

- Modern computing devices aren't ad hoc constructions; a rich theory underpins their design and operation.
- Focusing on computer architecture specifically, Boolean algebra is central to more or less everything:
 - 1. in 1840s Boole unified concepts in logic and set theory, predating what we now know as **abstract algebra**, which then
 - 2. enabled Shannon to design analyse and design electrical circuits via logic gates in seminal 1937s work.

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the temperature is $20^{\circ}C$

this statement is false the temperature is too hot

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whose meaning

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- 1. can be evaluated to give a truth value, i.e., false or true,
- 2. must be unambiguous,

A proposition is basically a statement

the temperature is $20^{\circ}C$ the temperature is $x^{\circ}C$ this statement is false the temperature is too hot

- 1. can be evaluated to give a truth value, i.e., false or true,
- 2. must be unambiguous,
- 3. can include free variables, and

A proposition is basically a statement

```
f = the temperature is 20^{\circ}C

g(x) = the temperature is x^{\circ}C

this statement is false

the temperature is too hot
```

- 1. can be **evaluated** to give a **truth value**, i.e., **false** or **true**,
- 2. must be unambiguous,
- 3. can include free variables, and
- can be represented using a short-hand variable or function, whereby free variables must be bound to concrete arguments before evaluation.

► Single statements can be combined using various **connectives**, e.g.,

the temperature is not $20^{\circ}C$

adding parentheses where needed to add clarity, so that

1. "not x" is denoted $\neg x$,

► Single statements can be combined using various **connectives**, e.g.,

¬(the temperature is $20^{\circ}C$)

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1. "not x" is denoted $\neg x$,

► Single statements can be combined using various **connectives**, e.g.,

the temperature is 20°C and it is sunny

- 1. "not x" is denoted $\neg x$,
- 2. "x and y" is denoted $x \wedge y$,

► Single statements can be combined using various **connectives**, e.g.,

(the temperature is $20^{\circ}C$) \land (it is sunny)

- 1. "not x" is denoted $\neg x$,
- 2. "x and y" is denoted $x \wedge y$,

► Single statements can be combined using various **connectives**, e.g.,

the temperature is $20^{\circ}C$ or it is sunny

- 1. "not x" is denoted $\neg x$,
- 2. "x and y" is denoted $x \wedge y$,
- 3. "x or y" is denoted $x \lor y$, and usually called inclusive-or,

► Single statements can be combined using various **connectives**, e.g.,

(the temperature is $20^{\circ}C$) \vee (it is sunny)

- 1. "not x" is denoted $\neg x$,
- 2. "x and y" is denoted $x \wedge y$,
- 3. "x or y" is denoted $x \lor y$, and usually called inclusive-or,

► Single statements can be combined using various **connectives**, e.g.,

either the temperature is 20°C or it is sunny ,but not both

- 1. "not x" is denoted $\neg x$,
- 2. "x and y" is denoted $x \wedge y$,
- 3. "x or y" is denoted $x \vee y$, and usually called inclusive-or,
- 4. "x or y but not x and y" is denoted $x \oplus y$, and usually called exclusive-or,

► Single statements can be combined using various **connectives**, e.g.,

(the temperature is $20^{\circ}C$) \oplus (it is sunny)

- 1. "not x" is denoted $\neg x$,
- 2. "x and y" is denoted $x \wedge y$,
- 3. "x or y" is denoted $x \vee y$, and usually called inclusive-or,
- 4. "x or y but not x and y" is denoted $x \oplus y$, and usually called exclusive-or,

Single statements can be combined using various connectives, e.g.,

the temperature being 20°C implies that it is sunny

- 1. "not x" is denoted $\neg x$,
- 2. "x and y" is denoted $x \wedge y$,
- 3. "x or y" is denoted $x \lor y$, and usually called inclusive-or,
- 4. "x or y but not x and y" is denoted $x \oplus y$, and usually called exclusive-or,
- 5. "x implies y" is denoted $x \Rightarrow y$, and sometimes written "if x then y", and

► Single statements can be combined using various **connectives**, e.g.,

(the temperature is $20^{\circ}C$) \Rightarrow (it is sunny)

- 1. "not x" is denoted $\neg x$,
- 2. "x and y" is denoted $x \wedge y$,
- 3. "x or y" is denoted $x \lor y$, and usually called inclusive-or,
- 4. "x or y but not x and y" is denoted $x \oplus y$, and usually called exclusive-or,
- 5. "x implies y" is denoted $x \Rightarrow y$, and sometimes written "if x then y", and

Single statements can be combined using various connectives, e.g.,

the temperature is 20°C is equivalent to it being sunny

- 1. "not x" is denoted $\neg x$,
- 2. "x and y" is denoted $x \wedge y$,
- 3. "x or y" is denoted $x \lor y$, and usually called inclusive-or,
- 4. "x or y but not x and y" is denoted $x \oplus y$, and usually called exclusive-or,
- 5. "x implies y" is denoted $x \Rightarrow y$, and sometimes written "if x then y", and
- 6. "x is equivalent to y" is denoted $x \equiv y$, and sometimes written "x if and only if y" or "x iff. y".

Single statements can be combined using various connectives, e.g.,

(the temperature is $20^{\circ}C$) \equiv (it is sunny)

- 1. "not x" is denoted $\neg x$,
- 2. "x and y" is denoted $x \wedge y$,
- 3. "x or y" is denoted $x \lor y$, and usually called inclusive-or,
- 4. "x or y but not x and y" is denoted $x \oplus y$, and usually called exclusive-or,
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- 6. "x is equivalent to y" is denoted $x \equiv y$, and sometimes written "x if and only if y" or "x iff. y".
- ► You might also hear more formal terms for these connectives:
 - ► ¬ is often termed logical **compliment** (or **negation**),
 - ► ∧ is often termed logical **conjunction**,
 - v is often termed logical (inclusive) disjunction,
 - ► ⊕ is often termed logical (exclusive) **disjunction**,
 - ightharpoonup \Rightarrow is often termed logical **implication**, and
 - ightharpoonup \equiv is often termed logical **equivalence**.

You can think of the same thing diagrammatically, i.e.,

$$r = \text{(the temperature is } 20^{\circ}\text{C}) \land \text{(it is sunny)}$$

=

the temperature is
$$20^{\circ}C \rightarrow \bigwedge \qquad \uparrow \qquad r$$
 it is sunny $\rightarrow \qquad \uparrow \qquad \uparrow$

but either way, the question is how do we **evaluate** the (compound) proposition (or **expression**) to produce a truth value?

 Since each statement can evaluate to true or false only, we can enumerate the possible outcomes in a truth table, e.g., if

x =the temperature is $20^{\circ}C$

y = it is sunny

 $r = \text{(the temperature is } 20^{\circ}C) \land \text{(it is sunny)}$

then

inp	outs	output
x	у	r
false	false	false
false	true	false
true	false	false
true	true	true

• With n inputs, the truth table will have 2^n rows: each row details the output(s) associated with a given assignment to the inputs.

Definition (logical connectives)

x	у	$\neg x$	$x \wedge y$	$x \vee y$	$x \oplus y$	$x \Rightarrow y$	$x \equiv y$
false	false	true	false	false	false	true	true
false	true	true	false	true	true	true	false
true	false	false	false	true	true	false	false
true	true	false	true	true	false	true	true

Example

Imagine that now

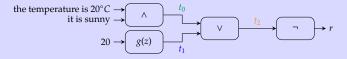
x = the temperature is $20^{\circ}C$

y = it is sunny

g(z) = the temperature is $z^{\circ}C$

 $r = \neg(((\text{the temperature is } 20^{\circ}C) \land (\text{it is sunny})) \lor (\text{the temperature is } z^{\circ}C))$

which we translate into the diagrammatic form



An example evaluation might be as follows:

inputs		int	output		
					$\overline{}$
х	у	t_0	t_1	t_2	r
false	false	false	false	false	true
false	true	false	false	false	true
true	false	false	true	true	false
true	true	true	true	true	false

- If you look closely, some commonalities between propositional logic and other concepts in Mathematics start to emerge ...
- ... among many possibilities, consider *one* example:
 - 1. In **propositional logic**, for some truth value x we have that

$$x \vee \mathbf{false} = x$$

and

$$x \wedge \mathbf{true} = x$$
.

2. In **set theory**, for some set *x* we have that

$$x \cup \emptyset = x$$

and

$$x \cap \mathcal{U} = x$$
.

3. In **elementary algebra**, for some number *x* we have that

$$x + 0 = x$$

and

$$x \cdot 1 = x$$
.

Thou must

- work with the set B = {0,1} of binary digits, using 0 and 1 instead of false and true,
- 2. shorten every statement into either a **variable** *or* **function**,
- use the unary operator ¬ (or NOT) and the binary operators ∧, ∨ and ⊕ (or AND, OR and XOR) to form expressions,
- 4. manipulate said expressions according to some axioms (or rules)

then call the result Boolean algebra.



 One benefit of Boolean axiomatisation is that we can now manipulate expressions without evaluating them, e.g., via

Definition								
Name	Axiom(s)							
commutativity association distribution	$x \wedge y (x \wedge y) \wedge z x \wedge (y \vee z)$	= = =	$y \wedge x$ $x \wedge (y \wedge z)$ $(x \wedge y) \vee (x \wedge z)$	$x \lor y (x \lor y) \lor z x \lor (y \land z)$	=	$y \lor x$ $x \lor (y \lor z)$ $(x \lor y) \land (x \lor z)$		

- ► Why?!
 - 1. to prove equivalence (without brute-force enumeration), or
 - 2. to perform simplification (e.g., less operators).

 One benefit of Boolean axiomatisation is that we can now manipulate expressions without evaluating them, e.g., via

Definition							
Name	Axiom(s)						
identity null idempoten inverse	$ \begin{array}{ccc} $		$x \lor 0$ $x \lor 1$ $x \lor x$ $x \lor \neg x$	= x = 1 = x = 1			

- ► Why?!
 - 1. to prove equivalence (without brute-force enumeration), or
 - 2. to perform simplification (e.g., less operators).

 One benefit of Boolean axiomatisation is that we can now manipulate expressions without evaluating them, e.g., via

Definition							
Name	Axiom(s)						
absorption de Morgan	$x \wedge (x \vee y)$ $\neg (x \wedge y)$	= =	$x \\ \neg x \lor \neg y$	$x \lor (x \land y)$ $\neg (x \lor y)$	≡ ≡	$x \rightarrow x \land \neg y$	

- ► Why?!
 - 1. to prove equivalence (without brute-force enumeration), or
 - 2. to perform simplification (e.g., less operators).

 One benefit of Boolean axiomatisation is that we can now manipulate expressions without evaluating them, e.g., via

Definition	Definition						
Name	Axiom(s)						
equivalence implication involution	$ \begin{array}{cccc} x \equiv y & \equiv & (x \Rightarrow y) \land (y \Rightarrow x) \\ x \Rightarrow y & \equiv & \neg x \lor y \\ \neg \neg x & \equiv & x \end{array} $						

- ► Why?!
 - 1. to prove equivalence (without brute-force enumeration), or
 - 2. to perform simplification (e.g., less operators).

Definition (principle of duality)

The fact there are AND and OR forms of most axioms hints at a more general underlying principle. Consider a Boolean expression e: the **dual expression** e^D is formed by

- 1. leaving each variable as is,
- 2. swapping each ∧ with ∨ and vice versa, and
- 3. swapping each 0 with 1 and vice versa.

Of course e and e^D are different expressions, and clearly not equivalent; if we start with some $e \equiv f$ however, then we do still get $e^D \equiv f^D$.

Example

Consider axioms for

1. distribution, e.g., if

$$e = x \wedge (y \vee z) \equiv (x \wedge y) \vee (x \wedge z)$$

then

$$e^D = x \lor (y \land z) \equiv (x \lor y) \land (x \lor z)$$

and

2. identity, e.g., if
$$e = x \land 1 \equiv x$$

then

$$e^D = x \vee 0 \equiv x$$
.

Definition (principle of complements)

The de Morgan axiom can be turned into a more general principle. Consider a Boolean expression e: the **complement** expression $\neg e$ is formed by

- 1. swapping each variable x with the complement $\neg x$,
- 2. swapping each ∧ with ∨ and vice versa, and
- 3. swapping each 0 with 1 and vice versa.

Example

Consider that if

$$e = x \wedge y \wedge z$$
,

then by the above we should find

$$f = \neg e = (\neg x) \vee (\neg y) \vee (\neg z).$$

Proof:

ı								-
ı	x	у	Z	$\neg x$	$\neg y$	$\neg z$	е	Ĵ
	0	0	0	1	1	1	0	1
ı	0	0	1	1	1	0	0	1
ı	0	1	0	1	0	1	0	1
	0	1	1	1	0	0	0	1
	1	0	0	0	1	1	0	1
ı	1	0	1	0	1	0	0	1
ı	1	1	0	0	0	1	0	1
	1	1	1	0	0	0	1	0

Definition (standard forms)

Consider a Boolean expression:

 When the expression is written as a sum (i.e., OR) of terms which each comprise the product (i.e., AND) of variables, e.g.,

$$(a \wedge b \wedge c) \vee (d \wedge e \wedge f),$$

minterm

it is said to be in **disjunctive normal form** or **Sum of Products (SoP)** form; the terms are called the **minterms**. Note that each variable can exist as-is *or* complemented using NOT, meaning

$$(\neg a \wedge b \wedge c) \vee (d \wedge \neg e \wedge f),$$

minterm

is also a valid SoP expression.

When the expression is written as a product (i.e., AND) of terms which each comprise the sum (i.e., OR) of variables, e.g.,

$$(a \vee b \vee c) \wedge (d \vee e \vee f),$$

maxterm

it is said to be in **conjunctive normal form** or **Product of Sums (PoS)** form; the terms are called the **maxterms**. As above each variable can exist as-is *or* complemented using NOT.

Conclusions

► Take away points:

- 1. In essence, Boolean algebra is a (somewhat) cosmetic extension of what you already know: however, keep in mind that
 - any Boolean function f which can be expressed by a truth table can be computed using a Boolean
 expression, so
 - if we can construct *physical* implementations of NOT, AND and OR we can build something to actually compute f.
- So the pay-off for understanding Boolean algebra is that it explains how computation (and hence computers) work in practice.
- 3. The point is that from here on is that you just learn
 - the truth tables for each Boolean operator, and
 - the set of axioms for Boolean algebra,

and focus on applying the theory to practical challenges and tasks, rather than the theory itself.

References and Further Reading

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