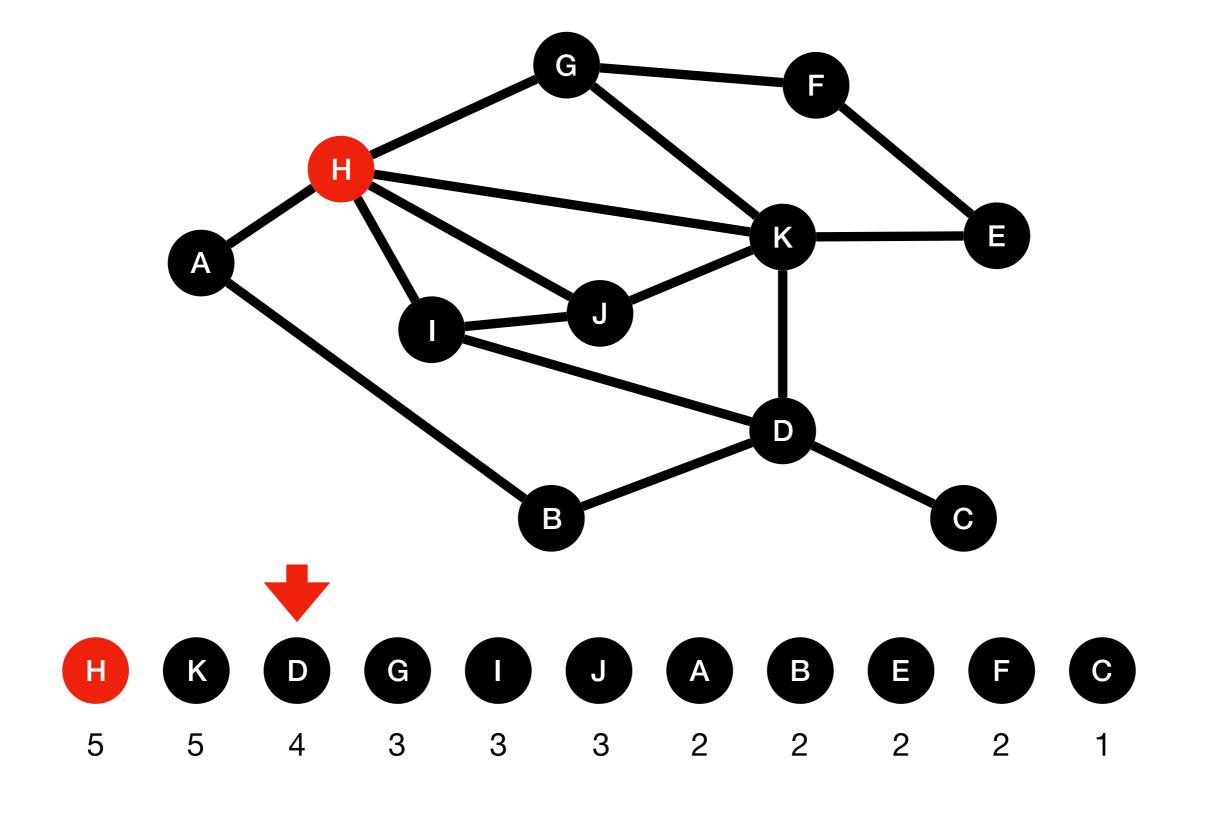


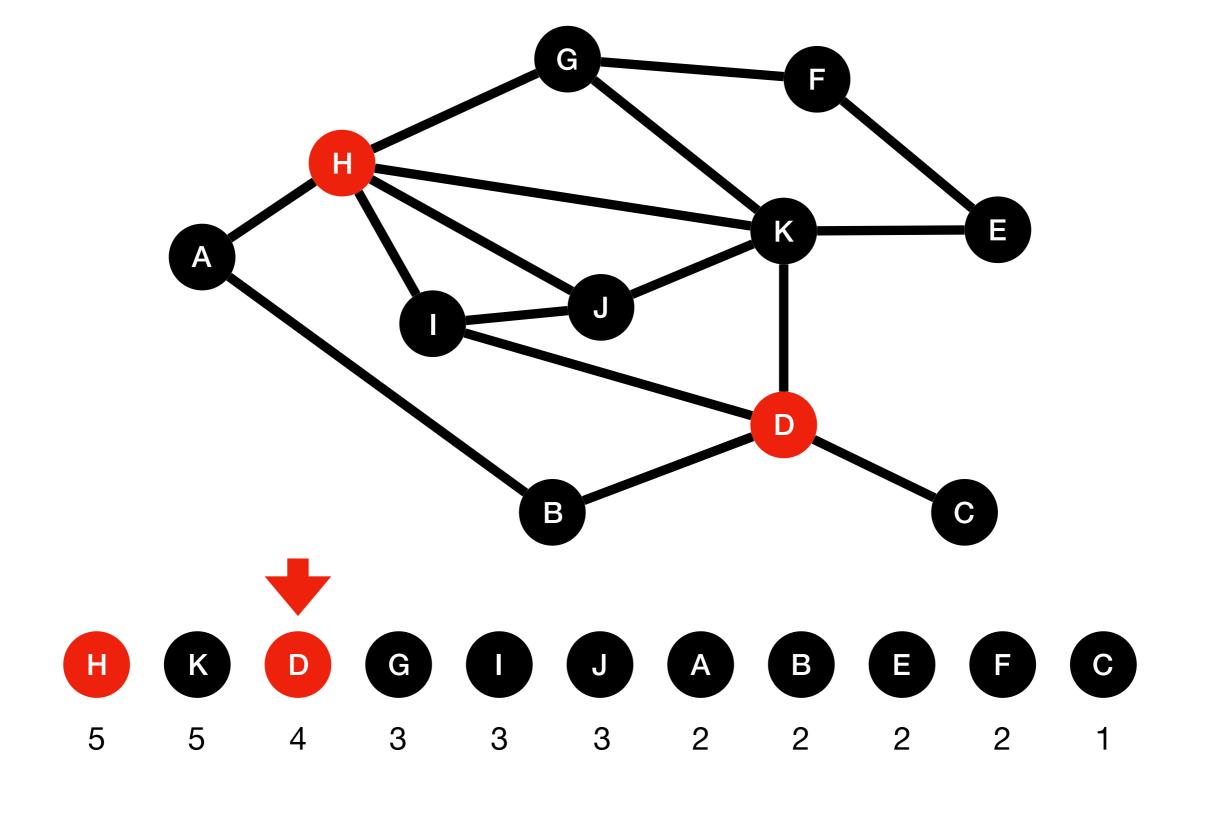


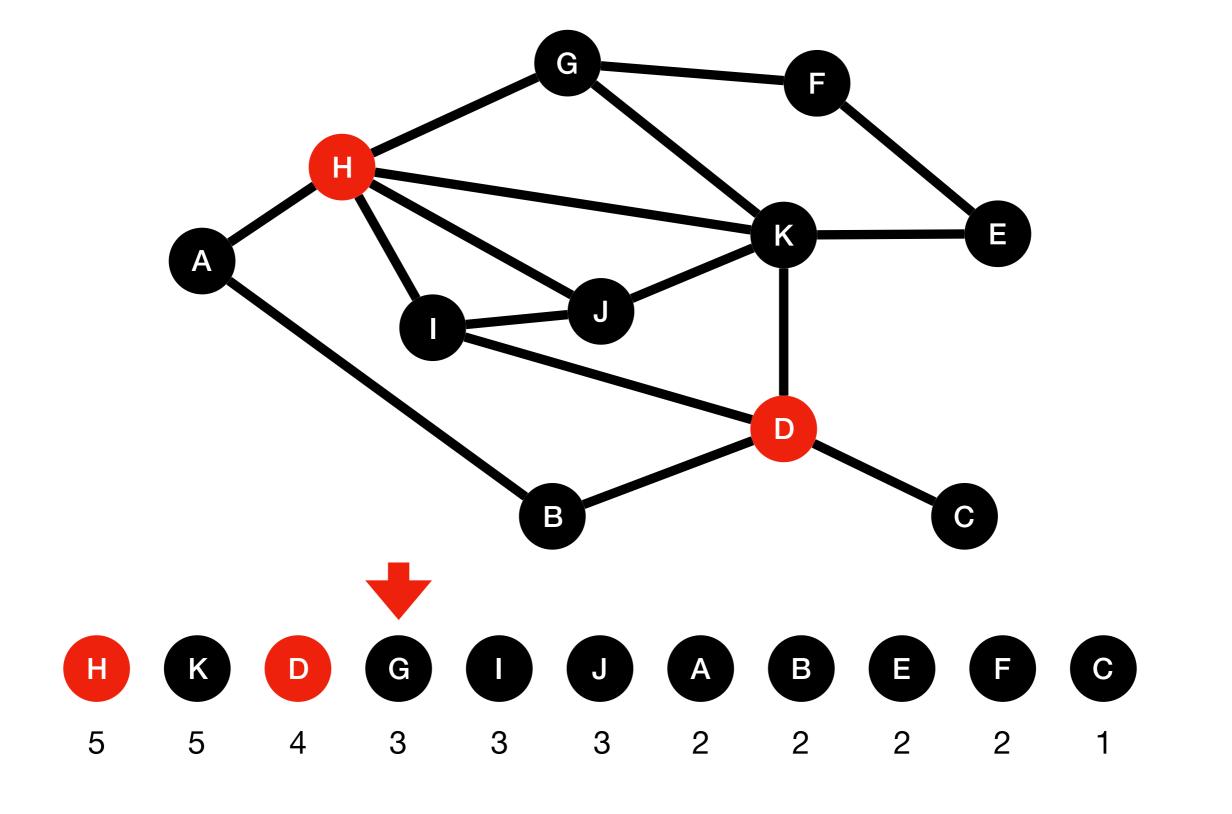


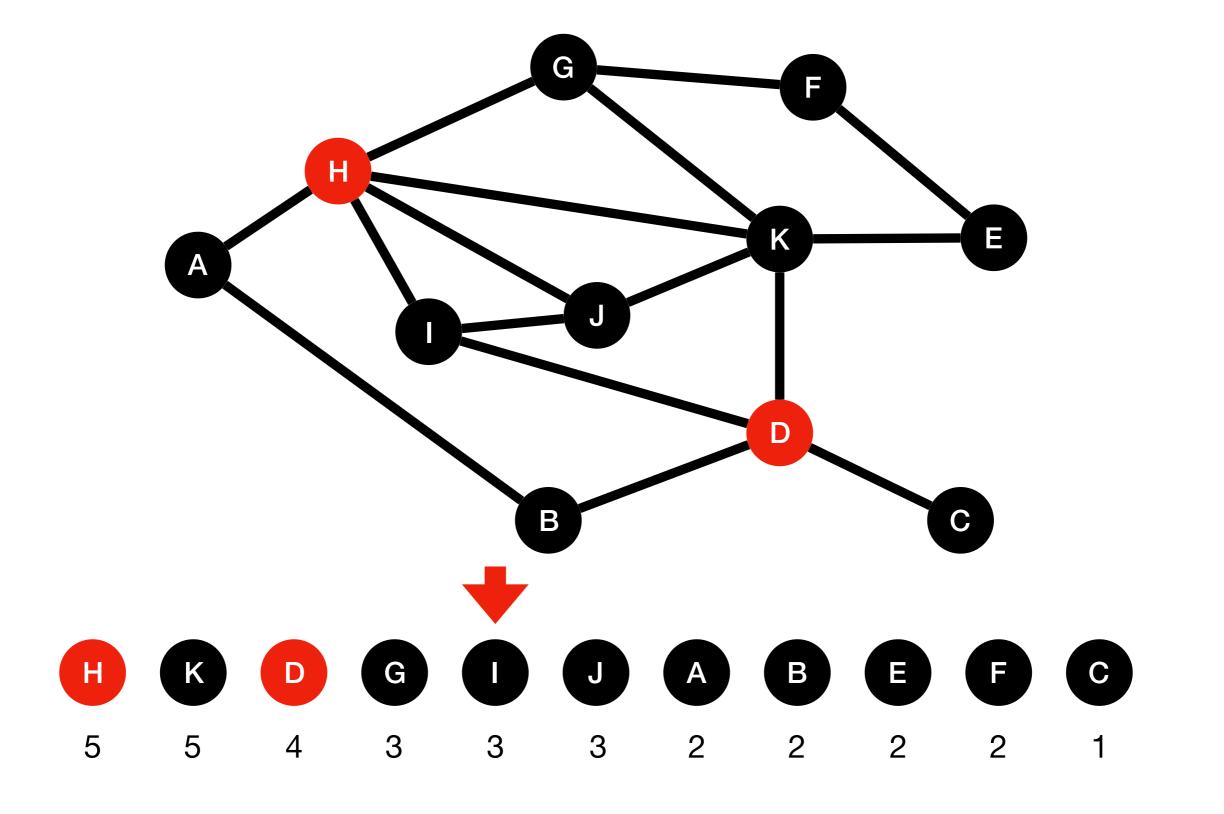


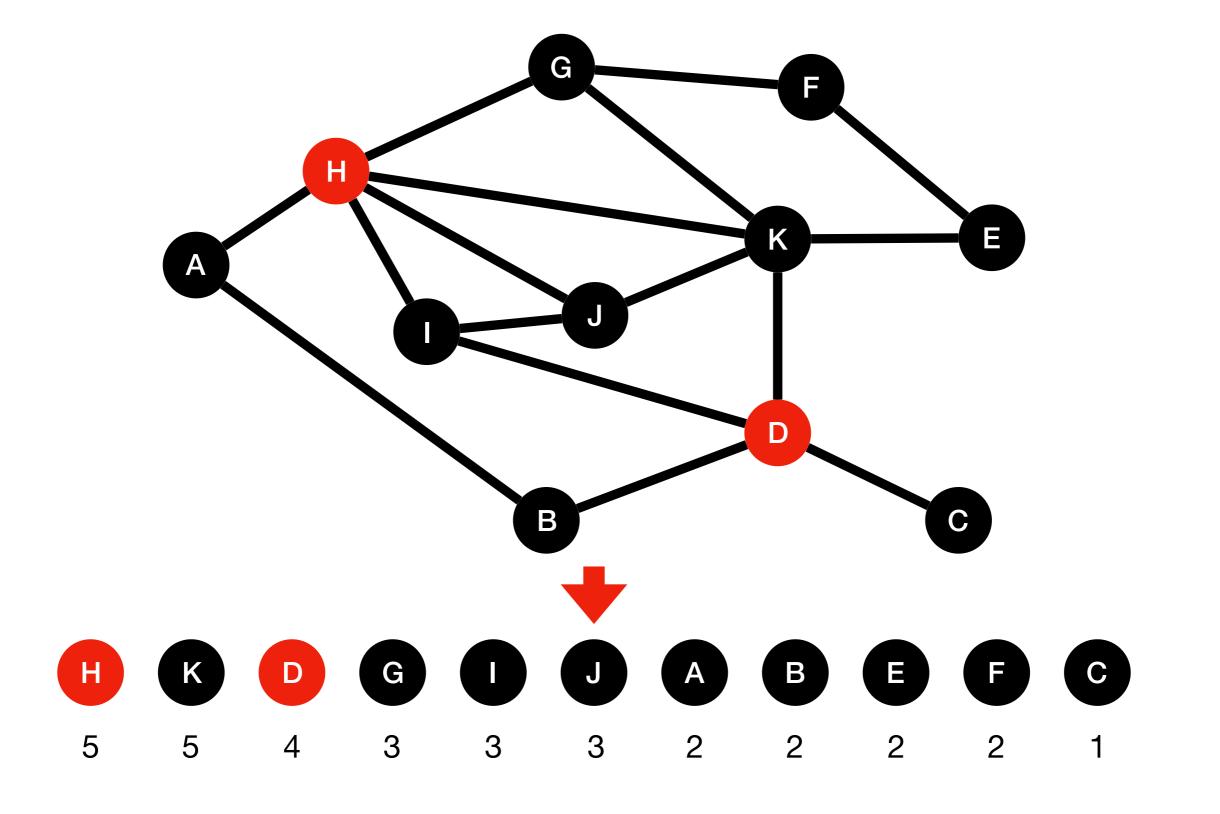


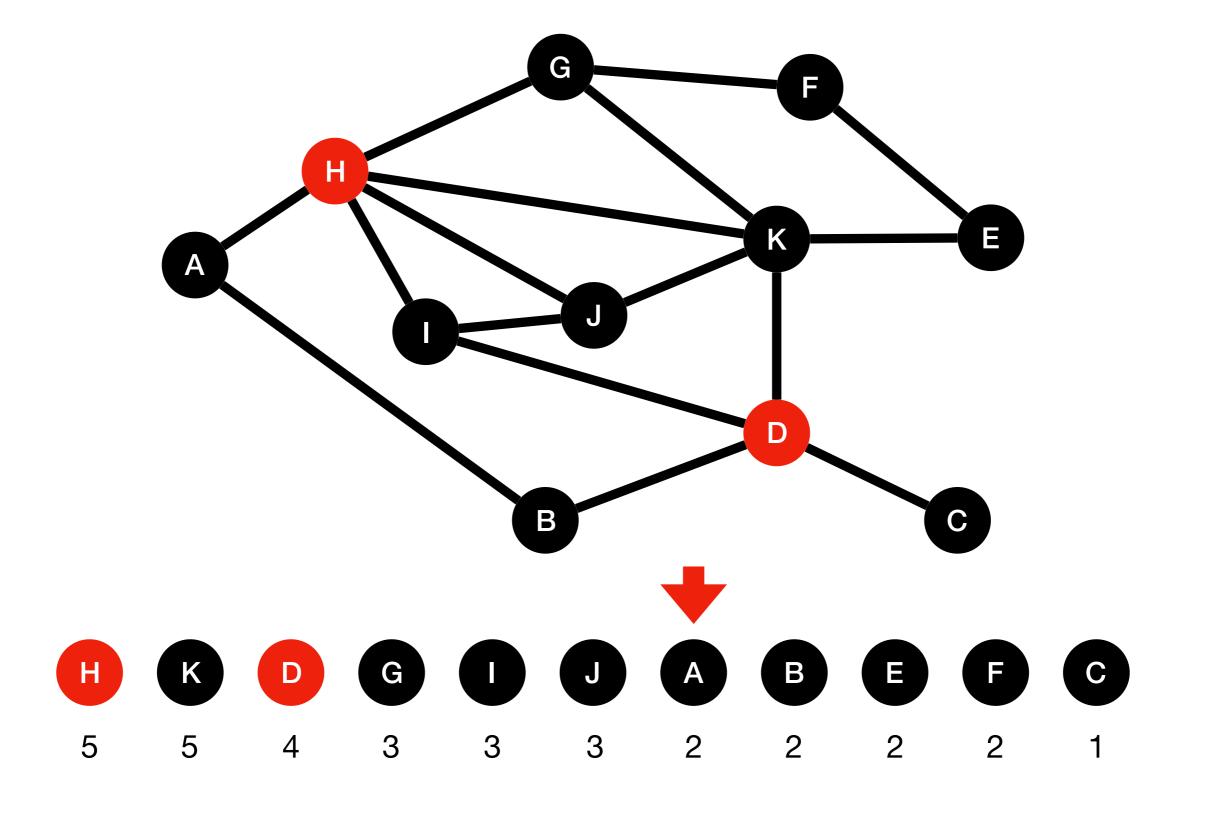


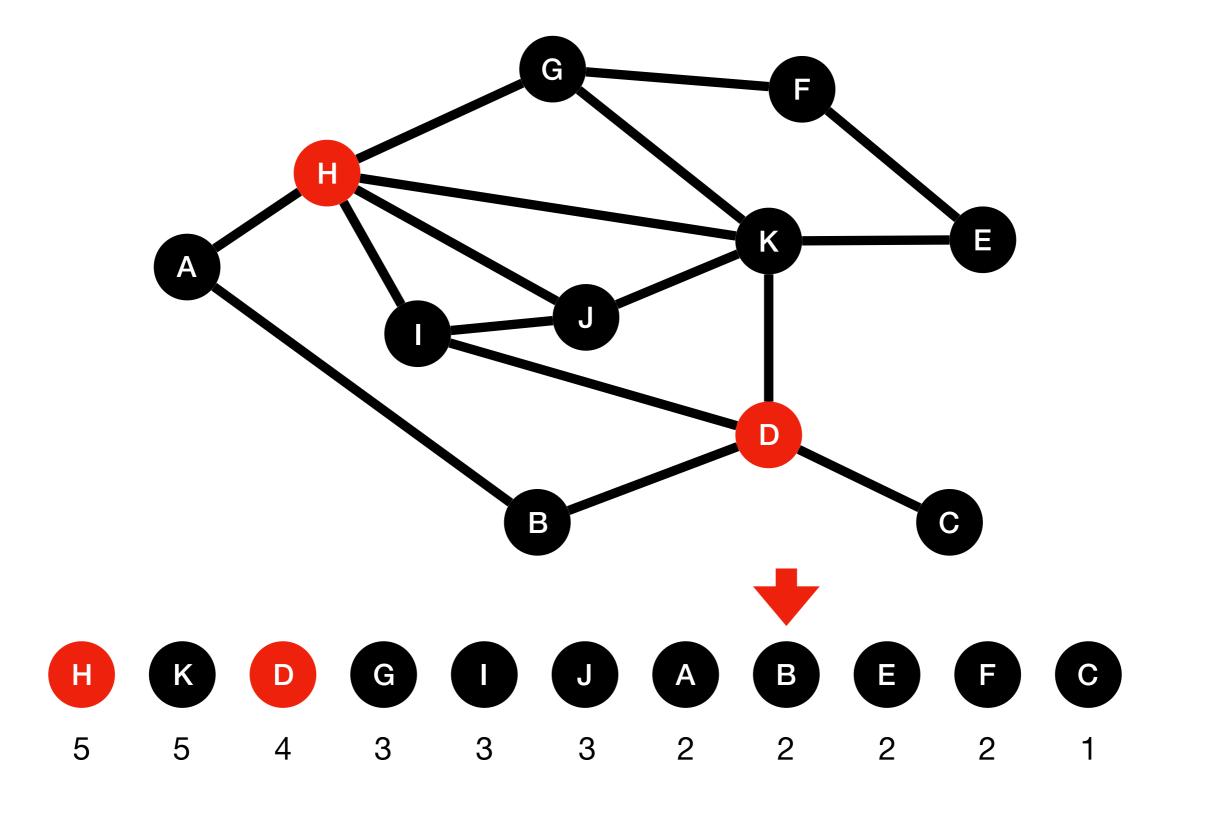


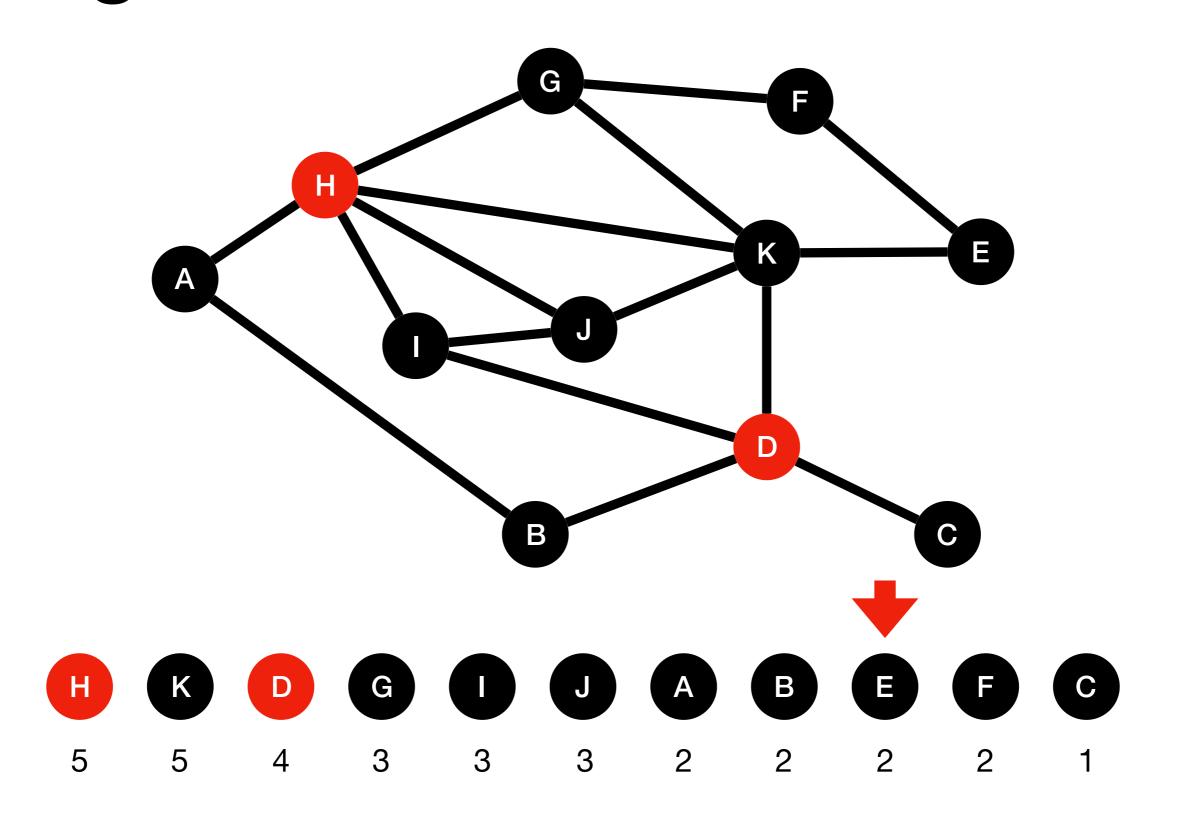


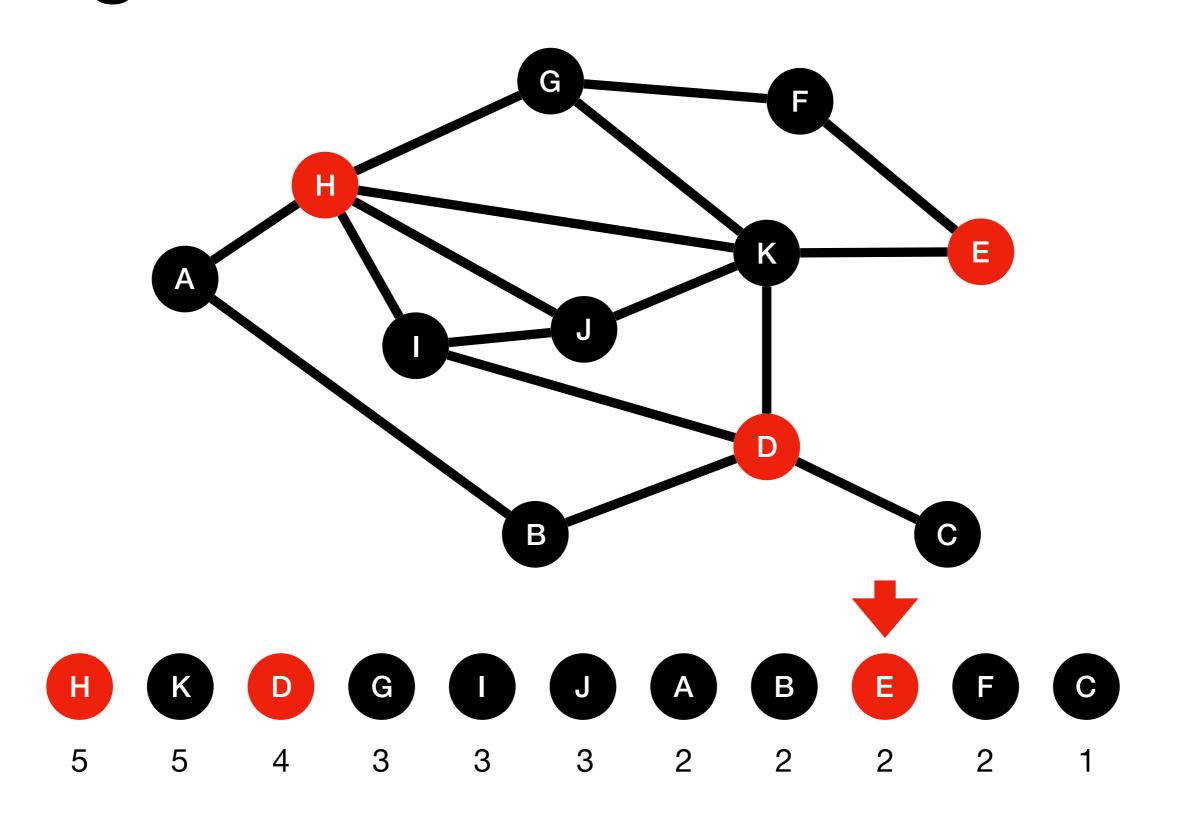


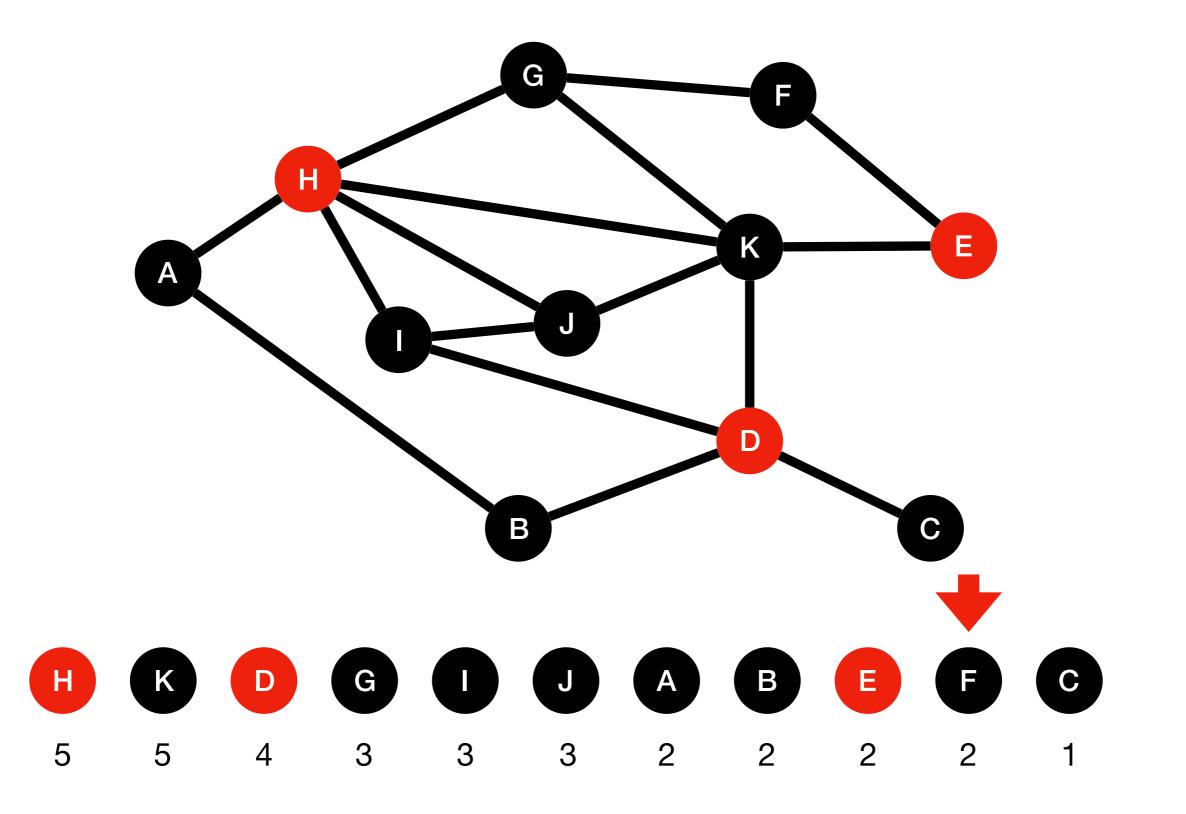


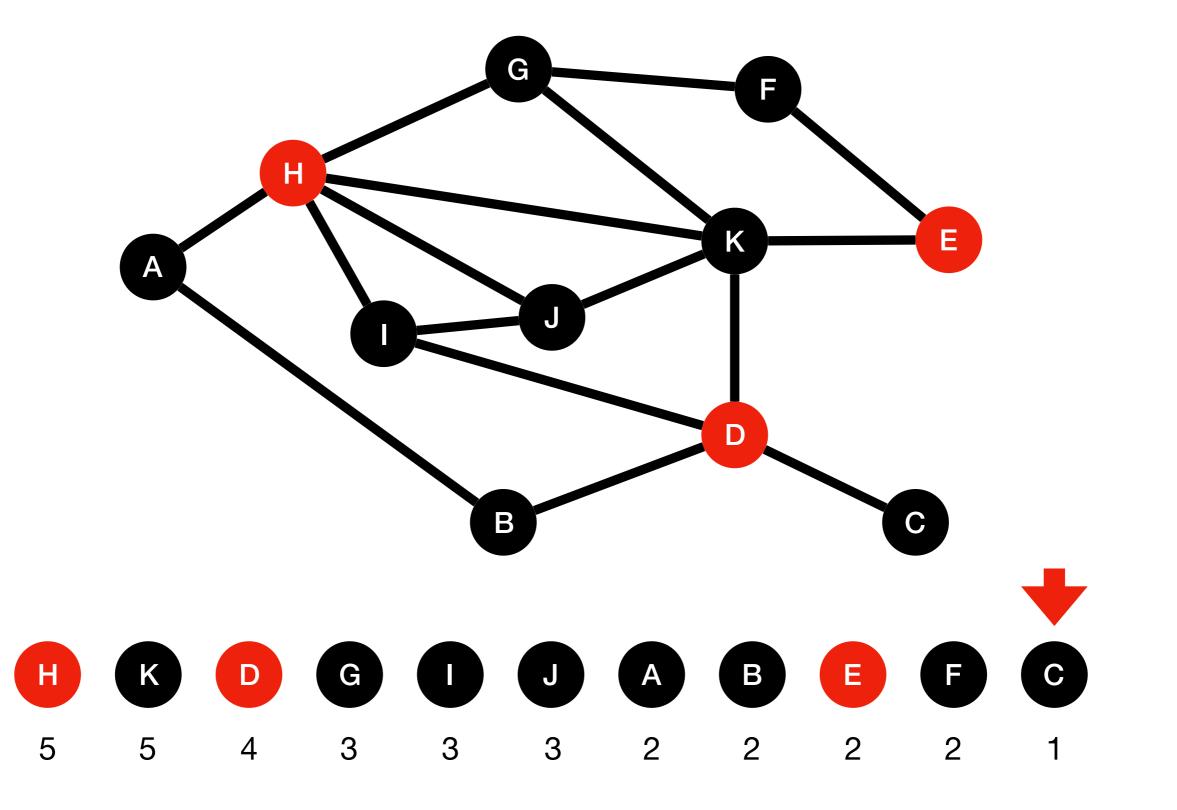


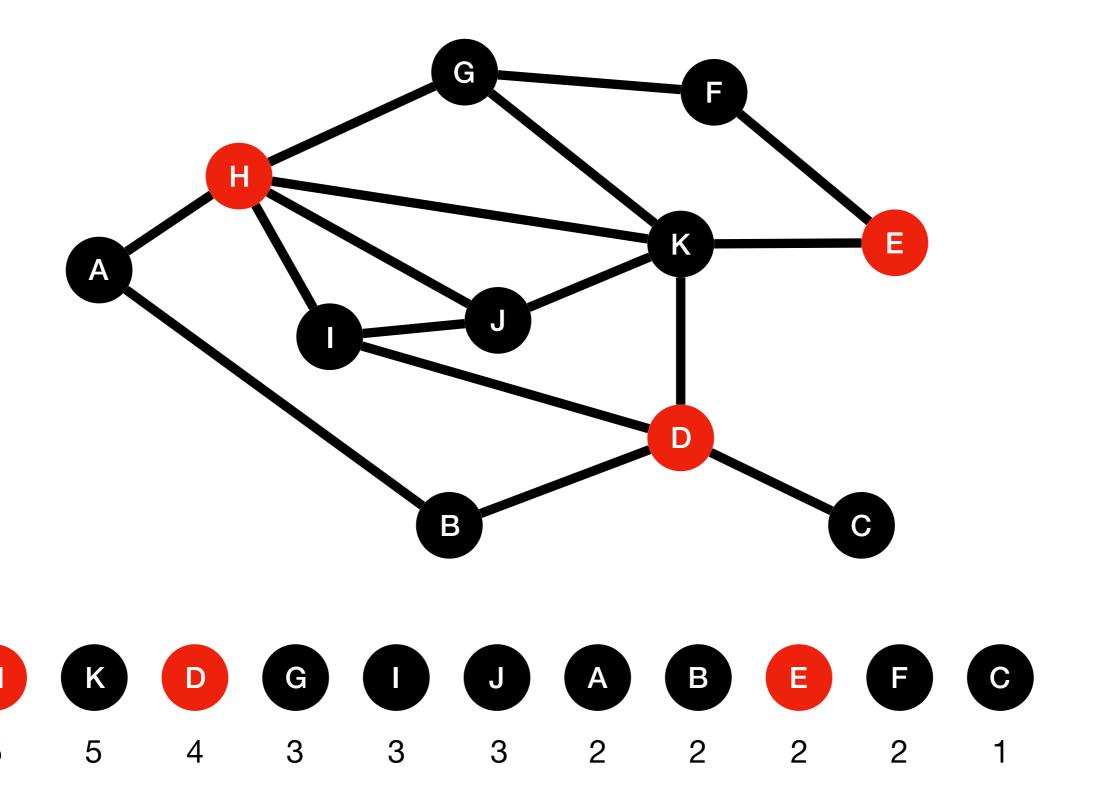


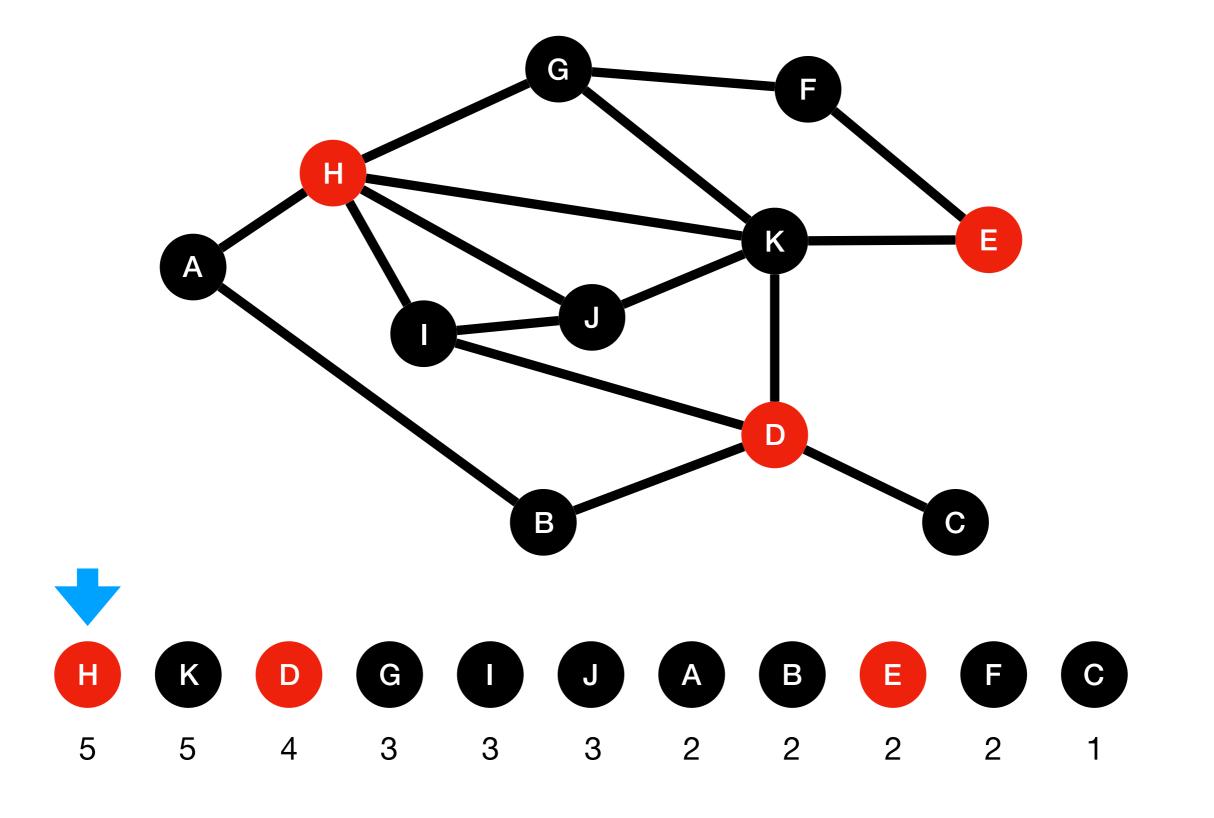


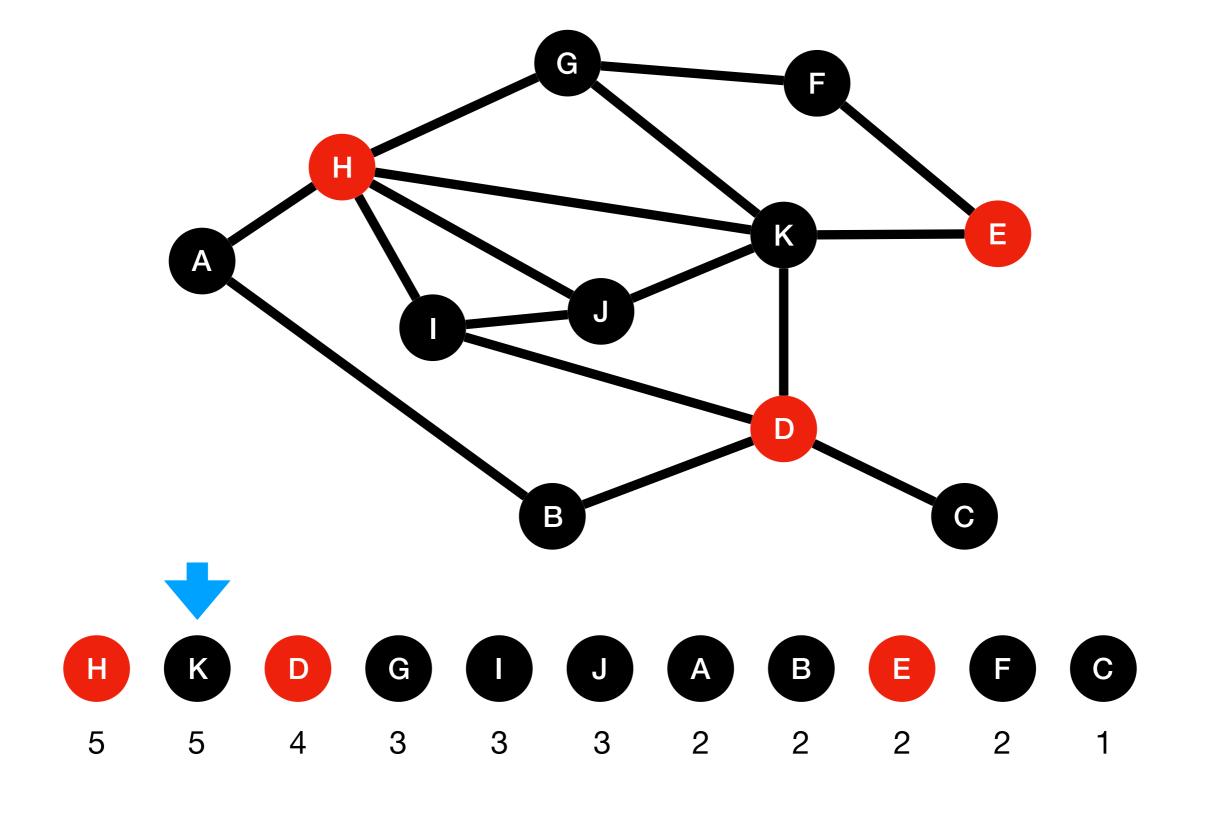


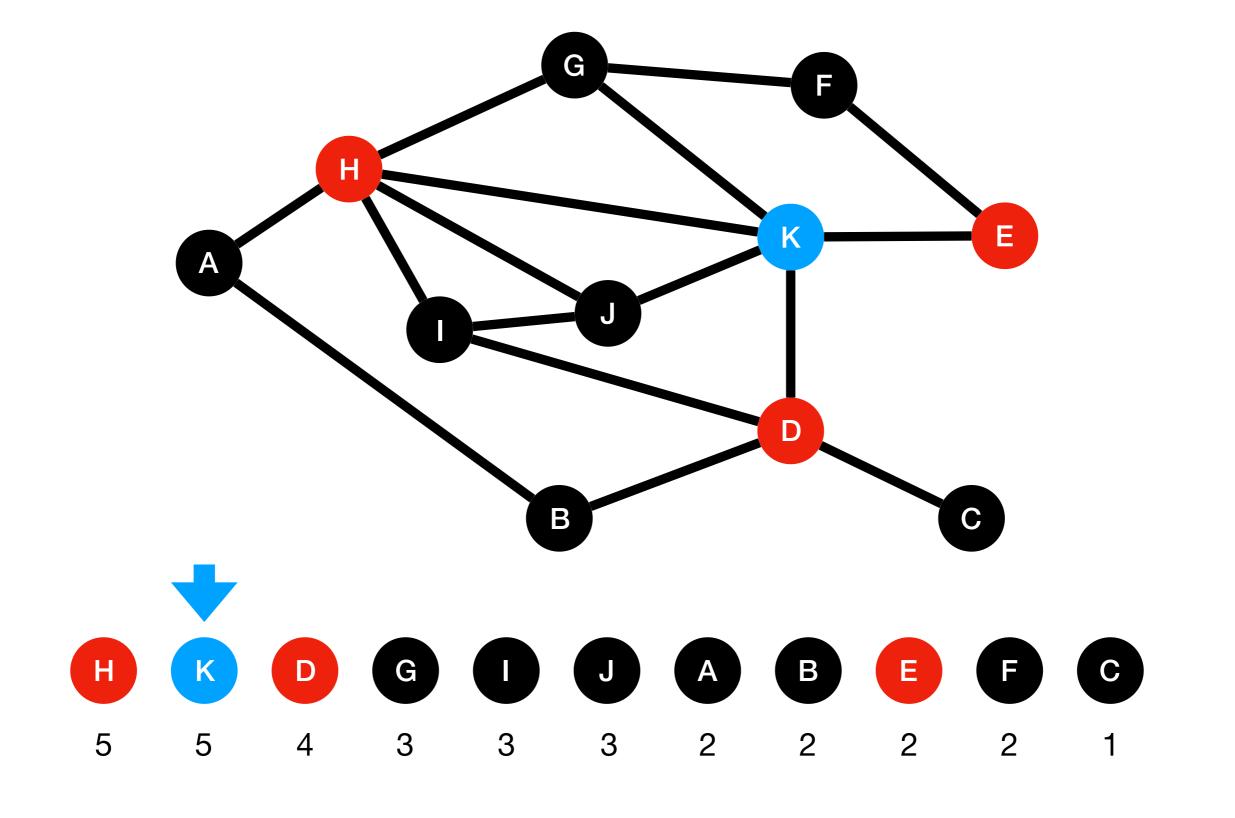


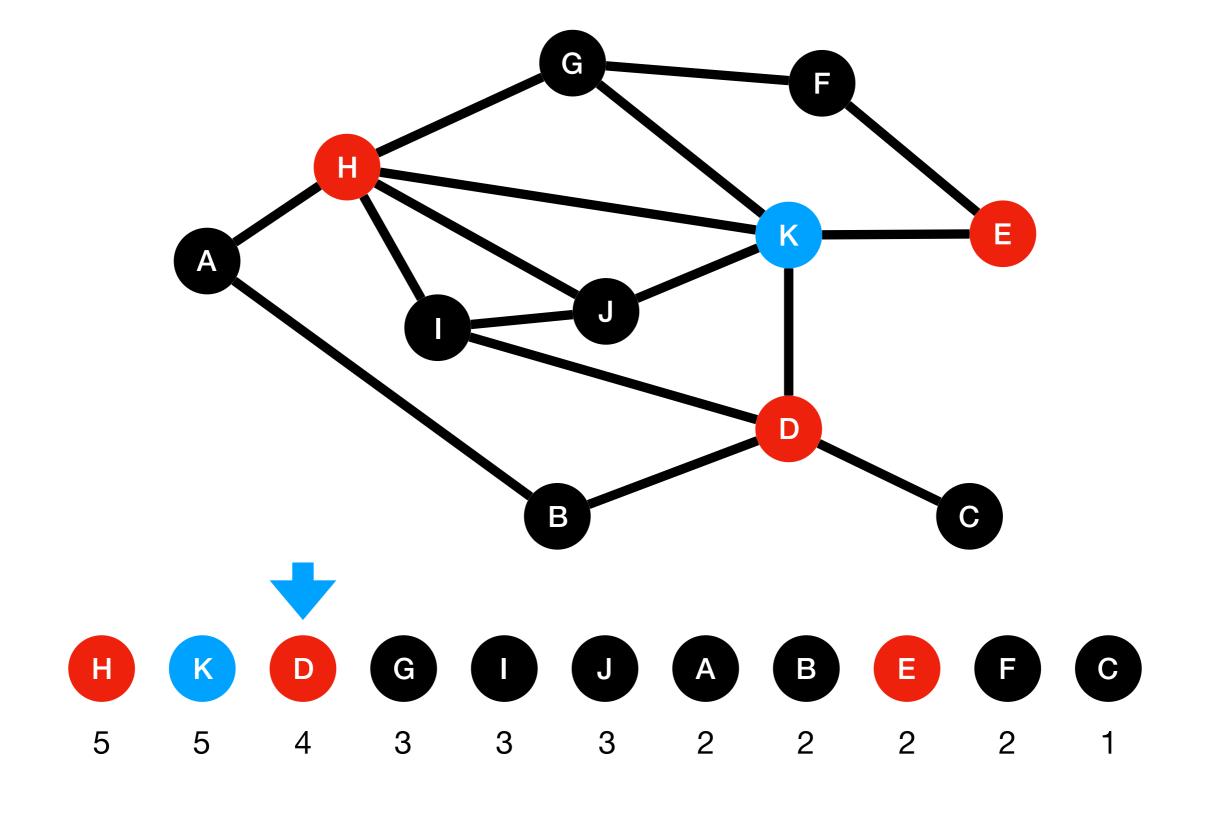


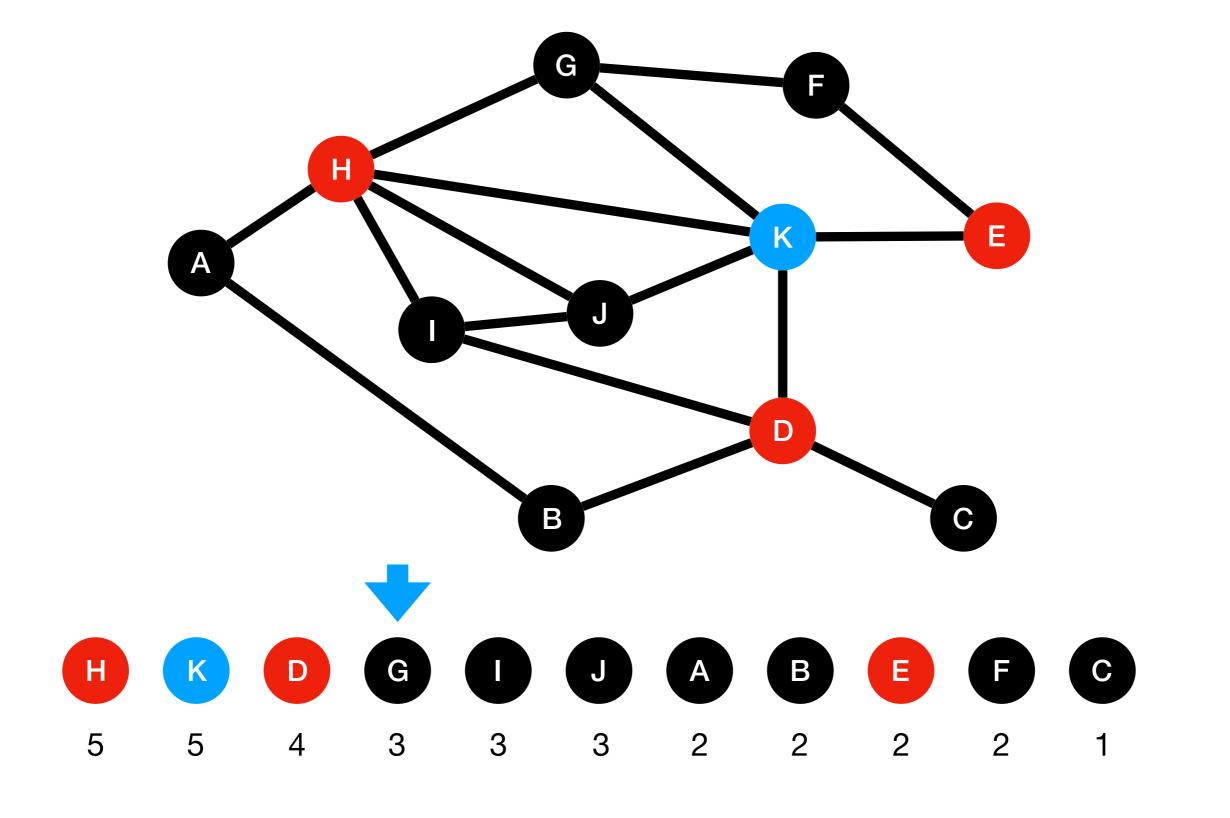


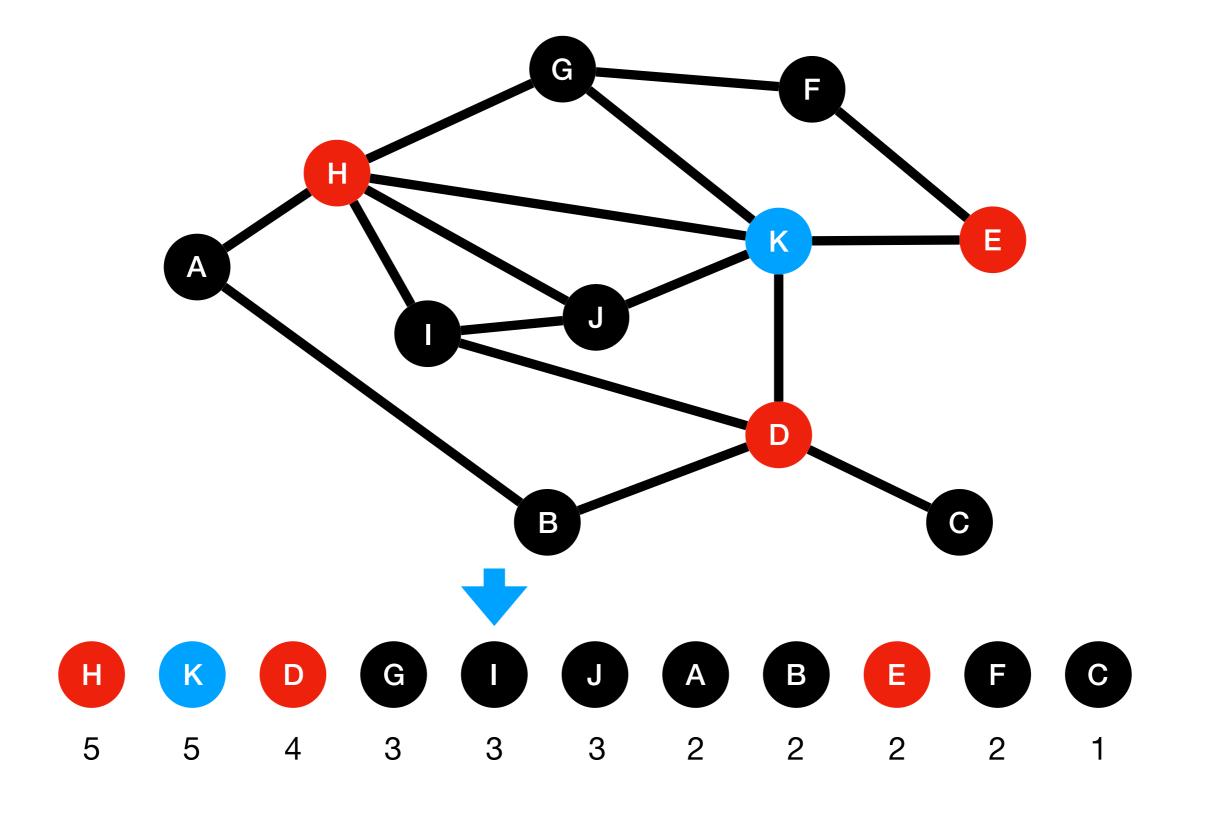


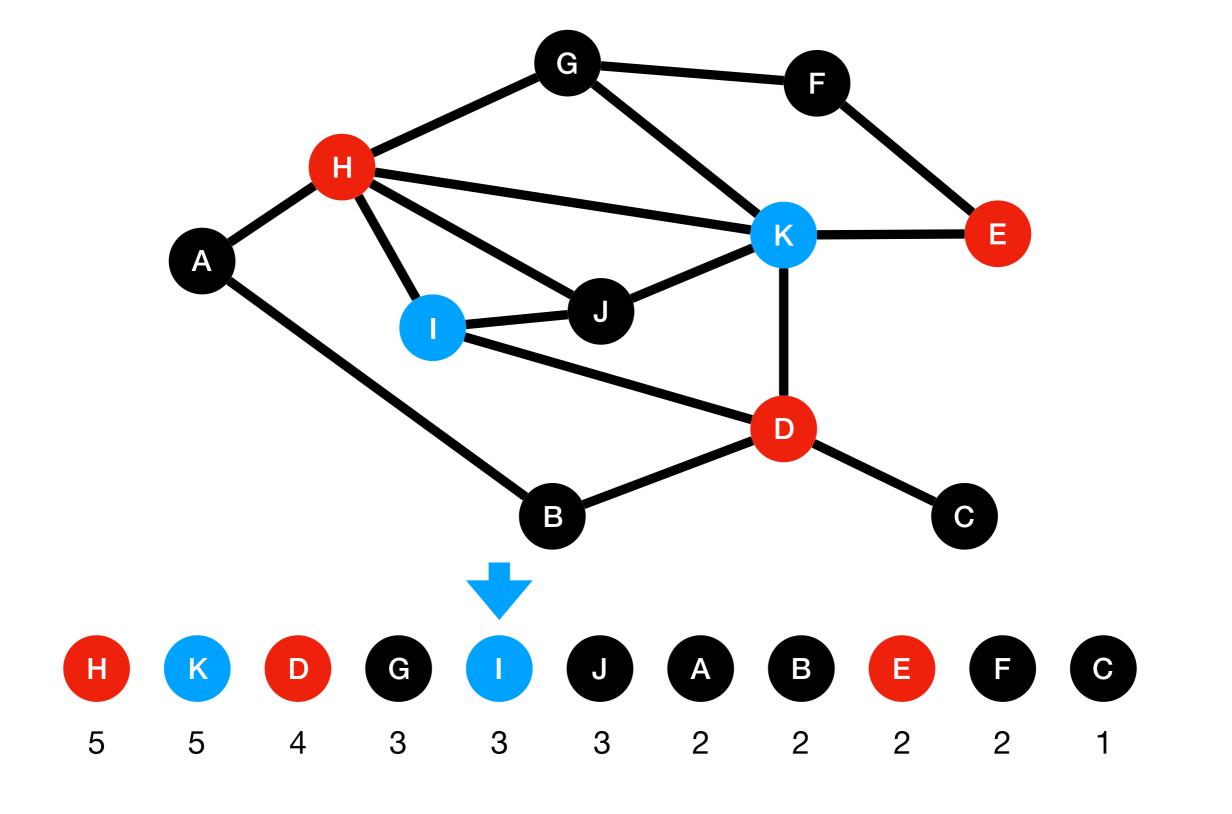


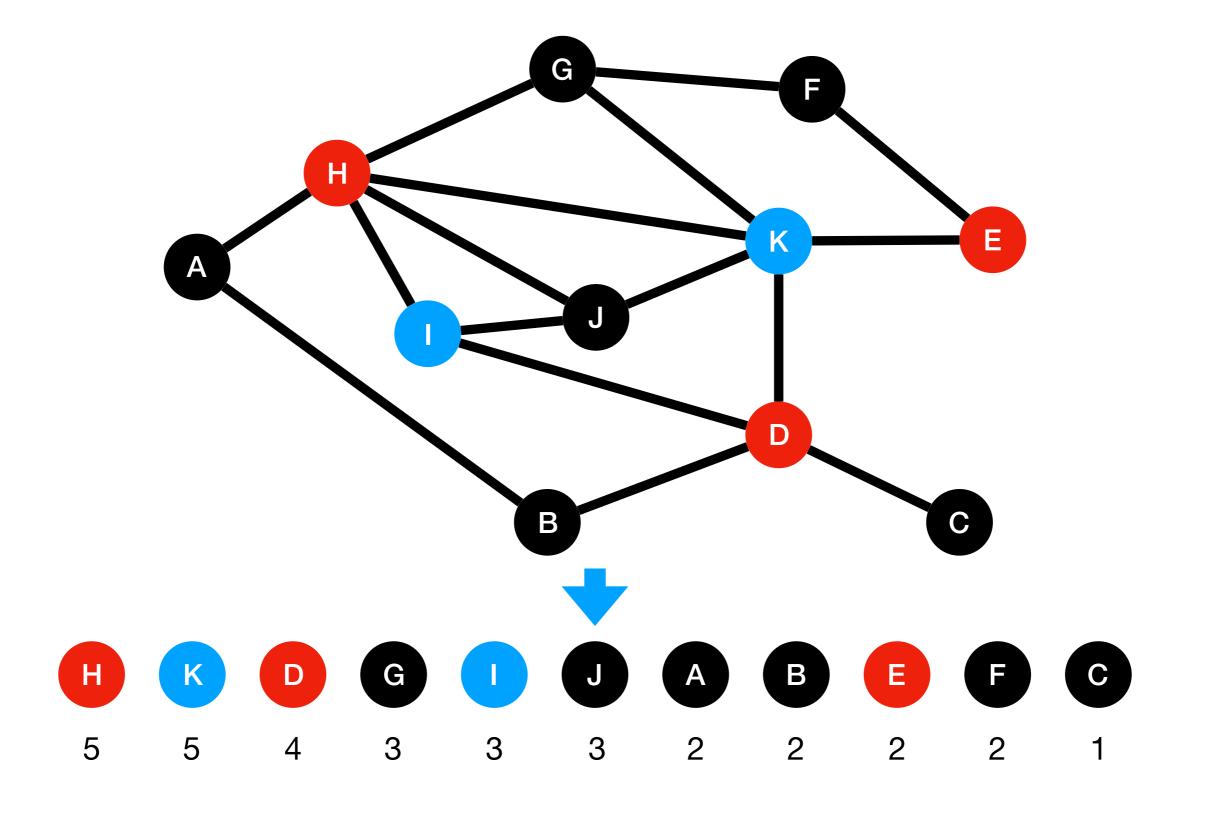


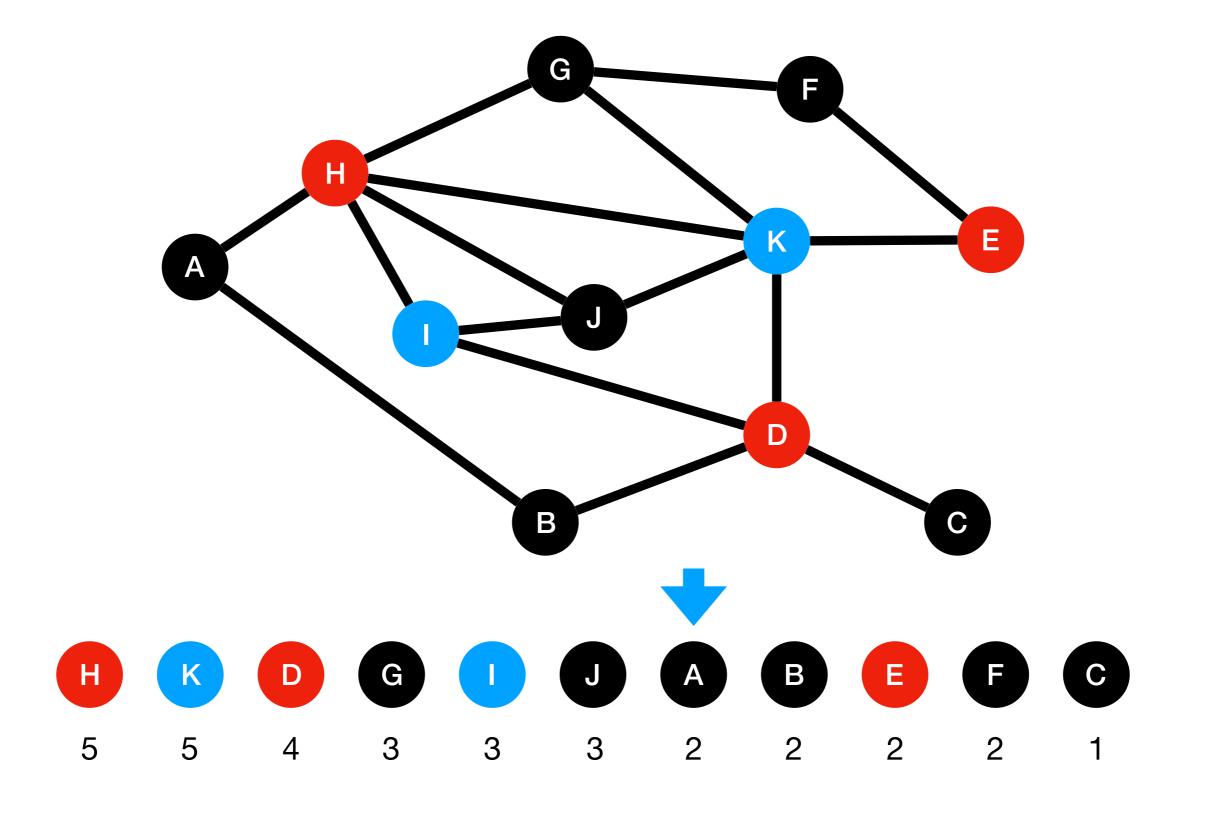


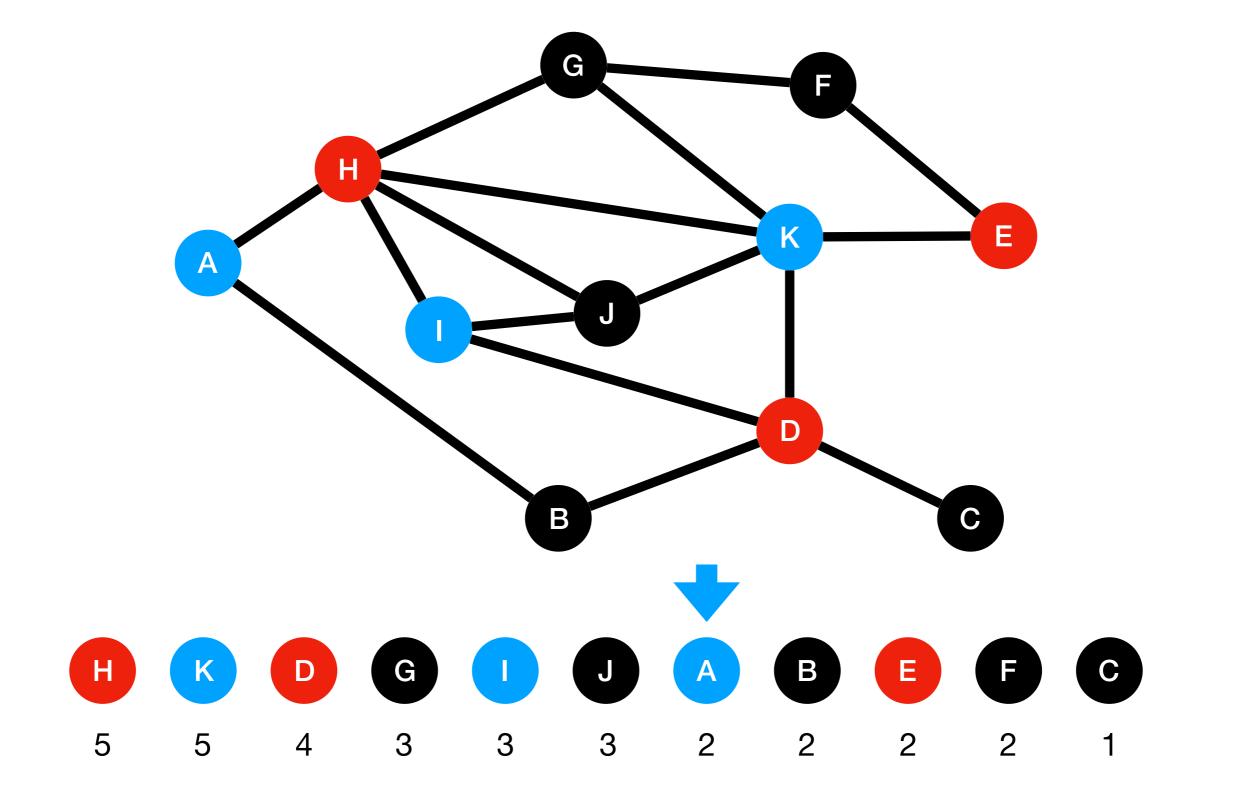


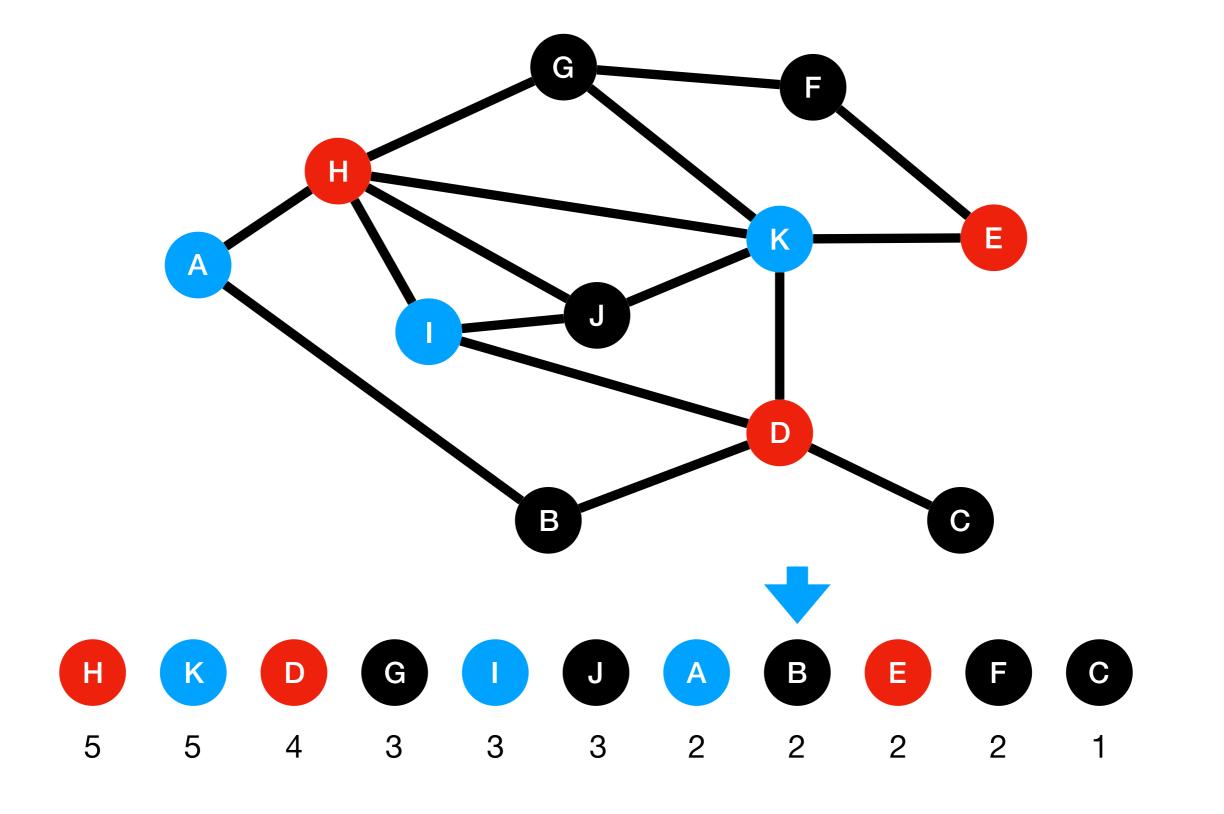


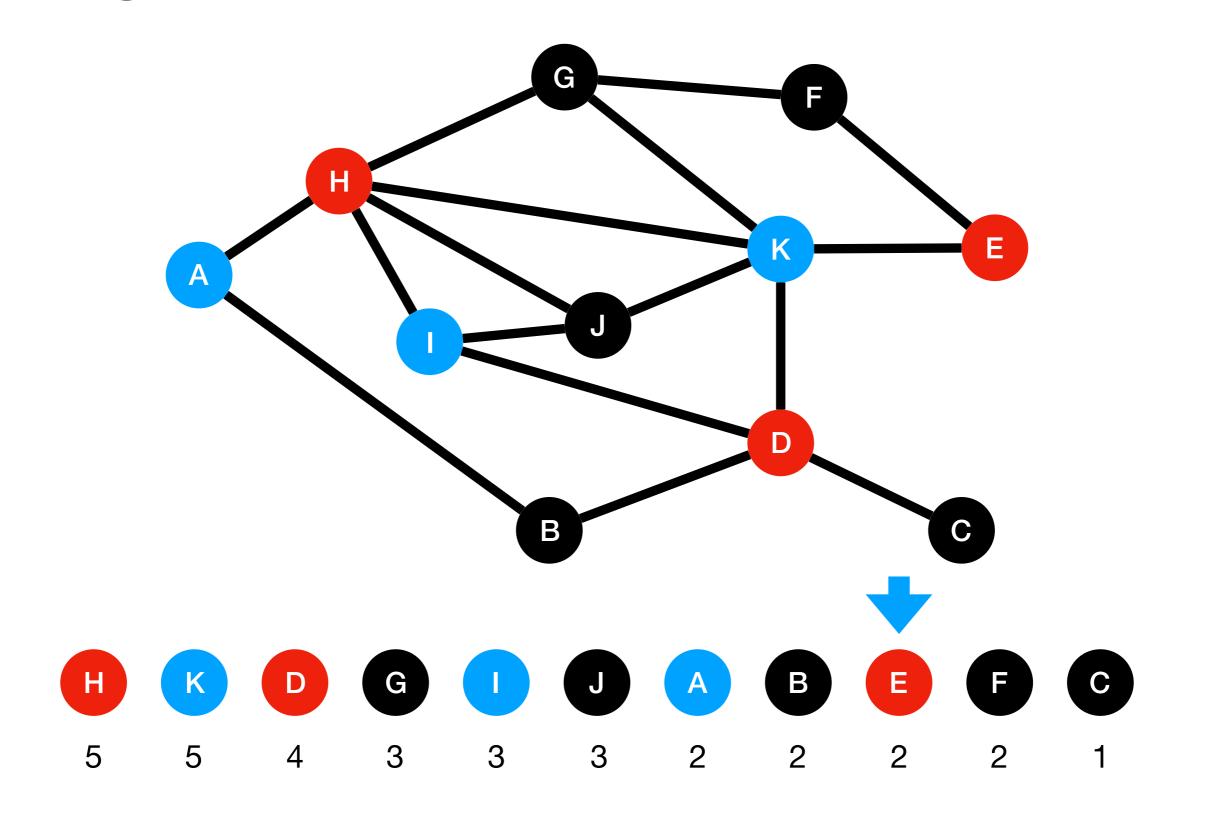


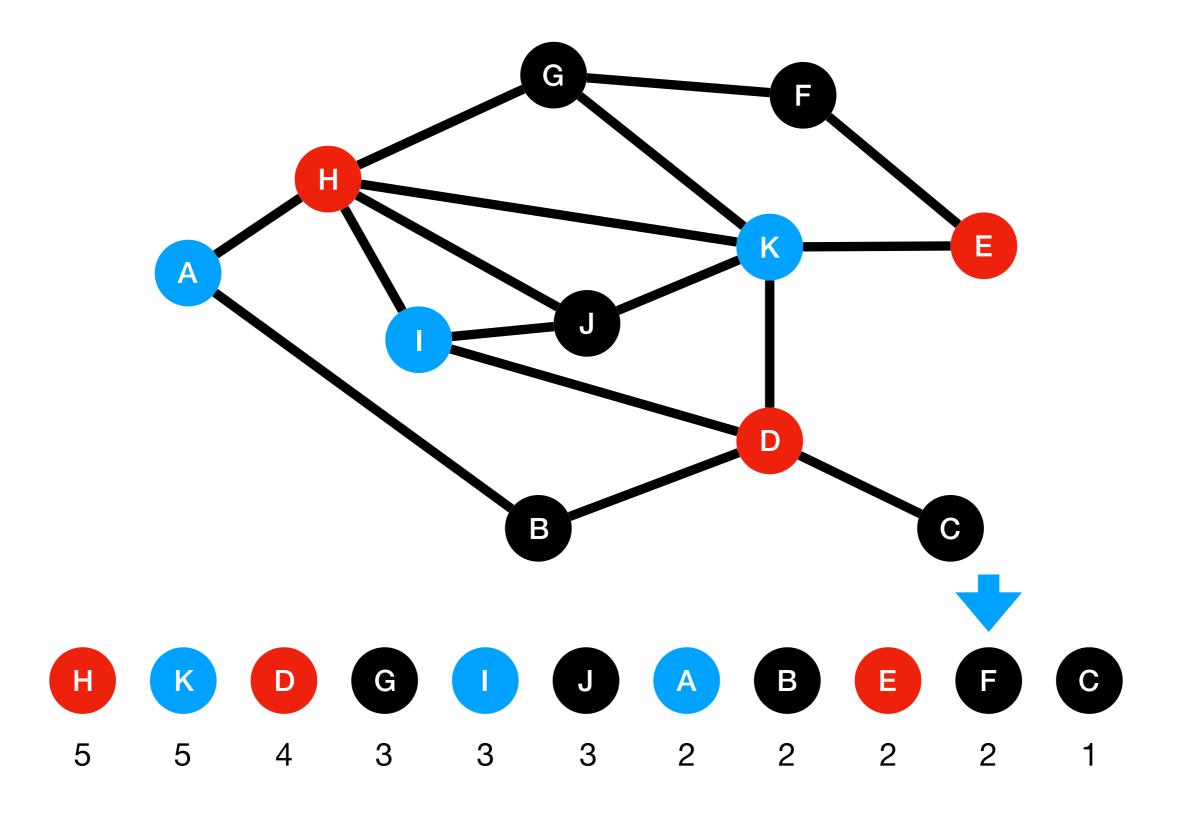


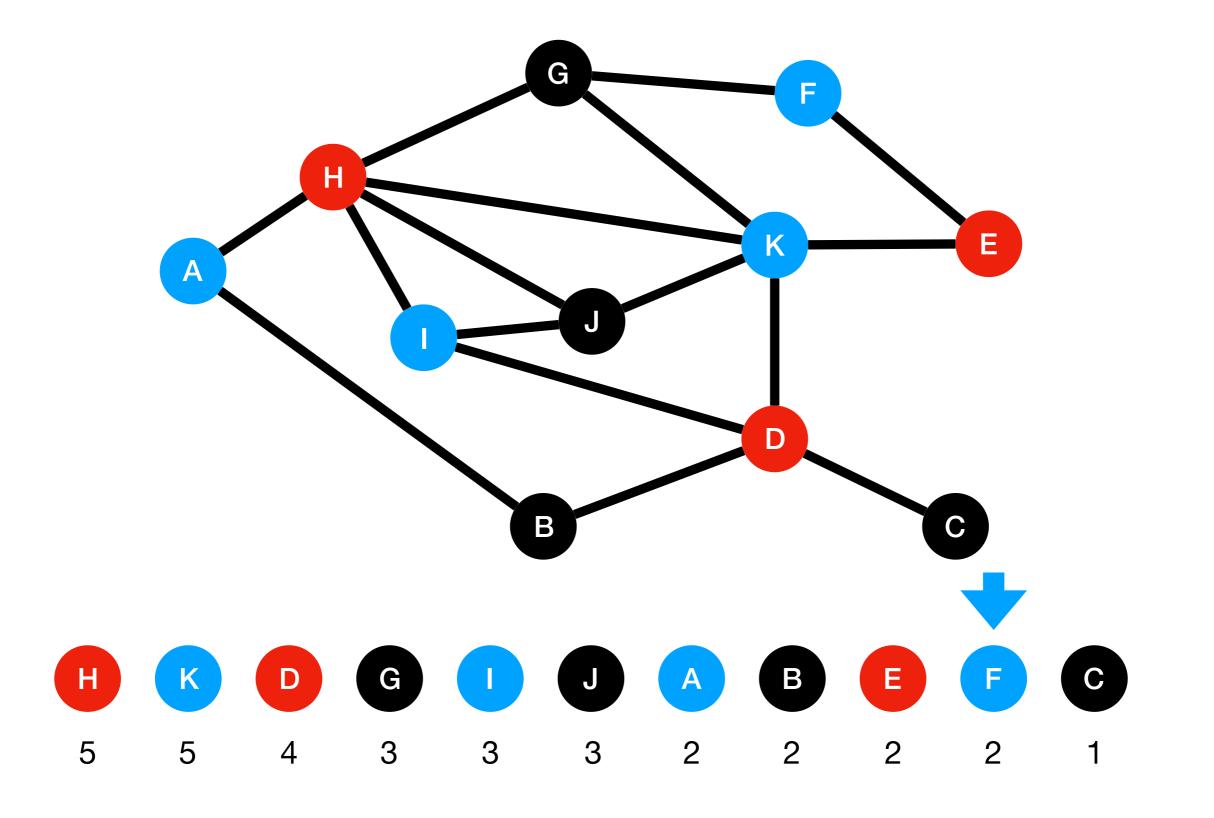


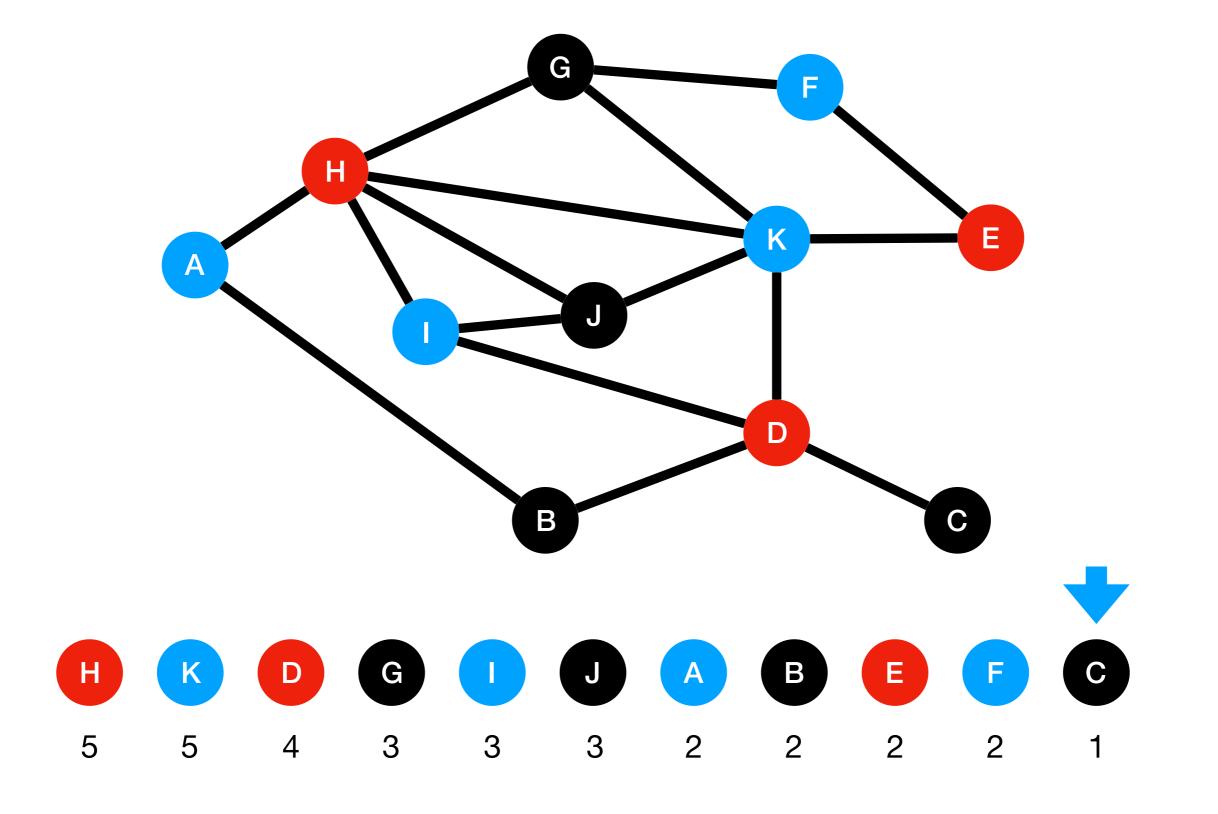


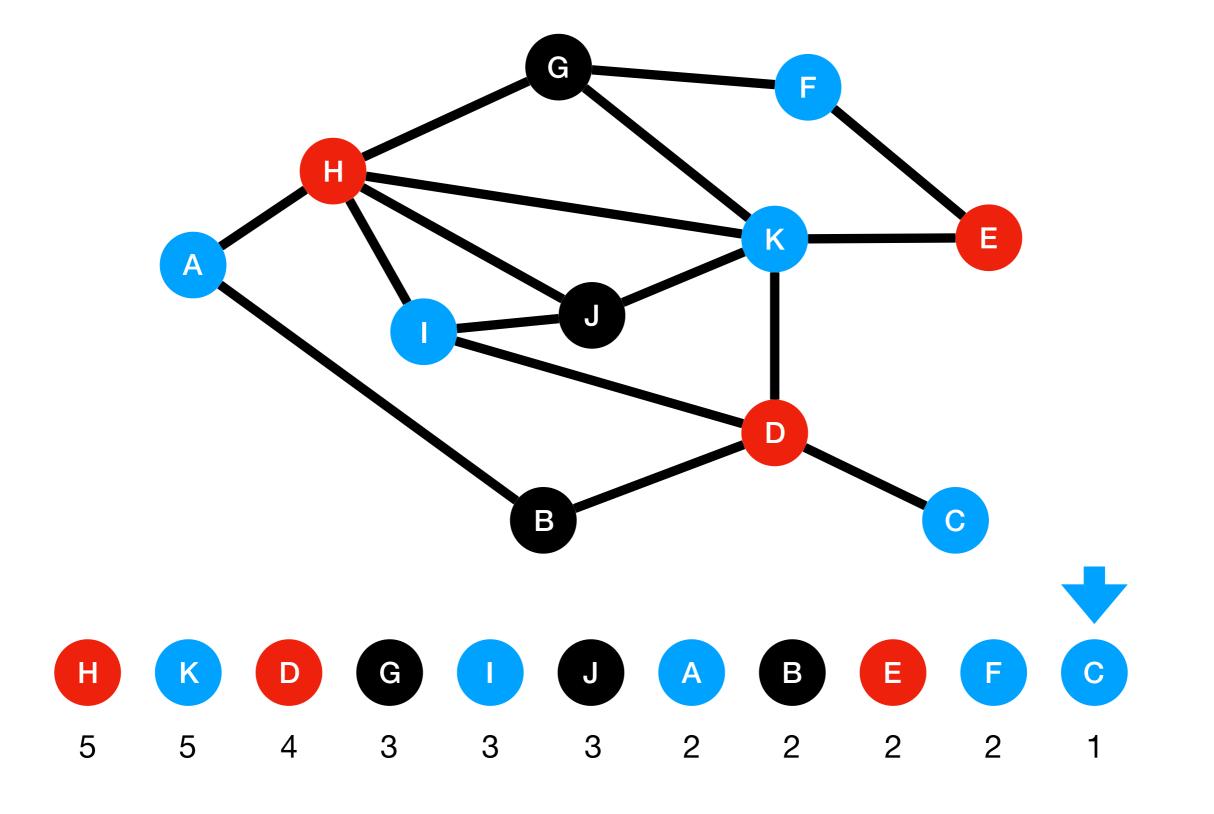


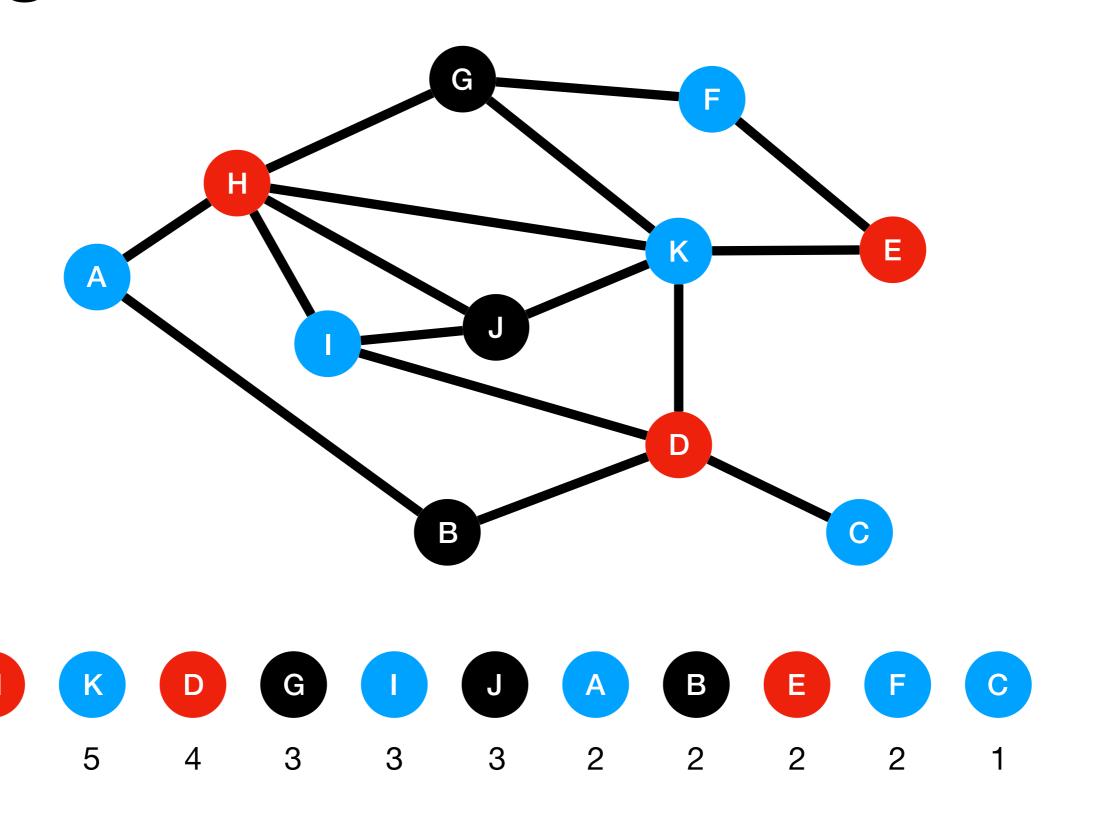


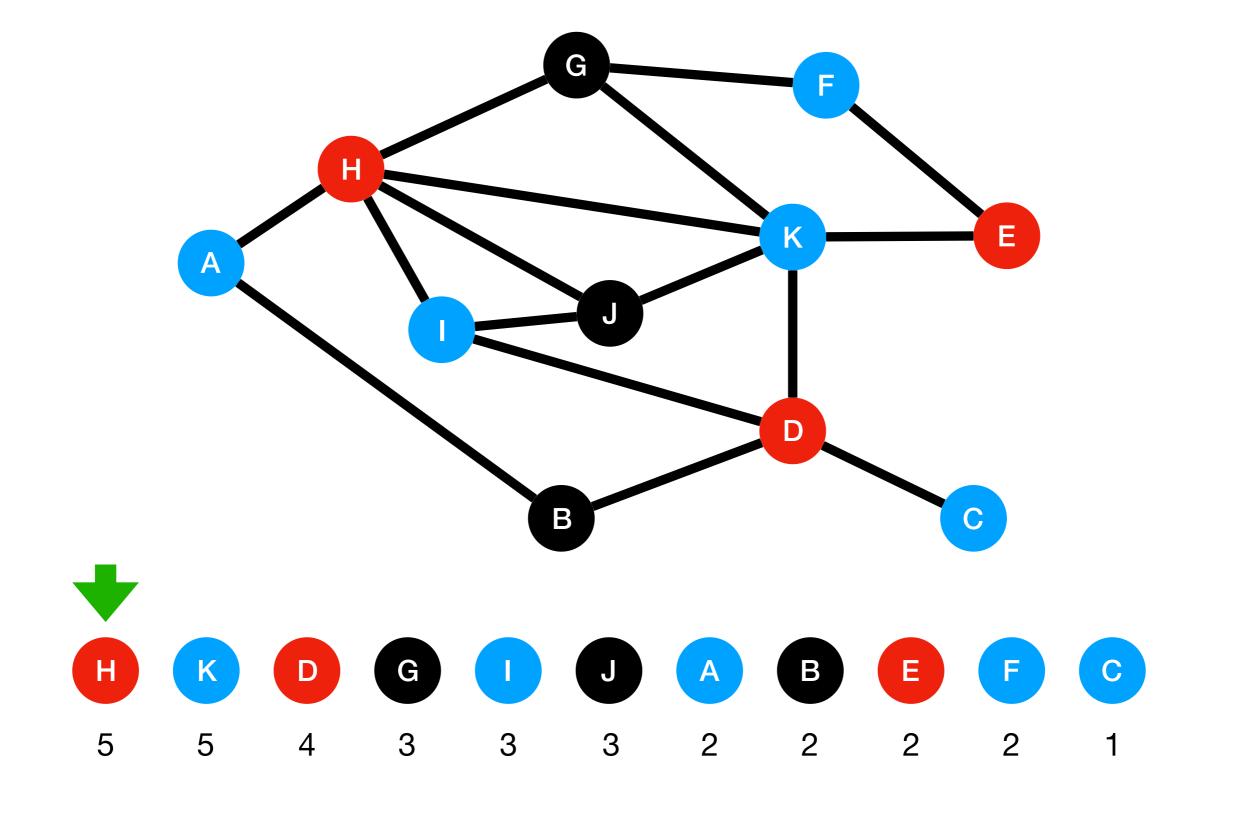


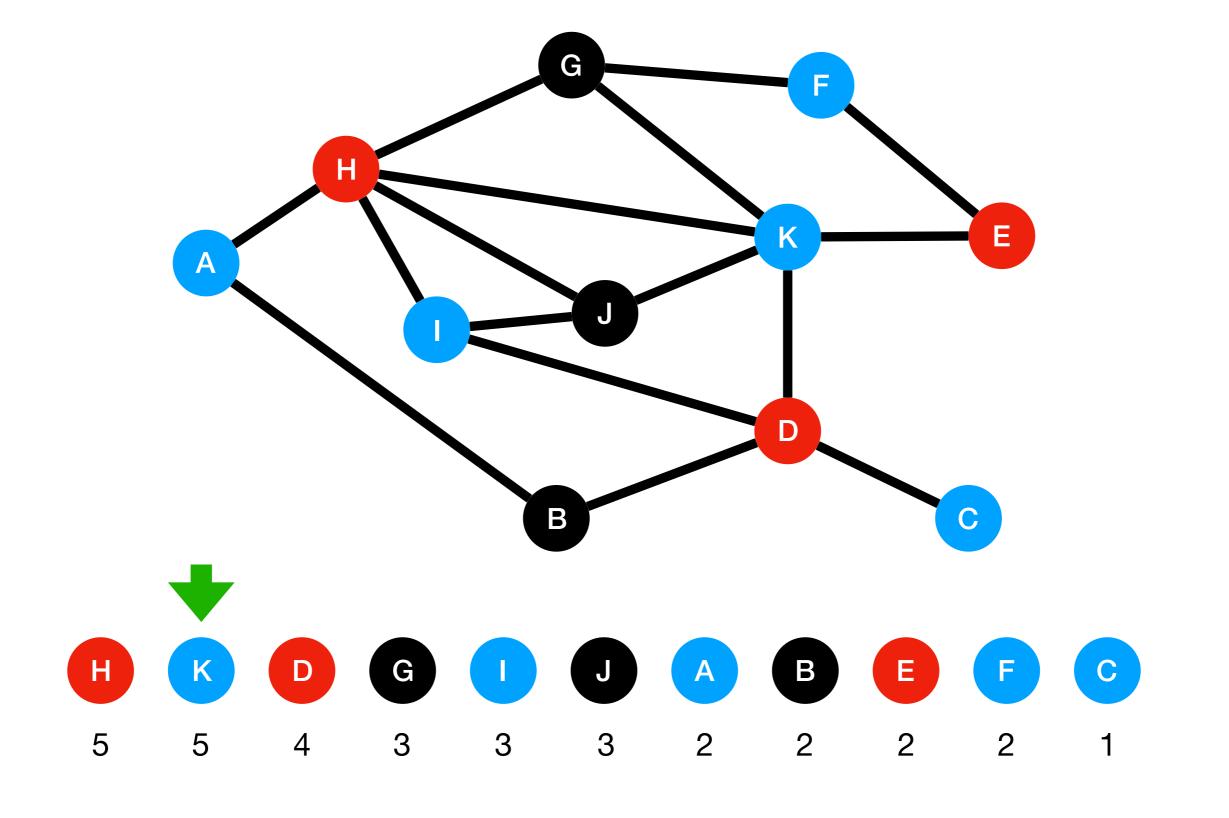


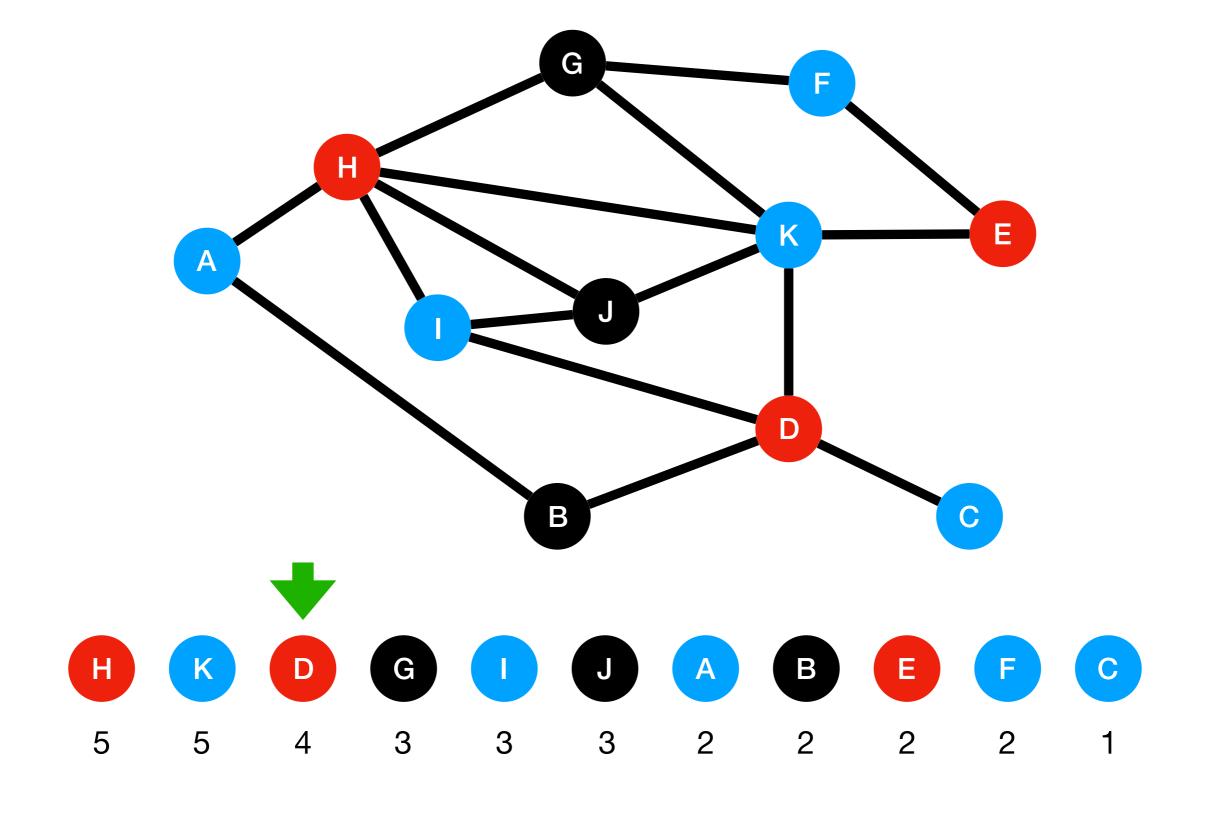


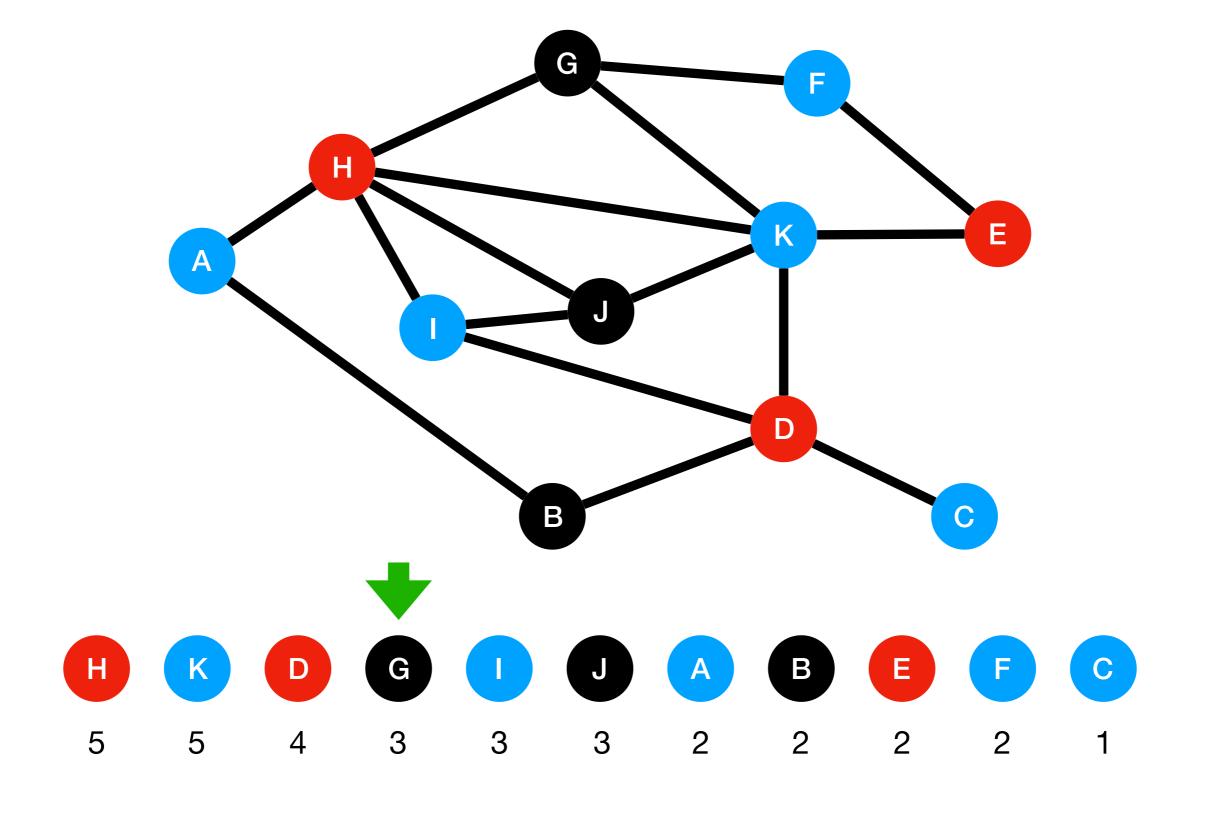


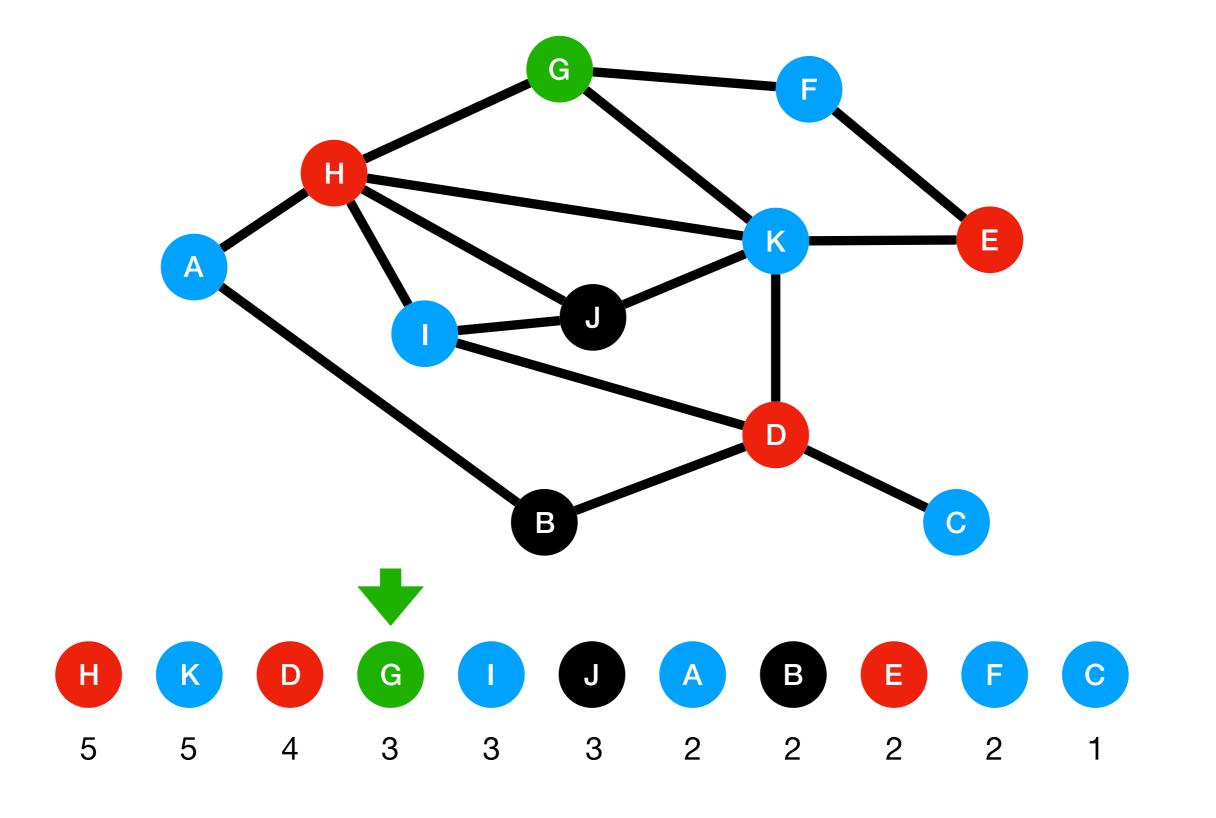


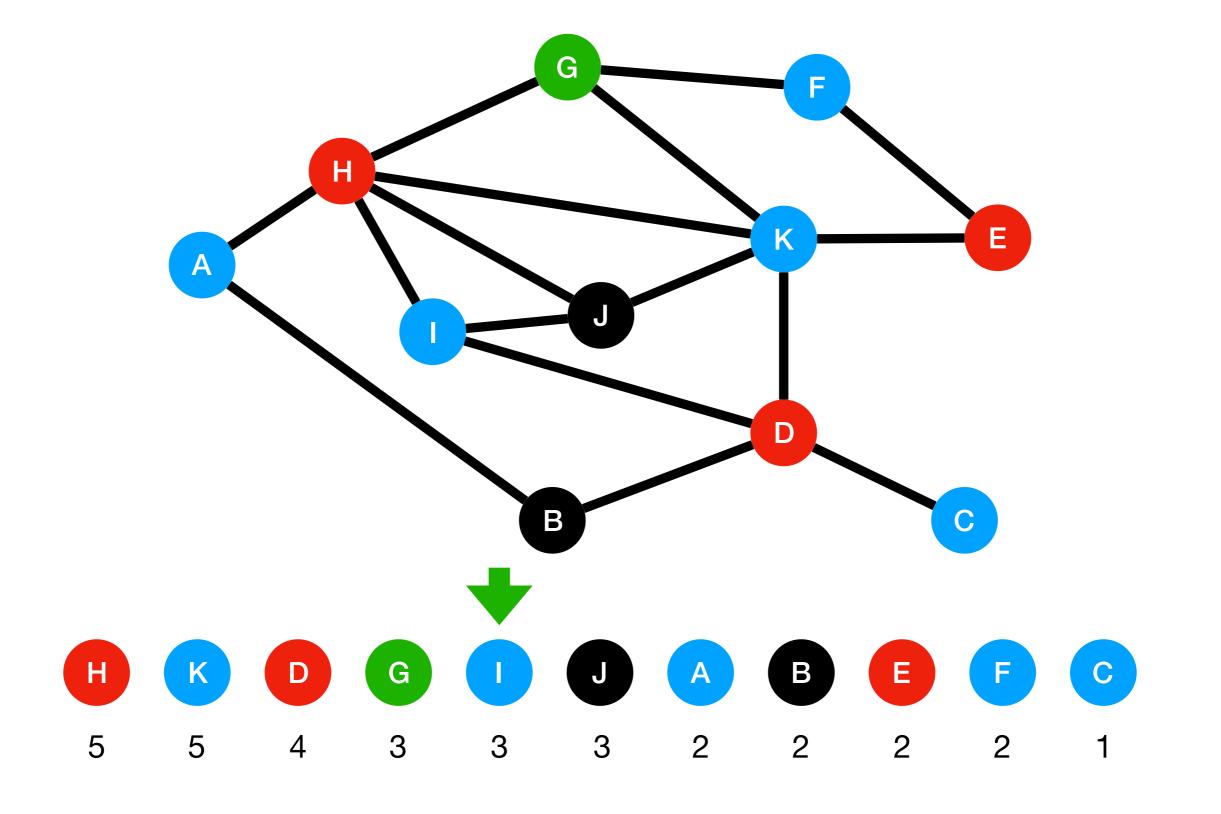


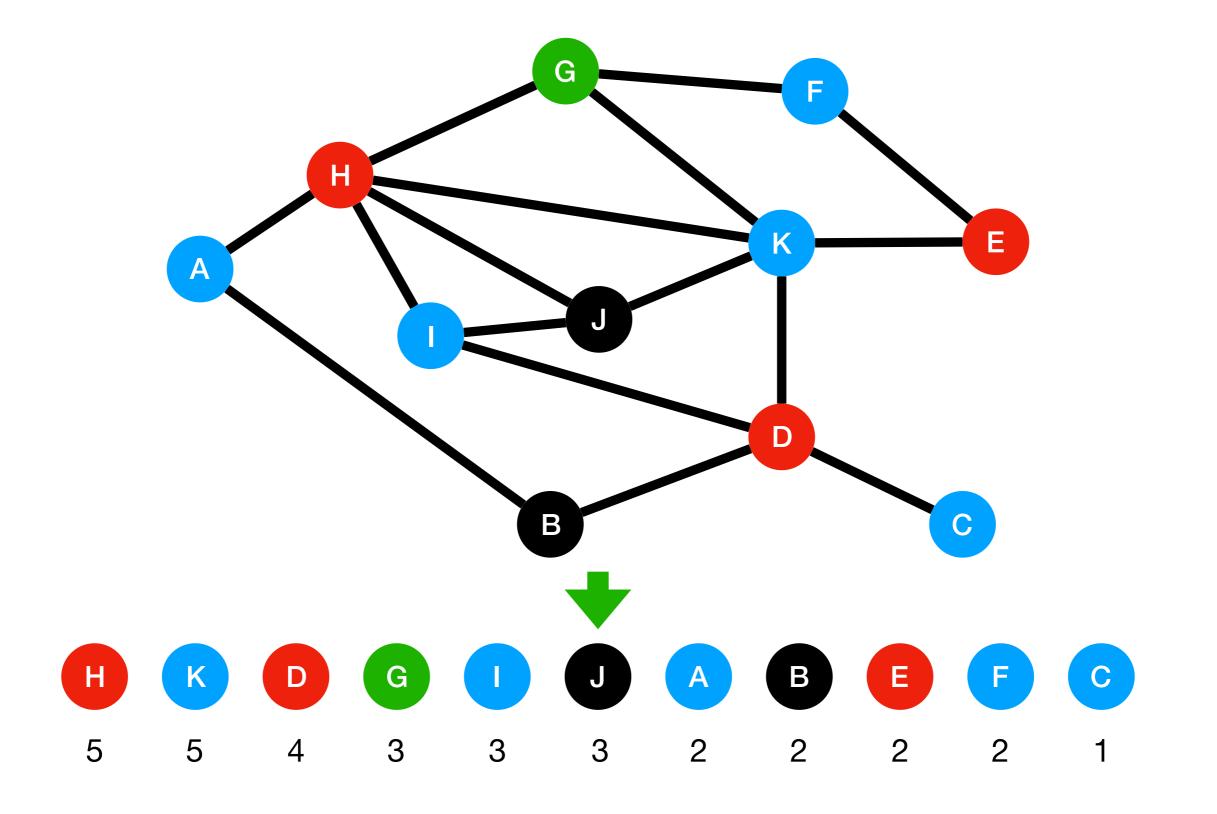


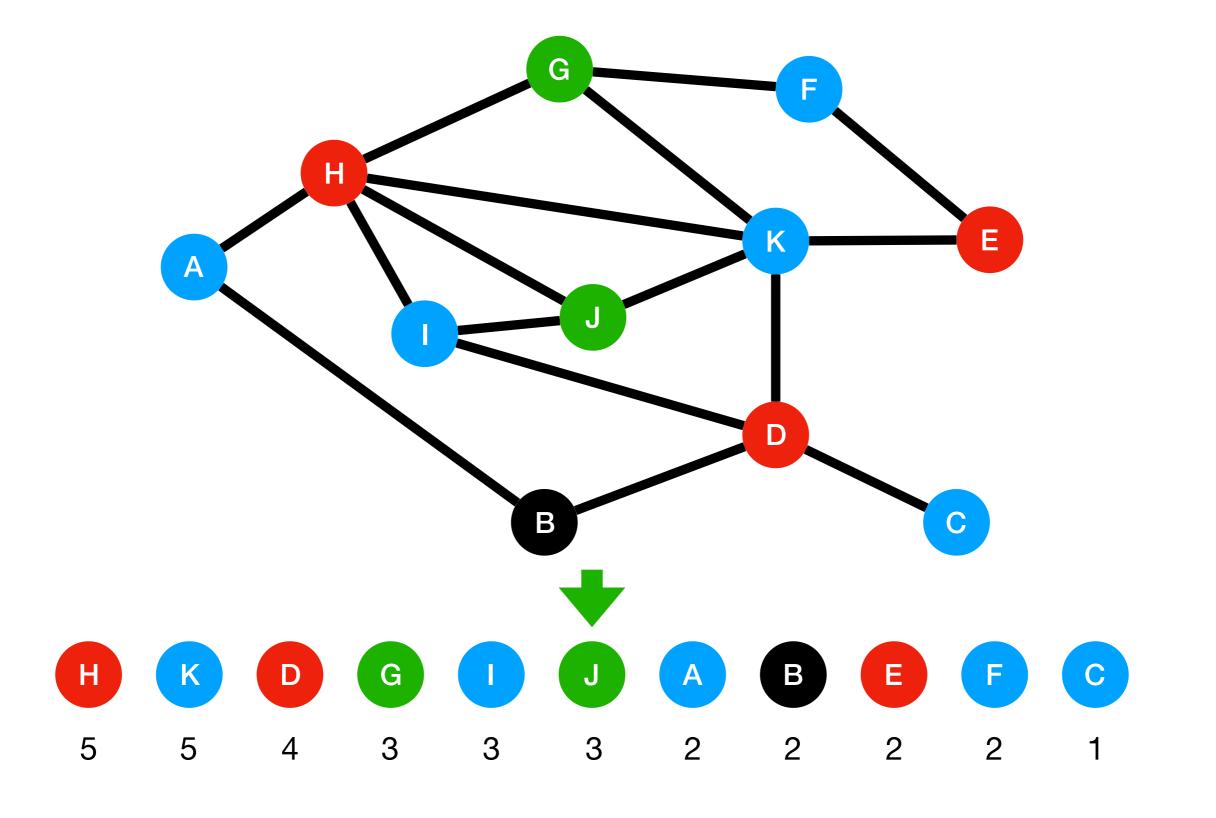


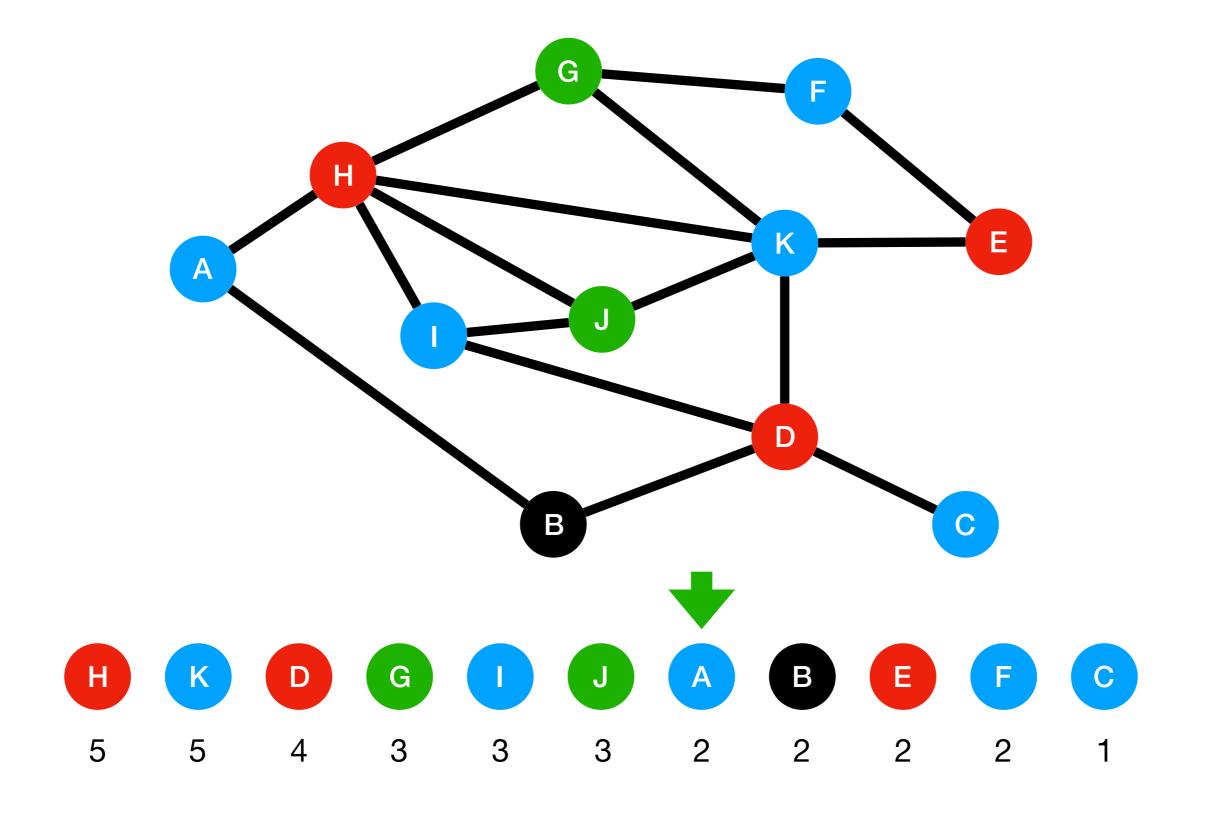


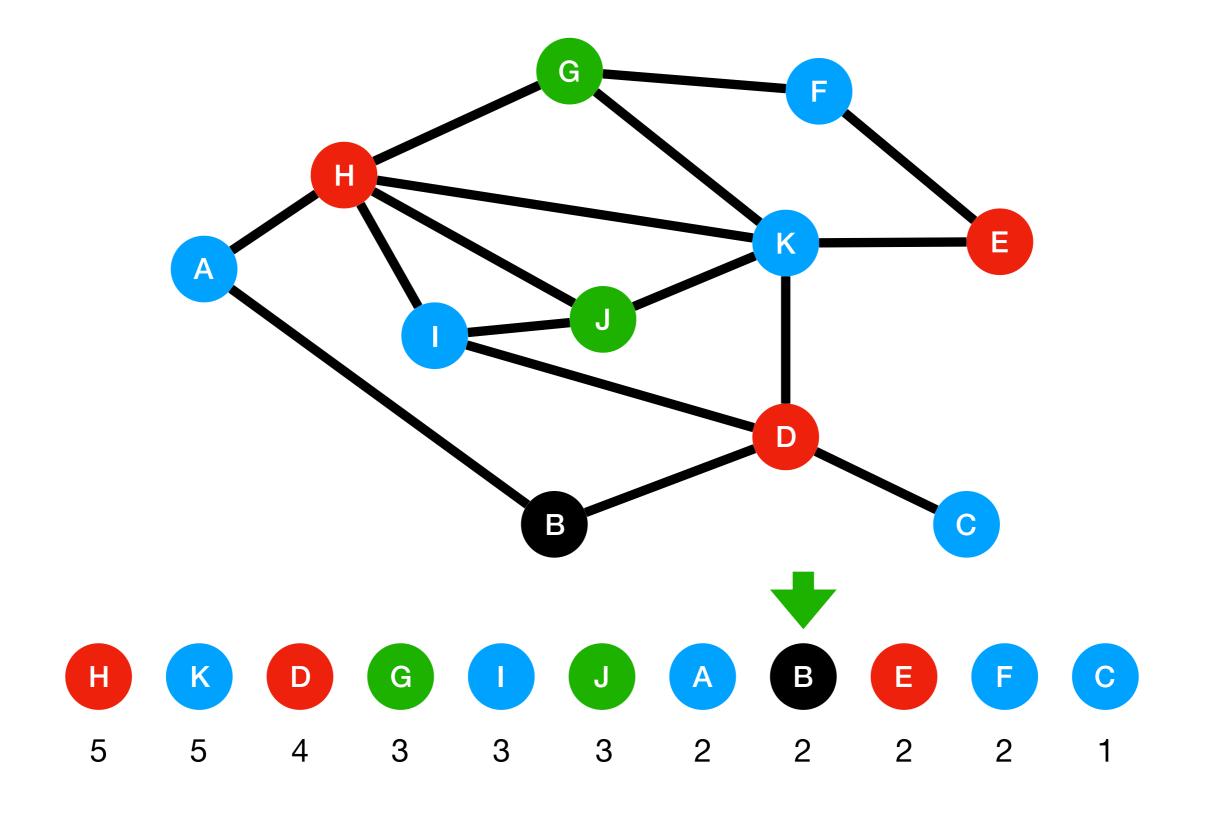


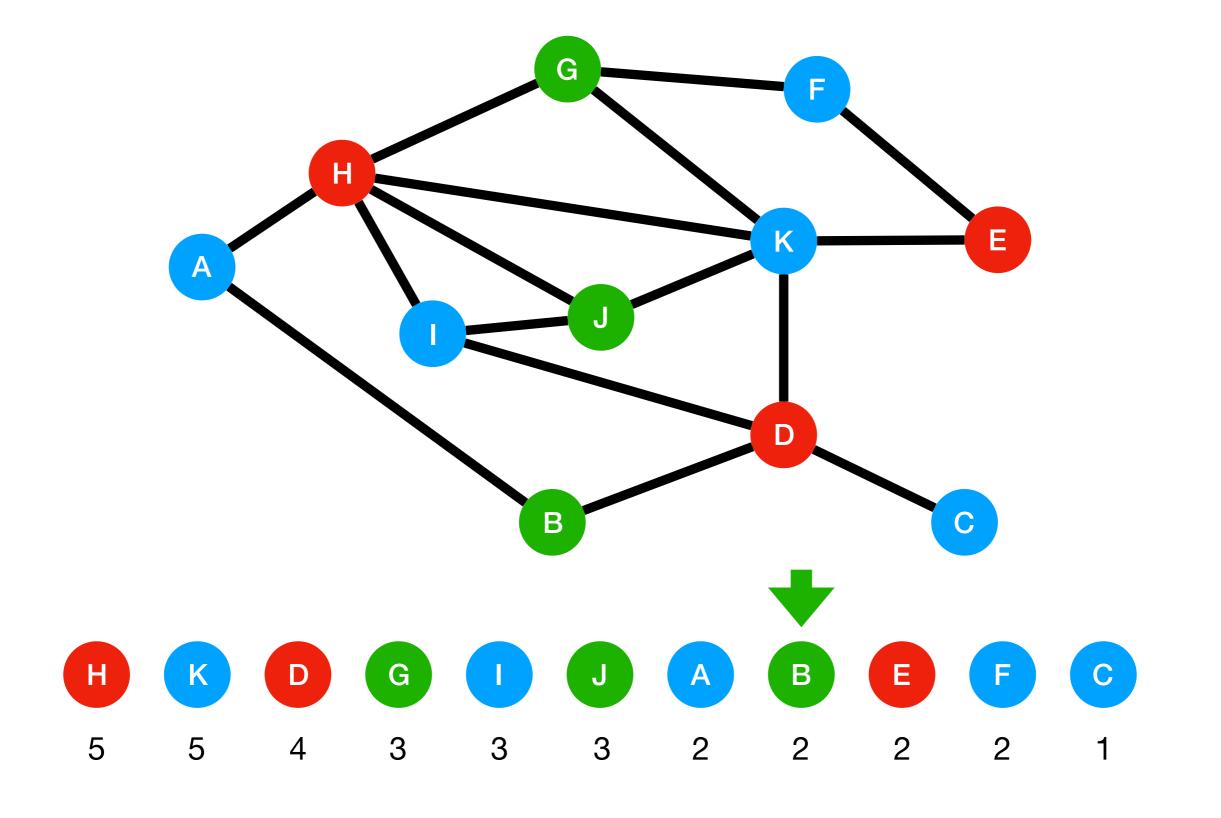


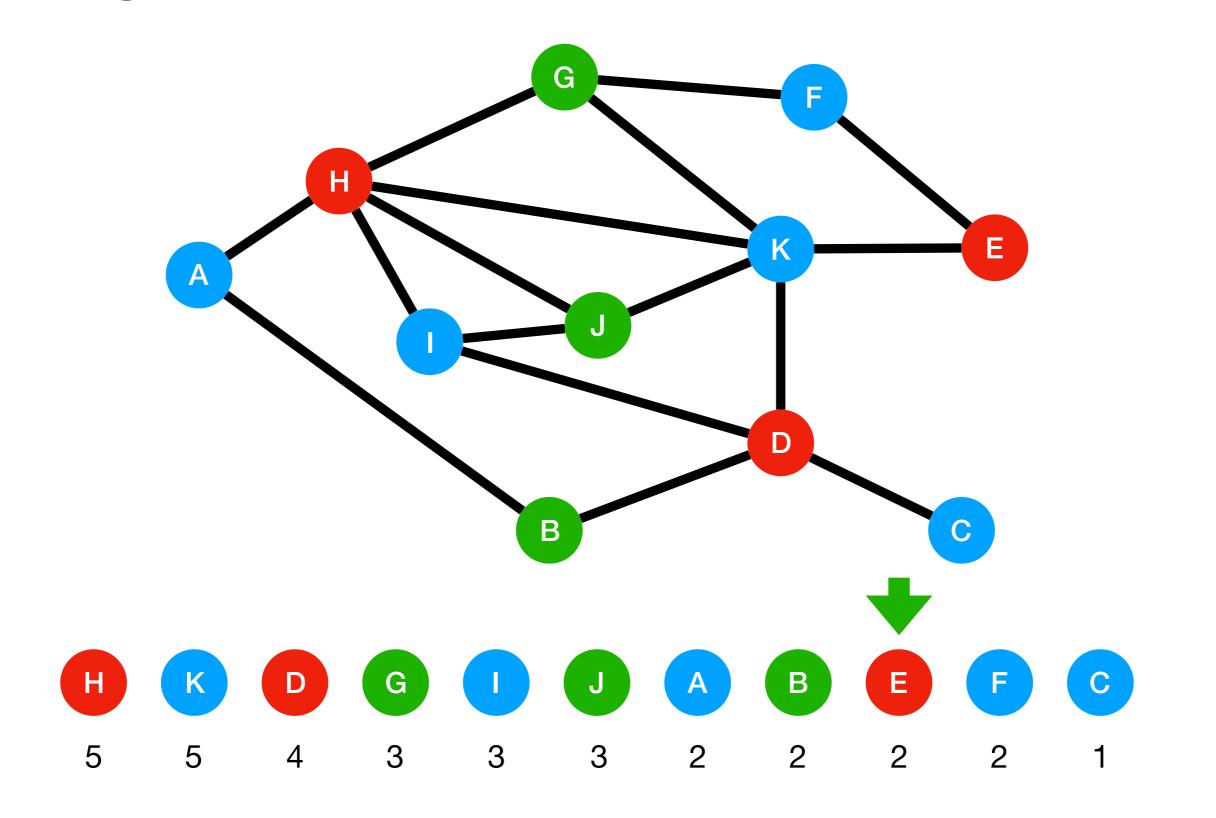


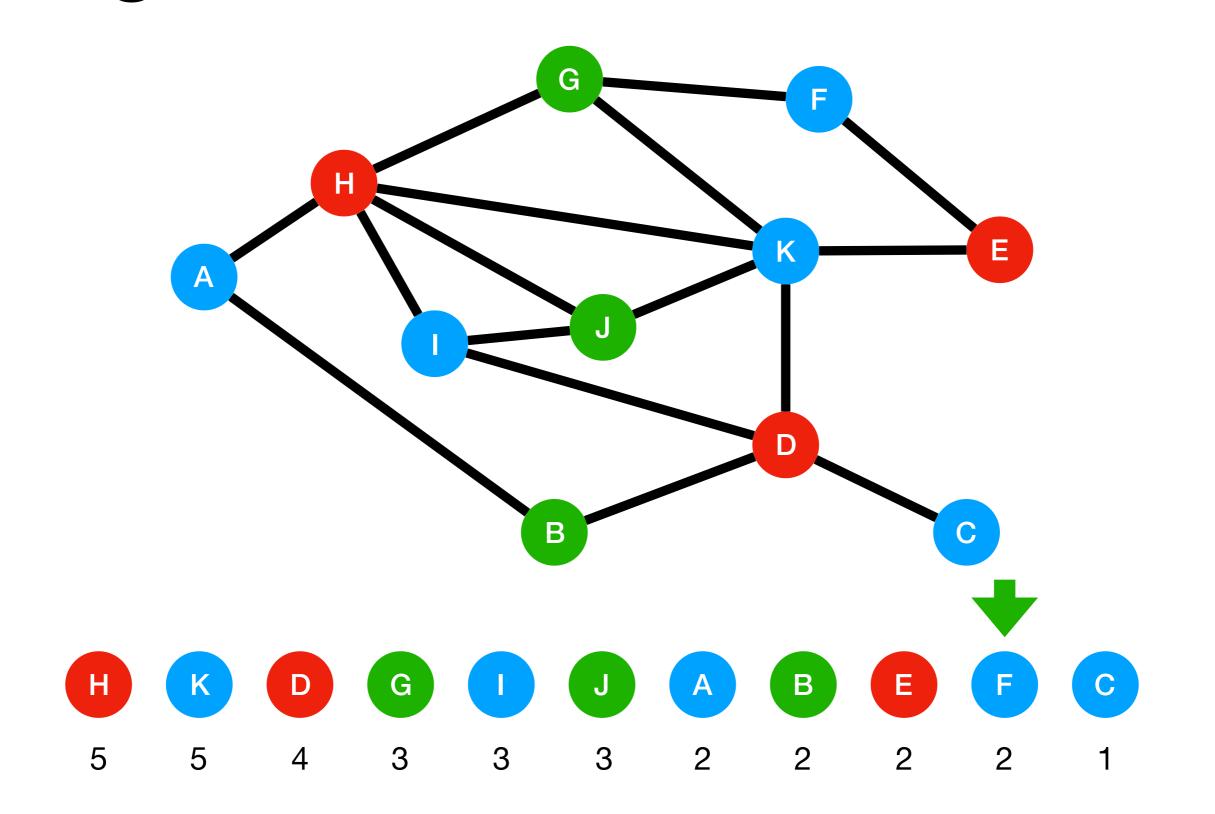


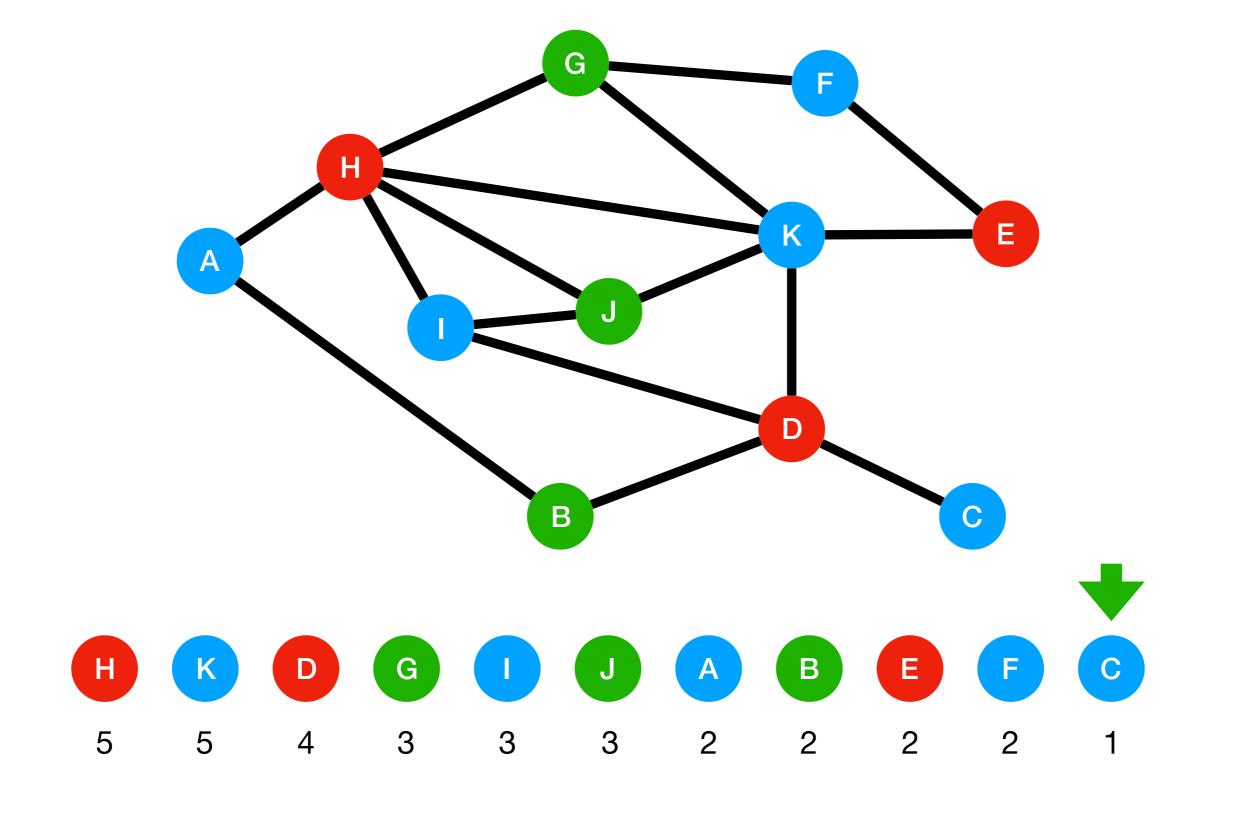


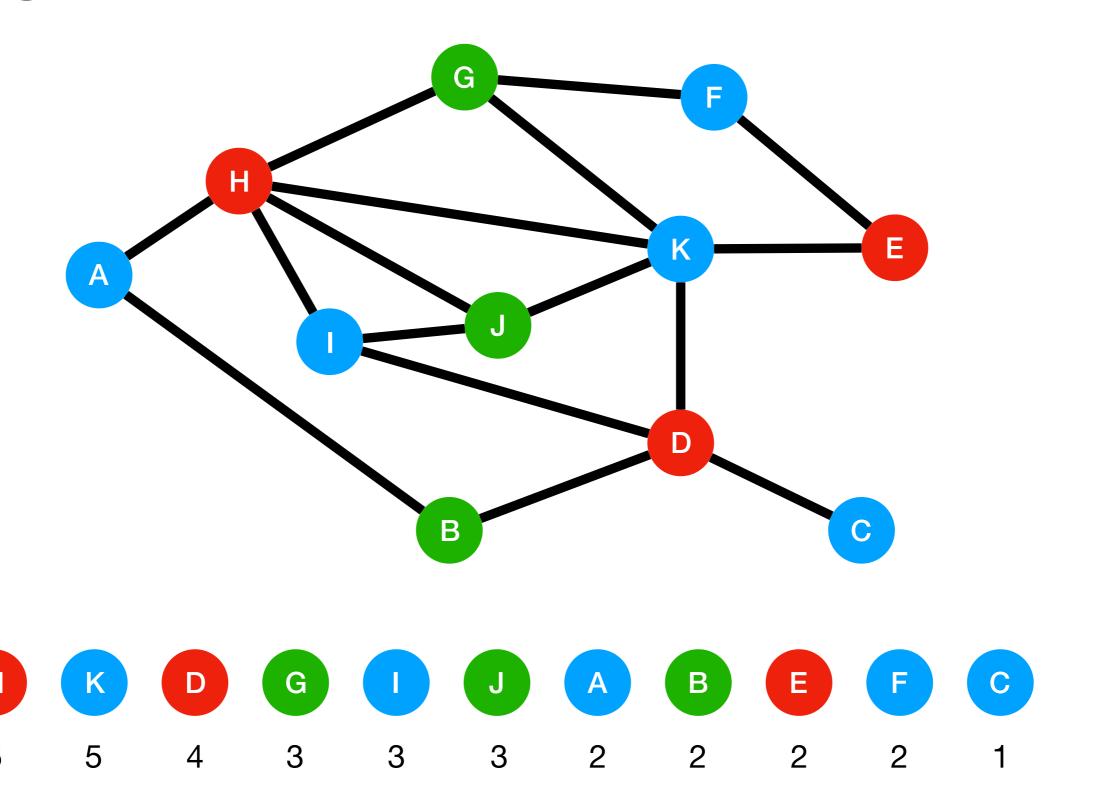


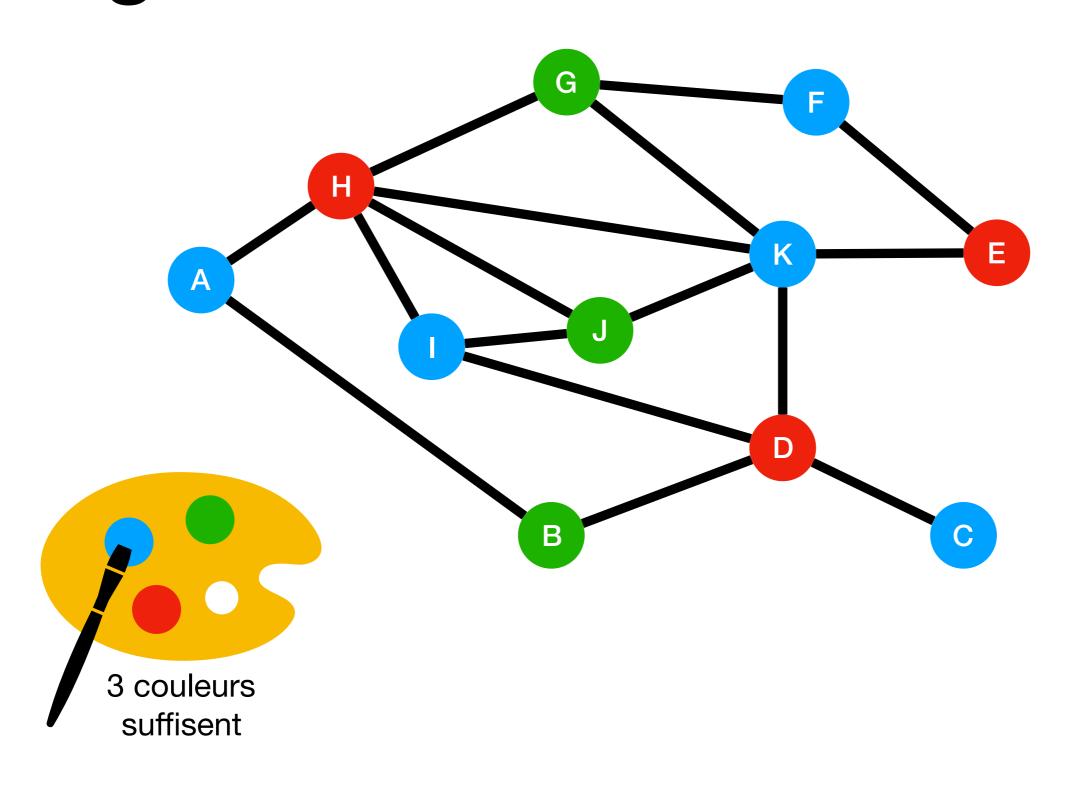


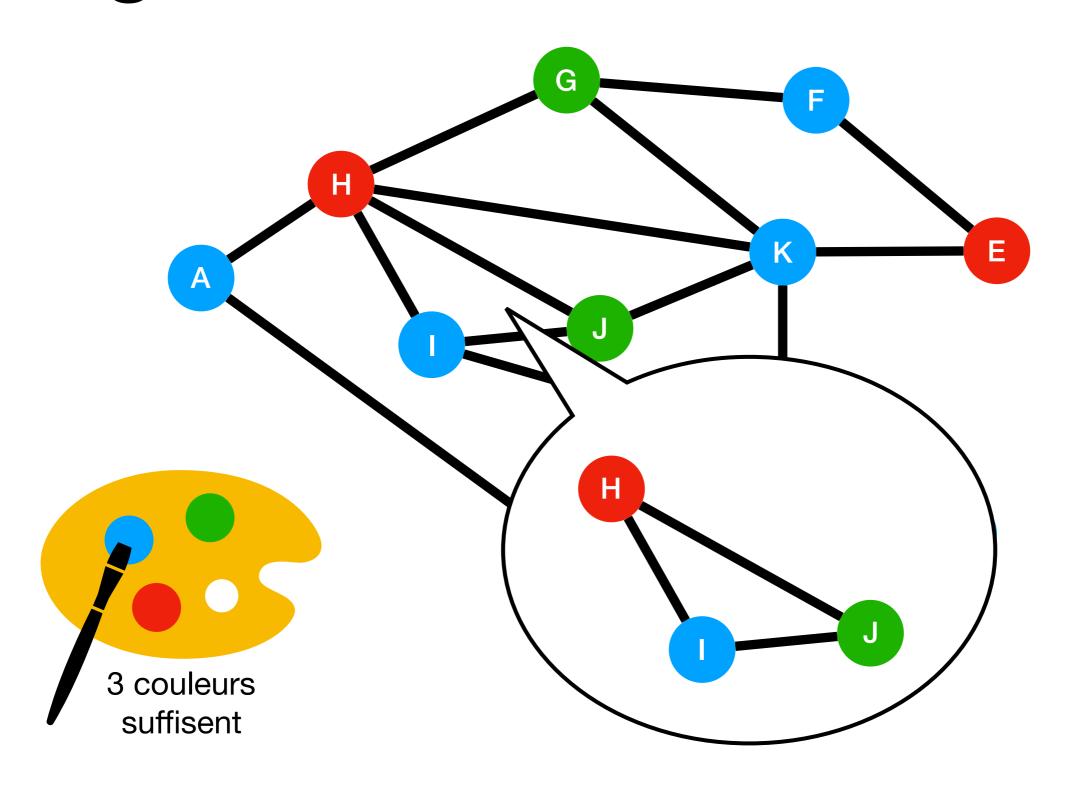


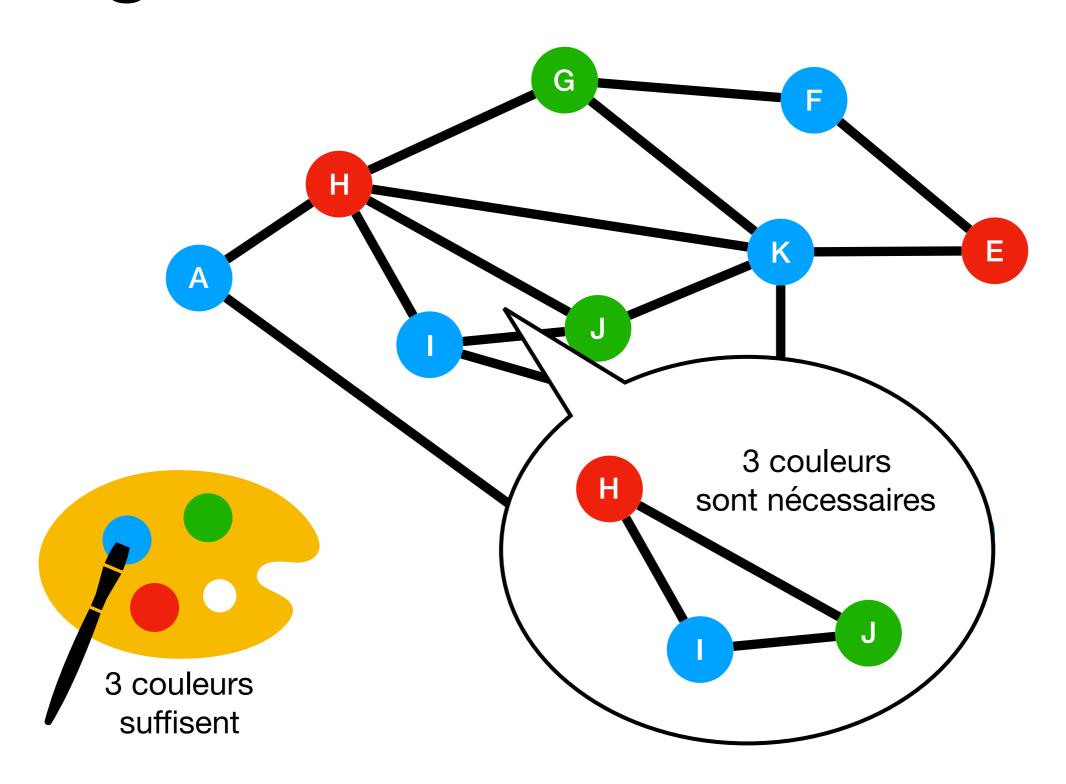












- Trier le sommets du graphe par ordre de degré (nombre de voisins) décroissant
- couleur ≔ rouge
- Tant qu'il y a encore des sommets en noir faire
 - Parcourir la liste triée des sommets et colorer en couleur les sommets en noir qui ne sont pas connectés à d'autres sommets de la même couleur
 - Choisir une nouvelle couleur

