

# Data Types

Haskell  $\rightarrow$  MinHS

```
data Direction = East | West | North | South
data Foo = Foo Int Bool Int
data Tree = Node Tree Tree | Leaf
```

Each nullary data constr. is isomorphic to 1

Recursive occurrences?

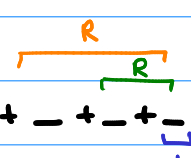
MinHS types isomorphic to these types?

$$\text{Direction} \simeq 1 + (1 + (1 + 1))$$

$$\text{Foo} \simeq 1 \times (\text{Int} \times (\text{Bool} \times \text{Int})) \cong \text{Int} \times (\text{Bool} \times \text{Int})$$

$$\text{Tree} \simeq \text{rec } t. (t \times t + 1)$$

MinHS terms of these types?

$$\text{North} \simeq \text{In}_R (\text{In}_R (\text{In}_L ()))$$


$$\text{Foo } 3161 \text{ False } 3141 \simeq (3161, (\text{False}, 3141))$$

$$\text{Node Leaf Leaf} \simeq \text{roll} (\text{In}_L (\text{roll} (\text{In}_R ()), \text{roll} (\text{In}_R ()))) \quad ??$$

Last term is a bit tricky. Notice that when using  $\text{In}_L$  and  $\text{In}_R$ , one of the types in the sum type  $T_1 + T_2$  is **free/arbitrary**. So

choose this as  $T_1$

$$\begin{aligned} \text{In}_R () &: (\text{rec } t. (t \times t + 1), \text{rec } t. (t \times t + 1)) + 1 \\ &\cong (t \times t + 1) [t := \text{rec } t. (t \times t + 1)] \end{aligned}$$

$$\Rightarrow \text{roll} (\text{In}_R ()): \text{rec } t. (t \times t + 1) \longleftarrow \text{This is how we write Leaf}$$

$$\Rightarrow (\text{roll} (\text{In}_R ()), \text{roll} (\text{In}_R ())) : (\text{rec } t. (t \times t + 1)) \times (\text{rec } t. (t \times t + 1))$$

$$\Rightarrow \text{In}_L ((\text{roll} (\text{In}_R ()), \text{roll} (\text{In}_R ()))) : (\text{rec } t. (t \times t + 1)) \times (\text{rec } t. (t \times t + 1)) + 1$$

$\simeq (t \times t + 1) [t := \text{rec } t. (t \times t + 1)]$  choose this as  $T_2$

$$\Rightarrow \text{roll} (\text{In}_L (...)) : \text{rec } t. (t \times t + 1)$$

↑ This is how we write Node Leaf Leaf.

## Recursive Types

$\text{rec } t. (\text{Int} + t):$

$$\frac{\frac{\text{Int} \vdash \text{Int}}{\text{InL } 3|6| : \text{Int} + (\text{rec } t. (\text{Int} + t))}}{\text{roll}(\text{InL } 3|6|) : \text{rec } t. (\text{Int} + t)}$$

$t_2$  is free in the typing rule for  $\text{InL}$  so I can let it be whatever I want.

How about  $\text{rec } t. (\text{Int} \times t)$ ? Does not have any finite terms!  
But can construct using Recfuns:

$$\text{recfun } f \ x = \text{roll } (x, \underbrace{f \ x}_{: (\text{Int}, \text{rec } t. (\text{Int} \times t))}) \quad \text{where } x : \text{Int}$$

## Isomorphism Equivalence

Which are isomorphic assuming a strict semantics?  $\Rightarrow$  Finite terms only!

Impossible to construct any finite terms

$\nwarrow$   $\text{rec } t. (t \times t)$

$0$

$1$

$1 \times 0$

$\cong 0$

$\text{rec } t. (t \times 1)$

$\text{rec } t. (t + 1)$

$\rightarrow$  Impossible to construct any finite terms  
 $\searrow$  only finite term is  $\text{roll}(\text{InR } ()) \xleftrightarrow{(\text{bij.})} ()$

## Curry-Howard Isomorphism (programs $\cong$ proofs!)

$\nearrow$  Type is  $A + B \rightarrow B + A$

Prove  $A \vee B \Rightarrow B \vee A$  by writing a suitable MinTS term:

$$\begin{aligned} \text{recfun } f \ x = \text{case } x \text{ of } & \text{InL } x_A \rightarrow \text{InR } x_A : B + A \\ & \text{InR } x_B \rightarrow \text{InL } x_B : B + A \end{aligned}$$

This program is a PROOF that  $A \vee B \Rightarrow B \vee A$  !