# Universidad de Buenos Aires Facultad de Ciencias Exactas y Naturales Departamento de Computación

Teoría de Lenguajes Primer Cuatrimestre de 2013

## Trabajo práctico

Micro HTML Prettyprint

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#### 1. Introducción

En el presente trabajo práctico definiremos la gramática para generar un lenguaje que interprete una versión simplificada de HTML.

#### 2. Asunciones y aclaraciones

A continuación detallaremos una serie de asunciones que hemos tomado al respecto del trabajo:

- Asumimos que la sección head tiene que ir necesariamente antes que la sección body.
- Asumimos que los tags < title >< /title > solo puede aparecer una sola vez y el tag
   < script >< /script > puede aparecer varias veces además de que ambos son opcionales
   y pueden estar en cualquier orden.
- Asumimos que dentro de la sección title puede haber cualquier texto sin tags definidos por el lenguaje.
- Asumimos que dentro de la sección script no pueden contener otras subsecciones script
- Asumimos que los espacios entre tags son inicialmente filtrados por el analizador léxico y que no llegan a la gramática así como también los tags de comentarios.
- Asumimos en las regular expressions que los nombres de los tags son todos en minúsculas y sin ningún espacio.
- En los gráficos de árboles de derivación los ejemplos se basan en código de html para mejor legibilidad, pero en realidad la derivación se realizará en base a los tokens traducidos por el analizador léxico.

#### 3. Gramática

```
G = \langle V_N, V_T, P, BEGIN \rangle donde
   V_N = \{ BEGIN, HTML, HEAD, LEER_HEAD, TITLE, SCRIPTS, MORE_SCRIPTS, BODY,
           LEER_BODY }
   V_T = \{ \text{ initHtml, endHtml, initHead, endHead, initTitle, endTitle, initScript, endScript, } \}
           noScripts, initBody, endBody, initDiv, endDiv, initH1, endH1, initP, endP, br,
   y P está dada por:
          BEGIN \longrightarrow
                          initHtml HTML endHtml
           \mathrm{HTML} \,\longrightarrow\,
                          HEAD BODY
           \text{HEAD} \longrightarrow
                          \lambda | initHead LEER_HEAD endHead
                          SCRIPTS TITLE MORE_SCRIPTS |
    \text{LEER\_HEAD} \longrightarrow
                          TITLE MORE_SCRIPTS | MORE_SCRIPTS
           \mathrm{TITLE} \longrightarrow
                          initTitle noTags endTitle
        {\rm SCRIPTS} \longrightarrow
                          initScript noScripts endScripts MORE_SCRIPTS
MORE\_SCRIPTS \longrightarrow
                          \lambda | initBody LEER_BODY endBody
           BODY \longrightarrow
                          \lambda | initBody LEER_BODY endBody
    LEER\_BODY \longrightarrow
                          \lambda | noTags LEER_BODY | br LEER_BODY |
                          initDiv LEER_BODY endDiv LEER_BODY |
                          initP LEER_BODY endP LEER_BODY |
                          initH1 LEER_BODY endH1 LEER_BODY
```

#### 4. Expresiones regulares

Terminal	Expresión regular
initHtml	<html></html>
endHtml	
initHead	<head></head>
endHead	
initBody	<body></body>
endBody	
initTitle	<title>&lt;/td&gt;&lt;/tr&gt;&lt;tr&gt;&lt;td&gt;endTitle&lt;/td&gt;&lt;td&gt;</title>
initScript	<script></td></tr><tr><td>endScript</td><td></script>
initDiv	<div></div>
endDiv	
initH1	<h1></h1>
endH1	
initP	
endP	
br	
noTags	^(((?!(TAGS_GRAMATICA)).)*\$
noScripts	^((?!( <script>   </script> )).)*\$

$$\label{eq:condeta} Donde\ TAGS\_GRAMATICA =  [|](html)  | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |  [|  \[|  \\[|  \\\[|  \\\\[|  \\\\\[|  \\\\\\[|  \\\\\\\[|  \\\\\\\\[|  \\\\\\\\\[|  \\\\\\\\\\[|  \\\\\\\\\\\[|  \\\\\\\\\\\\[|  \\\\\\\\\\\\\[|  \\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|  \\[|  \\\[|\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\]\\\(html\\\)\\]\\(html\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\\]\\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\\)\\\\\\\\\\\\\\]\\\\\\\\\\\\\\(html\\\\\\\\\\\\\\)\\\\\\\\\\\\\]\\\\\\\\\\\\\(html\\\\\\\\\\\\\)\\\\\\\\\\\\]\\\\\\\\\\\\(html\\\\\\\\\\\\)\\\\\\\\\\\]\\\\\\\\\\\(html\\\\\\\\\\\)\\\\\\\\\\]\\\\\\\\\\(html\\\\\\\\\\)\\\\\\\\\]\\\\\\\\\(html\\\\\\\\\)\\\\\\\\]\\\\\\\\(html\\\\\\\\)\\\\\\\]\\\\\\\(html\\\\\\\)\\\\\\]\\\\\\(html\\\\\\)\\\\\]\\\\\(html\\\\\)\\\\]\\\\(html\\\\)\\\]\\\(html\\\)\\]\\(html\\)\]\(html\)](html)$$

El tag no Tags se representa con una expresión regular que identifica cualquier texto sin tags definidos por la gramática.

Por otro lado, la expresión regular para identificar los comentarios es la siguiente:

$$comment =$$

Los diferentes símbolos utilizados en las expresiones denotan:

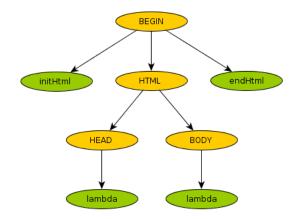
- $\longrightarrow$  Indica disyunción entre dos expresiones.
- $\hat{\ }\longrightarrow$  Matchea la expresión al comienzo de una línea.
- $\$ \longrightarrow Matchea la expresión al final de una línea.$
- ?! --> No machea la aparición denotada a su derecha.
- $* \longrightarrow$  Indica una repetición de 0 o más veces.

### 5. Árboles de derivación

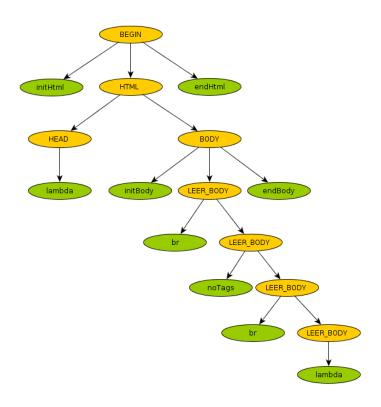
A continuación mostramos varios ejemplos de derivaciones válidas:

• Ejemplo 1: Html Empty.

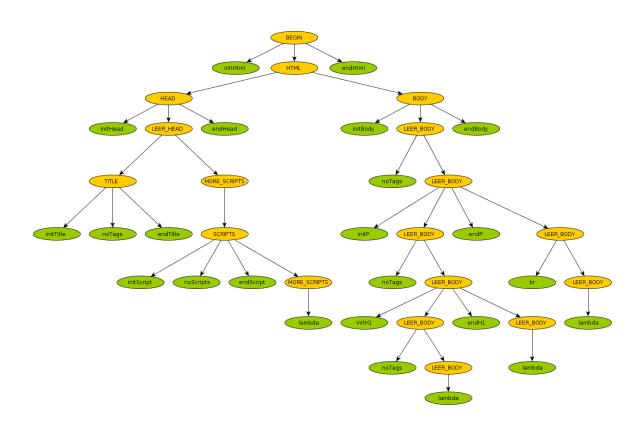
<html>



• Ejemplo 2: Empty head body with text.

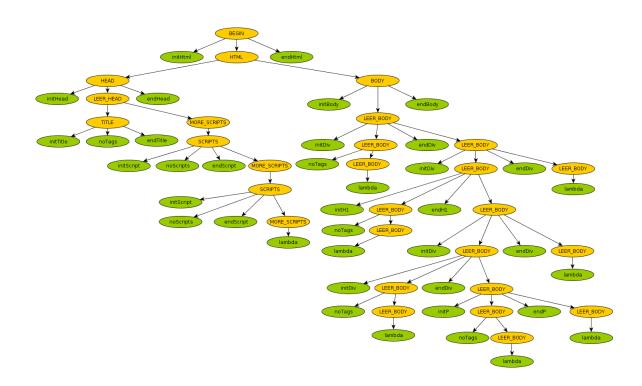


• Ejemplo 3: Tp example.

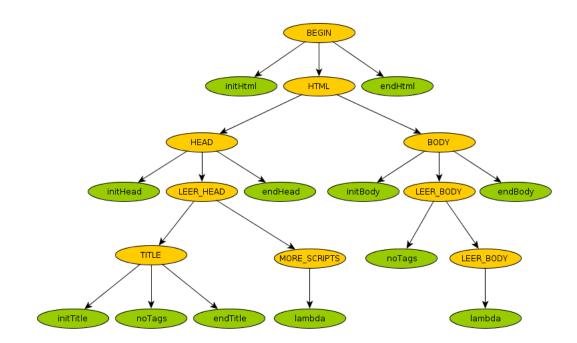


■ Ejemplo 4: Huge Body.

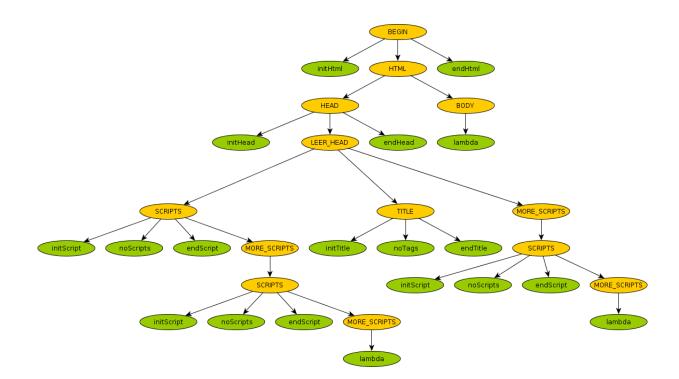
```
<html>
   <head>
     <title> Something's Gone Terribly Wrong </title>
     <script> </script>
     <script> bN_cfg = {p: {"dL_ch":"us.hpmguncat",
               "dL_dpt":"error","cobrand":"HuffPost"}};
      </script>
   </head>
   <body>
     <div> </div>
     <div>
        <h1> Oh, Noes! A 404! </h1>
        <div>
            <div> </div>
            or 
        </div>
     </div>
   </body>
</html>
```



■ Ejemplo 5: Just title empty body.



• Ejemplo 6: Several scripts middle title no body.



• Ejemplo 7: Just scripts nested tags and text.

