Scala for TAPL'ers 3

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Parsing

Parsing

- Check whether input has a given structure
 - Typically do more than just checking
 - Produce AST
- Acceptable structure is defined by a grammar

Grammar (in BNF)

```
term = '\' ident '.' term
| ident
| term term
```

Recursive descent

```
def ident(input: String): Boolean = ...
def term(input: String): Boolean
  = (input(0) == '\\' && ident(input.substring(1)) &&
      input(i) == '.' && term(input.substring(i+1)))
      Il ident(input)
      Il term(input) && term(inputRest)
```

Recursive descent

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def ident(input: String): Boolean = ...
def term(input: String): Boolean
    = (input(0) == '\\' && ident(input.substring(1)) &&
        input(i) == '.' && term(input.substring(i+1)))
        Il ident(input)
        Il term(input) && term(inputRest)
```

Don't know i and inputRest
Can you spot other (more serious) problems?

Essence of this problem:

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twoIdents = ident ident

def twoIdents(input: String) = ident(input) && ident(rest)
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Beside success, track how much input was consumed:

```
def twoIdents(input: String): Option[String]
```

Parsing failed: result is None (false in previous slide) else, result is Some(x) where x is the rest of the input

```
twoIdents = ident ident

def twoIdents(input: String) = ident(input) && ident(rest)
```

```
twoIdents = ident ident

def twoIdents(input: String) = seq(ident, ident)(input)
```

Factoring out input-management

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Factoring out input-management

```
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```

This is the essence of combinator parsing!

Censored

Combinator Parsers

- A parser is first-class: a function (and thus an object)
- Build complex parsers by combining simpler ones
- Combinators are just methods that take parsers and produce a new one

Parser

One more piece of information: result value

```
trait Result[+T]

case class Success[+T](result: T, rest: Input) extends Result[T]
case class Failure(errMsg: String, rest: Input) extends Result[Nothing]
```

Parser with seq

```
abstract class Parser[+T] extends (Input => Result[T]) {
 def ~ [U](p: Parser[U]): Parser[Pair[T, U]]
   = Parser { in: Input =>
      this(in) match {
        case Success(x, rest) =>
         p(rest) match {
           case Success(y, rest2) => Success((x, y), rest2)
           case Failure(e, r) => Failure(e, r)
        case Failure(e, r) => Failure(e, r)
```

Alternation

```
def | [U >: T](p: Parser[U]): Parser[U]
= Parser { in: Input =>
        this(in) match {
        case s@Success(_, _) => s
        case _ => p(in)
      }
}
```

Alternation

```
def | [U >: T](p: Parser[U]): Parser[U]
= Parser { in: Input =>
        this(in) match {
        case s@Success(_, _) => s
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```

```
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def | [U >: T](p: Parser[U]): Parser[U]
= Parser { in: Input =>
        this(in) match {
        case s@Success(_, _) => s
        case _ => p(in)
      }
}
```

```
def term =
    ( '\\' ~ ident ~ '.' ~ term
    | ident
    | term ~ term
    )
```

Implicit accept

Implicit accept

```
def term = accept('\\').~(ident).~(accept('.')).~(term).l
(ident).l(term.~(term))
```

Problem 2: Ordered Choice

```
term = '\' ident '.' term
| ident
| term term
```

- In BNF, the order of the alternatives does not matter.
- With our implementation of |, it does!
 - Commits to first successful branch

Problem 2: Ordered Choice

```
def term =
   ( term ~ term
   | '\\' ~ ident ~ '.' ~ term
   | ident
   )
```

- In BNF, the order of the alternatives does not matter.
- With our implementation of |, it does!
 - Commits to first successful branch

```
def term =
   ( term ~ term
   | '\\' ~ ident ~ '.' ~ term
   | ident
   )
```

term is recursive

- term is recursive
 - harmless
 - left-recursion

Harmless Cycles

```
def term =
   ( term ~ term)
   | '\\' ~ ident ~ '.' ~ term)
   | ident
   )
```

```
def ~ [U](p: => Parser[U]): Parser[Pair[T, U]]
def | [U >: T](p: => Parser[U]): Parser[U]
```

 combinators ~ and | take call-by-name arguments

Left-recursion

```
def term =
   (term ~ term
   | '\\' ~ ident ~ '.' ~ term
   | ident
   )
```

- Resulting parser immediately tries itself on the same input
 - (ordering due to semantics of |)
- CBN would not solve anything

Rewriting left-recursion

```
def term0 =
    ( '\\' ~ ident ~ '.' ~ term
    | ident
    )

def term = rep1(term0)
def rep1[T](p: Parser[T]) = p ~ rep1(p) | p
```

Lifting

```
def ^{(V)}[U](f: T \Rightarrow U): Parser[U]
  = Parser { in: Input =>
       this(in) match {
          case Success(x, r) \Rightarrow Success(f(x), r)
          case Failure(e, r) => Failure(e, r)
def ~> [U](p: => Parser[U]): Parser[U]
  = this \sim p \land \land \{case(x, y) \Rightarrow y\}
def <~ [U](p: => Parser[U]): Parser[T]
  = this \sim p \land \land \{case(x, y) \Rightarrow x\}
```

Producing an AST

```
def term0 : Parser[Term] =
  (('\\' ~> ident) ~ ('.' ~> term) ^^ {case (i, b) => Abs(i, b)}
  I ident ^{n} = Var(n)
def term = chainl1(term0, ' ' ^{^{\prime}} {x => App(_: Term, _: Term)})
def ident: Parser[String] = ...
class Term
case class Var(name: String) extends Term
case class Abs(name: String, body: Term) extends Term
case class App(fun: Term, arg: Term) extends Term
```

Lexical Parsing

- Typically, parsing happens in phases:
 - lexical parsing: take stream of characters and produce stream of tokens
 - syntactical parsing: take stream of tokens, produce AST, ...
- Lexical parsing
 - Take care of low-level issues
 - Does not need full power of parsing

Integrated lex'ing

```
def term0 : Parser[Term] =
   (('\\' ~> wss(ident)) ~ ('.' ~> wss(term)) ^^ {
                                      case (i, b) \Rightarrow Abs(i, b)
   I ident ^{n} = Var(n)
def term = chainl1(term0, ws ^{x} = App(_: Term, _: Term))
def ident = rep1(letter, letter | digit) ^^ {_.mkString("")}
def letter = acceptIf(_.isLetter)
def digit = acceptIf(_.isDigit)
def ws = rep1(accept(' '))
def wss[T](p: Parser[T]): Parser[T] = opt(ws) ~> p <~ opt(ws)</pre>
```

Homework

- Get to know the combinators
 - Report&code available (will keep improving it until project assignment is released)
 - Experiment!
 - Suggestion: make your own language for screen scraping:
 - rep(row containing {bold.class("title") ~ div.class("descr")})