Ecole Polytechnique Federale de Lausanne

MASTER THESIS

Efficient Deoptimization

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in the

LARA Computer Science

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Declaration of Authorship

I, Adrien GHOSN, declare that this thesis titled, "Efficient Deoptimization" and the work presented in it are my own. I confirm that:

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- Where I have consulted the published work of others, this is always clearly attributed.
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- I have acknowledged all main sources of help.
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"Thanks to my solid academic training, today I can write hundreds of words on virtually any topic without possessing a shred of information, which is how I got a good job in journalism."
Dave Barry

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Abstract

Faculty Name Computer Science

Master in Computer Science

Efficient Deoptimization

by Adrien GHOSN

The Thesis Abstract is written here (and usually kept to just this page). The page is kept centered vertically so can expand into the blank space above the title too...

Acknowledgements

The acknowledgements and the people to thank go here, don't forget to include your project advisor...

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 ω angular frequency rad

For/Dedicated to/To my...

Introduction

Related Work

2.1 On Stack Replacement, General Principle

2.1.1 Definition & Overview

On-Stack replacement (OSR) is a set of techniques that consist in dynamically transferring the execution, at run time, between different pieces of code. The action of transferring the execution to another code artefact is called an OSR transition.

On-Stack replacement can be viewed, at a high level, as a mechanism that allows to transform the currently executing code, into another version of itself. This transformation mechanism has been used to allow the bi-directional transition between different levels of code optimisations. We can therefore reduce it to two main purposes: transforming an executing piece of code into a more optimised version of itself, and undoing transformations that were previously performed. While similar, these two types of transformation have very different goals.

In several virtual machines (CITE PAPERS), some of which will be presented in (REFERENCE), On-Stack replacement has been used to improve the performance of long running functions. When the VM identifies a piece of code as being "hot", i.e., it hogs the execution, it suspends its execution, recompiles it to a higher level of optimisation, and transfers the execution to the newly generated version of the function. This differs from a simple Just-In-Time (JIT) compiler, since the recompilation takes place during the execution of the function, rather than just before its execution. However, both techniques rely on run time profiling data to uncover new optimisation opportunities.

In this case, OSR is used to improve performance.

On-Stack replace allows to

- 2.1.2 The origins: SELF debugging
- 2.1.3 Why is OSR interesting?
- 2.2 On Stack Replacement & Virtual Machines
- 2.2.1 In Java
- 2.2.2 LLVM
- 2.3 A Description of Existing Implementations
- 2.3.1 The OSR points
- 2.3.2 The Transition Mechanism
- 2.3.3 Constraints and Limitations
- 2.3.4 Generating on the Fly VS Caching
- 2.3.5 Discussion

Theoretical Model

- 3.1 The OSR points
- 3.2 The Transition Mechanism
- 3.3 Constaints

Implementation

Appendix A

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