Learning to program with F#

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Part I

F# basics

### Chapter 4

## Quick-start guide

Programming is the art of solving problems by writing a program to be executed by a computer. For example, to solve the following problem,

What is the sum of 357 and 864?

we have written the following program in F#,

```
let a = 357
let b = 864
let c = a + b
printfn "%A" c

1221
Listing 4.1: quickStartSum.fsx - A script to add 2 numbers and print the result to the console.
```

In box the above, we see our program was saved as a script in a file called quickStartSum.fsx, and in the console we executed the program by typing the command fsharpi quickStartSum.fsx. The result is then printed in the console to be 1221.

To solve the program we made program consisting of several lines, where each line was a statement. The first statement let a = 357 used the let keyword to bind the value 357 to the name a. Likewise, we bound the value 864 to the name b, but to the name c we bound the result of evaluating the expression a + b. That is, first the value a + b was calculated by substituting the names of a and b with their values to give the expression, 357 + 864, then this expression was evaluated by adding the values to give, 1221, and this value was finally bound to the name c. The last line printed the value of c to the console followed by a LF (line feed, see Appendix B.1) with the printfn function. Here printfn is a function of 2 arguments: "%A" and c. Notice, that in contrast to many other languages, F# does not use parentheses to frame the list of arguments, nor does it use commas to separate them. In general, the printfn function always has 1 or more arguments, and the first is a format string. A string is a sequence of characters starting and ending with double quotation marks. E.g., let s = "this is a string of characters" binds the string "this is..." to the name s. For the printfn function, the format string may be any string, but if it contains format character sequences, such as %A, then the values following the format string are substituted. The format string must match the value type, that is, here c is of type integer, whereas the format string %A matches any type.

Types are a central concept in F#. In the script 4.1 we bound values of types int and string to names. The values were not *declared* to have these types, instead the types were *inferred* by F#. Had we typed these statements line by line in an interactive session, then we would have seen the inferred types:

 $\cdot$  statement

·let

 $\cdot$  keyword

· binding

 $\cdot$  expression

 $\cdot$  format string

 $\cdot$  string

· type

· type declaration

· type inference

```
> let a = 357;;
val a : int = 357
> let b = 864;;
val b : int = 864
> let c = a + b;;
val c : int = 1221
> printfn "%A" c;;
1221
val it : unit = ()
Listing 4.2: fsharpi
```

The an interactive session displays the type using the *val* keyword. Since the value is also responded, then the last printfn statement is superfluous. However, it is ill adviced to design programs to be run in an interactive session, since the scripts needs to be manually copied every time it is to be run, and since the starting state may be unclear.

·val

Advice!

Were we to solve a slightly different problem,

What is the sum of 357.6 and 863.4?

then we would have to use floating point arithmatic instead of integers, and the program would look like,

```
let a = 357.6
let b = 863.4
let c = a + b
printfn "%A" c

1221.0
Listing 4.3: quickStartSumFloat.fsx - Floating point types and arithmatic.
```

On the surface, this could appear as an almost negligible change, but the set of integers and the set of real numbers (floats) require quite different representations, in order to be effective on a computer, and as a consequence, the implementation of their operations such as addition are very different. Thus, although the response is an integer, it has type float, which is indicated by 1221.0 which is not the same as 1221. F# is very picky about types, and generally does not allow types to be mixed. E.g., in an interactive session,

```
> let a = 357;;
val a : int = 357
> let b = 863.4;;
val b : float = 863.4
> let c = a + b;;
```

```
let c = a + b;;
-----
/Users/sporring/Desktop/fsharpNotes/stdin(3,13): error FS0001: The
    type 'float' does not match the type 'int'
Listing 4.4: fsharpi
```

we see that binding a name to a number without a decimal point is inferred to be integer, while when binding to a number with a decimal point, then the type is inferred to be a float, and when trying to add values of integer and floating point, then we get an error.

F# is a functional first programming language, and one implication is that names are constant and cannot be changed. If attempted, then F# will return an error as, e.g.,

However, if the same was performed in an interactive session,

```
> let a = 357;;
val a : int = 357
> let a = 864;;
val a : int = 864
Listing 4.6: fsharpi
```

then apparently rebinding is legal. The difference is that the ;; token defines a new nested scope. A token is a letter or a word, which the F# considers as an atomic unit. A scope is an area in a program, where a binding is valid. Scopes can be nested, and in F# a binding may reuse names in a nested scope, in which case the previous value is overshadowed. I.e., attempting the same without ;; between the two let statements results in an error, e.g.,

```
A ·;;
um, ·token
ted ·scope
een ·nested scope
· overshadow
```

```
> let a = 357
- let a = 864;;
let a = 864;;
----^
/Users/sporring/Desktop/fsharpNotes/stdin(2,5): error FS0037:
    Duplicate definition of value 'a'
Listing 4.7: fsharpi
```

Scopes can be visualized as nested squares as shown in Figure 4.1.

In F# functions are also values, and defining a function sum as part of the solution to the above  $\cdot$  function program gives,

```
let a = 357;
let a = 864;;
(a) Illegal let a = 357;;
let a = 864;;
```

Figure 4.1: Binding of the the same name in the same scope is illegal in F# 2, but legal in a different scopes. In (a) the two bindings are in the same scope, which is illegal, while in (b) the bindings are in separate scopes by the extra ;; token, which is legal.

```
let sum x y = x + y
let c = sum 357 864
printfn "%A" c

1221
Listing 4.8: quickStartSumFct.fsx - A script to add 2 numbers using a user defined function.
```

Entering the function into an interactive session will illustrate the inferred type, the function sum has: val sum: x:int \* y:int -> int, by which is meant that sum is a mapping from the set product of integers with integers into integers. Type inference in F# may cause problems, since the type of a function is inferred in the context, in which it is defined. E.g., in an interactive session, defining the sum in one scope on a single line will default the types to integers, F#'s favorite type, which will give an error, if it in a nested scope is to be used for floats,

A remedy is to either define the function in the same scope as its use,

```
> let sum x y = x + y
- let c = sum 357.6 863.4;;

val sum : x:float -> y:float -> float
val c : float = 1221.0
Listing 4.10: fsharpi
```

In this chapter, we have scratched the surface of learning how to program by concentrating on a number of key programming concepts and how they are expressed in the F# language. In the following chapters, we will expand the description of F# with features used in all programming approaches.

### Chapter 5

## Numbers, Characters, and Strings

All programs rely on processing of data, and an essential property of data is its type. F# contains a number of built-in types, and it designed such that it is easy to define new types. The simplest types are called  $primitive\ types$ , and a table of some of the most commonly used primitive types are shown in Table 5.1. A literal is a fixed value such as "3", and F# supports  $literal\ types$  which are indicated by a suffix in most cases and as shown in the table. A name is bound to a value by the syntax,

 $\cdot$  type

- · primitive types
- $\cdot \, literal$
- $\cdot$  literal type

```
"let" [ "mutable" ] ident [":" type] "=" expr ["in" expr]
```

That is, the *let* keyword indicates that the following is a binding of a name with an expression, and that the type may be specified with the : token. The binding may be mutable, which will be discussed in Section 5.4, and the binding may only be for the last expression as indicated by the *in*. The simplest example of an expression is a *literal*, i.e., a constant such as the number 3 or a function. Functions will be discussed in detail in Chapter 6. Examples of let statements with *literals* 

```
· let
· :
· in
· literal
· literals
```

```
> let a = 3
- let b = 4u
- let c = 5.6
- let d = 7.9f
- let e = 'A'
- let f = 'B'B
- let g = "ABC"
- let h = ();;

val a : int = 3
val b : uint32 = 4u
val c : float = 5.6
val d : float32 = 7.9000001f
val e : char = 'A'
val f : byte = 66uy
val g : string = "ABC"
val h : unit = ()
Listing 5.1: fsharpi
```

Here  $a, b, \ldots, h$  are names that we have chosen, and which by the binding operation are made equivalent to the corresponding values. Note that we did not specify the type of the name, and that F# interpreted the type from the literal form of the right-hand-side. Types matter, since the operations that can be performed on integers are quite different from those that can be performed on characters and strings. I.e.,

Metatype	Type name	Suffix	Literal	Description
Boolean	bool	none	true	Boolean values true or false
Integer	int	none or l	3	Integer values from -2,147,483,648 to
				2,147,483,647
	byte	uy	3uy	Integer values from 0 to 255
	sbyte	У	3y	Integer values from -128 to 127
	int16	s	3s	Integer values from -32768 to 32767
	uint16	us	3us	Integer values from 0 to 65535
	int32	none or l	3	Synonymous with int
	uint32	u or ul	3u	Integer values from 0 to 4,294,967,295
	int64	L	3L	Integer values from -9,223,372,036,854,775,808
				to 9,223,372,036,854,775,807
	uint64	UL	3UL	Integer values from 0 to
				18,446,744,073,709,551,615
	bigint	I	3I	Integer not limited to 64 bits
	nativeint	n	3n	A native pointer as a signed integer
	unativeint	un	3un	A native pointer as an unsigned integer
Real	float	none	3.0	64-bit IEEE 754 floating point value from
				$-\infty$ to $\infty$
	double	none	3.0	Synonymous with float
	single	F or f	3.0f	A 32-bit floating point type
	float32	F or f	3.0f	Synonymous with single
	decimal	M or m	$3.0 \mathrm{m}$	A floating point data type that has at least
				28 significant digits
Character	char	none	'c'	Unicode character
	byte	В	'c'B	ASCII character
	string	none	"abc"	Unicode sequence of characters
	byte[]	В	"abc"B	Unicode sequence of characters
	string or byte[]	@	@"\n"	Verbatim string
None	unit	none	()	No value denoted

Table 5.1: List of primitive types and the corresponding literal. The most commonly used types are highlighted in bold. Note that string verbatim uses a prefix instead of suffix notation. For at description of floating point numbers see Appendix A.2 and for ASCII and Unicode characters see Appendix B.

```
> let a = 3
- let b = 3.0
- let c = '3'
- let d = "3";;

val a : int = 3
val b : float = 3.0
val c : char = '3'
val d : string = "3"
Listing 5.2: fsharpi
```

the variables a, b, c, and d all represent the number 3, but their types are different, and hence they are quite different values. When specifying the type, then the type and the literal form must match, i.e., mixing types and literals gives an error,

```
> let a : float = 3;;
let a : float = 3;;
------
/Users/sporring/repositories/fsharpNotes/stdin(50,17): error FS0001
    : This expression was expected to have type
    float
but here has type
    int
Listing 5.3: fsharpi
```

since the left-hand-side is a name of type float while the right-hand-side is a literal of type integer. Many primitive types are compatible and may changed to each other by  $type\ casting$ . E.g.,

· type casting

```
> let a = float 3;;

val a : float = 3.0

Listing 5.4: fsharpi
```

where the left-hand-side is inferred to by of type float, since the integer number 3 is casted to float resulting in a similar floating point value, in this case the float point number 3.0. As a particular note, the boolean values are often treated as the integer values 0 and 1, however casting can only be performed with built-in functions, e.g.,

```
> let a = System.Convert.ToBoolean 1
- let b = System.Convert.ToBoolean 0
- let c = System.Convert.ToInt32 true
- let d = System.Convert.ToInt32 false;;

val a : bool = true
val b : bool = false
val c : int = 1
val d : int = 0
Listing 5.5: fsharpi
```

Here System.Convert.ToBoolean is the name of a function ToBoolean, which is a member of the · member

class Convert that is included in the namespace System. Namespaces, classes, and members are all part of Structured programming to be discussed in Part IV. For more on functions see Section 6 Typecasting is often a destructive operation, e.g., typecasting a float to int removes the part after the decimal point without rounding,

```
· class
· namespace
```

```
let a = 357.6
let b = int a
printfn "%A -> %A" a b

357.6 -> 357
Listing 5.6: quickStartDownCast.fsx - Fractional part is removed by downcasting.
```

Here we typecasted to a lesser type, in the sense that integers is a subset of floats, which is called downcasting. The opposite is called upcasting is often non-destructive, as Listing 5.4 showed, where an integer was casted to a float while retaining its value. As a side note, rounding a number y.x, where y is the whole part and x is the fractional part, is the operation of mapping numbers in the interval  $y.x \in [y.0, y.5)$  to y and  $y.x \in [y.5, y+1)$  to y+1. This can be performed by downcasting as follows,

```
· downcasting
```

- ·upcasting
- $\cdot$  rounding

```
let a = 357.6
let b = int (a + 0.5)
printfn "%A -> %A" a b

357.6 -> 358
Listing 5.7: rounding.fsx - The rounding function may be obtained by downcasting.
```

since if  $y.x \in [y.0, y.5)$ , then  $y.x + 0.5 \in [y.5, y + 1)$ , from which downcasting removes the fractional part resulting in y. And if  $y.x \in [y.5, y + 1)$ , then  $y.x + 0.5 \in [y + 1, y + 1.5)$ , from which downcasting removes the fractional part resulting in y + 1. Hence, the result is rounding.

If parentheses are ommitted in Listing 5.7, then F# will interpret the expression as (int a) + 0.5, which is erroneous, since addition of an integer with a float is undefined. This is an example of precedence, i.e., typecasting takes precedence over addition meaning that it is performed before addition. Consider the arithmetic expression, whose result is bound to a by

```
> let a = 3 + 4 * 5;;

val a : int = 23
Listing 5.8: fsharpi
```

Here, the addition and multiplication functions are shown in *infix notation* with the *operator* tokens + and \*. To arrive at the resulting value 23, F# has to decide in which order to perform the calculation. There are 2 possible orders, 3 + (4 \* 5) or (3 + 4) \* 5, which gives different results. For integer arithmetic, the correct order is of course to multiply before addition, and we say that multiplication takes *precedence* over addition. Every atomic operation that F# can perform is ordered in terms of its precedences, and for some common built-in operators shown in Table 5.2, the precedence is shown by the order they are given in the table. Associativity implies the order in which calculations are performed for operators of same precedence. For some operators and type combinations association matters little, e.g., multiplication associates to the left and exponentiation associates to the right, e.g., in

```
· infix notation · operator
```

 $\cdot \ precedence$ 

```
> let a = 3.0*4.0*5.0

- let b = (3.0*4.0)*5.0
```

Operator	Associativity	Example	Description
f x	Left	f 3	Function evaluation
+op, -op	Left	-3	Unary operator
&&	Left	true && true	Boolean and
11	Left	true    true	Boolean or
op ** op	Right	3.0 ** 2.0	Exponent
op * op, op / op, op % op	Left	3.0 / 2.0	Multiplication, division and re-
			mainder
op + op, op - op	Left	3.0 + 2.0	Addition and subtraction binary
			operators

Table 5.2: Some common operators, their precedence, and their associativity. Rows are ordered from highest to lowest precedences, such that op \* op has higher precedence than op + op. Operators in the same row has same precedence.

```
- let c = 3.0*(4.0*5.0);;

val a : float = 60.0
val b : float = 60.0
val c : float = 60.0

> let d = 4.0 ** 3.0 ** 2.0
- let e = (4.0 ** 3.0) ** 2.0
- let f = 4.0 ** (3.0 ** 2.0);;

val d : float = 262144.0
val e : float = 4096.0
val f : float = 262144.0
Listing 5.9: fsharpi
```

the expression for  ${\tt a}$  is interpreted as  ${\tt b}$  but gives the same results as  ${\tt c}$  since association does not matter for multiplication of numbers, but the expression for  ${\tt d}$  is interpreted as  ${\tt f}$  which is quite different from  ${\tt e}$ .

A less common notation is to define bindings for expressions using the *in* keyword, e.g.,

```
let p = 2.0 in printfn "%A" (3.0 ** p)

9.0
Listing 5.10: numbersIn.fsx - The name p is only bound in the nested scope following the keyword in.
```

Here p is only bound in the *scope* of the expression following the in keyword, in this the printfn · scope statement, and p is unbound in lines that follows.

·in

### 5.1 Integers and Reals

The set of integers and reals are infinitely large, and since all computers have limited resources, it is not possible to represent these sets in their entirety. The various integer and floating point types listed in Table 5.1 are finite subset where the integer types have been reduced by limiting their ranges and the floating point types have been reduced by sampling the space of reals. An in-depth description of integer and floating point implementations can be found in Appendix ??. The int and float are the most common types.

For integers the following arithmetic operators are defined:

+op, -op: These are unary plus and minus operators, and plus has no effect, but minus changes the sign, e.g.,

```
> let a = 5
- let b = -a;;

val a : int = 5
val b : int = -5
Listing 5.11: fsharpi
```

op + op, op - op, op \* op: These are binary operators, where addition, subtraction and multiplication performs the usual operations,

```
> let a = 7 + 3
- let b = 7 - 3
- let c = 7 * 3;;

val a : int = 10
val b : int = 4
val c : int = 21
Listing 5.12: fsharpi
```

op / op, op % op: These are binary operators, and division performs integer division, where the fractional part is discarded after division, and the \% is the remainder operator, which calculates the remainder after integer division,

```
> let a = 7 / 3
- let b = 7 % 3;;

val a : int = 2
val b : int = 1
Listing 5.13: fsharpi
```

Together integer division and remainder is a lossless representation of the original number as,

```
> let x = 7
- let whole = x / 3
- let remainder = x % 3
- let y = whole * 3 + remainder;;

val x : int = 7
val whole : int = 2
val remainder : int = 1
val y : int = 7
Listing 5.14: fsharpi
```

And we see that x and y is bound to the same value.

Integer exponentiation is not defined as an operator, but this is available the built-in function pown, e.g.,

```
> pown 2 5;;
val it : int = 32
Listing 5.15: fsharpi
```

which is equal to  $2^5$ . Note that when no let statement is used in conjunction with an expression then F# automatically binds the result to the it name, i.e., the above is equal to

·it

```
> let it = pown 2 5;;

val it : int = 32
Listing 5.16: fsharpi
```

Rumor has it, that the name it is an abreviation for 'irrelevant'.

Performing arithmetic operations on int types requires extra care, since the result may cause overflow, underflow, and even exceptions, e.g., the range of the integer type sbyte is [-128...127], which causes problems in the following example,

```
\begin{array}{l} \cdot \ overflow \\ \cdot \ underflow \end{array}
```

```
> let a = 100y - let b = 30y - let c = a+b;;

val a : sbyte = 100y
val b : sbyte = 30y
val c : sbyte = -126y
Listing 5.17: fsharpi
```

Here 100+30=130, which is larger than the biggest **sbyte**, and the result is an overflow. Similarly, we get an underflow, when the arithmetic result falls below the smallest value storable in an **sbyte**,

```
> let a = -100y
- let b = -30y
- let c = a+b;;

val a : sbyte = -100y
val b : sbyte = -30y
val c : sbyte = 126y
Listing 5.18: fsharpi
```

Notice that neither overflow nor underflow error gave rise to an error message, which is why such bugs are difficult to find. Dividing any non-zero number with 0 is infinite, which is also outside the domain of any of the integer types, but in this case, F# casts an exception,

· exception

```
Stopped due to error
Listing 5.19: fsharpi
```

The output looks daunting at first sight, but the first and last line of the error message are the most important parts, which tells us what exception was cast and why the program stopped. The middle are technical details concerning which part of the program caused this, and can be ignored for the time being. Exceptions are a type of *run-time error*, and are treated in Chapter ??

 $\cdot$  run-time error

Integers can also be written in binary, octal, or hexadecimal format using the prefixes 0b, 0o, and 0x, e.g.,

```
> let a = 0b1011
- let b = 0o13
- let c = 0xb;;

val a : int = 11
val b : int = 11
val c : int = 11
Listing 5.20: fsharpi
```

For a description of binary representations see Appendix A.1. The overflow error in Listing 5.17 can be understood in terms of the binary representation of integers: In binary,  $130 = 10000010_2$ , and this binary pattern is interpreted differently as byte and sbyte,

```
> let a = 0b10000010uy
- let b = 0b10000010y;;

val a : byte = 130uy
val b : sbyte = -126y
Listing 5.21: fsharpi
```

That is, for signed bytes, the left-most bit is used to represent the sign, and since the addition of  $100 = 01100100_2$  and  $30 = 00011110_b$  is  $130 = 10000010_2$  causes the left-most bit to be used, then this is wrongly interpreted as a negative number, when stored in an sbyte.

For floating point numbers the following arithmetic operators are defined:

+op, -op: These are unary plus and minus operators, and plus has no effect, but minus changes the sign, e.g.,

```
> let a = 5.0
- let b = -a;;

val a : float = 5.0
val b : float = -5.0
Listing 5.22: fsharpi
```

op + op, op - op, op \* op, op / op: These are binary operators, where addition, subtraction, multiplication, and division performs the usual operations,

```
> let a = 7.0 + 3.0
- let b = 7.0 - 3.0
- let c = 7.0 * 3.0
- let d = 7.0 / 3.0;;
```

op % op: The binary remainder operator, and division performs integer division, where the fractional part is discarded after division, and the \% is the remainder operator, which calculates the remainder after integer division,

```
> let a = 7.0 / 3.0
- let b = 7.0 % 3.0;;

val a : int = 2.0
val b : int = 1.0
Listing 5.24: fsharpi
```

The remainder for floating point numbers can be fractional, but division, rounding, and remainder is still a lossless representation of the original number as,

```
> let x = 7.0
- let division = x / 3.2
- let whole = float (int (division + 0.5))
- let remainder = x % 3.2
- let y = whole * 3.2 + remainder;;

val x : float = 7.0
val division : float = 2.1875
val whole : float = 2.0
val remainder : float = 0.6
val y : float = 7.0
Listing 5.25: fsharpi
```

And we see that x and y is bound to the same value.

op \*\* op: In spite of an unusual notation, the binary exponentiation operator performs the usual calculation,

```
> let a = 2.0 ** 5.0;;

val a : float = 32.0
Listing 5.26: fsharpi
```

which is equal to  $2^5$ .

Arithmetic using float will not cause over- and underflow problems, since the IEEE 754 standard includes the special numbers  $\pm \infty$  and NaN. E.g.,

```
> let a = 1.0/0.0
- let b = 0.0/0.0;;

val a : float = infinity
```

a	b	$a \cdot b$	a+b	$\bar{a}$
0	0	0	0	1
0	1	0	1	1
1	0	0	1	0
1	1	1	1	0

Table 5.3: Truth table for boolean 'and', 'or', and 'not' operators. Value 0 is false and 1 is true.

```
val b : float = nan
Listing 5.27: fsharpi
```

However, the **float** type has limite precision, since there is only a finite number of numbers that can be stored in a float. E.g.,

```
> let a = 357.8
- let b = a+0.1
- let c = b+0.1
- let d = c - 358.0;;

val a : float = 357.8
val b : float = 357.9
val c : float = 358.0
val d : float = 5.684341886e-14
Listing 5.28: fsharpi
```

Hence, although c appears to be correctly calculated, by the subtraction we see, that the value bound in c is not exactly the same as 358.0, and the reason is that the neither 357.8 nor 0.1 are exactly representable as a float, which is why the repeated addition accumulates a small representation error.

#### 5.2 Booleans

Boolean arithmatic is the basis of almost all computers and particularly important for controling program flow, which will be discussed in Chapter ??. Boolean values are one of 2 possible values, true or false, which is also sometimes written as 1 and 0. Two basic operations on boolean values are 'and' often also written as multiplication, and 'or' often written as addition, and 'not' often written as a bar above the value. All possible combination of input on these values can be written on tabular form, known as a truth table, shown in Table 5.3. That is, the multiplication and addition are good mnemonics for remembering the result of the 'and' and 'or' operators. In F# the values true and false are used, and the operators && for 'and',  $|\cdot|$  for 'or', and the function not for 'not', such that the above table is reproduced by,

```
\cdot and \cdot or \cdot not
```

· truth table

```
> let t = true
 let f
          false
          "%A %A %A %A %A"
                            f f (f && f)
                                          (f || f) (not f)
          "%A %A
                   %A %A
                          % A ''
                               f t (f && t) (f || t) (not f)
          " % A
                %A %A %A
                          %A" t f (t && f) (t || f) (not t)
- printfn
                        % A
                    % A
                             %A" t t (t && t) (t || t) (not t);;
          " % A
                % A
             a*b
                   a+b
                         not a
false false false
false true
            false true
                         true
      false false true
```

```
true true true false

val t : bool = true
val f : bool = false
val it : unit = ()
Listing 5.29: fsharpi
```

Careful spacing in the format string of the printfn function was used to align columns. Next section will discuss more elegant formatting options.

#### 5.3 Chars and Strings

#### 5.4 Mutable Data

The most common syntax for a value definition is

```
( "let'' [ "mutable" ] ident [":" type] "=" expr "in" expr
```

or alternatively

```
"let" ["mutable"] ident [":" type] "=" expr ["in"] expr
```

In the above, ident may be replaced with a more complicated pattern, but this is outside the scope of this text. If a value has been defined as mutable, then it's value may be changed using the following syntax,

```
expr "<-" expr
```

*Mutable data* is synonymous with the term *variable*. A variable an area in the computers working memory associated with a name and a type, and this area may be read from and written to during program execution. For example,

· Mutable data · variable

```
let mutable x = Unchecked.defaultof<int> // Declare a variable x of
   type int and assign the corresponding default value to it.
printfn "%d" x
x <- 5 // Assign a new value 5 to x
printfn "%d" x
x <- -3 // Assign a new value -3 to x
printfn "%d" x</pre>
Using 5.30: mutableAssignReassing.fsx -
```

Here a area in memory was denoted x, declared as type integer and assigned a default value. Later, this value of of x was replaced with another integer and yet another integer. The operator '|<-|' is used to distinguish the statement from the mathematical concept of equality. A short-hand for the above is available as,

```
let mutable x = 5 // Declare a variable x and assign the value 5 to
   it
printfn "%d" x
x <- -3 // Assign a new value -3 to x
printfn "%d" x

5
-3
Listing 5.31: mutableAssignReassingShort.fsx -</pre>
```

where the assignment of the default value was skipped, and the type was inferred from the assignment operation. However, it's important to note, that when the variable is declared, then the '|=|' operator must be used, while later reassignments must use the '|<-|' operator. Type mismatches will result in an error,

```
let mutable x = 5 // Declare a variable x and assign the value 5 to
   it
printfn "%d" x
x <- -3.0 // This is illegal, since -3.0 is a float, while x is of
   type int
printfn "%d" x

/Users/sporring/repositories/fsharpNotes/
   mutableAssignReassingTypeError.fsx(3,6): error FS0001: This
   expression was expected to have type
   int
but here has type
   float
Listing 5.32: mutableAssignReassingTypeError.fsx -</pre>
```

A typical variable is a counter of type integer, and a typical use of counters is to increment them, i.e., erasing a new value to be one more that its previous value. For example,

```
let mutable x = 5 // Declare a variable x and assign the value 5 to
   it
printfn "%d" x
x <- x + 1 // Assign a new value -3 to x
printfn "%d" x

5
6
Listing 5.33: mutableAssignIncrement.fsx -</pre>
```

An function that elegantly implements the incrementation operation may be constructed as,

```
let incr =
  let mutable counter = 0
  fun () ->
```

```
counter <- counter + 1
    counter
printfn "%d" (incr ())
printfn "%d" (incr ())

printfn "%d" (incr ())

1
2
3
Listing 5.34: mutableAssignIncrementEncapsulation.fsx -</pre>
```

<sup>1</sup> Here the output of incr is an anonymous function, that takes no argument, increments the variable of incr and returns the new value of the counter. This construction is called *encapsulation*, since the variable counter is hidden by the function incr from the user, i.e., the user need not be concerned with how the increment operator is implemented and the variable name used by incr does not clutter the scope where it is used.

 $\cdot \ encapsulation$ 

Variables cannot be returned from functions, that is,

```
let g () =
   let x = 0
   x
printfn "%d" (g ())

Clisting 5.35: mutableAssignReturnValue.fsx -
```

declares a function that has no arguments and returns the value 0, while the same for a variable is illegal,

```
let g () =
  let mutual x = 0
  x
printfn "%d" (g ())

/Users/sporring/repositories/fsharpNotes/
  mutableAssignReturnVariable.fsx(3,3): error FS0039: The value
  or constructor 'x' is not defined
Listing 5.36: mutableAssignReturnVariable.fsx -
```

There is a workaround for this by using reference cells by the build-in function ref and operators |!| · reference cells and |:=|,

```
let g () =
  let x = ref 0
  x
let y = g ()
  printfn "%d" !y
  y := 3
```

<sup>&</sup>lt;sup>1</sup>Explain why this works!

```
printfn "%d" !y

0
3
Listing 5.37: mutableAssignReturnRefCell.fsx -
```

That is, the ref function creates a reference variable, the '|!|' and the '|:=|' operators reads and writes its value. Reference cells are in some language called pointers, and their use is strongly discouraged, since they may cause *side-effects*, which is the effect that one function changes the state of another, such as the following example demonstrates,

 $\cdot$  side-effects

```
let updateFactor factor =
  factor := 2

let multiplyWithFactor x =
  let a = ref 1
  updateFactor a
  !a * x

printfn "%d" (multiplyWithFactor 3)

6
Listing 5.38: mutableAssignReturnSideEffect.fsx -
```

In the example, the function updateFactor changes a variable in the scope of multiplyWithFactor, which is prone to errors, since the style of programming does not follow the usual assignment syntax. Better style of programming is,

```
let updateFactor () =
   2

let multiplyWithFactor x =
   let a = ref 1
   a := updateFactor ()
   !a * x

printfn "%d" (multiplyWithFactor 3)

6
Listing 5.39: mutableAssignReturnWithoutSideEffect.fsx -
```

Here there can be no doubt in multiplyWithFactor that the value of 'a' is changing. Side-effects do have their use, but should in general be avoided at almost all costs, and in general it is advised to refrain from using ref cells.

# Part II Imperative programming

# Part III Declarative programming

# Part IV Structured programming

# $\begin{array}{c} {\rm Part~V} \\ {\rm Appendix} \end{array}$

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