EDF PROJECT REPORT

Analytical methods calculations	
Hyperperiod2	
CPU load3	
Schedulability using URM3	
By Time demand analysis4	
Simso offline simulator5	
Configuration5	
Tasks5	
Results6	
Chart simulation7	
Keil simulator in run-tim	10
Implementing trace and Tick macro hooks10	
Results15	
Logic analyzer16	
Using FreeRTOS Statistics23	
Conclusion25	

1. Using analytical methods calculate the following for the given set of tasks:

Calculate the system hyperperiod

Our system has 6 tasks as follow

Task	Period	Execution time	Deadline
Button_1_Monitor	50	0.0233	50
Button_2_Monitor	50	0.0233	50
Periodic_Transmitter	100	0.16	100
Uart_Receiver	20	WCE = 4.92	20
Load_1_Simulation	10	5.01	10
Load_2_Simulation	100	12.075	100

These execution times was extracted from real time simulation from both trace hooks counter calculations and logic analyzer of (GPIO) pins.

According to periodicity we can see that all periods can be a fraction of the largest periodic task .

Largest periodic function is 100 ms.

Button_1_Monitor and Button_2_Monitor occur 2 times in 100 ms.

Uart_Receiver occur 5 times in 100 ms.

Load_1_Simulation occur 10 times in 100 ms.

So the hyperperiod is 100 ms in which all tasks have to execute at this time.

Calculate the CPU load:

	Occurrence in	Execution time	Total execution	Percentage
	100 ms		time in hyper	
			period	
Button_1_Monitor	2	0.0233	0.0466	0%
Button_2_Monitor	2	0.0233	0.0466	0%
Periodic_Transmitter	1	0.16	0.16	0%
Uart_Receiver	5	WCE = 4.9	24.5	24.5%
		ms		
Load_1_Simulation	10	5	50	50%
Load_2_Simulation	1	12	12	12%
Total			In WCE = 87	87%

Note* this calculation depends on the worst case execution time (WCET) as the function Uart_Receiver not always takes 4.9 ms at the worst case the function will wait until all queue "3" messages will be sent ,and if there is no messages in QUEUE it will not wait to send all messages and it will return immediate

During real time simulation the average CPU load is 73%.

Check system schedulability using URM and time demand analysis techniques (Assuming the given set of tasks are scheduled using a fixed priority rate -monotonic scheduler)

1- By URM

The total utilization (U) = Sum(Execution time / Period)
=
$$(0.0233/50) + (0.0233/50) + (0.16/100) + (4.9/20) + (5/10) + (12/100)$$

= $0.000466 + 0.000466 + 0.0016 + 0.245 + 0.5 + 0.12 = 0.86753$ (it is < 1)

URM (Ulub) =
$$6(2^{(1/6)} - 1) = 0.734$$

According to the calculation the total utilization time (0.86799) > URM (0.734) so the system (may or may not be schedulable) and cannot be grantee as schedulable and we need more analysis

2- By Time demand analysis

Schedulable system is grantee when Time provided > Time demand

Time Provided is the hyperperiod = 100 ms

The time demand for each task according to priority "However priority depends on deadline"

1- Load 1 Simulation

W(10)(Load_1_Simulation) = 5 + 0 = 5 ms... < 10 ms so the system is feasible for Load_1_Simulation task

2- Uart_Receiver

W(20)(Uart_Receiver) = 4.9 + (20/10)*5 = 14.9ms <20ms so the system is feasible for Uart_Receiver task

3- Button_1_Monitor

W(50)(Button_1_Monitor) = 0.0233+(50/50)*0.0233+(50/10)*5+(50/20)*4.9 = 0.0466+5*5+2.5->3*4.9 = 39.466 ms.... < 50 ms so the system is feasible for Button 1 Monitor task

4- Button 2 Monitor

W(50)(Button_2_Monitor) = 0.0233+(50/50)*0.0233+(50/10)*5+(50/20)*4.9 = 0.0466+5*5+2.5->3*4.9 = 39.466 ms.... < 50 ms so the system is feasible for Button 2 Monitor task

5- Periodic_Transmitter

 $W(100)(\ \text{Periodic_Transmitter}) = 0.16 + (100/100)*12 + (100/50)*0.0233 + (100/50)*0.0233 + (100/10)*5 + (100/20)*4.9 = 0.16 + 12 + 0.0466 + 0.0466 + 50 + 24.5 = 86.753 \text{ms}.... < 100 \text{ ms} \quad \text{so the system is feasible for Periodic_Transmitter task}$

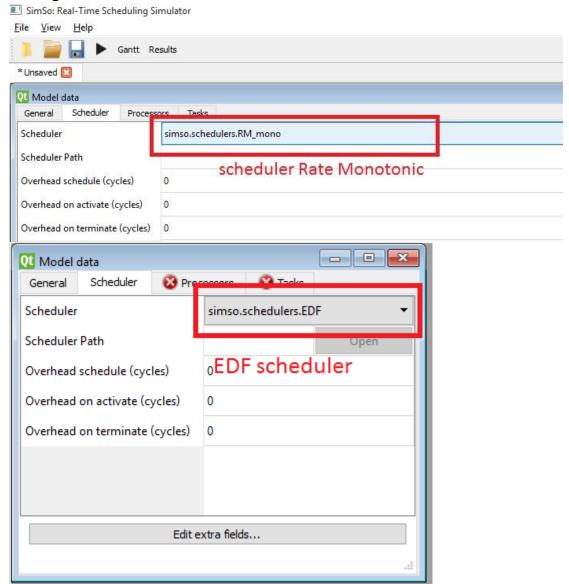
6- Load 2 Simulation

 $W(100)(\ Load_2_Simulation) = 12 + (100/100)*0.16 + (100/50)*0.0233 + (100/50)*0.0233 + (100/10)*5 + (100/20)*4.9 = 12 + 0.16 + 0.0466 + 0.0466 + 50 + 24.5 = 86.753\ ms.... < 100\ ms \ so\ the$ system is feasible for Load_2_Simulation task

In conclusion by analytic method the system is schedulable

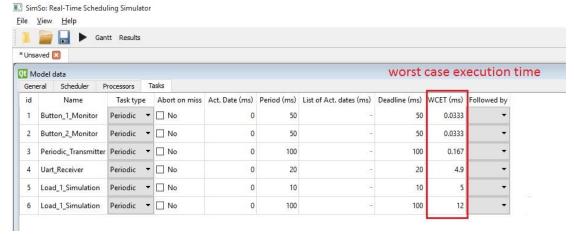
2. Using Simso offline simulator, simulate the given set of tasks assuming: Fixed priority rate monotonic scheduler

1- Configuration



The scheduler we choose either Rate Monotonic or EDF both schedulers give the same results.

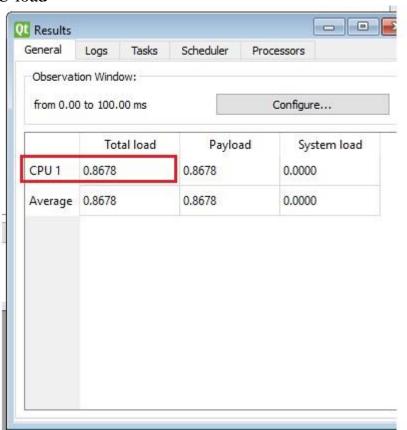
2- Tasks



We entered the task parameters as shown and we entered worst execution time "WCET" according to real time simulation (data extracted from keil simulation from logic analyser mainly and other methods so please refer to section >> 3. Using Keil simulator in run-time and the given set of tasks)

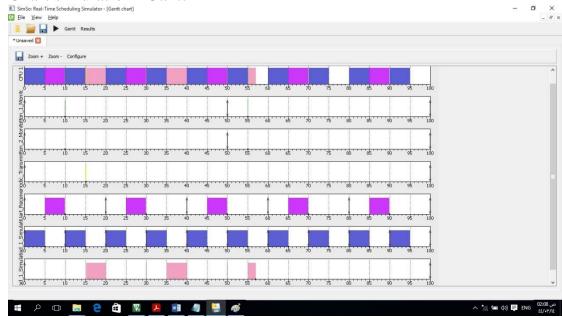
3- Result

a- CPU load



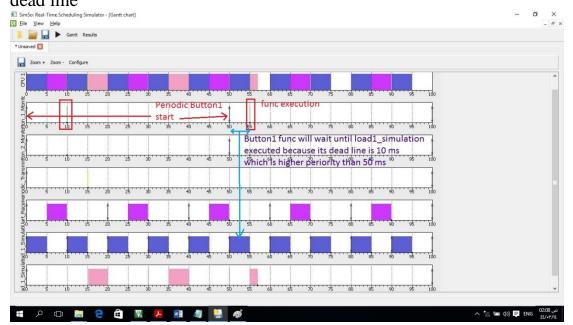
As we can see that CPU load is in the range of 86.8% in WCET of all tasks.

b- Chart simulation "Gantt"



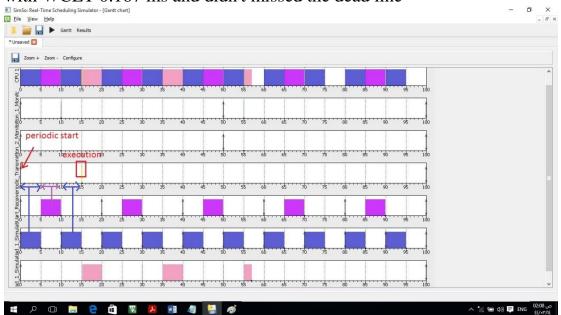
As we can see all tasks don't break their deadlines

*As regard Button1 and Button2 tasks they are very short period occurring 2 times in 100 ms scale and didn't breaks its dead line



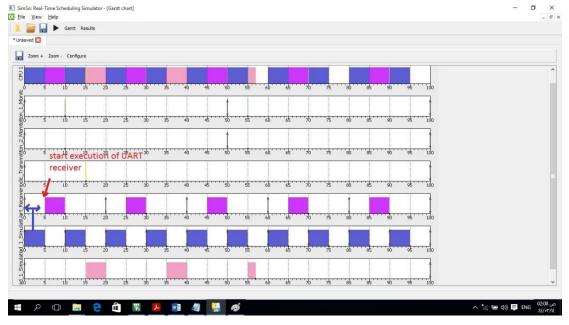
From the chart its execution will be delayed until load1simulation task end its work which will take 5 ms because its dead line is 10 ms < 50 ms of the button task deadline.

*as regard periodic_transmitter it occurs at period of 100 ms with WCET 0.167 ms and didn't missed the dead line

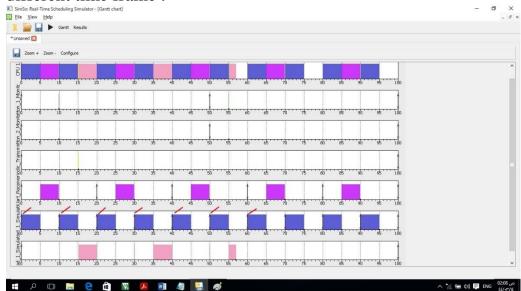


As we can see the periodic transmitter func execution will be delayed for 15 ms this is because its deadline is 100 ms seconds so it will be delayed until load1 func (10 ms deadline) executed and finished then UART receiver func (20 ms dedline) executed and fished, again load1 func.

*as regard UART receiver func (20 ms deadline) it will be executed immediately (5 ms) after load1 func (10 ms deadline) because the priority depends on deadline and it didn't break its deadline

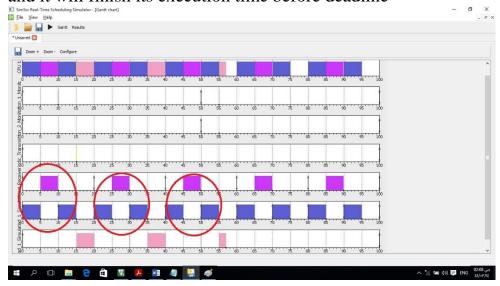


*as regard load1 func (10 ms deadline) it will preempt any other function and will delay all functions until its execution and its very clear preemption of the load2 func (100 ms deadline) at different time frame.



As we can see no delay in execution of the function every time because it is the least deadline

*as regard load2 func (100 ms deadline and execution time of 12 ms) it will be delayed and preempted by all other functions and it will finish its execution time before deadline



In conclusion according to SIMSO simulation program all tasks didn't break their deadlines and the nearest deadline will preempt other tasks with high cpu load (86.8%) and the system is schedulable.

- 3. Using Keil simulator in run-time and the given set of tasks:
 - 1- Implementing trace and Tick macro hooks and GPIO

In the FreeRTOSConfig.h file

```
main.c serial.c GPIO.h GPIO.c tasks.c GPIO_cfg.c FreeRTOSConfig.h task.h list.c list.h po
       * See http://www.freertos.org/a00110.html
  46
      #define configUSE_PREEMPTION
  49
      #define configUSE IDLE HOOK
  50 #define configUSE TICK HOOK
      #define configCPU_CLOCK_HZ
#define configTICK_RATE_HZ
                                        unsigned long ) 600000000 ) /\star =12.0MHz xtal multiplied by 5 usin
  52
                                     ( ( TickType_t ) 1000 )
      #define configMAX PRIORITIES
      #define configMINIMAL_STACK_SIZE ( ( unsigned short ) 90 )
      #define configTOTAL HEAP SIZE ( ( size t ) 13 * 1024 )
#define configMAX TASK NAME LEN ( 8 )
  #define configUSE_16_BIT_TICKS
     #define configIDLE_SHOULD_YIELD
                                      200 // this is the idle period and must be larger than any other
      #define IDLE_PERIOD
```

We will set the macros configUSE_TICK_HOOK and configUSE_TRACE_FACILITY to 1 to enable idle hook and traceability

Also we will set the (configUSE_APPLICATION_TASK_TAG) to 1 to give tasks name tags to trace each task

```
78 | #define traceTASK_SWITCHED_IN()
79 =
                                                                       {switch((int)pxCurrentTCB->pxTaskTag){\
                                                                          case Button_1_Monitor_Tag:\
    GPIO_write(TaskMonitor,Button_1_Monitor_PIN,PIN_IS_HIGH);\
    TaskTraceArr[Button_1_Monitor_TimeStampIn] = T1TC;\
82
83
84
85
86
87
88
89
90
91
92
93
94
95
96
97
98
99
                                                                          case Button_2_Monitor_Tag:\
   GPIO_write(TaskMonitor,Button_2_Monitor_PIN,PIN_IS_HIGH);\
   TaskTraceArr[Button_2_Monitor_TimeStampIn] = T1TC;\
                                                                             break; \
                                                                          case Periodic_Transmitter_Tag:\
   GPIO_write(TaskMonitor, Periodic_Transmitter_PIN, PIN_IS_HIGH);\
                                                                             TaskTraceArr[Periodic_Transmitter_TimeStampIn] = T1TC;\
                                                                             GPIO_write(TaskMonitor, Uart_Receiver_PIN, PIN_IS_HIGH);\
                                                                             TaskTraceArr[Uart_Receiver_TimeStampIn] = T1TC;
                                                                          break;\
case Load_1_Simulation_Tag:\
                                                                             GPIO write (TaskMonitor, Load 1 Simulation PIN, PIN IS HIGH);\
                                                                             TaskTraceArr[Load_1_Simulation_TimeStampIn] = T1TC;
                                                                          case Load 2 Simulation Tag:\
101
102
103
                                                                             GPIO_write(TaskMonitor,Load_2_Simulation_PIN,PIN_IS_HIGH);\
TaskTraceArr[Load_2_Simulation_TimeStampIn] = T1TC;\
104
105
106
                                                                          case Idle_Tag:\
GPIO_write(TaskMonitor,Idle_PIN,PIN_IS_HIGH);\
TaskTraceArr[Idle_TimeStampIn] = T1TC;\
                                                                             break; \
                                                                             }while(0)
```

```
| Section | Sect
```

We defined both trace TASK_SWITCHED_IN() & traceTASK_SWITCHED_OUT() . this function like macros are prone for errors and it is better to define it as a function call however we want to make the least time delay to accurately calculate the performance of the system

```
78 #define traceTASK_SWITCHED_IN()
                                                         {switch((int)pxCurrentTCB->pxTaskTag){\ going through task tags
                        if the function tag is ...
80
                    if the function tag is ... ____case Button_1 Monitor_Tag:\
make gpio specifc pin high _____GPIO_write(TaskMonitor, Button_1 Monitor_PIN, PIN_IS_HIGH);\
                                                          TaskTraceArr[Button_1_Monitor_TimeStampIn] = T1TC;\
           store the current timer1 counter ---
83
                                                              break:\
           inside the trace array
                                                            case Button_2_Monitor_Tag:\
85
                                                              GPIO_write(TaskMonitor,Button_2_Monitor_PIN,PIN_IS_HIGH);\
86
                                                               TaskTraceArr[Button_2_Monitor_TimeStampIn] = T1TC;\
                                                               break; \
                                                            case Periodic_Transmitter_Tag:\
111 | #define traceTASK_SWITCHED_OUT() | do\
112 | if current taks tag is ..
114 | make specifc pin low
                                          ad current timer1 counts

alculate difference

accumulate this calacultions

break;

case Button 2 Monitor Tags(\)

case Button 2 Monitor Tags(\)
          read current timer1 counts
```

Inside FreeRTOSConfig.h after we have to include my "TraceDef.h" which have defintions .. TaskTraceArr[] present inside main.c

In Main.c file

Defining TaskTraceArr[] , Button_1_Monitor_ExcutionTime ,Button_2_Monitor_ExcutionTime,

And CPU_Load,Button_1_Monitor_LoadPercent, Button_2_Monitor_LoadPercent,

Variables to monitor them during run-time of the system both we will obtain absolute and percentage values

Also we define WCET_Monitor variable which stands for worst case execution time to monitor specific task WCET according to hash defined macros inside the "TraceDef.h"

We will give each task its specific tag number

```
307 // used to calculate worest case excution time and periodic excution time
308 = #if (configTaskExcutionTimePeriodic ==1)
309 = {
310 = #ifdef Task1_WCET
311 if (WCET counter >= WCETcounterThreshold) {
312
       if (WCET Monitor < TaskTraceArr[Button 1 Monitor TotalEcutionTime]) {
         WCET Monitor = TaskTraceArr[Button 1 Monitor TotalEcutionTime];
313
         WCET counter = WCETcounterThreshold;
314
    WCET_counter++;
315
316
317
     TaskTraceArr[Button 1 Monitor TotalEcutionTime] =0;
319
320
321 #endif
322
323
324
    }
```

To calculate WCET for specific task you have to set "configTaskExcutionTimePeriodic" and "Task1_WCET" macros to 1 inside "TraceDef.h" file

As you can see we will not sample WCET except after counting n number of task execution because during start execution of simulator it will pass very large number to WCET_Monitor variable which is not correct

Inside "TraceDef.h"

```
#define MAX_TRACE_NUMB 35 defining numb of elements in trace array
10
                                we will monitor tasks through PORTO
   #define TaskMonitor PORT 0
11
12
14 // INPUTS BUTTONS
                             Defining input pins
   #define Button1 in PIN0
15
   #define Button2 in PIN1
16
17
   18 // OUTPUT MONITORING
   #define Button 1 Monitor PIN PIN2
19
                                            defining output pins
                                 PIN3
PIN4
PIN5
20
   #define Button_2_Monitor_PIN
21
   #define Periodic Transmitter PIN
22 #define Uart_Receiver_PIN
23 #define Load 1 Simulation PIN
24 | #define Load 2 Simulation PIN
                                   PIN7
25 | #define Idle_PIN
                                   PIN8
26
   #define Tick PIN
                                    PIN9
27
28
30 // Task Tags that will be traced during running
31 #define Idle Tag
                               0
32
   #define Button 1 Monitor Tag
                                   1
                                        defining tag number for
33 #define Button_2_Monitor_Tag 2
                                          each task
34 | #define Periodic Transmitter Tag 3
35 #define Uart Receiver Tag
36 #define Load 1 Simulation Tag
37 #define Load 2 Simulation Tag
// Array elements which will trace the time of tasks excution
42
   // present in TaskTraceArr[MAX TRACE NUMB]
   // Task1
43
44
   #define Button 1 Monitor TimeStampIn
45 #define Button 1 Monitor TimeStampOut
                                              1
46 #define Button 1 Monitor EcutionTime
47
   #define Button_1_Monitor_TotalEcutionTime
   // Task2
48
    #define Button 2 Monitor TimeStampIn
   #define Button 2 Monitor TimeStampOut
50
51 #define Button 2 Monitor EcutionTime
52 | #define Button_2_Monitor_TotalEcutionTime
53
   // Task3
54
    #define Periodic Transmitter TimeStampIn
55
   #define Periodic Transmitter TimeStampOut
                                             10
56 #define Periodic Transmitter EcutionTime
   #define Periodic Transmitter TotalEcutionTime 11
57
58
   //Task4
59
    #define Uart_Receiver_TimeStampIn
60
   #define Uart Receiver TimeStampOut
                                              13
61 #define Uart Receiver EcutionTime
                                             14
```

```
94
    #define configTaskExcutionTimePeriodic 1
95
    96
97
    // uncomment only (one) definition to monitor the variable WCET Monitor
    // according to task number the code to get the worst case excustion time will added
98
99
    // to enable this variable you must define configTaskExcutionTimePeriodic = 1
100
    //#define Task1_WCET
//#define Task2 WCET
101
102
    #define Task3 WCET
103
104
    //#define Task4 WCET
105
    //#define Task5 WCET
106
    //#define Task6 WCET
107
    //#define Task7 WCET
```

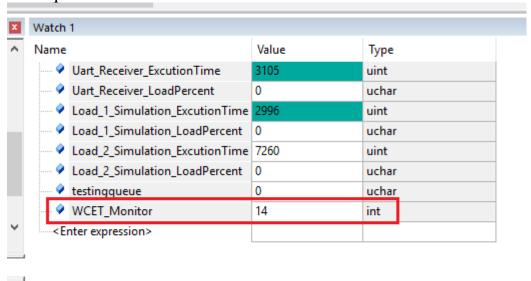
To calculate task specific WCET to be monitored by value of WCET Monitor.

For accurate results we will set TIMER1 devision to 100 instead of 1000

Then inside "TraceDef.h"

Set configTaskExcutionTimePeriodic to 1 and one by one define $Taskn_WCET$

Then in Watch Variable window add the WCET_Monitor variable to read the required task WCET

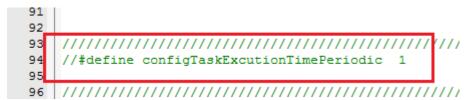


In this case we displayed WCET OF Button1 task which is 14 and to convert it to ms we will divide it by 600 (14/600) = 0.0233333 >> 0.023 ms or 23 us

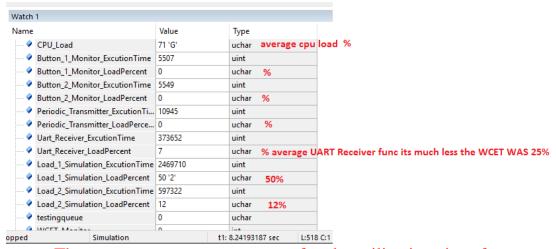
Results

Tasks	WCET Read	Result
Button_1_Monitor	14	0.0233
Button_2_Monitor	14	0.0233
Periodic_Transmitter	96	0.16
Uart_Receiver	2953	4.92
Load_1_Simulation	3006	5.01
Load_2_Simulation	7246	12.075

Average results percentage

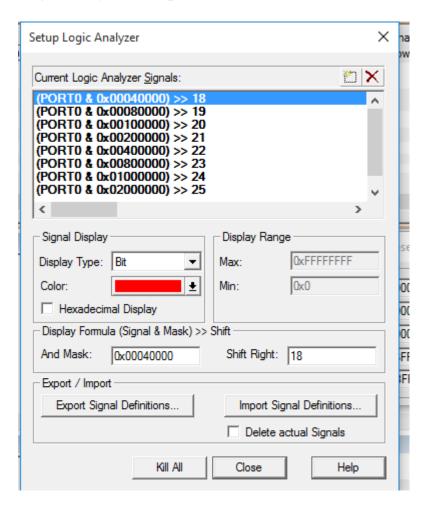


Undefine configTaskExcutionTimePeriodic then run simulation



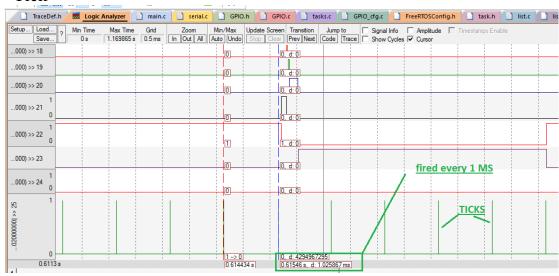
These are average percentage of tasks utilization time from trace macros

2- Logic analyzer interpretation



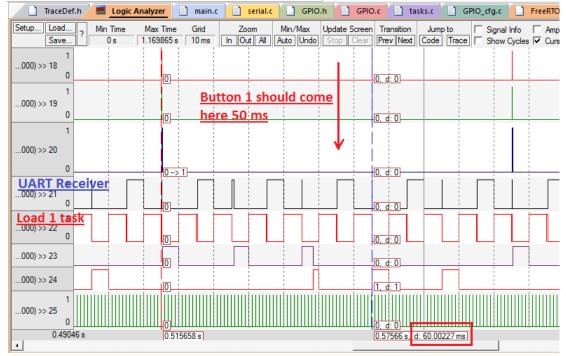
Button_1_Monitor	$pin \gg 18$
Button_2_Monitor	pin >> 19
Periodic_Transmitter	pin >> 20
Uart_Receiver	pin >> 21
Load_1_Simulation	pin >> 22
Load_2_Simulation	pin >> 23
Idle	pin >> 24
Tick	pin >> 25

*Tick

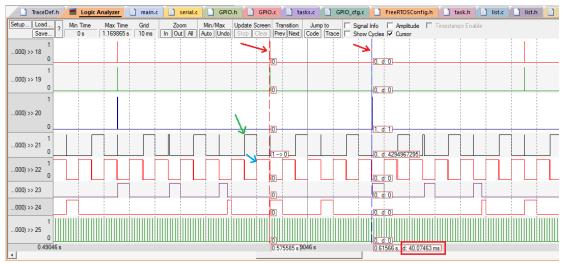


By logic analyzer the ticks will be fired every 1 ms

Button_1_Monitor & Button_2_Monitor

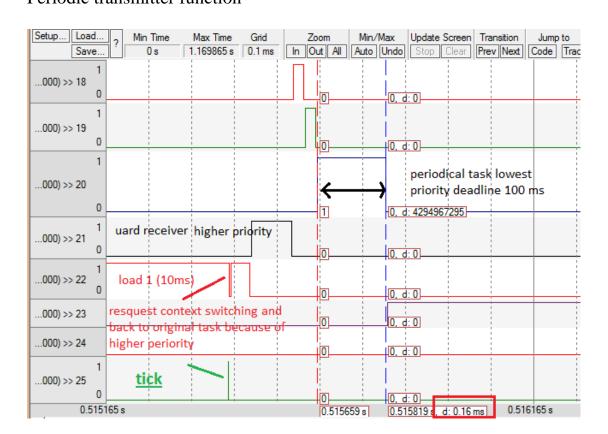


Button_1_Monitor & Button_2_Monitor should come every 50 ms . you can see the are delayed to 60 ms why??



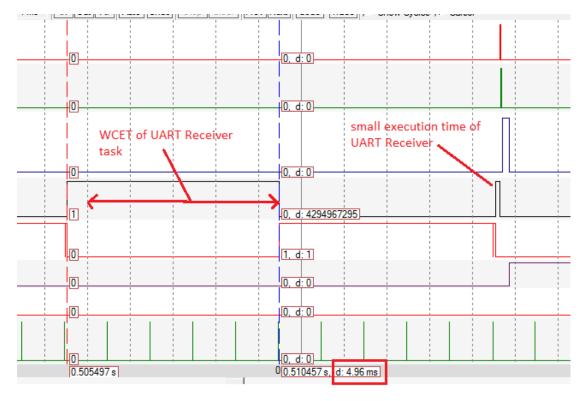
Because load1 simulation task (10 ms deadline) blue arrow and UART receiver task (20 ms deadline) Green arrow executed before it because their priorities are higher delaying it by 10 ms





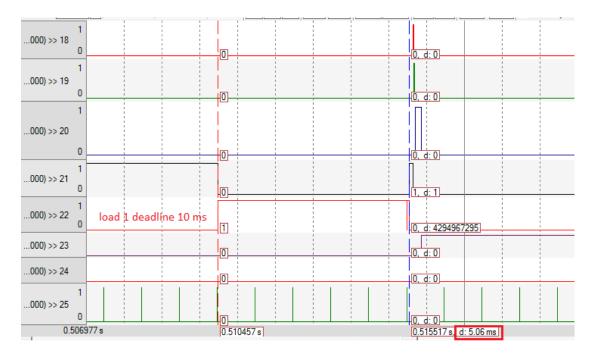
After the tick and insertion of the new tasks which were in delayed list the system added them to EDF Ready list and request for context switching and because the load1 task were running and didn't finished its execution and its deadline is 10 ms, context switching will get back to load1 task making this spike . then after this the UART receiver task will be executed because it is higher priority (20 ms deadline) and after its execution the other tasks (button 1 and 2) then at the last periodically transmitter task execute.

Its execution time is 0.16 ms

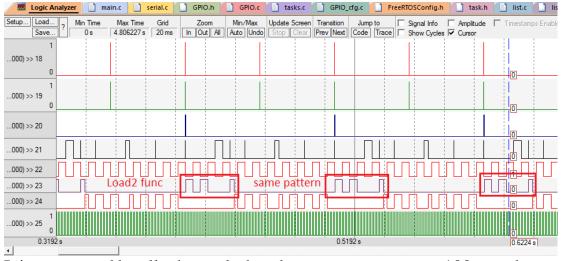


The worst case execution time of Uart_Receiver is 4.96 ms this is because at some point the queue contains many messages to be sent and it will wait UART module to send those strings.

Load_1_simulation

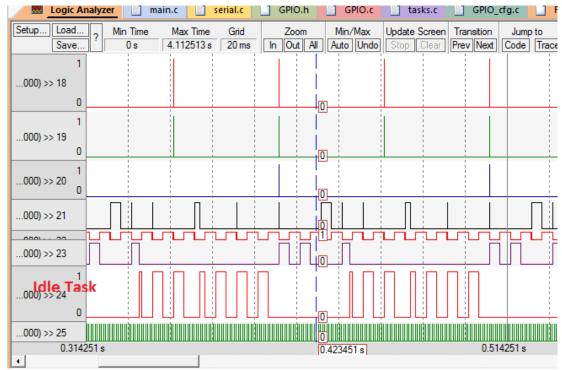


execution time of load 1 is 5 ms Load_2_simulation



It is preempted by all other tasks has the same pattern every 100 ms = hyper period it is about 12 ms

IDLE task



It is preempted by any other tasks

3- Using FreeRTOS Statistics

Implementation

In FreeRTOS.h

```
#define configMAX TASK NAME LEN ( 8 )
#define configUSE TRACE FACILITY 1 ////
#define configUSE 16 BIT TICKS 0
#define configUDLE SHOULD YIELD 0
```

Set configUSE_TRACE_FACILITY to 1

Add those macros

In Tasks.c

```
2720
2721 = #if ( configUSE TRACE FACILITY == 1 )
2722
       UBaseType t uxTaskGetSystemState( TaskStatus t * const pxTaskStatusArray, const UBa
2723
2724 🖨 - {
2725
2726
        UBaseType t uxTask = 0, uxQueue = configMAX PRIORITIES;
2727
2728
          vTaskSuspendAll();
2729 🗀
2730
            /\star Is there a space in the array for each task in the system? \star/
2731
            if( uxArraySize >= uxCurrentNumberOfTasks )
2732 📋
         /* Fill in an TaskStatus_t structure with information on each
```

Inside uxTaskGetSystemState function

We will change to this by adding (List_t *) &(xReadyTasksListEDF) instead of &(pxReadyTasksLists[uxQueue])

Results

				I
DLE	57063	10%	Idle 10% so cpu load 90%	
Load2	68543	12%	load2 12 ms = 12%	
UartR	138964	25%	WCET by using for loop load 5ms	
Button1	316	<1%		
Button2	313	<1%		
Periodi	969	<1%		
Load1	277762	51%	load1 5 ms every 10 ms = 5%	

You can see that the UART load in WCET is 25%

	N 100						
UART#2							
Button1	is HIGH		Button2 is HIGH	Message	from	Periodic	Func
Button1	is HIGH		Button2 is HIGH	I			
DLE	235281	28%					
Load2	108967	13%					
UartR	62432	7%					
Button1	475	<1%					
Button2	478	<1%					
Periodi	1451	<1%					
Load1	425894	50%	UART Receiver printing it	ts values			
Button1	is HIGH		Button2 is HIGH	Message	from	Periodic	Func
Button1	is HIGH		Button2 is HIGH	I			
DLE	252083	28%					
Load2	116751	13%					
L	66893	7%	this is the average UART Rece	eiver utilization			
UartR	00050						
UartR Button1		<1%					
	508	<1% <1%					
Button1	508 511						
Button1 Button2	508 511 1548	<1%					
Button1 Button2 Periodi	508 511 1548	<1% <1%					
Button1	508						

When we just make the UART Receiver func do its job the average utilization will be 7% those results are consistent with our implementation

In conclusion the EDF implementation meets the requirements of no task will break its deadline and the system overall is schedulable with low margin as in WCET the cpu load will be 90% and we have only 10% margine to play around. All results from our implementation are consistent with FreeRTOS statistics function even more reliable than it.