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Simplified_PageRank Documentation

Data Structures used

- One vector of unordered maps, Four unordered maps
- edges– vector<unordered_map<int, double>>
 - Use – Store all of the edges encountered, specifically tracking the incoming edges to any site along with the weight of the edge. Also used to do the pseudo-matrix multiplication in the page rank calculation.
 - Index – The integer equivalent of the referenced site.
 - Element – unordered_map (key: integer equivalent to the referencing site, value: weight of the edge ie 1/outdegree)
 - Why – The vector allowed for $O(1)$ insertions and lookups, and the unordered_map uses less space than a vector of vectors while still providing $O(1)$ lookups, along with a second spot (the value) for the weights to be.
- edgeOutDegrees – unordered_map<int, int>
 - Use – An unordered_map used to track the outdegree of each site.
 - Key – site's integer representation
 - Value – outdegree of key
 - Why – The map wasn't necessary, a vector could have worked just as well. Unordered to provide $O(1)$ average lookups.
- intToSite– unordered_map<int, string>
 - Use – provides a mapping between integers and site names. Each site is mapped to an integer when they are first encountered, and this is the complement to siteToInt.
 - Key – site's integer representation
 - Value – site's name
 - Why – The necessity of this map was due to the $O(n)$ nature of cpp's hashing function with strings. Had I used strings as the keys to all of the other unordered_maps, I would have been unable to take advantage of the $O(1)$ lookup time due to the hash function being linear to the site name length.
- SiteToInt – map<string, int>
 - Use – provides a mapping between a site name and an integer. Used to map a site to an integer, complemented by intToSite.
 - Key – site name
 - Value – site's integer value.
 - Why – Once again used to avoid the $O(n)$ relationship between the hash function and string length. This function is only used while adding edges (at $O(\text{length of site name})$), and afterwards sites are only referred to by their integer mapping.

Time Complexities: All cases are in the same order if not explicitly mentioned. Average is equal to best in all cases.

AddEdge: $O(\text{length of site name})$ Average, $O(\text{length of site name} * V)$ Worst

- No Loops
- Edges - at: $O(1)$, map operations: $O(1)$ Average, $O(V)$ Worst
- siteToInt - all operations: $O(\text{length of site name})$ Average, $O(\text{length of site name} * V)$ Worst
- intToSite - all operations: $O(1)$ Average, $O(V)$ Worst
- edgeOutDegrees - all operations: $O(1)$ Average, $O(V)$ Worst

calcEdgeValues: $O(V^2)$ Avg, $O(V^3)$ Worst

- Loop through all vertices: $O(V)$
 - Page_ranks.emplace: $O(1)$ Average, $O(V)$ Worst
 - Edges - at - size: $O(1)$
 - Loop through edges connected to first vertex: $O(V)$
 - edgeOutDegrees.at - $O(1)$ Avg, $O(V)$ Worst

pageRanks: $O(P * V^2)$ Avg, $O(P * V^3)$ Worst

- Loop # of iterations times: $O(P)$
 - Loop through all of the sites: $O(V)$
 - Loop through all sites connected to first vertex: $O(V)$
 - PageRanks - at: $O(1)$ Avg, $O(V)$ Worst
 - NewRanks -.emplace: $O(1)$ Avg, $O(V)$ Worst

printEdges: $O(V^2 * \text{length of site name})$ - used for testing

- Loop through all of the sites: $O(V)$
 - Cout: $O(\text{length of site name})$
 - Loop through sites connected to the to site: $O(V)$
 - Cout: $O(\text{length of site name})$

printOutDegrees: $O(V * \text{length of site name})$ Avg, $O(V^2 * \text{length of site name})$ Worst- used for testing

- Loop through all of the sites: $O(V)$
 - intToSite - at: $O(1)$ Avg, $O(V)$ Worst
 - Cout: $O(\text{length of site name})$

printRanks: $O(V * \text{length of site name})$ Average, $O(V^2 * \text{length of site name})$ Worst

- Loop through all sites: $O(V)$
 - Ordered_ranks -.emplace: $O(\text{length of site name})$ Avg, $O(\text{length of site name} * V)$ Worst
- Loop through all sites: $O(V)$
 - Cout - $O(\text{length of site name})$

main:

Avg: $O((E * \text{Length of site name}) + V^2 + (P * V^2) + (V * \text{Length of site name}))$

Worst: $O((E * \text{Length of site name}) + V^3 + (P * V^3) + (V^2 * \text{Length of site name}))$

As E is on the same order as V (graph is given to be sparse) –

Avg: $O((P * V^2) + (V * \text{Length of site name}))$

Worst: $O((P * V^3) + (V^2 * \text{length of site name}))$

- Loop through number of edges: $O(E)$
 - AddEdge: $O(\text{length of site name})$ Average, $O(\text{length of site name} * V)$ Worst
- CalcEdgeValues: $O(V^2)$ Avg, $O(V^3)$ Worst
- pageRanks: $O(P * V^2)$ Avg, $O(P * V^3)$ Worst
- printRanks: $O(V * \text{length of site name})$ Average, $O(V^2 * \text{length of site name})$ Worst

Reflection:

If I were to redo this assignment, I would take a couple more minutes to think through the necessities of the project.

I redid this assignment after realizing that I had made a mistake that significantly increased my time complexity. Specifically, I didn't have the `intToSite` and `siteToInt` maps in my first attempt at the project. After running some stress tests, I realized that my code was too slow, mainly due to the $O(n)$ relationship between site name length and the hash function that C++ uses. After realizing this, I almost completely redid the project, which was unfortunate. So, I should have taken time to realize that the integer representation was necessary to reduce the time complexity.

Additionally, I would have used a vector of vectors to contain all of the edge out degrees, as this would have made it possible to handle parallel edges.