

MODULE *Reno*

Implements simple *AIMD TCP* protocol. It can use any pthe model, provided that it implements the specification of *SimplePathModel*.

EXTENDS *TLC*, *Sequences*, *Integers*

*MAX\_ARRIVAL* : The rate at which packets can arrive. Helps lower runtime.

*SSTHRESH\_START*: Starting value of Slow-Start Threshold.

*MAX\_WINDOW* : Maximum allowed value of *cwnd*.

*added* : The number of packets arriving.

*arrivals* : Total number of packets arrived.

*buffered* : Number of packets buffered for service.

*cwnd* : Congestion window.

*ssthresh* : Slow-Start threshold

CONSTANT *C*, *MAX\_T*, *MAX\_ARRIVAL*, *DROP\_ACK*,  
*SSTHRESH\_START*, *MAX\_WINDOW*

VARIABLES *t*, *ticked*,  
*nAck*, *inFlight*, *timeout*,  
*added*, *arrivals*, *buffered*, *cwnd*, *ssthresh*

*timeVars*  $\triangleq$   $\langle t, \textit{ticked} \rangle$

*pathVars*  $\triangleq$   $\langle nAck, inFlight, timeout \rangle$

*tcpVars*  $\triangleq$   $\langle arrivals, buffered, cwnd, ssthresh, added \rangle$

*vars*  $\triangleq$   $\langle t, \textit{ticked}, nAck, inFlight, timeout, arrivals, added, buffered, cwnd, ssthresh \rangle$

*time*  $\triangleq$  INSTANCE *Time* WITH  $t \leftarrow t$ ,  $\textit{ticked} \leftarrow \textit{ticked}$ ,  $\textit{MAX\_T} \leftarrow \textit{MAX\_T}$

Here, “*SimplePathModel*” is used. Any other path model used must implement:

- 1) “*ExcessivePacketDropIsEnabled*”: A boolean formula that specifies if path capacity has been exceeded. This MUST always trigger a timeout and thus disables the *TCP* until that timeout happens.
- 2) “*Init*” and “*Next*” predicates.
- 3) “*Fairness*” for fairness conditions if required.
- 4) Variables  $\langle nAck, inFlight, timeout \rangle$  at the very least.

*path*  $\triangleq$  INSTANCE *SimplePathModel* WITH  $nAck \leftarrow nAck$ ,  $inFlight \leftarrow inFlight$ ,  
 $timeout \leftarrow timeout$ ,  $C \leftarrow C$ ,  
 $\textit{MAX\_ARRIVAL} \leftarrow \textit{MAX\_ARRIVAL}$ ,  
 $\textit{DROP\_ACK} \leftarrow \textit{DROP\_ACK}$

ASSUME  $\wedge \textit{SSTHRESH\_START} > 1$   
 $\wedge \textit{MAX\_WINDOW} > \textit{SSTHRESH\_START}$   
 $\wedge \textit{MAX\_T} > 2$

A note about “*added*”:

“added” signifies that some packets have arrived and should go into the input buffer. If added is nonzero, all *TCP* window actions are disabled until we add the packets into the buffer.

This does not change the possible behaviors, but it helps prevent some repeated states when those packets need to be sent.

$$\begin{aligned}
TypeOK \triangleq & \wedge path!TypeOK \\
& \wedge arrivals \in Nat \\
& \wedge arrivals \geq 0 \\
& \wedge buffered \in Nat \\
& \wedge buffered \geq 0 \\
& \wedge added \in Nat \\
& \wedge added \geq 0 \\
& \wedge cwnd \in Nat \\
& \wedge cwnd \geq 1 \\
& \wedge ssthresh \in Nat \\
& \wedge ssthresh \geq 2
\end{aligned}$$

These are the 3 basic conditions that the specification must satisfy. It gurantees the safety of instantiated specifications.

$$Termination \triangleq path!Termination$$

$$Finished \triangleq path!Finished$$

$$RateLimited \triangleq path!RateLimited$$

$$\begin{aligned}
Init \triangleq & \wedge path!Init \\
& \wedge arrivals = 0 \\
& \wedge cwnd = 1 \\
& \wedge ssthresh = SSTHRESH\_START \\
& \wedge buffered = 0 \\
& \wedge added = 0
\end{aligned}$$

*EnableIncreaseWindow*:

Only enabled when no timeout has happened and we have not exceeded *MAX\_WINDOW*. If still in slow-start, one *ACK* at least should be present, if during congestion avoidance, we need a whole *cwnd* worth of acks to increase the window.

*IncreaseWindow*:

The action itself doubles *cwnd* if during slow-start or adds one to it if during congestion avoidance. Appropriate amounts of *ACKs* are consumed as well ...

$$\begin{aligned}
EnableIncreaseWindow \triangleq & \wedge timeout = 0 \\
& \wedge cwnd < MAX\_WINDOW \\
& \wedge IF \ cwnd < ssthresh \\
& \quad THEN \ nAck > 0 \\
& \quad ELSE \ nAck \geq cwnd \\
& \wedge Finished = FALSE
\end{aligned}$$

$$\begin{aligned} & \wedge \neg path!ExcessivePacketDropIsEnabled \\ & \wedge added = 0 \end{aligned}$$

$$\begin{aligned} IncreaseWindow & \triangleq \wedge EnableIncreaseWindow \\ & \wedge IF \ cwnd < ssthresh \\ & \quad THEN \ \wedge \ cwnd' = 2 * cwnd \\ & \quad \quad \wedge \ nAck' = nAck - 1 \\ & \quad ELSE \ \wedge \ nAck' = nAck - cwnd \\ & \quad \quad \wedge \ cwnd' = cwnd + 1 \\ & \wedge UNCHANGED \ \langle ssthresh, timeout, inFlight, added, \\ & \quad arrivals, buffered, t, ticked \rangle \end{aligned}$$

*EnableDecreaseWindow:*

Only enabled during timeouts.

*DecreaseWindow:*

It resets the number of received ACKs, sets timeout to zero and cuts the window in half if slow-start is finished, else *cwnd* is set to 1. Half of current *cwnd* is added to slow-start threshold.

$$\begin{aligned} EnableDecreaseWindow & \triangleq \wedge timeout = 1 \\ & \wedge Finished = FALSE \\ & \wedge \neg path!ExcessivePacketDropIsEnabled \\ & \wedge added = 0 \end{aligned}$$

$$\begin{aligned} DecreaseWindow & \triangleq \wedge EnableDecreaseWindow \\ & \wedge IF \ cwnd \geq 4 \\ & \quad THEN \ ssthresh' = cwnd \div 2 \\ & \quad ELSE \ ssthresh' = 2 \\ & \wedge IF \ cwnd < ssthresh \\ & \quad THEN \ \wedge \ cwnd' = 1 \\ & \quad ELSE \ \wedge \ cwnd' = cwnd \div 2 \\ & \wedge timeout' = 0 \\ & \wedge nAck' = 0 \\ & \wedge UNCHANGED \ \langle inFlight, added, arrivals, buffered, t, \\ & \quad ticked \rangle \end{aligned}$$

$$Max(a, b) \triangleq IF \ a > b \ THEN \ a \ ELSE \ b$$

$$\begin{aligned} newPacketsAllowed(timePassed, packetsArrived) & \triangleq Max( \\ & \quad timePassed * MAX\_ARRIVAL - packetsArrived - 1, 0 \\ & ) \end{aligned}$$

$$\begin{aligned} getRandomArrival(timePassed, packetsArrived) & \triangleq RandomElement( \\ & \quad 0 \dots newPacketsAllowed(timePassed, packetsArrived) \\ & ) \end{aligned}$$

*PacketAcceptIsEnabled:*

Enabled only when timeout is zero and we have not exceeded the maximum arrival rate yet.

*AcceptPackets:*

Draws a random number of packets up to a maximum value that gurantees arrival rate is not exceeded, and assigns it to “added”. If there are packets to add, all *TCP* actions are disabled until these new packets are added to the input buffer.

$$\begin{aligned} \text{PacketAcceptIsEnabled} \triangleq & \quad \wedge \text{ticked} = 0 \\ & \quad \wedge \text{timeout} = 0 \\ & \quad \wedge \text{Finished} = \text{FALSE} \\ & \quad \wedge (t * \text{MAX\_ARRIVAL} - \text{arrivals}) > 0 \\ & \quad \wedge \neg \text{path!ExcessivePacketDropIsEnabled} \end{aligned}$$

$$\begin{aligned} \text{AcceptPackets} \triangleq & \quad \wedge \text{PacketAcceptIsEnabled} \\ & \quad \wedge \text{added}' = \text{getRandomArrival}(t, \text{arrivals}) \\ & \quad \wedge \text{UNCHANGED } \langle t, \text{ticked}, n\text{Ack}, in\text{Flight}, \text{timeout}, cwnd, \\ & \quad \quad ssthresh, \text{arrivals}, buffered \rangle \end{aligned}$$

*PacketArrivalIsEnabled:*

Only enabled when added is non-zero and the rest of packet acceptance conditions also hold.

*PacketArrival:*

Adds the value of “added” to total arrivals and buffered values, and then resets “added” to zero, enabling furthur *TCP* actions.

$$\begin{aligned} \text{PacketArrivalIsEnabled} \triangleq & \quad \wedge \text{added} > 0 \\ & \quad \wedge \text{PacketAcceptIsEnabled} \end{aligned}$$

$$\begin{aligned} \text{PacketArrival} \triangleq & \quad \wedge \text{PacketArrivalIsEnabled} \\ & \quad \wedge \text{arrivals}' = \text{arrivals} + \text{added} \\ & \quad \wedge \text{buffered}' = \text{buffered} + \text{added} \\ & \quad \wedge \text{added}' = 0 \\ & \quad \wedge \text{UNCHANGED } \langle t, \text{ticked}, n\text{Ack}, in\text{Flight}, \text{timeout}, cwnd, \\ & \quad \quad ssthresh \rangle \end{aligned}$$

*PacketSendIsEnabled:*

Only enabled when there are buffered packets to deliver and  $in\text{Flight} < cwnd$ , meaning that the window has empty space.

*SendNewPackets:*

If there are enough buffered packets, the window is filled completely, if not, the window is filled until the buffer is empty.

$$\begin{aligned} \text{PacketSendIsEnabled} \triangleq & \quad \wedge \text{ticked} = 0 \\ & \quad \wedge \text{Finished} = \text{FALSE} \\ & \quad \wedge \text{buffered} > 0 \\ & \quad \wedge in\text{Flight} < cwnd \end{aligned}$$

$$\begin{aligned}
& \wedge \text{timeout} = 0 \\
& \wedge \neg \text{path!ExcessivePacketDropIsEnabled} \\
& \wedge \text{added} = 0 \\
\text{SendNewPackets} & \triangleq \wedge \text{PacketSendIsEnabled} \\
& \wedge \text{IF } \text{cwnd} - \text{inFlight} > \text{buffered} \\
& \quad \text{THEN } \wedge \text{inFlight}' = \text{inFlight} + \text{buffered} \\
& \quad \wedge \text{buffered}' = 0 \\
& \quad \text{ELSE } \wedge \text{buffered}' = \text{buffered} - (\text{cwnd} - \text{inFlight}) \\
& \quad \wedge \text{inFlight}' = \text{cwnd} \\
& \wedge \text{UNCHANGED } \langle t, \text{ticked}, n\text{Ack}, \text{timeout}, \text{cwnd}, \text{added}, \\
& \quad \text{ssthresh}, \text{arrivals} \rangle \\
\text{Next} & \triangleq \vee \text{IncreaseWindow} \\
& \vee \text{DecreaseWindow} \\
& \vee \text{SendNewPackets} \\
& \vee \text{AcceptPackets} \\
& \vee \text{PacketArrival} \\
& \vee \wedge \text{path!Next} \\
& \wedge \text{UNCHANGED } \text{tcpVars} \\
\text{Fairness} & \triangleq \wedge \text{path!Fairness} \\
\text{Spec} & \triangleq \text{Init} \wedge \Box[\text{Next}]_{\text{vars}} \wedge \text{Fairness}
\end{aligned}$$


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