



인하공업전문대학
INHA TECHNICAL COLLEGE

TCP/IP 네트워크 프로그래밍 4주차

인하공업전문대학 컴퓨터 정보과
김한결 강사

- 네트워크 기초 복습
 - ✓ 패킷이란?
 - ✓ 비트와 바이트란?
 - ✓ 문자코드
 - ✓ 랜과 WAN의 차이
 - ✓ 가정에서의 네트워크 구성
 - ✓ 회사에서의 네트워크 구성

네트워크 기초 복습

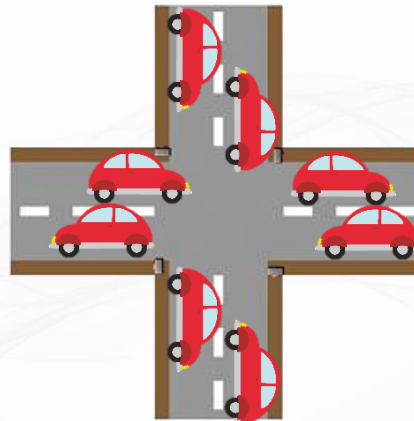
• 패킷이란?

- 네트워크를 통해 전송되는 데이터의 작은 조각
- 큰 데이터가 있더라도 작게 나누어서 보내는 게 규칙

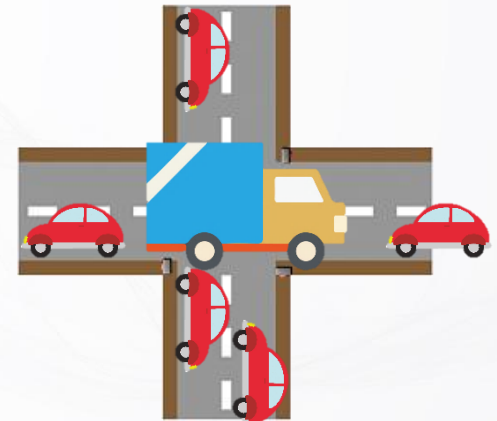
패킷 = 택배



흐름 원할



흐름 정체



대역폭 -> 일반적으로 네트워크에서 이용 가능한 최대 전송 속도로 정보를 전송할 수 있는 단위 시간당 전송량을 말함.

• 비트와 바이트란?

- 비트 -> 정보를 나타내는 최소 단위
- 2진수 '0' 과 '1' 집합인 디지털 데이터

= 비트와 바이트 =

비트 (binary digit, bit)	바이트 (Byte)
0과 1, 두 가지 값만 가질 수 있는 측정 단위	여덟 개의 비트로 구성된 데이터의 양을 나타내는 단위
 0 OFF FALSE  1 ON TRUE	

삼성반도체이야기
samsungsemiconstory.com

출처 : <https://www.samsungsemiconstory.com/2440>

네트워크 기초 복습

• 문자코드

- ASCII 코드
- 네트워크 통신, 메일 메시지 에서 사용되는 기본 문자

제어 문자			공백 문자			구두점			숫자			알파벳		
10진	16진	문자	10진	16진	문자	10진	16진	문자	10진	16진	문자	10진	16진	문자
0	0x00	NUL	32	0x20	SP	64	0x40	@	96	0x60	a			
1	0x01	SOH	33	0x21	!	65	0x41	A	97	0x61	a			
2	0x02	STX	34	0x22	"	66	0x42	B	98	0x62	b			
3	0x03	ETX	35	0x23	#	67	0x43	C	99	0x63	c			
4	0x04	EOT	36	0x24	\$	68	0x44	D	100	0x64	d			
5	0x05	ENQ	37	0x25	%	69	0x45	E	101	0x65	e			
6	0x06	ACK	38	0x26	&	70	0x46	F	102	0x66	f			
7	0x07	BEL	39	0x27	'	71	0x47	G	103	0x67	g			
8	0x08	BS	40	0x28	(72	0x48	H	104	0x68	h			
9	0x09	HT	41	0x29)	73	0x49	I	105	0x69	i			
10	0x0A	LF	42	0x2A	*	74	0x4A	J	106	0x6A	j			
11	0x0B	VT	43	0x2B	+	75	0x4B	K	107	0x6B	k			
12	0x0C	FF	44	0x2C	,	76	0x4C	L	108	0x6C	l			
13	0x0D	CR	45	0x2D	-	77	0x4D	M	109	0x6D	m			
14	0x0E	SO	46	0x2E	.	78	0x4E	N	110	0x6E	n			
15	0x0F	SI	47	0x2F	/	79	0x4F	O	111	0x6F	o			
16	0x10	DLE	48	0x30	0	80	0x50	P	112	0x70	p			
17	0x11	DC1	49	0x31	1	81	0x51	Q	113	0x71	q			
18	0x12	DC2	50	0x32	2	82	0x52	R	114	0x72	r			
19	0x13	DC3	51	0x33	3	83	0x53	S	115	0x73	s			
20	0x14	DC4	52	0x34	4	84	0x54	T	116	0x74	t			
21	0x15	NAK	53	0x35	5	85	0x55	U	117	0x75	u			
22	0x16	SYN	54	0x36	6	86	0x56	V	118	0x76	v			
23	0x17	ETB	55	0x37	7	87	0x57	W	119	0x77	w			
24	0x18	CAN	56	0x38	8	88	0x58	X	120	0x78	x			
25	0x19	EM	57	0x39	9	89	0x59	Y	121	0x79	y			
26	0x1A	SUB	58	0x3A	:	90	0x5A	Z	122	0x7A	z			
27	0x1B	ESC	59	0x3B	;	91	0x5B	[123	0x7B	{			
28	0x1C	FS	60	0x3C	<	92	0x5C	\	124	0x7C				
29	0x1D	GS	61	0x3D	=	93	0x5D]	125	0x7D	}			
30	0x1E	RS	62	0x3E	>	94	0x5E	^	126	0x7E	~			
31	0x1F	US	63	0x3F	?	95	0x5F	_	127	0x7F	DEL			

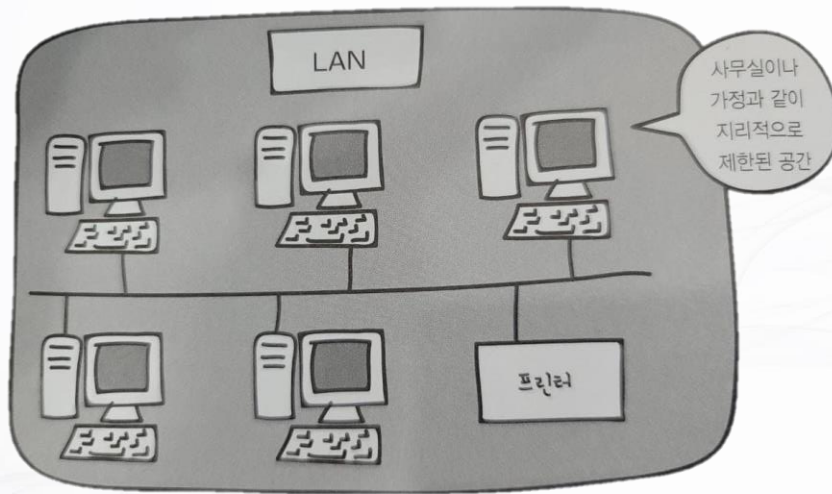
출처 :

<https://m.blog.naver.com/PostView.nhn?blogId=kyuhgmi&logNo=220073551241&proxyReferer=https:%2F%2Fwww.google.com%2F>

네트워크 기초 복습

• 랜(Lan)과 웬(Wan)의 차이

- Local Area Network (근거리 통신망)
- 사무실, 가정 같이 지리적으로 제한된 공간

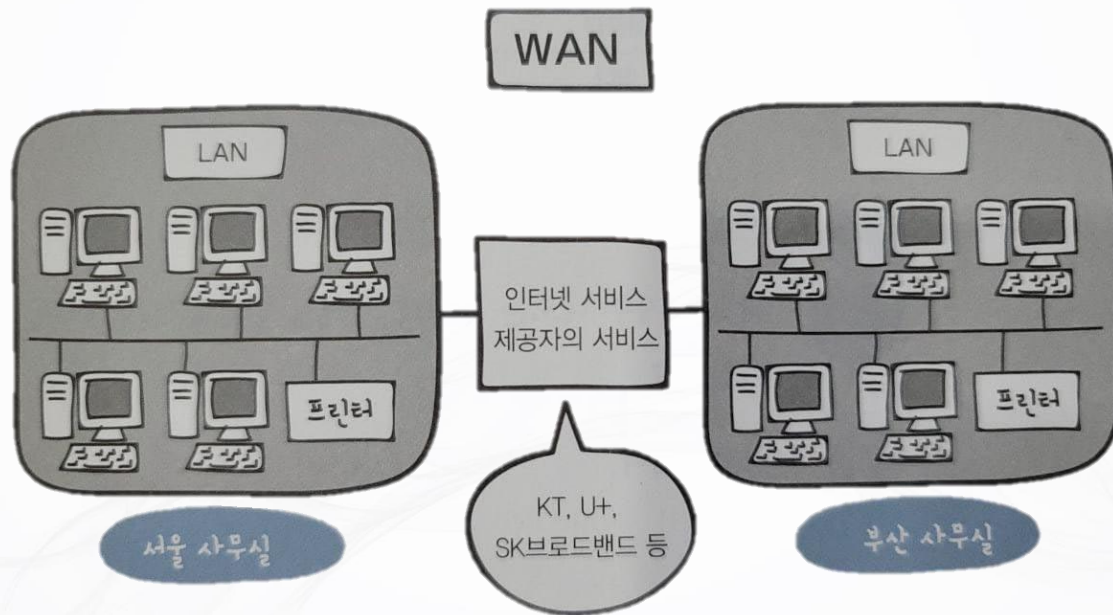


출처 : 미즈구치 카츠야 지음, 모두의 네트워크

네트워크 기초 복습

• 랜(Lan)과 Wan(Wan)의 차이

- Wide Area Network(광역 통신망)
- ISP, Internet Service Provider 인터넷 서비스 제공자



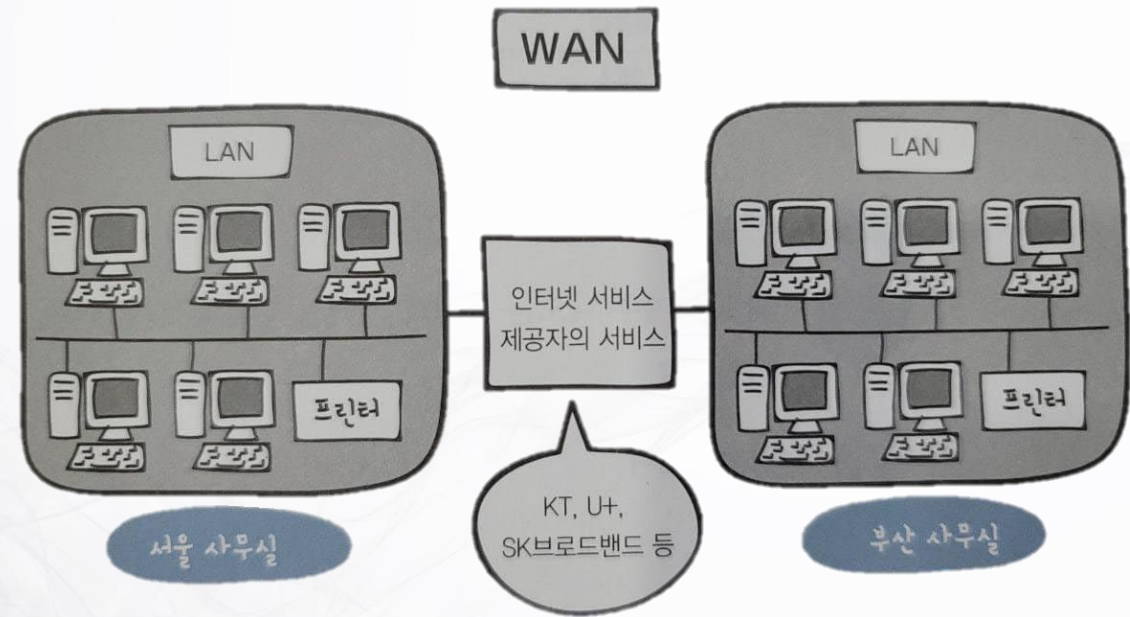
출처 : 미즈구치 카츠야 지음, 모두의 네트워크

네트워크 기초 복습

랜(Lan)과 Wan(Wan)의 차이

- Wide Area Network(광역 통신망)
- ISP, Internet Service Provider
인터넷 서비스 제공자

	LAN	WAN
범위	좁다(건물이나 특정지역)	넓다(랜과 랜을 연결)
속도	빠르다	느리다
오류	적다	많다

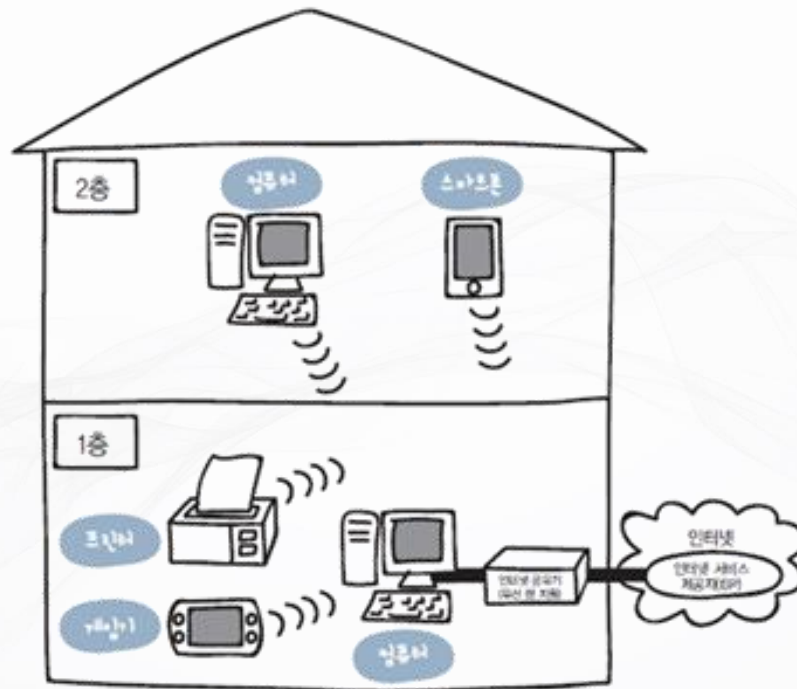


출처 : 미즈구치 카츠야 지음, 모두의 네트워크

네트워크 기초 복습

• 가정에서의 네트워크 구성

- 인터넷 서비스 제공자, 인터넷 회선을 결정 하고 계약
- 네트워크 장비 -> 유무선 공유기 사용
- 접속 방식 -> 유선 랜 방식, 무선 랜 방식

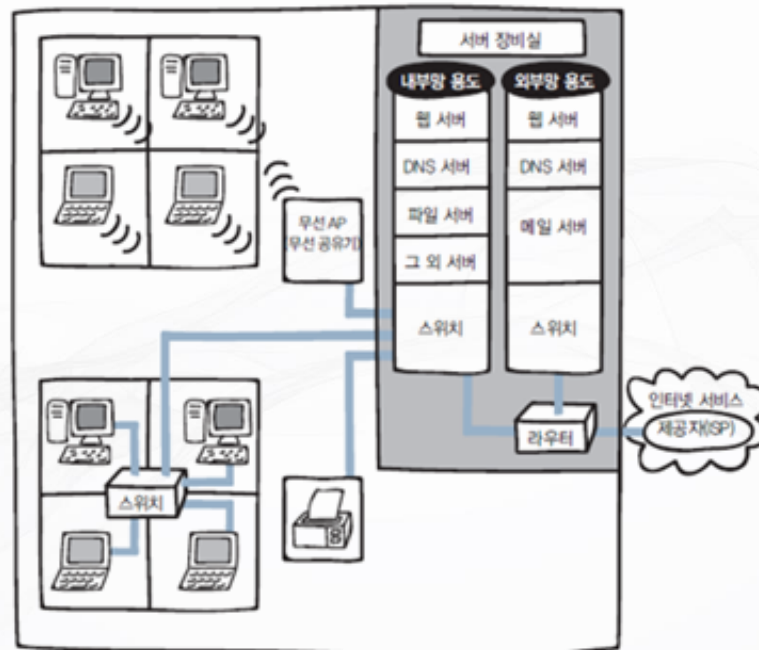


출처 : <https://gourmet-eundong.tistory.com/16>

네트워크 기초 복습

• 회사에서의 네트워크 구성

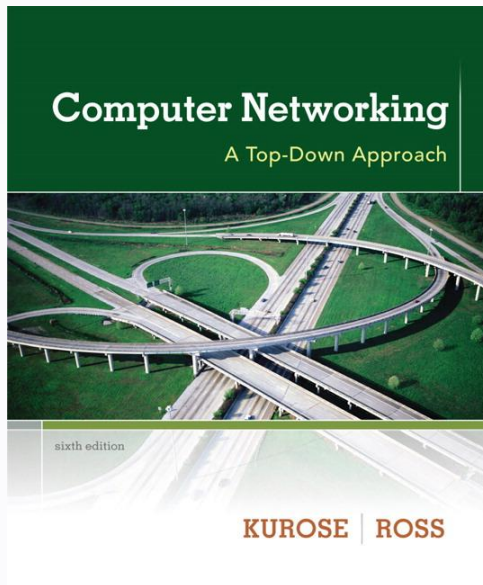
- DMZ, Demilitarized Zone -> 외부에 공개하기 위한 네트워크
- 외부에 공개하는 서버, 웹서버, DNS 서버, 메일서버 등
- 회사의 서버는 온프레미스나 클라우드로 운영



출처 : <https://gourmet-eundong.tistory.com/16>

• 연습문제

1. 네트워크에서 전송되는 작은 데이터 조각을 (**패킷**) 이라고 한다
2. 컴퓨터는 (**0**) 과 (**1**)만 이해한다
3. 정보를 표시하는 최소의 단위를 (**bit**)라고 한다.
4. 특정 건물이나 지역을 범위로 하고 속도가 빠르며 오류 발생 확률이 낮은 네트워크를 (**LAN**) 라고 한다.
5. 전기 통신 사업자가 제공하는 서비스를 사용하여 구축된 속도가 느리고 오류가 발생하기 쉬운 네트워크를 (**WAN**)이라고 한다.
6. 인터넷에 연결하려면 우선 (**ISP**) 와 인터넷 회선을 결정하고 계약한다.
7. 외부에 공개하기 위한 네트워크를 (**DMZ**)라고 한다.
8. 기업의 서버는 (**온프레미스**)나 클라우드 중 하나로 운영되고 있다.



Computer Networking: A Top Down Approach

6th edition

Jim Kurose, Keith Ross

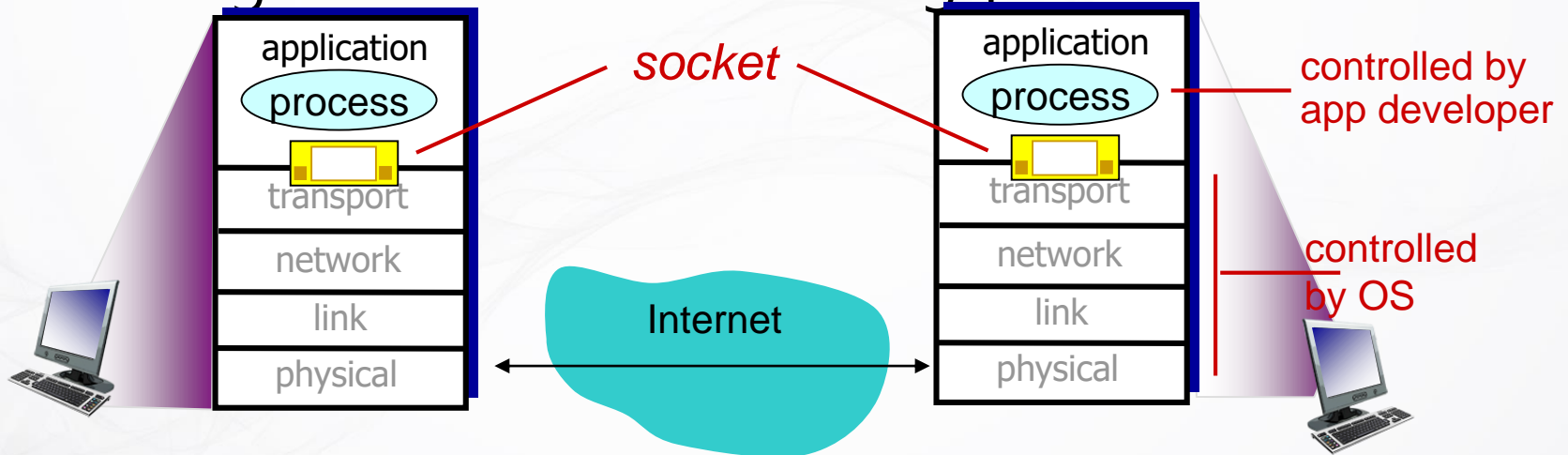
Addison-Wesley

March 2012

- sockets
- Addressing processes
- protocol
- Transport protocol
- TCP UDP
- http
- Non-persistent http

Sockets

- process sends/receives messages to/from its **socket**
- socket analogous to door
 - sending process shoves message out door
 - sending process relies on transport infrastructure on other side of door to deliver message to socket at receiving process



Addressing processes

- to receive messages, process must have *identifier*
- host device has unique 32-bit IP address
- *Q:* does IP address of host on which process runs suffice for identifying the process?
 - *A:* no, *many* processes can be running on same host
- *identifier* includes both **IP address** and **port numbers** associated with process on host.
- example port numbers:
 - HTTP server: 80
 - mail server: 25
- to send HTTP message to gaia.cs.umass.edu web server:
 - **IP address:** 128.119.245.12
 - **port number:** 80
- more shortly...

App-layer Protocol Defines

- types of messages exchanged,
 - e.g., request, response
- message syntax:
 - what fields in messages & how fields are delineated
- message semantics
 - meaning of information in fields
- rules for when and how processes send & respond to messages

open protocols:

- defined in RFCs
- allows for interoperability
- e.g., HTTP, SMTP

proprietary protocols:

- e.g., Skype

What transport service does an app need?

data integrity

- some apps (e.g., file transfer, web transactions) require 100% reliable data transfer
- other apps (e.g., audio) can tolerate some loss

timing

- some apps (e.g., Internet telephony, interactive games) require low delay to be “effective”

throughput

- ❖ some apps (e.g., multimedia) require minimum amount of throughput to be “effective”
- ❖ other apps (“elastic apps”) make use of whatever throughput they get

security

- ❖ encryption, data integrity, ...

Transport service requirements: common apps

application	data loss	throughput	time sensitive
file transfer	no loss	elastic	no
e-mail	no loss	elastic	no
Web documents	no loss	elastic	no
real-time audio/video	loss-tolerant	audio: 5kbps-1Mbps video:10kbps-5Mbps c	yes, 100' s mse
stored audio/video	loss-tolerant	same as above	
interactive games	loss-tolerant	few kbps up	yes, few secs
text messaging	no loss	elastic	yes, 100' s mse c yes and no

TCP service:

- *reliable transport* between sending and receiving process
- *flow control*: sender won't overwhelm receiver
- *congestion control*: throttle sender when network overloaded
- *does not provide*: timing, minimum throughput guarantee, security
- *connection-oriented*: setup required between client and server processes

UDP service:

- *unreliable data transfer* between sending and receiving process
- *does not provide*: reliability, flow control, congestion control, timing, throughput guarantee, security, or connection setup,

Q: why bother? Why is there a UDP?

Internet apps: application transport protocols

application	application layer protocol	underlying transport protocol
e-mail	SMTP [RFC 2821]	TCP
remote terminal access	Telnet [RFC 854]	TCP
Web	HTTP [RFC 2616]	TCP
file transfer	FTP [RFC 959]	TCP
streaming multimedia	HTTP (e.g., YouTube), RTP [RFC 1889]	TCP or UDP
Internet telephony	SIP, RTP, proprietary (e.g., Skype)	TCP or UDP

- *web page* consists of *objects*
- object can be HTML file, JPEG image, Java applet, audio file,...
- web page consists of *base HTML-file* which includes *several referenced objects*
- each object is addressable by a *URL*, e.g.,

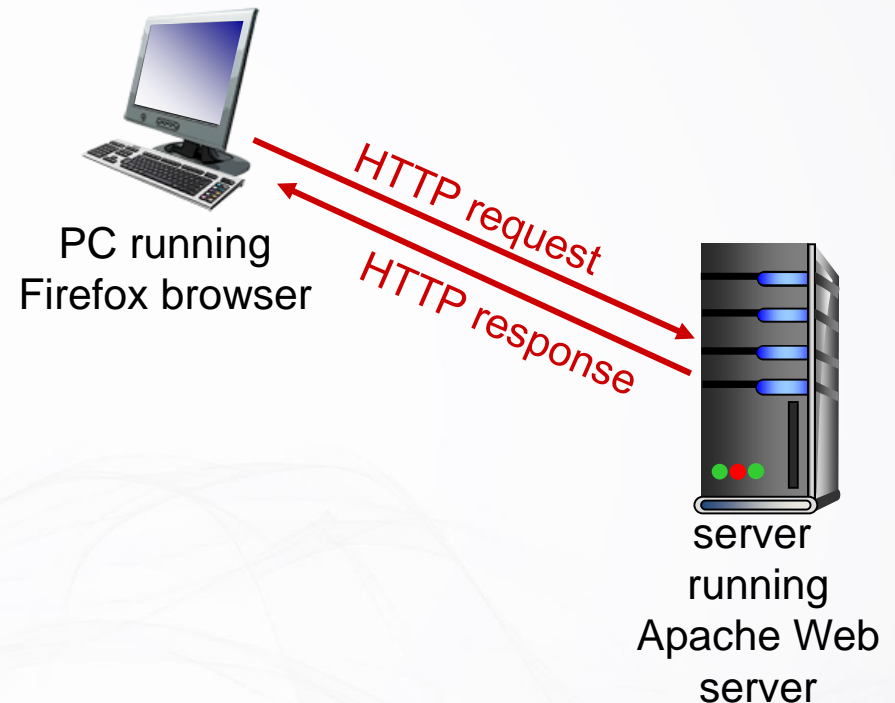
`www.inhatec.ac.kr/someDept/pic.gif`

host name

path name

HTTP: hypertext transfer protocol

- Web's application layer protocol
- client/server model
 - *client*: browser that requests, receives, (using HTTP protocol) and “displays” Web objects
 - *server*: Web server sends (using HTTP protocol) objects in response to requests



uses TCP:

- client initiates TCP connection (creates socket) to server, port 80
- server accepts TCP connection from client
- HTTP messages (application-layer protocol messages) exchanged between browser (HTTP client) and Web server (HTTP server)
- TCP connection closed

HTTP is “stateless”

- server maintains no information about past client requests

aside
protocols that maintain “state” are complex!

- ❖ past history (state) must be maintained
- ❖ if server/client crashes, their views of “state” may be inconsistent, must be reconciled

non-persistent HTTP

- at most one object sent over TCP connection
 - connection then closed
- downloading multiple objects required multiple connections

persistent HTTP

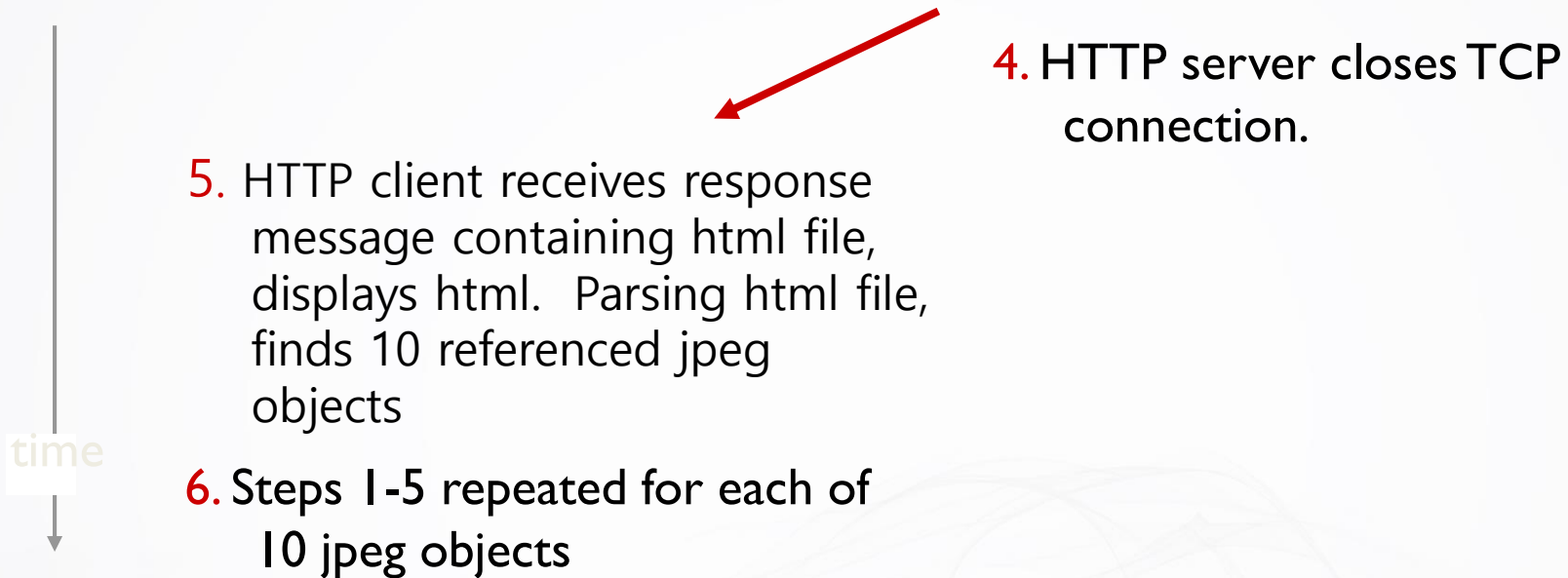
- multiple objects can be sent over single TCP connection between client, server

Non-persistent HTTP

-
- ```
graph LR; 1a[1a. HTTP client initiates TCP connection to HTTP server (process) at www.inhatc.ac.kr on port 80] --> 1b[1b. HTTP server at host www.inhatc.ac.kr waiting for TCP connection at port 80. "accepts" connection, notifying client]; 1b --> 2[2. HTTP client sends HTTP request message (containing URL) into TCP connection socket. Message indicates that client wants object someDepartment/home.index]; 2 --> 3[3. HTTP server receives request message, forms response message containing requested object, and sends message into its socket];
```
- 1a. HTTP client initiates TCP connection to HTTP server (process) at `www.inhatc.ac.kr` on port 80
- 1b. HTTP server at host `www.inhatc.ac.kr` waiting for TCP connection at port 80. “accepts” connection, notifying client
2. HTTP client sends HTTP *request message* (containing URL) into TCP connection socket. Message indicates that client wants object `someDepartment/home.index`
3. HTTP server receives request message, forms *response message* containing requested object, and sends message into its socket

time

## Non-persistent HTTP (cont.)



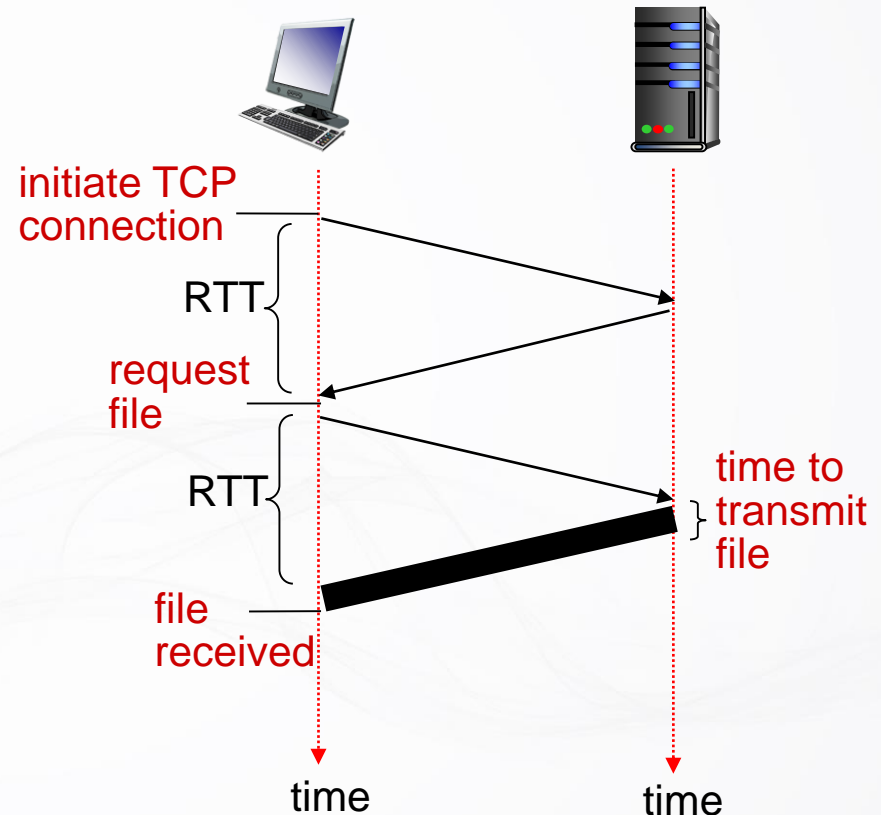
## Non-persistent HTTP : response time

**RTT (definition):** time for a small packet to travel from client to server and back

**HTTP response time:**

- one RTT to initiate TCP connection
- one RTT for HTTP request and first few bytes of HTTP response to return
- file transmission time
- non-persistent HTTP response time =

$2\text{RTT} + \text{file transmission time}$



**4주차 수업이 끝났습니다**

**고생하셨습니다.**

