Cloud and Cluster Data Management



THANKS TO M. GROSSNIKLAUS

Facebook: Motivation for Hive

Growth of the Facebook data warehouse

2007: 15TB of net data

2010: 700TB of net data

– 2011: >30PB of net data

– 2012: >100PB of net data

- Scalable data analysis used across the company
 - ad hoc analysis
 - business intelligence
 - Insights for the Facebook Ad Network
 - analytics for page owners
 - system monitoring

Motivation for Hive (continued)

- Original Facebook data processing infrastructure
 - built using a commercial RDBMS prior to 2008
 - became inadequate as daily data processing jobs took longer than a day
- Hadoop was selected as a replacement
 - pros: petabyte scale and use of commodity hardware
 - cons: using it was not easy for end user not familiar with map-reduce
 - "Hadoop lacked the expressiveness of [..] query languages like SQL and users ended up spending hours (if not days) to write programs for even simple analysis."

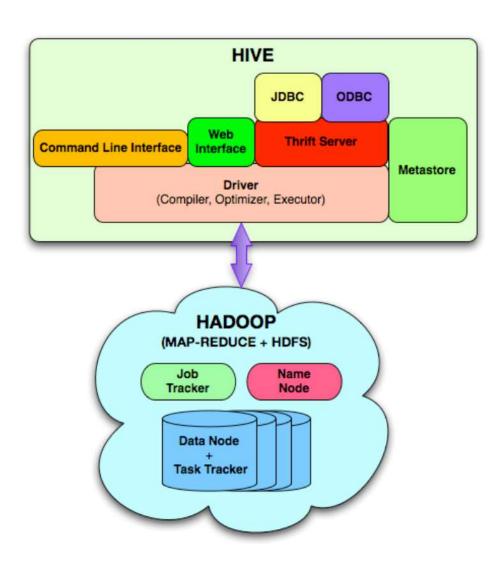
Motivation for Hive (continued)

- Hive is intended to address this problem by bridging the gap between RDBMS and Hadoop
 - "Our vision was to bring the familiar concepts of tables, columns, partitions and a subset of SQL to the unstructured world of Hadoop"

• Hive provides:

- tools to enable easy data extract/transform/load (ETL)
- a mechanism to impose structure on a variety of data formats
- access to files stored either directly in HDFS or in other data storage systems such as Hbase, Cassandra, MongoDB, and Google Spreadsheets
- a simple SQL-like query language
- query execution via MapReduce

Hive Architecture



- Clients use command line interface (CLI), Web UI, or JDBC/ODBC driver
- HiveServer provides Thrift and JDBC/ODBC interfaces
- Metastore stores system catalogue and metadata about tables, columns, partitions etc.
- Driver manages lifecycle of HiveQL statement as it moves through Hive

Data Model

- Schemas are required in Hive
- Hive structures data into well-understood database concepts like tables, columns, rows, and partitions
- Primitive types
 - Integers: bigint (8 bytes), int (4 bytes), smallint (2 bytes), tinyint (1 byte)
 - Floating point: float (single precision), double (double precision)
 - String

Complex Types

- Complex types
 - Associative arrays: map<key-type, value-type>
 - Lists: list<element-type>
 - Structs: struct<field-name: field-type, ...>
- Complex types are templated and can be composed to create types of arbitrary complexity
 - li list<map<string, struct<p1:int, p2:int>>>

Complex Types

- Complex types
 - Associative arrays: map<key-type, value-type>
 - Lists: list<element-type>
 - Structs: struct<field-name: field-type, ...>
- Accessors
 - Associative arrays: m['key']
 - Lists: li[0]
 - Structs: s.field-name
- Example:
 - li list<map<string, struct<p1:int, p2:int>>
 - t1.li[0] ['key'].p1 gives the p1 field of the struct associated with the key of the first array of the list li

Participation Question 1

- From the list of data models below, what the most similar data model to hive? Briefly explain your answer.
 - Relational
 - Key-value
 - Document
 - Column family
 - Graph

Query Language

- HiveQL is a subset of SQL plus some extensions
 - FROM clause sub-queries
 - various types of joins: inner, left outer, right outer and outer joins
 - Cartesian products
 - group by and aggregation
 - union all
 - create table as select

Limitations

- only equality joins
- joins need to be written using ANSI join syntax (not in WHERE clause)
- no support for inserts in existing table or data partition
- all inserts overwrite existing data

Query Language

- Hive supports user defined functions written in java
- Three types of UDFs
 - UDF: user defined function
 - Input: single row
 - Output: single row
 - UDAF: user defined aggregate function
 - Input: multiple rows
 - Output: single row
 - UDTF: user defined table function
 - Input: single row
 - Output: multiple rows (table)

Creating Tables

- Tables are created using the **CREATE TABLE** DDL statement
- Example:

```
CREATE TABLE t1(
    st string,
    fl float,
    li list<map<string, struct<p1:int, p2:int>>
);
```

- Tables may be partitioned or non-partitioned (we'll see more about this later)
- Partitioned tables are created using the PARTITIONED BY statement

```
CREATE TABLE test_part(c1 string, c2 string)
PARTITIONED BY (ds string, hr int);
```

Inserting Data

Example

```
INSERT OVERWRITE TABLE t2
SELECT t3.c2, COUNT(1)
FROM t3
WHERE t3.c1 <= 20
GROUP BY t3.c2;</pre>
```

- OVERWRITE (instead of INTO) keyword to make semantics of insert statement explicit
- Lack of INSERT INTO, UPDATE, and DELETE enable simple mechanisms to deal with reader and writer concurrency
- At Facebook, these restrictions have not been a problem
 - data is loaded into warehouse daily or hourly
 - each batch is loaded into a new partition of the table that corresponds to that day or hour

Inserting Data

- Hive supports inserting data into HDFS, local directories, or directly into partitions (more on that later)
- Inserting into HDFS

```
INSERT OVERWRITE DIRECTORY '/output_dir'
SELECT t3.c2, AVG(t3.c1)
FROM t3
WHERE t3.c1 > 20 AND t3.c1 <= 30
GROUP BY t3.c2;</pre>
```

Inserting into local directory

```
INSERT OVERWRITE LOCAL DIRECTORY '/home/dir'
SELECT t3.c2, SUM(t3.c1)
FROM t3
WHERE t3.c1 > 30
GROUP BY t3.c2;
```

Inserting Data

- Hive supports inserting data into multiple tables/files from a single source given multiple transformations
- Example (corrected from paper):

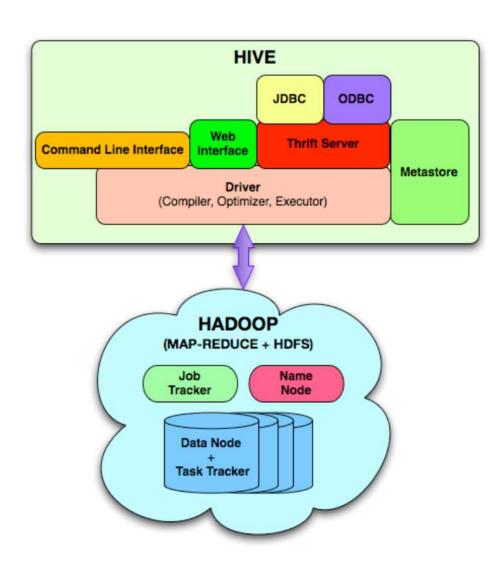
```
FROM +1
    INSERT OVERWRITE TABLE t2
    SELECT t1.c2, count(1)
    WHERE t1.c1 <= 20
    GROUP BY t1.c2;
    INSERT OVERWRITE DIRECTORY '/output dir'
    SELECT t1.c2, AVG(t1.c1)
    WHERE t1.c1 > 20 AND t1.c1 <= 30
    GROUP BY t1.c2;
    INSERT OVERWRITE LOCAL DIRECTORY '/home/dir'
    SELECT t1.c2, SUM(t1.c1)
    WHERE t1.c1 > 30
    GROUP BY t1.c2;
```

We Gotta Have Map-Reduce!

- HiveQL has extensions to express map-reduce programs
- Example

- MAP clause indicates how the input columns are transformed by the mapper UDF (and supplies schema)
- CLUSTER BY clause specifies output columns that are hashed and distributed to reducers
- REDUCE clause specifies the UDF to be used by the reducers

Hive Architecture



- Clients use command line interface, Web UI, or JDBC/ODBC driver
- HiveServer provides Thrift and JDBC/ODBC interfaces
- Metastore stores system catalogue and metadata about tables, columns, partitions etc.
- Driver manages lifecycle of HiveQL statement as it moves through Hive

Metastore

- Stores system catalog and metadata about tables, columns, partitions, etc.
- Uses a traditional RDBMS "as this information needs to be served fast"
- Backed up regularly (since everything depends on this)
- Needs to be able to scale with the number of submitted queries (we don't want thousands of Hadoop workers hitting this DB for every task)
- Only Query Compiler talks to Metastore (metadata is then sent to Hadoop workers in XML plans at runtime)

Data Storage

- Table metadata associates data in a table to HDFS directories
 - tables: represented by a top-level directory in HDFS
 - table partitions: stored as a sub-directory of the table directory
 - buckets: stores the actual data and resides in the sub-directory that corresponds to the bucket's partition, or in the top-level directory if there are no partitions
- Tables are stored under the Hive root directory

```
CREATE TABLE test_table (...);
```

Creates a directory like

 <warehouse_root_directory>/test_table
 where <warehouse_root_directory> is determined by the Hive configuration

Partitions

 Partitioned tables are created using the PARTITIONED BY clause in the CREATE TABLE statement

```
CREATE TABLE test_part(c1 string, c2 int)
PARTITIONED BY (ds string, hr int);
```

- Note that partitioning columns are not part of the table data
- New partitions can be created through an INSERT statement or an ALTER statement that adds a partition to a table

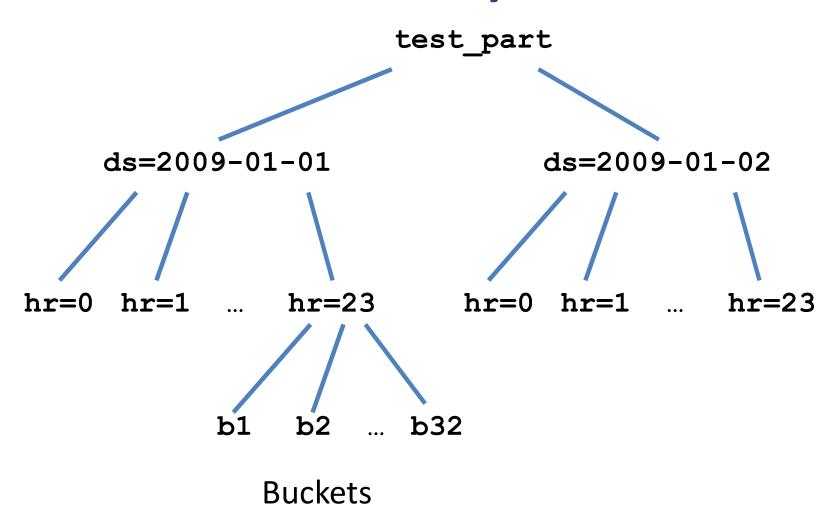
Partition Example

```
INSERT OVERWRITE TABLE test_part
PARTITION(ds='2009-01-01', hr=12)
SELECT * FROM t;
ALTER TABLE test_part
ADD PARTITION(ds='2009-02-02', hr=11);
```

- Each of these statements creates a new directory
 - /.../test_part/ds=2009-01-01/hr=12
 - /.../test_part/ds=2009-02-02/hr=11
- HiveQL compiler uses this information to prune directories that need to be scanned to evaluate a query

```
SELECT * FROM test_part WHERE ds='2009-01-01';
SELECT * FROM test_part
WHERE ds='2009-02-02' AND hr=11;
```

Directory Structure



Participation Question 2

How is partitioning in Hive different from indexing in a DBMS?

Buckets

- A bucket is a file in the leaf level directory of a table or partition
- Users specify number of buckets and column on which to bucket data using the CLUSTERED BY clause

```
CREATE TABLE test_part(c1 string, c2 int)
PARTITIONED BY (ds string, hr int)
CLUSTERED BY (c1) INTO 32 BUCKETS;
```

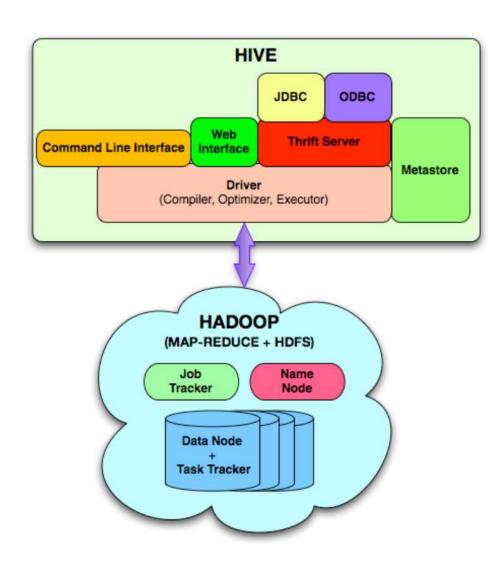
Buckets

- Bucket information is then used to prune data in the case the user runs queries on a sample of data
- Example:

```
SELECT * FROM test_part TABLESAMPLE (2 OUT OF 32);
```

 This query will only use 1/32 of the data as a sample from the second bucket in each partition

Hive Architecture



- Clients use command line interface, Web UI, or JDBC/ODBC driver
- HiveServer provides Thrift and JDBC/ODBC interfaces
- Metastore stores system catalogue and metadata about tables, columns, partitions etc.
- Driver manages lifecycle of HiveQL statement as it moves through Hive

Query Compiler

- Parses HiveQL using Antlr to generate an abstract syntax tree
- Type checks and performs semantic analysis based on Metastore information
- Naïve rule-based optimizations
- Compiles HiveQL into a directed acyclic graph of MapReduce tasks

Optimizations

Column Pruning

 Ensures that only columns needed in query expressions are deserialized and used by the execution plan

Predicate Pushdown

Filters out rows in the first scan if possible

Partition Pruning

Ensures that only partitions needed by the query plan are used

Optimizations

Map side joins

- If one table in a join is very small it can be replicated in all of the mappers and joined with other tables
- User must know ahead of time which are the small tables and provide hints to Hive

```
SELECT /*+ MAPJOIN(t2) */ t1.c1, t2.c1
FROM t1 JOIN t2 ON(t1.c2 = t2.c2);
```

Join reordering

 Smaller tables are kept in memory and larger tables are streamed in reducers ensuring that the join does not exceed memory limits

Optimizations

- GROUP BY repartitioning
 - If data is skewed in GROUP BY columns the user can specify hints as with MAPJOIN

```
set hive.groupby.skewindata=true;
SELECT t1.c1, sum(t1.c2)
FROM t1
GROUP BY t1;
```

- Hashed based partial aggregations in mappers
 - Hive enables users to control the amount of memory used on mappers to hold rows in a hash table
 - As soon as that amount of memory is used, partial aggregates are sent to reducers.

Execution Engine

- MapReduce tasks are executed in the order of their dependencies
- Independent tasks can be executed in parallel

Hive Usage at Facebook

- Data processing task
 - more than 50% of the workload are ad-hoc queries
 - remaining workload produces data for reporting dashboards
 - range from simple summarization tasks to generate rollups and cubes to more advanced machine learning algorithms
- Hive is used by novice and expert users
- Types of Applications:
 - Summarization
 - Eg: Daily/Weekly aggregations of impression/click counts
 - Ad hoc Analysis
 - Eg: how many group admins broken down by state/country
 - Data Mining (Assembling training data)
 - Eg: User Engagement as a function of user attributes

Hive Usage Elsewhere

CNET

Hive used for data mining, internal log analysis and ad hoc queries

eHarmony

Hive used as a source for reporting/analytics and machine learning

Grooveshark

Hive used for user analytics, dataset cleaning, and machine learning
 R&D

Last.fm

Hive used for various ad hoc queries

Scribd

- Hive used for machine learning, data mining, ad-hoc querying, and both internal and user-facing analytics
- Also: Netflix, Tata Consultancy, Hortonworks, Quoble

References

- C. Olston, B. Reed, U. Srivastava, R. Kumar, and A. Tomkins: Pig Latin: A Not-So-Foreign Language for Data Processing. Proc. Intl. Conf. on Management of Data (SIGMOD), pp. 1099-1110, 2008.
- A. Thusoo, J. S. Sarma, N. Jain, Z. Shao, P. Chakka, N. Zhang, S. Anthony, H. Liu, R. Murthy: Hive A Petabyte Scale Data Warehouse Using Hadoop. Proc. Intl. Conf. on Data Engineering (ICDE), pp. 996-1005, 2010.