

LeetCode

A project dedicated to DS&A

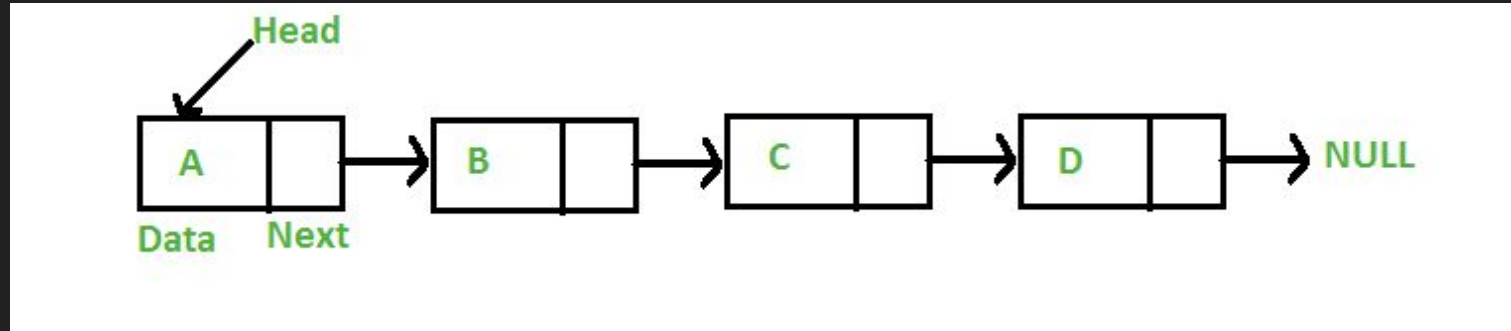
Agenda

1. Quick Review
2. Binary Trees
3. DFS Fundamentals
4. LeetCode Problems

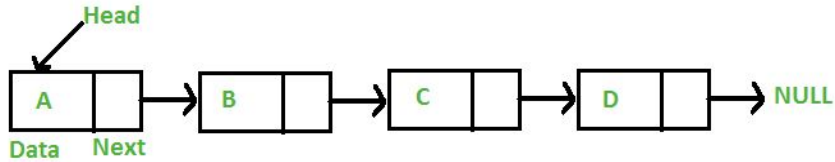
Quick Review

Linked Lists

- A **linked list** is a **linear data structure** that is not stored contiguously
 - i.e. the elements in the list are not stored next to each other in memory
- In a singly linked list, **each node** has data and a pointer to the next node

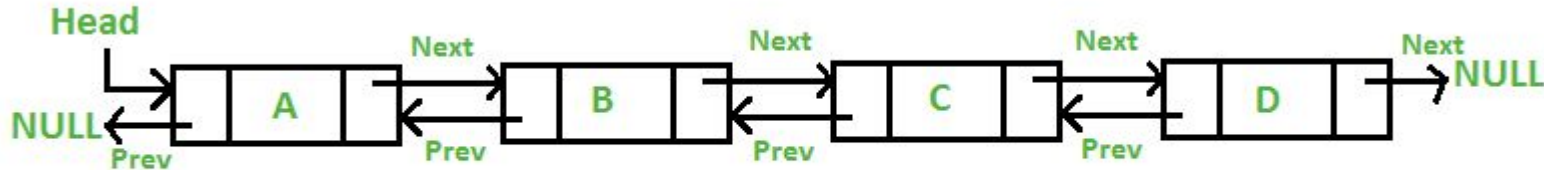


Types of Linked Lists



Single Linked Lists

Circular Linked Lists

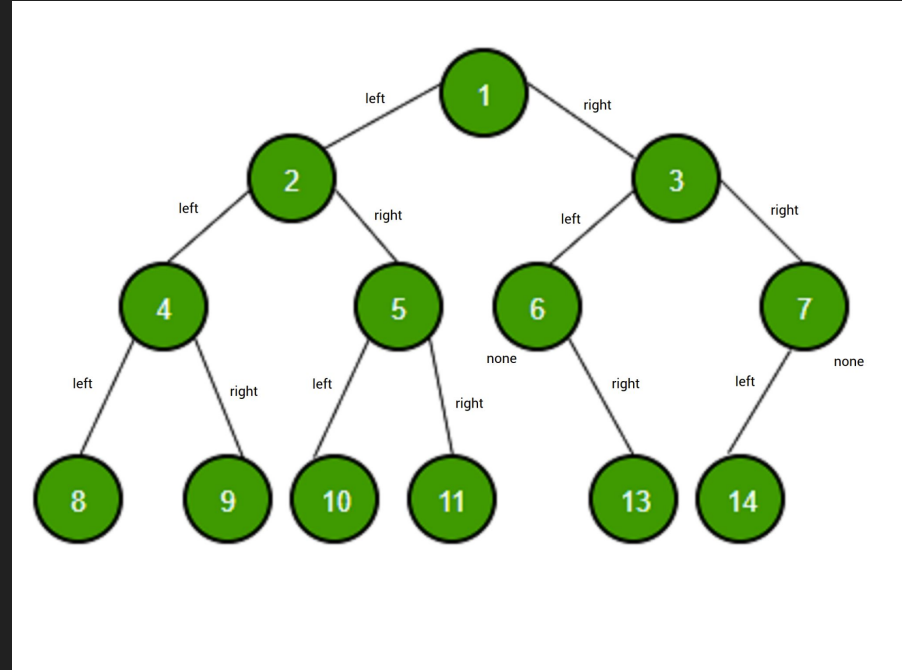


Doubly Linked Lists

Binary Trees

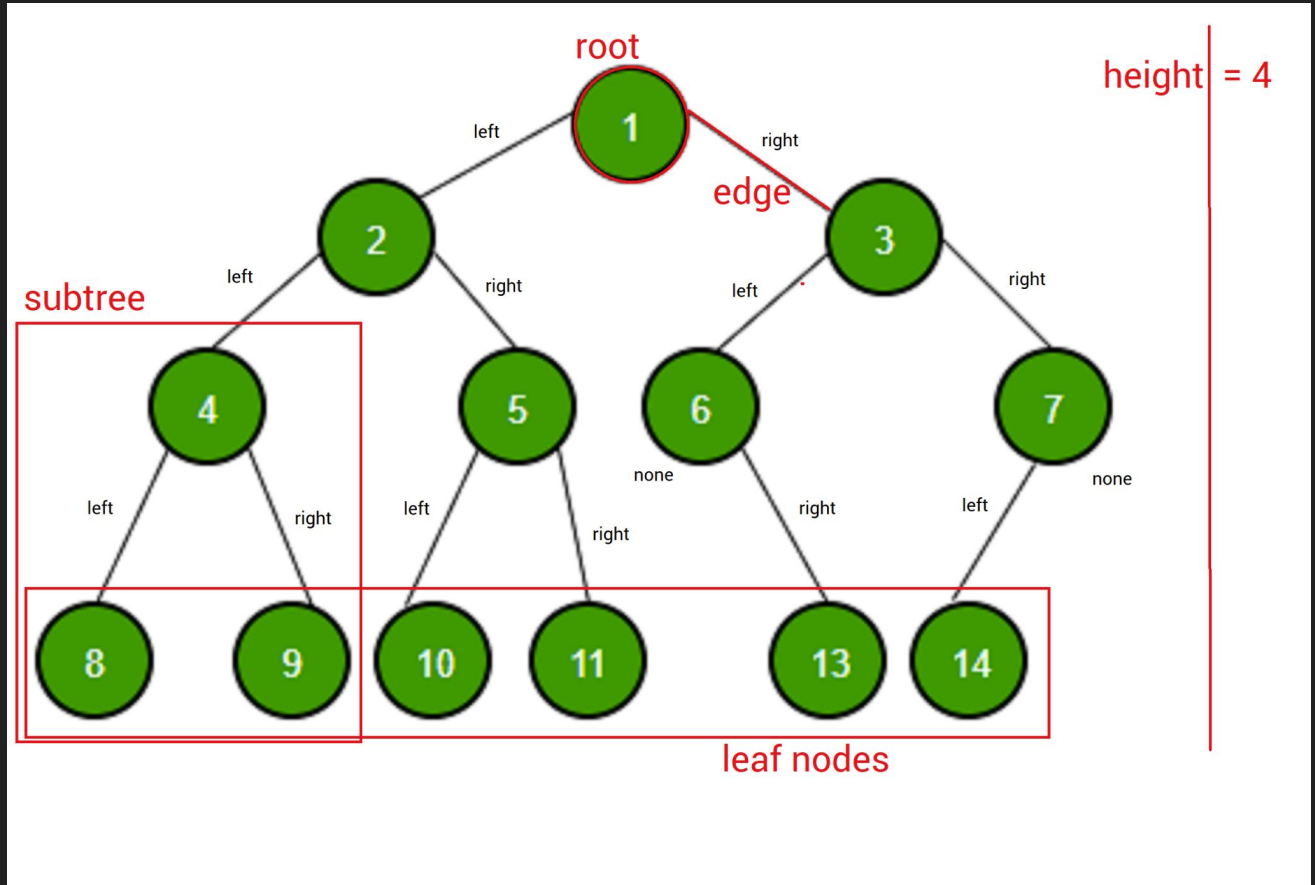
Binary Trees

- A **binary tree** is represented by a pointer to the topmost **node** of the tree
- If the tree is **empty**, then the value of the root is **None**



Tree Terminology

- Root
- Edge
- Leaves
- Subtree
- Height

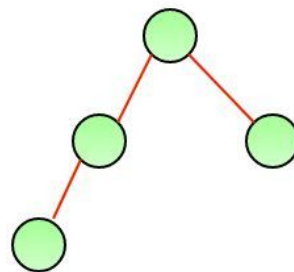


Full, Complete, and Perfect Binary Trees

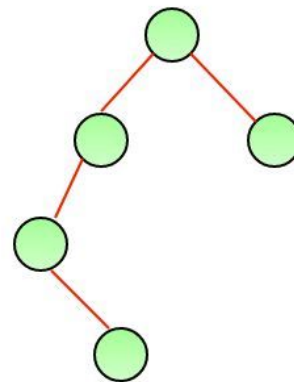
- Full Binary Tree
 - Each node has 0 or 2 children
- Complete Binary Tree
 - All levels are completely filled except for the last row
 - The last row's nodes are all as left as possible
- Perfect Binary Tree
 - All levels are completely filled
 - Total Number of Nodes = $2 * \text{Number of Leaf Nodes} - 1$

Balanced Binary Tree

- The **height difference** between the left and right subtree of the node is **not more than 1**.
- This property is important for ensuring the effectiveness of our algorithms...



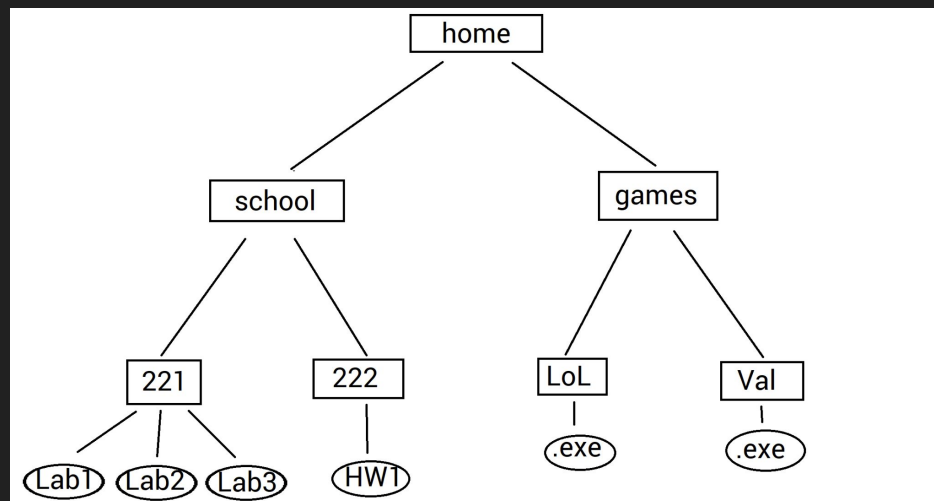
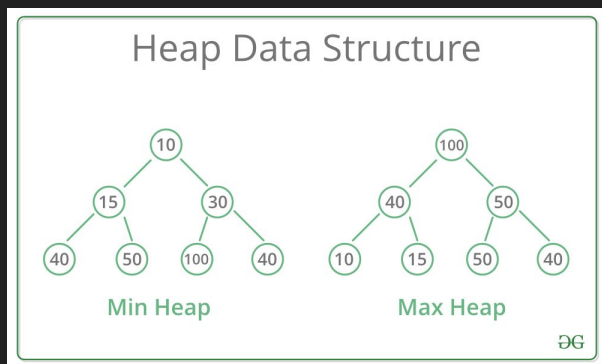
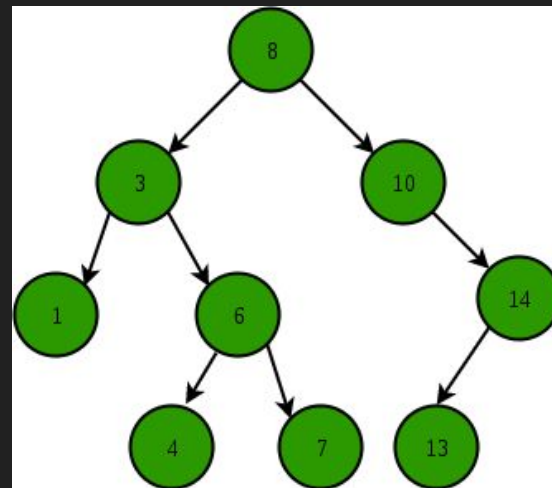
A height balanced tree



Not a height balanced tree

When should I use a tree?

- When making “decisions” in an algorithm
 - I.e. less than or greater than
- Storing hierarchical data
 - I.e. a file system
- Inside of other data structures
 - Heaps
 - Priority Queues



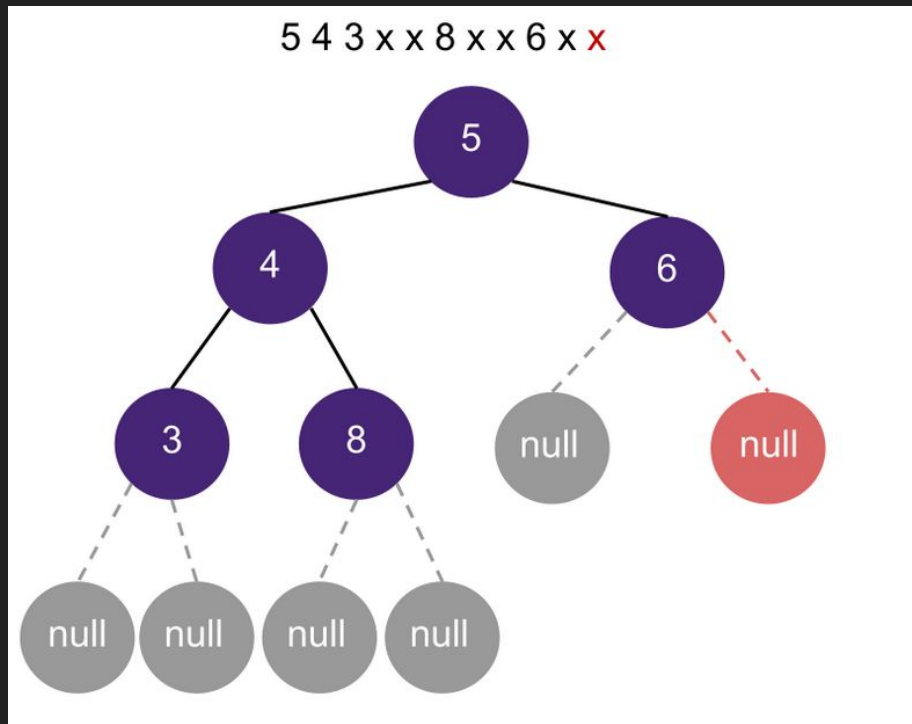
Time Complexity

- **Searching:** Since we haven't introduced special types of trees, searching is worst case $O(n)$
- **Insertion:** Insertion also has a worst case of $O(n)$, though it is $O(h)$ if the tree is balanced
- **Deletion:** Since we have to search before deleting, it is also $O(n)$
- These time complexities can be improved significantly when we introduce special tree types, such as Binary Search Trees.

Local Setup

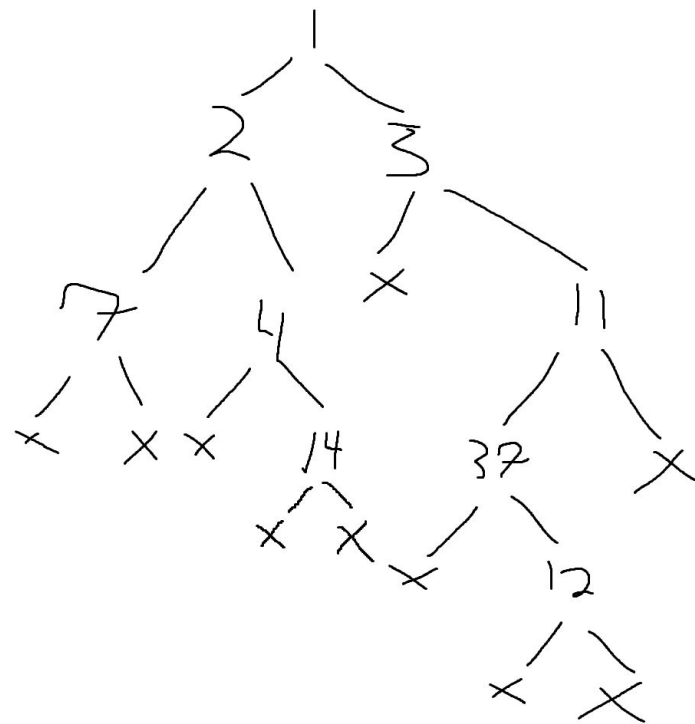
- Get code from our new GitHub:
<https://tx.ag/LCBinaryTree>
- Replace “yourFunction” with your function name
- Replace “yourFunction” with your function name in the last line
- Input will be given in the form shown on the right —>

Example Input



Another Example of Input

- This is just the way that I have encoded input in the GitHub...
 - You can do this differently if you write your own!

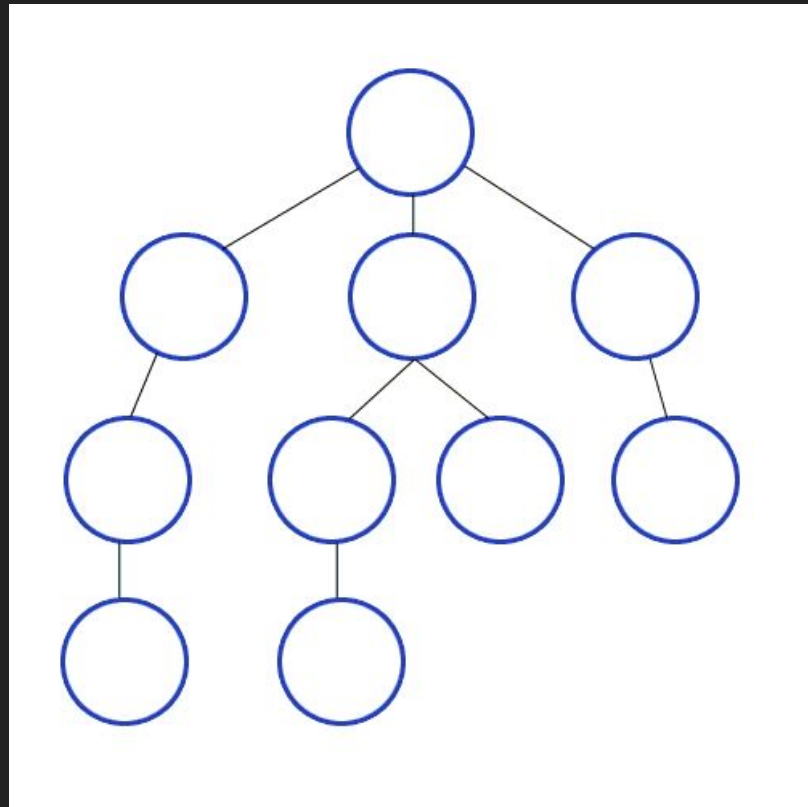


1 2 7 x x 4 x 14 x x 3 x 11 37 x 12 x x x

DFS Fundamentals

Depth First Search (DFS)

- Depth First means we go as deep as we can to look for a value
 - I.e. we go down the rabbit hole as far as possible, then come back up and try again with another rabbit hole
- Two different options
 - Iterative
 - Recursive



Recursive DFS Boilerplate Code

```
def dfs(root, target):  
    if root is None:  
        return None  
    if root.val == target:  
        return root  
    # return non-null return value from the recursive calls  
    left = dfs(root.left, target)  
    if left is not None:  
        return left  
  
    # at this point, we know left is null, and right could be null or non-null  
    # we return right child's recursive call result directly because  
    # - if it's non-null we should return it  
    # - if it's null, then both left and right are null, we want to return null  
    return dfs(root.right, target)  
# the code can be shortened to: return dfs(root.left, target) or dfs(root.right, target)
```

Iterative DFS Code

```
def dfsIterative(root: TreeNode, target):  
    """Iterative Depth First Search (DFS)"""  
    if root is None:  
        return  
    stack = []  
    curr = root  
    prev = None  
    while stack or curr is not None: # While stack isn't empty OR curr isn't None  
        if curr is not None:  
            if curr.val == target: # if curr node is target, return it  
                return curr  
            stack.append(curr) # append curr node to stack  
            curr = curr.left # update curr node to curr.left  
        else:  
            prev = stack.pop() # pop node off of stack  
            curr = prev.right # update curr to prev.right  
    return # target is not in the tree
```

LeetCode Problems

Example Problem #1 - LeetCode 104.

- <https://leetcode.com/problems/maximum-depth-of-binary-tree/>
- Can we implement what we have learned to solve this problem?
 - DFS will work
 - You can use an iterative or recursive solution
 - Other algorithms will also work, but we haven't talked about them yet.
- Time Complexity?
 - $O(n)$
- Space Complexity?
 - $O(1)$

Group Work (if there is time)

Until next time...

Keep practicing

Practice Problems

Easy (supposedly):

- <https://leetcode.com/problems/invert-binary-tree/>
- <https://leetcode.com/problems/subtree-of-another-tree/>

A bit more difficult:

- <https://leetcode.com/problems/binary-tree-maximum-path-sum/>

Quite difficult (but also cool):

- <https://leetcode.com/problems/serialize-and-deserialize-binary-tree/>

Questions?