

Setup

Grab a GPU node:

```
$ source /projects/hpc/bootcamp/login-gpu.sh
```

Once on a node, load the GPU programming package and download today's examples:

```
$ module load CUDA
$ git clone https://github.com/agrippa/hpc-bootcamp $HOME/bootcamp-gpu
```

It also may be useful to have the CUDA API references open:

<http://docs.nvidia.com/cuda/cuda-runtime-api/index.html>

Slides are accessible in your browser at: <https://github.com/agrippa/hpc-bootcamp>

GPU Accelerated Computing

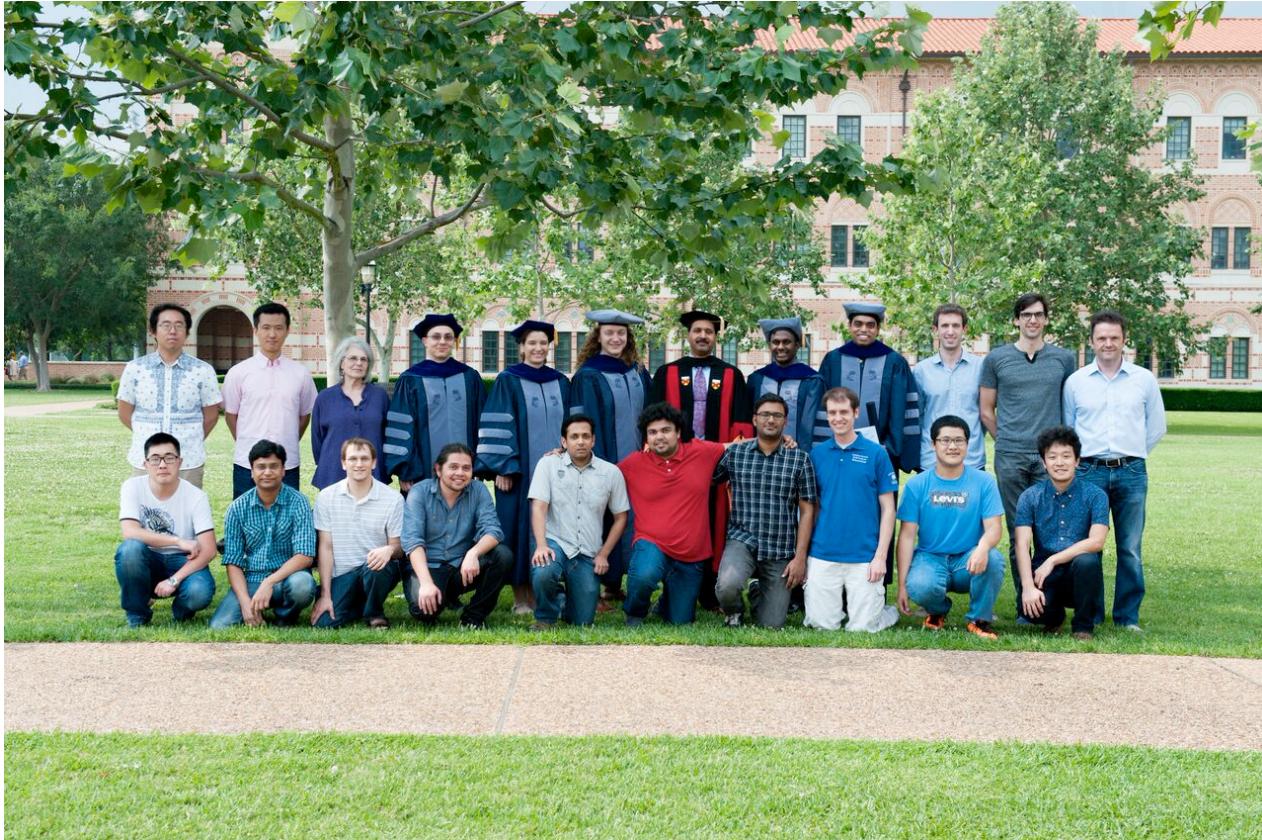
Max Grossman

Habanero Extreme Scale Software Research Group, Rice University
Author, Professional CUDA C Programming

Outline

1. A Brief History of Computing
2. GPU Architecture Overview
3. An Overview of CUDA
4. Hands-On with a Scientific App
5. Related Topics:
 1. CUDA Libraries
 2. Alternative GPU Programming Models
 3. Alternatives to GPUs

About the Speaker



Group home page:
habanero.rice.edu

Group Github:
github.com/habanero-rice

Programming resources:
hc.rice.edu

About the Speaker

Application Development:

- NASA JPL: Atmospheric simulators
- Repsol R&D: GPU Enablement
- Halliburton, Shell R&D: New application development

Runtime Systems for HPC

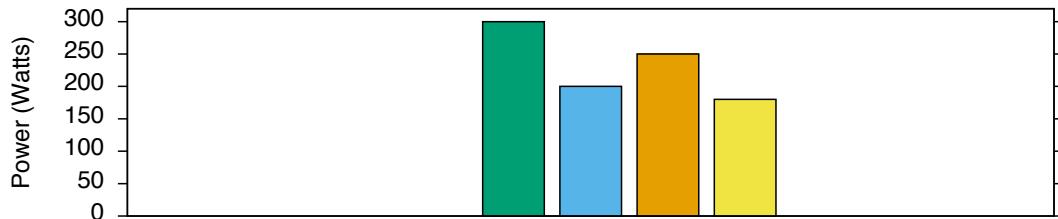
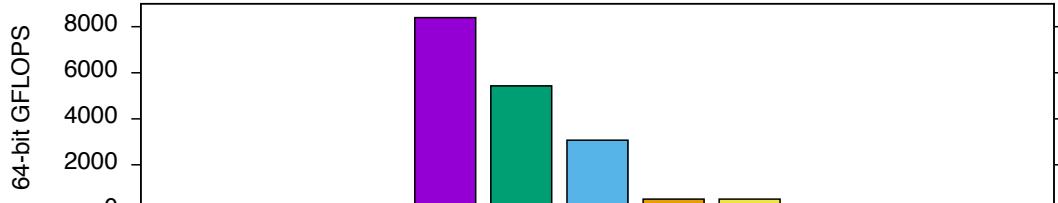
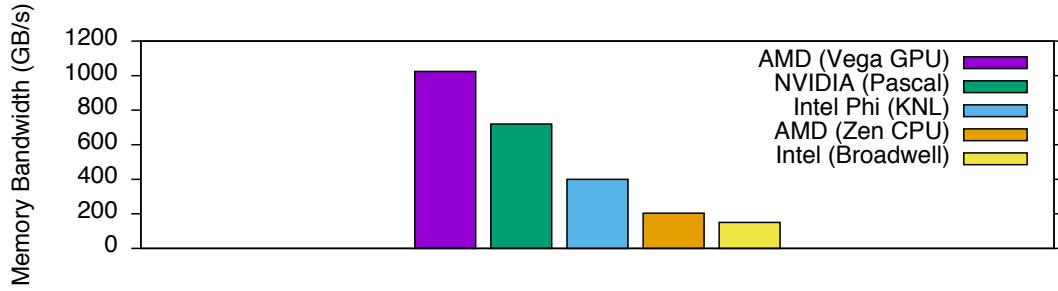
Data Analytics (“Big Data”)

Compilers and Programmer Tools



Why GPUs?

Upcoming Hardware Generations



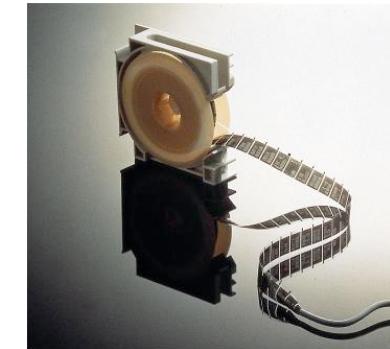
Is your hardware holding you back, or are you holding back your hardware?

A (Brief) History of Early Personal Computing

1971: Intel 4004

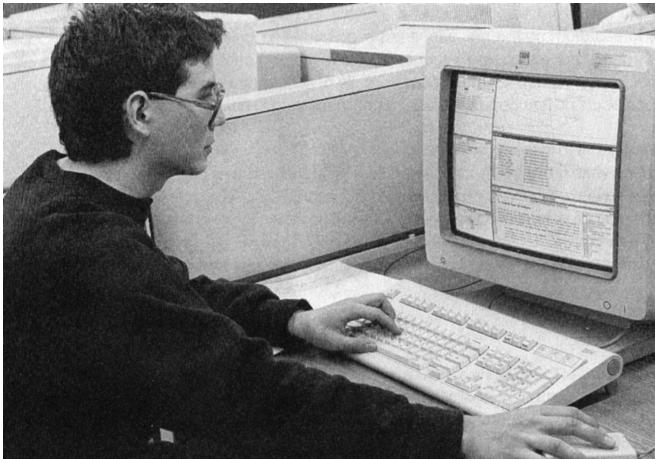


1974: Intel 8080



A (Brief) History of Early Personal Computing

1982: AMD licensed to sell Intel 8086/8088 processors, because of demand for IBM workstations



1987: First Sun SPARC released



A (Brief) History of Early Personal Computing



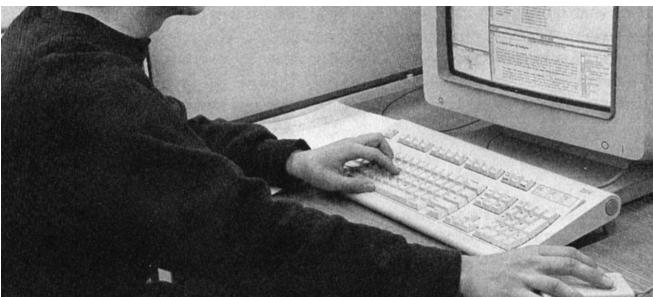
A (Brief) History of Early Personal Computing



What do all of these applications have in common?

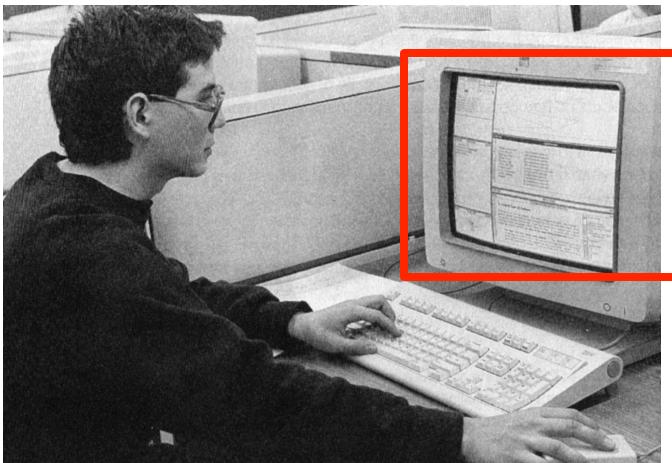


Application performance is entirely dependent on straight-line performance.



A (Brief) History of Early Personal Computing

First dedicated Graphics Processing Units (GPUs) introduced in the 1990s.



A (Brief) History of Early Personal Computing

Only became programmable in the early 2000s.

Game Developers Conference

Standard “Grid” Computation

- Initialize “view” (pixels:texels::1:1)

```
glMatrixMode(GL_MODELVIEW);
glLoadIdentity();
glMatrixMode(GL_PROJECTION);
glLoadIdentity();
glOrtho(0, 1, 0, 1, 0, 1);
glViewport(0, 0, gridResX, gridResY);
```
- For each algorithm step:
 - Activate render-to-texture
 - Setup input textures, fragment program
 - Draw a full-screen quad (1 unit x 1 unit)



Game Developers Conference

GPU: high performance growth

- CPU
 - Annual growth ~1.5x → decade growth ~ 60x
 - Moore's law
- GPU
 - Annual growth > 2.0x → decade growth > 1000x
 - Much faster than Moore's law



A (Brief) History of Early Personal Computing

| | |
|---|--|
| 1 | Tianhe-2 [MilkyWay-2] - TH-IVB-FEP Cluster, Intel Xeon E5-2692 12C 2.200GHz, TH Express-2, Intel Xeon Phi 31S1P NUDT |
| 2 | Titan - Cray XK7 , Opteron 6274 16C 2.200GHz, Cray Gemini interconnect, NVIDIA K20x Cray Inc. |
| 3 | Sequoia - BlueGene/Q, Power BQC 16C 1.60 GHz, Custom IBM |
| 4 | K computer, SPARC64 VIIIfx 2.0GHz, Tofu interconnect Fujitsu |
| 5 | Mira - BlueGene/Q, Power BQC 16C 1.60GHz, Custom IBM |
| 6 | Trinity - Cray XC40, Xeon E5-2698v3 16C 2.3GHz, Aries interconnect Cray Inc. |
| 7 | Piz Daint - Cray XC30, Xeon E5-2670 8C 2.600GHz, Aries interconnect , NVIDIA K20x Cray Inc. |

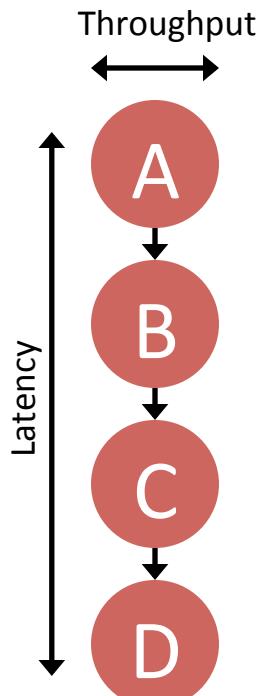
And a decade later...

| Green500 Rank | MFLOPS/W | Site* | Computer* | Total Power (kW) |
|---------------|----------|--|--|------------------|
| 1 | 7,031.58 | Institute of Physical and Chemical Research (RIKEN) | Shoubu - ExaScaler-1.4 80Brick, Xeon E5-2618Lv3 8C 2.3GHz, Infiniband FDR, PEZY-SC | 50.32 |
| 2 | 5,331.79 | GSIC Center, Tokyo Institute of Technology | TSUBAME-KFC/DL - LX 1U-4GPU/104Re-1G Cluster, Intel Xeon E5-2620v2 8C 2.1GHz, Infiniband FDR, NVIDIA Tesla K80 | 51.13 |
| 3 | 5,271.81 | GSI Helmholtz Center | ASUS ESC4000 FDR/G2S, Intel Xeon E5-2690v2 10C 3GHz, Infiniband FDR, AMD FirePro S9150 | 57.15 |
| 4 | 4,778.46 | Institute of Modern Physics (IMP), Chinese Academy of Sciences | Sugon Cluster W780i, Xeon E5-2640v3 8C 2.6GHz, Infiniband QDR, NVIDIA Tesla K80 | 65.00 |
| 5 | 4,112.11 | Stanford Research Computing Center | XStream - Cray CS-Storm, Intel Xeon E5-2680v2 10C 2.8GHz, Infiniband FDR, Nvidia K80 | 190.00 |
| 6 | 3,856.90 | IT Company | Inspur TS10000 HPC Server, Xeon E5-2620v3 8C 2.4GHz, 10G Ethernet, NVIDIA Tesla K40 | 58.00 |
| 7 | 3,775.45 | Internet Service | Inspur TS10000 HPC Server, Intel Xeon E5-2620v2 8C 2.1GHz, 10G Ethernet, NVIDIA Tesla K40 | 110.00 |
| 8 | 3,775.45 | Internet Service | Inspur TS10000 HPC Server, Intel Xeon E5-2620v2 8C 2.1GHz, 10G Ethernet, NVIDIA Tesla K40 | 110.00 |
| 9 | 3,775.45 | Internet Service | Inspur TS10000 HPC Server, Intel Xeon E5-2620v2 8C 2.1GHz, 10G Ethernet, NVIDIA Tesla K40 | 110.00 |
| 10 | 3,775.45 | Internet Service | Inspur TS10000 HPC Server, Intel Xeon E5-2620v2 8C 2.1GHz, 10G Ethernet, NVIDIA Tesla K40 | 110.00 |

| TITAN | SUMMIT |
|----------------------------------|---------------------------------------|
| 18,688 | ~3,400 |
| (1) 16-core AMD Opteron per node | (Multiple) IBM POWER 9s per node |
| (1) NVIDIA Kepler K20x per node | (Multiple) NVIDIA Volta GPUs per node |
| 32GB (DDR3) | >512GB (HBM+DDR4) |
| PCI Gen2 | NVLINK (5-12x PCIe3) |
| Gemini | Dual Rail EDR-IB (23 GB/s) |
| 9 MW | 10 MW |

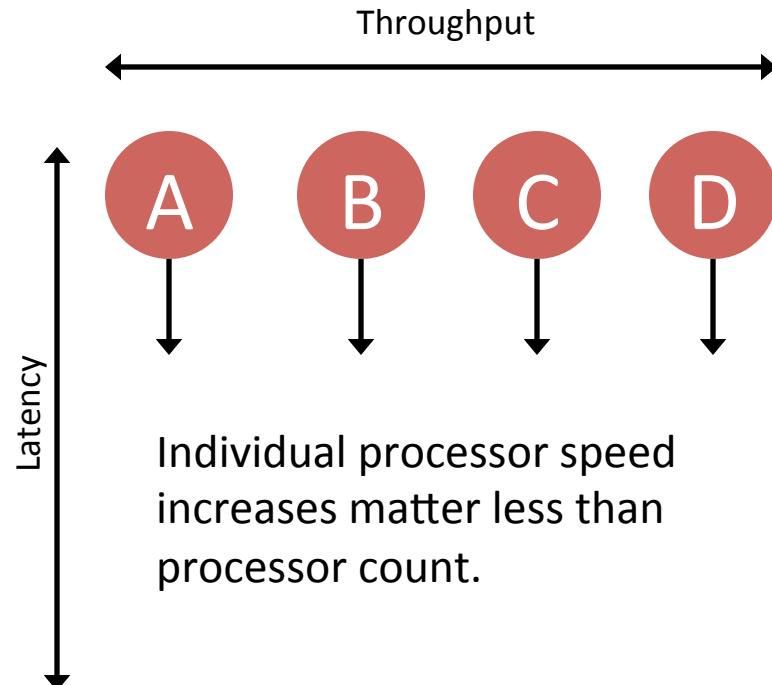
Latency vs. Throughput in Apps & Datasets

Latency-bound



Increasing the # of processors doesn't matter, only straight-line performance.

Throughput-bound



Individual processor speed increases matter less than processor count.

Latency-Optimized Processors

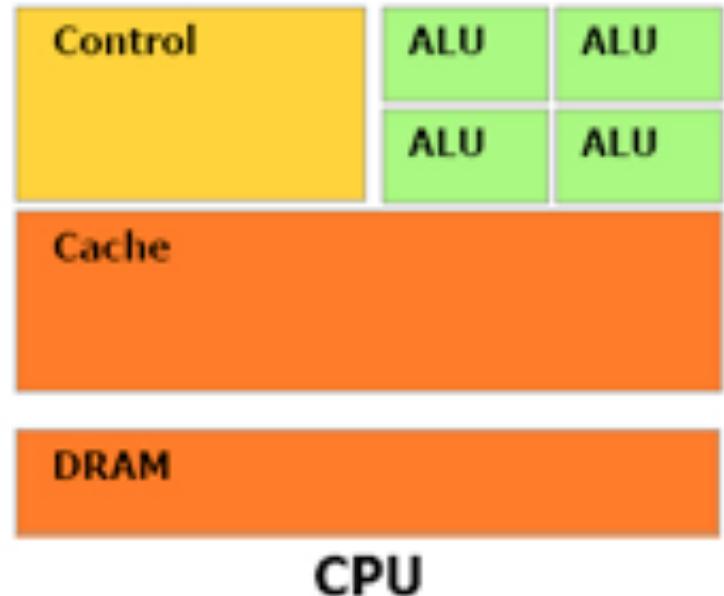
For most consumer (i.e. interactive) applications, latency is the #1 performance metric.

↑ clock frequencies

↑ control logic (e.g. branch prediction)

↑ cache per core

↓ parallelism/throughput



Throughput-Optimized Processors

For many scientific applications, throughput is the #1 performance metric.

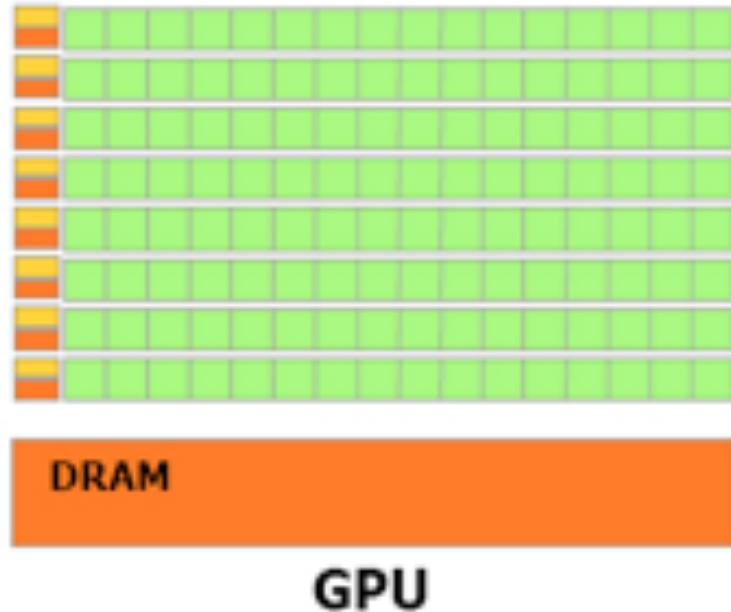
↑ cores/concurrency

↑ memory bandwidth

↓ clock frequency

↓ cache

↓ control logic



Throughput-Bound Applications

While many applications are latency-bound, the important ones (scientific, finance, medical, etc.) tend to be throughput bound.

TABLE 6: Hardware Configuration

| Device | Type | Number of cores | Memory size |
|---------|-----------------------------|-----------------|--------------------|
| 4 x GPU | GeForce GTX Titan 0.876 GHz | 2688 | 6 GB Global Memory |
| CPU | Intel Core i7-4770 3.40 GHz | 4 | 32 GB |

TABLE 7: Results

| Number of Elements | Simulation Time (s) | | |
|--------------------|---------------------|-----------------------|-------------------------|
| | Chronos 4 GPUs | Well Known FE Program | Performance Improvement |
| 200.000 | 21 | 516 (8.6 min) | 24,57 x |
| 500.000 | 43 | 3407 (56.78 min) | 79,23 x |
| 1.000.000 | 83 | Insufficient Memory | x |
| 2.000.000 | 168 | Insufficient Memory | x |

“Speeding up a Finite Element Computation on GPU”, GTC 15

| Reconstruction Stage | Single CPU time | Single GPU time | Single GPU speedup | Two GPU time | Two GPU speedup |
|------------------------------------|-----------------|-----------------|--------------------|--------------|-----------------|
| Preprocessing and Support Function | 34.09 | 9.36 | 3.64X | 7.72 | 4.42X |
| Refraction-Corrected Ray Tracing | 1899.98 | 63.29 | 30.02X | 45.53 | 41.73X |
| Compounding Views | 39.33 | 0.84 | 46.71X | 0.84 | 46.71X |
| Entire Reflection Reconstruction | 2108.40 | 79.16 | 26.63X | 54.57 | 38.64X |

“Multi-GPU Accelerated Refraction-Corrected Reflection Image Reconstruction for 3D Ultrasound Breast Imaging”, GTC 15

Overview of (NVIDIA) GPU Architecture

Architectural Generations

Discussion will focus on characteristics shared across architectural generations.

When important, differences between generations will be noted.

- Fermi → Kepler → Maxwell → Pascal

Much of the material in this section can be found in the white papers below.



http://www.nvidia.com/content/pdf/fermi_white_papers/nvidia_fermi_compute_architecture_whitepaper.pdf

<https://www.nvidia.com/content/PDF/kepler/NVIDIA-Kepler-GK110-Architecture-Whitepaper.pdf>

<https://developer.nvidia.com/maxwell-compute-architecture>

<http://images.nvidia.com/content/pdf/tesla/whitepaper/pascal-architecture-whitepaper.pdf>

The CUDA Core

The CUDA Core is the finest granularity of execution on NVIDIA GPUs.

In a single cycle, executes either:

- 32-bit integer arithmetic instruction
- 32-bit floating point arithmetic instruction

CUDA
Core

CUDA core != x86 core.

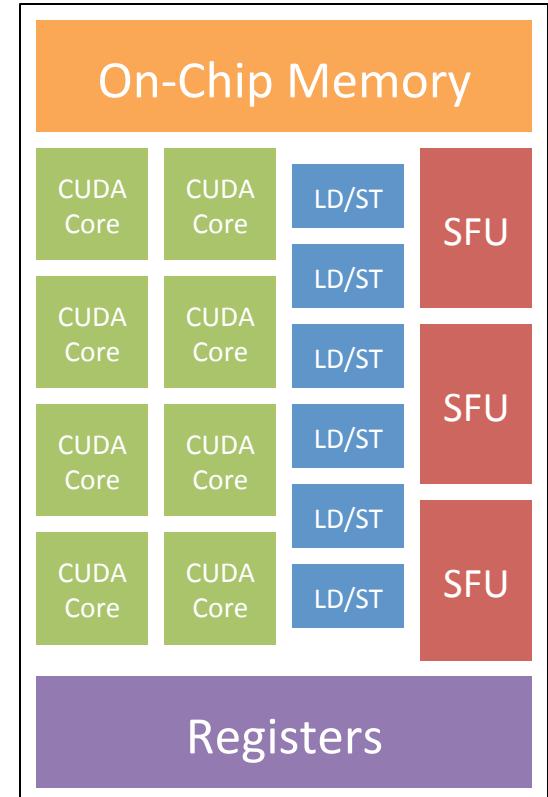
The Streaming Multiprocessor (SM)

SM is analogous to “core” in x86 processors

Consists of many CUDA cores

- 32 for Fermi, 192 for Kepler, 128 for Maxwell, 64 for Pascal
- All CUDA cores in the same SM execute in lock step
- Registers + on-chip memory are shared among threads running on this SM

Think “vector core” but with more flexibility.

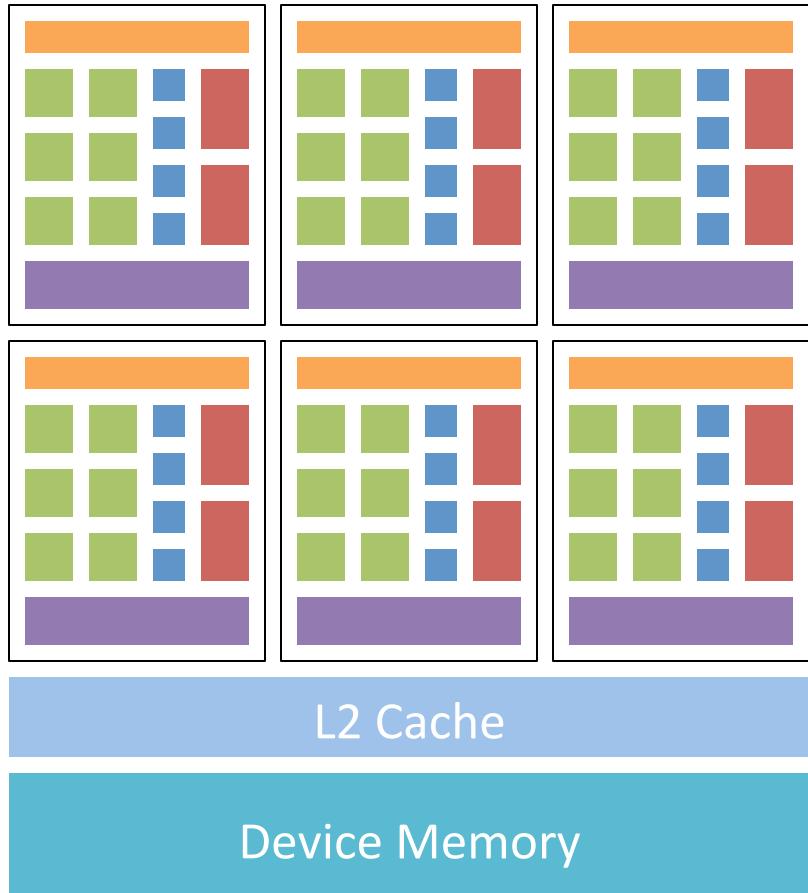


The GPU

Consists of $O(10)$ SMs

L2 cache, device memory (DRAM) shared among all SMs

Off-chip latency $\sim 100x$ worse than on-chip



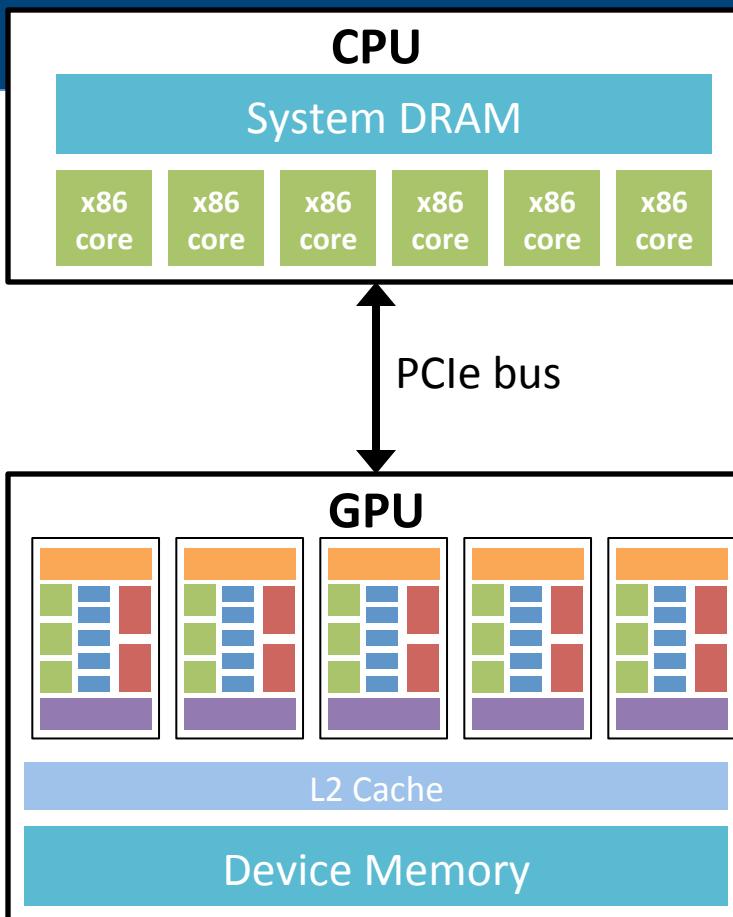
The Full System

GPUs are not standalone processors, require a CPU to manage them.

Commonly refer to CPU as “host” and GPU as “device”.

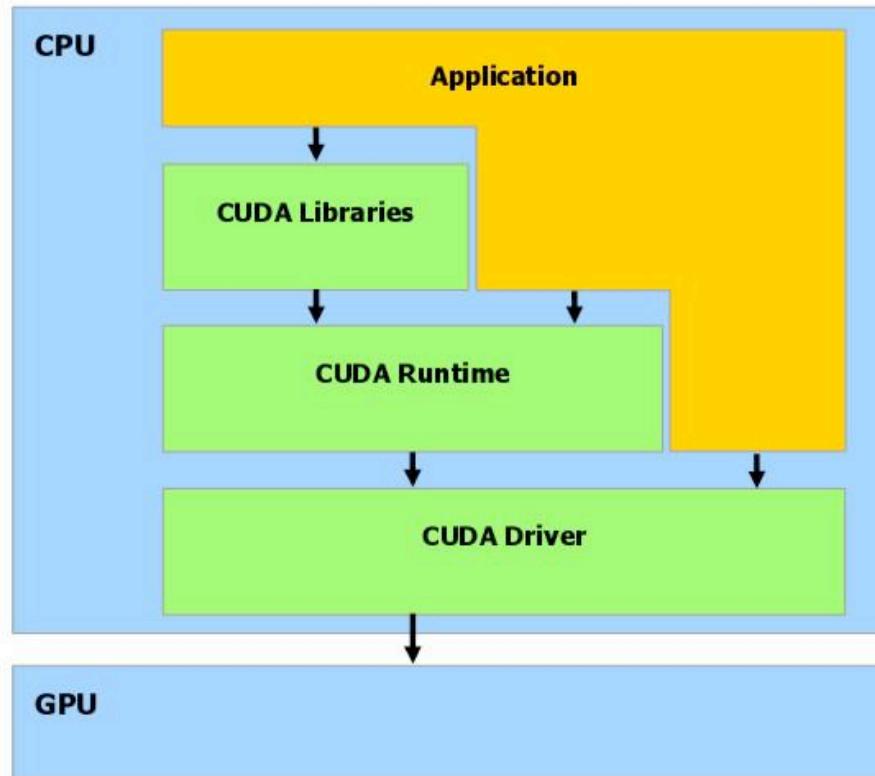
Connected by very high latency PCIe bus.

Analogous to two MPI processes, versus OpenMP/Cilk everything shared model.



Hands-On With CUDA, a Simple Example

CUDA Software Stack



A Simple CUDA Example

Kernel is single-threaded.

A single CUDA thread per CUDA core.

Let's start by considering vector addition: $C = B + A$

```
void vector_add(int *C, int *B, int *A, int N) {  
    int i;  
    for (i = 0; i < N; i++) {  
        C[i] = A[i] + B[i];  
    }  
}
```



CUDA
Core

A Simple CUDA Example

To enable GPU execution of this function, simply add **__global__**:

```
__global__ void vector_add(int *C, int *B, int *A,  
    int N) {  
    int i;  
    for (i = 0; i < N; i++) {  
        C[i] = A[i] + B[i];  
    }  
}
```



CUDA
Core

A Simple CUDA Example

CUDA function type qualifiers:

| Qualifier | Executes on... | Callable from... |
|-------------------------|----------------|------------------|
| <code>__global__</code> | device | host |
| <code>__host__</code> | host | host |
| <code>__device__</code> | device | device |

```
__device__ void bar(...) { ... }

__global__ void foo(...) {
    bar(...);
}
```



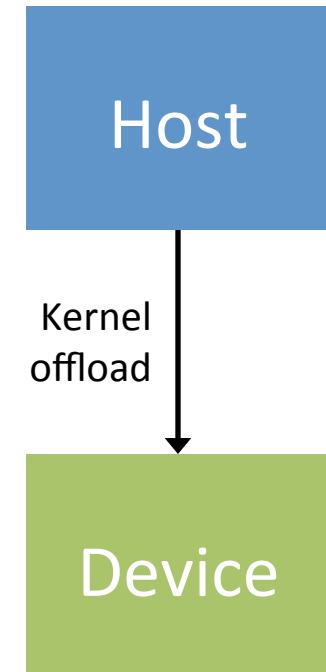
CUDA
Core

A Simple CUDA Example

Calling this kernel from the host requires special syntax to indicate that the function call targets the device:

```
__global__ void vector_add(...) {  
    ...  
}  
  
vector_add<<<1, 1>>>(C, A, B, N);
```

Run one thread on the GPU
(more on this later)



A Simple CUDA Example

One final step: the data.

Recall that host and device are physically discrete processors → separate address spaces, a pointer on one has no meaning on the other.

| Host | Device |
|---|--|
| int *A = (int *)malloc(N * sizeof(int)); int *B = (int *)malloc(N * sizeof(int)); int *C = (int *)malloc(N * sizeof(int)); vector_add<<<1, 1>>>(C, A, B, N); | __global__ void vector_add(int *C, int *A, int *B, int N) { int i; for (i = 0; i < N; i++) { C[i] = A[i] + B[i]; } } |

A Simple CUDA Example

Host

```
int *A = (int *)malloc(N * sizeof(int));  
int *B = (int *)malloc(N * sizeof(int));  
int *C = (int *)malloc(N * sizeof(int));  
  
vector_add<<<1, 1>>>(C, A, B, N);
```



Host

Device

```
__global__ void vector_add(int *C,  
                           int *A, int *B, int N) {  
    int i;  
    for (i = 0; i < N, i++) {  
        C[i] = A[i] + B[i];  
    }  
}
```

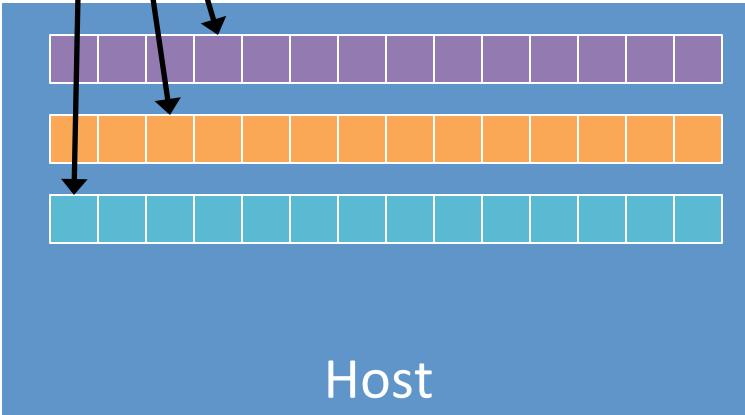
PCIe

Device

A Simple CUDA Example

Host

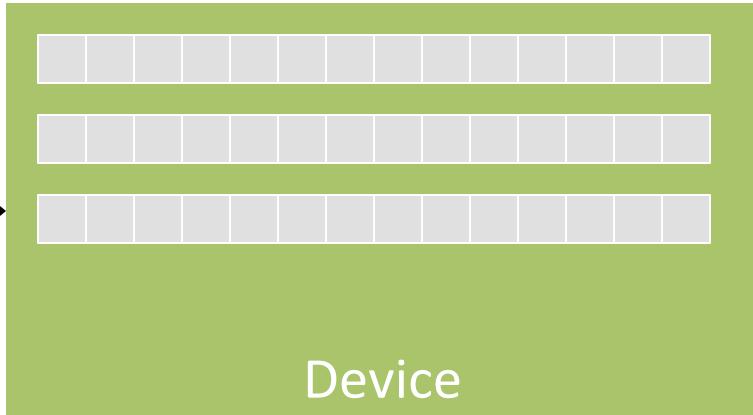
```
int *A = (int *)malloc(N * sizeof(int));  
int *B = (int *)malloc(N * sizeof(int));  
int *C = (int *)malloc(N * sizeof(int));  
  
vector_add<<<1, 1>>>(C, A, B, N);
```



Device

```
__global__ void vector_add(int *C,  
    int *A, int *B, int N) {  
    int i;  
    for (i = 0; i < N; i++) {  
        C[i] = A[i] + B[i];  
    }  
}
```

}



A Simple CUDA Example

```
Host                               Device  
int *A = (int *)malloc(N * sizeof(int)); __global__ void vector_add(int *C,  
int *B = (int *)malloc(N * sizeof(int));     int *A, int *B, int N) {  
int *C = (int *)malloc(N * sizeof(int));     int i;  
cudaError_t cudaMalloc(void **devPtr, size_t size);         C[i] = A[i] + B[i];  
vector_add<<<1, 1>>>(C, A, B, N); }  
}
```

Allocate size bytes of memory on GPU.

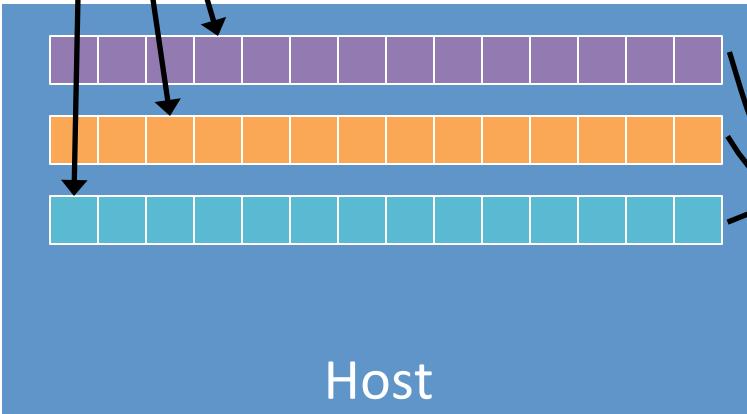
Store address at *devPtr.



A Simple CUDA Example

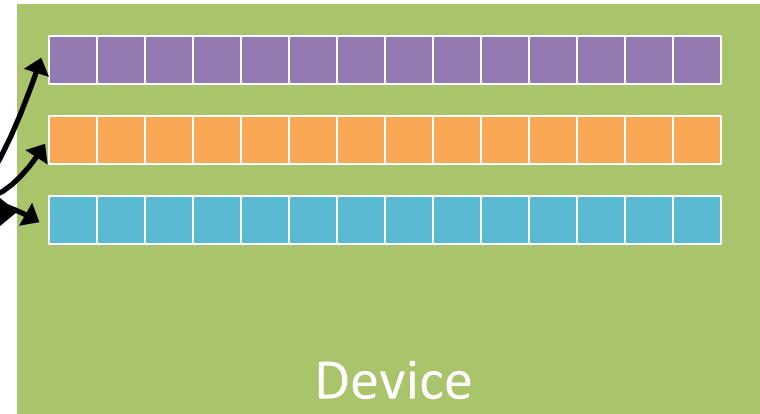
Host

```
int *A = (int *)malloc(N * sizeof(int));  
int *B = (int *)malloc(N * sizeof(int));  
int *C = (int *)malloc(N * sizeof(int));  
  
vector_add<<<1, 1>>>(C, A, B, N);
```



Device

```
_global_ void vector_add(int *C,  
                           int *A, int *B, int N) {  
    int i;  
    for (i = 0; i < N; i++) {  
        C[i] = A[i] + B[i];  
    }  
}
```



A Simple CUDA Example

```
Host                               Device  
int *A = (int *)malloc(N * sizeof(int)); global void vector_add(int *C,  
int *B = (int *)malloc(N * sizeof(int));           for (i = 0; i < N; i++) {  
int *C = (int *)malloc(N * sizeof(int));           C[i] = A[i] + B[i];  
  
vector_add<<<1, 1>>>(C, A, B, N);  
  
Transfer size bytes from src to dst.
```

Address spaces must obey direction, can be:

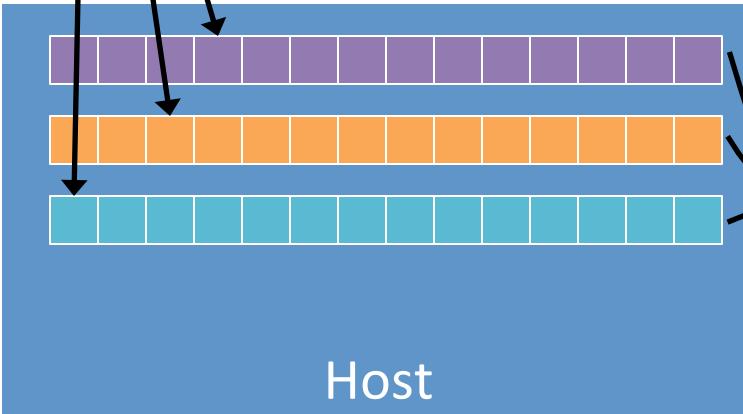
- cudaMemcpyHostToHost
- cudaMemcpyDeviceToDevice
- cudaMemcpyHostToDevice
- cudaMemcpyDeviceToHost



A Simple CUDA Example

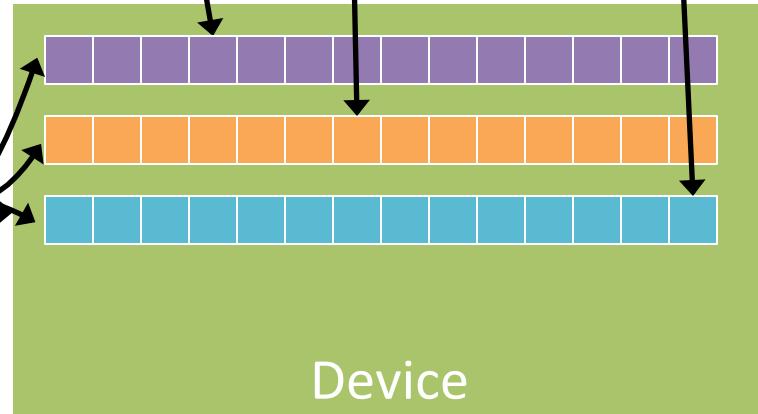
Host

```
int *A = (int *)malloc(N * sizeof(int));  
int *B = (int *)malloc(N * sizeof(int));  
int *C = (int *)malloc(N * sizeof(int));  
  
vector_add<<<1, 1>>>(C, A, B, N);
```



Device

```
__global__ void vector_add(int *C,  
                           int *A, int *B, int N) {  
    int i;  
    for (i = 0; i < N; i++) {  
        C[i] = A[i] + B[i];  
    }  
}
```



A Simple CUDA Example

```
int *A, *d_A;  
A = (int *)malloc(N * sizeof(int));  
cudaMalloc((void **) &d_A, N * sizeof(int));  
cudaMemcpy(A, d_A, N * sizeof(int), cudaMemcpyHostToDevice);  
  
...  
  
vector_add<<<1, 1>>>(C, A, B, N);
```

```
cudaError_t cudaMalloc(void **devPtr, size_t size);  
  
cudaError_t cudaMemcpy(void *dst, const void *src,  
size_t size, enum cudaMemcpyKind direction);  
  
cudaError_t cudaFree(void *addr);
```

Host

Device

A Simple CUDA Example

Building CUDA programs requires **nvcc**. Supports all of the usual flags you expect from a compiler.

```
$ cd ~/bootcamp-gpu/src/00_vecadd  
$ nvcc -arch=sm_20 vecadd.cu -o vecadd
```

```
↑          ↑          ↑          ↑  
NVIDIA's GPU  Compile and Input CUDA  Output  
compiler      optimize for the file    executable  
                  Fermi architecture
```

*Makefiles will be provided for your convenience

Hands On - 00_vecadd

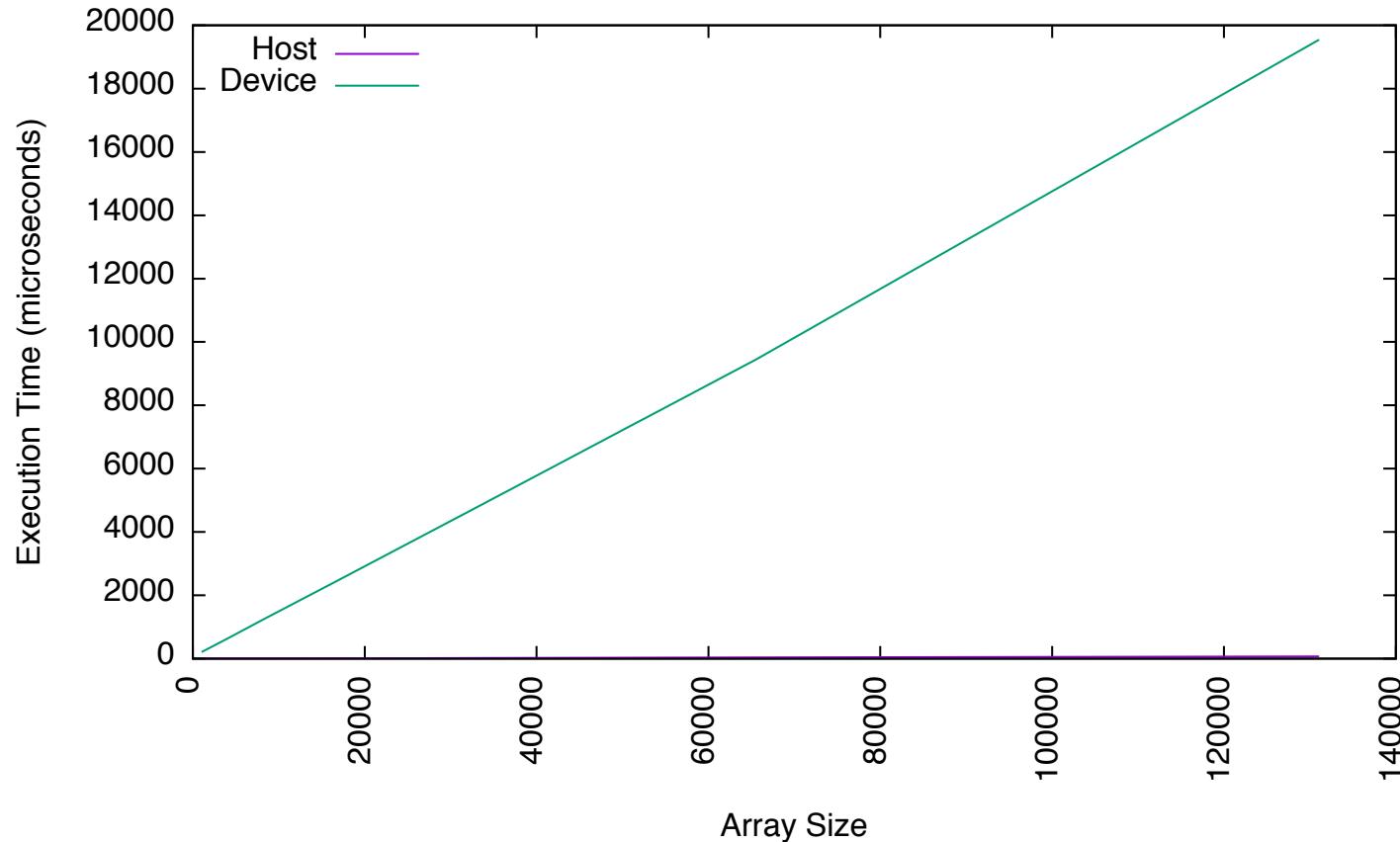
Don't forget to follow the instructions on slide 1.

```
$ cd ~/bootcamp-gpu/src/00_vecadd  
$ make  
$ vim/emacs vecadd.cu # Get familiar with the code  
$ ./vecadd 131072 # Vectors with 131,072 elements
```

What kind of performance do you observe?

Does it change relatively with different input sizes?

Hands On - 00_vecadd



Review - Basic CUDA

Running a function on the GPU requires 1) `__global__` annotation, and 2) kernel launch syntax (`<<<...>>>`).

CUDA forces you to think about discrete address spaces and manage coherency.

```
cudaError_t cudaMalloc(void **devPtr, size_t size);

cudaError_t cudaMemcpy(void *dst, const void *src,
                      size_t size, enum cudaMemcpyKind direction);

cudaError_t cudaFree(void *addr);
```

Copying and pasting code is not the path to performance in CUDA.

nvprof

nvprof is a command-line profiler for CUDA applications.

```
$ nvprof --help
Usage: nvprof [options] [CUDA-application] [application-arguments]
...
...
```

Can be used in a number of modes:

- Default/Summary Mode
- API Trace Mode
- Event/Metric Summary Mode
- Event/Metric Trace Mode

Hands On – nvprof Summary mode

A good starting point is usually Default/Summary mode.

```
$ nvprof ./vecadd 131072
```

Hands On – nvprof Summary mode

A good starting point is usually Default/Summary mode.

```
$ nvprof ./vecadd 131072
```

```
==17036== NVPROF is profiling process 17036, command: ./vecadd 131072
Finished! All 131072 elements validate.
```

```
Took 621 microseconds on the host
```

```
Took 19725 microseconds on the device, 0.03148x speedup
```

```
==17036== Profiling application: ./vecadd 131072
```

```
==17036== Profiling result:
```

| Time(%) | Time | Calls | Avg | Min | Max | Name |
|---------------|-----------------|----------|-----------------|-----------------|-----------------|--|
| 98.61% | 18.656ms | 1 | 18.656ms | 18.656ms | 18.656ms | vector_add(int*, int*, int*, int) |
| 0.95% | 180.03us | 2 | 90.015us | 89.951us | 90.079us | [CUDA memcpy HtoD] |
| 0.44% | 82.655us | 1 | 82.655us | 82.655us | 82.655us | [CUDA memcpy DtoH] |

nvprof

Diving deeper requires Event/Metric Summary Mode.

```
$ nvprof --query-events # List all enabled events  
$ nvprof --query-metrics # List all enabled metrics
```

Events are raw data collected by the hardware, e.g. # instructions (`inst_issued`)

Metrics are derived based on events, e.g. # instructions per cycle (`issued_ipc`)

Hands On – nvprof Metrics

Diving deeper requires Event/Metric Summary Mode.

1. ipc: # of instructions executed per cycle on the GPU
2. sm_efficiency: % of time an SM on the GPU is doing useful work
3. alu_fu_utilization: indicates how well the CUDA cores are being used

```
$ nvprof --metrics ipc,sm_efficiency,alu_fu_utilization ./vecadd 131072
```

Hands On – nvprof Metrics

Diving deeper requires Event/Metric Summary Mode.

```
$ nvprof --metrics ipc,sm_efficiency,alu_fu_utilization ./vecadd 131072
```

```
==30413== Warning: Some kernel(s) will be replayed on device 0 in order to collect all events/metrics.  
Finished! All 131072 elements validate.
```

```
Took 621 microseconds on the host
```

```
Took 174666 microseconds on the device, 0.00356x speedup
```

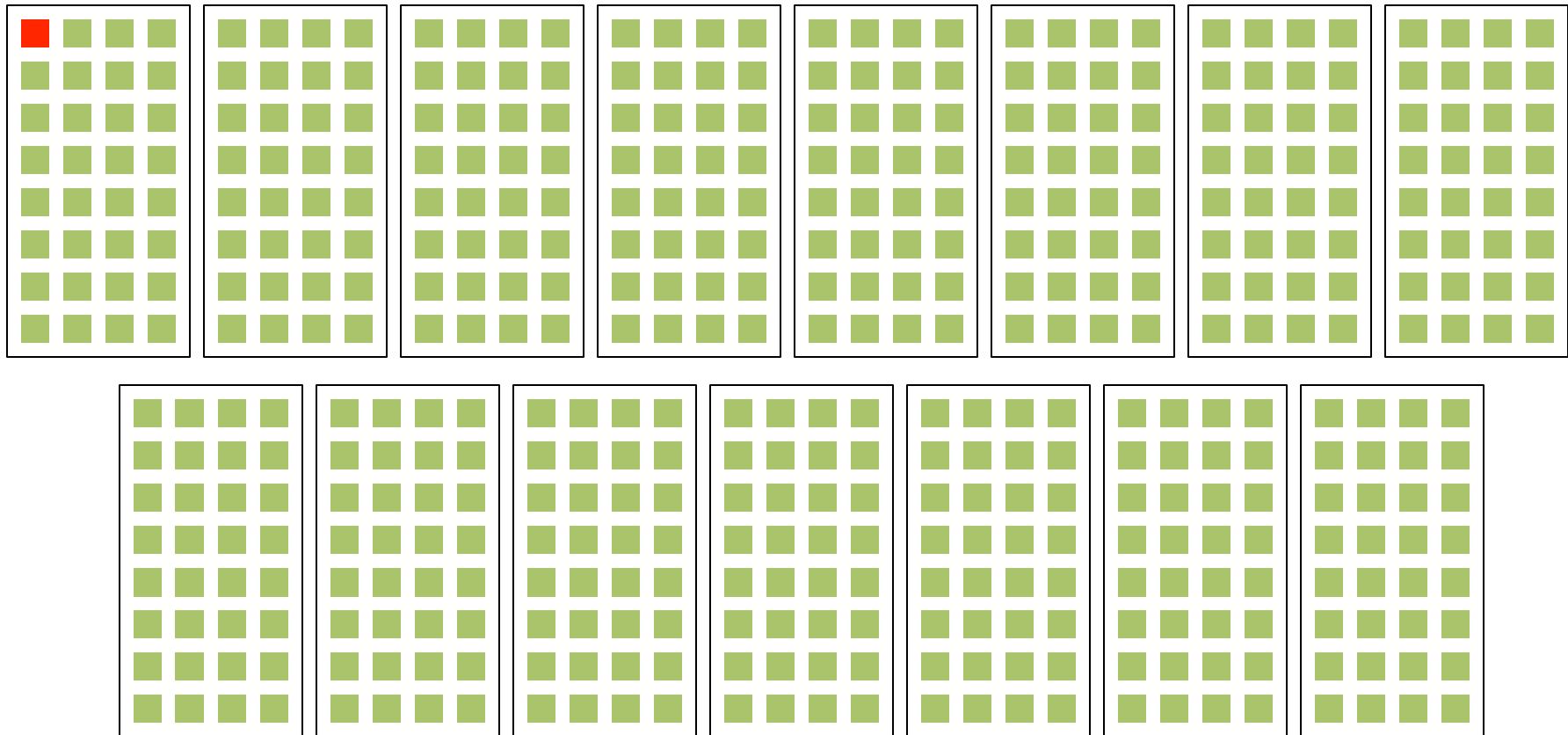
```
==30413== Profiling application: ./vecadd 131072
```

```
==30413== Profiling result:
```

```
==30413== Metric result:
```

| Invocations | Metric Name | Metric Description | Min ... |
|---|--------------------|--------------------------------------|--------------|
| Device "Tesla M2050 (0)" | | | |
| Kernel: vector_add(int*, int*, int*, int) | | | |
| 1 | sm_efficiency | Multiprocessor Activity | 7.14% ... |
| 1 | ipc | Executed IPC | 0.103982 ... |
| 1 | alu_fu_utilization | Arithmetic Function Unit Utilization | Low (1) ... |

nvprof



CUDA Execution Model

Or How to Write Parallel CUDA Programs

Flynn's Taxonomy

| | Single Instruction | Multiple Instructions |
|---------------|--------------------|-----------------------|
| Single Data | SISD | MISD |
| Multiple Data | SIMD | MIMD |

Single Instruction, Single Data stream (SISD)

A sequential computer which exploits no parallelism in either the instruction or data streams. e.g., old single processor PC

Single Instruction, Multiple Data streams (SIMD)

A computer which exploits multiple data streams against a single instruction stream to perform operations which may be naturally parallelized. e.g. graphics processing unit

Multiple Instruction, Single Data stream (MISD)

Multiple instructions operate on a single data stream. Uncommon architecture which is generally used for fault tolerance. Heterogeneous systems operate on the same data stream and must agree on the result. e.g. the Space Shuttle flight control computer.

Multiple Instruction, Multiple Data streams (MIMD)

Multiple autonomous processors simultaneously executing different instructions on different data. e.g. a PC cluster memory space.

CUDA Execution Model

Recall:

- Programmer writes single-threaded kernel
- **1 CUDA thread per CUDA core.**



CUDA
Core

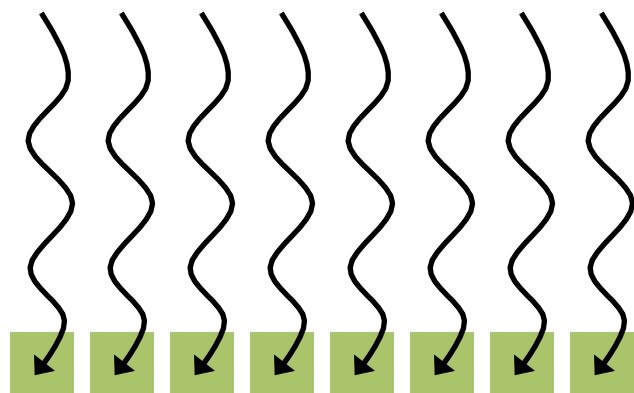
CUDA Execution Model

Recall:

- Programmer writes single-threaded kernel
- **1 CUDA thread per CUDA core.**

Threads are scheduled in groups of 32, called a **warp**.

- Warps execute in lock step (same instruction at same cycle)
- All threads in warp on same SM



CUDA Execution Model

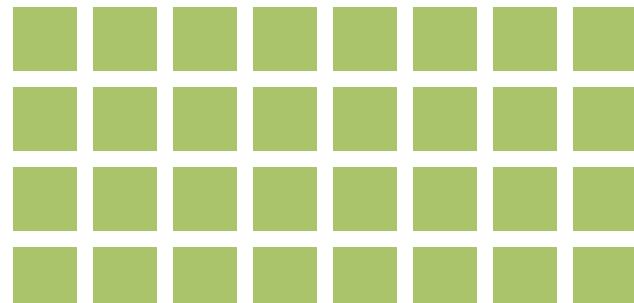
Recall:

- Programmer writes single-threaded kernel
- **1 CUDA thread** per CUDA core.

Threads are scheduled in groups of 32, called a **warp**.

Warps are grouped into **thread blocks**.

- Warps in a block execute on same SM
- Warp execution in same thread block can be interleaved with each other on SM
- SM resources are partitioned among all threads in resident thread blocks



e.g. 4 warps in one thread block

CUDA Execution Model

Recall:

- Programmer writes single-threaded kernel
- **1 CUDA thread** per CUDA core.

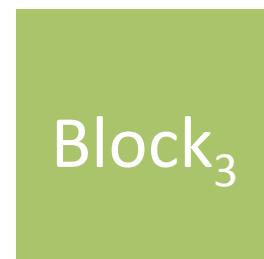
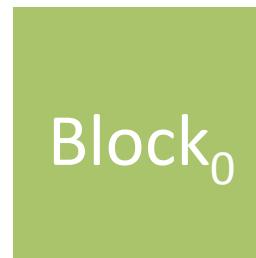
Threads are scheduled in groups of 32, called a **warp**.

Warp s are grouped into **thread blocks**.

A single kernel **grid** can include many thread blocks.

- Kernel call is not complete until all thread blocks within it complete.

`vector_add<<<...>>>(...)`



CUDA Execution Model

Blocks and grids are configurable by kernel launch arguments:

```
vector_add<<<nblocks, nthreads_per_block>>>(...)
```

```
vector_add<<<1, 1>>>(...); // 1 block with 1 thread inside it
```

```
vector_add<<<1, 32>>>(...); // 1 block with 1 full warp
```

```
vector_add<<<1, 256>>>(...); // 1 block with 256 threads (8 warps)
```

```
vector_add<<<N, 256>>>(...); // N blocks with 256 threads/block
```

CUDA Execution Model

Blocks and grids can also be multi-dimensional:

```
// Equivalent to vector_add<<<1, 256>>>(...);  
dim3 1d_block(256, 1, 1);  
vector_add<<<1, 1d_block>>>(...);  
  
dim3 block(16, 16); // 2D 16x16 block of threads, 256 total  
  
dim3 grid(4, 4); // 2D 4x4 grid of thread blocks, 16 total  
  
kernel<<<grid, block>>>(...);
```

CUDA Execution Model

Special CUDA kernel variables can be used to check the coordinates of a thread in an executing grid.

| Variable Name | Description |
|---|--|
| <code>threadIdx.x</code> , <code>threadIdx.y</code> , ... | Offset of a thread within a thread block |
| <code>blockIdx.x</code> , <code>blockIdx.y</code> , ... | Offset of a block within a grid |
| <code>blockDim.x</code> , ... | # of threads per block |
| <code>gridDim.x</code> , ... | # of blocks in grid |
| <code>blockIdx.x * blockDim.x + threadIdx.x</code> | One way to calculate globally unique thread ID for 1D blocks and grids |

CUDA Execution Model

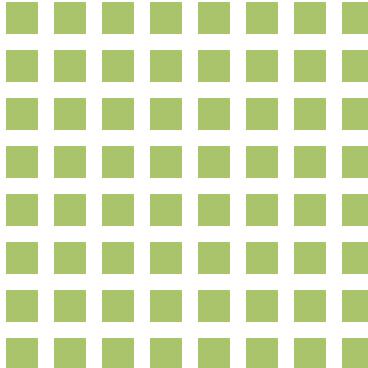
```
__global__ void matrix_add(...) {
    const int row = blockIdx.x * blockDim.x + threadIdx.x;
    const int col = blockIdx.y * blockDim.y + threadIdx.y;
    ...
}

dim3 block(8, 8); // 2D 8x8 block of threads, 64 total
dim3 grid(2, 2); // 2D 2x2 grid of thread blocks, 4 total

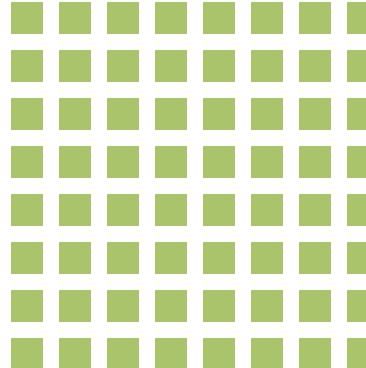
vector_add<<<grid, block>>>(...);
```

CUDA Execution Model

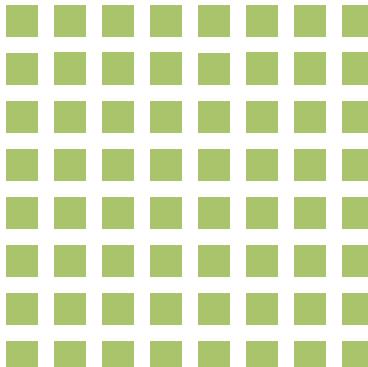
Block (0, 0)



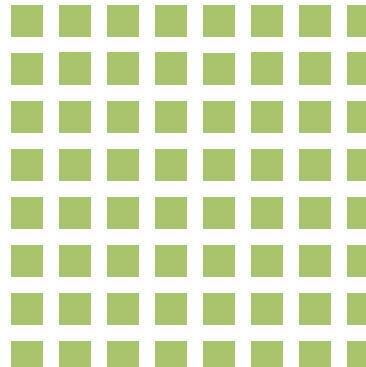
Block (0, 1)



Block (1, 0)



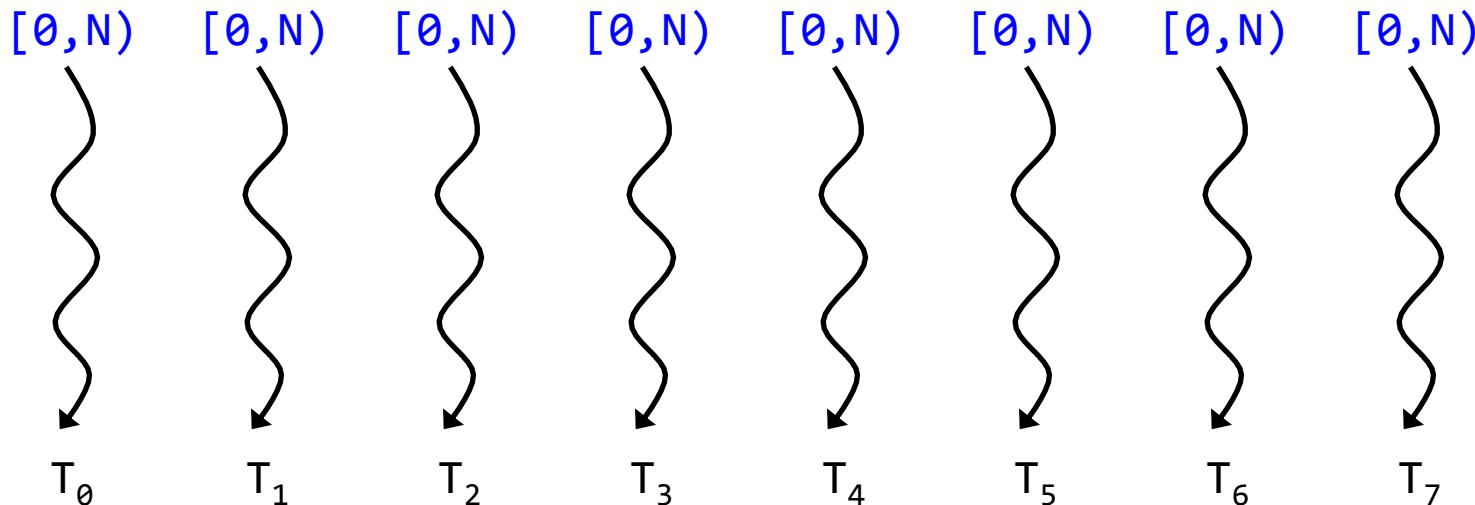
Block (1, 1)



CUDA Execution Model

Launch configuration alone is insufficient for our `vector_add` to run in parallel.

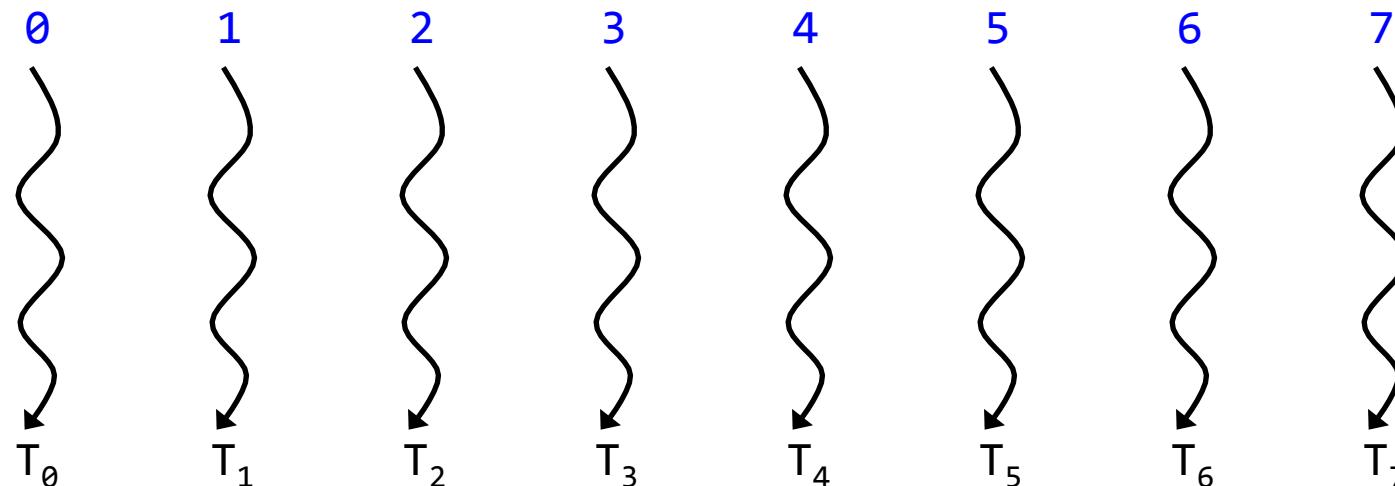
```
vector_add<<<N / 256, 256>>>(C, A, B, N);
```



CUDA Execution Model

Must also change the work performed by the kernel.

```
__global__ void vector_add(int *C, int *B, int *A, int N) {  
    const int i = blockIdx.x * blockDim.x + threadIdx.x;  
    C[i] = A[i] + B[i];  
}
```



Hands On – Adding CUDA Parallelism

Using the previous slides as a guide, add parallelism to the vecadd example:

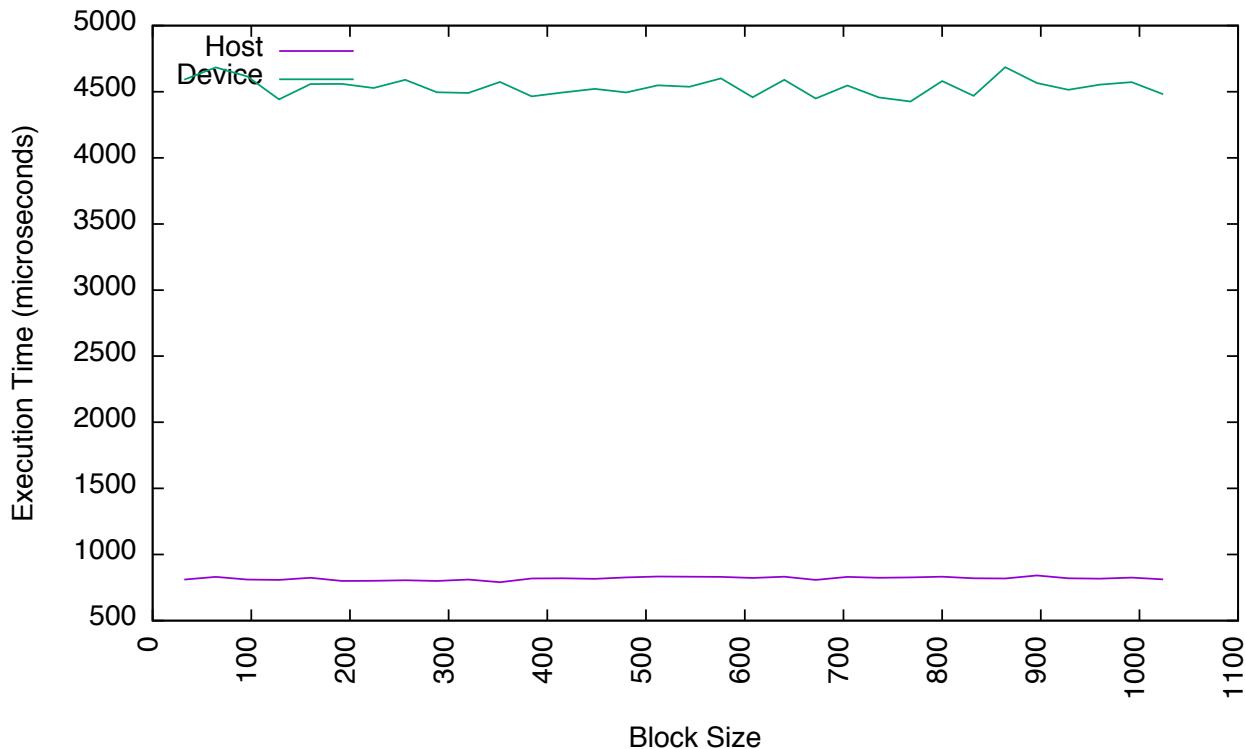
1. Change the <<<...>>> launch configuration to use B threads per block, N/B blocks
 - Be careful that the N you run is evenly divisible by B
 - CUDA supports up to 1,024 threads per block
2. Change the kernel to process a single data point

Does performance improve?

```
$ ./vecadd 1048576
```

Does block size affect performance?

Hands On – Adding CUDA Parallelism



More similar host and device performance (was ~150,000 us).

No clear trends in performance as a function of block size (yet).

Still not great performance, time to go back to nvprof.

Hands On – Adding CUDA Parallelism

\$HOME/bootcamp-gpu/src/01_vecadd contains an example solution that allows configuring threads per block from the command line:

```
$ cd $HOME/bootcamp-gpu/src/01_vecadd  
$ ./vecadd <N> <threads-per-block>
```

Try:

```
$ nvprof --metrics ipc,sm_efficiency,alu_fu_utilization \  
./vecadd 1048576 256  
  
$ nvprof ./vecadd 1048576 256
```

Hands On – Adding CUDA Parallelism

```
$ nvprof --metrics ipc,sm_efficiency,alu_fu_utilization ./vecadd 1048576
```

```
==32216== NVPROF is profiling process 32216, command: ./vecadd 1048576
==32216== Warning: Some kernel(s) will be replayed on device 0 in order to collect all events/metrics.
Finished! All 1048576 elements validate using 256 threads per block.
```

```
Took 5032 microseconds on the host
```

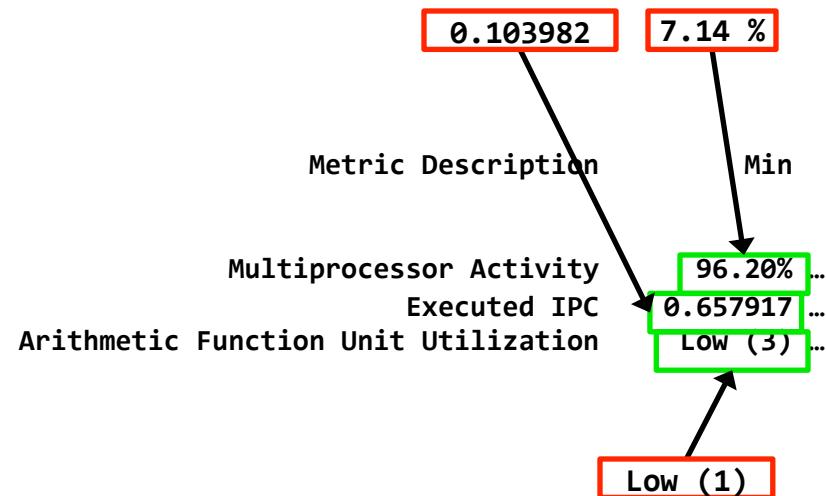
```
Took 21110 microseconds on the device, 0.23837x speedup
```

```
==32216== Profiling application: ./vecadd 1048576
```

```
==32216== Profiling result:
```

```
==32216== Metric result:
```

| Invocations | Metric Name |
|---|--------------------|
| Device "Tesla M2050 (0)" | |
| Kernel: vector_add(int*, int*, int*, int) | |
| 1 | sm_efficiency |
| 1 | ipc |
| 1 | alu_fu_utilization |



Hands On – Adding CUDA Parallelism

```
$ nvprof ./vecadd 1048576 256
```

```
[jmg3@gpu-014 01_vecadd]$ /opt/apps/software/Compiler/GCC/4.4.7/CUDA/6.5.14/bin/nvprof ./vecadd  
1048576 32  
==32199== NVPROF is profiling process 32199, command: ./vecadd 1048576 32  
Finished! All 1048576 elements validate using 32 threads per block.  
Took 5049 microseconds on the host  
Took 4541 microseconds on the device, 1.11187x speedup  
==32199== Profiling application: ./vecadd 1048576 32  
==32199== Profiling result:  


| Time(%) | Time     | Calls | Avg      | Min      | Max      | Name                              |
|---------|----------|-------|----------|----------|----------|-----------------------------------|
| 44.85%  | 1.5758ms | 1     | 1.5758ms | 1.5758ms | 1.5758ms | [CUDA memcpy DtoH]                |
| 43.93%  | 1.5436ms | 2     | 771.81us | 767.70us | 775.93us | [CUDA memcpy HtoD]                |
| 10.26%  | 360.38us | 1     | 360.38us | 360.38us | 360.38us | vector_add(int*, int*, int*, int) |
| 0.96%   | 33.856us | 1     | 33.856us | 33.856us | 33.856us | [CUDA memset]                     |


```

Review – CUDA Execution Model

nvcc – Compile CUDA applications, shares many flags with x86 compilers, use –arch to target specific GPU generations.

nvprof – Tool for analyzing CUDA applications, profile-driven optimization!

CUDA thread hierarchy: **thread**, **warp**, **thread block**, and **grid**

Access thread coordinates from running kernel via **blockIdx.x**, **threadIdx.x**, **blockDim.x**

Good to maximize parallelism of CUDA kernels over input, PCIe transfers can quickly become bottlenecks

CUDA Asynchronous Execution

CUDA Asynchronous Execution

Most APIs are synchronous by default.

- When you call `malloc`, you expect its return value to point to immediately addressable memory.

Commonly, APIs for high latency operations support some level of asynchrony:

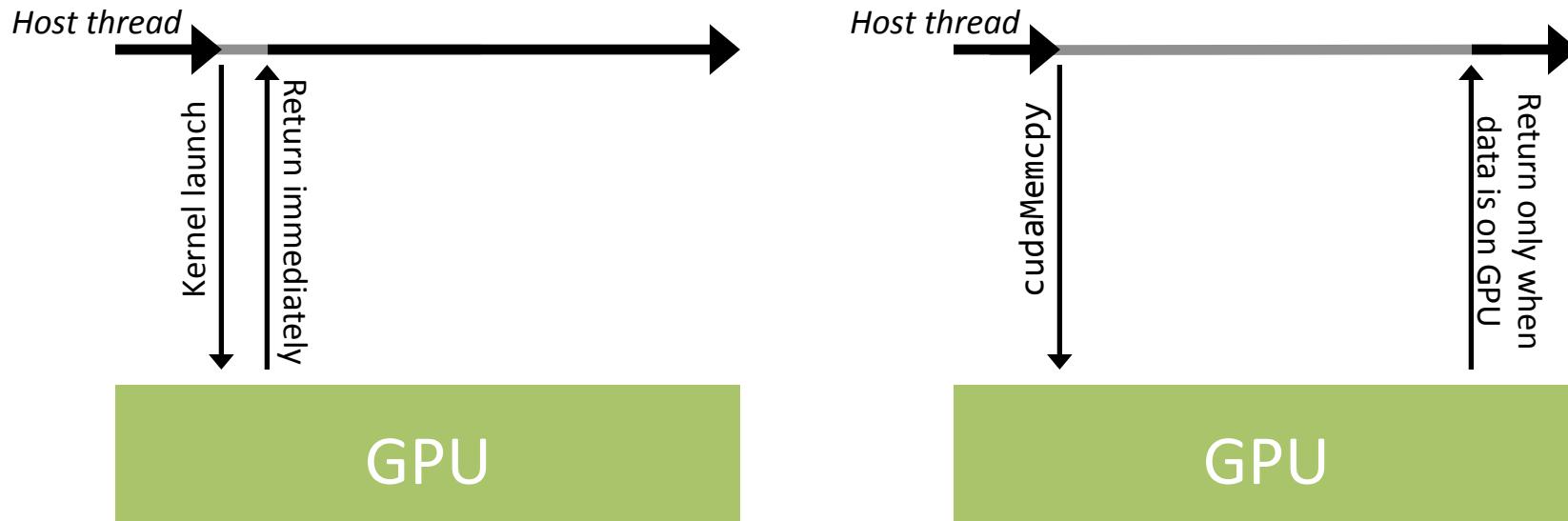
- POSIX AIO
- NodeJS/Javascript callbacks
- E-mail

Overlapping useful work on-core with asynchronous work off-core can benefit performance.

CUDA Asynchronous Execution

CUDA supports both blocking and asynchronous APIs.

A kernel launch <<<...>>> is asynchronous. A cudaMemcpy is blocking.



CUDA Asynchronous Execution

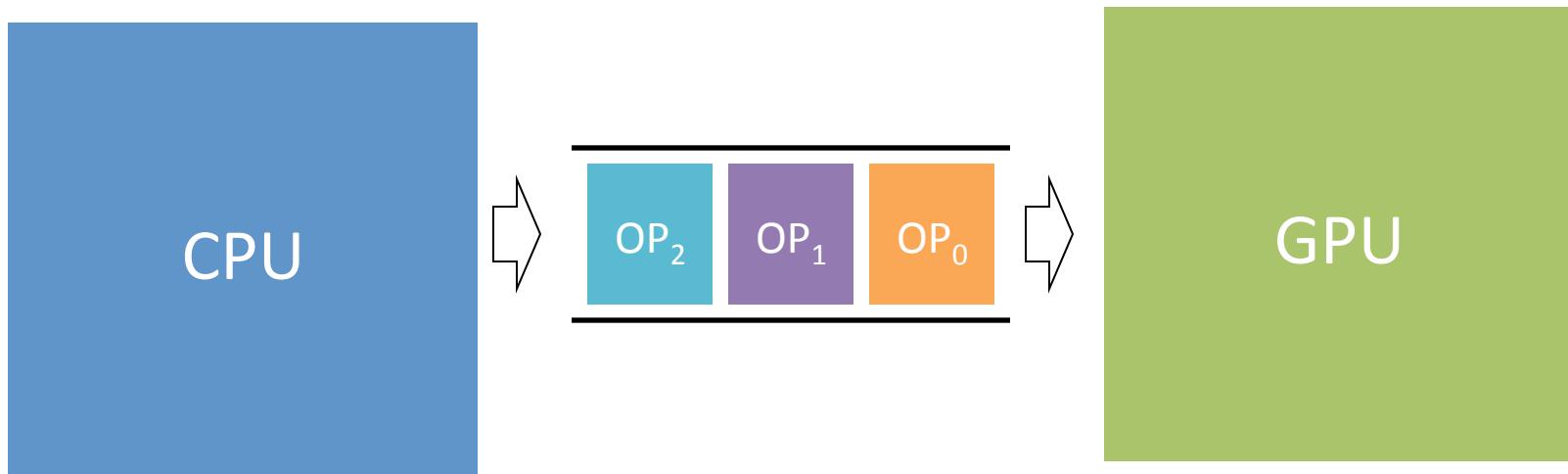
Some blocking APIs have asynchronous equivalents. Some don't.

| Blocking API | Asynchronous API |
|--------------|---------------------------|
| cudaMalloc | |
| cudaMemcpy | cudaMemcpyAsync |
| | Kernel launch (<<<...>>>) |
| cudaFree | |

CUDA Asynchronous Execution

The core of CUDA asynchrony is **CUDA streams**.

A CUDA stream defines a sequence of CUDA operations (e.g. kernel launch, memory copy, etc.) to be performed in order.



CUDA Asynchronous Execution

Streams must be explicitly created by the programmer:

```
cudaStream_t stream;  
cudaError_t cudaStreamCreate(cudaStream_t *stream);  
cudaError_t cudaStreamDestroy(cudaStream_t *stream);
```

CUDA Asynchronous Execution

Every asynchronous operation is associated with a CUDA stream.

If the CUDA stream is not explicitly set, a **default stream** is used.

Implicit default stream

```
kernel<<<...>>>(...);  
cudaMemcpyAsync(...);
```

Explicit stream

```
kernel<<<..., stream>>>(...);  
cudaMemcpyAsync(..., stream);
```

CUDA Asynchronous Execution

Stream status can be queried or blocked on.

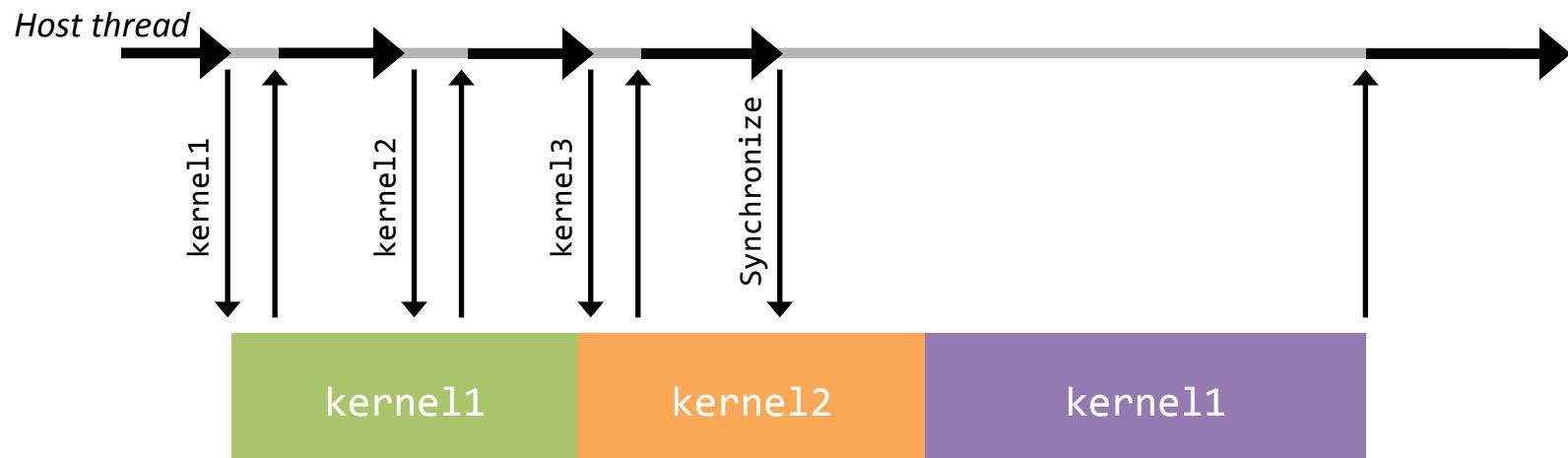
Implicitly includes all past operations in that stream.

```
cudaStream_t stream;  
cudaStreamCreate(&stream);  
kernel<<<..., stream>>>(...);  
cudaStreamSynchronize(stream);
```

```
cudaError_t cudaStreamSynchronize(cudaStream_t stream);  
cudaError_t cudaStreamQuery(cudaStream_t stream);
```

CUDA Asynchronous Execution

```
kernel1<<<..., stream>>>(...);  
kernel2<<<..., stream>>>(...);  
kernel3<<<..., stream>>>(...);  
cudaStreamSynchronize(stream);
```

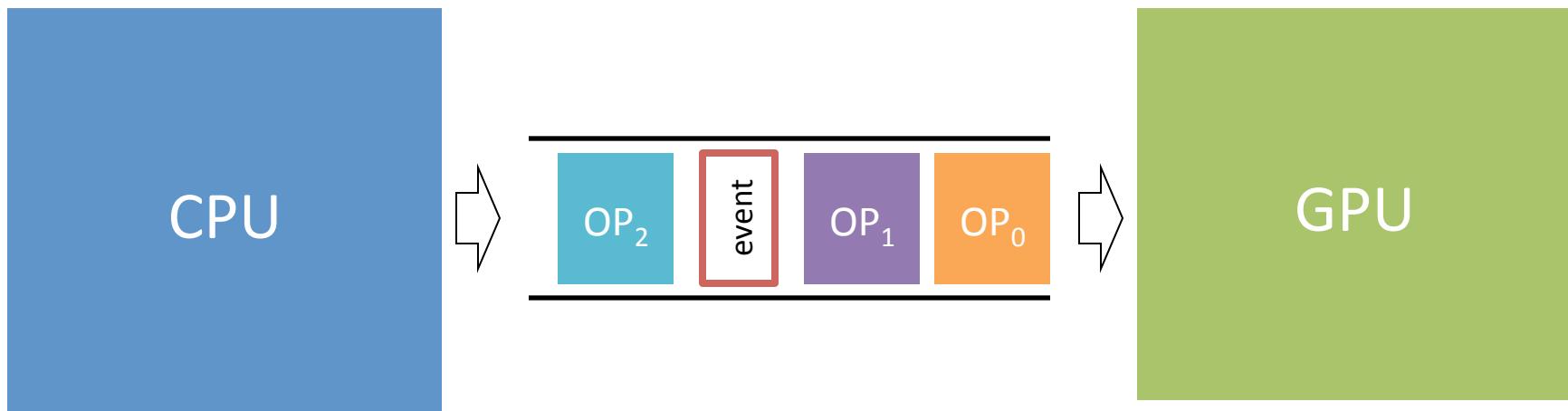


CUDA Asynchronous Execution

Also possible to query specific points in time within a stream using **CUDA Events**.

Events are inserted in to streams and satisfied as they exit the stream.

Enable querying of all operations inserted in to the stream prior to the event.



CUDA Asynchronous Execution

Like streams, events can be created, destroyed, queried, and blocked on.

```
cudaEvent_t event;
cudaError_t cudaEventCreate(cudaEvent_t *event,
                           unsigned int flags);
cudaError_t cudaEventDestroy(cudaEvent_t event);
cudaError_t cudaEventQuery(cudaEvent_t event);
cudaError_t cudaEventSynchronize(cudaEvent_t event);
```

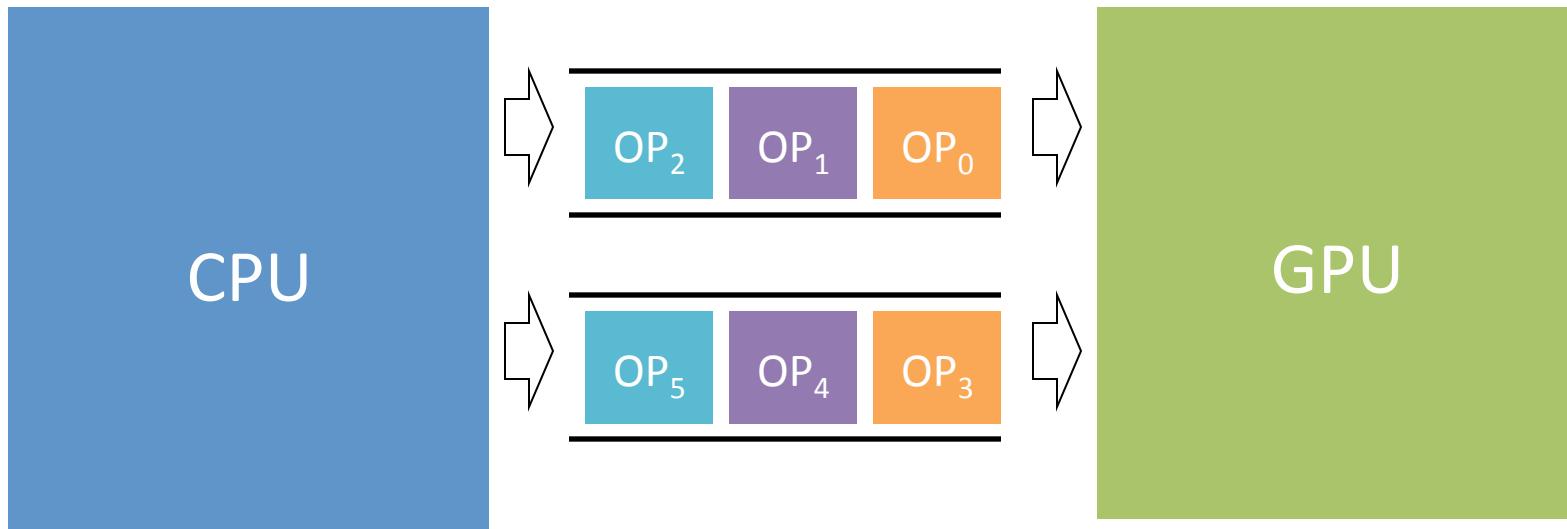
Placing an event in a stream:

```
cudaError_t cudaEventRecord(cudaEvent_t event,
                           cudaStream_t stream);
```

CUDA Asynchronous Execution

Synchronizing on all operations issued to a given device is possible with `cudaDeviceSynchronize`.

```
cudaError_t cudaDeviceSynchronize();
```

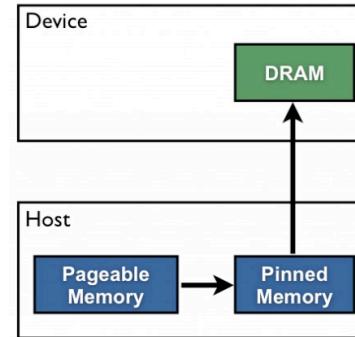


Special Note on cudaMemcpyAsync

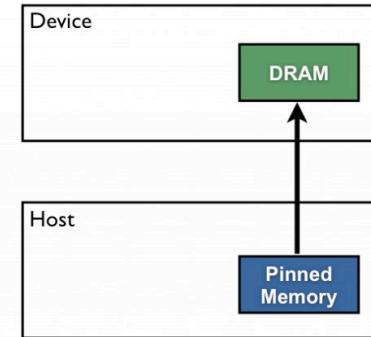
Host allocations transferred through cudaMemcpyAsync must be page-locked/pinned.

CUDA provides cudaMallocHost for page-locked host allocations.

Pageable Data Transfer



Pinned Data Transfer



```
cudaError_t cudaMallocHost(void **ptr, size_t size);
```

CUDA Asynchronous Execution

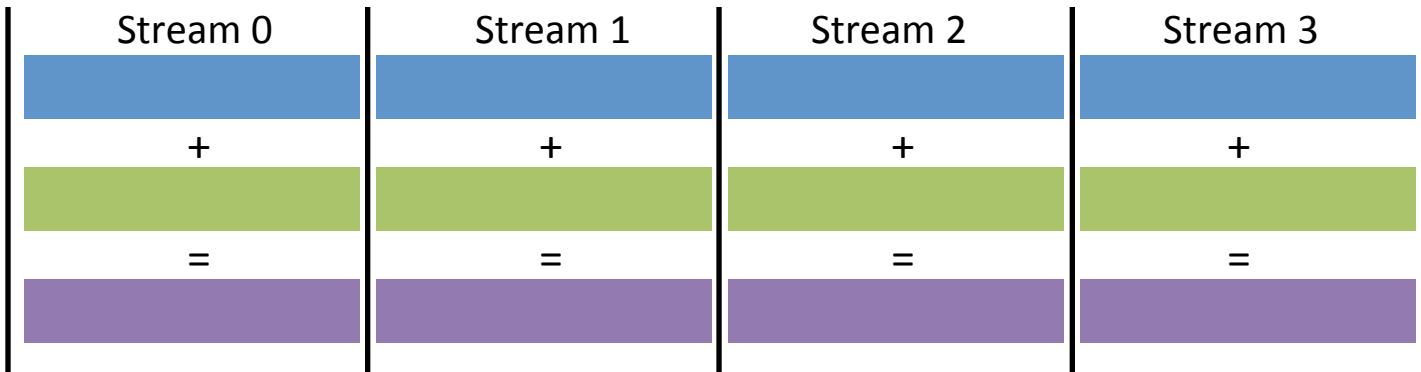
\$HOME/bootcamp-gpu/src/02_vecadd_streams/vecadd.cu contains an example of chunked vector addition.



```
$ ./vecadd <N> <nchunks>
```

Hands On – CUDA Streams

Complete the template in `$HOME/bootcamp-gpu/src/02_vecadd_streams/vecadd.cu` based on the TODOs contained to use a different stream on each vector chunk.



```
$ ./vecadd 4194304 4 # Vector of 4,194,304 elements, 4 chunks
```

Do multiple streams significantly improve performance for this example?

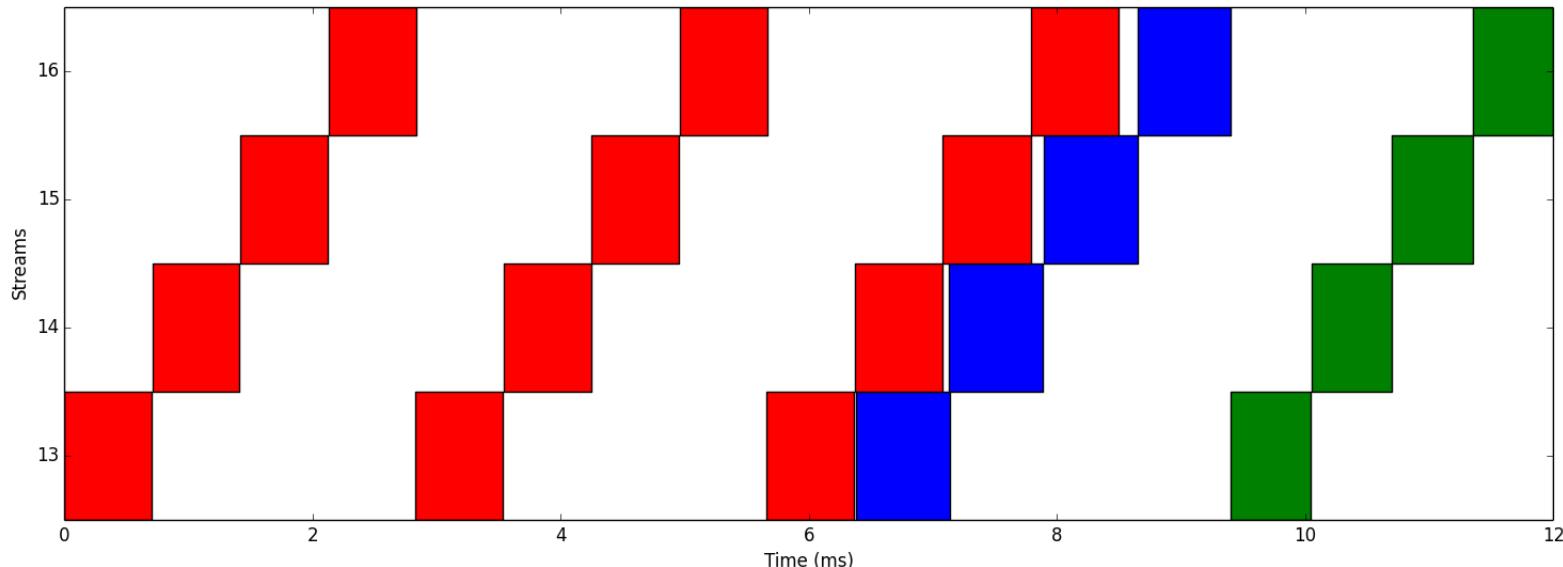
Hands On – CUDA Streams

Try using nvprof's GPU tracing mode to understand stream interleaving.

```
$ nvprof --print-gpu-trace ./vecadd_solution 4194304 1
```

```
$ nvprof --print-gpu-trace ./vecadd_solution 4194304 4
```

Hands On – CUDA Streams



Review – CUDA Asynchrony

A **CUDA stream** defines a sequence of CUDA operations (e.g. kernel launch, memory copy, etc.) to be performed in order.

CUDA events are inserted in to streams and satisfied as they exit the stream, mark a specific point-in-time in the processing of the operations in that stream.

Using asynchronous APIs requires an understanding of streams (and events), can benefit performance by enabling host-device, host-copy, device-copy overlaps.

```
cudaStreamCreate, cudaStreamDestroy, cudaStreamSynchronize,  
cudaStreamQuery, cudaEventCreate, cudaEventDestroy,  
cudaEventQuery, cudaEventSynchronize, cudaEventRecord,  
cudaDeviceSynchronize, cudaMallocHost
```

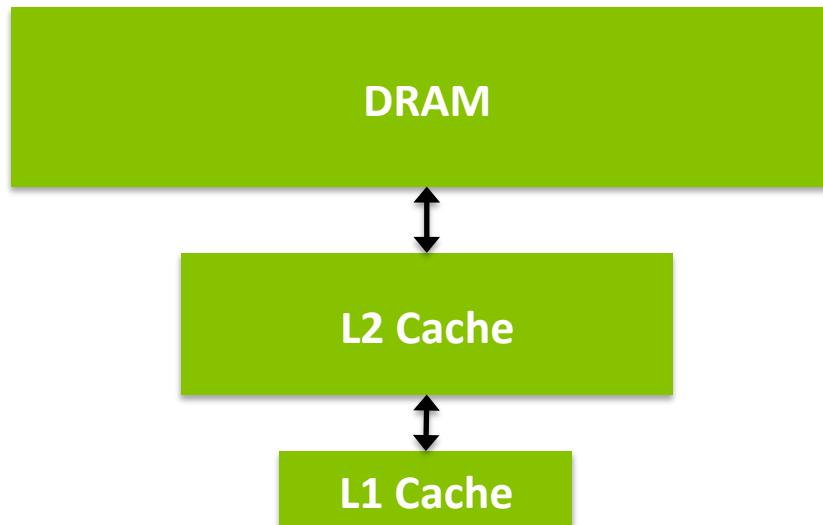
CUDA Memory Hierarchy

Memory Hierarchies

Up until now, we've glossed over efficient data access from CUDA kernels.

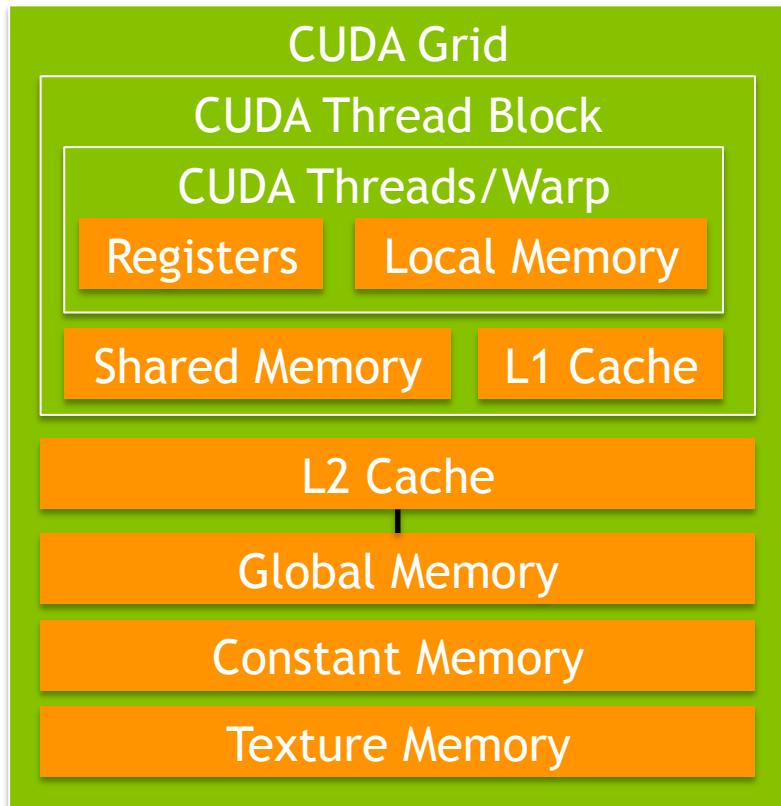
On CPU, this generally involves optimizing for cache line locality.

- A memory hierarchy emulates a large amount of low-latency memory



CUDA Memory Hierarchy

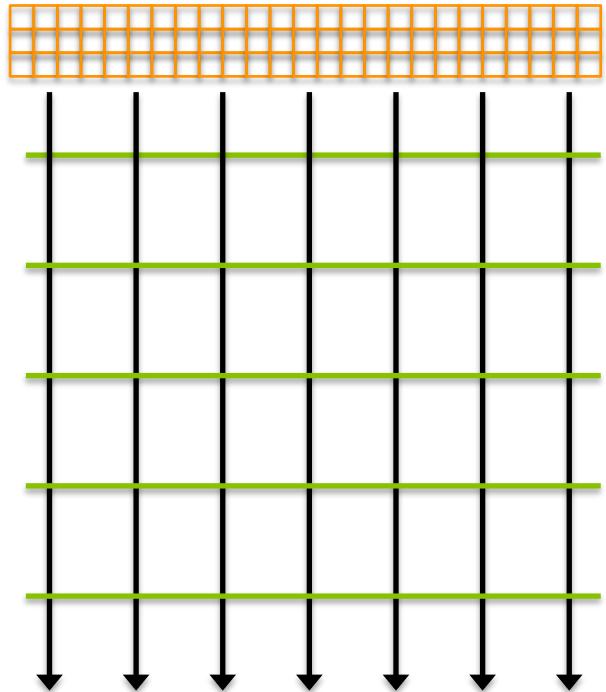
- The CUDA Memory Hierarchy is more complex than the CPU's
 - Many different types of memory, each with special-purpose characteristics
 - More explicit control over data movement



CUDA Memory Hierarchy

- **Registers**

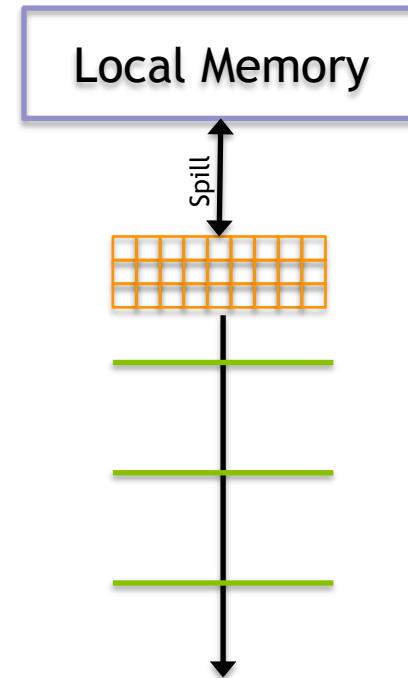
- Lowest latency memory space on the GPU
- Private to each CUDA thread
- Constant pool of registers per-SM divided among threads in resident thread blocks
- Architecture-dependent limit on number of registers per thread
- Registers are not explicitly used by the programmer, implicitly allocated by the compiler
- `nvcc -maxrregcount` allows you to limit # registers per thread



CUDA Memory Hierarchy

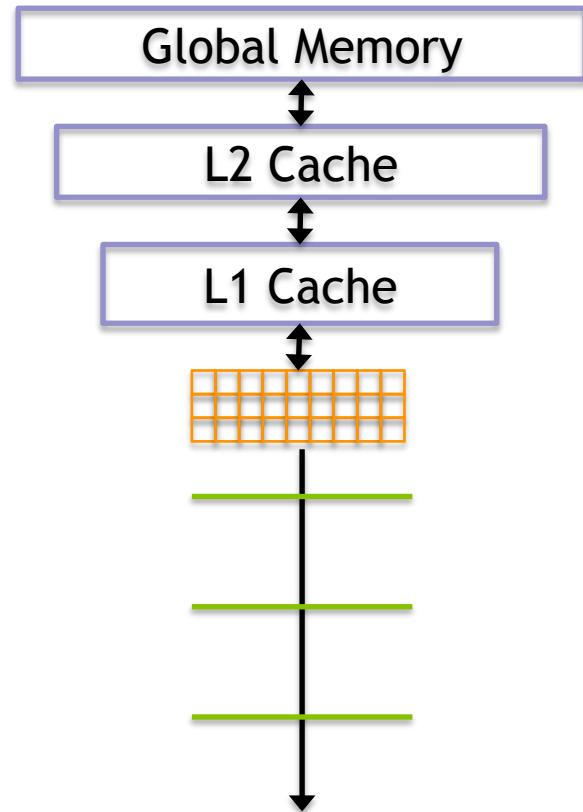
- **Local Memory**

- When registers are exhausted, variables spill to local memory
- Variables likely to be placed in local memory: large local structures or arrays, local arrays whose indices cannot be determined at compile-time
- Local memory is not physical, variables stored in local memory are spilled to Global Memory, L1 cache, or L2 cache
- Not explicitly controlled by programmer



CUDA Memory Hierarchy

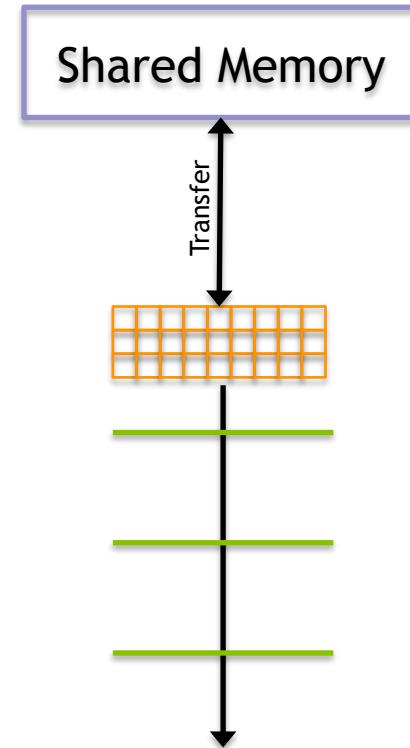
- **GPU Caches**
 - Behavior of GPU caches is architecture-dependent
 - Per-SM L1 cache stored on-chip
 - Per-GPU L2 cache stored off-chip, caches values for all SMs
 - Due to parallelism of accesses, GPU caches can be difficult to reason about



CUDA Memory Hierarchy

- **Shared Memory**

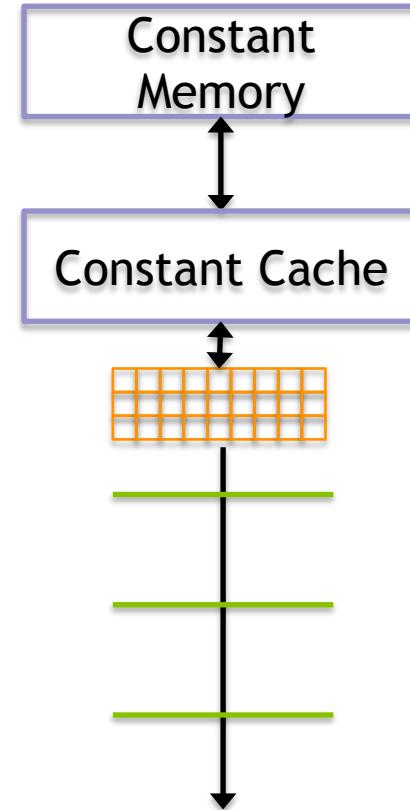
- Declared with the `__shared__` keyword
- Low-latency, high bandwidth
- Shared by all threads in a thread block
- Explicitly allocated and managed by the programmer, manual L1 cache
- Stored on-SM, same physical memory as the GPU L1 cache
- On-SM memory is statically partitioned between L1 cache and shared memory



CUDA Memory Hierarchy

- **Constant Memory**

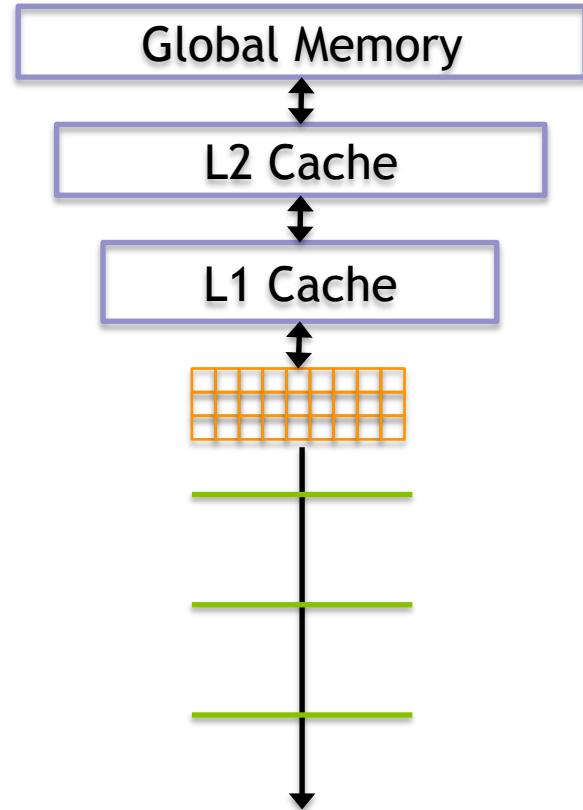
- Declared with the `__constant__` keyword
- Read-only
- Limited in size: 64KB
- Stored in device memory (same physical location as Global Memory)
- Cached in a per-SM constant cache
- Optimized for all threads in a warp accessing the same memory cell



CUDA Memory Hierarchy

- **Global Memory**

- Large, high-latency memory
- Stored in device memory (along with constant and texture memory)
- Can be declared statically with `__device__`
- Can be allocated dynamically with `cudaMalloc`
- Explicitly managed by the programmer
- Optimized for all threads in a warp accessing neighbouring memory cells



Review - CUDA Memory Spaces

| MEMORY | ON/OFF CHIP | CACHED | ACCESS | SCOPE | LIFETIME |
|----------|-------------|--------|--------|----------------------|-----------------|
| Register | On | n/a | R/W | 1 thread | Thread |
| Local | Off | † | R/W | 1 thread | Thread |
| Shared | On | n/a | R/W | All threads in block | Block |
| Global | Off | † | R/W | All threads + host | Host allocation |
| Constant | Off | Yes | R | All threads + host | Host allocation |
| Texture | Off | Yes | R | All threads + host | Host allocation |

CUDA Kernel Synchronization

Kernel Synchronization

Synchronization in CUDA is closely tied with the CUDA memory hierarchy.

Relative to CPU programming models, synchronization options are constrained in CUDA.

- CUDA is designed to be highly scalable and portable across generations
- Lots of synchronization negates both of these principles

There is no global synchronization in CUDA (i.e. no `#pragma omp barrier`)

Instead, CUDA offers local synchronization and some global memory fences.

- For performance reasons, use of these should still be minimized

Kernel Synchronization

| Function | Action | Scope | Notes |
|-------------------------------------|--------------|--------------|---|
| <code>__syncthreads()</code> | Barrier | Thread block | Most commonly useful function. |
| <code>__threadfence_block()</code> | Memory fence | Thread block | Applies to shared and global mem. |
| <code>__threadfence()</code> | Memory fence | Device | Same as <code>__threadfence_block()</code> , but also a global memory write fence relative to all threads on the device. |
| <code>__threadfence_system()</code> | Memory fence | Whole system | Same as <code>__threadfence_block()</code> , but also a global memory write fence relative to all threads on the device and host and read fence for calling thread. |

Review - Kernel Synchronization

Good chance you'll only ever use `__syncthreads()`

Avoid synchronization at all costs.

Global synchronization by kernel termination.

Optimizing CUDA Memory Accesses

CPUs rely heavily on hardware-managed caches for performance.

GPU caching is a more challenging problem

- Thousands of threads cooperating on a problem
- Difficult to predict the next round of accesses for all threads

For efficient global memory access, GPUs instead rely on:

1. Large device memory bandwidth
2. Aligned and coalesced memory access patterns
3. Maintaining sufficient pending I/O operations to keep the memory bus saturated and hide global memory latency

Optimizing CUDA Memory Accesses

Aligned and coalesced global memory accesses are key to optimizing an application's use of global memory bandwidth.

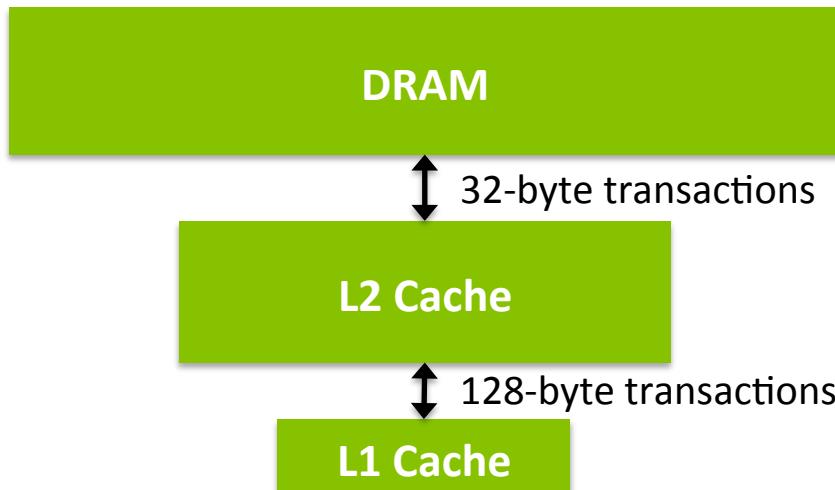
1. **Coalesced:** threads in a warp reference memory addresses that can all be serviced by a single global memory transaction (think of a memory transaction as the process of bringing a cache line into the cache)
2. **Aligned:** the global memory accesses by threads within a warp start at an address boundary that is an even multiple of the size of a global memory transaction

Optimizing CUDA Memory Accesses

Global memory transactions are either 32 or 128 contiguous bytes.

Size of a memory transaction depends on which caches it passes through.

- **L1 + L2 - 128 bytes, L2 only - 32 bytes**



Optimizing CUDA Memory Accesses

Which caches a global memory transaction passes through depends on GPU architecture and the type of access (read vs. write).

| Architecture | Reads Cacheable in L1? | Cached by Default? |
|----------------------|------------------------|--------------------|
| Fermi | Yes | Yes |
| Kepler | No | No |
| Kepler K40 and later | Yes | No |

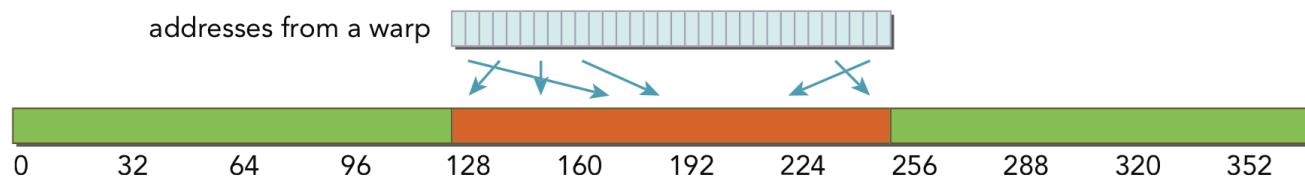
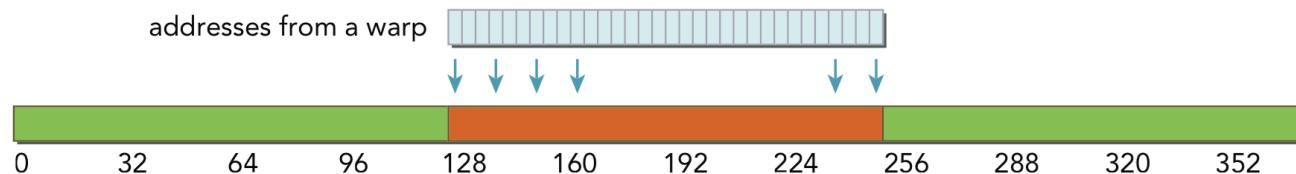
Global memory reads cached in L2 if not in L1

| Architecture | Writes Cacheable in L1? | Cached by Default? |
|----------------------|-------------------------|--------------------|
| Fermi | No | No |
| Kepler | No | No |
| Kepler K40 and later | No | No |

Global memory writes always cached in L2

Optimizing CUDA Memory Accesses

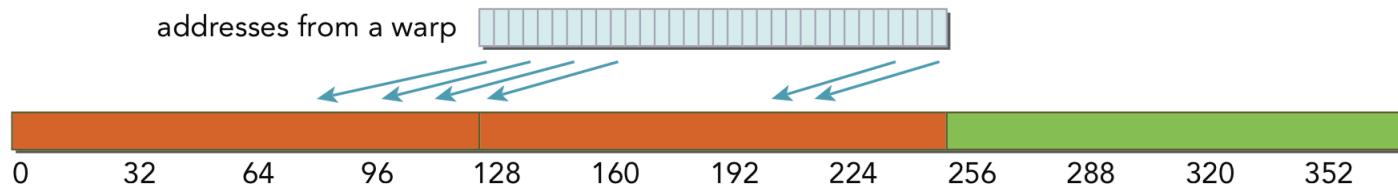
Aligned and Coalesced Memory Access (*w/ L1 cache*)



With 128-byte transactions, a single transaction is required and all of the loaded bytes are used

Optimizing CUDA Memory Accesses

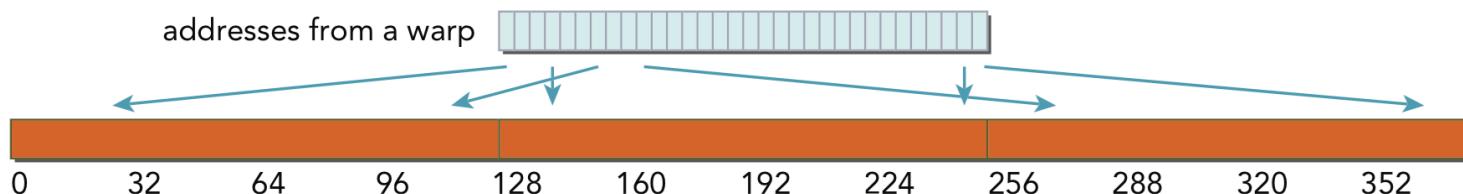
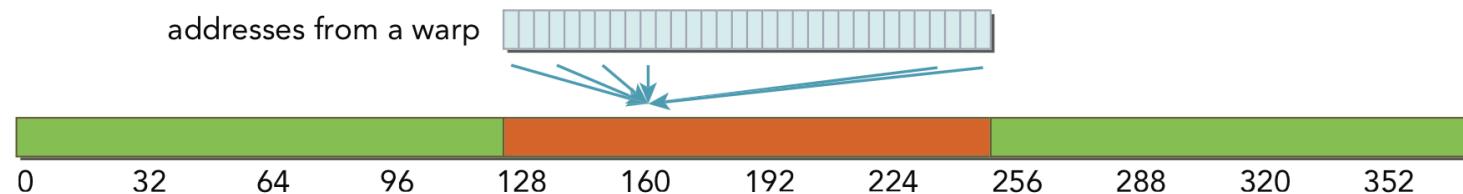
Misaligned and Coalesced Memory Access (*w/ L1 cache*)



With 128-byte transactions, two memory transactions are required to load all requested bytes. Only half of the loaded bytes are used.

Optimizing CUDA Memory Accesses

Misaligned and Uncoalesced Memory Access (*w/ L1 cache*)



With uncoalesced loads, many more bytes loaded than requested.

Optimizing CUDA Memory Accesses

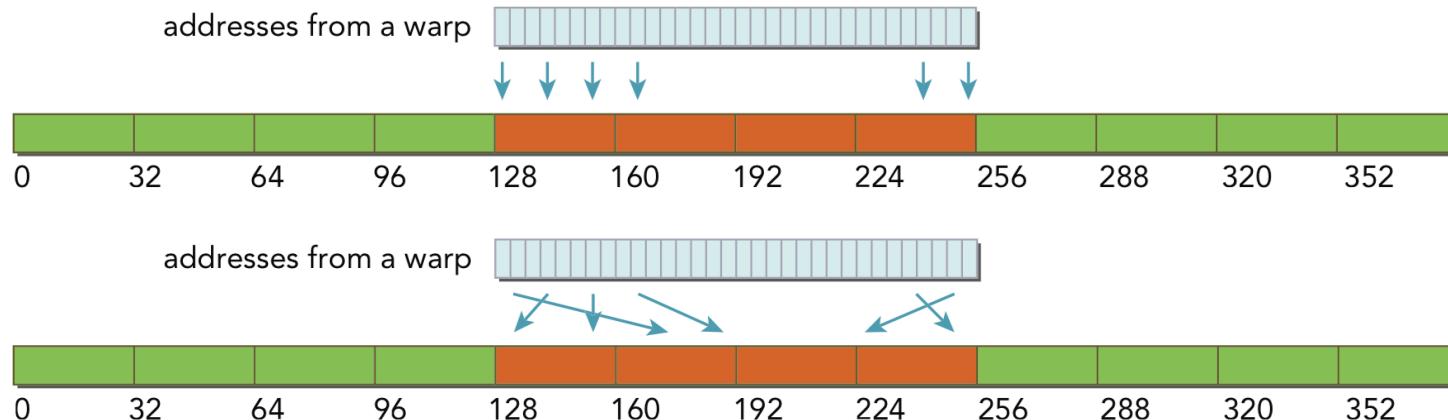
Memory accesses that are not cached in L1 are serviced by 32-byte transactions

This can improve memory bandwidth utilization

However, the L2 cache is device-wide, higher latency than L1, and still relatively small → many applications may take a performance hit if L1 cache is not used for reads

Optimizing CUDA Memory Accesses

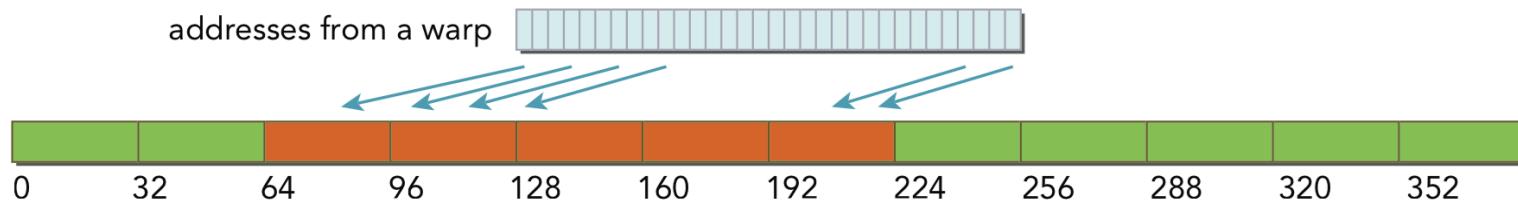
Aligned and Coalesced Memory Access (*w/o L1 cache*)



With 32-byte transactions, four transactions are required and all of the loaded bytes are used

Optimizing CUDA Memory Accesses

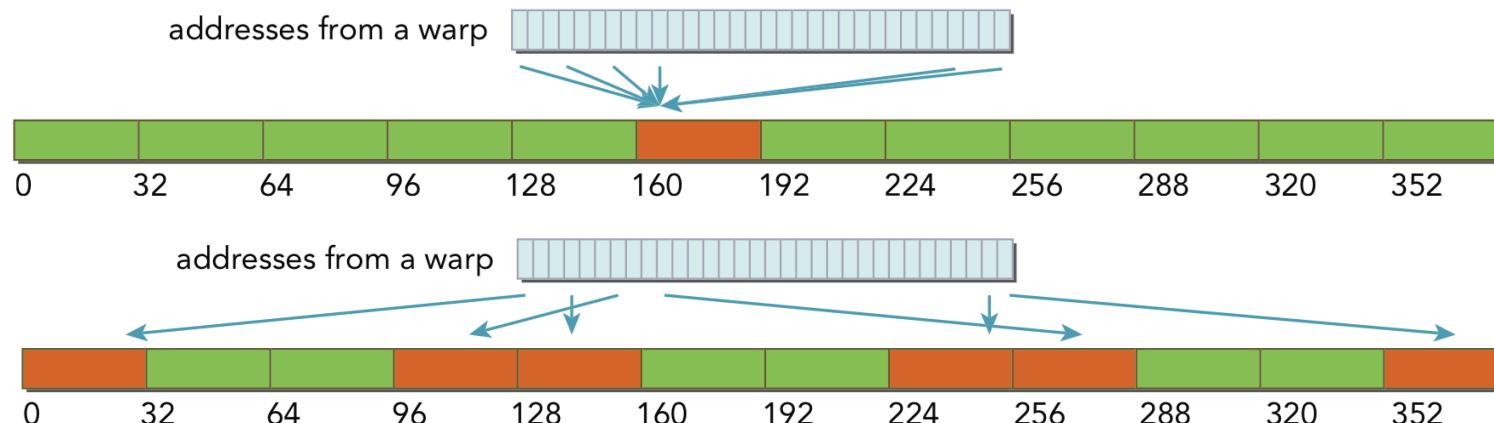
Misaligned and Coalesced Memory Access (*w/o L1 cache*)



With 32-byte transactions, extra memory transactions are still required to load all requested bytes but the number of wasted bytes is likely reduced compared to 128-byte transactions.

Optimizing CUDA Memory Accesses

Misaligned and Uncoalesced Memory Access (*w/o L1 cache*)

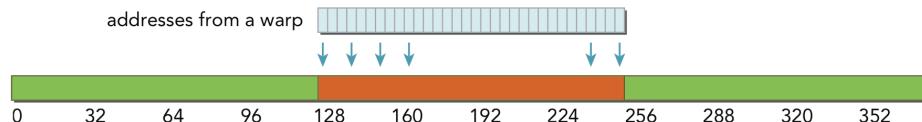


With uncoalesced loads, more bytes loaded than requested but better efficiency than with 128-byte transactions.

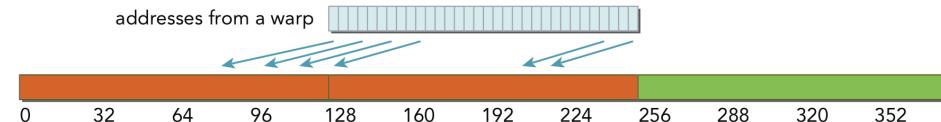
Hands On – Memory Access Optimization

To illustrate these points, we'll experiment with a variation on the vector add micro-benchmark under `$HOME/bootcamp-gpu/src/03_vecadd_misaligned`.

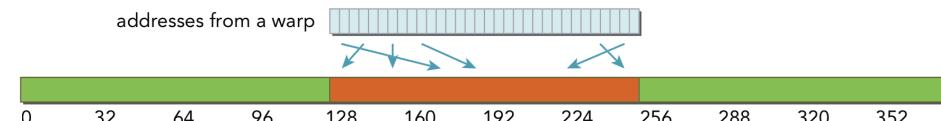
`vector_add`



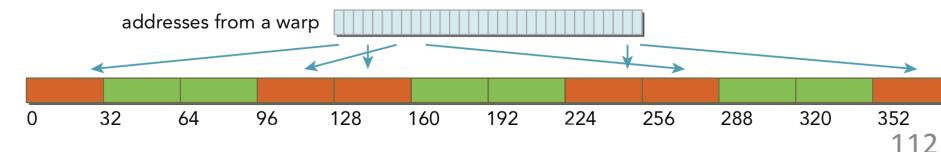
`vector_add_read_offset`



`vector_add_write_offset`



`vector_add_weirdly_coalesced`



`vector_add_not_coalesced`

Hands On – Memory Access Optimization

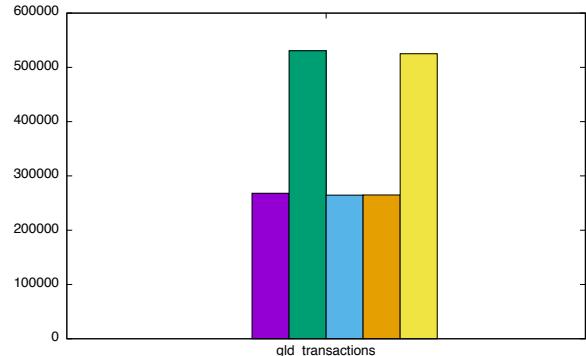
Experiment with several new metrics from nvprof:

1. gld_transactions: # of global memory load transactions
2. gst_transactions: # of global memory store transactions
3. gld_transactions_per_request: mean # of gld txs issued to satisfy a warp's request
4. gst_transactions_per_request: mean # of gst txs issued to satisfy a warp's request
5. gld_efficiency: gld_requests / gld_transactions
6. gst_efficiency: gst_requests / gst_transactions

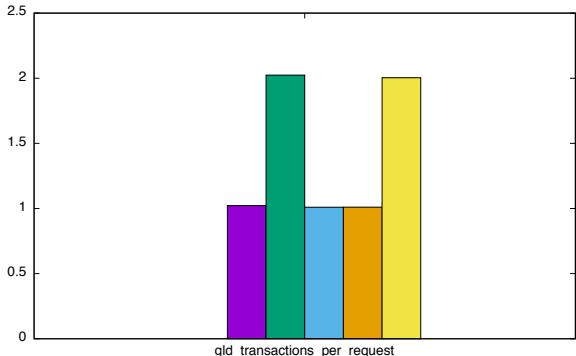
```
$ nvprof --metrics gld_transactions ./vecadd 4194304  
$ nvprof --metrics gst_transactions ./vecadd 4194304  
$ ...
```

What do you observe?

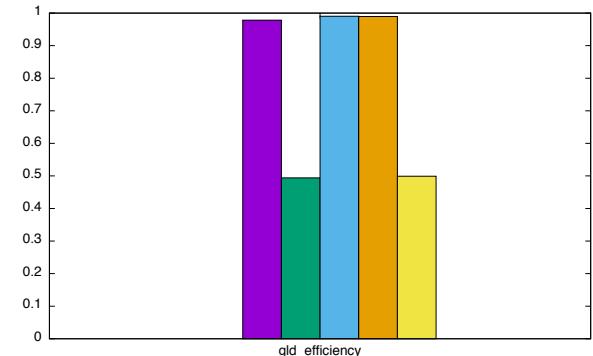
Hands On – Memory Access Optimization



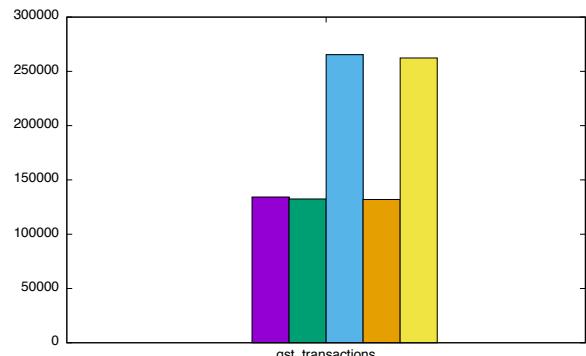
Misaligned, uncoalesced reads \uparrow `gld_transactions`



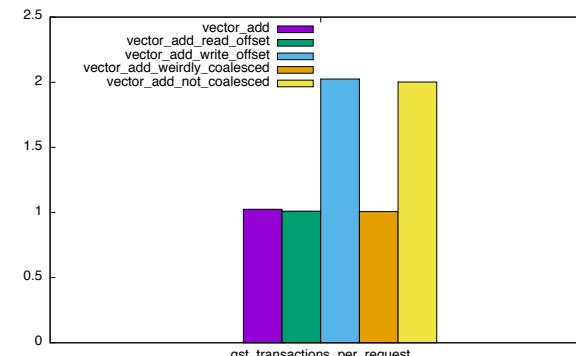
Misaligned, uncoalesced reads \uparrow `gld_transactions_per_request`



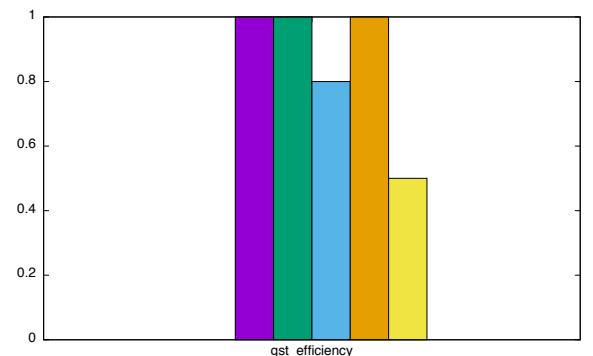
Misaligned, uncoalesced reads \downarrow `gld_efficiency`



Misaligned, uncoalesced writes \uparrow `gst_transactions`



Misaligned, uncoalesced writes \uparrow `gst_transactions_per_request`



Misaligned, uncoalesced writes \downarrow `gst_efficiency`

Optimizing CUDA Memory Accesses

| | | | | | | |
|--|---|------------------------------|---------------------------------------|-----------------|-----------------|-----------------|
| Kernel: <code>vector_add(int*, int*, int*, int)</code> | 1 | gld_transactions | Global Load Transactions | 267988 | 267988 | 267988 |
| | 1 | gst_transactions | Global Store Transactions | 134288 | 134288 | 134288 |
| | 1 | gld_transactions_per_request | Global Load Transactions Per Request | 1.022293 | 1.022293 | 1.022293 |
| | 1 | gst_transactions_per_request | Global Store Transactions Per Request | 1.024536 | 1.024536 | 1.024536 |
| | 1 | gld_efficiency | Global Memory Load Efficiency | 97.82% | 97.82% | 97.82% |
| | 1 | gst_efficiency | Global Memory Store Efficiency | 100.00% | 100.00% | 100.00% |
| Kernel: <code>vector_add_read_offset(int*, int*, int*, int, int)</code> | 1 | gld_transactions | Global Load Transactions | 530593 | 530593 | 530593 |
| | 1 | gst_transactions | Global Store Transactions | 132440 | 132440 | 132440 |
| | 1 | gld_transactions_per_request | Global Load Transactions Per Request | 2.024052 | 2.024052 | 2.024052 |
| | 1 | gst_transactions_per_request | Global Store Transactions Per Request | 1.010437 | 1.010437 | 1.010437 |
| | 1 | gld_efficiency | Global Memory Load Efficiency | 49.41% | 49.41% | 49.41% |
| | 1 | gst_efficiency | Global Memory Store Efficiency | 100.00% | 100.00% | 100.00% |
| Kernel: <code>vector_add_write_offset(int*, int*, int*, int, int)</code> | 1 | gld_transactions | Global Load Transactions | 264684 | 264684 | 264684 |
| | 1 | gst_transactions | Global Store Transactions | 265384 | 265384 | 265384 |
| | 1 | gld_transactions_per_request | Global Load Transactions Per Request | 1.009689 | 1.009689 | 1.009689 |
| | 1 | gst_transactions_per_request | Global Store Transactions Per Request | 2.024719 | 2.024719 | 2.024719 |
| | 1 | gld_efficiency | Global Memory Load Efficiency | 99.04% | 99.04% | 99.04% |
| | 1 | gst_efficiency | Global Memory Store Efficiency | 80.00% | 80.00% | 80.00% |
| Kernel: <code>vector_add_weirdly_coalesced(int*, int*, int*, int)</code> | 1 | gld_transactions | Global Load Transactions | 264908 | 264908 | 264908 |
| | 1 | gst_transactions | Global Store Transactions | 132020 | 132020 | 132020 |
| | 1 | gld_transactions_per_request | Global Load Transactions Per Request | 1.010544 | 1.010544 | 1.010544 |
| | 1 | gst_transactions_per_request | Global Store Transactions Per Request | 1.007233 | 1.007233 | 1.007233 |
| | 1 | gld_efficiency | Global Memory Load Efficiency | 98.96% | 98.96% | 98.96% |
| | 1 | gst_efficiency | Global Memory Store Efficiency | 200.00% | 200.00% | 200.00% |
| Kernel: <code>vector_add_not_coalesced(int*, int*, int*, int)</code> | 1 | gld_transactions | Global Load Transactions | 525168 | 525168 | 525168 |
| | 1 | gst_transactions | Global Store Transactions | 262360 | 262360 | 262360 |
| | 1 | gld_transactions_per_request | Global Load Transactions Per Request | 2.003357 | 2.003357 | 2.003357 |
| | 1 | gst_transactions_per_request | Global Store Transactions Per Request | 2.001648 | 2.001648 | 2.001648 |
| | 1 | gld_efficiency | Global Memory Load Efficiency | 49.92% | 49.92% | 49.92% |
| | 1 | gst_efficiency | Global Memory Store Efficiency | 50.00% | 50.00% | 50.00% |

Review - Optimizing CUDA Memory Access

Seek **aligned** and **coalesced** global memory accesses when optimizing CUDA kernel performance.

The way in which a load/store is serviced depends on the cache(s) it passes through.

Aligned and coalesced accesses reduce the number of transactions necessary and the efficiency of each of those transactions (i.e. # bytes requested / # bytes loaded).

CUDA Atomic Operations (Briefly)

Atomic Operations

Atomic operations are a special class of mathematical operations in computing

- An atomic operation is performed uninterruptedly, so that there is no interference from other threads
- When a thread's atomic operation has completed, it can be certain its requested changes have been made without interference from other threads
- Atomic operations are particularly useful on the massively parallel GPU

Atomic operations may improve correctness, but can have a detrimental impact on performance

- Out of thousands of threads, only one can succeed in accessing a shared variable
- Atomicity requires reading and writing DRAM, no caching allowed

Atomic Operations

```
__global__ void sumAll(int *in_array, int N, int *out_scalar) {
    int i = blockIdx.x * blockDim.x + threadIdx.x;
    if (i < N) {
        // out_scalar += out_array[i] would cause massive data race
        atomicAdd(out_scalar, in_array[i]);
    }
}
```

Atomic Operations Summary

CUDA supports a variety of atomic operations. CAS can be used to build your own.

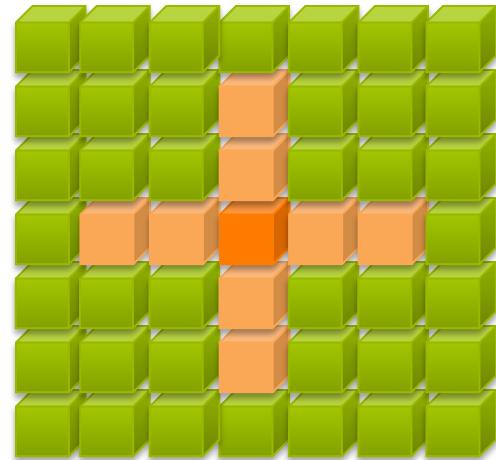
| OPERATION | FUNCTION | SUPPORTED TYPES |
|--------------------|------------|--|
| Addition | atomicAdd | int, unsigned int, unsigned long long int, float |
| Subtraction | atomicSub | int, unsigned int |
| Unconditional Swap | atomicExch | int, unsigned int, unsigned long long int, float |
| Minimum | atomicMin | int, unsigned int, unsigned long long int |
| Maximum | atomicMax | int, unsigned int, unsigned long long int |
| Increment | atomicInc | unsigned int |
| Decrement | atomicDec | unsigned int |
| Compare-And-Swap | atomicCAS | int, unsigned int, unsigned long long int |
| And | atomicAnd | int, unsigned int, unsigned long long int |
| Or | atomicOr | int, unsigned int, unsigned long long int |
| Xor | atomicXor | int, unsigned int, unsigned long long int |

Hands on with a Real World App

A More Complex Application

2D waveform propagation (stencil)

```
for (int y = 0; y < ny; y++) {  
    for (int x = 0; x < nx; x++) {  
  
        for (int d = 1; d <= radius; d++) {  
            div += c_coeff[d] * (curr[y_pos_offset] +  
                curr[y_neg_offset] + curr[x_pos_offset] +  
                curr[x_neg_offset]);  
        }  
  
        next[this_offset] = temp + div * vsq[this_offset];  
    }  
}
```



We'll use iso2d as a more realistic example to explore optimizations.

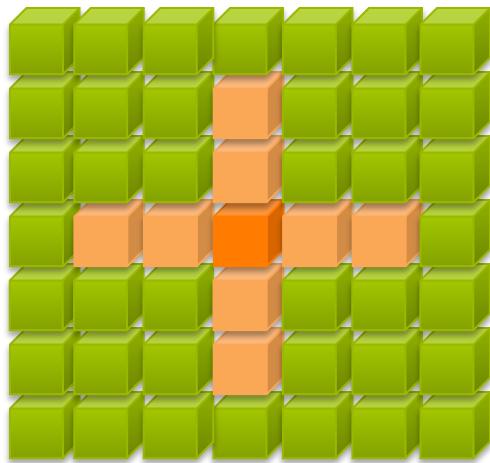
A More Complex Application

Provided reference code:

- `04_iso2d_seq/`: sequential C implementation

```
$ cd $HOME/bootcamp-gpu/src/04_iso2d_seq  
$ make
```

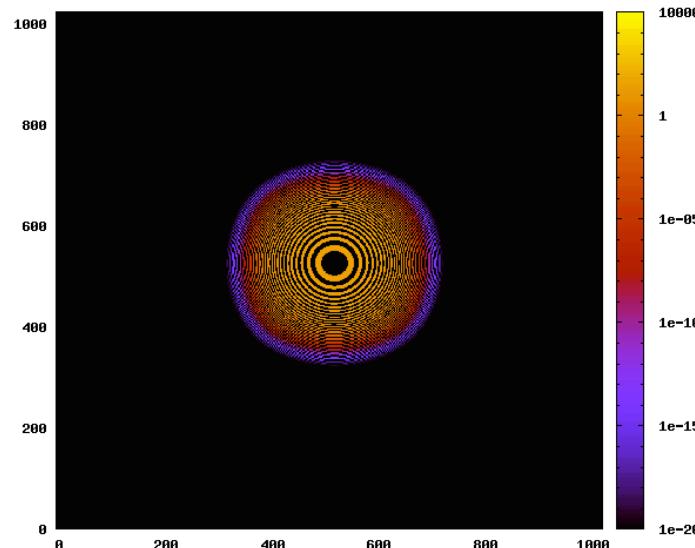
```
$ ./iso2d -x 1024 -y 1024 -i 800 -t
```



A More Complex Example

iso2d's -t flag will produce an output file to visualize for correctness.

```
$ ./iso2d_common/iso.sh -x 1024 -y 1024 -i snap.text  
$ eog iso.png # Must use ssh -X to connect to DAVINCI
```



Hands On – A Quick Review of OpenMP

Let's start by examining the parallelism possible in iso2d using OpenMP.

Where could you add a `#pragma omp parallel for` in iso2d.cpp for the most parallelism?

Compare the performance of your OpenMP implementation to the provided sequential code by modifying iso2d.cpp in 04_iso2d_seq.

Visualize the results to verify the correctness of your OpenMP implementation.

(If you're having trouble you can cheat and look in 05_iso2d_omp/)

Hands On – OpenMP to CUDA

Next, your task is to port the OpenMP version of `iso2d` to CUDA, and check its performance.

`06_iso2d_cuda/iso2d.cu` contains TODOs to help in the port.

Use `nvprof`'s Summary Mode to check where your implementation is spending time.

(If you're having trouble you can cheat and look in `06_iso2d_cuda/iso2d_solution.cu`)

Hands On – OpenMP to CUDA

Use nvprof's Summary Mode to check where your implementation is spending time.

```
iso_r4_2x: 0.4557509422 s total, 0.0005638975 s/step, 1859.52 Mcells/s/step
==19313== Profiling application: ./iso2d_cuda_solution -x 1024 -y 1024 -i 800
==19313== Profiling result:
Time(%)      Time      Calls      Avg      Min      Max      Name
 98.94%  432.51ms        800  540.64us  537.07us  546.56us  fwd_kernel
   0.85%  3.7121ms        804   4.6160us  1.2150us  955.81us  [CUDA memcpy HtoD]
   0.21%  911.11us         1  911.11us  911.11us  911.11us  [CUDA memcpy DtoH]

==19313== API calls:
```

| Time(%) | Time | Calls | Avg | Min | Max | Name |
|---------|----------|-------|----------|----------|----------|-------------------|
| 82.61% | 444.60ms | 805 | 552.29us | 9.3120us | 1.2511ms | cudaMemcpy |
| 15.56% | 83.736ms | 4 | 20.934ms | 144.61us | 83.285ms | cudaMalloc |
| 1.18% | 6.3430ms | 800 | 7.9280us | 7.1660us | 28.247us | cudaLaunch |
| 0.29% | 1.5825ms | 6400 | 247ns | 207ns | 8.7500us | cudaSetupArgument |

Hands On – OpenMP to CUDA

Dig deeper using nvprof metrics. We'll start by just looking at the ones we've tried in past examples.

\$ nvprof --metrics <metric> -x 1024 -y 1024 -i 800 -t with:

1. ipc
2. sm_efficiency
3. alu_fu_utilization
4. gld_transactions
5. gst_transactions
6. gld_transactions_per_request
7. gst_transactions_per_request
8. gld_efficiency
9. gst_efficiency

Hands On – OpenMP to CUDA

Dig deeper using nvprof metrics. We'll start by just looking at the ones we've tried in past examples.

| Invocations | Metric Name | Metric Description | Avg |
|--------------------------|------------------------------|---------------------------------------|----------|
| Device "Tesla M2050 (0)" | | | |
| Kernel: fwd_kernel | | | |
| 800 | sm_efficiency | Multiprocessor Activity | 98.08% |
| 800 | ipc | Executed IPC | 1.128842 |
| 800 | gld_transactions | Global Load Transactions | 1238112 |
| 800 | gst_transactions | Global Store Transactions | 65994 |
| 800 | gld_transactions_per_request | Global Load Transactions Per Request | 1.988641 |
| 800 | gst_transactions_per_request | Global Store Transactions Per Request | 2.013987 |
| 800 | gld_efficiency | Global Memory Load Efficiency | 50.29% |
| 800 | gst_efficiency | Global Memory Store Efficiency | 80.00% |
| 800 | alu_fu_utilization | Arithmetic Function Unit Utilization | Mid (5) |

What does this kernel load?

| Accessed | Coalesced | Aligned? | Frequency of Access |
|--------------------|-----------|-------------|---------------------|
| curr[this_offset] | Yes | Maybe | A few threads |
| next[this_offset] | Yes | Maybe | A few threads |
| curr[y_pos_offset] | Yes | Maybe | A few threads |
| curr[y_neg_offset] | Yes | Maybe | A few threads |
| curr[x_pos_offset] | Yes | Maybe | A few threads |
| curr[x_neg_offset] | Yes | Maybe | A few threads |
| vsq[this_offset] | Yes | Maybe | A few threads |
| c_coeff[d] | No | Usually Not | All threads |

iso2d

Constant memory is optimized for broadcast accesses (all threads read same location).

Variables in CUDA constant memory must be declared statically:

```
__constant__ TYPE const_c_coeff[NUM_COEFF];
```

Can be referenced as normal arrays from CUDA kernels:

```
const_c_coeff[d]
```

Initialized using a special-purpose cudaMemcpy API:

```
cudaMemcpyToSymbol(const_c_coeff, c_coeff, NUM_COEFF * sizeof(TYPE));
```

Hands On – Constant Memory

Try your hand at optimizing the iso2d CUDA version using constant memory.

Re-analyze the overall performance and the same profiler metrics.

(Example in 07_iso2d_cuda_cmem)

Hands On – Constant Memory

Try your hand at optimizing the iso2d CUDA version using constant memory.

Re-analyze the overall performance and the same profiler metrics.

| Version | Performance (Mcells/s/step) | Improvement |
|------------|-----------------------------|-------------|
| Sequential | 33.56 | |
| OpenMP | 167.71 | 5.00 |
| CUDA | 2,016.59 | 12.02 |
| + cmem | 2,104.14 | 1.04 |

Hands On – Constant Memory

Try your hand at optimizing the iso2d CUDA version using constant memory.

Re-analyze the overall performance and the same profiler metrics.

| Invocations | Metric Name | Metric Description | Avg |
|--------------------------|------------------------------|---------------------------------------|----------|
| Device "Tesla M2050 (0)" | | | |
| Kernel: fwd_kernel | | | |
| 800 | sm_efficiency | Multiprocessor Activity | 97.99% |
| 800 | ipc | Executed IPC | 1.115025 |
| 800 | gld_transactions | Global Load Transactions | 1238327 |
| 800 | gst_transactions | Global Store Transactions | 66072 |
| 800 | gld_transactions_per_request | Global Load Transactions Per Request | 1.988987 |
| 800 | gst_transactions_per_request | Global Store Transactions Per Request | 2.016384 |
| 800 | gld_efficiency | Global Memory Load Efficiency | 50.28% |
| 800 | gst_efficiency | Global Memory Store Efficiency | 80.00% |
| 800 | alu_fu_utilization | Arithmetic Function Unit Utilization | Mid (5) |

No significant change in metrics

Hands On – Constant Memory

Back to the drawing board.

| Accessed | Coalesced | Aligned? | Frequency of Access |
|--------------------|-----------|-------------|---------------------|
| curr[this_offset] | Yes | Maybe | A few threads |
| next[this_offset] | Yes | Maybe | A few threads |
| curr[y_pos_offset] | Yes | Maybe | A few threads |
| curr[y_neg_offset] | Yes | Maybe | A few threads |
| curr[x_pos_offset] | Yes | Maybe | A few threads |
| curr[x_neg_offset] | Yes | Maybe | A few threads |
| vsq[this_offset] | Yes | Maybe | A few threads |
| c_coeff[d] | No | Usually Not | All threads |

Hands On – Access Alignment

Depending on the width of the 2D grid we're applying our stencil to, there may be many accesses that are mis-aligned.

By padding rows in the grid to be a length evenly divisible by 128 bytes, we can improve the # of aligned accesses (thereby reducing transactions).

Try writing a version of your constant memory version of `iso2d` that pads all rows of `curr`, `next`, and `vsq` out to 128 bytes.

Re-analyze the overall performance and the same profiler metrics.

Hands On – Access Alignment

| Version | Performance (Mcells/s/step) | Improvement |
|------------|-----------------------------|-------------|
| Sequential | 33.56 | |
| OpenMP | 167.71 | 5.00 |
| CUDA | 2,016.59 | 12.02 |
| + cmem | 2,104.14 | 1.04 |
| + aligned | 2,806.49 | 1.33 |

Hands On – Access Alignment

| Metric | Old Value | New Value |
|------------------------------|-----------|-----------|
| sm_efficiency | 98.08% | 98.16% |
| ipc | 1.128842 | 1.594335 |
| gld_transactions | 1238112 | 887798 |
| gst_transactions | 65994 | 32884 |
| gld_transactions_per_request | 1.988641 | 1.425972 |
| gst_transactions_per_request | 2.013987 | 1.003541 |
| gld_efficiency | 50.29% | 70.13% |
| gst_efficiency | 80.00% | 100.00% |
| alu_fu_utilization | Mid (5) | High (7) |

Hands On – nvprof Throughput Metrics

Let's collect some new metrics:

1. gld_throughput: Achieved throughput for global memory accesses
2. dram_read_throughput: Throughput between DRAM and L2
3. l2_read_throughput: Throughput through L2 cache

Hands On – nvprof Throughput Metrics

1. gld_throughput: Achieved throughput for global memory accesses
 2. dram_read_throughput: Throughput between DRAM and L2
 3. l2_read_throughput: Throughput through L2 cache

```

Kernel: fwd_kernel(float*, float*, float*, int, int, int, int)
    gld_throughput          Global Load Throughput
    dram_read_throughput   Device Memory Read Throughput
    l2_read_throughput     L2 Throughput (Reads)

```

Massive discrepancy?

| |
|------------|
| 288.79GB/s |
| 49.043GB/s |
| 70.209GB/s |

iso2d

Throughput metrics suggest that we're hitting a lot in L1 cache.

Can confirm using the `l1_cache_global_hit_rate` metric.

| Metric Name | Metric Description | Avg |
|---------------------------------------|--------------------|---------------|
| Device "Tesla M2050 (0)" | | |
| Kernel: fwd_kernel | | |
| <code>l1_cache_global_hit_rate</code> | L1 Global Hit Rate | 78.41% |

Good: we're getting lots of locality benefits.

Bad: L1 cache may be a bottleneck, losing some of our on-chip memory to shared memory.

iso2d

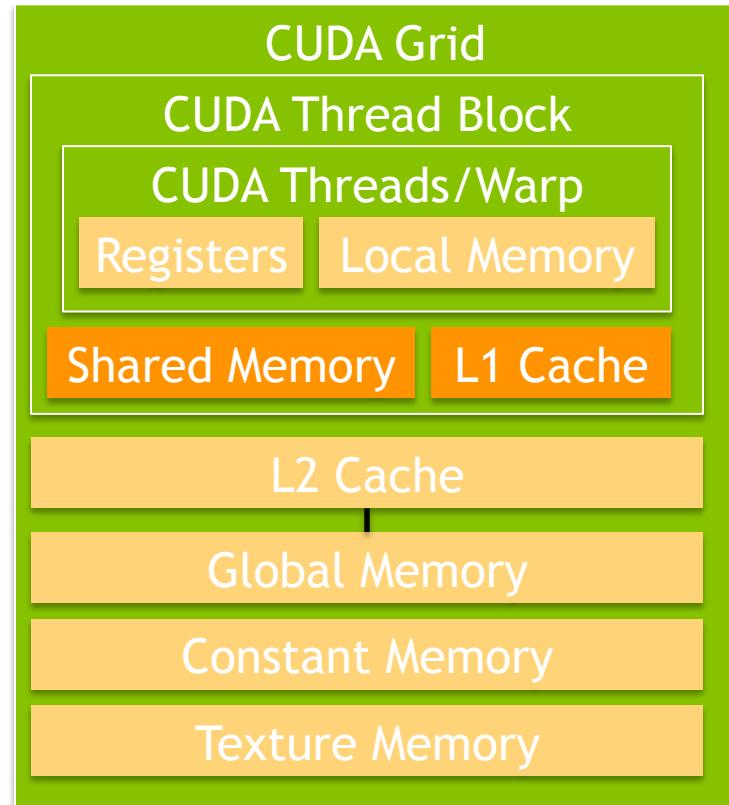
Original discussion of shared memory:

- On-chip, low latency, high bandwidth
- Same physical memory as L1 cache

Reconfigure the partitioning of on-chip memory between L1 and shared memory:

```
cudaError_t cudaThreadSetCacheConfig(  
    enum cudaFuncCache cacheConfig);
```

- `cudaFuncCachePreferNone` (default)
- `cudaFuncCachePreferShared`
- `cudaFuncCachePreferL1`



Hands On – On-Chip Memory Configuration

How does `cudaThreadSetCacheConfig` affect the performance of `iso2d`?

Try changing the call to it at the start of the `08_iso2d_cuda_aligned` example to be either:

```
cudaThreadSetCacheConfig(cudaFuncCachePreferShared);
```

```
cudaThreadSetCacheConfig(cudaFuncCachePreferL1);
```

Then, measure any change in overall performance and L1 hit ratios.

Hands On – On-Chip Memory Configuration

| Version | Performance (Mcells/s/step) | Improvement |
|-----------------------------|-----------------------------|-------------|
| Sequential | 33.56 | |
| OpenMP | 167.71 | 5.00 |
| CUDA | 2,016.59 | 12.02 |
| + cmem | 2,104.14 | 1.04 |
| + aligned | 2,806.49 | 1.33 |
| + cudaFuncCachePreferShared | 2,805.55 | 1.00 |
| + cudaFuncCachePreferL1 | 2,825.09 | 1.01 |

w/ `cudaFuncCachePreferNone`:

78.41% L1 Global Hit Rate

w/ `cudaFuncCachePreferShared`:

78.41% L1 Global Hit Rate

w/ `cudaFuncCachePreferL1`:

79.73% L1 Global Hit Rate

Hands On – On-Chip Memory Configuration

| Version | Performance (Mcells/s/step) | Improvement |
|------------------------------|-----------------------------|-------------|
| Sequential | 33.56 | |
| OpenMP | 167.71 | 5.00 |
| CUDA | 2,016.59 | 12.02 |
| CUDA w/ cmem | 2,104.14 | 1.04 |
| CUDA w/ cmem and aligned | 2,806.49 | 1.33 |
| w/ cudaFuncCachePreferShared | 2,805.55 | 1.00 |
| w/ cudaFuncCachePreferL1 | 2,825.09 | 1.01 |

w/ cudaFuncCachePreferNone:

78.41% L1 Global Hit Rate

w/ cudaFuncCachePreferShared:

78.41% L1 Global Hit Rate

w/ cudaFuncCachePreferL1:

79.73% L1 Global Hit Rate

**Appears to be some improvement possible by increasing
on-chip locality.**

Time to look in to shared memory.

CUDA Shared Memory

Shared memory is shared by threads in the same block.

- On-SM memory

Shared memory can be allocated statically or dynamically

Statically Allocated Shared Memory

- Size is fixed at compile-time
- Can declare many statically allocated shared memory variables
- Can be declared globally or inside a device function
- Can be multi-dimensional

```
__shared__ int s_arr[256][256];
```

CUDA Shared Memory

Dynamically Allocated Shared Memory

- Size in bytes is set at kernel launch with a third kernel launch configurable
- Can only have one dynamically allocated shared memory array per kernel
- Must be one-dimensional array

```
__global__ void kernel(...) {
    extern __shared__ int s_arr[];
    ...
}

kernel<<<nblocks, threads_per_block, shared_memory_bytes>>>(...);
```

CUDA Shared Memory

What is reused in iso2d? i.e. what can benefit from improved on-chip locality?

```
for (int d = 1; d <= radius; d++) {
    int y_pos_offset = POINT_OFFSET(x, y + d, dimx, radius);
    int y_neg_offset = POINT_OFFSET(x, y - d, dimx, radius);
    int x_pos_offset = POINT_OFFSET(x + d, y, dimx, radius);
    int x_neg_offset = POINT_OFFSET(x - d, y, dimx, radius);
    div += const_c_coeff[d] * (curr[y_pos_offset] +
                               curr[y_neg_offset] + curr[x_pos_offset] +
                               curr[x_neg_offset]);
}
```

Hands On – Shared Memory

Take a look at using shared memory to tile curr.

This is the toughest transformation we'll try today, so there is also a reference implementation under `09_iso2d_cuda_smem/`.

You can start by declaring a dynamic shared memory allocation in the kernel, and figuring out how large it needs to be from the launch:

```
extern __shared__ TYPE cache[];  
fwd_kernel<<<..., SHAREDY(conf.radius) * SHAREDX(conf.radius) *  
        sizeof(TYPE)>>>(...);
```

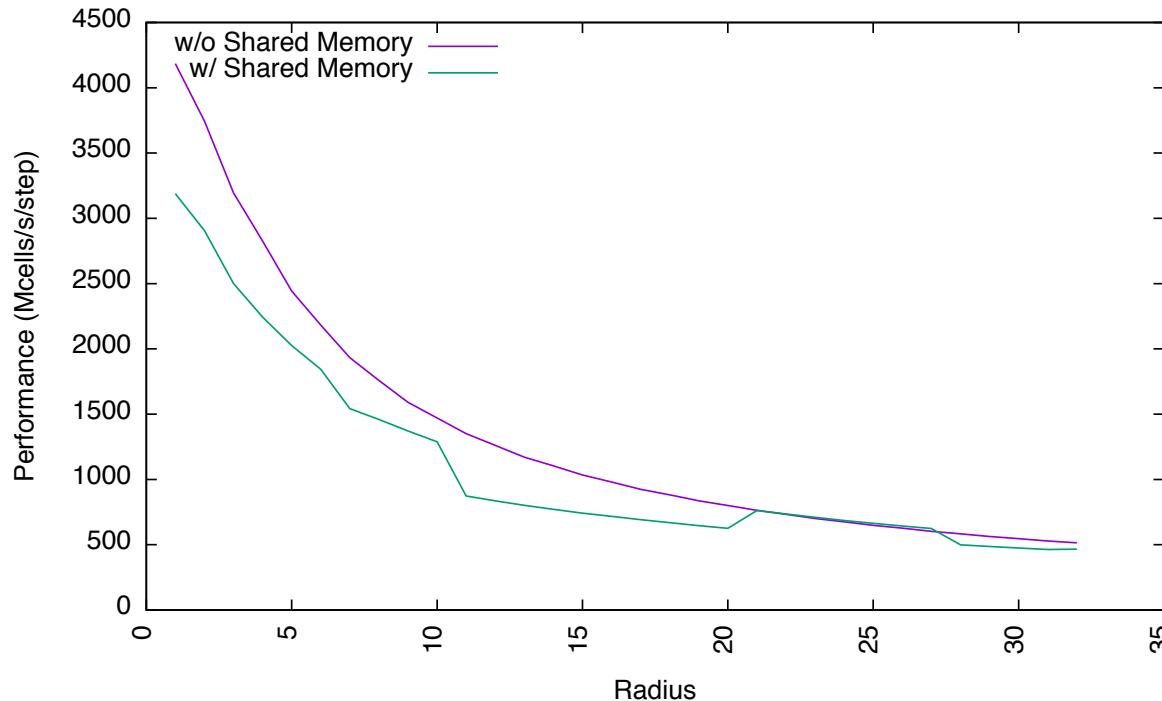
Give it a shot!

Hands On – Shared Memory

| Version | Performance (Mcells/s/step) | Improvement |
|-------------------------|-----------------------------|-------------|
| Sequential | 33.56 | |
| OpenMP | 167.71 | 5.00 |
| CUDA | 2,016.59 | 12.02 |
| + cmem | 2,104.14 | 1.04 |
| + aligned | 2,806.49 | 1.33 |
| + cudaFuncCachePreferL1 | 2,825.09 | 1.01 |
| + smem | 2,242.79 | 0.79 |

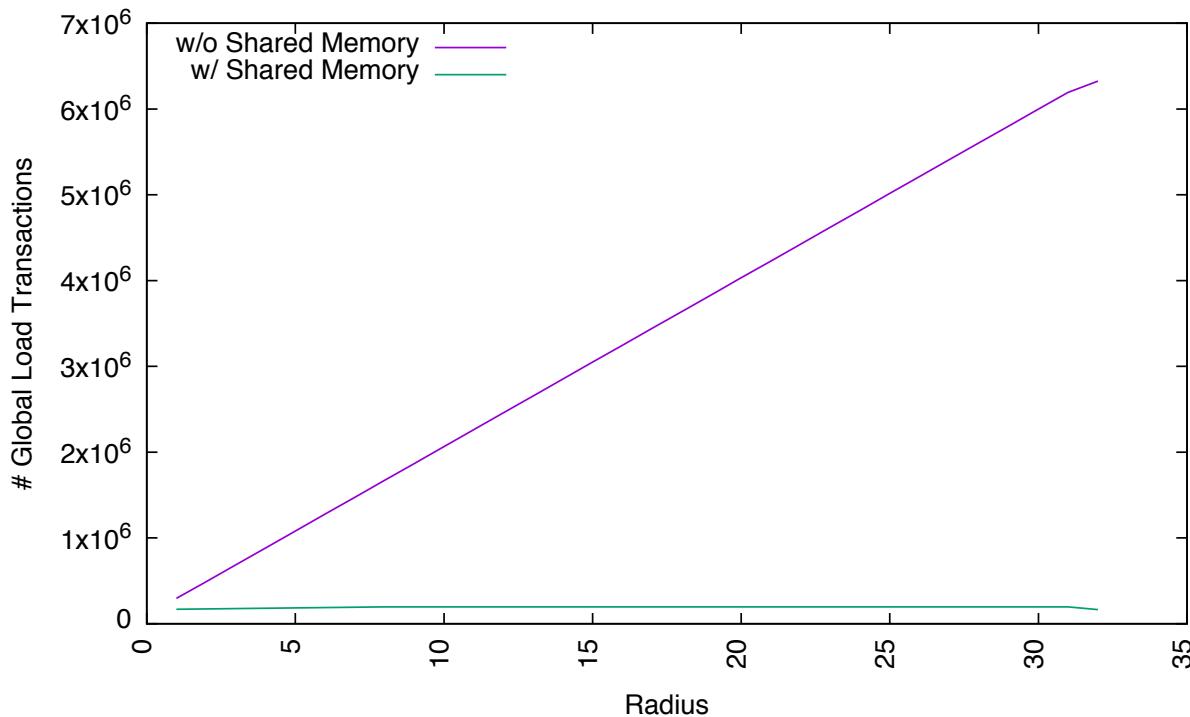
Hands On – Shared Memory

Increasing radius causes more reuse, explicit management helps.



Hands On – Shared Memory

Very effective at reducing global load transactions.



Hands On – Shared Memory

Still don't see any overall performance benefit because our manual management of shared memory is doing no better than the automatic management of L1, but incurs additional overheads (instructions executed, registers, etc).

Without shared memory, L1 Global Hit Rate = 80.85%

With shared memory, L1 Global Hit Rate = 4.49%

iso2d Summary

| Version | Performance (Mcells/s/step) | Improvement |
|-------------------------|-----------------------------|-------------|
| Sequential | 33.56 | |
| OpenMP | 167.71 | 5.00 |
| CUDA | 2,016.59 | 12.02 |
| + cmem | 2,104.14 | 1.04 |
| + aligned | 2,806.49 | 1.33 |
| + cudaFuncCachePreferL1 | 2,825.09 | 1.01 |
| + smem | 2,242.79 | 0.79 |

CUDA Libraries

CUDA Libraries

CUDA Libraries offer high-level, high performance, domain-specific APIs for GPUs.

Examples:

1. CUSPARSE: Sparse linear algebra
2. CUBLAS: Dense linear algebra
3. CUDNN: Deep Neural Networks
4. CUFFT: Fast Fourier Transforms

Take advantage of expert-tuned kernels without writing a line of kernel code.

Many CUDA libraries emulate their host counterparts (e.g. CUBLAS and BLAS).

Common Library Workflow

1. Create a library-specific handle that manages contextual information useful for the library's operation.
 - Many CUDA Libraries have the concept of a handle which stores opaque library-specific information on the host which many library functions access
 - Programmer's responsibility to manage this handle
 - For example: `cublasHandle_t`, `cufftHandle`, `cusparseHandle_t`, `curandGenerator_t`
2. Allocate device memory for inputs and outputs to the library function.
 - Use `cudaMalloc` as usual

Common Library Workflow

3. If inputs are not already in a library-supported format, convert them to be accessible by the library.
 - Many CUDA Libraries only accept data in a specific format
 - For example: column-major vs. row-major arrays
4. Populate the pre-allocated device memory with inputs in a supported format.
 - In many cases, this step simply implies a `cudaMemcpy` or one of its variants to make the data accessible on the GPU
 - Some libraries provide custom transfer functions, for example: `cublasSetVector` optimizes strided copies for the CUBLAS library

Common Library Workflow

5. Configure the library computation to be executed.
 - In some libraries, this is a no-op
 - Others require additional metadata to execute library computation correctly
 - In some cases this configuration takes the form of extra parameters passed to library functions, others set fields in the library handle
6. Execute a library call that offloads the desired computation to the GPU.
 - No GPU-specific knowledge required

Common Library Workflow

7. Retrieve the results of that computation from device memory, possibly in a library-determined format.
 - Again, this may be as simple as a `cudaMemcpy` or require a library-specific function
8. If necessary, convert the retrieved data to the application's native format.
 - If a conversion to a library-specific format was necessary, this step ensures the application can now use the calculated data
 - In general, it is best to keep the application format and library format the same, reducing overhead from repeated conversions

Common Library Workflow

9. Release CUDA resources.
 - Includes the usual CUDA cleanup (`cudaFree`, `cudaStreamDestroy`, etc) plus any library-specific cleanup
10. Continue with the remainder of the application.

Common Library Workflow

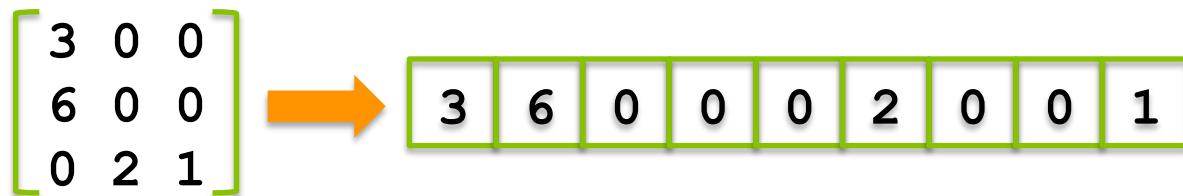
- Not all libraries follow this workflow, and not all libraries require every step in this workflow
 - In fact, for many libraries many steps are skipped
 - Keeping this workflow in mind will help give you context on what the library might be doing behind the scenes and where you are in the process
- Next, we'll look at one simple illustration of this workflow

cuBLAS

- cuBLAS is a port of a popular linear algebra library, BLAS
- cuBLAS (like BLAS) splits its subroutines into multiple levels based on data types processed:
 - Level 1: vector-only operations (e.g. vector addition)
 - Level 2: matrix-vector operations (e.g. matrix-vector multiplication)
 - Level 3: matrix-matrix operations (e.g. matrix multiplication)

cuBLAS Ideosyncracies

- For legacy compatibility, cuBLAS operates on column-major matrices



- cuBLAS also has a legacy API which was dropped in CUDA 4.0, we'll be looking at the new cuBLAS API
 - If you find cuBLAS code that doesn't quite match up, you may be looking at the old cuBLAS API

cuBLAS Data Management

Device memory in cuBLAS is allocated as you're used to: `cudaMalloc`

Transferring data to the device uses cuBLAS-specific functions:
`cublasGetVector/cublasSetVector` and `cublasGetMatrix/cublasSetMatrix`

cuBLAS Data Management

- Example:

```
cublasStatus_t cublasSetVector(int n, int elemSize,  
                               const void *x, int incx, void *y, int incy);
```

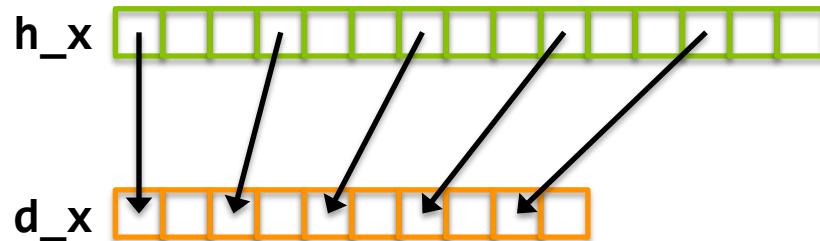
where:

- **n** is the number of elements to transfer to the GPU
- **elemSize** is the size of each element (e.g. `sizeof(int)`)
- **x** is the vector on the host to copy from
- **incx** is a stride in **x** of the array cells to transfer to
- **y** is the vector on the GPU to copy to
- **incy** is a stride in **y** of the array cells to transfer to

cuBLAS Data Management

- Example:

```
cublasSetVector(5, sizeof(int), h_x, 3, d_x, 2);
```



cuBLAS Data Management

- Similarly:

```
cublasStatus_t cublasSetMatrix(int rows, int cols,  
    int elemSize, const void *A, int lda, void *B, int ldb);
```

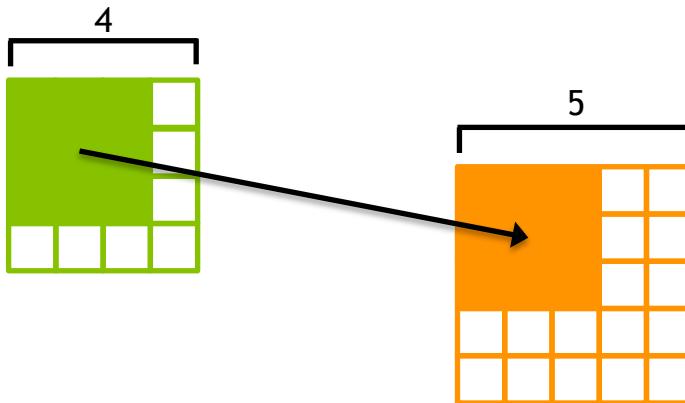
where:

- rows is the number of rows in a matrix to copy
- cols is the number of cols in a matrix to copy
- elemSize is the size of each cell in the matrix (e.g. sizeof(int))
- A is the source matrix on the host
- lda is the number of rows in the underlying array for A
- B is the destination matrix on the GPU
- ldb is the number of rows in the underlying array for B

cuBLAS Data Management

- Similarly:

```
cublasSetMatrix(3, 3, sizeof(int), h_A, 4, d_A, 5);
```



cuBLAS Example

Let's take a look at a simple cuBLAS example that performs a matrix-vector multiplication. This example uses 6 of the 10 steps in the common library workflow:

1. Create a cuBLAS handle using `cublasCreateHandle`
2. Allocate device memory for inputs and outputs using `cudaMalloc`
3. Populate device memory using `cublasSetVector`, `cublasSetMatrix`
4. Call `cublasSgemv` to run matrix-vector multiplication on the GPU
5. Retrieve results from the GPU using `cublasGetVector`
6. Release CUDA and cuBLAS resources using `cudaFree`, `cublasDestroy`

cuBLAS Example

```
cublasCreate(&handle);
cudaMalloc((void **) &dA, sizeof(float) * M * N);
cudaMalloc((void **) &dX, sizeof(float) * N);
cudaMalloc((void **) &dY, sizeof(float) * M);

cublasSetVector(N, sizeof(float), X, 1, dX, 1);
cublasSetVector(M, sizeof(float), Y, 1, dY, 1);
cublasSetMatrix(M, N, sizeof(float), A, M, dA, M);

cublasSgemv(handle, CUBLAS_OP_N, M, N, &alpha, dA, M, dX, 1, &beta, dY, 1);

cublasGetVector(M, sizeof(float), dY, 1, Y, 1);
```

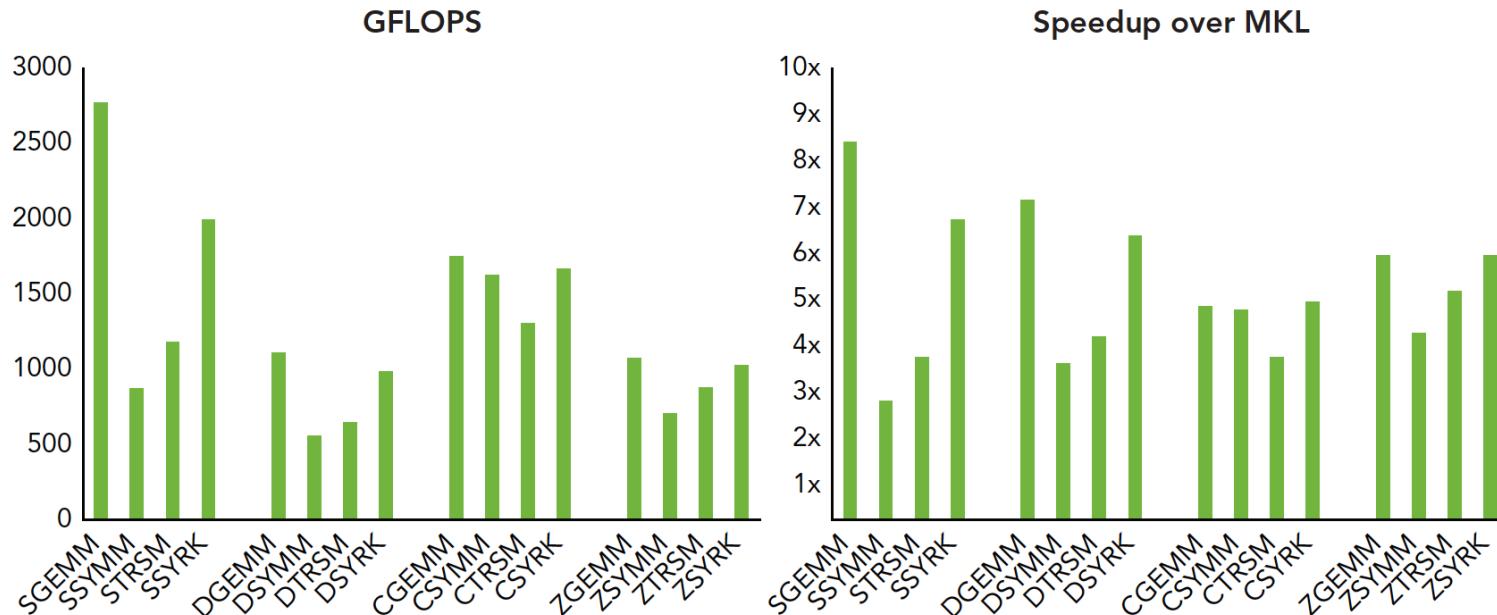
cuBLAS Portability

Porting to cuBLAS from BLAS is a straightforward process. In general, it requires:

- Adding device memory allocation/freeing (`cudaMalloc`, `cudaFree`)
- Adding device transfer functions (`cublasSetVector`,
`cublasSetMatrix`, etc)
- Transform library routine calls from BLAS to cuBLAS (e.g. `cbLAS_sgemv`
→ `cublasSgemv`)

cuBLAS Performance

cuBLAS Level 3: >1 TFLOPS double-precision



- MKL 10.3.6 on Intel SandyBridge E5-2687W @ 3.10GHz
- CUBLAS 5.0.30 on K20X, input and output data on device

Review – CUDA Libraries

CUDA Libraries offer high-level, high performance, domain-specific APIs for GPUs.

Deliberately designed to facilitate porting from legacy, host-only libraries.

| LIBRARY NAME | DOMAIN |
|---------------------------------|--|
| NVIDIA cuFFT | Fast Fourier Transforms |
| NVIDIA cuBLAS | Linear Algebra (BLAS Library) |
| CULA Tools | Linear Algebra |
| MAGMA | Next-gen Linear Algebra |
| IMSL Fortran Numerical Library | Mathematics and Statistics |
| NVIDIA cuSPARSE | Sparse Linear Algebra |
| NVIDIA CUSP | Sparse Linear Algebra and Graph Computations |
| AccelerEyes ArrayFire | Mathematics, Signal and Image Processing, and Statistics |
| NVIDIA cuRAND | Random Number Generation |
| NVIDIA NPP | Image and Signal Processing |
| NVIDIA CUDA Math Library | Mathematics |
| Thrust | Parallel Algorithms and Data Structures |
| HiPLAR | Linear Algebra in R |
| Geometry Performance Primitives | Computational Geometry |
| Paralution | Sparse Iterative Methods |
| AmgX | Core Solvers |

Alternatives to CUDA

OpenCL

A more open standard for parallel programming of a variety of processors.

- Supports GPUs, CPUs, FPGAs

Much of the OpenCL API was based on the CUDA APIs:

| CUDA API | OpenCL API |
|-----------------------|------------------------|
| cudaMalloc | clCreateBuffer |
| cudaFree | clReleaseMemObject |
| cudaMemcpyAsync | clEnqueueWriteBuffer |
| cudaMemcpyAsync | clEnqueueReadBuffer |
| kernel<<<...>>>(...); | clEnqueueNDRangeKernel |

OpenCL

| Pros | Cons |
|--|---|
| Portable API, run one kernel everywhere. | More verbose, explicit API |
| Open source | Weaker tooling than CUDA |
| Not tied to a single vendor | Smaller community than CUDA, less established |
| Collaboration among many companies, constantly pushing the standard forward. | |

OpenMP Accelerators



target

Create a device data environment and execute the construct on the same device.

| C/C++ |
|--|
| #pragma omp target [clause[.,] clause],... new-line parallel-loop-construct parallel-sections-construct |

| C/C++ |
|---|
| !\$omp target [clause[.,] clause],... parallel-loop-construct parallel-sections-construct !\$omp end target |

| Fortran |
|---------|
|---------|

Clauses

- device(*integer-expression*)
- map(*list*)
- mapto(*list*)
- mapfrom(*list*)
- scratch(*list*)
- num_threads(*list*)
- if(*scalar-expression*)

```
sum = 0;  
#pragma omp target device(acc0) map(B,C)  
#pragma omp parallel for reduction(+:sum)  
for (i=0; i<N; i++)  
    sum += B[i] * C[i]
```

OpenMP Accelerators

| Pros | Cons |
|---|--|
| High-level, easy to work with | Relatively new, implementations are still coming along, community is still growing, tooling is nonexistent |
| Supported by a large open-source consortium | Higher abstractions mean less control over optimizations |
| Single parallel programming model for host and device | |

OpenACC

Pragma-based (ala OpenMP) programming for GPUs, primarily supported by PGI (a.k.a. NVIDIA).

```
#pragma acc parallel loop gang deviceptr(d_A, d_B, d_C)
    for (i = 0; i < M; i++)
#pragma acc loop worker vector
    for (j = 0; j < P; j++) {
        float sum = 0.0f;

        for (k = 0; k < N; k++)
            sum += d_A[i * N + k] * d_B[k * P + j];

        d_C[i * P + j] = sum;
    }
```

OpenACC

| Pros | Cons |
|---|--|
| High-level, easy to work with | With OpenMP 4.0, future of OpenACC is unclear |
| Earliest example of pragma-based GPU programming, so likely the most stable | Higher abstractions mean less control over optimizations |
| Supported by NVIDIA/PGI | |

Trilinos Project

“The Trilinos Project is an effort to develop algorithms and enabling technologies within an object-oriented software framework for the solution of large-scale, complex multi-physics engineering and scientific problems.”

Over 50 packages: <https://trilinos.org/packages/>

- Includes packages for BLAS, Preconditioners, Linear Solvers, Nonlinear Solvers, Eigensolvers, Automatic Differentiation, Domain Decomposition, Partitioning, Mesh Manipulation, ...

Supports CUDA, OpenMP as backends

Provided free and open source by Sandia National Lab

Trilinos Project

| Pros | Cons |
|--|---|
| Like a super CUDA library | Long compile times due to heavy use of C++ templating |
| Open source and free, developed/ supported by the national labs | Performance can be hit or miss |

GPU Acceleration of the JVM

Many projects looking at JVM acceleration:

- APARAPI: <https://github.com/aparapi/aparapi>
- Rootbeer: <https://github.com/pcpratts/rootbeer1>
- SWAT: <https://github.com/agrippa/spark-swat>
- IBM's J9 JVM:
http://ahayashi.blogs.rice.edu/files/2013/07/IBM_Java8_GPU_PACT15_cameraready-sj0tik.pdf

| Pros | Cons |
|---|---|
| Broader set of GPU applications | Overheads from JIT-ing, serialization, etc may make speedup harder to achieve |
| High-level JVM programming languages | Most solutions are still in the research stage |
| Speedup relative to JVM-based apps may be > speedup relative to native apps | Little or no control over optimization |

HClib

Open source, unified, library-based programming system for x86 and GPU.

Lightweight work-stealing runtime.

finish-async, futures, parallel-for, atomics, accelerators, integration with MPI/
OpenSHMEM.

Emphasize portability from OpenMP, MPI, OpenSHMEM, other legacy APIs.

<http://hc.rice.edu/>

https://github.com/habanero-rice/hclib/tree/resource_workers

HClib

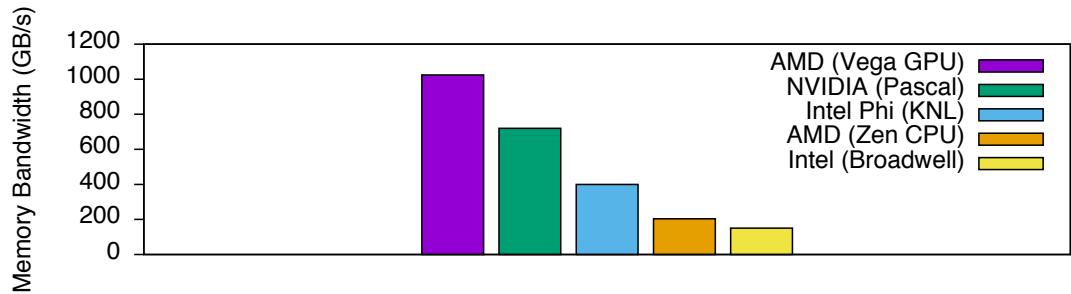
| Pros | Cons |
|---|---|
| Unified scheduling of x86, GPU, communication tasks | No backing by organizations other than Rice |
| High performance runtime | Only support C/C++ |
| Support for a wide range of parallel programming constructs | |
| Simplified porting from other programming models | |

<http://hc.rice.edu/>

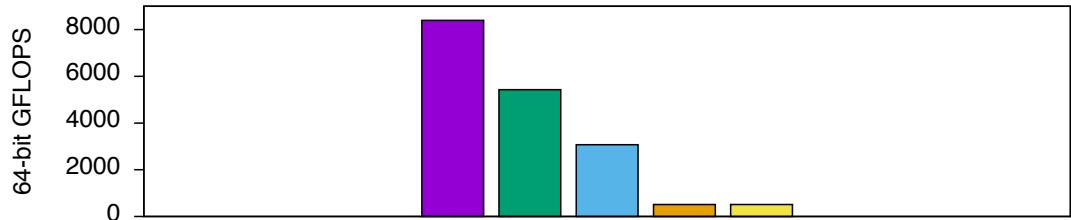
https://github.com/habanero-rice/hclib/tree/resource_workers

Alternatives to GPUs

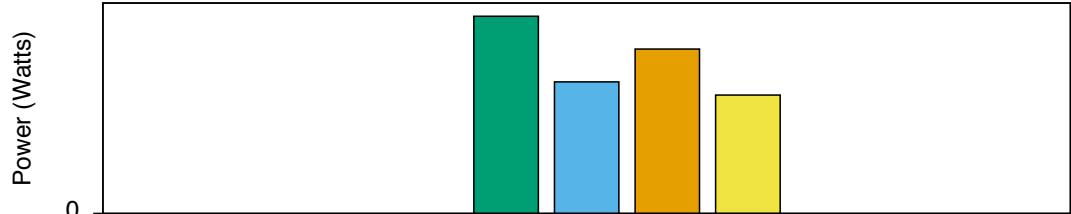
FPGA, KNL, x86



Is your hardware holding you back, or are you holding back your hardware?



GPUs have higher hardware peaks, but also have performance characteristics that you actually stand a chance of reasoning about.



Summary

Topics Not Covered

Multi-GPU CUDA Programming

Shared Bank Conflicts

CUDA-Aware MPI

Most of the CUDA libraries

Advanced Stream Usage

Dynamic Parallelism

Advanced Instruction Optimization

Warp Shuffle Instructions

Unified Memory

cuda-gdb

Texture Memory

...

Topics Covered

GPU Architecture

Optimizing CUDA Data Access

Alternatives to CUDA

nvprof

Hands On with a Non-Trivial CUDA

CUDA Execution Model

Kernel Synchronization

Application

Contact:

Max Grossman

jmaxg3@gmail.com

CUDA Asynchrony

CUDA Atomic Operations

CUDA Memory Hierarchy

CUDA Libraries