Counting primitive subsets and other statistics of the divisor graph of $\{1, 2, \dots n\}$

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- $\mathcal{P}_0 = \{1\}.$



Theorem (Ahlswede, Khatchatrian, Sárközy, 1999)

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If we restrict to primitive subsets of $\{1, 2, \dots n\}$ then the set $\left\{ \left\lfloor \frac{n}{2} \right\rfloor + 1, \dots n \right\}$ is primitive, and has size $\left\lfloor \frac{n}{2} \right\rfloor$.

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What do primitive subsets of the first few integers look like?

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How many primitive subsets of $\{1, 2 \dots n\}$ are there?

Counting primitive sets

Let Q(n) count the primitive sets with largest element at most n.

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              Number of primitive subsequences of \{1, 2, ..., n\}.
   1, 2, 3, 5, 7, 13, 17, 33, 45, 73, 103, 205, 253, 505, 733, 1133, 1529, 3057, 3897,
   7793, 10241, 16513, 24593, 49185, 59265, 109297, 163369, 262489, 355729, 711457, 879937,
   1759873, 2360641, 3908545, 5858113, 10534337, 12701537, 25403073, 38090337, 63299265,
   81044097, 162088193, 205482593, 410965185, 570487233, 855676353 (list; graph; refs; listen; history; text;
   internal format)
   OFFSET
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   COMMENTS
                 a(n) counts all subsequences of {1, ..., n} in which no term divides any
                   other. If n is a prime a(n) = 2*a(n-1)-1 because for each subsequence s
                   counted by a(n-1) two different subsequences are counted by a(n): s and
                   s.n. There is only one exception: 1.n is not a primitive subsequence
                   because 1 divides n. For all n>1: a(n) < 2*a(n-1). - Alois P. Heinz, Mar
                   07 2011
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Every subset of $\left(\frac{n}{2},n\right]$ is primitive. There are $2^{\lceil\frac{n}{2}\rceil} \geq \sqrt{2}^n$ such subsets.

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Proof is not effective: Gives no insight on the value of this constant.



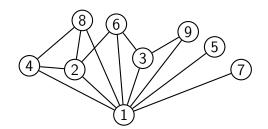
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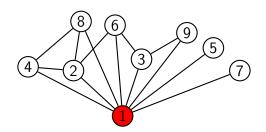
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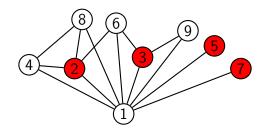
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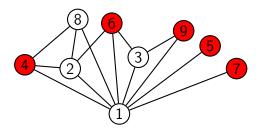
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Proof: Graphs, and smooth number estimates.



Numerical Bound



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A general theorem

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Theorem (M., 2019)

Fix $\epsilon > 0$, $A \ge 0$. Suppose f(k, n) depends only on the connected component of k in the divisor graph of [k, n] and $|f(k, n)| \le A$. Then

$$\sum_{a=1}^{n} f(a, n) = nC_f + O_A \left(n \exp \left(-\sqrt{\left(\frac{1}{6} - \epsilon\right) \log n \log \log n} \right) \right)$$

where

$$C_f = \sum_{i=1}^{\infty} \sum_{\substack{d \ p+(d) < i}} \sum_{t \in [id,(i+1)d)} \left(rac{f(d,t)}{t(t+1)} \prod_{p \leq i} rac{p-1}{p}
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Corollary

Let V(n) denote the median size of primitive subsets of [1, n]. Then

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Question

Is $V(n) \sim vn$ for some v? If so, is v computable?

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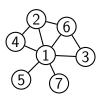
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Example: C(7) = 2



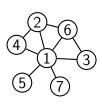
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The divisor graph of [1,7] can be covered by $\{7,1,5\}$ and $\{3,6,2,4\}$ but it is not possible to use a single path.



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The constant ρ is effectively computable and 0.1916 $< \nu <$ 0.2143.

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Thank you!

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