Universal source coding and the Online Bayes algorithm

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Combining experts in the log loss framework

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The online Bayes Algorithm

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Generalization to larger sets of distributions

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- ► [L_A^T] is the code length if *A* is combined with arithmetic coding.

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- Goal: minimize regret:

$$-\sum_{t=1}^{T} \log p_A^t(c^t) + \min_{i=1,\dots,N} \left(-\sum_{t=1}^{T} \log p_i^t(c^t) \right)$$

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▶ Dividing by T we get $\frac{L_{T}^{T}}{T} = \min_{i} \frac{L_{T}^{T}}{T} + \frac{\log N}{T}$

Upper bound on $\sum_{i=1}^{N} w_i^{T+1}$ for **Hedge** (η)

Lemma (upper bound)

For any sequence of loss vectors ℓ^1, \dots, ℓ^T we have

$$\ln\left(\sum_{i=1}^N w_i^{T+1}\right) \leq -(1-e^{-\eta})L_{\mathsf{Hedge}(\eta)}.$$

Tuning η as a function of T

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per iteration we get:

$$\frac{L_{\mathsf{Hedge}(\eta)}}{T} \leq \min_{i} \frac{L_{i}}{T} + \sqrt{\frac{2 \ln N}{T}} + \frac{\ln N}{T}$$

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- More generally, the regret is smaller if many of the experts perform well.

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 - ► For number of mistakes Bayesian method cannot be "fixed". Requires variable learning rate.

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 - Markov Chain Monte Carlo Sample the posterior. Can sometimes be done efficiently. Efficient sampling relates to mixing rate of markov chain whose limit dist is the posterior dist.

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- Most sequences do not correspond to valid prediction algorithms.

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 - ► A prediction for the next bit $p(\vec{X}) \in [0, 1]$.
 - ► To ensure *p* has a finite description. Restrict to rational numbers *n*/*m*
- Any online prediction algorithm can be represented as code $\vec{b}(E)$ for U. The code length is $|\vec{b}(E)|$.
- Most sequences do not correspond to valid prediction algorithms.
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- Fix a universal Turing machine U.
- ▶ An online prediction algorithm *E* is a program that
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- ▶ Run the Bayes algorithm over "all" prediction algorithms.
- ▶ technical details: On iteration t, $|\vec{X}| = t$. Use the predictions of programs \vec{b} such that $|\vec{b}| \le t$ and for which $V(\vec{b}, \vec{X}, 2^t) = 1$. Assing the remaining mass the prediction 1/2 (insuring a loss of 1)

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- We don't pay a penalty for copies.
- More generally, the regret is smaller if many of the experts perform well.

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- ► Can we still get a meaningful bound?

Bayes Algorithm for biased coins

▶ Replace the initial weight by a density measure

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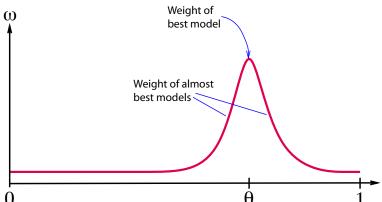
▶ We need a new lower bound on the final total weight

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$$\begin{aligned} L_A - L_{\min} & \leq & \ln \int_0^1 w(\theta) e^{-L_{\theta}} d\theta - \ln e^{L_{\min}} \\ & = & \ln \int_0^1 w(\theta) e^{-(L_{\theta} - L_{\min})} d\theta \\ & = & \ln \int_0^1 w(\theta) e^{T(g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}))} d\theta \end{aligned}$$

► Taylor expansion of $g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta})$ around $\theta = \hat{\theta}$.

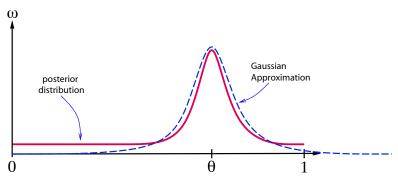
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Laplace Approximation (details)

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$$= w(\hat{\theta}) \sqrt{\frac{-2\pi}{T \frac{d^{2}}{d\theta^{2}} \Big|_{\theta = \hat{\theta}} (g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}))}} + O(T^{-3/2})$$

Choosing the optimal prior

▶ Choose $w(\theta)$ to maximize the worst-case final total weight

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▶ Make bound equal for all $\hat{\theta} \in [0, 1]$ by choosing

$$w^*(\hat{\theta}) = \frac{1}{Z} \sqrt{\frac{\frac{g^2}{d\theta^2}\Big|_{\theta=\hat{\theta}} (g(\hat{\theta},\theta) - g(\hat{\theta},\hat{\theta}))}{-2\pi}},$$

where **Z** is the normalization factor:

$$Z = \sqrt{rac{1}{2\pi}} \int_0^1 \sqrt{rac{d^2}{d heta^2}}igg|_{ heta=\hat{ heta}} (g(\hat{ heta},\hat{ heta}) - g(\hat{ heta}, heta)) \ d\hat{ heta}$$

The bound for the optimal prior

Plugging in we get

$$L_{A} - L_{\min} \leq \ln \int_{0}^{1} w^{*}(\theta) e^{T(g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}))} d\theta$$

$$= \ln \left(\sqrt{\frac{2\pi Z}{T}} + O(T^{-3/2}) \right)$$

$$= \frac{1}{2} \ln \frac{T}{2\pi} - \frac{1}{2} \ln Z + O(1/T) .$$

Solving for log-loss

▶ The exponent in the integral is

$$g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}) = \hat{\theta} \ln \frac{\hat{\theta}}{\theta} + (1 - \hat{\theta}) \ln \frac{1 - \hat{\theta}}{1 - \theta} = D_{KL}(\hat{\theta}||\theta)$$

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▶ The optimal prior:

$$w^*(\hat{\theta}) = \frac{1}{\pi \sqrt{\hat{\theta}(1-\hat{\theta})}}$$

Known in general as Jeffrey's prior. And, in this case, the Dirichlet-(1/2, 1/2) prior.

The cumulative log loss of Bayes using Jeffrey's prior

$$L_A - L_{min} \le \frac{1}{2} \ln(T+1) + \frac{1}{2} \ln \frac{\pi}{2} + O(1/T)$$

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► This is called the Trichevsky Trofimov prediction rule.

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▶ Suffers larger regret when $\hat{\theta}$ is far from 1/2

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The constant C is optimal.

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Generalization to larger sets of distributions

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► [Haussler and Opper] show that the coefficient in front of In *T* is optimal for distribution families where the metric entropy is up to

$$N(1/\epsilon) = O(e^{\epsilon^{-\alpha}})$$

For all $\alpha \leq 5/2$.