

## Contents

<b>1 Misc</b>	<b>2</b>		
1.1 Contest	2	4.2.1 Matroid Intersection	16
1.1.1 Makefile	2	4.2.2 De Bruijn Sequence	16
1.2 How Did We Get Here?	2	4.2.3 Multinomial	17
1.2.1 Macros	2	4.3 Algebra	17
1.2.2 Fast I/O	2	4.3.1 Formal Power Series	17
1.2.3 constexpr	2	4.4 Theorems	18
1.2.4 Bump Allocator	2	4.4.1 Kirchhoff's Theorem	18
1.3 Tools	3	4.4.2 Tutte's Matrix	18
1.3.1 Floating Point Binary Search	3	4.4.3 Cayley's Formula	18
1.3.2 SplitMix64	3	4.4.4 Erdős–Gallai Theorem	18
1.3.3 <random>	3	4.4.5 Burnside's Lemma	18
1.3.4 x86 Stack Hack	3	<b>5 Numeric</b>	<b>19</b>
1.4 Algorithms	3	5.1 Barrett Reduction	19
1.4.1 Bit Hacks	3	5.2 Long Long Multiplication	19
1.4.2 Aliens Trick	3	5.3 Fast Fourier Transform	19
1.4.3 Hilbert Curve	3	5.4 Fast Walsh-Hadamard Transform	19
1.4.4 Infinite Grid Knight Distance	3	5.5 Subset Convolution	19
1.4.5 Poker Hand	3	5.6 Linear Recurrences	19
1.4.6 Longest Increasing Subsequence	3	5.6.1 Berlekamp-Massey Algorithm	19
1.4.7 Mo's Algorithm on Tree	4	5.6.2 Linear Recurrence Calculation	19
<b>2 Data Structures</b>	<b>5</b>	5.7 Matrices	20
2.1 GNU PBDS	5	5.7.1 Determinant	20
2.2 Segment Tree (ZKW)	5	5.7.2 Inverse	20
2.3 Line Container	5	5.7.3 Characteristic Polynomial	20
2.4 Li-Chao Tree	5	5.7.4 Solve Linear Equation	21
2.5 Heavy-Light Decomposition	5	5.8 Polynomial Interpolation	21
2.6 Wavelet Matrix	6	5.9 Simplex Algorithm	21
2.7 Link-Cut Tree	6	<b>6 Geometry</b>	<b>23</b>
<b>3 Graph</b>	<b>8</b>	6.1 Point	23
3.1 Modeling	8	6.1.1 Quarternion	23
3.2 Matching/Flows	8	6.1.2 Spherical Coordinates	23
3.2.1 Dinic's Algorithm	8	6.2 Segments	23
3.2.2 Minimum Cost Flow	8	6.3 Convex Hull	23
3.2.3 Gomory-Hu Tree	9	6.3.1 3D Hull	23
3.2.4 Global Minimum Cut	9	6.4 Angular Sort	24
3.2.5 Bipartite Minimum Cover	9	6.5 Convex Polygon Minkowski Sum	24
3.2.6 Edmonds' Algorithm	9	6.6 Point In Polygon	24
3.2.7 Minimum Weight Matching	10	6.6.1 Convex Version	24
3.2.8 Stable Marriage	10	6.6.2 Offline Multiple Points Version	24
3.2.9 Kuhn-Munkres algorithm	11	6.7 Closest Pair	25
3.3 Shortest Path Faster Algorithm	11	6.8 Minimum Enclosing Circle	25
3.4 Strongly Connected Components	11	6.9 Delaunay Triangulation	25
3.4.1 2-Satisfiability	12	6.9.1 Slower Version	26
3.5 Biconnected Components	12	6.10 Half Plane Intersection	26
3.5.1 Articulation Points	12	<b>7 Strings</b>	<b>27</b>
3.5.2 Bridges	12	7.1 Knuth-Morris-Pratt Algorithm	27
3.6 Triconnected Components	12	7.2 Aho-Corasick Automaton	27
3.7 Centroid Decomposition	12	7.3 Suffix Array	27
3.8 Minimum Mean Cycle	13	7.4 Suffix Tree	27
3.9 Directed MST	13	7.5 Cocke-Younger-Kasami Algorithm	28
3.10 Maximum Clique	13	7.6 Z Value	28
3.11 Dominator Tree	14	7.7 Manacher's Algorithm	28
3.12 Manhattan Distance MST	14	7.8 Minimum Rotation	28
<b>4 Math</b>	<b>15</b>	7.9 Palindromic Tree	29
4.1 Number Theory	15	<b>8 Debug List</b>	<b>30</b>
4.1.1 Mod Struct	15		
4.1.2 Miller-Rabin	15		
4.1.3 Linear Sieve	15		
4.1.4 Get Factors	15		
4.1.5 Binary GCD	15		
4.1.6 Extended GCD	15		
4.1.7 Chinese Remainder Theorem	15		
4.1.8 Baby-Step Giant-Step	15		
4.1.9 Pollard's Rho	16		
4.1.10 Tonelli-Shanks Algorithm	16		
4.1.11 Chinese Sieve	16		
4.1.12 Rational Number Binary Search	16		
4.1.13 Farey Sequence	16		
4.2 Combinatorics	16		

## 1. Misc

### 1.1. Contest

#### 1.1.1. Makefile

```
1 .PRECIOUS: ./p%
3 %: p%
4     ulimit -s unlimited && ./s<
5 p%: p%.cpp
6     g++ -o $@ $< -std=c++17 -Wall -Wextra -Wshadow \
7         -fsanitize=address,undefined
```

### 1.2. How Did We Get Here?

#### 1.2.1. Macros

Use vectorizations and math optimizations at your own peril.  
For gcc≥9, there are `[[likely]]` and `[[unlikely]]` attributes.  
Call gcc with `-fopt-info-optimized-missed-optall` for optimization info.

```
1 #define _GLIBCXX_DEBUG 1 // for debug mode
2 #define _GLIBCXX_SANITIZE_VECTOR 1 // for asan on vectors
3 #pragma GCC optimize("O3", "unroll-loops")
4 #pragma GCC optimize("fast-math")
5 #pragma GCC target("avx,avx2,abm,bmi,bmi2") // tip: `lscpu`
6 // before a loop
7 #pragma GCC unroll 16 // 0 or 1 -> no unrolling
8 #pragma GCC ivdep
```

#### 1.2.2. Fast I/O

```
1 struct scanner {
2     static constexpr size_t LEN = 32 << 20;
3     char *buf, *buf_ptr, *buf_end;
4     scanner()
5         : buf(new char[LEN]), buf_ptr(buf + LEN),
6           buf_end(buf + LEN) {}
7     ~scanner() { delete[] buf; }
8     char getc() {
9         if (buf_ptr == buf_end) [[unlikely]]
10            buf_end = buf + fread_unlocked(buf, 1, LEN, stdin);
11            buf_ptr = buf;
12            return *(buf_ptr++);
13        }
14        char seek(char del) {
15            char c;
16            while ((c = getc()) < del) {}
17            return c;
18        }
19        void read(int &t) {
20            bool neg = false;
21            char c = seek('-');
22            if (c == '-') neg = true, t = 0;
23            else t = c ^ '0';
24            while ((c = getc()) >= '0') t = t * 10 + (c ^ '0');
25            if (neg) t = -t;
26        }
27    };
28    struct printer {
29        static constexpr size_t CPI = 21, LEN = 32 << 20;
30        char *buf, *buf_ptr, *buf_end, *tbuf;
31        char *int_buf, *int_buf_end;
32        printer()
33            : buf(new char[LEN]), buf_ptr(buf),
34              buf_end(buf + LEN), int_buf(new char[CPI + 1]),
35              int_buf_end(int_buf + CPI - 1) {}
36        ~printer() {
37            flush();
38            delete[] buf, delete[] int_buf;
39        }
40        void flush() {
41            fwrite_unlocked(buf, 1, buf_ptr - buf, stdout);
42            buf_ptr = buf;
43        }
44        void write(const char &c) {
45            *buf_ptr = c;
46            if (++buf_ptr == buf_end) [[unlikely]]
47                flush();
48        }
49        void write(const char *s) {
50            for (; *s != '\0'; ++s) write(*s);
51        }
52        void write(int x) {
53            if (x < 0) write('-'), x = -x;
54            if (x == 0) [[unlikely]]
55                return write('0');
56            return write('0');
57            for (tbuf = int_buf_end; x != 0; --tbuf, x /= 10)
58                *tbuf = '0' + char(x % 10);
```

```
59     write(++tbuf);
60 }
61 };
```

## Kotlin

```
1 import java.io.*
2 import java.util.*
3
4 @JvmField val cin = System.`in`.bufferedReader()
5 @JvmField val cout = PrintWriter(System.out, false)
6 @JvmField var tokenizer: StringTokenizer = StringTokenizer("")
7 fun nextLine() = cin.readLine()!!
8 fun read(): String {
9     while(!tokenizer.hasMoreTokens())
10        tokenizer = StringTokenizer(nextLine())
11        return tokenizer.nextToken()
12    }
13
14 // example
15 fun main() {
16     val n = read().toInt()
17     val a = DoubleArray(n) { read().toDouble() }
18     cout.println("omg hi")
19     cout.flush()
20 }
```

#### 1.2.3. constexpr

Some default limits in gcc (7.x - trunk):

- constexpr recursion depth: 512
- constexpr loop iteration per function: 262144
- constexpr operation count per function: 33554432
- template recursion depth: 900 (gcc *might* segfault first)

```
1 constexpr array<int, 10> fibonacci{[] {
2     array<int, 10> a{};
3     a[0] = a[1] = 1;
4     for (int i = 2; i < 10; i++) a[i] = a[i - 1] + a[i - 2];
5     return a;
6 }}();
7 static_assert(fibonacci[9] == 55, "CE");
8
9 template <typename F, typename INT, INT... S>
10 constexpr void for_constexpr(integer_sequence<INT, S...>,
11                               F &&func) {
12     int _[] = {(func(integral_constant<INT, S>{}), 0)...};
13 }
14 // example
15 template <typename... T> void print_tuple(tuple<T...> t) {
16     for_constexpr(make_index_sequence<sizeof...(T)>{}),
17         [&](auto i) { cout << get<i>(t) << '\n'; };
18 }
```

#### 1.2.4. Bump Allocator

```
1
2
3 // global bump allocator
4 char mem[256 << 20]; // 256 MB
5 size_t rsp = sizeof mem;
6 void *operator new(size_t s) {
7     assert(s < rsp); // MLE
8     return (void *)&mem[rsp -= s];
9 }
10 void operator delete(void *) {}
11
12 // bump allocator for STL / pbds containers
13 char mem[256 << 20];
14 size_t rsp = sizeof mem;
15 template <typename T> struct bump {
16     typedef T value_type;
17     bump() {}
18     template <typename U> bump(U, ...) {}
19     T *allocate(size_t n) {
20         rsp -= n * sizeof(T);
21         rsp &= 0 - alignof(T);
22         return (T *)(&mem + rsp);
23     }
24     void deallocate(T *, size_t n) {}
25 };
```

### 1.3. Tools

#### 1.3.1. Floating Point Binary Search

```

1 union di {
2     double d;
3     ull i;
4 };
5 bool check(double);
6 // binary search in [L, R) with relative error 2^-eps
7 double binary_search(double L, double R, int eps) {
8     di l = {L}, r = {R}, m;
9     while (r.i - l.i > 1LL << (52 - eps)) {
10         m.i = (l.i + r.i) >> 1;
11         if (check(m.d)) r = m;
12         else l = m;
13     }
14     return l.d;
15 }

```

#### 1.3.2. SplitMix64

```

1 using ull = unsigned long long;
2 inline ull splitmix64(ull x) {
3     // change to `static ull x = SEED;` for DRBG
4     ull z = (x += 0x9E3779B97F4A7C15);
5     z = (z ^ (z >> 30)) * 0xBF58476D1CE4E5B9;
6     z = (z ^ (z >> 27)) * 0x94D049BB133111EB;
7     return z ^ (z >> 31);
8 }

```

#### 1.3.3. <random>

```

1 #ifdef __unix__
2     random_device rd;
3     mt19937_64 RNG(rd());
4 #else
5     const auto SEED = chrono::high_resolution_clock::now()
6         .time_since_epoch()
7         .count();
8     mt19937_64 RNG(SEED);
9 #endif
10 // random uint_fast64_t: RNG();
11 // uniform random of type T (int, double, ...) in [l, r]:
12 // uniform_int_distribution<T> dist(l, r); dist(RNG);

```

#### 1.3.4. x86 Stack Hack

```

1 constexpr size_t size = 200 << 20; // 200MiB
2 int main() {
3     register long rsp asm("rsp");
4     char *buf = new char[size];
5     asm("movq %0, %%rsp\n" :: "r"(buf + size));
6     // do stuff
7     asm("movq %0, %%rsp\n" :: "r"(rsp));
8     delete[] buf;
9 }

```

### 1.4. Algorithms

#### 1.4.1. Bit Hacks

```

1 // next permutation of x as a bit sequence
2 ull next_bits_permutation(ull x) {
3     ull c = __builtin_ctzll(x), r = x + (1ULL << c);
4     return (r ^ x) >> (c + 2) | r;
5 }
6 // iterate over all (proper) subsets of bitset s
7 void subsets(ull s) {
8     for (ull x = s; x;) { --x &= s; /* do stuff */ }
9 }

```

#### 1.4.2. Aliens Trick

```

1 // min dp[i] value and its i (smallest one)
2 pll get_dp(int cost);
3 ll aliens(int k, int l, int r) {
4     while (l != r) {
5         int m = (l + r) / 2;
6         auto [f, s] = get_dp(m);
7         if (s == k) return f - m * k;
8         if (s < k) r = m;
9         else l = m + 1;
10     }
11     return get_dp(l).first - l * k;
12 }

```

#### 1.4.3. Hilbert Curve

```

1 ll hilbert(ll n, int x, int y) {
2     ll res = 0;
3     for (ll s = n; s /= 2;) {
4         int rx = !!((x & s), ry = !!((y & s);
5         res += s * s * ((3 * rx) ^ ry);
6         if (ry == 0) {
7             if (rx == 1) x = s - 1 - x, y = s - 1 - y;
8             swap(x, y);
9         }
10    }
11    return res;
12 }

```

#### 1.4.4. Infinite Grid Knight Distance

```

1 ll get_dist(ll dx, ll dy) {
2     if (++(dx = abs(dx)) > ++(dy = abs(dy))) swap(dx, dy);
3     if (dx == 1 && dy == 2) return 3;
4     if (dx == 3 && dy == 3) return 4;
5     ll lb = max(dy / 2, (dx + dy) / 3);
6     return ((dx ^ dy ^ lb) & 1) ? ++lb : lb;
7 }

```

#### 1.4.5. Poker Hand

```

1
2
3
4
5
6
7 using namespace std;
8
9 struct hand {
10     static constexpr auto rk = [] {
11         array<int, 256> x{};
12         auto s = "23456789TJQKACDHS";
13         for (int i = 0; i < 17; i++) x[s[i]] = i % 13;
14         return x;
15     }();
16     vector<pair<int, int>> v;
17     vector<int> cnt, vf, vs;
18     int type;
19     hand() : cnt(4), type(0) {}
20     void add_card(char suit, char rank) {
21         ++cnt[rk[suit]];
22         for (auto &[f, s] : v)
23             if (s == rk[rank]) return ++f, void();
24         v.emplace_back(1, rk[rank]);
25     }
26     void process() {
27         sort(v.rbegin(), v.rend());
28         for (auto [f, s] : v) vf.push_back(f), vs.push_back(s);
29         bool str = 0, flu = find(all(cnt), 5) != cnt.end();
30         if ((str = v.size() == 5))
31             for (int i = 1; i < 5; i++)
32                 if (vs[i] != vs[i - 1] + 1) str = 0;
33         if (vs == vector<int>{12, 3, 2, 1, 0})
34             str = 1, vs = {3, 2, 1, 0, -1};
35         if (str && flu) type = 9;
36         else if (vf[0] == 4) type = 8;
37         else if (vf[0] == 3 && vf[1] == 2) type = 7;
38         else if (str || flu) type = 5 + flu;
39         else if (vf[0] == 3) type = 4;
40         else if (vf[0] == 2) type = 2 + (vf[1] == 2);
41         else type = 1;
42     }
43     bool operator<(const hand &b) const {
44         return make_tuple(type, vf, vs) <
45             make_tuple(b.type, b.vf, b.vs);
46     }
47 };

```

#### 1.4.6. Longest Increasing Subsequence

```

1
2
3 template <class I> vi lis(const vector<I> &S) {
4     if (S.empty()) return {};
5     vi prev(sz(S));
6     typedef pair<I, int> p;
7     vector<p> res;
8     rep(i, 0, sz(S)) {
9         // change 0 -> i for longest non-decreasing subsequence
10        auto it = lower_bound(all(res), p[S[i], 0]);
11        if (it == res.end())
12            res.emplace_back(), it = res.end() - 1;
13        *it = {S[i], i};
14        prev[i] = it == res.begin() ? 0 : (it - 1) ->second;
15    }
16 }

```

```

15 }
16 int L = sz(res), cur = res.back().second;
17 vi ans(L);
18 while (L--) ans[L] = cur, cur = prev[cur];
19 return ans;
20 }

```

#### 1.4.7. Mo's Algorithm on Tree

```

1 void MoAlgoOnTree() {
2     Dfs(0, -1);
3     vector<int> euler(tk);
4     for (int i = 0; i < n; ++i) {
5         euler[tin[i]] = i;
6         euler[tout[i]] = i;
7     }
8     vector<int> l(q), r(q), qr(q), sp(q, -1);
9     for (int i = 0; i < q; ++i) {
10         if (tin[u[i]] > tin[v[i]]) swap(u[i], v[i]);
11         int z = GetLCA(u[i], v[i]);
12         sp[i] = z[i];
13         if (z == u) l[i] = tin[u[i]], r[i] = tin[v[i]];
14         else l[i] = tout[u[i]], r[i] = tin[v[i]];
15         qr[i] = i;
16     }
17     sort(qr.begin(), qr.end(), [&](int i, int j) {
18         if (l[i] / kB == l[j] / kB) return r[i] < r[j];
19         return l[i] / kB < l[j] / kB;
20     });
21     vector<bool> used(n);
22     // Add(v): add/remove v to/from the path based on used[v]
23     for (int i = 0, tl = 0, tr = -1; i < q; ++i) {
24         while (tl < l[qr[i]]) Add(euler[tl++]);
25         while (tl > l[qr[i]]) Add(euler[--tl]);
26         while (tr > r[qr[i]]) Add(euler[tr--]);
27         while (tr < r[qr[i]]) Add(euler[++tr]);
28         // add/remove LCA(u, v) if necessary
29     }
30 }

```



## 2.6. Wavelet Matrix

```

1
3
5 #pragma GCC target("popcnt,bmi2")
5 #include <immintrin.h>

7 // T is unsigned. You might want to compress values first
7 template <typename T> struct wavelet_matrix {
9     static_assert(is_unsigned_v<T>, "only unsigned T");
9     struct bit_vector {
11         static constexpr uint W = 64;
11         uint n, cnt0;
13         vector<ull> bits;
13         vector<uint> sum;
15         bit_vector(uint n_)
15             : n(n_), bits(n / W + 1), sum(n / W + 1) {}
17         void build() {
17             for (uint j = 0; j != n / W; ++j)
19                 sum[j + 1] = sum[j] + _mm_popcnt_u64(bits[j]);
21             cnt0 = rank0(n);
23         void set_bit(uint i) { bits[i / W] |= 1ULL << i % W; }
23         bool operator[](uint i) const {
25             return !!(bits[i / W] & 1ULL << i % W);
25         }
27         uint rank1(uint i) const {
27             return sum[i / W] +
29                 _mm_popcnt_u64(_bzhi_u64(bits[i / W], i % W));
29         }
31         uint rank0(uint i) const { return i - rank1(i); }
31     };
33     vector<bit_vector> b;
33     wavelet_matrix(const vector<T> &a) : n(a.size()) {
35         lg =
37             _lg(max(*max_element(a.begin(), a.end()), T(1))) + 1;
37         b.assign(lg, n);
39         vector<T> cur = a, nxt(n);
39         for (int h = lg; h--;) {
41             for (uint i = 0; i < n; ++i)
43                 if (cur[i] & (T(1) << h)) b[h].set_bit(i);
43             b[h].build();
45             int il = 0, ir = b[h].cnt0;
45             for (uint i = 0; i < n; ++i)
47                 nxt[(b[h][i] ? ir : il)++] = cur[i];
47             swap(cur, nxt);
49         }
49         T operator[](uint i) const {
51             T res = 0;
51             for (int h = lg; h--;)
53                 if (b[h][i])
55                     i += b[h].cnt0 - b[h].rank0(i), res |= T(1) << h;
55             return res;
57         }
57         // query k-th smallest (0-based) in a[l, r)
57         T kth(uint l, uint r, uint k) const {
59             T res = 0;
59             for (int h = lg; h--;) {
61                 uint tl = b[h].rank0(l), tr = b[h].rank0(r);
63                 if (k >= tr - tl) {
65                     k -= tr - tl;
65                     l += b[h].cnt0 - tl;
65                     r += b[h].cnt0 - tr;
67                     res |= T(1) << h;
67                 } else l = tl, r = tr;
69             }
69             return res;
71         }
71         // count of i in [l, r) with a[i] < u
71         uint count(uint l, uint r, T u) const {
73             if (u >= T(1) << lg) return r - l;
73             uint res = 0;
75             for (int h = lg; h--;) {
77                 uint tl = b[h].rank0(l), tr = b[h].rank0(r);
77                 if (u & (T(1) << h)) {
79                     l += b[h].cnt0 - tl;
79                     r += b[h].cnt0 - tr;
81                     res += tr - tl;
81                 } else l = tl, r = tr;
83             }
83             return res;
85 };

```

## 2.7. Link-Cut Tree

```

1
3 const int MXN = 100005;
3 const int MEM = 100005;
5
5 struct Splay {
7     static Splay nil, mem[MEM], *pmem;
7     Splay *ch[2], *f;
9     int val, rev, size;
9     Splay() : val(-1), rev(0), size(0) {
11         f = ch[0] = ch[1] = &nil;
13     }
13     Splay(int _val) : val(_val), rev(0), size(1) {
15         f = ch[0] = ch[1] = &nil;
15     }
17     bool isr() {
17         return f->ch[0] != this && f->ch[1] != this;
19     }
19     int dir() { return f->ch[0] == this ? 0 : 1; }
19     void setCh(Splay *c, int d) {
21         ch[d] = c;
21         if (c != &nil) c->f = this;
23         pull();
25     }
25     void push() {
27         if (rev) {
27             swap(ch[0], ch[1]);
27             if (ch[0] != &nil) ch[0]->rev ^= 1;
29             if (ch[1] != &nil) ch[1]->rev ^= 1;
29             rev = 0;
31         }
33     }
33     void pull() {
35         size = ch[0]->size + ch[1]->size + 1;
35         if (ch[0] != &nil) ch[0]->f = this;
37         if (ch[1] != &nil) ch[1]->f = this;
39     }
39     Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::mem;
39     Splay *nil = &Splay::nil;
41     void rotate(Splay *x) {
43         Splay *p = x->f;
43         int d = x->dir();
45         if (!p->isr()) p->f->setCh(x, p->dir());
45         else x->f = p->f;
47         p->setCh(x->ch[!d], d);
47         x->setCh(p, !d);
49         p->pull();
49         x->pull();
51     }
51     vector<Splay*> splayVec;
53     void splay(Splay *x) {
55         splayVec.clear();
55         for (Splay *q = x;; q = q->f) {
57             splayVec.push_back(q);
57             if (q->isr()) break;
59         }
59         reverse(begin(splayVec), end(splayVec));
59         for (auto it : splayVec) it->push();
61         while (!x->isr()) {
63             if (x->f->isr()) rotate(x);
63             else if (x->dir() == x->f->dir())
65                 rotate(x->f), rotate(x);
65             else rotate(x), rotate(x);
67         }
69     }
69     Splay *access(Splay *x) {
71         Splay *q = nil;
71         for (; x != nil; x = x->f) {
73             splay(x);
73             x->setCh(q, 1);
75             q = x;
77         }
77         return q;
79     }
79     void evert(Splay *x) {
79         access(x);
79         splay(x);
81         x->rev ^= 1;
81         x->push();
83         x->pull();
85     }
85     void link(Splay *x, Splay *y) {
87         // evert(x);
87         access(x);
87         splay(x);
89         evert(y);
89         x->setCh(y, 1);
91     }
91     void cut(Splay *x, Splay *y) {
93         // evert(x);

```

```
95     access(y);
96     splay(y);
97     y->push();
98     y->ch[0] = y->ch[0]->f = nil;
99 }
100
101 int N, Q;
102 Splay *vt[MXN];
103
104 int ask(Splay *x, Splay *y) {
105     access(x);
106     access(y);
107     splay(x);
108     int res = x->f->val;
109     if (res == -1) res = x->val;
110     return res;
111 }
112
113 int main(int argc, char **argv) {
114     scanf("%d%d", &N, &Q);
115     for (int i = 1; i <= N; i++)
116         vt[i] = new (Splay::pmem++) Splay(i);
117     while (Q--) {
118         char cmd[105];
119         int u, v;
120         scanf("%s", cmd);
121         if (cmd[1] == 'i') {
122             scanf("%d%d", &u, &v);
123             link(vt[v], vt[u]);
124         } else if (cmd[0] == 'c') {
125             scanf("%d", &v);
126             cut(vt[1], vt[v]);
127         } else {
128             scanf("%d%d", &u, &v);
129             int res = ask(vt[u], vt[v]);
130             printf("%d\n", res);
131         }
132     }
133 }
```



### 3. Graph

#### 3.1. Modeling

- Maximum/Minimum flow with lower bound / Circulation problem
  - Construct super source  $S$  and sink  $T$ .
  - For each edge  $(x, y, l, u)$ , connect  $x \rightarrow y$  with capacity  $u - l$ .
  - For each vertex  $v$ , denote by  $in(v)$  the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - If  $in(v) > 0$ , connect  $S \rightarrow v$  with capacity  $in(v)$ , otherwise, connect  $v \rightarrow T$  with capacity  $-in(v)$ .
    - To maximize, connect  $t \rightarrow s$  with capacity  $\infty$  (skip this in circulation problem), and let  $f$  be the maximum flow from  $S$  to  $T$ . If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from  $s$  to  $t$  is the answer.
    - To minimize, let  $f$  be the maximum flow from  $S$  to  $T$ . Connect  $t \rightarrow s$  with capacity  $\infty$  and let the flow from  $S$  to  $T$  be  $f'$ . If  $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise,  $f'$  is the answer.
  - The solution of each edge  $e$  is  $l_e + f_e$ , where  $f_e$  corresponds to the flow of edge  $e$  on the graph.
- Construct minimum vertex cover from maximum matching  $M$  on bipartite graph  $(X, Y)$ 
  - Redirect every edge:  $y \rightarrow x$  if  $(x, y) \in M$ ,  $x \rightarrow y$  otherwise.
  - DFS from unmatched vertices in  $X$ .
  - $x \in X$  is chosen iff  $x$  is unvisited.
  - $y \in Y$  is chosen iff  $y$  is visited.
- Minimum cost cyclic flow
  - Construct super source  $S$  and sink  $T$
  - For each edge  $(x, y, c)$ , connect  $x \rightarrow y$  with  $(cost, cap) = (c, 1)$  if  $c > 0$ , otherwise connect  $y \rightarrow x$  with  $(cost, cap) = (-c, 1)$
  - For each edge with  $c < 0$ , sum these cost as  $K$ , then increase  $d(y)$  by 1, decrease  $d(x)$  by 1
  - For each vertex  $v$  with  $d(v) > 0$ , connect  $S \rightarrow v$  with  $(cost, cap) = (0, d(v))$
  - For each vertex  $v$  with  $d(v) < 0$ , connect  $v \rightarrow T$  with  $(cost, cap) = (0, -d(v))$
  - Flow from  $S$  to  $T$ , the answer is the cost of the flow  $C + K$
- Maximum density induced subgraph
  - Binary search on answer, suppose we're checking answer  $T$
  - Construct a max flow model, let  $K$  be the sum of all weights
  - Connect source  $s \rightarrow v$ ,  $v \in G$  with capacity  $K$
  - For each edge  $(u, v, w)$  in  $G$ , connect  $u \rightarrow v$  and  $v \rightarrow u$  with capacity  $w$
  - For  $v \in G$ , connect it with sink  $v \rightarrow t$  with capacity  $K + 2T - (\sum_{e \in E(v)} w(e)) - 2w(v)$
  - $T$  is a valid answer if the maximum flow  $f < K|V|$
- Minimum weight edge cover
  - For each  $v \in V$  create a copy  $v'$ , and connect  $u' \rightarrow v'$  with weight  $w(u, v)$ .
  - Connect  $v \rightarrow v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to  $v$ .
  - Find the minimum weight perfect matching on  $G'$ .
- Project selection problem
  - If  $p_v > 0$ , create edge  $(s, v)$  with capacity  $p_v$ ; otherwise, create edge  $(v, t)$  with capacity  $-p_v$ .
  - Create edge  $(u, v)$  with capacity  $w$  with  $w$  being the cost of choosing  $u$  without choosing  $v$ .
  - The mincut is equivalent to the maximum profit of a subset of projects.
- 0/1 quadratic programming

$$\sum_x c_x x + \sum_y c_y \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y}')$$

can be minimized by the mincut of the following graph:

- Create edge  $(x, t)$  with capacity  $c_x$  and create edge  $(s, y)$  with capacity  $c_y$ .
- Create edge  $(x, y)$  with capacity  $c_{xy}$ .
- Create edge  $(x, y)$  and edge  $(x', y')$  with capacity  $c_{xyx'y'}$ .

### 3.2. Matching/Flows

#### 3.2.1. Dinic's Algorithm

```

1 struct Dinic {
2     struct edge {
3         int to, cap, flow, rev;
4     };
5     static constexpr int MAXN = 10000, MAXF = 1e9;
6     vector<edge> v[MAXN];
7     int top[MAXN], deep[MAXN], side[MAXN], s, t;
8     void make_edge(int s, int t, int cap) {
9         v[s].push_back({t, cap, 0, (int)v[t].size()});
10        v[t].push_back({s, 0, 0, (int)v[s].size() - 1});
11    }
12    int dfs(int a, int flow) {
13        if (a == t || !flow) return flow;
14        for (int &i = top[a]; i < v[a].size(); i++) {
15            edge &e = v[a][i];
16            if (deep[a] + 1 == deep[e.to] && e.cap - e.flow) {
17                int x = dfs(e.to, min(e.cap - e.flow, flow));
18                if (x) {
19                    e.flow += x, v[e.to][e.rev].flow -= x;
20                    return x;
21                }
22            }
23        }
24        deep[a] = -1;
25        return 0;
26    }
27    bool bfs() {
28        queue<int> q;
29        fill_n(deep, MAXN, 0);
30        q.push(s), deep[s] = 1;
31        int tmp;
32        while (!q.empty()) {
33            tmp = q.front(), q.pop();
34            for (edge e : v[tmp])
35                if (!deep[e.to] && e.cap != e.flow)
36                    deep[e.to] = deep[tmp] + 1, q.push(e.to);
37        }
38        return deep[t];
39    }
40    int max_flow(int _s, int _t) {
41        s = _s, t = _t;
42        int flow = 0, tflow;
43        while (bfs()) {
44            fill_n(top, MAXN, 0);
45            while ((tflow = dfs(s, MAXF))) flow += tflow;
46        }
47        return flow;
48    }
49    void reset() {
50        fill_n(side, MAXN, 0);
51        for (auto &i : v) i.clear();
52    }
53 };

```

#### 3.2.2. Minimum Cost Flow

```

1 struct MCF {
2     struct edge {
3         ll to, from, cap, flow, cost, rev;
4     } * fromE[MAXN];
5     vector<edge> v[MAXN];
6     ll n, s, t, flows[MAXN], dis[MAXN], pi[MAXN], flowlim;
7     void make_edge(int s, int t, ll cap, ll cost) {
8         if (!cap) return;
9         v[s].pb(edge{t, s, cap, 0LL, cost, v[t].size()});
10        v[t].pb(edge{s, t, 0LL, 0LL, -cost, v[s].size() - 1});
11    }
12    bitset<MAXN> vis;
13    void dijkstra() {
14        vis.reset();
15        __gnu_pbds::priority_queue<pair<ll, int>> q;
16        vector<decltype(q)::point_iterator> its(n);
17        q.push({0LL, s});
18        while (!q.empty()) {
19            int now = q.top().second;
20            q.pop();
21            if (vis[now]) continue;
22            vis[now] = 1;
23            ll ndis = dis[now] + pi[now];
24            for (edge &e : v[now]) {
25                if (e.flow == e.cap || vis[e.to]) continue;
26                if (dis[e.to] > ndis + e.cost - pi[e.to]) {
27                    dis[e.to] = ndis + e.cost - pi[e.to];
28                    flows[e.to] = min(flows[now], e.cap - e.flow);
29                    fromE[e.to] = &e;
30                    if (its[e.to] == q.end())
31                        its[e.to] = q.push({-dis[e.to], e.to});
32                    else q.modify(its[e.to], {-dis[e.to], e.to});
33                }
34            }
35        }
36    }
37 };

```



```

33     }
34 }
35 }
36 }
37 bool AP(ll &flow) {
38     fill_n(dis, n, INF);
39     fromE[s] = 0;
40     dis[s] = 0;
41     flows[s] = flowlim - flow;
42     dijkstra();
43     if (dis[t] == INF) return false;
44     flow += flows[t];
45     for (edge *e = fromE[t]; e; e = fromE[e->from]) {
46         e->flow += flows[t];
47         v[e->to][e->rev].flow -= flows[t];
48     }
49     for (int i = 0; i < n; i++)
50         pi[i] = min(pi[i] + dis[i], INF);
51     return true;
52 }
53 pll solve(int _s, int _t, ll _flowlim = INF) {
54     s = _s, t = _t, flowlim = _flowlim;
55     pll re;
56     while (re.F != flowlim && AP(re.F))
57         ;
58     for (int i = 0; i < n; i++)
59         for (edge &e : v[i])
60             if (e.flow != 0) re.S += e.flow * e.cost;
61     re.S /= 2;
62     return re;
63 }
64 void init(int _n) {
65     n = _n;
66     fill_n(pi, n, 0);
67     for (int i = 0; i < n; i++) v[i].clear();
68 }
69 void setpi(int s) {
70     fill_n(pi, n, INF);
71     pi[s] = 0;
72     for (ll it = 0, flag = 1, tdis; flag && it < n; it++) {
73         flag = 0;
74         for (int i = 0; i < n; i++)
75             if (pi[i] != INF)
76                 for (edge &e : v[i])
77                     if (e.cap && (tdis = pi[i] + e.cost) < pi[e.to])
78                         pi[e.to] = tdis, flag = 1;
79     }
80 }
81 };

```

### 3.2.3. Gomory-Hu Tree

Requires: Dinic's Algorithm

```

1
3 int e[MAXN][MAXN];
4 int p[MAXN];
5 Dinic D; // original graph
6 void gomory_hu() {
7     fill(p, p + n, 0);
8     fill(e[0], e[n], INF);
9     for (int s = 1; s < n; s++) {
10         int t = p[s];
11         Dinic F = D;
12         int tmp = F.max_flow(s, t);
13         for (int i = 1; i < s; i++)
14             e[s][i] = e[i][s] = min(tmp, e[t][i]);
15         for (int i = s + 1; i <= n; i++)
16             if (p[i] == t && F.side[i]) p[i] = s;
17     }
18 }

```

### 3.2.4. Global Minimum Cut

```

1
3 // weights is an adjacency matrix, undirected
4 pair<int, vi> getMinCut(vector<vi> &weights) {
5     int N = sz(weights);
6     vi used(N), cut, best_cut;
7     int best_weight = -1;
8
9     for (int phase = N - 1; phase >= 0; phase--) {
10         vi w = weights[0], added = used;
11         int prev, k = 0;
12         rep(i, 0, phase) {
13             prev = k;
14             k = -1;
15             rep(j, 1, N) if (!added[j] &&
16                             (k == -1 || w[j] > w[k])) k = j;

```

```

17         if (i == phase - 1) {
18             rep(j, 0, N) weights[prev][j] += weights[k][j];
19             rep(j, 0, N) weights[j][prev] = weights[prev][j];
20             used[k] = true;
21             cut.push_back(k);
22             if (best_weight == -1 || w[k] < best_weight) {
23                 best_cut = cut;
24                 best_weight = w[k];
25             }
26         } else {
27             rep(j, 0, N) w[j] += weights[k][j];
28             added[k] = true;
29         }
30     }
31     return {best_weight, best_cut};
32 }

```

### 3.2.5. Bipartite Minimum Cover

Requires: Dinic's Algorithm

```

1
3 // maximum independent set = all vertices not covered
4 // x : [0, n), y : [0, m]
5 struct Bipartite_vertex_cover {
6     Dinic D;
7     int n, m, s, t, x[maxn], y[maxn];
8     void make_edge(int x, int y) { D.make_edge(x, y + n, 1); }
9     int matching() {
10         int re = D.max_flow(s, t);
11         for (int i = 0; i < n; i++)
12             for (Dinic::edge &e : D.v[i])
13                 if (e.to != s && e.flow == 1) {
14                     x[i] = e.to - n, y[e.to - n] = i;
15                     break;
16                 }
17         return re;
18     }
19     // init() and matching() before use
20     void solve(vector<int> &vx, vector<int> &vy) {
21         bitset<maxn * 2 + 10> vis;
22         queue<int> q;
23         for (int i = 0; i < n; i++)
24             if (x[i] == -1) q.push(i), vis[i] = 1;
25         while (!q.empty()) {
26             int now = q.front();
27             q.pop();
28             if (now < n) {
29                 for (Dinic::edge &e : D.v[now])
30                     if (e.to != s && e.to - n != x[now] && !vis[e.to])
31                         vis[e.to] = 1, q.push(e.to);
32             } else {
33                 if (!vis[y[now - n]])
34                     vis[y[now - n]] = 1, q.push(y[now - n]);
35             }
36         }
37         for (int i = 0; i < n; i++)
38             if (!vis[i]) vx.pb(i);
39         for (int i = 0; i < m; i++)
40             if (vis[i + n]) vy.pb(i);
41     }
42     void init(int _n, int _m) {
43         n = _n, m = _m, s = n + m, t = s + 1;
44         for (int i = 0; i < n; i++)
45             x[i] = -1, D.make_edge(s, i, 1);
46         for (int i = 0; i < m; i++)
47             y[i] = -1, D.make_edge(i + n, t, 1);
48     }
49 };

```

### 3.2.6. Edmonds' Algorithm

```

1
3 struct Edmonds {
4     int n, T;
5     vector<vector<int>> g;
6     vector<int> pa, p, used, base;
7     Edmonds(int n) :
8         n(n), T(0), g(n), pa(n, -1), p(n), used(n),
9         base(n) {}
10     void add(int a, int b) {
11         g[a].push_back(b);
12         g[b].push_back(a);
13     }
14     int getBase(int i) {
15         while (i != base[i])
16             base[i] = base[base[i]], i = base[i];
17         return i;

```

```

}
vector<int> toJoin;
void mark_path(int v, int x, int b, vector<int> &path) {
    for (; getBase(v) != b; v = p[x]) {
        p[v] = x, x = pa[v];
        toJoin.push_back(v);
        toJoin.push_back(x);
        if (!used[x]) used[x] = ++T, path.push_back(x);
    }
}
bool go(int v) {
    for (int x : g[v]) {
        int b, bv = getBase(v), bx = getBase(x);
        if (bv == bx) {
            continue;
        } else if (used[x]) {
            vector<int> path;
            toJoin.clear();
            if (used[bx] < used[bv])
                mark_path(v, x, b = bx, path);
            else mark_path(x, v, b = bv, path);
            for (int z : toJoin) base[getBase(z)] = b;
            for (int z : path)
                if (go(z)) return 1;
        } else if (p[x] == -1) {
            p[x] = v;
            if (pa[x] == -1) {
                for (int y; x != -1; x = v)
                    y = p[x], v = pa[y], pa[x] = y, pa[y] = x;
                return 1;
            }
            if (!used[pa[x]]) {
                used[pa[x]] = ++T;
                if (go(pa[x])) return 1;
            }
        }
    }
    return 0;
}
void init_dfs() {
    for (int i = 0; i < n; i++)
        used[i] = 0, p[i] = -1, base[i] = i;
}
bool dfs(int root) {
    used[root] = ++T;
    return go(root);
}
void match() {
    int ans = 0;
    for (int v = 0; v < n; v++)
        for (int x : g[v])
            if (pa[v] == -1 && pa[x] == -1) {
                pa[v] = x, pa[x] = v, ans++;
                break;
            }
    init_dfs();
    for (int i = 0; i < n; i++)
        if (pa[i] == -1 && dfs(i)) ans++, init_dfs();
    cout << ans * 2 << "\n";
    for (int i = 0; i < n; i++)
        if (pa[i] > i)
            cout << i + 1 << " " << pa[i] + 1 << "\n";
}
};

```

### 3.2.7. Minimum Weight Matching

```

1 struct Graph {
2     static const int MAXN = 105;
3     int n, e[MAXN][MAXN];
4     int match[MAXN], d[MAXN], onstk[MAXN];
5     vector<int> stk;
6     void init(int _n) {
7         n = _n;
8         for (int i = 0; i < n; i++)
9             for (int j = 0; j < n; j++)
10                // change to appropriate infinity
11                // if not complete graph
12                e[i][j] = 0;
13    }
14    void add_edge(int u, int v, int w) {
15        e[u][v] = e[v][u] = w;
16    }
17    bool SPFA(int u) {
18        if (onstk[u]) return true;
19        stk.push_back(u);
20        onstk[u] = 1;
21        for (int v = 0; v < n; v++) {
22            if (u != v && match[u] != v && !onstk[v]) {
23                int m = match[v];
24                if (d[m] > d[u] - e[v][m] + e[u][v]) {

```

```

25                d[m] = d[u] - e[v][m] + e[u][v];
26                onstk[v] = 1;
27                stk.push_back(v);
28                if (SPFA(m)) return true;
29                stk.pop_back();
30                onstk[v] = 0;
31            }
32        }
33    }
34    onstk[u] = 0;
35    stk.pop_back();
36    return false;
37 }
38 int solve() {
39     for (int i = 0; i < n; i += 2) {
40         match[i] = i + 1;
41         match[i + 1] = i;
42     }
43     while (true) {
44         int found = 0;
45         for (int i = 0; i < n; i++) onstk[i] = d[i] = 0;
46         for (int i = 0; i < n; i++) {
47             stk.clear();
48             if (!onstk[i] && SPFA(i)) {
49                 found = 1;
50                 while (stk.size() >= 2) {
51                     int u = stk.back();
52                     stk.pop_back();
53                     int v = stk.back();
54                     stk.pop_back();
55                     match[u] = v;
56                     match[v] = u;
57                 }
58             }
59             if (!found) break;
60         }
61         int ret = 0;
62         for (int i = 0; i < n; i++) ret += e[i][match[i]];
63         ret /= 2;
64         return ret;
65     }
66 } graph;

```

### 3.2.8. Stable Marriage

```

1 // normal stable marriage problem
2 /* input:
3 3
4 Albert Laura Nancy Marcy
5 Brad Marcy Nancy Laura
6 Chuck Laura Marcy Nancy
7 Laura Chuck Albert Brad
8 Marcy Albert Chuck Brad
9 Nancy Brad Albert Chuck
10 */
11
12 using namespace std;
13 const int MAXN = 505;
14
15 int n;
16 int favor[MAXN][MAXN]; // favor[boy_id][rank] = girl_id;
17 int order[MAXN][MAXN]; // order[girl_id][boy_id] = rank;
18 int current[MAXN]; // current[boy_id] = rank;
19 // boy_id will pursue current[boy_id] girl.
20 int girl_current[MAXN]; // girl[girl_id] = boy_id;
21
22 void initialize() {
23     for (int i = 0; i < n; i++) {
24         current[i] = 0;
25         girl_current[i] = n;
26         order[i][n] = n;
27     }
28 }
29
30 map<string, int> male, female;
31 string bname[MAXN], gname[MAXN];
32 int fit = 0;
33
34 void stable_marriage() {
35     queue<int> que;
36     for (int i = 0; i < n; i++) que.push(i);
37     while (!que.empty()) {
38         int boy_id = que.front();
39         que.pop();
40
41         int girl_id = favor[boy_id][current[boy_id]];
42         current[boy_id]++;

```

```

47     if (order[girl_id][boy_id] <
        order[girl_id][girl_current[girl_id]]) {
49         if (girl_current[girl_id] < n)
            que.push(girl_current[girl_id]);
        girl_current[girl_id] = boy_id;
51     } else {
        que.push(boy_id);
53     }
55 }

57 int main() {
    cin >> n;

59     for (int i = 0; i < n; i++) {
        string p, t;
        cin >> p;
61         male[p] = i;
        bname[i] = p;
63         for (int j = 0; j < n; j++) {
            cin >> t;
65             if (!female.count(t)) {
                gname[fit] = t;
                female[t] = fit++;
67             }
            favor[i][j] = female[t];
71         }
73     }

75     for (int i = 0; i < n; i++) {
        string p, t;
        cin >> p;
77         for (int j = 0; j < n; j++) {
            cin >> t;
79             order[female[p]][male[t]] = j;
81         }
83     }

    initialize();
    stable_marriage();

85     for (int i = 0; i < n; i++) {
        cout << bname[i] << " "
87         << gname[favor[i][current[i] - 1]] << endl;
89     }
91 }

```

### 3.2.9. Kuhn-Munkres algorithm

```

1 // Maximum Weight Perfect Bipartite Matching
// Detect non-perfect-matching:
3 // 1. set all edge[i][j] as INF
// 2. if solve() >= INF, it is not perfect matching.

5 typedef long long ll;
7 struct KM {
    static const int MAXN = 1050;
9     static const ll INF = 1LL << 60;
    int n, match[MAXN], vx[MAXN], vy[MAXN];
11     ll edge[MAXN][MAXN], lx[MAXN], ly[MAXN], slack[MAXN];
    void init(int _n) {
        n = _n;
13         for (int i = 0; i < n; i++)
            for (int j = 0; j < n; j++) edge[i][j] = 0;
15     }
17     void add_edge(int x, int y, ll w) { edge[x][y] = w; }
    bool DFS(int x) {
19         vx[x] = 1;
        for (int y = 0; y < n; y++) {
21             if (vy[y]) continue;
            if (lx[x] + ly[y] > edge[x][y]) {
23                 slack[y] =
                    min(slack[y], lx[x] + ly[y] - edge[x][y]);
25             } else {
                vy[y] = 1;
27                 if (match[y] == -1 || DFS(match[y])) {
                    match[y] = x;
29                     return true;
31                 }
33             }
            return false;
35     }
    ll solve() {
        fill(match, match + n, -1);
        fill(lx, lx + n, -INF);
        fill(ly, ly + n, 0);
37         for (int i = 0; i < n; i++)
            for (int j = 0; j < n; j++)
39                 lx[i] = max(lx[i], edge[i][j]);
41         for (int i = 0; i < n; i++) {

```

```

43         fill(slack, slack + n, INF);
        while (true) {
45             fill(vx, vx + n, 0);
            fill(vy, vy + n, 0);
47             if (DFS(i)) break;
            ll d = INF;
49             for (int j = 0; j < n; j++)
                if (!vy[j]) d = min(d, slack[j]);
51             for (int j = 0; j < n; j++) {
                if (vx[j]) lx[j] -= d;
53                 if (vy[j]) ly[j] += d;
                else slack[j] -= d;
55             }
57         }
        ll res = 0;
59         for (int i = 0; i < n; i++) {
            res += edge[match[i]][i];
61         }
        return res;
63     }
} graph;

```

### 3.3. Shortest Path Faster Algorithm

```

1 struct SPFA {
    static const int maxn = 1010, INF = 1e9;
3     int dis[maxn];
    bitset<maxn> inq, inneg;
5     queue<int> q, tq;
    vector<pii> v[maxn];
7     void make_edge(int s, int t, int w) {
        v[s].emplace_back(t, w);
9     }
    void dfs(int a) {
11         inneg[a] = 1;
        for (pii i : v[a])
13             if (!inneg[i.F]) dfs(i.F);
15     }
    bool solve(int n, int s) { // true if have neg-cycle
        for (int i = 0; i <= n; i++) dis[i] = INF;
17         dis[s] = 0, q.push(s);
        for (int i = 0; i < n; i++) {
19             inq.reset();
            int now;
21             while (!q.empty()) {
                now = q.front(), q.pop();
23                 for (pii &i : v[now]) {
                    if (dis[i.F] > dis[now] + i.S) {
25                         dis[i.F] = dis[now] + i.S;
                        if (!inq[i.F]) tq.push(i.F), inq[i.F] = 1;
27                     }
                }
29             }
            q.swap(tq);
31         }
        bool re = !q.empty();
33         inneg.reset();
        while (!q.empty()) {
35             if (!inneg[q.front()]) dfs(q.front());
            q.pop();
37         }
        return re;
39     }
    void reset(int n) {
41         for (int i = 0; i <= n; i++) v[i].clear();
43     }
};

```

### 3.4. Strongly Connected Components

```

1 struct TarjanScc {
    int n, step;
3     vector<int> time, low, instk, stk;
    vector<vector<int>> e, scc;
5     TarjanScc(int n_)
        : n(n_), step(0), time(n), low(n), instk(n), e(n) {}
7     void add_edge(int u, int v) { e[u].push_back(v); }
    void dfs(int x) {
9         time[x] = low[x] = ++step;
        stk.push_back(x);
        instk[x] = 1;
11         for (int y : e[x])
            if (!time[y]) {
13                 dfs(y);
                low[x] = min(low[x], low[y]);
15             } else if (instk[y]) {
                low[x] = min(low[x], time[y]);
17             }
        if (time[x] == low[x]) {
19             scc.emplace_back();

```

```

21     for (int y = -1; y != x;) {
22         y = stk.back();
23         stk.pop_back();
24         instk[y] = 0;
25         scc.back().push_back(y);
26     }
27 }
28
29 void solve() {
30     for (int i = 0; i < n; i++)
31         if (!time[i]) dfs(i);
32     reverse(scc.begin(), scc.end());
33     // scc in topological order
34 }
35 };

```

### 3.4.1. 2-Satisfiability

Requires: Strongly Connected Components

```

1
2 // 1 based, vertex in SCC = MAXN * 2
3 // (not i) is i + n
4
5 struct two_SAT {
6     int n, ans[MAXN];
7     SCC S;
8     void imply(int a, int b) { S.make_edge(a, b); }
9     bool solve(int _n) {
10         n = _n;
11         S.solve(n * 2);
12         for (int i = 1; i <= n; i++) {
13             if (S.scc[i] == S.scc[i + n]) return false;
14             ans[i] = (S.scc[i] < S.scc[i + n]);
15         }
16         return true;
17     }
18     void init(int _n) {
19         n = _n;
20         fill_n(ans, n + 1, 0);
21         S.init(n * 2);
22     }
23 } SAT;

```

## 3.5. Biconnected Components

### 3.5.1. Articulation Points

```

1 void dfs(int x, int p) {
2     tin[x] = low[x] = ++t;
3     int ch = 0;
4     for (auto u : g[x])
5         if (u.first != p) {
6             if (!ins[u.second])
7                 st.push(u.second), ins[u.second] = true;
8             if (tin[u.first]) {
9                 low[x] = min(low[x], tin[u.first]);
10                continue;
11            }
12            ++ch;
13            dfs(u.first, x);
14            low[x] = min(low[x], low[u.first]);
15            if (low[u.first] >= tin[x]) {
16                cut[x] = true;
17                ++sz;
18                while (true) {
19                    int e = st.top();
20                    st.pop();
21                    bcc[e] = sz;
22                    if (e == u.second) break;
23                }
24            }
25        }
26    if (ch == 1 && p == -1) cut[x] = false;
27 }

```

### 3.5.2. Bridges

```

1 // if there are multi-edges, then they are not bridges
2 void dfs(int x, int p) {
3     tin[x] = low[x] = ++t;
4     st.push(x);
5     for (auto u : g[x])
6         if (u.first != p) {
7             if (tin[u.first]) {
8                 low[x] = min(low[x], tin[u.first]);
9                 continue;
10            }
11            dfs(u.first, x);
12            low[x] = min(low[x], low[u.first]);
13            if (low[u.first] == tin[u.first]) br[u.second] = true;

```

```

15     }
16     if (tin[x] == low[x]) {
17         ++sz;
18         while (st.size()) {
19             int u = st.top();
20             st.pop();
21             bcc[u] = sz;
22             if (u == x) break;
23         }
24     }
25 }

```

## 3.6. Triconnected Components

```

1
2 // requires a union-find data structure
3
4 struct ThreeEdgeCC {
5     int V, ind;
6     vector<int> id, pre, post, low, deg, path;
7     vector<vector<int>> components;
8     UnionFind uf;
9     template <class Graph>
10     void dfs(const Graph &G, int v, int prev) {
11         pre[v] = ++ind;
12         for (int w : G[v])
13             if (w != v) {
14                 if (w == prev) {
15                     prev = -1;
16                     continue;
17                 }
18                 if (pre[w] != -1) {
19                     if (pre[w] < pre[v]) {
20                         deg[v]++;
21                         low[v] = min(low[v], pre[w]);
22                     } else {
23                         deg[v]--;
24                         int &u = path[v];
25                         for (; u != -1 && pre[u] <= pre[w] &&
26                             pre[w] <= post[u];) {
27                             uf.join(v, u);
28                             deg[v] += deg[u];
29                             u = path[u];
30                         }
31                     }
32                 }
33                 continue;
34             }
35         dfs(G, w, v);
36         if (path[w] == -1 && deg[w] <= 1) {
37             deg[v] += deg[w];
38             low[v] = min(low[v], low[w]);
39             continue;
40         }
41         if (deg[w] == 0) w = path[w];
42         if (low[v] > low[w]) {
43             low[v] = min(low[v], low[w]);
44             swap(w, path[v]);
45         }
46         for (; w != -1; w = path[w]) {
47             uf.join(v, w);
48             deg[v] += deg[w];
49         }
50     }
51     post[v] = ind;
52 }
53 template <class Graph>
54 ThreeEdgeCC(const Graph &G)
55     : V(G.size()), ind(-1), id(V, -1), pre(V, -1),
56       post(V), low(V, INT_MAX), deg(V, 0), path(V, -1),
57       uf(V) {
58     for (int v = 0; v < V; v++)
59         if (pre[v] == -1) dfs(G, v, -1);
60     components.reserve(uf.cnt);
61     for (int v = 0; v < V; v++)
62         if (uf.find(v) == v) {
63             id[v] = components.size();
64             components.emplace_back(1, v);
65             components.back().reserve(uf.getSize(v));
66         }
67     for (int v = 0; v < V; v++)
68         if (id[v] == -1)
69             components[id[v] = id[uf.find(v)]] .push_back(v);
70 }
71 };

```

## 3.7. Centroid Decomposition

```

1 void get_center(int now) {
2     v[now] = true;
3     vtx.push_back(now);

```



```

35     C[1].clear(), C[2].clear();
36     for (auto v : T) {
37         int k = 1;
38         auto f = [&](int i) { return e[v.i][i]; };
39         while (any_of(all(C[k]), f)) k++;
40         if (k > mxk) mxk = k, C[mxk + 1].clear();
41         if (k < mnk) T[j++].i = v.i;
42         C[k].push_back(v.i);
43     }
44     if (j > 0) T[j - 1].d = 0;
45     rep(k, mnk, mxk + 1) for (int i : C[k]) T[j].i = i,
46                             T[j++].d = k;
47     expand(T, lev + 1);
48 } else if (sz(q) > sz(qmax)) qmax = q;
49 q.pop_back(), R.pop_back();
50 }
51 }
52 vi maxClique() {
53     init(V), expand(V);
54     return qmax;
55 }
56 Maxclique(vb conn)
57 : e(conn), C(sz(e) + 1), S(sz(C)), old(S) {
58     rep(i, 0, sz(e)) V.push_back({i});
59 }
60 };

```

### 3.11. Dominator Tree

```

1 // idom[n] is the unique node that strictly dominates n but
2 // does not strictly dominate any other node that strictly
3 // dominates n. idom[n] = 0 if n is entry or the entry
4 // cannot reach n.
5 struct DominatorTree {
6     static const int MAXN = 200010;
7     int n, s;
8     vector<int> g[MAXN], pred[MAXN];
9     vector<int> cov[MAXN];
10    int dfn[MAXN], nfd[MAXN], ts;
11    int par[MAXN];
12    int sdom[MAXN], idom[MAXN];
13    int mom[MAXN], mn[MAXN];
14
15    inline bool cmp(int u, int v) { return dfn[u] < dfn[v]; }
16
17    int eval(int u) {
18        if (mom[u] == u) return u;
19        int res = eval(mom[u]);
20        if (cmp(sdom[mn[mom[u]]], sdom[mn[u]]))
21            mn[u] = mn[mom[u]];
22        return mom[u] = res;
23    }
24
25    void init(int _n, int _s) {
26        n = _n;
27        s = _s;
28        REP1(i, 1, n) {
29            g[i].clear();
30            pred[i].clear();
31            idom[i] = 0;
32        }
33    }
34    void add_edge(int u, int v) {
35        g[u].push_back(v);
36        pred[v].push_back(u);
37    }
38    void DFS(int u) {
39        ts++;
40        dfn[u] = ts;
41        nfd[ts] = u;
42        for (int v : g[u])
43            if (dfn[v] == 0) {
44                par[v] = u;
45                DFS(v);
46            }
47    }
48    void build() {
49        ts = 0;
50        REP1(i, 1, n) {
51            dfn[i] = nfd[i] = 0;
52            cov[i].clear();
53            mom[i] = mn[i] = sdom[i] = i;
54        }
55        DFS(s);
56        for (int i = ts; i >= 2; i--) {
57            int u = nfd[i];
58            if (u == 0) continue;
59            for (int v : pred[u])
60                if (dfn[v]) {
61                    eval(v);

```

```

63            if (cmp(sdom[mn[v]], sdom[u]))
64                sdom[u] = sdom[mn[v]];
65        }
66        cov[sdom[u]].push_back(u);
67        mom[u] = par[u];
68        for (int w : cov[par[u]]) {
69            eval(w);
70            if (cmp(sdom[mn[w]], par[u])) idom[w] = mn[w];
71            else idom[w] = par[u];
72        }
73        cov[par[u]].clear();
74    }
75    REP1(i, 2, ts) {
76        int u = nfd[i];
77        if (u == 0) continue;
78        if (idom[u] != sdom[u]) idom[u] = idom[idom[u]];
79    }
80 } dom;

```

### 3.12. Manhattan Distance MST

```

1
2
3 // returns [(dist, from, to), ...]
4 // then do normal mst afterwards
5 typedef Point<int> P;
6 vector<array<int, 3>> manhattanMST(vector<P> ps) {
7     vi id(sz(ps));
8     iota(all(id), 0);
9     vector<array<int, 3>> edges;
10    rep(k, 0, 4) {
11        sort(all(id), [&](int i, int j) {
12            return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;
13        });
14        map<int, int> sweep;
15        for (int i : id) {
16            for (auto it = sweep.lower_bound(-ps[i].y);
17                 it != sweep.end(); sweep.erase(it++)) {
18                int j = it->second;
19                P d = ps[i] - ps[j];
20                if (d.y > d.x) break;
21                edges.push_back({d.y + d.x, i, j});
22            }
23            sweep[-ps[i].y] = i;
24        }
25        for (P &p : ps)
26            if (k & 1) p.x = -p.x;
27            else swap(p.x, p.y);
28    }
29    return edges;
30 }

```



## 4. Math

### 4.1. Number Theory

#### 4.1.1. Mod Struct

A list of safe primes: 26003, 27767, 28319, 28979, 29243, 29759, 30467, 910927547, 919012223, 947326223, 990669467, 1007939579, 1019126699, 929760389146037459, 975500632317046523, 989312547895528379

NTT prime $p$	$p - 1$	primitive root
65537	$1 \ll 16$	3
998244353	$119 \ll 23$	3
2748779069441	$5 \ll 39$	3
1945555039024054273	$27 \ll 56$	5

Requires: Extended GCD

```

1
3 template <typename T> struct M {
4     static T MOD; // change to constexpr if already known
5     T v;
6     M(T x = 0) {
7         v = (-MOD <= x && x < MOD) ? x : x % MOD;
8         if (v < 0) v += MOD;
9     }
10    explicit operator T() const { return v; }
11    bool operator==(const M &b) const { return v == b.v; }
12    bool operator!=(const M &b) const { return v != b.v; }
13    M operator-(const M &b) { return M(-v); }
14    M operator+(const M &b) { return M(v + b.v); }
15    M operator-(const M &b) { return M(v - b.v); }
16    M operator*(const M &b) { return M((__int128)v * b.v % MOD); }
17    M operator/(const M &b) { return *this * (b ^ (MOD - 2)); }
18    // change above implementation to this if MOD is not prime
19    M inv() {
20        auto [p, _, g] = extgcd(v, MOD);
21        return assert(g == 1), p;
22    }
23    friend M operator^(M a, ll b) {
24        M ans(1);
25        for (; b >= 1, a *= a)
26            if (b & 1) ans *= a;
27        return ans;
28    }
29    friend M &operator+=(M &a, M b) { return a = a + b; }
30    friend M &operator-=(M &a, M b) { return a = a - b; }
31    friend M &operator*=(M &a, M b) { return a = a * b; }
32    friend M &operator/=(M &a, M b) { return a = a / b; }
33 };
34 using Mod = M<int>;
35 template <> int Mod::MOD = 1'000'000'007;
36 int &MOD = Mod::MOD;

```

#### 4.1.2. Miller-Rabin

Requires: Mod Struct

```

1
3 // checks if Mod::MOD is prime
4 bool is_prime() {
5     if (MOD < 2 || MOD % 2 == 0) return MOD == 2;
6     Mod A[] = {2, 7, 61}; // for int values (< 2^31)
7     // ll: 2, 325, 9375, 28178, 450775, 9780504, 1795265022
8     int s = __builtin_ctzll(MOD - 1), i;
9     for (Mod a : A) {
10        Mod x = a ^ (MOD >> s);
11        for (i = 0; i < s && (x + 1).v > 2; i++) x *= x;
12        if (i && x != -1) return 0;
13    }
14    return 1;
15 }

```

#### 4.1.3. Linear Sieve

```

1 constexpr ll MAXN = 1000000;
2 bitset<MAXN> is_prime;
3 vector<ll> primes;
4 ll mpf[MAXN], phi[MAXN], mu[MAXN];
5
6 void sieve() {
7     is_prime.set();
8     is_prime[1] = 0;
9     mu[1] = phi[1] = 1;
10    for (ll i = 2; i < MAXN; i++) {
11        if (is_prime[i]) {
12            mpf[i] = i;
13            primes.push_back(i);
14            phi[i] = i - 1;
15            mu[i] = -1;
16        }
17    }
18 }

```

```

17 for (ll p : primes) {
18     if (p > mpf[i] || i * p >= MAXN) break;
19     is_prime[i * p] = 0;
20     mpf[i * p] = p;
21     mu[i * p] = -mu[i];
22     if (i % p == 0)
23         phi[i * p] = phi[i] * p, mu[i * p] = 0;
24     else phi[i * p] = phi[i] * (p - 1);
25 }
26 }
27 }

```

#### 4.1.4. Get Factors

Requires: Linear Sieve

```

1
3 vector<ll> all_factors(ll n) {
4     vector<ll> fac = {1};
5     while (n > 1) {
6         const ll p = mpf[n];
7         vector<ll> cur = {1};
8         while (n % p == 0) {
9             n /= p;
10            cur.push_back(cur.back() * p);
11        }
12        vector<ll> tmp;
13        for (auto x : fac)
14            for (auto y : cur) tmp.push_back(x * y);
15        tmp.swap(fac);
16    }
17    return fac;
18 }

```

#### 4.1.5. Binary GCD

```

1 // returns the gcd of non-negative a, b
2 ull bin_gcd(ull a, ull b) {
3     if (!a || !b) return a + b;
4     int s = __builtin_ctzll(a | b);
5     a >>= __builtin_ctzll(a);
6     while (b) {
7         if ((b >>= __builtin_ctzll(b)) < a) swap(a, b);
8         b >>= __builtin_ctzll(b);
9     }
10    return a << s;
11 }

```

#### 4.1.6. Extended GCD

```

1 // returns (p, q, g): p * a + q * b == g == gcd(a, b)
2 // g is not guaranteed to be positive when a < 0 or b < 0
3 tuple<ll, ll, ll> extgcd(ll a, ll b) {
4     ll s = 1, t = 0, u = 0, v = 1;
5     while (b) {
6         ll q = a / b;
7         swap(a -= q * b, b);
8         swap(s -= q * t, t);
9         swap(u -= q * v, v);
10    }
11    return {s, u, a};
12 }

```

#### 4.1.7. Chinese Remainder Theorem

Requires: Extended GCD

```

1 // for 0 <= a < m, 0 <= b < n, returns the smallest x >= 0
2 // such that x % m == a and x % n == b
3 ll crt(ll a, ll m, ll b, ll n) {
4     if (n > m) swap(a, b), swap(m, n);
5     auto [x, y, g] = extgcd(m, n);
6     assert((a - b) % g == 0); // no solution
7     x = ((b - a) / g * x) % (n / g) * m + a;
8     return x < 0 ? x + m / g * n : x;
9 }

```

#### 4.1.8. Baby-Step Giant-Step

Requires: Mod Struct

```

1
3 // returns x such that a ^ x = b where x \in [l, r)
4 ll bsgs(Mod a, Mod b, ll l = 0, ll r = MOD - 1) {
5     int m = sqrt(r - l) + 1, i;
6     unordered_map<ll, ll> tb;
7     Mod d = (a ^ l) / b;
8     for (i = 0, d = (a ^ l) / b; i < m; i++, d *= a)
9         if (d == 1) return l + i;
10    else tb[(ll)d] = l + i;
11    Mod c = Mod(1) / (a ^ m);
12    for (i = 0, d = 1; i < m; i++, d *= c)
13        if (auto j = tb.find((ll)d); j != tb.end())
14            return j->second + i * m;
15    return assert(0), -1; // no solution
16 }

```



#### 4.1.9. Pollard's Rho

```

1 ll f(ll x, ll mod) { return (x * x + 1) % mod; }
  // n should be composite
2 ll pollard_rho(ll n) {
3     if (!(n & 1)) return 2;
4     while (1) {
5         ll y = 2, x = RNG() % (n - 1) + 1, res = 1;
6         for (int sz = 2; res == 1; sz *= 2) {
7             for (int i = 0; i < sz && res <= 1; i++) {
8                 x = f(x, n);
9                 res = __gcd(abs(x - y), n);
10            }
11            y = x;
12        }
13        if (res != 0 && res != n) return res;
14    }
15 }

```

#### 4.1.10. Tonelli-Shanks Algorithm

Requires: Mod Struct

```

1
2
3 int legendre(Mod a) {
4     if (a == 0) return 0;
5     return (a ^ ((MOD - 1) / 2)) == 1 ? 1 : -1;
6 }
7 Mod sqrt(Mod a) {
8     assert(legendre(a) != -1); // no solution
9     ll p = MOD, s = p - 1;
10    if (a == 0) return 0;
11    if (p == 2) return 1;
12    if (p % 4 == 3) return a ^ ((p + 1) / 4);
13    int r, m;
14    for (r = 0; !(s & 1); r++) s >>= 1;
15    Mod n = 2;
16    while (legendre(n) != -1) n += 1;
17    Mod x = a ^ ((s + 1) / 2), b = a ^ s, g = n ^ s;
18    while (b != 1) {
19        Mod t = b;
20        for (m = 0; t != 1; m++) t *= g;
21        Mod gs = g ^ (1LL << (r - m - 1));
22        g = gs * gs, x *= gs, b *= g, r = m;
23    }
24    return x;
25 }
  // to get sqrt(X) modulo p^k, where p is an odd prime:
26 // c = x^2 (mod p), c = X^2 (mod p^k), q = p^(k-1)
  // X = x^q * c^((p^k-2q+1)/2) (mod p^k)

```

#### 4.1.11. Chinese Sieve

```

1 const ll N = 1000000;
  // f, g, h multiplicative, h = f (dirichlet convolution) g
2 ll pre_g(ll n);
3 ll pre_h(ll n);
4 // preprocessed prefix sum of f
5 ll pre_f[N];
6 // prefix sum of multiplicative function f
7 ll solve_f(ll n) {
8     static unordered_map<ll, ll> m;
9     if (n < N) return pre_f[n];
10    if (m.count(n)) return m[n];
11    ll ans = pre_h(n);
12    for (ll l = 2, r; l <= n; l = r + 1) {
13        r = n / (n / l);
14        ans -= (pre_g(r) - pre_g(l - 1)) * djs_f(n / l);
15    }
16    return m[n] = ans;
17 }

```

#### 4.1.12. Rational Number Binary Search

```

1 struct QQ {
2     ll p, q;
3     QQ go(QQ b, ll d) { return {p + b.p * d, q + b.q * d}; }
4 };
5 bool pred(QQ);
  // returns smallest p/q in [lo, hi] such that
6 // pred(p/q) is true, and 0 <= p, q <= N
7 QQ frac_bs(ll N) {
8     QQ lo{0, 1}, hi{1, 0};
9     if (pred(lo)) return lo;
10    assert(pred(hi));
11    bool dir = 1, L = 1, H = 1;
12    for (; L || H; dir = !dir) {
13        ll len = 0, step = 1;
14        for (int t = 0; t < 2 && (t ? step /= 2 : step *= 2);)
15            if (QQ mid = hi.go(lo, len + step);
16                mid.p > N || mid.q > N || dir ^ pred(mid))

```

```

19         t++;
20         else len += step;
21         swap(lo, hi = hi.go(lo, len));
22         (dir ? L : H) = !!len;
23     }
24     return dir ? hi : lo;
25 }

```

#### 4.1.13. Farey Sequence

```

1 // returns (e/f), where (a/b, c/d, e/f) are
  // three consecutive terms in the order n farey sequence
2 // to start, call next_farey(n, 0, 1, 1, n)
3 pll next_farey(ll n, ll a, ll b, ll c, ll d) {
4     ll p = (n + b) / d;
5     return pll(p * c - a, p * d - b);
6 }

```

### 4.2. Combinatorics

#### 4.2.1. Matroid Intersection

This template assumes 2 weighted matroids of the same type, and that removing an element is much more expensive than checking if one can be added. **Remember to change the implementation details.**

The ground set is  $0, 1, \dots, n-1$ , where element  $i$  has weight  $w[i]$ . For the unweighted version, remove weights and change BF/SPFA to BFS.

```

1 constexpr int N = 100;
2 constexpr int INF = 1e9;
3
4 struct Matroid { // represents an independent set
5     Matroid(bitset<N>); // initialize from an independent set
6     bool can_add(int); // if adding will break independence
7     Matroid remove(int); // removing from the set
8 };
9
10 auto matroid_intersection(int n, const vector<int> &w) {
11     bitset<N> S;
12     for (int sz = 1; sz <= n; sz++) {
13         Matroid M1(S), M2(S);
14
15         vector<vector<pii>> e(n + 2);
16         for (int j = 0; j < n; j++)
17             if (!S[j]) {
18                 if (M1.can_add(j)) e[n].emplace_back(j, -w[j]);
19                 if (M2.can_add(j)) e[j].emplace_back(n + 1, 0);
20             }
21         for (int i = 0; i < n; i++)
22             if (S[i]) {
23                 Matroid T1 = M1.remove(i), T2 = M2.remove(i);
24                 for (int j = 0; j < n; j++)
25                     if (!S[j]) {
26                         if (T1.can_add(j)) e[i].emplace_back(j, -w[j]);
27                         if (T2.can_add(j)) e[j].emplace_back(i, w[i]);
28                     }
29             }
30
31         vector<pii> dis(n + 2, {INF, 0});
32         vector<int> prev(n + 2, -1);
33         dis[n] = {0, 0};
34         // change to SPFA for more speed, if necessary
35         bool upd = 1;
36         while (upd) {
37             upd = 0;
38             for (int u = 0; u < n + 2; u++)
39                 for (auto [v, c] : e[u]) {
40                     pii x{dis[u].first + c, dis[u].second + 1};
41                     if (x < dis[v]) dis[v] = x, prev[v] = u, upd = 1;
42                 }
43         }
44
45         if (dis[n + 1].first < INF)
46             for (int x = prev[n + 1]; x != n; x = prev[x])
47                 S.flip(x);
48         else break;
49
50         // S is the max-weighted independent set with size sz
51     }
52     return S;
53 }

```

#### 4.2.2. De Bruijn Sequence

```

1 int res[kN], aux[kN], a[kN], sz;
2 void Rec(int t, int p, int n, int k) {
3     if (t > n) {
4         if (n % p == 0)
5             for (int i = 1; i <= p; ++i) res[sz++] = aux[i];
6         else {

```

```

7   aux[t] = aux[t - p];
9   Rec(t + 1, p, n, k);
11  for (aux[t] = aux[t - p] + 1; aux[t] < k; ++aux[t])
13  Rec(t + 1, t, n, k);
15  }
17  int DeBruijn(int k, int n) {
19  // return cyclic string of length k^n such that every
21  // string of length n using k character appears as a
23  // substring.
25  if (k == 1) return res[0] = 0, 1;
27  fill(aux, aux + k * n, 0);
29  return sz = 0, Rec(1, 1, n, k), sz;
31  }

```

### 4.2.3. Multinomial

```

1  // ways to permute v[i]
3  ll multinomial(vi &v) {
5  ll c = 1, m = v.empty() ? 1 : v[0];
7  for (int i = 1; i < v.size(); i++)
9  for (int j = 0; i < v[i]; j++) c = c * ++m / (j + 1);
11 return c;
13 }

```

## 4.3. Algebra

### 4.3.1. Formal Power Series

```

1  template <typename mint>
3  struct FormalPowerSeries : vector<mint> {
5  using vector<mint>::vector;
7  using FPS = FormalPowerSeries;
9  FPS &operator+=(const FPS &r) {
11 if (r.size() > this->size()) this->resize(r.size());
13 for (int i = 0; i < (int)r.size(); i++)
15 (*this)[i] += r[i];
17 return *this;
19 }
21 FPS &operator+=(const mint &r) {
23 if (this->empty()) this->resize(1);
25 (*this)[0] += r;
27 return *this;
29 }
31 FPS &operator-=(const FPS &r) {
33 if (r.size() > this->size()) this->resize(r.size());
35 for (int i = 0; i < (int)r.size(); i++)
37 (*this)[i] -= r[i];
39 return *this;
41 }
43 FPS &operator-=(const mint &r) {
45 if (this->empty()) this->resize(1);
47 (*this)[0] -= r;
49 return *this;
51 }
53 FPS &operator*=(const mint &v) {
55 for (int k = 0; k < (int)this->size(); k++)
57 (*this)[k] *= v;
59 return *this;
61 }
63 FPS &operator/=(const FPS &r) {
65 if (this->size() < r.size()) {
67 this->clear();
69 return *this;
71 }
73 int n = this->size() - r.size() + 1;
75 if ((int)r.size() <= 64) {
77 FPS f(*this), g(r);
79 g.shrink();
81 mint coeff = g.back().inverse();
83 for (auto &x : g) x *= coeff;
85 int deg = (int)f.size() - (int)g.size() + 1;
87 int gs = g.size();
89 FPS quo(deg);
91 for (int i = deg - 1; i >= 0; i--) {
93 quo[i] = f[i + gs - 1];
95 for (int j = 0; j < gs; j++)
97 f[i + j] -= quo[i] * g[j];
99 }
101 *this = quo * coeff;

```

```

61 this->resize(n, mint(0));
63 return *this;
65 }
67 return *this = ((*this).rev().pre(n) * r.rev().inv(n))
69 .pre(n)
71 .rev();
73 }
75 FPS operator*(const FPS &r) const {
77 return FPS(*this) += r;
79 }
81 FPS operator+(const mint &v) const {
83 return FPS(*this) += v;
85 }
87 FPS operator-(const FPS &r) const {
89 return FPS(*this) -= r;
91 }
93 FPS operator-(const mint &v) const {
95 return FPS(*this) -= v;
97 }
99 FPS operator*(const FPS &r) const {
101 return FPS(*this) *= r;
103 }
105 FPS operator*(const mint &v) const {
107 return FPS(*this) *= v;
109 }
111 FPS operator/(const FPS &r) const {
113 return FPS(*this) /= r;
115 }
117 FPS operator%(const FPS &r) const {
119 return FPS(*this) %= r;
121 }
123 FPS operator-() const {
125 FPS ret(this->size());
127 for (int i = 0; i < (int)this->size(); i++)
129 ret[i] = -(*this)[i];
131 return ret;
133 }
135 void shrink() {
137 while (this->size() && this->back() == mint(0))
139 this->pop_back();
141 }
143 FPS rev() const {
145 FPS ret(*this);
147 reverse(begin(ret), end(ret));
149 return ret;
151 }
153 FPS dot(FPS r) const {
155 FPS ret(min(this->size(), r.size()));
157 for (int i = 0; i < (int)ret.size(); i++)
159 ret[i] = (*this)[i] * r[i];
161 return ret;
163 }
165 FPS pre(int sz) const {
167 return FPS(begin(*this),
169 begin(*this) + min((int)this->size(), sz));
171 }
173 FPS operator>>(int sz) const {
175 if ((int)this->size() <= sz) return {};
177 FPS ret(*this);
179 ret.erase(ret.begin(), ret.begin() + sz);
181 return ret;
183 }
185 FPS operator<<(int sz) const {
187 FPS ret(*this);
189 ret.insert(ret.begin(), sz, mint(0));
191 return ret;
193 }
195 FPS diff() const {
197 const int n = (int)this->size();
199 FPS ret(max(0, n - 1));
201 mint one(1), coeff(1);
203 for (int i = 1; i < n; i++) {
205 ret[i - 1] = (*this)[i] * coeff;
207 coeff += one;
209 }
211 return ret;
213 }

```

```

153 FPS integral() const {
154     const int n = (int)this->size();
155     FPS ret(n + 1);
156     ret[0] = mint(0);
157     if (n > 0) ret[1] = mint(1);
158     auto mod = mint::get_mod();
159     for (int i = 2; i <= n; i++)
160         ret[i] = (-ret[mod % i]) * (mod / i);
161     for (int i = 0; i < n; i++) ret[i + 1] *= (*this)[i];
162     return ret;
163 }
164
165 mint eval(mint x) const {
166     mint r = 0, w = 1;
167     for (auto &v : *this) r += w * v, w *= x;
168     return r;
169 }
170
171 FPS log(int deg = -1) const {
172     assert((*this)[0] == mint(1));
173     if (deg == -1) deg = (int)this->size();
174     return (this->diff() * this->inv(deg))
175         .pre(deg - 1)
176         .integral();
177 }
178
179 FPS pow(int64_t k, int deg = -1) const {
180     const int n = (int)this->size();
181     if (deg == -1) deg = n;
182     for (int i = 0; i < n; i++) {
183         if ((*this)[i] != mint(0)) {
184             if (i * k > deg) return FPS(deg, mint(0));
185             mint rev = mint(1) / (*this)[i];
186             FPS ret =
187                 (((*this * rev) >> i).log(deg) * k).exp(deg) *
188                 ((*this)[i].pow(k));
189             ret = (ret << (i * k)).pre(deg);
190             if ((int)ret.size() < deg) ret.resize(deg, mint(0));
191             return ret;
192         }
193     }
194     return FPS(deg, mint(0));
195 }
196
197 static void *ntt_ptr;
198 static void set_fft();
199 FPS &operator*=(const FPS &r);
200 void ntt();
201 void intt();
202 void ntt_doubling();
203 static int ntt_pr();
204 FPS inv(int deg = -1) const;
205 FPS exp(int deg = -1) const;
206 };
207 template <typename mint>
208 void *FormalPowerSeries<mint>::ntt_ptr = nullptr;

```

## 4.4. Theorems

### 4.4.1. Kirchhoff's Theorem

Denote  $L$  be a  $n \times n$  matrix as the Laplacian matrix of graph  $G$ , where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where  $c$  is the number of edge  $(i, j)$  in  $G$ .

- The number of undirected spanning in  $G$  is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at  $r$  in  $G$  is  $|\det(\tilde{L}_{rr})|$ .

### 4.4.2. Tutte's Matrix

Let  $D$  be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if  $i < j$  and  $(i, j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{\text{rank}(D)}{2}$  is the maximum matching on  $G$ .

### 4.4.3. Cayley's Formula

- Given a degree sequence  $d_1, d_2, \dots, d_n$  for each *labeled* vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)! \cdots (d_n-1)!}$$

spanning trees.

- Let  $T_{n,k}$  be the number of *labeled* forests on  $n$  vertices with  $k$  components, such that vertex  $1, 2, \dots, k$  belong to different components. Then  $T_{n,k} = kn^{n-k-1}$ .

### 4.4.4. Erdős–Gallai Theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \dots \geq d_n$  can be represented as the degree sequence of a finite simple graph on  $n$  vertices if and only if  $d_1 + d_2 + \dots + d_n$  is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

holds for all  $1 \leq k \leq n$ .

### 4.4.5. Burnside's Lemma

Let  $X$  be a set and  $G$  be a group that acts on  $X$ . For  $g \in G$ , denote by  $X^g$  the elements fixed by  $g$ :

$$X^g = \{x \in X \mid gx \in X\}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

## 5. Numeric

### 5.1. Barrett Reduction

```

1 using ull = unsigned long long;
2 using ul = __uint128_t;
3 // very fast calculation of a % m
4 struct reduction {
5     const ull m, d;
6     explicit reduction(ull m) : m(m), d(((ul)1 << 64) / m) {}
7     inline ull operator()(ull a) const {
8         ull q = (ull)((ul)d * a) >> 64;
9         return (a - q * m) >= m ? a - m : a;
10    }
11 };

```

### 5.2. Long Long Multiplication

```

1 using ull = unsigned long long;
2 using ll = long long;
3 using ld = long double;
4 // returns a * b % M where a, b < M < 2**63
5 ull mult(ull a, ull b, ull M) {
6     ll ret = a * b - M * ull(ld(a) * ld(b) / ld(M));
7     return ret + M * (ret < 0) - M * (ret >= (ll)M);
8 }

```

### 5.3. Fast Fourier Transform

```

1 template <typename T>
2 void fft_(int n, vector<T> &a, vector<T> &rt, bool inv) {
3     vector<int> br(n);
4     for (int i = 1; i < n; i++) {
5         br[i] = (i & 1) ? br[i - 1] + n / 2 : br[i / 2] / 2;
6         if (br[i] > i) swap(a[i], a[br[i]]);
7     }
8     for (int len = 2; len <= n; len *= 2)
9         for (int i = 0; i < n; i += len)
10             for (int j = 0; j < len / 2; j++) {
11                 int pos = n / len * (inv ? len - j : j);
12                 T u = a[i + j], v = a[i + j + len / 2] * rt[pos];
13                 a[i + j] = u + v, a[i + j + len / 2] = u - v;
14             }
15     if (T minv = T(1) / T(n); inv)
16         for (T &x : a) x *= minv;
17 }

```

Requires: Mod Struct

```

1
2
3 void ntt(vector<Mod> &a, bool inv, Mod primitive_root) {
4     int n = a.size();
5     Mod root = primitive_root ^ (MOD - 1) / n;
6     vector<Mod> rt(n + 1, 1);
7     for (int i = 0; i < n; i++) rt[i + 1] = rt[i] * root;
8     fft_(n, a, rt, inv);
9 }
10
11 void fft(vector<complex<double>> &a, bool inv) {
12     int n = a.size();
13     vector<complex<double>> rt(n + 1);
14     double arg = acos(-1) * 2 / n;
15     for (int i = 0; i < n; i++)
16         rt[i] = {cos(arg * i), sin(arg * i)};
17     fft_(n, a, rt, inv);
18 }

```

### 5.4. Fast Walsh-Hadamard Transform

Requires: Mod Struct

```

1
2
3 void fwht(vector<Mod> &a, bool inv) {
4     int n = a.size();
5     for (int d = 1; d < n; d <= 1)
6         for (int m = 0; m < n; m++)
7             if (!(m & d)) {
8                 inv ? a[m] -= a[m | d] : a[m] += a[m | d]; // AND
9                 inv ? a[m | d] -= a[m] : a[m | d] += a[m]; // OR
10                 Mod x = a[m], y = a[m | d]; // XOR
11                 a[m] = x + y, a[m | d] = x - y; // XOR
12             }
13     if (Mod iv = Mod(1) / n; inv) // XOR
14         for (Mod &i : a) i *= iv; // XOR
15 }

```

### 5.5. Subset Convolution

Requires: Mod Struct

```

1 #pragma GCC target("popcnt")
2 #include <immintrin.h>
3
4 void fwht(int n, vector<vector<Mod>> &a, bool inv) {
5     for (int h = 0; h < n; h++)
6         for (int i = 0; i < (1 << n); i++)
7             if (!(i & (1 << h)))
8                 for (int k = 0; k <= n; k++)
9                     inv ? a[i | (1 << h)][k] -= a[i][k]
10                        : a[i | (1 << h)][k] += a[i][k];
11 }
12 // c[k] = sum(popcnt(i & j) == sz && i | j == k) a[i] * b[j]
13 vector<Mod> subset_convolution(int n, int sz,
14                               const vector<Mod> &a,
15                               const vector<Mod> &b) {
16     int len = n + sz + 1, N = 1 << n;
17     vector<vector<Mod>> a(1 << n, vector<Mod>(len, 0)), b = a;
18     for (int i = 0; i < N; i++)
19         a[i][_mm_popcnt_u64(i)] = a[i];
20     b[i][_mm_popcnt_u64(i)] = b[i];
21     fwht(n, a, 0), fwht(n, b, 0);
22     for (int i = 0; i < N; i++) {
23         vector<Mod> tmp(len);
24         for (int j = 0; j < len; j++)
25             for (int k = 0; k <= j; k++)
26                 tmp[j] += a[i][k] * b[i][j - k];
27         a[i] = tmp;
28     }
29     fwht(n, a, 1);
30     vector<Mod> c(N);
31     for (int i = 0; i < N; i++)
32         c[i] = a[i][_mm_popcnt_u64(i) + sz];
33     return c;
34 }

```

### 5.6. Linear Recurrences

#### 5.6.1. Berlekamp-Massey Algorithm

```

1 template <typename T>
2 vector<T> berlekamp_massey(const vector<T> &s) {
3     int n = s.size(), l = 0, m = 1;
4     vector<T> r(n), p(n);
5     r[0] = p[0] = 1;
6     T b = 1, d = 0;
7     for (int i = 0; i < n; i++, m++, d = 0) {
8         for (int j = 0; j <= l; j++) d += r[j] * s[i - j];
9         if ((d /= b) == 0) continue; // change if T is float
10        auto t = r;
11        for (int j = m; j < n; j++) r[j] -= d * p[j - m];
12        if (l * 2 <= i) l = i + 1 - l, b = d, m = 0, p = t;
13    }
14    return r.resize(l + 1), reverse(r.begin(), r.end()), r;
15 }

```

#### 5.6.2. Linear Recurrence Calculation

```

1 template <typename T> struct lin_rec {
2     using poly = vector<T>;
3     poly mul(poly a, poly b, poly m) {
4         int n = m.size();
5         poly r(n);
6         for (int i = n - 1; i >= 0; i--) {
7             r.insert(r.begin(), 0), r.pop_back();
8             T c = r[n - 1] + a[n - 1] * b[i];
9             // c /= m[n - 1]; if m is not monic
10            for (int j = 0; j < n; j++)
11                r[j] += a[j] * b[i] - c * m[j];
12        }
13        return r;
14    }
15    poly pow(poly p, ll k, poly m) {
16        poly r(m.size());
17        r[0] = 1;
18        for (; k >= 1; p = mul(p, p, m))
19            if (k & 1) r = mul(r, p, m);
20        return r;
21    }
22    T calc(poly t, poly r, ll k) {
23        int n = r.size();
24        poly p(n);
25        p[1] = 1;
26        poly q = pow(p, k, r);
27        T ans = 0;
28        for (int i = 0; i < n; i++) ans += t[i] * q[i];
29        return ans;
30    }
31 };

```

## 5.7. Matrices

### 5.7.1. Determinant

Requires: Mod Struct

```

1
3 Mod det(vector<vector<Mod>> a) {
4     int n = a.size();
5     Mod ans = 1;
6     for (int i = 0; i < n; i++) {
7         int b = i;
8         for (int j = i + 1; j < n; j++)
9             if (a[j][i] != 0) {
10                b = j;
11                break;
12            }
13        if (i != b) swap(a[i], a[b]), ans = -ans;
14        ans *= a[i][i];
15        if (ans == 0) return 0;
16        for (int j = i + 1; j < n; j++) {
17            Mod v = a[j][i] / a[i][i];
18            if (v != 0)
19                for (int k = i + 1; k < n; k++)
20                    a[j][k] -= v * a[i][k];
21        }
22    }
23    return ans;
}

```

```

1 double det(vector<vector<double>> a) {
2     int n = a.size();
3     double ans = 1;
4     for (int i = 0; i < n; i++) {
5         int b = i;
6         for (int j = i + 1; j < n; j++)
7             if (fabs(a[j][i]) > fabs(a[b][i])) b = j;
8         if (i != b) swap(a[i], a[b]), ans = -ans;
9         ans *= a[i][i];
10        if (ans == 0) return 0;
11        for (int j = i + 1; j < n; j++) {
12            double v = a[j][i] / a[i][i];
13            if (v != 0)
14                for (int k = i + 1; k < n; k++)
15                    a[j][k] -= v * a[i][k];
16        }
17    }
18    return ans;
19 }

```

### 5.7.2. Inverse

```

1
3 // Returns rank.
4 // Result is stored in A unless singular (rank < n).
5 // For prime powers, repeatedly set
6 // A^{-1} = A^{-1} (2I - A*A^{-1}) (mod p^k)
7 // where A^{-1} starts as the inverse of A mod p,
8 // and k is doubled in each step.
9
10 int matInv(vector<vector<double>> &A) {
11     int n = sz(A);
12     vi col(n);
13     vector<vector<double>> tmp(n, vector<double>(n));
14     rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
15
16     rep(i, 0, n) {
17         int r = i, c = i;
18         rep(j, i, n)
19             rep(k, i, n) if (fabs(A[j][k]) > fabs(A[r][c])) r = j, c = k;
20
21         if (fabs(A[r][c]) < 1e-12) return i;
22         A[i].swap(A[r]);
23         tmp[i].swap(tmp[r]);
24         rep(j, 0, n) swap(A[j][i], A[j][c]),
25             swap(tmp[j][i], tmp[j][c]);
26         swap(col[i], col[c]);
27         double v = A[i][i];
28         rep(j, i + 1, n) {
29             double f = A[j][i] / v;
30             A[j][i] = 0;
31             rep(k, i + 1, n) A[j][k] -= f * A[i][k];
32             rep(k, 0, n) tmp[j][k] -= f * tmp[i][k];
33         }
34         rep(j, i + 1, n) A[i][j] /= v;
35         rep(j, 0, n) tmp[i][j] /= v;
36         A[i][i] = 1;
37     }
}

```

```

39
40 for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
41     double v = A[j][i];
42     rep(k, 0, n) tmp[j][k] -= v * tmp[i][k];
43 }
44
45 rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] = tmp[i][j];
46 return n;
47 }
48
49 int matInv_mod(vector<vector<ll>> &A) {
50     int n = sz(A);
51     vi col(n);
52     vector<vector<ll>> tmp(n, vector<ll>(n));
53     rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
54
55     rep(i, 0, n) {
56         int r = i, c = i;
57         rep(j, i, n) rep(k, i, n) if (A[j][k]) {
58             r = j;
59             c = k;
60             goto found;
61         }
62         return i;
63     found:
64         A[i].swap(A[r]);
65         tmp[i].swap(tmp[r]);
66         rep(j, 0, n) swap(A[j][i], A[j][c]),
67             swap(tmp[j][i], tmp[j][c]);
68         swap(col[i], col[c]);
69         ll v = modpow(A[i][i], mod - 2);
70         rep(j, i + 1, n) {
71             ll f = A[j][i] * v % mod;
72             A[j][i] = 0;
73             rep(k, i + 1, n) A[j][k] =
74                 (A[j][k] - f * A[i][k]) % mod;
75             rep(k, 0, n) tmp[j][k] =
76                 (tmp[j][k] - f * tmp[i][k]) % mod;
77         }
78         rep(j, i + 1, n) A[i][j] = A[i][j] * v % mod;
79         rep(j, 0, n) tmp[i][j] = tmp[i][j] * v % mod;
80         A[i][i] = 1;
81     }
82
83     for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
84         ll v = A[j][i];
85         rep(k, 0, n) tmp[j][k] =
86             (tmp[j][k] - v * tmp[i][k]) % mod;
87     }
88
89     rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] =
90         tmp[i][j] % mod + (tmp[i][j] < 0 ? mod : 0);
91     return n;
92 }

```

### 5.7.3. Characteristic Polynomial

```

1
3 // calculate det(a - xI)
4 template <typename T>
5 vector<T> CharacteristicPolynomial(vector<vector<T>> a) {
6     int N = a.size();
7
8     for (int j = 0; j < N - 2; j++) {
9         for (int i = j + 1; i < N; i++) {
10             if (a[i][j] != 0) {
11                 swap(a[j + 1], a[i]);
12                 for (int k = 0; k < N; k++)
13                     swap(a[k][j + 1], a[k][i]);
14                 break;
15             }
16         }
17         if (a[j + 1][j] != 0) {
18             T inv = T(1) / a[j + 1][j];
19             for (int i = j + 2; i < N; i++) {
20                 if (a[i][j] == 0) continue;
21                 T coe = inv * a[i][j];
22                 for (int l = j; l < N; l++)
23                     a[i][l] -= coe * a[j + 1][l];
24                 for (int k = 0; k < N; k++)
25                     a[k][j + 1] += coe * a[k][i];
26             }
27         }
28     }
29 }
30
31 vector<vector<T>> p(N + 1);
32 p[0] = {T(1)};
33 for (int i = 1; i <= N; i++) {
34     p[i].resize(i + 1);
35 }

```

```

35     for (int j = 0; j < i; j++) {
36         p[i][j + 1] -= p[i - 1][j];
37         p[i][j] += p[i - 1][j] * a[i - 1][i - 1];
38     }
39     T x = 1;
40     for (int m = 1; m < i; m++) {
41         x *= -a[i - m][i - m - 1];
42         T coe = x * a[i - m - 1][i - 1];
43         for (int j = 0; j < i - m; j++)
44             p[i][j] += coe * p[i - m - 1][j];
45     }
46 }
47 return p[N];
48 }

```

#### 5.7.4. Solve Linear Equation

```

1 // solves for x: A * x = b
2
3 typedef vector<double> vd;
4 const double eps = 1e-12;
5
6 // solves for x: A * x = b
7 int solveLinear(vector<vd> &A, vd &b, vd &x) {
8     int n = sz(A), m = sz(x), rank = 0, br, bc;
9     if (n) assert(sz(A[0]) == m);
10    vi col(m);
11    iota(all(col), 0);
12
13    rep(i, 0, n) {
14        double v, bv = 0;
15        rep(r, i, n) rep(c, i, m) if ((v = fabs(A[r][c])) > bv)
16            br = r, bc = c, bv = v;
17        if (bv <= eps) {
18            rep(j, i, n) if (fabs(b[j]) > eps) return -1;
19            break;
20        }
21        swap(A[i], A[br]);
22        swap(b[i], b[br]);
23        swap(col[i], col[bc]);
24        rep(j, 0, n) swap(A[j][i], A[j][bc]);
25        bv = 1 / A[i][i];
26        rep(j, i + 1, n) {
27            double fac = A[j][i] * bv;
28            b[j] -= fac * b[i];
29            rep(k, i + 1, m) A[j][k] -= fac * A[i][k];
30        }
31        rank++;
32    }
33
34    x.assign(m, 0);
35    for (int i = rank; i--;) {
36        b[i] /= A[i][i];
37        x[col[i]] = b[i];
38        rep(j, 0, i) b[j] -= A[j][i] * b[i];
39    }
40    return rank; // (multiple solutions if rank < m)
41 }

```

#### 5.8. Polynomial Interpolation

```

1 // returns a, such that a[0]x^0 + a[1]x^1 + a[2]x^2 + ...
2 // passes through the given points
3 typedef vector<double> vd;
4 vd interpolate(vd x, vd y, int n) {
5     vd res(n), temp(n);
6     rep(k, 0, n - 1) rep(i, k + 1, n) y[i] =
7         (y[i] - y[k]) / (x[i] - x[k]);
8     double last = 0;
9     temp[0] = 1;
10    rep(k, 0, n) rep(i, 0, n) {
11        res[i] += y[k] * temp[i];
12        swap(last, temp[i]);
13        temp[i] -= last * x[k];
14    }
15    return res;
16 }

```

#### 5.9. Simplex Algorithm

```

1 // Two-phase simplex algorithm for solving linear programs
2 // of the form
3 //
4 //      maximize    c^T x
5 //      subject to  Ax <= b
6 //                  x >= 0
7 //
8 // INPUT: A -- an m x n matrix

```

```

9 //      b -- an m-dimensional vector
10 //      c -- an n-dimensional vector
11 //      x -- a vector where the optimal solution will be
12 //           stored
13 //
14 // OUTPUT: value of the optimal solution (infinity if
15 // unbounded
16 //         above, nan if infeasible)
17 //
18 // To use this code, create an LPSolver object with A, b,
19 // and c as arguments. Then, call Solve(x).
20
21 typedef long double ld;
22 typedef vector<ld> vd;
23 typedef vector<vd> vvd;
24 typedef vector<int> vi;
25
26 const ld EPS = 1e-9;
27
28 struct LPSolver {
29     int m, n;
30     vi B, N;
31     vvd D;
32
33     LPSolver(const vvd &A, const vd &b, const vd &c)
34         : m(b.size()), n(c.size()), N(n + 1), B(m),
35           D(m + 2, vd(n + 2)) {
36         for (int i = 0; i < m; i++)
37             for (int j = 0; j < n; j++) D[i][j] = A[i][j];
38         for (int i = 0; i < m; i++) {
39             B[i] = n + i;
40             D[i][n] = -1;
41             D[i][n + 1] = b[i];
42         }
43         for (int j = 0; j < n; j++) {
44             N[j] = j;
45             D[m][j] = -c[j];
46         }
47         N[n] = -1;
48         D[m + 1][n] = 1;
49     }
50
51     void Pivot(int r, int s) {
52         double inv = 1.0 / D[r][s];
53         for (int i = 0; i < m + 2; i++)
54             if (i != r)
55                 for (int j = 0; j < n + 2; j++)
56                     if (j != s) D[i][j] -= D[r][j] * D[i][s] * inv;
57         for (int j = 0; j < n + 2; j++)
58             if (j != s) D[r][j] *= inv;
59         for (int i = 0; i < m + 2; i++)
60             if (i != r) D[i][s] *= -inv;
61         D[r][s] = inv;
62         swap(B[r], N[s]);
63     }
64
65     bool Simplex(int phase) {
66         int x = phase == 1 ? m + 1 : m;
67         while (true) {
68             int s = -1;
69             for (int j = 0; j <= n; j++) {
70                 if (phase == 2 && N[j] == -1) continue;
71                 if (s == -1 || D[x][j] < D[x][s] ||
72                     D[x][j] == D[x][s] && N[j] < N[s])
73                     s = j;
74             }
75             if (D[x][s] > -EPS) return true;
76             int r = -1;
77             for (int i = 0; i < m; i++) {
78                 if (D[i][s] < EPS) continue;
79                 if (r == -1 ||
80                     D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
81                     (D[i][n + 1] / D[i][s]) ==
82                     (D[r][n + 1] / D[r][s]) &&
83                     B[i] < B[r])
84                     r = i;
85             }
86             if (r == -1) return false;
87             Pivot(r, s);
88         }
89     }
90
91     ld Solve(vd &x) {
92         int r = 0;
93         for (int i = 1; i < m; i++)
94             if (D[i][n + 1] < D[r][n + 1]) r = i;
95         if (D[r][n + 1] < -EPS) {
96             Pivot(r, n);
97             if (!Simplex(1) || D[m + 1][n + 1] < -EPS)
98                 return -numeric_limits<ld>::infinity();
99             for (int i = 0; i < m; i++)

```



```

101     if (B[i] == -1) {
102         int s = -1;
103         for (int j = 0; j <= n; j++)
104             if (s == -1 || D[i][j] < D[i][s] ||
105                 D[i][j] == D[i][s] && N[j] < N[s])
106                 s = j;
107         Pivot(i, s);
108     }
109     if (!Simplex(2)) return numeric_limits<ld>::infinity();
110     x = vd(n);
111     for (int i = 0; i < m; i++)
112         if (B[i] < n) x[B[i]] = D[i][n + 1];
113     return D[m][n + 1];
114 }
115 };
116
117 int main() {
118
119     const int m = 4;
120     const int n = 3;
121     ld _A[m][n] = {
122         {6, -1, 0}, {-1, -5, 0}, {1, 5, 1}, {-1, -5, -1}};
123     ld _b[m] = {10, -4, 5, -5};
124     ld _c[n] = {1, -1, 0};
125
126     vvd A(m);
127     vd b(_b, _b + m);
128     vd c(_c, _c + n);
129     for (int i = 0; i < m; i++) A[i] = vd(_A[i], _A[i] + n);
130
131     LPSolver solver(A, b, c);
132     vd x;
133     ld value = solver.Solve(x);
134
135     cerr << "VALUE: " << value << endl; // VALUE: 1.29032
136     cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
137     for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
138     cerr << endl;
139     return 0;
140 }

```



## 6. Geometry

### 6.1. Point

```

1 template <typename T> struct P {
2     T x, y;
3     P(T x = 0, T y = 0) : x(x), y(y) {}
4     bool operator<(const P &p) const {
5         return tie(x, y) < tie(p.x, p.y);
6     }
7     bool operator==(const P &p) const {
8         return tie(x, y) == tie(p.x, p.y);
9     }
10    P operator-() const { return {-x, -y}; }
11    P operator+(P p) const { return {x + p.x, y + p.y}; }
12    P operator-(P p) const { return {x - p.x, y - p.y}; }
13    P operator*(T d) const { return {x * d, y * d}; }
14    P operator/(T d) const { return {x / d, y / d}; }
15    T dist2() const { return x * x + y * y; }
16    double len() const { return sqrt(dist2()); }
17    P unit() const { return *this / len(); }
18    friend T dot(P a, P b) { return a.x * b.x + a.y * b.y; }
19    friend T cross(P a, P b) { return a.x * b.y - a.y * b.x; }
20    friend T cross(P a, P b, P o) {
21        return cross(a - o, b - o);
22    }
23 };
using pt = P<ll>;

```

#### 6.1.1. Quaternion

```

1 constexpr double PI = 3.141592653589793;
2 constexpr double EPS = 1e-7;
3 struct Q {
4     using T = double;
5     T x, y, z, r;
6     Q(T r = 0) : x(0), y(0), z(0), r(r) {}
7     Q(T x, T y, T z, T r = 0) : x(x), y(y), z(z), r(r) {}
8     friend bool operator==(const Q &a, const Q &b) {
9         return (a - b).abs2() <= EPS;
10    }
11    friend bool operator!=(const Q &a, const Q &b) {
12        return !(a == b);
13    }
14    Q operator-() { return Q(-x, -y, -z, -r); }
15    Q operator+(const Q &b) const {
16        return Q(x + b.x, y + b.y, z + b.z, r + b.r);
17    }
18    Q operator-(const Q &b) const {
19        return Q(x - b.x, y - b.y, z - b.z, r - b.r);
20    }
21    Q operator*(const T &t) const {
22        return Q(x * t, y * t, z * t, r * t);
23    }
24    Q operator*(const Q &b) const {
25        return Q(r * b.x + x * b.r + y * b.z - z * b.y,
26                r * b.y - x * b.z + y * b.r + z * b.x,
27                r * b.z + x * b.y - y * b.x + z * b.r,
28                r * b.r - x * b.x - y * b.y - z * b.z);
29    }
30    Q operator/(const Q &b) const { return *this * b.inv(); }
31    T abs2() const { return r * r + x * x + y * y + z * z; }
32    T len() const { return sqrt(abs2()); }
33    Q conj() const { return Q(-x, -y, -z, r); }
34    Q unit() const { return *this * (1.0 / len()); }
35    Q inv() const { return conj() * (1.0 / abs2()); }
36    friend T dot(Q a, Q b) {
37        return a.x * b.x + a.y * b.y + a.z * b.z;
38    }
39    friend Q cross(Q a, Q b) {
40        return Q(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z,
41                a.x * b.y - a.y * b.x);
42    }
43    friend Q rotation_around(Q axis, T angle) {
44        return axis.unit() * sin(angle / 2) + cos(angle / 2);
45    }
46    Q rotated_around(Q axis, T angle) {
47        Q u = rotation_around(axis, angle);
48        return u * *this / u;
49    }
50    friend Q rotation_between(Q a, Q b) {
51        a = a.unit(), b = b.unit();
52        if (a == -b) {
53            // degenerate case
54            Q ortho = abs(a.y) > EPS ? cross(a, Q(1, 0, 0))
55                : cross(a, Q(0, 1, 0));
56            return rotation_around(ortho, PI);
57        }
58        return (a * (a + b)).conj();
59    }
60 };

```

### 6.1.2. Spherical Coordinates

```

1 struct car_p {
2     double x, y, z;
3 };
4 struct sph_p {
5     double r, theta, phi;
6 };
7 sph_p conv(car_p p) {
8     double r = sqrt(p.x * p.x + p.y * p.y + p.z * p.z);
9     double theta = asin(p.y / r);
10    double phi = atan2(p.y, p.x);
11    return {r, theta, phi};
12 }
13 car_p conv(sph_p p) {
14     double x = p.r * cos(p.theta) * sin(p.phi);
15     double y = p.r * cos(p.theta) * cos(p.phi);
16     double z = p.r * sin(p.theta);
17     return {x, y, z};
18 }

```

### 6.2. Segments

```

1 // for non-collinear ABCD, if segments AB and CD intersect
2 bool intersects(pt a, pt b, pt c, pt d) {
3     if (cross(b, c, a) * cross(b, d, a) > 0) return false;
4     if (cross(d, a, c) * cross(d, b, c) > 0) return false;
5     return true;
6 }
7 // the intersection point of lines AB and CD
8 pt intersect(pt a, pt b, pt c, pt d) {
9     auto x = cross(b, c, a), y = cross(b, d, a);
10    if (x == y) {
11        // if(abs(x, y) < 1e-8) {
12        // is parallel
13    } else {
14        return d * (x / (x - y)) - c * (y / (x - y));
15    }
16 }

```

### 6.3. Convex Hull

```

1 // returns a convex hull in counterclockwise order
2 // for a non-strict one, change cross >= to >
3 vector<pt> convex_hull(vector<pt> p) {
4     sort(ALL(p));
5     if (p[0] == p.back()) return {p[0]};
6     int n = p.size(), t = 0;
7     vector<pt> h(n + 1);
8     for (int _ = 2, s = 0; _--; s = --t, reverse(ALL(p)))
9         for (pt i : p) {
10             while (t > s + 1 && cross(i, h[t - 1], h[t - 2]) >= 0)
11                 t--;
12             h[t++] = i;
13         }
14     return h.resize(t), h;
15 }

```

#### 6.3.1. 3D Hull

```

1
2 typedef Point3D<double> P3;
3
4 struct PR {
5     void ins(int x) { (a == -1 ? a : b) = x; }
6     void rem(int x) { (a == x ? a : b) = -1; }
7     int cnt() { return (a != -1) + (b != -1); }
8     int a, b;
9 };
10
11 struct F {
12     P3 q;
13     int a, b, c;
14 };
15
16 vector<F> hull3d(const vector<P3> &A) {
17     assert(sz(A) >= 4);
18     vector<vector<PR>> E(sz(A), vector<PR>(sz(A), {-1, -1}));
19     #define E(x, y) E[f.x][f.y]
20     vector<F> FS;
21     auto mf = [&](int i, int j, int k, int l) {
22         P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
23         if (q.dot(A[l]) > q.dot(A[i])) q = q * -1;
24         F f{q, i, j, k};
25         E(a, b).ins(k);
26         E(a, c).ins(j);
27         E(b, c).ins(i);
28         FS.push_back(f);
29     };

```

```

31 rep(i, 0, 4) rep(j, i + 1, 4) rep(k, j + 1, 4)
mf(i, j, k, 6 - i - j - k);
33
rep(i, 4, sz(A)) {
35     rep(j, 0, sz(FS)) {
        F f = FS[j];
37         if (f.q.dot(A[i]) > f.q.dot(A[f.a])) {
            E(a, b).rem(f.c);
39             E(a, c).rem(f.b);
            E(b, c).rem(f.a);
41             swap(FS[j--], FS.back());
            FS.pop_back();
43         }
    }
45     int nw = sz(FS);
    rep(j, 0, nw) {
47         F f = FS[j];
        #define C(a, b, c)
49         if (E(a, b).cnt() != 2) mf(f.a, f.b, i, f.c);
        C(a, b, c);
51         C(a, c, b);
        C(b, c, a);
53     }
    for (F &it : FS)
55         if ((A[it.b] - A[it.a])
            .cross(A[it.c] - A[it.a])
            .dot(it.q) <= 0)
57             swap(it.c, it.b);
59     return FS;
61 };

```

#### 6.4. Angular Sort

```

1 auto angle_cmp = [](const pt &a, const pt &b) {
    auto btm = [](const pt &a) {
3         return a.y < 0 || (a.y == 0 && a.x < 0);
    };
5     return make_tuple(btm(a), a.y * b.x, abs2(a)) <
        make_tuple(btm(b), a.x * b.y, abs2(b));
7 };
void angular_sort(vector<pt> &p) {
9     sort(p.begin(), p.end(), angle_cmp);
}

```

#### 6.5. Convex Polygon Minkowski Sum

```

1 // O(n) convex polygon minkowski sum
// must be sorted and counterclockwise
3 vector<pt> minkowski_sum(vector<pt> p, vector<pt> q) {
    auto diff = [](vector<pt> &c) {
5         auto rcmp = [](pt a, pt b) {
            return pt{a.y, a.x} < pt{b.y, b.x};
7         };
        rotate(c.begin(), min_element(ALL(c), rcmp), c.end());
9         c.push_back(c[0]);
        vector<pt> ret;
11         for (int i = 1; i < c.size(); i++)
            ret.push_back(c[i] - c[i - 1]);
13         return ret;
    };
15     auto dp = diff(p), dq = diff(q);
    pt cur = p[0] + q[0];
17     vector<pt> d(dp.size() + dq.size(), ret = {cur};
    // include angle_cmp from angular-sort.cpp
19     merge(ALL(dp), ALL(dq), d.begin(), angle_cmp);
    // optional: make ret strictly convex (UB if degenerate)
21     int now = 0;
    for (int i = 1; i < d.size(); i++) {
23         if (cross(d[i], d[now]) == 0) d[now] = d[now] + d[i];
        else d[++now] = d[i];
25     }
    d.resize(now + 1);
27     // end optional part
    for (pt v : d) ret.push_back(cur = cur + v);
29     return ret.pop_back(), ret;
}

```

#### 6.6. Point In Polygon

```

1 bool on_segment(pt a, pt b, pt p) {
    return cross(a, b, p) == 0 && dot((p - a), (p - b)) <= 0;
3 }
// p can be any polygon, but this is O(n)
5 bool inside(const vector<pt> &p, pt a) {
    int cnt = 0, n = p.size();
7     for (int i = 0; i < n; i++) {
        pt l = p[i], r = p[(i + 1) % n];
9         // change to return 0; for strict version
        if (on_segment(l, r, a)) return 1;
11         cnt ^= ((a.y < l.y) - (a.y < r.y)) * cross(l, r, a) > 0;
    }
}

```

```

13     return cnt;
}

```

##### 6.6.1. Convex Version

```

1 // no preprocessing version
// p must be a strict convex hull, counterclockwise
3 // if point is inside or on border
bool is_inside(const vector<pt> &c, pt p) {
5     int n = c.size(), l = 1, r = n - 1;
    if (cross(c[0], c[1], p) < 0) return false;
7     if (cross(c[n - 1], c[0], p) < 0) return false;
    while (l < r - 1) {
9         int m = (l + r) / 2;
        T a = cross(c[0], c[m], p);
11         if (a > 0) l = m;
        else if (a < 0) r = m;
13         else return dot(c[0] - p, c[m] - p) <= 0;
    }
15     if (l == r) return dot(c[0] - p, c[l] - p) <= 0;
    else return cross(c[l], c[r], p) >= 0;
17 }

19 // with preprocessing version
vector<pt> vecs;
21 pt center;
// p must be a strict convex hull, counterclockwise
23 // BEWARE OF OVERFLOWS!!
void preprocess(vector<pt> p) {
25     for (auto &v : p) v = v * 3;
    center = p[0] + p[1] + p[2];
27     center.x /= 3, center.y /= 3;
    for (auto &v : p) v = v - center;
29     vecs = (angular_sort(p), p);
}

31 bool intersect_strict(pt a, pt b, pt c, pt d) {
    if (cross(b, c, a) * cross(b, d, a) > 0) return false;
33     if (cross(d, a, c) * cross(d, b, c) >= 0) return false;
    return true;
35 }
// if point is inside or on border
37 bool query(pt p) {
    p = p * 3 - center;
39     auto pr = upper_bound(ALL(vecs), p, angle_cmp);
    if (pr == vecs.end()) pr = vecs.begin();
41     auto pl = (pr == vecs.begin()) ? vecs.back() : *(pr - 1);
    return !intersect_strict({0, 0}, p, pl, *pr);
43 }

```

##### 6.6.2. Offline Multiple Points Version

Requires: Point, GNU PBDS

```

1
3
5
using Double = __float128;
7 using Point = pt<Double, Double>;

9 int n, m;
vector<Point> poly;
11 vector<Point> query;
vector<int> ans;

13 struct Segment {
15     Point a, b;
    int id;
17 };
vector<Segment> segs;

19 Double Xnow;
21 inline Double get_y(const Segment &u, Double xnow = Xnow) {
    const Point &a = u.a;
23     const Point &b = u.b;
    return (a.y * (b.x - xnow) + b.y * (xnow - a.x)) /
25         (b.x - a.x);
}

27 bool operator<(Segment u, Segment v) {
    Double yu = get_y(u);
29     Double yv = get_y(v);
    if (yu != yv) return yu < yv;
31     return u.id < v.id;
}

33 ordered_map<Segment> st;

35 struct Event {
    int type; // +1 insert seg, -1 remove seg, 0 query
37     Double x, y;
}

```

```

    int id;
};
bool operator<(Event a, Event b) {
    if (a.x != b.x) return a.x < b.x;
    if (a.type != b.type) return a.type < b.type;
    return a.y < b.y;
}
vector<Event> events;

void solve() {
    set<Double> xs;
    set<Point> ps;
    for (int i = 0; i < n; i++) {
        xs.insert(poly[i].x);
        ps.insert(poly[i]);
    }
    for (int i = 0; i < n; i++) {
        Segment s{poly[i], poly[(i + 1) % n], i};
        if (s.a.x > s.b.x ||
            (s.a.x == s.b.x && s.a.y > s.b.y)) {
            swap(s.a, s.b);
        }
        segs.push_back(s);

        if (s.a.x != s.b.x) {
            events.push_back({+1, s.a.x + 0.2, s.a.y, i});
            events.push_back({-1, s.b.x - 0.2, s.b.y, i});
        }
    }
    for (int i = 0; i < m; i++) {
        events.push_back({0, query[i].x, query[i].y, i});
    }
    sort(events.begin(), events.end());
    int cnt = 0;
    for (Event e : events) {
        int i = e.id;
        Xnow = e.x;
        if (e.type == 0) {
            Double x = e.x;
            Double y = e.y;
            Segment tmp = {{x - 1, y}, {x + 1, y}, -1};
            auto it = st.lower_bound(tmp);

            if (ps.count(query[i]) > 0) {
                ans[i] = 0;
            } else if (xs.count(x) > 0) {
                ans[i] = -2;
            } else if (it != st.end() &&
                get_y(*it) == get_y(tmp)) {
                ans[i] = 0;
            } else if (it != st.begin() &&
                get_y(*prev(it)) == get_y(tmp)) {
                ans[i] = 0;
            } else {
                int rk = st.order_of_key(tmp);
                if (rk % 2 == 1) {
                    ans[i] = 1;
                } else {
                    ans[i] = -1;
                }
            }
        } else if (e.type == 1) {
            st.insert(segs[i]);
            assert((int)st.size() == ++cnt);
        } else if (e.type == -1) {
            st.erase(segs[i]);
            assert((int)st.size() == --cnt);
        }
    }
}

```

## 6.7. Closest Pair

```

vector<pll> p; // sort by x first!
bool cmpy(const pll &a, const pll &b) const {
    return a.y < b.y;
}
ll sq(ll x) { return x * x; }
// returns (minimum dist)^2 in [l, r]
ll solve(int l, int r) {
    if (r - l <= 1) return 1e18;
    int m = (l + r) / 2;
    ll mid = p[m].x, d = min(solve(l, m), solve(m, r));
    auto pb = p.begin();
    inplace_merge(pb + l, pb + m, pb + r, cmpy);
    vector<pll> s;
    for (int i = l; i < r; i++)
        if (sq(p[i].x - mid) < d) s.push_back(p[i]);
    for (int i = 0; i < s.size(); i++)
        for (int j = i + 1;
            j < s.size() && sq(s[j].y - s[i].y) < d; j++)

```

```

    d = min(d, dis(s[i], s[j]));
    return d;
}

```

## 6.8. Minimum Enclosing Circle

```

1
2
3 typedef Point<double> P;
4 double ccRadius(const P &A, const P &B, const P &C) {
5     return (B - A).dist() * (C - B).dist() * (A - C).dist() /
6         abs((B - A).cross(C - A)) / 2;
7 }
8 P ccCenter(const P &A, const P &B, const P &C) {
9     P b = C - A, c = B - A;
10    return A + (b * c.dist2() - c * b.dist2()).perp() /
11        b.cross(c) / 2;
12 }
13 pair<P, double> mec(vector<P> ps) {
14    shuffle(all(ps), mt19937(time(0)));
15    P o = ps[0];
16    double r = 0, EPS = 1 + 1e-8;
17    rep(i, 0, sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
18        o = ps[i], r = 0;
19        rep(j, 0, i) if ((o - ps[j]).dist() > r * EPS) {
20            o = (ps[i] + ps[j]) / 2;
21            r = (o - ps[i]).dist();
22            rep(k, 0, j) if ((o - ps[k]).dist() > r * EPS) {
23                o = ccCenter(ps[i], ps[j], ps[k]);
24                r = (o - ps[i]).dist();
25            }
26        }
27    }
28    return {o, r};
29 }

```

## 6.9. Delaunay Triangulation

```

1
2
3 typedef Point<ll> P;
4 typedef struct Quad *Q;
5 typedef __int128_t lll; // (can be ll if coords are < 2e4)
6 P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
7
8 struct Quad {
9     bool mark;
10    Q o, rot;
11    P p;
12    P F() { return r()->p; }
13    Q r() { return rot->rot; }
14    Q prev() { return rot->o->rot; }
15    Q next() { return r()->prev(); }
16 };
17 bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
18    lll p2 = p.dist2(), A = a.dist2() - p2,
19        B = b.dist2() - p2, C = c.dist2() - p2;
20    return p.cross(a, b) * C + p.cross(b, c) * A +
21        p.cross(c, a) * B >
22        0;
23 }
24 Q makeEdge(P orig, P dest) {
25    Q q[] = {new Quad{0, 0, 0, orig}, new Quad{0, 0, 0, arb},
26        new Quad{0, 0, 0, dest}, new Quad{0, 0, 0, arb}};
27    rep(i, 0, 4) q[i]->o = q[(i - 3) % 4],
28        q[i]->rot = q[(i + 1) % 4];
29    return *q;
30 }
31 void splice(Q a, Q b) {
32    swap(a->o->rot->o, b->o->rot->o);
33    swap(a->o, b->o);
34 }
35 Q connect(Q a, Q b) {
36    Q q = makeEdge(a->F(), b->p);
37    splice(q, a->next());
38    splice(q->r(), b);
39    return q;
40 }
41
42 pair<Q, Q> rec(const vector<P> &s) {
43    if (sz(s) <= 3) {
44        Q a = makeEdge(s[0], s[1]),
45            b = makeEdge(s[1], s.back());
46        if (sz(s) == 2) return {a, a->r()};
47        splice(a->r(), b);
48        auto side = s[0].cross(s[1], s[2]);
49        Q c = side ? connect(b, a) : 0;
50        return {side < 0 ? c->r() : a, side < 0 ? c : b->r()};
51    }
52 }

```

```

#define H(e)      e->F(), e->p
#define valid(e) (e->F().cross(H(base)) > 0)
Q A, B, ra, rb;
int half = sz(s) / 2;
tie(ra, A) = rec({all(s) - half});
tie(B, rb) = rec({sz(s) - half + all(s)});
while ((B->p.cross(H(A)) < 0 && (A = A->next())) ||
        (A->p.cross(H(B)) > 0 && (B = B->r()->o)))
;
Q base = connect(B->r(), A);
if (A->p == ra->p) ra = base->r();
if (B->p == rb->p) rb = base;

#define DEL(e, init, dir)
Q e = init->dir;
if (valid(e))
    while (circ(e->dir->F(), H(base), e->F())) {
        Q t = e->dir;
        splice(e, e->prev());
        splice(e->r(), e->r()->prev());
        e = t;
    }
for (;;) {
    DEL(LC, base->r(), o);
    DEL(RC, base, prev());
    if (!valid(LC) && !valid(RC)) break;
    if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
        base = connect(RC, base->r());
    else base = connect(base->r(), LC->r());
}
return {ra, rb};
}

// returns [A_0, B_0, C_0, A_1, B_1, ...]
// where A_i, B_i, C_i are counter-clockwise triangles
vector<P> triangulate(vector<P> pts) {
    sort(all(pts));
    assert(unique(all(pts)) == pts.end());
    if (sz(pts) < 2) return {};
    Q e = rec(pts).first;
    vector<Q> q = {e};
    int qi = 0;
    while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
    #define ADD
    {
        Q c = e;
        do {
            c->mark = 1;
            pts.push_back(c->p);
            q.push_back(c->r());
            c = c->next();
        } while (c != e);
    }
    ADD;
    pts.clear();
    while (qi < sz(q))
        if (!(e = q[qi++])->mark) ADD;
    return pts;
}

```

### 6.9.1. Slower Version

```

1
3 template <class P, class F>
void delaunay(vector<P> &ps, F trfun) {
5     if (sz(ps) == 3) {
        int d = (ps[0].cross(ps[1], ps[2]) < 0);
        trfun(0, 1 + d, 2 - d);
    }
    vector<P3> p3;
    for (P p : ps) p3.emplace_back(p.x, p.y, p.dist2());
11    if (sz(ps) > 3)
        for (auto t : hull3d(p3))
            if ((p3[t.b] - p3[t.a])
                .cross(p3[t.c] - p3[t.a])
                .dot(P3(0, 0, 1)) < 0)
                trfun(t.a, t.c, t.b);
17 }

```

### 6.10. Half Plane Intersection

```

1 struct Line {
    Point P;
    Vector v;
3     bool operator<(const Line &b) const {
        return atan2(v.y, v.x) < atan2(b.v.y, b.v.x);
    }
};
7 bool OnLeft(const Line &L, const Point &p) {
    return Cross(L.v, p - L.P) > 0;
9

```

```

}
11 Point GetIntersection(Line a, Line b) {
    Vector u = a.P - b.P;
13     Double t = Cross(b.v, u) / Cross(a.v, b.v);
    return a.P + a.v * t;
15 }
int HalfplaneIntersection(Line *L, int n, Point *poly) {
17     sort(L, L + n);

19     int first, last;
    Point *p = new Point[n];
    Line *q = new Line[n];
    q[first = last = 0] = L[0];
23     for (int i = 1; i < n; i++) {
        while (first < last && !OnLeft(L[i], p[last - 1]))
            last--;
25         while (first < last && !OnLeft(L[i], p[first])) first++;
        q[++last] = L[i];
27         if (fabs(Cross(q[last].v, q[last - 1].v)) < EPS) {
            last--;
29             if (OnLeft(q[last], L[i].P)) q[last] = L[i];
        }
31         if (first < last)
            p[last - 1] = GetIntersection(q[last - 1], q[last]);
33     }
    while (first < last && !OnLeft(q[first], p[last - 1]))
        last--;
35     if (last - first <= 1) return 0;
    p[last] = GetIntersection(q[last], q[first]);
37
39     int m = 0;
    for (int i = first; i <= last; i++) poly[m++] = p[i];
41     return m;
43 }

```

## 7. Strings

### 7.1. Knuth-Morris-Pratt Algorithm

```

1 vector<int> pi(const string &s) {
2     vector<int> p(s.size());
3     for (int i = 1; i < s.size(); i++) {
4         int g = p[i - 1];
5         while (g && s[i] != s[g]) g = p[g - 1];
6         p[i] = g + (s[i] == s[g]);
7     }
8     return p;
9 }
10
11 vector<int> match(const string &s, const string &pat) {
12     vector<int> p = pi(pat + '\0' + s), res;
13     for (int i = p.size() - s.size(); i < p.size(); i++)
14         if (p[i] == pat.size())
15             res.push_back(i - 2 * pat.size());
16     return res;
17 }

```

### 7.2. Aho-Corasick Automaton

```

1 struct Aho_Corasick {
2     static const int maxc = 26, maxn = 4e5;
3     struct NODES {
4         int Next[maxc], fail, ans;
5     };
6     NODES T[maxn];
7     int top, qtop, q[maxn];
8     int get_node(const int &fail) {
9         fill_n(T[top].Next, maxc, 0);
10        T[top].fail = fail;
11        T[top].ans = 0;
12        return top++;
13    }
14    int insert(const string &s) {
15        int ptr = 1;
16        for (char c : s) { // change char id
17            c -= 'a';
18            if (!T[ptr].Next[c]) T[ptr].Next[c] = get_node(ptr);
19            ptr = T[ptr].Next[c];
20        }
21        return ptr;
22    } // return ans_last_place
23    void build_fail(int ptr) {
24        int tmp;
25        for (int i = 0; i < maxc; i++)
26            if (T[ptr].Next[i]) {
27                tmp = T[ptr].fail;
28                while (tmp != 1 && !T[tmp].Next[i])
29                    tmp = T[tmp].fail;
30                if (T[tmp].Next[i] != T[ptr].Next[i])
31                    if (T[tmp].Next[i]) tmp = T[tmp].Next[i];
32                T[T[ptr].Next[i]].fail = tmp;
33                q[qtop++] = T[ptr].Next[i];
34            }
35    }
36    void AC_auto(const string &s) {
37        int ptr = 1;
38        for (char c : s) {
39            while (ptr != 1 && !T[ptr].Next[c]) ptr = T[ptr].fail;
40            if (T[ptr].Next[c]) {
41                ptr = T[ptr].Next[c];
42                T[ptr].ans++;
43            }
44        }
45    }
46    void Solve(string &s) {
47        for (char &c : s) // change char id
48            c -= 'a';
49        for (int i = 0; i < qtop; i++) build_fail(q[i]);
50        AC_auto(s);
51        for (int i = qtop - 1; i > -1; i--)
52            T[T[q[i]].fail].ans += T[q[i]].ans;
53    }
54    void reset() {
55        qtop = top = q[0] = 1;
56        get_node(1);
57    }
58 } AC;
59 // usage example
60 string s, S;
61 int n, t, ans_place[50000];
62 int main() {
63     Tie cin >> t;
64     while (t--) {
65         AC.reset();
66         cin >> S >> n;

```

```

67         for (int i = 0; i < n; i++) {
68             cin >> s;
69             ans_place[i] = AC.insert(s);
70         }
71         AC.Solve(S);
72         for (int i = 0; i < n; i++)
73             cout << AC.T[ans_place[i]].ans << '\n';
74     }
75 }

```

### 7.3. Suffix Array

```

1 // sa[i]: starting index of suffix at rank i
2 // 0-indexed, sa[0] = n (empty string)
3 // lcp[i]: lcp of sa[i] and sa[i - 1], lcp[0] = 0
4 struct SuffixArray {
5     vector<int> sa, lcp;
6     SuffixArray(string &s,
7                 int lim = 256) { // or basic_string<int>
8         int n = sz(s) + 1, k = 0, a, b;
9         vector<int> x(all(s) + 1), y(n), ws(max(n, lim)),
10            rank(n);
11         sa = lcp = y, iota(all(sa), 0);
12         for (int j = 0, p = 0; p < n; j = max(1, j * 2), lim = p) {
13             p = j, iota(all(y), n - j);
14             for (int i = 0; i < n; i++)
15                 if (sa[i] >= j) y[p++] = sa[i] - j;
16             fill(all(ws), 0);
17             for (int i = 0; i < n; i++) ws[x[i]]++;
18             for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];
19             for (int i = n; i--;) sa[-ws[x[i]]] = y[i];
20             swap(x, y), p = 1, x[sa[0]] = 0;
21             for (int i = 1; i < n; i++)
22                 a = sa[i - 1], b = sa[i],
23                 x[b] = (y[a] == y[b] && y[a + j] == y[b + j])
24                     ? p - 1 : p++;
25         }
26         for (int i = 1; i < n; i++) rank[sa[i]] = i;
27         for (int i = 0, j; i < n - 1; lcp[rank[i+1]] = k)
28             for (k && k--, j = sa[rank[i] - 1];
29                  s[i + k] == s[j + k]; k++);
30     }
31 };

```

### 7.4. Suffix Tree

```

1 struct SAM {
2     static const int maxc = 26; // char range
3     static const int maxn = 10010; // string len
4     struct Node {
5         Node *green, *edge[maxc];
6         int max_len, in, times;
7     } * root, *last, reg[maxn * 2];
8     int top;
9     Node *get_node(int _max) {
10        Node *re = &reg[top++];
11        re->in = 0, re->times = 1;
12        re->max_len = _max, re->green = 0;
13        for (int i = 0; i < maxc; i++) re->edge[i] = 0;
14        return re;
15    }
16    void insert(const char c) { // c in range [0, maxc)
17        Node *p = last;
18        last = get_node(p->max_len + 1);
19        while (p && !p->edge[c])
20            p->edge[c] = last, p = p->green;
21        if (!p) last->green = root;
22        else {
23            Node *pot_green = p->edge[c];
24            if ((pot_green->max_len) == (p->max_len + 1))
25                last->green = pot_green;
26            else {
27                Node *wish = get_node(p->max_len + 1);
28                wish->times = 0;
29                while (p && p->edge[c] == pot_green)
30                    p->edge[c] = wish, p = p->green;
31                for (int i = 0; i < maxc; i++)
32                    wish->edge[i] = pot_green->edge[i];
33                wish->green = pot_green->green;
34                pot_green->green = wish;
35                last->green = wish;
36            }
37        }
38    }

```



```

39 Node *q[maxn * 2];
40 int ql, qr;
41 void get_times(Node *p) {
42     ql = 0, qr = -1, reg[0].in = 1;
43     for (int i = 1; i < top; i++) reg[i].green->in++;
44     for (int i = 0; i < top; i++)
45         if (!reg[i].in) q[++qr] = &reg[i];
46     while (ql <= qr) {
47         q[ql]->green->times += q[ql]->times;
48         if (!(--q[ql]->green->in)) q[++qr] = q[ql]->green;
49         ql++;
50     }
51 }
52 void build(const string &s) {
53     top = 0;
54     root = last = get_node(0);
55     for (char c : s) insert(c - 'a'); // change char id
56     get_times(root);
57 }
58 // call build before solve
59 int solve(const string &s) {
60     Node *p = root;
61     for (char c : s)
62         if (!(p = p->edge[c - 'a'])) // change char id
63             return 0;
64     return p->times;
65 }
};

```

## 7.5. Cocke-Younger-Kasami Algorithm

```

1
2
3 struct rule {
4     // s -> xy
5     // if y == -1, then s -> x (unit rule)
6     int s, x, y, cost;
7 };
8 int state;
9 // state (id) for each letter (variable)
10 // lowercase letters are terminal symbols
11 map<char, int> rules;
12 vector<rule> cnf;
13 void init() {
14     state = 0;
15     rules.clear();
16     cnf.clear();
17 }
18 // convert a cfg rule to cnf (but with unit rules) and add
19 // it
20 void add_to_cnf(char s, const string &p, int cost) {
21     if (!rules.count(s)) rules[s] = state++;
22     for (char c : p)
23         if (!rules.count(c)) rules[c] = state++;
24     if (p.size() == 1) {
25         cnf.push_back({rules[s], rules[p[0]], -1, cost});
26     } else {
27         // length >= 3 -> split
28         int left = rules[s];
29         int sz = p.size();
30         for (int i = 0; i < sz - 2; i++) {
31             cnf.push_back({left, rules[p[i]], state, 0});
32             left = state++;
33         }
34         cnf.push_back(
35             {left, rules[p[sz - 2]], rules[p[sz - 1]], cost});
36     }
37 }
38
39 constexpr int MAXN = 55;
40 vector<long long> dp[MAXN][MAXN];
41 // unit rules with negative costs can cause negative cycles
42 vector<bool> neg_INF[MAXN][MAXN];
43
44 void relax(int l, int r, rule c, long long cost,
45           bool neg_c = 0) {
46     if (!neg_INF[l][r][c.s] &&
47         (neg_INF[l][r][c.x] || cost < dp[l][r][c.s])) {
48         if (neg_c || neg_INF[l][r][c.x]) {
49             dp[l][r][c.s] = 0;
50             neg_INF[l][r][c.s] = true;
51         } else {
52             dp[l][r][c.s] = cost;
53         }
54     }
55 }
56 void bellman(int l, int r, int n) {
57     for (int k = 1; k <= state; k++)
58         for (rule c : cnf)
59             if (c.y == -1)
60                 relax(l, r, c, dp[l][r][c.x] + c.cost, k == n);

```

```

61 }
62 void cyk(const string &s) {
63     vector<int> tok;
64     for (char c : s) tok.push_back(rules[c]);
65     for (int i = 0; i < tok.size(); i++) {
66         for (int j = 0; j < tok.size(); j++) {
67             dp[i][j] = vector<long long>(state + 1, INT_MAX);
68             neg_INF[i][j] = vector<bool>(state + 1, false);
69         }
70         dp[i][i][tok[i]] = 0;
71         bellman(i, i, tok.size());
72     }
73     for (int r = 1; r < tok.size(); r++) {
74         for (int l = r - 1; l >= 0; l--) {
75             for (int k = l; k < r; k++)
76                 for (rule c : cnf)
77                     if (c.y != -1)
78                         relax(l, r, c,
79                             dp[l][k][c.x] + dp[k + 1][r][c.y] +
80                             c.cost);
81             bellman(l, r, tok.size());
82         }
83     }
84 }
85 // usage example
86 int main() {
87     init();
88     add_to_cnf('S', "aSc", 1);
89     add_to_cnf('S', "BBB", 1);
90     add_to_cnf('S', "SB", 1);
91     add_to_cnf('B', "b", 1);
92     cyk("abbbbc");
93     // dp[0][s.size() - 1][rules[start]] = min cost to
94     // generate s
95     cout << dp[0][5][rules['S']] << '\n'; // 7
96     cyk("acbc");
97     cout << dp[0][3][rules['S']] << '\n'; // INT_MAX
98     add_to_cnf('S', "S", -1);
99     cyk("abbbbc");
100     cout << neg_INF[0][5][rules['S']] << '\n'; // 1
101 }

```

## 7.6. Z Value

```

1 int z[n];
2 void zval(string s) {
3     // z[i] => longest common prefix of s and s[i:], i > 0
4     int n = s.size();
5     z[0] = 0;
6     for (int b = 0, i = 1; i < n; i++) {
7         if (z[b] + b <= i) z[i] = 0;
8         else z[i] = min(z[i - b], z[b] + b - i);
9         while (s[i + z[i]] == s[z[i]]) z[i]++;
10        if (i + z[i] > b + z[b]) b = i;
11    }
12 }

```

## 7.7. Manacher's Algorithm

```

1 int z[n];
2 void manacher(string s) {
3     // z[i] => longest odd palindrome centered at i is
4     // s[i - z[i]] ... i + z[i]]
5     // to get all palindromes (including even length),
6     // insert a '#' between each s[i] and s[i + 1]
7     int n = s.size();
8     z[0] = 0;
9     for (int b = 0, i = 1; i < n; i++) {
10        if (z[b] + b >= i)
11            z[i] = min(z[2 * b - i], b + z[b] - i);
12        else z[i] = 0;
13        while (i + z[i] + 1 < n && i - z[i] - 1 >= 0 &&
14            s[i + z[i] + 1] == s[i - z[i] - 1])
15            z[i]++;
16        if (z[i] + i > z[b] + b) b = i;
17    }
18 }

```

## 7.8. Minimum Rotation

```

1 int min_rotation(string s) {
2     int a = 0, n = s.size();
3     s += s;
4     for (int b = 0; b < n; b++) {
5         for (int k = 0; k < n; k++) {
6             if (a + k == b || s[a + k] < s[b + k]) {
7                 b += max(0, k - 1);
8                 break;
9             }
10            if (s[a + k] > s[b + k]) {

```

```

11     a = b;
12     break;
13 }
14 }
15 }
16 return a;
17 }

```

## 7.9. Palindromic Tree

```

1
3 struct palindromic_tree {
4     struct node {
5         int next[26], fail, len;
6         int cnt,
7         num; // cnt: appear times, num: number of pal. suf.
8         node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
9             for (int i = 0; i < 26; ++i) next[i] = 0;
10        }
11    };
12    vector<node> St;
13    vector<char> s;
14    int last, n;
15    palindromic_tree() : St(2), last(1), n(0) {
16        St[0].fail = 1, St[1].len = -1, s.pb(-1);
17    }
18    inline void clear() {
19        St.clear(), s.clear(), last = 1, n = 0;
20        St.pb(0), St.pb(-1);
21        St[0].fail = 1, s.pb(-1);
22    }
23    inline int get_fail(int x) {
24        while (s[n - St[x].len - 1] != s[n]) x = St[x].fail;
25        return x;
26    }
27    inline void add(int c) {
28        s.push_back(c - 'a'), ++n;
29        int cur = get_fail(last);
30        if (!St[cur].next[c]) {
31            int now = SZ(St);
32            St.pb(St[cur].len + 2);
33            St[now].fail = St[get_fail(St[cur].fail)].next[c];
34            St[cur].next[c] = now;
35            St[now].num = St[St[now].fail].num + 1;
36        }
37        last = St[cur].next[c], ++St[last].cnt;
38    }
39    inline void count() { // counting cnt
40        auto i = St.rbegin();
41        for (; i != St.rend(); ++i) {
42            St[i->fail].cnt += i->cnt;
43        }
44    }
45    inline int size() { // The number of diff. pal.
46        return SZ(St) - 2;
47    }
48 };

```



## 8. Debug List

- 1 - Pre-submit:
  - 3 - Did you make a typo when copying a template?
  - 3 - Test more cases if unsure.
    - Write a naive solution and check small cases.
  - 5 - Submit the correct file.
- 7 - General Debugging:
  - Read the whole problem again.
  - 9 - Have a teammate read the problem.
  - Have a teammate read your code.
    - 11 - Explain you solution to them (or a rubber duck).
  - Print the code and its output / debug output.
  - 13 - Go to the toilet.
- 15 - Wrong Answer:
  - Any possible overflows?
    - 17 - > `__int128` ?
      - Try `-ftrapv` or `#pragma GCC optimize("trapv")`
  - 19 - Floating point errors?
    - > `long double` ?
    - 21 - turn off math optimizations
      - check for `'=='`, `'>='`, `'acos(1.000000001)'`, etc.
  - 23 - Did you forget to sort or unique?
  - Generate large and worst "corner" cases.
  - 25 - Check your `'m' / 'n'`, `'i' / 'j'` and `'x' / 'y'`.
  - Are everything initialized or reset properly?
  - 27 - Are you sure about the STL thing you are using?
    - Read cppreference (should be available).
  - 29 - Print everything and run it on pen and paper.
- 31 - Time Limit Exceeded:
  - Calculate your time complexity again.
  - 33 - Does the program actually end?
    - Check for `'while(q.size())'` etc.
  - 35 - Test the largest cases locally.
  - Did you do unnecessary stuff?
    - 37 - e.g. pass vectors by value
    - e.g. `'memset'` for every test case
  - 39 - Is your constant factor reasonable?
- 41 - Runtime Error:
  - Check memory usage.
    - 43 - Forget to clear or destroy stuff?
      - > `'vector::shrink_to_fit()'`
  - 45 - Stack overflow?
  - Bad pointer / array access?
    - 47 - Try `'-fsanitize=address'`
  - Division by zero? NaN's?