Little Self-Replicating Programs

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Value.lhs 1

This module contains declarations of the basic types that we'll be using throughout the rest of the code. It also contains a few little helper functions that didn't have better homes. The following language extensions just make things a bit easier, letting us automatically derive a few typeclass instances and implement the MonadState typeclass.

```
{-# LANGUAGE GeneralizedNewtypeDeriving #-}
{-# LANGUAGE MultiParamTypeClasses #-}
module Value (
    Value(..),
    EvalError(..),
    Thread(..),
    WorldState(..),
    throw,
    pause,
    runThread,
    liftRandom,
) where
```

We'll be using transformers to build up the Thread monad, so we need the following imports:

```
import Control.Monad.Identity
import Control.Monad.Except
import Control.Monad.State
import Control.Monad.Coroutine
```

We also need the Rand monad to deal with mutations and random initialization. There is a RandT transformer, but since it's just a wrapper for StateT and since we're already using one of those to keep track of the WorldState, I thought it would be more straightforward to add a random generator to the WorldState and make a liftRandom helper function (see below).

```
import Control.Monad.Random
import System.Random
```

Basically everything is currently implemented with maps, even the universe of cells, which would more reasonably have been implemented as an array. This was just for simplicity. Array.Diff is still experimental, and it would have been annoying to wrap everything in ST or IO, so I just went with maps for everything. Future versions could use more efficient data structures, but since the point of this project was (1) to be a proof of concept and (2) to try to get a parallelization speedup, it seemed fine to have the sequential code be a bit inefficient.

```
import qualified Data. Map as Map
```

The Value type represents both code and data in our small interpreted language. It only has support for integers and functions, and the functions are all fexprs for simplicity.

```
data Value = IntVal Int
             PrimFunc String (Value 
ightarrow Thread Value)
             Lambda Int Value
             Variable Int
             FuncCall Value Value
instance Show Value where
    show (IntVal x) = show x
    show (PrimFunc name _) = name
```

```
show (Lambda var val) =
    "(lambda var:" ++ show var ++ " " ++ show val ++ ")"
show (Variable var) = "var:" ++ show var
show (FuncCall f a) = "(" ++ show f ++ " " ++ show a ++ ")"
```

The following is just a helpful type alias, because there are lots of maps from integers to values.

```
type ValueMap = Map.Map Int Value
```

The following is our error type, which causes the evaluation of a thread to halt prematurely. A different version of this code could have different EvalError constructors to allow for easier inspection of what the code is doing, but for this first version I went with a single constructor.

```
data EvalError = EvalError
```

The WorldState type contains all of the state data a single thread of execution needs in order to operate. The univMap is a read-only map from cell-number to value (which should be thought of as an array), and is the same across all threads. The univSize field is the size of the universe, i.e. the number of cells. The univEdits map contains the current thread's edits to the universe since the last time the different univMaps have been synchronized. When a cell's value is queried, it is first searched for in univEdits, and then in univMap. When a cell's value is written to, it is written in univEdits. The envMap contains the current thread's local scope. This is unique to each thread. This is mainly used for arguments and local variables. The randomGen field is the random generator for the current thread. The generators for different threads are initialized with different random seeds. Finally, the location of the cell the thread is currently evaluating is stored in cellPos.

A Thread is an identity coroutine (meaning it can be paused, but doesn't generate a value until it's finished) that can fail with an EvalError, and always has a WorldState, even if it has failed.

The following is a helper instance making it easier to access the internal WorldState.

```
instance MonadState WorldState Thread where get = Thread  lift  get put = Thread \circ lift \circ put
```

We don't need a MonadError instance since we never need to catch any errors, so this is essentially just the throwError method from the MonadError typeclass.

```
throw :: EvalError \rightarrow Thread a throw = Thread \circ lift \circ throwError
```

As the name suggests, the pause function pauses the current thread.

```
pause :: Thread ()
pause = Thread $ suspend $ Identity $ return ()
```

The runThread function just runs the whole monad transformer and gets it into a form we can work with directly. This is done at every step of execution.

```
type Unwrapped a = (Either EvalError (Either (Thread a) a), WorldState)
runThread :: WorldState → Thread a → Unwrapped a
runThread state (Thread t) =
    unwrapId ∘ runIdentity ∘ flip runStateT state ∘ runExceptT ∘ resume $ t
    where
        unwrapId (Right (Left (Identity t)), s) = (Right $ Left $ Thread t, s)
        unwrapId (Right (Right x), s) = (Right $ Right x, s)
        unwrapId (Left err, s) = (Left err, s)
```

The following is just a helper function that lifts an action from the Rand monad to the Thread monad.

```
\label{eq:liftRandom: Rand StdGen a } \begin{array}{l} \text{ IiftRandom :: Rand StdGen a } \rightarrow \text{ Thread a } \\ \text{ liftRandom rand = do} \\ \text{ state } \leftarrow \text{ get} \\ \text{ let } (\texttt{x}, \texttt{ g}) = \text{runRand rand \$ randomGen state} \\ \text{ put \$ state $\{$ randomGen = g $\}$} \\ \text{ return $x$} \end{array}
```

2 State.lhs

This module contains some helper functions for dealing with the WorldState.

```
module State (
     getVar,
     setVar,
     getCell,
     setCell,
     getCellPos,
     setCellPos,
     getSize,
) where
import Value
import Control.Monad.State
import System.Random
import qualified Data. Map as Map
The getVar function just gets a variable from the local execution scope, or throws an error if it's
not found.
\mathtt{getVar} :: \mathtt{Int} \to \mathtt{Thread} \ \mathtt{Value}
getVar x = do
     \mathtt{state} \, \leftarrow \, \mathtt{get}
     case envMap state Map.!? x of
          \texttt{Just} \ \texttt{y} \ \to \ \texttt{return} \ \texttt{y}
          Nothing \rightarrow throw EvalError
The setVar function just sets a variable in the local execution scope.
\mathtt{setVar} :: \mathtt{Int} \to \mathtt{Value} \to \mathtt{Thread} ()
setVar x v = do
     state \leftarrow get
     put $ state { envMap = Map.insert x v $ envMap state }
```

The getCell function just gets the value of a cell from the universe, looking first in the current thread's edits and then in the read-only univMap. It causes the program to crash if the requested cell is out of bounds.

```
\begin{tabular}{ll} \tt getCell :: Int $\to$ Thread Value \\ \tt getCell $x = do \\ \tt state &\leftarrow get \\ \tt return $ Map.union (univEdits state) (univMap state) Map.! $x$ \\ \end{tabular}
```

The setCell function just sets the value of a cell in the thread's local univEdits.

The following just gets the index of the cell the current thread is evaluating. This is used for some of the locality-sensitive builtin functions.

```
getCellPos :: Thread Int
getCellPos = do
    state ← get
    return $ cellPos state

The following just sets the index of the cell the current thread is evaluating.
setCellPos :: Int → Thread ()
setCellPos x = do
    state ← get
    put $ state { cellPos = x }

The following just gets the size of the universe, i.e. the total number of cells.
getSize :: Thread Int
getSize = do
    state ← get
    return $ univSize state
```

3 Eval.lhs

This module is very short, and only exists because it can't go anywhere else. It only contains one function, the eval function, which evaluates a value in the current thread.

```
module Eval (
     eval,
) where
import Value
import State
\mathtt{eval} \; :: \; \mathtt{Value} \; \to \; \mathtt{Thread} \; \; \mathtt{Value}
Values are all autoquoted, as in Lisp:
eval x@(IntVal _) = return x
eval x@(PrimFunc _ _) = return x
eval x@(Lambda _ _ ) = return x
Evaluating a variable just gets it from the local environment:
```

```
eval (Variable x) = getVar x
```

Evaluating a function is the only time the current thread gets paused, and it gets paused between when the result is evaluated and when it is returned. Attempting to evaluate something that isn't a function kills the thread.

```
eval (FuncCall f a) = do
     \texttt{f'} \leftarrow \texttt{eval} \ \texttt{f}
     case f' of
           PrimFunc \_ g \rightarrow do
                 y \leftarrow g a
                  pause
                  return y
           Lambda x v 
ightarrow do
                 setVar x a
                  y \leftarrow eval v
                 pause
                 return y

ightarrow throw EvalError
```

4 Builtins.lhs

This module contains all the builtin functions that can be used. Hopefully I didn't forget anything that prevents the interpreter from being Turing-complete.

```
module Builtins (
    primFuncs,
) where
import Value
import State
import Eval
There's just one export, the primFuncs export, which is just a list of PrimFuncs.
primFuncs :: [Value]
primFuncs = [macro3 "if" ifFunc,
             macro2 "define" define,
             func1 "peek" peek,
             func2 "poke" poke,
             func2 "+" $ intOp (+),
             func2 "-" $ intOp (-),
             func2 "*" $ intOp (*),
             func2 ">" $ intBoolOp (>),
             func2 "<" $ intBoolOp (<),</pre>
             func2 "=" $ intBoolOp (==),
             func2 "&&" $ boolOp (&&),
             func2 "||" $ boolOp (||),
             func1 "eval" eval,
             func1 "lambda-get-var" lambdaGetVar,
             func1 "lambda-get-val" lambdaGetVal,
             func2 "lambda-set-var" lambdaSetVar,
             func2 "lambda-set-val" lambdaSetVal,
             func1 "funccall-get-func" funcCallGetFunc,
             func1 "funccall-get-arg" funcCallGetArg,
             func2 "funccall-set-func" funcCallSetFunc,
             func2 "funccall-set-arg" funcCallSetArg]
```

The following are just some helper functions that make defining multi-parameter functions and macros easier. The difference is that functions automatically evaluate their parameters, but macros do not. Fundamentally they're both fexprs, the only reason there are both is to reduce repetition in function definitions. There's probably some crazy dependent-type way to make these helpers work for functions of any arity, but since there are only two of each type, it seemed fine to do it by hand.

```
\begin{array}{lll} \texttt{func1} & \texttt{:: String} \, \to \, \texttt{(Value} \, \to \, \texttt{Thread Value)} \, \to \, \texttt{Value} \\ \texttt{func1} & \texttt{name} \, \, \texttt{f} \, = \texttt{PrimFunc name} \, \, \$ \, \, \lambda \texttt{x} \, \to \, \texttt{do} \\ & \texttt{x'} \, \leftarrow \, \texttt{eval} \, \, \texttt{x} \\ & \texttt{f} \, \, \, \texttt{x'} \end{array}
```

The following is just an if macro. There aren't booleans in the language, so positive integers are treated as true and negative ones as false.

```
\begin{array}{lll} \text{ifFunc} & \text{:: Value} \ \rightarrow \ \text{Value} \ \rightarrow \ \text{Thread Value} \\ & \text{ifFunc b thenExpr elseExpr} = \text{do} \\ & \text{b'} \leftarrow \text{eval b} \\ & \text{case b' of} \\ & & \text{IntVal } \ \text{x} \ \rightarrow \ \text{if } \ \text{x} > 0 \\ & & \text{then eval thenExpr} \\ & & \text{else eval elseExpr} \\ & & & \rightarrow \ \text{throw EvalError} \end{array}
```

The following is a macro that sets a local variable (recall that these are unique to each thread of execution).

```
define :: Value → Value → Thread Value
define (Variable x) y = do
    y' ← eval y
    setVar x y'
    return y'
define _ _ = throw EvalError
```

The peek and poke functions read from and write to cells, respectively. The names are a reference to early BASIC machines.

```
peek :: Value → Thread Value
peek (IntVal x) = do
    n ← getSize
    y ← getCellPos
    getCell ((x + y) 'mod' n)
peek _ = throw EvalError

poke :: Value → Value → Thread Value
poke (IntVal x) val = do
    y ← getCellPos
    n ← getSize
    setCell ((x + y) 'mod' n) val
    return val
poke _ _ = throw EvalError
```

The intOp, intBoolOp, and boolOp functions are helpers for defining builtin binary operators on integers and "bools", which are also just integers. The operators that output booleans output 1 for true and 0 for false.

```
intOp _ _ _ = throw EvalError
{\tt intBoolOp} :: (Int 	o Int 	o Bool) 	o Value 	o Value 	o Thread Value
intBoolOp op (IntVal x) (IntVal y) = return $ IntVal $
     if op x y then 1 else 0
intBoolOp _ _ _ = throw EvalError
\texttt{boolOp} \ :: \ (\texttt{Bool} \ \rightarrow \ \texttt{Bool}) \ \rightarrow \ \texttt{Value} \ \rightarrow \ \texttt{Value} \ \rightarrow \ \texttt{Thread} \ \texttt{Value}
boolOp op (IntVal x) (IntVal y) = return $ IntVal $
     if op (x > 0) (y > 0) then 1 else 0
boolOp _ _ _ = throw EvalError
The next few functions are for metaprogramming, allowing the construction and deconstruction of
lambdas and function calls. They should be helpful in allowing actual self-replication.
{\tt lambdaGetVar} \ :: \ {\tt Value} \ \to \ {\tt Thread} \ {\tt Value}
lambdaGetVar (Lambda x _) = return $ Variable x
lambdaGetVar _ = throw EvalError
lambdaGetVal :: Value \rightarrow Thread Value
lambdaGetVal (Lambda _ y) = return y
lambdaGetVal _ = throw EvalError
{\tt lambdaSetVar} \ :: \ {\tt Value} \ \to \ {\tt Value} \ \to \ {\tt Thread} \ {\tt Value}
lambdaSetVar (Lambda _ y) (Variable x) = return $ Lambda x y
lambdaSetVar _ _ = throw EvalError
{\tt lambdaSetVal} \ :: \ {\tt Value} \ \to \ {\tt Value} \ \to \ {\tt Thread} \ {\tt Value}
lambdaSetVal (Lambda x _) y = return $ Lambda x y
lambdaSetVal _ _ = throw EvalError
\texttt{funcCallGetFunc} \; :: \; \texttt{Value} \; \rightarrow \; \texttt{Thread} \; \; \texttt{Value}
funcCallGetFunc (FuncCall f _) = return f
funcCallGetFunc _ = throw EvalError
funcCallGetArg :: Value \rightarrow Thread Value
funcCallGetArg (FuncCall _ a) = return a
funcCallGetArg _ = throw EvalError
\texttt{funcCallSetFunc} \; :: \; \texttt{Value} \; \rightarrow \; \texttt{Value} \; \rightarrow \; \texttt{Thread} \; \; \texttt{Value}
funcCallSetFunc (FuncCall _ a) f = return $ FuncCall f a
funcCallSetFunc _ _ = throw EvalError
{	t funcCallSetArg}:: {	t Value} 
ightarrow {	t Value} 
ightarrow {	t Thread} {	t Value}
funcCallSetArg (FuncCall f _) a = return $ FuncCall f a
funcCallSetArg _ _ = throw EvalError
```

intOp :: (Int \rightarrow Int \rightarrow Int) \rightarrow Value \rightarrow Value \rightarrow Thread Value intOp op (IntVal x) (IntVal y) = return \$ IntVal \$ op x y

5 Mutate.lhs

This module contains actions that mutate contents of cells and generate new random values. The actions are all in the Rand monad instead of in the Thread monad because they only need the random generator, not any other part of the WorldState, so it would have been overkill to give them an entire Thread coroutine. Also because they're not just used from within the Thread monad, but are also used for initializing the universe at the beginning of execution.

```
module Mutate (
    mutate,
    randomValue,
) where
import Value
import State
import Builtins
import System.Random
import Control.Monad.Random
The following is just a helper type alias for the monad we're going to be doing everything in.
type RandM = Rand StdGen
The following represents the probability that any mutation is going to occur this step.
mutateP :: Double
mutateP = 0.01
The following represents the probability that the parameter of a Lambda, as opposed to its body,
will be mutated.
mutateParP :: Double
mutateParP = 0.2
The following is the probability that a function, as opposed to its argument, will be mutated in a
function call.
mutateFuncP :: Double
mutateFuncP = 0.3
The following is the probability that an entirely new random value will be generated, as opposed
to the differential modifications that are otherwise performed.
mutateTypeP :: Double
mutateTypeP = 0.1
When an int is mutated, it is randomly incremented or decremented with equal probability.
\mathtt{mutateInt} \; :: \; \mathtt{Int} \; \rightarrow \; \mathtt{RandM} \; \; \mathtt{Int}
mutateInt x = do
    \texttt{b} \leftarrow \texttt{getRandom}
    return $ if b then x + 1 else x - 1
When a new integer is generated, it is selected from the range [-5, 5].
randInt :: RandM Int
randInt = getRandomR (-5, 5)
```

```
randIntVal :: RandM Value
randIntVal = randInt >>= return o IntVal
```

When a new primitive function is generated, it is selected at random from the list of primitive functions.

```
randPrimFunc :: RandM Value
randPrimFunc = do
   i ← getRandomR (0, length primFuncs - 1)
   return $ primFuncs !! i
```

When a new lambda is generated, its parameter is a random integer from the range [-5, 5], and its body is a randomly generated value.

```
randLambda :: RandM Value
randLambda = do
    x ← randInt
    v ← randomValue
    return $ Lambda x v
```

When a variable is generated, it is selected at random from the range [-5, 5].

```
randVariable :: RandM Value
randVariable = randInt >>= return o Variable
```

When a function call is generated, both the function and the argument are randomly generated values. If the function is an integer this will result in the thread crashing, but that's sufficiently low probability that it wouldn't have been worth making a separate random function generator.

```
\begin{tabular}{ll} randFuncCall :: RandM Value \\ randFuncCall = do \\ f \leftarrow randomValue \\ a \leftarrow randomValue \\ return \$ FuncCall f a \\ \end{tabular}
```

The mutateInplace function just puts together all of the above probabilities and mutators and applies them to an arbitrary value.

```
{\tt mutateInplace} :: {\tt Value} \to {\tt RandM} {\tt Value}
mutateInplace (IntVal x) = mutateInt x >>= return o IntVal
mutateInplace (PrimFunc _ _) = randPrimFunc
mutateInplace (Lambda x v) = do
    b \leftarrow getRandom
    if b < mutateParP then do
         x' \leftarrow mutateInt x
         return $ Lambda x' v
    else do
         v' \leftarrow mutateInplace v
         return $ Lambda x v'
mutateInplace (Variable x) = mutateInt x >>= return o Variable
mutateInplace (FuncCall f a) = do
    b \leftarrow getRandom
    if b < mutateFuncP then do
         f' \leftarrow mutateInplace f
         return $ FuncCall f' a
    else do
         a' \leftarrow mutateInplace a
         return $ FuncCall f a'
```

When a new random value is required, its type is selected at random from the 5 different possible types (ints, primitive functions, lambdas, variables, and function calls), with equal probability. A different version of the code could have different probabilities for each of the types, but there are already a lot of parameters to deal with, and it's not clear how these probabilities should deviate from uniform for better performance.

```
\begin{array}{lll} {\rm randomValue} & {\rm :: \ RandM \ Value} \\ {\rm randomValue} & = {\rm do} \\ {\rm b} & \leftarrow {\rm \ getRandomR} \ (0, \ 4) \\ {\rm case} & {\rm b} & {\rm :: \ Int} \ {\rm of} \\ & 0 & \rightarrow {\rm \ randIntVal} \\ & 1 & \rightarrow {\rm \ randPrimFunc} \\ & 2 & \rightarrow {\rm \ randLambda} \\ & 3 & \rightarrow {\rm \ randVariable} \\ & 4 & \rightarrow {\rm \ randFuncCall} \end{array}
```

The mutateValue function just generates a new random value with probability mutateTypeP, and mutates the existing value otherwise.

```
\begin{array}{ll} \texttt{mutateValue} :: \texttt{Value} \to \texttt{RandM} \ \texttt{Value} \\ \texttt{mutateValue} \ x = \texttt{do} \\ \texttt{b} \leftarrow \texttt{getRandom} \\ \texttt{if} \ \texttt{b} < \texttt{mutateTypeP} \ \texttt{then} \ \texttt{randomValue} \\ \texttt{else} \ \texttt{mutateInplace} \ x \end{array}
```

When we mutate a thread, we first check whether any mutations are going to occur (probability mutateP), and if so, we pick a random cell to mutate, and then mutate it.

```
mutate :: Thread ()
mutate = do
    b ← liftRandom getRandom
    when (b < mutateP) $ do
        n ← getSize
        i ← liftRandom $ getRandomR (0, n - 1)
        x ← getCell i
        x' ← liftRandom $ mutateValue x
        setCell i x'</pre>
```

6 Rep.lhs

```
module Rep (
    runStep,
    runN,
) where
import Value
import State
import Eval
import Mutate
{\tt import\ Control.Monad.Random}
import qualified Data. Map as Map
randomThread :: Thread Value
randomThread = do
    n \leftarrow getSize
    i \leftarrow liftRandom \$ getRandomR (0, n - 1)
    setCellPos i
    \texttt{cell} \, \leftarrow \, \texttt{getCell} \, \, \texttt{i}
    eval cell
runStep :: ([WorldState], [Thread Value]) → ([WorldState], [Thread Value])
runStep (states, threads) = (states'', threads'') where
    (threads', states') = unzip
         [runThread s (mutate >> t) | (s, t) \leftarrow zip states threads]
    restartThread (Left err) = randomThread
    restartThread (Right (Left t)) = t
    restartThread (Right (Right _)) = randomThread
    threads' = map restartThread threads'
    univ = univMap $ head states
    univ' = Map.union (Map.unions $ map univEdits states') univ
    updateState state = state { univMap = univ', univEdits = Map.empty }
    states' = map updateState states'
initialize :: Int 	o Int 	o Int 	o ([WorldState], [Thread Value])
initialize nCells nThreads seed = (states, threads) where
    rand = do
        \texttt{cells} \, \leftarrow \, \texttt{sequence} \, \, \texttt{\$ replicate nCells randomValue}
        seeds ← sequence $ replicate nThreads getRandom
        return (cells, seeds)
    (cells, seeds) = evalRand rand $ mkStdGen seed
    univ = Map.fromList $ zip [0..] cells
    makeState s = WorldState { univMap = univ,
                                  univSize = nCells,
                                  univEdits = Map.empty,
                                  envMap = Map.empty,
                                  randomGen = mkStdGen s,
                                  cellPos = 0 }
    states = [makeState s | s \leftarrow seeds]
    threads = replicate nThreads randomThread
```