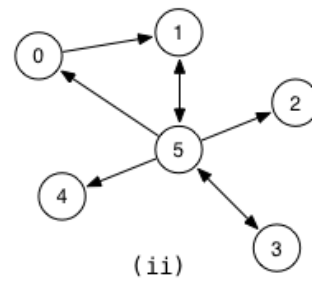
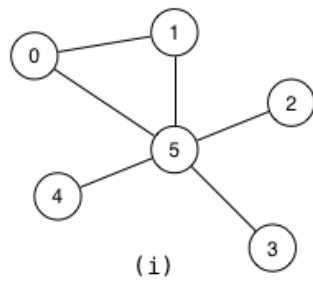


Q1.



First graph:

	0	1	2	3	4	5
0	0	1	0	0	0	1
1	1	0	0	0	0	1
2	0	0	0	0	0	1
3	0	0	0	0	0	1
4	0	0	0	0	0	1
5	1	1	1	1	1	0

Adj. Matrix

0	→ 1, 5
1	→ 0, 5
2	→ 5
3	→ 5
4	→ 5
5	→ 0, 1, 2, 3, 4

Adj. List

Second graph:

	0	1	2	3	4	5
0	0	1	0	0	0	0
1	0	0	0	0	0	1
2	0	0	0	0	0	0
3	0	0	0	0	0	1
4	0	0	0	0	0	0
5	1	1	1	1	1	0

Adj. Matrix

0	→ 1
1	→ 5
2	
3	→ 5
4	
5	→ 0, 1, 2, 3, 4

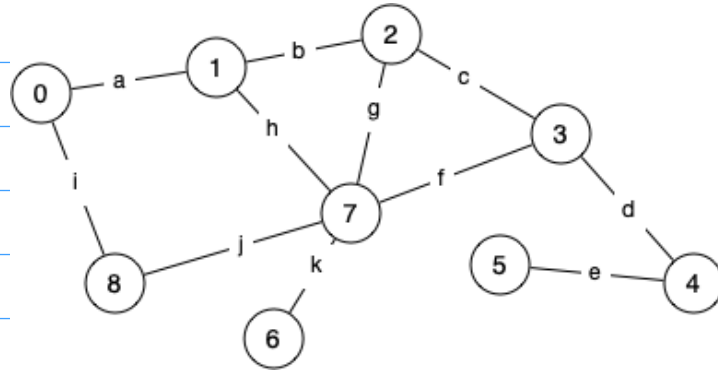
Adj. List

Note the space complexity differences:

Matrix $\Rightarrow O(V^2)$

List $\Rightarrow O(V+E)$

Q2.



Cycle: Path that starts + ends at the same node.
 Clique: Complete subgraphs (with $|V| \geq 3$)

# Edges	# Cycles	# Cliques
11	6	2

v	0	1	2	3	4	5	6	7	8
deg(v)	2	3	3	3	2	1	1	5	2

Longest path from $5 \xrightarrow{*} 8$: $5 \rightarrow 4 \rightarrow 3 \rightarrow 2 \rightarrow 7 \rightarrow 1 \rightarrow 0 \rightarrow 8$

Q4. Differences between adj. matrices of directed / undirected graphs:

Undirected graphs have symmetric adj. matrices

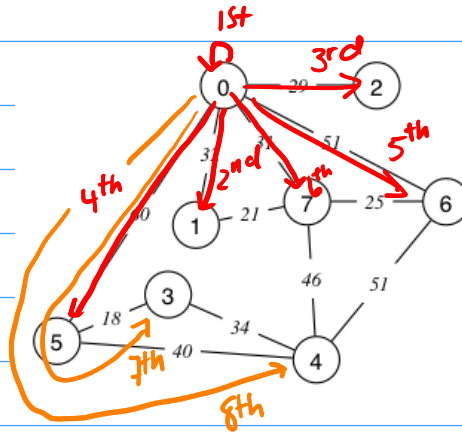
(M is symmetric $\Rightarrow M[i][j] = M[j][i]$ for all i, j)

Why?

$M[i][j] = 1 \quad \Leftrightarrow \quad \text{can go directly from } i \text{ to } j$
 $\Leftrightarrow \quad \text{can go directly from } j \text{ to } i \text{ too}$
 $\Leftrightarrow \quad M[j][i] = 1.$

Similarly for $M[i][j] = 0 \quad \Leftrightarrow \quad M[j][i] = 0$

Q6.

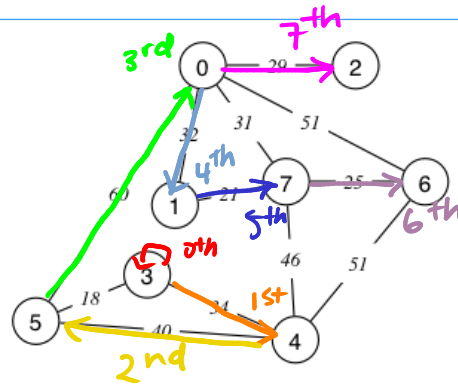


breadthFirst(g, 0);

<u>Iteration</u>	<u>Visited</u>	<u>Queue</u> (head, ..., tail)
0	0 1 2 3 4 5 6 7	0
1	0 1 2 3 4 5 6 7	1, 2, 5, 6, 7
2	0 1 2 3 4 5 6 7	2, 5, 6, 7, 7
3	0 1 2 3 4 5 6 7	5, 6, 7, 7
4	0 1 2 3 4 5 6 7	6, 7, 7, 3, 4
5	0 1 2 3 4 5 6 7	7, 7, 3, 4, 4, 7
6	0 1 2 3 4 5 6 7	7, 3, 4, 4, 7, 4
7	0 1 2 3 4 5 6 7	4, 4, 7, 4, 4
8	0 1 2 3 4 5 6 7	4, 7, 4, 4
9..12	0 1 2 3 4 5 6 7	(eventually becomes empty)

Output: 0, 1, 2, 5, 6, 7, 3, 4

BFS \Rightarrow visit neighbours, then neighbours of neighbours, etc.



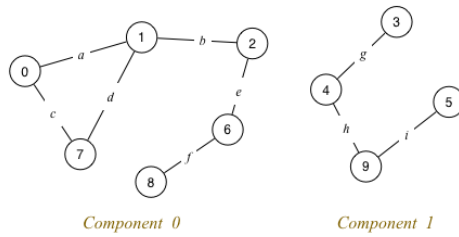
depthFirst(g, 3);

<u>Iteration</u>	<u>Visited</u>	<u>Stack</u> (top, ..., bottom)
0	0 1 2 3 4 5 6 7	3
1	0 1 2 3 4 5 6 7	4, 5
2	0 1 2 3 4 5 6 7	5, 6, 7, 5
3	0 1 2 3 4 5 6 7	0, 6, 7, 5
4	0 1 2 3 4 5 6 7	1, 2, 6, 7, 6, 7, 5
5	0 1 2 3 4 5 6 7	7, 2, 6, 7, 6, 7, 5
6	0 1 2 3 4 5 6 7	6, 2, 6, 7, 6, 7, 5
7	0 1 2 3 4 5 6 7	2, 6, 7, 6, 7, 5
8	0 1 2 3 4 5 6 7	6, 7, 6, 7, 5
9..12	0 1 2 3 4 5 6 7	(eventually becomes empty)

Output: 3, 4, 5, 0, 1, 7, 6, 2

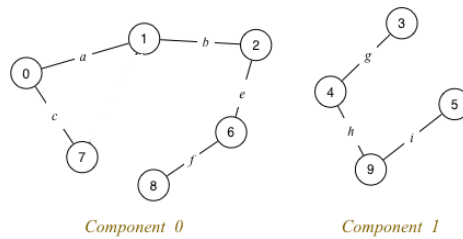
DFS \Rightarrow Follow a path as far as you can, then backtrack to find the next path to explore

Q8.



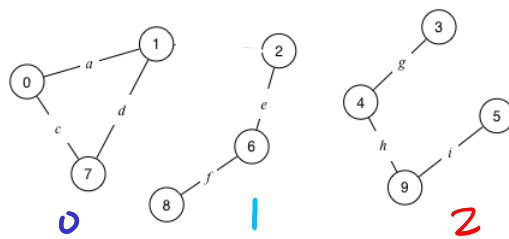
Bridges: e, f, g, h, i

Edge d removed:



connected components unchanged

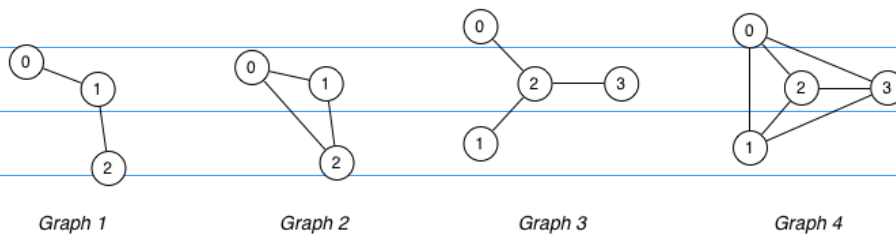
Edge b removed:



$cc[] = \{0, 0, 1, 2, 2, 2, 1, 0, 1, 2\}$

(1 more connected component)

Q9.



Euler path: visit each edge exactly once

Hamiltonian path: visit each vertex exactly once

E for Edge and Euler

Same thing for Euler/Hamilton circuits except you must end up where you started.

Lemmas:

(1) G has an Euler path

\Leftrightarrow There are exactly 2 vertices with odd degree

(2) G has an Euler circuit

\Leftrightarrow Every vertex has even degree

No similar lemmas for Hamiltonian paths/circuits. $\ddot{\smile}$

	Euler path	Euler circuit	Hamiltonian path	Hamiltonian circuit
Graph 1	✓	X	✓	X
Graph 2	✓	✓	✓	✓
Graph 3	X	X	X	X
Graph 4	X	X	✓	✓