```
# Assignment #10: dp & bfs
Updated 2 GMT+8 Nov 25, 2024
2024 fall, Complied by <mark>付耀贤, 信息管理系/mark>
## 1. 题目
### LuoguP1255 数楼梯
dp, bfs, https://www.luogu.com.cn/problem/P1255
思路:
DP 经典题目,没有问题。
代码:
n=int(input())
dp = [0] * (n + 3)
dp[1] = 1
dp[2] = 2
for i in range(3, n + 1):
   dp[i]=dp[i-1]+dp[i-2]
print(dp[n])
代码运行截图 <mark> (至少包含有"Accepted") </mark>
fushion
11-26 16:03:55
               Accepted
                           P1255 数楼梯
                                                   © 163ms / 

4.98MB / 

124B Python 3
### 27528: 跳台阶
dp, http://cs101.openjudge.cn/practice/27528/
思路:
dp 关系变为:dp[i] = dp[i-1] + dp[i-2] + ... + dp[0], 本质是一样的
代码:
n=int(input())
dp = [0] * (n+1)
dp[0] = 1
       dp[i]+=dp[j]
print(dp[n])
```

代码运行截图 == (至少包含有"Accepted") ==

状态: Accepted

```
      據代码
      #: 47408891

      n=int(input())
      题目: 27528

      dp = [0] * (n+1)
      提交人: beginner

      dp[0] = 1
      内存: 3612kB

      for i in range(1, n+1):
      时间: 29ms

      for j in range(i):
      语言: Python3

      dp[i]+=dp[j]
      提交时间: 2024-11-26 16:24:21
```

474D. Flowers

dp, https://codeforces.com/problemset/problem/474/D

思路:

K=2 时,可以推出来 n**2//4+1;但是超时。dp 关系找不出来,,看题解有点糊涂。

```
dp[0] = 1
       dp[n] = dp[n - 1] # 如果直接吃红花的选择
       prefix sum[i] = (prefix sum[i - 1] + dp[i]) % MOD
   # 返回查询结果
           result.append((prefix sum[b] - prefix sum[a - 1]) % MOD)
           result.append(prefix sum[b] % MOD)
import sys
input = sys.stdin.read
data = input().splitlines()
```

```
results = solve(t, k, queries)
sys.stdout.write('\n'.join(map(str, results)) + '\n')
```

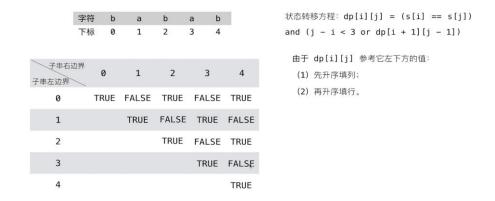
```
293341322 Nov/26/2024 18:30<sup>UTC+8</sup> aglint <u>D - Flowers</u> Python 3 Accepted 328 ms KB
```

LeetCode5.最长回文子串

dp, two pointers, string, https://leetcode.cn/problems/longestpalindromic-substring/

思路:

看的题解, 这张图很清楚!



```
class Solution:
    def longestPalindrome(self, s: str) -> str:
        n = len(s)
        if n < 2:
            return s
        max_len = 1
        begin = 0
        # dp[i][j] 表示 s[i..j] 是否是回文串
        dp = [[False] * n for _ in range(n)]
        for i in range(n):
            dp[i][i] = True

# 递推开始
# 先枚举子串长度
for L in range(2, n + 1):
        # 枚举左边界, 左边界的上限设置可以宽松一些
```

```
for i in range(n):
    # 由 L 和 i 可以确定右边界,即 j - i + 1 = L 得
    j = L + i - 1
    if j >= n:
        break
    if s[i] != s[j]:
        dp[i][j] = False
    else:
        if j - i < 3:
             dp[i][j] = True
        else:
             dp[i][j] = dp[i + 1][j - 1]

# 只要 dp[i][j] = true 成立,就表示子串 s[i..L] 是回文,此时记录回文长度和起始位置
    if dp[i][j] and j - i + 1 > max_len:
        max_len = j - i + 1
        begin = i
    return s[begin:begin + max_len]
```

通过

N Strange I3haskarahGJ 提交于 2024.11.26 18:06

```
### 12029: 水淹七军

bfs, dfs, http://cs101.openjudge.cn/practice/12029/

思路:
看题解的 bfs 实现法。
```

```
from collections import deque
import sys
input = sys.stdin.read

def is_valid(x, y, m, n):
    return 0 <= x < m and 0 <= y < n

def bfs(start_x, start_y, start_height, m, n, h, water_height):
    dx = [-1, 1, 0, 0]</pre>
```

```
dy = [0, 0, -1, 1]
    q = deque([(start x, start y, start height)])
   water height[start x][start y] = start height
   while q:
        x, y, height = q.popleft()
        for i in range(4):
            nx, ny = x + dx[i], y + dy[i]
            if is valid(nx, ny, m, n) and h[nx][ny] < height:</pre>
                if water_height[nx][ny] < height:</pre>
                    water height[nx][ny] = height
                    q.append((nx, ny, height))
def main():
   data = input().split()
    idx = 0
   k = int(data[idx])
   idx += 1
   results = []
    for in range(k):
        m, n = map(int, data[idx:idx + 2])
        idx += 2
        h = []
        for i in range(m):
            h.append(list(map(int, data[idx:idx + n])))
            idx += n
        water_height = [[0] * n for _ in range(m)]
        i, j = map(int, data[idx:idx + 2])
        idx += 2
        i, j = i - 1, j - 1
        p = int(data[idx])
        idx += 1
        for in range(p):
            x, y = map(int, data[idx:idx + 2])
            idx += 2
            x, y = x - 1, y - 1
            if h[x][y] <= h[i][j]:
                continue
            bfs(x, y, h[x][y], m, n, h, water_height)
```

```
results.append("Yes" if water_height[i][j] > 0 else "No")

sys.stdout.write("\n".join(results) + "\n")

if __name__ == "__main__":
    main()
```

```
      據代码
      基本信息

      from collections import deque import sys input = sys.stdin.read
      題目: 12029

      def is_valid(x, y, m, n): return 0 <= x < m and 0 <= y < n</td>
      语言: Python3

      提交时间: 2024-11-26 23:22:49
```

02802: 小游戏

bfs, http://cs101.openjudge.cn/practice/02802/

思路:

一眼做不出来,直接学习题解代码。

```
import heapq
num1=1
while True:
    w,h=map(int,input().split())
    if w==0 and h==0:
        break
    print(f"Board #{num1}:")
    martix=[[" "]*(w+2)]+[[" "]+list(input())+[" "] for _ in
range(h)]+[[" "]*(w+2)]
    dir=[(0,1),(0,-1),(1,0),(-1,0)]
    num2=1
    while True:
        x1,y1,x2,y2=map(int,input().split())
        if x1==0 and x2==0 and y1==0 and y2==0:
             break
        queue,flag=[],False
        vis=set()
        heapq.heappush(queue,(0,x1,y1,-1))
```

```
martix[y2][x2]=" "
    vis.add((-1,x1,y1))
    while queue:
        step,x,y,dirs=heapq.heappop(queue)
        if x==x2 and y==y2:
            flag=True
            break
        for i,(dx,dy) in enumerate(dir):
            px,py=x+dx,y+dy
            if 0<=px<=w+1 and 0<=py<=h+1 and (i,px,py) not in vis
and martix[py][px]!="X":
            vis.add((i,px,py))
            heapq.heappush(queue,(step+(dirs!=i),px,py,i))
        if flag:
            print(f"Pair {num2}: {step} segments.")
        else:
            print(f"Pair {num2}: impossible.")
        martix[y2][x2]="X"
        num2+=1
    print()
    num1+=1</pre>
```

状态: Accepted

```
    源代码
    #: 47417132

    import heapq
numl=1
    題目: 02802

    while True:
w,h-map(int,input().split())
if w==0 and h==0:
break
    内存: 4688kB

    世紀
    財间: 75ms

    语言: Python3
    提交时间: 2024-11-26 23:14:13
```

2. 学习总结和收获

我觉得要把学习策略转向了,很清楚一些难题自己很难理解,即使当时明白了考场上也很难做出来。不在难题上花费很多时间,大概能懂题解就行。省下来的时间开始做一些类似往年期末中等难度的题目和比较套路化、模板化的较难题目。