5

## 5.1

We have  $M = (Q, \Sigma, \delta)$  and  $\pi = \{H_i\}_{i \in I}$  a admissible partion of Q. If M is complete then for all  $i \in I$  and for all  $a \in \Sigma$  there exists **exactly one** partion such that for  $j \in I$   $H_i\delta_a \subseteq H_j$ . By definition of  $\pi$  (Lemma 3.48) there exists **at least one**  $j \in I$  with  $H_i\delta_a \subseteq H_j$ . Notice  $H_i\delta_a \neq \emptyset$  because M is complete.

We show that there only exists **exactly one**  $j \in I$ . Suppose there exists  $j, k \in I$  with  $H_i \delta_a \subseteq H_{j,k}$  for all  $a \in \Sigma$  and  $j \neq k$ . We choose an arbitrary  $q \in H_i$  then the following must hold:

$$q\delta_a = q_j \in H_j$$
$$q\delta_a = q_k \in H_k$$

Notice  $q_j \neq q_k$  because  $H_j \cap H_k = \emptyset$ . This is a contraction because  $q\delta_a$  is not right unique anymore.

## 5.2

We proove for a transformation semigroup transformation (Q, S) which is irreducable such that for all  $q \in Q$  either |qS| = 1 or qS = Q. Assume |qS| = 1 for any  $q \in Q$ . This means we find one arbitray but fixed  $q' \in Q$  such that qS = q'. Moreover,  $q = H_i$  and  $q' = H_j$  with  $i, j \in I$ , it is the trival partition of singleton classes. Now assume qS = Q for any  $q \in Q$ . This means for each  $s \in S$  with qs = q' we map to a different q' such that  $\bigcup q' = Q$ . Thus, we get the trivial partition of Q.

Finally suppose  $|qS| > 1 \land qS \neq Q$  for any  $q \in Q$ . Then qS maps into two or more patitions (up to |Q| - 1) which contradicts the condition of (Q, S) being irreducable.

## 5.3

1. Since, Aut(M) is the set of all state machine automorphisms, this means that f is a bijective function on  $Q \times Q$ , therefore, it only permutates the states of M. As,  $\Sigma$  is mapped to  $\Sigma$  by the identity function, the transactions  $\delta$  of M do not change.

Given Lemma 2.89, we know that given a set Q,  $(S_Q, \circ, id_Q)$  is a group. We have already established that Aut(M) only permutates states of M, and as Aut(M) includes 'all' state machine automorphisms, it also includes the identity function. This concludes the proof, that Aut(M) is a group.

2. Assume, that there exists a  $q \in Q$  such that f(q) = q.

## **5.4**