

Our network is a graph  $(S, F)$  with  $S$  a set of switches and  $F$  a set of fibers. We associate a binary variable  $x_{ij}$  to each pair of adjacent switches  $(i, j) \in S^2$ , such that

$$x_{ij} = \begin{cases} 1 & \text{if the fiber } (i, j) \in F \text{ belongs to the path} \\ 0 & \text{otherwise} \end{cases}$$

We note  $d_{ij}$  the distance between  $i$  and  $j$ . The length  $L$  of a path is

$$L = \sum_{(i,j) \in F} d_{ij} x_{ij}$$

Finding the shortest path comes down to minimizing  $L$ .

A few constraints must be introduced to ensure the flow conservation :

- The difference between the outgoing traffic and the incoming traffic at the source  $s$  must be equal to 1.

$$\sum_{j:(s,j) \in F} x_{sj} - \sum_{j:(j,s) \in F} x_{js} = 1$$

- For any switch other than the source and the destination, the outgoing traffic and the incoming traffic must be equal (flow conservation).

$$\sum_{j:(i,j) \in F} x_{ij} - \sum_{j:(j,i) \in F} x_{ji} = 0 \quad \forall i \in S \setminus \{s, d\}$$

- the  $x_{ij}$  are binary variables.

$$x_{ij} \in \{0, 1\} \quad \forall (i, j) \in F$$

Note : at the destination switch, we must have

$$\sum_{j:(d,j) \in F} x_{dj} - \sum_{j:(j,d) \in F} x_{jd} = -1$$

. However, it can be proven that this is a consequence of the two above-mentioned constraints combined : it doesn't need to be an explicit constraint.

$$\sum_{j:(s,j) \in F} x_{sj} - \sum_{j:(j,s) \in F} x_{js} = 1$$

We note  $G(P, E)$  the transformed graph which nodes are traffic paths  $p \in P$ , and we define  $\Lambda$  a set of wavelengths such that  $\Lambda = \{\lambda_1, \lambda_2, \dots, \lambda_{|\Lambda|}\}$ .

Given a wavelength  $\lambda$  and a path  $p$ , we define the binary variables

$$x_p^\lambda = \begin{cases} 1 & \text{if } \lambda \text{ is assigned to } p \\ 0 & \text{otherwise} \end{cases}$$

$$y_\lambda = \begin{cases} 1 & \text{if } \lambda \text{ is used at least once} \\ 0 & \text{otherwise} \end{cases}$$

Minimizing the number of assigned wavelengths in the optical network comes down to minimizing

$$\sum_{\lambda \in \Lambda} y_\lambda$$

To obtain a valid tour, we need to add the following constraints :

— For every path, there must one and only one wavelength assigned.

$$\sum_{\lambda \in \Lambda} x_p^\lambda = 1 \quad \forall p \in P$$

— If a wavelength  $\lambda$  is not used ( $y_\lambda = 0$ ), it cannot be assigned to any path ( $x_p^\lambda = 0 \quad \forall p \in P$ ), and if it is used ( $y_\lambda = 1$ ), it cannot be assigned to more than one path (i.e  $x_p^\lambda + x_{p'}^\lambda \leq 1 \quad \forall (p, p') \in P^2$ ). In other words,

$$x_p^\lambda + x_{p'}^\lambda \leq y_\lambda \quad \forall \lambda \in \Lambda, \quad \forall (p, p') \in P^2$$

— Wavelengths are assigned sequentially, in increasing order of indices.

$$y_{\lambda_k} \geq y_{\lambda_{k+1}} \quad \forall k \in \llbracket 1, |\Lambda| - 1 \rrbracket$$

— For every path  $p$  and every wavelength  $\lambda$ ,  $x_p^\lambda$  and  $y_\lambda$  are binary variables.

$$y_\lambda \in \{0, 1\} \quad \forall \lambda \in \Lambda$$

$$x_p^\lambda \in \{0, 1\} \quad \forall p \in P, \quad \forall \lambda \in \Lambda$$