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## 1. Misc

### 1.1. Contest

#### 1.1.1. Makefile

```

1 .PRECIOUS: ./p%
3 %: p%
5   ulimit -s unlimited && ./$<
5 p%: %.cpp
6   g++ -O $@ <$ -std=c++17 -Wall -Wextra -Wshadow \
7     -fsanitize=address,undefined

```

### 1.2. How Did We Get Here?

#### 1.2.1. Macros

Use vectorizations and math optimizations at your own peril.  
For gcc $\geq$ 9, there are `[[likely]]` and `[[unlikely]]` attributes.  
Call gcc with `-fopt-info-optimized-missed-optall` for optimization info.

```

1 #define _GLIBCXX_DEBUG 1 // for debug mode
2 #define _GLIBCXX_SANITIZE_VECTOR 1 // for asan on vectors
3 #pragma GCC optimize("O3", "unroll-loops")
4 #pragma GCC optimize("fast-math")
5 #pragma GCC target("avx,avx2,abm,bmi,bmi2") // tip: `lscpu`  

// before a loop
7 #pragma GCC unroll 16 // 0 or 1 -> no unrolling
# pragma GCC ivdep

```

#### 1.2.2. Fast I/O

```

1 struct scanner {
2     static constexpr size_t LEN = 32 << 20;
3     char *buf, *buf_ptr, *buf_end;
4     scanner() : buf(new char[LEN]), buf_ptr(buf + LEN),
5                  buf_end(buf + LEN) {}
6     ~scanner() { delete[] buf; }
7     char get() {
8         if (buf_ptr == buf_end) [[unlikely]]
9             buf_end = buf + fread_unlocked(buf, 1, LEN, stdin),
10            buf_ptr = buf;
11        return *(buf_ptr++);
12    }
13    char seek(char del) {
14        char c;
15        while ((c = getc()) < del) {}
16        return c;
17    }
18    void read(int &t) {
19        bool neg = false;
20        char c = seek('-');
21        if (c == '-') neg = true, t = 0;
22        else t = c ^ '0';
23        while ((c = getc()) >= '0') t = t * 10 + (c ^ '0');
24        if (neg) t = -t;
25    }
26    struct printer {
27        static constexpr size_t CPI = 21, LEN = 32 << 20;
28        char *buf, *buf_ptr, *buf_end, *tbuf;
29        char *int_buf, *int_buf_end;
30        printer() : buf(new char[LEN]), buf_ptr(buf),
31                    buf_end(buf + LEN), int_buf(new char[CPI + 1]()),
32                    int_buf_end(int_buf + CPI - 1) {}
33        ~printer() {
34            flush();
35            delete[] buf, delete[] int_buf;
36        }
37        void flush() {
38            fwrite_unlocked(buf, 1, buf_ptr - buf, stdout);
39            buf_ptr = buf;
40        }
41        void write_(const char &c) {
42            *buf_ptr = c;
43            if (++buf_ptr == buf_end) [[unlikely]]
44                flush();
45        }
46        void write_(const char *s) {
47            for (; *s != '\0'; ++s) write_(*s);
48        }
49        void write(int x) {
50            if (x < 0) write_('-'), x = -x;
51            if (x == 0) [[unlikely]]
52                return write_('0');
53            for (tbuf = int_buf_end; x != 0; --tbuf, x /= 10)
54                *tbuf = '0' + char(x % 10);
55            write_(++tbuf);
56        }
57    };
58 };

```

## Kotlin

```

1 import java.io.*
2 import java.util.*
3
4 @JvmField val cin = System.`in`.bufferedReader()
5 @JvmField val cout = PrintWriter(System.out, false)
6 @JvmField var tokenizer: StringTokenizer = StringTokenizer("")
7 fun nextLine() = cin.readLine()!!
8 fun read(): String {
9     while (!tokenizer.hasMoreTokens())
10        tokenizer = StringTokenizer(nextLine())
11    return tokenizer.nextToken()
12}
13
14 // example
15 fun main() {
16     val n = read().toInt()
17     val a = DoubleArray(n) { read().toDouble() }
18     cout.println("omg hi")
19     cout.flush()
20 }

```

#### 1.2.3. constexpr

Some default limits in gcc (7.x - trunk):

- constexpr recursion depth: 512
- constexpr loop iteration per function: 262144
- constexpr operation count per function: 33554432
- template recursion depth: 900 (gcc might segfault first)

```

1 constexpr array<int, 10> fibonacci[] {
2     array<int, 10> a{};
3     a[0] = a[1] = 1;
4     for (int i = 2; i < 10; i++) a[i] = a[i - 1] + a[i - 2];
5     return a;
6 }();
7 static_assert(fibonacci[9] == 55, "CE");
8
9 template <typename F, typename INT, INT... S>
10 constexpr void for_constexpr(integer_sequence<INT, S...>,  

11                             F &&func) {
12     int _[] = {{func(integral_constant<INT, S>{})}, 0}...
13 }
14 // example
15 template <typename... T> void print_tuple(tuple<T...> t) {
16     for_constexpr(make_index_sequence<sizeof...(T)>{},  

17                  [&](auto i) { cout << get<i>(t) << '\n'; });
18 }

```

#### 1.2.4. Bump Allocator

```

1
2
3 // global bump allocator
4 char mem[256 << 20]; // 256 MB
5 size_t rsp = sizeof(mem);
6 void *operator new(size_t s) {
7     assert(s < rsp); // MLE
8     return (void *)&mem[rsp -= s];
9 }
10 void operator delete(void *) {}
11
12 // bump allocator for STL / pbds containers
13 char mem[256 << 20];
14 size_t rsp = sizeof(mem);
15 template <typename T> struct bump {
16     typedef T value_type;
17     bump() {}
18     template <typename U> bump(U, ...) {}
19     T *allocate(size_t n) {
20         rsp -= n * sizeof(T);
21         rsp &= 0 - alignof(T);
22         return (T *) (mem + rsp);
23     }
24     void deallocate(T *, size_t n) {}
25 }

```

## 1.3. Tools

### 1.3.1. Floating Point Binary Search

```

1 union di {
2     double d;
3     ull i;
4 };
5 bool check(double);
6 // binary search in [L, R) with relative error 2^-eps
7 double binary_search(double L, double R, int eps) {
8     di l = {L}, r = {R}, m;
9     while (r.i - l.i > 1LL << (52 - eps)) {
10         m.i = (l.i + r.i) >> 1;
11     }
12 }

```

```

11     if (check(m.d)) r = m;
12     else l = m;
13 }
14 return l.d;
15 }
```

### 1.3.2. SplitMix64

```

1 using ull = unsigned long long;
2 inline ull splitmix64(ull x) {
3     // change to `static ull x = SEED;` for DRBG
4     ull z = (x += 0x9E3779B97F4A7C15);
5     z = (z ^ (z >> 30)) * 0xBF58476D1CE4E5B9;
6     z = (z ^ (z >> 27)) * 0x94D049BB133111EB;
7     return z ^ (z >> 31);
}
```

### 1.3.3. <random>

```

1 #ifdef __unix__
2 random_device rd;
3 mt19937_64 RNG(rd());
4 #else
5 const auto SEED = chrono::high_resolution_clock::now()
6         .time_since_epoch()
7         .count();
8 mt19937_64 RNG(SEED);
9 #endif
10 // random uint_fast64_t: RNG();
11 // uniform random of type T (int, double, ...) in [l, r];
12 // uniform_int_distribution<T> dist(l, r); dist(RNG);
```

### 1.3.4. x86 Stack Hack

```

1 constexpr size_t size = 200 << 20; // 200MiB
2 int main() {
3     register long rsp asm("rsp");
4     char *buf = new char[size];
5     asm("movq %0, %%rsp\n" :: "r"(buf + size));
6     // do stuff
7     asm("movq %0, %%rsp\n" :: "r"(rsp));
8     delete[] buf;
9 }
```

## 1.4. Algorithms

### 1.4.1. Bit Hacks

```

1 // next permutation of x as a bit sequence
2 ull next_bits_permutation(ull x) {
3     ull c = __builtin_ctzll(x), r = x + (1ULL << c);
4     return (r ^ x) >> (c + 2) | r;
5 }
6 // iterate over all (proper) subsets of bitset s
7 void subsets(ull s) {
8     for (ull x = s; x;) { --x &= s; /* do stuff */ }
9 }
```

### 1.4.2. Aliens Trick

```

1 // min dp[i] value and its i (smallest one)
2 pll get_dp(int cost);
3 ll aliens(int k, int l, int r) {
4     while (l != r) {
5         int m = (l + r) / 2;
6         auto [f, s] = get_dp(m);
7         if (s == k) return f - m * k;
8         if (s < k) r = m;
9         else l = m + 1;
10    }
11 return get_dp(l).first - l * k;
}
```

### 1.4.3. Hilbert Curve

```

1 ll hilbert(ll n, int x, int y) {
2     ll res = 0;
3     for (ll s = n; s /= 2;) {
4         int rx = !(x & s), ry = !(y & s);
5         res += s * s * ((3 * rx) ^ ry);
6         if (ry == 0) {
7             if (rx == 1) x = s - 1 - x, y = s - 1 - y;
8             swap(x, y);
9         }
10    }
11 return res;
}
```

### 1.4.4. Infinite Grid Knight Distance

```

1 ll get_dist(ll dx, ll dy) {
2     if (++(dx = abs(dx)) > ++(dy = abs(dy))) swap(dx, dy);
3     if (dx == 1 && dy == 2) return 3;
4     if (dx == 3 && dy == 3) return 4;
5     ll lb = max(dy / 2, (dx + dy) / 3);
6     return ((dx ^ dy ^ lb) & 1) ? ++lb : lb;
7 }
```

### 1.4.5. Poker Hand

```

1
2
3
4
5
6
7 using namespace std;
8
9 struct hand {
10     static constexpr auto rk = [] {
11         array<int, 256> x{};
12         auto s = "23456789TJQKACDHS";
13         for (int i = 0; i < 17; i++) x[s[i]] = i % 13;
14         }();
15     vector<pair<int, int>> v;
16     vector<int> cnt, vf, vs;
17     int type;
18     hand() : cnt(4), type(0) {}
19     void add_card(char suit, char rank) {
20         ++cnt[rk[suit]];
21         for (auto &[f, s] : v)
22             if (s == rk[rank]) return ++f, void();
23         v.emplace_back(1, rk[rank]);
24     }
25     void process() {
26         sort(v.rbegin(), v.rend());
27         for (auto [f, s] : v) vf.push_back(f), vs.push_back(s);
28         bool str = 0, flu = find(all(cnt), 5) != cnt.end();
29         if ((str = v.size() == 5))
30             for (int i = 1; i < 5; i++)
31                 if (vs[i] != vs[i - 1] + 1) str = 0;
32         if (vs == vector<int>{12, 3, 2, 1, 0})
33             str = 1, vs = {3, 2, 1, 0, -1};
34         if (str && flu) type = 9;
35         else if (vf[0] == 4) type = 8;
36         else if (vf[0] == 3 && vf[1] == 2) type = 7;
37         else if (str || flu) type = 5 + flu;
38         else if (vf[0] == 3) type = 4;
39         else if (vf[0] == 2) type = 2 + (vf[1] == 2);
40         else type = 1;
41     }
42     bool operator<(const hand &b) const {
43         return make_tuple(type, vf, vs) <
44                make_tuple(b.type, b.vf, b.vs);
45     }
46 }
```

### 1.4.6. Longest Increasing Subsequence

```

1
2
3 template <class I> vi lis(const vector<I> &S) {
4     if (S.empty()) return {};
5     vi prev(sz(S));
6     typedef pair<I, int> p;
7     vector<p> res;
8     rep(i, 0, sz(S)) {
9         // change 0 -> i for longest non-decreasing subsequence
10        auto it = lower_bound(all(res), p{S[i], 0});
11        if (it == res.end())
12            res.emplace_back(), it = res.end() - 1;
13        *it = {S[i], i};
14        prev[i] = it == res.begin() ? 0 : (it - 1)->second;
15    }
16    int L = sz(res), cur = res.back().second;
17    vi ans(L);
18    while (L--) ans[L] = cur, cur = prev[cur];
19 }
```

### 1.4.7. Mo's Algorithm on Tree

```

1 void MoAlgoOnTree() {
2     Dfs(0, -1);
3     vector<int> euler(tk);
4     for (int i = 0; i < n; ++i) {
5         euler[tin[i]] = i;
6         euler[tout[i]] = i;
7     }
8     vector<int> l(q), r(q), qr(q), sp(q, -1);
```

```
9 | for (int i = 0; i < q; ++i) {  
10|   if (tin[u[i]] > tin[v[i]]) swap(u[i], v[i]);  
11|   int z = GetLCA(u[i], v[i]);  
12|   sp[i] = z[i];  
13|   if (z == u) l[i] = tin[u[i]], r[i] = tin[v[i]];  
14|   else l[i] = tout[u[i]], r[i] = tin[v[i]];  
15|   qr[i] = i;  
16| }  
17| sort(qr.begin(), qr.end(), [&](int i, int j) {  
18|   if (l[i] / kB == l[j] / kB) return r[i] < r[j];  
19|   return l[i] / kB < l[j] / kB;  
20|});  
21| vector<bool> used(n);  
22| // Add(v): add/remove v to/from the path based on used[v]  
23| for (int i = 0, tl = 0, tr = -1; i < q; ++i) {  
24|   while (tl < l[qr[i]]) Add(euler[tl++]);  
25|   while (tl > l[qr[i]]) Add(euler[--tl]);  
26|   while (tr > r[qr[i]]) Add(euler[tr--]);  
27|   while (tr < r[qr[i]]) Add(euler[++tr]);  
28|   // add/remove LCA(u, v) if necessary  
29| }
```

## 2. Data Structures

### 2.1. GNU PBDS

```

1 #include <ext/pb_ds/assoc_container.hpp>
2 #include <ext/pb_ds/priority_queue.hpp>
3 #include <ext/pb_ds/tree_policy.hpp>
4 using namespace __gnu_pbds;
5
6 // most std::map + order_of_key, find_by_order, split, join
7 template <typename T, typename U = null_type>
8 using ordered_map = tree<T, U, std::less<>, rb_tree_tag,
9                         tree_order_statistics_node_update>;
10 // useful tags: rb_tree_tag, splay_tree_tag
11
12 template <typename T> struct myhash {
13     size_t operator()(T x) const; // splitmix, bswap(x*R), ...
14 };
15 // most of std::unordered_map, but faster (needs good hash)
16 template <typename T, typename U = null_type>
17 using hash_table = gp_hash_table<T, U, myhash<T>>;
18
19 // most std::priority_queue + modify, erase, split, join
20 using heap = priority_queue<int, std::less<>>;
21 // useful tags: pairing_heap_tag, binary_heap_tag,
22 //                 (rc_)?binomial_heap_tag, thin_heap_tag

```

### 2.2. 2D Partial Sums

```

1 using vvi = vector<vector<int>>;
2 using vvll = vector<vector<ll>>;
3 using vll = vector<ll>;
4
5 struct PrefixSum2D {
6     vvll pref; // 0-based 2-D prefix sum
7     void build(const vvll &v) { // creates a copy
8         int n = v.size(), m = v[0].size();
9         pref.assign(n, vll(m, 0));
10        for (int i = 0; i < n; i++) {
11            for (int j = 0; j < m; j++) {
12                pref[i][j] = v[i][j] + (i ? pref[i - 1][j] : 0) +
13                           (j ? pref[i][j - 1] : 0) -
14                           (i && j ? pref[i - 1][j - 1] : 0);
15            }
16        }
17    }
18    ll query(int ulx, int uly, int brx, int bry) const {
19        ll ans = pref[brx][bry];
20        if (ulx) ans -= pref[ulx - 1][bry];
21        if (uly) ans -= pref[brx][uly - 1];
22        if (ulx && uly) ans += pref[ulx - 1][uly - 1];
23        return ans;
24    }
25    ll query(int ulx, int uly, int size) const {
26        return query(ulx, uly, ulx + size - 1, uly + size - 1);
27    }
28}; // PartialSum2D : PrefixSum2D
29 struct PartialSum2D : PrefixSum2D {
30     vvll diff; // 0 based
31     int n, m;
32     PartialSum2D(int _n, int _m) : n(_n), m(_m) {
33         diff.assign(n + 1, vll(m + 1, 0));
34     }
35     // add c from [ulx,uly] to [brx,bry]
36     void update(int ulx, int uly, int brx, int bry, ll c) {
37         diff[ulx][uly] += c;
38         diff[ulx][bry + 1] -= c;
39         diff[brx + 1][uly] -= c;
40         diff[brx + 1][bry + 1] += c;
41     }
42     void update(int ulx, int uly, int size, ll c) {
43         int brx = ulx + size - 1;
44         int bry = uly + size - 1;
45         update(ulx, uly, brx, bry, c);
46     }
47     // process the grid using prefix sum
48     void process() { this->build(diff); }
49 };
50 // usage
51 PrefixSum2D pref;
52 pref.build(v); // takes 2d 0-based vector as input
53 pref.query(x1, y1, x2, y2); // sum of region
54
55 PartialSum2D part(n, m); // dimension of grid 0 based
56 part.update(x1, y1, x2, y2, 1); // add 1 in region
57 // must run after all updates
58 part.process(); // prefix sum on diff array
59 // only exists after processing
60 vvll &grid = part.pref; // processed diff array
61 part.query(x1, y1, x2, y2); // gives sum of region

```

### 2.3. Segment Tree (ZKW)

```

1 struct segtree {

```

```

3     using T = int;
4     T f(T a, T b) { return a + b; } // any monoid operation
5     static constexpr T ID = 0; // identity element
6     int n;
7     vector<T> v;
8     segtree(int n_) : n(n_), v(2 * n, ID) {}
9     segtree(vector<T> &a) : n(a.size()), v(2 * n, ID) {
10        copy_n(a.begin(), n, v.begin() + n);
11        for (int i = n - 1; i > 0; i--) {
12            v[i] = f(v[i * 2], v[i * 2 + 1]);
13        }
14    }
15    void update(int i, T x) {
16        for (v[i += n] = x; i /= 2;) {
17            v[i] = f(v[i * 2], v[i * 2 + 1]);
18        }
19    }
20    T query(int l, int r) {
21        T tl = ID, tr = ID;
22        for (l += n, r += n; l < r; l /= 2, r /= 2) {
23            if (l & 1) tl = f(tl, v[l++]);
24            if (r & 1) tr = f(v[--r], tr);
25        }
26        return f(tl, tr);
27    }

```

### 2.4. Line Container

```

1
3     struct Line {
4         mutable ll k, m, p;
5         bool operator<(const Line &o) const { return k < o.k; }
6         bool operator<(ll x) const { return p < x; }
7     };
8     // add: line y=kx+m, query: maximum y of given x
9     struct LineContainer : multiset<Line, less<>> {
10        // (for doubles, use inf = 1/.0, div(a,b) = a/b)
11        static const ll inf = LLONG_MAX;
12        ll div(ll a, ll b) { // floored division
13            return a / b - ((a ^ b) < 0 && a % b);
14        }
15        bool isect(iterator x, iterator y) {
16            if (y == end()) return x->p = inf, 0;
17            if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
18            else x->p = div(y->m - x->m, x->k - y->k);
19            return x->p >= y->p;
20        }
21        void add(ll k, ll m) {
22            auto z = insert({k, m, 0}), y = z++, x = y;
23            while (isect(y, z)) z = erase(z);
24            if (x != begin() && isect(--x, y))
25                isect(x, y = erase(y));
26            while ((y = x) != begin() && (--x)->p >= y->p)
27                isect(x, erase(y));
28        }
29        ll query(ll x) {
30            assert(!empty());
31            auto l = *lower_bound(x);
32            return l.k * x + l.m;
33        }

```

### 2.5. Li-Chao Tree

```

1     constexpr ll MAXN = 2e5, INF = 2e18;
2     struct Line {
3         ll m, b;
4         Line() : m(0), b(-INF) {}
5         Line(ll _m, ll _b) : m(_m), b(_b) {}
6         ll operator()(ll x) const { return m * x + b; }
7     };
8     struct Li_Chao {
9         Line a[MAXN * 4];
10        void insert(Line seg, int l, int r, int v = 1) {
11            if (l == r) {
12                if (seg(l) > a[v](l)) a[v] = seg;
13                return;
14            }
15            int mid = (l + r) >> 1;
16            if (a[v].m > seg.m) swap(a[v], seg);
17            if (a[v](mid) < seg(mid)) {
18                swap(a[v], seg);
19                insert(seg, l, mid, v << 1);
20            } else insert(seg, mid + 1, r, v << 1 | 1);
21        }
22        ll query(int x, int l, int r, int v = 1) {
23            if (l == r) return a[v](x);
24            int mid = (l + r) >> 1;
25            if (x <= mid)
26                return max(a[v](x), query(x, l, mid, v << 1));
27            else
28                return max(a[v](x), query(x, mid + 1, r, v << 1));
29        }

```

## 2.6. Heavy-Light Decomposition

```

1
3 template <bool VALS_EDGES> struct HLD {
4     int N, tim = 0;
5     vector<vi> adj;
6     vi par, siz, depth, rt, pos;
7     Node *tree;
8     HLD(vector<vi> adj_) {
9         : N(sz(adj_)), adj(adj_), par(N, -1), siz(N, 1),
10        depth(N), rt(N), pos(N), tree(new Node(0, N)) {
11            dfsSz(0);
12            dfsHld(0);
13        }
14        void dfsSz(int v) {
15            if (par[v] != -1)
16                adj[v].erase(find(all(adj[v]), par[v]));
17            for (int &u : adj[v]) {
18                par[u] = v, depth[u] = depth[v] + 1;
19                dfsSz(u);
20                siz[v] += siz[u];
21                if (siz[u] > siz[adj[v][0]]) swap(u, adj[v][0]);
22            }
23        }
24        void dfsHld(int v) {
25            pos[v] = tim++;
26            for (int u : adj[v]) {
27                rt[u] = (u == adj[v][0] ? rt[v] : u);
28                dfsHld(u);
29            }
30        }
31        template <class B> void process(int u, int v, B op) {
32            for (; rt[u] != rt[v]; v = par[rt[v]]) {
33                if (depth[rt[u]] > depth[rt[v]]) swap(u, v);
34                op(pos[rt[v]], pos[v] + 1);
35            }
36            if (depth[u] > depth[v]) swap(u, v);
37            op(pos[u] + VALS_EDGES, pos[v] + 1);
38        }
39        void modifyPath(int u, int v, int val) {
40            process(u, v,
41                     [&](int l, int r) { tree->add(l, r, val); });
42        }
43        int queryPath(int u,
44                      int v) { // Modify depending on problem
45            int res = -1e9;
46            process(u, v, [&](int l, int r) {
47                res = max(res, tree->query(l, r));
48            });
49            return res;
50        }
51        int querySubtree(int v) { // modifySubtree is similar
52            return tree->query(pos[v] + VALS_EDGES,
53                                 pos[v] + siz[v]);
54        }
55    };

```

## 2.7. Wavelet Matrix

```

1
3 #pragma GCC target("popcnt,bmi2")
4 #include <immintrin.h>
5
6 // T is unsigned. You might want to compress values first
7 template <typename T> struct wavelet_matrix {
8     static_assert(is_unsigned_v<T>, "only unsigned T");
9     struct bit_vector {
10         static constexpr uint W = 64;
11         uint n, cnt0;
12         vector<ull> bits;
13         vector<uint> sum;
14         bit_vector(uint n_) :
15             n(n_), bits(n / W + 1), sum(n / W + 1) {}
16         void build() {
17             for (uint j = 0; j != n / W; ++j)
18                 sum[j + 1] = sum[j] + _mm_popcnt_u64(bits[j]);
19             cnt0 = rank0(n);
20         }
21         void set_bit(uint i) { bits[i / W] |= 1ULL << i % W; }
22         bool operator[](uint i) const {
23             return !(bits[i / W] & 1ULL << i % W);
24         }
25         uint rank1(uint i) const {
26             return sum[i / W] +
27                   _mm_popcnt_u64(_bzhi_u64(bits[i / W], i % W));
28         }
29         uint rank0(uint i) const { return i - rank1(i); }
30     };
31     uint n, lg;
32     vector<bit_vector> b;
33     wavelet_matrix(const vector<T> &a) : n(a.size()) {

```

```

35     lg =
36         lg(max(*max_element(a.begin(), a.end()), T(1))) + 1;
37     b.assign(lg, n);
38     vector<T> cur = a, nxt(n);
39     for (int h = lg; h--;) {
40         for (uint i = 0; i < n; ++i)
41             if (cur[i] & (T(1) << h)) b[h].set_bit(i);
42         b[h].build();
43         int il = 0, ir = b[h].cnt0;
44         for (uint i = 0; i < n; ++i)
45             nxt[(b[h][i] ? ir : il)++] = cur[i];
46         swap(cur, nxt);
47     }
48     T operator[](uint i) const {
49         T res = 0;
50         for (int h = lg; h--;) {
51             if (b[h][i])
52                 i += b[h].cnt0 - b[h].rank0(i), res |= T(1) << h;
53             else i = b[h].rank0(i);
54         }
55         return res;
56     }
57     // query k-th smallest (0-based) in a[l, r]
58     T kth(uint l, uint r, uint k) const {
59         T res = 0;
60         for (int h = lg; h--;) {
61             uint tl = b[h].rank0(l), tr = b[h].rank0(r);
62             if (k >= tr - tl) {
63                 k -= tr - tl;
64                 l += b[h].cnt0 - tl;
65                 r += b[h].cnt0 - tr;
66                 res |= T(1) << h;
67             } else l = tl, r = tr;
68         }
69         return res;
70     }
71     // count of i in [l, r) with a[i] < u
72     uint count(uint l, uint r, T u) const {
73         if (u >= T(1) << lg) return r - l;
74         uint res = 0;
75         for (int h = lg; h--;) {
76             uint tl = b[h].rank0(l), tr = b[h].rank0(r);
77             if (u & (T(1) << h)) {
78                 l += b[h].cnt0 - tl;
79                 r += b[h].cnt0 - tr;
80                 res += tr - tl;
81             } else l = tl, r = tr;
82         }
83         return res;
84     }
85 }

```

## 2.8. Link-Cut Tree

```

1
3 const int MXN = 100005;
4 const int MEM = 100005;
5
6 struct Splay {
7     static Splay nil, mem[MEM], *pmem;
8     Splay *ch[2], *f;
9     int val, rev, size;
10    Splay() : val(-1), rev(0), size(0) {
11        f = ch[0] = ch[1] = &nil;
12    }
13    Splay(int _val) : val(_val), rev(0), size(1) {
14        f = ch[0] = ch[1] = &nil;
15    }
16    bool isr() {
17        return f->ch[0] != this && f->ch[1] != this;
18    }
19    int dir() { return f->ch[0] == this ? 0 : 1; }
20    void setCh(Splay *c, int d) {
21        ch[d] = c;
22        if (c != &nil) c->f = this;
23        pull();
24    }
25    void push() {
26        if (rev) {
27            swap(ch[0], ch[1]);
28            if (ch[0] != &nil) ch[0]->rev ^= 1;
29            if (ch[1] != &nil) ch[1]->rev ^= 1;
30            rev = 0;
31        }
32    }
33    void pull() {
34        size = ch[0]->size + ch[1]->size + 1;
35        if (ch[0] != &nil) ch[0]->f = this;
36        if (ch[1] != &nil) ch[1]->f = this;
37    }
38    Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::mem;
39    Splay *nil = &Splay::nil;

```

```

41 void rotate(Splay *x) {
42     Splay *p = x->f;
43     int d = x->dir();
44     if (!p->isr()) p->f->setCh(x, p->dir());
45     else x->f = p->f;
46     p->setCh(x->ch[!d], d);
47     x->setCh(p, !d);
48     p->pull();
49     x->pull();
50 }
51
52 vector<Splay *> splayVec;
53 void splay(Splay *x) {
54     splayVec.clear();
55     for (Splay *q = x;; q = q->f) {
56         splayVec.push_back(q);
57         if (q->isr()) break;
58     }
59     reverse(begin(splayVec), end(splayVec));
60     for (auto it : splayVec) it->push();
61     while (!x->isr()) {
62         if (x->f->isr()) rotate(x);
63         else if (x->dir() == x->f->dir())
64             rotate(x->f), rotate(x);
65         else rotate(x), rotate(x);
66     }
67 }
68
68 Splay *access(Splay *x) {
69     Splay *q = nil;
70     for (; x != nil; x = x->f) {
71         splay(x);
72         x->setCh(q, 1);
73         q = x;
74     }
75     return q;
76 }
77 void evert(Splay *x) {
78     access(x);
79     splay(x);
80     x->rev ^= 1;
81     x->push();
82     x->pull();
83 }
84 void link(Splay *x, Splay *y) {
85     // evert(x);
86     access(x);
87     splay(x);
88     evert(y);
89     x->setCh(y, 1);
90 }
91 void cut(Splay *x, Splay *y) {
92     // evert(x);
93     access(y);
94     splay(y);
95     y->push();
96     y->ch[0] = y->ch[0]->f = nil;
97 }
98
99 int N, Q;
100 Splay *vt[MXN];
101
102 int ask(Splay *x, Splay *y) {
103     access(x);
104     access(y);
105     splay(x);
106     int res = x->f->val;
107     if (res == -1) res = x->val;
108     return res;
109 }
110
111 int main(int argc, char **argv) {
112     scanf("%d%d", &N, &Q);
113     for (int i = 1; i <= N; i++)
114         vt[i] = new (Splay::pmem++) Splay(i);
115     while (Q--) {
116         char cmd[105];
117         int u, v;
118         scanf("%s", cmd);
119         if (cmd[1] == 'i') {
120             scanf("%d%d", &u, &v);
121             link(vt[v], vt[u]);
122         } else if (cmd[0] == 'c') {
123             scanf("%d", &v);
124             cut(vt[1], vt[v]);
125         } else {
126             scanf("%d%d", &u, &v);
127             int res = ask(vt[u], vt[v]);
128             printf("%d\n", res);
129         }
130     }
131 }

```

### 3. Graph

#### 3.1. Modeling

- Maximum/Minimum flow with lower bound / Circulation problem

1. Construct super source  $S$  and sink  $T$ .

2. For each edge  $(x, y, l, u)$ , connect  $x \rightarrow y$  with capacity  $u - l$ .

3. For each vertex  $v$ , denote by  $in(v)$  the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.

4. If  $in(v) > 0$ , connect  $S \rightarrow v$  with capacity  $in(v)$ , otherwise, connect  $v \rightarrow T$  with capacity  $-in(v)$ .

– To maximize, connect  $t \rightarrow s$  with capacity  $\infty$  (skip this in circulation problem), and let  $f$  be the maximum flow from  $S$  to  $T$ .

If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from  $s$  to  $t$  is the answer.

– To minimize, let  $f$  be the maximum flow from  $S$  to  $T$ . Connect  $t \rightarrow s$  with capacity  $\infty$  and let the flow from  $S$  to  $T$  be  $f'$ . If  $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise,  $f'$

is the answer.

5. The solution of each edge  $e$  is  $l_e + f_e$ , where  $f_e$  corresponds to the flow of edge  $e$  on the graph.

- Construct minimum vertex cover from maximum matching  $M$  on bipartite graph  $(X, Y)$

1. Redirect every edge:  $y \rightarrow x$  if  $(x, y) \in M$ ,  $x \rightarrow y$  otherwise.

2. DFS from unmatched vertices in  $X$ .

3.  $x \in X$  is chosen iff  $x$  is unvisited.

4.  $y \in Y$  is chosen iff  $y$  is visited.

- Minimum cost cyclic flow

1. Construct super source  $S$  and sink  $T$

2. For each edge  $(x, y, c)$ , connect  $x \rightarrow y$  with  $(cost, cap) = (c, 1)$  if  $c > 0$ , otherwise connect  $y \rightarrow x$  with  $(cost, cap) = (-c, 1)$

3. For each edge with  $c < 0$ , sum these cost as  $K$ , then increase  $d(y)$  by 1, decrease  $d(x)$  by 1

4. For each vertex  $v$  with  $d(v) > 0$ , connect  $S \rightarrow v$  with  $(cost, cap) = (0, d(v))$

5. For each vertex  $v$  with  $d(v) < 0$ , connect  $v \rightarrow T$  with  $(cost, cap) = (0, -d(v))$

6. Flow from  $S$  to  $T$ , the answer is the cost of the flow  $C + K$

- Maximum density induced subgraph

1. Binary search on answer, suppose we're checking answer  $T$

2. Construct a max flow model, let  $K$  be the sum of all weights

3. Connect source  $s \rightarrow v$ ,  $v \in G$  with capacity  $K$

4. For each edge  $(u, v, w)$  in  $G$ , connect  $u \rightarrow v$  and  $v \rightarrow u$  with capacity  $w$

5. For  $v \in G$ , connect it with sink  $v \rightarrow t$  with capacity  $K + 2T - (\sum_{e \in E(v)} w(e)) - 2w(v)$

6.  $T$  is a valid answer if the maximum flow  $f < K|V|$

- Minimum weight edge cover

1. For each  $v \in V$  create a copy  $v'$ , and connect  $u' \rightarrow v'$  with weight  $w(u, v)$ .

2. Connect  $v \rightarrow v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to  $v$ .

3. Find the minimum weight perfect matching on  $G'$ .

- Project selection problem

1. If  $p_v > 0$ , create edge  $(s, v)$  with capacity  $p_v$ ; otherwise, create edge  $(v, t)$  with capacity  $-p_v$ .

2. Create edge  $(u, v)$  with capacity  $w$  with  $w$  being the cost of choosing  $u$  without choosing  $v$ .

3. The mincut is equivalent to the maximum profit of a subset of projects.

- 0/1 quadratic programming

$$\sum_x c_{xx} + \sum_y c_y \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y'})$$

can be minimized by the mincut of the following graph:

1. Create edge  $(x, t)$  with capacity  $c_x$  and create edge  $(s, y)$  with capacity  $c_y$ .

2. Create edge  $(x, y)$  with capacity  $c_{xy}$ .

3. Create edge  $(x, y)$  and edge  $(x', y')$  with capacity  $c_{xyx'y'}$ .

#### 3.2. Matching/Flows

##### 3.2.1. Dinic's Algorithm

```

1 struct Dinic {
2     struct edge {
3         int to, cap, flow, rev;
4     };
5     static constexpr int MAXN = 1000, MAXF = 1e9;
6     vector<edge> v[MAXN];
7     int top[MAXN], deep[MAXN], side[MAXN], s, t;
8     void make_edge(int s, int t, int cap) {
9         v[s].push_back({t, cap, 0, (int)v[t].size()});
10        v[t].push_back({s, 0, 0, (int)v[s].size() - 1});
11    }
12    int dfs(int a, int flow) {
13        if (a == t || !flow) return flow;
14        for (int &i = top[a]; i < v[a].size(); i++) {
15            edge &e = v[a][i];
16            if (deep[a] + 1 == deep[e.to] && e.cap - e.flow) {
17                int x = dfs(e.to, min(e.cap - e.flow, flow));
18                if (x) {
19                    e.flow += x, v[e.to][e.rev].flow -= x;
20                    return x;
21                }
22            }
23        }
24        deep[a] = -1;
25        return 0;
26    }
27    bool bfs() {
28        queue<int> q;
29        fill_n(deep, MAXN, 0);
30        q.push(s), deep[s] = 1;
31        int tmp;
32        while (!q.empty()) {
33            tmp = q.front(), q.pop();
34            for (edge e : v[tmp])
35                if (!deep[e.to] && e.cap != e.flow)
36                    deep[e.to] = deep[tmp] + 1, q.push(e.to);
37        }
38        return deep[t];
39    }
40    int max_flow(int _s, int _t) {
41        s = _s, t = _t;
42        int flow = 0, tflow;
43        while (bfs())
44        fill_n(side, MAXN, 0);
45        while ((tflow = dfs(s, MAXF))) flow += tflow;
46    }
47    void reset() {
48        fill_n(side, MAXN, 0);
49        for (auto &i : v) i.clear();
50    }
51};
```

##### 3.2.2. Minimum Cost Flow

```

1 struct MCF {
2     struct edge {
3         ll to, from, cap, flow, cost, rev;
4     } *fromE[MAXN];
5     vector<edge> v[MAXN];
6     ll n, s, t, flows[MAXN], dis[MAXN], pi[MAXN], flowlim;
7     void make_edge(int s, int t, ll cap, ll cost) {
8         if (!cap) return;
9         v[s].pb(edge{t, s, cap, 0LL, cost, v[t].size()});
10        v[t].pb(edge{s, t, 0LL, 0LL, -cost, v[s].size() - 1});
11    }
12    bitset<MAXN> vis;
13    void dijkstra() {
14        vis.reset();
15        gnu_pbds::priority_queue<pair<ll, int>> q;
16        vector<decltype(q)::point_iterator> its(n);
17        q.push({0LL, s});
18        while (!q.empty()) {
19            int now = q.top().second;
20            q.pop();
21            if (vis[now]) continue;
22            vis[now] = 1;
23            ll ndis = dis[now] + pi[now];
24            for (edge &e : v[now]) {
25                if (e.flow == e.cap || vis[e.to]) continue;
26                if (dis[e.to] > ndis + e.cost - pi[e.to]) {
27                    dis[e.to] = ndis + e.cost - pi[e.to];
28                    flows[e.to] = min(flows[now], e.cap - e.flow);
29                    fromE[e.to] = &e;
30                    if (its[e.to] == q.end())
31                        its[e.to] = q.push({dis[e.to], e.to});
32                    else q.modify(its[e.to], {-dis[e.to], e.to});
33                }
34            }
35        }
36    }
37};
```

```

35     }
36   }
37   bool AP(ll &flow) {
38     fill_n(dis, n, INF);
39     fromE[s] = 0;
40     dis[s] = 0;
41     flows[s] = flowlim - flow;
42     dijkstra();
43     if (dis[t] == INF) return false;
44     flow += flows[t];
45     for (edge *e = fromE[t]; e; e = fromE[e->from]) {
46       e->flow += flows[t];
47       v[e->to][e->rev].flow -= flows[t];
48     }
49     for (int i = 0; i < n; i++)
50       pi[i] = min(pi[i] + dis[i], INF);
51     return true;
52   }
53   pll solve(int _s, int _t, ll flowlim = INF) {
54     s = _s, t = _t, flowlim = _flowlim;
55     pll re;
56     while (re.F != flowlim && AP(re.F));
57     for (int i = 0; i < n; i++)
58       for (edge &e : v[i])
59         if (e.flow != 0) re.S += e.flow * e.cost;
60     re.S /= 2;
61     return re;
62   }
63   void init(int _n) {
64     n = _n;
65     fill_n(pi, n, 0);
66     for (int i = 0; i < n; i++) v[i].clear();
67   }
68   void setpi(int s) {
69     fill_n(pi, n, INF);
70     pi[s] = 0;
71     for (ll it = 0, flag = 1, tdis; flag && it < n; it++) {
72       flag = 0;
73       for (int i = 0; i < n; i++)
74         if (pi[i] != INF)
75           for (edge &e : v[i])
76             if (e.cap && (tdis = pi[i] + e.cost) < pi[e.to])
77               pi[e.to] = tdis, flag = 1;
78     }
79   }
}

```

### 3.2.3. Gomory-Hu Tree

Requires: Dinic's Algorithm

```

1
3 int e[MAXN][MAXN];
4 int p[MAXN];
5 Dinic D; // original graph
6 void gomory_hu() {
7   fill(p, p + n, 0);
8   fill(e[0], e[n], INF);
9   for (int s = 1; s < n; s++) {
10     int t = p[s];
11     Dinic F = D;
12     int tmp = F.max_flow(s, t);
13     for (int i = 1; i < s; i++)
14       e[s][i] = e[i][s] = min(tmp, e[t][i]);
15     for (int i = s + 1; i <= n; i++)
16       if (p[i] == t && F.side[i]) p[i] = s;
17   }
}

```

### 3.2.4. Global Minimum Cut

```

1
3 // weights is an adjacency matrix, undirected
4 pair<int, vi> getMinCut(vector<vi> &weights) {
5   int N = sz(weights);
6   vi used(N), cut, best_cut;
7   int best_weight = -1;
8
9   for (int phase = N - 1; phase >= 0; phase--) {
10     vi w = weights[0], added = used;
11     int prev, k = 0;
12     rep(i, 0, phase) {
13       prev = k;
14       k = -1;
15       rep(j, 1, N) if (!added[j] &&
16                     (k == -1 || w[j] > w[k])) k = j;
17       if (i == phase - 1) {
18         rep(j, 0, N) weights[prev][j] += weights[k][j];
19         rep(j, 0, N) weights[j][prev] = weights[prev][j];
20         used[k] = true;
21         cut.push_back(k);
22         if (best_weight == -1 || w[k] < best_weight) {

```

```

23           best_cut = cut;
24           best_weight = w[k];
25         }
26       } else {
27         rep(j, 0, N) w[j] += weights[k][j];
28         added[k] = true;
29       }
30     }
31   }
32   return {best_weight, best_cut};
33 }

```

### 3.2.5. Bipartite Minimum Cover

Requires: Dinic's Algorithm

```

1
3 // maximum independent set = all vertices not covered
4 // x : [0, n), y : [0, m]
5 struct Bipartite_vertex_cover {
6   Dinic D;
7   int n, m, s, t, x[maxn], y[maxn];
8   void make_edge(int x, int y) { D.make_edge(x, y + n, 1); }
9   int matching() {
10     int re = D.max_flow(s, t);
11     for (int i = 0; i < n; i++)
12       for (Dinic::edge &e : D.v[i])
13         if (e.to != s && e.flow == 1) {
14           x[i] = e.to - n, y[e.to - n] = i;
15           break;
16         }
17     return re;
18   }
19   // init() and matching() before use
20   void solve(vector<int> &vx, vector<int> &vy) {
21     bitset<maxn * 2 + 10> vis;
22     queue<int> q;
23     for (int i = 0; i < n; i++)
24       if (x[i] == -1) q.push(i), vis[i] = 1;
25     while (!q.empty()) {
26       int now = q.front();
27       q.pop();
28       if (now < n) {
29         for (Dinic::edge &e : D.v[now])
30           if (e.to != s && e.to - n != x[now] && !vis[e.to])
31             vis[e.to] = 1, q.push(e.to);
32       } else {
33         if (!vis[y[now - n]])
34           vis[y[now - n]] = 1, q.push(y[now - n]);
35       }
36     }
37     for (int i = 0; i < n; i++)
38       if (!vis[i]) vx.pb(i);
39     for (int i = 0; i < m; i++)
40       if (vis[i + n]) vy.pb(i);
41   }
42   void init(int _n, int _m) {
43     n = _n, m = _m, s = n + m, t = s + 1;
44     for (int i = 0; i < n; i++)
45       x[i] = -1, D.make_edge(s, i, 1);
46     for (int i = 0; i < m; i++)
47       y[i] = -1, D.make_edge(i + n, t, 1);
48   }
49 }

```

### 3.2.6. Edmonds' Algorithm

```

1
3 struct Edmonds {
4   int n, T;
5   vector<vector<int>> g;
6   vector<int> pa, p, used, base;
7   Edmonds(int n)
8     : n(n), T(0), g(n), pa(n, -1), p(n), used(n),
9      base(n) {}
10  void add(int a, int b) {
11    g[a].push_back(b);
12    g[b].push_back(a);
13  }
14  int getBase(int i) {
15    while (i != base[i])
16      base[i] = base[base[i]], i = base[i];
17    return i;
18  }
19  vector<int> toJoin;
20  void mark_path(int v, int x, int b, vector<int> &path) {
21    for (; getBase(v) != b; v = pa[v]) {
22      p[v] = x, x = pa[v];
23      toJoin.push_back(v);
24      toJoin.push_back(x);
25      if (!used[x]) used[x] = ++T, path.push_back(x);
26    }
}

```

```

27 }
28 bool go(int v) {
29     for (int x : g[v]) {
30         int b, bv = getBase(v), bx = getBase(x);
31         if (bv == bx) {
32             continue;
33         } else if (used[x]) {
34             vector<int> path;
35             toJoin.clear();
36             if (used[bx] < used[bv])
37                 mark_path(v, x, b = bx, path);
38             else mark_path(x, v, b = bv, path);
39             for (int z : toJoin) base[getBase(z)] = b;
40             for (int z : path)
41                 if (go(z)) return 1;
42         } else if (p[x] == -1) {
43             p[x] = v;
44             if (pa[x] == -1) {
45                 for (int y; x != -1; x = v)
46                     y = p[x], v = pa[y], pa[x] = y, pa[y] = x;
47                 return 1;
48             }
49             if (!used[pa[x]]) {
50                 used[pa[x]] = ++T;
51                 if (go(pa[x])) return 1;
52             }
53         }
54     }
55     return 0;
56 }
57 void init_dfs() {
58     for (int i = 0; i < n; i++)
59         used[i] = 0, p[i] = -1, base[i] = i;
60 }
61 bool dfs(int root) {
62     used[root] = ++T;
63     return go(root);
64 }
65 void match() {
66     int ans = 0;
67     for (int v = 0; v < n; v++)
68         for (int x : g[v])
69             if (pa[v] == -1 && pa[x] == -1) {
70                 pa[v] = x, pa[x] = v, ans++;
71                 break;
72             }
73     init_dfs();
74     for (int i = 0; i < n; i++)
75         if (pa[i] == -1 && dfs(i)) ans++, init_dfs();
76     cout << ans * 2 << "\n";
77     for (int i = 0; i < n; i++)
78         if (pa[i] > i)
79             cout << i + 1 << " " << pa[i] + 1 << "\n";
80 }
81 }

```

### 3.2.7. Minimum Weight Matching

```

1 struct Graph {
2     static const int MAXN = 105;
3     int n, e[MAXN][MAXN];
4     int match[MAXN], d[MAXN], onstk[MAXN];
5     vector<int> stk;
6     void init(int _n) {
7         n = _n;
8         for (int i = 0; i < n; i++)
9             for (int j = 0; j < n; j++)
10                // change to appropriate infinity
11                // if not complete graph
12                e[i][j] = 0;
13    }
14    void add_edge(int u, int v, int w) {
15        e[u][v] = e[v][u] = w;
16    }
17    bool SPFA(int u) {
18        if (onstk[u]) return true;
19        stk.push_back(u);
20        onstk[u] = 1;
21        for (int v = 0; v < n; v++) {
22            if (u != v && match[u] != v && !onstk[v]) {
23                int m = match[v];
24                if (d[m] > d[u] - e[v][m] + e[u][v]) {
25                    d[m] = d[u] - e[v][m] + e[u][v];
26                    onstk[v] = 1;
27                    stk.push_back(v);
28                    if (SPFA(m)) return true;
29                    stk.pop_back();
30                    onstk[v] = 0;
31                }
32            }
33        }
34        onstk[u] = 0;
35        stk.pop_back();
36        return false;
37    }

```

```
37 }
38 int solve() {
39     for (int i = 0; i < n; i += 2) {
40         match[i] = i + 1;
41         match[i + 1] = i;
42     }
43     while (true) {
44         int found = 0;
45         for (int i = 0; i < n; i++) onstk[i] = d[i] = 0;
46         for (int i = 0; i < n; i++) {
47             stk.clear();
48             if (!onstk[i] && SPFA(i)) {
49                 found = 1;
50                 while (stk.size() >= 2) {
51                     int u = stk.back();
52                     stk.pop_back();
53                     int v = stk.back();
54                     stk.pop_back();
55                     match[u] = v;
56                     match[v] = u;
57                 }
58             }
59             if (!found) break;
60         }
61         int ret = 0;
62         for (int i = 0; i < n; i++) ret += e[i][match[i]];
63         ret /= 2;
64         return ret;
65     }
66 } graph;
```

### 3.2.8. Stable Marriage

```

1 // normal stable marriage problem
/* input:
3
Albert Laura Nancy Marcy
5 Brad Marcy Nancy Laura
Chuck Laura Marcy Nancy
7 Laura Chuck Albert Brad
Marcy Albert Chuck Brad
9 Nancy Brad Albert Chuck
*/
11

13 using namespace std;
const int MAXN = 505;
15
16 int n;
17 int favor[MAXN][MAXN]; // favor[boy_id][rank] = girl_id;
18 int order[MAXN][MAXN]; // order[girl_id][boy_id] = rank;
19 int current[MAXN]; // current[boy_id] = rank;
// boy_id will pursue current[boy_id] girl.
20 int girl_current[MAXN]; // girl[girl_id] = boy_id;

22 void initialize() {
    for (int i = 0; i < n; i++) {
        current[i] = 0;
        girl_current[i] = n;
        order[i][n] = n;
    }
29 }

31 map<string, int> male, female;
32 string bname[MAXN], gname[MAXN];
33 int fit = 0;

35 void stable_marriage() {
36
    queue<int> que;
    for (int i = 0; i < n; i++) que.push(i);
39    while (!que.empty()) {
        int boy_id = que.front();
        que.pop();

        int girl_id = favor[boy_id][current[boy_id]];
        current[boy_id]++;
45
        if (order[girl_id][boy_id] <
            order[girl_id][girl_current[girl_id]]) {
            if (girl_current[girl_id] < n)
                que.push(girl_current[girl_id]);
            girl_current[girl_id] = boy_id;
51        } else {
            que.push(boy_id);
53        }
55    }
56
57 int main() {
58     cin >> n;
59
60     for (int i = 0; i < n; i++) {

```

```

61     string p, t;
63     cin >> p;
64     male[p] = i;
65     bname[i] = p;
66     for (int j = 0; j < n; j++) {
67         cin >> t;
68         if (!female.count(t)) {
69             gname[fit] = t;
70             female[t] = fit++;
71         }
72         favor[i][j] = female[t];
73     }
74
75     for (int i = 0; i < n; i++) {
76         string p, t;
77         cin >> p;
78         for (int j = 0; j < n; j++) {
79             cin >> t;
80             order[female[p]][male[t]] = j;
81         }
82     }
83
84     initialize();
85     stable_marriage();
86
87     for (int i = 0; i < n; i++) {
88         cout << bname[i] << " "
89         << gname[favor[i][current[i] - 1]] << endl;
90     }
91 }

```

### 3.2.9. Kuhn-Munkres algorithm

```

1 // Maximum Weight Perfect Bipartite Matching
2 // Detect non-perfect-matching:
3 // 1. set all edge[i][j] as INF
4 // 2. if solve() >= INF, it is not perfect matching.
5
6     typedef long long ll;
7     struct KM {
8         static const int MAXN = 1050;
9         static const ll INF = 1LL << 60;
10        int n, match[MAXN], vx[MAXN], vy[MAXN];
11        ll edge[MAXN][MAXN], lx[MAXN], ly[MAXN], slack[MAXN];
12        void init(int _n) {
13            n = _n;
14            for (int i = 0; i < n; i++)
15                for (int j = 0; j < n; j++) edge[i][j] = 0;
16        }
17        void add_edge(int x, int y, ll w) { edge[x][y] = w; }
18        bool DFS(int x) {
19            vx[x] = 1;
20            for (int y = 0; y < n; y++) {
21                if (vy[y]) continue;
22                if (lx[x] + ly[y] > edge[x][y]) {
23                    slack[y] =
24                        min(slack[y], lx[x] + ly[y] - edge[x][y]);
25                } else {
26                    vy[y] = 1;
27                    if (match[y] == -1 || DFS(match[y])) {
28                        match[y] = x;
29                        return true;
30                    }
31                }
32            }
33            return false;
34        }
35        ll solve() {
36            fill(match, match + n, -1);
37            fill(lx, lx + n, -INF);
38            fill(ly, ly + n, 0);
39            for (int i = 0; i < n; i++)
40                for (int j = 0; j < n; j++)
41                    lx[i] = max(lx[i], edge[i][j]);
42            for (int i = 0; i < n; i++) {
43                fill(slack, slack + n, INF);
44                while (true) {
45                    fill(vx, vx + n, 0);
46                    fill(vy, vy + n, 0);
47                    if (DFS(i)) break;
48                    ll d = INF;
49                    for (int j = 0; j < n; j++)
50                        if (!vy[j]) d = min(d, slack[j]);
51                    for (int j = 0; j < n; j++) {
52                        if (vx[j]) lx[j] -= d;
53                        if (vy[j]) ly[j] += d;
54                        else slack[j] -= d;
55                    }
56                }
57            }
58            ll res = 0;
59            for (int i = 0; i < n; i++) {

```

```

61         }
62         return res;
63     }
64 }

```

### 3.3. Shortest Path Faster Algorithm

```

1 struct SPFA {
2     static const int maxn = 1010, INF = 1e9;
3     int dis[maxn];
4     bitset<maxn> inq, inneg;
5     queue<int> q, tq;
6     vector<pii> v[maxn];
7     void make_edge(int s, int t, int w) {
8         v[s].emplace_back(t, w);
9     }
10    void dfs(int a) {
11        inneg[a] = 1;
12        for (pii i : v[a])
13            if (!inneg[i.F]) dfs(i.F);
14    }
15    bool solve(int n, int s) { // true if have neg-cycle
16        for (int i = 0; i <= n; i++) dis[i] = INF;
17        dis[s] = 0, q.push(s);
18        for (int i = 0; i < n; i++) {
19            inq.reset();
20            int now;
21            while (!q.empty()) {
22                now = q.front(), q.pop();
23                for (pii i &i : v[now]) {
24                    if (dis[i.F] > dis[now] + i.S) {
25                        dis[i.F] = dis[now] + i.S;
26                        if (!inq[i.F]) tq.push(i.F), inq[i.F] = 1;
27                    }
28                }
29            }
30            q.swap(tq);
31        }
32        bool re = !q.empty();
33        inneg.reset();
34        while (!q.empty()) {
35            if (!inneg[q.front()]) dfs(q.front());
36            q.pop();
37        }
38        return re;
39    }
40    void reset(int n) {
41        for (int i = 0; i <= n; i++) v[i].clear();
42    }
43 }

```

### 3.4. Strongly Connected Components

```

1 struct TarjanScc {
2     int n, step;
3     vector<int> time, low, instk, stk;
4     vector<vector<int>> e, scc;
5     TarjanScc(int n_) : n(n_), step(0), time(n), low(n), instk(n), e(n) {}
6     void add_edge(int u, int v) { e[u].push_back(v); }
7     void dfs(int x) {
8         time[x] = low[x] = ++step;
9         stk.push_back(x);
10        instk[x] = 1;
11        for (int y : e[x]) {
12            if (!time[y]) {
13                dfs(y);
14                low[x] = min(low[x], low[y]);
15            } else if (instk[y]) {
16                low[x] = min(low[x], time[y]);
17            }
18            if (time[x] == low[x]) {
19                scc.emplace_back();
20                for (int y = -1; y != x;) {
21                    y = stk.back();
22                    stk.pop_back();
23                    instk[y] = 0;
24                    scc.back().push_back(y);
25                }
26            }
27        }
28    }
29    void solve() {
30        for (int i = 0; i < n; i++)
31            if (!time[i]) dfs(i);
32        reverse(scc.begin(), scc.end());
33        // scc in topological order
34    }
35 }

```

#### 3.4.1. 2-Satisfiability

Requires: Strongly Connected Components

```

1
3 // 1 based, vertex in SCC = MAXN * 2
// (not i) is i + n
5 struct two_SAT {
    int n, ans[MAXN];
    SCC S;
    void imply(int a, int b) { S.make_edge(a, b); }
    bool solve(int _n) {
        n = _n;
        S.solve(n * 2);
        for (int i = 1; i <= n; i++) {
            if (S.scc[i] == S.scc[i + n]) return false;
            ans[i] = (S.scc[i] < S.scc[i + n]);
        }
        return true;
    }
    void init(int _n) {
        n = _n;
        fill_n(ans, n + 1, 0);
        S.init(n * 2);
    }
} SAT;

```

### 3.5. Biconnected Components

#### 3.5.1. Articulation Points

```

1 void dfs(int x, int p) {
    tin[x] = low[x] = ++t;
    int ch = 0;
    for (auto u : g[x])
        if (u.first != p) {
            if (!ins[u.second])
                st.push(u.second), ins[u.second] = true;
            if (tin[u.first])
                low[x] = min(low[x], tin[u.first]);
            continue;
        }
        ++ch;
    dfs(u.first, x);
    low[x] = min(low[x], low[u.first]);
    if (low[u.first] >= tin[x]) {
        cut[x] = true;
        ++sz;
        while (true) {
            int e = st.top();
            st.pop();
            bcc[e] = sz;
            if (e == u.second) break;
        }
    }
    if (ch == 1 && p == -1) cut[x] = false;
}

```

#### 3.5.2. Bridges

```

1 // if there are multi-edges, then they are not bridges
void dfs(int x, int p) {
    tin[x] = low[x] = ++t;
    st.push(x);
    for (auto u : g[x])
        if (u.first != p) {
            if (!tin[u.first])
                low[x] = min(low[x], tin[u.first]);
            continue;
        }
    dfs(u.first, x);
    low[x] = min(low[x], low[u.first]);
    if (low[u.first] == tin[u.first]) br[u.second] = true;
}
if (tin[x] == low[x]) {
    ++sz;
    while (st.size())
        int u = st.top();
        st.pop();
        bcc[u] = sz;
        if (u == x) break;
}

```

### 3.6. Triconnected Components

```

1
3 // requires a union-find data structure
5 struct ThreeEdgeCC {
    int V, ind;
    vector<int> id, pre, post, low, deg, path;
}

```

```

9     vector<vector<int>> components;
10    UnionFind uf;
11    template <class Graph>
12    void dfs(const Graph &G, int v, int prev) {
13        pre[v] = ++ind;
14        for (int w : G[v])
15            if (w != v) {
16                if (w == prev) {
17                    prev = -1;
18                    continue;
19                }
20                if (pre[w] == -1) {
21                    if (pre[w] < pre[v]) {
22                        deg[v]++;
23                        low[v] = min(low[v], pre[w]);
24                    } else {
25                        deg[v]--;
26                        int &u = path[v];
27                        for (; u != -1 && pre[u] <= pre[w] &&
28                            pre[w] <= post[u];) {
29                            uf.join(v, u);
30                            deg[v] += deg[u];
31                            u = path[u];
32                        }
33                    }
34                }
35            }
36        dfs(G, w, v);
37        if (path[w] == -1 && deg[w] <= 1) {
38            deg[v] += deg[w];
39            low[v] = min(low[v], low[w]);
40            continue;
41        }
42        if (deg[w] == 0) w = path[w];
43        if (low[v] > low[w]) {
44            low[v] = min(low[v], low[w]);
45            swap(w, path[v]);
46        }
47        for (; w != -1; w = path[w]) {
48            uf.join(v, w);
49            deg[v] += deg[w];
50        }
51        post[v] = ind;
52    }
53    template <class Graph>
54    ThreeEdgeCC(const Graph &G)
55        : V(G.size()), ind(-1), id(V, -1), pre(V, -1),
56        post(V), low(V, INT_MAX), deg(V, 0), path(V, -1),
57        uf(V) {
58        for (int v = 0; v < V; v++)
59            if (pre[v] == -1) dfs(G, v, -1);
60        components.reserve(uf.cnt);
61        for (int v = 0; v < V; v++)
62            if (uf.find(v) == v) {
63                id[v] = components.size();
64                components.emplace_back(1, v);
65                components.back().reserve(uf.getSize(v));
66            }
67        for (int v = 0; v < V; v++)
68            if (id[v] == -1)
69                components[id[v]] = id[uf.find(v)].push_back(v);
70    };
71 }

```

### 3.7. Centroid Decomposition

```

1 void get_center(int now) {
2     v[now] = true;
3     vtx.push_back(now);
4     sz[now] = 1;
5     mx[now] = 0;
6     for (int u : G[now])
7         if (!v[u]) {
8             get_center(u);
9             mx[now] = max(mx[now], sz[u]);
10            sz[now] += sz[u];
11        }
12    }
13 void get_dis(int now, int d, int len) {
14    dis[d][now] = cnt;
15    v[now] = true;
16    for (auto u : G[now])
17        if (!v[u.first]) { get_dis(u, d, len + u.second); }
18    }
19 void dfs(int now, int fa, int d) {
20    get_center(now);
21    int c = -1;
22    for (int i : vtx) {
23        if (max(mx[i], (int)vtx.size() - sz[i]) <=
24            (int)vtx.size() / 2)
25            c = i;
26        v[i] = false;
27    }
}

```

```

29  get_dis(c, d, 0);
30  for (int i : vtx) v[i] = false;
31  v[c] = true;
32  vtx.clear();
33  dep[c] = d;
34  p[c] = fa;
35  for (auto u : G[c])
36    if (u.first != fa && !v[u.first]) {
37      dfs(u.first, c, d + 1);
38    }
39 }

```

### 3.8. Minimum Mean Cycle

```

1
3 // d[i][j] == 0 if {i,j} !in E
4 long long d[1003][1003], dp[1003][1003];
5
6 pair<long long, long long> MMWC() {
7  memset(dp, 0x3f, sizeof(dp));
8  for (int i = 1; i <= n; ++i) dp[0][i] = 0;
9  for (int i = 1; i <= n; ++i) {
10    for (int j = 1; j <= n; ++j) {
11      for (int k = 1; k <= n; ++k) {
12        dp[i][k] = min(dp[i - 1][j] + d[j][k], dp[i][k]);
13      }
14    }
15  }
16  long long au = 1ll << 31, ad = 1;
17  for (int i = 1; i <= n; ++i) {
18    if (dp[n][i] == 0x3f3f3f3f3f3f3f3f) continue;
19    long long u = 0, d = 1;
20    for (int j = n - 1; j >= 0; --j) {
21      if ((dp[n][i] - dp[j][i]) * d > u * (n - j)) {
22        u = dp[n][i] - dp[j][i];
23        d = n - j;
24      }
25    }
26    if (u * ad < au * d) au = u, ad = d;
27  }
28  long long g = __gcd(au, ad);
29  return make_pair(au / g, ad / g);
}

```

### 3.9. Directed MST

```

1 template <typename T> struct DMST {
2   T g[maxn][maxn], fw[maxn];
3   int n, fr[maxn];
4   bool vis[maxn], inc[maxn];
5   void clear() {
6     for (int i = 0; i < maxn; ++i) {
7       for (int j = 0; j < maxn; ++j) g[i][j] = inf;
8       vis[i] = inc[i] = false;
9     }
10  }
11  void addedge(int u, int v, T w) {
12    g[u][v] = min(g[u][v], w);
13  }
14  T operator()(int root, int _n) {
15    n = _n;
16    if (!dfs(root) != n) return -1;
17    T ans = 0;
18    while (true) {
19      for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] = i;
20      for (int i = 1; i <= n; ++i)
21        if (!inc[i]) {
22          for (int j = 1; j <= n; ++j) {
23            if (!inc[j] && i != j && g[j][i] < fw[i]) {
24              fw[i] = g[j][i];
25              fr[i] = j;
26            }
27          }
28        }
29        int x = -1;
30        for (int i = 1; i <= n; ++i)
31          if (i != root && !inc[i]) {
32            int j = i, c = 0;
33            while (j != root && fr[j] != i && c <= n)
34              ++c, j = fr[j];
35            if (j == root || c > n) continue;
36            else {
37              x = i;
38              break;
39            }
40        }
41        if (!~x) {
42          for (int i = 1; i <= n; ++i)
43            if (i != root && !inc[i]) ans += fw[i];
44        }
45        int y = x;
}

```

```

47  for (int i = 1; i <= n; ++i) vis[i] = false;
48  do {
49    ans += fw[y];
50    y = fr[y];
51    vis[y] = inc[y] = true;
52  } while (y != x);
53  inc[x] = false;
54  for (int k = 1; k <= n; ++k)
55    if (vis[k]) {
56      for (int j = 1; j <= n; ++j)
57        if (!vis[j]) {
58          if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
59          if (g[j][k] < inf && g[j][k] - fw[k] < g[j][x])
60            g[j][x] = g[j][k] - fw[k];
61        }
62    }
63  }
64  return ans;
}
int dfs(int now) {
  int r = 1;
  vis[now] = true;
  for (int i = 1; i <= n; ++i)
    if (g[now][i] < inf && !vis[i]) r += dfs(i);
  return r;
}

```

### 3.10. Maximum Clique

```

1 // source: KACTL
2
3 typedef vector<bitset<200>> vb;
4 struct Maxclique {
5   double limit = 0.025, pk = 0;
6   struct Vertex {
7     int i, d = 0;
8   };
9   typedef vector<Vertex> vv;
10  vb e;
11  vv V;
12  vector<vi> C;
13  vi qmax, q, S, old;
14  void init(vv &r) {
15    for (auto &v : r) v.d = 0;
16    for (auto &v : r)
17      for (auto j : r) v.d += e[v.i][j.i];
18    sort(all(r), [](auto a, auto b) { return a.d > b.d; });
19    int mxD = r[0].d;
20    rep(i, 0, sz(r)) r[i].d = min(i, mxD) + 1;
21  }
22  void expand(vv &R, int lev = 1) {
23    S[lev] += S[lev - 1] - old[lev];
24    old[lev] = S[lev - 1];
25    while (sz(R)) {
26      if (sz(q) + R.back().d <= sz(qmax)) return;
27      q.push_back(R.back().i);
28      vv T;
29      for (auto v : R)
30        if (e[R.back().i][v.i]) T.push_back({v.i});
31      if (sz(T)) {
32        if (S[lev]++ / ++pk < limit) init(T);
33        int j = 0, mxk = 1,
34        mnk = max(sz(qmax) - sz(q) + 1, 1);
35        C[1].clear(), C[2].clear();
36        for (auto v : T) {
37          int k = 1;
38          auto f = [&](int i) { return e[v.i][i]; };
39          while (any_of(all(C[k]), f)) k++;
40          if (k > mxk) mxk = k, C[mxk + 1].clear();
41          if (k < mnk) T[j++].i = v.i;
42          C[k].push_back(v.i);
43        }
44        if (j > 0) T[j - 1].d = 0;
45        rep(k, mnk, mxk + 1) for (int i : C[k]) T[j++].i = i,
46                                    T[j++].d = k;
47        expand(T, lev + 1);
48      } else if (sz(q) > sz(qmax)) qmax = q;
49      q.pop_back(), R.pop_back();
50    }
51  }
52  vi maxClique() {
53    init(V), expand(V);
54    return qmax;
55  }
56  Maxclique(vb conn)
57    : e(conn), C(sz(e) + 1), S(sz(C)), old(S) {
58      rep(i, 0, sz(e)) V.push_back({i});
59    }
60 }

```

### 3.11. Dominator Tree

```

1 // idom[n] is the unique node that strictly dominates n but
2 // does not strictly dominate any other node that strictly
3 // dominates n. idom[n] = 0 if n is entry or the entry
4 // cannot reach n.
5 struct DominatorTree {
6     static const int MAXN = 200010;
7     int n, s;
8     vector<int> g[MAXN], pred[MAXN];
9     vector<int> cov[MAXN];
10    int dfn[MAXN], nfd[MAXN], ts;
11    int par[MAXN];
12    int sdom[MAXN], idom[MAXN];
13    int mom[MAXN], mn[MAXN];
14
15    inline bool cmp(int u, int v) { return dfn[u] < dfn[v]; }
16
17    int eval(int u) {
18        if (mom[u] == u) return u;
19        int res = eval(mom[u]);
20        if (cmp(sdom[mn[mom[u]]], sdom[mn[u]]))
21            mn[u] = mn[mom[u]];
22        return mom[u] = res;
23    }
24
25    void init(int _n, int _s) {
26        n = _n;
27        s = _s;
28        REP1(i, 1, n) {
29            g[i].clear();
30            pred[i].clear();
31            idom[i] = 0;
32        }
33    }
34    void add_edge(int u, int v) {
35        g[u].push_back(v);
36        pred[v].push_back(u);
37    }
38    void DFS(int u) {
39        ts++;
40        dfn[u] = ts;
41        nfd[ts] = u;
42        for (int v : g[u])
43            if (dfn[v] == 0) {
44                par[v] = u;
45                DFS(v);
46            }
47    }
48    void build() {
49        ts = 0;
50        REP1(i, 1, n) {
51            dfn[i] = nfd[i] = 0;
52            cov[i].clear();
53            mom[i] = mn[i] = sdom[i] = i;
54        }
55        DFS(s);
56        for (int i = ts; i >= 2; i--) {
57            int u = nfd[i];
58            if (u == 0) continue;
59            for (int v : pred[u])
60                if (dfn[v]) {
61                    eval(v);
62                    if (cmp(sdom[mn[v]], sdom[u]))
63                        sdom[u] = sdom[mn[v]];
64                }
65            cov[sdom[u]].push_back(u);
66            mom[u] = par[u];
67            for (int w : cov[par[u]]) {
68                eval(w);
69                if (cmp(sdom[mn[w]], par[u])) idom[w] = mn[w];
70                else idom[w] = par[u];
71            }
72            cov[par[u]].clear();
73        }
74        REP1(i, 2, ts) {
75            int u = nfd[i];
76            if (u == 0) continue;
77            if (idom[u] != sdom[u]) idom[u] = idom[idom[u]];
78        }
79    }
80 } dom;

```

```

11    rep(k, 0, 4) {
12        sort(all(id), [&](int i, int j) {
13            return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;
14        });
15        map<int, int> sweep;
16        for (int i : id) {
17            for (auto it = sweep.lower_bound(-ps[i].y);
18                 it != sweep.end(); sweep.erase(it++)) {
19                int j = it->second;
20                P d = ps[i] - ps[j];
21                if (d.y > d.x) break;
22                edges.push_back({d.y + d.x, i, j});
23            }
24            sweep[-ps[i].y] = i;
25        }
26        for (P &p : ps)
27            if (k & 1) p.x = -p.x;
28            else swap(p.x, p.y);
29    }
30    return edges;
31 }

```

### 3.12. Manhattan Distance MST

```

1
3 // returns [(dist, from, to), ...]
4 // then do normal mst afterwards
5 typedef Point<int> P;
6 vector<array<int, 3>> manhattanMST(vector<P> ps) {
7     vi id(sz(ps));
8     iota(all(id), 0);
9     vector<array<int, 3>> edges;

```

## 4. Math

### 4.1. Number Theory

#### 4.1.1. Mod Struct

A list of safe primes: 26003, 27767, 28319, 28979, 29243, 29759, 30467, 910927547, 919012223, 947326223, 990669467, 1007939579, 1019126699, 929760389146037459, 975500632317046523, 989312547895528379

NTT prime $p$	$p - 1$	primitive root
65537	$1 \lll 16$	3
998244353	$119 \lll 23$	3
2748779069441	$5 \lll 39$	3
1945555039024054273	$27 \lll 56$	5

Requires: Extended GCD

```

1
3 template <typename T> struct M {
4     static T MOD; // change to constexpr if already known
5     T v;
6     M(T x = 0) {
7         v = (-MOD <= x && x < MOD) ? x : x % MOD;
8         if (v < 0) v += MOD;
9     }
10    explicit operator T() const { return v; }
11    bool operator==(const M &b) const { return v == b.v; }
12    bool operator!=(const M &b) const { return v != b.v; }
13    M operator-() { return M(-v); }
14    M operator+(M b) { return M(v + b.v); }
15    M operator-(M b) { return M(v - b.v); }
16    M operator*(M b) { return M((__int128)v * b.v % MOD); }
17    M operator/(M b) { return *this * (b ^ (MOD - 2)); }
18    // change above implementation to this if MOD is not prime
19    M inv() {
20        auto [p, _, g] = extgcd(v, MOD);
21        return assert(g == 1), p;
22    }
23    friend M operator^(M a, ll b) {
24        M ans(1);
25        for (; b; b >>= 1, a *= a)
26            if (b & 1) ans *= a;
27        return ans;
28    }
29    friend M &operator+=(M &a, M b) { return a = a + b; }
30    friend M &operator-=(M &a, M b) { return a = a - b; }
31    friend M &operator*=(M &a, M b) { return a = a * b; }
32    friend M &operator/=(M &a, M b) { return a = a / b; }
33};
34 using Mod = M<int>;
35 template <> int Mod::MOD = 1'000'000'007;
36 int &MOD = Mod::MOD;

```

#### 4.1.2. Miller-Rabin

Requires: Mod Struct

```

1
3 // checks if Mod::MOD is prime
4 bool is_prime() {
5     if (MOD < 2 || MOD % 2 == 0) return MOD == 2;
6     Mod A[] = {2, 7, 61}; // for int values (< 2^31)
7     // ll: 2, 325, 9375, 28178, 450775, 9780504, 1795265022
8     int s = __builtin_ctzll(MOD - 1), i;
9     for (Mod a : A) {
10         Mod x = a ^ (MOD >> s);
11         for (i = 0; i < s && (x + 1).v > 2; i++) x *= x;
12         if (i && x != -1) return 0;
13     }
14     return 1;
15 }

```

#### 4.1.3. Linear Sieve

```

1 constexpr ll MAXN = 1000000;
2 bitset<MAXN> is_prime;
3 vector<ll> primes;
4 ll mpf[MAXN], phi[MAXN], mu[MAXN];
5
6 void sieve() {
7     is_prime.set();
8     is_prime[1] = 0;
9     mu[1] = phi[1] = 1;
10    for (ll i = 2; i < MAXN; i++) {
11        if (is_prime[i]) {
12            mpf[i] = i;
13            primes.push_back(i);
14            phi[i] = i - 1;
15            mu[i] = -1;
16        }
17        for (ll p : primes) {
18            if (p > mpf[i] || i * p >= MAXN) break;
19            is_prime[i * p] = 0;
20            mpf[i * p] = p;
21            mu[i * p] = -mu[i];
22            if (i % p == 0)
23                phi[i * p] = phi[i] * p, mu[i * p] = 0;
24            else phi[i * p] = phi[i] * (p - 1);
25        }
26    }
27 }

```

```

19    is_prime[i * p] = 0;
20    mpf[i * p] = p;
21    mu[i * p] = -mu[i];
22    if (i % p == 0)
23        phi[i * p] = phi[i] * p, mu[i * p] = 0;
24    else phi[i * p] = phi[i] * (p - 1);
25 }
26 }
27 }

```

#### 4.1.4. Get Factors

Requires: Linear Sieve

```

1
3 vector<ll> all_factors(ll n) {
4     vector<ll> fac = {1};
5     while (n > 1) {
6         const ll p = mpf[n];
7         vector<ll> cur = {1};
8         while (n % p == 0) {
9             n /= p;
10            cur.push_back(cur.back() * p);
11        }
12        vector<ll> tmp;
13        for (auto x : fac)
14            for (auto y : cur) tmp.push_back(x * y);
15        tmp.swap(fac);
16    }
17    return fac;
18 }

```

#### 4.1.5. Binary GCD

```

1 // returns the gcd of non-negative a, b
2 ull bin_gcd(ull a, ull b) {
3     if (!a || !b) return a + b;
4     int s = __builtin_ctzll(a | b);
5     a >>= __builtin_ctzll(a);
6     while (b) {
7         if ((b >>= __builtin_ctzll(b)) < a) swap(a, b);
8         b -= a;
9     }
10    return a << s;
11 }

```

#### 4.1.6. Extended GCD

```

1 // returns (p, q, g): p * a + q * b == g == gcd(a, b)
2 // g is not guaranteed to be positive when a < 0 or b < 0
3 tuple<ll, ll, ll> extgcd(ll a, ll b) {
4     ll s = 1, t = 0, u = 0, v = 1;
5     while (b) {
6         ll q = a / b;
7         swap(a -= q * b, b);
8         swap(s -= q * t, t);
9         swap(u -= q * v, v);
10    }
11    return {s, u, a};
12 }

```

#### 4.1.7. Chinese Remainder Theorem

Requires: Extended GCD

```

1 // for 0 <= a < m, 0 <= b < n, returns the smallest x >= 0
2 // such that x % m == a and x % n == b
3 ll crt(ll a, ll m, ll b, ll n) {
4     if (n > m) swap(a, b), swap(m, n);
5     auto [x, y, g] = extgcd(m, n);
6     assert((a - b) % g == 0); // no solution
7     x = ((b - a) / g * x) % (n / g) * m + a;
8     return x < 0 ? x + m / g * n : x;
9 }

```

#### 4.1.8. Baby-Step Giant-Step

Requires: Mod Struct

```

1
3 // returns x such that a ^ x = b where x in [l, r)
4 ll bsgs(Mod a, Mod b, ll l = 0, ll r = MOD - 1) {
5     int m = sqrt(r - l) + 1, i;
6     unordered_map<ll, ll> tb;
7     Mod d = (a ^ l) / b;
8     for (i = 0, d = (a ^ l) / b; i < m; i++, d *= a)
9         if (d == 1) return l + i;
10        else tb[(ll)d] = l + i;
11    Mod c = Mod(1) / (a ^ m);
12    for (i = 0, d = 1; i < m; i++, d *= c)
13        if (auto j = tb.find((ll)d); j != tb.end())
14            return j->second + i * m;
15    return assert(0), -1; // no solution
16 }

```

#### 4.1.9. Pollard's Rho

```

1 ll f(ll x, ll mod) { return (x * x + 1) % mod; }
// n should be composite
3 ll pollard_rho(ll n) {
    if (!(n & 1)) return 2;
    while (1) {
        ll y = 2, x = RNG() % (n - 1) + 1, res = 1;
        for (int sz = 2; res == 1; sz *= 2) {
            for (int i = 0; i < sz && res <= 1; i++) {
                x = f(x, n);
                res = __gcd(abs(x - y), n);
            }
            y = x;
        }
        if (res != 0 && res != n) return res;
    }
}

```

#### 4.1.10. Tonelli-Shanks Algorithm

Requires: Mod Struct

```

1
3 int legendre(Mod a) {
    if (a == 0) return 0;
    return (a ^ ((MOD - 1) / 2)) == 1 ? 1 : -1;
}
7 Mod sqrt(Mod a) {
    assert(legendre(a) != -1); // no solution
    ll p = MOD, s = p - 1;
    if (a == 0) return 0;
    if (p == 2) return 1;
    if (p % 4 == 3) return a ^ ((p + 1) / 4);
    int r, m;
    for (r = 0; !(s & 1); r++) s >>= 1;
    Mod n = 2;
    while (legendre(n) != -1) n += 1;
    Mod x = a ^ ((s + 1) / 2), b = a ^ s, g = n ^ s;
    while (b != 1) {
        Mod t = b;
        for (m = 0; t != 1; m++) t *= t;
        Mod gs = g ^ (1LL << (r - m - 1));
        g = gs * gs, x *= gs, b *= g, r = m;
    }
    return x;
}
// to get sqrt(X) modulo p^k, where p is an odd prime:
// c = x^2 (mod p), c = X^2 (mod p^k), q = p^(k-1)
// X = x^q * c^((p^k-2q+1)/2) (mod p^k)

```

#### 4.1.11. Chinese Sieve

```

1 const ll N = 1000000;
// f, g, h multiplicative, h = f (dirichlet convolution) g
3 ll pre_g(ll n);
ll pre_h(ll n);
// preprocessed prefix sum of f
ll pre_f[N];
// prefix sum of multiplicative function f
ll solve_f(ll n) {
    static unordered_map<ll, ll> m;
    if (n < N) return pre_f[n];
    if (m.count(n)) return m[n];
    ll ans = pre_h(n);
    for (ll l = 2, r; l <= n; l = r + 1) {
        r = n / (n / l);
        ans -= (pre_g(r) - pre_g(l - 1)) * djs_f(n / l);
    }
    return m[n] = ans;
}

```

#### 4.1.12. Rational Number Binary Search

```

1 struct QQ {
    ll p, q;
3     QQ go(QQ b, ll d) { return {p + b.p * d, q + b.q * d}; }
};
5 bool pred(QQ);
// returns smallest p/q in [lo, hi] such that
7 // pred(p/q) is true, and 0 <= p, q <= N
QQ frac_bs(ll N) {
    QQ lo{0, 1}, hi{1, 0};
    if (pred(lo)) return lo;
    assert(pred(hi));
    bool dir = 1, L = 1, H = 1;
    for (; L || H; dir = !dir) {
        ll len = 0, step = 1;
        for (int t = 0; t < 2 && (t ? step /= 2 : step *= 2);)
            if (QQ mid = hi.go(lo, len + step));
                mid.p > N || mid.q > N || dir ^ pred(mid))
                    t++;
            else len += step;
    }
}

```

```

21     swap(lo, hi = hi.go(lo, len));
22     (dir ? L : H) = !len;
23 }
return dir ? hi : lo;
}

```

#### 4.1.13. Farey Sequence

```

1 // returns (e/f), where (a/b, c/d, e/f) are
// three consecutive terms in the order n farey sequence
3 // to start, call next_farey(n, 0, 1, 1, n)
pll next_farey(ll n, ll a, ll b, ll c, ll d) {
5     ll p = (n + b) / d;
    return pll(p * c - a, p * d - b);
7 }

```

#### 4.2. Combinatorics

##### 4.2.1. Matroid Intersection

This template assumes 2 weighted matroids of the same type, and that removing an element is much more expensive than checking if one can be added. **Remember to change the implementation details.**

The ground set is  $0, 1, \dots, n - 1$ , where element  $i$  has weight  $w[i]$ . For the unweighted version, remove weights and change BF/SPFA to BFS.

```

1 constexpr int N = 100;
constexpr int INF = le9;
3
struct Matroid { // represents an independent set
5     Matroid(bitset<N>); // initialize from an independent set
    bool can_add(int); // if adding will break independence
    Matroid Remove(int); // removing from the set
};
9
auto matroid_intersection(int n, const vector<int> &w) {
11    bitset<N> S;
    for (int sz = 1; sz <= n; sz++) {
        Matroid M1(S), M2(S);
15    vector<vector<pii>> e(n + 2);
        for (int j = 0; j < n; j++) {
            if (!S[j]) {
                if (M1.can_add(j)) e[n].emplace_back(j, -w[j]);
                if (M2.can_add(j)) e[j].emplace_back(n + 1, 0);
            }
        }
21    for (int i = 0; i < n; i++) {
            if (S[i]) {
                Matroid T1 = M1.remove(i), T2 = M2.remove(i);
                for (int j = 0; j < n; j++) {
                    if (!S[j]) {
                        if (T1.can_add(j)) e[i].emplace_back(j, -w[j]);
                        if (T2.can_add(j)) e[j].emplace_back(i, w[i]);
                    }
                }
29    }
31    vector<pii> dis(n + 2, {INF, 0});
    vector<int> prev(n + 2, -1);
    dis[0] = {0, 0};
    // change to SPFA for more speed, if necessary
35    bool upd = 1;
    while (upd) {
37        upd = 0;
        for (int u = 0; u < n + 2; u++) {
            for (auto [v, c] : e[u]) {
                pii x(dis[u].first + c, dis[u].second + 1);
                if (x < dis[v]) dis[v] = x, prev[v] = u, upd = 1;
            }
43    }
45    if (dis[n + 1].first < INF)
        for (int x = prev[n + 1]; x != n; x = prev[x])
            S.flip(x);
        else break;
49    // S is the max-weighted independent set with size sz
51    }
53    return S;
}

```

##### 4.2.2. De Bruijn Sequence

```

1 int res[kN], aux[kN], a[kN], sz;
void Rec(int t, int p, int n, int k) {
3     if (t > n) {
        if (n % p == 0)
            for (int i = 1; i <= p; ++i) res[sz++] = aux[i];
    } else {
        aux[t] = aux[t - p];
        Rec(t + 1, p, n, k);
        for (aux[t] = aux[t - p] + 1; aux[t] < k; ++aux[t])
            Rec(t + 1, t, n, k);
    }
}

```

```

11 }
13 int DeBruijn(int k, int n) {
14     // return cyclic string of length  $k^n$  such that every
15     // string of length  $n$  using  $k$  character appears as a
16     // substring.
17     if (k == 1) return res[0] = 0, 1;
18     fill(aux, aux + k * n, 0);
19     return sz = 0, Rec(1, 1, n, k), sz;
}

```

#### 4.2.3. Multinomial

```

1
3 // ways to permute v[i]
4 ll multinomial(vi &v) {
5     ll c = 1, m = v.empty() ? 1 : v[0];
6     for (int i = 1; i < v.size(); i++)
7         for (int j = 0; i < v[i]; j++) c = c * ++m / (j + 1);
8     return c;
9 }

```

#### 4.3. Algebra

##### 4.3.1. Formal Power Series

```

1
3 template <typename mint>
4 struct FormalPowerSeries : vector<mint> {
5     using vector<mint>::vector;
6     using FPS = FormalPowerSeries;
7
8     FPS &operator+=(const FPS &r) {
9         if (r.size() > this->size()) this->resize(r.size());
10        for (int i = 0; i < (int)r.size(); i++)
11            (*this)[i] += r[i];
12        return *this;
13    }
15
16    FPS &operator+=(const mint &r) {
17        if (this->empty()) this->resize(1);
18        (*this)[0] += r;
19        return *this;
20    }
21
22    FPS &operator-=(const FPS &r) {
23        if (r.size() > this->size()) this->resize(r.size());
24        for (int i = 0; i < (int)r.size(); i++)
25            (*this)[i] -= r[i];
26        return *this;
27    }
28
29    FPS &operator-=(const mint &r) {
30        if (this->empty()) this->resize(1);
31        (*this)[0] -= r;
32        return *this;
33    }
34
35    FPS &operator*=(const mint &v) {
36        for (int k = 0; k < (int)this->size(); k++)
37            (*this)[k] *= v;
38        return *this;
39    }
40
41    FPS &operator/=(const FPS &r) {
42        if (this->size() < r.size()) {
43            this->clear();
44            return *this;
45        }
46        int n = this->size() - r.size() + 1;
47        if ((int)r.size() <= 64) {
48            FPS f(*this), g(r);
49            g.shrink();
50            mint coeff = g.back().inverse();
51            for (auto &x : g) x *= coeff;
52            int deg = (int)f.size() - (int)g.size() + 1;
53            int gs = g.size();
54            FPS quo(deg);
55            for (int i = deg - 1; i >= 0; i--) {
56                quo[i] = f[i + gs - 1];
57                for (int j = 0; j < gs; j++)
58                    f[i + j] -= quo[i] * g[j];
59            }
60            *this = quo * coeff;
61            this->resize(n, mint(0));
62            return *this;
63        }
64        return *this = ((*this).rev().pre(n) * r.rev().inv(n))
65            .pre(n)
66            .rev();
}

```

```

67 }
68
69 FPS &operator%=(const FPS &r) {
70     *this -= *this / r * r;
71     shrink();
72     return *this;
73 }
74
75 FPS operator+(const FPS &r) const {
76     return FPS(*this) += r;
77 }
78
79 FPS operator-(const mint &v) const {
80     return FPS(*this) -= v;
81 }
82
83 FPS operator-(const FPS &r) const {
84     return FPS(*this) -= r;
85 }
86
87 FPS operator*(const FPS &r) const {
88     return FPS(*this) *= r;
89 }
90
91 FPS operator*(const mint &v) const {
92     return FPS(*this) *= v;
93 }
94
95 FPS operator/(const FPS &r) const {
96     return FPS(*this) /= r;
97 }
98
99 FPS operator%=(const FPS &r) const {
100    return FPS(*this) %= r;
101 }
102
103 FPS operator-() const {
104     FPS ret(this->size());
105     for (int i = 0; i < (int)this->size(); i++)
106         ret[i] = -(*this)[i];
107     return ret;
108 }
109
110 void shrink() {
111     while (this->size() && this->back() == mint(0))
112         this->pop_back();
113 }
114
115 FPS rev() const {
116     FPS ret(*this);
117     reverse(begin(ret), end(ret));
118     return ret;
119 }
120
121 FPS dot(FPS r) const {
122     FPS ret(min(this->size(), r.size()));
123     for (int i = 0; i < (int)ret.size(); i++)
124         ret[i] = (*this)[i] * r[i];
125     return ret;
126 }
127
128 FPS pre(int sz) const {
129     return FPS(begin(*this),
130               begin(*this) + min((int)this->size(), sz));
131 }
132
133 FPS operator>>(int sz) const {
134     if ((int)this->size() <= sz) return {};
135     FPS ret(*this);
136     ret.erase(ret.begin(), ret.begin() + sz);
137     return ret;
138 }
139
140 FPS operator<<(int sz) const {
141     FPS ret(*this);
142     ret.insert(ret.begin(), sz, mint(0));
143     return ret;
144 }
145
146 FPS diff() const {
147     const int n = (int)this->size();
148     FPS ret(max(0, n - 1));
149     mint one(1), coeff(1);
150     for (int i = 1; i < n; i++) {
151         ret[i - 1] = (*this)[i] * coeff;
152         coeff += one;
153     }
154     return ret;
155 }
156
157 FPS integral() const {
158     const int n = (int)this->size();
159     FPS ret(n + 1);
160     ret[0] = mint(0);
161     if (n > 0) ret[1] = mint(1);
162     auto mod = mint::get_mod();
163     for (int i = 2; i <= n; i++)
164         ret[i] = (-ret[mod % i]) * (mod / i);
165 }

```

```

161     for (int i = 0; i < n; i++) ret[i + 1] *= (*this)[i];
162     return ret;
163 }
164
165 mint eval(mint x) const {
166     mint r = 0, w = 1;
167     for (auto &v : *this) r += w * v, w *= x;
168     return r;
169 }
170
171 FPS log(int deg = -1) const {
172     assert((*this)[0] == mint(1));
173     if (deg == -1) deg = (int)this->size();
174     return (this->diff() * this->inv(deg))
175         .pre(deg - 1)
176         .integral();
177 }
178
179 FPS pow(int64_t k, int deg = -1) const {
180     const int n = (int)this->size();
181     if (deg == -1) deg = n;
182     for (int i = 0; i < n; i++) {
183         if ((*this)[i] != mint(0)) {
184             if (i * k > deg) return FPS(deg, mint(0));
185             mint rev = mint(1) / (*this)[i];
186             FPS ret =
187                 (((*this * rev) >> i).log(deg) * k).exp(deg) *
188                 ((*this)[i].pow(k));
189             ret = (ret << (i * k)).pre(deg);
190             if ((int)ret.size() < deg) ret.resize(deg, mint(0));
191             return ret;
192         }
193     }
194     return FPS(deg, mint(0));
195 }
196
197 static void *ntt_ptr;
198 static void set_fft();
199 FPS &operator*=(const FPS &r);
200 void ntt();
201 void intt();
202 void ntt_doubling();
203 static int ntt_pr();
204 FPS inv(int deg = -1) const;
205 FPS exp(int deg = -1) const;
206 };
207 template <typename mint>
208 void *FormalPowerSeries<mint>::ntt_ptr = nullptr;

```

## 4.4. Theorems

### 4.4.1. Kirchhoff's Theorem

Denote  $L$  be a  $n \times n$  matrix as the Laplacian matrix of graph  $G$ , where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where  $c$  is the number of edge  $(i, j)$  in  $G$ .

- The number of undirected spanning in  $G$  is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at  $r$  in  $G$  is  $|\det(\tilde{L}_{rr})|$ .

### 4.4.2. Tutte's Matrix

Let  $D$  be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if  $i < j$  and  $(i, j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{\text{rank}(D)}{2}$  is the maximum matching on  $G$ .

### 4.4.3. Cayley's Formula

- Given a degree sequence  $d_1, d_2, \dots, d_n$  for each labeled vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

- Let  $T_{n,k}$  be the number of labeled forests on  $n$  vertices with  $k$  components, such that vertex  $1, 2, \dots, k$  belong to different components. Then  $T_{n,k} = kn^{n-k-1}$ .

### 4.4.4. Erdős–Gallai Theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \dots \geq d_n$  can be represented as the degree sequence of a finite simple graph on  $n$  vertices if and only if  $d_1 + d_2 + \dots + d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

holds for all  $1 \leq k \leq n$ .

### 4.4.5. Burnside's Lemma

Let  $X$  be a set and  $G$  be a group that acts on  $X$ . For  $g \in G$ , denote by  $X^g$  the elements fixed by  $g$ :

$$X^g = \{x \in X \mid gx \in X\}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

## 5. Numeric

### 5.1. Barrett Reduction

```

1 using ull = unsigned long long;
2 using uL = __uint128_t;
3 // very fast calculation of a % m
4 struct reduction {
5     const ull m, d;
6     explicit reduction(ull m) : m(m), d(((uL)1 << 64) / m) {}
7     inline ull operator()(ull a) const {
8         ull q = (ull)((uL)d * a) >> 64;
9         return (a - q * m) >= m ? a - m : a;
10    }
11 };

```

### 5.2. Long Long Multiplication

```

1 using ull = unsigned long long;
2 using ll = long long;
3 using ld = long double;
4 // returns a * b % M where a, b < M < 2**63
5 ull mult(ull a, ull b, ull M) {
6     ll ret = a * b - M * ull(ld(a) * ld(b) / ld(M));
7     return ret + M * (ret < 0) - M * (ret >= (ll)M);
}

```

### 5.3. Fast Fourier Transform

```

1 template <typename T>
2 void fft_(int n, vector<T> &a, vector<T> &rt, bool inv) {
3     vector<int> br(n);
4     for (int i = 1; i < n; i++) {
5         br[i] = (i & 1) ? br[i - 1] + n / 2 : br[i / 2] / 2;
6         if (br[i] > i) swap(a[i], a[br[i]]);
7     }
8     for (int len = 2; len <= n; len *= 2)
9         for (int i = 0; i < n; i += len)
10            for (int j = 0; j < len / 2; j++) {
11                int pos = n / len * (inv ? len - j : j);
12                T u = a[i + j], v = a[i + j + len / 2] * rt[pos];
13                a[i + j] = u + v, a[i + j + len / 2] = u - v;
14            }
15     if (T minv = T(1) / T(n); inv)
16         for (T &x : a) x *= minv;
17 }

```

Requires: Mod Struct

```

1
3 void ntt(vector<Mod> &a, bool inv, Mod primitive_root) {
4     int n = a.size();
5     Mod root = primitive_root ^ (MOD - 1) / n;
6     vector<Mod> rt(n + 1, 1);
7     for (int i = 0; i < n; i++) rt[i + 1] = rt[i] * root;
8     fft_(n, a, rt, inv);
9 }
10 void fft(vector<complex<double>> &a, bool inv) {
11     int n = a.size();
12     vector<complex<double>> rt(n + 1);
13     double arg = acos(-1) * 2 / n;
14     for (int i = 0; i <= n; i++)
15         rt[i] = {cos(arg * i), sin(arg * i)};
16     fft_(n, a, rt, inv);
17 }

```

### 5.4. Fast Walsh-Hadamard Transform

Requires: Mod Struct

```

1
3 void fwht(vector<Mod> &a, bool inv) {
4     int n = a.size();
5     for (int d = 1; d < n; d <= 1)
6         for (int m = 0; m < n; m++)
7             if (!(m & d)) {
8                 inv ? a[m] -= a[m | d] : a[m] += a[m | d]; // AND
9                 inv ? a[m | d] -= a[m] : a[m | d] += a[m]; // OR
10                Mod x = a[m], y = a[m | d]; // XOR
11                a[m] = x + y, a[m | d] = x - y; // XOR
12            }
13     if (Mod iv = Mod(1) / n; inv) // XOR
14         for (Mod &i : a) i *= iv; // XOR
15 }

```

### 5.5. Subset Convolution

Requires: Mod Struct

```

1 #pragma GCC target("popcnt")
2 #include <immintrin.h>
3
4 void fwht(int n, vector<vector<Mod>> &a, bool inv) {
5     for (int h = 0; h < n; h++)
6         for (int i = 0; i < (1 << n); i++)
7             if (!(i & (1 << h)))
8                 for (int k = 0; k <= n; k++)
9                     inv ? a[i | (1 << h)][k] -= a[i][k]
10                        : a[i | (1 << h)][k] += a[i][k];
11 }
12 // c[k] = sum(popcnt(i & j) == sz && i | j == k) a[i] * b[j]
13 vector<Mod> subset_convolution(int n, int sz,
14                                 const vector<Mod> &a,
15                                 const vector<Mod> &b_) {
16     int len = n + sz + 1, N = 1 << n;
17     vector<vector<Mod>> a(1 << n, vector<Mod>(len, 0)), b = a;
18     for (int i = 0; i < N; i++)
19         a[i][__mm_popcnt_u64(i)] = a[i];
20     b[i][__mm_popcnt_u64(i)] = b[i];
21     fwht(n, a, 0), fwht(n, b, 0);
22     for (int i = 0; i < N; i++) {
23         vector<Mod> tmp(len);
24         for (int j = 0; j < len; j++)
25             for (int k = 0; k <= j; k++)
26                 tmp[j] += a[i][k] * b[i][j - k];
27         a[i] = tmp;
28     }
29     fwht(n, a, 1);
30     vector<Mod> c(N);
31     for (int i = 0; i < N; i++)
32         c[i] = a[i][__mm_popcnt_u64(i) + sz];
33     return c;
34 }
35 }

```

### 5.6. Linear Recurrences

#### 5.6.1. Berlekamp-Massey Algorithm

```

1 template <typename T>
2 vector<T> berlekamp_massey(const vector<T> &s) {
3     int n = s.size(), l = 0, m = 1;
4     vector<T> r(n), p(n);
5     r[0] = p[0] = 1;
6     T b = 1, d = 0;
7     for (int i = 0; i < n; i++, m++, d = 0) {
8         for (int j = 0; j <= l; j++) d += r[j] * s[i - j];
9         if ((d /= b) == 0) continue; // change if T is float
10        auto t = r;
11        for (int j = m; j < n; j++) r[j] -= d * p[j - m];
12        if (l * 2 <= i) l = i + 1 - l, b *= d, m = 0, p = t;
13    }
14    return r.resize(l + 1), reverse(r.begin(), r.end()), r;
15 }

```

#### 5.6.2. Linear Recurrence Calculation

```

1 template <typename T> struct lin_rec {
2     using poly = vector<T>;
3     poly mul(poly a, poly b, poly m) {
4         int n = m.size();
5         poly r(n);
6         for (int i = n - 1; i >= 0; i--) {
7             r.insert(r.begin(), 0), r.pop_back();
8             T c = r[n - 1] + a[n - 1] * b[i];
9             // c /= m[n - 1]; if m is not monic
10            for (int j = 0; j < n; j++)
11                r[j] += a[j] * b[i] - c * m[j];
12        }
13        return r;
14    }
15    poly pow(poly p, ll k, poly m) {
16        poly r(m.size());
17        r[0] = 1;
18        for (; k; k >= 1, p = mul(p, p, m))
19            if (k & 1) r = mul(r, p, m);
20        return r;
21    }
22    T calc(poly t, poly r, ll k) {
23        int n = r.size();
24        poly p(n);
25        p[1] = 1;
26        poly q = pow(p, k, r);
27        T ans = 0;
28        for (int i = 0; i < n; i++) ans += t[i] * q[i];
29        return ans;
30    }
31 }

```

## 5.7. Matrices

### 5.7.1. Determinant

Requires: Mod Struct

```

1
2 Mod det(vector<vector<Mod>> a) {
3     int n = a.size();
4     Mod ans = 1;
5     for (int i = 0; i < n; i++) {
6         int b = i;
7         for (int j = i + 1; j < n; j++) {
8             if (a[j][i] != 0) {
9                 b = j;
10                break;
11            }
12            if (i != b) swap(a[i], a[b]), ans = -ans;
13            ans *= a[i][i];
14            if (ans == 0) return 0;
15            for (int j = i + 1; j < n; j++) {
16                Mod v = a[j][i] / a[i][i];
17                if (v != 0)
18                    for (int k = i + 1; k < n; k++)
19                        a[j][k] -= v * a[i][k];
20        }
21    }
22    return ans;
23 }
```

```

1 double det(vector<vector<double>> a) {
2     int n = a.size();
3     double ans = 1;
4     for (int i = 0; i < n; i++) {
5         int b = i;
6         for (int j = i + 1; j < n; j++)
7             if (fabs(a[j][i]) > fabs(a[b][i])) b = j;
8         if (i != b) swap(a[i], a[b]), ans = -ans;
9         ans *= a[i][i];
10        if (ans == 0) return 0;
11        for (int j = i + 1; j < n; j++) {
12            double v = a[j][i] / a[i][i];
13            if (v != 0)
14                for (int k = i + 1; k < n; k++)
15                    a[j][k] -= v * a[i][k];
16    }
17    return ans;
18 }
```

### 5.7.2. Inverse

```

1
2 // Returns rank.
3 // Result is stored in A unless singular (rank < n).
4 // For prime powers, repeatedly set
5 // A^{-1} = A^{-1} (2I - A^*A^{-1}) (mod p^k)
6 // where A^{-1} starts as the inverse of A mod p,
7 // and k is doubled in each step.
8
9 int matInv(vector<vector<double>> &A) {
10    int n = sz(A);
11    vi col(n);
12    vector<vector<double>> tmp(n, vector<double>(n));
13    rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
14
15    rep(i, 0, n) {
16        int r = i, c = i;
17        rep(j, i, n)
18            rep(k, i, n) if (fabs(A[j][k]) > fabs(A[r][c])) r = j, c = k;
19        if (fabs(A[r][c]) < 1e-12) return i;
20        A[i].swap(A[r]);
21        tmp[i].swap(tmp[r]);
22        rep(j, 0, n) swap(A[j][i], A[j][c]);
23        swap(tmp[j][i], tmp[j][c]);
24        swap(col[i], col[c]);
25        double v = A[i][i];
26        rep(j, i + 1, n) {
27            double f = A[j][i] / v;
28            A[j][i] = 0;
29            rep(k, i + 1, n) A[j][k] -= f * A[i][k];
30            rep(k, 0, n) tmp[j][k] -= f * tmp[i][k];
31        }
32        rep(j, i + 1, n) A[i][j] /= v;
33        rep(j, 0, n) tmp[i][j] /= v;
34        A[i][i] = 1;
35    }
36    for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
37        double v = A[j][i];
38 }
```

```

43           rep(k, 0, n) tmp[j][k] -= v * tmp[i][k];
44     }
45     rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] = tmp[i][j];
46     return n;
47 }
48
49 int matInv_mod(vector<vector<ll>> &A) {
50     int n = sz(A);
51     vi col(n);
52     vector<vector<ll>> tmp(n, vector<ll>(n));
53     rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
54
55     rep(i, 0, n) {
56         int r = i, c = i;
57         rep(j, i, n) rep(k, i, n) if (A[j][k]) {
58             r = j;
59             c = k;
60             goto found;
61         }
62     }
63     return i;
64 found:
65     A[i].swap(A[r]);
66     tmp[i].swap(tmp[r]);
67     rep(j, 0, n) swap(A[j][i], A[j][c]);
68     swap(tmp[j][i], tmp[j][c]);
69     swap(col[i], col[c]);
70     ll v = modpow(A[i][i], mod - 2);
71     rep(j, i + 1, n) {
72         ll f = A[j][i] * v % mod;
73         A[j][i] = 0;
74         rep(k, i + 1, n) A[j][k] =
75             (A[j][k] - f * A[i][k]) % mod;
76         rep(k, 0, n) tmp[j][k] =
77             (tmp[j][k] - f * tmp[i][k]) % mod;
78     }
79     rep(j, i + 1, n) A[i][j] = A[i][j] * v % mod;
80     rep(j, 0, n) tmp[i][j] = tmp[i][j] * v % mod;
81     A[i][i] = 1;
82
83     for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
84         ll v = A[j][i];
85         rep(k, 0, n) tmp[j][k] =
86             (tmp[j][k] - v * tmp[i][k]) % mod;
87     }
88
89     rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] =
90         tmp[i][j] % mod + (tmp[i][j] < 0 ? mod : 0);
91     return n;
92 }
```

### 5.7.3. Characteristic Polynomial

```

1
2
3 // calculate det(a - xI)
4 template <typename T>
5 vector<T> CharacteristicPolynomial(vector<vector<T>> a) {
6     int N = a.size();
7
8     for (int j = 0; j < N - 2; j++) {
9         for (int i = j + 1; i < N; i++) {
10            if (a[i][j] != 0) {
11                swap(a[j + 1], a[i]);
12                for (int k = 0; k < N; k++)
13                    swap(a[k][j + 1], a[k][i]);
14                break;
15            }
16        }
17        if (a[j + 1][j] != 0) {
18            T inv = T(1) / a[j + 1][j];
19            for (int i = j + 2; i < N; i++) {
20                if (a[i][j] == 0) continue;
21                T coe = inv * a[i][j];
22                for (int l = j; l < N; l++)
23                    a[i][l] -= coe * a[j + 1][l];
24                for (int k = 0; k < N; k++)
25                    a[k][j + 1] += coe * a[k][i];
26            }
27        }
28    }
29
30    vector<vector<T>> p(N + 1);
31    p[0] = {T(1)};
32    for (int i = 1; i <= N; i++) {
33        p[i].resize(i + 1);
34        for (int j = 0; j < i; j++) {
35            p[i][j + 1] -= p[i - 1][j];
36            p[i][j] += p[i - 1][j] * a[i - 1][i - 1];
37        }
38        T x = 1;
39        for (int m = 1; m < i; m++) {
```

```

41     x *= -a[i - m][i - m - 1];
42     Tcoe = x * a[i - m - 1][i - 1];
43     for (int j = 0; j < i - m; j++)
44         p[i][j] += coe * p[i - m - 1][j];
45   }
46   return p[N];
}

```

#### 5.7.4. Solve Linear Equation

```

1

3 typedef vector<double> vd;
5 const double eps = 1e-12;

// solves for x: A * x = b
7 int solveLinear(vector<vd> &A, vd &b, vd &x) {
9     int n = sz(A), m = sz(x), rank = 0, br, bc;
11    if (n) assert(sz(A[0]) == m);
12    vi col(m);
13    iota(all(col), 0);

14    rep(i, 0, n) {
15        double v, bv = 0;
16        rep(r, i, n) rep(c, i, m) if ((v = fabs(A[r][c])) > bv)
17            br = r,
18            bc = c, bv = v;
19        if (bv <= eps) {
20            rep(j, i, n) if (fabs(b[j]) > eps) return -1;
21            break;
22        }
23        swap(A[i], A[br]);
24        swap(b[i], b[br]);
25        swap(col[i], col[bc]);
26        rep(j, 0, n) swap(A[j][i], A[j][bc]);
27        bv = 1 / A[i][i];
28        rep(j, i + 1, n) {
29            double fac = A[j][i] * bv;
30            b[j] -= fac * b[i];
31            rep(k, i + 1, m) A[j][k] -= fac * A[i][k];
32        }
33        rank++;
34    }

35    x.assign(m, 0);
36    for (int i = rank; i--) {
37        b[i] /= A[i][i];
38        x[col[i]] = b[i];
39        rep(j, 0, i) b[j] -= A[j][i] * b[i];
40    }
41    return rank; // (multiple solutions if rank < m)
}

```

#### 5.8. Polynomial Interpolation

```

1

3 // returns a, such that a[0]x^0 + a[1]x^1 + a[2]x^2 + ...
4 // passes through the given points
5 typedef vector<double> vd;
6 vd interpolate(vd x, vd y, int n) {
7     vd res(n), temp(n);
8     rep(k, 0, n - 1) rep(i, k + 1, n) y[i] =
9         (y[i] - y[k]) / (x[i] - x[k]);
10    double last = 0;
11    temp[0] = 1;
12    rep(k, 0, n) rep(i, 0, n) {
13        res[i] += y[k] * temp[i];
14        swap(last, temp[i]);
15        temp[i] -= last * x[k];
16    }
17    return res;
}

```

#### 5.9. Simplex Algorithm

```

1 // Two-phase simplex algorithm for solving linear programs
2 // of the form
3 //
4 //      maximize      c^T x
5 //      subject to    Ax <= b
6 //                      x >= 0
7 //
8 // INPUT: A -- an m x n matrix
9 //        b -- an m-dimensional vector
10 //        c -- an n-dimensional vector
11 //        x -- a vector where the optimal solution will be
12 //              stored
13 //
14 // OUTPUT: value of the optimal solution (infinity if
15 //        unbounded
16 //              above, nan if infeasible)

```

```

17 // // To use this code, create an LPSolver object with A, b,
18 // // and c as arguments. Then, call Solve(x).
19 // // typedef long double ld;
20 // typedef vector<ld> vd;
21 // typedef vector<vd> vvd;
22 // typedef vector<int> vi;
23

24 const ld EPS = 1e-9;

25 struct LPSolver {
26     int m, n;
27     vi B, N;
28     vvd D;
29

30     LPSolver(const vvd &A, const vd &b, const vd &c)
31         : m(b.size()), n(c.size()), N(n + 1), B(m),
32           D(m + 2, vd(n + 2)) {
33         for (int i = 0; i < m; i++) {
34             for (int j = 0; j < n; j++) D[i][j] = A[i][j];
35             for (int i = 0; i < m; i++) {
36                 B[i] = n + i;
37                 D[i][n] = -1;
38                 D[i][n + 1] = b[i];
39             }
40             for (int j = 0; j < n; j++) {
41                 N[j] = j;
42                 D[m][j] = -c[j];
43             }
44             N[n] = -1;
45             D[m + 1][n] = 1;
46         }
47     }

48     void Pivot(int r, int s) {
49         double inv = 1.0 / D[r][s];
50         for (int i = 0; i < m + 2; i++) {
51             if (i != r)
52                 for (int j = 0; j < n + 2; j++) {
53                     if (j != s) D[i][j] -= D[r][j] * D[i][s] * inv;
54                 }
55             if (j != s) D[r][j] *= inv;
56             for (int i = 0; i < m + 2; i++) {
57                 if (i != r) D[i][s] *= -inv;
58                 D[r][s] = inv;
59                 swap(B[r], N[s]);
60             }
61         }
62     }

63     bool Simplex(int phase) {
64         int x = phase == 1 ? m + 1 : m;
65         while (true) {
66             int s = -1;
67             for (int j = 0; j <= n; j++) {
68                 if (phase == 2 && N[j] == -1) continue;
69                 if (s == -1 || D[x][j] < D[x][s] ||
70                     D[x][j] == D[x][s] && N[j] < N[s])
71                     s = j;
72             }
73             if (D[x][s] > -EPS) return true;
74             int r = -1;
75             for (int i = 0; i < m; i++) {
76                 if (D[i][s] < EPS) continue;
77                 if (r == -1 ||
78                     D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
79                     (D[i][n + 1] / D[i][s]) ==
80                     (D[r][n + 1] / D[r][s]) &&
81                     B[i] < B[r])
82                     r = i;
83             }
84             if (r == -1) return false;
85             Pivot(r, s);
86         }
87     }

88     ld Solve(vd &x) {
89         int r = 0;
90         for (int i = 1; i < m; i++) {
91             if (D[i][n + 1] < D[r][n + 1]) r = i;
92         }
93         if (D[r][n + 1] < -EPS) {
94             Pivot(r, n);
95             if (!Simplex(1) || D[m + 1][n + 1] < -EPS)
96                 return numeric_limits<ld>::infinity();
97             for (int i = 0; i < m; i++) {
98                 if (B[i] == -1) {
99                     int s = -1;
100                    for (int j = 0; j <= n; j++) {
101                        if (s == -1 || D[i][j] < D[i][s] ||
102                            D[i][j] == D[i][s] && N[j] < N[s])
103                            s = j;
104                    }
105                    Pivot(i, s);
106                }
107            }
108            if (!Simplex(2)) return numeric_limits<ld>::infinity();
109            x = vd(n);
110        }
111    }

```

```
111     for (int i = 0; i < m; i++)
112         if (B[i] < n) x[B[i]] = D[i][n + 1];
113     return D[m][n + 1];
114 }
115 };
116
117 int main() {
118
119     const int m = 4;
120     const int n = 3;
121     ld _A[m][n] = {
122         {6, -1, 0}, {-1, -5, 0}, {1, 5, 1}, {-1, -5, -1}};
123     ld _b[m] = {10, -4, 5, -5};
124     ld _c[n] = {1, -1, 0};
125
126     vvd A(m);
127     vd b(_b, _b + m);
128     vd c(_c, _c + n);
129     for (int i = 0; i < m; i++) A[i] = vd(_A[i], _A[i] + n);
130
131     LPSolver solver(A, b, c);
132     vd x;
133     ld value = solver.Solve(x);
134
135     cerr << "VALUE: " << value << endl; // VALUE: 1.29032
136     cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
137     for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
138     cerr << endl;
139     return 0;
140 }
```

## 6. Geometry

### 6.1. Point

```

1 template <typename T> struct P {
2     T x, y;
3     P(T x = 0, T y = 0) : x(x), y(y) {}
4     bool operator<(const P &p) const {
5         return tie(x, y) < tie(p.x, p.y);
6     }
7     bool operator==(const P &p) const {
8         return tie(x, y) == tie(p.x, p.y);
9     }
10    P operator-() const { return {-x, -y}; }
11    P operator+(P p) const { return {x + p.x, y + p.y}; }
12    P operator-(P p) const { return {x - p.x, y - p.y}; }
13    P operator*(T d) const { return {x * d, y * d}; }
14    P operator/(T d) const { return {x / d, y / d}; }
15    T dist2() const { return x * x + y * y; }
16    double len() const { return sqrt(dist2()); }
17    P unit() const { return *this / len(); }
18    friend T dot(P a, P b) { return a.x * b.x + a.y * b.y; }
19    friend T cross(P a, P b) { return a.x * b.y - a.y * b.x; }
20    friend T cross(P a, P b, P o) {
21        return cross(a - o, b - o);
22    }
23};
```

using pt = P<ll>;

#### 6.1.1. Quaternion

```

1 constexpr double PI = 3.141592653589793;
2 constexpr double EPS = 1e-7;
3 struct Q {
4     using T = double;
5     T x, y, z, r;
6     Q(T r = 0) : x(0), y(0), z(0), r(r) {}
7     Q(T x, T y, T z, T r = 0) : x(x), y(y), z(z), r(r) {}
8     friend bool operator==(const Q &a, const Q &b) {
9         return (a - b).abs2() <= EPS;
10    }
11    friend bool operator!=(const Q &a, const Q &b) {
12        return !(a == b);
13    }
14    Q operator-() { return Q(-x, -y, -z, -r); }
15    Q operator+(const Q &b) const {
16        return Q(x + b.x, y + b.y, z + b.z, r + b.r);
17    }
18    Q operator-(const Q &b) const {
19        return Q(x - b.x, y - b.y, z - b.z, r - b.r);
20    }
21    Q operator*(const T &t) const {
22        return Q(x * t, y * t, z * t, r * t);
23    }
24    Q operator*(const Q &b) const {
25        return Q(r * b.x + x * b.r + y * b.z - z * b.y,
26                  r * b.y - x * b.z + y * b.r + z * b.x,
27                  r * b.z + x * b.y - y * b.x + z * b.r,
28                  r * b.r - x * b.x - y * b.y - z * b.z);
29    }
30    Q operator/(const Q &b) const { return *this * b.inv(); }
31    T abs2() const { return r * r + x * x + y * y + z * z; }
32    T len() const { return sqrt(abs2()); }
33    Q conj() const { return Q(-x, -y, -z, r); }
34    Q unit() const { return *this * (1.0 / len()); }
35    Q inv() const { return conj() * (1.0 / abs2()); }
36    friend T dot(Q a, Q b) {
37        return a.x * b.x + a.y * b.y + a.z * b.z;
38    }
39    friend Q cross(Q a, Q b) {
40        return Q(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z,
41                  a.x * b.y - a.y * b.x);
42    }
43    friend Q rotation_around(Q axis, T angle) {
44        return axis.unit() * sin(angle / 2) + cos(angle / 2);
45    }
46    Q rotated_around(Q axis, T angle) {
47        Q u = rotation_around(axis, angle);
48        return u * *this / u;
49    }
50    friend Q rotation_between(Q a, Q b) {
51        a = a.unit(), b = b.unit();
52        if (a == -b) {
53            // degenerate case
54            Q ortho = abs(a.y) > EPS ? cross(a, Q(1, 0, 0))
55                                      : cross(a, Q(0, 1, 0));
56            return rotation_around(ortho, PI);
57        }
58        return (a * (a + b)).conj();
59    }
};
```

### 6.1.2. Spherical Coordinates

```

1 struct car_p {
2     double x, y, z;
3 };
4 struct sph_p {
5     double r, theta, phi;
6 };
7
7 sph_p conv(car_p p) {
8     double r = sqrt(p.x * p.x + p.y * p.y + p.z * p.z);
9     double theta = asin(p.y / r);
10    double phi = atan2(p.y, p.x);
11    return {r, theta, phi};
12}
13 car_p conv(sph_p p) {
14    double x = p.r * cos(p.theta) * sin(p.phi);
15    double y = p.r * cos(p.theta) * cos(p.phi);
16    double z = p.r * sin(p.theta);
17    return {x, y, z};
18}
```

### 6.2. Segments

```

1 // for non-collinear ABCD, if segments AB and CD intersect
2 bool intersects(pt a, pt b, pt c, pt d) {
3     if (cross(b, c, a) * cross(b, d, a) > 0) return false;
4     if (cross(d, a, c) * cross(d, b, c) > 0) return false;
5     return true;
6 }
7 // the intersection point of lines AB and CD
8 pt intersect(pt a, pt b, pt c, pt d) {
9     auto x = cross(b, c, a), y = cross(b, d, a);
10    if (x == y) {
11        // if(abs(x, y) < 1e-8) {
12        // is parallel
13    } else {
14        return d * (x / (x - y)) - c * (y / (x - y));
15    }
16 }
```

### 6.3. Convex Hull

```

1 // returns a convex hull in counterclockwise order
2 // for a non-strict one, change cross >= to >
3 vector<pt> convex_hull(vector<pt> p) {
4     sort(ALL(p));
5     if (p[0] == p.back()) return {p[0]};
6     int n = p.size(), t = 0;
7     vector<pt> h(n + 1);
8     for (int _ = 2, s = 0; _--; s = --t, reverse(ALL(p)))
9         for (pt i : p) {
10             while (t > s + 1 && cross(i, h[t - 1], h[t - 2]) >= 0)
11                 t--;
12             h[t++] = i;
13         }
14     return h.resize(t), h;
15 }
```

#### 6.3.1. 3D Hull

```

1
3 typedef Point3D<double> P3;
5 struct PR {
6     void ins(int x) { (a == -1 ? a : b) = x; }
7     void rem(int x) { (a == x ? a : b) = -1; }
8     int cnt() { return (a != -1) + (b != -1); }
9     int a, b;
10 }
11
12 struct F {
13     P3 q;
14     int a, b, c;
15 };
16
17 vector<F> hull3d(const vector<P3> &A) {
18     assert(sz(A) >= 4);
19     vector<vector<PR>> E(sz(A), vector<PR>(sz(A), {-1, -1}));
20     #define E(x, y) E[f.x][f.y]
21     vector<F> FS;
22     auto mf = [&](int i, int j, int k, int l) {
23         P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
24         if (q.dot(A[l]) > q.dot(A[i])) q = q * -1;
25         F f{q, i, j, k};
26         E(a, b).ins(k);
27         E(a, c).ins(j);
28         E(b, c).ins(i);
29         FS.push_back(f);
30     };
31     rep(i, 0, 4) rep(j, i + 1, 4) rep(k, j + 1, 4)
32         mf(i, j, k, 6 - i - j - k);
33 }
```

```

33
35     rep(i, 4, sz(A)) {
36         rep(j, 0, sz(FS)) {
37             F f = FS[j];
38             if (f.q.dot(A[i]) > f.q.dot(A[f.a])) {
39                 E(a, b).rem(f.c);
40                 E(a, c).rem(f.b);
41                 E(b, c).rem(f.a);
42                 swap(FS[j--], FS.back());
43                 FS.pop_back();
44             }
45             int nw = sz(FS);
46             rep(j, 0, nw) {
47                 F f = FS[j];
48 #define C(a, b, c)
49                 if (E(a, b).cnt() != 2) mf(f.a, f.b, i, f.c);
50                 C(a, b, c);
51                 C(a, c, b);
52                 C(b, c, a);
53             }
54             for (F &it : FS)
55                 if ((A[it.b] - A[it.a])
56                     .cross(A[it.c] - A[it.a])
57                     .dot(it.q) <= 0)
58                     swap(it.c, it.b);
59             return FS;
60     };
61 }
```

#### 6.4. Angular Sort

```

1 auto angle_cmp = [](const pt &a, const pt &b) {
2     auto btm = [](const pt &a) {
3         return a.y < 0 || (a.y == 0 && a.x < 0);
4     };
5     return make_tuple(btm(a), a.y * b.x, abs2(a)) <
6         make_tuple(btm(b), a.x * b.y, abs2(b));
7 };
8 void angular_sort(vector<pt> &p) {
9     sort(p.begin(), p.end(), angle_cmp);
10 }
```

#### 6.5. Convex Polygon Minkowski Sum

```

1 // O(n) convex polygon minkowski sum
2 // must be sorted and counterclockwise
3 vector<pt> minkowski_sum(vector<pt> p, vector<pt> q) {
4     auto diff = [](vector<pt> &c) {
5         auto rcmp = [](pt a, pt b) {
6             return pt{a.y, a.x} < pt{b.y, b.x};
7         };
8         rotate(c.begin(), min_element(ALL(c), rcmp), c.end());
9         c.push_back(c[0]);
10        vector<pt> ret;
11        for (int i = 1; i < c.size(); i++)
12            ret.push_back(c[i] - c[i - 1]);
13        return ret;
14    };
15    auto dp = diff(p), dq = diff(q);
16    pt cur = p[0] + q[0];
17    vector<pt> d(dp.size() + dq.size()), ret = {cur};
18 // include angle_cmp from angular-sort.cpp
19    merge(ALL(dp), ALL(dq), d.begin(), angle_cmp);
20 // optional: make ret strictly convex (UB if degenerate)
21    int now = 0;
22    for (int i = 1; i < d.size(); i++) {
23        if (cross(d[i], d[now]) == 0) d[now] = d[now] + d[i];
24        else d[++now] = d[i];
25    }
26    d.resize(now + 1);
27 // end optional part
28    for (pt v : d) ret.push_back(cur = cur + v);
29    return ret.pop_back(), ret;
30 }
```

#### 6.6. Point In Polygon

```

1 bool on_segment(pt a, pt b, pt p) {
2     return cross(a, b, p) == 0 && dot((p - a), (p - b)) <= 0;
3 }
4 // p can be any polygon, but this is O(n)
5 bool inside(const vector<pt> &p, pt a) {
6     int cnt = 0, n = p.size();
7     for (int i = 0; i < n; i++) {
8         pt l = p[i], r = p[(i + 1) % n];
9         // change to return 0; for strict version
10        if (on_segment(l, r, a)) return 1;
11        cnt ^= ((a.y < l.y) - (a.y < r.y)) * cross(l, r, a) > 0;
12    }
13    return cnt;
14 }
```

#### 6.6.1. Convex Version

```

1 // no preprocessing version
2 // p must be a strict convex hull, counterclockwise
3 // if point is inside or on border
4 bool is_inside(const vector<pt> &c, pt p) {
5     int n = c.size(), l = 1, r = n - 1;
6     if (cross(c[0], c[1], p) < 0) return false;
7     if (cross(c[n - 1], c[0], p) < 0) return false;
8     while (l < r - 1) {
9         int m = (l + r) / 2;
10        T a = cross(c[0], c[m], p);
11        if (a > 0) l = m;
12        else if (a < 0) r = m;
13        else return dot(c[0] - p, c[m] - p) <= 0;
14    }
15    if (l == r) return dot(c[0] - p, c[l] - p) <= 0;
16    else return cross(c[l], c[r], p) >= 0;
17 }

19 // with preprocessing version
20 vector<pt> vecs;
21 pt center;
22 // p must be a strict convex hull, counterclockwise
23 // BEWARE OF OVERFLOWS!!
24 void preprocess(vector<pt> p) {
25     for (auto &v : p) v = v * 3;
26     center = p[0] + p[1] + p[2];
27     center.x /= 3, center.y /= 3;
28     for (auto &v : p) v = v - center;
29     vecs = (angular_sort(p), p);
30 }
31 bool intersect_strict(pt a, pt b, pt c, pt d) {
32     if (cross(b, c, a) * cross(b, d, a) > 0) return false;
33     if (cross(d, a, c) * cross(d, b, c) >= 0) return false;
34     return true;
35 }
36 // if point is inside or on border
37 bool query(pt p) {
38     p = p * 3 - center;
39     auto pr = upper_bound(ALL(vecs), p, angle_cmp);
40     if (pr == vecs.end()) pr = vecs.begin();
41     auto pl = (pr == vecs.begin()) ? vecs.back() : *(pr - 1);
42     return !intersect_strict({0, 0}, p, pl, *pr);
43 }
```

#### 6.6.2. Offline Multiple Points Version

Requires: GNU PBDS, Point

```

1
2
3
4
5 using Double = __float128;
6 using Point = pt<Double, Double>;
7
8 int n, m;
9 vector<Point> poly;
10 vector<Point> query;
11 vector<int> ans;
12
13 struct Segment {
14     Point a, b;
15     int id;
16 };
17 vector<Segment> segs;
18
19 Double Xnow;
20 inline Double get_y(const Segment &u, Double xnow = Xnow) {
21     const Point &a = u.a;
22     const Point &b = u.b;
23     return (a.y * (b.x - xnow) + b.y * (xnow - a.x)) /
24         (b.x - a.x);
25 }
26
27 bool operator<(Segment u, Segment v) {
28     Double yu = get_y(u);
29     Double yv = get_y(v);
30     if (yu != yv) return yu < yv;
31     return u.id < v.id;
32 }
33 ordered_map<Segment> st;
34
35 struct Event {
36     int type; // +1 insert seg, -1 remove seg, 0 query
37     Double x, y;
38     int id;
39 };
40
41 bool operator<(Event a, Event b) {
42     if (a.x != b.x) return a.x < b.x;
43     if (a.type != b.type) return a.type < b.type;
44     return a.y < b.y;
45 }
```

```

45 | vector<Event> events;
46 |
47 | void solve() {
48 |     set<Double> xs;
49 |     set<Point> ps;
50 |     for (int i = 0; i < n; i++) {
51 |         xs.insert(poly[i].x);
52 |         ps.insert(poly[i]);
53 |     }
54 |     for (int i = 0; i < n; i++) {
55 |         Segment s{poly[i], poly[(i + 1) % n], i};
56 |         if (s.a.x > s.b.x ||
57 |             (s.a.x == s.b.x && s.a.y > s.b.y)) {
58 |             swap(s.a, s.b);
59 |         }
60 |         segs.push_back(s);
61 |
62 |         if (s.a.x != s.b.x) {
63 |             events.push_back({+1, s.a.x + 0.2, s.a.y, i});
64 |             events.push_back({-1, s.b.x - 0.2, s.b.y, i});
65 |         }
66 |     }
67 |     for (int i = 0; i < m; i++) {
68 |         events.push_back({0, query[i].x, query[i].y, i});
69 |     }
70 |     sort(events.begin(), events.end());
71 |     int cnt = 0;
72 |     for (Event e : events) {
73 |         int i = e.id;
74 |         Xnow = e.x;
75 |         if (e.type == 0) {
76 |             Double x = e.x;
77 |             Double y = e.y;
78 |             Segment tmp = {{x - 1, y}, {x + 1, y}, -1};
79 |             auto it = st.lower_bound(tmp);
80 |
81 |             if (ps.count(query[i]) > 0) {
82 |                 ans[i] = 0;
83 |             } else if (xs.count(x) > 0) {
84 |                 ans[i] = -2;
85 |             } else if (it != st.end() &&
86 |                        get_y(*it) == get_y(tmp)) {
87 |                 ans[i] = 0;
88 |             } else if (it != st.begin() &&
89 |                        get_y(*prev(it)) == get_y(tmp)) {
90 |                 ans[i] = 0;
91 |             } else {
92 |                 int rk = st.order_of_key(tmp);
93 |                 if (rk % 2 == 1) {
94 |                     ans[i] = 1;
95 |                 } else {
96 |                     ans[i] = -1;
97 |                 }
98 |             }
99 |         } else if (e.type == 1) {
100 |             st.insert(segs[i]);
101 |             assert((int)st.size() == ++cnt);
102 |         } else if (e.type == -1) {
103 |             st.erase(segs[i]);
104 |             assert((int)st.size() == --cnt);
105 |         }
106 |     }
107 | }

```

## 6.7. Closest Pair

```

1 | vector<pll> p; // sort by x first!
2 | bool cmpy(const pll &a, const pll &b) const {
3 |     return a.y < b.y;
4 | }
5 | ll sq(ll x) { return x * x; }
6 | // returns (minimum dist)^2 in [l, r)
7 | ll solve(int l, int r) {
8 |     if (r - l <= 1) return 1e18;
9 |     int m = (l + r) / 2;
10 |    ll mid = p[m].x, d = min(solve(l, m), solve(m, r));
11 |    auto pb = p.begin();
12 |    inplace_merge(pb + l, pb + m, pb + r, cmpy);
13 |    vector<pll> s;
14 |    for (int i = l; i < r; i++)
15 |        if (sq(p[i].x - mid) < d) s.push_back(p[i]);
16 |    for (int i = 0; i < s.size(); i++)
17 |        for (int j = i + 1;
18 |              j < s.size() && sq(s[j].y - s[i].y) < d; j++)
19 |            d = min(d, dis(s[i], s[j]));
20 |    return d;
21 |

```

## 6.8. Minimum Enclosing Circle

```

1 |
2 |
3 | typedef Point<double> P;

```

```

5 | double ccRadius(const P &A, const P &B, const P &C) {
6 |     return (B - A).dist() * (C - B).dist() * (A - C).dist() /
7 |             abs((B - A).cross(C - A)) / 2;
8 |
9 | P ccCenter(const P &A, const P &B, const P &C) {
10 |     P b = C - A, c = B - A;
11 |     return A + (b * c.dist2() - c * b.dist2()).perp() /
12 |             b.cross(c) / 2;
13 |
14 | pair<P, doublevector<P> ps) {
15 |     shuffle(all(ps), mt19937(time(0)));
16 |     P o = ps[0];
17 |     double r = 0, EPS = 1 + 1e-8;
18 |     rep(i, 0, sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
19 |         o = ps[i], r = 0;
20 |         rep(j, 0, i) if ((o - ps[j]).dist() > r * EPS) {
21 |             o = (ps[i] + ps[j]) / 2;
22 |             r = (o - ps[i]).dist();
23 |             rep(k, 0, j) if ((o - ps[k]).dist() > r * EPS) {
24 |                 o = ccCenter(ps[i], ps[j], ps[k]);
25 |                 r = (o - ps[i]).dist();
26 |             }
27 |         }
28 |     }
29 |     return {o, r};

```

## 6.9. Delaunay Triangulation

```

1 |
2 |
3 | typedef Point<ll> P;
4 | typedef struct Quad *Q;
5 | typedef _int128_t lll; // (can be ll if coords are < 2e4)
6 | P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
7 |
8 | struct Quad {
9 |     bool mark;
10 |     Q o, rot;
11 |     P p;
12 |     P F() { return r() ->p; }
13 |     Q r() { return rot->rot; }
14 |     Q prev() { return rot->o->rot; }
15 |     Q next() { return r() ->prev(); }
16 |
17 |     bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
18 |         lll p2 = p.dist2(), A = a.dist2() - p2,
19 |             B = b.dist2() - p2, C = c.dist2() - p2;
20 |         return p.cross(a, b) * C + p.cross(b, c) * A +
21 |                 p.cross(c, a) * B >
22 |                         0;
23 |
24 |     Q makeEdge(P orig, P dest) {
25 |         Q q[] = {new Quad{0, 0, 0, orig}, new Quad{0, 0, 0, arb},
26 |                  new Quad{0, 0, 0, dest}, new Quad{0, 0, 0, arb}};
27 |         rep(i, 0, 4) q[i]->o = q[i & 3],
28 |             q[i]->rot = q[(i + 1) & 3];
29 |         return *q;
30 |     }
31 |     void splice(Q a, Q b) {
32 |         swap(a->o->rot->o, b->o->rot->o);
33 |         swap(a->o, b->o);
34 |     }
35 |     Q connect(Q a, Q b) {
36 |         Q q = makeEdge(a->F(), b->p);
37 |         splice(q, a->next());
38 |         splice(q->r(), b);
39 |         return q;
40 |     }
41 |
42 |     pair<Q, Q> rec(const vector<P> &s) {
43 |         if (sz(s) <= 3) {
44 |             Q a = makeEdge(s[0], s[1]),
45 |                 b = makeEdge(s[1], s.back());
46 |             if (sz(s) == 2) return {a, a->r()};
47 |             splice(a->r(), b);
48 |             auto side = s[0].cross(s[1], s[2]);
49 |             Q c = side ? connect(b, a) : 0;
50 |             return {side < 0 ? c->r() : a, side < 0 ? c : b->r()};
51 |         }
52 |
53 |         #define H(e) e->F(), e->p
54 |         #define valid(e) (e->F().cross(H(base)) > 0)
55 |         Q A, B, ra, rb;
56 |         int half = sz(s) / 2;
57 |         tie(ra, A) = rec({all(s) - half});
58 |         tie(B, rb) = rec({sz(s) - half + all(s)});
59 |         while ((B->p).cross(H(A)) < 0 && (A = A->next()) ||
60 |                 (A->p).cross(H(B)) > 0 && (B = B->r()->o)));
61 |         Q base = connect(B->r(), A);
62 |         if (A->p == ra->p) ra = base->r();
63 |         if (B->p == rb->p) rb = base;
64 |
65 |

```

```

#define DEL(e, init, dir)
    Q e = init->dir;
    if (valid(e))
        while (circ(e->dir->F(), H(base), e->F())) {
            Q t = e->dir;
            splice(e, e->prev());
            splice(e->r(), e->r()->prev());
            e = t;
        }
    for (;;) {
        DEL(LC, base->r(), o);
        DEL(RC, base, prev());
        if (!valid(LC) && !valid(RC)) break;
        if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
            base = connect(RC, base->r());
        else base = connect(base->r(), LC->r());
    }
    return {ra, rb};
}

// returns [A_0, B_0, C_0, A_1, B_1, ...]
// where A_i, B_i, C_i are counter-clockwise triangles
vector<P> triangulate(vector<P> pts) {
    sort(all(pts));
    assert(unique(all(pts)) == pts.end());
    if (sz(pts) < 2) return {};
    Q e = rec(pts).first;
    vector<Q> q = {e};
    int qi = 0;
    while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
#define ADD
{
    Q c = e;
    do {
        c->mark = 1;
        pts.push_back(c->p);
        q.push_back(c->r());
        c = c->next();
    } while (c != e);
}
ADD;
pts.clear();
while (qi < sz(q))
    if (!(e = q[qi++])->mark) ADD;
return pts;
}

```

### 6.9.1. Slower Version

```

1
3 template <class P, class F>
void delaunay(vector<P> &ps, F trifun) {
5     if (sz(ps) == 3) {
        int d = (ps[0].cross(ps[1], ps[2]) < 0);
        trifun(0, 1 + d, 2 - d);
    }
9     vector<P3> p3;
11    for (P p : ps) p3.emplace_back(p.x, p.y, p.dist2());
13    if (sz(ps) > 3)
        for (auto t : hull3d(p3))
15        if ((p3[t.b] - p3[t.a])
            .cross(p3[t.c] - p3[t.a])
            .dot(P3(0, 0, 1)) < 0)
            trifun(t.a, t.c, t.b);
17 }

```

### 6.10. Half Plane Intersection

```

1 struct Line {
2     Point P;
3     Vector v;
4     bool operator<(const Line &b) const {
5         return atan2(v.y, v.x) < atan2(b.v.y, b.v.x);
6     }
7 };
8 bool OnLeft(const Line &L, const Point &p) {
9     return Cross(L.v, p - L.P) > 0;
}
10 Point GetIntersection(Line a, Line b) {
11     Vector u = a.P - b.P;
12     Double t = Cross(b.v, u) / Cross(a.v, b.v);
13     return a.P + a.v * t;
}
14 int HalfplaneIntersection(Line *L, int n, Point *poly) {
15     sort(L, L + n);
16
17     int first, last;
18     Point *p = new Point[n];
19     Line *q = new Line[n];
20     q[first = last = 0] = L[0];
21     for (int i = 1; i < n; i++) {
22         while (first < last && !OnLeft(L[i], p[last - 1]))
23             last--;
24         while (first < last && !OnLeft(L[i], p[first])) first++;
25         q[++last] = L[i];
26         if (fabs(Cross(q[last].v, q[last - 1].v)) < EPS) {
27             last--;
28             if (OnLeft(q[last], L[i].P)) q[last] = L[i];
29         }
30         if (first < last)
31             p[last - 1] = GetIntersection(q[last - 1], q[last]);
32     }
33     while (first < last && !OnLeft(q[first], p[last - 1]))
34         last--;
35     if (last - first <= 1) return 0;
36     p[last] = GetIntersection(q[last], q[first]);
37
38     int m = 0;
39     for (int i = first; i <= last; i++) poly[m++] = p[i];
40     return m;
41 }
42
43 }

```

## 7. Strings

### 7.1. Knuth-Morris-Pratt Algorithm

```

1 vector<int> pi(const string &s) {
2     vector<int> p(s.size());
3     for (int i = 1; i < s.size(); i++) {
4         int g = p[i - 1];
5         while (g && s[i] != s[g]) g = p[g - 1];
6         p[i] = g + (s[i] == s[g]);
7     }
8     return p;
9 }
10 vector<int> match(const string &s, const string &pat) {
11     vector<int> p = pi(pat + '\0' + s), res;
12     for (int i = p.size() - s.size(); i < p.size(); i++)
13         if (p[i] == pat.size())
14             res.push_back(i - 2 * pat.size());
15     return res;
}

```

### 7.2. Z Value

```

1 int z[n];
2 void zval(string s) {
3     // z[i] => longest common prefix of s and s[i:], i > 0
4     int n = s.size();
5     z[0] = 0;
6     for (int b = 0, i = 1; i < n; i++) {
7         if (z[b] + b <= i) z[i] = 0;
8         else z[i] = min(z[i - b], z[b] + b - i);
9         while (s[i + z[i]] == s[z[i]]) z[i]++;
10        if (i + z[i] > b + z[b]) b = i;
11    }
}

```

### 7.3. Manacher's Algorithm

```

1 int z[n];
2 void manacher(string s) {
3     // z[i] => longest odd palindrome centered at i is
4     //      s[i - z[i] ... i + z[i]]
5     // to get all palindromes (including even length),
6     // insert a '#' between each s[i] and s[i + 1]
7     int n = s.size();
8     z[0] = 0;
9     for (int b = 0, i = 1; i < n; i++) {
10        if (z[b] + b >= i)
11            z[i] = min(z[2 * b - i], b + z[b] - i);
12        else z[i] = 0;
13        while (i + z[i] + 1 < n && i - z[i] - 1 >= 0 &&
14              s[i + z[i] + 1] == s[i - z[i] - 1])
15            z[i]++;
16        if (z[i] + i > z[b] + b) b = i;
17    }
}

```

### 7.4. Minimum Rotation

```

1 int min_rotation(string s) {
2     int a = 0, n = s.size();
3     s += s;
4     for (int b = 0; b < n; b++) {
5         for (int k = 0; k < n; k++) {
6             if (a + k == b || s[a + k] < s[b + k]) {
7                 b += max(0, k - 1);
8                 break;
9             }
10            if (s[a + k] > s[b + k]) {
11                a = b;
12                break;
13            }
14        }
15    }
16    return a;
17 }

```

### 7.5. Aho-Corasick Automaton

```

1 struct Aho_Corasick {
2     static const int maxc = 26, maxn = 4e5;
3     struct NODES {
4         int Next[maxc], fail, ans;
5     };
6     NODES T[maxn];
7     int top, qtop, q[maxn];
8     int get_node(const int &fail) {
9         fill_n(T[top].Next, maxc, 0);
10        T[top].fail = fail;
11        T[top].ans = 0;
12        return top++;
13    }
14    int insert(const string &s) {
15        int ptr = 1;
16        for (char c : s) { // change char id
17            c -= 'a';
18            if (!T[ptr].Next[c]) T[ptr].Next[c] = get_node(ptr);
19            ptr = T[ptr].Next[c];
20        }
21        return ptr;
22    } // return ans_last_place
23    void build_fail(int ptr) {
24        int tmp;
25        for (int i = 0; i < maxc; i++) {
26            if (T[ptr].Next[i]) {
27                tmp = T[ptr].fail;
28                while (tmp != 1 && !T[tmp].Next[i])
29                    tmp = T[tmp].fail;
30                if (T[tmp].Next[i] != T[ptr].Next[i])
31                    if (T[tmp].Next[i]) tmp = T[tmp].Next[i];
32                T[T[ptr].Next[i]].fail = tmp;
33                q[qtop++] = T[ptr].Next[i];
34            }
35        }
36        void AC_auto(const string &s) {
37            int ptr = 1;
38            for (char c : s) {
39                while (ptr != 1 && !T[ptr].Next[c]) ptr = T[ptr].fail;
40                if (T[ptr].Next[c]) {
41                    ptr = T[ptr].Next[c];
42                    T[ptr].ans++;
43                }
44            }
45        }
46        void Solve(string &s) {
47            for (char &c : s) // change char id
48                c -= 'a';
49            for (int i = 0; i < qtop; i++) build_fail(q[i]);
50            AC_auto(s);
51            for (int i = qtop - 1; i > -1; i--)
52                T[T[q[i]].fail].ans += T[q[i]].ans;
53        }
54        void reset() {
55            qtop = top = q[0] = 1;
56            get_node(1);
57        }
58    } AC;
59    // usage example
60    string s, S;
61    int n, t, ans_place[50000];
62    int main() {
63        Tie cin >> t;
64        while (t--) {
65            AC.reset();
66            cin >> S >> n;
67            for (int i = 0; i < n; i++) {
68                cin >> s;
69                ans_place[i] = AC.insert(s);
70            }
71            AC.Solve(S);
72            for (int i = 0; i < n; i++)
73                cout << AC.T[ans_place[i]].ans << '\n';
74        }
75    }
}

```

## 8. Debug List

- 1 - Pre-submit:
  - Did you make a typo when copying a template?
  - Test more cases if unsure.
    - Write a naive solution and check small cases.
  - Submit the correct file.
- 7 - General Debugging:
  - Read the whole problem again.
  - Have a teammate read the problem.
  - Have a teammate read your code.
    - Explain your solution to them (or a rubber duck).
  - Print the code and its output / debug output.
  - Go to the toilet.
- 15 - Wrong Answer:
  - Any possible overflows?
    - > `\_\_int128` ?
    - Try `-ftrapv` or `#pragma GCC optimize("trapv")`
  - Floating point errors?
    - > `long double` ?
    - turn off math optimizations
    - check for `==`, `>=`, `acos(1.0000000001)`, etc.
  - Did you forget to sort or unique?
  - Generate large and worst "corner" cases.
  - Check your `m` / `n`, `i` / `j` and `x` / `y`.
  - Are everything initialized or reset properly?
  - Are you sure about the STL thing you are using?
    - Read cppreference (should be available).
  - Print everything and run it on pen and paper.
- 31 - Time Limit Exceeded:
  - Calculate your time complexity again.
  - Does the program actually end?
    - Check for `while(q.size())` etc.
  - Test the largest cases locally.
  - Did you do unnecessary stuff?
    - e.g. pass vectors by value
    - e.g. `memset` for every test case
  - Is your constant factor reasonable?
- 41 - Runtime Error:
  - Check memory usage.
    - Forget to clear or destroy stuff?
      - > `vector::shrink\_to\_fit()`
  - Stack overflow?
  - Bad pointer / array access?
    - Try `-fsanitize=address`
  - Division by zero? NaN's?