

# Advanced Coding Theory and Cryptography

*Notes by: Alex Pellegrini*

# Contents

1	An introduction to Gröbner bases	2
---	----------------------------------	---

# Chapter 1

## An introduction to Gröbner bases

**Theorem 2.1.10** (Hilbert's Basis Theorem)

*Proof.* We proceed by induction on the number of variables. Let  $I \subset A[X]$  be an ideal not finitely generated, we may assume it can be constructed by an infinite sequence  $(f_i)_{i \in \mathbb{N}}$  of independent polynomials of minimal degree. "Independent" means that  $f_i \in I \setminus J_i$  where we set  $J_i := \langle f_0, \dots, f_{i-1} \rangle$ . Now let  $a_i := lc(f_i)$  be the leading coefficient of  $f_i$  and consider  $J := \langle a_0, a_1, \dots \rangle \subset A$ . We know that  $J$  can be a basis for an ideal in  $A$  but since  $A$  is a Noetherian ring we have that there exists a finite basis for such ideal, say  $J = \langle a_1, \dots, a_N \rangle$ . We claim that  $I = \langle f_1, \dots, f_N \rangle =: I'$ . Suppose by contrary that this is not true then take a polynomial  $f_{N+1} \in I$ , we want to show that it is a linear combination of elements of  $I'$ :

$$a_{N+1} = u_1 a_1 + u_2 a_2 + \dots + u_N a_N$$

Consider

$$g := \sum_{i=1}^N u_i f_i x^{\deg(f_{N+1}) - \deg(f_i)} \in I'$$

it has the same degree and same leading coefficient as  $f_{N+1}$ . Now  $f_{N+1} - g \notin I'$  and has degree strictly less than  $f_{N+1}$  contradicting its minimality. Therefore  $f_{N+1} - g$  must be 0 and  $f_{N+1} \in I'$ .

The induction follows since we can consider  $A[X_1, \dots, X_m] = A'[X_m]$  where  $A' := [X_1, \dots, X_{m-1}]$  which we know is a Noetherian ring.  $\square$

**Lemma 2.1.13** (Dickson's Lemma)

*Proof.* We proceed by induction on the number of variable, by first proving the case with one variable. So we are considering  $\mathcal{M} = \{X_1^\alpha \mid \alpha \in \mathbb{N}\}$ , and

$T \subset \mathcal{M}$  a semigroup ideal. Since every  $t_i \in T$  is of the form  $t = X_1^{\alpha_i}$  we consider  $\beta = \min\{\alpha_i | X_1^{\alpha_i} \in T\}$ . We claim that  $T = \langle X_1^\beta \rangle$ . Indeed let  $t_j \in T$  then it is of the form  $t_j = X_1^{\alpha_j}$  so  $\frac{t_j}{t_i} = X_1^{\alpha_j - \beta}$  is well defined where  $\alpha_j - \beta > 0$  by minimality of  $\beta$ . We can take  $\gamma = \alpha_j - \beta$  hence:

$$t_j = X_1^\beta \cdot X_1^\gamma \in \langle X_1^\beta \rangle = T$$

We prove the more general case so let be  $m \in \mathbb{N}$  arbitrary and assume the lemma proved for  $m - 1$ .

Let  $T \subset \mathcal{M} = \{X_1^{a_1} \cdots X_m^{a_m} \mid (a_1, \dots, a_m) \in \mathbb{N}^m\}$ . Consider also the projection map  $\pi(X_1^{a_1} \cdots X_m^{a_m}) = X_1^{a_1} \cdots X_{m-1}^{a_{m-1}}$ . By induction hypothesis  $\pi(T)$  is a finitely generated semigroup ideal so we can find a basis, say  $\pi(T) = \langle t_1, \dots, t_k \rangle$ . Now let:

$$A_i := \min\{a_m \mid X_m^{a_m} | t, t \in T, \pi(t) = t_i\} \quad \forall i = 1, \dots, k$$

and furthermore

$$A := \min\{a_m \mid X_m^{a_m} \in T\}$$

We claim that  $T = \langle t_1 X_m^{A_1}, \dots, t_k X_m^{A_k}, X_m^A \rangle$  which is a finite set.

So pick an arbitrary  $t \in T$  so  $t = \pi(t) X_m^{a_{m_t}}$  for some  $a_{m_t} \in \mathbb{N}$ , we know that  $\exists t_i$  such that  $\pi(t) = s \cdot t_i$ , therefore  $t = s \cdot t_i \cdot X_m^{a_{m_t}}$  and by minimality of  $A_i$  we obtain that for:

$$t = s \cdot t_i \cdot X_m^{a_{m_t}} = s \cdot t_i \cdot X_m^{A_i} \cdot X_m^\gamma$$

for  $\gamma = a_{m_t} - A_i$ . Now  $\forall t \in T$  we have proved that  $t \in \langle t_i \cdot X_m^{A_i} \rangle$  which is contained in  $\langle t_1 X_m^{A_1}, \dots, t_k X_m^{A_k}, X_m^A \rangle$   $\square$

### Theorem 2.1.14

*Proof.*  $\Rightarrow$  Let  $f \in I$  then we can write:

$$f = \sum_{i=1}^k f_i \cdot p_i = f_1 \cdot p_1 + f_2 \cdot p_2 + \dots + f_k \cdot p_k, \quad f_i \in \mathcal{P}$$

So evaluating  $f(A)$  means to evaluate every  $p_i$  so:

$$\begin{aligned} f(A) &= f_1 \cdot p_1(A) + f_2 \cdot p_2(A) + \dots + f_k \cdot p_k(A) = \\ &= f_1 \cdot 0 + f_2 \cdot 0 + \dots + f_k \cdot 0 = 0 \end{aligned}$$

$\Leftarrow$  Trivial by setting  $f = p_i \quad \forall i = 1, \dots, k$   $\square$