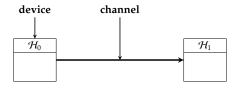
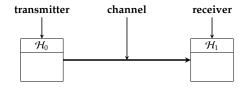
▶ We can *already* form a simple **point-to-point** communication channel



using TIA-232-F.

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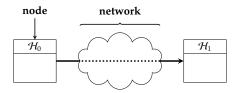
▶ We can *already* form a simple **point-to-point** communication channel



### using TIA-232-F, st.

- ► Good:
  - modular wrt. communication medium and protocol,
  - uses standardised components,
  - •
- Bad:
  - well defined, but quite limited functionality,
  - the organisation of components is fixed,
  - there are limits wrt. physical locality,
  - ▶ .

► We can *already* form a simple **point-to-point** communication channel



using TIA-232-F.

► Challenge: expand our remit to use of a **computer network** ...

COMS20001 lecture: week #19

▶ ... how?



COMS20001 lecture: week #19

- ▶ ... how?
  - 1. Requirements: what do we expect from a network?
    - supports a high degree of connectivity,
    - allows inter-operation between heterogeneous components,
    - uses appropriate level of abstraction to provide useful functionality,
    - satisfies any relevant quality metrics (e.g., efficiency, reliability),
    - can be (dynamically) scaled wrt. components and usage.

COMS20001 lecture: week #19

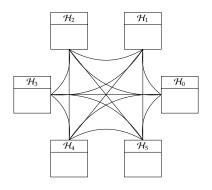
▶ ... how?

- 2. Architecture: how should we approach the design of a network?
  - (physical or logical) topology, i.e., the structure of components in the network, plus
  - control i.e., how do we use those components to communicate, and
  - standardisation, i.e., how do we ensure components which communicate can **inter-operate**.

▶ Idea #1: *fully*-connected, st.  $\mathcal{H}_i$  is connected to  $\mathcal{H}_i$  for  $j \in S = \{0, 1, ..., n-1\} \setminus \{i\}$ .

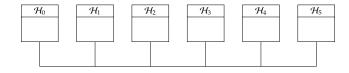
- ▶ Note that access to a given channel may be
  - 1. dedicated, or
  - shared, implying a need to control access but also the posibility to broadcast to multiple receiving nodes.

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- Example: a mesh topology, i.e.,



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- Example: a bus topology, i.e.,

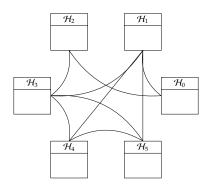


- ▶ Note that access to a given channel may be
  - 1. dedicated, or
  - shared, implying a need to control access but also the posibility to broadcast to multiple receiving nodes.

▶ Idea #2: *partially*-connected, st.  $\mathcal{H}_i$  is connected to  $\mathcal{H}_i$  for  $i \in S \subset \{0, 1, ..., n-1\} \setminus \{i\}$ .

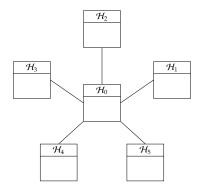
- Note that connectivity may now be either
  - 1. direct, or
  - 2. **indirect**, implying a need for intermediate **hops**, e.g., as realised via **store-and-forward** by (intermediate) switching nodes.

- ▶ Idea #2: partially-connected, st.  $\mathcal{H}_i$  is connected to  $\mathcal{H}_j$  for  $j \in S \subset \{0, 1, ..., n-1\} \setminus \{i\}$ .
- **Example:** a **mesh topology**, i.e.,



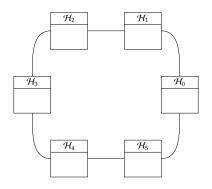
- Note that connectivity may now be either
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- ▶ Idea #2: *partially*-connected, st.  $\mathcal{H}_i$  is connected to  $\mathcal{H}_j$  for  $j \in S \subset \{0, 1, ..., n-1\} \setminus \{i\}$ .
- Example: a star topology, i.e.,



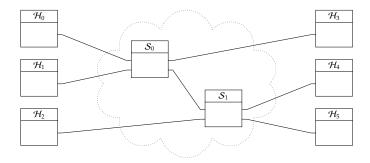
- Note that connectivity may now be either
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- Example: a ring topology, i.e.,

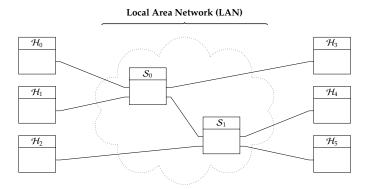


- Note that connectivity may now be either
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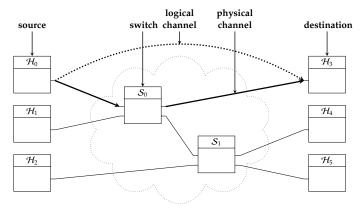
► We can *generalise* partially connected network topologies



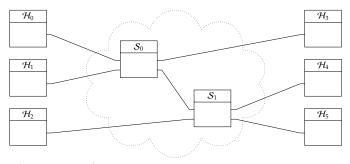
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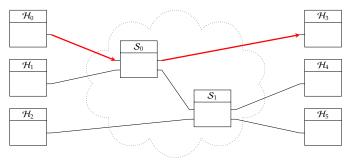


We can generalise partially connected network topologies



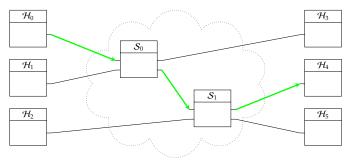
- circuit switching and
- packet switching (to support connection-based or connection-less channels).

We can generalise partially connected network topologies



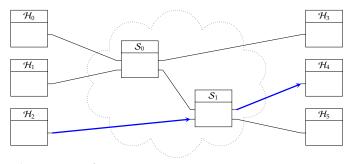
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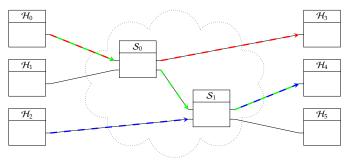
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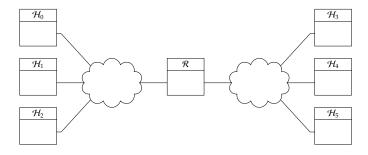
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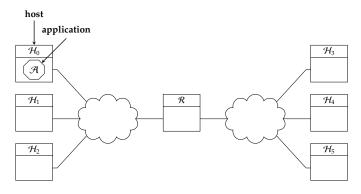
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### then, finally, noting

- a network provides connectivity between nodes, while
- an *inter*-network connects networks themselves together.

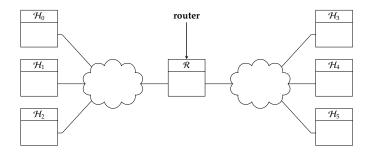
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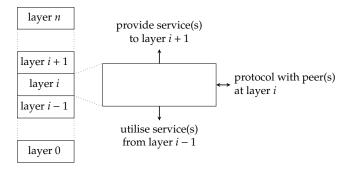
### then, finally, noting

- a network provides connectivity between nodes, while
- an *inter*-network connects networks themselves together.

- An inter-network architecture seems a good approach!
  - ► Good:
    - scalable by virtue of a flexible and modular design, and
    - offers resilience against failure.
  - Bad: we need to cope with the fact
    - ideally we support a multiplicity of communication mediums and protocols, and share access to both,
    - each host needs a globally unique address, and
    - router(s) need to know or discover how to communicate (or route) data from one host to another

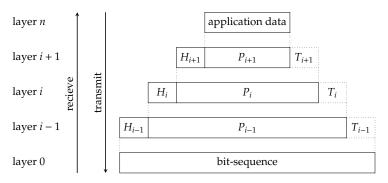
the difficulty of which is enhanced by a need to avoid impact on scalability (e.g., avoid centralised solutions).

▶ Step #1: rather than use one monolithic protocol, decompose control, e.g.,



and hence provide various advantages stemming from modularity ...

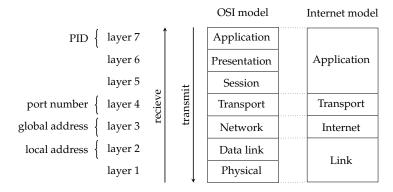
• ... then support inter-layer (peer or otherwise) communication using **encapsulation**:



► Step #2: specify (abstract) layers we need, e.g.,

			OSI model	Internet model
layer 7	1	`	Application	
layer 6			Presentation	Application
layer 5	e	īţ.	Session	
layer 4	recieve	transmit	Transport	Transport
layer 3	re	ţ.	Network	Internet
layer 2			Data link	Link
layer 1		$\downarrow$	Physical	LIIK

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					OSI model	Internet model
	layer 7	a			Application	
message {	layer 6				Presentation	Application
	layer 5		iţ		Session	
segment, datagram {	layer 4	recieve	transmit		Transport	Transport
packet {	layer 3	T.	Į.		Network	Internet
frame {	layer 2				Data link	Link
symbol (e.g., bit)	layer 1		$ \hspace{.1cm} $	.	Physical	LIIK

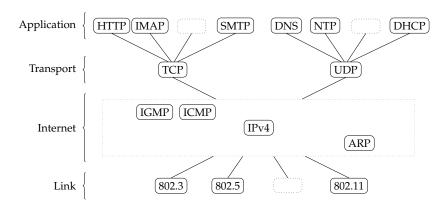
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					OSI model	Internet model
	layer 7	a	1		Application	
host	layer 6				Presentation	Application
nost	layer 5		transmit		Session	
	layer 4	recieve			Transport	Transport
router {	layer 3	re	Į,		Network	Internet
switch, bridge	layer 2				Data link	Link
NIC, repeater, hub	layer 1				Physical	LIIK

► Step #2: specify (abstract) layers we need, e.g.,

				OSI model		Internet model
user {	layer 7	1		Application		
	layer 6			Presentation		Application
	layer 5	, i		Session		
kernel {	layer 4	recieve		Transport		Transport
	layer 3	ti li		Network		Internet
hardware {	layer 2			Data link		Link
	layer 1			Physical		LINK

Step #3: populate layers with (concrete) protocols, e.g.,

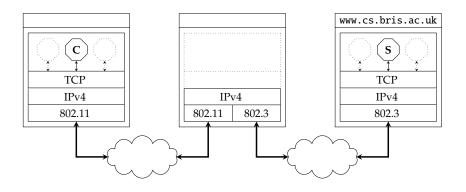


#### noting that

- a protocol graph is an abstract description of how protocols fit into the layered model, whereas
- \* a **protocol stack** is a concrete implementation of *one* top-to-bottom combination of protocols.

#### Conclusions

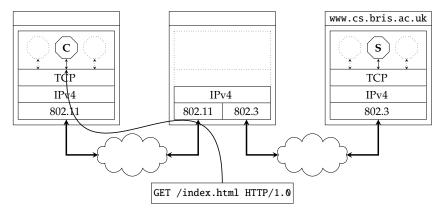
► As a final step, we can add some detail to our running example, e.g.,



with the rest of the unit aiming to further explain each layer.

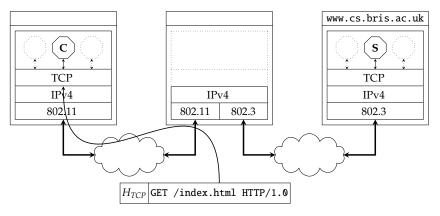
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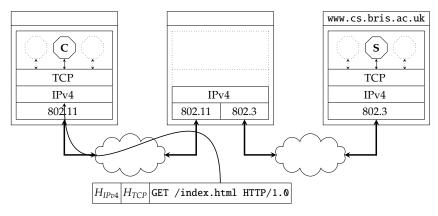


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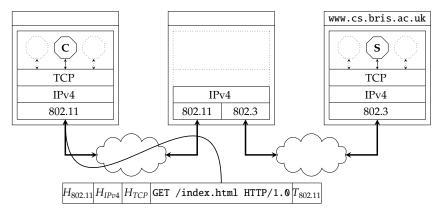
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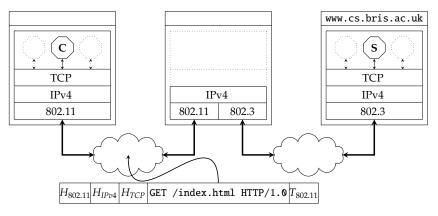
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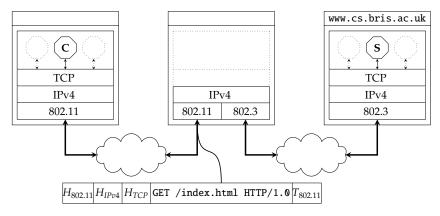
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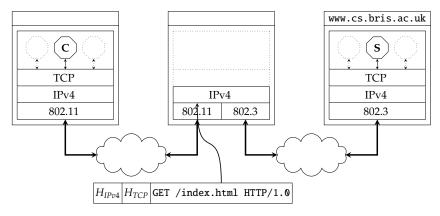
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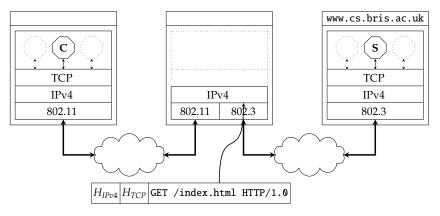
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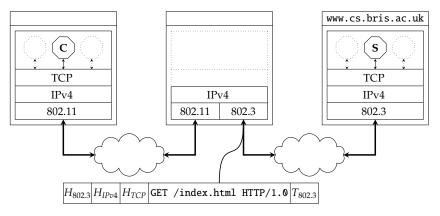
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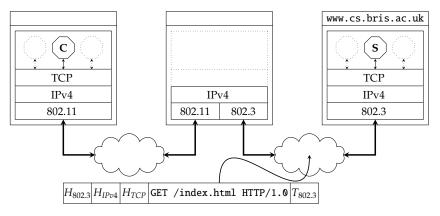
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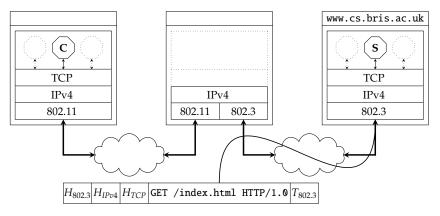
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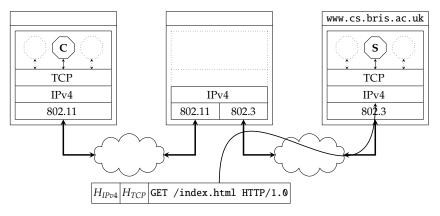
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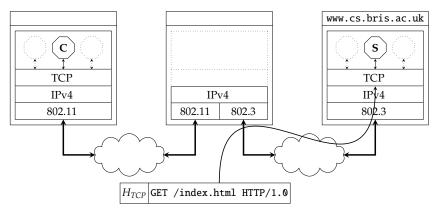
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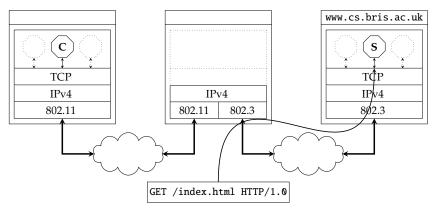
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#### Take away points:

- Although a lot of this content might seem very abstract, the point is that we now have a sensible high-level design.
- We'll fill in missing detail using a bottom-up approach by covering the
  - 1. link layer,
  - 2. internet layer,
  - 3. transport layer, and
  - 4. application layer

while emphasising the *general* concepts as applied in specific technologies ...

- ... keeping in mind the running, motivating example of HTTP.
- A central challenge throughout is design of solutions that work efficiently in the general-case, but guarentee correctness in (many) special- or corner-cases.

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