

4/1/23

Lab -7 - Proof by Resolution

Aim: To create a knowledge base using propositional logic and prove the given query using resolution.

Algorithm:

function PL-Resolution(KB, a) returns true or false
inputs: KB, the knowledge base, a sentence in propositional logic a;
the query, a sentence in propositional logic

Clause₁ - the set of clauses in the CNF representation of KB A - a new { }

loop do

for each pair of clauses C_i, C_j in Clause₁ do
 resolvents - PL-RESOLVE(C_i, C_j)

 if resolvents contains the empty clause
 then return true

 new - new \cup resolvents.

if new \subset Clause₁ then return false

Clause₁ - Clause₁ \cup new.

Code:

code:

import re

def negate(term):

return f'~{term}' if term[0] != '~' else
term[1]

def resolve(clause):

if len(clause) > 2:

t = split_term(clause)

return f'~{t[1]} v {t[0]}'

return ''

def split_term(term):

exp = '(~* [PQRS])'

term = re.findall(exp, term)

return term

def contradiction(query, clause):

contradictions = [f'~{query} v {negate(query)}',
f'~{negate(query)} v {query}']

return clause in contradictions or resolve(
clause) in contradictions

def resolve(kb, query):

temp = kb.get()

temp += [negate(query)]

steps = dict()

for rule in temp:

steps[rule] = 'given'

steps[negate(query)] = 'negated conclusion'

i = 0

while i < len(temp):

n = len(temp)

in while
↓

$J = (i+1) \cdot n$

$class = []$

while $j \neq i$:

term1 = split_term(comp[i])

term2 = split_term(comp[j])

for (in term1:

if negate() in term2:

$t_1 = [t \text{ for } t \text{ in term}_1$

if $t_1 = C$]

$t_2 = [t \text{ for } t \text{ in term}_2$

if $t_1 = \text{negate}(t_2)$]

$gen = t_1 + t_2$

if $\text{len}(gen) = -2$:

if $gen[0] = \text{negate}(gen[1])$

$class += [f' \{ gen[0]$

$\} \vee \{ gen[1] \}']$

else:

if contradiction(query,

$f' \{ gen[0] \} \vee \{ gen[1] \}']$:

temp.append (

$f' \{ gen[0] \} \vee \{ gen[1] \}']$)

steps[''] = f'

Resolved { temp[i] } and { temp[j] } to { temp[i-1] }

which is in turn null. \ \n A contradiction is found ~~then~~ when { negate(query) } is assumed as true. Hence, { query } is true. "

return steps

elif $\text{len}(gen) = -1$:

$class += [f' \{ gen[0] \}']$

else:

if contradiction(query, f'

$\{ term1[0] \} \vee \{ term2[0] \}']$:

temp.append (f' { term1

$\} \vee \{ term2[0] \}']$)

under if condition
↓

steps [''] = f " resolved
{tmp[i]} and {tmp[j]} to {tmp[-1]}, which
is in turn null. | WA contradiction is found when
{negate(query)} is assumed as true. Hence, {query}
is true".

under main while

return steps

for clause in clauses:

if clause not in temp and clause !=
reverse(clause) not in temp:

temp.append(clause)

steps[clause] = f " resolved from

{tmp[i]} and {tmp[j]}."

under main while

j = (j+1) % n.

def

i += 1

return steps

def resolution(kb, query):

kb = kb.split(' ')

steps = resolve(kb, query)

print('In step |t| clause |t| Derivation |t|')

print('!' * 30)

i = 1

for step in steps:

print(f' |i| |t| {step} |t| {steps[step]}
 |t|')

i += 1

def main():

print("Enter the kb: ")

kb = input()

print("Enter the query: ")

query = input()

resolution(kb, query)

main()

answer: Prove the KB:

$$\sim P \vee \sim Q \vee R \quad P \vee \sim P \vee Q \quad S$$

Prove the query:

R.

Step	Clause	Derivation
1.	$\sim P \vee \sim Q \vee R$	Given.
	P	Given.
	$\sim S \vee \sim T \vee Q$	Given.
	S	Given.
	$\sim R$	Assumed conclusion.
	$\sim Q \vee R$	Resolved from $\sim P \vee \sim Q \vee R$ and P .
	$\sim P \vee \sim Q$	Resolved from $\sim P \vee \sim Q \vee R$ and R .
	$\sim Q$	Resolved from P and $\sim P \vee \sim Q$.
	Q	Resolved from $\sim S \vee \sim T \vee Q$ and S .
	$\sim S \vee K$	Resolved from $\sim S \vee \sim T \vee Q$ and $\sim P$.
	$\sim S$	Resolved from $\sim S \vee \sim T \vee Q$ and $\sim P$.
	K	Resolved from S and $\sim S \vee K$.
	$\sim P$	Resolved from S and $\sim S \vee \sim P$.
	$\sim R$	Resolved from $\sim R$ and R to $R \vee \sim R$, which is a tautology.

A contradiction is found when $\sim R$ is assumed as true. Hence R is true.