

# The last, the least and the urgent...

A story of three policies

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# Outline

Introduction

A crash course on measure valued processes

Partial service queues and Earliest-Deadline-First

Deadline-oblivious policies

Simulations

Final remarks

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## Introduction

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# Motivation

A bit of history...

- Several queueing systems have service and **timing** requirements.
- Examples:
  - Computing tasks with real-time constraints.
  - Item delivery problems in logistics.
  - Emergency response.
  - etc. etc. etc.
- This has led to a long and rich history of research about **queues with abandonments** [Barrer, 1957; Stanford, 1979; Baccelli et al., 1984].

# Motivation

Recent developments...

One of the most used policies is **Earliest-Deadline-First (EDF)**

- Give priority to tasks with more urgent deadlines.

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Through fluid limits and diffusion approximations, establish performance:

- [Decreusefond and Moyal, 2008] establish EDF fluid limits in the single server case.
- [Kruk et al., 2011] provides diffusion approximations.
- [Moyal, 2013] establish some optimality properties of EDF.
- [Kang and Ramanan, 2010, 2012] analyze the many-server case.
- [Atar et al., 2018, 2023] establish asymptotic performance.

and many others...

## Common assumption

Customers renege *only* in the queue, and not during service.

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We call this the *call-center scenario*:

- Akin to waiting for the customer-help line to pick your call while you listen to annoying music.
- The underlying idea is that when a task reaches service, it will stay until completion.

**Key performance metric:** number of satisfied tasks (or reneging probability).



# Motivation

## Partial service queues

In several queueing systems:

- Tasks may abandon during service.
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We call this setting *queues with partial service*.

Some examples:

- Electrical vehicle charging: customers leave the system with a *partial charge*.
- LLM inference: longer computation times lead to better answers, but these may be interrupted to deliver a quick response.
- File transfers over the Internet, that can be resumed later.

# Key points of this talk

- Provide some suitable representation of the state space and dynamics of these partial service queues.
- Analyze several interesting policies under a suitable fluid model.
- Compute the main performance metric here: [attained work](#).
- *Last but not least*: show that the simple LCFS policy [exhibits the same performance](#) than EDF in this setting, without using deadline information.

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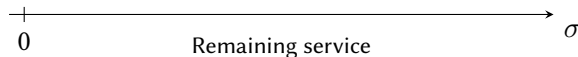
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# Measure valued stochastic processes

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- Tasks arrive as a Poisson process of intensity  $\lambda$ .
- Each task has a service requirement  $S \sim g(\sigma)$ .



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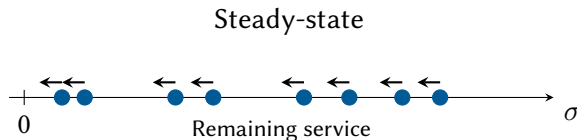




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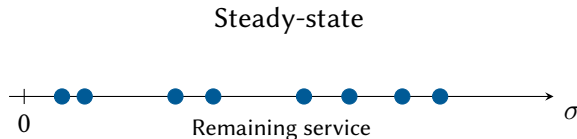
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State-descriptor:

$$\Phi_t = \sum_i \delta_{\sigma_i(t)}$$

a Point-process on the positive half-line.

- $\Phi_t$  is a measure-valued Markov process.
- Its dynamics can be characterized through its generator.
- In steady state:

$$\Phi \sim \text{Poisson Process with mean measure } \mu(d\sigma) = \lambda \bar{G}(\sigma) d\sigma$$

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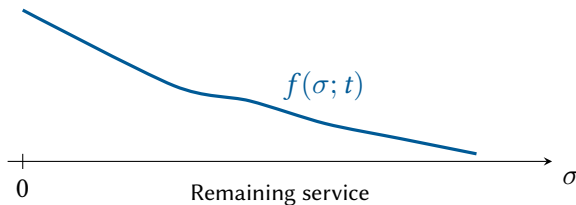
where  $\bar{G}$  is the CCDF of  $S$ .

### Interpretation:

- Write  $\mu(d\sigma) = \rho \left[ \frac{1}{E[S]} (1 - G(\sigma)) \right] d\sigma$ , with  $\rho = \lambda E[S]$ .
- Then  $\left[ \frac{1}{E[S]} (1 - G(\sigma)) \right] d\sigma$  is the *residual service time distribution* associated to  $G$ .
- In steady-state, the total number of customers  $\sim \text{Poisson}(\rho)$  and distributed in  $\sigma$  as the residual lifetime distribution.

## M/G/∞, fluid approximation.

Suppose that we can replace  $\Phi_t$  by a general measure  $\mu_t$  with density  $f(\sigma; t)$ .



- Mass is transported to the left at rate 1.
- New mass arrives at  $\sigma$  with intensity  $\lambda g(\sigma) d\sigma dt$ .

We can combine this in the following [transport equation](#):

$$\frac{\partial f}{\partial t} = -\frac{\partial f}{\partial \sigma} + \lambda g(\sigma).$$

## M/G/ $\infty$ , fluid approximation.

Imposing equilibrium and the boundary condition  $f(\sigma) \rightarrow 0$  as  $\sigma \rightarrow \infty$  we get:

$$\frac{\partial f}{\partial \sigma} + \lambda g(\sigma) = 0 \implies f(\sigma) = \lambda \int_{\sigma}^{\infty} g(u) du = \lambda \bar{G}(\sigma),$$

so the fluid approximation recovers the mean measure of  $\Phi$ .

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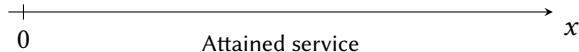
so the fluid approximation recovers the mean measure of  $\Phi$ .

- This is a deterministic measure, with total mass  $\rho$ ...
- ...distributed in the real line as the residual service distribution.
- Serves as an approximation of  $\Phi$  in a large scale system ( $\lambda \rightarrow \infty$ ).

## $M/G/\infty$ : take two

Attained service state descriptor

Here is another approach to model the same system [Kang and Ramanan, 2010]:





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New mass arrives, rate  $\lambda dt$

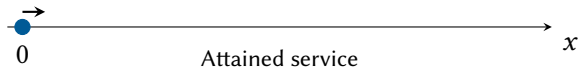


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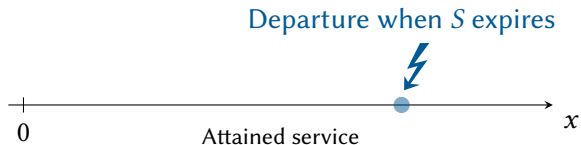
Drifts to the right at rate 1



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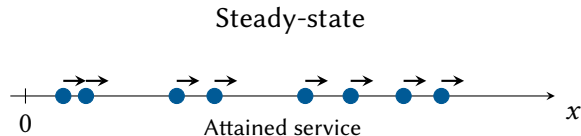
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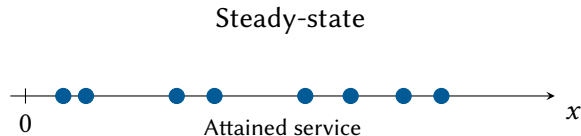
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State-descriptor:

$$\tilde{\Phi}_t = \sum_i \delta_{x_i(t)}$$

a Point-process on the positive half-line, where  $x_i(t)$  is the elapsed time in the system

# M/G/ $\infty$ , take two

## Steady-state

$\tilde{\Phi}_t$  is a measure-valued Markov process.

- Mass always arrive at 0 with rate  $\lambda dt$ .
- Transports to the right at rate 1.
- Leaves the system at rate  $h(x)$ , the **hazard rate function**:

$$h(x) = \lim_{dt \rightarrow 0} P(S \in [x, x + dt] \mid S > x) = \frac{g(x)}{\bar{G}(x)} = -\frac{\partial}{\partial x} \log \bar{G}(x).$$

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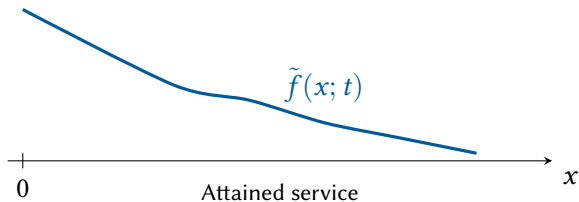
$$\tilde{\Phi} \sim \text{Poisson Process with mean measure } \nu(dx) = \lambda \bar{G}(x) dx$$

So the reversed representation has the same distribution, because in a random point in time the elapsed service and the remaining service have the same distribution.

# M/G/∞: take two

Fluid approximation.

Suppose that we can replace  $\tilde{\Phi}_t$  by a general measure  $\nu_t$  with density  $\tilde{f}(x; t)$ .

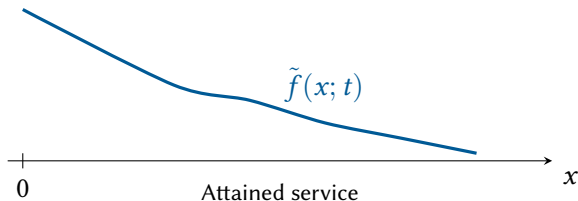




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Suppose that we can replace  $\tilde{\Phi}_t$  by a general measure  $\nu_t$  with density  $\tilde{f}(x; t)$ .



The corresponding transport equation is (informally):

$$\frac{\partial \tilde{f}}{\partial t} = -\frac{\partial \tilde{f}}{\partial x} - h(x)\tilde{f} + \lambda\delta_0.$$

## M/G/∞: take two

Fluid equilibrium.

Imposing equilibrium we get:

$$\frac{\partial \tilde{f}}{\partial x} = -h(x)\tilde{f} + \lambda\delta_0.$$

Solving (in a distribution sense) with the boundary condition  $\tilde{f}(\infty) = 0$  we get:

$$\tilde{f}(x) = \lambda e^{-\int_0^x h(u)du}.$$

But by definition  $\int_0^x h(u)du = -\log \bar{G}(x)$ , and thus:

$$\tilde{f}(x) = \lambda \bar{G}(x)$$

So the transport fluid equation recovers again the mean measure of the steady-state.

- We can model  $M/G$  systems by using two state descriptors:
  - The remaining service  $\Phi$ .
  - The attained service  $\tilde{\Phi}$ .
- Both admit reasonable fluid approximations, which correspond to transport equations.
- In fact this has been used in the literature to model abandonments (since they operate as  $M/G/\infty$  systems in some sense).

# Lessons learned

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**Question:** can we do more using this machinery of measure-valued processes?

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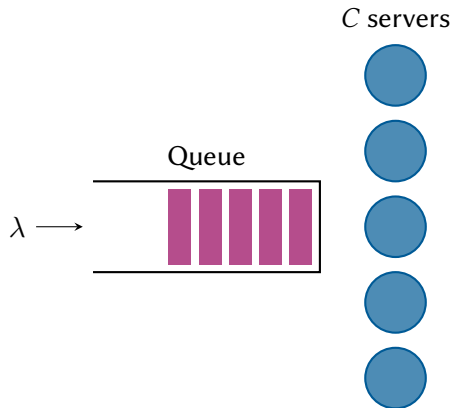
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# Partial service queues

## Setting

Consider an  $M/G/C$  system where:

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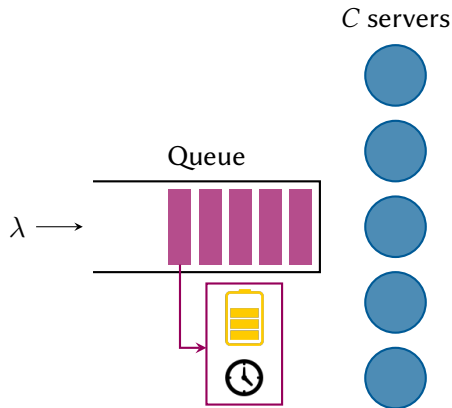


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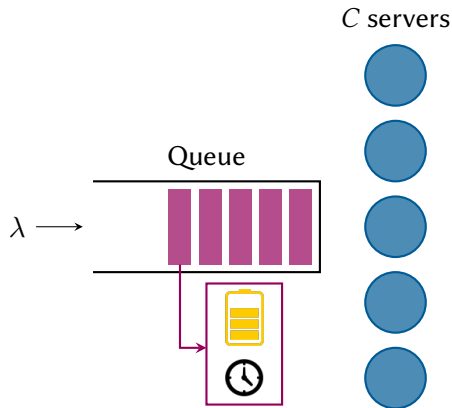


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  - $S_i$ : service time (at rate 1).
  - $T_i$ : sojourn time or deadline.
- $(S_i, T_i)$  are independent across jobs.
- Follow a common distribution  $G(\sigma, \tau)$ , possibly correlated.





# Partial service queues

## Definition

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- In particular, they may leave **during service**.
- Key performance metrics:
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  - Equivalently,  $S_r := S - S_a$ , amount of service **reneged**.

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- **Problem**: we have to keep track of remaining service and deadlines simultaneously!

- Before proceeding, it is useful to define the **system load**:

$$\rho := \lambda E[\min\{S, T\}].$$

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- **Interpretation**: the mean number of customers on a system with  $C = \infty$ .
- What we expect in a large scale fluid model:
  - If  $\rho < C$  (underload), all tasks can be served,  $S_a = \min\{S, T\}$ .
  - If  $\rho > C$  (overload), demand *curtailing* will occur. How? It depends on the policy...

# System evolution

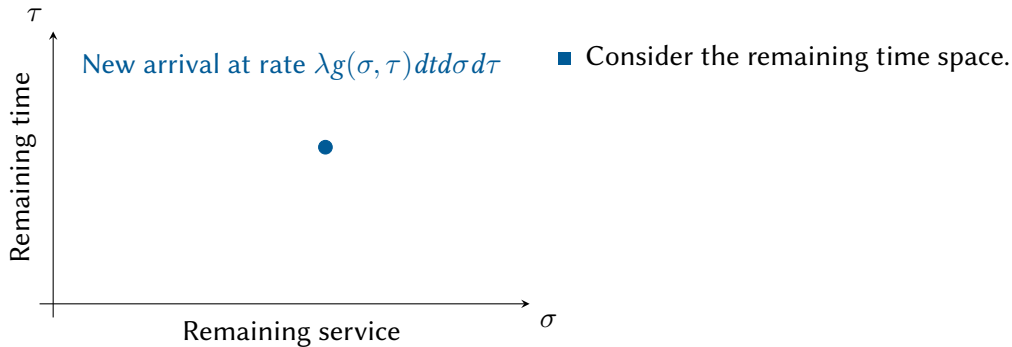
## Remaining service times



- Consider the remaining time space.

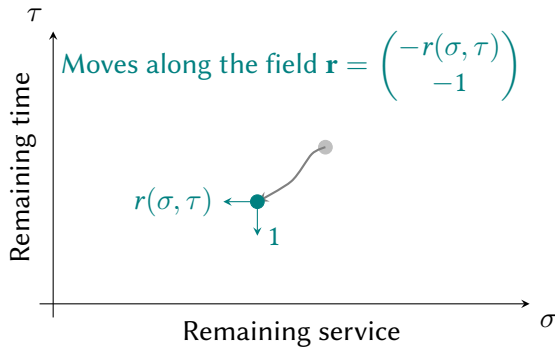
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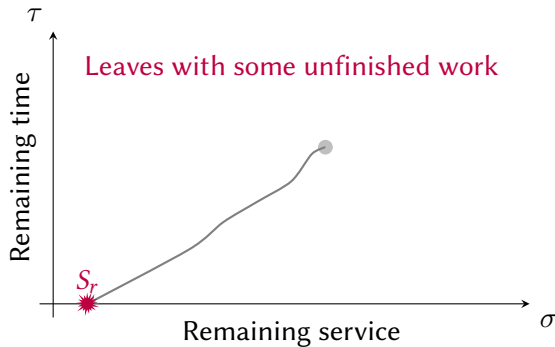


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- May depend on any combination of  $(\sigma, \tau)$ .



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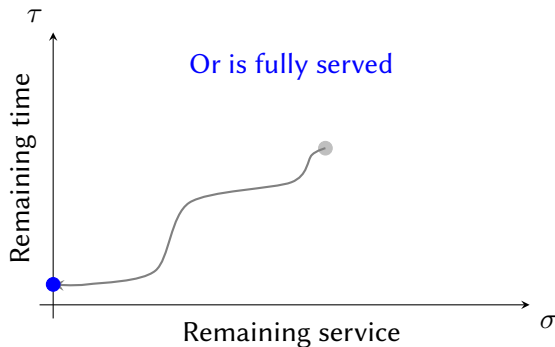
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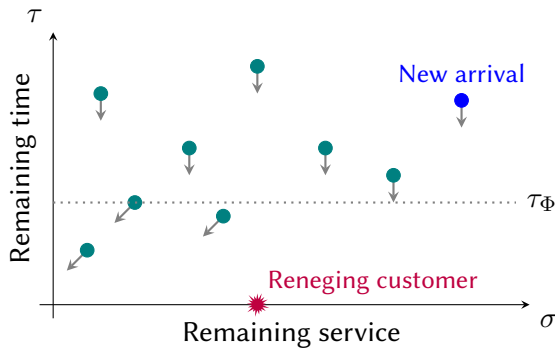
## Remaining service times



- Consider the remaining time space.
- **Policy** defines how tasks are served.
- May depend on any combination of  $(\sigma, \tau)$ .
- State descriptor:

$$\Phi_t = \sum_i \delta_{(\sigma_i(t), \tau_i(t))}$$

## Example: Earliest-deadline-first



New arrival

- Serve the  $C$  most urgent customers.
- Corresponds to taking:

$$r_\Phi(\sigma, \tau) = \mathbf{1}_{\{\tau \leq \tau_\Phi\}}$$

with

$$\tau_\Phi := \sup\{\tau \geq 0 : \Phi(\mathbb{R}_+ \times (0, \tau]) < C\}.$$

- Replace  $\Phi_t$  by a (fluid) measure  $\mu_t$ .
- Now mass drifts along the field:

$$\mathbf{r}_\mu(\sigma, \tau) = \begin{pmatrix} -r_\mu(\sigma, \tau) \\ -1 \end{pmatrix}$$

- With  $r_\mu$  satisfying:

$$0 \leq r_\mu \leq 1$$

$$\iint r_\mu(\sigma, \tau) \mu(d\sigma, d\tau) \leq \min\{\mu(\mathbb{R}_{++}^2), C\}.$$

We will describe these dynamics in terms of the projections

$$\langle \varphi, \mu \rangle := \iint \varphi(\sigma, \tau) \mu(d\sigma, d\tau)$$

of the state measure with respect to a test function  $\varphi : \mathbb{R}_{++}^2 \rightarrow \mathbb{R}$ , with continuous derivatives and compact support, i.e.  $\varphi \in \mathcal{C}_c^1(\mathbb{R}_{++}^2)$ .

We have:

$$\langle \varphi, \mu_{t+dt} \rangle = \iint \varphi(\sigma - r_{\mu_t} dt, \tau - dt) \mu_t(d\sigma, d\tau) + \lambda dt \iint \varphi(\sigma, \tau) g(\sigma, \tau) d\sigma d\tau + o(dt).$$

# Fluid model dynamics

## Weak formulation

$$\begin{aligned}\frac{\partial}{\partial t} \langle \varphi, \mu_t \rangle &= \lim_{dt \rightarrow 0} \iint \frac{1}{dt} [\varphi(\sigma - r_{\mu_t} dt, \tau - dt) - \varphi(\sigma, \tau)] \mu_t(d\sigma, d\tau) \\ &\quad + \lambda \iint \varphi(\sigma, \tau) g(\sigma, \tau) d\sigma d\tau \\ &= - \iint [r_{\mu_t}(\sigma, \tau) \varphi_\sigma(\sigma, \tau) + \varphi_\tau(\sigma, \tau)] \mu_t(d\sigma, d\tau) + \lambda \iint \varphi(\sigma, \tau) g(\sigma, \tau) d\sigma d\tau,\end{aligned}$$

Equivalently:

$$\begin{aligned} \langle \varphi, \mu_t \rangle = \langle \varphi, \mu_0 \rangle + \int_0^t \bigg[ & - \iint [r_{\mu_s}(\sigma, \tau) \varphi_\sigma(\sigma, \tau) + \varphi_\tau(\sigma, \tau)] \mu_t(d\sigma, d\tau) \\ & + \lambda \iint \varphi(\sigma, \tau) g(\sigma, \tau) d\sigma d\tau \bigg] ds, \end{aligned}$$

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for any  $\varphi \in \mathcal{C}_c^1(\mathbb{R}_{++}^2)$ .

Looks daunting, but is not that bad...



If  $\mu_t$  admits a density  $f(\sigma, \tau; t)$  with respect to the Lebesgue measure, it corresponds to:

$$\frac{\partial f}{\partial t} + \nabla \cdot [\mathbf{r}_{\mu_t} f] = \lambda g$$

a transport equation.

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Example: EDF

$$\frac{\partial f}{\partial t} = \frac{\partial f}{\partial \sigma} \mathbf{1}_{\{\tau < \tau_{\mu_t}\}} + \frac{\partial f}{\partial \tau} + \lambda g$$

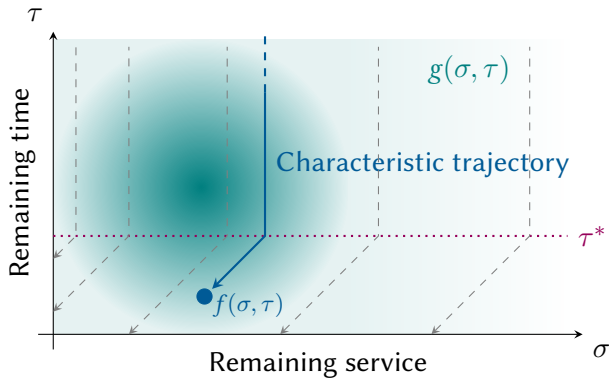
Imposing equilibrium we get:

- $\tau_{\mu^*} = \tau^*$  becomes a constant.
- The measure  $\mu^*$  must satisfy:

$$\frac{\partial f}{\partial \sigma} \mathbf{1}_{\{\tau < \tau^*\}} + \frac{\partial f}{\partial \tau} + \lambda g = 0.$$

- Linear PDE that can be easily solved by the method of characteristics.

# Solving the EDF transport equation



### Theorem

*Assume that  $\rho > C$  and the equation*

$$\lambda E[\min\{S, T, \tau^*\}] = C$$

*has a unique solution  $\tau^* > 0$ . Consider the measure  $\mu^*$  given by the following density:*

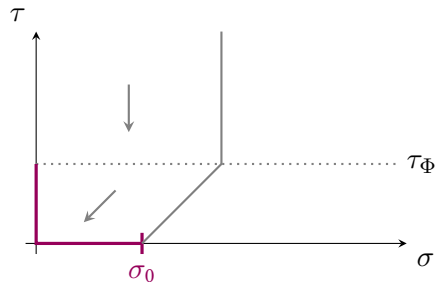
$$f(\sigma, \tau) = \lambda \left[ \int_0^{(\tau^* - \tau)^+} g(\sigma + u, \tau + u) du + \int_{(\tau^* - \tau)^+}^{\infty} g(\sigma + (\tau^* - \tau)^+, \tau + u) du \right].$$

*This measure is a fluid equilibrium for the EDF policy, and*

$$\tau^* = \sup \{ \tau \geq 0 : \mu^*(\mathbb{R}_{++} \times (0, \tau]) \leq C \}.$$

# EDF performance in equilibrium

- Let us compute the rate at which work is **reneged**.
- Compute the rate at which mass exits with  $S_r < \sigma_0$ .



## Proposition

$$\int_0^{\tau^*} f(0, \tau) d\tau + \int_0^{\sigma_0} f(\sigma, 0) d\sigma = \lambda P(S - \min\{S, T, \tau^*\} < \sigma_0).$$

$$\text{i.e. } S_a = S - S_r = \min\{S, T, \tau^*\}.$$

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# What if we do not know the deadlines?

- Deadlines are often hard to estimate in practice.
- Moreover, tasks may under-report their deadline to get priority!
- What about **deadline-oblivious** policies?
  - Can we model them?
  - What is their performance?



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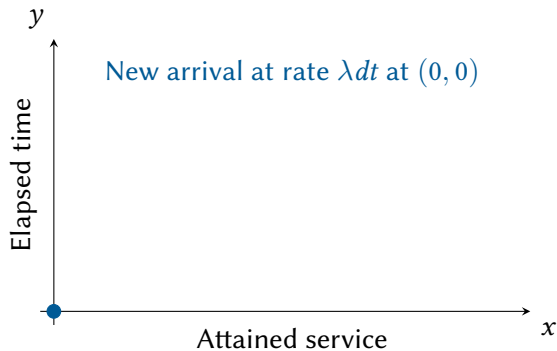
**Problem:** we need a new state-space...

# Attained service state descriptor



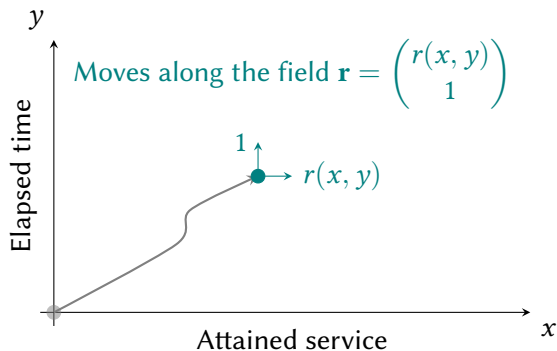
- Consider the elapsed time space.

# Attained service state descriptor



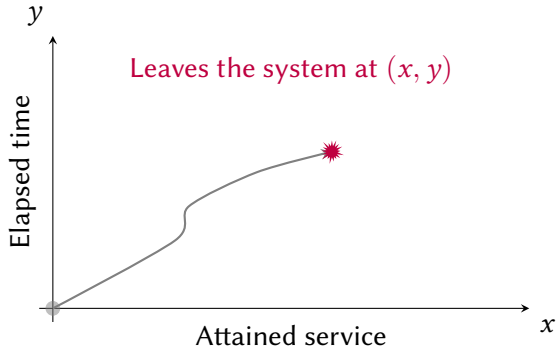
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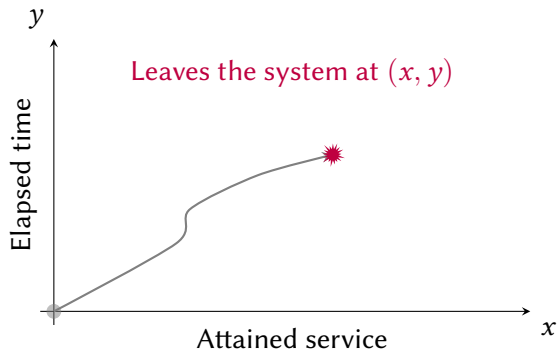
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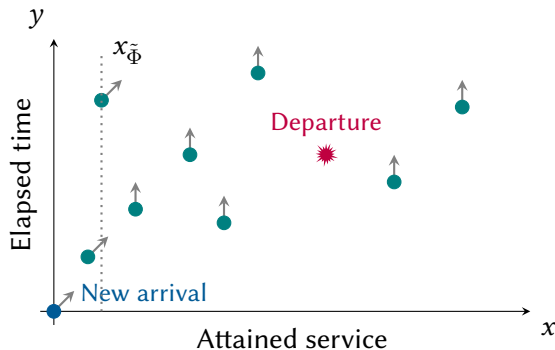
# Attained service state descriptor



- Consider the elapsed time space.
- **Policy** again defines how tasks are served.
- May depend on any combination of  $(x, y)$ .
- State descriptor:

$$\tilde{\Phi}_t = \sum_i \delta_{(x_i(t), y_i(t))}$$

## Example: Least-Attained-Service policy



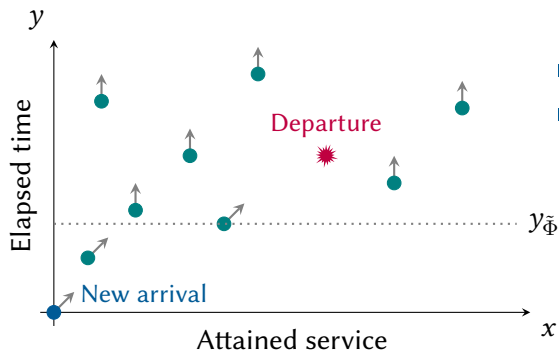
- Serve the  $C$  least-served tasks.
- Corresponds to taking:

$$r_{\tilde{\Phi}}(x, y) = \mathbf{1}_{\{x \leq x_{\tilde{\Phi}}\}}$$

with

$$x_{\tilde{\Phi}} := \sup\{x : \tilde{\Phi}([0, x] \times \mathbb{R}_+) \leq C\}.$$

## Example: Last-Come-First-Served policy



■ Serve the  $C$  more recent tasks.

■ Corresponds to taking:

$$r_{\tilde{\Phi}}(x, y) = \mathbf{1}_{\{y \leq y_{\tilde{\Phi}}\}}$$

with

$$y_{\tilde{\Phi}} := \sup\{y : \tilde{\Phi}(\mathbb{R}_+ \times [0, y]) \leq C\}$$



# The hazard rate field

We have a new problem: what is the rate at which users **leave** the system?

# The hazard rate field

We have a new problem: what is the rate at which users **leave** the system?

Let  $\bar{G}(x, y) = P(S > x, T > y)$  and define:

## Definition (Hazard rate field)

$$\mathbf{h}(x, y) = -\nabla \log \bar{G}(x, y) \quad \text{i.e.}$$

- $h^x(x, y) = P(S \in [x, x + dx], T > S \mid S > x, T > y)$
- $h^y(x, y) = P(T \in [y, y + dy], S > T \mid S > x, T > y)$

Interpretation:  $\mathbf{h}$  stores the rate at which  $\min\{S, T\}$  is attained due to  $S$  or  $T$  expiring.

- Replace  $\tilde{\Phi}_t$  by a (fluid) measure  $\nu_t$ .
- Now mass arrives at  $(0, 0)$  at rate  $\lambda$ .
- Drifts along the field:

$$\mathbf{r}_\nu(x, y) = \begin{pmatrix} r_\nu(x, y) \\ 1 \end{pmatrix}$$

- With  $r_\nu$  satisfying:

$$0 \leq r_\nu \leq 1$$

$$\iint r_\nu(x, y) \nu(dx, dy) \leq \min\{\nu(\mathbb{R}_+^2), C\}.$$

Now we have to compute the departure rate  $\eta_\nu(x, y)$ :

$$\eta_\nu(x, y) := \lim_{dt \rightarrow 0} \frac{1}{dt} P(\{S \in (x, x + r_{\tilde{\Phi}} dt)\} \cup \{T \in (y, y + dt)\} \mid S > x, T > y)$$

By the chain rule and some computations:

$$\eta_\nu(x, y) = \frac{1}{\bar{G}(x, y)} \left[ -\frac{\partial}{\partial x} \bar{G}(x, y) r_{\tilde{\Phi}}(x, y) - \frac{\partial}{\partial y} \bar{G}(x, y) \right]$$

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Therefore:

$$\eta_\nu(x, y) = h^x(x, y) r_\nu(x, y) + h^y(x, y) = \mathbf{r}_\nu(x, y) \cdot \mathbf{h}(x, y).$$

# Attained service transport equation

- We now have all ingredients to formulate the dynamics of the system.
- The transport equation in the elapsed service space is (informally):

$$\frac{\partial \bar{f}}{\partial t} + \nabla \cdot [\mathbf{r}_{\nu_t} \bar{f}] + [\mathbf{r}_{\nu_t} \cdot \mathbf{h}] \bar{f} = \lambda \delta_{(0,0)}.$$

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where  $\tilde{f}$  is the density of  $\nu_t$ .

- The above equation must be treated in weak form:
  - To account for the impulse mass at  $(0, 0)$  driving the system.
  - To allow solutions without a density as we shall see.

# Last come first served

## Fluid equilibrium

Recall that LCFS can be modeled by:

$$r_\nu(x, y) = \mathbf{1}_{\{y < y_\nu\}}$$

with

$$y_\nu = \sup \{y \geq 0 : \nu(\mathbb{R}_+ \times [0, y]) \leq C\}.$$



# Last come first served

## Fluid equilibrium

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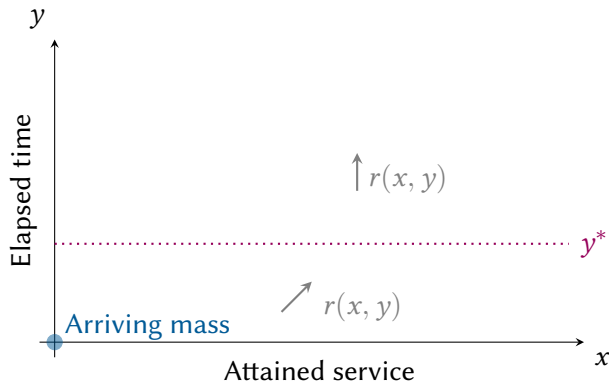
$$y_\nu = \sup \{y \geq 0 : \nu(\mathbb{R}_+ \times [0, y]) \leq C\}.$$

Imposing equilibrium,  $\nu^*$ ,  $y^*$  fixed, we have to solve:

$$\nabla \cdot [\mathbf{r}_{\nu^*} \bar{f}] + [\mathbf{r}_{\nu^*} \cdot \mathbf{h}] \bar{f} = \lambda \delta_{(0,0)}.$$

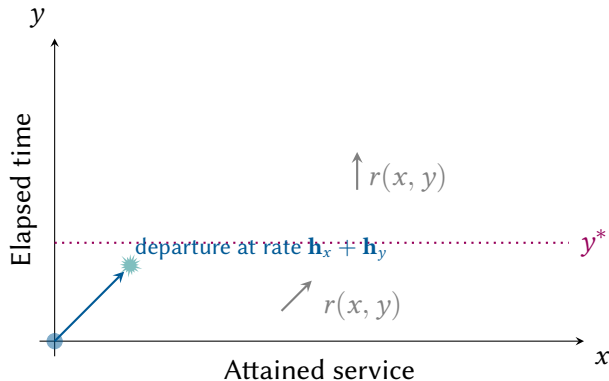
# Solving the transport equation

Last come first served case



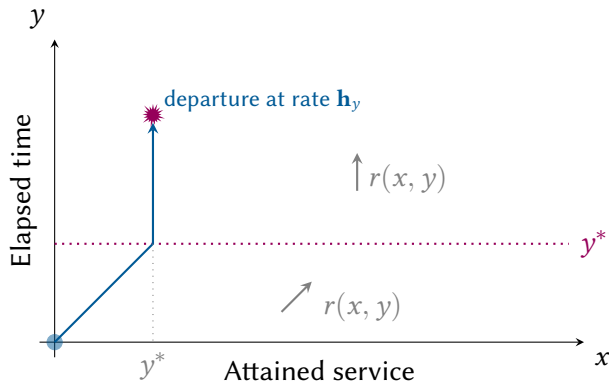
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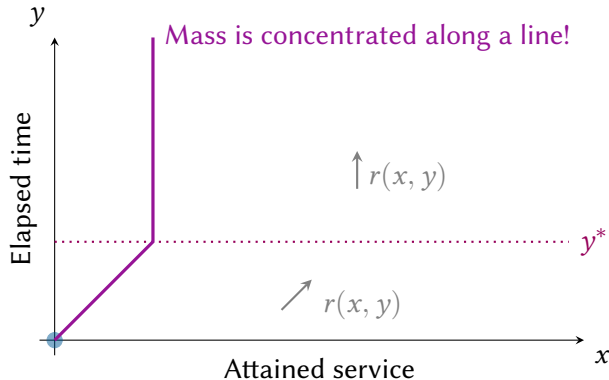
# Solving the transport equation

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# Solving the transport equation

Last come first served case



# Deadline-oblivious policies in overload

## Theorem

*Assume that  $\rho > C$  and the equation*

$$\lambda E[\min\{S, T, z^*\}] = C$$

*has a unique solution  $z^* > 0$ . Consider the measure  $\nu^*$  given by:*

$$\langle \varphi, \nu^* \rangle = \lambda \left[ \int_0^{z^*} \varphi(u, u) \bar{G}(u, u) du + \int_{z^*}^{\infty} \varphi(z^*, u) \bar{G}(z^*, u) du \right],$$

*for all  $\varphi \in C_c(\mathbb{R}_+^2)$ . Then this measure is the equilibrium measure for both the Least-Attained-Service and Last-Come-First-Served policies.*

## LAS/LCFS performance in equilibrium

Compute the rate at which mass leaves the system with less than  $x_0$  attained service:

$$\iint_{[0, x_0] \times \mathbb{R}_+} \eta_{\nu^*}(x, y) \nu^*(dx, dy).$$

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## Proposition

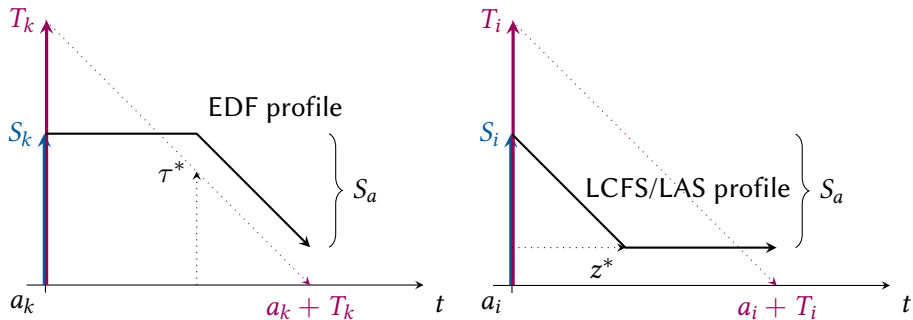
*Assume that  $\rho > C$ . Then*

$$\int_{[0, x_0] \times \mathbb{R}_+} [h^x(x, y) \mathbf{1}_{\{y < z^*\}} + h^y(x, y)] \nu^*(dx, dy) = \lambda P(\min\{S, T, z^*\} \leq x_0).$$

So again the attained work is  $S_a = \min\{S, T, z^*\}!!$



# Graphical explanation



Since  $\tau^* = x^* = y^* = z^*$ , performance is the same in all three policies!!!

# Outline

Introduction

A crash course on measure valued processes

Partial service queues and Earliest-Deadline-First

Deadline-oblivious policies

**Simulations**

Final remarks

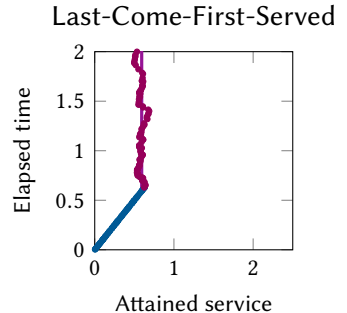
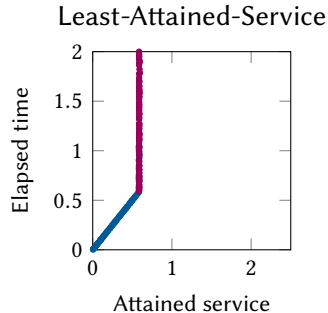
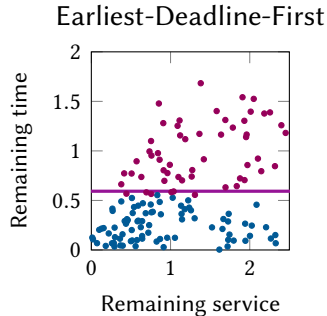
# Simulations with correlated $S$ and $T$

- We finally validate our fluid approximation by stochastic simulations
- In order to account for correlations, we take:

$$S = e^U \quad \text{and} \quad T = e^V \quad \text{with} \quad (U, V) \sim \mathcal{N} \left( \begin{pmatrix} 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 & 0.9 \\ 0.9 & 1 \end{pmatrix} \right).$$

- In particular, the random variables  $U$  and  $V$  are correlated with normal distributions, and therefore  $S$  and  $T$  are correlated with log-normal distributions.
- In this case,  $E[\min\{S, T\}] \approx 1.37$  can only be numerically estimated.
- We choose  $\lambda = 200$  and  $C = 100$ , then  $z^* \approx 0.593$ .

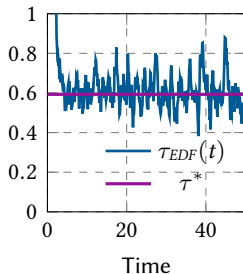
# State space snapshots



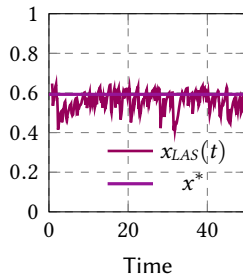
Blue dots are in service, red dots are not in service.

# Stochastic threshold evolution

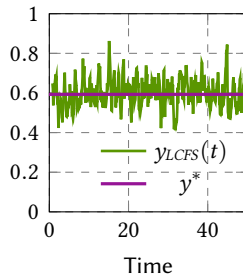
Earliest-Deadline-First



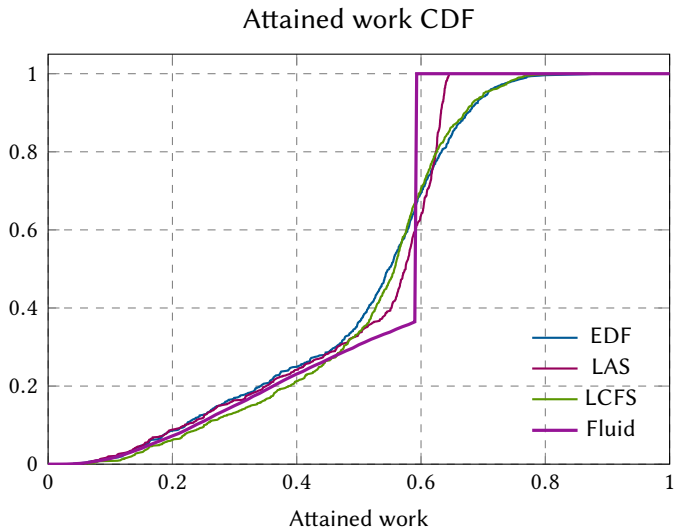
Least-Attained-Service



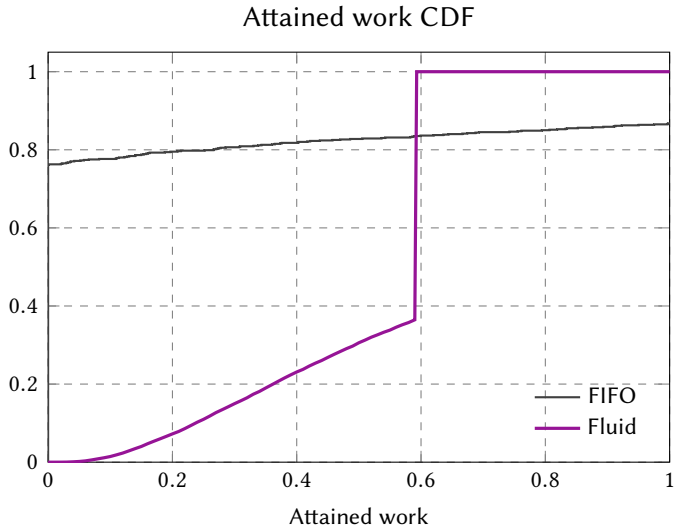
Last-Come-First-Served



# Attained work empirical CDF



# Comparison with FIFO



# Outline

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**Final remarks**



# Messages from the talk

- Measure-valued processes are a powerful tool to model general service queues.
- Partial service queues require two-dimensional measures.
- Our proposed dynamics for fluid models are tractable and approximate the real system.
- Last-but-not-least: in this setting, **deadline-oblivious** policies can be used without performance penalty!

- Analyze further policies using these tools (FCFS is easy for instance).
- Establish process-level convergence to the fluid models (long work...help needed...)
- Devise new policies and/or analyze different settings:
  - Tasks stay until service completion, but we want to measure the average *tardiness*, i.e. how late they depart.

# Gracias!

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<https://aferragu.github.io>

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