COMS10003 Workshop Sheet 7 - outline solutions.

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1. No solution given.

2. If $r_1 = a$ mod c then $a = n_1 c + r_1$, similarly if $r_2 = b$ mod c then $b = n_2 c + r_2$ so the right hand side is $r_1 + r_2$ mod c and the left hand side is $n_1 c + r_1 + n_2 c + r_2$ mod c and we can loose the $(n_1 + n_2)c$ since it's a multiple of c. For the second one the right hand side is $r_1 r_2$ mod c whereas the left hand side is $(r_1 + n_1 c)(r_2 + n_2 c)$ mod $c = (r_1 r_2 + \text{stuff} \cdot c)$ mod $c = r_1 r_2$ mod c.

3. $60 = 3 \cdot 4 \cdot 5$ has 12 factors; it is the smallest number with 12 factors, but 72, 84, 90 and 96 do as well. The idea is to quicky write out the table, but you can guess it is a number with all the low factors, so, for example, under 1000 the best is $840=3 \cdot 5 \cdot 7 \cdot 8$ with 32 factors.

4. Take an example first, $8 = 2^3$ and it is divided by 1, 2, 4 and 8. More generally $p^r / p^s = p^{r-s}$ and the answer is $r + 1$, taking in to account $s = 0$.

5. The lemma about division.

    (a) If $a|b$ then $b = ma$ and if $x|y$ then $y = nx$ so $by = mnax$ and hence $ax|by$.

    (b) So if $a|b$ then $b = ma$ and if $b|c$ then $c = nb$ hence $c = mna$ and $a|c$.

    (c) Well if $a|b$ then $b = na$ and $n \neq 0$ if $b \neq 0$ so $b >= a$.

    (d) In the usual way $b = ma$ and $c = na$ so $bx + cy = xma + yna$ so $a|(bx + cy)$.

6. If one was a prime then the integers wouldn't be a unique factorization domain and lots of theorems would be harder to state.

7. The first 99 values are given at `http://oeis.org/A000010`.

8. If $d|p^n$ it must be in the form p^s, so if $d|a$ then $a = mp$ for some m. Now, to work out the possible values of m, divide p^n by p, giving p^{n-1}.

9. Any number that isn't p is coprime with p so $\phi(p) = p - 1$.

10. From our calculation above, there are p^{r-1} numbers which are not co-prime with p^r so

$$\phi(p^r) = p^r - p^{r-1} = p^r \left(1 - \frac{1}{p}\right) \tag{1}$$

11. This follows from what we have done above

$$\phi(n) = n \prod \left(1 - \frac{1}{p_i}\right) \tag{2}$$

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12. So if $d|a$ and $d|b$ then $d|(a-b)$, conversely, if $d|(a-b)$ and $d|b$ then $d|a$ so a and b and $a-b$ and b have the same common divisors, so the have the same greatest divisor. The important point is you need to argue both ways to show the set of common divisors are the same and not just that one is contained in the other.

13. So the answer if $(\phi(n)-2)/2$; basically, if you skip on to the $i+k$th star each time, if $(n,k) \neq 1$ you get back to where you started without visiting all the points, $k=1$ is explicitely excluded in the question and if $(n,k)=1$ then $(n-k,n)=1$ but that gives the same star, so you have to divide by two.

14. No answer given; you can check your answers at `http://gcd.awardspace.com/`, though of course, it is best to write your own program to do this.

15. It is useful to decide first of all what 'slowest convergence' means; it should probably mean the most steps for the size of the number. Obviously then the slowest convergence will be numbers with gcd one, if $(a,b) = d$ then $(a/d, b/d) = 1$ will take the same number of steps, but with a/d and b/d smaller than a and b is $d \neq 1$. Now, work backwards along the Euclid algorithm so that the numbers are as small as possible. The last step for numbers with gcd one is

$$x_2 = k_1 \cdot 1 + 0 \tag{3}$$

where I am calling the last number in the algorithm before you reach the one and zero x_2; the idea being that $x_1 = 1$ and $x_0 = 0$. Now to make x_2 as small as possible $k_1 = 1$ and $x_2 = 1$. Now go back one step to the step that lead to this one

$$x_3 = k_2 x_2 + 1 \tag{4}$$

and again for x_3 to be as small as possible $k_2 = 1$ giving $x_3 = x_2 + 1 = 2$. Now keep going for as many terms as you are interested in, for the two numbers x_n and x_{n-1} to be as small as possible the k_is are all chosen to be one so that

$$x_n = x_{n-1} + x_{n-2} \tag{5}$$

and so on, that is, the Fibonacci sequence.