

Part II: Grammars

Submission: On Canvas, under Assignments

Deadline: Week 5, Thurs 16th Feb 2023

Notes on answering this question:

- **Pay attention to what kind of grammar I ask for** (*regular*, or *context free*, or some other type).
- For instance: if I ask for a **regular** grammar and you give me some other kind of grammar, then you may lose marks!
- In particular, if your regular grammar contains a rule of the form $S \rightarrow aby$ or $S \rightarrow D$ or $S \rightarrow \epsilon D$, then this is not a regular grammar and your answer contains an error.
- Always clearly specify the start symbol.
- The alphabet (set of tokens) of a language cannot contain ϵ . ϵ is the empty string, that is, an absence of tokens. If your answer contains something of the form $T = \{\epsilon, \dots\}$ (where T is supposed to be a set of tokens for your language) — then your answer is probably wrong and you're probably losing marks.

Some of you have not met set notation before, so here's a quick tutorial:

- The language determined by the regex $/a^*/$ is $\{a^n \mid n \in \mathbb{N}\}$ or equivalently $\{a^n \mid n \in \{0, 1, 2, \dots\}\}$ or equivalently $\{a^n \mid n \geq 0\}$.
- The language determined by the regex $/(a|b)?/$ is $\{a, b, \epsilon\}$.
- The language determined by the regex $/a^+b^+ /$ is $\{a^m b^n \mid m, n \geq 1\}$ or equivalently $\{a^m b^n \mid m \geq 1, n \geq 1\}$ or equivalently $\{a^m b^n \mid m, n \in \{1, 2, 3, \dots\}\}$.
- The language determined by the English description "any nonzero digit" is $\{1, 2, 3, 4, 5, 6, 7, 8, 9\}$ or equivalently $\{1, 2, \dots, 9\}$ or equivalently $\{n \mid 9 \geq n \geq 1\}$ or equivalently $\{n \mid 0 < n \leq 9\}$.

Now on to the questions:

1. Write the language determined by the regex $/a^*b^*/$

The language determined is $\{a^n b^m \mid n, m \geq 0\}$

2. Write a **regular** grammar to generate the language determined by the regex $/a^*b^*/$

$S ::= a \mid aS \mid aT \mid T$

$T ::= b \mid bT$

3. Write the language determined by the regex $/(ab)^*/$

The language determined is $\{(ab)^n \mid n \geq 0\}$

4. Write a regular grammar to generate the language determined by the regex $/(ab)^*/$

$S ::= \epsilon \mid ab \mid abS$

5. Write the language determined by $/Whiske?y/$. The alphabet is $\{W,h,i,s,k,e,y\}$ (W is a terminal symbol here).

The language determined is $\{Whiske^n y \mid 1 \geq n \geq 0\}$

6. Write a regular grammar to generate the language matched by $/Whiske?y/$

$S ::= Whiskey \mid Whisky$

7. Write a regular grammar to generate decimal numbers; the relevant regex is $/[1-9][0-9]^*(\.[0-9]^*[1-9])?/$.

You may find it useful to use notation resembling $D ::= 0 \mid 1 \mid \dots \mid 9$ to denote an set of ten production rules.

$S ::= A \mid A.BA \mid AB$

$A ::= C \mid CBAB$

$C ::= 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \mid \epsilon$

$B ::= 0 \mid \epsilon$

8. Give a **context free** grammar for the language $L = \{a^n b^n \mid n \in \mathbb{N}\}$.

$S ::= \epsilon \mid aSb$

9. Give a context free grammar for the set of properly bracketed sentences over alphabet $\{\emptyset, (,)\}$. So \emptyset and $((\emptyset))$ are fine, and $\emptyset\emptyset$ and $\emptyset)$ are not fine.

$S ::= \emptyset \mid (S)$

10. Write a context free grammar to generate possibly bracketed expressions of arithmetic with $+$ and $*$ and single digit numbers.

So \emptyset and (\emptyset) and $\emptyset+1+2$ and $(\emptyset+9)$ are fine, and $(\emptyset+)$ and 12 are not fine.

$S ::= \emptyset \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \mid S+S \mid S*S \mid (S)$

11. Give a grammar for palindromes over the alphabet $\{a, b\}$.

$S ::= a \mid b \mid aSa \mid bSb \mid \epsilon$

12. A parity-sequence is a sequence consisting of 0s and 1s that has an even number of ones. Give a grammar for parity-sequences.

$S ::= S1S1S \mid 0S \mid \epsilon$

13. Write a grammar for the set of numbers divisible by 4 (base 10; so the alphabet is $[0-9]$). You might like to read [this page on divisibility testing](#), first. You may use dots notation to represent a sequence, as in for example the sequence Z3,Z6,Z9,...,Z999 to mean “Z followed by some number divisible by three and strictly between 0 and 1000”.

$S ::= 0 \mid 4 \mid 8 \mid A00 \mid A04 \mid A08 \mid A12 \mid A16 \mid \dots \mid A96$
 $A ::= 0A \mid 1A \mid 2A \mid 3A \mid 4A \mid 5A \mid 6A \mid 7A \mid 8A \mid 9A \mid AA \mid \epsilon$

14. Write a grammar for the set of numbers divisible by 3 (base 10; so the alphabet is $[0-9]$). So for example 0, 003, and 120 should be in your language, and 1, 2, and 5 should not

$S \rightarrow 0A \mid 3A \mid 6A \mid 9A \mid 1B \mid 4B \mid 7B \mid 2C \mid 5C \mid 8C$
 $A \rightarrow 0A \mid 3A \mid 6A \mid 9A \mid 1B \mid 4B \mid 7B \mid 2C \mid 5C \mid 8C \mid \epsilon$
 $B \rightarrow 0B \mid 3B \mid 6B \mid 9B \mid 1C \mid 4C \mid 7C \mid 2A \mid 5A \mid 8A$
 $C \rightarrow 0C \mid 3C \mid 6C \mid 9C \mid 1A \mid 4A \mid 7A \mid 2B \mid 5B \mid 8B$

15. (Unmarked) Write a grammar for the set of numbers divisible by 11 (base 10; so the alphabet is $[0-9]$).

More on Grammars

16. State which of the following production rules are *left-regular*, *right-regular*, *left-recursive*, *right-recursive*, *context-free* (or more than one, or none, of these):

1. $Thsi \rightarrow This$ (a rewrite auto-applied by Microsoft Word).

Context Free Grammar

2. $Sentence \rightarrow Subject Verb Object$. (Here Sentence, Subject, Verb, and Object are non-terminal symbols.)

Context Free Grammar

3. $X \rightarrow Xa$.

Left Recursive because X is the non-terminal.

4. $X \rightarrow XaX$.

Context Free Grammar

17. What is the object language generated by $X \rightarrow Xa$ (see Lecture 2)? Explain your answer.

The object language generated by this can be $\{a, aa, aaa, aaaa, \dots\}$ as object language is the set of strings of terminals.

18. Construct context-free grammars that generate the following languages:

1. $(ab|ba)^*$,

$S ::= ab \mid ba \mid abS \mid baS \mid \epsilon$

2. $\{(ab)^n a^n \mid n \geq 1\}$,

$S ::= aba \mid abSa$

3. $\{w \in \{a,b\}^* \mid w \text{ is a palindrome}\}$

$S ::= a \mid b \mid aSa \mid bSb \mid \epsilon$

4. $\{w \in \{a,b\}^* \mid w \text{ contains exactly two bs and any number of as}\}$

$S ::= S'bS'bS'$

$S' ::= aS' \mid \epsilon$

5. $\{a^n b^m \mid 0 \leq n \leq m \leq 2n\}$

$S ::= abS \mid baS \mid abbS \mid bbaS \mid babS \mid \epsilon$

19. Consider the grammar $G = (\{S, A, B\}, \{a, b\}, P, S)$ with productions

$S \rightarrow SAB \mid \epsilon$

$A \rightarrow aA \mid \epsilon$

$B \rightarrow Bb \mid \epsilon$

1. Give a leftmost derivation for aababba.

$S \Rightarrow SAB$

$\Rightarrow SABAB$

$\Rightarrow SABABAB$

$\Rightarrow \epsilon ABABAB$

$\Rightarrow aABABAB$

$\Rightarrow aaABABAB$

$\Rightarrow aa\epsilon BABAB$

$\Rightarrow aaBbABAB$

$\Rightarrow aa\epsilon bABAB$

$\Rightarrow aabaABAB$

$\Rightarrow aaba\epsilon BAB$

$\Rightarrow aabaBbAB$

$\Rightarrow aabaBbbAB$

$\Rightarrow aaba\epsilon bbAB$

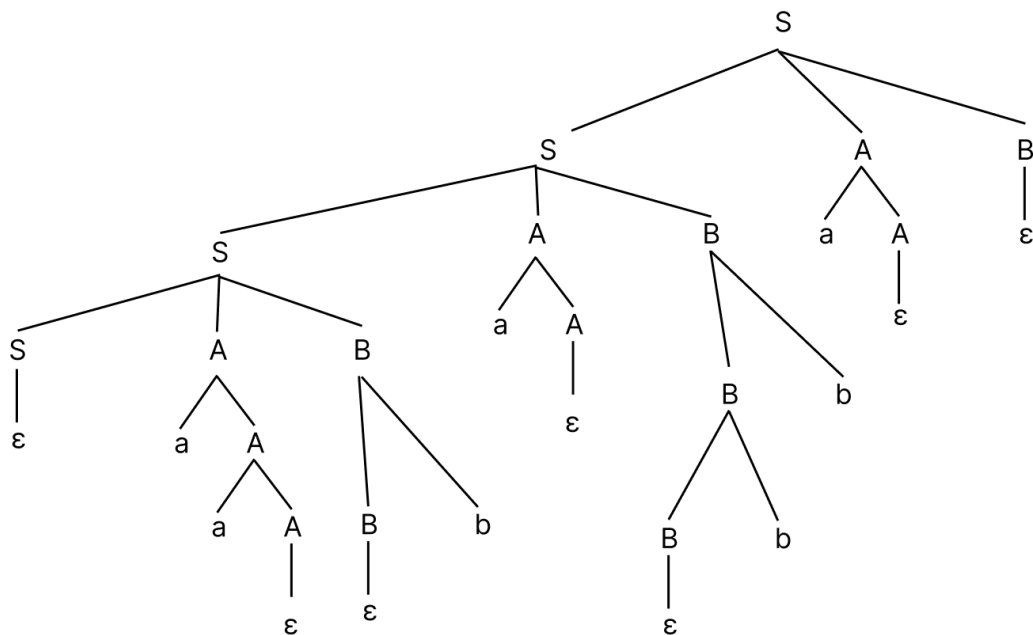
$\Rightarrow aababbaAB$

$\Rightarrow aababba\epsilon B$

$\Rightarrow aababba\epsilon$

$\Rightarrow aababba$

2. Draw the derivation tree corresponding to your derivation.



20. State with proof whether the following grammars are ambiguous or unambiguous:

1. $G = (\{S\}, \{a, b\}, P, S)$ with productions $S \rightarrow aSa \mid aSbSa \mid \epsilon$

The above production rule is ambiguous because when we try to generate string "aababaaa", this can be produced by two methods.

2. $G = (\{S\}, \{a, b\}, P, S)$ with productions $S \rightarrow aaSb \mid abSbS \mid \epsilon$

The above production rule is unambiguous because when we try to generate string "ababbb", this can be produced by only one method.

21. Rewrite the following grammar to eliminate left-recursion:

$\text{exp} \rightarrow \text{exp} + \text{term} \mid \text{exp} - \text{term} \mid \text{term}$

$\text{exp} ::= \text{term exp'}$

$\text{exp'} ::= + \text{term exp'} \mid - \text{term exp'} \mid \epsilon$

22. Write a grammar to parse arithmetic expressions (with + and ×) on integers (such as 0, 42, -1). This is a little harder than it first sounds because you will need to write a grammar to generate correctly-formatted positive and negative numbers; a grammar that generates -01 is unacceptable. You may use dots notation to indicate obvious replication of rules, as in “ $D \rightarrow 0 | \dots | 9$ ”, without comment. (Note that a parser is not an evaluator. This question is **not** asking you to write an evaluator. Also note that the grammar only needs to be unambiguous if the question asks you to provide an unambiguous grammar.)

```
S ::= -AB | AB
B ::= +AB
B ::= -AB
B ::= ε
A ::= 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | AD
D ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9
```

23. Write two distinct parse trees for $2+3*4$ and explain in intuitive terms the significance of the two different parses to their denotation.

Using the Grammar

$\langle \text{exp} \rangle ::= \langle \text{digit} \rangle \mid \langle \text{digit} \rangle + \langle \text{exp} \rangle \mid \langle \text{exp} \rangle * \langle \text{digit} \rangle$

$\langle \text{digit} \rangle ::= 2 \mid 3 \mid 4$

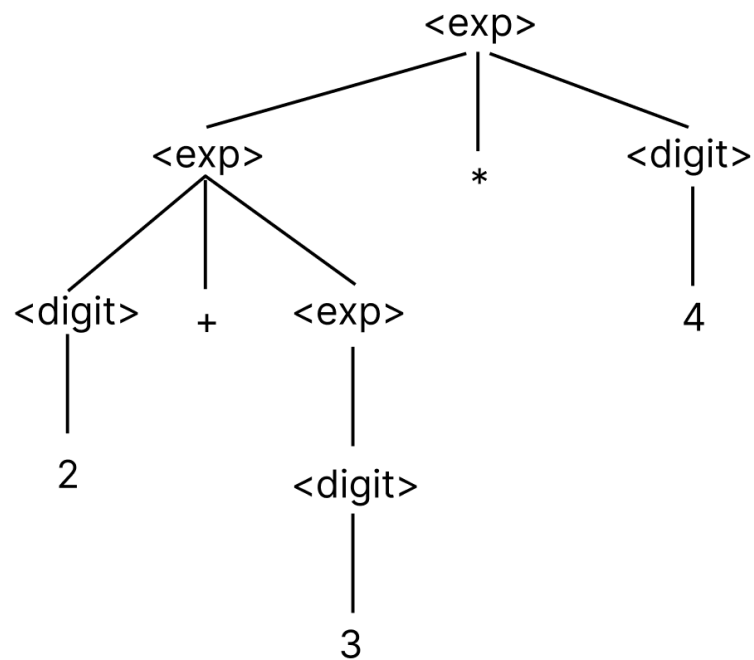


Figure 1: Figure for $2 + 3 * 4$

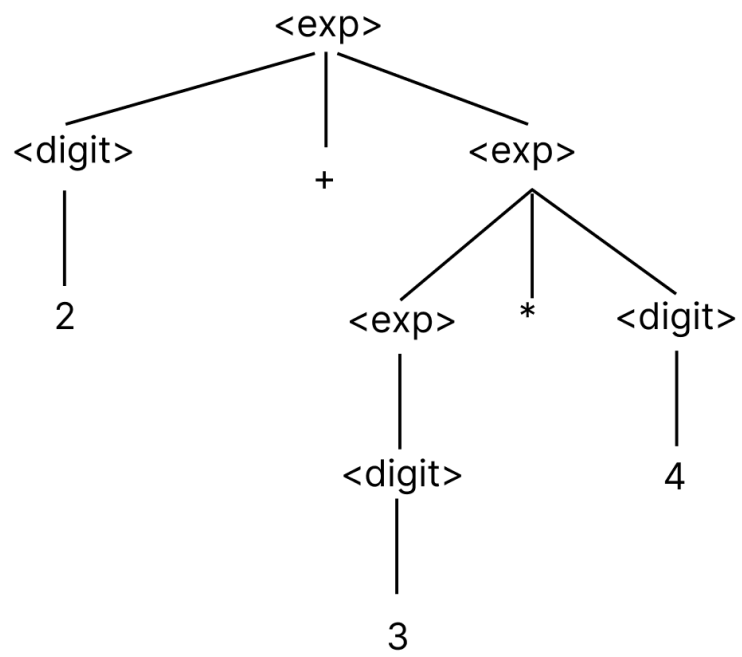


Figure 2: For $2 + 3 * 4$