A Framework for Automatic OpenMP Code Generation

Raghesh A (CS09M032)

Guide: Dr. Shankar Balachandran

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Outline

- Introduction
- The Polyhedral Model
- LLVM
- Polly
- OpenMP Code Generation in Polly
- Testing with PolyBench
- Conclusion and Future Work
- Setting up the environment
- Various Tools Used in Polyhedral Community

```
Source code
float A[1024];
int main()
  int i, j;
  for (i = 0; i < 1024; i++)
    for (j = 0; j < 5000000; j++)
      A[i] += i
```

LLVM-IR Sequential

```
define i32 @main() nounwind {
entry:
 %retval = alloca i32, align 4
 %i = alloca i32. align 4
 %j = alloca i32, align 4
  store i32 0, i32 * %retval
  store i32 0. i32 * %i. align 4
  br label %for.cond
for . cond:
 %tmp = load i32 * %i, align 4
 %cmp = icmp slt i32 %tmp, 1024
  br il %cmp, label %for.body,
              label %for.end12
for.body:
  store i32 0. i32 * %i. align 4
  br label %for.cond1
for . cond1:
 %tmp2 = load i32 * %j, align 4
 %cmp3 = icmp slt i32 %tmp2, 5000000
  br i1 %cmp3, label %for.body4,
               label %for end
```

LLVM-IR Sequential

```
for.bodv4:
 %tmp5 = load i32 * %i, align 4
 %conv = sitofp i32 %tmp5 to float
 \%tmp6 = load i32 * %i. align 4
 %arrayidx = getelementptr inbounds
              [1024 x float]*
              @A. i32 0. i32 %tmp6
 %tmp7 = load float * %arravidx
 %add = fadd float %tmp7, %conv
  store float %add, float* %arrayidx
  br label %for.inc
for .inc:
 %tmp8 = load i32* %j, align 4
 \%inc = add nsw i32 \%tmp8. 1
  store i32 %inc. i32 * %i. align 4
  br label %for.cond1
for end.
  br label %for inc9
for inc9:
 %tmp10 = load i32* \%i, align 4
 \%inc11 = add nsw i32 \%tmp10. 1
  store i32 %inc11, i32* %i, align 4
  br label %for cond
for end12.
 %0 = load i32 * %retval
  ret i32 %0
```

Source code with OpenMP pragmas

```
float A[1024];
int main()
  int i, j;
 #pragma omp parallel for \
  schedule(runtime) private(j)
  for (i = 0; i < 1024; i++)
    for (j = 0; j < 5000000; j++)
     A[i] += i:
```

LLVM-IR Manual

```
define i32 @main() nounwind {
entry:
 %retval = alloca i32
 \%i = alloca i32
 %i = alloca i32
 "alloca point" = bitcast i32 0 to i32
  call void
       @GOMP_parallel_loop_runtime_start (
       void (i8*)* @main.omp_fn.0,
       i8* null, i32 0, i32 0,
       i32 1024, i32 1) nounwind
  call void
       @main.omp_fn.0(i8* null) nounwind
  call void
       @GOMP_parallel_end() nounwind
  br label %return
return ·
: preds = %entry
 %retval1 = load i32* %retval
; < i32 > [\#uses=1]
  ret i32 %retval1
```

LLVM-IR Manual

```
define internal void @main.omp_fn.0(
        i8 * %.omp_data_i) nounwind {
entry:
  <some initializations here>
  store i8 * %.omp_data_i,
        i8 ** %.omp_data_i_addr
  br label %bb
bb:
  %1 = call zeroext i8
       @GOMP_loop_runtime_next(
        i32 * %.istart0.3,
        i32 * %.iend0.4) nounwind
  %toBool = icmp ne i8 %1.0
  br i1 %toBool, label %bb2, label %bb1
hh1 ·
  call void @GOMP_loop_end_nowait()
  nounwind
  br label %return
hh2·
 <body of the loop>
```

LLVM-IR Automatic

```
@A = common global
     [1024 x float]
     zeroinitializer, align 4
define i32 @main() nounwind {
<some initialization >
pollyBB:
  %insertInst = zext i1 true to i16
  %omp.userContext = alloca
     %main.omp_subfn.omp.userContext
  %0 = getelementptr inbounds
       %main.omp_subfn.omp.userContext*
       %omp.userContext, i32 0, i32 0
  store [1024 x float]*
        @A, [1024 x float]** %0
  call void
      @GOMP_parallel_loop_runtime_start(
      void (i8*)* @main.omp_subfn.
      i8 * %omp_data, i32 0, i32 0,
      i32 1024, i32 1)
  call void @main.omp_subfn(i8*%omp_data) }
  call void @GOMP_parallel_end()
  br label %for.end12.region
```

LLVM-IR Automatic

```
define internal void
       @main.omp_subfn(
       i8 * %omp.userContext) {
omp.setup:
  <some initialization >
omp.exit:
  call void @GOMP_loop_end_nowait()
  ret void
omp.checkNext:
  \%2 = call i8
       @GOMP_loop_runtime_next(
       i32 * %omp. lowerBoundPtr.
       i32 * %omp.upperBoundPtr)
omp.loadIVBounds:
<body of the loop>
```

Necessary Background

- Parallelism in programs
 - · Parallelism and locality
 - Realizing parallelism
- Auto parallelization
- The polyhedral model
- LLVM
- Polly

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Workdone: "OpenMP Code Generation in Polly"

- Examples for transformations with polyhedral model
 - Transformation for improving data locality

- Examples for transformations with polyhedral model
 - Transformation for improving data locality

```
for(i = 1; i <= 10; i++)
A[i] = 10;
for(j = 6; j <= 15; j++)
A[j] = 15;
```

- Examples for transformations with polyhedral model
 - Transformation for improving data locality

```
\begin{array}{lll} \text{for} \, (\, i \, = \, 1; \, \, i \, <= \, 10; \, \, i \, ++) \\ A [\, i \, ] \, = \, 10; \\ \text{for} \, (\, j \, = \, 6; \, \, j \, <= \, 15; \, \, j \, ++) \\ A [\, j \, ] \, = \, 15; \end{array}
```

```
\begin{array}{lll} & \text{for} \left( \begin{array}{l} i & = 1; & i <= 5; & i++ \right) \\ & A \left[ \begin{array}{l} i \end{array} \right] & = 10; \\ & \text{for} \left( \begin{array}{l} j & = 6; & j <= 15; & j++ \right) \\ & A \left[ \begin{array}{l} j \end{array} \right] & = 15; \end{array}
```

- Examples for transformations with polyhedral model
 - Transformation for improving data locality

```
\begin{array}{lll} \text{for} \, (\, i \, = \, 1; \, \, i \, <= \, 10; \, \, i \, ++) \\ A \, [\, i \, ] \, = \, 10; \\ \text{for} \, (\, j \, = \, 6; \, \, j \, <= \, 15; \, \, j \, ++) \\ A \, [\, j \, ] \, = \, 15; \end{array}
```

```
\begin{array}{lll} \text{for}\,(\,i\,=\,1;\,\,i\,<=\,5;\,\,i++) \\ & \text{A[i]}\,=\,10;\\ \text{for}\,(\,j\,=\,6;\,\,j\,<=\,15;\,\,j++) \\ & \text{A[j]}\,=\,15; \end{array}
```

- Examples for transformations with polyhedral model
 - Transformation for improving data locality

```
\begin{array}{lll} \text{for} \, (\, i \, = \, 1; \; i \, <= \, 10; \; i + +) \\ A [\, i \,] \, = \, 10; \\ \text{for} \, (\, j \, = \, 6; \; j \, <= \, 15; \; j + +) \\ A [\, j \,] \, = \, 15; \end{array}
```

```
for (i = 0; i < 8; i++)
sum += A[i];
```

- Examples for transformations with polyhedral model
 - Transformation for improving data locality

```
\begin{array}{lll} & \text{for} \ (i = 1; \ i <= 10; \ i++) \\ & \ A[\ i\ ] = 10; \\ & \ \text{for} \ (j = 6; \ j <= 15; \ j++) \\ & \ A[\ j\ ] = 15; \end{array}
```

```
for (i = 0; i < 8; i++)
sum += A[i];
```

```
\begin{array}{lll} \text{for}\,(\,i\,=\,1;\,\,i\,<=\,5;\,\,i++) \\ & \text{A[i]}\,=\,10; \\ \text{for}\,(\,j\,=\,6;\,\,j\,<=\,15;\,\,j++) \\ & \text{A[j]}\,=\,15; \end{array}
```

```
<create and initialize an array 'tmp'>
for (i = 0; i < 8; i++)
  tmp[i % 4] += A[i];
sum = tmp[0] + tmp[1] + tmp[2] + tmp[3];</pre>
```

- Examples for transformations with polyhedral model
 - Transformation for improving data locality

```
\begin{array}{lll} \text{for} (\ i &= 1; \ i <= 5; \ i++) \\ A[\ i ] &= 10; \\ \text{for} (\ j &= 6; \ j <= 15; \ j++) \\ A[\ j ] &= 15; \end{array}
```

```
for ( i = 0; i < 8; i++) sum += A[i];
```

```
<create and initialize an array 'tmp'>
for (i = 0; i < 8; i++)
    tmp[i % 4] += A[i];
sum = tmp[0] + tmp[1] + tmp[2] + tmp[3];</pre>
```

```
\begin{array}{l} {\sf parfor} \ (\, \text{ii} \, = \, 0; \ \ \text{ii} \, < \, 4; \ \ \text{ii} \, + +) \\ {\sf tmp[\, \text{ii} \, ]} \, = \, 0; \\ {\sf for} \ (\, \text{i} \, = \, \text{ii} \, * \, 2; \ \ \text{i} \, < \, (\, \text{ii} \, + 1) \, * \, 2; \ \ \text{i} \, + +) \\ {\sf tmp[\, \text{ii} \, ]} \, \, + \, {\sf A[\, \text{i} \, ]}; \\ {\sf sum} \, = \, {\sf tmp[\, 0]} \, + \, {\sf tmp[\, 1]} \, + \, {\sf tmp[\, 2]} \, + \, {\sf tmp[\, 3]}; \end{array}
```

Polyhedral representation of programs

- Iteration domain
- Schedule
- Access function

Dynamic instances of each statement is represented as an integer point in statement's polyhedron

- Why not AST?
 - Dynamic instances of statements not captured
 - Rigid data structure
 - Less expressive than polyhedral model

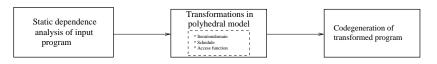


Figure: Transformation in polyhedral model

Iteration domain

Iteration domain

for (int
$$i=2$$
; $i <= 6$; $i++$)
for (int $j=2$; $j <= 6$; $j++$)
if ($i <= j$)
A[i] = 10; // S2

Iteration domain for ${\bf S}1$ is

$$D_{S1} \; = \; \{(i,j) \; \epsilon \; Z^2 \; | \; 2 \; \leq \; i \; \leq \; N \; \wedge \; 2 \; \leq \; j \; \leq \; N \}$$

Iteration domain for \$2 is

$$D_{S2} = \{(i,j) \in Z^2 \mid 2 \leq i \leq 6 \land 2 \leq j \leq 6 \land i \leq j\}$$

Iteration domain

for (int
$$i=2;\ i <= N;\ i++)$$

for (int $j=2;\ j <= N;\ j++)A[i]=10; // S1$

$$\begin{array}{lll} \text{for (int $i=2$; $i<=6$; $i++$)} \\ \text{for (int $j=2$; $j<=6$; $j++$)} \\ \text{if ($i<=j$)} \\ \text{A[$i]} = 10$; $//$ S2 \\ \end{array}$$

Iteration domain for S1 is

$$D_{S1} \ = \ \{(i,j) \ \epsilon \ Z^2 \ | \ 2 \ \leq \ i \ \leq \ N \ \land \ 2 \ \leq \ j \ \leq \ N\}$$

Iteration domain for \$2 is

$$D_{S2} \; = \; \{(i,j) \; \epsilon \; Z^2 \; | \; 2 \; \leq \; i \; \leq \; 6 \; \wedge \; 2 \; \leq \; j \; \leq \; 6 \; \wedge \; i \; \leq \; j \}$$

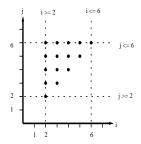


Figure: Graphical representation of iteration domain(S2)

Schedule

• Scattering function

Schedule

Scattering function

```
for (int i = 2; i <= 4; i++)
for (int j = 2; j <= 4; j++)
P[i][j] = A[i] * B[j] ; // S3
```

Examples:

$$\phi_{S3}(i,j) = (i,j)
\phi_{S3}(i,j) = (j,i)$$

Schedule

Scattering function

```
for (int i = 2; i <= 4; i++)
for (int j = 2; j <= 4; j++)
P[i][j] = A[i] * B[j] ; // S3
```

Examples:

```
\phi_{S3}(i,j) = (i,j)
\phi_{S3}(i,j) = (j,i)
```

Code generated by Cloog for $\phi_{S3}(i,j) = (j,i)$

```
for (t1 = 2; t1 <= 4; t1++) {
  for (t2 = 2; t2 <= 4; t2++) {
    i = t2; j = t1;
    P[i+j] += A[i] + B[j];
}
```

Loops are interchanged here by applying this transformation

Access function

$$A[i+j][i+N]$$

Array access function: $F_A(i,j) = (i+j, i+N)$

Change array access function for better locality

SCoP - Static Control Part

Example for SCoP

```
for (i = 0; i < 5*N; i++)
for (j = N; j < 3*i + 5*N + 6; j++)
A[i-j] = A[i];
if (i < N - 10)
A[i + 20] = j;
```

- Structured control flow
 - Regular for loops
 - Conditions
- Affine expressions in:
 - · Loop bounds
 - Conditions
 - Access functions
- Side effect free(Pure functions)



LLVM

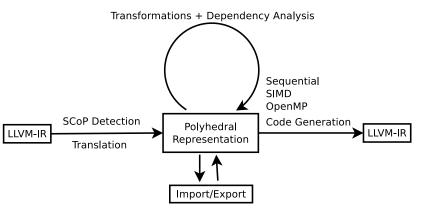
- LLVM (Low Level Virtual Machine)
 - Framework for implementing compilers
 - Common low level code repersentation
 - Lifelong analysis and transformation of programs
- LLVM relaxes SCoP constraints -> more SCoPs are detected
 - Regular for loops -> Anything that acts like a regular for loop
 - Affine experssions -> Expressions that calculates an affine result
 - Side effect free known
 - Memory access through arrays arrays + pointers
- Independent of programming language

Polly I

- Polly (Polyhedral Optimization in LLVM)
 - Implementing Polyhedral Optimization in LLVM
 - Effort towards Auto Parallelism in programs.
- Implementation
 - LLVM-IR to polyhedral model
 - Region-based SCoP detection
 - Semantic SCoPs
 - Polyhedral model
 - The integer set library
 - Composable polyhedral transformations
 - Export/Import
 - Polyhedral model to LLVM-IR
- Related work
 - gcc Graphite



Polly II



External Optimizers / Manual Optimizations

Figure: Architecture of Polly

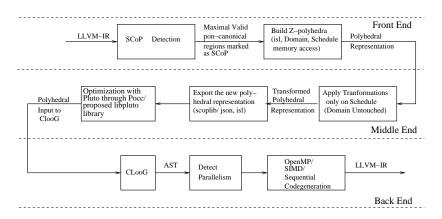


Figure: Detailed control flow in Polly

- Code generation pass in Polly
- Detecting parallelism in Polly
- Generating OpenMP library calls

```
for (int i = 0; i \le N; i++) A[i] = 1;
```

PollyBB

polly.loop_body:

```
%indvar=phi i32{0,%entry}
                        br label %polly BB
%insertInst=Zext i1 true to i16
%omp.userConyext=alloca %foo.omp_subfn.omp.userContext
%o=getelementptr inbounds %foo.omp_subfn.omp.userContext * %omp.userContext
* %omp.userContext, i32 0,i32 0
store [100 x float]*@A,[100xfloat]**%0
%omp_data=bitcast %foo.omp_subfn.omp.userContext * %omp.userContext to i8*
call void @GOMP_parallel_foop_runtime_start
(void(i8 *)*@foo.omp_subfn.i3* %omp_data,i32 0,i32 0,i32 100,i32 1) call void@foo.omp_subfn8*%omp_data)
call void@GOMP_parallel_end()
br label %polly.after loop.region
```

polly.after_loop.region: br label %polly.after_loop

Figure: CFG showing sequence of OpenMP library calls

Support for inner loops

```
for (int i = 0; i < M; i++)
for (int j = 0; j < N; j++)
A[i][j] = A[i-1][j] + B[i-1][j];
```

Surrounding induction variables and parameters need to be passed to the subfunction

Support for inner loops

```
for (int i = 0; i < M; i++)
for (int j = 0; j < N; j++)
A[i][j] = A[i-1][j] + B[i-1][j];
```

Surrounding induction variables and parameters need to be passed to the subfunction

Dealing with memory references

Enabling OpenMP code generation in Polly

```
export LIBPOLLY=<path to cmake>/lib/LLVMPolly.so
pollycc -fpolly -fparallel a.c

OR

# Generate the LLVM-IR files from source code.
clang -S -emit-Ilvm a.c
alias opt="opt -load $LIBPOLLY
# Apply optimizations to prepare code for polly
opt -S -mem2reg -loop-simplify -indvars a.c -o a.preopt.II
# Generate OpenMP code with Polly
opt -S -polly-codegen -enable-polly-openmp a.preopt.II -o a.II
# Link with libgomp
Ilc a.II -o a.s
Ilvm-gcc a.s -lgomp
```

- OpenMP testcases
 - Polly follows LLVM testing infrastrcutre

Testing with PolyBench

- PolyBench
 Benchmarks from
 - linear algebra
 - datamining
 - stencil computation
 - solver and manipulation algorithms operating on matrices

Experimental results

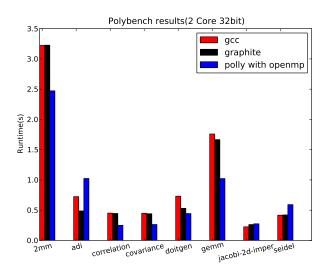


Figure: Performance comparison(2 core 32 bit)

Experimental results

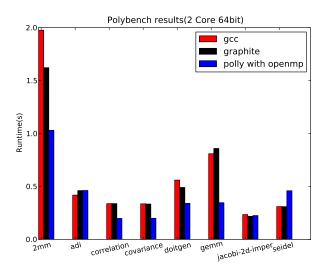


Figure: Performance comparison(2 core 64bit)

Experimental results

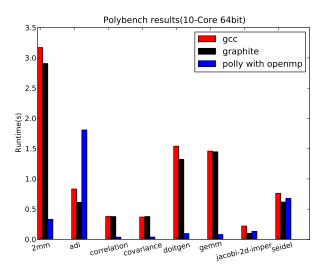


Figure: Performance comparison(10-core 64 bit)

Conclusion and Future Work

- Conclusion
- Support for memory access transformations in Polly
- Increasing coverage of Polly
 - Increasing SCoP coverage
 - Increasing the system coverage
- Integrating profile guided optimization into Polly

Setting up the environment

- CLooG
- PoCC
- Scoplib
- Building LLVM with Polly

Various Tools Used in Polyhedral Community

- ClooG
- PLUTO
- VisualPolylib