Introduction à l'informatique CM5

Antonio E. Porreca https://aeporreca.org/introinfo

Rappel A Vendredi 18 octobre on a le partiel!

Algorithmes de tri

Algorithmes de tri pour accélérer la recherche dans un tableau

- La recherche dans un tableau non trié prend temps O(n) avec la recherche séquentielle (ou linéaire)
- Par contre, on peut faire une recherche dichotomique dans un tableau trié en temps O(log₂ n)
- Donc ça vaut la peine de trier le tableau si on a beaucoup de recherches à faire

Algorithmes de tri dans le commerce électronique

amazonie.fr





Chercher:

Le Petit Prince

Trier par:

Résultats 1-20 sur 928572785 pour « Le Petit Prince »

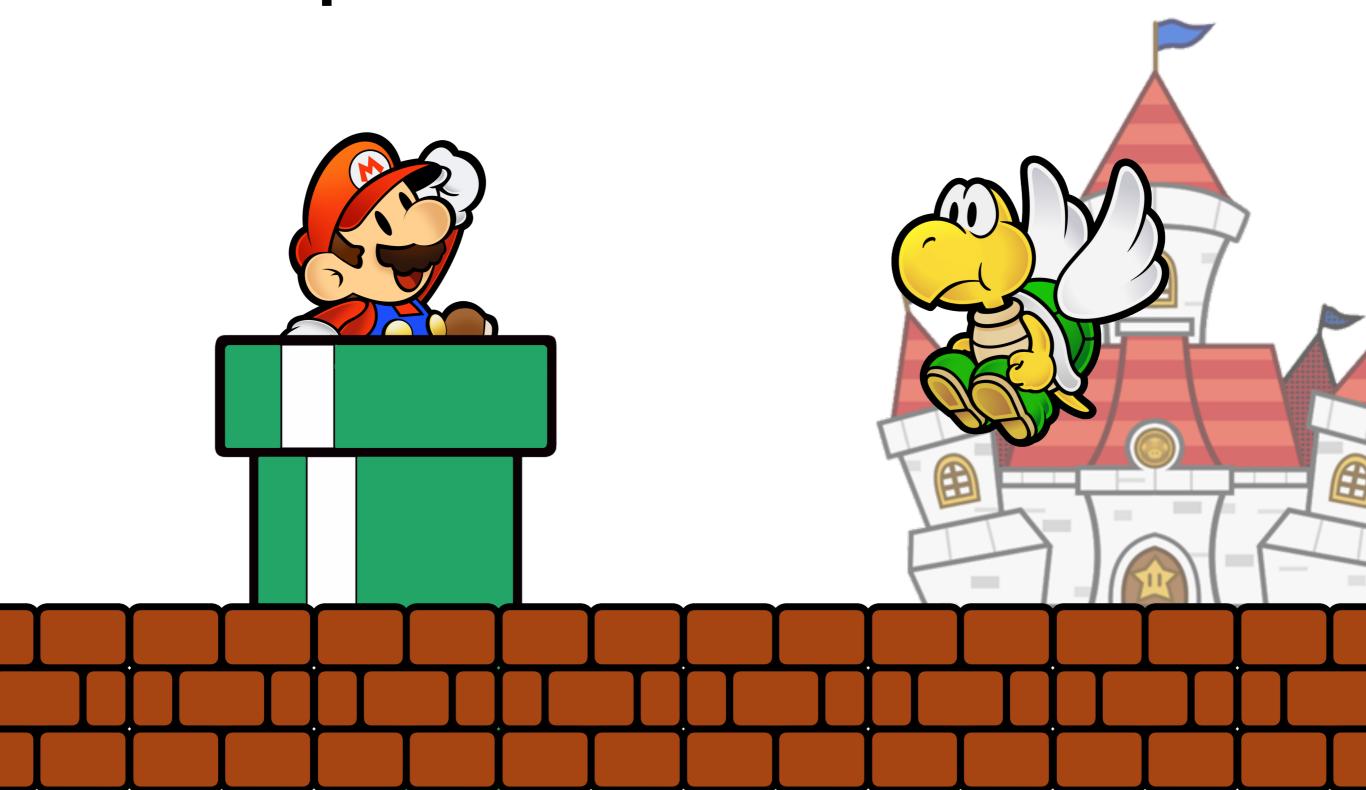


Le Petit Prince de Antoine de Saint-Exur

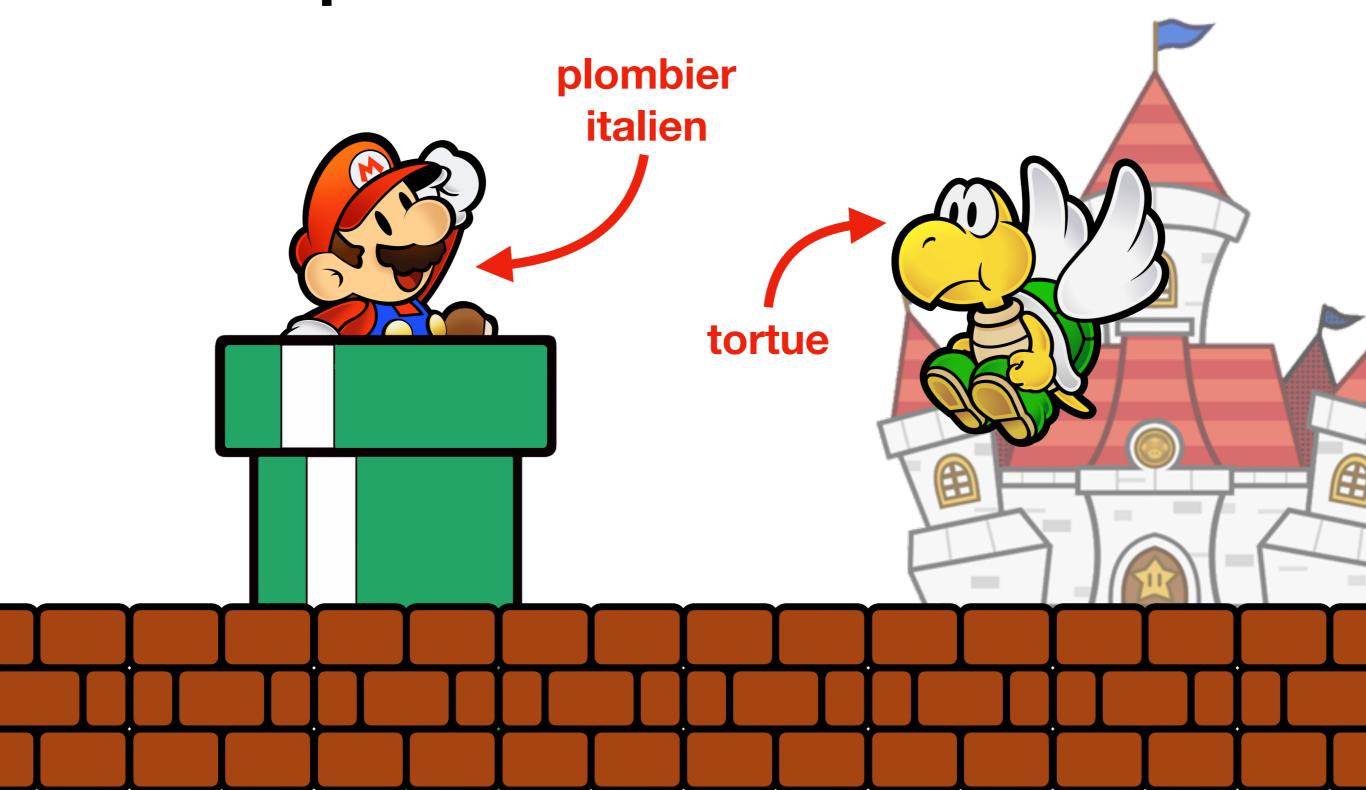
Format poche 6,90 € Format Kinder 6,49 €

prix croissant
prix décroissant
note moyenne
nouveauté

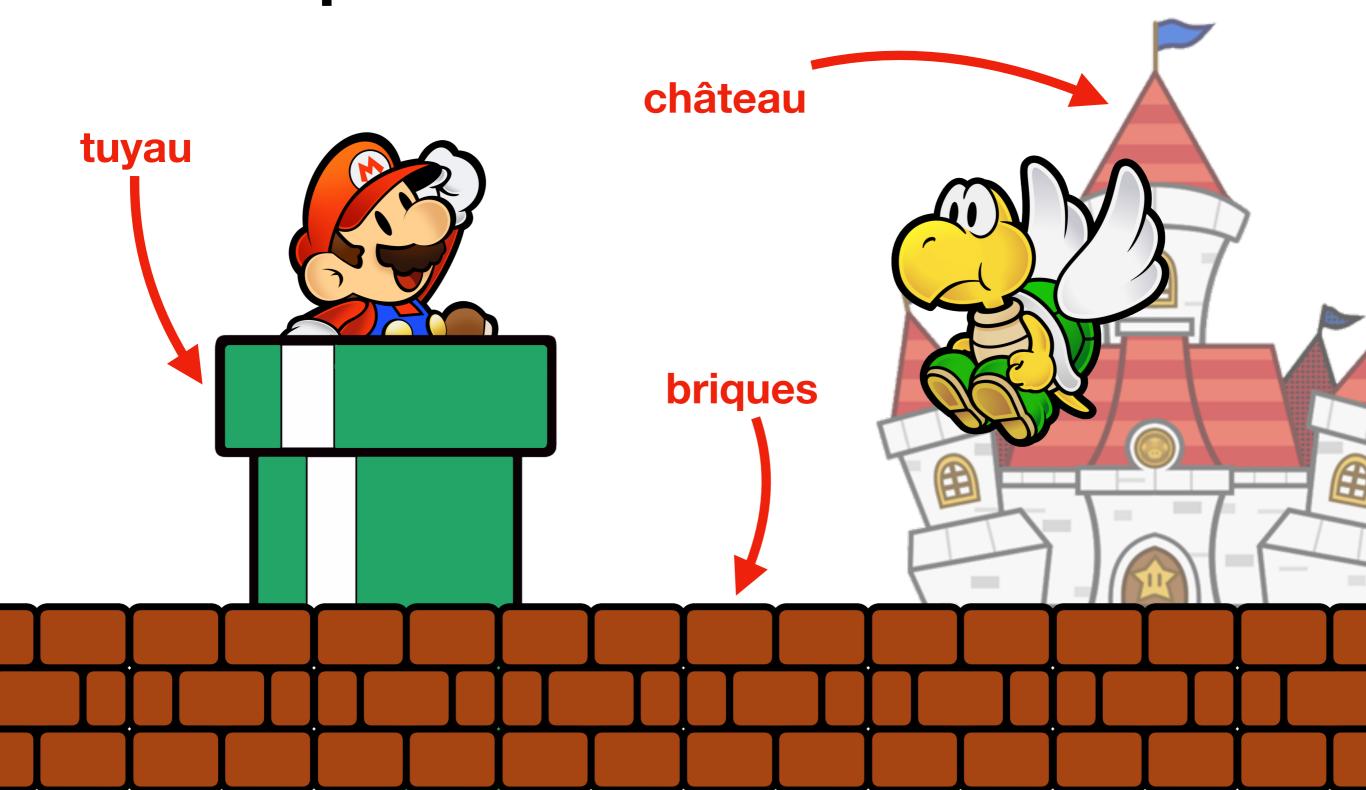
Algos de tri dans le jeux vidéo « Super Plombiers Italiens »



Algos de tri dans le jeux vidéo « Super Plombiers Italiens »

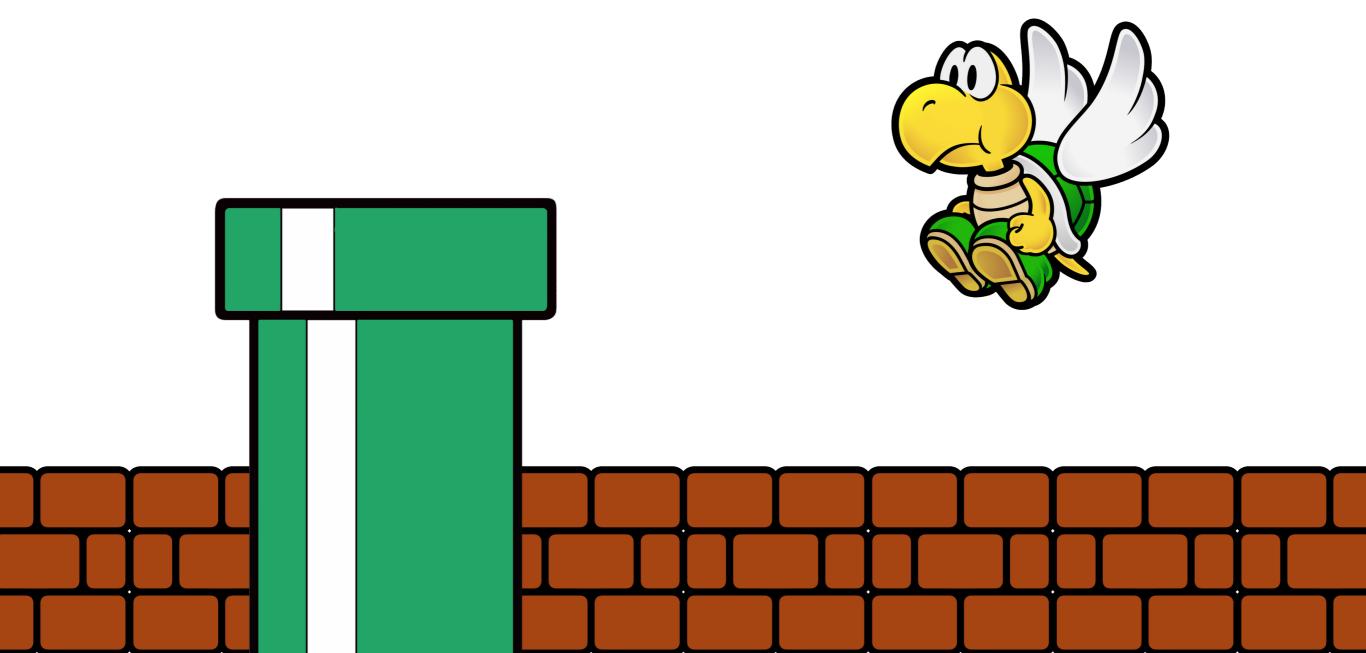


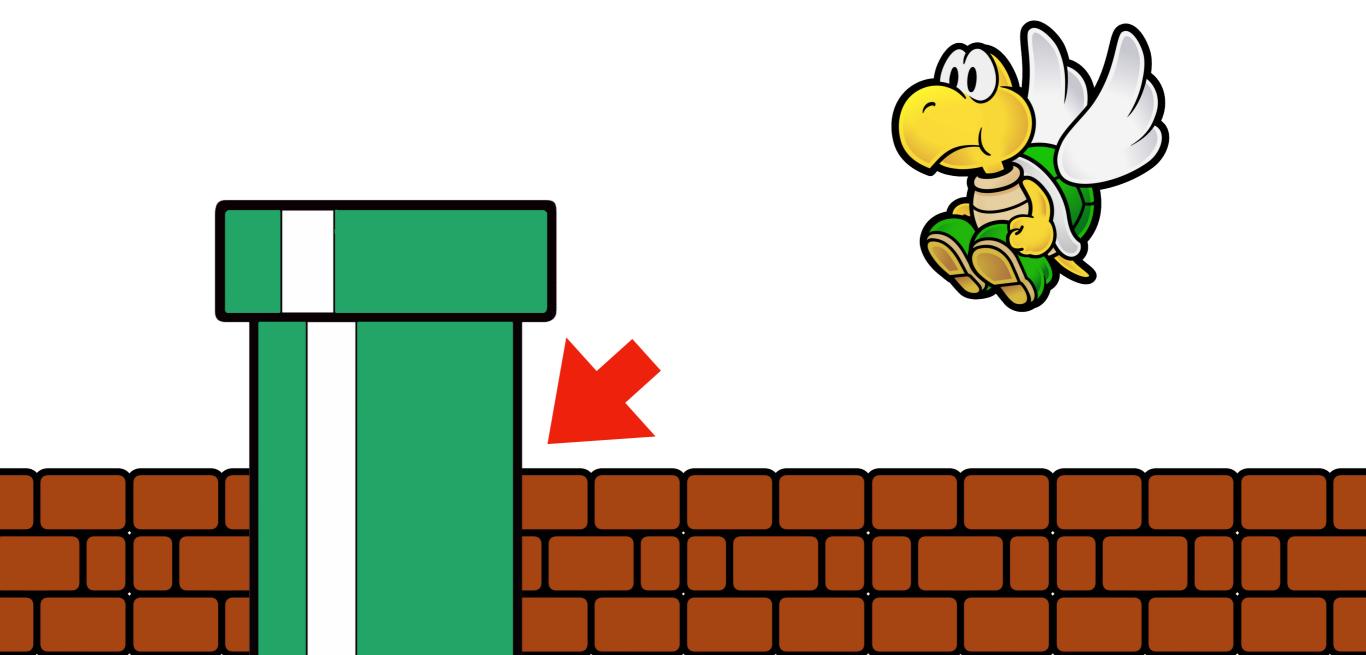
Algos de tri dans le jeux vidéo « Super Plombiers Italiens »

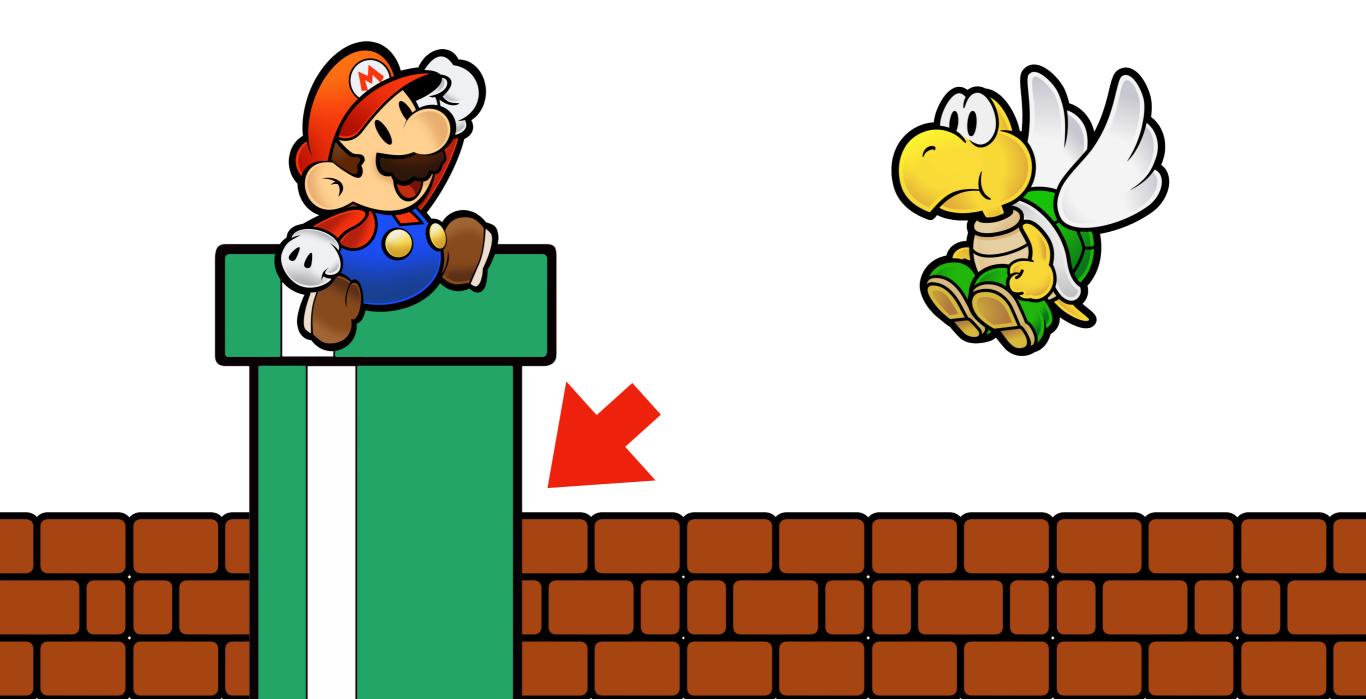


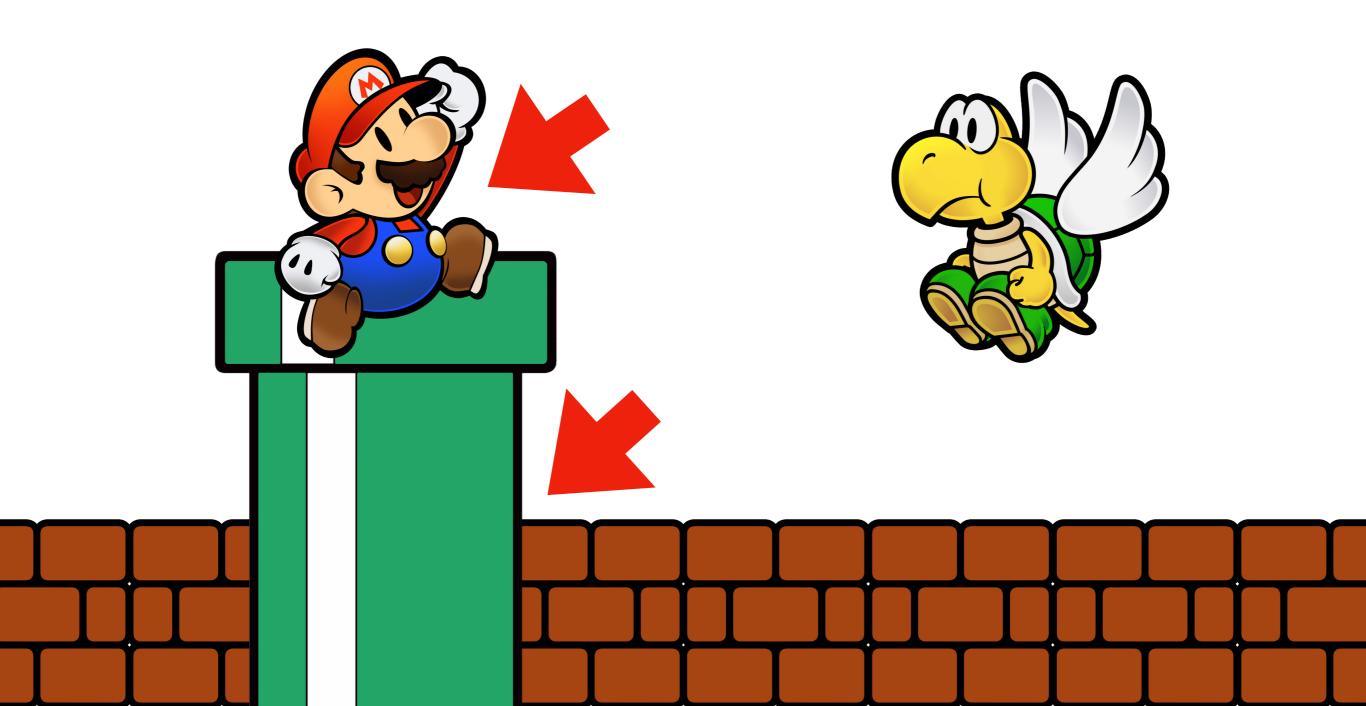


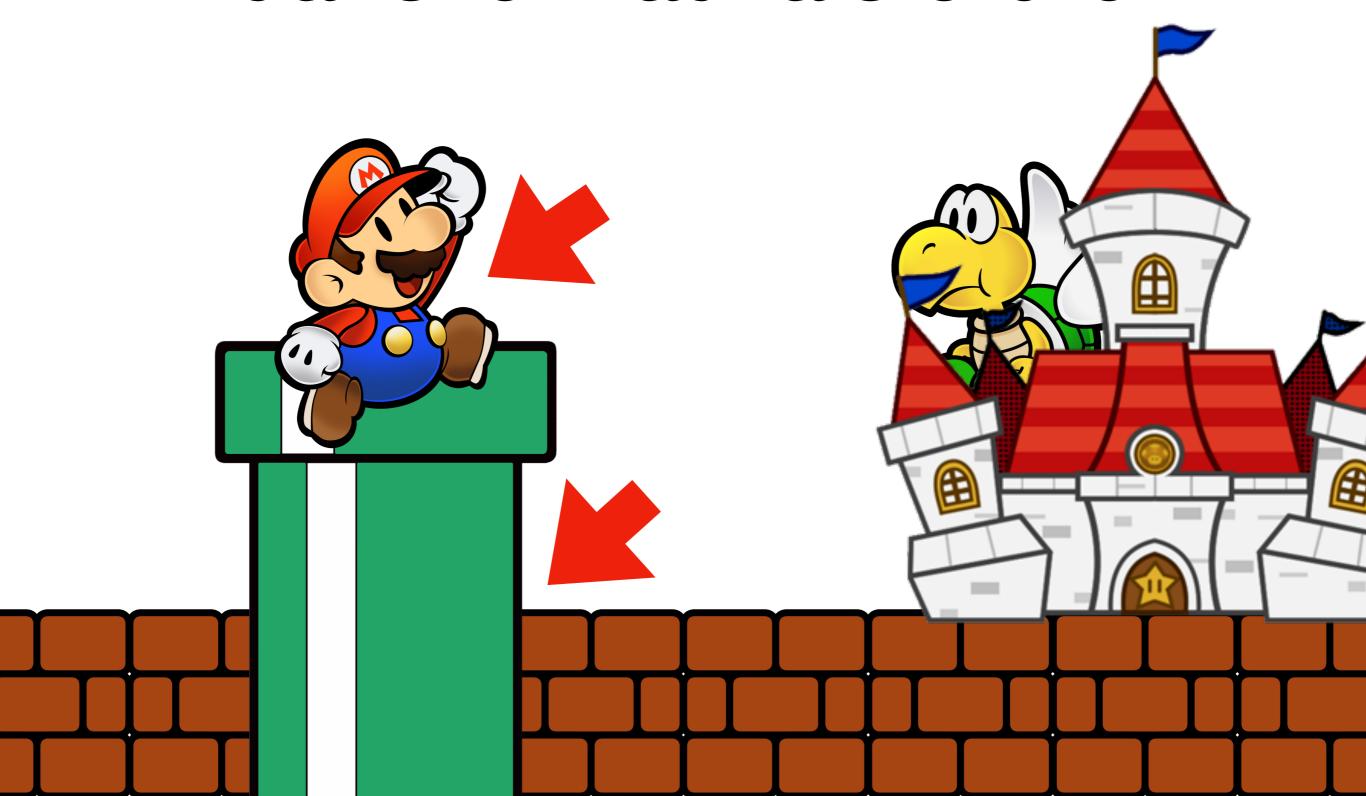


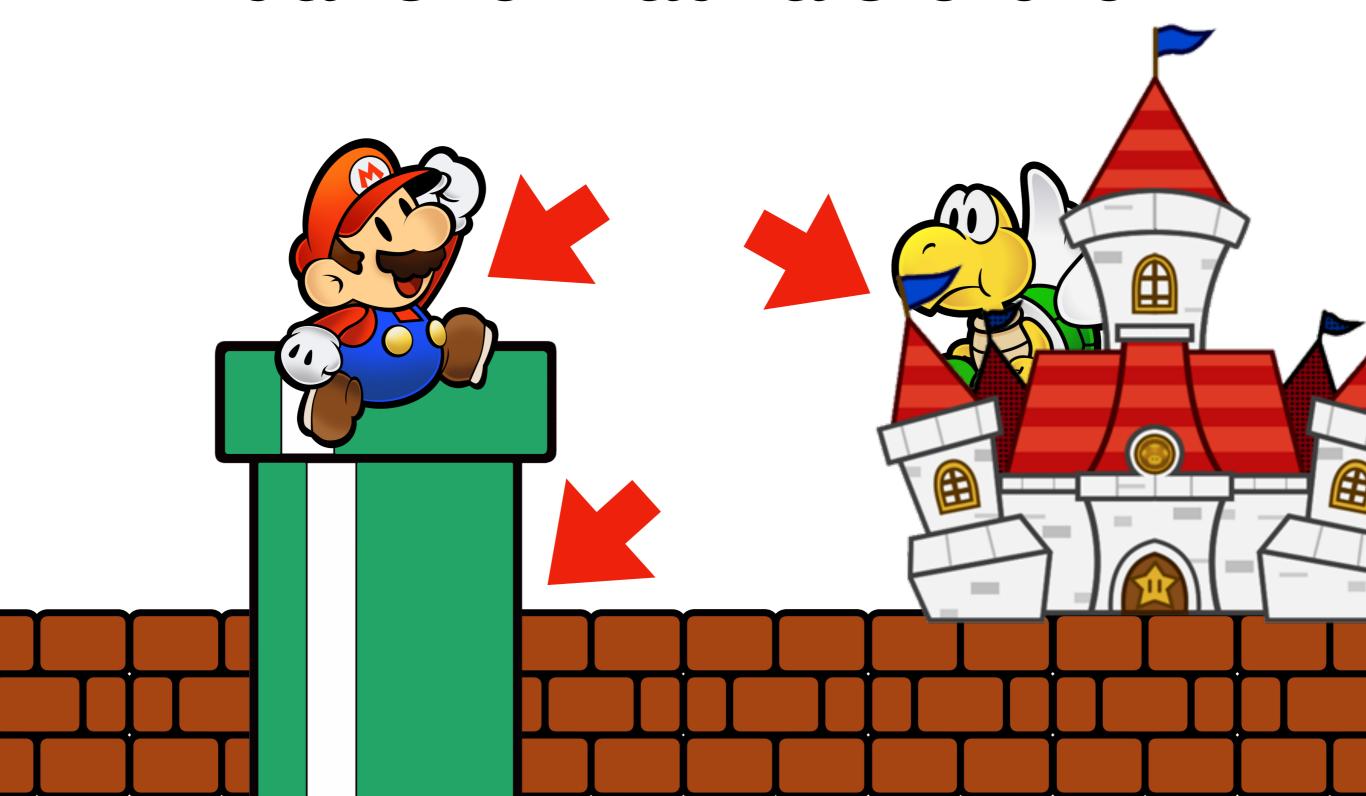


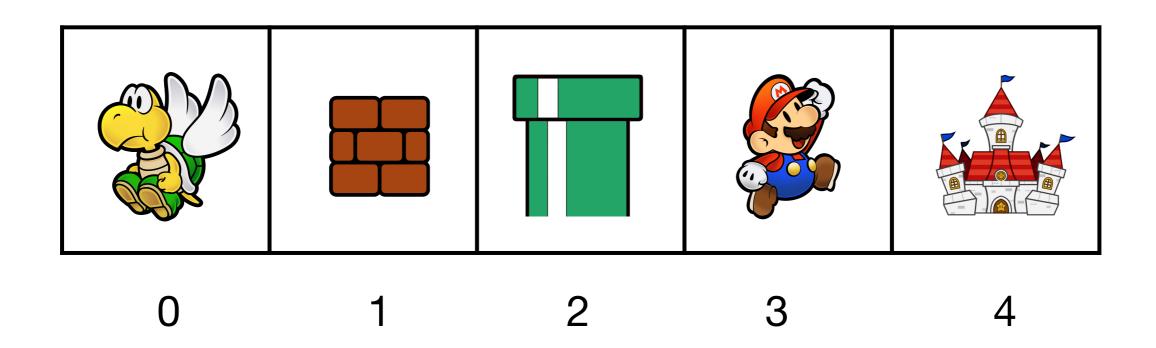


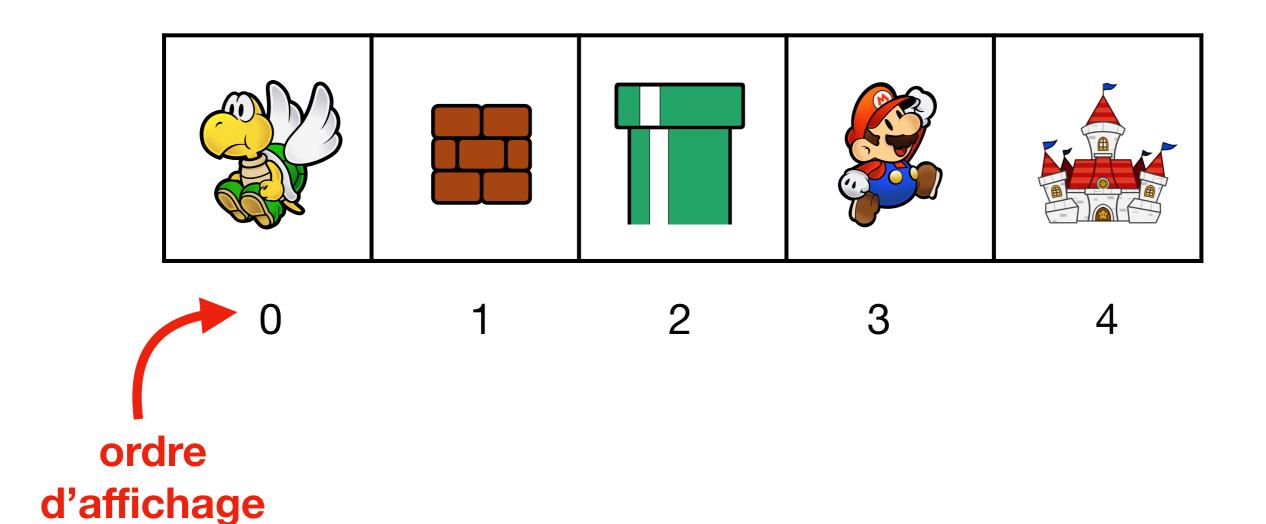


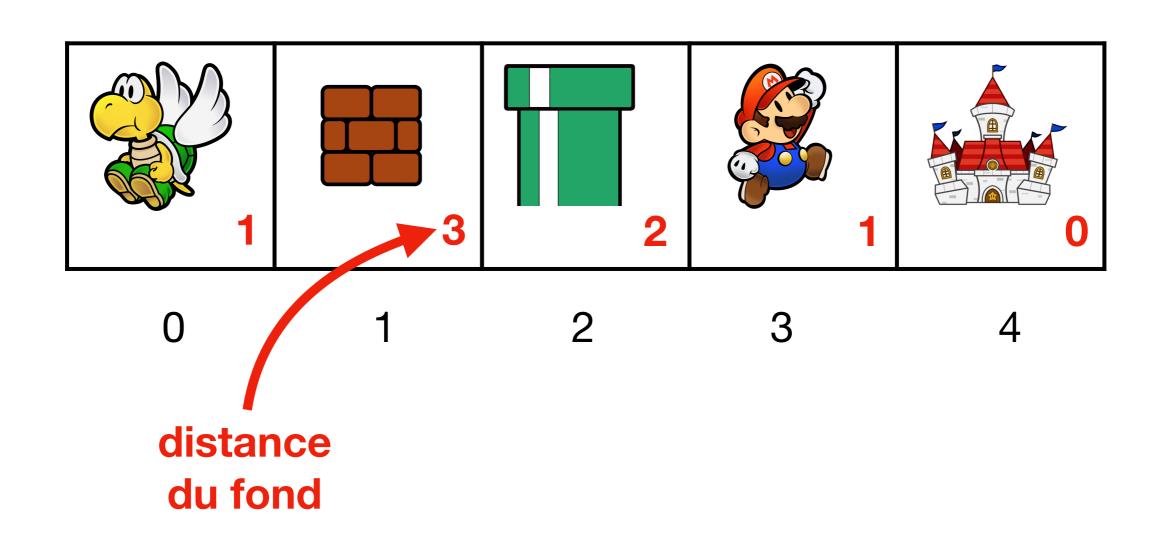


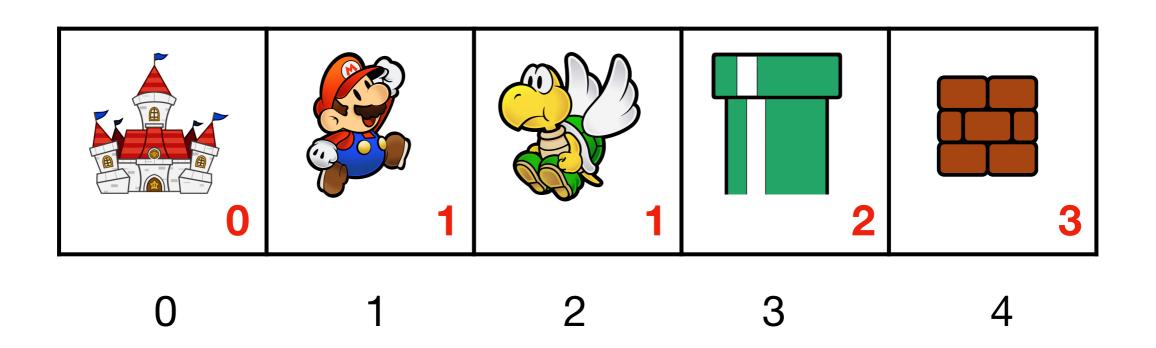












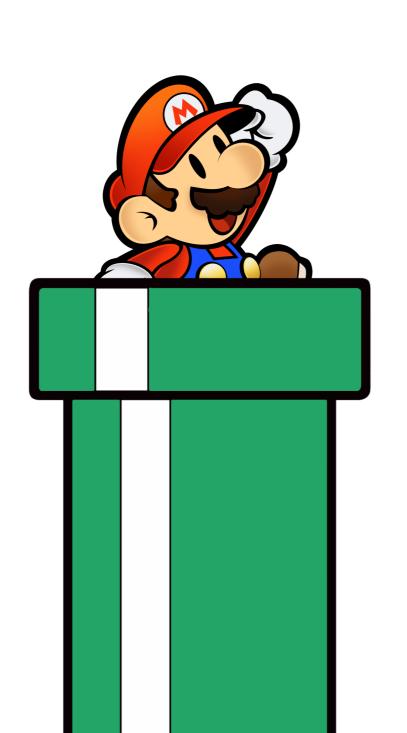


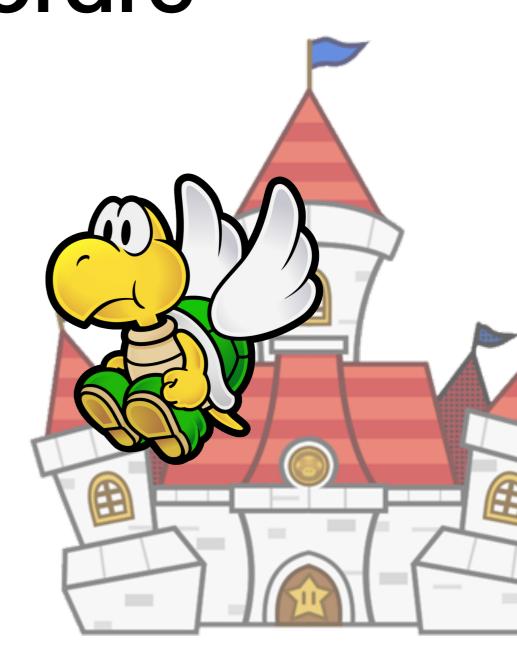


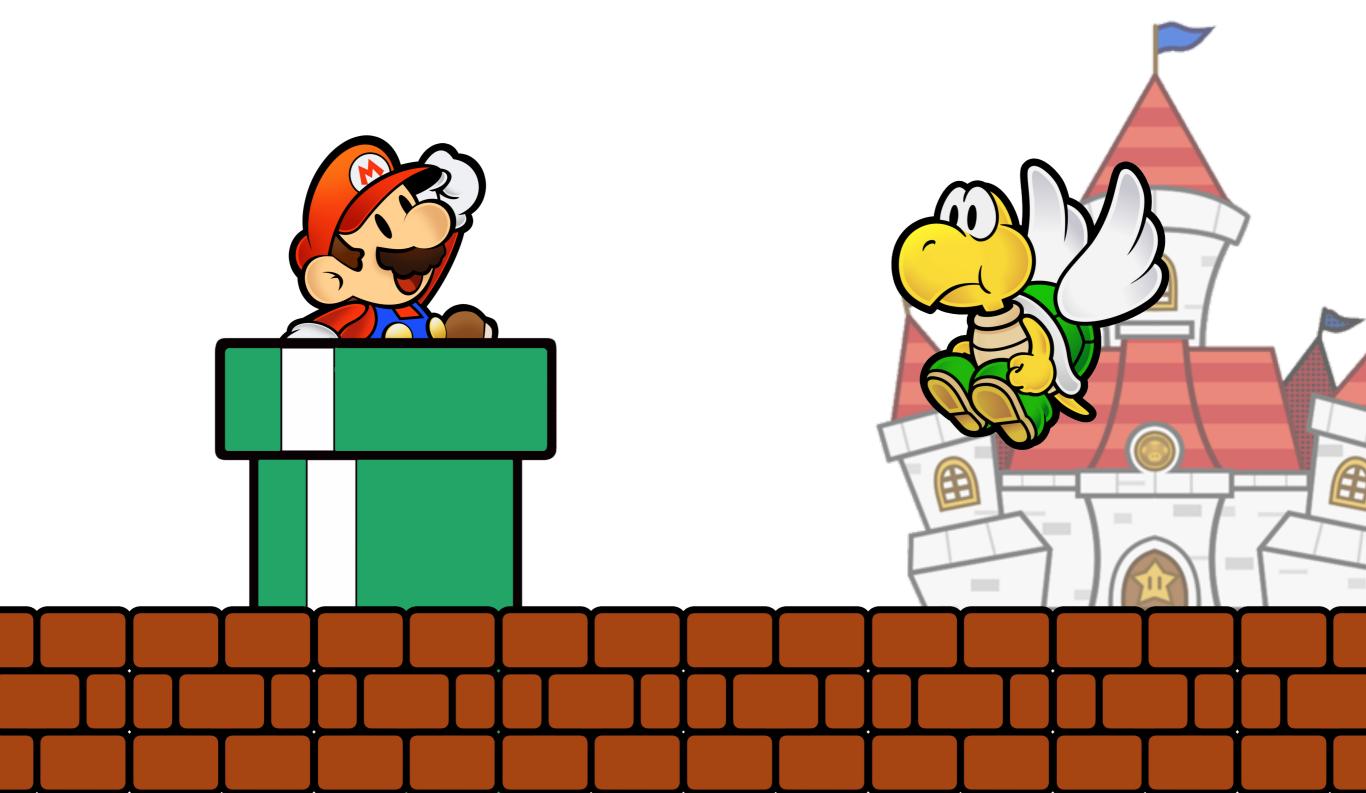


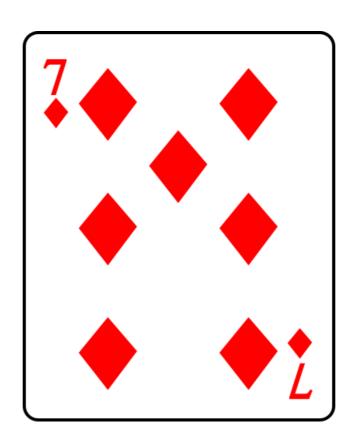


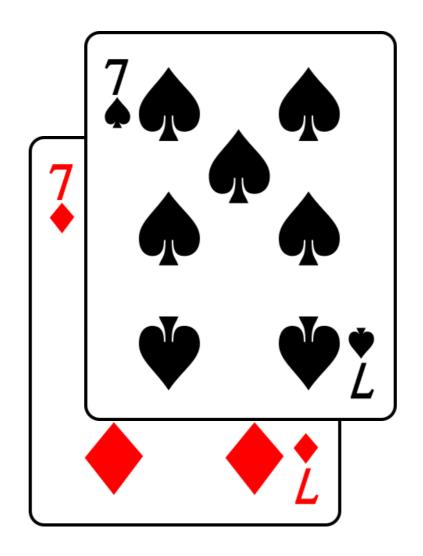


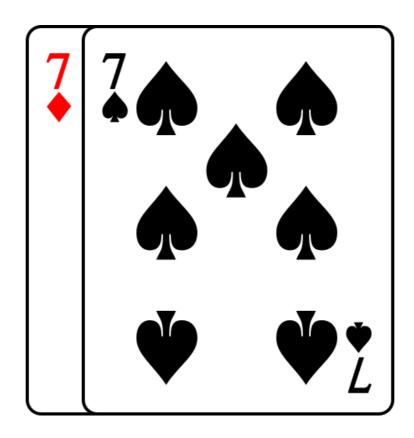


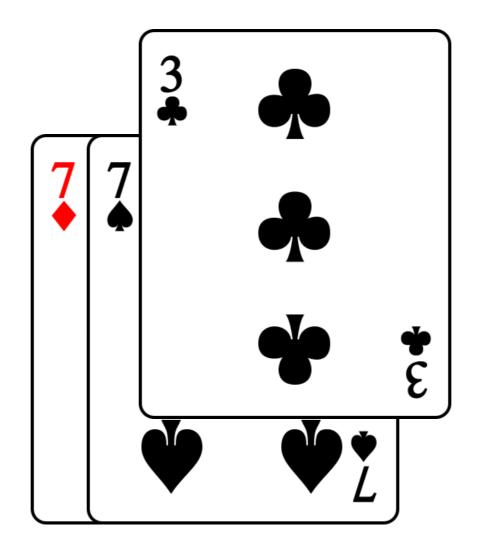


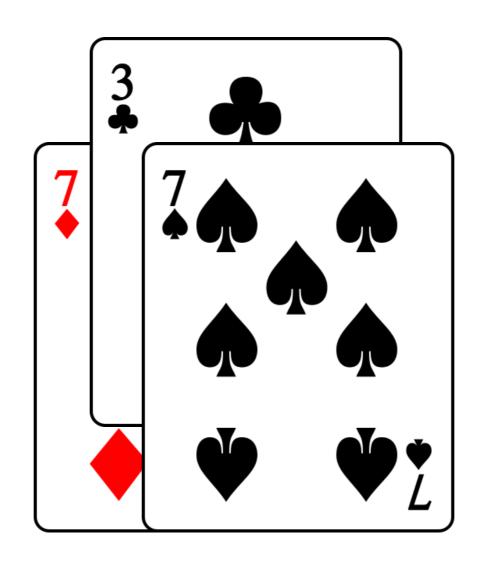


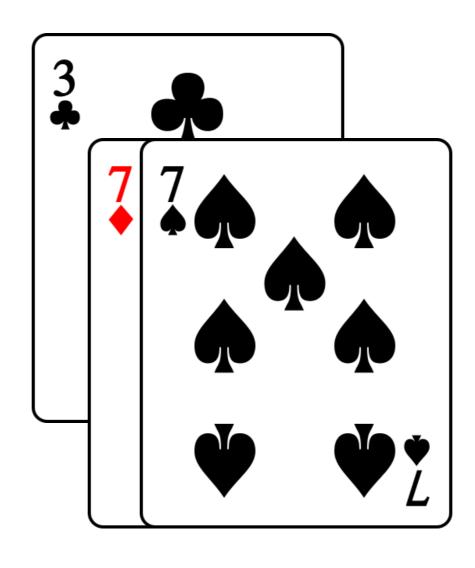


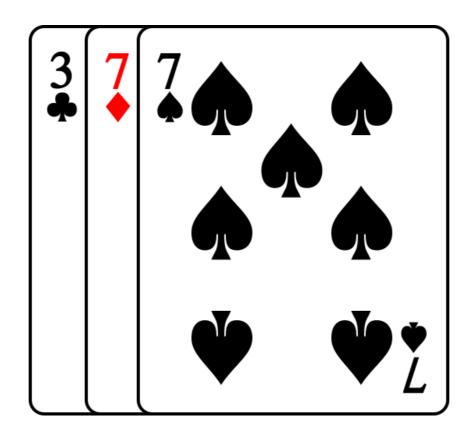


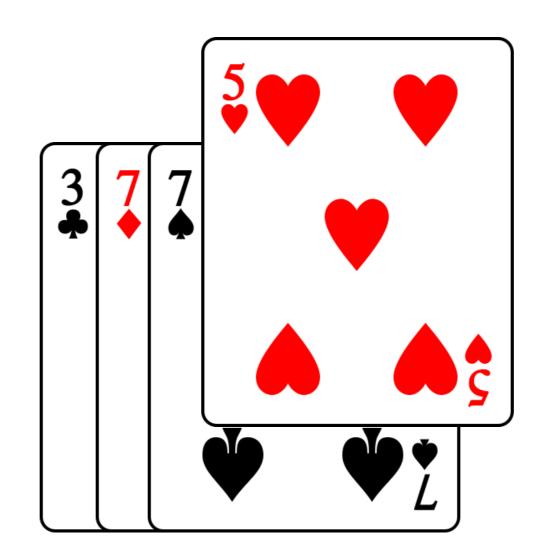


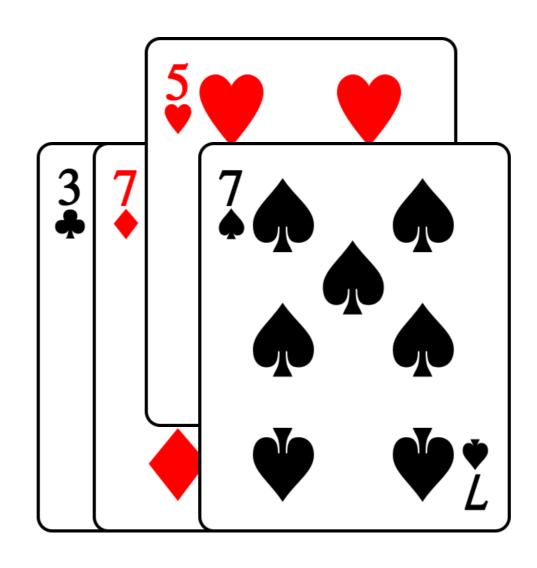


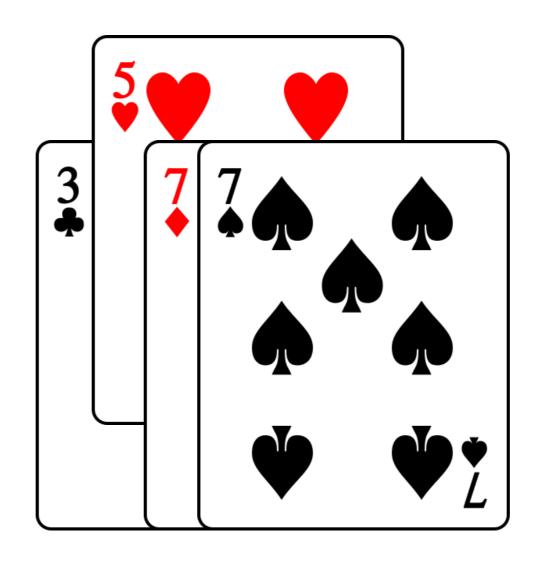


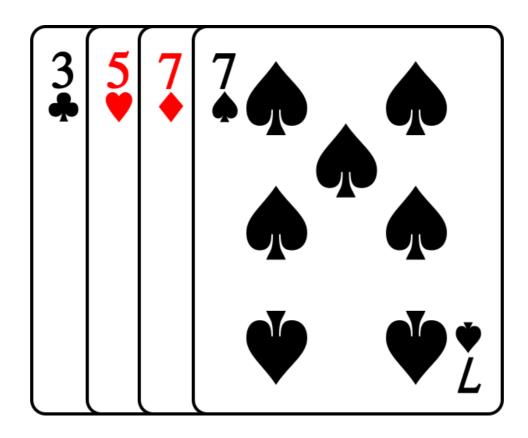


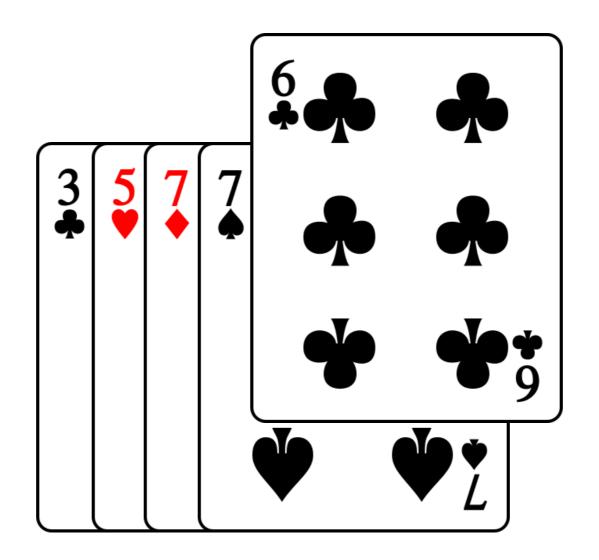


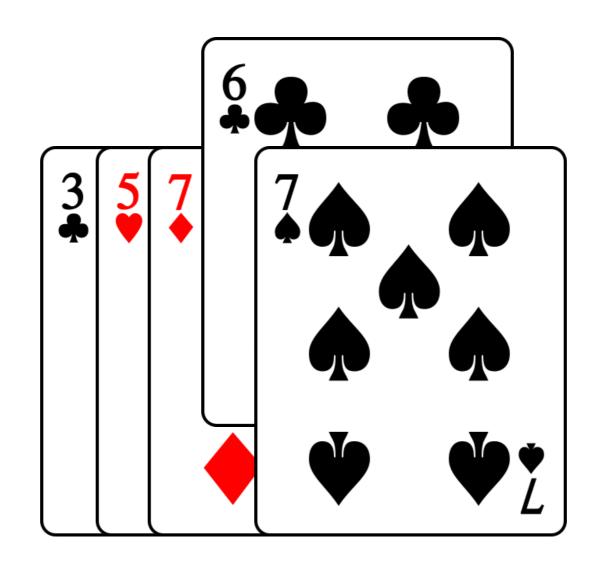


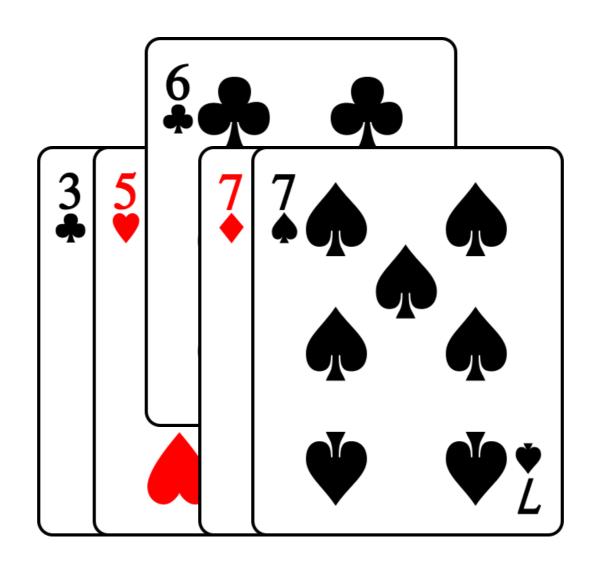


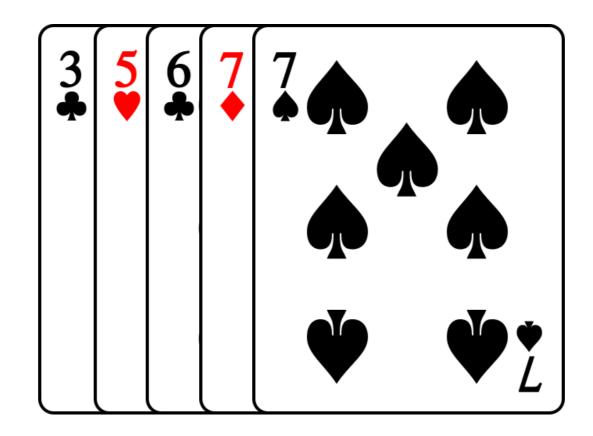


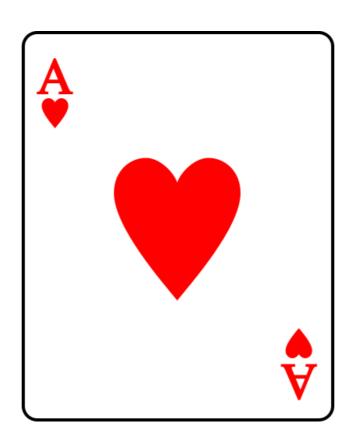


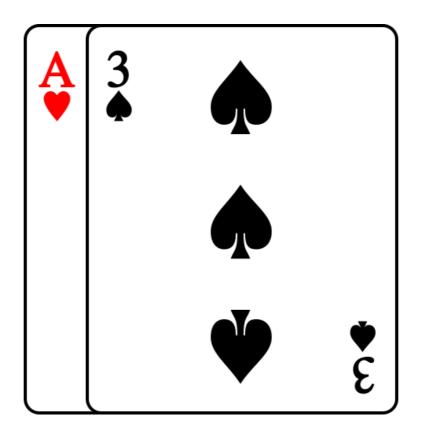


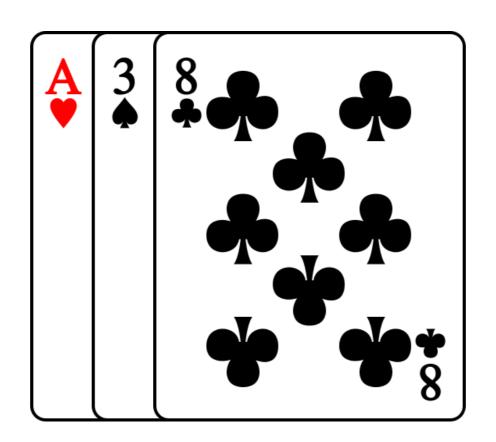


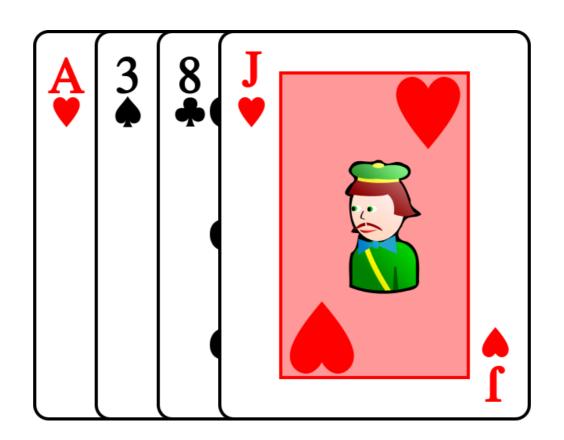


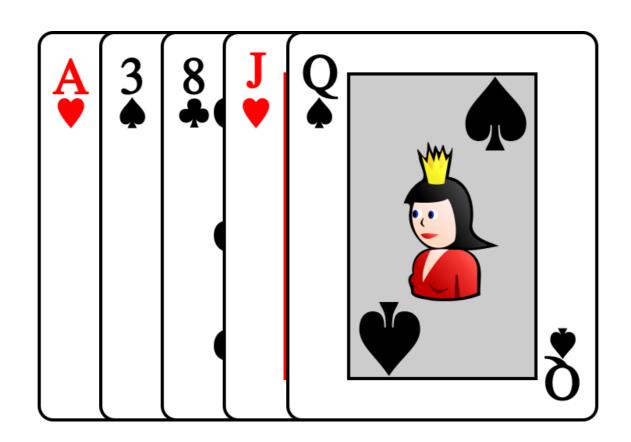




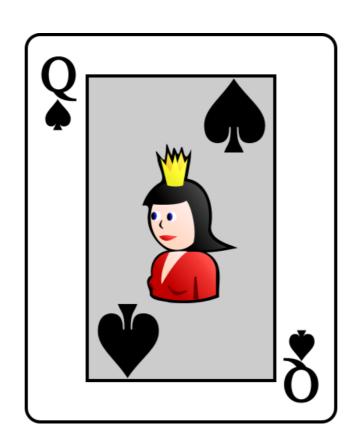


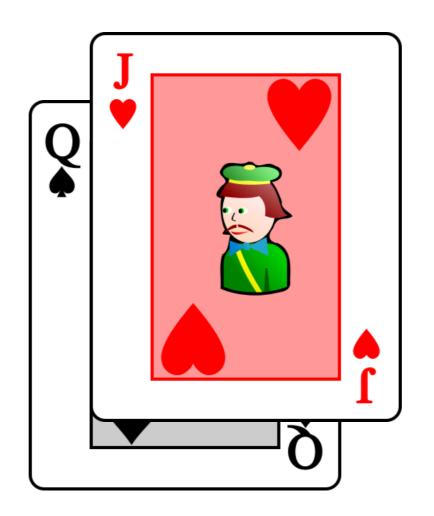


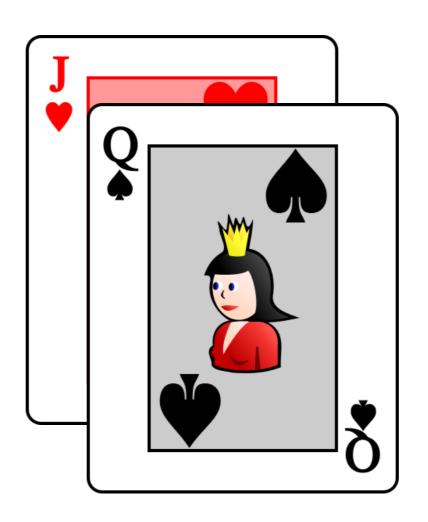


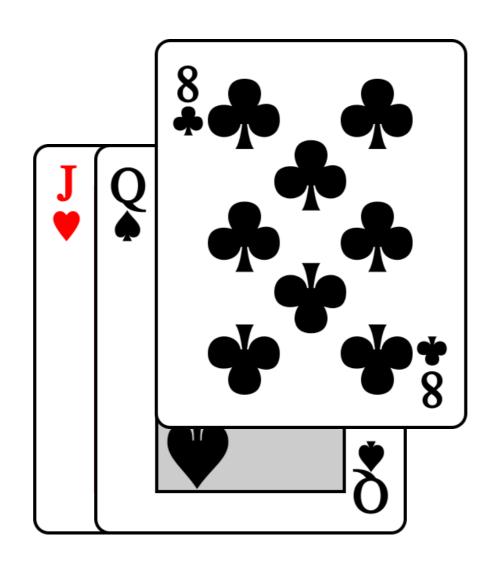


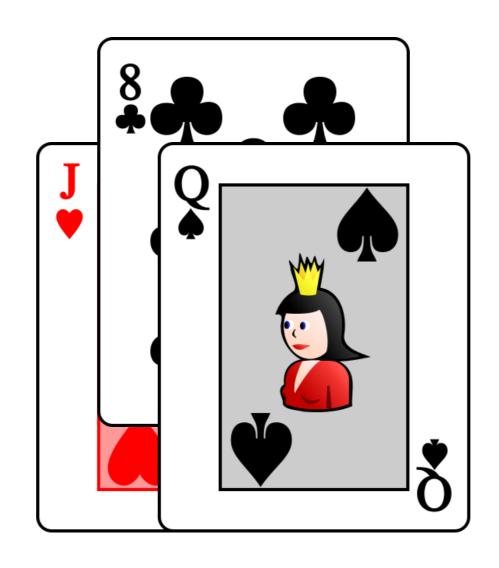
- Les cartes arrivent déjà triées
- On fait n opérations (déplacements de cartes)

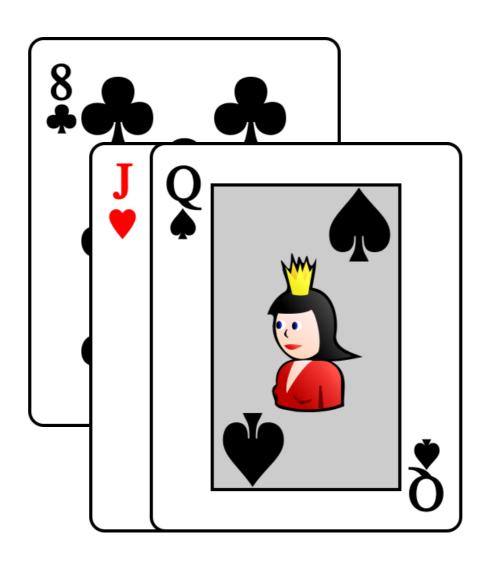


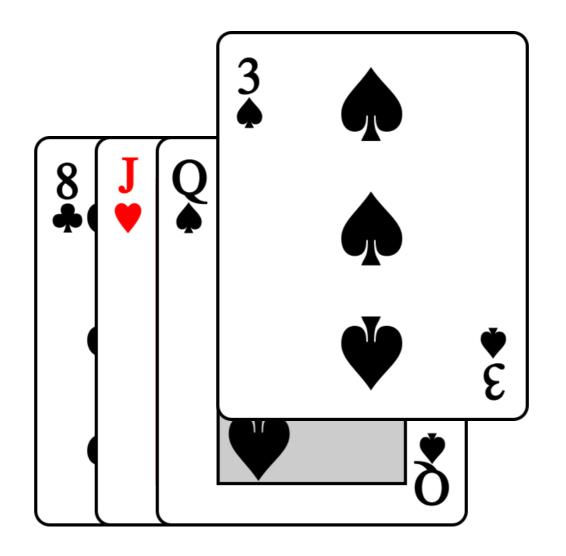


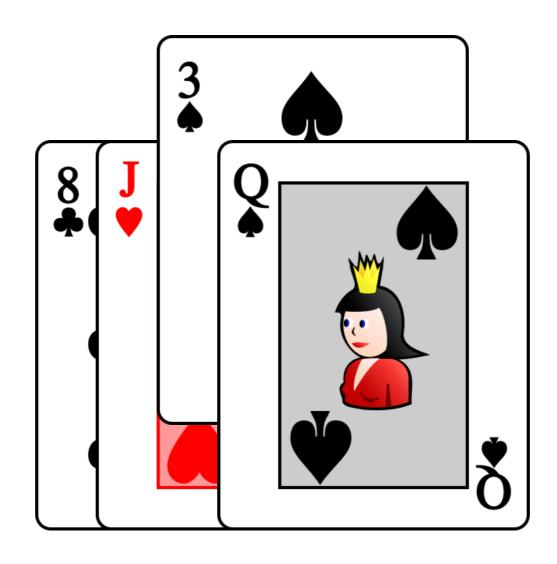


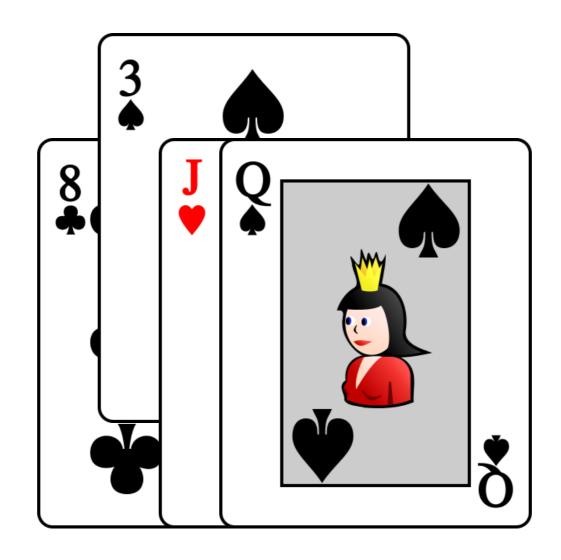


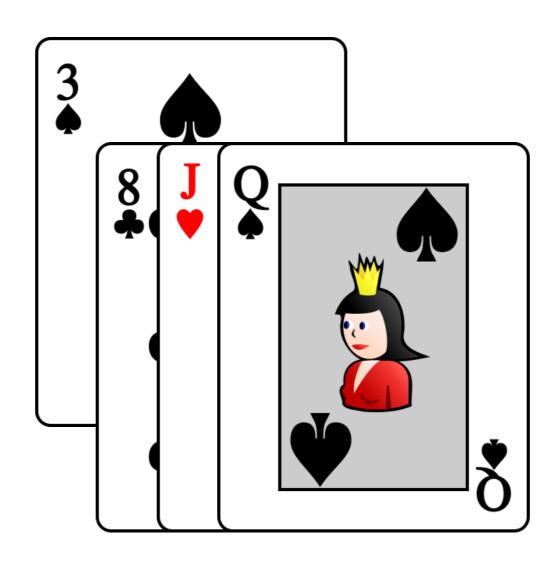


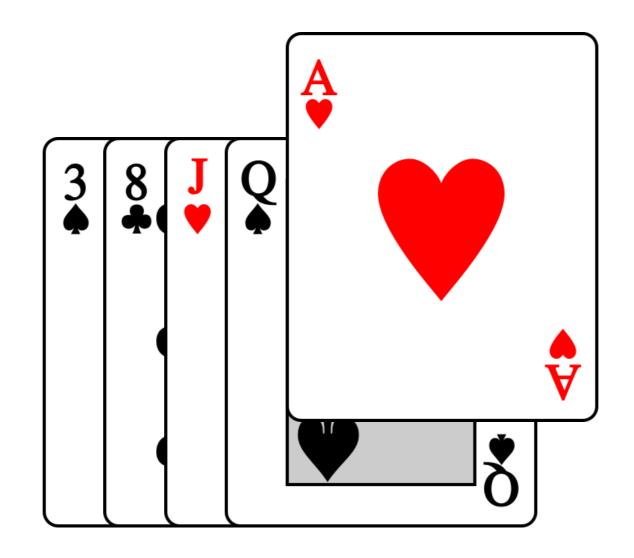


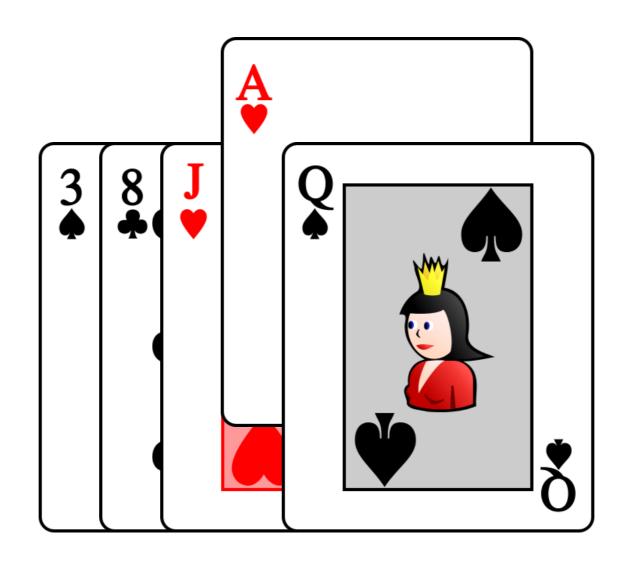


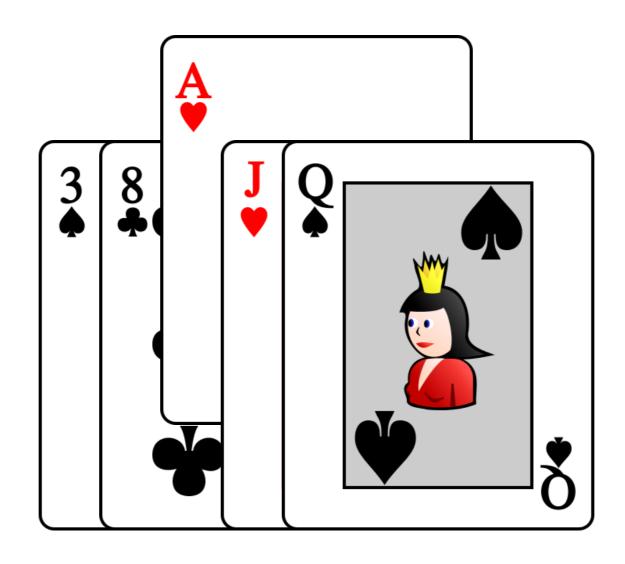


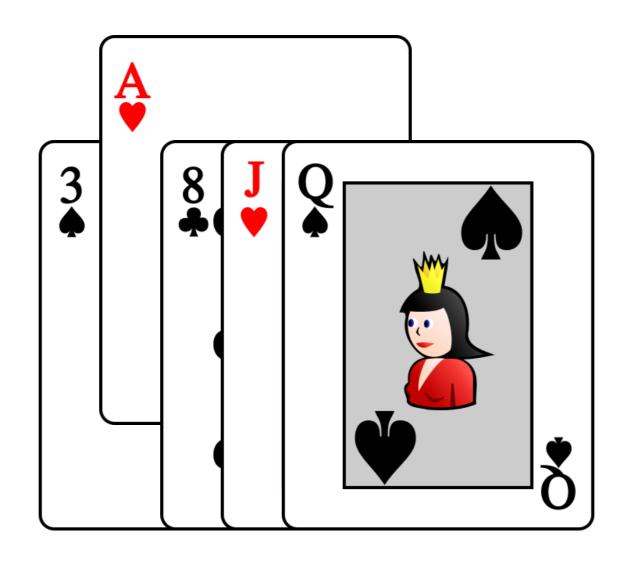


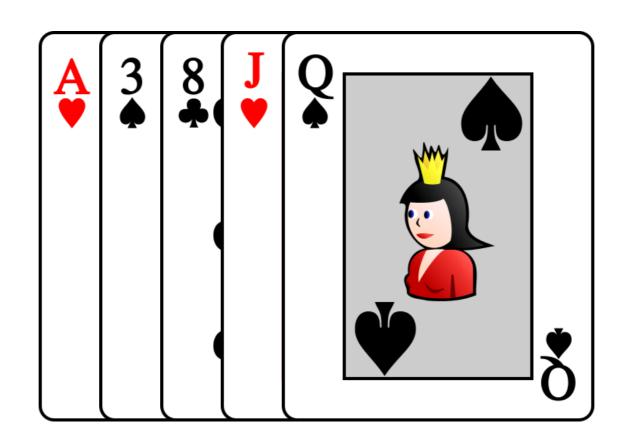












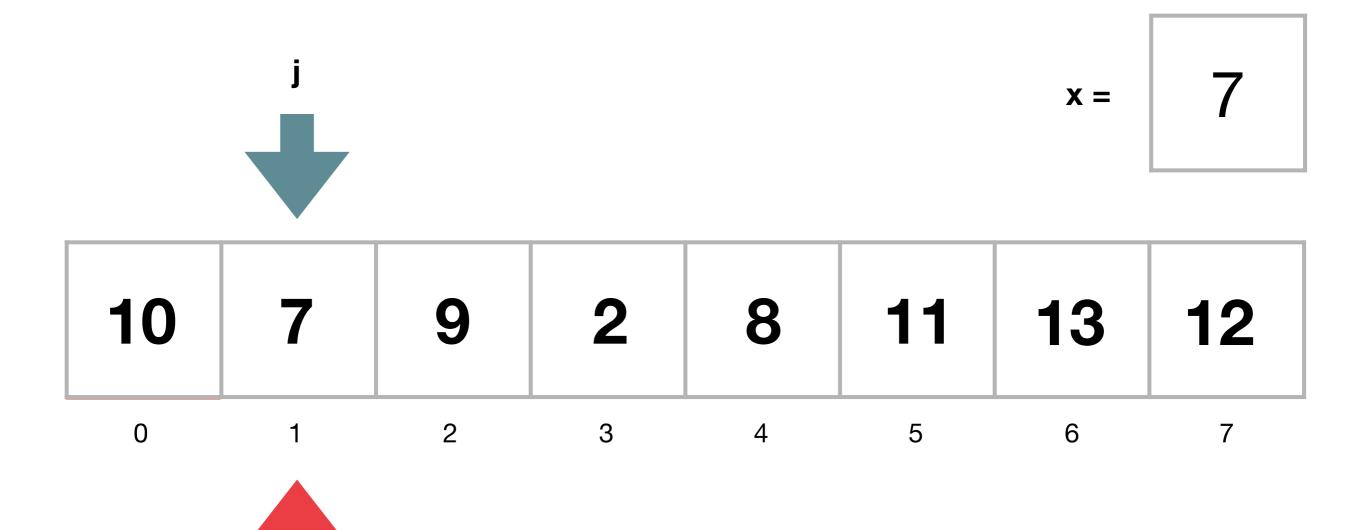
- Les cartes arrivent en ordre décroissant
- On fait i opérations pour la i-ème carte
- Le nombre totale est $1 + 2 + 3 + \cdots + n$

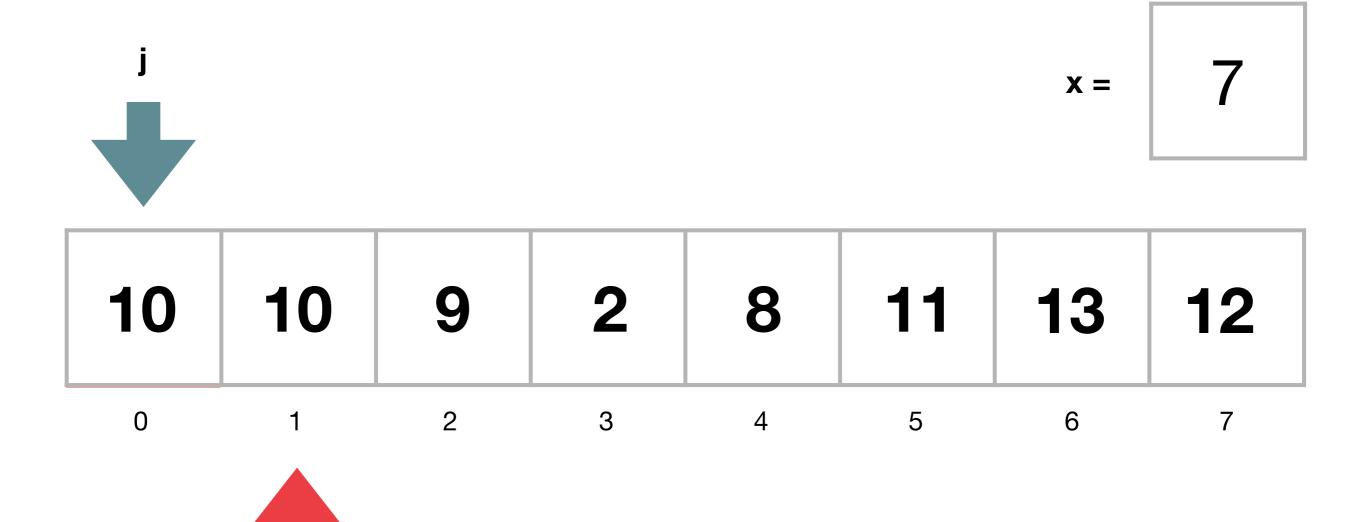
$$\sum_{i=1}^{n} i = \frac{1}{2}n(n+1) = \frac{1}{2}(n^2 + n) \in O(n^2)$$

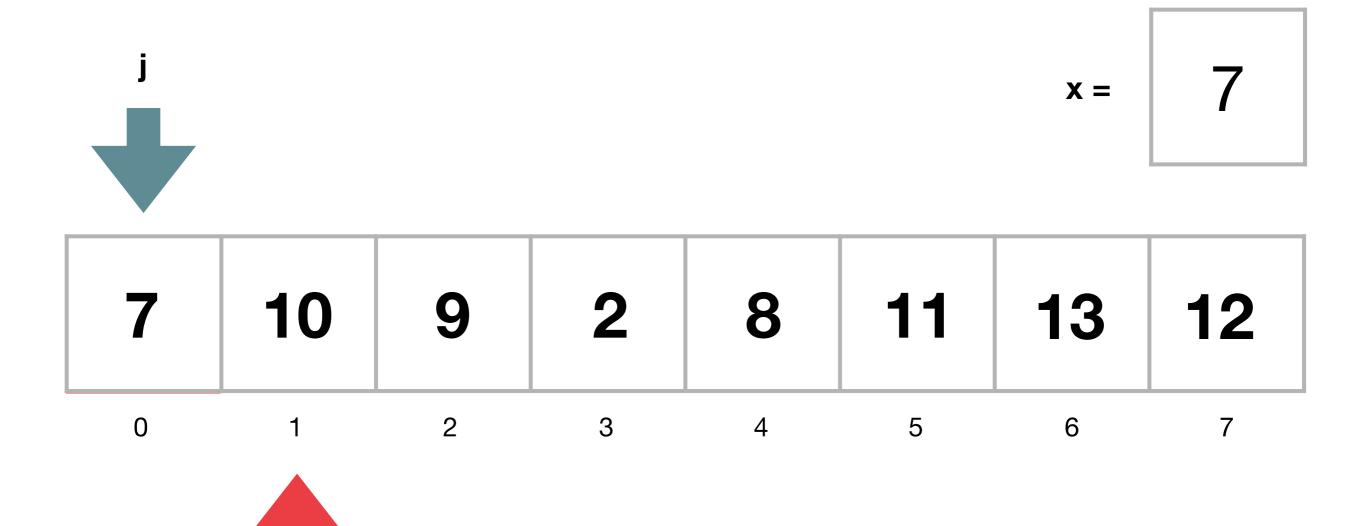
Tri d'un tableau par insertion

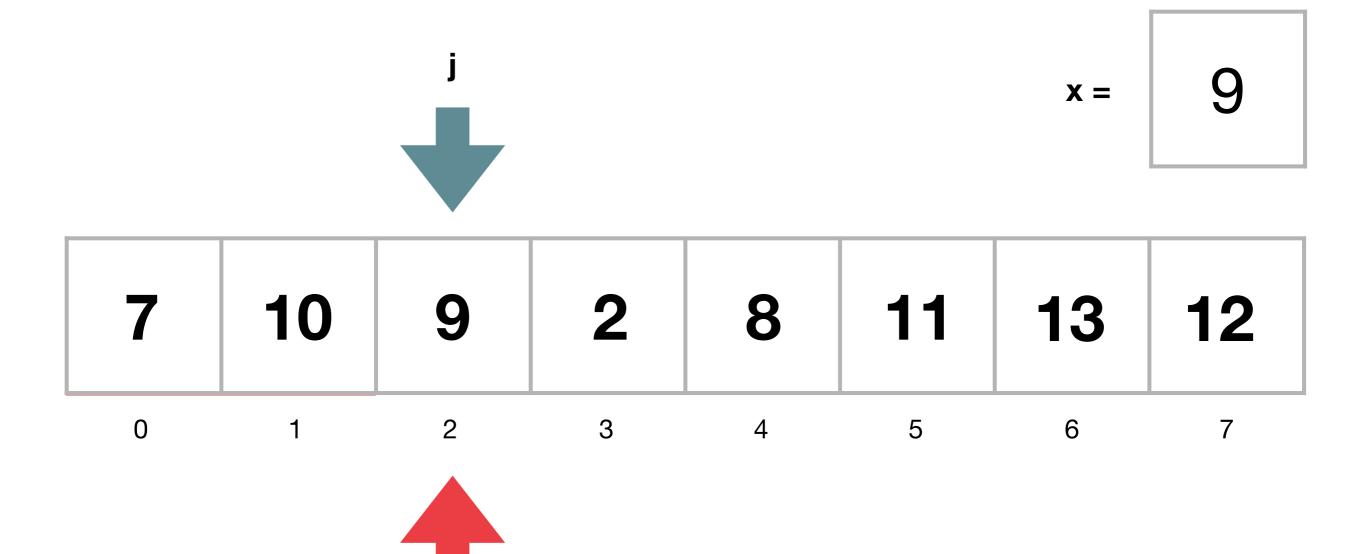
```
procedure trier-par-insertion(T)
   n = longueur(T)
   pour i = 1 \hat{a} n - 1 \text{ faire } (on saute le 0)
       x = T[i]
       i := i
       tant que j > 0 et x < T[j - 1] faire
           (décaler d'un élément)
           T[i] := T[i - 1]
           j ≔ j – 1
       fin tant que
       (ici x \ge T[j-1] ou bien j=0)
       T[i] = X
   fin pour
fin procedure
```

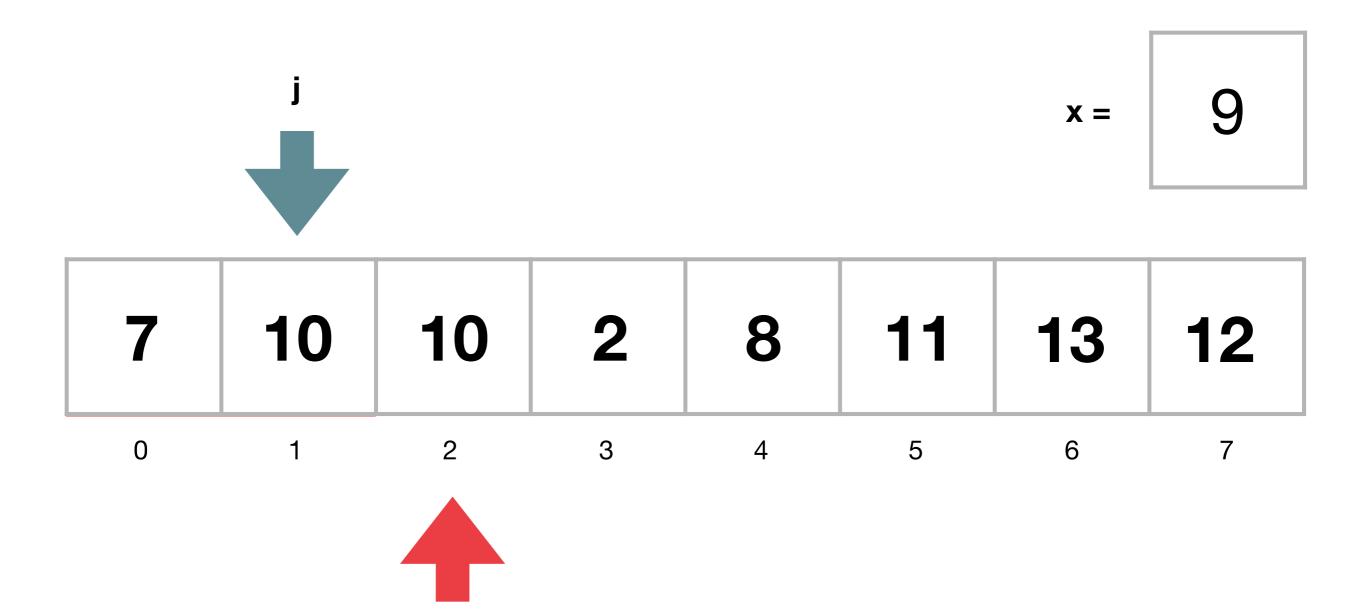
10	7	9	2	8	11	13	12
0	1	2	3	4	5	6	7

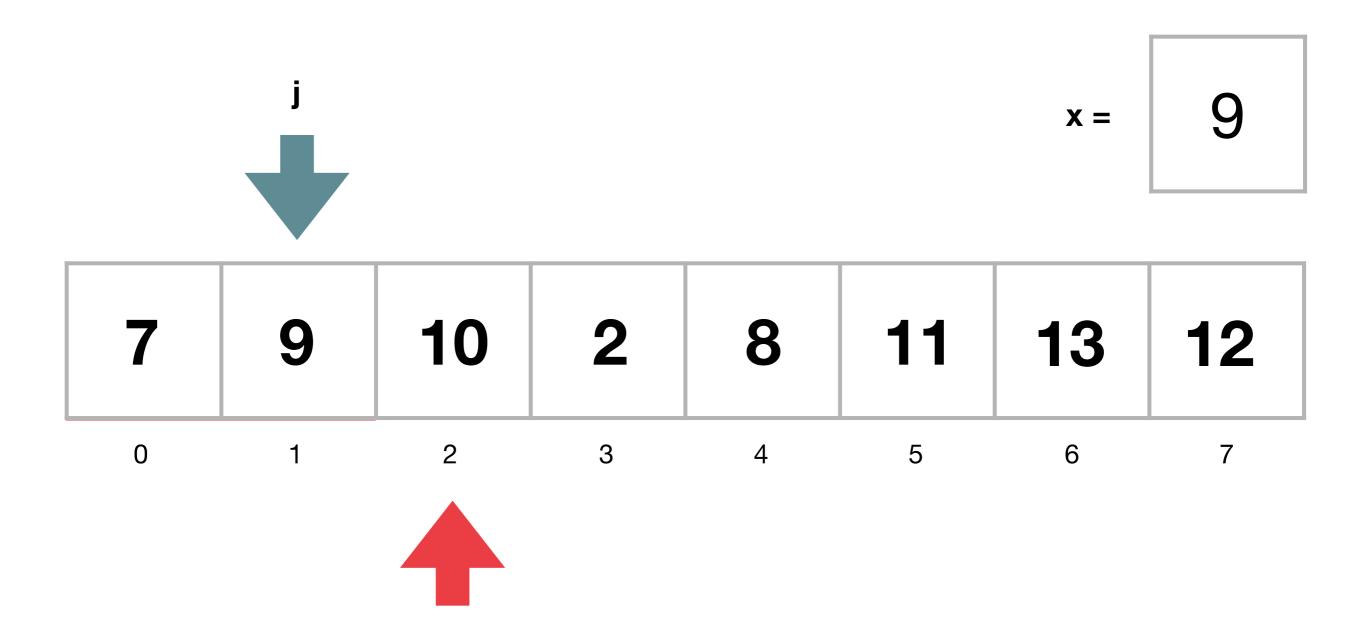


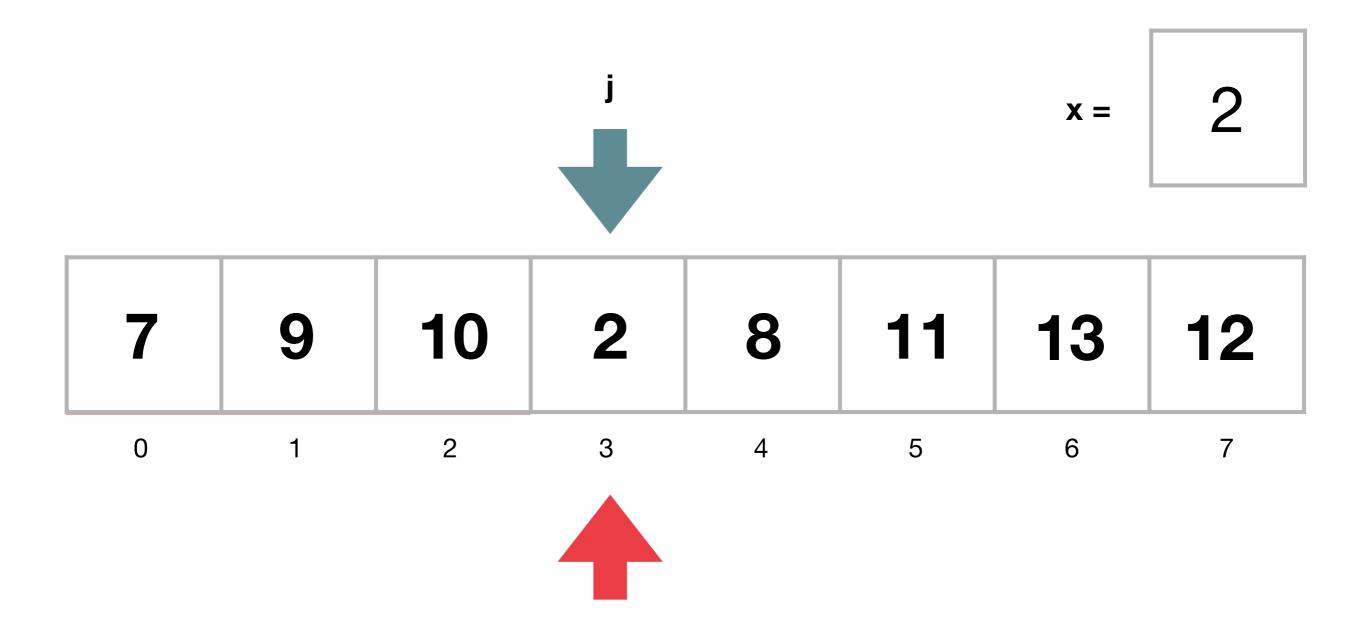


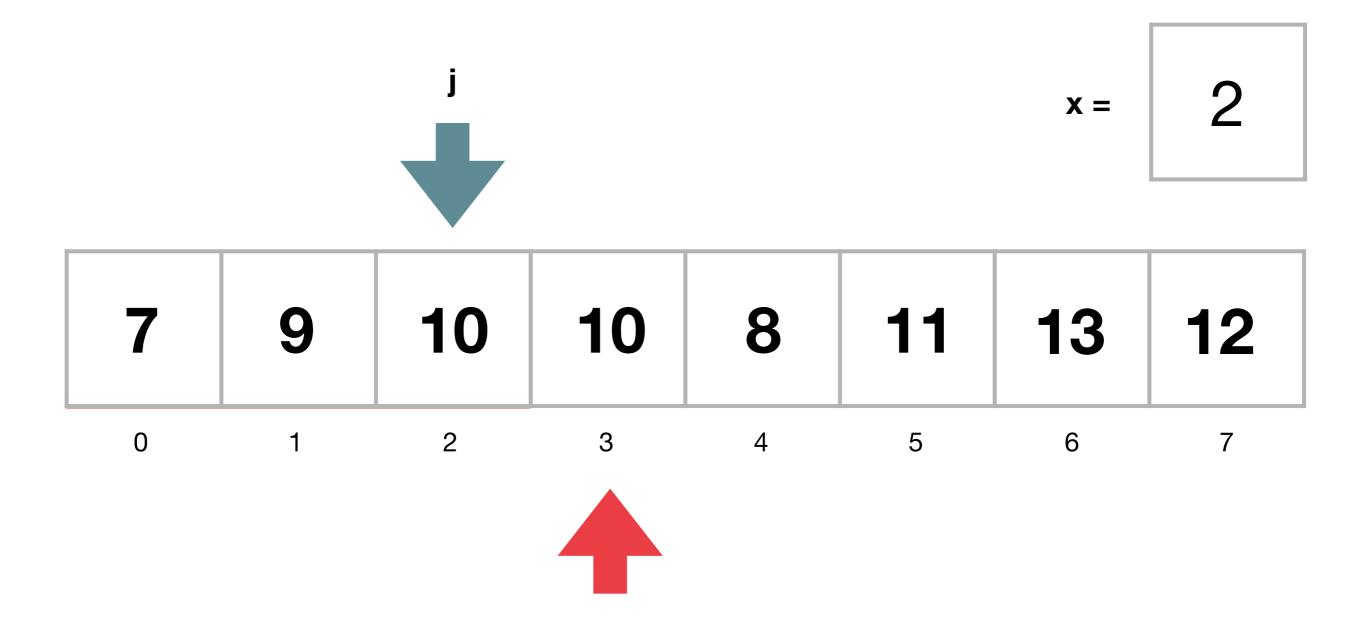


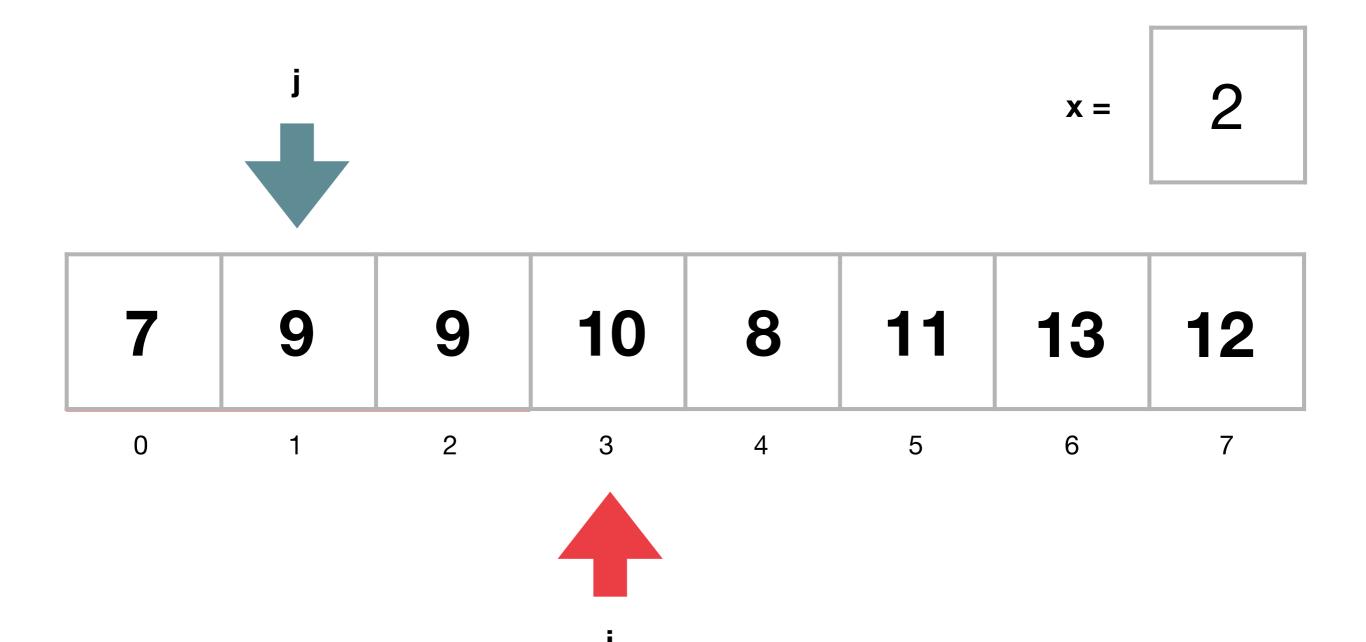


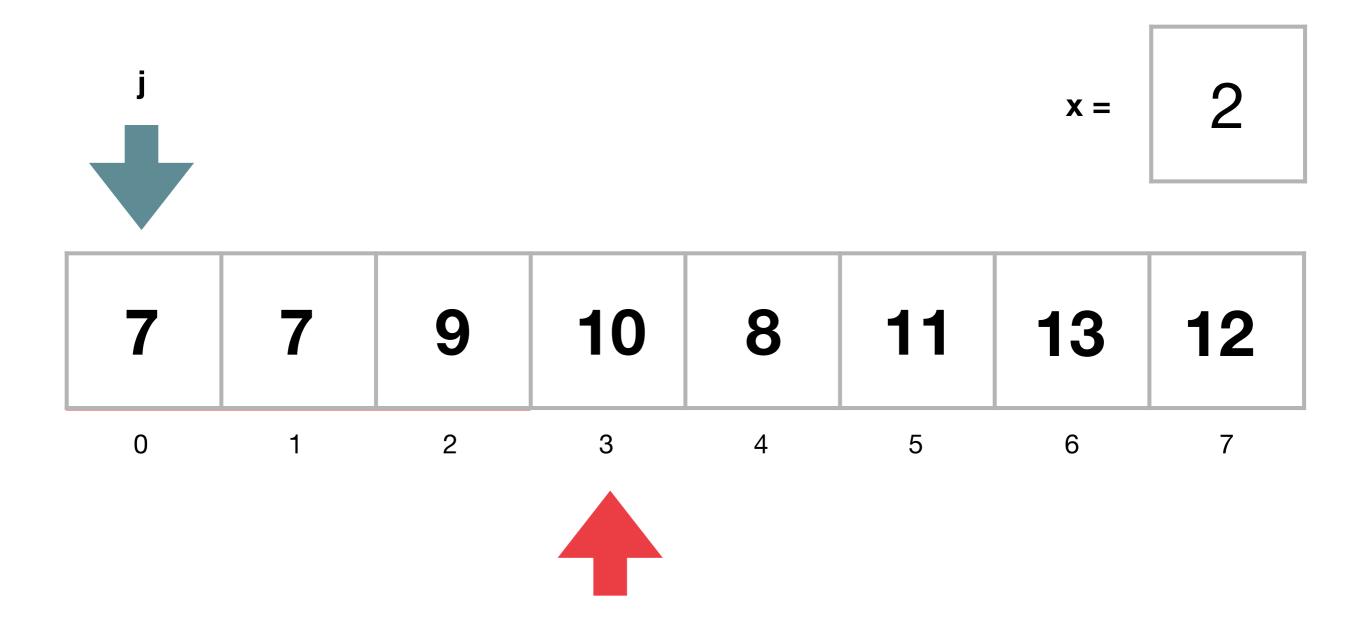


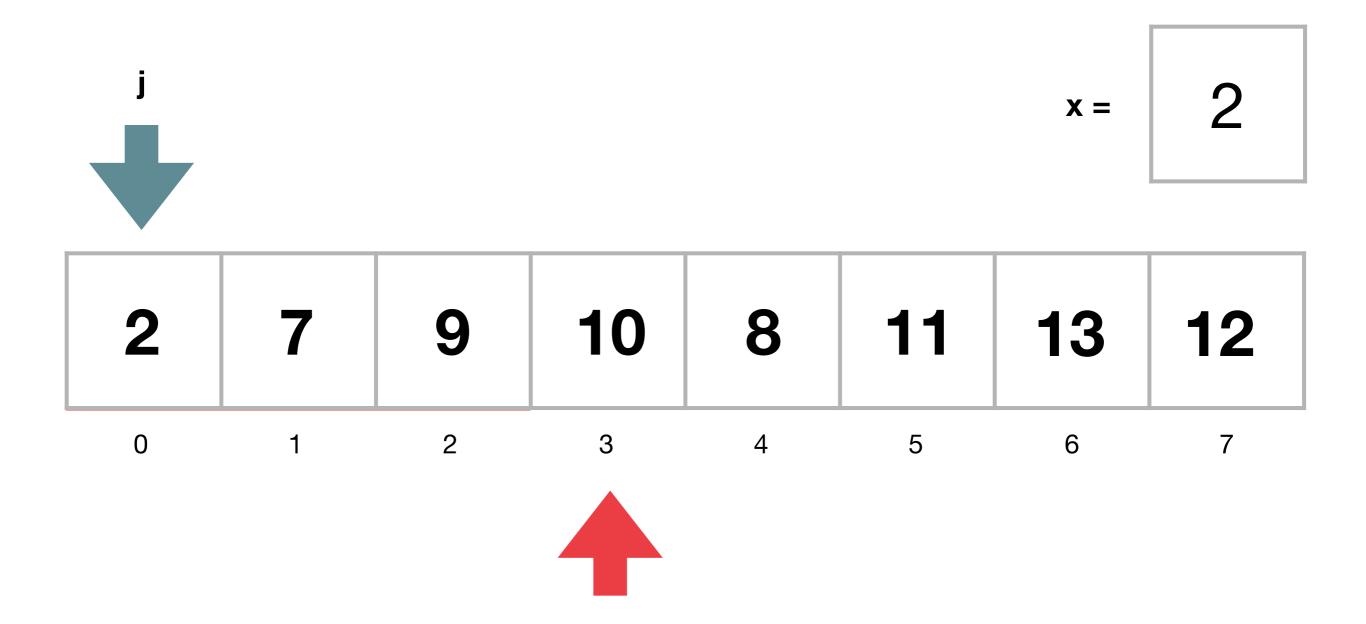


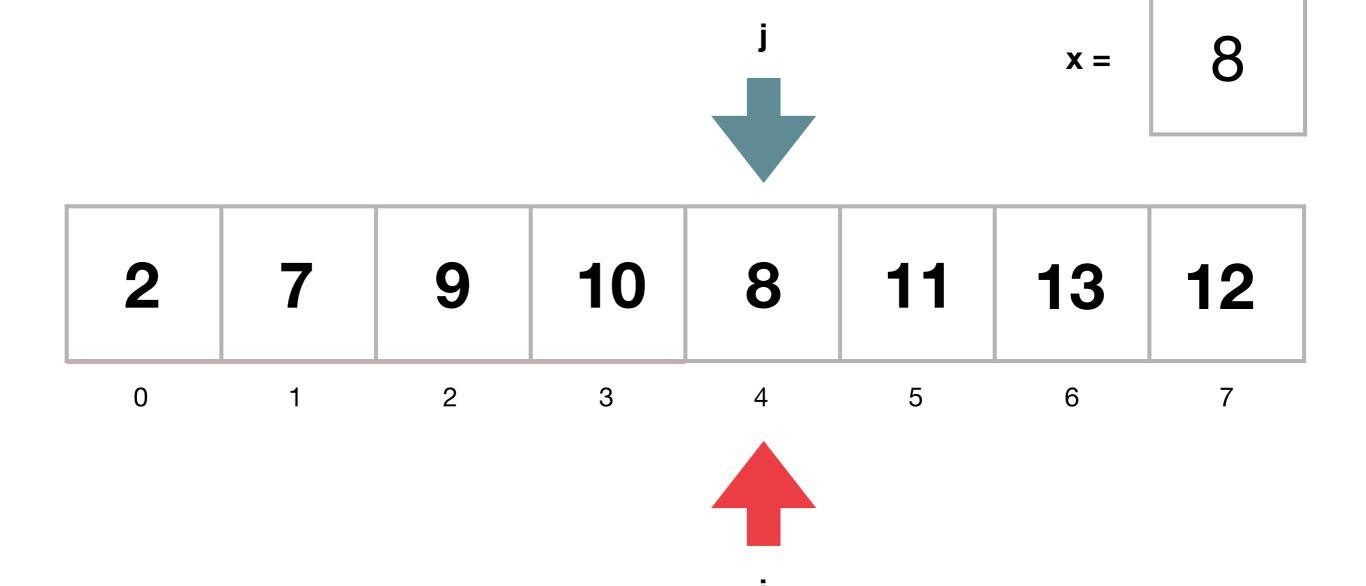


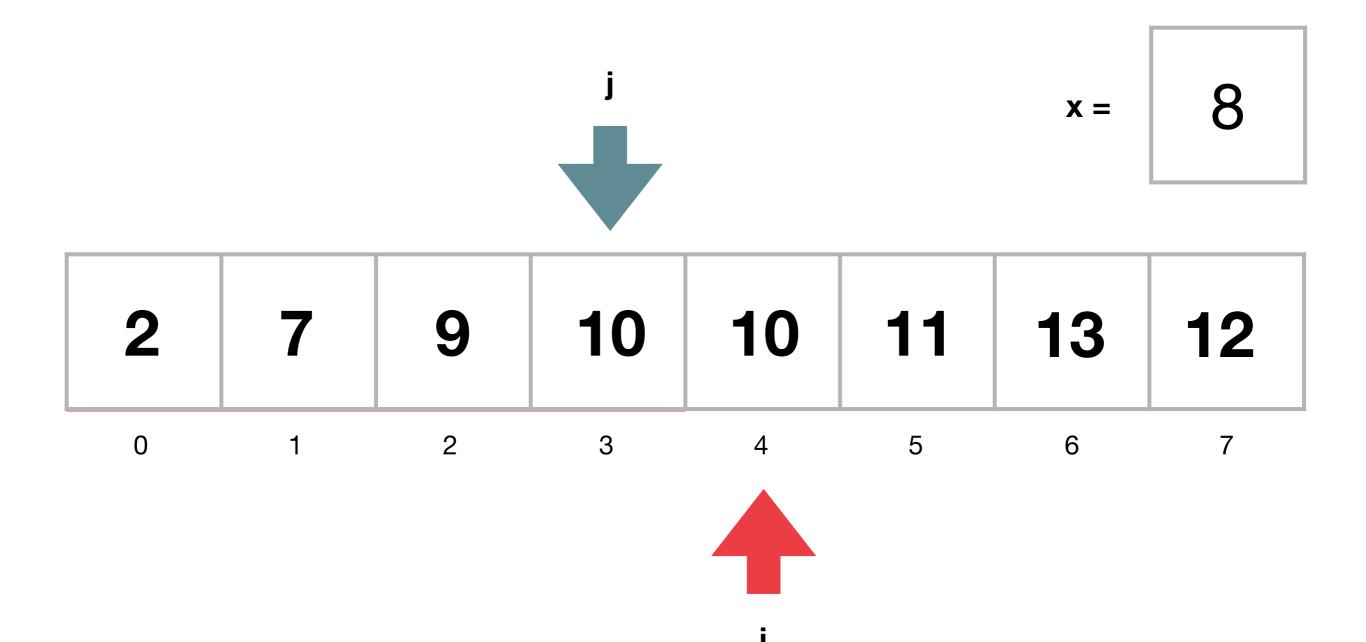


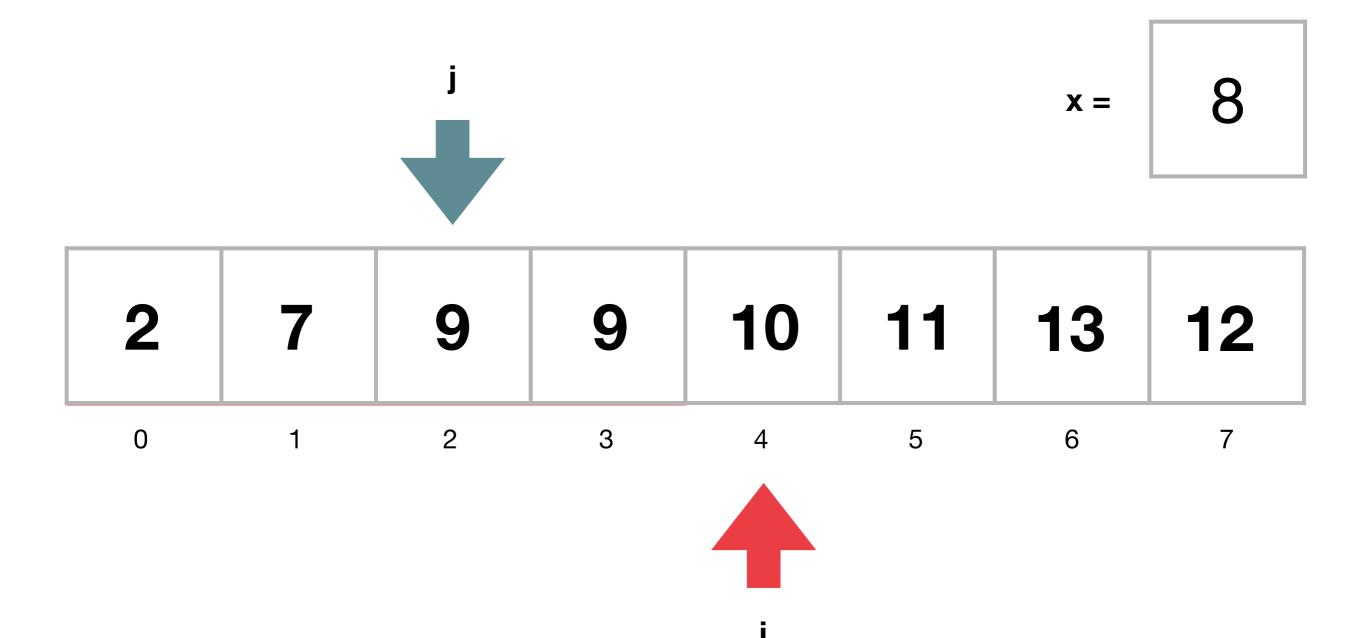


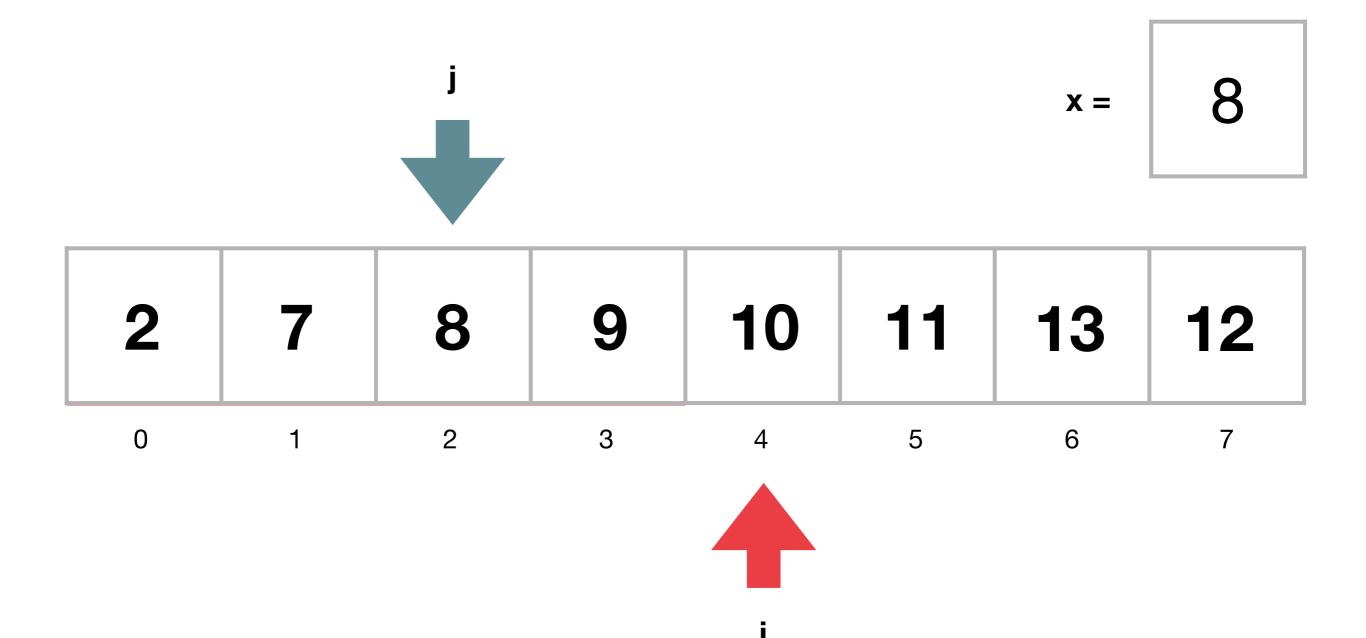




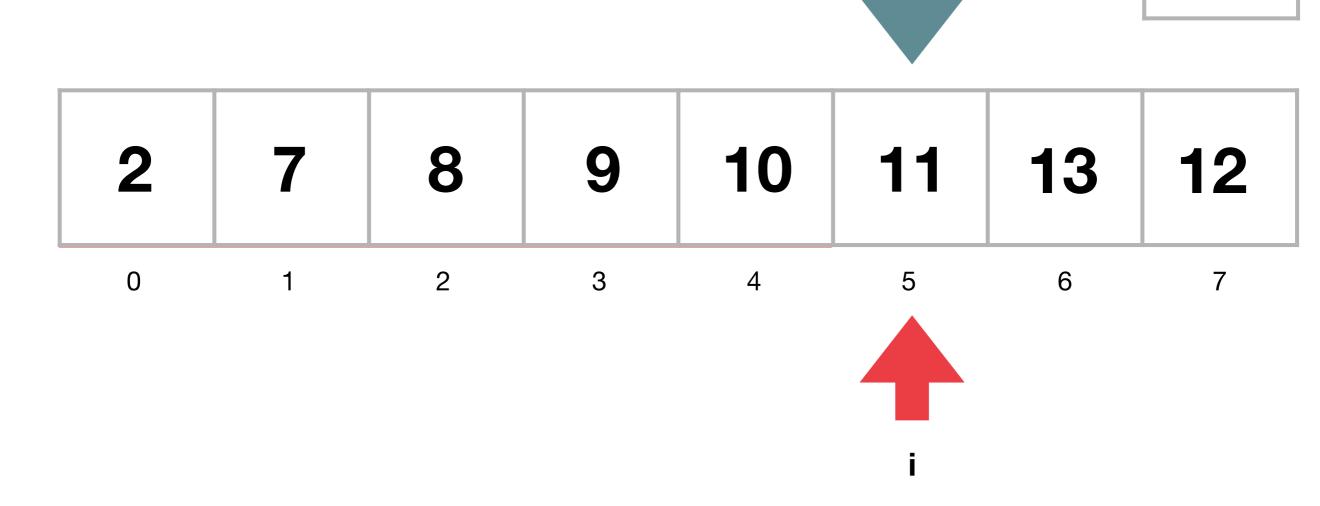


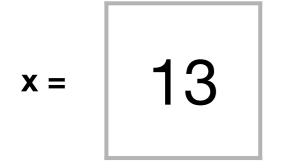






x =



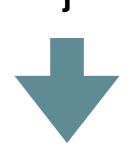




2	7	8	9	10	11	13	12
0	4	0	0	4	5	0	7







2	7	8	9	10	11	13	12
0	1	2	3	4	5	6	7

2 3 4 5 6 7

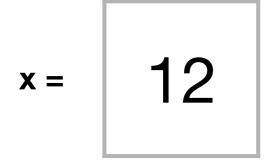






2	7	8	9	10	11	13	13
0	1	2	3	4	5	6	7

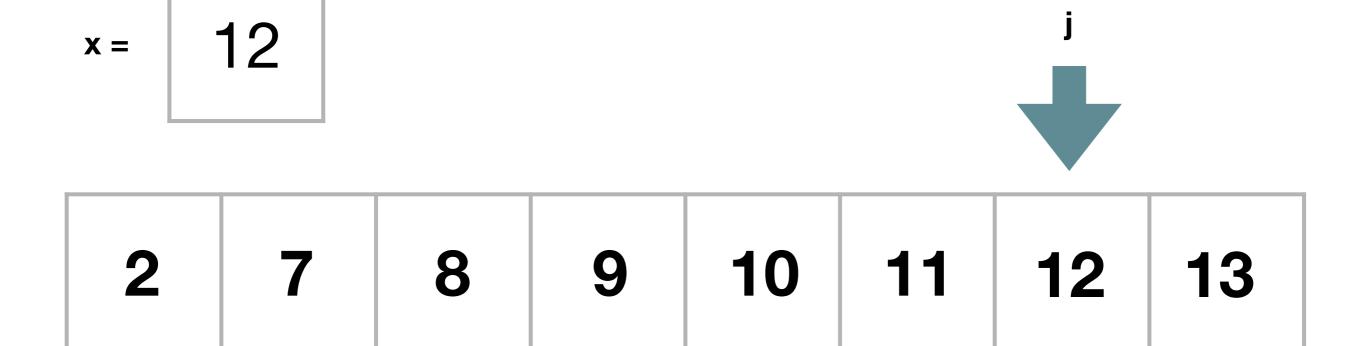






2	7	8	9	10	11	12	13
0	1	2	3	4	5	6	7





0

Terminaison

```
procedure trier-par-insertion(T)
   n = longueur(T)
    pour i = 1 \hat{a} n - 1 faire (on saute le 0)
       x = T[i]
       i := i
        tant que j > 0 et x < T[j - 1] faire
            (décaler d'un élément)
            T[i] = T[i-1]
           j = j - 1
        fin tant que
        (ici x \ge T[j-1] ou bien j=0)
       T[i] = x
   fin pour
fin procedure
```

- La boucle pour termine toujours
- La boucle tant que termine (au pire) quand j = 0

Correction

```
procedure trier-par-insertion(T)
    n = longueur(T)
    pour i = 1 \hat{a} n - 1 \text{ faire } (on saute le 0)
        x = T[i]
        i := i
        tant que j > 0 et x < T[j-1] faire
            (décaler d'un élément)
            T[i] = T[i-1]
            j = j - 1
        fin tant que
        (ici x \ge T[j-1] ou bien j=0)
        T[i] = x
    fin pour
fin procedure
```

Le sous-tableau
 T[0, ..., i – 1] est trié au début de la boucle pour

Efficacité

```
procedure trier-par-insertion(T)
    n = longueur(T)
    pour i = 1 \hat{a} n - 1 \text{ faire } (on saute le 0)
        x = T[i]
        i := i
        tant que j > 0 et x < T[j - 1] faire
            (décaler d'un élément)
            T[i] = T[i-1]
            j = j - 1
        fin tant que
        (ici x \ge T[i-1] ou bien j=0)
        T[i] = x
    fin pour
fin procedure
```

- O(n) opérations dans le meilleur des cas
- O(n²) opérations dans le pire des cas