

Primitive and Geometric Progression Free Sets Without Large Gaps

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Gaps between prime numbers

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In theory: At most 10^{53} . (Best we can prove)

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Summary: The largest gap between primes up to x has size between

$$\frac{\log x \log \log x \log \log \log \log x}{\log \log \log x} \quad \text{and} \quad x^{0.525}$$

but we think the truth is $\log^2 x$.

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Definition

A set $S \subset \mathbb{N}$ is **primitive** if no element of the set divides another: if $m, n \in S$ are distinct then $m \nmid n$.

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For example, $\mathcal{P}_2 = \{4, 6, 9, 10, 14, 15, 21, 22, \dots\}$.

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Theorem (Erdős)

If S is a primitive set then

$$\sum_{n \in S} \frac{1}{n \log n} < \infty.$$

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Conjecture (Erdős)

If S is a primitive set then

$$\sum_{n \in S} \frac{1}{n \log n} < \sum_{p \in \mathcal{P}} \frac{1}{p \log p} = 1.63 \dots$$

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Due to the nature of the construction of these sets, it isn't possible to get good upper bounds for the largest gap in such a set.

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Theorem (Linear Sieve of Rosser and Iwaniec)

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$$P(X_i = m) = \frac{C_\epsilon}{m^{1+\epsilon/4}}$$

with $C_\epsilon = \frac{1}{\zeta(1+\epsilon/4)}$ chosen to normalize the distribution.

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consist of precisely those integers, n , where the total number of prime factors of n agrees with the variable X_i for **all** its prime divisors, p_i .

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as $x \rightarrow \infty$.

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Using that $y = \exp \left(\sqrt{(2 + \epsilon) \log x \log \log x} \right)$ The inner exponent is

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The probability that no integer in $[x - cy, x]$ is included in Q is at most $\exp\left(-\exp\left(C'_\epsilon \sqrt{\log x \log \log x}\right)\right)$.

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Thus there exists a set, Q satisfying the properties of the theorem.

Pairwise coprime sets

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Theorem (M.)

For every $\epsilon > 0$ there exists a pairwise coprime set $b_1, b_2, b_3 \dots$, such that

$$b_n - b_{n-1} \ll b_n^{\alpha+\epsilon}$$

where $\alpha = \frac{5-\sqrt{17}}{2} = 0.43845\dots$

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The squarefree integers s_1, s_2, \dots satisfy

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We can find larger sets that avoid progressions, but can't say much about the gaps of these sets.

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Note V avoids 3-term progressions: If (a, ar, ar^2) is a progression, a prime in r appears to different powers in at least two of these terms.

Thank you!