

Logic Programming

Outline

- Pemrograman logika
- FOL
- Logika Predikat
- Horn clause
- Prolog

What is Logic Programming (LP)?

“Say what you want, not how you want it done”

- Programming paradigm which has its foundations in mathematical logic.
- Expresses the logic of a computation without describing its control flow → declarative
- Composed of a set of axioms, or rules which define relationships between objects, not sequences of instructions.

First Order Logic (FOL)

FOL : Symbolized reasoning in which each sentence, or statement, is broken down into a subject and a predicate (relation).

Terms

- **a constant**: single individual or concept
- **a variable** that stands for different individuals

Predicates/Relations

- **a relation** that maps n terms to a truth value true (T) or false (F).
- **a function**: a mapping that maps n terms to a term

Quantifiers

- Variables are used in conjunction with quantifiers
- “there exist”(\exists) and “for all”(\forall)

Connectives

- $\neg, \rightarrow, \leftrightarrow, \wedge, \vee$

Sentence in FOL

- An atomic sentence is simply a predicate applied to a set of terms.
 - $\text{Owns}(\text{John}, \text{Car1})$
- The standard propositional connectives ($\vee \neg \wedge \Rightarrow \Leftrightarrow$) can be used to construct complex sentences:
 - $\text{Owns}(\text{John}, \text{Car1}) \vee \text{Owns}(\text{Fred}, \text{Car1})$

- Universal quantifier: $\forall x$

Asserts that a sentence is **true for all** values of variable **x**

$\forall x [\text{Crow}(x) \rightarrow \text{Black}(x)]$ All crows are black

- Existential quantifier: \exists

Asserts that a sentence is **true for at least one** value of a variable **x**.

$\exists x [\text{Squirrel}(x) \wedge \text{Black}(x)]$ Some squirrels are black

Common Mistakes to Avoid

$$\forall x \text{ human}(x) \Rightarrow \text{mammal}(x)$$

- It's a big AND: Equivalent to the **conjunction** of all the instantiations of variable x :

$$\begin{aligned} &(\text{human}(\text{Jerry}) \Rightarrow \text{mammal}(\text{Jerry})) \wedge \\ &(\text{human}(\text{Jing}) \Rightarrow \text{mammal}(\text{Jing})) \wedge \\ &(\text{human}(\text{laptop}) \Rightarrow \text{mammal}(\text{laptop})) \wedge \dots \end{aligned}$$

- Common **mistake** is to use \wedge as main connective

$$\forall x \text{ human}(x) \wedge \text{mammal}(x)$$

- This means everything is human and a mammal!

$$\begin{aligned} &(\text{human}(\text{Jerry}) \wedge \text{mammal}(\text{Jerry})) \wedge \\ &(\text{human}(\text{Jing}) \wedge \text{mammal}(\text{Jing})) \wedge \\ &(\text{human}(\text{laptop}) \wedge \text{mammal}(\text{laptop})) \wedge \dots \end{aligned}$$

Common Mistakes to Avoid

$$\exists x \text{ human}(x) \wedge \text{male}(x)$$

- It's a big OR: Equivalent to the **disjunction** of all the instantiations of variable x :

$$(\text{human}(\text{Jerry}) \wedge \text{male}(\text{Jerry})) \vee$$

$$(\text{human}(\text{Jing}) \wedge \text{male}(\text{Jing})) \vee$$

$$(\text{human}(\text{laptop}) \wedge \text{male}(\text{laptop})) \vee \dots$$

- Common **mistake** is to use \Rightarrow as main connective
 - “Some pig can fly”

$$\exists x \text{ pig}(x) \Rightarrow \text{fly}(x) \text{ (wrong)}$$

- This is true if there is something not a pig!

$$(\text{pig}(\text{Jerry}) \Rightarrow \text{fly}(\text{Jerry})) \vee$$

$$(\text{pig}(\text{laptop}) \Rightarrow \text{fly}(\text{laptop})) \vee \dots$$

Nesting Quantifiers

- The order of quantifiers of the same type doesn't matter

$\forall x \forall y (\text{Parent}(x,y) \wedge \text{Male}(y) \Rightarrow \text{Son}(y,x))$

$\exists x \exists y (\text{Loves}(x,y) \wedge \text{Loves}(y,x))$

- The order of mixed quantifiers does matter:

$\forall x \exists y (\text{Loves}(x,y)) \rightarrow$ everybody loves somebody,

$\exists y \forall x (\text{Loves}(x,y)) \rightarrow$ there is someone who is loved by everyone

$\forall y \exists x (\text{Loves}(x,y)) \rightarrow$ everyone has someone who loves them.

$\exists x \forall y (\text{Loves}(x,y)) \rightarrow$ there is someone who loves everyone in the universe

Relation Between Quantifiers

- General Identities

$$\forall x \neg P \Leftrightarrow \neg \exists x P$$

$$\neg \forall x P \Leftrightarrow \exists x \neg P$$

$$\forall x P \Leftrightarrow \neg \exists x \neg P$$

$$\exists x P \Leftrightarrow \neg \forall x \neg P$$

Equality

- Identity relation
- Mary owns two cats. Inequality needed to insure x and y are distinct.
- $\exists x \exists y (\text{Owns}(\text{Mary}, x) \wedge \text{Cat}(x) \wedge \text{Owns}(\text{Mary}, y) \wedge \text{Cat}(y) \wedge \neg(x=y))$

Translate Sentence to FOL

"All *Ps* are *Qs*."

$\forall x. (P(x) \rightarrow Q(x))$

"Some *Ps* are *Qs*."

$\exists x. (P(x) \wedge Q(x))$

"No *Ps* are *Qs*."

$\forall x. (P(x) \rightarrow \neg Q(x))$

"Some *Ps* aren't *Qs*."

$\exists x. (P(x) \wedge \neg Q(x))$

Identifies objects, variables, predicates, function,
in the sentence

Translate Sentence to FOL

Every orange cat is fluffy.

Variable : 1 variable

Predicates : orange, cat, fluffy

Quantifier : \forall

"All Ps are Qs."

$\forall x. (P(x) \rightarrow Q(x))$

$\forall x. (x \text{ is an orange cat} \rightarrow x \text{ is fluffy})$

$\forall x. (x \text{ is an orange cat} \rightarrow x \text{ is fluffy})$

$\forall x. (x \text{ is an orange cat} \rightarrow \text{Fluffy}(x))$

$\forall x. (\text{Orange}(Cat(x)) \rightarrow \text{Fluffy}(x))$

$\forall x. (x \text{ is orange and } x \text{ is a cat} \rightarrow \text{Fluffy}(x))$

$\forall x. (\text{Orange}(x) \wedge \text{Cat}(x) \rightarrow \text{Fluffy}(x))$

Translate Sentence to FOL

There's a corgi that loves everyone.

Variable : 2 variables

Predicates : corgi, person, loves

Quantifier : \exists , \forall

"Some Ps are Qs."

$\exists x. (P(x) \wedge Q(x))$

$\exists x. (x \text{ is a corgi} \wedge x \text{ loves everyone})$

$\exists x. (Corgi(x) \wedge x \text{ loves everyone})$

$\exists x. (Corgi(x) \wedge x \text{ loves every person } y)$

$\exists x. (Corgi(x) \wedge \text{every person } y \text{ is loved by } x)$

"All Ps are Qs."

$\forall x. (P(x) \rightarrow Q(x))$

$\exists x. (Corgi(x) \wedge \forall y. (y \text{ is a person} \rightarrow y \text{ is loved by } x))$

$\exists x. (Corgi(x) \wedge \forall y. (Person(y) \rightarrow Loves(x, y)))$

Translate Sentence to FOL

Everybody loves at least one corgi.

"All Ps are Qs."

$\forall x. (P(x) \rightarrow Q(x))$

Variable : 2 variables

Predicates : corgi, person, loves

Quantifier : \exists, \forall

$\forall x. (x \text{ is a person} \rightarrow x \text{ loves at least one corgi})$

$\forall x. (Person(x) \rightarrow x \text{ loves at least one corgi})$

$\forall x. (Person(x) \rightarrow x \text{ loves at least one corgi } y)$

$\forall x. (Person(x) \rightarrow \text{there is a corgi } y \text{ that is loved by } x)$

"Some Ps are Qs."

$\exists x. (P(x) \wedge Q(x))$

$\forall x. (Person(x) \rightarrow \exists y. (y \text{ is a corgi} \wedge y \text{ is loved by } x))$

$\forall x. (Person(x) \rightarrow \exists y. (Corgi(y) \wedge Loves(x, y)))$

Horn Clause

- Perhatikan pernyataan: **if** $(P1 \wedge P2 \wedge \dots \wedge Pn)$ **then** Q .
- Dapat juga ditulis sebagai: $Q \leftarrow (P1 \wedge P2 \wedge \dots \wedge Pn)$
dibaca Q hanya jika $P1$ dan $P2$ dan ... dan Pn
- Pernyataan Q akan benar (True) jika semua pernyataan $P1, P2, \dots, Pn$ secara simultan benar.
- Ingat: $A \rightarrow B$ setara (memiliki nilai kebenaran yang sama) dengan : $\neg A \vee B$.
- Oleh karena itu, pernyataan implikasi tadi dapat dinyatakan dalam bentuk *disjunctive normal* sebagai:
 $Q \vee \neg P1 \vee \neg P2 \dots \vee \neg Pn$
- Ekspresi ini disebut Horn Clause

Horn Clause

- Prolog:

```
c:- a, b.  
a.  
b.
```

- Horn formula:

$$[c \vee \neg a \vee \neg b] \wedge a \wedge b$$
$$[c, \neg a, \neg b] \quad [a] \quad [b]$$

Proving Horn Clause

- Prolog is based on computing with Horn clauses
- Let us attempt to prove c by contradiction
- Hence, the goal clause is $\neg c$

$[c, \neg a, \neg b] \quad [a] \quad [b] \quad [\neg c]$

Resolution Inference Rule

$(A \vee B \vee C)$

$(\neg A)$

$\therefore (B \vee C)$

$[c, \neg a, \neg b] \quad [a] \quad [b] \quad [\neg c]$

$[c, \neg b] \quad [b] \quad [\neg c]$

$[c] \quad [\neg c]$

\square

Prolog

- Conversion FOL to Prolog
 - $(\forall x) (\forall y) (\forall z) (\text{father}(x,z) \wedge \text{parent}(z,y) \rightarrow \text{grandfather}(x,y))$
 $\text{grandfather}(X,Y) \text{ :- father}(X,Z), \text{parent}(Z,Y)$
- **Atoms:** string karakter yang dimulai dengan huruf kecil, menyatakan nama dari obyek atau relasi. Misal: joko, ibu, lelaki
- **Term:**
 - **Constants:** atoms
 - **Variables:** string karakter yang dimulai dengan huruf besar. Misal: X
 - **Structure:** predikat dengan nama dan sejumlah argumen yang telah *fixed*. Misal: ibu(shinta,X).

Rules & Facts

- Technically, we don't create a Prolog program; we create a Prolog database.
- In the database are two types of clauses: Facts and Rules.
- **Facts** (fakta): single piece of information
 - `blue(sky).`
 - `mammal(rabbit).`
 - `plays(john, hockey).`
- **Rules** (aturan): logic expression "hanya jika", Horn clauses.
 - generate new information from facts, other rules, and even themselves
 - Rules have the form **head :- body.**
 - the head and the body are clauses that typically use variables instead of constants (because rules are used to define general relationships).
 - Contoh:
 - `habisDibagiDua(X) :- genap(X).`
 - Pernyataan tersebut setara dengan pernyataan kalkulus predikat:
$$(\forall x) (\text{genap}(x) \rightarrow \text{habisDibagiDua}(x))$$

Conjunction & Disjunction

- *Conjunction* dari predikat direpresentasikan sebagai deretan struktur yang dipisahkan oleh tanda koma (,).
- *Disjunction* dari predikat direpresentasikan sebagai deretan struktur yang dipisahkan oleh tanda titik koma (;).
 - sibling(X,Y):- kakak(X,Y) ; adik(X,Y).
 - sibling(X,Y):- kakak(X,Y).
 - sibling(X,Y):- adik(X,Y).
- *Negation* predikat untuk negasi adalah not.
 - ganjil(X):- not genap(X).

How Prolog Works

- A Prolog program is initiated by a **query** - a predicate or a sequence of predicates to be proved.
- The predicate to be proved is called a **goal**.
- Prolog tries to **match** the goal to the head of some fact or some rule.
- If a match is found with a rule, Prolog continues with proving each predicate in the body of the rule.

Example

```
man(adam) .  
man(peter) .  
man(paul) .  
woman(marry) .  
woman(eve) .  
parent(adam,peter) .    % means adam is parent of peter  
parent(eve,peter) .  
parent(adam,paul) .  
parent(marry,paul) .
```

Facts

```
father(F,C):-man(F),parent(F,C) .  
mother(M,C):-woman(M),parent(M,C) .  
is_father(F):-father(F,_) .  
is_mother(M):-mother(M,_) .
```

Rules

```
?-father(X,paul) .
```

Query

```
X=Adam
```

Example 2

X grandmother of Z

mother(pam,bob).

mother(pat,jim).

father(tom,bob).

father(tom,liz).

father(bob, ann).

father(bob,pat).

grandmother(X,Z) :- mother(X,Y), mother(Y,Z).

grandmother(X,Z) :- mother(X,Y), father(Y,Z).

grandfather(X,Z) :- father(X,Y), mother(Y,Z).

grandfather(X,Z):-father(X,Y), father(Y,Z)

Example 3

X professors Y

```
studies(charlie, csc135).  
studies(olivia, csc135).  
studies(jack, csc131).  
studies(arthur, csc134).
```

```
teaches(kirke, csc135).  
teaches(collins, csc131).  
teaches(collins, csc171).  
teaches(juniper, csc134).
```

```
professor(X, Y) :- teaches(X, C), studies(Y, C).
```


Exercise 1

- Any two pancakes x and y taste similar
- Everyone knows at least two people

Exercise 2

- Buatlah klausa Prolog untuk persoalan berikut:
 1. Hewan buas adalah hewan yang berwarna gelap, berbadan besar, dan gigi bertaring.
 2. Hewan jinak adalah hewan yang berwarna terang, gigi tak bertaring dan berbadan kecil.
 3. Faktorial dari 0 adalah 1.
 4. Faktorial dari suatu bilangan bulat positif.

```
factorial(0) => 1  
factorial(n) => factorial(n-1) * n;
```

5. Fungsi f didefinisikan sbb:

$$f(x, y) = \begin{cases} x & , x > y \\ y & , \text{selainnya} \end{cases}$$