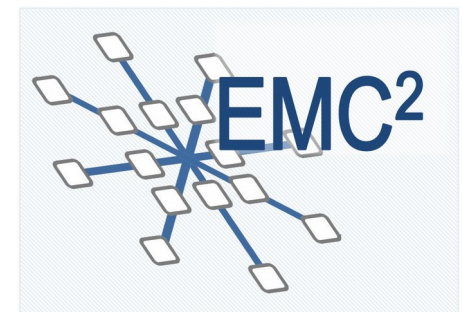


# Safety-critical control in mixed criticality embedded systems

Emil Hjelm



# Presentation outline

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- Background
- Research question and problem statement
- What has been done in the field?
- Demonstrator system overview
- Results
- Conclusion

# Background



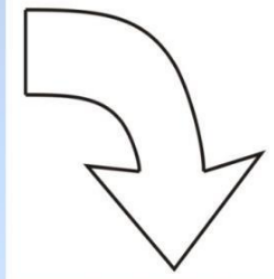
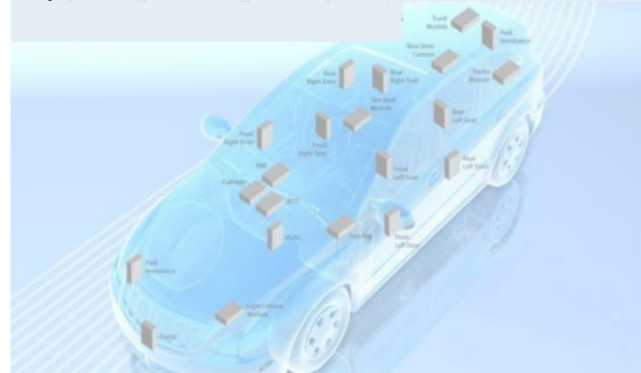
Today, modern cars have around 100 ECUs, this means:

- Inefficient resource usage
- Clear separation between functions

Moving towards MCS would mean:

- Efficient resource usage
  - Less power consumed
  - Less weight
  - Less volume
  - Lower production cost
- Difficult to separate functions of different criticality
- Fewer but more complex systems

Many heterogeneous single-core systems, specialized for the individual criticality levels



Multi-core systems for mixed criticalities



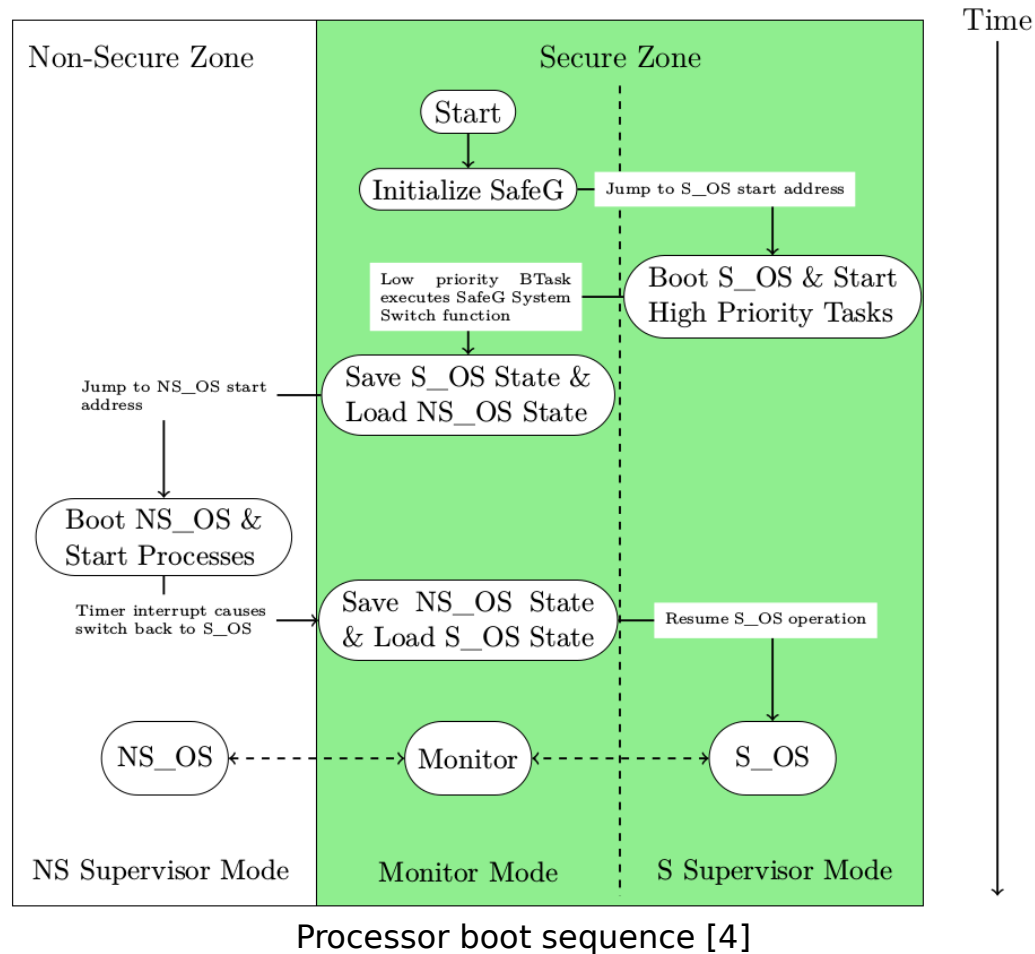
Illustration of combining applications to reduce ECUs [1]

## Four different Automotive Safety Integrity Levels, ASILs

Severity	Probability	Controllability		
		C1	C2	C3
S1	E1	QM	QM	QM
	E2	QM	QM	QM
	E3	QM	QM	A
	E4	QM	A	B
S2	E1	QM	QM	QM
	E2	QM	QM	A
	E3	QM	A	B
	E4	A	B	C
S3	E1	QM	QM	A
	E2	QM	A	B
	E3	A	B	C
	E4	B	C	D

[3] Road vehicles – functional safety – part 9: Automotive safety integrity level (asil)-oriented and safety-oriented analyses. Standard, International Organization for Standardization, Geneva, CH, November 2011.

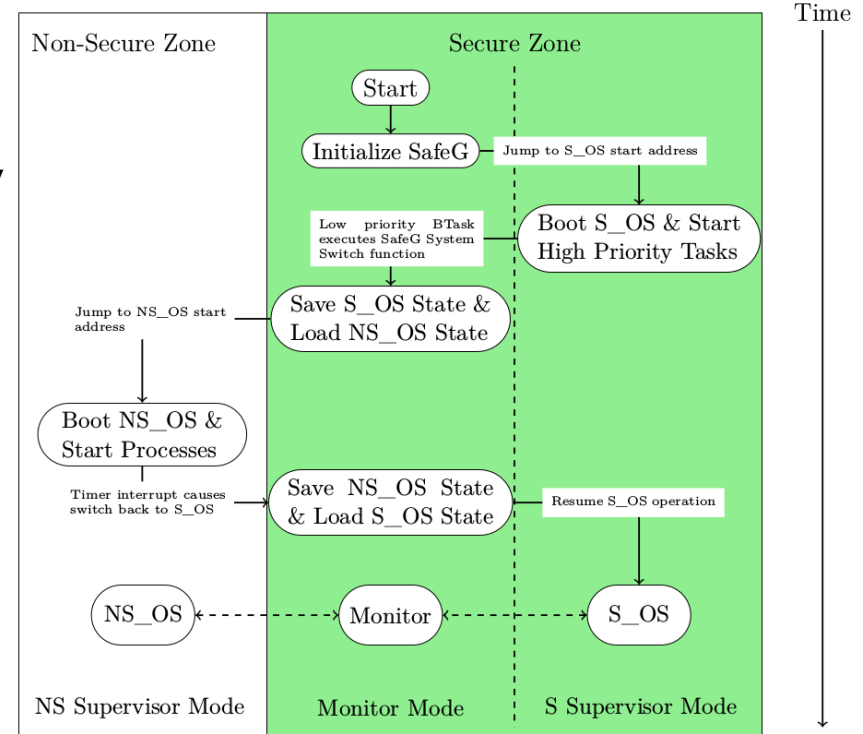
# Alten MCS - Processor switching state



# Alten MCS - Processor switching state



- Mode-switch takes  $\sim 2 \mu\text{s}$
- Theoretical maximum frequency for any control algorithm  $f_{\text{max}} < 250 \text{ kHz}$



Processor boot sequence [4]



Is virtualization an efficient approach when trying to reconcile the conflicting requirements of partitioning for safety assurance and sharing for efficient resource usage when implementing a safety-critical control system?

## Problem statement

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Implement distance keeping control algorithm as a safety critical application on a mixed criticality system. Evaluate performance of processor

- Missed deadlines?
- Robustness of hypervisor
- CPU utilization



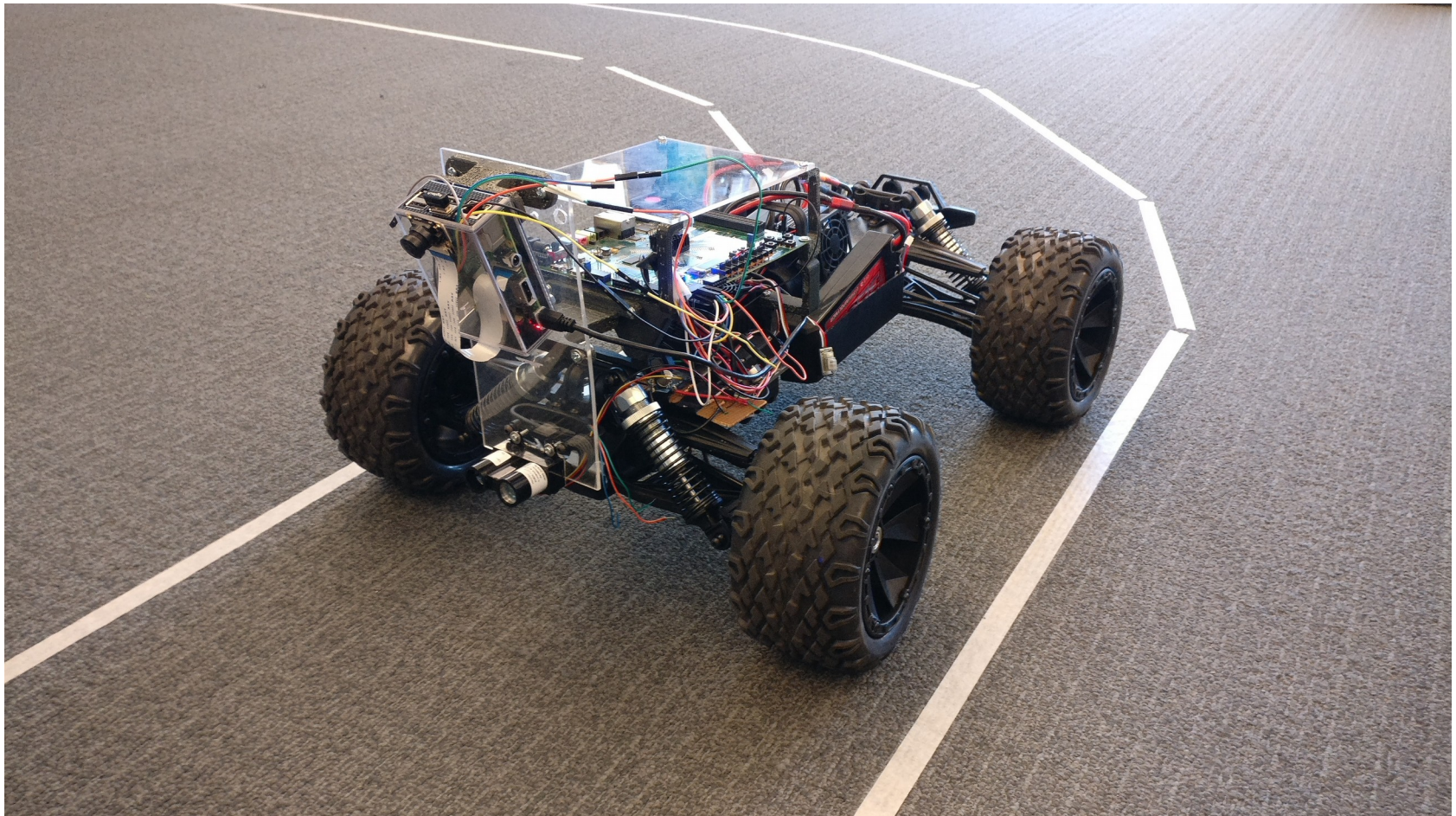
- Steve Vestal showed that neither RM nor DM was optimal when criticality levels are introduced [6]
- Sanjoy Baruah and Steve Vestal showed that EDF does not dominate FP in mixed criticality systems [7]

[6] Steve Vestal. Preemptive scheduling of multi-criticality systems with varying degrees of execution time assurance, 2007.

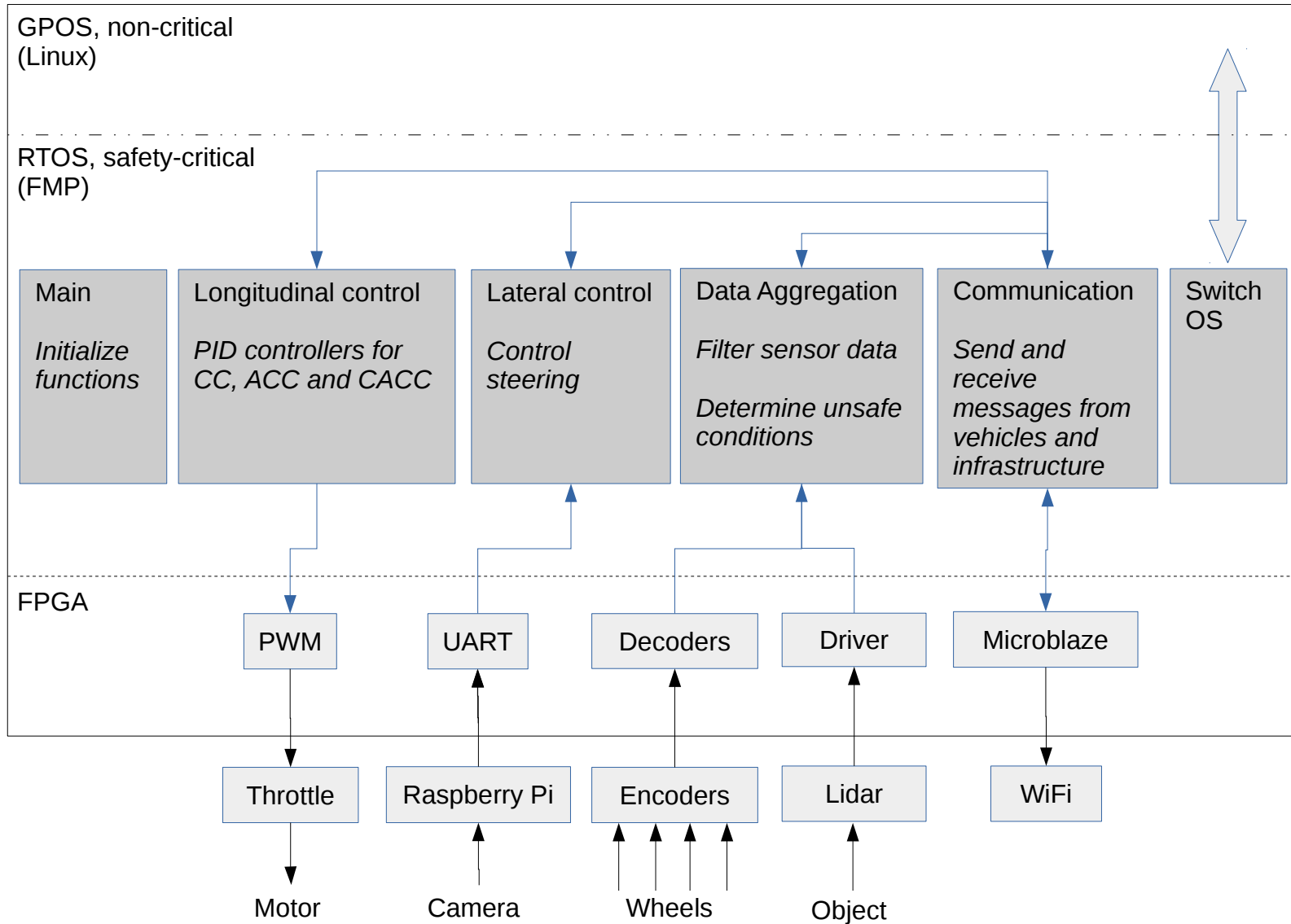
[7] Sanjoy Baruah and Steve Vestal. Schedulability analysis of sporadic tasks with multiple criticality specifications, 2008.

# System overview

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# System overview



# Processor scheduling



Data aggregation

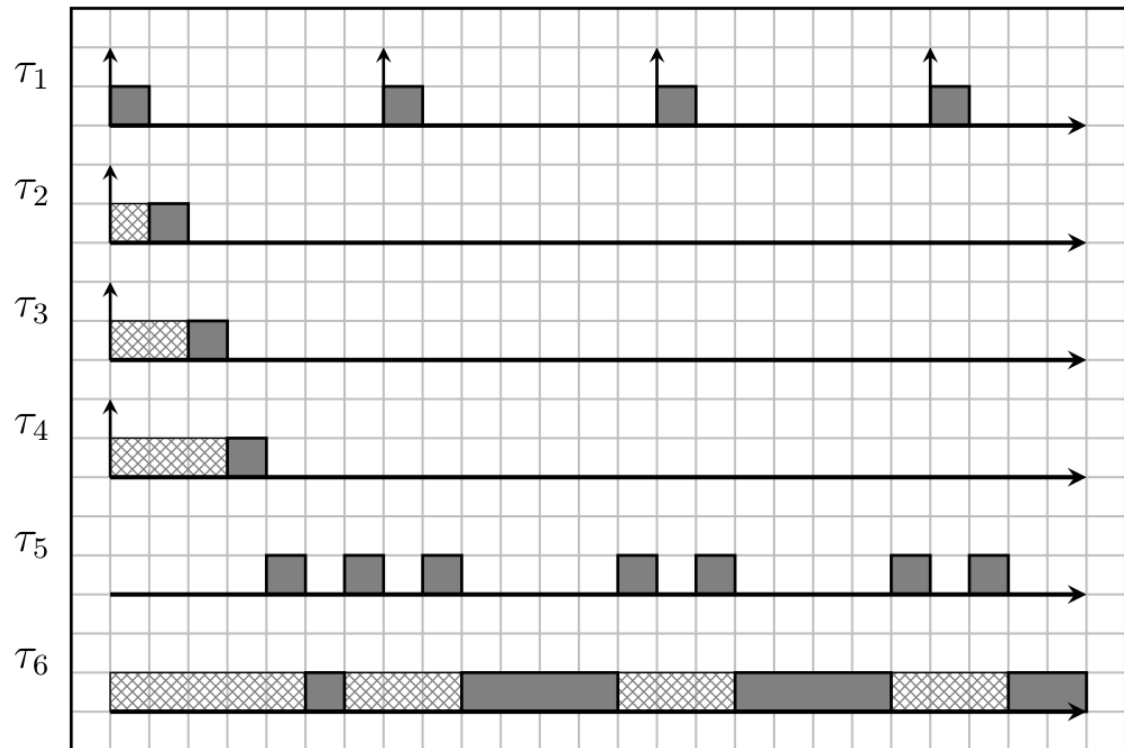
Communication

Lateral control

Longitudinal control

OS switch

GPOS

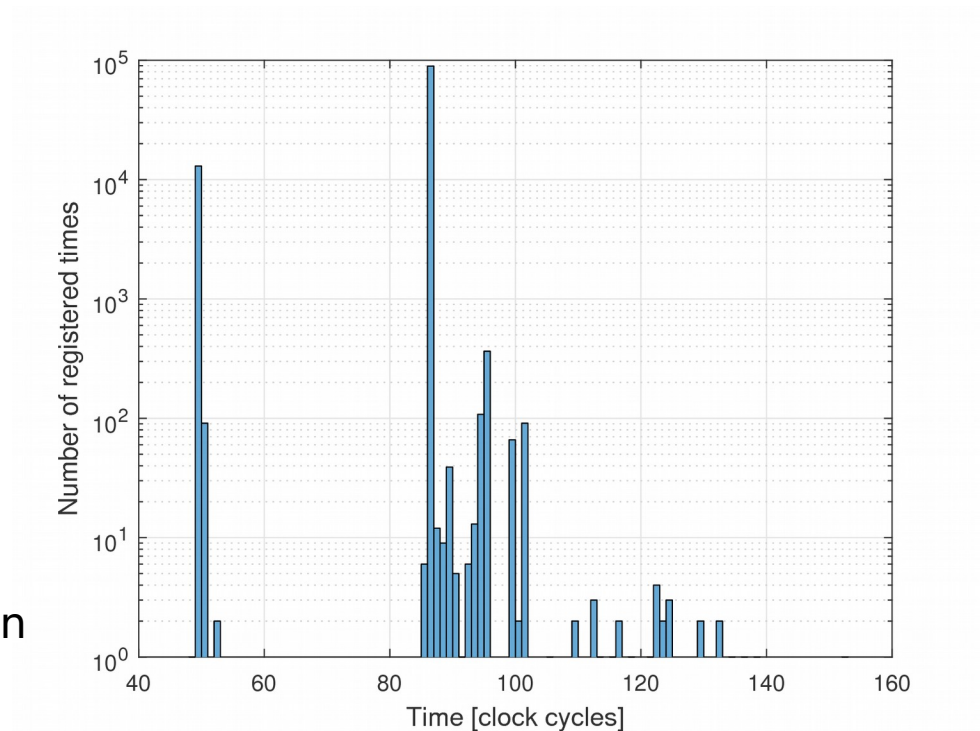


Time not to scale

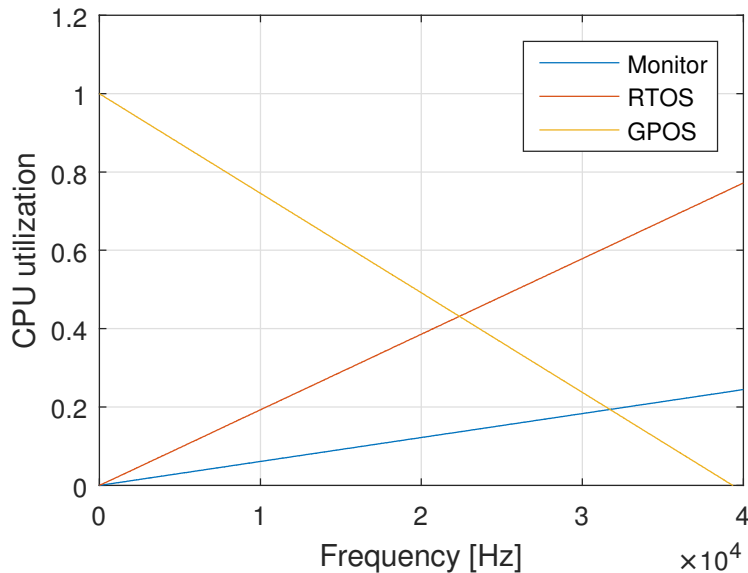
# OS switch execution times



- WCET:  
153 clock cycles  
3.06  $\mu$ s
- Median execution time:  
86 clock cycles  
1.72
- 1 kHz switching frequency:
  - RTOS CPU utilization 0.005%
  - Hypervisor overhead 0.6%
  - Increased total CPU utilization to potentially 99.4%



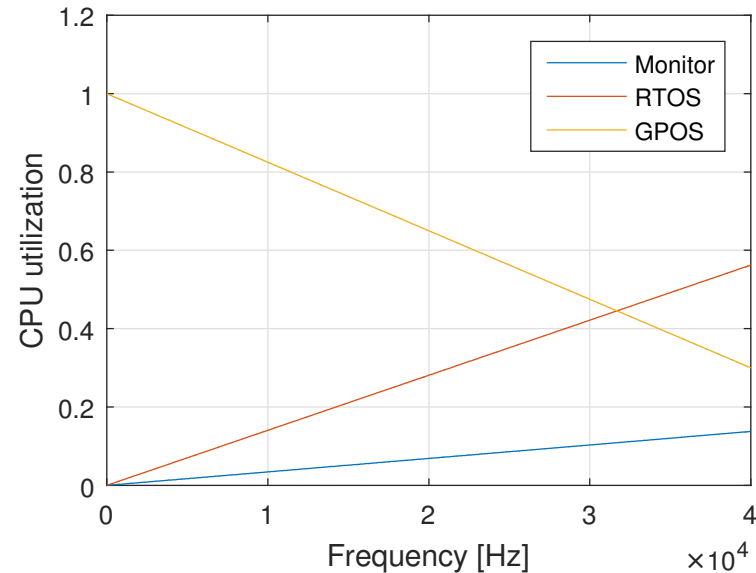
## WCET



100% CPU utilization by RTOS  
and monitor at 40kHz

(Meaning 0% time for GPOS)

## Median execution time



34% CPU utilization by GPOS  
at 40kHz if RTOS and monitor  
runs for median execution  
time



Is virtualization an efficient approach when trying to reconcile the conflicting requirements of partitioning for safety assurance and sharing for efficient resource usage when implementing a safety-critical control system?

Yes!

# Questions

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