

Tutorial 4: Dynamic Programming

This tutorial helps you develop skills in the learning outcome of the course: “Able to design algorithms using suitable strategies (dynamic programming, etc) to solve a problem, able to analyse the efficiencies of different algorithms for problems like optimal sequencing for matrix multiplication, the longest common subsequence, etc”.

1. Find the length of the longest common subsequence and a longest common subsequence of CAGAG and ACTGG by the dynamic programming algorithm in the lecture notes.

2. The H-number $H(n)$ is defined as follows:

$H(0) = 1$, and for $n > 0$:

$H(n) = H(n-1) + H(n-3) + H(n-5) + \dots + H(0)$ when n is odd

$H(n) = H(n-2) + H(n-4) + H(n-6) + \dots + H(0)$ when n is even.

- a) Give a recursive algorithm to compute $H(n)$ for an arbitrary n as suggested by the recurrence equation given for $H(n)$. Draw the tree that represents the recursive calls made when $H(8)$ is computed.
- b) Draw the subproblem graph for $H(8)$ and $H(9)$.
- c) Write an iterative algorithm using the dynamic programming approach (bottom-up). What are the time and space required?

3. The binomial coefficients can be defined by the recurrence equation:

$$C(n, k) = C(n-1, k-1) + C(n-1, k) \quad \text{for } n > 0 \text{ and } k > 0$$

$$C(n, 0) = 1 \quad \text{for } n \geq 0$$

$$C(0, k) = 0 \quad \text{for } k > 0$$

$C(n, k)$ is also called “ n choose k ”. This is the number of ways to choose k distinct objects from a set of n objects.

- (a) Give a recursive algorithm as suggested by the recurrence equation given for $C(n, k)$.
- (b) Draw the subproblem graph for $C(5, 3)$.
- (c) Write a recursive algorithm using the dynamic programming approach (top-down) stating the data structure used for the dictionary. What is the space and time complexity respectively?
- (d) Write an iterative algorithm using the dynamic programming approach (bottom-up). What is the space and time complexity respectively?

4. Construct an example with only three or four matrices where the worst multiplication order does at least 100 times as many element-wise multiplications as the best order.
5. Suppose the dimensions of the matrices A , B , C , and D are 20×2 , 2×15 , 15×40 , and 40×4 , respectively, and we want to know how best to compute $A \times B \times C \times D$. Show the arrays **cost** and **last** computed by Algorithm `matrixOrder()` in the lecture notes.
6. We have a knapsack of size 10 and 4 objects. The sizes and the profits of the objects are given by the table below. Find a subset of the objects that fits in the knapsack that maximizes the total profit by the dynamic programming algorithm in the lecture notes.

p	10	40	30	50
s	5	4	6	3

7. $S1$ is a sequence of $n1$ characters and $S2$ is a sequence of $n2$ characters. All characters are from the set $\{'a', 'c', 'g', 't'\}$. An alignment is defined by inserting any number of character $'_'$ (the underscore character) into $S1$ and $S2$ so that the resulting sequences $S1'$ and $S2'$ are of equal length. Each character in $S1'$ has to be aligned with the same character or an underscore in the same position in $S2'$ and vice versa. The cost of an alignment of $S1$ and $S2$ is defined as the number of underscore characters inserted in $S1$ and $S2$. For example, $S1 = \text{"ctatg"}$ and $S2 = \text{"ttaagc"}$. One possible alignment is

$S1' = \text{"ct_at_g_"} \text{ and}$
 $S2' = \text{"_tta_agc"}$

Both $S1'$ and $S2'$ have length 8 and the cost is 5. We want to find the minimum cost of aligning two sequences, denoted as $\text{alignment}(n1, n2)$.

- (a) Give a recursive definition of $\text{alignment}(n1, n2)$.
- (b) Draw the subproblem graph for $\text{alignment}(3, 4)$.
- (c) Design a dynamic programming algorithm of $\text{alignment}(n1, n2)$ using the bottom-up approach.