

Contents

1 Misc	1		
1.1 Contest	1	4.2.1 Matroid Intersection	14
1.1.1 Makefile	1	4.2.2 De Bruijn Sequence	14
1.2 How Did We Get Here?	1	4.2.3 Multinomial	15
1.2.1 Macros	1	4.3 Algebra	15
1.2.2 Fast I/O	2	4.3.1 Formal Power Series	15
1.2.3 constexpr	2	4.4 Theorems	16
1.2.4 Bump Allocator	2	4.4.1 Kirchhoff's Theorem	16
1.3 Tools	2	4.4.2 Tutte's Matrix	16
1.3.1 Floating Point Binary Search	2	4.4.3 Cayley's Formula	16
1.3.2 SplitMix64	2	4.4.4 Erdős–Gallai Theorem	16
1.3.3 <random>	3	4.4.5 Burnside's Lemma	16
1.3.4 x86 Stack Hack	3	5 Numeric	16
1.4 Algorithms	3	5.1 Barrett Reduction	16
1.4.1 Bit Hacks	3	5.2 Long Long Multiplication	16
1.4.2 Aliens Trick	3	5.3 Fast Fourier Transform	16
1.4.3 Hilbert Curve	3	5.4 Fast Walsh-Hadamard Transform	17
1.4.4 Infinite Grid Knight Distance	3	5.5 Subset Convolution	17
1.4.5 Poker Hand	3	5.6 Linear Recurrences	17
1.4.6 Longest Increasing Subsequence	3	5.6.1 Berlekamp-Massey Algorithm	17
1.4.7 Mo's Algorithm on Tree	3	5.6.2 Linear Recurrence Calculation	17
2 Data Structures	4	5.7 Matrices	17
2.1 GNU PBDS	4	5.7.1 Determinant	17
2.2 Segment Tree (ZKW)	4	5.7.2 Inverse	17
2.3 Line Container	4	5.7.3 Characteristic Polynomial	18
2.4 Li-Chao Tree	4	5.7.4 Solve Linear Equation	18
2.5 Heavy-Light Decomposition	4	5.8 Polynomial Interpolation	18
2.6 Wavelet Matrix	5	5.9 Simplex Algorithm	19
2.7 Link-Cut Tree	5	6 Geometry	19
3 Graph	6	6.1 Point	19
3.1 Modeling	6	6.1.1 Quarternion	20
3.2 Matching/Flows	6	6.1.2 Spherical Coordinates	20
3.2.1 Dinic's Algorithm	6	6.2 Segments	20
3.2.2 Minimum Cost Flow	7	6.3 Convex Hull	20
3.2.3 Gomory-Hu Tree	7	6.3.1 3D Hull	20
3.2.4 Global Minimum Cut	7	6.4 Angular Sort	21
3.2.5 Bipartite Minimum Cover	7	6.5 Convex Polygon Minkowski Sum	21
3.2.6 Edmonds' Algorithm	8	6.6 Point In Polygon	21
3.2.7 Minimum Weight Matching	8	6.6.1 Convex Version	21
3.2.8 Stable Marriage	8	6.6.2 Offline Multiple Points Version	21
3.2.9 Kuhn-Munkres algorithm	9	6.7 Closest Pair	22
3.3 Shortest Path Faster Algorithm	9	6.8 Minimum Enclosing Circle	22
3.4 Strongly Connected Components	10	6.9 Delaunay Triangulation	22
3.4.1 2-Satisfiability	10	6.9.1 Slower Version	23
3.5 Biconnected Components	10	6.10 Half Plane Intersection	23
3.5.1 Articulation Points	10	7 Strings	23
3.5.2 Bridges	10	7.1 Knuth-Morris-Pratt Algorithm	23
3.6 Triconnected Components	10	7.2 Aho-Corasick Automaton	23
3.7 Centroid Decomposition	11	7.3 Suffix Array	24
3.8 Minimum Mean Cycle	11	7.4 Suffix Tree	24
3.9 Directed MST	11	7.5 Cocke-Younger-Kasami Algorithm	24
3.10 Maximum Clique	11	7.6 Z Value	25
3.11 Dominator Tree	12	7.7 Manacher's Algorithm	25
3.12 Manhattan Distance MST	12	7.8 Minimum Rotation	25
4 Math	13	7.9 Palindromic Tree	25
4.1 Number Theory	13	8 Debug List	25
4.1.1 Mod Struct	13	1. Misc	
4.1.2 Miller-Rabin	13	1.1. Contest	
4.1.3 Linear Sieve	13	1.1.1. Makefile	
4.1.4 Get Factors	13		
4.1.5 Binary GCD	13		
4.1.6 Extended GCD	13		
4.1.7 Chinese Remainder Theorem	13		
4.1.8 Baby-Step Giant-Step	13		
4.1.9 Pollard's Rho	14		
4.1.10 Tonelli-Shanks Algorithm	14		
4.1.11 Chinese Sieve	14		
4.1.12 Rational Number Binary Search	14		
4.1.13 Farey Sequence	14		
4.2 Combinatorics	14		

```

1 .PRECIOUS: ./p%
3 %: p%
4 ulimit -s unlimited && ./p%
5 p%: p%.cpp
7 g++ -o $@ $< -std=c++17 -Wall -Wextra -Wshadow \
  -fsanitize=address,undefined

```

1.2. How Did We Get Here?

1.2.1. Macros

Use vectorizations and math optimizations at your own peril.
For gcc \geq 9, there are `[[likely]]` and `[[unlikely]]` attributes.

Call gcc with `-fopt-info-optimized-missed-optall` for optimization info.

```
1 #define _GLIBCXX_DEBUG 1 // for debug mode
2 #define _GLIBCXX_SANITIZE_VECTOR 1 // for asan on vectors
3 #pragma GCC optimize("O3", "unroll-loops")
4 #pragma GCC optimize("fast-math")
5 #pragma GCC target("avx,avx2,abm,bmi,bmi2") // tip: `lscpu`
6 // before a loop
7 #pragma GCC unroll 16 // 0 or 1 -> no unrolling
8 #pragma GCC ivdep
```

1.2.2. Fast I/O

```
1 struct scanner {
2     static constexpr size_t LEN = 32 << 20;
3     char *buf, *buf_ptr, *buf_end;
4     scanner()
5         : buf(new char[LEN]), buf_ptr(buf + LEN),
6           buf_end(buf + LEN) {}
7     ~scanner() { delete[] buf; }
8     char getc() {
9         if (buf_ptr == buf_end) [[unlikely]]
10             buf_end = buf + fread_unlocked(buf, 1, LEN, stdin),
11             buf_ptr = buf;
12         return *(buf_ptr++);
13     }
14     char seek(char del) {
15         char c;
16         while ((c = getc()) < del) {}
17         return c;
18     }
19     void read(int &t) {
20         bool neg = false;
21         char c = seek('-');
22         if (c == '-') neg = true, t = 0;
23         else t = c ^ '0';
24         while ((c = getc()) >= '0') t = t * 10 + (c ^ '0');
25         if (neg) t = -t;
26     }
27 };
28 struct printer {
29     static constexpr size_t CPI = 21, LEN = 32 << 20;
30     char *buf, *buf_ptr, *buf_end, *tbuf;
31     char *int_buf, *int_buf_end;
32     printer()
33         : buf(new char[LEN]), buf_ptr(buf),
34           buf_end(buf + LEN), int_buf(new char[CPI + 1]),
35           int_buf_end(int_buf + CPI + 1) {}
36     ~printer() {
37         flush();
38         delete[] buf, delete[] int_buf;
39     }
40     void flush() {
41         fwrite_unlocked(buf, 1, buf_ptr - buf, stdout);
42         buf_ptr = buf;
43     }
44     void write(const char &c) {
45         *buf_ptr = c;
46         if (++buf_ptr == buf_end) [[unlikely]]
47             flush();
48     }
49     void write(const char *s) {
50         for (; *s != '\0'; ++s) write(*s);
51     }
52     void write(int x) {
53         if (x < 0) write('-', x = -x);
54         if (x == 0) [[unlikely]]
55             return write('0');
56         for (tbuf = int_buf_end; x != 0; --tbuf, x /= 10)
57             *tbuf = '0' + char(x % 10);
58         write(++tbuf);
59     }
60 };
```

Kotlin

```
1 import java.io.*
2 import java.util.*
3
4 @JvmField val cin = System.`in`.bufferedReader()
5 @JvmField val cout = PrintWriter(System.out, false)
6 @JvmField var tokenizer: StringTokenizer = StringTokenizer("")
7 fun nextLine() = cin.readLine()!!
8 fun read(): String {
9     while(!tokenizer.hasMoreTokens())
10         tokenizer = StringTokenizer(nextLine())
11     return tokenizer.nextToken()
12 }
```

```
// example
15 fun main() {
16     val n = read().toInt()
17     val a = DoubleArray(n) { read().toDouble() }
18     cout.println("omg hi")
19     cout.flush()
20 }
```

1.2.3. constexpr

Some default limits in gcc (7.x - trunk):

- constexpr recursion depth: 512
- constexpr loop iteration per function: 262144
- constexpr operation count per function: 33554432
- template recursion depth: 900 (gcc *might* segfault first)

```
1 constexpr array<int, 10> fibonacci{[] {
2     array<int, 10> a{};
3     a[0] = a[1] = 1;
4     for (int i = 2; i < 10; i++) a[i] = a[i - 1] + a[i - 2];
5     return a;
6 }}();
7 static_assert(fibonacci[9] == 55, "CE");
8
9 template <typename F, typename INT, INT... S>
10 constexpr void for_constexpr(integer_sequence<INT, S...>,
11                             F &&func) {
12     int _[] = {(func(integral_constant<INT, S>{}), 0)...};
13 }
14 // example
15 template <typename... T> void print_tuple(tuple<T...> t) {
16     for_constexpr(make_index_sequence<sizeof...(T)>{}),
17                 [&](auto i) { cout << get<i>(t) << '\n'; }};
```

1.2.4. Bump Allocator

```
1 // global bump allocator
2
3 char mem[256 << 20]; // 256 MB
4 size_t rsp = sizeof mem;
5 void *operator new(size_t s) {
6     assert(s < rsp); // MLE
7     return (void *)&mem[rsp -= s];
8 }
9 void operator delete(void *) {}
10
11 // bump allocator for STL / pbds containers
12 char mem[256 << 20];
13 size_t rsp = sizeof mem;
14 template <typename T> struct bump {
15     typedef T value_type;
16     bump() {}
17     template <typename U> bump(U, ...) {}
18     T *allocate(size_t n) {
19         rsp -= n * sizeof(T);
20         rsp &= 0 - alignof(T);
21         return (T *)(&mem + rsp);
22     }
23     void deallocate(T *, size_t n) {}
24 };
```

1.3. Tools

1.3.1. Floating Point Binary Search

```
1 union di {
2     double d;
3     ull i;
4 };
5 bool check(double);
6 // binary search in [L, R) with relative error 2^-eps
7 double binary_search(double L, double R, int eps) {
8     di l = {L}, r = {R}, m;
9     while (r.i - l.i > 1LL << (52 - eps)) {
10         m.i = (l.i + r.i) >> 1;
11         if (check(m.d)) r = m;
12         else l = m;
13     }
14     return l.d;
15 }
```

1.3.2. SplitMix64

```
1 using ull = unsigned long long;
2 inline ull splitmix64(ull x) {
3     // change to `static ull x = SEED;` for DRBG
4     ull z = (x += 0x9E3779B97F4A7C15);
5     z = (z ^ (z >> 30)) * 0xBF58476D1CE4E5B9;
6     z = (z ^ (z >> 27)) * 0x94D049B133111EB;
7     return z ^ (z >> 31);
8 }
```

1.3.3. <random>

```
1 #ifdef __unix__
  random_device rd;
  mt19937_64 RNG(rd());
2 #else
  const auto SEED = chrono::high_resolution_clock::now()
    .time_since_epoch()
    .count();
  mt19937_64 RNG(SEED);
3 #endif
  // random uint_fast64_t: RNG();
  // uniform random of type T (int, double, ...) in [l, r]:
  // uniform_int_distribution<T> dist(l, r); dist(RNG);
```

1.3.4. x86 Stack Hack

```
1 constexpr size_t size = 200 << 20; // 200MiB
2 int main() {
  register long rsp asm("rsp");
  char *buf = new char[size];
  asm("movq %0, %%rsp\n" :: "r"(buf + size));
  // do stuff
  asm("movq %0, %%rsp\n" :: "r"(rsp));
  delete[] buf;
3 }
```

1.4. Algorithms

1.4.1. Bit Hacks

```
1 // next permutation of x as a bit sequence
  ull next_bits_permutation(ull x) {
  2 ull c = __builtin_ctzll(x), r = x + (1ULL << c);
    return (r ^ x) >> (c + 2) | r;
  3 }
  // iterate over all (proper) subsets of bitset s
  void subsets(ull s) {
  4 for (ull x = s; x;) { --x &= s; /* do stuff */ }
  5 }
```

1.4.2. Aliens Trick

```
1 // min dp[i] value and its i (smallest one)
  pll get_dp(int cost);
  ll aliens(int k, int l, int r) {
  2 while (l != r) {
    int m = (l + r) / 2;
    auto [f, s] = get_dp(m);
    if (s == k) return f - m * k;
    if (s < k) r = m;
    else l = m + 1;
  3 }
  return get_dp(l).first - l * k;
  4 }
```

1.4.3. Hilbert Curve

```
1 ll hilbert(ll n, int x, int y) {
  ll res = 0;
  for (ll s = n; s /= 2;) {
  2 int rx = !(x & s), ry = !(y & s);
    res += s * s * ((3 * rx) ^ ry);
    if (ry == 0) {
  3 if (rx == 1) x = s - 1 - x, y = s - 1 - y;
      swap(x, y);
    }
  4 }
  return res;
  5 }
```

1.4.4. Infinite Grid Knight Distance

```
1 ll get_dist(ll dx, ll dy) {
  if (++(dx = abs(dx)) > ++(dy = abs(dy))) swap(dx, dy);
  2 if (dx == 1 && dy == 2) return 3;
  if (dx == 3 && dy == 3) return 4;
  ll lb = max(dy / 2, (dx + dy) / 3);
  return ((dx ^ dy ^ lb) & 1) ? ++lb : lb;
  3 }
```

1.4.5. Poker Hand

```
1
2
3
4
5
6
7 using namespace std;
```

```
9 struct hand {
  static constexpr auto rk = [] {
  10 array<int, 256> x{};
    auto s = "23456789TJQKACDHS";
    for (int i = 0; i < 17; i++) x[s[i]] = i % 13;
    return x;
  11 }();
  vector<pair<int, int>> v;
  vector<int> cnt, vf, vs;
  int type;
  hand() : cnt(4), type(0) {}
  void add_card(char suit, char rank) {
  12 ++cnt[rk[suit]];
    for (auto &[f, s] : v)
  13 if (s == rk[rank]) return ++f, void();
    v.emplace_back(1, rk[rank]);
  14 }
  void process() {
  15 sort(v.rbegin(), v.rend());
    for (auto [f, s] : v) vf.push_back(f), vs.push_back(s);
  16 bool str = 0, flu = find(all(cnt), 5) != cnt.end();
    if ((str = v.size() == 5))
  17 for (int i = 1; i < 5; i++)
    if (vs[i] != vs[i - 1] + 1) str = 0;
  18 if (vs == vector<int>{12, 3, 2, 1, 0})
    str = 1, vs = {3, 2, 1, 0, -1};
  19 if (str && flu) type = 9;
    else if (vf[0] == 4) type = 8;
  20 else if (vf[0] == 3 && vf[1] == 2) type = 7;
    else if (str || flu) type = 5 + flu;
  21 else if (vf[0] == 3) type = 4;
    else if (vf[0] == 2) type = 2 + (vf[1] == 2);
  22 else type = 1;
  23 }
  bool operator<(const hand &b) const {
  24 return make_tuple(type, vf, vs) <
    make_tuple(b.type, b.vf, b.vs);
  25 }
  26 };
```

1.4.6. Longest Increasing Subsequence

```
1
2
3 template <class I> vi lis(const vector<I> &S) {
  if (S.empty()) return {};
  4 vi prev(sz(S));
  typedef pair<I, int> p;
  vector<p> res;
  rep(i, 0, sz(S)) {
  5 // change 0 -> i for longest non-decreasing subsequence
    auto it = lower_bound(all(res), p[S[i], 0]);
    if (it == res.end())
  6 res.emplace_back(), it = res.end() - 1;
    *it = {S[i], i};
    prev[i] = it == res.begin() ? 0 : (it - 1)->second;
  7 }
  int L = sz(res), cur = res.back().second;
  vi ans(L);
  while (L--) ans[L] = cur, cur = prev[cur];
  8 return ans;
  9 }
```

1.4.7. Mo's Algorithm on Tree

```
1 void MoAlgoOnTree() {
  Dfs(0, -1);
  vector<int> euler(tk);
  for (int i = 0; i < n; ++i) {
  2 euler[tin[i]] = i;
    euler[tout[i]] = i;
  3 }
  vector<int> l(q), r(q), qr(q), sp(q, -1);
  for (int i = 0; i < q; ++i) {
  4 if (tin[u[i]] > tin[v[i]]) swap(u[i], v[i]);
    int z = GetLCA(u[i], v[i]);
    sp[i] = z[i];
  5 if (z == u) l[i] = tin[u[i]], r[i] = tin[v[i]];
    else l[i] = tout[u[i]], r[i] = tin[v[i]];
    qr[i] = i;
  6 }
  sort(qr.begin(), qr.end(), [&](int i, int j) {
  7 if (l[i] / kB == l[j] / kB) return r[i] < r[j];
    return l[i] / kB < l[j] / kB;
  8 });
  vector<bool> used(n);
  // Add(v): add/remove v to/from the path based on used[v]
  for (int i = 0, tl = 0, tr = -1; i < q; ++i) {
  9 while (tl < l[qr[i]]) Add(euler[tl++]);
    while (tl > l[qr[i]]) Add(euler[--tl]);
    while (tr > r[qr[i]]) Add(euler[tr--]);
  10 }
```

```

27     while (tr < r[qr[i]]) Add(euler[++tr]);
28     // add/remove LCA(u, v) if necessary
29 }

```

2. Data Structures

2.1. GNU PBDS

```

1 #include <ext/pb_ds/assoc_container.hpp>
2 #include <ext/pb_ds/priority_queue.hpp>
3 #include <ext/pb_ds/tree_policy.hpp>
4 using namespace __gnu_pbds;
5
6 // most std::map + order_of_key, find_by_order, split, join
7 template <typename T, typename U = null_type>
8 using ordered_map = tree<T, U, std::less<>, rb_tree_tag,
9     tree_order_statistics_node_update>;
10 // useful tags: rb_tree_tag, splay_tree_tag
11
12 template <typename T> struct myhash {
13     size_t operator()(T x) const; // splitmix, bswap(x*R), ...
14 };
15 // most of std::unordered_map, but faster (needs good hash)
16 template <typename T, typename U = null_type>
17 using hash_table = gp_hash_table<T, U, myhash<T>>;
18
19 // most std::priority_queue + modify, erase, split, join
20 using heap = priority_queue<int, std::less<>>;
21 // useful tags: pairing_heap_tag, binary_heap_tag,
22 //             (rc_)?binomial_heap_tag, thin_heap_tag

```

2.2. Segment Tree (ZKW)

```

1 struct segtree {
2     using T = int;
3     T f(T a, T b) { return a + b; } // any monoid operation
4     static constexpr T ID = 0; // identity element
5     int n;
6     vector<T> v;
7     segtree(int n_) : n(n_), v(2 * n, ID) {}
8     segtree(vector<T> &a) : n(a.size()), v(2 * n, ID) {
9         copy_n(a.begin(), n, v.begin() + n);
10        for (int i = n - 1; i > 0; i--)
11            v[i] = f(v[i * 2], v[i * 2 + 1]);
12    }
13    void update(int i, T x) {
14        for (v[i += n] = x; i /= 2;)
15            v[i] = f(v[i * 2], v[i * 2 + 1]);
16    }
17    T query(int l, int r) {
18        T tl = ID, tr = ID;
19        for (l += n, r += n; l < r; l /= 2, r /= 2) {
20            if (l & 1) tl = f(tl, v[l++]);
21            if (r & 1) tr = f(v[--r], tr);
22        }
23        return f(tl, tr);
24    }
25 };

```

2.3. Line Container

```

1
2
3 struct Line {
4     mutable ll k, m, p;
5     bool operator<(const Line &o) const { return k < o.k; }
6     bool operator<(ll x) const { return p < x; }
7 };
8 // add: line y=kx+m, query: maximum y of given x
9 struct LineContainer : multiset<Line, less<>> {
10     // (for doubles, use inf = 1/.0, div(a,b) = a/b)
11     static const ll inf = LLONG_MAX;
12     ll div(ll a, ll b) { // floored division
13         return a / b - ((a ^ b) < 0 && a % b);
14     }
15     bool isect(iterator x, iterator y) {
16         if (y == end()) return x->p = inf, 0;
17         if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
18         else x->p = div(y->m - x->m, x->k - y->k);
19         return x->p >= y->p;
20     }
21     void add(ll k, ll m) {
22         auto z = insert({k, m, 0}), y = z++, x = y;
23         while (isect(y, z)) z = erase(z);
24         if (x != begin() && isect(--x, y))
25             isect(x, y = erase(y));
26         while ((y = x) != begin() && (--x)->p >= y->p)
27             isect(x, erase(y));
28     }
29 };

```

```

29 ll query(ll x) {
30     assert(!empty());
31     auto l = *lower_bound(x);
32     return l.k * x + l.m;
33 }
34 };

```

2.4. Li-Chao Tree

```

1 constexpr ll MAXN = 2e5, INF = 2e18;
2 struct Line {
3     ll m, b;
4     Line() : m(0), b(-INF) {}
5     Line(ll _m, ll _b) : m(_m), b(_b) {}
6     ll operator()(ll x) const { return m * x + b; }
7 };
8 struct LiChao {
9     Line a[MAXN * 4];
10    void insert(Line seg, int l, int r, int v = 1) {
11        if (l == r) {
12            if (seg(l) > a[v](l)) a[v] = seg;
13            return;
14        }
15        int mid = (l + r) >> 1;
16        if (a[v].m > seg.m) swap(a[v], seg);
17        if (a[v](mid) < seg(mid)) {
18            swap(a[v], seg);
19            insert(seg, l, mid, v << 1);
20        } else insert(seg, mid + 1, r, v << 1 | 1);
21    }
22    ll query(int x, int l, int r, int v = 1) {
23        if (l == r) return a[v](x);
24        int mid = (l + r) >> 1;
25        if (x <= mid)
26            return max(a[v](x), query(x, l, mid, v << 1));
27        else
28            return max(a[v](x), query(x, mid + 1, r, v << 1 | 1));
29    }
30 };

```

2.5. Heavy-Light Decomposition

```

1
2
3 template <bool VALS_EDGES> struct HLD {
4     int N, tim = 0;
5     vector<vi> adj;
6     vi par, siz, depth, rt, pos;
7     Node *tree;
8     HLD(vector<vi> adj_)
9         : N(sz(adj_)), adj(adj_), par(N, -1), siz(N, 1),
10         depth(N), rt(N), pos(N), tree(new Node(0, N)) {}
11     dfsSz(0);
12     dfsHld(0);
13 }
14 void dfsSz(int v) {
15     if (par[v] != -1)
16         adj[v].erase(find(all(adj[v]), par[v]));
17     for (int &u : adj[v]) {
18         par[u] = v, depth[u] = depth[v] + 1;
19         dfsSz(u);
20         siz[v] += siz[u];
21         if (siz[u] > siz[adj[v][0]]) swap(u, adj[v][0]);
22     }
23 }
24 void dfsHld(int v) {
25     pos[v] = tim++;
26     for (int u : adj[v]) {
27         rt[u] = (u == adj[v][0] ? rt[v] : u);
28         dfsHld(u);
29     }
30 }
31 template <class B> void process(int u, int v, B op) {
32     for (; rt[u] != rt[v]; v = par[rt[v]]) {
33         if (depth[rt[u]] > depth[rt[v]]) swap(u, v);
34         op(pos[rt[v]], pos[v] + 1);
35     }
36     if (depth[u] > depth[v]) swap(u, v);
37     op(pos[u] + VALS_EDGES, pos[v] + 1);
38 }
39 void modifyPath(int u, int v, int val) {
40     process(u, v, [&](int l, int r) { tree->add(l, r, val); });
41 }
42 int queryPath(int u,
43     int v) { // Modify depending on problem
44     int res = -1e9;
45     process(u, v, [&](int l, int r) {
46         res = max(res, tree->query(l, r));
47     });
48     return res;
49 };

```

```

}
51 int querySubtree(int v) { // modifySubtree is similar
    return tree->query(pos[v] + VALS_EDGES,
53         pos[v] + siz[v]);
}
55 };

```

2.6. Wavelet Matrix

```

1
3
#pragma GCC target("popcnt,bmi2")
#include <immintrin.h>

7 // T is unsigned. You might want to compress values first
template <typename T> struct wavelet_matrix {
9     static_assert(is_unsigned_v<T>, "only unsigned T");
    struct bit_vector {
11         static constexpr uint W = 64;
        uint n, cnt0;
        vector<ull> bits;
        vector<uint> sum;
        bit_vector(uint n_)
13             : n(n_), bits(n / W + 1), sum(n / W + 1) {}
        void build() {
15             for (uint j = 0; j != n / W; ++j)
                sum[j + 1] = sum[j] + _mm_popcnt_u64(bits[j]);
                cnt0 = rank0(n);
21         }
        void set_bit(uint i) { bits[i / W] |= 1ULL << i % W; }
23         bool operator[](uint i) const {
            return !(bits[i / W] & 1ULL << i % W);
25         }
        uint rank1(uint i) const {
27             return sum[i / W] +
                _mm_popcnt_u64(_bzh_u64(bits[i / W], i % W));
29         }
        uint rank0(uint i) const { return i - rank1(i); }
31     };
    uint n, lg;
    vector<bit_vector> b;
    wavelet_matrix(const vector<T> &a) : n(a.size()) {
33         lg =
            lg(max(*max_element(a.begin(), a.end()), T(1))) + 1;
35         b.assign(lg, n);
        vector<T> cur = a, nxt(n);
        for (int h = lg; h--;) {
37             for (uint i = 0; i < n; ++i)
                if (cur[i] & (T(1) << h)) b[h].set_bit(i);
                b[h].build();
                int il = 0, ir = b[h].cnt0;
                for (uint i = 0; i < n; ++i)
41                 nxt[(b[h][i] ? ir : il)++] = cur[i];
                swap(cur, nxt);
43             }
        }
45     }
    T operator[](uint i) const {
        T res = 0;
51         for (int h = lg; h--;)
            if (b[h][i])
                i += b[h].cnt0 - b[h].rank0(i), res |= T(1) << h;
                else i = b[h].rank0(i);
53         return res;
55     }
    // query k-th smallest (0-based) in a[l, r]
    T kth(uint l, uint r, uint k) const {
57         T res = 0;
        for (int h = lg; h--;) {
            uint tl = b[h].rank0(l), tr = b[h].rank0(r);
            if (k >= tr - tl) {
61                 k -= tr - tl;
                l += b[h].cnt0 - tl;
                r += b[h].cnt0 - tr;
                res |= T(1) << h;
                } else l = tl, r = tr;
63             }
        }
        return res;
65     }
    // count of i in [l, r] with a[i] < u
    uint count(uint l, uint r, T u) const {
71         if (u >= T(1) << lg) return r - l;
        uint res = 0;
        for (int h = lg; h--;) {
            uint tl = b[h].rank0(l), tr = b[h].rank0(r);
            if (u & (T(1) << h)) {
73                 l += b[h].cnt0 - tl;
                r += b[h].cnt0 - tr;
                res += tr - tl;
                } else l = tl, r = tr;
75             }
        }
        return res;
77     }
    // count of i in [l, r] with a[i] < u
    uint count(uint l, uint r, T u) const {
79         if (u >= T(1) << lg) return r - l;
        uint res = 0;
        for (int h = lg; h--;) {
            uint tl = b[h].rank0(l), tr = b[h].rank0(r);
            if (u & (T(1) << h)) {
81                 l += b[h].cnt0 - tl;
                r += b[h].cnt0 - tr;
                res += tr - tl;
                } else l = tl, r = tr;
            }
        }
        return res;
83     }

```

```

83     return res;
    }
85 };

```

2.7. Link-Cut Tree

```

1
3 const int MXN = 100005;
const int MEM = 100005;

5 struct Splay {
7     static Splay nil, mem[MEM], *pmem;
    Splay *ch[2], *f;
    int val, rev, size;
    Splay() : val(-1), rev(0), size(0) {
11         f = ch[0] = ch[1] = &nil;
    }
    Splay(int _val) : val(_val), rev(0), size(1) {
13         f = ch[0] = ch[1] = &nil;
    }
    bool isr() {
15         return f->ch[0] != this && f->ch[1] != this;
    }
    int dir() { return f->ch[0] == this ? 0 : 1; }
    void setCh(Splay *c, int d) {
21         ch[d] = c;
        if (c != &nil) c->f = this;
        pull();
23     }
    void push() {
25         if (rev) {
            swap(ch[0], ch[1]);
            if (ch[0] != &nil) ch[0]->rev ^= 1;
            if (ch[1] != &nil) ch[1]->rev ^= 1;
            rev = 0;
27         }
    }
    void pull() {
33         size = ch[0]->size + ch[1]->size + 1;
        if (ch[0] != &nil) ch[0]->f = this;
        if (ch[1] != &nil) ch[1]->f = this;
35     }
} Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::mem;
Splay *nil = &Splay::nil;

41 void rotate(Splay *x) {
    Splay *p = x->f;
    int d = x->dir();
    if (!p->isr()) p->f->setCh(x, p->dir());
    else x->f = p->f;
    p->setCh(x->ch[!d], d);
    x->setCh(p, !d);
    p->pull();
    x->pull();
43 }

45 vector<Splay*> splayVec;
void splay(Splay *x) {
    splayVec.clear();
    for (Splay *q = x; q = q->f) {
51         splayVec.push_back(q);
        if (q->isr()) break;
    }
    reverse(begin(splayVec), end(splayVec));
    for (auto it : splayVec) it->push();
    while (!x->isr()) {
61         if (x->f->isr()) rotate(x);
        else if (x->dir() == x->f->dir())
            rotate(x->f), rotate(x);
        else rotate(x), rotate(x);
        }
    }

63 Splay *access(Splay *x) {
    Splay *q = nil;
    for (; x != nil; x = x->f) {
71         splay(x);
        x->setCh(q, 1);
        q = x;
    }
    return q;
73 }

75 void evert(Splay *x) {
    access(x);
    splay(x);
    x->rev ^= 1;
    x->push();
    x->pull();
    }

81 void link(Splay *x, Splay *y) {
    access(x);
    splay(x);
    x->ch[1] = y;
    y->f = x;
    y->pull();
    }

```



```

87 // evert(x);
88 access(x);
89 splay(x);
90 evert(y);
91 x->setCh(y, 1);
92 }
93 void cut(Splay *x, Splay *y) {
94 // evert(x);
95 access(y);
96 splay(y);
97 y->push();
98 y->ch[0] = y->ch[0]->f = nil;
99 }
100 int N, Q;
101 Splay *vt[MXN];
102
103 int ask(Splay *x, Splay *y) {
104 access(x);
105 access(y);
106 splay(x);
107 int res = x->f->val;
108 if (res == -1) res = x->val;
109 return res;
110 }
111
112 int main(int argc, char **argv) {
113 scanf("%d%d", &N, &Q);
114 for (int i = 1; i <= N; i++)
115   vt[i] = new (Splay::pmem++) Splay(i);
116 while (Q--) {
117   char cmd[105];
118   int u, v;
119   scanf("%s", cmd);
120   if (cmd[1] == 'i') {
121     scanf("%d%d", &u, &v);
122     link(vt[v], vt[u]);
123   } else if (cmd[0] == 'c') {
124     scanf("%d", &v);
125     cut(vt[1], vt[v]);
126   } else {
127     scanf("%d%d", &u, &v);
128     int res = ask(vt[u], vt[v]);
129     printf("%d\n", res);
130   }
131 }

```

3. Graph

3.1. Modeling

- Maximum/Minimum flow with lower bound / Circulation problem
 - Construct super source S and sink T .
 - For each edge (x, y, l, u) , connect $x \rightarrow y$ with capacity $u - l$.
 - For each vertex v , denote by $in(v)$ the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - If $in(v) > 0$, connect $S \rightarrow v$ with capacity $in(v)$, otherwise, connect $v \rightarrow T$ with capacity $-in(v)$.
 - To maximize, connect $t \rightarrow s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T . If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T . Connect $t \rightarrow s$ with capacity ∞ and let the flow from S to T be f' . If $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
 - Redirect every edge: $y \rightarrow x$ if $(x, y) \in M$, $x \rightarrow y$ otherwise.
 - DFS from unmatched vertices in X .
 - $x \in X$ is chosen iff x is unvisited.
 - $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
 - Construct super source S and sink T .
 - For each edge (x, y, c) , connect $x \rightarrow y$ with $(cost, cap) = (c, 1)$ if $c > 0$, otherwise connect $y \rightarrow x$ with $(cost, cap) = (-c, 1)$.
 - For each edge with $c < 0$, sum these cost as K , then increase $d(y)$ by 1, decrease $d(x)$ by 1.
 - For each vertex v with $d(v) > 0$, connect $S \rightarrow v$ with $(cost, cap) = (0, d(v))$.
 - For each vertex v with $d(v) < 0$, connect $v \rightarrow T$ with $(cost, cap) = (0, -d(v))$.
 - Flow from S to T , the answer is the cost of the flow $C + K$.
- Maximum density induced subgraph
 - Binary search on answer, suppose we're checking answer T .
 - Construct a max flow model, let K be the sum of all weights.
 - Connect source $s \rightarrow v$, $v \in G$ with capacity K .

- For each edge (u, v, w) in G , connect $u \rightarrow v$ and $v \rightarrow u$ with capacity w .
 - For $v \in G$, connect it with sink $v \rightarrow t$ with capacity $K + 2T - (\sum_{e \in E(v)} w(e)) - 2w(v)$.
 - T is a valid answer if the maximum flow $f < K|V|$.
- Minimum weight edge cover
 - For each $v \in V$ create a copy v' , and connect $u' \rightarrow v'$ with weight $w(u, v)$.
 - Connect $v \rightarrow v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v .
 - Find the minimum weight perfect matching on G' .
 - Project selection problem
 - If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$.
 - Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v .
 - The mincut is equivalent to the maximum profit of a subset of projects.
 - 0/1 quadratic programming

$$\sum_x c_x x + \sum_y c_y \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y}')$$

can be minimized by the mincut of the following graph:

- Create edge (x, t) with capacity c_x and create edge (s, y) with capacity c_y .
- Create edge (x, y) with capacity c_{xy} .
- Create edge (x, y) and edge (x', y') with capacity $c_{xyx'y'}$.

3.2. Matching/Flows

3.2.1. Dinic's Algorithm

```

1 struct Dinic {
2   struct edge {
3     int to, cap, flow, rev;
4   };
5   static constexpr int MAXN = 1000, MAXF = 1e9;
6   vector<edge> v[MAXN];
7   int top[MAXN], deep[MAXN], side[MAXN], s, t;
8   void make_edge(int s, int t, int cap) {
9     v[s].push_back({t, cap, 0, (int)v[t].size()});
10    v[t].push_back({s, 0, 0, (int)v[s].size() - 1});
11  }
12  int dfs(int a, int flow) {
13    if (a == t || !flow) return flow;
14    for (int &i = top[a]; i < v[a].size(); i++) {
15      edge &e = v[a][i];
16      if (deep[a] + 1 == deep[e.to] && e.cap - e.flow) {
17        int x = dfs(e.to, min(e.cap - e.flow, flow));
18        if (x) {
19          e.flow += x, v[e.to][e.rev].flow -= x;
20          return x;
21        }
22      }
23    }
24    deep[a] = -1;
25    return 0;
26  }
27  bool bfs() {
28    queue<int> q;
29    fill_n(deep, MAXN, 0);
30    q.push(s), deep[s] = 1;
31    int tmp;
32    while (!q.empty()) {
33      tmp = q.front(), q.pop();
34      for (edge e : v[tmp])
35        if (!deep[e.to] && e.cap != e.flow)
36          deep[e.to] = deep[tmp] + 1, q.push(e.to);
37    }
38    return deep[t];
39  }
40  int max_flow(int _s, int _t) {
41    s = _s, t = _t;
42    int flow = 0, tflow;
43    while (bfs()) {
44      fill_n(top, MAXN, 0);
45      while ((tflow = dfs(s, MAXF))) flow += tflow;
46    }
47    return flow;
48  }
49  void reset() {
50    fill_n(side, MAXN, 0);
51    for (auto &i : v) i.clear();
52  }
53 };

```

3.2.2. Minimum Cost Flow

```

1 struct MCF {
2     struct edge {
3         ll to, from, cap, flow, cost, rev;
4     } * fromE[MAXN];
5     vector<edge> v[MAXN];
6     ll n, s, t, flows[MAXN], dis[MAXN], pi[MAXN], flowlim;
7     void make_edge(int s, int t, ll cap, ll cost) {
8         if (!cap) return;
9         v[s].pb(edge{t, s, cap, 0LL, cost, v[t].size()});
10        v[t].pb(edge{s, t, 0LL, 0LL, -cost, v[s].size() - 1});
11    }
12    bitset<MAXN> vis;
13    void dijkstra() {
14        vis.reset();
15        __gnu_pbds::priority_queue<pair<ll, int>> q;
16        vector<decltype(q)::point_iterator> its(n);
17        q.push({0LL, s});
18        while (!q.empty()) {
19            int now = q.top().second;
20            q.pop();
21            if (vis[now]) continue;
22            vis[now] = 1;
23            ll ndis = dis[now] + pi[now];
24            for (edge &e : v[now]) {
25                if (e.flow == e.cap || vis[e.to]) continue;
26                if (dis[e.to] > ndis + e.cost - pi[e.to]) {
27                    dis[e.to] = ndis + e.cost - pi[e.to];
28                    flows[e.to] = min(flows[now], e.cap - e.flow);
29                    fromE[e.to] = &e;
30                    if (its[e.to] == q.end())
31                        its[e.to] = q.push({-dis[e.to], e.to});
32                    else q.modify(its[e.to], {-dis[e.to], e.to});
33                }
34            }
35        }
36    }
37    bool AP(ll &flow) {
38        fill_n(dis, n, INF);
39        fromE[s] = 0;
40        dis[s] = 0;
41        flows[s] = flowlim - flow;
42        dijkstra();
43        if (dis[t] == INF) return false;
44        flow += flows[t];
45        for (edge *e = fromE[t]; e; e = fromE[e->from]) {
46            e->flow += flows[t];
47            v[e->to][e->rev].flow -= flows[t];
48        }
49        for (int i = 0; i < n; i++)
50            pi[i] = min(pi[i] + dis[i], INF);
51        return true;
52    }
53    pll solve(int _s, int _t, ll _flowlim = INF) {
54        s = _s, t = _t, flowlim = _flowlim;
55        pll re;
56        while (re.F != flowlim && AP(re.F))
57            ;
58        for (int i = 0; i < n; i++)
59            for (edge &e : v[i])
60                if (e.flow != 0) re.S += e.flow * e.cost;
61        re.S /= 2;
62        return re;
63    }
64    void init(int _n) {
65        n = _n;
66        fill_n(pi, n, 0);
67        for (int i = 0; i < n; i++) v[i].clear();
68    }
69    void setpi(int s) {
70        fill_n(pi, n, INF);
71        pi[s] = 0;
72        for (ll it = 0, flag = 1, tdis; flag && it < n; it++) {
73            flag = 0;
74            for (int i = 0; i < n; i++)
75                if (pi[i] != INF)
76                    for (edge &e : v[i])
77                        if (e.cap && (tdis = pi[i] + e.cost) < pi[e.to])
78                            pi[e.to] = tdis, flag = 1;
79        }
80    }
81 };

```

3.2.3. Gomory-Hu Tree

Requires: Dinic's Algorithm

```

1
2
3 int e[MAXN][MAXN];
4 int p[MAXN];

```

```

5 Dinic D; // original graph
6 void gomory_hu() {
7     fill(p, p + n, 0);
8     fill(e[0], e[n], INF);
9     for (int s = 1; s < n; s++) {
10         int t = p[s];
11         Dinic F = D;
12         int tmp = F.max_flow(s, t);
13         for (int i = 1; i < s; i++)
14             e[s][i] = e[i][s] = min(tmp, e[t][i]);
15         for (int i = s + 1; i <= n; i++)
16             if (p[i] == t && F.side[i]) p[i] = s;
17     }
18 }

```

3.2.4. Global Minimum Cut

```

1
2 // weights is an adjacency matrix, undirected
3 pair<int, vi> getMinCut(vector<vi> &weights) {
4     int N = sz(weights);
5     vi used(N), cut, best_cut;
6     int best_weight = -1;
7
8     for (int phase = N - 1; phase >= 0; phase--) {
9         vi w = weights[0], added = used;
10        int prev, k = 0;
11        rep(i, 0, phase) {
12            prev = k;
13            k = -1;
14            rep(j, 1, N) if (!added[j] &&
15                            (k == -1 || w[j] > w[k])) k = j;
16            if (i == phase - 1) {
17                rep(j, 0, N) weights[prev][j] += weights[k][j];
18                rep(j, 0, N) weights[j][prev] = weights[prev][j];
19                used[k] = true;
20                cut.push_back(k);
21                if (best_weight == -1 || w[k] < best_weight) {
22                    best_cut = cut;
23                    best_weight = w[k];
24                }
25            } else {
26                rep(j, 0, N) w[j] += weights[k][j];
27                added[k] = true;
28            }
29        }
30    }
31    return {best_weight, best_cut};
32 }

```

3.2.5. Bipartite Minimum Cover

Requires: Dinic's Algorithm

```

1
2 // maximum independent set = all vertices not covered
3 // x : [0, n], y : [0, m]
4 struct Bipartite_vertex_cover {
5     Dinic D;
6     int n, m, s, t, x[maxn], y[maxn];
7     void make_edge(int x, int y) { D.make_edge(x, y + n, 1); }
8     int matching() {
9         int re = D.max_flow(s, t);
10        for (int i = 0; i < n; i++)
11            for (Dinic::edge &e : D.v[i])
12                if (e.to != s && e.flow == 1) {
13                    x[i] = e.to - n, y[e.to - n] = i;
14                    break;
15                }
16        return re;
17    }
18    // init() and matching() before use
19    void solve(vector<int> &vx, vector<int> &vy) {
20        bitset<maxn * 2 + 10> vis;
21        queue<int> q;
22        for (int i = 0; i < n; i++)
23            if (x[i] == -1) q.push(i), vis[i] = 1;
24        while (!q.empty()) {
25            int now = q.front();
26            q.pop();
27            if (now < n) {
28                for (Dinic::edge &e : D.v[now])
29                    if (e.to != s && e.to - n != x[now] && !vis[e.to])
30                        vis[e.to] = 1, q.push(e.to);
31            } else {
32                if (!vis[y[now - n]])
33                    vis[y[now - n]] = 1, q.push(y[now - n]);
34            }
35        }
36    }
37 }

```

```

37     for (int i = 0; i < n; i++)
38         if (!vis[i]) vx.pb(i);
39     for (int i = 0; i < m; i++)
40         if (vis[i + n]) vy.pb(i);
41 }
42 void init(int _n, int _m) {
43     n = _n, m = _m, s = n + m, t = s + 1;
44     for (int i = 0; i < n; i++)
45         x[i] = -1, D.make_edge(s, i, 1);
46     for (int i = 0; i < m; i++)
47         y[i] = -1, D.make_edge(i + n, t, 1);
48 }
49 };

```

3.2.6. Edmonds' Algorithm

```

1 struct Edmonds {
2     int n, T;
3     vector<vector<int>> g;
4     vector<int> pa, p, used, base;
5     Edmonds(int n) : n(n), T(0), g(n), pa(n, -1), p(n), used(n),
6         base(n) {}
7     void add(int a, int b) {
8         g[a].push_back(b);
9         g[b].push_back(a);
10    }
11    int getBase(int i) {
12        while (i != base[i])
13            base[i] = base[base[i]], i = base[i];
14        return i;
15    }
16    vector<int> toJoin;
17    void mark_path(int v, int x, int b, vector<int> &path) {
18        for (; getBase(v) != b; v = p[x]) {
19            p[v] = x, x = pa[v];
20            toJoin.push_back(v);
21            toJoin.push_back(x);
22            if (!used[x]) used[x] = ++T, path.push_back(x);
23        }
24    }
25    bool go(int v) {
26        for (int x : g[v]) {
27            int b, bv = getBase(v), bx = getBase(x);
28            if (bv == bx) continue;
29            } else if (used[x]) {
30                vector<int> path;
31                toJoin.clear();
32                if (used[bx] < used[bv])
33                    mark_path(v, x, b = bx, path);
34                else mark_path(x, v, b = bv, path);
35                for (int z : toJoin) base[getBase(z)] = b;
36                for (int z : path)
37                    if (go(z)) return 1;
38            } else if (p[x] == -1) {
39                p[x] = v;
40                if (pa[x] == -1) {
41                    for (int y; x != -1; x = v)
42                        y = p[x], v = pa[y], pa[x] = y, pa[y] = x;
43                    return 1;
44                }
45                if (!used[pa[x]]) {
46                    used[pa[x]] = ++T;
47                    if (go(pa[x])) return 1;
48                }
49            }
50        }
51        return 0;
52    }
53    void init_dfs() {
54        for (int i = 0; i < n; i++)
55            used[i] = 0, p[i] = -1, base[i] = i;
56    }
57    bool dfs(int root) {
58        used[root] = ++T;
59        return go(root);
60    }
61    void match() {
62        int ans = 0;
63        for (int v = 0; v < n; v++)
64            for (int x : g[v])
65                if (pa[v] == -1 && pa[x] == -1) {
66                    pa[v] = x, pa[x] = v, ans++;
67                    break;
68                }
69        init_dfs();
70        for (int i = 0; i < n; i++)
71            if (pa[i] == -1 && dfs(i)) ans++, init_dfs();
72    }
73 };

```

```

74     cout << ans * 2 << "\n";
75     for (int i = 0; i < n; i++)
76         if (pa[i] > i)
77             cout << i + 1 << " " << pa[i] + 1 << "\n";
78 }
79 };

```

3.2.7. Minimum Weight Matching

```

1 struct Graph {
2     static const int MAXN = 105;
3     int n, e[MAXN][MAXN];
4     int match[MAXN], d[MAXN], onstk[MAXN];
5     vector<int> stk;
6     void init(int _n) {
7         n = _n;
8         for (int i = 0; i < n; i++)
9             for (int j = 0; j < n; j++)
10                 // change to appropriate infinity
11                 // if not complete graph
12                 e[i][j] = 0;
13    }
14    void add_edge(int u, int v, int w) {
15        e[u][v] = e[v][u] = w;
16    }
17    bool SPFA(int u) {
18        if (onstk[u]) return true;
19        stk.push_back(u);
20        onstk[u] = 1;
21        for (int v = 0; v < n; v++) {
22            if (u != v && match[u] != v && !onstk[v]) {
23                int m = match[v];
24                if (d[m] > d[u] - e[v][m] + e[u][v]) {
25                    d[m] = d[u] - e[v][m] + e[u][v];
26                    onstk[v] = 1;
27                    stk.push_back(v);
28                    if (SPFA(m)) return true;
29                    stk.pop_back();
30                    onstk[v] = 0;
31                }
32            }
33        }
34        onstk[u] = 0;
35        stk.pop_back();
36        return false;
37    }
38    int solve() {
39        for (int i = 0; i < n; i += 2) {
40            match[i] = i + 1;
41            match[i + 1] = i;
42        }
43        while (true) {
44            int found = 0;
45            for (int i = 0; i < n; i++) onstk[i] = d[i] = 0;
46            for (int i = 0; i < n; i++) {
47                stk.clear();
48                if (!onstk[i] && SPFA(i)) {
49                    found = 1;
50                    while (stk.size() >= 2) {
51                        int u = stk.back();
52                        stk.pop_back();
53                        int v = stk.back();
54                        stk.pop_back();
55                        match[u] = v;
56                        match[v] = u;
57                    }
58                }
59            }
60            if (!found) break;
61        }
62        int ret = 0;
63        for (int i = 0; i < n; i++) ret += e[i][match[i]];
64        ret /= 2;
65        return ret;
66    }
67 } graph;

```

3.2.8. Stable Marriage

```

1 // normal stable marriage problem
2 /* input:
3 3
4 Albert Laura Nancy Marcy
5 Brad Marcy Nancy Laura
6 Chuck Laura Marcy Nancy
7 Laura Chuck Albert Brad
8 Marcy Albert Chuck Brad
9 Nancy Brad Albert Chuck
10 */
11

```



```

13 using namespace std;
14 const int MAXN = 505;
15
16 int n;
17 int favor[MAXN][MAXN]; // favor[boy_id][rank] = girl_id;
18 int order[MAXN][MAXN]; // order[girl_id][boy_id] = rank;
19 int current[MAXN]; // current[boy_id] = rank;
20 // boy_id will pursue current[boy_id] girl.
21 int girl_current[MAXN]; // girl[girl_id] = boy_id;
22
23 void initialize() {
24     for (int i = 0; i < n; i++) {
25         current[i] = 0;
26         girl_current[i] = n;
27         order[i][n] = n;
28     }
29 }
30
31 map<string, int> male, female;
32 string bname[MAXN], gname[MAXN];
33 int fit = 0;
34
35 void stable_marriage() {
36     queue<int> que;
37     for (int i = 0; i < n; i++) que.push(i);
38     while (!que.empty()) {
39         int boy_id = que.front();
40         que.pop();
41
42         int girl_id = favor[boy_id][current[boy_id]];
43         current[boy_id]++;
44
45         if (order[girl_id][boy_id] <
46             order[girl_id][girl_current[girl_id]]) {
47             if (girl_current[girl_id] < n)
48                 que.push(girl_current[girl_id]);
49             girl_current[girl_id] = boy_id;
50         } else {
51             que.push(boy_id);
52         }
53     }
54 }
55
56 int main() {
57     cin >> n;
58
59     for (int i = 0; i < n; i++) {
60         string p, t;
61         cin >> p;
62         male[p] = i;
63         bname[i] = p;
64         for (int j = 0; j < n; j++) {
65             cin >> t;
66             if (!female.count(t)) {
67                 gname[fit] = t;
68                 female[t] = fit++;
69             }
70             favor[i][j] = female[t];
71         }
72     }
73
74     for (int i = 0; i < n; i++) {
75         string p, t;
76         cin >> p;
77         for (int j = 0; j < n; j++) {
78             cin >> t;
79             order[female[p]][male[t]] = j;
80         }
81     }
82
83     initialize();
84     stable_marriage();
85
86     for (int i = 0; i < n; i++) {
87         cout << bname[i] << " "
88              << gname[favor[i][current[i] - 1]] << endl;
89     }
90 }

```

3.2.9. Kuhn-Munkres algorithm

```

1 // Maximum Weight Perfect Bipartite Matching
2 // Detect non-perfect-matching:
3 // 1. set all edge[i][j] as INF
4 // 2. if solve() >= INF, it is not perfect matching.
5
6 typedef long long ll;
7 struct KM {
8     static const int MAXN = 1050;
9     static const ll INF = 1LL << 60;

```

```

11 int n, match[MAXN], vx[MAXN], vy[MAXN];
12 ll edge[MAXN][MAXN], lx[MAXN], ly[MAXN], slack[MAXN];
13 void init(int n) {
14     n = n;
15     for (int i = 0; i < n; i++)
16         for (int j = 0; j < n; j++) edge[i][j] = 0;
17
18 void add_edge(int x, int y, ll w) { edge[x][y] = w; }
19 bool DFS(int x) {
20     vx[x] = 1;
21     for (int y = 0; y < n; y++) {
22         if (vy[y]) continue;
23         if (lx[x] + ly[y] > edge[x][y]) {
24             slack[y] =
25                 min(slack[y], lx[x] + ly[y] - edge[x][y]);
26         } else {
27             vy[y] = 1;
28             if (match[y] == -1 || DFS(match[y])) {
29                 match[y] = x;
30                 return true;
31             }
32         }
33     }
34     return false;
35 }
36 ll solve() {
37     fill(match, match + n, -1);
38     fill(lx, lx + n, -INF);
39     fill(ly, ly + n, 0);
40     for (int i = 0; i < n; i++)
41         for (int j = 0; j < n; j++)
42             lx[i] = max(lx[i], edge[i][j]);
43     for (int i = 0; i < n; i++) {
44         fill(slack, slack + n, INF);
45         while (true) {
46             fill(vx, vx + n, 0);
47             fill(vy, vy + n, 0);
48             if (DFS(i)) break;
49             ll d = INF;
50             for (int j = 0; j < n; j++)
51                 if (!vy[j]) d = min(d, slack[j]);
52             for (int j = 0; j < n; j++) {
53                 if (vx[j]) lx[j] -= d;
54                 if (vy[j]) ly[j] += d;
55                 else slack[j] -= d;
56             }
57         }
58     }
59     ll res = 0;
60     for (int i = 0; i < n; i++) {
61         res += edge[match[i]][i];
62     }
63     return res;
64 }
65 } graph;

```

3.3. Shortest Path Faster Algorithm

```

1 struct SPFA {
2     static const int maxn = 1010, INF = 1e9;
3     int dis[maxn];
4     bitset<maxn> inq, inneg;
5     queue<int> q, tq;
6     vector<pii> v[maxn];
7     void make_edge(int s, int t, int w) {
8         v[s].emplace_back(t, w);
9     }
10    void dfs(int a) {
11        inneg[a] = 1;
12        for (pii i : v[a])
13            if (!inneg[i.F]) dfs(i.F);
14    }
15    bool solve(int n, int s) { // true if have neg-cycle
16        for (int i = 0; i <= n; i++) dis[i] = INF;
17        dis[s] = 0, q.push(s);
18        for (int i = 0; i < n; i++) {
19            inq.reset();
20            int now;
21            while (!q.empty()) {
22                now = q.front(), q.pop();
23                for (pii &i : v[now]) {
24                    if (dis[i.F] > dis[now] + i.S) {
25                        dis[i.F] = dis[now] + i.S;
26                        if (!inq[i.F]) tq.push(i.F), inq[i.F] = 1;
27                    }
28                }
29            }
30            q.swap(tq);
31        }
32        bool re = !q.empty();
33        inneg.reset();

```

```

35 while (!q.empty()) {
36     if (!inneg[q.front()]) dfs(q.front());
37     q.pop();
38 }
39 return re;
40 }
41 void reset(int n) {
42     for (int i = 0; i <= n; i++) v[i].clear();
43 };

```

3.4. Strongly Connected Components

```

1 struct TarjanScc {
2     int n, step;
3     vector<int> time, low, instk, stk;
4     vector<vector<int>> e, scc;
5     TarjanScc(int n_)
6         : n(n_), step(0), time(n), low(n), instk(n), e(n) {}
7     void add_edge(int u, int v) { e[u].push_back(v); }
8     void dfs(int x) {
9         time[x] = low[x] = ++step;
10        stk.push_back(x);
11        instk[x] = 1;
12        for (int y : e[x])
13            if (!time[y]) {
14                dfs(y);
15                low[x] = min(low[x], low[y]);
16            } else if (instk[y]) {
17                low[x] = min(low[x], time[y]);
18            }
19        if (time[x] == low[x]) {
20            scc.emplace_back();
21            for (int y = -1; y != x; ) {
22                y = stk.back();
23                stk.pop_back();
24                instk[y] = 0;
25                scc.back().push_back(y);
26            }
27        }
28    }
29    void solve() {
30        for (int i = 0; i < n; i++)
31            if (!time[i]) dfs(i);
32        reverse(scc.begin(), scc.end());
33        // scc in topological order
34    }
35 };

```

3.4.1. 2-Satisfiability

Requires: Strongly Connected Components

```

1 // 1 based, vertex in SCC = MAXN * 2
2 // (not i) is i + n
3 struct two_SAT {
4     int n, ans[MAXN];
5     SCC S;
6     void imply(int a, int b) { S.make_edge(a, b); }
7     bool solve(int _n) {
8         n = _n;
9         S.solve(n * 2);
10        for (int i = 1; i <= n; i++) {
11            if (S.scc[i] == S.scc[i + n]) return false;
12            ans[i] = (S.scc[i] < S.scc[i + n]);
13        }
14        return true;
15    }
16    void init(int _n) {
17        n = _n;
18        fill_n(ans, n + 1, 0);
19        S.init(n * 2);
20    }
21 } SAT;

```

3.5. Biconnected Components

3.5.1. Articulation Points

```

1 void dfs(int x, int p) {
2     tin[x] = low[x] = ++t;
3     int ch = 0;
4     for (auto u : g[x])
5         if (u.first != p) {
6             if (!ins[u.second])
7                 st.push(u.second), ins[u.second] = true;
8             if (tin[u.first]) {
9                 low[x] = min(low[x], tin[u.first]);
10                continue;

```

```

11            }
12            ++ch;
13            dfs(u.first, x);
14            low[x] = min(low[x], low[u.first]);
15            if (low[u.first] >= tin[x]) {
16                cut[x] = true;
17                ++sz;
18                while (true) {
19                    int e = st.top();
20                    st.pop();
21                    bcc[e] = sz;
22                    if (e == u.second) break;
23                }
24            }
25        }
26        if (ch == 1 && p == -1) cut[x] = false;
27    }

```

3.5.2. Bridges

```

1 // if there are multi-edges, then they are not bridges
2 void dfs(int x, int p) {
3     tin[x] = low[x] = ++t;
4     st.push(x);
5     for (auto u : g[x])
6         if (u.first != p) {
7             if (tin[u.first]) {
8                 low[x] = min(low[x], tin[u.first]);
9                 continue;
10            }
11            dfs(u.first, x);
12            low[x] = min(low[x], low[u.first]);
13            if (low[u.first] == tin[u.first]) br[u.second] = true;
14        }
15    if (tin[x] == low[x]) {
16        ++sz;
17        while (st.size()) {
18            int u = st.top();
19            st.pop();
20            bcc[u] = sz;
21            if (u == x) break;
22        }
23    }
24 }

```

3.6. Triconnected Components

```

1 // requires a union-find data structure
2 struct ThreeEdgeCC {
3     int V, ind;
4     vector<int> id, pre, post, low, deg, path;
5     vector<vector<int>> components;
6     UnionFind uf;
7     template <class Graph>
8     void dfs(const Graph &G, int v, int prev) {
9         pre[v] = ++ind;
10        for (int w : G[v])
11            if (w != v) {
12                if (w == prev) {
13                    prev = -1;
14                    continue;
15                }
16                if (pre[w] != -1) {
17                    if (pre[w] < pre[v]) {
18                        deg[v]++;
19                        low[v] = min(low[v], pre[w]);
20                    } else {
21                        deg[v]--;
22                        int &u = path[v];
23                        for (; u != -1 && pre[u] <= pre[w] &&
24                             pre[w] <= post[u];) {
25                            uf.join(v, u);
26                            deg[v] += deg[u];
27                            u = path[u];
28                        }
29                    }
30                }
31                continue;
32            }
33        dfs(G, w, v);
34        if (path[w] == -1 && deg[w] <= 1) {
35            deg[v] += deg[w];
36            low[v] = min(low[v], low[w]);
37            continue;
38        }
39        if (deg[w] == 0) w = path[w];
40        if (low[v] > low[w]) {
41            low[v] = min(low[v], low[w]);
42            swap(w, path[v]);

```

```

45     }
46     for (; w != -1; w = path[w]) {
47         uf.join(v, w);
48         deg[v] += deg[w];
49     }
50 }
51 post[v] = ind;
52 }
53 template <class Graph>
54 ThreeEdgeCC(const Graph &G)
55 : V(G.size()), ind(-1), id(V, -1), pre(V, -1),
56   post(V), low(V, INT_MAX), deg(V, 0), path(V, -1),
57   uf(V) {
58     for (int v = 0; v < V; v++)
59         if (pre[v] == -1) dfs(G, v, -1);
60     components.reserve(uf.cnt);
61     for (int v = 0; v < V; v++)
62         if (uf.find(v) == v) {
63             id[v] = components.size();
64             components.emplace_back(1, v);
65             components.back().reserve(uf.getSize(v));
66         }
67     for (int v = 0; v < V; v++)
68         if (id[v] == -1)
69             components[id[v] = id[uf.find(v)]] .push_back(v);
70 }
71 };

```

3.7. Centroid Decomposition

```

1 void get_center(int now) {
2     v[now] = true;
3     vtx.push_back(now);
4     sz[now] = 1;
5     mx[now] = 0;
6     for (int u : G[now])
7         if (!v[u]) {
8             get_center(u);
9             mx[now] = max(mx[now], sz[u]);
10            sz[now] += sz[u];
11        }
12    }
13 void get_dis(int now, int d, int len) {
14     dis[d][now] = cnt;
15     v[now] = true;
16     for (auto u : G[now])
17         if (!v[u.first]) { get_dis(u, d, len + u.second); }
18 }
19 void dfs(int now, int fa, int d) {
20     get_center(now);
21     int c = -1;
22     for (int i : vtx) {
23         if (max(mx[i], (int)vtx.size() - sz[i]) <=
24             (int)vtx.size() / 2)
25             c = i;
26         v[i] = false;
27     }
28     get_dis(c, d, 0);
29     for (int i : vtx) v[i] = false;
30     v[c] = true;
31     vtx.clear();
32     dep[c] = d;
33     p[c] = fa;
34     for (auto u : G[c])
35         if (u.first != fa && !v[u.first]) {
36             dfs(u.first, c, d + 1);
37         }
38 }

```

3.8. Minimum Mean Cycle

```

1 // d[i][j] == 0 if {i,j} !in E
2 long long d[1003][1003], dp[1003][1003];
3
4 pair<long long, long long> MMWC() {
5     memset(dp, 0x3f, sizeof(dp));
6     for (int i = 1; i <= n; ++i) dp[0][i] = 0;
7     for (int i = 1; i <= n; ++i) {
8         for (int j = 1; j <= n; ++j) {
9             for (int k = 1; k <= n; ++k) {
10                dp[i][k] = min(dp[i - 1][j] + d[j][k], dp[i][k]);
11            }
12        }
13    }
14    long long au = 1ll << 31, ad = 1;
15    for (int i = 1; i <= n; ++i) {
16        if (dp[n][i] == 0x3f3f3f3f3f3f3f3f) continue;
17        long long u = 0, d = 1;
18        for (int j = n - 1; j >= 0; --j) {

```

```

21         if ((dp[n][i] - dp[j][i]) * d > u * (n - j)) {
22             u = dp[n][i] - dp[j][i];
23             d = n - j;
24         }
25     }
26     if (u * ad < au * d) au = u, ad = d;
27 }
28 long long g = __gcd(au, ad);
29 return make_pair(au / g, ad / g);
30 }

```

3.9. Directed MST

```

1 template <typename T> struct DMST {
2     T g[maxn][maxn], fw[maxn];
3     int n, fr[maxn];
4     bool vis[maxn], inc[maxn];
5     void clear() {
6         for (int i = 0; i < maxn; ++i) {
7             for (int j = 0; j < maxn; ++j) g[i][j] = inf;
8             vis[i] = inc[i] = false;
9         }
10    }
11    void addedge(int u, int v, T w) {
12        g[u][v] = min(g[u][v], w);
13    }
14    T operator()(int root, int _n) {
15        n = _n;
16        if (dfs(root) != n) return -1;
17        T ans = 0;
18        while (true) {
19            for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] = i;
20            for (int i = 1; i <= n; ++i)
21                if (!inc[i]) {
22                    for (int j = 1; j <= n; ++j) {
23                        if (!inc[j] && i != j && g[j][i] < fw[i]) {
24                            fw[i] = g[j][i];
25                            fr[i] = j;
26                        }
27                    }
28                }
29            int x = -1;
30            for (int i = 1; i <= n; ++i)
31                if (i != root && !inc[i]) {
32                    int j = i, c = 0;
33                    while (j != root && fr[j] != i && c <= n)
34                        ++c, j = fr[j];
35                    if (j == root || c > n) continue;
36                    else {
37                        x = i;
38                        break;
39                    }
40                }
41            if (!x) {
42                for (int i = 1; i <= n; ++i)
43                    if (i != root && !inc[i]) ans += fw[i];
44                return ans;
45            }
46            int y = x;
47            for (int i = 1; i <= n; ++i) vis[i] = false;
48            do {
49                ans += fw[y];
50                y = fr[y];
51                vis[y] = inc[y] = true;
52            } while (y != x);
53            inc[x] = false;
54            for (int k = 1; k <= n; ++k)
55                if (vis[k]) {
56                    for (int j = 1; j <= n; ++j)
57                        if (!vis[j]) {
58                            if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
59                            if (g[j][k] < inf &&
60                                g[j][k] - fw[k] < g[j][x])
61                                g[j][x] = g[j][k] - fw[k];
62                        }
63                }
64            }
65        return ans;
66    }
67    int dfs(int now) {
68        int r = 1;
69        vis[now] = true;
70        for (int i = 1; i <= n; ++i)
71            if (g[now][i] < inf && !vis[i]) r += dfs(i);
72        return r;
73    }
74 };

```

3.10. Maximum Clique

```
1 // source: KACTL
```

```

3 typedef vector<bitset<200>> vb;
  struct MaxClique {
5     double limit = 0.025, pk = 0;
    struct Vertex {
7         int i, d = 0;
    };
    typedef vector<Vertex> vv;
    vb e;
    vv V;
    vector<vi> C;
    vi qmax, q, S, old;
    void init(vv &r) {
13         for (auto &v : r) v.d = 0;
        for (auto &v : r)
15             for (auto j : r) v.d += e[v.i][j.i];
        sort(all(r), [](auto a, auto b) { return a.d > b.d; });
17         int mxD = r[0].d;
        rep(i, 0, sz(r)) r[i].d = min(i, mxD) + 1;
    }
21 void expand(vv &R, int lev = 1) {
    S[lev] += S[lev - 1] - old[lev];
    old[lev] = S[lev - 1];
    while (sz(R)) {
23         if (sz(q) + R.back().d <= sz(qmax)) return;
        q.push_back(R.back().i);
        vv T;
25         for (auto v : R)
            if (e[R.back().i][v.i] T.push_back({v.i});
        if (sz(T)) {
27             if (S[lev]++ / ++pk < limit) init(T);
            int j = 0, mxk = 1,
                mnk = max(sz(qmax) - sz(q) + 1, 1);
            C[1].clear(), C[2].clear();
            for (auto v : T) {
31                 int k = 1;
                auto f = [&](int i) { return e[v.i][i]; };
                while (any_of(all(C[k]), f)) k++;
                if (k > mxk) mxk = k, C[mxk + 1].clear();
                if (k < mnk) T[j++] = v.i;
                C[k].push_back(v.i);
            }
            if (j > 0) T[j - 1].d = 0;
            rep(k, mnk, mxk + 1) for (int i : C[k]) T[j].i = i,
33                 T[j++] = k;
            expand(T, lev + 1);
        } else if (sz(q) > sz(qmax)) qmax = q;
        q.pop_back(), R.pop_back();
    }
51 }
vi maxClique() {
    init(V), expand(V);
    return qmax;
}
53 MaxClique(vb conn)
    : e(conn), C(sz(e) + 1), S(sz(C)), old(S) {
55     rep(i, 0, sz(e)) V.push_back({i});
}
57 }
61 };

```

3.11. Dominator Tree

```

1 // idom[n] is the unique node that strictly dominates n but
  // does not strictly dominate any other node that strictly
  // dominates n. idom[n] = 0 if n is entry or the entry
  // cannot reach n.
  struct DominatorTree {
5     static const int MAXN = 200010;
    int n, s;
    vector<int> g[MAXN], pred[MAXN];
    vector<int> cov[MAXN];
    int dfn[MAXN], nfd[MAXN], ts;
    int par[MAXN];
    int sdom[MAXN], idom[MAXN];
    int mom[MAXN], mn[MAXN];
13     inline bool cmp(int u, int v) { return dfn[u] < dfn[v]; }
    int eval(int u) {
15         if (mom[u] == u) return u;
        int res = eval(mom[u]);
        if (cmp(sdom[mn[mom[u]]], sdom[mn[u]]))
            mn[u] = mn[mom[u]];
        return mom[u] = res;
    }
23 void init(int _n, int _s) {
    n = _n;
    s = _s;
    REPl(i, 1, n) {
25         g[i].clear();
    }
    pred[i].clear();
    idom[i] = 0;
}

```

```

    pred[i].clear();
    idom[i] = 0;
}
31 void add_edge(int u, int v) {
    g[u].push_back(v);
    pred[v].push_back(u);
}
33 void DFS(int u) {
    ts++;
    dfn[u] = ts;
    nfd[ts] = u;
    for (int v : g[u])
    35         if (dfn[v] == 0) {
            par[v] = u;
            DFS(v);
        }
    }
    void build() {
    37         ts = 0;
        REPl(i, 1, n) {
            dfn[i] = nfd[i] = 0;
            cov[i].clear();
            mom[i] = mn[i] = sdom[i] = i;
        }
        DFS(s);
        for (int i = ts; i >= 2; i--) {
            int u = nfd[i];
            if (u == 0) continue;
            for (int v : pred[u])
            39                 if (dfn[v]) {
                    eval(v);
                    if (cmp(sdom[mn[v]], sdom[u]))
                        sdom[u] = sdom[mn[v]];
                }
            cov[sdom[u]].push_back(u);
            mom[u] = par[u];
            for (int w : cov[par[u]]) {
                eval(w);
                if (cmp(sdom[mn[w]], par[u])) idom[w] = mn[w];
                else idom[w] = par[u];
            }
            cov[par[u]].clear();
        }
        REPl(i, 2, ts) {
            int u = nfd[i];
            if (u == 0) continue;
            if (idom[u] != sdom[u]) idom[u] = idom[idom[u]];
        }
    }
    } dom;
}

```

3.12. Manhattan Distance MST

```

1 // returns [(dist, from, to), ...]
  // then do normal mst afterwards
  typedef Point<int> P;
  vector<array<int, 3>> manhattanMST(vector<P> ps) {
    vi id(sz(ps));
    iota(all(id), 0);
    vector<array<int, 3>> edges;
    rep(k, 0, 4) {
        sort(all(id), [&](int i, int j) {
            return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;
        });
        map<int, int> sweep;
        for (int i : id) {
            for (auto it = sweep.lower_bound(-ps[i].y);
                it != sweep.end(); sweep.erase(it++)) {
                int j = it->second;
                P d = ps[i] - ps[j];
                if (d.y > d.x) break;
                edges.push_back({d.y + d.x, i, j});
            }
            sweep[-ps[i].y] = i;
        }
        for (P &p : ps)
            if (k & 1) p.x = -p.x;
            else swap(p.x, p.y);
    }
    return edges;
}

```

4. Math

4.1. Number Theory

4.1.1. Mod Struct

A list of safe primes: 26003, 27767, 28319, 28979, 29243, 29759, 30467, 910927547, 919012223, 947326223, 990669467, 1007939579, 1019126699, 929760389146037459, 975500632317046523, 989312547895528379

NTT prime p	$p-1$	primitive root
65537	$1 \ll 16$	3
998244353	$119 \ll 23$	3
2748779069441	$5 \ll 39$	3
1945555039024054273	$27 \ll 56$	5

Requires: Extended GCD

```

1
3 template <typename T> struct M {
4     static T MOD; // change to constexpr if already known
5     T v;
6     M(T x = 0) {
7         v = (-MOD <= x && x < MOD) ? x : x % MOD;
8         if (v < 0) v += MOD;
9     }
10    explicit operator T() const { return v; }
11    bool operator==(const M &b) const { return v == b.v; }
12    bool operator!=(const M &b) const { return v != b.v; }
13    M operator-() { return M(-v); }
14    M operator+(M b) { return M(v + b.v); }
15    M operator-(M b) { return M(v - b.v); }
16    M operator*(M b) { return M((__int128)v * b.v % MOD); }
17    M operator/(M b) { return *this * (b ^ (MOD - 2)); }
18    // change above implementation to this if MOD is not prime
19    M inv() {
20        auto [p, _, g] = extgcd(v, MOD);
21        return assert(g == 1), p;
22    }
23    friend M operator^(M a, ll b) {
24        M ans(1);
25        for (; b >= 1, a *= a)
26            if (b & 1) ans *= a;
27        return ans;
28    }
29    friend M &operator+=(M &a, M b) { return a = a + b; }
30    friend M &operator-=(M &a, M b) { return a = a - b; }
31    friend M &operator*=(M &a, M b) { return a = a * b; }
32    friend M &operator/=(M &a, M b) { return a = a / b; }
33 };
34 using Mod = M<int>;
35 template <> int Mod::MOD = 1'000'000'007;
36 int &MOD = Mod::MOD;

```

4.1.2. Miller-Rabin

Requires: Mod Struct

```

1
3 // checks if Mod::MOD is prime
4 bool is_prime() {
5     if (MOD < 2 || MOD % 2 == 0) return MOD == 2;
6     Mod A[] = {2, 7, 61}; // for int values (< 2^31)
7     // ll: 2, 325, 9375, 28178, 450775, 9780504, 1795265022
8     int s = __builtin_ctzll(MOD - 1), i;
9     for (Mod a : A) {
10        Mod x = a ^ (MOD >> s);
11        for (i = 0; i < s && (x + 1).v > 2; i++) x *= x;
12        if (i && x != -1) return 0;
13    }
14    return 1;
15 }

```

4.1.3. Linear Sieve

```

1 constexpr ll MAXN = 1000000;
2 bitset<MAXN> is_prime;
3 vector<ll> primes;
4 ll mpf[MAXN], phi[MAXN], mu[MAXN];
5
6 void sieve() {
7     is_prime.set();
8     is_prime[1] = 0;
9     mu[1] = phi[1] = 1;
10    for (ll i = 2; i < MAXN; i++) {
11        if (is_prime[i]) {
12            mpf[i] = i;
13            primes.push_back(i);
14            phi[i] = i - 1;
15            mu[i] = -1;
16        }
17    }
18 }

```

```

17 for (ll p : primes) {
18     if (p > mpf[i] || i * p >= MAXN) break;
19     is_prime[i * p] = 0;
20     mpf[i * p] = p;
21     mu[i * p] = -mu[i];
22     if (i % p == 0)
23         phi[i * p] = phi[i] * p, mu[i * p] = 0;
24     else phi[i * p] = phi[i] * (p - 1);
25 }
26 }
27 }

```

4.1.4. Get Factors

Requires: Linear Sieve

```

1
3 vector<ll> all_factors(ll n) {
4     vector<ll> fac = {1};
5     while (n > 1) {
6         const ll p = mpf[n];
7         vector<ll> cur = {1};
8         while (n % p == 0) {
9             n /= p;
10            cur.push_back(cur.back() * p);
11        }
12        vector<ll> tmp;
13        for (auto x : fac)
14            for (auto y : cur) tmp.push_back(x * y);
15        tmp.swap(fac);
16    }
17    return fac;
18 }

```

4.1.5. Binary GCD

```

1 // returns the gcd of non-negative a, b
2 ull bin_gcd(ull a, ull b) {
3     if (!a || !b) return a + b;
4     int s = __builtin_ctzll(a | b);
5     a >>= __builtin_ctzll(a);
6     while (b) {
7         if ((b >>= __builtin_ctzll(b)) < a) swap(a, b);
8         b -= a;
9     }
10    return a << s;
11 }

```

4.1.6. Extended GCD

```

1 // returns (p, q, g): p * a + q * b == g == gcd(a, b)
2 // g is not guaranteed to be positive when a < 0 or b < 0
3 tuple<ll, ll, ll> extgcd(ll a, ll b) {
4     ll s = 1, t = 0, u = 0, v = 1;
5     while (b) {
6         ll q = a / b;
7         swap(a -= q * b, b);
8         swap(s -= q * t, t);
9         swap(u -= q * v, v);
10    }
11    return {s, u, a};
12 }

```

4.1.7. Chinese Remainder Theorem

Requires: Extended GCD

```

1 // for 0 <= a < m, 0 <= b < n, returns the smallest x >= 0
2 // such that x % m == a and x % n == b
3 ll crt(ll a, ll m, ll b, ll n) {
4     if (n > m) swap(a, b), swap(m, n);
5     auto [x, y, g] = extgcd(m, n);
6     assert((a - b) % g == 0); // no solution
7     x = ((b - a) / g * x) % (n / g) * m + a;
8     return x < 0 ? x + m / g * n : x;
9 }

```

4.1.8. Baby-Step Giant-Step

Requires: Mod Struct

```

1
3 // returns x such that a ^ x = b where x in [l, r)
4 ll bsqs(Mod a, Mod b, ll l = 0, ll r = MOD - 1) {
5     int m = sqrt(r - l) + 1, i;
6     unordered_map<ll, ll> tb;
7     Mod d = (a ^ l) / b;
8     for (i = 0, d = (a ^ l) / b; i < m; i++, d *= a)
9         if (d == 1) return l + i;
10    else tb[(ll)d] = l + i;
11    Mod c = Mod(1) / (a ^ m);
12    for (i = 0, d = 1; i < m; i++, d *= c)
13        if (auto j = tb.find((ll)d); j != tb.end())
14            return j->second + i * m;
15    return assert(0), -1; // no solution
16 }

```


4.1.9. Pollard's Rho

```

1 ll f(ll x, ll mod) { return (x * x + 1) % mod; }
  // n should be composite
2 ll pollard_rho(ll n) {
3     if (!(n & 1)) return 2;
4     while (1) {
5         ll y = 2, x = RNG() % (n - 1) + 1, res = 1;
6         for (int sz = 2; res == 1; sz *= 2) {
7             for (int i = 0; i < sz && res <= 1; i++) {
8                 x = f(x, n);
9                 res = __gcd(abs(x - y), n);
10            }
11            y = x;
12        }
13        if (res != 0 && res != n) return res;
14    }
15 }

```

4.1.10. Tonelli-Shanks Algorithm

Requires: Mod Struct

```

1
2
3 int legendre(Mod a) {
4     if (a == 0) return 0;
5     return (a ^ ((MOD - 1) / 2)) == 1 ? 1 : -1;
6 }
7 Mod sqrt(Mod a) {
8     assert(legendre(a) != -1); // no solution
9     ll p = MOD, s = p - 1;
10    if (a == 0) return 0;
11    if (p == 2) return 1;
12    if (p % 4 == 3) return a ^ ((p + 1) / 4);
13    int r, m;
14    for (r = 0; !(s & 1); r++) s >>= 1;
15    Mod n = 2;
16    while (legendre(n) != -1) n += 1;
17    Mod x = a ^ ((s + 1) / 2), b = a ^ s, g = n ^ s;
18    while (b != 1) {
19        Mod t = b;
20        for (m = 0; t != 1; m++) t *= g;
21        Mod gs = g ^ (1LL << (r - m - 1));
22        g = gs * gs, x *= gs, b *= g, r = m;
23    }
24    return x;
25 }
  // to get sqrt(X) modulo p^k, where p is an odd prime:
26 // c = x^2 (mod p), c = X^2 (mod p^k), q = p^(k-1)
27 // X = x^q * c^((p^k-2q+1)/2) (mod p^k)

```

4.1.11. Chinese Sieve

```

1 const ll N = 1000000;
  // f, g, h multiplicative, h = f (dirichlet convolution) g
2 ll pre_g(ll n);
3 ll pre_h(ll n);
4 // preprocessed prefix sum of f
5 ll pre_f[N];
6 // prefix sum of multiplicative function f
7 ll solve_f(ll n) {
8     static unordered_map<ll, ll> m;
9     if (n < N) return pre_f[n];
10    if (m.count(n)) return m[n];
11    ll ans = pre_h(n);
12    for (ll l = 2, r; l <= n; l = r + 1) {
13        r = n / (n / l);
14        ans -= (pre_g(r) - pre_g(l - 1)) * djs_f(n / l);
15    }
16    return m[n] = ans;
17 }

```

4.1.12. Rational Number Binary Search

```

1 struct QQ {
2     ll p, q;
3     QQ go(QQ b, ll d) { return {p + b.p * d, q + b.q * d}; }
4 };
5 bool pred(QQ);
  // returns smallest p/q in [lo, hi] such that
6 // pred(p/q) is true, and 0 <= p, q <= N
7 QQ frac_bs(ll N) {
8     QQ lo{0, 1}, hi{1, 0};
9     if (pred(lo)) return lo;
10    assert(pred(hi));
11    bool dir = 1, L = 1, H = 1;
12    for (; L || H; dir = !dir) {
13        ll len = 0, step = 1;
14        for (int t = 0; t < 2 && (t ? step /= 2 : step *= 2);)
15            if (QQ mid = hi.go(lo, len + step);
16                mid.p > N || mid.q > N || dir ^ pred(mid))

```

```

19         t++;
20         else len += step;
21         swap(lo, hi = hi.go(lo, len));
22         (dir ? L : H) = !len;
23     }
24     return dir ? hi : lo;
25 }

```

4.1.13. Farey Sequence

```

1 // returns (e/f), where (a/b, c/d, e/f) are
  // three consecutive terms in the order n farey sequence
2 // to start, call next_farey(n, 0, 1, 1, n)
3 pll next_farey(ll n, ll a, ll b, ll c, ll d) {
4     ll p = (n + b) / d;
5     return pll(p * c - a, p * d - b);
6 }

```

4.2. Combinatorics

4.2.1. Matroid Intersection

This template assumes 2 weighted matroids of the same type, and that removing an element is much more expensive than checking if one can be added. **Remember to change the implementation details.**

The ground set is $0, 1, \dots, n-1$, where element i has weight $w[i]$. For the unweighted version, remove weights and change BF/SPFA to BFS.

```

1 constexpr int N = 100;
2 constexpr int INF = 1e9;
3
4 struct Matroid { // represents an independent set
5     Matroid(bitset<N>); // initialize from an independent set
6     bool can_add(int); // if adding will break independence
7     Matroid remove(int); // removing from the set
8 };
9
10 auto matroid_intersection(int n, const vector<int> &w) {
11     bitset<N> S;
12     for (int sz = 1; sz <= n; sz++) {
13         Matroid M1(S), M2(S);
14
15         vector<vector<pii>> e(n + 2);
16         for (int j = 0; j < n; j++)
17             if (!S[j]) {
18                 if (M1.can_add(j)) e[n].emplace_back(j, -w[j]);
19                 if (M2.can_add(j)) e[j].emplace_back(n + 1, 0);
20             }
21         for (int i = 0; i < n; i++)
22             if (S[i]) {
23                 Matroid T1 = M1.remove(i), T2 = M2.remove(i);
24                 for (int j = 0; j < n; j++)
25                     if (!S[j]) {
26                         if (T1.can_add(j)) e[i].emplace_back(j, -w[j]);
27                         if (T2.can_add(j)) e[j].emplace_back(i, w[i]);
28                     }
29             }
30
31         vector<pii> dis(n + 2, {INF, 0});
32         vector<int> prev(n + 2, -1);
33         dis[n] = {0, 0};
34         // change to SPFA for more speed, if necessary
35         bool upd = 1;
36         while (upd) {
37             upd = 0;
38             for (int u = 0; u < n + 2; u++)
39                 for (auto [v, c] : e[u]) {
40                     pii x{dis[u].first + c, dis[u].second + 1};
41                     if (x < dis[v]) dis[v] = x, prev[v] = u, upd = 1;
42                 }
43         }
44
45         if (dis[n + 1].first < INF)
46             for (int x = prev[n + 1]; x != n; x = prev[x])
47                 S.flip(x);
48         else break;
49
50         // S is the max-weighted independent set with size sz
51     }
52     return S;
53 }

```

4.2.2. De Bruijn Sequence

```

1 int res[kN], aux[kN], a[kN], sz;
2 void Rec(int t, int p, int n, int k) {
3     if (t > n) {
4         if (n % p == 0)
5             for (int i = 1; i <= p; ++i) res[sz++] = aux[i];
6         else {

```

```

7   aux[t] = aux[t - p];
9   Rec(t + 1, p, n, k);
11  for (aux[t] = aux[t - p] + 1; aux[t] < k; ++aux[t])
13  Rec(t + 1, t, n, k);
15  }
17  int DeBruijn(int k, int n) {
19  // return cyclic string of length k^n such that every
21  // string of length n using k character appears as a
23  // substring.
25  if (k == 1) return res[0] = 0, 1;
27  fill(aux, aux + k * n, 0);
29  return sz = 0, Rec(1, 1, n, k), sz;
31  }

```

4.2.3. Multinomial

```

1  // ways to permute v[i]
3  ll multinomial(vi &v) {
5  ll c = 1, m = v.empty() ? 1 : v[0];
7  for (int i = 1; i < v.size(); i++)
9  for (int j = 0; i < v[i]; j++) c = c * ++m / (j + 1);
11 return c;
13 }

```

4.3. Algebra

4.3.1. Formal Power Series

```

1  template <typename mint>
3  struct FormalPowerSeries : vector<mint> {
5  using vector<mint>::vector;
7  using FPS = FormalPowerSeries;
9  FPS &operator+=(const FPS &r) {
11 if (r.size() > this->size()) this->resize(r.size());
13 for (int i = 0; i < (int)r.size(); i++)
15 (*this)[i] += r[i];
17 return *this;
19 }
21 FPS &operator+=(const mint &r) {
23 if (this->empty()) this->resize(1);
25 (*this)[0] += r;
27 return *this;
29 }
31 FPS &operator-=(const FPS &r) {
33 if (r.size() > this->size()) this->resize(r.size());
35 for (int i = 0; i < (int)r.size(); i++)
37 (*this)[i] -= r[i];
39 return *this;
41 }
43 FPS &operator-=(const mint &r) {
45 if (this->empty()) this->resize(1);
47 (*this)[0] -= r;
49 return *this;
51 }
53 FPS &operator*=(const mint &v) {
55 for (int k = 0; k < (int)this->size(); k++)
57 (*this)[k] *= v;
59 return *this;
61 }
63 FPS &operator/=(const FPS &r) {
65 if (this->size() < r.size()) {
67 this->clear();
69 return *this;
71 }
73 int n = this->size() - r.size() + 1;
75 if ((int)r.size() <= 64) {
77 FPS f(*this), g(r);
79 g.shrink();
81 mint coeff = g.back().inverse();
83 for (auto &x : g) x *= coeff;
85 int deg = (int)f.size() - (int)g.size() + 1;
87 int gs = g.size();
89 FPS quo(deg);
91 for (int i = deg - 1; i >= 0; i--) {
93 quo[i] = f[i + gs - 1];
95 for (int j = 0; j < gs; j++)
97 f[i + j] -= quo[i] * g[j];
99 }
101 *this = quo * coeff;

```

```

61 this->resize(n, mint(0));
63 return *this;
65 }
67 return *this = ((*this).rev().pre(n) * r.rev().inv(n))
69 .pre(n)
71 .rev();
73 }
75 FPS &operator%=(const FPS &r) const {
77 *this -= *this / r * r;
79 shrink();
81 return *this;
83 }
85 FPS operator+(const FPS &r) const {
87 return FPS(*this) += r;
89 }
91 FPS operator+(const mint &v) const {
93 return FPS(*this) += v;
95 }
97 FPS operator-(const FPS &r) const {
99 return FPS(*this) -= r;
101 }
103 FPS operator-(const mint &v) const {
105 return FPS(*this) -= v;
107 }
109 FPS operator*(const FPS &r) const {
111 return FPS(*this) *= r;
113 }
115 FPS operator*(const mint &v) const {
117 return FPS(*this) *= v;
119 }
121 FPS operator/(const FPS &r) const {
123 return FPS(*this) /= r;
125 }
127 FPS operator/(const FPS &r) const {
129 return FPS(*this) /= r;
131 }
133 FPS operator%(const FPS &r) const {
135 return FPS(*this) %= r;
137 }
139 FPS operator-() const {
141 FPS ret(this->size());
143 for (int i = 0; i < (int)this->size(); i++)
145 ret[i] = -(*this)[i];
147 return ret;
149 }
151 void shrink() {
153 while (this->size() && this->back() == mint(0))
155 this->pop_back();
157 }
159 FPS rev() const {
161 FPS ret(*this);
163 reverse(begin(ret), end(ret));
165 return ret;
167 }
169 FPS dot(FPS r) const {
171 FPS ret(min(this->size(), r.size()));
173 for (int i = 0; i < (int)ret.size(); i++)
175 ret[i] = (*this)[i] * r[i];
177 return ret;
179 }
181 FPS pre(int sz) const {
183 return FPS(begin(*this),
185 begin(*this) + min((int)this->size(), sz));
187 }
189 FPS operator>>(int sz) const {
191 if ((int)this->size() <= sz) return {};
193 FPS ret(*this);
195 ret.erase(ret.begin(), ret.begin() + sz);
197 return ret;
199 }
201 FPS operator<<(int sz) const {
203 FPS ret(*this);
205 ret.insert(ret.begin(), sz, mint(0));
207 return ret;
209 }
211 FPS diff() const {
213 const int n = (int)this->size();
215 FPS ret(max(0, n - 1));
217 mint one(1), coeff(1);
219 for (int i = 1; i < n; i++) {
221 ret[i - 1] = (*this)[i] * coeff;
223 coeff += one;
225 }
227 return ret;
229 }

```

```

153 FPS integral() const {
154     const int n = (int)this->size();
155     FPS ret(n + 1);
156     ret[0] = mint(0);
157     if (n > 0) ret[1] = mint(1);
158     auto mod = mint::get_mod();
159     for (int i = 2; i <= n; i++)
160         ret[i] = (-ret[mod % i]) * (mod / i);
161     for (int i = 0; i < n; i++) ret[i + 1] *= (*this)[i];
162     return ret;
163 }
164
165 mint eval(mint x) const {
166     mint r = 0, w = 1;
167     for (auto &v : *this) r += w * v, w *= x;
168     return r;
169 }
170
171 FPS log(int deg = -1) const {
172     assert((*this)[0] == mint(1));
173     if (deg == -1) deg = (int)this->size();
174     return (this->diff() * this->inv(deg))
175         .pre(deg - 1)
176         .integral();
177 }
178
179 FPS pow(int64_t k, int deg = -1) const {
180     const int n = (int)this->size();
181     if (deg == -1) deg = n;
182     for (int i = 0; i < n; i++) {
183         if ((*this)[i] != mint(0)) {
184             if (i * k > deg) return FPS(deg, mint(0));
185             mint rev = mint(1) / (*this)[i];
186             FPS ret =
187                 ((*this * rev) >> i).log(deg) * k.exp(deg) *
188                 ((*this)[i].pow(k));
189             ret = (ret << (i * k)).pre(deg);
190             if ((int)ret.size() < deg) ret.resize(deg, mint(0));
191             return ret;
192         }
193     }
194     return FPS(deg, mint(0));
195 }
196
197 static void *ntt_ptr;
198 static void set_fft();
199 FPS &operator*=(const FPS &r);
200 void ntt();
201 void intt();
202 void ntt_doubling();
203 static int ntt_pr();
204 FPS inv(int deg = -1) const;
205 FPS exp(int deg = -1) const;
206 };
207 template <typename mint>
208 void *FormalPowerSeries<mint>::ntt_ptr = nullptr;

```

4.4. Theorems

4.4.1. Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G , where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i, j) in G .

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

4.4.2. Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if $i < j$ and $(i, j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{\text{rank}(D)}{2}$ is the maximum matching on G .

4.4.3. Cayley's Formula

- Given a degree sequence d_1, d_2, \dots, d_n for each labeled vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1, 2, \dots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.

4.4.4. Erdős–Gallai Theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \dots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + d_2 + \dots + d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

holds for all $1 \leq k \leq n$.

4.4.5. Burnside's Lemma

Let X be a set and G be a group that acts on X . For $g \in G$, denote by X^g the elements fixed by g :

$$X^g = \{x \in X \mid gx \in X\}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

5. Numeric

5.1. Barrett Reduction

```

1 using ull = unsigned long long;
2 using ul = __uint128_t;
3 // very fast calculation of a % m
4 struct reduction {
5     const ull m, d;
6     explicit reduction(ull m) : m(m), d(((ul)1 << 64) / m) {}
7     inline ull operator()(ull a) const {
8         ull q = (ull)((a * d) >> 64);
9         return (a - q * m) >= m ? a - m : a;
10    }
11 };

```

5.2. Long Long Multiplication

```

1 using ull = unsigned long long;
2 using ll = long long;
3 using ld = long double;
4 // returns a * b % M where a, b < M < 2**63
5 ull mult(ull a, ull b, ull M) {
6     ll ret = a * b - M * ull(ld(a) * ld(b) / ld(M));
7     return ret + M * (ret < 0) - M * (ret >= (ll)M);
8 }

```

5.3. Fast Fourier Transform

```

1 template <typename T>
2 void fft(int n, vector<T> &a, vector<T> &rt, bool inv) {
3     vector<int> br(n);
4     for (int i = 1; i < n; i++) {
5         br[i] = (i & 1) ? br[i - 1] + n / 2 : br[i / 2] / 2;
6         if (br[i] > i) swap(a[i], a[br[i]]);
7     }
8     for (int len = 2; len <= n; len *= 2)
9         for (int i = 0; i < n; i += len)
10             for (int j = 0; j < len / 2; j++) {
11                 int pos = n / len * (inv ? len - j : j);
12                 T u = a[i + j], v = a[i + j + len / 2] * rt[pos];
13                 a[i + j] = u + v, a[i + j + len / 2] = u - v;
14             }
15     if (T minv = T(1) / T(n); inv)
16         for (T &x : a) x *= minv;
17 }

```

Requires: Mod Struct

```

1
2 void ntt(vector<Mod> &a, bool inv, Mod primitive_root) {
3     int n = a.size();
4     Mod root = primitive_root ^ (MOD - 1) / n;
5     vector<Mod> rt(n + 1, 1);
6     for (int i = 0; i < n; i++) rt[i + 1] = rt[i] * root;
7     fft_(n, a, rt, inv);
8 }
9
10 void fft(vector<complex<double>> &a, bool inv) {
11     int n = a.size();
12     vector<complex<double>> rt(n + 1);
13     double arg = acos(-1) * 2 / n;
14     for (int i = 0; i <= n; i++)
15         rt[i] = {cos(arg * i), sin(arg * i)};
16     fft_(n, a, rt, inv);
17 }

```

5.4. Fast Walsh-Hadamard Transform

Requires: Mod Struct

```

1 void fwht(vector<Mod> &a, bool inv) {
2     int n = a.size();
3     for (int d = 1; d < n; d <= 1)
4         for (int m = 0; m < n; m++)
5             if (!(m & d)) {
6                 inv ? a[m] -= a[m | d] : a[m] += a[m | d]; // AND
7                 inv ? a[m | d] -= a[m] : a[m | d] += a[m]; // OR
8                 Mod x = a[m], y = a[m | d]; // XOR
9                 a[m] = x + y, a[m | d] = x - y; // XOR
10            }
11    if (Mod iv = Mod(1) / n; inv) // XOR
12        for (Mod &i : a) i *= iv; // XOR
13 }

```

5.5. Subset Convolution

Requires: Mod Struct

```

1 #pragma GCC target("popcnt")
2 #include <immintrin.h>
3
4 void fwht(int n, vector<vector<Mod>> &a, bool inv) {
5     for (int h = 0; h < n; h++)
6         for (int i = 0; i < (1 << n); i++)
7             if (!(i & (1 << h)))
8                 for (int k = 0; k <= n; k++)
9                     inv ? a[i | (1 << h)][k] -= a[i][k] :
10                        a[i | (1 << h)][k] += a[i][k];
11 }
12 // c[k] = sum(popcnt(i & j) == sz && i | j == k) a[i] * b[j]
13 vector<Mod> subset_convolution(int n, int sz,
14                               const vector<Mod> &a_,
15                               const vector<Mod> &b_) {
16     int len = n + sz + 1, N = 1 << n;
17     vector<vector<Mod>> a(1 << n, vector<Mod>(len, 0)), b = a;
18     for (int i = 0; i < N; i++)
19         a[i][_mm_popcnt_u64(i)] = a_[i],
20         b[i][_mm_popcnt_u64(i)] = b_[i];
21     fwht(n, a, 0), fwht(n, b, 0);
22     for (int i = 0; i < N; i++) {
23         vector<Mod> tmp(len);
24         for (int j = 0; j < len; j++)
25             for (int k = 0; k <= j; k++)
26                 tmp[j] += a[i][k] * b[i][j - k];
27         a[i] = tmp;
28     }
29     fwht(n, a, 1);
30     vector<Mod> c(N);
31     for (int i = 0; i < N; i++)
32         c[i] = a[i][_mm_popcnt_u64(i) + sz];
33     return c;
34 }

```

5.6. Linear Recurrences

5.6.1. Berlekamp-Massey Algorithm

```

1 template <typename T>
2 vector<T> berlekamp_massey(const vector<T> &s) {
3     int n = s.size(), l = 0, m = 1;
4     vector<T> r(n), p(n);
5     r[0] = p[0] = 1;
6     T b = 1, d = 0;
7     for (int i = 0; i < n; i++, m++, d = 0) {
8         for (int j = 0; j <= l; j++) d += r[j] * s[i - j];
9         if ((d /= b) == 0) continue; // change if T is float
10        auto t = r;
11        for (int j = m; j < n; j++) r[j] -= d * p[j - m];
12        if (l * 2 <= i) l = i + 1 - l, b *= d, m = 0, p = t;
13    }
14    return r.resize(l + 1), reverse(r.begin(), r.end()), r;
15 }

```

5.6.2. Linear Recurrence Calculation

```

1 template <typename T> struct lin_rec {
2     using poly = vector<T>;
3     poly mul(poly a, poly b, poly m) {
4         int n = m.size();
5         poly r(n);
6         for (int i = n - 1; i >= 0; i--) {
7             r.insert(r.begin(), 0), r.pop_back();
8             T c = r[n - 1] + a[n - 1] * b[i];
9             // c /= m[n - 1]; if m is not monic
10            for (int j = 0; j < n; j++)

```

```

11        r[j] += a[j] * b[i] - c * m[j];
12    }
13    return r;
14 }
15 poly pow(poly p, ll k, poly m) {
16     poly r(m.size());
17     r[0] = 1;
18     for (; k >= 1; p = mul(p, p, m))
19         if (k & 1) r = mul(r, p, m);
20     return r;
21 }
22 T calc(poly t, poly r, ll k) {
23     int n = r.size();
24     poly p(n);
25     p[1] = 1;
26     poly q = pow(p, k, r);
27     T ans = 0;
28     for (int i = 0; i < n; i++) ans += t[i] * q[i];
29     return ans;
30 }
31 };

```

5.7. Matrices

5.7.1. Determinant

Requires: Mod Struct

```

1 Mod det(vector<vector<Mod>> a) {
2     int n = a.size();
3     Mod ans = 1;
4     for (int i = 0; i < n; i++) {
5         int b = i;
6         for (int j = i + 1; j < n; j++)
7             if (a[j][i] != 0) {
8                 b = j;
9                 break;
10            }
11        if (i != b) swap(a[i], a[b]), ans = -ans;
12        ans *= a[i][i];
13        if (ans == 0) return 0;
14        for (int j = i + 1; j < n; j++) {
15            Mod v = a[j][i] / a[i][i];
16            if (v != 0)
17                for (int k = i + 1; k < n; k++)
18                    a[j][k] -= v * a[i][k];
19        }
20    }
21    return ans;
22 }

```

```

1 double det(vector<vector<double>> a) {
2     int n = a.size();
3     double ans = 1;
4     for (int i = 0; i < n; i++) {
5         int b = i;
6         for (int j = i + 1; j < n; j++)
7             if (fabs(a[j][i]) > fabs(a[b][i])) b = j;
8         if (i != b) swap(a[i], a[b]), ans = -ans;
9         ans *= a[i][i];
10        if (ans == 0) return 0;
11        for (int j = i + 1; j < n; j++) {
12            double v = a[j][i] / a[i][i];
13            if (v != 0)
14                for (int k = i + 1; k < n; k++)
15                    a[j][k] -= v * a[i][k];
16        }
17    }
18    return ans;
19 }

```

5.7.2. Inverse

```

1 // Returns rank.
2 // Result is stored in A unless singular (rank < n).
3 // For prime powers, repeatedly set
4 // A^{-1} = A^{-1} (2I - A*A^{-1}) (mod p^k)
5 // where A^{-1} starts as the inverse of A mod p,
6 // and k is doubled in each step.
7
8 int matInv(vector<vector<double>> &A) {
9     int n = sz(A);
10    vi col(n);
11    vector<vector<double>> tmp(n, vector<double>(n));
12    rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
13
14    rep(i, 0, n) {

```

```

17  int r = i, c = i;
18  rep(j, i, n)
19  rep(k, i, n) if (fabs(A[j][k]) > fabs(A[r][c])) r = j,
                                     c = k;

21  if (fabs(A[r][c]) < 1e-12) return i;
22  A[i].swap(A[r]);
23  tmp[i].swap(tmp[r]);
24  rep(j, 0, n) swap(A[j][i], A[j][c]);
25  swap(tmp[j][i], tmp[j][c]);
26  swap(col[i], col[c]);
27  double v = A[i][i];
28  rep(j, i + 1, n) {
29      double f = A[j][i] / v;
30      A[j][i] = 0;
31      rep(k, i + 1, n) A[j][k] -= f * A[i][k];
32      rep(k, 0, n) tmp[j][k] -= f * tmp[i][k];
33  }
34  rep(j, i + 1, n) A[i][j] /= v;
35  rep(j, 0, n) tmp[i][j] /= v;
36  A[i][i] = 1;
37  }

39  for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
40      double v = A[j][i];
41      rep(k, 0, n) tmp[j][k] -= v * tmp[i][k];
42  }

43  rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] = tmp[i][j];
44  return n;
45  }

46  int matInv_mod(vector<vector<ll>> &A) {
47      int n = sz(A);
48      vi col(n);
49      vector<vector<ll>> tmp(n, vector<ll>(n));
50      rep(i, 0, n) tmp[i][i] = 1, col[i] = i;

51  rep(i, 0, n) {
52      int r = i, c = i;
53      rep(j, i, n) rep(k, i, n) if (A[j][k]) {
54          r = j;
55          c = k;
56          goto found;
57      }
58      return i;
59  found:
60      A[i].swap(A[r]);
61      tmp[i].swap(tmp[r]);
62      rep(j, 0, n) swap(A[j][i], A[j][c]);
63      swap(tmp[j][i], tmp[j][c]);
64      swap(col[i], col[c]);
65      ll v = modpow(A[i][i], mod - 2);
66      rep(j, i + 1, n) {
67          ll f = A[j][i] * v % mod;
68          A[j][i] = 0;
69          rep(k, i + 1, n) A[j][k] =
70              (A[j][k] - f * A[i][k]) % mod;
71          rep(k, 0, n) tmp[j][k] =
72              (tmp[j][k] - f * tmp[i][k]) % mod;
73      }
74      rep(j, i + 1, n) A[i][j] = A[i][j] * v % mod;
75      rep(j, 0, n) tmp[i][j] = tmp[i][j] * v % mod;
76      A[i][i] = 1;
77  }

78  for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
79      ll v = A[j][i];
80      rep(k, 0, n) tmp[j][k] =
81          (tmp[j][k] - v * tmp[i][k]) % mod;
82  }

83  rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] =
84      tmp[i][j] % mod + (tmp[i][j] < 0 ? mod : 0);
85  return n;
86  }

```

5.7.3. Characteristic Polynomial

```

1  // calculate det(a - xI)
2  template <typename T>
3  vector<T> CharacteristicPolynomial(vector<vector<T>> a) {
4      int N = a.size();

5  for (int j = 0; j < N - 2; j++) {
6      for (int i = j + 1; i < N; i++) {
7          if (a[i][j] != 0) {
8              swap(a[j + 1], a[i]);

```

```

13  for (int k = 0; k < N; k++)
14      swap(a[k][j + 1], a[k][i]);
15  break;
16  }
17  }
18  if (a[j + 1][j] != 0) {
19      T inv = T(1) / a[j + 1][j];
20      for (int i = j + 2; i < N; i++) {
21          if (a[i][j] == 0) continue;
22          T coe = inv * a[i][j];
23          for (int l = j; l < N; l++)
24              a[i][l] -= coe * a[j + 1][l];
25          for (int k = 0; k < N; k++)
26              a[k][j + 1] += coe * a[k][i];
27      }
28  }
29  }

30  vector<vector<T>> p(N + 1);
31  p[0] = {T(1)};
32  for (int i = 1; i <= N; i++) {
33      p[i].resize(i + 1);
34      for (int j = 0; j < i; j++) {
35          p[i][j + 1] -= p[i - 1][j];
36          p[i][j] += p[i - 1][j] * a[i - 1][i - 1];
37      }
38      T x = 1;
39      for (int m = 1; m < i; m++) {
40          x *= -a[i - m][i - m - 1];
41          T coe = x * a[i - m - 1][i - 1];
42          for (int j = 0; j < i - m; j++)
43              p[i][j] += coe * p[i - m - 1][j];
44      }
45  }
46  return p[N];
47  }

```

5.7.4. Solve Linear Equation

```

1  // solves for x: A * x = b
2  typedef vector<double> vd;
3  const double eps = 1e-12;

4  // solves for x: A * x = b
5  int solveLinear(vector<vd> &A, vd &b, vd &x) {
6      int n = sz(A), m = sz(x), rank = 0, br, bc;
7      if (n) assert(sz(A[0]) == m);
8      vi col(m);
9      iota(all(col), 0);

10  rep(i, 0, n) {
11      double v, bv = 0;
12      rep(r, i, n) rep(c, i, m) if ((v = fabs(A[r][c])) > bv)
13          br = r, bc = c, bv = v;
14      if (bv <= eps) {
15          rep(j, i, n) if (fabs(b[j]) > eps) return -1;
16          break;
17      }
18      swap(A[i], A[br]);
19      swap(b[i], b[br]);
20      swap(col[i], col[bc]);
21      rep(j, 0, n) swap(A[j][i], A[j][bc]);
22      bv = 1 / A[i][i];
23      rep(j, i + 1, n) {
24          double fac = A[j][i] * bv;
25          b[j] -= fac * b[i];
26          rep(k, i + 1, m) A[j][k] -= fac * A[i][k];
27      }
28      rank++;
29  }

30  x.assign(m, 0);
31  for (int i = rank; i--;) {
32      b[i] /= A[i][i];
33      x[col[i]] = b[i];
34      rep(j, 0, i) b[j] -= A[j][i] * b[i];
35  }
36  return rank; // (multiple solutions if rank < m)
37  }

```

5.8. Polynomial Interpolation

```

1  // returns a, such that a[0]x^0 + a[1]x^1 + a[2]x^2 + ...
2  // passes through the given points
3  typedef vector<double> vd;
4  vd interpolate(vd x, vd y, int n) {
5      vd res(n), temp(n);

```



```

rep(k, 0, n - 1) rep(i, k + 1, n) y[i] =
(y[i] - y[k]) / (x[i] - x[k]);
double last = 0;
temp[0] = 1;
rep(k, 0, n) rep(i, 0, n) {
    res[i] += y[k] * temp[i];
    swap(last, temp[i]);
    temp[i] -= last * x[k];
}
return res;
}

```

5.9. Simplex Algorithm

```

// Two-phase simplex algorithm for solving linear programs
// of the form
//
//      maximize      c^T x
//      subject to    Ax <= b
//                  x >= 0
//
// INPUT: A -- an m x n matrix
//        b -- an m-dimensional vector
//        c -- an n-dimensional vector
//        x -- a vector where the optimal solution will be
//        stored
//
// OUTPUT: value of the optimal solution (infinity if
// unbounded
//         above, nan if infeasible)
//
// To use this code, create an LPSolver object with A, b,
// and c as arguments. Then, call Solve(x).

typedef long double ld;
typedef vector<ld> vd;
typedef vector<vd> vvd;
typedef vector<int> vi;

const ld EPS = 1e-9;

struct LPSolver {
    int m, n;
    vi B, N;
    vvd D;

    LPSolver(const vvd &A, const vd &b, const vd &c)
        : m(b.size()), n(c.size()), N(n + 1), B(m),
          D(m + 2, vd(n + 2)) {
        for (int i = 0; i < m; i++)
            for (int j = 0; j < n; j++) D[i][j] = A[i][j];
        for (int i = 0; i < m; i++) {
            B[i] = n + i;
            D[i][n] = -1;
            D[i][n + 1] = b[i];
        }
        for (int j = 0; j < n; j++) {
            N[j] = j;
            D[m][j] = -c[j];
        }
        N[n] = -1;
        D[m + 1][n] = 1;
    }

    void Pivot(int r, int s) {
        double inv = 1.0 / D[r][s];
        for (int i = 0; i < m + 2; i++)
            if (i != r)
                for (int j = 0; j < n + 2; j++)
                    if (j != s) D[i][j] -= D[r][j] * D[i][s] * inv;
        for (int j = 0; j < n + 2; j++)
            if (j != s) D[r][j] *= inv;
        for (int i = 0; i < m + 2; i++)
            if (i != r) D[i][s] *= -inv;
        D[r][s] = inv;
        swap(B[r], N[s]);
    }

    bool Simplex(int phase) {
        int x = phase == 1 ? m + 1 : m;
        while (true) {
            int s = -1;
            for (int j = 0; j <= n; j++) {
                if (phase == 2 && N[j] == -1) continue;
                if (s == -1 || D[x][j] < D[x][s] ||
                    D[x][j] == D[x][s] && N[j] < N[s])
                    s = j;
            }
            if (D[x][s] > -EPS) return true;
            int r = -1;
            for (int i = 0; i < m; i++) {

```

```

                if (D[i][s] < EPS) continue;
                if (r == -1 ||
                    D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
                    (D[i][n + 1] / D[i][s]) ==
                    (D[r][n + 1] / D[r][s]) &&
                    B[i] < B[r])
                    r = i;
            }
            if (r == -1) return false;
            Pivot(r, s);
        }
    }

    ld Solve(vd &x) {
        int r = 0;
        for (int i = 1; i < m; i++)
            if (D[i][n + 1] < D[r][n + 1]) r = i;
        if (D[r][n + 1] < -EPS) {
            Pivot(r, n);
            if (!Simplex(1) || D[m + 1][n + 1] < -EPS)
                return numeric_limits<ld>::infinity();
            for (int i = 0; i < m; i++)
                if (B[i] == -1) {
                    int s = -1;
                    for (int j = 0; j <= n; j++)
                        if (s == -1 || D[i][j] < D[i][s] ||
                            D[i][j] == D[i][s] && N[j] < N[s])
                            s = j;
                    Pivot(i, s);
                }
            if (!Simplex(2)) return numeric_limits<ld>::infinity();
            x = vd(n);
            for (int i = 0; i < m; i++)
                if (B[i] < n) x[B[i]] = D[i][n + 1];
            return D[m][n + 1];
        }
    }

    int main() {
        const int m = 4;
        const int n = 3;
        ld _A[m][n] = {
            {6, -1, 0}, {-1, -5, 0}, {1, 5, 1}, {-1, -5, -1}};
        ld _b[m] = {10, -4, 5, -5};
        ld _c[n] = {1, -1, 0};

        vvd A(m);
        vd b(_b, _b + m);
        vd c(_c, _c + n);
        for (int i = 0; i < m; i++) A[i] = vd(_A[i], _A[i] + n);

        LPSolver solver(A, b, c);
        vd x;
        ld value = solver.Solve(x);

        cerr << "VALUE: " << value << endl; // VALUE: 1.29032
        cerr << "SOLUTION: "; // SOLUTION: 1.74194 0.451613 1
        for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
        cerr << endl;
        return 0;
    }
}

```

6. Geometry

6.1. Point

```

template <typename T> struct P {
    T x, y;
    P(T x = 0, T y = 0) : x(x), y(y) {}
    bool operator<(const P &p) const {
        return tie(x, y) < tie(p.x, p.y);
    }
    bool operator==(const P &p) const {
        return tie(x, y) == tie(p.x, p.y);
    }
    P operator-() const { return {-x, -y}; }
    P operator+(P p) const { return {x + p.x, y + p.y}; }
    P operator-(P p) const { return {x - p.x, y - p.y}; }
    P operator*(T d) const { return {x * d, y * d}; }
    P operator/(T d) const { return {x / d, y / d}; }
    T dist2() const { return x * x + y * y; }
    double len() const { return sqrt(dist2()); }
    P unit() const { return *this / len(); }
    friend T dot(P a, P b) { return a.x * b.x + a.y * b.y; }
    friend T cross(P a, P b) { return a.x * b.y - a.y * b.x; }
    friend T cross(P a, P b, P o) {
        return cross(a - o, b - o);
    }
}

```

```
23 };
    using pt = P<ll>;
```

6.1.1. Quarternion

```
1 constexpr double PI = 3.141592653589793;
  constexpr double EPS = 1e-7;
3 struct Q {
    using T = double;
5     T x, y, z, r;
    Q(T r = 0) : x(0), y(0), z(0), r(r) {}
7     Q(T x, T y, T z, T r = 0) : x(x), y(y), z(z), r(r) {}
    friend bool operator==(const Q &a, const Q &b) {
        return (a - b).abs2() <= EPS;
9     }
    friend bool operator!=(const Q &a, const Q &b) {
        return !(a == b);
11    }
    Q operator-() { return Q(-x, -y, -z, -r); }
13    Q operator+(const Q &b) const {
        return Q(x + b.x, y + b.y, z + b.z, r + b.r);
15    }
    Q operator-(const Q &b) const {
        return Q(x - b.x, y - b.y, z - b.z, r - b.r);
17    }
    Q operator*(const T &t) const {
        return Q(x * t, y * t, z * t, r * t);
19    }
    Q operator*(const Q &b) const {
        return Q(r * b.x + x * b.r + y * b.z - z * b.y,
21                r * b.y - x * b.z + y * b.r + z * b.x,
                r * b.z + x * b.y - y * b.x + z * b.r,
23                r * b.r - x * b.x - y * b.y - z * b.z);
25    }
    Q operator/(const Q &b) const { return *this * b.inv(); }
31    T abs2() const { return r * r + x * x + y * y + z * z; }
    T len() const { return sqrt(abs2()); }
33    Q conj() const { return Q(-x, -y, -z, r); }
    Q unit() const { return *this * (1.0 / len()); }
35    Q inv() const { return conj() * (1.0 / abs2()); }
    friend T dot(Q a, Q b) {
        return a.x * b.x + a.y * b.y + a.z * b.z;
37    }
    friend Q cross(Q a, Q b) {
        return Q(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z,
41                a.x * b.y - a.y * b.x);
43    }
    friend Q rotation_around(Q axis, T angle) {
        return axis.unit() * sin(angle / 2) + cos(angle / 2);
45    }
    Q rotated_around(Q axis, T angle) {
        Q u = rotation_around(axis, angle);
47        return u * *this / u;
49    }
    friend Q rotation_between(Q a, Q b) {
        a = a.unit(), b = b.unit();
51        if (a == -b) {
            // degenerate case
            Q ortho = abs(a.y) > EPS ? cross(a, Q(1, 0, 0))
53                : cross(a, Q(0, 1, 0));
            return rotation_around(ortho, PI);
55        }
        return (a * (a + b)).conj();
57    }
};
```

6.1.2. Spherical Coordinates

```
1 struct car_p {
    double x, y, z;
3 };
  struct sph_p {
    double r, theta, phi;
5 };
  sph_p conv(car_p p) {
9     double r = sqrt(p.x * p.x + p.y * p.y + p.z * p.z);
    double theta = asin(p.y / r);
    double phi = atan2(p.y, p.x);
11    return {r, theta, phi};
13 }
  car_p conv(sph_p p) {
15     double x = p.r * cos(p.theta) * sin(p.phi);
    double y = p.r * cos(p.theta) * cos(p.phi);
17     double z = p.r * sin(p.theta);
    return {x, y, z};
19 }
```

6.2. Segments

```
1 // for non-collinear ABCD, if segments AB and CD intersect
bool intersects(pt a, pt b, pt c, pt d) {
3     if (cross(b, c, a) * cross(b, d, a) > 0) return false;
    if (cross(d, a, c) * cross(d, b, c) > 0) return false;
5     return true;
6 }
  // the intersection point of lines AB and CD
pt intersect(pt a, pt b, pt c, pt d) {
9     auto x = cross(b, c, a), y = cross(b, d, a);
    if (x == y) {
11        // if(abs(x, y) < 1e-8) {
        // is parallel
13    } else {
        return d * (x / (x - y)) - c * (y / (x - y));
15    }
}
```

6.3. Convex Hull

```
1 // returns a convex hull in counterclockwise order
  // for a non-strict one, change cross >= to >
vector<pt> convex_hull(vector<pt> p) {
3     sort(ALL(p));
    if (p[0] == p.back()) return {p[0]};
5     int n = p.size(), t = 0;
    vector<pt> h(n + 1);
7     for (int _ = 2, s = 0; _--; s = --t, reverse(ALL(p)))
        for (pt i : p) {
9             while (t > s + 1 && cross(i, h[t - 1], h[t - 2]) >= 0)
                t--;
11            h[t++] = i;
13        }
    return h.resize(t), h;
15 }
```

6.3.1. 3D Hull

```
1
3 typedef Point3D<double> P3;
4
5 struct PR {
    void ins(int x) { (a == -1 ? a : b) = x; }
7     void rem(int x) { (a == x ? a : b) = -1; }
    int cnt() { return (a != -1) + (b != -1); }
9     int a, b;
10 };
11
12 struct F {
13     P3 q;
    int a, b, c;
15 };
16
17 vector<F> hull3d(const vector<P3> &A) {
    assert(sz(A) >= 4);
19     vector<vector<PR>> E(sz(A), vector<PR>(sz(A), {-1, -1}));
    #define E(x, y) E[f.x][f.y]
21     vector<F> FS;
    auto mf = [&](int i, int j, int k, int l) {
23         P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
        if (q.dot(A[l]) > q.dot(A[i])) q = q * -1;
25         F f{q, i, j, k};
        E(a, b).ins(k);
27         E(a, c).ins(j);
        E(b, c).ins(i);
29         FS.push_back(f);
30     };
    rep(i, 0, 4) rep(j, i + 1, 4) rep(k, j + 1, 4)
31     mf(i, j, k, 6 - i - j - k);
32
33     rep(i, 4, sz(A)) {
        rep(j, 0, sz(FS)) {
35             F f = FS[j];
            if (f.q.dot(A[i]) > f.q.dot(A[f.a])) {
37                 E(a, b).rem(f.c);
                E(a, c).rem(f.b);
39                 E(b, c).rem(f.a);
                swap(FS[j--], FS.back());
41                 FS.pop_back();
42             }
43         }
        int nw = sz(FS);
45         rep(j, 0, nw) {
            F f = FS[j];
47             #define C(a, b, c)
            if (E(a, b).cnt() != 2) mf(f.a, f.b, i, f.c);
49             C(a, b, c);
            C(a, c, b);
51             C(b, c, a);
        }
```

```

53     }
54 }
55 for (F &it : FS)
56     if ((A[it.b] - A[it.a])
57         .cross(A[it.c] - A[it.a])
58         .dot(it.q) <= 0)
59         swap(it.c, it.b);
60 return FS;
61 };

```

6.4. Angular Sort

```

1 auto angle_cmp = [](const pt &a, const pt &b) {
2     auto btm = [](const pt &a) {
3         return a.y < 0 || (a.y == 0 && a.x < 0);
4     };
5     return make_tuple(btm(a), a.y * b.x, abs2(a)) <
6         make_tuple(btm(b), a.x * b.y, abs2(b));
7 };
8 void angular_sort(vector<pt> &p) {
9     sort(p.begin(), p.end(), angle_cmp);
10 }

```

6.5. Convex Polygon Minkowski Sum

```

1 // O(n) convex polygon minkowski sum
2 // must be sorted and counterclockwise
3 vector<pt> minkowski_sum(vector<pt> p, vector<pt> q) {
4     auto diff = [](vector<pt> &c) {
5         auto rcmp = [](pt a, pt b) {
6             return pt{a.y, a.x} < pt{b.y, b.x};
7         };
8         rotate(c.begin(), min_element(ALL(c), rcmp), c.end());
9         c.push_back(c[0]);
10        vector<pt> ret;
11        for (int i = 1; i < c.size(); i++)
12            ret.push_back(c[i] - c[i - 1]);
13        return ret;
14    };
15    auto dp = diff(p), dq = diff(q);
16    pt cur = p[0] + q[0];
17    vector<pt> d(dp.size() + dq.size()), ret = {cur};
18    // include angle_cmp from angular_sort.cpp
19    merge(ALL(dp), ALL(dq), d.begin(), angle_cmp);
20    // optional: make ret strictly convex (UB if degenerate)
21    int now = 0;
22    for (int i = 1; i < d.size(); i++) {
23        if (cross(d[i], d[now]) == 0) d[now] = d[now] + d[i];
24        else d[++now] = d[i];
25    }
26    d.resize(now + 1);
27    // end optional part
28    for (pt v : d) ret.push_back(cur = cur + v);
29    return ret.pop_back(), ret;
30 }

```

6.6. Point In Polygon

```

1 bool on_segment(pt a, pt b, pt p) {
2     return cross(a, b, p) == 0 && dot((p - a), (p - b)) <= 0;
3 }
4 // p can be any polygon, but this is O(n)
5 bool inside(const vector<pt> &p, pt a) {
6     int cnt = 0, n = p.size();
7     for (int i = 0; i < n; i++) {
8         pt l = p[i], r = p[(i + 1) % n];
9         // change to return 0; for strict version
10        if (on_segment(l, r, a)) return 1;
11        cnt ^= ((a.y < l.y) - (a.y < r.y)) * cross(l, r, a) > 0;
12    }
13    return cnt;
14 }

```

6.6.1. Convex Version

```

1 // no preprocessing version
2 // p must be a strict convex hull, counterclockwise
3 // if point is inside or on border
4 bool is_inside(const vector<pt> &c, pt p) {
5     int n = c.size(), l = 1, r = n - 1;
6     if (cross(c[0], c[1], p) < 0) return false;
7     if (cross(c[n - 1], c[0], p) < 0) return false;
8     while (l < r - 1) {
9         int m = (l + r) / 2;
10        T a = cross(c[0], c[m], p);
11        if (a > 0) l = m;
12        else if (a < 0) r = m;
13        else return dot(c[0] - p, c[m] - p) <= 0;
14    }
15    if (l == r) return dot(c[0] - p, c[l] - p) <= 0;

```

```

17 } else return cross(c[l], c[r], p) >= 0;
18 }
19 // with preprocessing version
20 vector<pt> vecs;
21 pt center;
22 // p must be a strict convex hull, counterclockwise
23 // BEWARE OF OVERFLOWS!!
24 void preprocess(vector<pt> p) {
25     for (auto &v : p) v = v * 3;
26     center = p[0] + p[1] + p[2];
27     center.x /= 3, center.y /= 3;
28     for (auto &v : p) v = v - center;
29     vecs = (angular_sort(p), p);
30 }
31 bool intersect_strict(pt a, pt b, pt c, pt d) {
32     if (cross(b, c, a) * cross(b, d, a) > 0) return false;
33     if (cross(d, a, c) * cross(d, b, c) >= 0) return false;
34     return true;
35 }
36 // if point is inside or on border
37 bool query(pt p) {
38     p = p * 3 - center;
39     auto pr = upper_bound(ALL(vecs), p, angle_cmp);
40     if (pr == vecs.end()) pr = vecs.begin();
41     auto pl = (pr == vecs.begin()) ? vecs.back() : *(pr - 1);
42     return !intersect_strict({0, 0}, p, pl, *pr);
43 }

```

6.6.2. Offline Multiple Points Version

Requires: GNU PBDS, Point

```

1
2
3
4
5 using Double = __float128;
6 using Point = pt<Double, Double>;
7
8
9 int n, m;
10 vector<Point> poly;
11 vector<Point> query;
12 vector<int> ans;
13
14 struct Segment {
15     Point a, b;
16     int id;
17 };
18 vector<Segment> segs;
19
20 Double Xnow;
21 inline Double get_y(const Segment &u, Double xnow = Xnow) {
22     const Point &a = u.a;
23     const Point &b = u.b;
24     return (a.y * (b.x - xnow) + b.y * (xnow - a.x)) /
25         (b.x - a.x);
26 }
27 bool operator<(Segment u, Segment v) {
28     Double yu = get_y(u);
29     Double yv = get_y(v);
30     if (yu != yv) return yu < yv;
31     return u.id < v.id;
32 }
33 ordered_map<Segment> st;
34
35 struct Event {
36     int type; // +1 insert seg, -1 remove seg, 0 query
37     Double x, y;
38     int id;
39 };
40 bool operator<(Event a, Event b) {
41     if (a.x != b.x) return a.x < b.x;
42     if (a.type != b.type) return a.type < b.type;
43     return a.y < b.y;
44 }
45 vector<Event> events;
46
47 void solve() {
48     set<Double> xs;
49     set<Point> ps;
50     for (int i = 0; i < n; i++) {
51         xs.insert(poly[i].x);
52         ps.insert(poly[i]);
53     }
54     for (int i = 0; i < n; i++) {
55         Segment s{poly[i], poly[(i + 1) % n], i};
56         if (s.a.x > s.b.x ||
57             (s.a.x == s.b.x && s.a.y > s.b.y)) {
58             swap(s.a, s.b);
59         }

```

```

    segs.push_back(s);
61
    if (s.a.x != s.b.x) {
63        events.push_back({+1, s.a.x + 0.2, s.a.y, i});
        events.push_back({-1, s.b.x - 0.2, s.b.y, i});
65    }
}
67 for (int i = 0; i < m; i++) {
    events.push_back({0, query[i].x, query[i].y, i});
69 }
sort(events.begin(), events.end());
71 int cnt = 0;
for (Event e : events) {
73     int i = e.id;
    Xnow = e.x;
75     if (e.type == 0) {
        Double x = e.x;
77         Double y = e.y;
        Segment tmp = {{x - 1, y}, {x + 1, y}, -1};
79         auto it = st.lower_bound(tmp);

81         if (ps.count(query[i]) > 0) {
            ans[i] = 0;
83         } else if (xs.count(x) > 0) {
            ans[i] = -2;
85         } else if (it != st.end() &&
            get_y(*it) == get_y(tmp)) {
87             ans[i] = 0;
89         } else if (it != st.begin() &&
            get_y(*prev(it)) == get_y(tmp)) {
            ans[i] = 0;
91         } else {
            int rk = st.order_of_key(tmp);
93             if (rk % 2 == 1) {
                ans[i] = 1;
95             } else {
                ans[i] = -1;
97             }
99         }
        } else if (e.type == 1) {
            st.insert(segs[i]);
101             assert((int)st.size() == ++cnt);
        } else if (e.type == -1) {
            st.erase(segs[i]);
103             assert((int)st.size() == --cnt);
105         }
    }
107 }

```

6.7. Closest Pair

```

1 vector<pll> p; // sort by x first!
bool cmpy(const pll &a, const pll &b) const {
3     return a.y < b.y;
}
5 ll sq(ll x) { return x * x; }
// returns (minimum dist)^2 in [l, r)
7 ll solve(int l, int r) {
    if (r - l <= 1) return l * l;
9     int m = (l + r) / 2;
    ll mid = p[m].x, d = min(solve(l, m), solve(m, r));
11     auto pb = p.begin();
    inplace_merge(pb + l, pb + m, pb + r, cmpy);
13     vector<pll> s;
    for (int i = l; i < r; i++)
15         if (sq(p[i].x - mid) < d) s.push_back(p[i]);
    for (int i = 0; i < s.size(); i++)
17         for (int j = i + 1;
            j < s.size() && sq(s[j].y - s[i].y) < d; j++)
19             d = min(d, dis(s[i], s[j]));
21     return d;
}

```

6.8. Minimum Enclosing Circle

```

1
3 typedef Point<double> P;
double ccRadius(const P &A, const P &B, const P &C) {
5     return (B - A).dist() * (C - B).dist() * (A - C).dist() /
        abs((B - A).cross(C - A)) / 2;
7 }
P ccCenter(const P &A, const P &B, const P &C) {
9     P b = C - A, c = B - A;
    return A + (b * c.dist2() - c * b.dist2()).perp() /
11         b.cross(c) / 2;
}
13 pair<P, double> mec(vector<P> ps) {
    shuffle(all(ps), mt19937(time(0)));
15     P o = ps[0];
    double r = 0, EPS = 1 + 1e-8;

```

```

17     rep(i, 0, sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
        o = ps[i], r = 0;
19         rep(j, 0, i) if ((o - ps[j]).dist() > r * EPS) {
            o = (ps[i] + ps[j]) / 2;
21             r = (o - ps[i]).dist();
            rep(k, 0, j) if ((o - ps[k]).dist() > r * EPS) {
23                 o = ccCenter(ps[i], ps[j], ps[k]);
                r = (o - ps[i]).dist();
25             }
27         }
        return {o, r};
29     }
}

```

6.9. Delaunay Triangulation

```

1
3 typedef Point<ll> P;
typedef struct Quad *Q;
5 typedef __int128 t ll; // (can be ll if coords are < 2e4)
P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
7
9 struct Quad {
    bool mark;
    Q o, rot;
    P p;
11     P F() { return r()->p; }
    Q r() { return rot->rot; }
    Q prev() { return rot->o->rot; }
    Q next() { return r()->prev(); }
13 };
15 bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
    ll p2 = p.dist2(), A = a.dist2() - p2,
17     B = b.dist2() - p2, C = c.dist2() - p2;
    return p.cross(a, b) * C + p.cross(b, c) * A +
21     p.cross(c, a) * B > 0;
23 }
25 Q makeEdge(P orig, P dest) {
    Q q[] = {new Quad{0, 0, 0, orig}, new Quad{0, 0, 0, arb},
27         new Quad{0, 0, 0, dest}, new Quad{0, 0, 0, arb}};
    rep(i, 0, 4) q[i]->o = q[i-1] & 3,
29     q[i]->rot = q[(i + 1) & 3];
    return *q;
31 }
33 void splice(Q a, Q b) {
    swap(a->o->rot->o, b->o->rot->o);
    swap(a->o, b->o);
35 }
37 Q connect(Q a, Q b) {
    Q q = makeEdge(a->F(), b->p);
    splice(q, a->next());
    splice(q->r(), b);
39     return q;
41 }
43 pair<Q, Q> rec(const vector<P> &s) {
    if (sz(s) <= 3) {
45         Q a = makeEdge(s[0], s[1]),
            b = makeEdge(s[1], s.back());
47         if (sz(s) == 2) return {a, a->r()};
        splice(a->r(), b);
49         auto side = s[0].cross(s[1], s[2]);
        Q c = side ? connect(b, a) : 0;
51         return {side < 0 ? c->r() : a, side < 0 ? c : b->r()};
53     }
55 #define H(e) e->F(), e->p
    #define valid(e) (e->F().cross(H(base)) > 0)
    Q A, B, ra, rb;
57     int half = sz(s) / 2;
    tie(ra, A) = rec({all(s) - half});
    tie(B, rb) = rec({sz(s) - half + all(s)});
59     while ((B->p.cross(H(A)) < 0 && (A = A->next())) ||
        (A->p.cross(H(B)) > 0 && (B = B->r()->o)))
61         ;
    Q base = connect(B->r(), A);
    if (A->p == ra->p) ra = base->r();
63     if (B->p == rb->p) rb = base;
65
67 #define DEL(e, init, dir)
    Q e = init->dir;
69     if (valid(e))
        while (circ(e->dir->F(), H(base), e->F())) {
71         Q t = e->dir;
            splice(e, e->prev());
            splice(e->r(), e->r()->prev());
73         e = t;
75     }
}

```

```

77     DEL(LC, base->r(), o);
       DEL(RC, base, prev());
79     if (!valid(LC) && !valid(RC)) break;
       if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
81         base = connect(RC, base->r());
       else base = connect(base->r(), LC->r());
83 }
85 return {ra, rb};
87 }
// returns [A_0, B_0, C_0, A_1, B_1, ...]
// where A_i, B_i, C_i are counter-clockwise triangles
89 vector<P> triangulate(vector<P> pts) {
    sort(all(pts));
    assert(unique(all(pts)) == pts.end());
    if (sz(pts) < 2) return {};
    Q e = rec(pts).first;
    vector<Q> q = {e};
    int qi = 0;
    while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
97 #define ADD
    {
99         Q c = e;
        do {
101             c->mark = 1;
            pts.push_back(c->p);
103             q.push_back(c->r());
            c = c->next();
105         } while (c != e);
    }
107 ADD;
    pts.clear();
109 while (qi < sz(q))
    if (!(e = q[qi++])->mark) ADD;
111 return pts;
}

```

6.9.1. Slower Version

```

1
3 template<class P, class F>
void delaunay(vector<P> &ps, F trifun) {
5     if (sz(ps) == 3) {
        int d = (ps[0].cross(ps[1], ps[2]) < 0);
        trifun(0, 1 + d, 2 - d);
    }
    vector<P3> p3;
    for (P p : ps) p3.emplace_back(p.x, p.y, p.dist2());
    if (sz(ps) > 3)
        for (auto t : hull3d(p3))
            if ((p3[t.b] - p3[t.a])
                .cross(p3[t.c] - p3[t.a])
                .dot(P3(0, 0, 1)) < 0)
                trifun(t.a, t.c, t.b);
17 }

```

6.10. Half Plane Intersection

```

1 struct Line {
    Point P;
    Vector v;
    bool operator<(const Line &b) const {
        return atan2(v.y, v.x) < atan2(b.v.y, b.v.x);
    }
};
9 bool OnLeft(const Line &L, const Point &p) {
    return Cross(L.v, p - L.P) > 0;
}
11 Point GetIntersection(Line a, Line b) {
    Vector u = a.P - b.P;
    Double t = Cross(b.v, u) / Cross(a.v, b.v);
    return a.P + a.v * t;
15 }
int HalfplaneIntersection(Line *L, int n, Point *poly) {
    sort(L, L + n);
19     int first, last;
    Point *p = new Point[n];
    Line *q = new Line[n];
    q[first = last = 0] = L[0];
    for (int i = 1; i < n; i++) {
        while (first < last && !OnLeft(L[i], p[last - 1]))
            last--;
        while (first < last && !OnLeft(L[i], p[first])) first++;
        q[++last] = L[i];
        if (fabs(Cross(q[last].v, q[last - 1].v)) < EPS) {
            last--;
            if (OnLeft(q[last], L[i].P)) q[last] = L[i];
31     }
}

```

```

       if (first < last)
           p[last - 1] = GetIntersection(q[last - 1], q[last]);
    }
    while (first < last && !OnLeft(q[first], p[last - 1]))
        last--;
    if (last - first <= 1) return 0;
    p[last] = GetIntersection(q[last], q[first]);
39     int m = 0;
    for (int i = first; i <= last; i++) poly[m++] = p[i];
    return m;
43 }

```

7. Strings

7.1. Knuth-Morris-Pratt Algorithm

```

1
3 vector<int> pi(const string &s) {
    vector<int> p(s.size());
    for (int i = 1; i < s.size(); i++) {
        int g = p[i - 1];
        while (g && s[i] != s[g]) g = p[g - 1];
        p[i] = g + (s[i] == s[g]);
    }
    return p;
11 }
vector<int> match(const string &s, const string &pat) {
    vector<int> p = pi(pat + '\0' + s), res;
    for (int i = p.size() - s.size(); i < p.size(); i++)
        if (p[i] == pat.size())
            res.push_back(i - 2 * pat.size());
17     return res;
}

```

7.2. Aho-Corasick Automaton

```

1 struct Aho_Corasick {
    static const int maxc = 26, maxn = 4e5;
    struct NODES {
        int Next[maxc], fail, ans;
    };
    NODES T[maxn];
    int top, qtop, q[maxn];
    int get_node(const int &fail) {
        fill_n(T[top].Next, maxc, 0);
        T[top].fail = fail;
        T[top].ans = 0;
        return top++;
    }
    int insert(const string &s) {
        int ptr = 1;
        for (char c : s) { // change char id
            c -= 'a';
            if (!T[ptr].Next[c]) T[ptr].Next[c] = get_node(ptr);
            ptr = T[ptr].Next[c];
        }
        return ptr;
    } // return ans_last_place
    void build_fail(int ptr) {
        int tmp;
        for (int i = 0; i < maxc; i++)
            if (T[ptr].Next[i]) {
                tmp = T[ptr].fail;
                while (tmp != 1 && !T[tmp].Next[i])
                    tmp = T[tmp].fail;
                if (T[tmp].Next[i] != T[ptr].Next[i])
                    if (T[tmp].Next[i]) tmp = T[tmp].Next[i];
                T[T[ptr].Next[i]].fail = tmp;
                q[qtop++] = T[ptr].Next[i];
            }
    }
    void AC_auto(const string &s) {
        int ptr = 1;
        for (char c : s) {
            while (ptr != 1 && !T[ptr].Next[c]) ptr = T[ptr].fail;
            if (T[ptr].Next[c]) {
                ptr = T[ptr].Next[c];
                T[ptr].ans++;
            }
        }
    }
    void Solve(string &s) {
        for (char &c : s) // change char id
            c -= 'a';
        for (int i = 0; i < qtop; i++) build_fail(q[i]);
        AC_auto(s);
        for (int i = qtop - 1; i > -1; i--)
            T[T[q[i]].fail].ans += T[q[i]].ans;
51     }
}

```



```

53 }
54 void reset() {
55     qtop = top = q[0] = 1;
56     get_node(1);
57 }
58 } AC;
59 // usage example
60 string s, S;
61 int n, t, ans_place[50000];
62 int main() {
63     Tie cin >> t;
64     while (t--) {
65         AC.reset();
66         cin >> S >> n;
67         for (int i = 0; i < n; i++) {
68             cin >> s;
69             ans_place[i] = AC.insert(s);
70         }
71         AC.Solve(S);
72         for (int i = 0; i < n; i++)
73             cout << AC.T[ans_place[i]].ans << '\n';
74     }
75 }

```

7.3. Suffix Array

```

1 // sa[i]: starting index of suffix at rank i
2 // 0-indexed, sa[0] = n (empty string)
3 // lcp[i]: lcp of sa[i] and sa[i - 1], lcp[0] = 0
4 struct SuffixArray {
5     vector<int> sa, lcp;
6     SuffixArray(string &s,
7         int lim = 256) { // or basic_string<int>
8         int n = sz(s) + 1, k = 0, a, b;
9         vector<int> x(all(s) + 1), y(n), ws(max(n, lim)),
10             rank(n);
11         sa = lcp = y, iota(all(sa), 0);
12         for (int j = 0, p = 0; p < n;
13             j = max(1, j * 2), lim = p) {
14             p = j, iota(all(y), n - j);
15             for (int i = 0; i < n; i++)
16                 if (sa[i] >= j) y[p++] = sa[i] - j;
17             fill(all(ws), 0);
18             for (int i = 0; i < n; i++) ws[x[i]]++;
19             for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];
20             for (int i = n; i--;) sa[-ws[x[i]]] = y[i];
21             swap(x, y), p = 1, x[sa[0]] = 0;
22             for (int i = 1; i < n; i++)
23                 a = sa[i - 1], b = sa[i],
24                 x[b] = (y[a] == y[b] && y[a + j] == y[b + j])
25                     ? p - 1 : p++;
26         }
27         for (int i = 1; i < n; i++) rank[sa[i]] = i;
28         for (int i = 0, j; i < n - 1; lcp[rank[i+1]] = k)
29             for (k && k--, j = sa[rank[i] - 1];
30                 s[i + k] == s[j + k]; k++);
31     };
32 };

```

7.4. Suffix Tree

```

1 struct SAM {
2     static const int maxc = 26; // char range
3     static const int maxn = 10010; // string len
4     struct Node {
5         Node *green, *edge[maxc];
6         int max_len, in, times;
7     } * root, *last, reg[maxn * 2];
8     int top;
9     Node *get_node(int _max) {
10         Node *re = &reg[top++];
11         re->in = 0, re->times = 1;
12         re->max_len = _max, re->green = 0;
13         for (int i = 0; i < maxc; i++) re->edge[i] = 0;
14         return re;
15     }
16     void insert(const char c) { // c in range [0, maxc)
17         Node *p = last;
18         last = get_node(p->max_len + 1);
19         while (p && !p->edge[c])
20             p->edge[c] = last, p = p->green;
21         if (!p) last->green = root;
22         else {
23             Node *pot_green = p->edge[c];
24             if ((pot_green->max_len) == (p->max_len + 1))

```

```

25         last->green = pot_green;
26     }
27     else {
28         Node *wish = get_node(p->max_len + 1);
29         wish->times = 0;
30         while (p && p->edge[c] == pot_green)
31             p->edge[c] = wish, p = p->green;
32         for (int i = 0; i < maxc; i++)
33             wish->edge[i] = pot_green->edge[i];
34         wish->green = pot_green->green;
35         pot_green->green = wish;
36         last->green = wish;
37     }
38 }
39 Node *q[maxn * 2];
40 int ql, qr;
41 void get_times(Node *p) {
42     ql = 0, qr = -1, reg[0].in = 1;
43     for (int i = 1; i < top; i++) reg[i].green->in++;
44     for (int i = 0; i < top; i++)
45         if (!reg[i].in) q[++qr] = &reg[i];
46     while (ql <= qr) {
47         q[ql]->green->times += q[ql]->times;
48         if (!(--q[ql]->green->in)) q[++qr] = q[ql]->green;
49         ql++;
50     }
51 }
52 void build(const string &s) {
53     top = 0;
54     root = last = get_node(0);
55     for (char c : s) insert(c - 'a'); // change char id
56     get_times(root);
57 }
58 // call build before solve
59 int solve(const string &s) {
60     Node *p = root;
61     for (char c : s)
62         if (!(p = p->edge[c - 'a'])) // change char id
63             return 0;
64     return p->times;
65 }

```

7.5. Cocke-Younger-Kasami Algorithm

```

1 struct rule {
2     // s -> xy
3     // if y == -1, then s -> x (unit rule)
4     int s, x, y, cost;
5 };
6 int state;
7 // state (id) for each letter (variable)
8 // lowercase letters are terminal symbols
9 map<char, int> rules;
10 vector<rule> cnf;
11 void init() {
12     state = 0;
13     rules.clear();
14     cnf.clear();
15 }
16 // convert a cfg rule to cnf (but with unit rules) and add
17 // it
18 void add_to_cnf(char s, const string &p, int cost) {
19     if (!rules.count(s)) rules[s] = state++;
20     for (char c : p)
21         if (!rules.count(c)) rules[c] = state++;
22         if (p.size() == 1) {
23             cnf.push_back({rules[s], rules[p[0]], -1, cost});
24         } else {
25             // length >= 3 -> split
26             int left = rules[s];
27             int sz = p.size();
28             for (int i = 0; i < sz - 2; i++) {
29                 cnf.push_back({left, rules[p[i]], state, 0});
30                 left = state++;
31             }
32             cnf.push_back(
33                 {left, rules[p[sz - 2]], rules[p[sz - 1]], cost});
34         }
35 }
36 }
37 }
38 constexpr int MAXN = 55;
39 vector<long long> dp[MAXN][MAXN];
40 // unit rules with negative costs can cause negative cycles
41 vector<bool> neg_INF[MAXN][MAXN];
42 void relax(int l, int r, rule c, long long cost,
43     bool neg_c = 0) {
44     if (!neg_INF[l][r][c.s] &&

```

```

47     (neg_INF[l][r][c.x] || cost < dp[l][r][c.s])) {
49     if (neg_c || neg_INF[l][r][c.x]) {
        dp[l][r][c.s] = 0;
        neg_INF[l][r][c.s] = true;
51     } else {
        dp[l][r][c.s] = cost;
53     }
55 }
57 void bellman(int l, int r, int n) {
    for (int k = 1; k <= state; k++)
        for (rule c : cnf)
            if (c.y == -1)
                relax(l, r, c, dp[l][r][c.x] + c.cost, k == n);
61 }
63 void cyk(const string &s) {
    vector<int> tok;
    for (char c : s) tok.push_back(rules[c]);
    for (int i = 0; i < tok.size(); i++) {
        for (int j = 0; j < tok.size(); j++) {
65             dp[i][j] = vector<long long>(state + 1, INT_MAX);
            neg_INF[i][j] = vector<bool>(state + 1, false);
69         }
        dp[i][i][tok[i]] = 0;
        bellman(i, i, tok.size());
71     }
    for (int r = 1; r < tok.size(); r++) {
        for (int l = r - 1; l >= 0; l--) {
73             for (int k = l; k < r; k++)
                for (rule c : cnf)
                    if (c.y != -1)
75                         relax(l, r, c,
                                dp[l][k][c.x] + dp[k + 1][r][c.y] +
                                c.cost);
77             bellman(l, r, tok.size());
81         }
83     }
85 }
87 // usage example
89 int main() {
    init();
    add_to_cnf('S', "aSc", 1);
    add_to_cnf('S', "BBB", 1);
    add_to_cnf('S', "SB", 1);
    add_to_cnf('B', "b", 1);
    cyk("abbbbc");
    // dp[0][s.size() - 1][rules[start]] = min cost to
    // generate s
    cout << dp[0][5][rules['S']] << '\n'; // 7
    cyk("acbc");
    cout << dp[0][3][rules['S']] << '\n'; // INT_MAX
    add_to_cnf('S', "S", -1);
    cyk("abbbbc");
    cout << neg_INF[0][5][rules['S']] << '\n'; // 1
101 }

```

7.6. Z Value

```

1 int z[n];
2 void zval(string s) {
3     // z[i] => longest common prefix of s and s[i:], i > 0
    int n = s.size();
    z[0] = 0;
    for (int b = 0, i = 1; i < n; i++) {
        if (z[b] + b <= i) z[i] = 0;
        else z[i] = min(z[i - b], z[b] + b - i);
        while (s[i + z[i]] == s[z[i]]) z[i]++;
        if (i + z[i] > b + z[b]) b = i;
9     }
11 }

```

7.7. Manacher's Algorithm

```

1 int z[n];
2 void manacher(string s) {
3     // z[i] => longest odd palindrome centered at i is
    // s[i - z[i]] ... i + z[i]
    // to get all palindromes (including even length),
    // insert a '#' between each s[i] and s[i + 1]
    int n = s.size();
    z[0] = 0;
    for (int b = 0, i = 1; i < n; i++) {
        if (z[b] + b >= i)
            z[i] = min(z[2 * b - i], b + z[b] - i);
        else z[i] = 0;
        while (i + z[i] + 1 < n && i - z[i] - 1 >= 0 &&
            s[i + z[i] + 1] == s[i - z[i] - 1])
            z[i]++;
        if (z[i] + i > z[b] + b) b = i;
13     }
15 }

```

7.8. Minimum Rotation

```

1 int min_rotation(string s) {
    int a = 0, n = s.size();
    s += s;
    for (int b = 0; b < n; b++) {
        for (int k = 0; k < n; k++) {
            if (a + k == b || s[a + k] < s[b + k]) {
                b = max(0, k - 1);
                break;
            }
            if (s[a + k] > s[b + k]) {
                a = b;
                break;
            }
        }
        return a;
17 }

```

7.9. Palindromic Tree

```

1
2
3 struct palindromic_tree {
    struct node {
        int next[26], fail, len;
        int cnt,
        num; // cnt: appear times, num: number of pal. suf.
        node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
            for (int i = 0; i < 26; ++i) next[i] = 0;
        }
    };
    vector<node> St;
    vector<char> s;
    int last, n;
    palindromic_tree() : St(2), last(1), n(0) {
        St[0].fail = 1, St[1].len = -1, s.pb(-1);
    }
    inline void clear() {
        St.clear(), s.clear(), last = 1, n = 0;
        St.pb(0), St.pb(-1);
        St[0].fail = 1, s.pb(-1);
    }
    inline int get_fail(int x) {
        while (s[n - St[x].len - 1] != s[n]) x = St[x].fail;
        return x;
    }
    inline void add(int c) {
        s.push_back(c - 'a'), ++n;
        int cur = get_fail(last);
        if (!St[cur].next[c]) {
            int now = SZ(St);
            St.pb(St[cur].len + 2);
            St[now].fail = St[get_fail(St[cur].fail)].next[c];
            St[cur].next[c] = now;
            St[now].num = St[St[now].fail].num + 1;
        }
        last = St[cur].next[c], ++St[last].cnt;
    }
    inline void count() { // counting cnt
        auto i = St.rbegin();
        for (; i != St.rend(); ++i) {
            St[i->fail].cnt += i->cnt;
        }
    }
    inline int size() { // The number of diff. pal.
        return SZ(St) - 2;
    }
};

```

8. Debug List

- 1 - Pre-submit:
 - Did you make a typo when copying a template?
 - Test more cases if unsure.
 - Write a naive solution and check small cases.
 - Submit the correct file.
- 7 - General Debugging:
 - Read the whole problem again.
 - Have a teammate read the problem.
 - Have a teammate read your code.
 - Explain your solution to them (or a rubber duck).
 - Print the code and its output / debug output.
 - Go to the toilet.
- 15 - Wrong Answer:
 - Any possible overflows?
 - > `'__int128'`?

```
19 - Try `-ftrapv` or `#pragma GCC optimize("trapv")`
- Floating point errors?
- > `long double` ?
21 - turn off math optimizations
- check for `==`, `>=`, `acos(1.000000001)`, etc.
23 - Did you forget to sort or unique?
- Generate large and worst "corner" cases.
25 - Check your `m` / `n`, `i` / `j` and `x` / `y`.
- Are everything initialized or reset properly?
27 - Are you sure about the STL thing you are using?
- Read cppreference (should be available).
29 - Print everything and run it on pen and paper.

31 - Time Limit Exceeded:
- Calculate your time complexity again.
33 - Does the program actually end?
- Check for `while(q.size())` etc.
35 - Test the largest cases locally.
- Did you do unnecessary stuff?
37 - e.g. pass vectors by value
- e.g. `memset` for every test case
39 - Is your constant factor reasonable?

41 - Runtime Error:
- Check memory usage.
43 - Forget to clear or destroy stuff?
- > `vector::shrink_to_fit()`
45 - Stack overflow?
- Bad pointer / array access?
47 - Try `-fsanitize=address`
- Division by zero? NaN's?
```