

## 2

### 2.1

Since the group is finite, we can generate the powers of an element  $s \in S$  infinitely. However there is only a finite amount of elements in  $S$ . Thus, there exists an integer  $x$  such that  $s^x = s$ . If  $x = 2$  we are done, else we multiply both sides with  $s^{x-2}$ . Then we substitute  $s^{(x-1)}$  with  $y$  and the result is  $y^2 = y'$ .

$$\begin{aligned} s^2 \cdot s^{x-2} &= s \cdot s^{x-2} \\ (s^x \cdot s^{-1}) \cdot (s^x \cdot s^{-1}) &= s^{x-1} \\ s^{x-1} \cdot s^{x-1} &= s^{x-1} \\ (s^{x-1})^2 &= s^{x-1} \end{aligned}$$

□

### 2.2

First we proof  $g, g' \in G, g \sim_H g'$  if  $g = hg'$  is an equivalence relation by showing it is reflexive, symmetric and transitive.

Reflexivity: For all  $g \in G$  we have to prove that  $g = hg'$ , since  $H$  is a Subgroup of  $G$ , it contains the neutral Element of  $G$  therefore, we can always choose the neutral element for  $h$  with  $g = hg'$ . Thus  $g = g'$  and  $g \sim_H g$ .

Symmetry: Assume  $g \sim_H g'$  with  $g = hg'$ . Then we can add in inverse element  $h^{-1}$ , which is also in  $H$  because it is a group, resulting in  $h^{-1}g' = g$ . Hence,  $g' \sim_H g$  is also true.

Transitivity: Assume  $g \sim_H g'$  and  $g' \sim_H g''$  with  $g'' \in G$  and  $h_1, h_2 \in H : g = h_1g', g' = h_2g''$ . We can substitute  $g'$  resulting in  $g = h_1h_2g''$ . Since  $H$  is closed  $h_1h_2$  is also in  $H$ . Hence,  $g \sim_H g''$  is also true.

Since we proofed that  $\sim_H$  is a equivalence relation,  $G/H$  is partition induced by  $\sim_H$ . □

## 2.3

We proof if  $G$  is a permutation group on  $Q$  then  $G$  acts on  $Q$ . Therefore we show that  $G$  fullfills the two conditions

$$\forall q \in Q, g_1, g_2 \in G : q(g_1 g_2) = (q g_1) g_2 \quad (1)$$

$$g_1, g_2 \in G : \text{if } q g_1 = q g_2 \text{ for all } q \in Q \text{ then } g_1 = g_2 \quad (2)$$

Hence,  $G$  is a group by Lemma 2.89 it is associative and satifies (1). Also permutations are bijective which satifies (2).  $\square$

## 2.4

$f$  is a monoid morphismus as you can first can combine two words of  $\Sigma^*$  and count the length or count first the length and then add them up. For a given word  $w \in \Sigma^*$  its equivalence class is  $[w]$  with all other words with the same length. Since the words of  $\Sigma^*$  can be arbitrary long, the order of the quotion set is infinte.  $\square$