

hw 2b) requires $(A \otimes B)(C \otimes D) = AC \otimes BD$

should prove this identity on the hw

back to optimization...

A algorithms for finding minima of
 $f(x)$ $f: \mathbb{R}^n \rightarrow \mathbb{R}$

Structure of most algorithms goes like this:

Start w/ some x_0

for $k = 0, 1, 2, \dots$ until convergence

- compute $g_k = \nabla f(\vec{x}_k)$

Newton's
method

{ maybe compute
the hessian (Newton's method)
 $H_k = \nabla^2 f(\vec{x}_k)$ & solve
 $H_k d_k = -g_k$ for d_k

deriv. of f wrt
each component
of f . ie, x is
a vector

steepest
descent

OR
{ $d_k = -g_k$

$x_{k+1} = x_k + \alpha_k d_k$, for some α_k T.B.D.

how to determine "convergence"?

also, note that w/ hessian, (newton's method) need to check if $g_k^T d_k < 0 \Rightarrow$ we indeed have a descent direction. (it's guaranteed / by defⁿ for steepest descent).

if we know we're working w/ a descent direction, then for a reasonable choice of alpha we're going to go down.

convergence test: $\|g_k\|$ is "small enough".

small enough might be abs. or relative error.

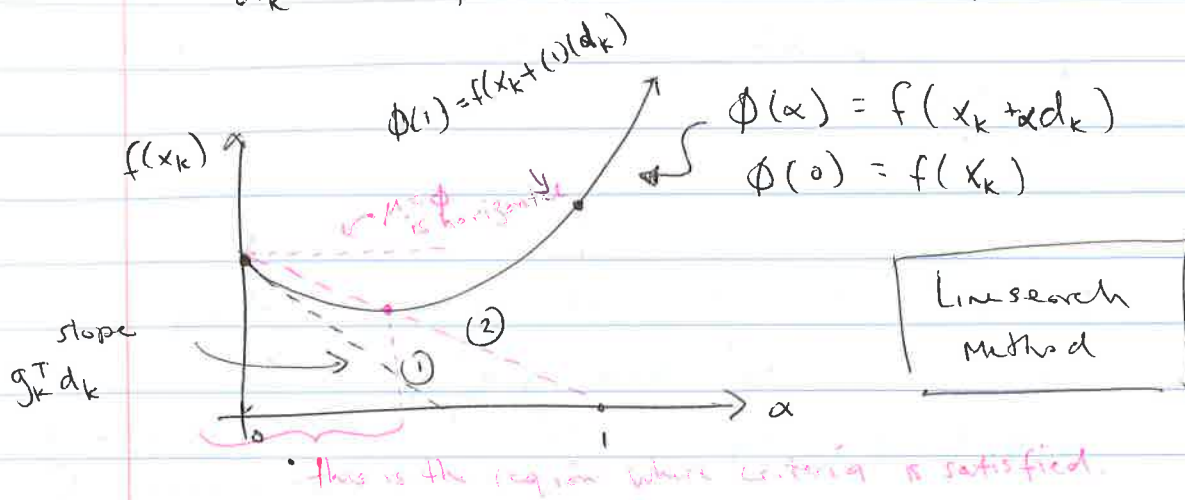
ie. $\frac{\|g_k\|}{\|f(x_k)\|} \leq \tau$, tolerance, which depends on the problem / context.

Cost - assume we have an analytic formula that includes expressions for p_f & p_f^2 .

there are n iterations for steepest descent, n^2 for newton (? i think this is what he said...).

~~scribble~~

So, how to choose d_k ? (assume we have a d_k already - by whatever method).



in this example $\phi(1) > f(x_k)$, so we want to pick another version of α . Strategy?

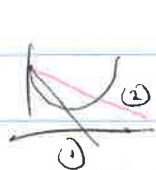
Taylor Series:

$$\phi(\alpha) = f(x_k + \alpha d_k) = f(x_k) + \alpha g_k^T d_k + \underbrace{\theta(d^2)}_{\text{ignore}}$$

Line through
 $(0, f(x_k))$ w/
 slope $g_k^T d_k (< 0)$

since we know $d_k < 0$, since Line (1) is below function f + tangent to f @ 0, then if we create a new line w/ a larger (but still -ive) slope, then it is guaranteed

to be above $f^* = f$ for at least some period of time. (even if you had something like



where f goes below line (1)).

the pink dashed line is a line through

$(0, f(x_k))$ w/ slope $\mu g_k^T d_k$, for

$\mu \in (0, 1)$. eg. $\mu = 1/2$.

$\mu = 0 \Rightarrow$ horizontal
 $\mu = 1 \Rightarrow$ slope is $g_k^T d_k$, same as original line.

Req't on α_k :

ex. 1

$$f(x_k + \alpha_k d_k) \leq f(x_k) + \mu \alpha_k g_k^T d_k$$

value of ~~black~~ f^* line at α_k

value of pink line at α_k

$\mu = 1/2$ convention.
 people research this.

(oo, fun!)

try $\alpha_k = 1$. if req't is satisfied, then accept.

"simple-minded approach"

otherwise, try $\alpha_k \leftarrow \alpha_k / 2$ (reduce α_k)

iterate...

* note: \exists still more book keeping to do, to ensure we have a global min.

for n -dimensions, think of the 2D graph as a cross-section in the direction of \vec{d}_k (vector).

Typical convergence thm

For a descent method applied to a function f that is bounded below, for which gradient g is Lipschitz continuous, then $\|g_k\| \rightarrow 0$

↓ Lipschitz continuous

$$\equiv \|g(x) - g(y)\| \leq L\|x - y\|$$

for some $L > 0$

$\forall x, y$ in the region of interest.

3 better ways to choose α ? viz. simple minded approach discussed earlier - want a diff way that could be smarter ... ie, if we can use our knowledge of f to reduce stupid guesses, we are happier.

in particular, we'll have to compute eqⁿ (1) less often.

Cost of computing ①? well,

$f(x_k)$ - free

$\mu \alpha_k g_k^T d_k$ - cheap

but $f(x_k + \alpha_k d_k)$ (LHS) might be
arbitrarily complex $f \approx f$.

So, an alternative approach to meet the req't
of eq'n ①:

start w/ $\alpha_{k_1} = 1$

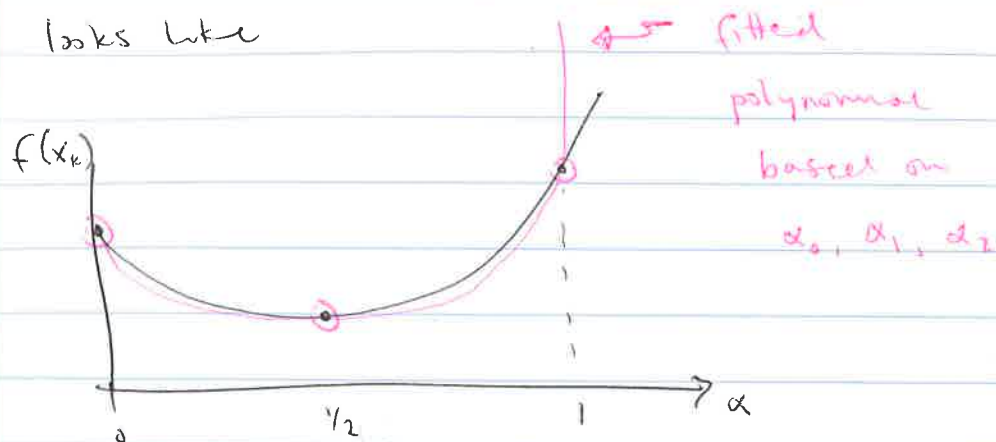
$\alpha_{k_2} =$ something like
 $\frac{1}{2} \alpha_{k_1}$

Now, we have candidate values:

$\alpha_k = 0$	$f(x_k)$	$\alpha_k = 1$	$f(x_k + d_k)$
$\alpha = \frac{1}{2}$	$f(x_k + \frac{1}{2} d_k)$		

Let $p(\alpha) =$ quadratic polynomial that

interpolates 3 points. Then, find α
that minimizes $p(\alpha)$. call this
 α_{k_3} .



the fitted polynomial makes use of values we've already computed.