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Question 1:

a) i) $\{ \Rightarrow \}$ It is not functionally complete

Every formula can be expressed in DNF. Therefore $\{ \vee, \wedge, \neg \}$ is functionally complete. We have to express every connective in the set using \Rightarrow .

$p \vee q \equiv \neg p \Rightarrow q$. ' \vee ' cannot be expressed using ' \Rightarrow ' because we also need ' \neg '.

$$\begin{aligned} p \wedge q &\equiv \neg(\neg p \vee \neg q) \quad (\text{De Morgan law}) \\ &\equiv \neg(p \Rightarrow \neg q) \end{aligned}$$

' \wedge ' cannot be expressed using ' \Rightarrow ' because we also need ' \neg '

$$p \Leftrightarrow q \equiv (p \Rightarrow q) \wedge (q \Rightarrow p)$$

$$\begin{aligned} &\equiv \neg(\neg(p \Rightarrow q) \vee \neg(q \Rightarrow p)) \quad (\text{De Morgan}) \\ &\equiv \neg((p \Rightarrow q) \Rightarrow \neg(q \Rightarrow p)) \end{aligned}$$

\Leftrightarrow cannot be expressed using ' \Rightarrow ' because we need ' \neg '

ii) $\{ \Rightarrow, \neg \}$ We know that $\{ \neg, \vee, \wedge \}$ is functionally

$$p \vee q \equiv \neg p \Rightarrow q \quad \checkmark \quad \text{complete.}$$

$$p \wedge q \equiv \neg(p \Rightarrow \neg q) \quad \checkmark$$

$$p \Leftrightarrow q \equiv (p \Rightarrow q) \wedge (q \Rightarrow p)$$

$$\equiv \neg(\neg(p \Rightarrow q) \vee \neg(q \Rightarrow p)) \quad (\text{De Morgan})$$

$$\equiv \neg((p \Rightarrow q) \Rightarrow \neg(q \Rightarrow p)) \quad \checkmark$$

$\{ \Rightarrow, \neg \}$ is functionally complete because we've expressed $\{ \neg, \vee, \wedge \}$ using just $\{ \Rightarrow, \neg \}$

Question 1

b)

A	B	C	$\neg A$	$\neg B$	$\neg C$	$\neg A \Rightarrow B$	$A \wedge \neg C$	$(A \wedge \neg C) \Leftrightarrow \neg B$	$(\neg A \Rightarrow B) \wedge ((A \wedge \neg C) \Leftrightarrow \neg B)$
T	T	T	F	F	F	T	F	T	T \leftarrow
T	T	F	F	F	T	T	T	F	F
T	F	T	F	T	F	T	F	F	F
T	F	F	F	T	T	T	T	T	T \leftarrow
F	T	T	T	F	F	T	F	T	T \leftarrow
F	T	F	T	F	T	T	F	T	T \leftarrow
F	F	T	T	T	F	F	F	F	F
F	F	F	T	T	T	F	F	F	F

DNF: We find in the last column all entries with the value T. Those are called the minterms

DNF is the disjunction of the minterms:

$$(A \wedge B \wedge C) \vee (A \wedge \neg B \wedge \neg C) \vee (\neg A \wedge B \wedge C) \vee (\neg A \wedge \neg B \wedge \neg C)$$

Since this is not in minimised form yet, is called principal disjunctive normal form.

$\begin{matrix} AB \\ C \end{matrix}$	$\bar{A}\bar{B}$	$\bar{A}B$	AB	$A\bar{B}$
C	F	T	T	F
\bar{C}	F	T	F	T

$$(B \wedge C) \vee (\neg A \wedge B) \vee (A \wedge \neg B \wedge \neg C)$$

We use Karnaugh map to find the minimised form, by observing which variables stay the same

CNF: We find in the last column all entries with the value F. Those are called maxterms

CNF is the conjunction of the maxterms (Conjunction of disjunctions)

$$\neg[(A \wedge B \wedge \neg C) \vee (A \wedge \neg B \wedge C) \vee (\neg A \wedge \neg B \wedge \neg C) \vee (\neg A \wedge \neg B \wedge C)]$$

$$(1A \vee 1B \vee C) \wedge (1A \vee B \vee \neg C) \wedge (A \vee B \vee \neg C) \wedge (A \vee B \vee C)$$

$\begin{matrix} AB \\ C \end{matrix}$	$\bar{A}\bar{B}$	$\bar{A}B$	AB	$A\bar{B}$
C	F	T	T	F
\bar{C}	F	T	F	T

$$\neg[(\neg A \wedge \neg B) \vee (\neg B \wedge C) \vee (A \wedge B \wedge \neg C)]$$

$$\neg[\neg(A \wedge B) \wedge \neg(B \wedge C) \wedge \neg(A \wedge B \wedge \neg C)]$$

$$(A \vee B) \wedge (B \vee \neg C) \wedge (1A \vee 1B \vee C)$$

Question 2

a) Proof by exhaustion

$$n = 5k+1 \quad n^4 - 1$$

$$= (5k+1)^4 - 1$$

$$= 25(5k^2+2k)^2 + 10(5k^2+2k) + 1 - 1 \quad \text{Divisible by 5 by direct proof}$$

$$= 5[5(5k^2+2k)^2 + 2(5k^2+2k)]$$

$$n = 5k+2 \quad n^4 - 1$$

$$= (5k+2)^4 - 1$$

$$= 25(5k^2+4k)^2 + 40(5k^2+4k) + 16 - 1$$

$$= 25(5k^2+4k)^2 + 40(5k^2+4k) + 15$$

Divisible by 5 by direct proof

$$= 5[5(5k^2+4k)^2 + 8(5k^2+4k) + 3]$$

$$n = 5k+3$$

$$n^4 - 1$$

$$= (5k+3)^4 - 1$$

$$= 25(5k^2+6k)^2 + 90(5k^2+6k) + 81 - 1$$

$$= 25(5k^2+6k)^2 + 90(5k^2+6k) + 80$$

Divisible by 5 by direct proof

$$= 5[5(5k^2+6k)^2 + 18(5k^2+6k) + 16]$$

$$n = 5k+4$$

$$n^4 - 1$$

$$= (5k+4)^4 - 1$$

$$= 25(5k^2+8k)^2 + 160(5k^2+8k) + 256 - 1$$

Divisible by 5 by direct proof

$$= 25(5k^2+8k)^2 + 160(5k^2+8k) + 255$$

$$= 5[5(5k^2+8k)^2 + 32(5k^2+8k) + 51]$$

Question 2

b) $P: x+y$ is even

$Q: x$ and y have the same parity

$$P \rightarrow Q \equiv \neg Q \rightarrow \neg P$$

If x and y have different parity, $x+y$ is odd

Case 1: x : odd y : even

$$x: 2k+1 \quad y: 2j$$

$$2k+1+2j = 2(k+j)+1$$

$x+y$ is odd by direct proof

Case 2: x : even y : odd

$$x: 2k \quad y: 2j+1$$

$$2k+2j+1 = 2(k+j)+1$$

$x+y$ is odd by direct proof

Since we've proved that $\neg Q \rightarrow \neg P$ and $\neg Q \rightarrow \neg P \equiv P \rightarrow Q$, we've proved that $P \rightarrow Q$ hold using proof by contrapositive

Question 3

a) Because $(40, 11) = 1$ there exists an inverse

$$40 = 3 \cdot 11 + 7$$

$$11 = 1 \cdot 7 + 4$$

$$7 = 1 \cdot 4 + 3$$

$$4 = 1 \cdot 3 + 1$$

$$1 = 4 - 3$$

$$1 = 4 - (7 - 4)$$

$$1 = 2 \cdot 4 - 7$$

$$1 = 2(11 - 7) - 7$$

$$1 = 2 \cdot 11 - 3 \cdot 7$$

$$1 = 2 \cdot 11 - 3 \cdot 40 + 9 \cdot 11$$

$$1 = \boxed{11} \cdot 11 - 3 \cdot 40$$

$$11^{-1} \equiv 11 \pmod{40}$$

b) $x'' \pmod{41} = 10$ $e = 11$ $p = 41$

$$d \equiv e^{-1} \pmod{p-1}$$

$$d \equiv 11^{-1} \pmod{40}$$

From part a) $d \equiv 11 \pmod{40}$

$$11 \cdot 11 \equiv 1 \pmod{40}$$

$$11 \cdot 11 = 1 + 40k$$

$$(x'')^{11} \equiv x^{1+40k} \equiv x \cdot x^{40k} \equiv x \pmod{41}$$

Using Fermat's Little Theorem

$$a^{p-1} \equiv 1 \pmod{p} \text{ if } p \nmid a$$

$$x = 10'' \pmod{41}$$

$$10^{11} = 10^{1+2+8} = 10 \cdot 10^2 \cdot 10^8$$

$$10^2 = 100 \equiv 18 \pmod{41}$$

$$10^4 \equiv 18^2 \equiv -4 \pmod{41}$$

$$10^8 \equiv (-4)^2 \equiv 16 \pmod{41}$$

$$10^{11} = 10 \cdot 10^2 \cdot 10^8 \equiv 10 \cdot 18 \cdot 16 \equiv 10 \pmod{41}$$

$$\boxed{x = 10}$$

Question 4

$$e=5$$

$$\phi(n) = \phi(pq) = \phi(p) \cdot \phi(q) = (p-1)(q-1)$$

$$n=111 = 3 \cdot 37$$

$$\phi(111) = \phi(3) \cdot \phi(37) = 2 \cdot 36 = 72$$

$$d \equiv e^{-1} \pmod{\phi(n)} \quad ed \equiv 1 \pmod{\phi(n)} \quad ed = 1 + k\phi(n)$$

$$C(M)^d = M^{ed} = M \cdot (M^{\phi(n)})^k \equiv M \pmod{n}$$

Because $(5, 72) = 1$ there is an inverse of e

$$72 = 14 \cdot 5 + 2$$

$$29 = (11101)_2$$

$$5 = 2 \cdot 2 + 1$$

$$1 = 5 - 2 \cdot 2$$

$$1 = 5 - 2 \cdot (72 - 14 \cdot 5)$$

$$1 = 5 - 2 \cdot 72 + 28 \cdot 5$$

$$1 = 29 \cdot 5 - 2 \cdot 72$$

$$d \equiv 29 \pmod{72}$$

$$101^2 \equiv 100 \pmod{111}$$

$$101^4 \equiv 100^2 \equiv 10 \pmod{111}$$

$$101^8 \equiv 10^2 \equiv 100 \pmod{111}$$

$$101^{16} \equiv 100^2 \equiv 10 \pmod{111}$$

$$101^{29} \equiv 101 \cdot 10 \cdot 100 \cdot 10 \equiv 11 \pmod{111}$$

$$81^2 \equiv 12 \pmod{111}$$

$$81^4 \equiv 12^2 \equiv 33 \pmod{111}$$

$$81^8 \equiv 33^2 \equiv 90 \equiv -21 \pmod{111}$$

$$81^{16} \equiv (-21)^2 \equiv -3 \pmod{111}$$

$$81^{29} \equiv 81 \cdot 33 \cdot (-21) \cdot (-3) \equiv 12 \pmod{111}$$

$$25^2 \equiv -41 \pmod{111}$$

$$25^4 \equiv (-41)^2 \equiv 16 \pmod{111}$$

$$25^8 \equiv 16^2 \equiv 34 \pmod{111}$$

$$25^{16} \equiv 34^2 \equiv 46 \pmod{111}$$

$$25^{29} \equiv 25 \cdot 16 \cdot 34 \cdot 46 \equiv 4 \pmod{111}$$

$$32^2 \equiv 25 \pmod{111}$$

$$32^4 \equiv 25^2 \equiv -41 \pmod{111}$$

$$32^8 \equiv (-41)^2 \equiv 16 \pmod{111}$$

$$32^{16} \equiv 16^2 \equiv 34 \pmod{111}$$

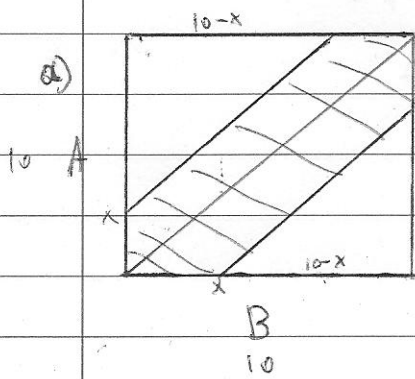
$$32^{29} \equiv 32 \cdot 34 \cdot 16 \cdot (-41) \equiv 2 \pmod{111}$$

The key $n=111$ is a 3-digit number, therefore the block size is 3.

Also in the encoded message there are 33 digits and $3 \mid 33$.

Everything greater than a 3-digit number would be greater than our key which is not possible.

Question 5



$$0.5 = \frac{(10 \cdot 10) - \frac{(10-x)^2}{2} \cdot 2}{(10 \cdot 10)}$$

$$50 = 100 - (100 - 20x + x^2)$$

$$A = B + x$$

$$50 = 100 - 100 + 20x - x^2$$

$$A = B - x$$

$$0 = -x^2 + 20x - 50$$

$$0 = x^2 - 20x + 50$$

$$x_{1,2} = \frac{20 \pm \sqrt{20^2 - 4 \cdot 1 \cdot 50}}{2}$$

$$= \frac{20 \pm 10\sqrt{2}}{2} = 5 \cdot (2 \pm \sqrt{2})$$

$$x_1 = 5(2 + \sqrt{2}) = 17.07$$

$$x_2 = 5(2 - \sqrt{2}) = 2.929$$

$$b) P(128|A) = 0.25 \quad P(A) = \frac{1}{2}$$

$$P(256|A) = 0.75 \quad P(B) = \frac{1}{2}$$

$$P(128|B) = 0.36 \quad P(256) = 0.75 \cdot \frac{1}{2} + 0.64 \cdot \frac{1}{2} = 0.695$$

$$P(256|B) = 0.64 \quad P(128) = 0.25 \cdot \frac{1}{2} + 0.36 \cdot \frac{1}{2} = 0.305$$

E_1 : 256 followed by 128

$$P(E_1) = (0.695) \cdot (0.305) = 0.211975$$

$$P(A|E_1) = \frac{P(E_1|A)P(A)}{P(E_1)} = \frac{(0.75)(0.36) \cdot \left(\frac{1}{2}\right) \cdot \left(\frac{1}{2}\right)}{0.211975} = \frac{2700}{8479} = 0.318$$