

# SURIM 2018 - Quantitative and Local Central Limit Theorems

## 1 General Background

### 1.1 Probability Theory

CLT:

$$|P(X_n < t) - P(Z < t)| \xrightarrow{n \rightarrow \infty} 0$$

LLT:

$$P(X_n = k) = \Phi(k) + o(1) \text{ where } \Phi(k) = \frac{1}{\sqrt{2\pi}\sigma_n} \exp(-(k - \mu_n)/\sigma_n)^2/2)$$

Berry-Esseen Theorem: If the third moment is finite, given i.i.d. random variables, gives a quantitative bound for CLT

Levy's Continuity Theorem: If a sequence of characteristic functions converge pointwise to  $\phi(t)$ ,  $X_n$  converges to random variable  $X$ , whose characteristic function is  $\phi(t)$ .

Stein's theorem (1986): If finite third and fourth moments, the dependency graph has finite degree  $d < \infty$ , we have a general CLT

### 1.2 Fourier Analysis

Fourier transform: physically, a complex valued function of frequency, magnitude represents amount of frequency present in original function

$$\hat{f}(\xi) := \int_{-\infty}^{+\infty} f(x) e^{-2\pi i x \xi} dx$$

Inverse transform:

$$f(x) := \int_{-\infty}^{+\infty} \hat{f}(\xi) e^{-2\pi i x \xi} d\xi$$

Dirac delta function: loosely,  $\delta(x) = \begin{cases} +\infty & x = 0 \\ 0 & x \neq 0 \end{cases}$

Parseval's Theorem: Fourier transform is unitary (loosely, integral of square of function equals the integral of square of its transform)

### 1.3 Measure Theory

Dominated Convergence Theorem (used in Gilmer Kopparty small  $t$ ): almost everywhere convergence of a sequence of functions implies convergence in the  $L^1$  norm.

## 2 Papers

### Gilmer and Kopparty (2014)

- Establishes that a local limit theorem for the number of triangles in a random graph exists
- More specifically, shows  $\int_{-\pi\sigma_n}^{\pi\sigma_n} |\phi_n(t) - e^{-t^2/2}| dt = o(1)$  (tends to 0).

- Proof splits cases into small, intermediate, and big  $t$ . Small comes relatively easily from CLT but authors provide no bound on rate of convergence; intermediate: splits the  $\binom{n}{2}$  edges up into those that are contained in a perfect matching  $M$  and those that are not; big  $t$  has a good  $t^{-50}$  bound.

## Berkowitz (2017)

- Provides a quantitative bound on how far the distribution of the number of triangles in a random graph can vary from a discrete normal distribution
- Shows  $\int_{-\pi\sigma_n}^{\pi\sigma_n} \left| \phi_n(t) - e^{-t^2/2} \right| dt = O(n^{-1/2+\epsilon})$

## Possibly related papers

### 2.0.1 Chatterjee and Diaconis (2017) A Central Limit Theorem for a New Statistic on Permutations

- Establishes a CLT for the random variable  $T(\pi) := D(\pi) + D(\pi^{-1})$
- Gives 6 approaches for establishing CLT for descents  $D(\pi)$ .

### 2.0.2 Chatterjee (2014) A Short Survey of Stein's Method

- More ideas for possible random variables for which CLT's haven't been shown: Euclidean salesman problem (choose set of  $n$  random points from a unit square, let  $P$  be the shortest path that visits all points and ends with the same point,  $P$  should satisfy a CLT), minimum matching problem (suppose  $n$  is even, pair the points into  $n/2$  pairs such that the sum  $S$  of the pairwise distances is minimized,  $S$  should satisfy a CLT), analogous problems with weights on edges

## 3 New Random Variables

### 3.1 Arithmetic Progressions

Let  $A_n$  be the number of arithmetic progressions in a random subset of  $F_n$ .

$A_n = \frac{1}{2} \sum_{k \in \mathbb{F}_n} \sum_{j \neq 0 \in \mathbb{F}_n} X_k X_{k+j} X_{k+2j}$  for  $n$  prime and not 2 or 3.

Given a fixed  $k \in \mathbb{F}_n$ , there are  $\frac{3}{2}(n-1)$  arithmetic progressions that contain  $k$ .

LLT Bound:

For small  $t$ : for  $|t| < \sqrt{n}$ ,

$$\left| \phi_Z(t) - e^{-t^2/2} \right| \leq O\left( \frac{t^3 e^{-t^2/3}}{\sqrt{n}} + \frac{t}{n^{5/4}} \right)$$

.

### 3.2 Number of Descents in a Permutation

Let  $D(\pi) = \#\{i : 1 \leq i \leq n-1, \pi(i) > \pi(i+1)\}$ .

$$\mathbb{E}[D(\pi)] = \frac{n-1}{2}; \text{Var}(D(\pi)) = \frac{n+1}{12}$$

Since dependency graph for  $D(\pi)$  has degree  $2 < \infty$ , by Stein's method, we have a generic CLT for  $D(\pi)$ .

Let  $X_i = \begin{cases} 1 & \pi(i) > \pi(i+1) \\ 0 & \text{else} \end{cases}$

Probability

$$P(X_i X_{i+1} \dots X_{i+k}) = \frac{1}{(k+2)!}$$

Eulerian numbers:  $A(n, m)$  is the number of permutations of the numbers from 1 to  $n$  that have exactly  $m$  descents (ascents).

Possible strategy: Consider the following decomposition:  $D(\pi) = \text{Odd}(\pi) + \text{Even}(\pi)$  where

$$\begin{aligned}\text{Odd}(\pi) &= \#\{i : i \text{ odd}, 1 \leq i \leq n-1, \pi(i) > \pi(i+1)\} \\ \text{Even}(\pi) &= \#\{i : i \text{ even}, 1 \leq i \leq n-1, \pi(i) > \pi(i+1)\}\end{aligned}$$

Then we essentially have a random variable conditioned on a binary string of the values of the odd  $X_i$ , with each string equally likely since  $\mathbb{E}(X_i) = 1/2$ .

$$\begin{aligned}P(X_{i+1} = 1 \mid X_i = 1, X_{i+2} = 1) &= 1/6 \\ P(X_{i+1} = 1 \mid X_i = 1, X_{i+2} = 0) &= 1/2 \\ P(X_{i+1} = 1 \mid X_i = 0, X_{i+2} = 1) &= 1/2 \\ P(X_{i+1} = 1 \mid X_i = 0, X_{i+2} = 0) &= 5/6\end{aligned}$$

If we condition on the values of  $X_i$  for odd  $i$ , the only remaining variables are the  $X_i$  for even  $i$ , which are all independent. Furthermore, the distribution of an even  $X_i$  is determined entirely by the odd  $X_{i-1}$  and  $X_{i+1}$  surrounding it, so the only important data from the odd  $X_i$ , aside from the total number that are 1, is the number  $N_{11}$  of subsequences 11 of the binary string, the numbers  $N_{10}$  of subsequences 10 of the binary string, etc. Once the odd  $X_i$  are determined and these are fixed,  $D$  is the sum of four binomial distributions (plus a constant).

To establish a local limit theorem, we want to prove a quantitative bound for  $\int_{-\pi\sigma_n}^{\pi\sigma_n} |\phi_n(t) - e^{-t^2/2}| dt$ .

$$\left| \mathbb{E} e^{itD/\sigma} \right| = \left| \mathbb{E} \exp \left( it \left( \sum_{j=1}^{N_{11}} B_j + \dots \right) / \sigma \right) \right|$$

## 4 AP Stuff small t or whatever

We define a random variable  $\mathcal{A}$  to be the number of 3-term arithmetic progressions in a randomly chosen subset of  $\mathbb{F}_n$ , where each element in  $\mathbb{F}_n$  has probability  $p$  of appearing in the subset. The elements of the probability space can thus be described as  $\mathbf{x} \in \{0, 1\}^n$  where  $\mathbf{x}_k$  is 1 with probability  $p$  and 0 with probability  $1 - p$ . Further, the random variable  $\mathcal{A}$  can be thought of as the function  $\mathcal{A} : \{0, 1\}^n \rightarrow \mathbb{N}$ . Although  $\mathcal{A}$  depends on both  $n$  and  $p$ , we take  $p$  to be fixed and constrain our analysis to  $n$ .

We define a  $p$ -biased Fourier basis for functions on the probability space. In order to do this we first define  $\chi_k : \{0, 1\}^n \rightarrow \mathbb{R}$  by

$$\chi_k := \chi_k(\mathbf{x}) := \frac{\mathbf{x}_k - p}{\sqrt{p(1-p)}} = \begin{cases} -\sqrt{\frac{p}{1-p}} & \text{if } \mathbf{x}_p = 0 \\ \sqrt{\frac{1-p}{p}} & \text{if } \mathbf{x}_p = 1 \end{cases}$$

so that  $\chi_k$  is a normalized version of  $x_k$ . Further, we can extend this to define, for an arbitrary set  $S \subset \mathbb{F}_n$ ,

$$\chi_S := \prod_{k \in S} \chi_k$$

Note that if we take the inner product of two functions  $f, g : \{0, 1\}^n \rightarrow \mathbb{R}$  to be  $\mathbb{E}[fg]$ , then  $\{\chi_S : S \subset \mathbb{F}_n\}$  forms an orthonormal basis for functions on our probability space.

So if we define the Fourier transform  $\hat{f} : \{0, 1\}^n \rightarrow \mathbb{R}$  of an arbitrary function  $f : \{0, 1\}^n \rightarrow \mathbb{R}$  by

$$\hat{f}(S) := \mathbb{E}[f(\mathbf{x}) \chi_S(\mathbf{x})]$$

then from the orthonormality of our basis we get

$$f(\mathbf{x}) = \sum_{S \subset \mathbb{F}_n} \hat{f}(S) \chi_S(\mathbf{x})$$

We will use this expansion to calculate the variance of  $\mathcal{A}$  and bound the pointwise distance of the characteristic function from that of the discrete Gaussian.

It will now be useful to normalize the random variable  $\mathcal{A}$ . We take the mean of  $\mathcal{A}$  to be  $\mu := \mathbb{E}[\mathcal{A}]$ . We write the variance of  $\mathcal{A}$  as  $\sigma^2 := \mathbb{E}[\mathcal{A}^2] - \mathbb{E}[\mathcal{A}]^2$ . So we define  $Z : \{0, 1\}^n \rightarrow \mathbb{R}$  by

$$Z := \frac{\mathcal{A} - \mu}{\sigma}$$

and we will often refer to the characteristic function of  $Z$  defined by  $\phi_Z(t) := \mathbb{E}[e^{itZ}]$ .

Before moving on to calculate the Fourier coefficients  $\hat{\mathcal{A}}(S)$ , we will first note that there are  $\binom{n}{2}$  possible (non-trivial) arithmetic progressions in  $\mathbb{F}_n$ . There are first  $n$  choices for the start of the arithmetic progression, then  $n-1$  choices for a non-trivial separation distance  $d$ , and finally both  $d$  and  $-d$  will have counted the same arithmetic progression from different starting points, so we divide by 2 to yield  $\frac{n(n-1)}{2} = \binom{n}{2}$ . Additionally, each 3-term arithmetic progression occurs with probability  $p^3$  (each of the three terms in the progression occur independently with probability  $p$ ). This allows us to calculate

$$\mu = \mathbb{E} \left[ \sum_{\Lambda} 1_{\Lambda} \right] = \sum_{\Lambda} \mathbb{E}[1_{\Lambda}] = \sum_{\Lambda} p^3 = p^3 \binom{n}{2}$$

where  $1_{\Lambda}$  is the indicator function for a particular 3-term arithmetic progression  $\Lambda$  in  $\mathbb{F}_n$ .

Furthermore, the fact that our basis for functions on the probability space is orthonormal allows us to calculate variance according to this formula, commonly known as Parseval's Theorem.

$$\sigma^2 = \sum_{S \neq \emptyset} \hat{\mathcal{A}}(S)^2$$

We now work with the arithmetic progression indicator functions  $1_{\Lambda}$  a bit more. We use  $k \in \Lambda$  to denote that  $k$  is a term in the arithmetic progression  $\Lambda$ . Therefore, the indicator function can be expressed as

$$\begin{aligned} 1_{\Lambda}(\mathbf{x}) &= \prod_{k \in \Lambda} \mathbf{x}_k = \prod_{k \in \Lambda} \left( \sqrt{p(1-p)} \chi_k + p \right) \\ &= p^3 + p^2 \sqrt{p(1-p)} \sum_{k \in \Lambda} \chi_k + p^2(1-p) \sum_{k_1 \neq k_2 \in \Lambda} \chi_{\{k_1, k_2\}} + p^{3/2}(1-p)^{3/2} \end{aligned}$$

Note that any two elements of  $\mathbb{F}_n$  appear in exactly 3 arithmetic progressions and any one element appears in exactly  $\frac{3}{2}(n-1)$  arithmetic progressions. Hence, by summing over all arithmetic progressions, we have

$$\mathcal{A} = p^3 \binom{n}{2} + \frac{3}{2}(n-1)p^2 \sqrt{p(1-p)} \sum_{k \in \mathbb{F}_n} \chi_k + 3p^2(1-p) \sum_{\{k_1, k_2\} \in \mathbb{F}_n} \chi_{\{k_1, k_2\}} + \sum_{\Lambda} p^{3/2}(1-p)^{3/2} \chi_{\Lambda}$$

Thus, we have the Fourier Transform of  $\mathcal{A}$ :

$$\hat{\mathcal{A}}(S) = \begin{cases} p^3 \binom{n}{2} & \text{if } S = \emptyset \\ \frac{3}{2}(n-1)p^2 \sqrt{p(1-p)} & \text{if } |S| = 1 \\ 3p^2(1-p) & \text{if } S = \{k_1, k_2\} \\ p^{3/2}(1-p)^{3/2} & \text{if } S = \Lambda \\ 0 & \text{else} \end{cases}$$

and we can use Parseval's Theorem to give us the variance of  $\mathcal{A}$  from this

$$\begin{aligned}
\sigma^2 &= \sum_{\substack{S \subset \mathbb{F}_n \\ S \neq \emptyset}} \hat{\mathcal{A}}(S)^2 \\
&= \sum_{k \in \mathbb{F}_n} \left( \frac{3}{2}(n-1)p^2\sqrt{p(1-p)} \right)^2 + \sum_{\{k_1, k_2\} \in \mathbb{F}_n} (3p^2(1-p))^2 + \sum_{\Lambda} \left( p^{3/2}(1-p)^{3/2} \right)^2 \\
&= \frac{9}{4}n(n-1)^2p^5(1-p) + 9\binom{n}{2}p^4(1-p)^2 + \binom{n}{2}p^3(1-p)^3 \\
&= \Theta(n^3)
\end{aligned}$$

We also derive the fact that  $\sigma = \Theta(n^{3/2})$  from this calculation. We proceed to show the following bound for small values of  $t$ :

**Proposition 4.1.** *For  $|t| \ll \sqrt{n}$ ,*

$$|\phi_Z(t) - e^{-t^2/2}| \leq O\left(\frac{t^3 e^{-t^2/3}}{\sqrt{n}} + \frac{t}{\sqrt{n}}\right).$$

*Proof.* To begin, we decompose  $Z = X + Y$ , where

$$Q = \frac{1}{\sqrt{n}}, X = \sum_{k \in \mathbb{F}_n} Q\chi_k, Y = \sum_{k \in \mathbb{F}_n} (\hat{Z}(k) - Q)\chi_k + \sum_{|S| \geq 2} \hat{Z}(S)\chi_S.$$

We first bound the distance from the characteristic function of  $X$  to the normal distribution: Since we normalized  $X$  using  $Q$ ,  $X$  is a random variable with mean 0 and variance 1. Since

$$L_n := n\mathbb{E}[|Q\chi_k|^3] = O\left(\frac{1}{\sqrt{n}}\right) < \infty,$$

then by Berry-Esseen, if  $t \leq \frac{1}{4L_n}$ , then

$$|\mathbb{E}[e^{itX}] - e^{-t^2/2}| \leq 16L_n t^3 e^{-t^2/3}.$$

Next we consider  $Y$ . By Cauchy-Schwarz, orthogonality of our basis, and Parseval's Theorem, we have

$$\mathbb{E}|Y| \leq \sqrt{\mathbb{E}|Y|^2} = \sqrt{\text{var}(Y)} = \sqrt{\sum_{k \in \mathbb{F}_n} (\hat{Z}(k) - Q)^2 + \sum_{|S| \geq 2} \hat{Z}(S)^2}.$$

Now using our above calculation of the variance  $\sigma^2$  of  $\mathcal{A}$ ,

$$\sum_{|S| \geq 2} \hat{Z}(S)^2 = \frac{9\binom{n}{2}p^4(1-p)^2 + \binom{n}{2}p^3(1-p)^3}{\sigma^2} = \Theta\left(\frac{1}{n}\right).$$

In addition, using our calculation of  $\hat{\mathcal{A}}$ ,

$$\begin{aligned}
n\sigma\hat{Z}^2(k) - \sigma^2 &= n\hat{\mathcal{A}}^2(k) - \sigma^2 = O(n^2) \\
n\hat{Z}^2(k) - 1 &= O\left(\frac{1}{n}\right) \\
\hat{Z}^2(k) - \frac{1}{n} &= O\left(\frac{1}{n^2}\right).
\end{aligned}$$

Since  $\hat{Z}(k) + Q = O\left(\frac{1}{\sqrt{n}}\right)$  and  $Q^2 = \frac{1}{n}$ ,

$$\left|\hat{Z}(k) - Q\right| \leq \left|\frac{\hat{Z}(k) - \frac{1}{n}}{\hat{Z}(k) + Q}\right| = O\left(\frac{1}{n^{3/2}}\right).$$

Hence, we have  $\mathbb{E}[Y] = O\left(\frac{1}{\sqrt{n}}\right)$ . Thus, we conclude that if  $t < \frac{1}{4L_n} = O(\sqrt{n})$ ,

$$\begin{aligned} \left|\phi_Z(t) - e^{-t^2/2}\right| &= \left|\mathbb{E}[e^{itZ} - e^{-t^2/2}]\right| = \left|\mathbb{E}[e^{it(X+Y)} - e^{-t^2/2}]\right| \\ &\leq \left|\mathbb{E}[e^{it(X+Y)} - e^{itX}]\right| + \left|\mathbb{E}[e^{itX} - e^{-t^2/2}]\right| \\ &\leq \mathbb{E}[tY] + 16L_n t^3 e^{-t^2/3} \quad (\text{using the Mean Value Theorem on } e^{itX}) \\ &= O\left(\frac{t^3 e^{-t^2/3}}{\sqrt{n}} + \frac{t}{\sqrt{n}}\right). \end{aligned}$$

□