Combining infinite sets of experts

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Outline

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The online Bayes Algorithm

► Total loss of expert i

$$L_i^t = -\sum_{i=1}^t \log p_i^s(c^s); \quad L_i^0 = 0$$

Weight of expert i

$$w_i^t = w_i^1 e^{-L_i^{t-1}} = w_i^1 \prod_{i=1}^{t-1} p_i^s(c^s)$$

Freedom to choose initial weights.

$$w_t^1 > 0$$
, $\sum_{i=1}^n w_i^1 = 1$

► Prediction of algorithm A

$$\mathbf{p}_A^t = \frac{\sum_{i=1}^N w_i^t \mathbf{p}_i^t}{\sum_{i=1}^N w_i^t}$$

Cumulative loss vs. Final total weight

Total weight:
$$W^t \doteq \sum_{i=1}^N w_i^t$$

$$\frac{W^{t+1}}{W^t} = \frac{\sum_{i=1}^{N} w_i^t e^{\log p_i^t(c^t)}}{\sum_{i=1}^{N} w_i^t} = \frac{\sum_{i=1}^{N} w_i^t p_i^t(c^t)}{\sum_{i=1}^{N} w_i^t} = p_A^t(c^t)$$
$$-\log \frac{W^{t+1}}{W^t} = -\log p_A^t(c^t)$$

$$-\log W^{T+1} = -\log \frac{W^{T+1}}{W^1} = -\sum_{t=1}^{T} \log p_A^t(c^t) = L_A^T$$

EQUALITY not bound!

Simple Bound

- ▶ Use non-uniform initial weights $\sum_i w_i^1 = 1$
- Total Weight is at least the weight of the best expert.

$$\begin{split} L_A^T &= -\log W^{T+1} = -\log \sum_{i=1}^N w_i^{T+1} \\ &= -\log \sum_{i=1}^N w_i^1 e^{-L_i^T} \le -\log \max_i \left(w_i^1 e^{-L_i^T} \right) \\ &= \min_i \left(L_i^T - \log w_i^1 \right) \end{split}$$

Standardizing online prediction algorithms

- Fix a universal Turing machine U.
- ▶ An online prediction algorithm *E* is a program that
 - given as input The past $\vec{X} \in \{0, 1\}^t$
 - runs finite time and outputs
 - ▶ A prediction for the next bit $p(\vec{X}) \in [0, 1]$.
 - To ensure p has a finite description. Restrict to rational numbers n/m
- Any online prediction algorithm can be represented as code $\vec{b}(E)$ for U. The code length is $|\vec{b}(E)|$.
- Most sequences do not correspond to valid prediction algorithms.
- ▶ $V(\vec{b}, \vec{X}, t) = 1$ if the program \vec{b} , given \vec{X} as input, halts within t steps and outputs a well-formed prediction. Otherwise $V(\vec{b}, \vec{X}, t) = 0$
- $V(\vec{b}, \vec{X}, t)$ is computable (recursively enumerable).

A universal prediction machine

- Assign to the code \vec{b} the initial weight $w_{\vec{b}}^1 = 2^{-|\vec{b}| \log_2 |\vec{b}|}$.
- ► The total initial weight over all finite binary sequences is one.
- ► Run the Bayes algorithm over "all" prediction algorithms.
- ▶ technical details: On iteration t, $|\vec{X}| = t$. Use the predictions of programs \vec{b} such that $|\vec{b}| \le t$ and for which $V(\vec{b}, \vec{X}, 2^t) = 1$. Assing the remaining mass the prediction 1/2 (insuring a loss of 1)

Performance of the universal prediction algorithm

- ▶ Using $L_A \leq \min_i (L_i \log w_i^1)$
- Assume E is a prediction algorithm which generates the tth prediction in time smaller than 2^t
- ▶ When $t \le |\vec{b}(E)|$ the algorithm is not used and thus it's loss is 1
- ▶ We get that the loss of the Universal algorithm is at most $2|\vec{b}(E)| + \log_2 |\vec{b}(E)| + L_E$
- More careful analysis can reduce $2|\vec{b}(E)| + \log_2|\vec{b}(E)|$ to $|\vec{b}(E)|$
- Code length is arbitrarily close to the Kolmogorov Complexity of the sequence.
- Ridiculously bad running time.

Bayes coding is better than two part codes

- Simple bound as good as bound for two part codes (MDL) but enables online compression
- ► Suppose we have K copies of each expert.
- ► Two part code has to point to one of the KN experts $L_A \le \log NK + \min_i L_i^T = \log NK + \min_i L_i^T$
- ▶ If we use Bayes predictor + arithmetic coding we get:

$$L_A = -\log W^{T+1} \le \log K \max_i \frac{1}{NK} e^{-L_i^T} = \log N + \min_i L_i^T$$

- We don't pay a penalty for copies.
- More generally, the regret is smaller if many of the experts perform well.

The biased coins set of experts

- ► Each expert corresponds to a biased coin, predicts with a fixed $\theta \in [0, 1]$.
- Set of experts is uncountably infinite.
- Only countably many experts can be assigned non-zero weight.
- Instead, we assign the experts a Density Measure.
- ▶ $L_A \le \min_i (L_i \log w_i^1)$ is meaningless.
- Can we still get a meaningful bound?

Bayes Algorithm for biased coins

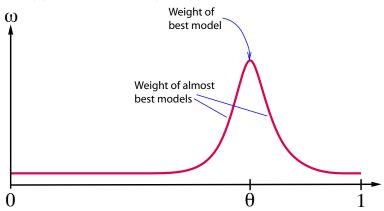
- ► Replace the initial weight by a density measure $w(\theta) = w^{1}(\theta), \int_{0}^{1} w(\theta) d\theta = 1$
- Relationship between final total weight and total log loss remains unchanged:

$$L_A = \ln \int_0^1 w(\theta) e^{-L_{\theta}^{T+1}} d\theta$$

We need a new lower bound on the final total weight

Main Idea

If $\mathbf{w}^t(\theta)$ is large then $\mathbf{w}^t(\theta + \epsilon)$ is also large.



Expanding the exponent around the peak

 \blacktriangleright For log loss the best θ is empirical distribution of the seq.

$$\hat{\theta} = \frac{\#\{x^t = 1; \ 1 \le t \le T\}}{T}$$

The total loss scales with T

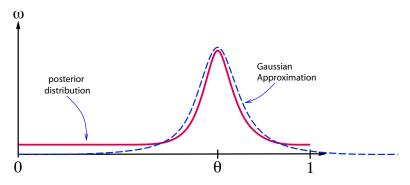
$$L_{\theta} = T \cdot (\hat{\theta}\ell(\theta, 1) + (1 - \hat{\theta})\ell(\theta, 0)) \doteq T \cdot g(\hat{\theta}, \theta)$$

$$\begin{array}{lcl} \textit{L}_{\textit{A}} - \textit{L}_{\text{min}} & \leq & \ln \int_{0}^{1} w(\theta) e^{-\textit{L}_{\theta}} d\theta - \ln e^{\textit{L}_{\text{min}}} \\ \\ & = & \ln \int_{0}^{1} w(\theta) e^{-(\textit{L}_{\theta} - \textit{L}_{\text{min}})} d\theta \\ \\ \textit{pause} & = & \ln \int_{0}^{1} w(\theta) e^{\textit{T}(g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}))} d\theta \end{array}$$

Laplace Approximation

Laplace approximation (idea)

- ► Taylor expansion of $g(\hat{\theta}, \theta) g(\hat{\theta}, \hat{\theta})$ around $\theta = \hat{\theta}$.
- First and second terms in the expansion are zero.
- ► Third term gives a quadratic expression in the exponent
- ightharpoonup \Rightarrow a gaussian approximation of the posterior.



Bayes using Jeffrey's prior
 Laplace Approximation

Laplace Approximation (details)

$$\int_{0}^{1} w(\theta) e^{T(g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}))} d\theta$$

$$= w(\hat{\theta}) \sqrt{\frac{-2\pi}{T \frac{d^{2}}{d\theta^{2}} \Big|_{\theta = \hat{\theta}} (g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}))}} + O(T^{-3/2})$$

Choosing the optimal prior

Choosing the optimal prior

▶ Choose $w(\theta)$ to maximize the worst-case final total weight

$$\min_{\hat{\theta}} w(\hat{\theta}) \sqrt{\frac{-2\pi}{T \left. \frac{d^2}{d\theta^2} \right|_{\theta = \hat{\theta}} (g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}))}}$$

▶ Make bound equal for all $\hat{\theta} \in [0, 1]$ by choosing

$$w^*(\hat{\theta}) = \frac{1}{Z} \sqrt{\frac{\frac{g^2}{d\theta^2}\Big|_{\theta=\hat{\theta}} (g(\hat{\theta},\theta) - g(\hat{\theta},\hat{\theta}))}{-2\pi}},$$

where **Z** is the normalization factor:

$$Z=\sqrt{rac{1}{2\pi}} \int_0^1 \left. \sqrt{rac{d^2}{d heta^2}}
ight|_{ heta=\hat{ heta}} (g(\hat{ heta},\hat{ heta})-g(\hat{ heta}, heta)) \left. d\hat{ heta}
ight|$$

Choosing the optimal prior

The bound for the optimal prior

Plugging in we get

$$L_{A} - L_{\min} \leq \ln \int_{0}^{1} w^{*}(\theta) e^{T(g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}))} d\theta$$

$$= \ln \left(\sqrt{\frac{2\pi Z}{T}} + O(T^{-3/2}) \right)$$

$$= \frac{1}{2} \ln \frac{T}{2\pi} - \frac{1}{2} \ln Z + O(1/T) .$$

Solving for log-loss

► The exponent in the integral is

$$g(\hat{ heta}, heta) - g(\hat{ heta}, \hat{ heta}) = \hat{ heta} \ln rac{\hat{ heta}}{ heta} + (1 - \hat{ heta}) \ln rac{1 - \hat{ heta}}{1 - heta} = D_{ extit{KL}}(\hat{ heta}|| heta)$$

The second derivative

$$\left. \frac{d^2}{d\theta^2} \right|_{\theta=\hat{\theta}} D_{KL}(\hat{\theta}||\theta) = \hat{\theta}(1-\hat{\theta})$$

Is called the empirical Fisher information

The optimal prior:

$$w^*(\hat{\theta}) = \frac{1}{\pi \sqrt{\hat{\theta}(1-\hat{\theta})}}$$

Known in general as Jeffrey's prior. And, in this case, the Dirichlet-(1/2, 1/2) prior.

Bayes using Jeffrey's prior
Choosing the optimal prior

The cumulative log loss of Bayes using Jeffrey's prior

$$L_A - L_{\min} \le \frac{1}{2} \ln(T+1) + \frac{1}{2} \ln \frac{\pi}{2} + O(1/T)$$

But what is the prediction rule?

- ► As luck would have it the Dirichlet prior is the conjugate prior for the Binomial distribution.
- ▶ Observed t bits, n of which were 1. The posterior is:

$$\frac{1}{Z\sqrt{\theta(1-\theta)}}\theta^{n}(1-\theta)^{t-n} = \frac{1}{Z}\theta^{n-1/2}(1-\theta)^{t-n-1/2}$$

The posterior average is:

$$\frac{\int_0^1 \theta^{n+1/2} (1-\theta)^{t-n-1/2} d\theta}{\int_0^1 \theta^{n-1/2} (1-\theta)^{t-n-1/2} d\theta} = \frac{n+1/2}{t+1}$$

▶ This is called the Trichevsky Trofimov prediction rule.

Laplace Rule of Succession

- Laplace suggested using the uniform prior, which is also a conjugate prior.
- In this case the posterior average is:

$$\frac{\int_{0}^{1} \theta^{n+1} (1-\theta)^{t-n} d\theta}{\int_{0}^{1} \theta^{n} (1-\theta)^{t-n} d\theta} = \frac{n+1}{t+2}$$

► The bound on the cumulative log loss is worse:

$$L_A - L_{\min} = \ln T + O(1)$$

▶ Suffers larger regret when $\hat{\theta}$ is far from 1/2

Shtarkov Lower bound

What is the optimal prediction when T is know in advance?

$$L_*^T - \min_{\theta} L_{\theta}^T \geq \frac{1}{2} \ln(T+1) + \frac{1}{2} \ln \frac{\pi}{2} - O(\frac{1}{\sqrt{T}})$$

Multinomial Distributions

- ► For a distribution over k elements (Multinomial) [Xie and Barron]
- ▶ Use the add 1/2 rule (KT).

$$p(i) = \frac{n_i + 1/2}{t + k/2}$$

Bound is

$$L_A - L_{\min} \leq \frac{k-1}{2} \ln T + C + o(1)$$

The constant C is optimal.

Exponential Distributions

- ► For any set of distributions from the exponential family defined by *k* parameters (Some technical conditions on closure of set??) [Rissanen??]
- Use Bayes Algorithm with Jeffrey's prior:

$$w^*(\hat{\theta}) = \frac{1}{Z} \frac{1}{\sqrt{\left|\mathbf{H}\left(D_{KL}(\hat{\theta}||\theta)\right)\right|_{\theta=\hat{\theta}}}}$$

H denotes the Hessian.

$$L_A - L_{\min} \le \frac{k-1}{2} \ln T - \ln Z + o(1)$$

General Distributions

- Characterize distribution family by metric entropy.
- Fixed parameter set usually corresponds to polynomial metrix entropy

$$N(1/\epsilon) = O\left(\frac{1}{\epsilon^d}\right)$$

d is the number of parameters.

► [Haussler and Opper] show that the coefficient in front of In *T* is optimal for distribution families where the metric entropy is up to

$$N(1/\epsilon) = O(e^{\epsilon^{-\alpha}})$$

For all $\alpha \leq 5/2$.

next Class

- Variable-length markov models a set of distributions with increasing number of parameters.
- THe context algorithm: An efficient implementation of the Bayes algorithm which achieves close-to-optimal worst case bounds.