

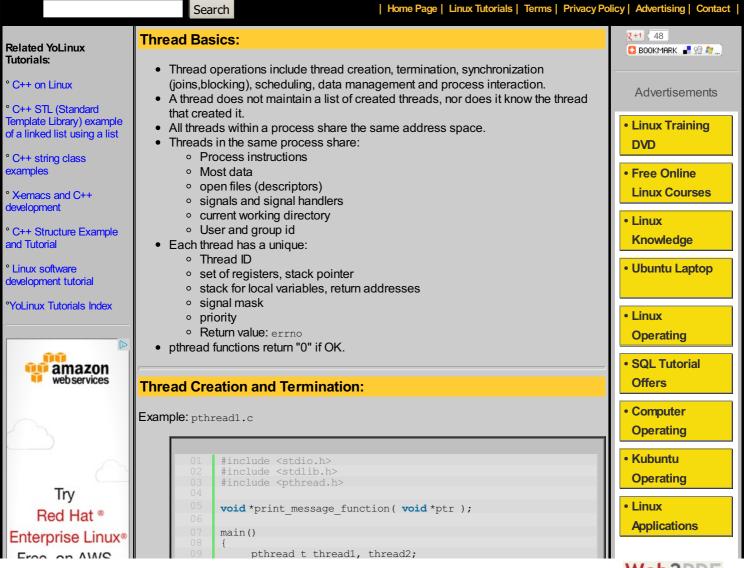
# **POSIX thread (pthread) libraries**

The POSIX thread libraries are a standards based thread API for C/C++. It allows one to spawn a new concurrent process flow. It is most effective on multi-processor or multi-core systems where the process flow can be scheduled to run on another processor thus gaining speed through parallel or distributed processing. Threads require less overhead than "forking" or spawning a new process because the system does not initialize a new system virtual memory space and environment for the process. While most effective on a multiprocessor system, gains are also found on uniprocessor systems which exploit latency in I/O and other system functions which may halt process execution. (One thread may execute while another is waiting for I/O or some other system latency.) Parallel programming technologies such as MPI and PVM are used in a distributed computing environment while threads are limited to a single computer system. All threads within a process share the same address space. A thread is spawned by defining a function and its arguments which will be processed in the thread. The purpose of using the POSIX thread library in your software is to execute software faster.

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```
const char *message1 = "Thread 1";
      const char *message2 = "Thread 2";
      int iret1, iret2;
     /* Create independent threads each of which will execute
function */
      iret1 = pthread create( &thread1, NULL,
print message function, (void*) message1);
      iret2 = pthread_create( &thread2, NULL,
print_message_function, (void*) message2);
      /* Wait till threads are complete before main continues.
Unless we
      ^{\prime \star} wait we run the risk of executing an exit which will nate ^{\star\prime}
terminate
     /* the process and all threads before the threads have leted. */
completed.
     pthread_join( thread1, NULL);
pthread_join( thread2, NULL);
     printf("Thread 1 returns: %d\n",iret1);
printf("Thread 2 returns: %d\n",iret2);
      exit(0);
}
void *print_message_function( void *ptr )
      char *message;
     message = (char *) ptr;
     printf("%s \n", message);
```

### Compile:

- C compiler: cc -pthread pthread1.c (Or cc -lpthread pthread1.c)
   or
- C++ compiler: g++ -pthread pthread1.c (or g++ -lpthread pthread1.c)

The GNU compiler now has the command line option "-pthread" while older versions of the compiler specify the pthread library explicitly with "-lpthread".

Run: ./a.out

## Results:

```
Thread 1
Thread 2
Thread 1 returns: 0
Thread 2 returns: 0
```

## Details:

- In this example the same function is used in each thread. The arguments are different. The functions need not be the same.
- Threads terminate by explicitly calling pthread\_exit, by letting the function return, or by a call to the function exit which will terminate the process including any threads.
- Function call: <a href="mailto:pthread\_create">pthread\_create</a> create a new thread

## Arguments:

- thread returns the thread id. (unsigned long int defined in bits/pthreadtypes.h)
- attr Set to NULL if default thread attributes are used. (else define members of the struct pthread\_attr\_t defined in bits/pthreadtypes.h)
   Attributes include:
  - detached state (joinable? Default: PTHREAD\_CREATE\_JOINABLE. Other option: PTHREAD\_CREATE\_DETACHED)
  - scheduling policy (real-time?
  - PTHREAD INHÉRIT SCHED, PTHREAD EXPLICIT SCHED, SCHED OTHER)
  - scheduling parameter
  - inheritsched attribute (Default: PTHREAD\_EXPLICIT\_SCHED Inherit from parent thread: PTHREAD INHERIT SCHED)



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- scope (Kernel threads: PTHREAD\_SCOPE\_SYSTEM User threads: PTHREAD SCOPE PROCESS Pick one or the other not both.)
- guard size
- stack address (See unistd.h and bits/posix\_opt.h POSIX\_THREAD\_ATTR\_STACKADDR)
- stack size (default minimum PTHREAD\_STACK\_SIZE set in pthread.h),
- void \* (\*start\_routine) pointer to the function to be threaded. Function has a single argument: pointer to void.
- \*arg pointer to argument of function. To pass multiple arguments, send a pointer to a structure.
- Function call: pthread\_join wait for termination of another thread

```
int pthread_join(pthread_t th, void **thread_return);
```

#### Arguments:

- th thread suspended until the thread identified by th terminates, either by calling pthread exit() or by being cancelled.
- thread\_return If thread\_return is not NULL, the return value of th is stored in the location pointed to by thread\_return.
- Function call: pthread\_exit terminate the calling thread

```
void pthread_exit(void *retval);
```

#### Arguments:

retval - Return value of thread.

This routine kills the thread. The pthread\_exit function never returns. If the thread is not detached, the thread id and return value may be examined from another thread by using pthread join.

Note: the return pointer \*retval, must not be of local scope otherwise it would cease to exist once the thread terminates.

• [C++ pitfalls]: The above sample program will compile with the GNU C and C++ compiler g++. The following function pointer representation below will work for C but not C++. Note the subtle differences and avoid the pitfall below:

```
void print_message_function( void *ptr );
...
iret1 = pthread_create( &thread1, NULL,
    (void*) &print_message_function, (void*) message1);
...
...
```

# Thread Synchronization:

The threads library provides three synchronization mechanisms:

- mutexes Mutual exclusion lock: Block access to variables by other threads. This
  enforces exclusive access by a thread to a variable or set of variables.
- joins Make a thread wait till others are complete (terminated).
- condition variables data type pthread cond t

## **Mutexes:**

Mutexes are used to prevent data inconsistencies due to operations by multiple threads upon the same memory area performed at the same time or to prevent race conditions where an order of operation upon the memory is expected. A contention or race condition often occurs when two or more threads need to perform operations on the same memory area, but the results of computations depends on the order in which these operations are performed. Mutexes are used for serializing shared resources such as memory. Anytime a global resource is accessed by more than one thread the resource should have a Mutex associated with it. One can apply a mutex to protect a segment of memory ("critical region") from other threads. Mutexes can be applied only to threads in a single process and do not work between processes as do semaphores.

Example threaded function:

Without Mutex	With Mutex





```
int
counter=0;

/* Function
C */
void
functionC()

counter++

    counter++

    counter++
    }

/* Note scope of variable and mutex are
the same */
pthread mutex t mutex1 =
PTHREAD_MUTEX_INITIALIZER;
int counter=0;

/* Function C */
void functionC()
{
    pthread_mutex_lock(&mutex1);
    counter++
    pthread_mutex_unlock(&mutex1);
}
```

Possible execution sequence						
Thread 1	Thread 2	Thread 1	Thread 2			
counter = 0	counter = 0	counter = 0	counter = 0			
counter =	counter =	counter =	Thread 2 locked out. Thread 1 has exclusive use of variable counter			
			counter = 2			

If register load and store operations for the incrementing of variable <code>counter</code> occurs with unfortunate timing, it is theoretically possible to have each thread increment and overwrite the same variable with the same value. Another possibility is that thread two would first increment <code>counter</code> locking out thread one until complete and then thread one would increment it to 2.

Sequence	Thread 1	Thread 2
1	counter = 0	counter=0
	Thread 1 locked out. Thread 2 has exclusive use of variable counter	counter = 1
3	counter = 2	

Code listing: mutex1.c

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>
void *functionC();
pthread_mutex_t mutex1 = PTHREAD_MUTEX_INITIALIZER;
int counter = 0;
main()
   int rc1, rc2;
   pthread_t thread1, thread2;
    /* Create independent threads each of which will execute
if( (rc1=pthread_create( &thread1, NULL, &functionC, NULL)) )
      printf("Thread creation failed: %d\n", rc1);
if( (rc2=pthread_create( &thread2, NULL, &functionC, NULL)) )
      printf("Thread creation failed: %d\n", rc2);
   /* Wait till threads are complete before main continues.
Unless we
   /\star wait we run the risk of executing an exit which will
terminate */
   /* the process and all threads before the threads have completed. */
   pthread_join( thread1, NULL);
pthread_join( thread2, NULL);
   exit(0);
void *functionC()
```

```
37
38
39
counter++;
40
41
42
}
pthread_mutex_lock( &mutex1 );
counter value: %d\n",counter);
pthread_mutex_unlock( &mutex1 );
}
```

Compile: cc -pthread mutex1.c (or cc -lpthread mutex1.c for older versions of the GNU compiler which explicitly reference the library)

Run: ./a.out Results:

Counter value: 1
Counter value: 2

When a mutex lock is attempted against a mutex which is held by another thread, the thread is blocked until the mutex is unlocked. When a thread terminates, the mutex does not unless explicitly unlocked. Nothing happens by default.

#### Man Pages:

- pthread\_mutex\_lock() acquire a lock on the specified mutex variable. If the mutex
  is already locked by another thread, this call will block the calling thread until the
  mutex is unlocked.
- pthread\_mutex\_unlock() unlock a mutex variable. An error is returned if mutex is already unlocked or owned by another thread.
- pthread\_mutex\_trylock() attempt to lock a mutex or will return error code if busy.
   Useful for preventing deadlock conditions.

## Joins:

A join is performed when one wants to wait for a thread to finish. A thread calling routine may launch multiple threads then wait for them to finish to get the results. One waits for the completion of the threads with a join.

Sample code: join1.c

```
#include <stdio.h>
#include <pthread.h>
#define NTHREADS 10
void *thread function(void *);
pthread_mutex_t mutex1 = PTHREAD_MUTEX_INITIALIZER;
int counter = 0;
   pthread t thread id[NTHREADS];
   inti, j;
    for(i=0; i < NTHREADS; i++)</pre>
      pthread create ( &thread id[i], NULL, thread function,
NULL );
 for(j=0; j < NTHREADS; j++)</pre>
       pthread join( thread id[j], NULL);
    /* Now that all threads are complete I can print the final
result.
   /* Without the join I could be printing a value before all
the threads */
/* have been
completed.
   printf("Final counter value: %d\n", counter);
void *thread function(void *dummyPtr)
   printf("Thread number %ld\n", pthread self());
   pthread mutex lock( &mutex1 );
   pthread mutex unlock( &mutex1 );
```



```
Compile: cc -pthread join1.c (or cc -lpthread join1.c for older versions of the GNU compiler which explicitly reference the library)
Run: ./a.out
Results:

Thread number 1026
Thread number 2051
Thread number 3076
Thread number 4101
Thread number 5126
Thread number 5126
Thread number 6151
Thread number 6151
Thread number 7176
Thread number 9226
Thread number 9226
Thread number 10251
Final counter value: 10
```

## Man Pages:

- pthread\_create() create a new thread
- pthread\_join() wait for termination of another thread
- pthread self() return identifier of current thread

# **Condition Variables:**

A condition variable is a variable of type pthread\_cond\_t and is used with the appropriate functions for waiting and later, process continuation. The condition variable mechanism allows threads to suspend execution and relinquish the processor until some condition is true. A condition variable must always be associated with a mutex to avoid a race condition created by one thread preparing to wait and another thread which may signal the condition before the first thread actually waits on it resulting in a deadlock. The thread will be perpetually waiting for a signal that is never sent. Any mutex can be used, there is no explicit link between the mutex and the condition variable.

Man pages of functions used in conjunction with the condition variable:

- · Creating/Destroying:
  - o pthread cond init
  - o pthread\_cond\_t cond = PTHREAD\_COND\_INITIALIZER;
  - pthread\_cond\_destroy
- · Waiting on condition:
  - pthread\_cond\_wait unlocks the mutex and waits for the condition variable cond to be signaled.
  - pthread cond timedwait place limit on how long it will block.
- Waking thread based on condition:
  - pthread\_cond\_signal restarts one of the threads that are waiting on the condition variable cond.
  - pthread\_cond\_broadcast wake up all threads blocked by the specified condition variable.

Example code: cond1.c

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>

pthread_mutex t count mutex = PTHREAD_MUTEX_INITIALIZER;
pthread_cond_t condition_var = PTHREAD_COND_INITIALIZER;

void *functionCount1();

void *functionCount2();
int count = 0;
#define COUNT_DONE 10
#define COUNT_HALT1 3
#define COUNT_HALT2 6

main()
{
pthread_t thread1, thread2;

pthread_create(&thread1, NULL, &functionCount1, NULL);
pthread_create(&thread2, NULL, &functionCount2, NULL);

pthread_join(thread1, NULL);
pthread_join(thread2, NULL);

printf("Final count: %d\n",count);
```

```
exit(0);
// Write numbers 1-3 and 8-10 as permitted by
functionCount2()
void *functionCount1()
   for(;;)
       // Lock mutex and then wait for signal to relase mutex
      pthread_mutex_lock( &count_mutex );
      // Wait while functionCount2() operates on count
// mutex unlocked if condition variable in
functionCount2() signaled.
      pthread cond wait( &condition var, &count mutex );
      printf("Counter value functionCount1: %d\n",count);
      pthread mutex unlock( &count mutex );
      if(count >= COUNT DONE) return(NULL);
}
// Write numbers 4-7
void *functionCount2()
    for(;;)
       pthread mutex lock( &count mutex );
       if( count < COUNT HALT1 || count > COUNT HALT2 )
           // Condition of if statement has been met.
           // Signal to free waiting thread by freeing the
mutex.
          // Note: functionCount1() is now permitted to
modify "count"
          pthread_cond_signal( &condition var );
       else
          count++;
          printf("Counter value functionCount2: %d\n", count);
       pthread mutex unlock ( &count mutex );
       if(count >= COUNT_DONE) return(NULL);
```

Compile: cc -pthread cond1.c (or cc -lpthread cond1.c for older versions of the GNU compiler which explicitly reference the library)

Run: ./a.out Results:

```
Counter value functionCount1: 1
Counter value functionCount1: 2
Counter value functionCount1: 3
Counter value functionCount2: 4
Counter value functionCount2: 5
Counter value functionCount2: 7
Counter value functionCount1: 8
Counter value functionCount1: 9
Counter value functionCount1: 10
Final count: 10
```

Note that functionCount1() was halted while count was between the values COUNT\_HALT1 and COUNT\_HALT2. The only thing that has been ensures is that functionCount2 will increment the count between the values COUNT\_HALT1 and COUNT\_HALT2. Everything else is random.

The logic conditions (the "if" and "while" statements) must be chosen to insure that the "signal" is executed if the "wait" is ever processed. Poor software logic can also lead to a deadlock condition.

Note: Race conditions abound with this example because count is used as the condition and can't be locked in the while statement without causing deadlock.



# Thread Scheduling:

When this option is enabled, each thread may have its own scheduling properties. Scheduling attributes may be specified:

- · during thread creation
- by dynamically by changing the attributes of a thread already created
- by defining the effect of a mutex on the thread's scheduling when creating a mutex
- by dynamically changing the scheduling of a thread during synchronization operations.

The threads library provides default values that are sufficient for most cases.

# **Thread Pitfalls:**

- Race conditions: While the code may appear on the screen in the order you wish
  the code to execute, threads are scheduled by the operating system and are
  executed at random. It cannot be assumed that threads are executed in the order
  they are created. They may also execute at different speeds. When threads are
  executing (racing to complete) they may give unexpected results (race condition).
  Mutexes and joins must be utilized to achieve a predictable execution order and
  outcome.
- Thread safe code: The threaded routines must call functions which are "thread safe". This means that there are no static or global variables which other threads may clobber or read assuming single threaded operation. If static or global variables are used then mutexes must be applied or the functions must be rewritten to avoid the use of these variables. In C, local variables are dynamically allocated on the stack. Therefore, any function that does not use static data or other shared resources is thread-safe. Thread-unsafe functions may be used by only one thread at a time in a program and the uniqueness of the thread must be ensured. Many non-reentrant functions return a pointer to static data. This can be avoided by returning dynamically allocated data or using caller-provided storage. An example of a non-thread safe function is strtok which is also not re-entrant. The "thread safe" version is the re-entrant version strtok r.
- Mutex Deadlock: This condition occurs when a mutex is applied but then not "unlocked". This causes program execution to halt indefinitely. It can also be caused by poor application of mutexes or joins. Be careful when applying two or more mutexes to a section of code. If the first pthread\_mutex\_lock is applied and the second pthread\_mutex\_lock fails due to another thread applying a mutex, the first mutex may eventually lock all other threads from accessing data including the thread which holds the second mutex. The threads may wait indefinitely for the resource to become free causing a deadlock. It is best to test and if failure occurs, free the resources and stall before retrying.

The order of applying the mutex is also important. The following code segment illustrates a potential for deadlock:



```
pthread_mutex_lock(&lock1);
...

pthread_mutex_lock(&lock2);
pthread_mutex_lock(&lock2);
pthread_mutex_lock(&lock1);
...
pthread_mutex_lock(&lock1);
pthread_mutex_lock(&lock1);
pthread_mutex_lock(&lock2);
...

main()

main()

main()

pthread_create(&thread1, NULL, function1, NULL);
pthread_create(&thread2, NULL, function2, NULL);
...
}
```

If function1 acquires the first mutex and function2 acquires the second, all resources are tied up and locked.

Condition Variable Deadlock: The logic conditions (the "if" and "while" statements)
must be chosen to insure that the "signal" is executed if the "wait" is ever
processed.

# Thread Debugging:

- GDB:
  - Debugging Programs with Multiple Threads
  - GDB: Stopping and starting multi-thread programs
  - GDB/MI: Threads commands
- DDD:
  - · Examining Threads

# Thread Man Pages:

- pthread atfork register handlers to be called at fork(2) time
- pthread\_attr\_destroy [pthread\_attr\_init] thread creation attributes
- pthread\_attr\_getdetachstate [pthread\_attr\_init] thread creation attributes
- pthread\_attr\_getguardsize get the guardsize attribute in the attr object.
- pthread\_attr\_getinheritsched [pthread\_attr\_init] thread creation attributes
- pthread\_attr\_getschedparam [pthread\_attr\_init] thread creation attributes
- pthread\_attr\_getschedpolicy [pthread\_attr\_init] thread creation attributes
- pthread\_attr\_getscope [pthread\_attr\_init] thread creation attributes
- pthread\_attr\_getstack get the thread creation stack attributes stackaddr and stacksize in the attr object.
- pthread\_attr\_getstackaddr get the thread creation stackaddr attributes stackaddr attribute in the attr object.
- pthread\_attr\_getstacksize get the thread creation stacksize attribute in the attr object.
- pthread\_attr\_init thread creation attributes
- pthread\_attr\_setdetachstate [pthread\_attr\_init] thread creation attributes
- pthread\_attr\_setguardsize set the guardsize attribute in the attr object.
- pthread\_attr\_setinheritsched [pthread\_attr\_init] thread creation attributes
- pthread\_attr\_setschedparam [pthread\_attr\_init] thread creation attributes
- pthread\_attr\_setschedpolicy [pthread\_attr\_init] thread creation attributes
- pthread\_attr\_setscope [pthread\_attr\_init] thread creation attributes
- pthread\_attr\_setstack set the thread creation stack attributes stackaddr and stacksize in the attr object.
- pthread\_attr\_setstackaddr set the thread creation stackaddr attributes stackaddr attribute in the attr object.
- pthread\_attr\_setstacksize set the thread creation stacksize attribute in the attr object.
- pthread\_cancel thread cancellation
- pthread\_cleanup\_pop [pthread\_cleanup\_push] install and remove cleanup handlers
- pthread\_cleanup\_pop\_restore\_np [pthread\_cleanup\_push] install and remove cleanup handlers
- pthread\_cleanup\_push install and remove cleanup handlers



- pthread\_cleanup\_push\_defer\_np [pthread\_cleanup\_push] install and remove cleanup handlers
- pthread\_condattr\_destroy [pthread\_condattr\_init] condition creation attributes
- pthread\_condattr\_init condition creation attributes
- pthread\_cond\_broadcast [pthread\_cond\_init] operations on conditions
- pthread cond destroy [pthread cond init] operations on conditions
- pthread\_cond\_init operations on conditions
- pthread\_cond\_signal [pthread\_cond\_init] operations on conditions
- pthread\_cond\_timedwait [pthread\_cond\_init] operations on conditions
- pthread cond wait [pthread cond init] operations on conditions
- pthread create create a new thread
- pthread detach put a running thread in the detached state
- pthread equal compare two thread identifiers
- pthread\_exit terminate the calling thread
- pthread\_getschedparam [pthread\_setschedparam] control thread scheduling parameters
- pthread\_getspecific [pthread\_key\_create] management of thread-specific data
- pthread\_join wait for termination of another thread
- pthread\_key\_create management of thread-specific data
- pthread\_key\_delete [pthread\_key\_create] management of thread-specific data
- pthread\_kill\_other\_threads\_np terminate all threads in program except calling thread
- pthread kill [pthread sigmask] handling of signals in threads
- pthread\_mutexattr\_destroy [pthread\_mutexattr\_init] mutex creation attributes
- pthread\_mutexattr\_getkind\_np [pthread\_mutexattr\_init] mutex creation attributes
- pthread\_mutexattr\_init mutex creation attributes
- pthread\_mutexattr\_setkind\_np [pthread\_mutexattr\_init] mutex creation attributes
- pthread mutex destroy [pthread mutex init] operations on mutexes
- pthread\_mutex\_init operations on mutexes
- pthread mutex lock [pthread mutex init] operations on mutexes
- pthread mutex trylock [pthread mutex init] operations on mutexes
- pthread\_mutex\_unlock [pthread\_mutex\_init] operations on mutexes
- pthread once once-only initialization
- pthread\_self return identifier of current thread
- pthread\_setcancelstate [pthread\_cancel] thread cancellation
- pthread\_setcanceltype [pthread\_cancel] thread cancellation
- pthread\_setschedparam control thread scheduling parameters
- pthread\_setspecific [pthread\_key\_create] management of thread-specific data
- pthread\_sigmask handling of signals in threads
- pthread testcancel [pthread cancel] thread cancellation

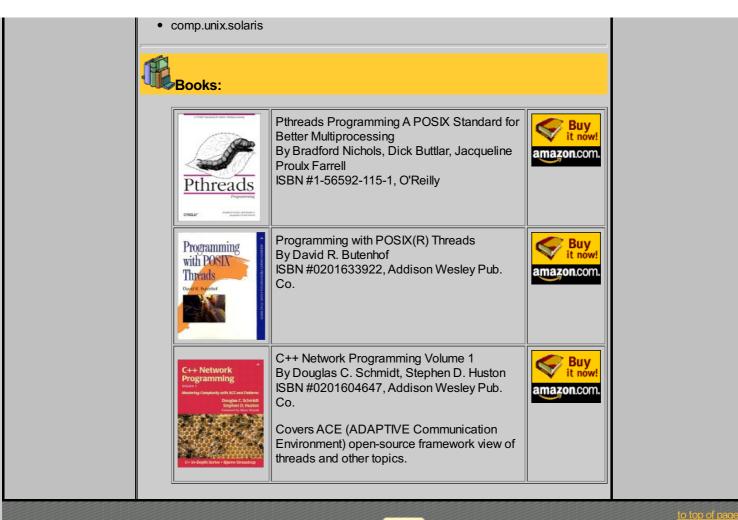
### Links:

- Fundamentals Of Multithreading Paul Mazzucco
- Native Posix Thread Library for Linux
- Posix threads for MS/Win32: [Announcement / description] sourceforge home page
- Introduction to Programming Threads
- GNU Portable Threads
- Comparison of thread implementations
- comp.programming.threads FAQ
- Examples
- Pthreads tutorial and examples of thread problems by Andrae Muys
- Helgrind: Valgrind KDE thread checker
- Sun's Multithreaded Programming Guide Not Linux but a good reference.
- Platform independent threads:
  - Gnome GLib 2 threads Thread abstraction; including mutexes, conditions and thread private data. [example]
  - OmniORB (CORBA) Thread Library
  - zThreads
- · C++ Thread classes:
  - GNU: Common C++ support for threading, sockets, file access, daemons, persistence, serial I/O, XML parsing and system services
  - ACE: Adaptive Communication Environment C++ interface
    - ACE programmers guide: [pdf] (see page 29 for threads)
  - C++ Thread classes sourceforge
  - QpThread

# News Groups:

comp.programming.threads













Windows Virtualization

www.Symantec.com/Virtualization Secure, Backup & Recover Virtual Machines with Symantec V-Ray.

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