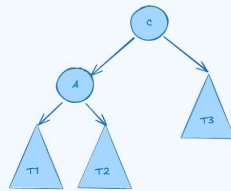
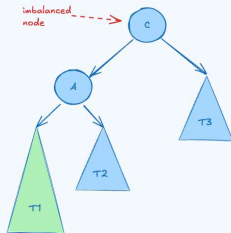


STARTING TREE for LL and LR



Before an insertion, height(C.left) and height(C.right) differ by 1.

Case 1: Left-Left Imbalance (Single Rotation)



Then, there is an insertion somewhere in T1 that causes A's height to increase by 1. Now, height(C.left) and height(C.right) differ by 2.

Important Observations:

- All values in T1 are smaller than both A and C.
- All values in T2 are bigger than A but smaller than C.
- All values in T3 are bigger than C.

→ Since all values in T2 are bigger than A but smaller than C, that subtree could live to the right of A or the left of C

```

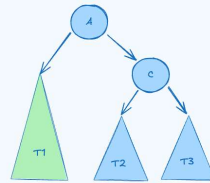
fun rotateWithLeftChild(C, parentOfC):
  A = C.left
  C.left = A.right
  A.right = C
  
```

```

if C == root of tree:
  root of tree = A
else:
  if parentOfC.left = C:
    parentOfC.left = A
  else
    parentOfC.right = A
  
```

```

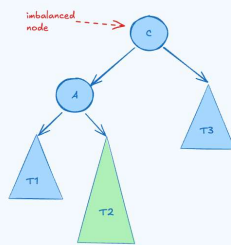
updateHeight(C)
updateHeight(A)
  
```



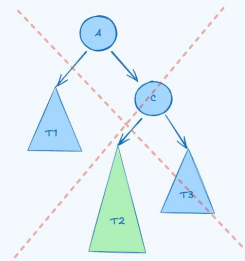
We can adjust for this imbalance with a SINGLE ROTATION by rotating the node of imbalance with its left child.

Notice after the re-balancing, A has the same height as C did in the original tree.

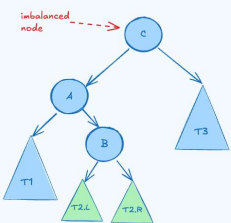
Case 2: Left-Right Imbalance (Double Rotation)



An insertion in T2 causes C to become the node of imbalance.



Notice that a Single Rotation wouldn't fix things.



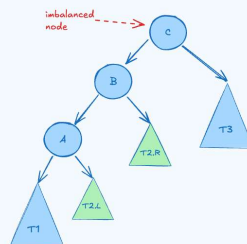
B represents the root of T2. Inserting into T2 cannot be fixed by a single rotation (shown above). T2.L and T2.R are shown as half way to the next level because only 1 (where the insertion happened) would extend down to the red --- level.

Observations:

- All values in T2.L fall between A and B.
- So, T2.L could be the left child of B or the right child of A.
- All values in T2.R fall between B and C.
- So, T2.R could be the right child of B or the left child of C.

To fix:

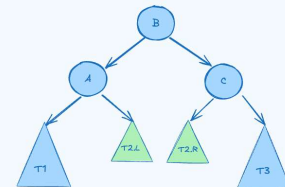
- Step 1: Rotate A with its right child (B). (same rotation as RR imbalance)
- Step 2: Rotate C with its left child (now it is B after step 1).



Step 1: Rotate A with Right Child B.

```

fun doubleRotateWithLeftChild(C, parentOfC):
  rotateWithRightChild(A, parentOfC)
  rotateWithLeftChild(C, parentOfC)
  
```

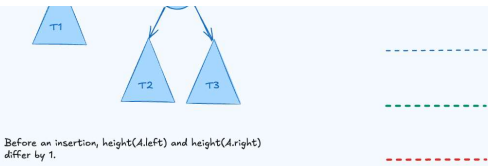


Step 2: Rotate C with Right Child B.

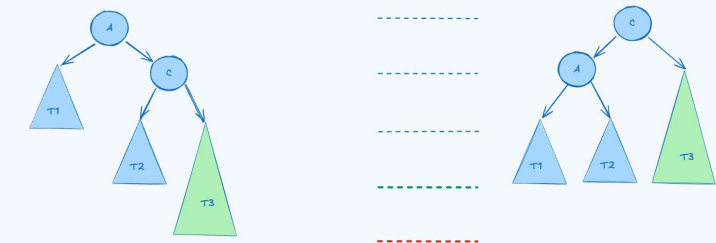
After this rotation, either T2.L or T2.R will be as deep as T1 and T3, but not both.

STARTING TREE For RR and RL





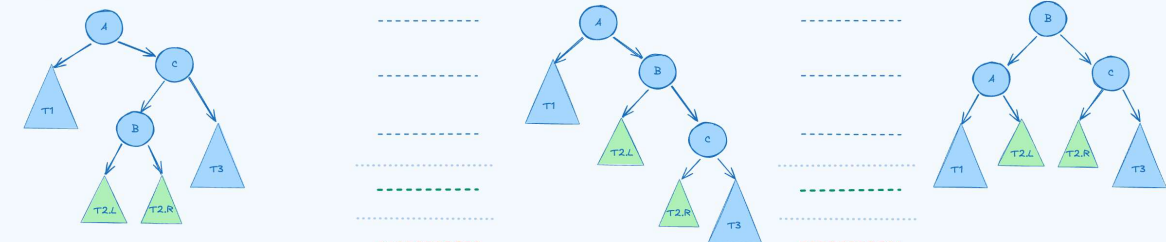
Case 4: Right Right Imbalance (Single Rotation)



Insertion into T3 causes it to be 2 levels deeper than T1, making A the node of imbalance.

Remember: All values in T2 fall between A and C. So T2 could be connected to the left of C or the Right of A.

Case 3: Right Left Imbalance (Double Rotation)



As with Case 2, an insertion into T2 cannot be solved with a single rotation. So, we need to explicitly consider the root of T2, labeled as B here.