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FACULTY OF INFORMATION ENGINEERING,
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Mathematical Logic for Computer Science

Lecture notes integrated with the book TODO

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Contents

Information and Contacts	1
1 Homeworks	2
1.1 Homework 1	2
1.2 Homework 2	6

Information and Contacts

Personal notes and summaries collected as part of the *Mathematical Logic for Computer Science* course offered by the degree in Computer Science of the University of Rome "La Sapienza".

Further information and notes can be found at the following link:

<https://github.com/aflaag-notes>. Anyone can feel free to report inaccuracies, improvements or requests through the Issue system provided by GitHub itself or by contacting the author privately:

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Suggested prerequisites:

TODO

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1

Homeworks

1.1 Homework 1

Exercise 1.4 Let $\mathcal{L} = \{E(x, y)\}$ be the language of graphs.

1. For each fixed $n \in \mathbb{N}$, write a sentence C_n such that for any graph \mathcal{G} , $\mathcal{G} \models C_n$ if and only if \mathcal{G} contains a cycle of length n .
2. Prove using Compactness that the property of being a *cycle* is not expressible by a theory in \mathcal{L} over the class of graphs.

Solution. Let $\mathcal{L} = \{E(x, y)\}$ be the language of graphs.

1. The property “ \mathcal{G} contains a cycle of length n ” can be written as follows

$$C_n := \exists x_1 \dots \exists x_n \left(\bigwedge_{\substack{1 \leq i, j \leq n \\ i \neq j}} \neg(x_i = x_j) \right) \wedge \left(\bigwedge_{1 \leq i \leq n-1} E(x_i, x_{i+1}) \wedge E(x_n, x_1) \right)$$

In fact, the first conjunction implies that x_1, \dots, x_n are *distinct*, and the second conjunction describes the existence of the n -long *cycle* itself.

2. Consider the property $P_n :=$ “ \mathcal{G} is a *cycle* of length n ”. This property can be expressed by *extending* C_n as follows:

$$\begin{aligned} V_n &:= \forall y \bigvee_{1 \leq j \leq n} (y = x_j) \\ E_n &:= \bigwedge_{1 \leq i \leq n-1} \bigwedge_{\substack{1 \leq j \leq n: \\ j \neq i+1}} \neg E(x_i, x_j) \wedge \bigwedge_{2 \leq j \leq n} \neg E(x_n, x_j) \\ C'_n &:= \exists x_1 \dots \exists x_n \quad C_n \wedge V_n \wedge E_n \end{aligned}$$

where we have that

- V_n ensures that \mathcal{G} has *exactly* n vertices
- E_n ensures that the only edges present in \mathcal{G} are the ones that describe the cycle graph of n vertices
- C'_n describes our property P_n

Now, consider the property $P :=$ “ \mathcal{G} is a *cycle*”, and in particular $\neg P :=$ “ \mathcal{G} is not a *cycle*”. We observe that we can build the following infinite theory

$$T^{\neg P} := \{\neg C'_n \mid n \in \mathbb{N}_{\geq 3}\}$$

for which it is easy to see that

$$\mathcal{G} \models \neg P \iff \neg P(\mathcal{G}) \text{ holds}$$

meaning that $\neg P$ is expressible through $T^{\neg P}$.

Claim: $T^{\neg P} \in \text{FINSAT}$.

Proof of the Claim. Fix $T_0 \subseteq_{fin} T^{\neg P}$. We observe that $T_0 := \{\neg C'_{i_1}, \dots, \neg C'_{i_k}\}$

for some $i_1, \dots, i_k \in \mathbb{N}$. Now, if we consider $i^* := \max_{j \in [k]} i_j$, then the cycle graph that has $i^* + 1$ vertices is clearly a structure that satisfies T_0 . \square

Claim: P is not expressible by a theory in \mathcal{L} over the class of graphs.

Proof of the Claim. By way of contradiction, suppose that P is expressible, i.e. there is a theory T^P for which P can be expressed. Then, consider the theory $T := T^P \cup T^{\neg P}$. By the previous claim, we have that $T \in \text{FINSAT}$, and by Compactness this is true if and only if $T \in \text{SAT}$. However, this is a contradiction, because a graph cannot be and not be a cycle at the same time. \square

Finally, this last claim concludes the proof. \square

Exercise 2.1 Consider the following two structures \mathcal{G}_1 and \mathcal{G}_2 for the languages of graphs:



Write at least two sentences distinguishing the two structures. Discuss the EF-game played on these structures: for what k can the Duplicator win the k -rounds game? For what k can the Spoiler win?

Solution. Some properties that can distinguish these two structures are the following:

1. “ \mathcal{G} contains a vertex of degree 3”, which is represented by the following sentence of rank 5

$$\exists x_1 \exists x_2 \exists x_3 \exists x_4 \left(\bigwedge_{\substack{1 \leq i, j \leq 4 \\ i \neq j}} \neg(x_i = x_j) \right) \wedge \left(\bigwedge_{2 \leq i \leq 4} E(x_1, x_i) \right) \wedge \left(\forall y \quad \neg E(x_1, y) \vee \bigvee_{2 \leq j \leq 4} (y = x_j) \right)$$

2. “ \mathcal{G} contains edges as \mathcal{G}_1 ”, which is represented by the following sentence of rank 5

$$\begin{aligned} \exists x_1 \exists x_2 \exists x_3 \exists x_4 \exists x_5 \quad & E(x_1, x_2) \wedge E(x_1, x_3) \wedge \\ & E(x_2, x_3) \wedge E(x_2, x_4) \wedge E(x_2, x_5) \wedge \\ & E(x_3, x_4) \wedge E(x_4, x_5) \wedge \\ & E(x_4, x_5) \end{aligned}$$

we observe that the edges of \mathcal{G}_2 are not sufficient to distinguish the two sentences, because \mathcal{G}_2 is a subgraph of \mathcal{G}_1

3. “ \mathcal{G} contains a cycle of length 5”, which is represented by C_5 of the previous exercise, and has rank 5
4. “ \mathcal{G} contains a cycle of length 4”, which is represented by C_4 of the previous exercise, and has rank 4
5. “ \mathcal{G} contains K_4 as subgraph”, which is represented by the following sentence having rank 4

$$\exists x_1 \exists x_2 \exists x_3 \exists x_4 \left(\bigwedge_{\substack{1 \leq i, j \leq 4 \\ i \neq j}} \neg(x_i = x_j) \right) \wedge \left(\bigwedge_{\substack{1 \leq i, j \leq 4 \\ i \neq j}} E(x_i, x_j) \right)$$

These sentences *may seem* to suggest that the two structures are 3-equivalent, meaning that there is no sentence of rank 3 that can distinguish \mathcal{G}_1 from \mathcal{G}_2 . For now, let's focus on proving that they are *at least* 2-equivalent.

Claim: The Duplicator wins $G_2(\mathcal{G}_1, \mathcal{G}_2)$.

Proof of the Claim. Let s_i and d_i be the i -th nodes chosen by the Spoiler and the Duplicator, respectively. Then, we can define the following strategy for the Duplicator:

- if $s_1 \in \{1, 4, 5\}$, then the Duplicator chooses $d_1 \in \{a, b, d, e\}$, otherwise if $s_1 \in \{2, 3\}$ then $d_1 = c$
- similarly, if $s_1 \in \{a, b, d, e\}$, then the Duplicator chooses $d_1 \in \{1, 4, 5\}$, otherwise if $s_1 = c$ then $d_1 \in \{2, 3\}$

Then, no matter the choice of s_2 , the Duplicator can always answer with a node d_2 that preserves the partial isomorphism, in fact:

- if $s_2 \sim s_1$, it is guaranteed that there is a vertex d_2 in the other structure such that $d_2 \sim d_1$ because $\delta(\mathcal{G}_1) = \delta(\mathcal{G}_2) = 2$ — and the same argument applies if $s_2 \sim d_1$ for finding a vertex $d_2 \sim s_1$
- if $s_2 \not\sim s_1$, the strategy that we provided for the Duplicator guarantees that there exists at least one vertex d_2 in the other structure such that $d_2 \not\sim d_1$ — and the same argument applies if $s_2 \not\sim d_1$ for finding a vertex $d_2 \not\sim s_1$

Thus, the Duplicator has a strategy to always win at least 2 rounds, therefore the Duplicator wins $G_2(\mathcal{G}_1, \mathcal{G}_2)$ by Ehrenfeucht's theorem. \square

Now that we proved that $\mathcal{G}_1 \equiv_2 \mathcal{G}_2$, is it true that they are also 3-equivalent? Unfortunately, the following claim proves that this is indeed false.

Claim: The Spoiler wins $G_3(\mathcal{G}_1, \mathcal{G}_2)$.

Proof of the Claim. The following is a strategy that guarantees the Spoiler to win in 3 rounds:

- let $s_1 \in \{4, 5\}$

- by the previous claim, we know that the strategy for the Duplicator to win at least 2 rounds is to choose $d_1 \in \{a, b, d, e\}$, thus we may assume that $d_1 \neq c$
- now, let $s_2 = 1$
- to preserve the partial isomorphism, we observe that
 - if $d_1 \in \{a, b\}$, then $d_2 \in \{d, e\}$
 - if $d_1 \in \{d, e\}$, then $d_2 \in \{a, b\}$
- now, it suffices for the Spoiler to choose s_3 in \mathcal{G}_2 such that $s_3 \sim d_2$ and $s_3 \neq c$: by construction of \mathcal{G}_2 , we see that $s_3 \approx d_1$, but all the vertices in $\{2, 3, 5\}$ are adjacent to s_1 , which would violate the partial isomorphism

□

In fact, we can actually find a property that distinguishes \mathcal{G}_1 from \mathcal{G}_2 which can be written through a sentence of rank 3: “there are two vertices x_1 and x_2 of \mathcal{G} such that for each third vertex x_3 there is a K_3 as subgraph of \mathcal{G} such that $V(K_3) = \{x_1, x_2, x_3\}$ ”

$$\exists x_1 \exists x_2 \forall x_3 \left(\bigwedge_{\substack{1 \leq i, j \leq 3 \\ i \neq j}} \neg(x_i = x_j) \right) \wedge E(x_1, x_2) \wedge E(x_2, x_3) \wedge E(x_3, x_1)$$

Let x_1, x_2 and x_3 be the three chosen vertices — and we may assume that $x_1 \sim x_2$ otherwise the sentence is trivially unsatisfied. Then, we observe that

- in \mathcal{G}_1 if $\{x_1, x_2\} = \{2, 3\}$, then for any other vertex $x_3 \in \{1, 4, 5\}$ we can always find a K_3 having x_1, x_2 and x_3 as its vertices
- in \mathcal{G}_2 we have two cases
 - if $\{x_1, x_2\} \subseteq \{a, b, c\}$, the property is unsatisfied for $x_3 \in \{d, e\}$
 - if $\{x_1, x_2\} \subseteq \{c, d, e\}$, the property is unsatisfied for $x_3 \in \{a, b\}$

In conclusion, we have that $\mathcal{G}_1 \equiv_2 \mathcal{G}_2$, and that $\mathcal{G}_1 \not\equiv_3 \mathcal{G}_2$.

□

1.2 Homework 2

Exercise 1.1 Let (W, R) be a *quasi-order*; that is, assume that R is transitive and reflexive. Define the binary relation \sim on W by putting $s \sim t \iff R(s, t) \wedge R(t, s)$.

(a) Show that \sim is an equivalence relation.

Let $[s]$ denote the equivalence class of s under this relation, and define the following relation on the collection of equivalence classes: $[s] \leq [t] \iff R(s, t)$.

(b) Show that this relation is well-defined.

(c) Show that \leq is a partial order.

Solution. We prove the statements as follows.

(a) To prove that \sim is an equivalence relation, it suffices to show that \sim has the following properties:

- *reflexivity*: $\forall s \in W \quad R(s, s)$ by reflexivity of R , therefore $s \sim s$
- *symmetry*: $\forall s, t \in W \quad s \sim t \iff R(s, t) \wedge R(t, s) \iff t \sim s$
- *transitivity*: $\forall s, t, u \in W \quad \begin{cases} s \sim t \iff R(s, t) \wedge R(t, s) \\ t \sim u \iff R(t, u) \wedge R(u, t) \end{cases}$ and by transitivity of R we have that

$$- R(s, t) \wedge R(t, u) \implies R(s, u)$$

$$- R(u, t) \wedge R(t, s) \implies R(u, s)$$

$$\text{and by definition } R(s, u) \wedge R(u, s) \iff s \sim u$$

(b) To prove that \leq is well-defined, we need to show that

$$\forall s, t, s', t' \quad s \sim s' \wedge t \sim t' \implies ([s] \leq [t] \iff [s'] \leq [t'])$$

We observe that

- $s \sim s' \iff R(s, s') \wedge R(s', s)$
- $t \sim t' \iff R(t, t') \wedge R(t', t)$

therefore, we have that

- $[s] \leq [t] \iff R(s, t)$, and by transitivity of R it holds that $R(s', s) \wedge R(s, t) \implies R(s', t)$; therefore, by transitivity of R again we have that $R(s', t) \wedge R(t, t') \implies R(s', t') \iff [s'] \leq [t']$
- $[s'] \leq [t'] \iff R(s', t')$, and by transitivity of R it holds that $R(s', t') \wedge R(t', t) \implies R(s', t)$; therefore, by transitivity of R again we have that $R(s, s') \wedge R(s', t) \implies R(s, t) \iff [s] \leq [t]$

(c) To prove that \leq is a partial order, it suffices to show that \leq has the following properties:

- *reflexivity*: $\forall s \in W \quad R(s, s)$ by reflexivity of R , and $R(s, s) \iff [s] \leq [s]$

- *antisymmetry*: $\forall s, t \in W \quad \left\{ \begin{array}{l} [s] \leq [t] \iff R(s, t) \\ [t] \leq [s] \iff R(t, s) \end{array} \right. \implies R(s, t) \wedge R(t, s) \iff s \sim t \iff [s] = [t]$
- *transitivity*: $\forall s, t, u \in W \quad \left\{ \begin{array}{l} [s] \leq [t] \iff R(s, t) \\ [t] \leq [u] \iff R(t, u) \end{array} \right. \implies R(s, t) \wedge R(t, u) \implies R(s, u)$ by transitivity of R , and $R(s, u) \iff [s] \leq [u]$

□

Exercise 2.2 Let $\mathcal{N} = (\mathbb{N}, S_1, S_2)$ and $\mathcal{B} = (\mathbb{B}, R_1, R_2)$ be the following frames for a modal similarity type with two diamonds \Diamond_1, \Diamond_2 . Here, \mathbb{N} is the set of natural numbers and \mathbb{B} is the set of strings of 0's and 1's, and the relations are defined by

$$\begin{aligned} S_1(m, n) &\iff n = m + 1 \\ S_2(m, n) &\iff m > n \\ R_1(s, t) &\iff t = s0 \vee t = s1 \\ R_2(s, t) &\iff t \sqsubset s \end{aligned}$$

where $t \sqsubset s$ if and only if t is a *proper prefix* of s — i.e. t is a prefix of s such that $t \neq s$ (thus t can be ε). Which of the following formulas are valid on \mathcal{N} and \mathcal{B} , respectively?

- (a) $(\Diamond_1 p \wedge \Diamond_2 q) \rightarrow \Diamond_1(p \wedge q)$
- (b) $(\Diamond_2 p \wedge \Diamond_2 q) \rightarrow \Diamond_2(p \wedge q)$
- (c) $(\Diamond_1 p \wedge \Diamond_1 q \wedge \Diamond_1 r) \rightarrow (\Diamond_1(p \wedge q) \vee \Diamond_1(p \wedge r) \vee \Diamond_1(q \wedge r))$
- (d) $p \rightarrow \Diamond_1 \Box_1 p$
- (e) $p \rightarrow \Diamond_2 \Box_1 p$
- (f) $p \rightarrow \Box_1 \Diamond_2 p$
- (g) $p \rightarrow \Box_2 \Diamond_1 p$

Solution. First, consider the following extension to the \wedge operator on the inductive definition of satisfiability of formulas.

Claim: Given a model $\mathfrak{M} = (W, R, V)$, and a state $w \in W$, it holds that $\mathfrak{M}, w \models \phi \wedge \psi \iff \mathfrak{M}, w \models \phi \wedge \mathfrak{M}, w \models \psi$.

Proof of the Claim. By using De Morgan's law, we have that

$$\begin{aligned} \mathfrak{M}, w \models \phi \wedge \psi &= \neg(\neg\phi \vee \neg\psi) \iff \neg\mathfrak{M}, w \models \neg\phi \vee \neg\psi \\ &\iff \neg(\mathfrak{M}, w \models \neg\phi \vee \mathfrak{M}, w \models \neg\psi) \\ &\iff \neg(\neg\mathfrak{M}, w \models \phi \vee \neg\mathfrak{M}, w \models \psi) \\ &\iff \mathfrak{M}, w \models \phi \wedge \mathfrak{M}, w \models \psi \end{aligned}$$

□

For all the following propositions, we will assume that $\mathfrak{M} = (\mathbb{N}, S_1, S_2, V)$ and $\mathfrak{M}' = (\mathbb{B}, R_1, R_2, V)$ are two models over \mathcal{N} and \mathcal{B} respectively.

(a) $(\Diamond_1 p \wedge \Diamond_1 q) \rightarrow \Diamond_1(p \wedge q)$

- By the claim, for any $m \in \mathbb{N}$ it holds that

$$\begin{aligned}
 \mathfrak{M}, m \models \Diamond_1 p \wedge \Diamond_1 q &\iff \mathfrak{M}, m \models \Diamond_1 p \wedge \mathfrak{M}, m \models \Diamond_1 q \\
 &\iff \begin{cases} \exists n_p \in \mathbb{N} & S_1(m, n_p) \wedge \mathfrak{M}, n_p \models p \\ \exists n_q \in \mathbb{N} & S_1(m, n_q) \wedge \mathfrak{M}, n_q \models q \end{cases} \\
 &\iff \begin{cases} \exists n_p \in \mathbb{N} & n_p = m + 1 \wedge n_p \in V(p) \\ \exists n_q \in \mathbb{N} & n_q = m + 1 \wedge n_q \in V(q) \end{cases} \\
 &\iff m + 1 \in V(p) \wedge m + 1 \in V(q) \\
 &\iff m + 1 \in V(p) \cap V(q)
 \end{aligned}$$

and again, by the claim we have that

$$\begin{aligned}
 \mathfrak{M}, m \models \Diamond_1(p \wedge q) &\iff \exists n \in \mathbb{N} \quad S_1(m, n) \wedge \mathfrak{M}, n \models p \wedge q \\
 &\iff \exists n \in \mathbb{N} \quad n = m + 1 \wedge (\mathfrak{M}, n \models p \wedge \mathfrak{M}, n \models q) \\
 &\iff \exists n \in \mathbb{N} \quad n = m + 1 \wedge (n \in V(p) \wedge n \in V(q)) \\
 &\iff \exists n \in \mathbb{N} \quad n = m + 1 \wedge n \in V(p) \cap V(q) \\
 &\iff m + 1 \in V(p) \cap V(q)
 \end{aligned}$$

from which we conclude that

$$\mathfrak{M}, m \models \Diamond_1 p \wedge \Diamond_1 q \iff m + 1 \in V(p) \cap V(q) \iff \mathfrak{M}, m \models \Diamond_1(p \wedge q)$$

implying that the formula is valid on \mathcal{N} .

- By the claim, for any $s \in \mathbb{B}$ it holds that

$$\begin{aligned}
 \mathfrak{M}', s \models \Diamond_1 p \wedge \Diamond_1 q &\iff \mathfrak{M}', s \models \Diamond_1 p \wedge \mathfrak{M}', s \models \Diamond_1 q \\
 &\iff \begin{cases} \exists t_p \in \mathbb{B} & R_1(s, t_p) \wedge \mathfrak{M}', t_p \models p \\ \exists t_q \in \mathbb{B} & R_1(s, t_q) \wedge \mathfrak{M}', t_q \models q \end{cases} \\
 &\iff \begin{cases} \exists t_p \in \mathbb{B} & (t_p = s0 \vee t_p = s1) \wedge t_p \in V(p) \\ \exists t_q \in \mathbb{B} & (t_q = s0 \vee t_q = s1) \wedge t_q \in V(q) \end{cases} \\
 &\iff \begin{cases} s0 \in V(p) \vee s1 \in V(p) \\ s0 \in V(q) \vee s1 \in V(q) \end{cases} \\
 &\iff \{s0, s1\} \cap V(p) \neq \emptyset \wedge \{s0, s1\} \cap V(q) \neq \emptyset
 \end{aligned}$$

and again, by the claim we have that

$$\begin{aligned}
 \mathfrak{M}', s \models \Diamond_1(p \wedge q) &\iff \exists t \in \mathbb{B} \quad R_1(s, t) \wedge \mathfrak{M}', t \models p \wedge q \\
 &\iff \exists t \in \mathbb{B} \quad (t = s0 \vee t = s1) \wedge (\mathfrak{M}', t \models p \wedge \mathfrak{M}', t \models q) \\
 &\iff \exists t \in \mathbb{B} \quad (t = s0 \vee t = s1) \wedge (t \in V(p) \wedge t \in V(q)) \\
 &\iff \exists t \in \mathbb{B} \quad (t = s0 \vee t = s1) \wedge t \in V(p) \cap V(q) \\
 &\iff s0 \in V(p) \cap V(q) \vee s1 \in V(p) \cap V(q) \\
 &\iff \{s0, s1\} \cap V(p) \cap V(q) \neq \emptyset
 \end{aligned}$$

Now suppose $V(p) = \{s0\}$ and $V(q) = \{s1\}$; then we have that $\{s0, s1\} \cap V(p) = \{s0\} \neq \emptyset \wedge \{s0, s1\} \cap V(q) = \{s1\} \neq \emptyset \iff \mathfrak{M}', s \models \Diamond_1 p \wedge \Diamond_1 q$ although $\{s0, s1\} \cap V(p) \cap V(q) = \{s0, s1\} \cap \emptyset = \emptyset \iff \mathfrak{M}', s \not\models \Diamond_1(p \wedge q)$, implying that the formula is not valid on \mathcal{B} .

(b) $(\Diamond_2 p \wedge \Diamond_2 q) \rightarrow \Diamond_2(p \wedge q)$

- By definition, for any $m \in \mathbb{N}$ it holds that

$$\begin{aligned} \mathfrak{M}, m \models \Diamond_2 p \wedge \Diamond_2 q &\iff \mathfrak{M}, m \models \Diamond_2 p \wedge \mathfrak{M}, m \models \Diamond_2 q \\ &\iff \begin{cases} \exists n_p \in \mathbb{N} & S_2(m, n_p) \wedge \mathfrak{M}, n_p \models p \\ \exists n_q \in \mathbb{N} & S_2(m, n_q) \wedge \mathfrak{M}, n_q \models q \end{cases} \\ &\iff \begin{cases} \exists n_p \in \mathbb{N} & m > n_p \wedge n_p \in V(p) \\ \exists n_q \in \mathbb{N} & m > n_q \wedge n_q \in V(q) \end{cases} \end{aligned}$$

and again, by the claim we have that

$$\begin{aligned} \mathfrak{M}, m \models \Diamond_2(p \wedge q) &\iff \exists n \in \mathbb{N} \quad S_2(m, n) \wedge \mathfrak{M}, n \models p \wedge q \\ &\iff \exists n \in \mathbb{N} \quad m > n \wedge (\mathfrak{M}, n \models p \wedge \mathfrak{M}, n \models q) \\ &\iff \exists n \in \mathbb{N} \quad m > n \wedge (n \in V(p) \wedge n \in V(q)) \\ &\iff \exists n \in \mathbb{N} \quad m > n \wedge n \in V(p) \cap V(q) \end{aligned}$$

Now take an $n \geq 2$, and consider $n_p, n_q \in \mathbb{N}$ such that $n_p \neq n_q \wedge n > n_p, n_q$, and suppose that $V(p) = \{n_p\}$ and $V(q) = \{n_q\}$; then we have that $\begin{cases} \exists n_p \in \mathbb{N} & n > n_p \wedge n_p \in V(p) \\ \exists n_q \in \mathbb{N} & n > n_q \wedge n_q \in V(q) \end{cases} \iff \mathfrak{M}, m \models \Diamond_2 p \wedge \Diamond_2 q$ although $n_p \neq n_q \implies V(p) \cap V(q) = \emptyset \implies \nexists n \in \mathbb{N} \quad m > n \wedge n \in V(p) \cap V(q) \iff \mathfrak{M}, m \not\models \Diamond_2(p \wedge q)$, implying that the formula is not valid on \mathcal{N} .

- By the claim, for any $s \in \mathbb{B}$ it holds that

$$\begin{aligned} \mathfrak{M}', s \models \Diamond_2 p \wedge \Diamond_2 q &\iff \mathfrak{M}', s \models \Diamond_2 p \wedge \mathfrak{M}', s \models \Diamond_2 q \\ &\iff \begin{cases} \exists t_p \in \mathbb{B} & R_2(s, t_p) \wedge \mathfrak{M}', t_p \models p \\ \exists t_q \in \mathbb{B} & R_2(s, t_q) \wedge \mathfrak{M}', t_q \models q \end{cases} \\ &\iff \begin{cases} \exists t_p \in \mathbb{B} & t_p \sqsubset s \wedge t_p \in V(p) \\ \exists t_q \in \mathbb{B} & t_q \sqsubset s \wedge t_q \in V(q) \end{cases} \end{aligned}$$

and again, by the claim we have that

$$\begin{aligned} \mathfrak{M}', s \models \Diamond_2(p \wedge q) &\iff \exists t \in \mathbb{B} \quad t \sqsubset s \wedge (\mathfrak{M}', t \models p \wedge \mathfrak{M}', t \models q) \\ &\iff \exists t \in \mathbb{B} \quad t \sqsubset s \wedge (t \in V(p) \wedge t \in V(q)) \\ &\iff \exists t \in \mathbb{B} \quad t \sqsubset s \wedge t \in V(p) \cap V(q) \end{aligned}$$

Now take $s = 000$, consider $t_p = 0$ and $t_q = 0$, and suppose that $V(p) = \{t_p\} = \{0\}$ and $V(q) = \{t_q\} = \{00\}$; we observe that $t_p = 0 \sqsubset 000 = s$ and $t_q = 00 \sqsubset 000 = s$, therefore $\begin{cases} \exists t_p \in \mathbb{B} & t_p \sqsubset s \wedge t_p \in V(p) \\ \exists t_q \in \mathbb{B} & t_q \sqsubset s \wedge t_q \in V(q) \end{cases} \iff \mathfrak{M}', s \models \Diamond_2 p \wedge \Diamond_2 q$

although $V(p) \cap V(q) = \{t_p\} \cap \{t_q\} = \{0\} \cap \{00\} = \emptyset \implies \nexists t \in \mathbb{B} \quad t \sqsubset s \wedge t \in V(p) \cap V(q) \iff \mathfrak{M}', s \not\models \Diamond_2(p \wedge q)$, implying that the formula is not valid on \mathcal{B} .

$$(c) (\Diamond_1 p \wedge \Diamond_1 q \wedge \Diamond_1 r) \rightarrow (\Diamond_1(p \wedge q) \vee \Diamond_1(p \wedge r) \vee \Diamond_1(q \wedge r))$$

- By the claim, for any $m \in \mathbb{N}$ it holds that

$$\begin{aligned} \mathfrak{M}, m \models \Diamond_1 p \wedge \Diamond_1 q \wedge \Diamond_1 r &\iff \mathfrak{M}, m \models \Diamond_1 p \wedge \mathfrak{M}, m \models \Diamond_1 q \wedge \mathfrak{M}, m \models \Diamond_1 r \\ &\iff \begin{cases} \exists n_p \in \mathbb{N} & S_1(m, n_p) \wedge \mathfrak{M}, n_p \models \Diamond_1 p \\ \exists n_q \in \mathbb{N} & S_1(m, n_q) \wedge \mathfrak{M}, n_q \models \Diamond_1 q \\ \exists n_r \in \mathbb{N} & S_1(m, n_r) \wedge \mathfrak{M}, n_r \models \Diamond_1 r \end{cases} \\ &\iff \begin{cases} \exists n_p \in \mathbb{N} & n_p = m + 1 \wedge n_p \in V(p) \\ \exists n_q \in \mathbb{N} & n_q = m + 1 \wedge n_q \in V(q) \\ \exists n_r \in \mathbb{N} & n_r = m + 1 \wedge n_r \in V(r) \end{cases} \\ &\iff \begin{cases} m + 1 \in V(p) \\ m + 1 \in V(q) \\ m + 1 \in V(r) \end{cases} \\ &\iff m + 1 \in V(p) \cap V(q) \cap V(r) \end{aligned}$$

and again, by the claim we have that

$$\begin{aligned} \mathfrak{M}, m \models \Diamond_1(p \wedge q) \vee \Diamond_1(p \wedge r) \vee \Diamond_1(q \wedge r) &\iff (\mathfrak{M}, m \models \Diamond_1(p \wedge q)) \\ &\quad \vee (\mathfrak{M}, m \models \Diamond_1(p \wedge r)) \\ &\quad \vee (\mathfrak{M}, m \models \Diamond_1(q \wedge r)) \\ &\iff (\exists n_1 \in \mathbb{N} \quad S_1(m, n_1) \wedge \mathfrak{M}, n_1 \models p \wedge q) \\ &\quad \vee (\exists n_2 \in \mathbb{N} \quad S_1(m, n_2) \wedge \mathfrak{M}, n_2 \models p \wedge r) \\ &\quad \vee (\exists n_3 \in \mathbb{N} \quad S_1(m, n_3) \wedge \mathfrak{M}, n_3 \models q \wedge r) \\ &\iff (\exists n_1 \in \mathbb{N} \quad n_1 = m + 1 \wedge n_1 \in V(p) \cap V(q)) \\ &\quad \vee (\exists n_2 \in \mathbb{N} \quad n_2 = m + 1 \wedge n_2 \in V(p) \cap V(r)) \\ &\quad \vee (\exists n_3 \in \mathbb{N} \quad n_3 = m + 1 \wedge n_3 \in V(q) \cap V(r)) \\ &\iff (m + 1 \in V(p) \cap V(q)) \\ &\quad \vee (m + 1 \in V(p) \cap V(r)) \\ &\quad \vee (m + 1 \in V(q) \cap V(r)) \end{aligned}$$

Hence, we see that

$$\begin{aligned} \mathfrak{M}, m \models \Diamond_1 p \wedge \Diamond_1 q \wedge \Diamond_1 r &\iff m + 1 \in V(p) \cap V(q) \cap V(r) \\ &\implies \begin{cases} m + 1 \in V(p) \cap V(q) \\ m + 1 \in V(p) \cap V(r) \\ m + 1 \in V(q) \cap V(r) \end{cases} \\ &\iff \mathfrak{M}, m \models \Diamond_1(p \wedge q) \vee \Diamond_1(p \wedge r) \vee \Diamond_1(q \wedge r) \end{aligned}$$

implying that the formula is valid in \mathcal{N} .

- By the claim, for any $s \in \mathbb{B}$ it holds that

$$\begin{aligned}
\mathfrak{M}', s \models \Diamond_1 p \wedge \Diamond_1 q \wedge \Diamond_1 r &\iff \mathfrak{M}', s \models \Diamond_1 p \wedge \mathfrak{M}', s \models \Diamond_1 q \wedge \mathfrak{M}', s \models \Diamond_1 r \\
&\iff \begin{cases} \exists t_p \in \mathbb{B} & R_1(s, t_p) \wedge \mathfrak{M}', t_p \models p \\ \exists t_q \in \mathbb{B} & R_1(s, t_q) \wedge \mathfrak{M}', t_q \models q \\ \exists t_r \in \mathbb{B} & R_1(s, t_r) \wedge \mathfrak{M}', t_r \models r \end{cases} \\
&\iff \begin{cases} \exists t_p \in \mathbb{B} & (t_p = s0 \vee t_p = s1) \wedge t_p \in V(p) \\ \exists t_q \in \mathbb{B} & (t_q = s0 \vee t_q = s1) \wedge t_q \in V(q) \\ \exists t_r \in \mathbb{B} & (t_r = s0 \vee t_r = s1) \wedge t_r \in V(r) \end{cases} \\
&\iff \begin{cases} \{s0, s1\} \cap V(p) \neq \emptyset \\ \{s0, s1\} \cap V(q) \neq \emptyset \\ \{s0, s1\} \cap V(r) \neq \emptyset \end{cases}
\end{aligned}$$

and again, by the claim we have that

$$\begin{aligned}
\mathfrak{M}', s \models \Diamond_1(p \wedge q) \vee \Diamond_1(p \wedge r) \vee \Diamond_1(q \wedge r) &\iff (\mathfrak{M}', s \models \Diamond_1(p \wedge q)) \\
&\vee (\mathfrak{M}', s \models \Diamond_1(p \wedge r)) \\
&\vee (\mathfrak{M}', s \models \Diamond_1(q \wedge r)) \\
&\iff (\exists t_1 \in \mathbb{B} \quad R_1(s, t_1) \wedge \mathfrak{M}', t_1 \models p \wedge q) \\
&\vee (\exists t_2 \in \mathbb{B} \quad R_1(s, t_2) \wedge \mathfrak{M}', t_2 \models p \wedge r) \\
&\vee (\exists t_3 \in \mathbb{B} \quad R_1(s, t_3) \wedge \mathfrak{M}', t_3 \models q \wedge r) \\
&\iff (\exists t_1 \in \mathbb{B} \quad (t_1 = s0 \vee t_1 = s1) \wedge t_1 \in V(p) \cap V(q)) \\
&\vee (\exists t_2 \in \mathbb{B} \quad (t_2 = s0 \vee t_2 = s1) \wedge t_2 \in V(p) \cap V(r)) \\
&\vee (\exists t_3 \in \mathbb{B} \quad (t_3 = s0 \vee t_3 = s1) \wedge t_3 \in V(q) \cap V(r)) \\
&\iff \{s0, s1\} \cap V(p) \cap V(q) \neq \emptyset \\
&\vee \{s0, s1\} \cap V(p) \cap V(r) \neq \emptyset \\
&\vee \{s0, s1\} \cap V(q) \cap V(r) \neq \emptyset
\end{aligned}$$

Now suppose that $\mathfrak{M}', s \models \Diamond_1 p \wedge \Diamond_1 q \wedge \Diamond_1 r$, which happens if and only if

$$\begin{cases} \{s0, s1\} \cap V(p) \neq \emptyset \\ \{s0, s1\} \cap V(q) \neq \emptyset \\ \{s0, s1\} \cap V(r) \neq \emptyset \end{cases} \text{ as proved previously; by the pigeonhole principle,}$$

since there are 2 strings in $\{s0, s1\}$ and we have 3 sets $V(p)$, $V(q)$ and $V(r)$, there must be at least one string $x \in \{s0, s1\}$ such that $x \in V(a) \cap V(b)$, where $a, b \in \{p, q, r\}$ distinct. Without loss of generality, suppose that $x = s0$ and $a = p$ and $b = q$; then we have that

$$\begin{aligned}
x = s0 \in V(p) \cap V(q) &\implies x \in \{s0, s1\} \cap V(p) \cap V(q) \\
&\implies \{s0, s1\} \cap V(p) \cap V(q) \neq \emptyset \\
&\implies \mathfrak{M}', s \models \Diamond_1(p \wedge q) \vee \Diamond_1(p \wedge r) \vee \Diamond_1(q \wedge r)
\end{aligned}$$

implying that the formula is valid on \mathcal{B} .

(d) $p \rightarrow \Diamond_1 \Box_2 p$

- By definition, for any $m \in \mathbb{N}$ it holds that

$$\begin{aligned}
\mathfrak{M}, m \models \Diamond_1 \Box_2 p &\iff \exists n \in \mathbb{N} \quad S_1(m, n) \wedge \mathfrak{M}, n \models \Box_2 p \\
&\iff \exists n \in \mathbb{N} \quad n = m + 1 \wedge (\forall k \in \mathbb{N} \quad S_2(n, k) \implies \mathfrak{M}, k \models p) \\
&\iff \exists n \in \mathbb{N} \quad n = m + 1 \wedge (\forall k \in \mathbb{N} \quad n > k \implies k \in V(p)) \\
&\iff \forall k \in \mathbb{N} \quad m + 1 > k \implies k \in V(p) \\
&\iff V(p) = \{k \in \mathbb{N} \mid m + 1 > k\}
\end{aligned}$$

Now take $m = 1 \in \mathbb{N}$, and suppose $V(p) = \{m\} = \{1\}$; then $m \in V(p) \iff \mathfrak{M}, m \models p$, however for instance $k = 0 \in \mathbb{N}$ is such that $m + 1 = 1 + 1 = 2 > 0 = k$ although $k = 0 \notin V(p)$, therefore $\exists k \in \mathbb{N} \quad m + 1 > k \wedge k \notin V(p) \implies V(p) \neq \{k \in \mathbb{N} \mid m + 1 > k\} \iff \mathfrak{M}, m \not\models \Diamond_1 \Box_2 p$ which implies that the formula is not valid on \mathcal{N} .

- By definition, for any $s \in \mathbb{B}$ it holds that

$$\begin{aligned}
\mathfrak{M}', s \models \Diamond_1 \Box_2 p &\iff \exists t \in \mathbb{B} \quad R_1(s, t) \wedge \mathfrak{M}', t \models \Box_2 p \\
&\iff \exists t \in \mathbb{B} \quad (t = s0 \vee t = s1) \wedge (\forall u \in \mathbb{B} \quad R_2(t, u) \implies \mathfrak{M}', u \models p) \\
&\iff \exists t \in \mathbb{B} \quad (t = s0 \vee t = s1) \wedge (\forall u \in \mathbb{B} \quad u \sqsubset t \implies u \in V(p))
\end{aligned}$$

Now take $s = 00 \in \mathbb{B}$, and suppose $V(p) = \{s\} = \{00\}$; then $s \in V(p) \iff \mathfrak{M}', s \models p$, however if $t = s0$ or $t = s1$, there still is $u = 0$ such that $u = 0 \sqsubset 00 = t$ although $u = 0 \notin V(p)$, therefore $(t = s0 \vee t = s1) \implies (\exists u \in \mathbb{B} \quad u \sqsubset t \wedge u \notin V(p)) \iff \mathfrak{M}', s \not\models \Diamond_1 \Box_2 p$ which implies that the formula is not valid on \mathcal{B} .

(e) $p \rightarrow \Diamond_2 \Box_1 p$

- By definition, for any $m \in \mathbb{N}$ it holds that

$$\begin{aligned}
\mathfrak{M}, m \models \Diamond_2 \Box_1 p &\iff \exists n \in \mathbb{N} \quad S_2(m, n) \wedge \mathfrak{M}, n \models \Box_1 p \\
&\iff \exists n \in \mathbb{N} \quad m > n \wedge (\forall k \in \mathbb{N} \quad S_1(n, k) \implies \mathfrak{M}, k \models p) \\
&\iff \exists n \in \mathbb{N} \quad m > n \wedge (\forall k \in \mathbb{N} \quad k = n + 1 \implies k \in V(p)) \\
&\iff \exists n \in \mathbb{N} \quad m > n \wedge n + 1 \in V(p)
\end{aligned}$$

Now take $m = 0 \in \mathbb{N}$, and suppose $V(p) = \{m\} = \{0\}$; then $m \in V(p) \iff \mathfrak{M}, m \models p$, however there is no $n \in \mathbb{N}$ such that $m = 0 > n$, therefore $\nexists n \in \mathbb{N} \quad m > n \wedge n + 1 \in V(p) \iff \mathfrak{M}, m \not\models \Diamond_2 \Box_1 p$ which implies that the formula is not valid on \mathcal{N} .

- By definition, for any $s \in \mathbb{B}$ it holds that

$$\begin{aligned}
\mathfrak{M}', s \models \Diamond_2 \Box_1 p &\iff \exists t \in \mathbb{B} \quad R_2(s, t) \wedge \mathfrak{M}', t \models \Box_1 p \\
&\iff \exists t \in \mathbb{B} \quad t \sqsubset s \wedge (\forall u \in \mathbb{B} \quad R_1(t, u) \implies \mathfrak{M}', u \models p) \\
&\iff \exists t \in \mathbb{B} \quad t \sqsubset s \wedge (\forall u \in \mathbb{B} \quad (u = t0 \vee u = t1) \implies u \in V(p)) \\
&\iff \exists t \in \mathbb{B} \quad t \sqsubset s \wedge (t0 \in V(p) \vee t1 \in V(p))
\end{aligned}$$

Now take $s = \varepsilon \in \mathbb{B}$, and suppose $V(p) = \{s\} = \{\varepsilon\}$; then $s \in V(p) \iff \mathfrak{M}', s \models p$, however there is no $t \in \mathbb{B}$ such that $t \sqsubset s$, therefore $\nexists t \in \mathbb{B} \quad t \sqsubset s \wedge (t0 \in V(p) \vee t1 \in V(p)) \iff \mathfrak{M}', s \not\models \Diamond_2 \Box_1 p$ which implies that the formula is not valid on \mathcal{B} .

(f) $p \rightarrow \Box_1 \Diamond_2 p$

- By definition, for any $m \in \mathbb{N}$ it holds that

$$\begin{aligned}
\mathfrak{M}, m \models \Box_1 \Diamond_2 p &\iff \forall n \in \mathbb{N} \quad S_1(m, n) \implies \mathfrak{M}, n \models \Diamond_2 p \\
&\iff \forall n \in \mathbb{N} \quad n = m + 1 \implies (\exists k \in \mathbb{N} \quad S_2(n, k) \wedge \mathfrak{M}, k \models p) \\
&\iff \forall n \in \mathbb{N} \quad n = m + 1 \implies (\exists k \in \mathbb{N} \quad n > k \wedge k \in V(p)) \\
&\iff \exists k \in \mathbb{N} \quad m + 1 > k \wedge k \in V(p)
\end{aligned}$$

Now suppose $m \in V(p) \iff \mathfrak{M}, m \models p$; we observe that for every $m \in \mathbb{N}$ it holds that $m + 1 > m$, therefore $m + 1 > m \wedge m \in V(p) \iff \exists k \in \mathbb{N} \quad m + 1 > k \wedge k \in V(p) \iff \mathfrak{M}, m \models \Box_1 \Diamond_2 p$, which implies that the formula is valid on \mathcal{N} .

- By definition, for any $s \in \mathbb{B}$ it holds that

$$\begin{aligned}
\mathfrak{M}', s \models \Box_1 \Diamond_2 p &\iff \forall t \in \mathbb{B} \quad R_1(s, t) \implies \mathfrak{M}', t \models \Diamond_2 p \\
&\iff \forall t \in \mathbb{B} \quad (t = s0 \vee t = s1) \implies (\exists u \in \mathbb{B} \quad R_2(t, u) \wedge \mathfrak{M}', u \models p) \\
&\iff \forall t \in \mathbb{B} \quad (t = s0 \vee t = s1) \implies (\exists u \in \mathbb{B} \quad u \sqsubset t \wedge u \in V(p))
\end{aligned}$$

Now suppose $s \in V(p) \iff \mathfrak{M}', s \models p$; we observe that for every $s \in \mathbb{B}$ it holds that $s \sqsubset s0, s1$, therefore $\exists u \in \mathbb{B} \quad (u \sqsubset s0 \vee u \sqsubset s1) \wedge u \in V(p) \iff \mathfrak{M}', s \models \Box_1 \Diamond_2 p$, which implies that the formula is valid on \mathcal{B} .

(g) $p \rightarrow \Box_2 \Diamond_1 p$

- By definition, for any $m \in \mathbb{N}$ it holds that

$$\begin{aligned}
\mathfrak{M}, m \models \Box_2 \Diamond_1 p &\iff \forall n \in \mathbb{N} \quad S_2(m, n) \implies \mathfrak{M}, n \models \Diamond_1 p \\
&\iff \forall n \in \mathbb{N} \quad m > n \implies (\exists k \in \mathbb{N} \quad S_1(n, k) \wedge \mathfrak{M}, k \models p) \\
&\iff \forall n \in \mathbb{N} \quad m > n \implies (\exists k \in \mathbb{N} \quad k = n + 1 \wedge k \in V(p)) \\
&\iff \forall n \in \mathbb{N} \quad m > n \implies n + 1 \in V(p)
\end{aligned}$$

Now take $m = 2 \in \mathbb{N}$, and suppose $V(p) = \{m\} = \{2\}$; then $m \in V(p) \iff \mathfrak{M}, m \models p$, however for instance $n = 0 \in \mathbb{N}$ is such that $m = 2 > 0 = n$ and $n + 1 = 0 + 1 = 1 \notin V(p)$, therefore $\exists n \in \mathbb{N} \quad m > n \wedge n + 1 \notin V(p) \iff \mathfrak{M}, m \not\models \Box_2 \Diamond_1 p$, which implies that the formula is not valid on \mathcal{N} .

- By definition, for any $s \in \mathbb{B}$ it holds that

$$\begin{aligned}
\mathfrak{M}', s \models \Box_2 \Diamond_1 p &\iff \forall t \in \mathbb{B} \quad R_2(s, t) \implies \mathfrak{M}', t \models \Diamond_1 p \\
&\iff \forall s \in \mathbb{B} \quad t \sqsubset s \implies (\exists u \in \mathbb{B} \quad R_1(t, u) \wedge \mathfrak{M}', u \models p) \\
&\iff \forall s \in \mathbb{B} \quad t \sqsubset s \implies (\exists u \in \mathbb{B} \quad (u = t0 \vee u = t1) \implies u \in V(p))
\end{aligned}$$

Now take $s = 000 \in \mathbb{B}$, and suppose $V(p) = \{s\} = \{000\}$; then $s \in V(p) \iff \mathfrak{M}', s \models p$, however for instance $t = 0 \in \mathbb{B}$ is such that $t = 0 \sqsubset 000 = s$ although there is no $u = t0 = 00$ or $u = t1 = 01$ such that $u \in V(p)$, therefore $\exists t \in \mathbb{B} \quad t \sqsubset s \wedge (\nexists u \in \mathbb{B} \quad (u = t0 \vee u = t1) \wedge u \in V(p)) \iff \mathfrak{M}', s \not\models \Box_2 \Diamond_1 p$ which implies that the formula is not valid on \mathcal{B} .

□

Exercise 3.2 Consider the basic modal language, and the tuple $\mathfrak{f} = (\mathbb{N}, <, A)$ where A is the collection of finite and co-finite subsets of \mathbb{N} . Show that \mathfrak{f} is a general frame.

Solution. First, consider the following two claims.

Claim 1: If $X \subseteq \mathbb{N}$ is finite, and $Y \subseteq \mathbb{N}$ is co-finite, then $X \cup Y$ is co-finite.

Proof of the Claim. Since X is finite, $\mathbb{N} - X$ is finite, and since Y is co-finite, $\mathbb{N} - Y$ is finite; this implies that

$$\mathbb{N} - (X \cup Y) = (\mathbb{N} - X) \cap (\mathbb{N} - Y)$$

is the intersection of a co-finite and a finite set. In particular, we observe that

- such intersection will be a subset of $\mathbb{N} - Y$ by definition of intersection
- $\mathbb{N} - Y$ is finite
- a subset of a finite set is always finite

concluding that such intersection must be finite as well. Lastly, by definition we have that $\mathbb{N} - (X \cup Y)$ is finite if and only if $X \cup Y$ is co-finite. □

Claim 2: If $X, Y \subseteq \mathbb{N}$ are co-finite, then $X \cup Y$ is co-finite.

Proof of the Claim. By repeating the same argument of the previous claim, we have that $\mathbb{N} - (X \cup Y) = (\mathbb{N} - X) \cap (\mathbb{N} - Y)$ except that in this case both $\mathbb{N} - X$ and $\mathbb{N} - Y$ are finite, which implies that their intersection must be finite, hence $\mathbb{N} - (X \cup Y)$ is co-finite by definition. □

To prove that \mathfrak{f} is a general frame, it suffices to prove that the set A is closed under the following operations

- *union:* fix two sets $X, Y \in A$; then, by definition of A , we have that
 - if both X and Y are finite, then $X \cup Y$ is finite, hence $X \cup Y \in A$
 - without loss of generality, if X is finite and Y is co-finite, by Claim 1 $X \cup Y$ is co-finite, therefore $X \cup Y \in A$
 - if both X and Y are finite, then $X \cup Y$ is co-finite by Claim 2, therefore $X \cup Y \in A$
- *relative complement:* fix a set $X \in A$; then, by definition of A we trivially have that
 - if X is finite, then $\mathbb{N} - X$ is co-finite, hence $\mathbb{N} - X \in A$
 - if X is co-finite, then $\mathbb{N} - X$ is finite, hence $\mathbb{N} - X \in A$
- *modal operations:* assume that “ $<$ ” is the relation referring to a unary modal operator $\langle < \rangle$, and fix a set $X \in A$; by definition, we have that

$$m_{\langle < \rangle}(X) = \{n \in \mathbb{N} \mid \exists x \in X \quad n < x\}$$

therefore, we have that

- if X is finite, then

$$m_{(<)}(X) = \{n \in \mathbb{N} \mid n < \max(X)\}$$

therefore $m_{(<)}(X)$ is an “initial segment of \mathbb{N} ”, implying that it is finite, hence $m_{(<)}(X) \in A$

- if X is co-finite, then $\mathbb{N} - X$ is finite, implying that X is infinite; this implies that $\max(X)$ is not defined, therefore

$$m_{(<)}(X) = \mathbb{N} \implies \mathbb{N} - m_{(<)}(X) = \mathbb{N} - \mathbb{N} = \emptyset$$

and since \emptyset is finite, we conclude that $m_{(<)}(X)$ is co-finite, thus $m_{(<)}(X) \in A$

□

Exercise 4.3 Let Σ be a set of formulas in the basic modal language, and let \mathbf{F} denote the class of all frames. Show that $\Sigma \models_{\mathbf{F}}^g \phi$ if and only if $\{\Box^n \sigma \mid \sigma \in \Sigma, n \in \mathbb{N}\} \models_{\mathbf{F}} \phi$.

Solution. Let $\Pi := \{\Box^n \sigma \mid \sigma \in \Sigma, n \in \mathbb{N}\}$. We are going to split the two directions of the proof into two claims.

Claim: If $\Sigma \models_{\mathbf{F}}^g \phi$, then $\Pi \models_{\mathbf{F}} \phi$.

Proof of the Claim. Assume that $\Sigma \models_{\mathbf{F}}^g \phi$; fix a frame $\mathcal{F} \in \mathbf{F}$ defined over a world W , a model \mathfrak{M} over \mathcal{F} , and a world $w \in W$. Suppose that $\mathfrak{M}, w \models \Pi$; then, by the claim of Exercise 2.2 we know that

$$\begin{aligned} & \forall \sigma \in \Sigma, n \in \mathbb{N} \quad \mathfrak{M}, w \models \Box^n \sigma \\ \iff & \forall \sigma \in \Sigma, n \in \mathbb{N}, x_1, \dots, x_n \in W \quad w \xrightarrow{R} x_1 \xrightarrow{R} \dots \xrightarrow{R} x_n \implies \mathfrak{M}, x_n \models \sigma \end{aligned}$$

where $a \xrightarrow{R} b \iff R(a, b)$.

Now, consider the following restriction of W

$$W_w := \{v \in W_w \mid \exists n \in \mathbb{N}, x_1, \dots, x_{n-1} \in W \quad w \xrightarrow{R} x_1 \xrightarrow{R} \dots \xrightarrow{R} x_{n-1} \xrightarrow{R} v\}$$

where $v \in W_w$ if and only if v can be R -reached from w through a sequence of R -accessible elements. Moreover, consider the following restriction of R

$$R_w := (W_w \times W_w) \cap R$$

in which we consider the tuples of R that connect elements of W_w . Then, since \mathbf{F} is the class of all frames, we know that $\mathcal{F}_w = (W_w, R_w) \in \mathbf{F}$. Lastly, consider a model \mathfrak{M}_w over \mathcal{F}_w such that $\mathfrak{M}_w = (W_w, R_w, V_w)$ where

$$V_w = W_w \cap V$$

Consider some $x_1, \dots, x_n \in W$ such that $w \xrightarrow{R} x_1 \xrightarrow{R} \dots \xrightarrow{R} x_n$; by definition of W_w , this implies that all x_1, \dots, x_n are R -reachable from w , which implies that $x_1, \dots, x_n \in W_w$; this means that

$$\begin{aligned} & \forall \sigma \in \Sigma, n \in \mathbb{N}, x_1, \dots, x_n \in W \quad w \xrightarrow{R} x_1 \xrightarrow{R} \dots \xrightarrow{R} x_n \implies \mathfrak{M}, x_n \models \sigma \\ \implies & \forall \sigma \in \Sigma, n \in \mathbb{N}, x_1, \dots, x_n \in W_W \quad w \xrightarrow{R} x_1 \xrightarrow{R} \dots \xrightarrow{R} x_n \implies \mathfrak{M}, x_n \models \sigma \end{aligned}$$

Moreover, since x_1, \dots, x_n are elements of W_w , by definition of R_w it holds that

$$\begin{aligned} & \forall \sigma \in \Sigma, n \in \mathbb{N}, x_1, \dots, x_n \in W_W \quad w \xrightarrow{R} x_1 \xrightarrow{R} \dots \xrightarrow{R} x_n \implies \mathfrak{M}, x_n \models \sigma \\ \implies & \forall \sigma \in \Sigma, n \in \mathbb{N}, x_1, \dots, x_n \in W_W \quad w \xrightarrow{R_w} x_1 \xrightarrow{R_w} \dots \xrightarrow{R_w} x_n \implies \mathfrak{M}, x_n \models \sigma \end{aligned}$$

Furthermore, since $W_w \subseteq W$ and $R_w \subseteq R$, by definition of \mathfrak{M}_w we get that

$$\begin{aligned} & \forall \sigma \in \Sigma, n \in \mathbb{N}, x_1, \dots, x_n \in W_W \quad w \xrightarrow{R_w} x_1 \xrightarrow{R_w} \dots \xrightarrow{R_w} x_n \implies \mathfrak{M}, x_n \models \sigma \\ \implies & \forall \sigma \in \Sigma, n \in \mathbb{N}, x_1, \dots, x_n \in W_W \quad w \xrightarrow{R_w} x_1 \xrightarrow{R_w} \dots \xrightarrow{R_w} x_n \implies \mathfrak{M}_w, x_n \models \sigma \end{aligned}$$

This observation concludes that

$$\begin{aligned} & \forall \sigma \in \Sigma, n \in \mathbb{N}, x_1, \dots, x_n \in W \quad w \xrightarrow{R} x_1 \xrightarrow{R} \dots \xrightarrow{R} x_n \implies \mathfrak{M}, x_n \models \sigma \\ \implies & \forall \sigma \in \Sigma, n \in \mathbb{N}, x_1, \dots, x_n \in W_W \quad w \xrightarrow{R_w} x_1 \xrightarrow{R_w} \dots \xrightarrow{R_w} x_n \implies \mathfrak{M}_w, x_n \models \sigma \end{aligned}$$

Now, since for any $v \in W_w$ there are $y_1, \dots, y_k \in W$ such that $w \xrightarrow{R_w} y_1 \xrightarrow{R_w} \dots \xrightarrow{R_w} y_k \xrightarrow{R} v$ by definition of W_w , the previous observation implies that

$$\begin{aligned} & \forall \sigma \in \Sigma \in \mathbb{N} \quad \mathfrak{M}, w \models \Box^n \sigma \\ \iff & \forall \sigma \in \Sigma, n \in \mathbb{N}, x_1, \dots, x_n \in W \quad w \xrightarrow{R} x_1 \xrightarrow{R} \dots \xrightarrow{R} x_n \implies \mathfrak{M}, x_n \models \sigma \\ \implies & \forall \sigma \in \Sigma, n \in \mathbb{N}, x_1, \dots, x_n \in W_W \quad w \xrightarrow{R_w} x_1 \xrightarrow{R_w} \dots \xrightarrow{R_w} x_n \implies \mathfrak{M}_w, x_n \models \sigma \\ \implies & \forall \sigma \in \Sigma, v \in W_w \quad \mathfrak{M}_w, v \models \sigma \\ \implies & \forall v \in W_w, \sigma \in \Sigma \quad \mathfrak{M}_w, v \models \sigma \\ \iff & \forall v \in W_w \quad \mathfrak{M}_w, v \models \Sigma \\ \implies & \forall v \in W_w \quad \mathfrak{M}_w, v \models \phi \quad (\Sigma \models_F^g \phi) \\ \iff & \forall v \in W_w \quad w \in V_w(\phi) \\ \implies & W_w \subseteq V_w(\phi) \subseteq V(\phi) \end{aligned}$$

Lastly, we observe that $w \in W_w$, and since $W_w \subseteq V(\phi)$ we have that $w \in V(\phi)$, which happens if and only if $\mathfrak{M}, w \models \phi$.

This proves that for any frame $\mathcal{F} \in \mathbf{F}$ defined over a world W , any model \mathfrak{M} over \mathcal{F} , and every world $w \in W$, it holds that

$$\mathfrak{M}, w \models \Pi \implies \mathfrak{M}, w \models \phi$$

which implies that $\Pi \models_F \phi$ by definition. □

Claim: If $\Pi \models_F \phi$, then $\Sigma \models_F^g \phi$.

Proof of the Claim. Assume that $\Pi \models_F \phi$; fix a frame $\mathcal{F} \in F$ defined over a world W , and a model \mathfrak{M} over \mathcal{F} . Suppose that $\forall w \in W \quad \mathfrak{M}, w \models \Sigma$; then, by the claim of Exercise 2.2 we obtain the following

$$\begin{aligned}
 & \forall w \in W \quad \mathfrak{M}, w \models \Sigma \\
 \iff & \forall w \in W, \sigma \in \Sigma \quad \mathfrak{M}, w \models \sigma \\
 \implies & \forall n \in \mathbb{N}, w, x_1, \dots, x_n \in W, \sigma \in \Sigma \quad w \xrightarrow{R} x_1 \xrightarrow{R} \dots \xrightarrow{R} x_n \implies \mathfrak{M}, x_n \models \sigma \\
 \implies & \forall n \in \mathbb{N}, w \in W, \sigma \in \Sigma \quad \mathfrak{M}, w \models \Box^n \sigma \\
 \iff & \forall w \in W \quad \mathfrak{M}, w \models \Pi \\
 \implies & \forall w \in W \quad \mathfrak{M}, w \models \phi \quad (\Pi \models_F \phi)
 \end{aligned}$$

This proves that for any frame $\mathcal{F} \in F$ defined over a world W , and any model \mathfrak{M} over \mathcal{F} it holds that

$$\forall w \in W \quad \mathfrak{M}, w \models \Sigma \implies \forall w \in W \quad \mathfrak{M}, w \models \phi$$

which implies that $\Sigma \models_F^g \phi$ by definition. □

Finally, the two claims conclude the exercise. □

Exercise 5.1 Give K -proofs of $(\Box p \wedge \Diamond q) \rightarrow \Diamond(p \wedge q)$ and $\Diamond(p \vee q) \leftrightarrow (\Diamond p \vee \Diamond q)$.

Solution. In the first section of the solution, we are going to prove some useful derivations that will be extensively used in the actual K -proof of the two propositions. The right side of each line will be one of the following:

- (K): the K axiom
- (D): the Dual axiom
- (T): a propositional Tautology
- (MP(i, j)): the Modus Ponens rule applied on lines i and j
- (S(i)): the Substitution rule applied on line i
- (G(i)): the Generalization rule applied on line i
- (C $_k$ (i_1, \dots, i_n)): the k -th Claim applied on lines i_1, \dots, i_n — $k \in [7]$ and n depends on the number of lines the Claim refers to

Claim 1: If $p \rightarrow q$ can be K -proved, and $q \rightarrow r$ can be K -proved, then $p \rightarrow r$ can be K -proved in 4 steps.

Proof of the Claim. Consider a K -proof in which $p \rightarrow q$ is proved at step i , and $q \rightarrow r$ is

proved at step j — without loss of generality suppose that $i < j$; then we have that

$$\begin{array}{ll}
 \dots & \\
 i. \vdash p \rightarrow q & \\
 \dots & \\
 j. \vdash q \rightarrow r & \\
 j+1. \vdash (a \rightarrow b) \rightarrow ((b \rightarrow c) \rightarrow (a \rightarrow c)) & (T) \\
 j+2. \vdash (p \rightarrow q) \rightarrow ((q \rightarrow r) \rightarrow (p \rightarrow r)) & (S(j+1)) \\
 j+3. \vdash (q \rightarrow r) \rightarrow (p \rightarrow r) & (MP(i, j+2)) \\
 j+4. \vdash p \rightarrow r & (MP(j, j+3))
 \end{array}$$

□

Claim 2: If $p \rightarrow q$ can be K -proved, then $\Box p \rightarrow \Box q$ can be K -proved in 4 steps.

Proof of the Claim. Consider a K -proof in which $p \rightarrow q$ is proved at step i ; then, we have that

$$\begin{array}{ll}
 \dots & \\
 i. \vdash p \rightarrow q & \\
 i+1. \vdash \Box(p \rightarrow q) & (G(i)) \\
 i+2. \vdash \Box(a \rightarrow b) \rightarrow (\Box a \rightarrow \Box b) & (K) \\
 i+3. \vdash \Box(p \rightarrow q) \rightarrow (\Box p \rightarrow \Box q) & (S(i+2)) \\
 i+4. \vdash \Box p \rightarrow \Box q & (MP(i+1, i+3))
 \end{array}$$

□

Claim 3: If $p \rightarrow q$ can be K -proved, and $p \rightarrow r$ can be K -proved, then $p \rightarrow q \wedge r$ can be K -proved in 4 steps.

Proof of the Claim. Consider a K -proof in which $p \rightarrow q$ is proved at step i , and $p \rightarrow r$ is proved at step j — without loss of generality suppose $i < j$; then, we have that

$$\begin{array}{ll}
 \dots & \\
 i. \vdash p \rightarrow q & \\
 \dots & \\
 j. \vdash p \rightarrow r & \\
 j+1. \vdash (a \rightarrow b) \rightarrow ((a \rightarrow c) \rightarrow (a \rightarrow b \wedge c)) & (T) \\
 j+2. \vdash (p \rightarrow q) \rightarrow ((p \rightarrow r) \rightarrow (p \rightarrow q \wedge r)) & (S(j+1)) \\
 j+3. \vdash (p \rightarrow r) \rightarrow (p \rightarrow q \wedge r) & (MP(i, j+2)) \\
 j+4. \vdash p \rightarrow q \wedge r & (MP(j, j+3))
 \end{array}$$

□

Claim 4: If $p \rightarrow q$ can be K -proved, then $\neg q \rightarrow \neg p$ can be K -proved in 3 steps. Moreover, if $p \rightarrow \neg q$ can be K -proved, then $q \rightarrow \neg p$ can be K -proved in 3 steps.

Proof of the Claim. Consider a K -proof in which $p \rightarrow q$ is proved at step i ; then, we have that

$$\begin{array}{ll}
 \dots & \\
 i. \vdash p \rightarrow q & \\
 i + 1. \vdash (a \rightarrow b) \rightarrow (\neg b \rightarrow \neg a) & (T) \\
 i + 2. \vdash (p \rightarrow q) \rightarrow (\neg q \rightarrow \neg p) & (S(i + 1)) \\
 i + 3. \vdash \neg q \rightarrow \neg p & (MP(i, i + 2))
 \end{array}$$

The same K -proof can be used to prove the rest of the claim by using the propositional tautology $(a \rightarrow \neg b) \rightarrow (b \rightarrow \neg a)$. \square

Claim 5: If $p \leftrightarrow q$ can be K -proved, then $p \rightarrow q$ and $q \rightarrow p$ can be K -proved in 3 steps.

Proof of the Claim. Consider a K -proof in which $p \leftrightarrow q$ is proved at step i ; then, we have that

$$\begin{array}{ll}
 \dots & \\
 i. \vdash p \leftrightarrow q & \\
 i + 1. \vdash (a \leftrightarrow b) \rightarrow (a \rightarrow b) & (T) \\
 i + 2. \vdash (p \leftrightarrow q) \rightarrow (p \rightarrow q) & (S(i + 1)) \\
 i + 3. \vdash p \rightarrow q & (MP(i, i + 2))
 \end{array}$$

The case for $q \rightarrow p$ can be proved analogously by using the propositional tautology $(a \leftrightarrow b) \rightarrow (b \rightarrow a)$. \square

Claim 6: If $p \rightarrow (q \rightarrow r)$ can be K -proved, then $p \wedge q \rightarrow r$ can be K -proved in 3 steps.

Proof of the Claim. Consider a K -proof in which $p \rightarrow (q \rightarrow r)$ is proved at step i ; then, we have that

$$\begin{array}{ll}
 \dots & \\
 i. \vdash p \rightarrow (q \rightarrow r) & \\
 i + 1. \vdash (a \rightarrow (b \rightarrow c)) \rightarrow (a \wedge b \rightarrow c) & (T) \\
 i + 2. \vdash (p \rightarrow (q \rightarrow r)) \rightarrow (p \wedge q \rightarrow r) & (S(i + 1)) \\
 i + 3. \vdash p \wedge q \rightarrow r & (MP(i, i + 2))
 \end{array}$$

\square

Claim 7: If $p \rightarrow q$ can be K -proved, and $q \rightarrow p$ can be K -proved, then $p \leftrightarrow q$ can be proved in 4 steps.

Proof of the Claim. Consider a K -proof in which $p \rightarrow q$ is proved at step i , and $q \rightarrow p$ can be proved at step j — without loss of generality suppose $i < j$; then, we have that

$$\begin{array}{ll}
\ldots & \\
i. \vdash p \rightarrow q & \\
\ldots & \\
j. \vdash q \rightarrow p & \\
j+1. \vdash (a \rightarrow b) \rightarrow ((b \rightarrow a) \rightarrow (a \leftrightarrow b)) & (T) \\
j+2. \vdash (p \rightarrow q) \rightarrow ((q \rightarrow p) \rightarrow (p \leftrightarrow q)) & (S(j+1)) \\
j+3. \vdash (q \rightarrow p) \rightarrow (p \leftrightarrow q) & (MP(i, j+2)) \\
j+4. \vdash p \leftrightarrow q & (MP(j, j+3))
\end{array}$$

□

Now that we proved some preliminary claims, we can K -prove the two given propositions.

Claim 8: $(\Box p \wedge \Diamond q) \rightarrow \Diamond(p \wedge q)$ is K -provable.

Proof of the Claim.

$$\begin{array}{ll}
1. \vdash \Box(p \rightarrow q) \rightarrow (\Box p \rightarrow \Box q) & (K) \\
2. \vdash (\neg a \vee \neg b) \rightarrow (a \rightarrow \neg b) & (T) \\
3. \vdash (\neg p \vee \neg q) \rightarrow (p \rightarrow \neg q) & (S(2)) \\
\ldots & \\
7. \vdash \Box(\neg p \vee \neg q) \rightarrow \Box(p \rightarrow \neg q) & (C_2(3)) \\
8. \vdash \Box(p \rightarrow \neg q) \rightarrow (\Box p \rightarrow \Box \neg q) & (S(1)) \\
\ldots & \\
12. \vdash \Box(\neg p \vee \neg q) \rightarrow (\Box p \rightarrow \Box \neg q) & (C_1(7, 8)) \\
13. \vdash (a \rightarrow b) \rightarrow (\neg a \vee b) & (T) \\
14. \vdash (\Box p \rightarrow \Box \neg q) \rightarrow (\neg \Box p \vee \Box \neg q) & (S(13)) \\
\ldots & \\
18. \vdash \Box(\neg p \vee \neg q) \rightarrow (\neg \Box p \vee \Box \neg q) & (C_1(12, 14)) \\
19. \vdash (\neg a \vee b) \rightarrow \neg(a \wedge \neg b) & (T) \\
20. \vdash (\neg \Box p \vee \Box \neg q) \rightarrow \neg(\Box p \wedge \neg \Box \neg q) & (S(19)) \\
\ldots & \\
24. \vdash \Box(\neg p \vee \neg q) \rightarrow \neg(\Box p \wedge \neg \Box \neg q) & (C_1(18, 20)) \\
\ldots & \\
27. \vdash (\Box p \wedge \neg \Box \neg q) \rightarrow \neg \Box(\neg p \vee \neg q) & (C_4(24)) \\
28. \vdash \neg(a \wedge b) \rightarrow (\neg a \vee \neg b) & (T) \\
29. \vdash \neg(p \wedge q) \rightarrow (\neg p \vee \neg q) & (S(27)) \\
\ldots &
\end{array}$$

- ...
33. $\vdash \Box \neg(p \wedge q) \rightarrow \Box(\neg p \vee \neg q)$ $(C_2(28))$
- ...
36. $\vdash \neg \Box(\neg p \vee \neg q) \rightarrow \neg \Box \neg(p \wedge q)$ $(C_4(32))$
- ...
40. $\vdash (\Box p \wedge \neg \Box \neg q) \rightarrow \neg \Box \neg(p \wedge q)$ $(C_1(27, 36))$
41. $\vdash \Diamond a \leftrightarrow \neg \Box \neg a$ (D)
42. $\vdash \Diamond(p \wedge q) \leftrightarrow \neg \Box \neg(p \wedge q)$ $(S(41))$
- ...
45. $\vdash \neg \Box \neg(p \wedge q) \rightarrow \Diamond(p \wedge q)$ $(C_5(42))$
- ...
49. $\vdash (\Box p \wedge \neg \Box \neg q) \rightarrow \Diamond(p \wedge q)$ $(C_1(40, 45))$
50. $\vdash \Diamond q \leftrightarrow \neg \Box \neg q$ $(S(41))$
- ...
53. $\vdash \Diamond q \rightarrow \neg \Box \neg q$ $(C_5(50))$
54. $\vdash (b \rightarrow c) \rightarrow (a \wedge b \rightarrow a \wedge c)$ (T)
55. $\vdash (\Diamond q \rightarrow \neg \Box \neg q) \rightarrow (\Box p \wedge \Diamond q \rightarrow \Box p \wedge \neg \Box \neg q)$ $(S(54))$
56. $\vdash \Box p \wedge \Diamond q \rightarrow \Box p \wedge \neg \Box \neg q$ $(MP(53, 55))$
- ...
60. $\vdash (\Box p \wedge \Diamond q) \rightarrow \Diamond(p \wedge q)$ $(C_1(56, 49))$

□

This claim concludes the K -proof of the first proposition. To K -prove the second proposition, we are going to split the K -proof into 3 claims.

Claim 9: $(\Diamond p \vee \Diamond q) \rightarrow \Diamond(p \vee q)$ is K -provable.

Proof of the Claim.

1. $\vdash \neg p \wedge \neg q \rightarrow \neg p$ (T)
- ...
5. $\vdash \Box(\neg p \wedge \neg q) \rightarrow \Box \neg p$ $(C_2(1))$
6. $\vdash \neg p \wedge \neg q \rightarrow \neg q$ (T)
- ...
10. $\vdash \Box(\neg p \wedge \neg q) \rightarrow \Box \neg q$ $(C_2(2))$
- ...
14. $\vdash \Box(\neg p \wedge \neg q) \rightarrow \Box \neg p \wedge \Box \neg q$ $(C_3(10))$
- ...

- ...
17. $\vdash \neg(\Box\neg p \wedge \Box\neg q) \rightarrow \neg\Box(\neg p \wedge \neg q)$ (C₄(14))
18. $\vdash \neg a \vee \neg b \rightarrow \neg(a \wedge b)$ (T)
19. $\vdash \neg\Box\neg p \vee \neg\Box\neg q \rightarrow \neg(\Box\neg p \wedge \Box\neg q)$ (S(18))
- ...
23. $\vdash \neg\Box\neg p \vee \neg\Box\neg q \rightarrow \neg\Box(\neg p \wedge \neg q)$ (C₁(22, 17))
24. $\vdash \neg(a \vee b) \rightarrow \neg a \wedge \neg b$ (T)
25. $\vdash \neg(p \vee q) \rightarrow \neg p \wedge \neg q$ (S(24))
- ...
29. $\vdash \Box\neg(p \vee q) \rightarrow \Box(\neg p \wedge \neg q)$ (C₂(25))
- ...
32. $\vdash \neg\Box(\neg p \wedge \neg q) \rightarrow \neg\Box\neg(p \vee q)$ (C₄(29))
- ...
36. $\vdash \neg\Box\neg p \vee \neg\Box\neg q \rightarrow \neg\Box\neg(p \vee q)$ (C₁(23, 32))
37. $\vdash \Diamond a \leftrightarrow \neg\Box\neg a$ (D)
- ...
40. $\vdash \neg\Box\neg a \rightarrow \Diamond a$ (C₅(37))
41. $\vdash \neg\Box\neg(p \vee q) \rightarrow \Diamond(p \vee q)$ (S(40))
- ...
45. $\vdash \neg\Box\neg p \vee \neg\Box\neg q \rightarrow \Diamond(p \vee q)$ (C₁(36, 41))
- ...
49. $\vdash \Diamond a \rightarrow \neg\Box\neg a$ (C₅(37))
50. $\vdash \Diamond p \rightarrow \neg\Box\neg p$ (S(49))
51. $\vdash \Diamond q \rightarrow \neg\Box\neg q$ (S(49))
52. $\vdash (a \rightarrow c) \rightarrow ((b \rightarrow d) \rightarrow (a \vee b \rightarrow c \vee d))$ (T)
53. $\vdash (\Diamond p \rightarrow \neg\Box\neg p) \rightarrow ((\Diamond q \rightarrow \neg\Box\neg q) \rightarrow (\Diamond p \vee \Diamond q \rightarrow \neg\Box\neg p \vee \neg\Box\neg q))$ (S(52))
54. $\vdash (\Diamond q \rightarrow \neg\Box\neg q) \rightarrow (\Diamond p \vee \Diamond q \rightarrow \neg\Box\neg p \vee \neg\Box\neg q)$ (MP(50, 53))
55. $\vdash \Diamond p \vee \Diamond q \rightarrow \neg\Box\neg p \vee \neg\Box\neg q$ (MP(51, 54))
- ...
59. $\vdash \Diamond p \vee \Diamond q \rightarrow \Diamond(p \vee q)$ (C₁(55, 48))

□

Claim 10: $\Diamond(p \vee q) \rightarrow (\Diamond p \vee \Diamond q)$ is K -provable.

Proof of the Claim.

1. $\vdash \neg p \rightarrow (\neg q \rightarrow \neg p \wedge \neg q)$ (T)
- ...
5. $\vdash \Box\neg p \rightarrow \Box(\neg q \rightarrow \neg p \vee \neg q)$ (C₂(1))
- ...

- ...
6. $\vdash \Box(a \rightarrow b) \rightarrow (\Box a \rightarrow \Box b)$ (K)
7. $\vdash \Box(\neg q \rightarrow \neg p \wedge \neg q) \rightarrow (\Box \neg q \rightarrow \Box(\neg p \wedge \neg q))$ (S(6))
- ...
11. $\vdash \Box \neg p \rightarrow (\Box \neg q \rightarrow \Box(\neg p \wedge \neg q))$ ($C_1(5, 7)$)
- ...
14. $\vdash \Box \neg p \wedge \Box \neg q \rightarrow \Box(\neg p \wedge \neg q)$ ($C_6(11)$)
- ...
17. $\vdash \neg \Box(\neg p \wedge \neg q) \rightarrow \neg(\Box \neg p \wedge \Box \neg q)$ ($C_4(14)$)
18. $\vdash \neg(a \wedge b) \rightarrow \neg a \vee \neg b$ (T)
19. $\vdash \neg(\Box \neg p \wedge \Box \neg q) \rightarrow \neg \Box \neg p \vee \neg \Box \neg q$ (S(18))
- ...
23. $\vdash \neg \Box(\neg p \wedge \neg q) \rightarrow \neg \Box \neg p \vee \neg \Box \neg q$ ($C_1(17, 19)$)
24. $\vdash (\neg a \wedge \neg b) \rightarrow \neg(a \vee b)$ (T)
25. $\vdash (\neg p \wedge \neg q) \rightarrow \neg(p \vee q)$ (S(24))
- ...
29. $\vdash \Box(\neg p \wedge \neg q) \rightarrow \Box \neg(p \vee q)$ ($C_2(25)$)
- ...
32. $\vdash \neg \Box \neg(p \vee q) \rightarrow \neg \Box(\neg p \wedge \neg q)$ ($C_4(29)$)
- ...
36. $\vdash \neg \Box \neg(p \vee q) \rightarrow \neg \Box \neg p \vee \neg \Box \neg q$ ($C_1(32, 23)$)
37. $\vdash \Diamond a \leftrightarrow \neg \Box \neg a$ (D)
- ...
41. $\vdash \neg \Box \neg a \rightarrow \Diamond a$ ($C_5(37)$)
42. $\vdash \neg \Box \neg p \rightarrow \Diamond p$ (S(41))
43. $\vdash \neg \Box \neg q \rightarrow \Diamond q$ (S(41))
44. $\vdash (a \rightarrow c) \rightarrow ((b \rightarrow d) \rightarrow (a \vee b \rightarrow c \vee d))$ (T)
45. $\vdash (\neg \Box \neg p \rightarrow \Diamond p) \rightarrow ((\neg \Box \neg q \rightarrow \Diamond q) \rightarrow (\neg \Box \neg p \vee \neg \Box \neg q \rightarrow \Diamond p \vee \Diamond q))$ (S(44))
46. $\vdash (\neg \Box \neg q \rightarrow \Diamond q) \rightarrow (\neg \Box \neg p \vee \neg \Box \neg q \rightarrow \Diamond p \vee \Diamond q)$ (MP(41, 45))
47. $\vdash \neg \Box \neg p \vee \neg \Box \neg q \rightarrow \Diamond p \vee \Diamond q$ (MP(42, 46))
- ...
51. $\vdash \Diamond a \rightarrow \neg \Box \neg a$ ($C_5(37)$)
52. $\vdash \Diamond(p \vee q) \rightarrow \neg \Box \neg(p \vee q)$ (S(51))
- ...
56. $\vdash \Diamond(p \vee q) \rightarrow (\neg \Box \neg p \vee \neg \Box \neg q)$ ($C_1(52, 36)$)
- ...
60. $\vdash \Diamond(p \vee q) \rightarrow \Diamond p \vee \Diamond q$ ($C_1(56, 47)$)

□

Claim 11: $\Diamond(p \vee q) \leftrightarrow \Diamond(p \vee \Diamond q)$ is K -provable.

Proof of the Claim.

$$\begin{array}{rcl}
 & \dots & \\
 59. & \vdash \Diamond p \vee \Diamond q \rightarrow \Diamond(p \vee q) & (C_9) \\
 & \dots & \\
 119. & \vdash \Diamond(p \vee q) \rightarrow \Diamond p \vee \Diamond q & (C_{10}) \\
 & \dots & \\
 123. & \vdash \Diamond(p \vee q) \leftrightarrow \Diamond p \vee \Diamond q & (C_7(2, 1))
 \end{array}$$

□

This last claim concludes that the second proposition is K -provable as well.

□