

# Loop optimizations

**Most programs spend most time executing loops.**

**So loop optimization is important:**

- **Code motion: avoid recomputing in loop.**
- **Induction variables: reduce number of loop counters.**
- **Loop unfolding: reduce amount of branching.**

# Example

```

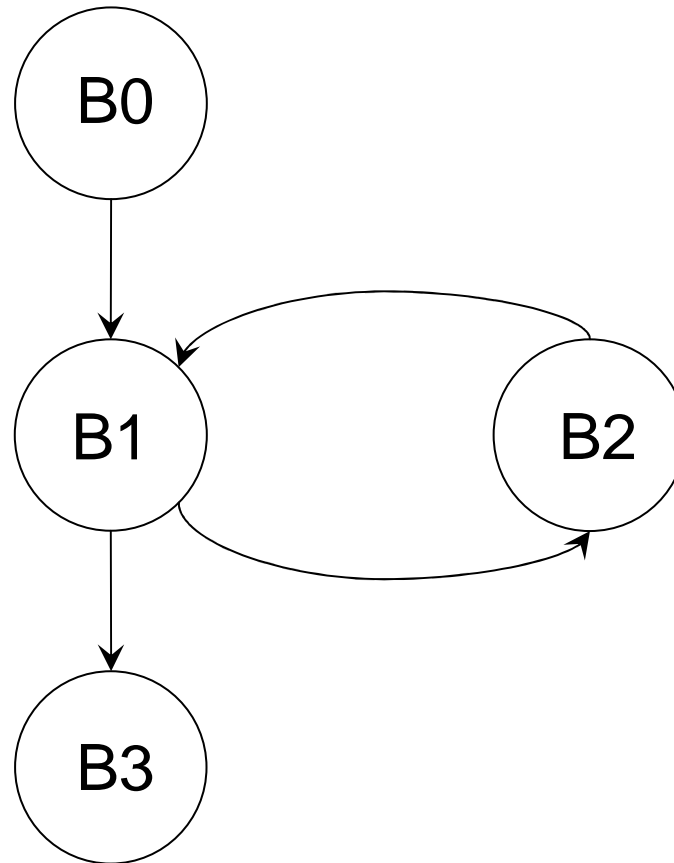
size = 2;
i = 0;
while (i <= size*5-1) {
    sum = sum + a[i];
    i = i + 1;
}
write(sum);

```

## Quadruples:

size = 2 i = 0 <b>goto 3</b>	B0
3: t0 = size * 5 t1 = t0 - 1 if (i <= t1) goto 6 else goto 11	B1
6: t2 = i * 4 t3 = M[a+t2] sum = sum + t3 i = i + 1 goto 3	B2
11: write(sum)	B3

## Flow graph:



# Code motion

Code in a loop that always computes same value can be moved (*hoisted*) to before the loop.

## Example:

### Quadruples

```
size = 2
i = 0
3:  t0 = size * 5
    t1 = t0 - 1
    if (i <= t1) goto 6 else goto 11
6:  t2 = i * 4
    t3 = M[a+t2]
    sum = sum + t3
    i = i + 1
    goto 3
11: write(sum)
```

### Optimized

```
size = 2
i = 0
t0 = size * 5
t1 = t0 - 1
5:  if (i <= t1) goto 6 else goto 11
6:  t2 = i * 4
    t3 = M[a+t2]
    sum = sum + t3
    i = i + 1
    goto 5
11: write(sum)
```

# Loop-invariant computations

*Loop-invariant computation:*

A quadruple

$d: \quad t = a_1 \text{ op } a_2$

where for each  $a_i$ :

- $a_i$  is a constant
- *or* all definitions of  $a_i$  reaching  $d$  are outside the loop
- *or* only one definition of  $a_i$ ,  $e$ , reaches  $d$ , and  $e$  is loop-invariant.

# Hoisting

## Loop-invariant computation

$d: \quad t = a_1 \text{ op } a_2$

can be moved before the loop only if:

1.  $t$  is not live on entry to the loop
2. *and*  $d$  is the only definition of  $t$  in the loop
3. *and*  $d$  “dominates” all loop exits at which  $t$  is live