

Longest common substring (DP)

```
# Time c. O(nm), space c. O(nm)
def longest_common_substring(s1, s2):
    m = [[0] * (1+len(s2)) for i in \
          range(1+len(s1))]
    max_len, end_idx = 0, 0
    for i in range(1, 1 + len(s1)):
        for j in range(1, 1 + len(s2)):
            if s1[i - 1] == s2[j - 1]:
                m[i][j] = m[i-1][j-1] + 1
                if m[i][j] > max_len:
                    max_len = m[i][j]
                    end_idx = i
            else:
                m[i][j] = 0
    return s1[end_idx-max_len : end_idx]
```

Longest common subsequence (DP)

```
# time c. O(nm), space c.
def lcs(s1, s2):
    matrix = [ [ '' for x in range(len(s2))]\
               for y in range(len(s1)) ]
    for i in range(len(s1)):
        for j in range(len(s2)):
            if s1[i] == s2[j]:
                if i == 0 or j == 0:
                    matrix[i][j] = s1[i]
                else:
                    matrix[i][j] = matrix[i-1][j-1] +\
                                   s1[i]
            else:
                matrix[i][j] = max( matrix[i-1][j],
                                    matrix[i][j-1],
                                    key=len)
    return matrix[-1][-1]
```

Make sentences with dictionary

```
def make_sent(string, dictionaries):
    global count
    if len(string) == 0:
        return True
    for i in range(0, len(string) + 1):
        prefix,suffix=string[:i],string[i:]
        if prefix in dictionaries:
            if suffix in dictionaries or\
               make_sent(suffix, dictionaries):
                count += 1
    return True

count = 0
string1 = "applet"
dictionary1 = {'app', 'let', 'apple', 't', 'applet'}
print( make_sent(string1, dictionary1) )
print( count )
```

True

3

Longest common substring in array

```
# time c. O(n*n*(n + n-1 + n-2 ... etc.)?
def longest_substr(arr):

    if len(arr) <= 1:
        return ''
    res      = ''
    reference = arr[0]
    for i in range(len(reference)):
        for j in range(len(reference)-i+1):
            candidate = reference[i:i+j]
            if j > len(res) and\
               all(candidate in x for x in arr):
                res = candidate
    return res
```

```
def max_subarray_sum(arr):
    if len(arr)==0: return 0
    max_sum = curr_sum = arr[0] # if neg
    start, tstart, end = 0, 0, 0
    for i in range(1, len(arr)):
        if arr[i] > curr_sum + arr[i]:
            curr_sum = arr[i]
            tstart = i
        else:
            curr_sum += arr[i]
        if curr_sum > max_sum:
            max_sum = curr_sum
            start = tstart
            end = i
    return max_sum, arr[ start:end+1 ]
```

```
# Time c. O(n), space c. O(n)
from collections import defaultdict
def findSubarraysWithSumK(arr, k):
    hash_map = defaultdict(list) # key=sum, value=end indices
    res = [] # subarrays
    curr_sum = 0
    for i in range( len(arr) ):
        curr_sum += arr[i]
        if curr_sum == k:
            res.append(arr[:i+1]) # or res.append((0, i+1))
        if curr_sum-k in hash_map:
            for value in hash_map[ curr_sum-k ]:
                res.append( arr[ value+1:i+1 ] )
                # or res.append((value+1, i+1))
        hash_map[ curr_sum ].append(i)
    return res
```

All non-contig. pairs sum up to k (unsorted)

```
def pair_sum(arr, k):
    if len(arr) < 2: return
    seen, output = set(), set()
    for num in arr:
        if k-num not in seen:
            seen.add(num)
        else:
            output.add( (min(num,k-num), max(num,k-num)) )
    return '\n'.join(map(str,list(output)))
```

Non-contig pair, sum closest to k (sorted)

```
# time c. O(n)
MAX_VAL = 1000000000
def closest_sum(arr, n, k):
    res = 0, 0
    l, r, diff = 0, n-1, MAX_VAL
    while l < r:
        if abs(arr[l] + arr[r] - k) < diff:
            res = l, r
            diff = abs(arr[l] + arr[r] - k)
        if arr[l] + arr[r] > k:
            r -= 1
        else:
            l += 1
    return arr[res[0]], arr[res[1]]
```

Partition - 2 subsets w/equal sum?

```
# t.c.= O(n*sum), space=O(sum)
def isSubsetSum_dp(arr, n):
    sum_ = 0
    for i in range(n): sum_ += arr[i]
    if sum_ % 2 != 0: return 0
    part = [0] * ((sum_ // 2) + 1) # 0 (False)
    # part[j]=true if exists subset w/sum=j
    for i in range(n):
        for j in range(sum_ // 2, arr[i]-1, -1):
            if (part[j - arr[i]] == 1 or j == arr[i]):
                part[j] = 1
    return part[sum_ // 2]
```

Is there subset whose sum=sum_?

```
# t.c.=2**n - each elem included or excluded
#space=O(n)-call stack in rec.=depth up to n
def isSubsetSum(arr, n, sum_):
    if sum_ == 0: # base case 1 - always possible
        return True
    if n == 0 and sum_ != 0:
        return False # base c. 2 - no elems left
    if arr[n-1] > sum_:
        return isSubsetSum(arr, n-1, sum_)
    return isSubsetSum(arr, n-1, sum_) or\
        isSubsetSum(arr, n-1, sum_-arr[n-1])
```

Smallest +int != sum of a subset

- Sort arr. If arr[i] > res => found a gap
- Elems after arr[i] > res. res = solution

$$[1, 2, 5] \Rightarrow 4$$

```
# time c. O(n) - sorted array
def findSmallest(arr, n):
    res = 1
    for i in range(n):
        if arr[i] <= res:
            res = res + arr[i]
        else:
            break
    return res
```

Partition - two subsets w/min diff. sums

$n * \text{sumTotal}$ w/memoization

```
# t.c.= O(2**n)-each step, 2 cases. Space=max depth of rec.=n
def findMinRec(arr, i, sumCalculated, sumTotal):
    # If last elem, sum of one subset = sumCalculated,
    # sum of the other=(sumTotal-sumCalculated). Take abs diff
    if i == 0:
        return abs( (sumTotal-sumCalculated) - sumCalculated )
    # For each arr[i-1], (1) we can exclude it from first set,
    # (2) we can include it. return min of two choices
    return min( findMinRec( arr, i-1, sumCalculated + arr[i-1], sumTotal),
                findMinRec( arr, i-1, sumCalculated, sumTotal ) )
    include memo in each call
```

```
def findmin_recursive(arr, n):
    # total sum of arr
    sumTotal = 0
    for i in range(n):
        sumTotal += arr[i]
    return findMinRec(arr, n, 0, sumTotal)
```

Balanced parenth. check

```
def balance_check(s):
    if len(s)%2 != 0: return False
    opening = set('([{')
    # matching pairs
    matches = set([ ('(', ')'),
                   ('[', ']'),
                   ('{', '}'), ]))
    stack = []
    for paren in s:
        if paren in opening:
            stack.append(paren)
        else:
            if len(stack) == 0:
                return False
            last_open = stack.pop()
            if (last_open,paren) not in matches:
                return False
    return len(stack) == 0
```

All non-contig. triplets w/ sum < k

```
# Time c. O(n^2); sort() + smart iteration
def count_triplets(arr, k):
    arr.sort()
    n = len(arr)
    count = 0
    for i in range(n-2):
        start = i + 1
        end = n - 1
        while start < end:
            if (arr[i] + arr[start] + arr[end]) >= k:
                end -= 1
            else:
                count += (end-start) # array sorted
                start += 1
    return count
```

Count paths

```
# Time c. = O(c**n) exponential
def count_paths_rec(m, n):
    if m == 1 or n == 1: return 1
    return count_paths_rec(m-1, n) + count_paths_rec(m, n-1)
```

```
# Time & space = O(mn)
def count_paths(m, n):
    if m < 1 or n < 1: return -1
    count = [[None for j in range(n)] for i in range(m)]
    for i in range(n): count[0][i] = 1
    for j in range(m): count[j][0] = 1
    for i in range(1, m):
        for j in range(1, n):
            count[i][j] = count[i - 1][j] + count[i][j - 1]
    return count[m - 1][n - 1]
```

```
# Time O(mn), space O(n)
def count_paths_dp(m, n):
    dp = [1 for i in range(n)]
    for i in range(1, m):
        for j in range(1, n):
            dp[j] = dp[j] + dp[j - 1]
    return dp[n - 1]
```

All non-contig. triplets w/sum=0

```
# Time c. O(n^2), space c. O(1)
def findTriplets(arr):
    arr.sort()
    n = len(arr)
    count = 0
    for i in range(n-2):
        start = i + 1
        end = n - 1
        while start < end:
            if arr[i] + arr[start] + arr[end] == 0:
                count += 1
                start += 1
                end -= 1
            elif (arr[i] + arr[start] + arr[end] < 0):
                start += 1
            else:
                end -= 1
    return count
```

Get all permutations of string

```
def permute(s):
    out = []
    if len(s) == 1:
        out = [s]
    else:
        for i, char in enumerate(s):
            for perm in permute(s[:i] + s[i+1:]):
                out += [char + perm]
    return out
```

```
permute('abc')
```

```
['abc', 'acb', 'bac', 'bca', 'cab', 'cba']
```

Longest Substr. w/K Chars

```
# Time c. O(n), space c. O(1)
# At most k distinct characters
def LongestSubstring(s, k):
    n = len(s)
    if n*k==0:
        return 0
    l, r = 0, 0
    mapp = dict()
    max_len = 0
    for c in s:
        mapp[c] = r
        r += 1
        if len(mapp) == k+1:
            del_idx = min(mapp.values())
            del mapp[s[del_idx]]
            l = del_idx + 1
            max_len = max(max_len, r - l)
    return max_len
```

Longest Increasing Path in Matrix

```
# time and space = O(mn)
def longestIncreasingPath(matrix):
    def dfs(i, j):
        if not dp[i][j]:
            val = matrix[i][j]
            dp[i][j] = 1 + max(
                dfs(i - 1, j) if i and val > matrix[i - 1][j] else 0,
                dfs(i + 1, j) if i < M - 1 and val > matrix[i + 1][j] else 0,
                dfs(i, j - 1) if j and val > matrix[i][j - 1] else 0,
                dfs(i, j + 1) if j < N - 1 and val > matrix[i][j + 1] else 0)
        return dp[i][j]
    if not matrix or not matrix[0]: return 0
    M, N = len(matrix), len(matrix[0])
    dp = [[0] * N for i in range(M)]
    return max(dfs(x, y) for x in range(M) for y in range(N))
```

Edit Levenshtein Distance

```
def levenshtein(s1, s2):
    if len(s1) > len(s2):
        s1, s2 = s2, s1
    distances = range(len(s1) + 1)
    for idx2, char2 in enumerate(s2):
        distances_ = [idx2 + 1]
        for idx1, char1 in enumerate(s1):
            if char1 == char2:
                distances_.append(distances[idx1])
            else:
                distances_.append(1 + min((distances[idx1], \
                    distances[idx1 + 1], distances_[-1])))
        distances = distances_
    return distances[-1]
```

Multiply strings

```
# GET NUMBER FROM CHAR: ord(char) - ord('0')
def multiply_strings(num1, num2):
    res = 0
    for i, c1 in enumerate(num1[::-1]):
        for j, c2 in enumerate(num2[::-1]):
            res += (ord(c1) - ord('0')) * \
                (ord(c2) - ord('0')) * \
                (10 ** (i + j))
    return str(res)
num1, num2 = '123', '456'
multiply_strings(num1, num2)
```

Longest Substr. w/Unique Ch

```
def lengthOfLongestSubstring2(s):
    mapp = {c: 0 for c in s}
    left = 0
    max_len = 0
    for right, c in enumerate(s):
        mapp[c] += 1
        while mapp[c] > 1:
            mapp[s[left]] -= 1
            left += 1
        max_len = max(max_len, right - left + 1)
    return max_len
```

Complexity classes

- O(1) constant-time algo - **no dependency on input size**. E.g. **formula** calculating an answer OR **list[idx]**.
- O(logn) algo **halves / reduces input size at each step**. Logarithmic because $\log_2 n = \# \text{ times to divide } n \text{ by 2 to get 1}$.
- O(sqrt(n)) *slower than O(logn), but faster than O(n)*; $\sqrt{n} = \sqrt{n} / n$, so \sqrt{n} lies **in the middle of input**.
- O(n) **iteration over the input** - accessing each input element at least once before reporting the answer.
- O(nlogn) often indicates that algo **sorts the input** OR algo uses **data structure** where *each operation takes O(logn) time*.
- O(n^2) two **nested** loops
- O(n^3) three **nested** loops.
- O(2^n) algo iterates through **all subsets of input elements**. E.g. subsets of {1,2,3}: {1}, {2}, {3}, {1,2}, {1,3}, {2,3} and {1,2,3}.
- O(n!) algo iterates through **all permutations of input elements**. E.g. permutations of {1,2,3}: (1,2,3), (1,3,2), (2,1,3), (2,3,1), (3,1,2) and (3,2,1).

Polynomial algo - time complexity of **O(n^k)** where k is const; if k small - algo efficient. All above except O(2^n) and O(n!) are polynomial. There are many important problems with **no known polynomial (=efficient) algo**, e.g. **NP-hard** problems

133. Clone Connected Undirected Graph



```
# t.c. O(n+m), n=nodes, m=edges s.c. O(n)
# Approach: BFS where visited is dict[curr_node] = cloned_node
class Node:
    def __init__(self, val=0, neighbors=None):
        self.val = val
        self.neighbors = neighbors if neighbors is not None else []
def cloneGraph(start: Node) -> Node:
    if not start:
        return start
    visited, queue = {}, [start] # Dict[visited node] = clone
    visited[start] = Node(start.val, []) # Clone start
    while queue:
        vertex = queue.pop(0)
        for neighbor in vertex.neighbors: # Iterate neighbors
            if neighbor not in visited:
                visited[neighbor] = Node(neighbor.val, [])#Clone neighbor
                queue.append(neighbor)
            # Append cloned neighbor
            visited[vertex].neighbors.append(visited[neighbor])
    return visited[start]
```

Merge Sort (mid)

Usage: sorting **linked lists**, inversion count in nearly sorted arr ($i < j$, but $A[i] > A[j]$), external sort (data too big for memory)

```
# time c. O(nLogn), space c. O(n)
def merge_sort(arr):
    # arr of length 1 - returned as is
    if len(arr) > 1:
        mid = len(arr) // 2
        left = merge_sort(arr[ :mid])
        right = merge_sort(arr[mid: ])
        i, j, k = 0, 0, 0
        while i < len(left) and j < len(right):
            if left[i] < right[j]:
                arr[k] = left[i]
                i += 1
            else:
                arr[k] = right[j]
                j += 1
            k += 1
        while i < len(left):
            arr[k] = left[i]
            i += 1
            k += 1
        while j < len(right):
            arr[k] = right[j]
            j += 1
            k += 1
    return arr
```

Quick Sort (pivot)

Pick **pivot** element (*first, last, random, median*), and **partition** array: all smaller elems before pivot, greater elements after pivot
Usage: sorting arrays

```
# O(nLogn), worst case O(n^2)
# time c. O(logn), qualifies as in-place
def quick_sort(arr, low, high):
    if low < high:
        pi = partition(arr, low, high)
        quick_sort(arr, low, pi-1)
        quick_sort(arr, pi+1, high)

def partition(arr, low, high):
    pivot = arr[high]
    i = low - 1
    for j in range(low, high):
        if arr[j] < pivot:
            i += 1
            arr[i], arr[j] = arr[j], arr[i]
    # place pivot in the middle
    arr[i+1], arr[high] = arr[high], arr[i+1]
    return i+1
```

Heap Sort

Usage: sort nearly sorted array, k largest (smallest) elems

Complete binary tree = every level filled, except possibly last, and nodes are in far left

Binary Heap = cbt where parent $>(<)$ children

Procedure: build max heap; replace max w/ last elem; reduce heap by 1; heapify root; repeat Array repr:

root = arr[0];

for any i-th node arr[i]:

a) arr[i-1]/2 = parent node

b) arr[(2*i)+1] = left child

c) arr[(2*i)+2] = right child

time O(nLogn), space in-place

```
def heap_sort(arr):
    n = len(arr)
    for i in range(n, -1, -1): # build heap
        heapify(arr, n, i)
    for i in range(n-1, 0, -1):
        arr[i], arr[0] = arr[0], arr[i]
        heapify(arr, i, 0) # extract elems
```

def heapify(arr, n, i):
 largest = i # root

l = 2 * i + 1

r = 2 * i + 2

if (l < n and arr[i] < arr[l]):

largest = l # if left child exists

if (r < n and arr[largest] < arr[r]):

largest = r # if right exists

if (largest != i):

arr[i], arr[largest] = arr[largest], arr[i]

heapify(arr, n, largest)

Binary Search

```
# time O(Logn), space (O1)
def binary_search(arr, value):
    if len(arr) == 0: return None
    min_idx, max_idx = 0, len(arr)
    while min_idx < max_idx:
        mid = (min_idx + max_idx) // 2
        if arr[mid] == value:
            return mid
        elif arr[mid] < value:
            min_idx = mid + 1
        else: max_idx = mid
    return None
```

Queue with two stacks

```
class Queue2Stacks(object):
    def __init__(self):
        self.instack = []
        self.outstack = []

    def enqueue(self, element):
        self.instack.append(element)

    def dequeue(self):
        if not self.outstack:
            while self.instack:
                self.outstack.append(
                    self.instack.pop())
        )
        return self.outstack.pop()
```

Next greater elem w/stack

[11, 14, 21] => 11:14, 14:21, 21:-1

```
# time O(n)
def printNGE(arr):
    stack = []
    stack.append(arr[0])
    for i in range(1, len(arr)):
        next_ = arr[i]
        if stack:
            elem = stack.pop()
            while elem < next_:
                print(str(elem) + " : " + str(next_))
                if not stack:
                    break
                elem = stack.pop()
            if elem > next_:
                stack.append(elem)
        stack.append(next_)
    while stack:
        elem = stack.pop()
        next_ = -1
        print(str(elem) + " : " + str(next_))

arr = [11, 14, 21, 3]
printNGE(arr)
```

Queue - FIFO, stack - LIFO

Dequeue – double-ended queue

Stack using two queues

```
from queue import Queue
class Stack:
    def __init__(self):
        self.q1 = []
        self.q2 = []
        self.curr_size = 0

    def push(self, x):
        self.curr_size += 1
        self.q2.append(x)
        while self.q1:
            self.q2.append(self.q1.pop(0))
        self.q = self.q1
        self.q1 = self.q2
        self.q2 = self.q

    def pop(self):
        if not self.q1:
            return
        self.curr_size -= 1
        return self.q1.pop(0)

    def peek(self):
        if not self.q1:
            return -1
        return self.q1[0]

    def size(self):
        return self.curr_size
```

Graphs

- vertices / nodes w/names (keys) and payloads (additional info);
- unordered = **undirected graph** vs. **directed graph** (digraph);
- weight on edge - cost from 1 vertex to the other (distance betw. cities)
- Path - # vertices connected by edges.
- Path length - # edges in unweighted g. (n-1) OR sum of weights in weighted g.
- Cycle in directed g. - path starts / ends at same vertex.
- If no cycles - **acyclic graph**; if also directed - directed acyclic graph (**DAG**)

Adjacency matrix = 2D matrix size $n \times n$ if graph has n vertices. Each cell $[i, j]$ indicates whether there is an edge between vertex i and vertex j (if edge - value is 1 or the weight, and 0 if no edge) - $\{[0, 1, 1], [1, 0, 0], [0, 1, 0]\}$

- Quick edge lookup t.c. $O(1)$ to see if edge present, add/remove edge;
- Simplicity of matrix, good for matrix operations (insights);
- Slow iteration - $O(n)$ to get vertex's _adj. list, $O(V^2)$ add vertex
- memory hog: space c. $O(n^2)$ => good for **dense graphs**
- adjacency lists are better for most tasks
- better for **weighted graphs**;
- Undirected graphs - adj. matrix symmetric

Adjacency list r.: edge_list = $\{[0, 1], [0, 2]\}$ or dict[list] (node: verts) or dict[dicts] (if weights)

- Efficient neighb. iteration and time c. $O(1)$ get vertex's adj. list (array indexing)
- Efficient storage (space c. = $O(2E)$)
- Slow edge lookup $O(E)$ ($E = \# v-s$) if edge (i, j) exists

Objects/pointer r.

- time c. $O(n)$ access a node;
- space c. $O(n)$ if nodes only, $O(n^2)$ w/pointers;
- better for **directed graphs** (pointers);
- high search time c.

Incidence matrix

rows = vertices, columns = edges, entries = is vertex incident to edge (bool)

BFS2

More efficient – enqueues node once

```
def bfs2(graph, start):
    visited, queue = {start}, [start]
    while queue:
        vertex = queue.pop(0)
        print(str(vertex) + " ", end="")
        for neighbor in graph[vertex]:
            if neighbor not in visited:
                visited.add(neighbor)
                queue.append(neighbor)

graph = {0: [1, 2], 1: [2], 2: [3], 3: [1, 2]}
bfs2(graph, 0)
```

Graphs

DFS

queue.pop(0) for BFS

```
# time c. O(V+E), space c. O(V)
def dfs(graph, start):
    visited, stack = set(), [start]
    while stack:
        vertex = stack.pop()
        if vertex not in visited:
            print(vertex, end=' ')
            visited.add(vertex)
            stack.extend(graph[vertex] - visited)
    return visited

graph = { 0: set([1, 2]), 1: set([2]),
          2: set([3]), 3: set([1, 2]) }
dfs(graph, 0)
```

0 2 3 1

Shortest path

Always returned first by BFS

```
def shortest_path(graph, start, goal):
    try:
        return next(bfs_paths(graph, start, goal))
    except StopIteration:
        return None
```

```
graph = {'A': set(['B', 'C']),
         'B': set(['A', 'D', 'E']),
         'C': set(['A', 'F']),
         'D': set(['B']),
         'E': set(['B', 'F']),
         'F': set(['C', 'E'])}

shortest_path(graph, 'A', 'F')
['A', 'C', 'F']
```

```
def kahn_toposort( graph ):
    topo = []
    in_degree = { k: 0 for k in graph }
    for k in graph:
        for v in graph[ k ]:
            in_degree[v] += 1
    q = [k for k,v in in_degree.items() if v==0]
    while q:
        k = q.pop()
        topo.append( k )
        for v in graph[ k ]:
            in_degree[ v ] -= 1
            if in_degree[ v ] == 0:
                q.append( v )
    return topo if len(topo)==len(graph) else []
```

Articulation point

(vertex) or
bridge (edge) –
disconnects graph
when removed.
Connectedness verified
by DFS

All Paths from Vertex A to Vertex B

```
def bfs_paths(graph, start, goal):
    queue = [(start, [start])]
    while queue:
        (vertex, path) = queue.pop(0)
        for nxt in graph[vertex] - set(path):
            if nxt == goal:
                yield path + [nxt]
            else:
                queue.append((nxt, path + [nxt]))
```

200. Number of Islands

2D grid of '1's (land) and '0's (water). Return # islands.
 Adjacent '1' land plots connected horiz. or vertically.
Time c.= $O(m^*n)$, **Space c.**= $O(m^*n)$ due to recursion stack.
 Each recursive call contains:

- func's local vars (indices i and j).
- pointer to the grid.
- return address, etc.

```
def numIslands(grid):
    if not grid:
        return 0
    count = 0
    for i in range(len(grid)):
        for j in range(len(grid[0])):
            # island found
            if grid[i][j] == '1':
                count += 1
                # delete found island
                dfs(grid, i, j)
    return count

def dfs(grid, i, j):
    ''' Delete found island '''
    if i<0 or j<0 or i>=len(grid) or\
    j>=len(grid[0]) or grid[i][j] != '1':
        return
    grid[i][j] = '#' #anything, but '1'
    dfs(grid, i+1, j)
    dfs(grid, i-1, j)
    dfs(grid, i, j+1)
    dfs(grid, i, j-1)

grid = [ ["1","1","0","0","0"],
         ["1","1","0","0","0"],
         ["0","0","1","0","0"],
         ["0","0","0","1","1"], ]
print(numIslands(grid))
```

269. Alien Dictionary (FB!)

- Alien language uses Eng ABC w/unknown order of letters
- You have list of strings - words sorted lexicographically.
- Return string of unique letters sorted lexicographically increasing
- No solution - return "". Multiple solutions - return any.
- 'abc' < 'abd', 'abc' < 'abcd'
- Solution:** a) BFS + get letter dependencies, b) toposort
- Time c.** $O(C)$ where C - total length of all words in input list
- Space c.** $O(1)$ or $O(U + \min(U^2, N))O(U + \min(U$

```
def alienOrder(words):
    # BFS + letter dependencies as adj_list
    adj_list = defaultdict(set) # adj_list for each letter
    # in_degree of each unique letter
    in_degree = Counter({c:0 for word in words for c in word})
    for first_word, second_word in zip(words, words[1:]):
        for c, d in zip(first_word, second_word):
            if c != d: # c comes before d
                if d not in adj_list[c]:
                    adj_list[c].add(d)
                    in_degree[d] += 1
                break
        else:
            # against rules - shorter word should be first
            if len(second_word) < len(first_word): return ''
    # TOPOLOGICAL SORT
    output = []
    queue = deque([c for c in in_degree if in_degree[c] == 0])
    while queue:
        c = queue.popleft()
        output.append(c)
        for d in adj_list[c]:
            in_degree[d] -= 1
            if in_degree[d] == 0:
                queue.append(d)
    # not all letters in output - cycle => no valid ordering
    if len(output) < len(in_degree):
        return ''
    return ''.join(output)
```

2a. Depth First Traversals (Binary Tree) BFS (iterative)

```
# time O(n), space O(h), h=height
class Node:
    def __init__(self, val):
        self.val = val
        self.left = None
        self.right = None

# Root Left Right
def Preorder(root):
    if root:
        print(root.val, end=' ')
        Preorder(root.left)
        Preorder(root.right)

# Left Right Root
def Postorder(root):
    if root:
        Postorder(root.left)
        Postorder(root.right)
        print(root.val, end=' '),

# Left Root Right
def Inorder(root):
    if root:
        Inorder(root.left)
        print(root.val, end=' ')
        Inorder(root.right)
```

```
root          = Node(1)
root.left     = Node(2)
root.right    = Node(3)
root.left.left = Node(4)
root.left.right = Node(5)
print("Preorder traversal:")
Preorder(root)
```

```
# time O(n), space O(n)
class Node:
    # create a new node
    def __init__(self, key):
        self.val = key
        self.left = None
        self.right = None

# print height
def LevelOrder(root):
    # base case
    if root is None:
        return
    queue = []
    queue.append(root)
    while(len(queue) > 0):
        node = queue.pop(0)
        print(node.val, end=' ')
        if node.left is not None:
            queue.append(node.left)
        if node.right is not None:
            queue.append(node.right)
```

Array representation of bin. tree:
node i
parent: $(i-1)/2$
left child: $2*i + 1$
right child: $2*i + 2$

TYPES OF TREES:

- **m-ary tree** - rooted tree, each node has no more than m children
- **Binary tree (BT)** - each node has *at most two children*
- **Complete BT** = bin tree, **every level is full**, except possibly last (far left).
- **BST** - parent's value greater than values in left subtree & less values in right subtree
- **Binary Heap** (max / min) = cbt, parent $>(<)$ children
- **Full (proper or plane) BT** - every node has *0 or 2 children*
- **Perfect BT** - all interior nodes have two children and all leaves have same depth
- **Infinite complete BT** - every node has two children (infinite levels)
- **Balanced BT** - L and R subtrees of each node differ in height by no more than 1
- **Degenerate (pathological) tree** - each parent has only one child = **linked list**
- **Representation:** list, dict, object

BFS (recursive) incl. height

```
# Recursive - time O(n^2), space O(n)
def LevelOrder_rec(root):
    h = height(root)
    for i in range(1, h+1):
        printGivenLevel(root, i)

# print nodes at a given level
def printGivenLevel(root, level):
    if root is None:
        return
    if level == 1:
        print(root.val, end=" ")
    elif level > 1:
        printGivenLevel(root.left, level-1)
        printGivenLevel(root.right, level-1)

# compute height - number of nodes from root
def height(node):
    if node is None:
        return 0
    else:
        # Compute the height of each subtree
        lheight = height(node.left)
        rheight = height(node.right)

        #Use the larger one
        if lheight > rheight:
            return lheight+1
        else:
            return rheight+1
```

```
root = Node(1)
root.left = Node(2)
root.right = Node(3)
root.left.left = Node(4)
root.left.right = Node(5)
LevelOrder_rec(root)
```

Check if BT is balanced

```
# time O(n), space c. O(h)-rec.stack: O(logn) or O(n)
def is_balanced(root):
    return -1 != get_depth(root)
# return 0 if unbalance. else depth+1
def get_depth(root):
    if not root:
        return 0
    left = get_depth(root.left)
    right = get_depth(root.right)
    if abs(left-right) > 1 or \
        left == -1 or right == -1:
        return -1
    return 1 + max(left, right)
```

Trees

Lowest common ancestor (any BT) Check if tree is BST

Time c. $O(n\log n)$ can be $O(n)$, space c. $O(n)$ -store tree_vals

```
def lca(root, p, q):
    if not root or root is p or root is q:
        return root
    left = lca(root.left, p, q)
    right = lca(root.right, p, q)
    if left and right:
        return root
    return left if left else right
```

Time c. $O(n)$, space c. $O(h)$ - recurs. stack:
 $O(\log n)$ balanced tree or $O(n)$ skewed tree

```
# If BST=>inorder traversal = sorted list
def inorder( root ):
    if tree != None:
        inorder(root.left)
        tree_vals.append(root.val)
        inorder(root.right)

def sort_check( tree_vals ):
    return tree_vals == sorted(tree_vals)

tree_vals = []
inorder(tree)
sort_check(tree_vals)
```

Invert a binary tree

Time c. $O(n)$, space c. $O(h)$ -rec.stack: best $O(\log n)$ or worst $O(n)$

```
def reverse(root):
    if not root: return
    root.left,root.right=root.right,root.left
    if root.left:
        reverse(root.left)
    if root.right:
        reverse(root.right)
```

Trim BST

Time c. $O(n)$, space c. $O(h)$ -rec.stack: best $O(\log n)$ or worst $O(n)$

```
def trimBST(tree, minVal, maxVal):
    if not tree: return
    tree.left=trimBST(tree.left, minVal, maxVal)
    tree.right=trimBST(tree.right, minVal, maxVal)
    if minVal<=tree.val<=maxVal:
        return tree
    if tree.val < minVal:
        return tree.right
    if tree.val > maxVal:
        return tree.left
```

BST node w/value closest to k

Time c. $O(h)$ -1 recurs. call per level, space c. $O(h)$ -rec.stack

```
def closest_value(root, k):
    a = root.val
    child = root.left if k < a else root.right
    if not child:
        return a
    b = closest_value(child, k)
    return min(k-a, k-b)|
```

Rotate matrix clockwise, in place

```
# time O(2M) -> O(M), space O(1)
class Solution(object): where M=# elem. in matrix = n^2
    def rotate(self, matrix):
        self.transpose(matrix)
        self.reflect(matrix)

    def transpose(self, matrix):
        n = len(matrix)
        for i in range(n):
            for j in range(i+1, n):
                matrix[i][j], matrix[j][i] =\
                    matrix[j][i], matrix[i][j]

    def reflect(self, matrix):
        n = len(matrix)
        for i in range(n):
            for j in range(n//2):
                matrix[i][j], matrix[i][-j-1] =\
                    matrix[i][-j-1], matrix[i][j]
```

Maximum path sum

#Time $O(n)$, space $O(h)$: balanced $O(\log n)$, skewed $O(n)$

```
def max_path_sum(root):
    max_sum = float('-inf')
    def gain(root):
        nonlocal max_sum
        if not root: return 0
        # only pos values
        left = max(gain(root.left), 0)
        right = max(gain(root.right), 0)
        # update global max_sum
        new_path = root.val+left+right
        max_sum = max(max_sum, new_path)
        # return max gain if continue same path up
        return root.val + max(left, right)
    gain(root)
    return max_sum
```

BST nodes in given range

Time c. $O(h+k)$ k=#nodes in range low-high,
space c. $O(h)$ -rec.stack: best $O(\log n)$ or worst $O(n)$

```
def getCount(root, low, high):
    if root == None: return 0
    if root.data==high and root.data==low:
        return 1
    if root.data <= high and root.data >= low:
        return (1 + getCount(root.left, low, high) +\
                getCount(root.right, low, high))
    elif root.data < low:
        return getCount(root.right, low, high)
    else:
        return getCount(root.left, low, high)
```

938. Range Sum of BST

```
# t.c.O(n), s.c.O(n) best O(logn) if balanced
def range_sum(root, low, high):
    if root is None: return 0
    res = 0
    if root.val < low:
        res += range_sum(root.right,
                          low, high)
    elif root.val > high:
        res += range_sum(root.left,
                          low, high)
    else:
        res += root.val +\
            range_sum(root.left, low, high) +\
            range_sum(root.right, low, high)
    return res
```

199. BT Right Side View (Top2Bottom)

```
# t.c.O(n), s.c.O(h)=rec.stack
# average logn, worst n
def rightSideView(root):
    if root is None:
        return []
    def helper(node, level):
        if level == len(rightside):
            rightside.append(node.val)
        for child in [node.right,
                      node.left]:
            if child:
                helper(child, level+1)
    rightside = []
    helper(root, 0)
    return rightside
```

Insert Number into BST

```
# t.c.=s.c.=O(n) worst, O(logn) avg
def insert(root, data):
    # Base case found where to insert
    if root is None:
        return TreeNode(data)
    if data < root.data:
        root.left = insert(root.left, data)
    else:
        root.right = insert(root.right, data)
    # Return (unchanged) root pointer
    return root
```

257. BT: All Root-to-Leaf Paths

Leaf = node with no children. Space c. O(n)
 T.c. O(n) traversing + str concat which in worst case - skewed tree = each path has length n so total t.c. can be $O(n^2)$

```
# recursive
def binaryTreePaths(root):
    def get_path(root, path):
        if root:
            path += str(root.val)
            if not root.left and not root.right:
                paths.append(path) # reached leaf
            else:
                path += '->'
                get_path(root.left, path)
                get_path(root.right, path)
    paths = []
    get_path(root, '')
    return paths

# iterative
def binaryTreePaths(root):
    if not root: return []
    paths = []
    stack = [(root, str(root.val))]
    while stack:
        node, path = stack.pop()
        if not node.left and not node.right:
            paths.append(path)
        if node.left:
            stack.append((node.left, path + '->' +\
                         str(node.left.val)))
        if node.right:
            stack.append((node.right, path + '->' +\
                         str(node.right.val)))
    return paths
```

114. Flatten Binary Tree to Linked List

```
# recursive, t.c. O(n), s.c. O(n)
def flattenTree(node):
    if not node:
        return None
    if not node.left and not node.right:
        return node # if leaf node, return node
    left = flattenTree(node.left)
    right = flattenTree(node.right)
    if left: # If left subtree, modify connect-s
        left.right = node.right
        node.right = node.left
        node.left = None
    # once re-wired, return "rightmost" node
    return right if right else left

# modifies in-place
def flatten(self, root: TreeNode) -> None:
    flattenTree(root)
```

114. Flatten Binary Tree to Linked List

```
# Iterative in-place, t.c. O(n), s.c. O(1)
def flatten(self, root: TreeNode) -> None:
    if not root:
        return None
    node = root
    while node:
        if node.left:
            rightmost = node.left # find rightmost node
            while rightmost.right:
                rightmost = rightmost.right
            rightmost.right = node.right #rewire connect-s
            node.right = node.left
            node.left = None
        node = node.right # move on to right side of tree
```

426. Convert BST to Sorted Circular Doubly LL

15

(FB Favorite. Toposort, graph traversal r favs too)

In circular doubly linked list - predecessor of first elem is last element; successor of last elem is first elem.
Inorder trav.: L->node->R, linking all nodes into DLL

```
# t.c.=O(n),s.c.=O(h) (O(n) worst if tree skewed,
# but O(logn) on avg
def treeToDoublyList(root):
    def helper(node):
        nonlocal last, first
        if node:
            helper(node.left) #left
            if last: # node
                # link prev node (last) w/node
                last.right, node.left = node, last
            else:
                # memorize 1st smallest node
                first = node
            last = node
            helper(node.right) #right
    if not root:
        return None
    # smallest (1st) & largest (last) nodes
    first, last = None, None
    helper(root)
    last.right, first.left = first, last # close DLL
    return first
```

Build Bin Tree from Array

Use level-order traversal to convert back to array

```
# time = space = O(n)
class Node:
    def __init__(self, val):
        self.val = val
        self.left = self.right = None

def insertLevelOrder(arr, i, n):
    root = None
    if i < n:                      # base case
        root = Node(arr[i])
        root.left = insertLevelOrder(arr, 2*i+1, n)
        root.right = insertLevelOrder(arr, 2*i+2, n)
    return root

arr = [1,2,3,4,5,6,None,None,None,7,8]
n = len(arr)
root = insertLevelOrder(arr, 0, n)
```

Delete Nodes, Return Forest

```
def delNodes( root: Node,
              to_delete: List[int],
            ) -> List[Node]:
    res, to_delete = [], set(to_delete)
    def helper(root):
        if root:
            # next line exec. after recursion on way up
            root.left, root.right = (helper(root.left),
                                      helper(root.right))

            if root.val not in to_delete:
                return root
            res.append(root.left) # if root is deleted
            res.append(root.right) # if root is deleted
    res.append(helper(root))
    return([ i for i in res if i ])
```

Nodes in Compl. Bin. Tree < O(n)

16

```
# time=O(d^2)=O((logN)^2), d=depth, space = O(1)
def countNodes(root: TreeNode) -> int:
    if not root: return 0
    d = compute_depth(root)
    if d == 0: return 1
    # Last level nodes - 0 to 2**d-1 (left->right)
    # Bin search to check how many nodes exist
    left, right = 1, 2**d - 1
    while left <= right:
        pivot = left + (right - left) // 2
        if exists(pivot, d, root):
            left = pivot + 1
        else:
            right = pivot - 1
    # there are only left nodes on last level
    return (2**d - 1) + left

def compute_depth(node: Node) -> int:
    d = 0
    while node.left:
        node = node.left
        d += 1
    return d

def exists(idx: int, d: int, node: Node) -> bool:
    '''Return True if last level node idx exists'''
    left, right = 0, 2**d - 1
    for _ in range(d):
        pivot = left + (right - left) // 2
        if idx <= pivot:
            node = node.left
            right = pivot
        else:
            node = node.right
            left = pivot + 1
    return node is not None
```

Singly Linked List (SLL)

```
class Node:
    def __init__(self, value):
        self.value = value
        self.next = None
```

Doubly Linked List (DLL)

```
class DLL_Node(object):
    def __init__(self, value):
        self.value = value
        self.next = None
        self.prev = None
```

Compare 2 strings as linked lists

```
# 0=same, 1=1st str lexicographically >,
# -1 - 2nd
class Node:
    def __init__(self, key):
        self.c = key ;
        self.next = None

def compare(list1, list2):
    while (list1 and list2 and\
           list1.c == list2.c):
        list1 = list1.next
        list2 = list2.next
    if (list1 and list2):
        return 1 if list1.c > list2.c else -1
    if (list1 and not list2):
        return 1
    if (list2 and not list1):
        return -1
    return 0

list1 = Node('g')
list1.next = Node('o')
...
print(compare(list1, list2))
```

nth to the last element

```
def nth_to_last_node(head, n):
    left_pointer = head
    right_pointer = head
    for i in range(n-1):
        if not right_pointer.next:
            print('n > length list')
            return
        right_pointer = right_pointer.next
    while right_pointer.next:
        left_pointer = left_pointer.next
        right_pointer = right_pointer.next
    return left_pointer.value
```

Reverse SLL

```
# time O(n), space O(1)
def reverse_sll(head):
    current = head
    previous = None
    next_node = None
    while current:
        next_node = current.next
        current.next = previous
        previous = current
        current = next_node
    return previous
```

Break cycle

```
def removeLoop(head, loop_node):
    pointer1 = head
    while(1):
        pointer2 = loop_node
        while (pointer2.next != loop_node and\
               pointer2.next != pointer1):
            pointer2 = pointer2.next
        if pointer2.next == pointer1:
            break
        pointer1 = pointer1.next
        pointer2.next = None
```

Check if cycle

```
def cycle_check(node):
    marker1 = node
    marker2 = node
    while marker2 and marker2.next:
        marker1 = marker1.next
        marker2 = marker2.next.next
        if marker2 == marker1:
            return True
    return False
```

Length of linked list

```
def length_LL(LL, head):
    if (not head):
        return 0
    else:
        return 1 +\nlength_LL(LL, head.next)
```

Linked Lists

Reorder List (3 challenges in one)

$L_0 \rightarrow L_n \rightarrow L_1 \rightarrow L_{n-1} \rightarrow L_2 \rightarrow L_{n-2} \rightarrow \dots$

```
# in-place, time O(n), space O(1)
def reorderList(head: Node) -> None:
    if not head: return
    # FIND MIDDLE
    slow = fast = head
    while fast and fast.next:
        slow = slow.next
        fast = fast.next.next
    # REVERSE SECOND LIST IN-PLACE
    prev, curr = None, slow
    while curr:
        curr.next, prev, curr = prev, curr, curr.next
    # MERGE 1->6->2->5->3->4
    first, second = head, prev
    while second.next:
        first.next, first = second, first.next
        second.next, second = first, second.next
```

21. Merge Two Sorted Lists

```
# t.c. O(n+m), s.c. O(1)
from typing import Optional
def mergeTwoLists(l1: Optional[ListNode],
                  l2: Optional[ListNode],
                  ) -> Optional[ListNode]:
    prehead = ListNode(0)
    current = prehead
    while l1 and l2:
        if l1.val <= l2.val:
            current.next = l1 #entire node
            l1 = l1.next
        else:
            current.next = l2
            l2 = l2.next
        current = current.next
    # l1 or l2 can still have nodes
    current.next = l1 if l1 else l2
    return prehead.next
```

Merge k Sorted Linked Lists

```
# time c. O(Nlogk), k = # lists. Space c. O(1)
def mergeKLists(lists):
    n = len(lists)
    interval = 1
    while interval < n:
        for i in range(0, n-interval, interval*2):
            lists[i] = merge2Lists( lists[i], lists[i+interval] )
        interval *= 2
    return lists[0] if n > 0 else None

def merge2Lists(l1, l2):
    head = point = Node(0)
    while l1 and l2:
        if l1.val <= l2.val:
            point.next = l1
            l1 = l1.next
        else:
            point.next = l2
            l2 = l2.next
        point = point.next
    if not l1: point.next=l2
    else: point.next=l1
    return head.next
```

Merge k Sorted Arrays

```
# time c. O(kN*Logk) since
# using heap (N*Logk) k times;
# space c. O(N) - output array
from heapq import merge
def mergeK(arr, k):
    l = arr[0]
    for i in range(k-1):
        l = list(merge(l, arr[i + 1]))
    return l
```

2. Add Two Numbers (as linked list)

20

```
# t.c. = s.c. = O(max(n, m))
def addTwoNumbers(l1: Node, l2: Node) -> Node:
    prehead = Node(0)
    current = prehead
    carry = 0
    while l1 or l2 or carry:
        val1 = l1.val if l1 else 0
        val2 = l2.val if l2 else 0
        total = val1 + val2 + carry
        carry = total // 10
        new_digit = total % 10
        current.next = Node(new_digit)
        current = current.next
        if l1: l1 = l1.next
        if l2: l2 = l2.next
    return prehead.next

# A dummy head to simplify handling the result.
# Pointer to construct the new list.
# carry from each addition
# Traverse both lists until exhausted.
# If either list ran out of digits, use 0

# New carry
```

543. Diameter of Binary Tree

Longest path (# edges) betw. any two nodes (thru or not thru root)

```
# t.c.=O(n), s.c.=O(n)
def diameterOfBinaryTree(root: TreeNode) -> int:
    def longest_path(node):
        if not node: return 0
        nonlocal diameter
        left_path = longest_path(node.left)
        right_path = longest_path(node.right)
        diameter = max(diameter, left_path+right_path)
        # add 1 for connection to parent
        return max(left_path, right_path) + 1
    diameter = 0
    longest_path(root)
    return diameter
```

138. Copy List with Random Pointer

```
# t.c.=O(n), s.c.=O(n)
from collections import defaultdict
class Node:
    def __init__(self,x,next=None,random=None):
        self.val = int(x)
        self.next = next
        self.random = random
    def copyRandomList(head: Node) -> Node:
        # Nodes as keys, their copies as values
        node_map = defaultdict(lambda: None)
        curr = head #Fill node_map
        while curr:
            node_map[curr] = Node(curr.val)
            curr = curr.next
        # Update pointers of copied nodes
        curr = head
        while curr:
            dup = node_map[curr]
            dup.next = node_map[curr.next]
            dup.random = node_map[curr.random]
            curr = curr.next
        return node_map[head]
```

Backtracking - Words from Phone Number

21

```
# time = O(N * (4^N)), N = len(digits), space = O(n)
def letterCombinations(digits: str) -> List[str]:
    if len(digits) == 0: return []
    letters = {"2": "abc", "3": "def", "4": "ghi",
               "5": "jkl", "6": "mno", "7": "pqrs",
               "8": "tuv", "9": "wxyz"}
    def backtrack(index, path):
        if len(path) == len(digits): # base case
            combinations.append("".join(path))
            return
        # letters mapped to current digit
        possible_letters = letters[digits[index]]
        for letter in possible_letters:
            path.append(letter)
            backtrack(index + 1, path)
            path.pop() # remove before moving on
    combinations = []
    backtrack(0, [])
    return combinations
```

Num ways to cover distance

All ways to cover distance it w/1, 2, 3 steps

```
# time O(n), space O(n)
def count_dp(dist):
    count = [0] * (dist + 1)
    count[0] = 1
    count[1] = 1
    count[2] = 2
    for i in range(3, dist + 1):
        count[i] = (count[i-1] +
                    count[i-2] +
                    count[i-3])
    return count[dist]

dist = 20
print(count_dp(dist))
```

Recursion

t.&s.c. O(n)

```
def fact(n):
    if n < 2:
        return 1
    else:
        return n*fact(n-1)

fact(5)
```

Memoization

M. doesn't help t.c., adds to s.c.

```
def factorial(n):
    if n < 2:
        return 1
    if not n in memo:
        memo[n] = n*factorial(n-1)
    return memo[n]

memo = {}
factorial(1)
```

Fibonnaci dynamic

Unlike for factorial M. helps here

```
# memoization (dynamic) s.c. space + memo = 2O(n)
# time O(n) - counting sequentially
def fib_dyn(n):
    if n == 0 or n == 1:
        return n
    if not n in memo:
        memo[n] = fib_dyn(n-1) +
                  fib_dyn(n-2)
    return memo[n]
```

Fibonnaci iterative

```
# iterative - tuple unpacking
# time O(n) - counting sequentially
def fib_iter(n):
    a = 0
    b = 1
    for i in range(n):
        a, b = b, a + b
    return a

for i in range(20):
    print(fib_iter(i), end=', ')
```

Fibonnaci recursive

```
# time O(2^n) because 2 calls for each n. S.c. O(n) -
def fib_rec(n): depth of recursion stack
    if n == 0 or n == 1:
        return n
    else:
        return fib_rec(n-1) + fib_rec(n-2)

# print 1st 20 Fib. nums
for i in range(20):
    print(fib_rec(i), end=', ')
```

Iterative Factorial

```
# t.c. O(n), s.c. O(1)
def factorial(n):
    if n < 0:
        raise ValueError("m")
    result = 1
    for i in range(1, n + 1):
        result *= i
    return result
```

Dijkstra's Algorithm

Shortest path from one vertex to all others

```
# time c. O(V + E*logE), iterative, BFS
import heapq    # min heap of (distance, vertex) pairs
# control iteration order - to pick vertex w/smallest dist
def calculate_distances(graph, start):
    # cost from start to each destination
    distances = {vertex: float('infinity') for vertex in graph}
    distances[start] = 0
    pq = [(0, start)]           # min heap or priority queue
    while len(pq) > 0:
        # pop smallest, maintain heap
        current_distance, current_vertex = heapq.heappop(pq)
        # nodes can be added to pq multiple times; process node
        # first time it's removed from pq
        if current_distance > distances[current_vertex]:
            continue
        for neighbor, weight in graph[current_vertex].items():
            distance = current_distance + weight
            # only if new path is better
            if distance < distances[neighbor]:
                distances[neighbor] = distance
                # push onto heap, maintain heap
                heapq.heappush(pq, (distance, neighbor))
    return distances
```

Min. spanning tree - Prim's algo

22

S.t. for a graph $G = (V, E)$ is acyclic subset of E connecting all vertices in V (sum of edge weights minimized)

Most efficient info flow. There may be several spanning trees - we need to find the minimum one.

Using priority queue to select next vertex for growing graph

```
from collections import defaultdict
import heapq
def min_spanning_tree(graph, start):
    ''' Outputs mst - min. spanning tree '''
    mst    = defaultdict(set)
    visited = set([start])
    edges   = [ (cost, start, to)
               for to, cost in graph[start].items() ]
    heapq.heapify(edges)
    while edges:
        cost, frm, to = heapq.heappop(edges)
        if to not in visited:
            visited.add(to)
            mst[frm].add(to)
            for to_next, cost in graph[to].items():
                if to_next not in visited:
                    heapq.heappush(edges, (cost, to, to_next))
    return mst
```

Floyd Warshall Algorithm

Shortest path between all pairs of nodes

```
# time c. O(V^3)
def floyd_marshall(graph):
    dist = list(map(lambda i : list(map(lambda j : j , i)) , graph))
    for k in range(V):
        for i in range(V):          # all vertices as source
            for j in range(V):      # all vertices as destination
                # update dist[i][j] if vertex k on shortest path from i to j
                dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j])
    return dist
```

```
# recursion + memoization
# arr coins + target amnt => fewest coins as change
def coins_dp(target, coins, known_results):
    min_coins = target
    if target in coins:
        known_results[target] = 1
        return 1
    elif known_results[target] > 0:
        return known_results[target]
    else:
        for i in [c for c in coins if c <= target]:
            num_coins = 1 + coins_dp( target-i,
                                      coins,
                                      known_results,
                                      )
            if num_coins < min_coins:
                min_coins = num_coins
                known_results[target] = min_coins
    return min_coins

target = 74
coins = [1,5,10,25]
known_results = [0]*(target+1)
coins_dp(target, coins, known_results)
```

Knapsack DP

```
# max sum of values for subset of weights <= W
# O(N*W) time and space
def knapSack(W, weights, values, n):
    DP = [[0 for x in range(W + 1)] for x in range(n + 1)]
    for i in range(n + 1):
        for w in range(W + 1):
            if i == 0 or w == 0:
                DP[i][w] = 0
            elif weights[i-1] <= w:
                DP[i][w] = max(values[i-1] + DP[i-1]\[w - weights[i-1]], DP[i-1][w])
            else:
                DP[i][w] = DP[i-1][w]
    return DP[n][W]

values = [60, 100, 120]
weights = [10, 20, 30]
W = 50
print(knapSack(W, weights, values, len(values)))
```

```
# time O(log(min(N,M)), space O(1). BIN SEARCH
def median(nums1, nums2):
    # capture edge cases
    if len(nums2) < len(nums1):
        nums1, nums2 = nums2, nums1
    total = len(nums1) + len(nums2)
    half = total // 2
    l, r = 0, len(nums1)-1
    # median is guaranteed
    while True:
        i = (l + r) // 2      # for nums1
        # subtr. 2 - j starts at 0, i starts at 0
        j = half - i - 2     # for nums2
        # overflow of indices
        nums1_left = nums1[i] if i >= 0 else float("-inf")
        nums1_right = nums1[i+1] if (i+1) < len(nums1) else float("inf")
        nums2_left = nums2[j] if j >= 0 else float("-inf")
        nums2_right = nums2[j+1] if (j+1) < len(nums2) else float("inf")
        # if correct partition is found
        if nums1_left <= nums2_right and nums2_left <= nums1_right:
            if total % 2:
                return min(nums1_right, nums2_right)
            else:
                return (max(nums1_left, nums2_left) +\
                        min(nums1_right, nums2_right)) / 2
        # if no correct partition - arrays are in ascending order
        elif nums1_left > nums2_right:
            r = i - 1
        else:
            l = i + 1
```

Traveling Salesman Problem (TSP)

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NP hard problem. There is no known polynomial time solution

```
# Naive approach. Time c. n!
from sys import maxsize
V = 4
def travellingSalesmanProblem(graph, s):
    vertices = [] # all verteces, but source vtx
    for i in range(V):
        if i != s:
            vertices.append(i)
    min_pathweight = maxsize # min weight Hamiltonian Cycle
    while True:
        current_pathweight = 0 # current Path weight(cost)
        k = s # compute current path weight
        for i in range(len(vertices)):
            current_pathweight += graph[k][vertices[i]]
            k = vertices[i]
        current_pathweight += graph[k][s]
        min_pathweight = min(min_pathweight,
                             current_pathweight) #update min
        if not next_permutation(vertices):
            break
    return min_pathweight
```

```
def next_permutation(L):
    n = len(L)
    i = n - 2
    while i >= 0 and L[i] >= L[i + 1]:
        i -= 1
    if i == -1:
        return False
    j = i + 1
    while j < n and L[j] > L[i]:
        j += 1
    j -= 1
    L[i], L[j] = L[j], L[i]
    left = i + 1
    right = n - 1
    while left < right:
        L[left], L[right] = L[right], L[left]
        left += 1
        right -= 1
    return True

# matrix representation of graph
graph = [[0, 10, 15, 20], [10, 0, 35, 25],
          [15, 35, 0, 30], [20, 25, 30, 0]]
start = 0
print(travellingSalesmanProblem(graph, s))
```

Kth Largest Element in Array

```
# time = O(n), space = O(1)
def findKthLargest(nums, k):
    def partition(left, right, pivot_idx):
        pivot = nums[pivot_idx]
        # 1. move pivot to end
        nums[pivot_idx], nums[right] = nums[right], nums[pivot_idx]
        # 2. move all smaller elements to the left
        store_idx = left
        for i in range(left, right):
            if nums[i] < pivot:
                nums[store_idx], nums[i] = nums[i], nums[store_idx]
                store_idx += 1
        # 3. move pivot to its final place
        nums[right], nums[store_idx] = nums[store_idx], nums[right]
    return store_idx

def select_rec(left, right, k_smallest):
    if left == right:      # base case - 1 elem
        return nums[left]
    pivot_idx = random.randint(left, right)
    # find pivot pos in sorted list
    pivot_idx = partition(left, right, pivot_idx)
    # if pivot in final sorted position
    if k_smallest == pivot_idx:
        return nums[k_smallest]
    elif k_smallest < pivot_idx:    # go left
        return select_rec(left, pivot_idx-1, k_smallest)
    else:                      # go right
        return select_rec(pivot_idx+1, right, k_smallest)
# kth largest = (n - k)th smallest
return select_rec(0, len(nums)-1, len(nums)-k)
```

$n \log(n)$ solution –
sort, then get
kth element

Minimum Window Substring

```
# time/space c. O(n+m)
# min. contiguous substr. of s w/all chars from t
from collections import Counter
def minWindow(s, t):
    if not t or not s:
        return ''
    # unique chars in t and curr window
    dict_t = Counter(t)
    curr = {}
    len_t = len(dict_t)
    len_curr = 0
    l, r = 0, 0
    # length, l, r
    res = float('inf'), None, None
    while r < len(s):
        char = s[r]
        curr[char] = curr.get(char, 0) + 1
        if char in dict_t and\
            curr[char] == dict_t[char]:
            len_curr += 1
        while l <= r and len_curr == len_t:
            char = s[l]
            if r - l + 1 < res[0]:
                res = (r - l + 1, l, r)
            curr[char] -= 1
            if char in dict_t and\
                curr[char] < dict_t[char]:
                len_curr -= 1
            l += 1
        r += 1
    return '' if res[0]==float('inf') else\
        s[res[1]:res[2]+1]
```

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```
tips.assign(tip_rate = tips["tip"] / tips["total_bill"])
tips.loc[tips["tip"] < 2, "tip"] *= 2
```

GROUP BY

```
tips.groupby("sex").size()    # female = 87, male = 90  
tips.groupby("day").agg({"tip": "mean", "day": "size"})
```

tip	day	
day		# agg() on one column in each df from grouped df.groupby('Team')['Points'].agg(np.mean)
Fri	2.586842	19
Sat	2.986897	87

```
tips.groupby(["smoker", "day"]).agg({"tip": ["size", "mean"]})
```

		tip	df.groupby('Team').agg(np.size)				
		size	mean	Rank Year Points			
smoker	day	Team			# M df.		
No	Fri	4	2.562500	Angels	2	2	2
	Sat	45	3.052889	Kings	3	3	3

MULTIPLE AGGREGATION FUNCTIONS ON ONE COL

```
df.groupby('Team')['Points'].agg([np.sum, np.mean, np.std])
```

	sum	mean	std
Team			
Angels	1536	768.000000	134.350288
Kings	2285	761.666667	24.006943

Remove missing data

```
# Remove rows with ANY missing values
df_dropped_rows = df.dropna()
# Remove columns with ANY missing values
df_dropped_columns = df.dropna(axis=1)
# Remove rows with missing values in specific columns
df_dropped_specific1 = df.dropna(subset=['A'])
# Remove columns with missing values in specific rows
df_dropped_specific2 = df.dropna(subset=[1], axis=1)
```

Transformation: df.groupby().transform()

Applied to **group** or **column**, returns an obj w/same index size

```
score = lambda x: (x - x.mean()) / x.std()*10  
df.groupby('Team').transform(score)
```

	Rank	Year	Points
0	-15.000000	-11.618950	12.843272
1	5.000000	-3.872983	3.020286

Filtration: df.groupby().filter()

Returns subset of df, filtering data on a criteri

```
df.groupby('Team').filter(lambda x: len(x) >= 3)
```

	Team	Rank	Year	Points
0	Riders	1	2014	876
1	Riders	2	2015	789

```
# MULTIPLE AGGREGATION FUNCTIONS ON MANY COLS  
df.groupby('Team')[['Points', 'Rank']].agg([min, max])
```

Team	Points			Rank		
	sum	mean	std	sum	mean	std
Angels	1536	768.000000	134.350288	5	2.500000	0.707107
Kings	2285	761.666667	24.006943	5	1.666667	1.154701

```
 DataFrame.merge( right_df,
                  how='inner', # {'left', 'right', 'outer', 'inner', 'cross'}
                  on=None,       # join key(s) - column
                  left_on=None, # if join key(s) w/different names (but same)
                  right_on=None,
                  left_index=False, # use index as join key(s)
                  right_index=False,
                  sort=False,      # Sort join keys lexicographically
                  suffixes=('_x', '_y'),
                  copy=True,
                  indicator=False, # col _merge saying source of each row
                  validate=None.   # check merge keys if "1:1", "1:m", etc.
```

Fill missing data

```
DataFrame.fillna(
    value=None, # number/dict/Series/df
    method=None, # 'bfill', 'ffill', None
    limit=None, # # of fffills/bfills
    axis=None,
    inplace=False,)

# Fill with a constant value
df_filled_constant = df.fillna(250)
# propagate next valid observation forward
df_filled_ffill = df.fillna(method='ffill')
# propagating next valid observation backward
df_filled_bfill = df.fillna(method='bfill')
# Fill all NaNs in cols w/values
values = {"A": 100, "B": 200, "C": 300,}
print(df.fillna(value=values))
# Fill w/df - replacement at same cols & idx
print(df.fillna(df2))
# Fill w/mean of col
df_filled_mean = df.fillna(df.mean())
# Fill w/median of cols
df_filled_median = df.fillna(df.median())
# Fill with mode of cols (mode() returns df)
df_filled_mode = df.fillna(df.mode().iloc[0])
# Interpolation (default: linear)
# method: 'index','pad','nearest','zero', etc.
df_interpolated = df.interpolate()
print(df_interpolated)
```

Identifying outliers using the IQR (Interquartile Range)

```
# Calculate Q1 (25th perc.) & Q3 (75th perc.)
Q1 = df['A'].quantile(0.25)
Q3 = df['A'].quantile(0.75)
IQR = Q3 - Q1
# Define outlier boundaries
lower_bound = Q1 - 1.5 * IQR
upper_bound = Q3 + 1.5 * IQR
# Identify outliers
outliers = df[(df['A'] < lower_bound) | \
    (df['A'] > upper_bound)]
```

Handling outliers

```
# REMOVE
df_no_outliers = df[(df['A'] >= lower_bound) & (df['A'] <= upper_bound)]
# CAP
df['A'] = np.where(df['A'] > upper_bound, upper_bound, df['A'])
df['A'] = np.where(df['A'] < lower_bound, lower_bound, df['A'])
# IMPUTE
median = df['A'].median()
df['A'] = np.where(df['A'] > upper_bound, median, df['A'])
df['A'] = np.where(df['A'] < lower_bound, median, df['A'])
# SCALE - RobustScaler if robust to outliers
from sklearn.preprocessing import RobustScaler
scaler = RobustScaler()
df['A'] = scaler.fit_transform(df[['A']])
```

Transforming Data

Z-Score: # standard deviations a data point is from the mean.

$\text{z-score} = (\text{value} - \text{mean}) / \text{std_dev}$ => StandardScaler()

```
from scipy.stats import zscore
df['A_log'] = np.log(df['A'])
df['A_zscore'] = zscore(df['A'])
print(df)
```

	A	A_capped	A_log	A_zscore
0	10	10.000	2.302585	-0.425312
1	-50	-20.625	NaN	-2.115292

Isolation Forest - ML to identify and handle outliers

```
from sklearn.ensemble import IsolationForest
iso = IsolationForest(contamination=0.1)
df['outliers'] = iso.fit_predict(df[['A']])
outliers = df[df['outliers'] == -1]
```

	A	A_capped	A_log	A_zscore	A_median	A_scaled	outliers
0	10	10.000	2.302585	-0.425312	10.0	-0.473118	1
1	-50	-20.625	NaN	-2.115292	21.0	-3.053763	-1

Read CSV files

```
pandas.read_csv()
```

- **filepath, sep=',' / regex, header='infer' / None / row number** (numbers for multi-index),
- **names=col names**. If they already exist in file, pass header=0 to rename,
- **index_col=col(s)** for row labels (None), **usecols=subset of cols to use (None)**
- **dtype=dtypes** for cols (None), **skiprows=line numbers to skip (0-indexed) or # lines to skip (int)**
- **skipfooter=# lines at bottom of file to skip (0)**, **nrows=# rows to read (None)**,
- **na_values=additional strings to recognize as NaN**, (also **true_values** and **false_values**),
- **skip_blank_lines=True, parse_dates=None, infer_datetime_format=_NoDefault.no_default**,
- **keep_date_col= keep original date col, date_format=None**,
- **thousands=None, decimal='.'**, **quotechar='''**, **quoting=0** (csv.QUOTE_MINIMAL, etc.),
- **escapechar=None, encoding=None**,
- **encoding_errors='strict' / 'ignore' / 'replace'**, **on_bad_lines='error' / 'warn' / 'skip'**,
- **delim_whitespace=whether or not whitespace will be used as the sep (equ. to sep='\\s+')**

```
df.to_json(file_name)
df = pd.read_json(file_name)
```

Rank

Rank col values in descend or ascen order

Method - How rank group of records that have the same value (ties):

- average: average rank of group
- min: lowest rank in group
- max: highest rank in group
- first: in order ranks appear
- dense: = 'min', but rank increases by 1

Axending=False => max value has rank 1

```
df['rank'] = df['num_legs_1178'].rank(
    method='first',
    ascending=False).astype("int")
df[['num_legs_1178', 'rank']]
```

	num_legs_1178	rank
falcon	2	3
dog	4	2
spider	8	1
fish	0	4

Set value for a cell

```
df.at['C', 'a'] = 10 # BY VALUE
df.iat[5,10] = 2 # BY INDEX
```

Select columns

```
df.dtypes.value_counts()
# col by type
df.select_dtypes(include=['int'])
# col by name - has "num" in name
df.filter(like='num')
# has digit(s) in col name
df.filter(regex='\d')
```

num_legs_1178

falcon 2

Select w/"get_loc" and "index"

```
start = df.columns.get_loc('num_legs_1178')
end = df.columns.get_loc('num_specimen')
df.iloc[:,4, start:end+1]
```

```
start = df.index.get_loc('dog')
end = df.index.get_loc('fish')
df.iloc[start:end+1, 2:5]
```

same things

```
df[df['num_legs_1178'] > 2]
df.query('num_legs_1178 > 2')
```

Melt - unpivot df from wide to long format

	A	B	C	A variable	value		
0	a	1	2	0	a	B	1
1	b	3	4	1	b	B	3
2	c	5	6	2	c	B	5
				3	a	C	2
				4	b	C	4
				5	c	C	6

1a. Aggregations on entire dataframe

WINDOW FUNCTIONS - find trends in data graphically by smoothing the curve

- **df.rolling()** - rolling window calculations; **window=window size**, **min_periods=min num observations in window required to have a value**.
- **df.expanding()** - same as rolling, but uses all the data up to that point in time. Equivalent statements: [df.rolling(window=len(df), min_periods=1).mean()] = [df.expanding(min_periods=1).mean()]
- **df.ewm()** - exponentially weighted window similar to expanding window, but each prior point is exponentially weighted down relative to the current point

```
df['column'].rolling(2).sum()
df['column'].rolling(2, min_periods=1).sum()
df['column'].rolling(2).mean()
```

Other

```
# cross-tabulation of 2+ factors
pd.crosstab(
    df["genus_overall"],
    df["rating_overall"],
    margins=True,
    normalize=0,)
```

```
# SQLite
import sqlite3
query = 'SELECT * FROM dune_table'
conn = sqlite3.connect('dune.db')
dune_df = pd.read_sql(query, conn)
dune_df.to_sql('dune_table', conn,
               index=False)
```

```
# reshape w/new cols & values + aggr.
pivot_df = pd.pivot_table(
    df, index='Name',
    columns=['Pets', 'Stores'],
    values='Weight' # col for values
    aggfunc='mean')
```

	rating_overall	1.2	3.4	5.2	7.8
genus_overall	avian	1.00	0.00	0.00	0.00
	canine	0.00	1.00	0.00	0.00

Code testing

Types of tests

- **Unit test** - finding bugs in *individual functions*.
Should be considered together with integration test - when a system is comprehensively unit tested, it makes integration testing far easier
- **Integration test** - test *multiple components* of application at once, incl. interactions between the parts - identifies defects in the interfaces between disparate parts of the codebase.
- **Functional test** - *whole system test*. Sometimes it's aka integration test, sometimes - user test

Basics of unittest library

What you need for a unit test:

- **test fixture** - preparations for test: creating temporary or proxy databases, directories, or starting a server.
- **test case** - individual unit of testing, checks response to a particular set of inputs (base class, TestCase, to create new test cases).
- **test suite** - collection of test cases, test suites, or both: to aggregate tests together.
- **test runner** - orchestrates the execution of tests and provides the outcome.

Additional [¶](#)

Main assertions: `assertEqual()`, `assertTrue()`, `assertFalse()`, `assertRaises()`

`unittest.main()` - command-line interface to run the test script

Test discovery via CLI finds all unittest files in the directory structure

Unittest can skip *individual tests* and whole classes of tests, marking a test as an "expected failure," (shouldn't be counted as a failure on a TestResult)

Unittest can distinguish *small differences* in tests (e.g. in some parameters) using the `subTest()` context manager. Doctest can test docstrings

```
import unittest
# Class containing unit tests
class TestStringMethods(unittest.TestCase):
    def setUp(self):
        # Setup required resources
        self.string = "hello world"
        self.empty_string = ""
        self.whitespace_string = " "
        self.numeric_string = "12345"
    def tearDown(self):
        # Clean up resources
        pass
    def test_upper(self):
        self.assertEqual(self.string.upper(),
                        "HELLO WORLD")
    def test_isupper(self):
        self.assertFalse(self.string.isupper())
        self.assertTrue("HELLO".isupper())
    def test_islower(self):
        self.assertTrue(self.string.islower())
        self.assertFalse("Hello".islower())
    def test_split(self):
        self.assertEqual(self.string.split(),
                        ["hello", "world"])
        with self.assertRaises(TypeError):
            s.split(2)
    # Entry point for running the tests
if __name__ == '__main__':
    unittest.main()
```

```
import pytest
```

```
# Fixture for setup and teardown
@pytest.fixture
def string_data():
    data = {
        "string": "hello world",
        "empty_string": "",
        "whitespace_string": "    ",
        "numeric_string": "12345"
    }
    return data

# Test functions
def test_upper(string_data):
    assert string_data["string"].upper() == "HELLO WORLD"

def test_isupper(string_data):
    assert not string_data["string"].isupper()
    assert "HELLO".isupper()

def test_islower(string_data):
    assert string_data["string"].islower()
    assert not "Hello".islower()

def test_split(string_data):
    assert string_data["string"].split() == ["hello", "world"]
    assert string_data["string"].split('o') == ["hell", " w", "rld"]

def test_isspace(string_data):
    assert string_data["whitespace_string"].isspace()
    assert not string_data["string"].isspace()
    assert not string_data["empty_string"].isspace()

def test_isdigit(string_data):
    assert string_data["numeric_string"].isdigit()
    assert not string_data["string"].isdigit()

def test_empty_string(string_data):
    assert string_data["empty_string"].upper() == ""
    assert string_data["empty_string"].split() == []
    assert string_data["empty_string"].islower()

# Run the tests
if __name__ == '__main__':
    pytest.main()
```

Type Hints

```
from typing import Optional

# var of specified type or None
def get_user_email(user_id: int) -> Optional[str]:
    emails = {1: "alice@example.com", 4: "bob@example.com"}
    return emails.get(user_id, None)

email = get_user_email(1)
print(email)                      # Output: alice@example.com

email = get_user_email(2)
print(email)                      # Output: None

from typing import Any

# any type. E.g. Dict[str, Any]
def print_value(value: Any) -> None:
    print(f"Value: {value}")

print_value(42)                  # Output: Value: 42
print_value("Hello")             # Output: Value: Hello
print_value([1, 2, 3])           # Output: Value: [1, 2, 3]

from typing import Union

# var is one of several types
def process_data(data: Union[int, str]) -> str:
    if isinstance(data, int):
        return f"Processed integer: {data}"
    elif isinstance(data, str):
        return f"Processed string: {data}"

from typing import Callable

# var is func or method
def execute_function(
    func: Callable[[int, int], int],
    a: int,
    b: int,
) -> int:
    return func(a, b)

def add(x: int, y: int) -> int:
    return x + y

result = execute_function(add, 5, 3)
print(result) # Output: 8
```

```

import openai
import tiktoken
import backoff

@backoff.on_exception(backoff.expo, openai.error.RateLimitError, max_time=10)
def get_response(model, messages, temperature=0.5, max_tokens=None):
    """Send request, return response"""
    response = openai.ChatCompletion.create(
        model = model,
        messages = messages,
        temperature = temperature, # range(0,2), the more the less deterministic / focused
        top_p = 1, # top probability mass, e.g. 0.1 = only tokens from top 10% proba mass
        n = 3, # number of chat completions
        #max_tokens = max_tokens, # tokens to return
        stream = False, # sequence to stop generation (new line, end of text, etc.)
        stop=None,
    )
    content = [response['choices'][0]['message']['content'], response['choices'][1]['message']['content'], response['choices'][2]['message']['content']]
    #num_tokens_api = response['usage']['prompt_tokens']
    return content

def get_gpt_title(text_, prompt_):
    """Classify text_ using prompt_ and ChatGPT API"""

    # compose messages and reduce length of chapter if needed
    messages = [
        {"role": "system", "content": "You are a helpful title generator.", },
        {"role": "user", "content": prompt_, },
    ]

    while not verify_num_tokens(model, messages):
        text_len = len(text_)
        text_ = text_[text_len//2]
        messages = [
            {"role": "system", "content": "You are a helpful title generator.", },
            {"role": "user", "content": prompt_, },
        ]

    return get_response(model, messages, )

prompt_one = """
1. Please generate a concise title for the text delimited below with triple backticks
2. The maximum length of title should be 5 words
"""

Text:
```
{}```

s = 'This is a text sample'
print(prompt_one.format(s), '\n')

#openai.api_key = os.getenv("OPENAI_API_KEY")
model = 'gpt-3.5-turbo'

text_col = 'text_merged'
#res = dict()
count = 0
for text_ in df[text_col].tolist():
 if t in res:
 continue
 try:
 res[text_] = get_gpt_title(text_, prompt_one)
 except openai.error.RateLimitError:
 print(f'\nText: {text_}\nRate limit error\n')
 except Exception as e:
 print(f'\nText: {text_}\nError: {e}\n')
 count += 1
 if count % 10 == 0:
 print(f'Processing chapter #{count}!|'

```

```
from transformers import T5Tokenizer, T5ForConditionalGeneration
import torch

The GPU 15GB can hold the whole model but not the CPU 13GB
torch.set_default_tensor_type(torch.cuda.FloatTensor)

tokenizer = T5Tokenizer.from_pretrained("google/flan-t5-xl")
model_t5 = T5ForConditionalGeneration.from_pretrained("google/flan-t5-xl", device_map="auto")

def generate(transcript, instruction, n=3):
 '''Original prompt that came with this notebook'''
 transcript = '\n'.join(transcript)[:1500]
 input_text = f'## Transcript\n\n{transcript}\n\n## Directions\n\n{instruction}'
 input_ids = tokenizer(input_text, return_tensors="pt").input_ids
 outputs = model_t5.generate(input_ids, max_new_tokens=100, num_return_sequences=n, num_beams=n)
 return [tokenizer.decode(o).replace('</s>', '').replace('<pad>', '').strip() for o in outputs]

segment_transcript = ["Alice: Hello, this meeting is about notes.",
 "Bill: Yes, the notes taken during a meeting",
 "Alice: And we must summarize these notes",
 "Bill: Good morning, yes.",
 "Alice: Okay that's everything."]
instruction = "In 1 to 5 words, generate a title for the above transcript"
print(generate(segment_transcript, instruction))
```

