

# Geodesic Paths and Distances on Meshes: PDE and Computational Geometry Approaches

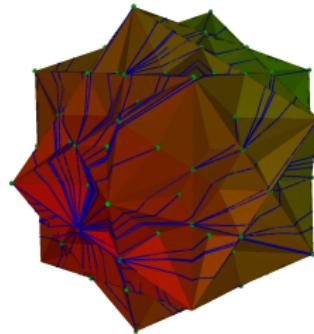
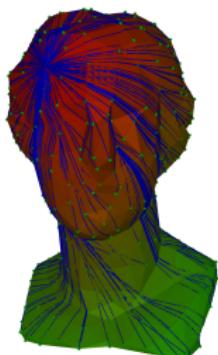
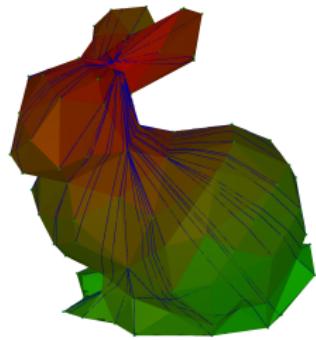
Matthieu Pierre Boyer & Antoine Groudiev

École Normale Supérieure

December 12, 2025

# Introduction

---



# Problem Positioning

---

Given a mesh of a 2-manifold embedded in  $\mathbb{R}^3$  and two points on the mesh, what is the shortest path between the points?

# Problem Positioning

---

Given a mesh of a 2-manifold embedded in  $\mathbb{R}^3$  and two points on the mesh, what is the shortest path between the points?

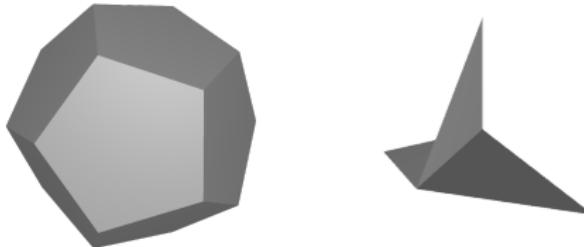
**2-Manifold:** a surface that looks locally like  $\mathbb{R}^2$ ;

# Problem Positioning

Given a mesh of a 2-manifold embedded in  $\mathbb{R}^3$  and two points on the mesh, what is the shortest path between the points?

**2-Manifold:** a surface that looks locally like  $\mathbb{R}^2$ ;

**Mesh:** a set  $\mathbb{V}$  of vertices, a set  $\mathcal{F}$  of faces in  $\mathbb{V}^3$ .



**Figure:** Examples of a 2-manifold (left) and non-manifold (right) mesh.

# Discrete Differential Geometry

The Laplace-Beltrami operator of a function  $u$  is given by:

$$(\Delta u)_i = \frac{1}{2A_i} \sum_{e=(i,j)} (\cot \alpha_{i,j} + \cot \beta_{i,j}) (u_i - u_j). \quad (1)$$

It quantifies the symmetric deviation of the variation of the value at a point.

# Physical Equations

---

We base ourselves on equations modelling the propagation of a phenomenon:

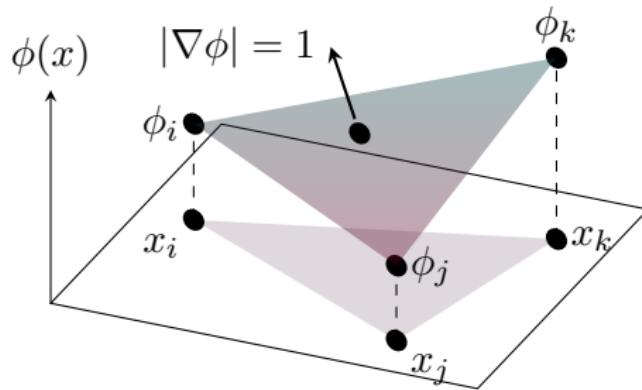
**Heat Equation**  $\delta u_t = \frac{d}{dt} u_0;$

**Poisson Equation**  $\Delta u = u_0$  for a fixed distribution  $u_0$ ;

**Spectral Embedding**  $\ell^2$  distance based on eigenspaces of the laplacian  $\Delta$ .

# Fast Marching

We compute a variation of Dijkstra's algorithm based on the Eikonal equation  $|\nabla u| = 1$  for wavefront propagation.



# Improved Chen-Han (ICH) Algorithm

**Idea:** Propagate **windows** accross edges of the mesh, encoding shortest paths from a source.

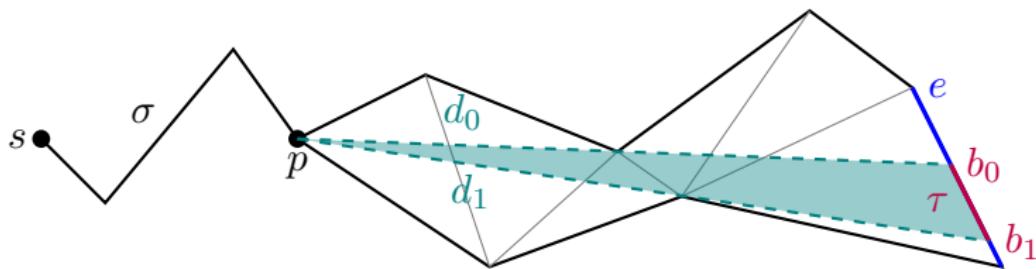
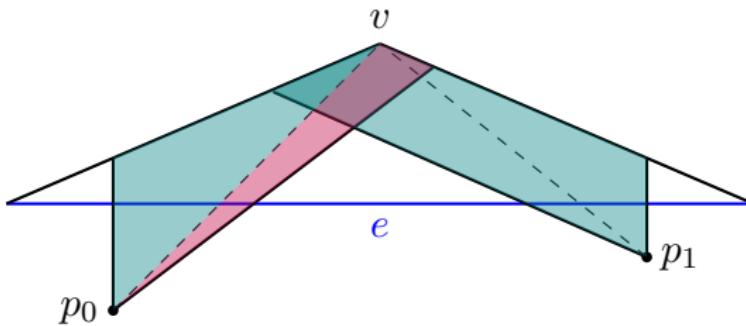


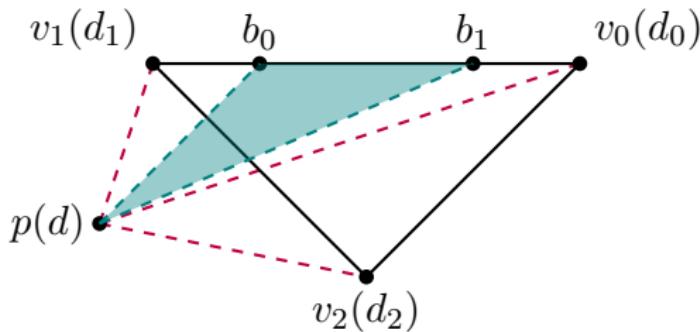
Figure: Illustration of a window  $w = (s, p, e, b_0, b_1, d_0, d_1, \sigma)$ .

# One Angle, One Split Rule



**Figure:** Illustration of the “one angle, one split” pruning rule: only keep three out of four windows.

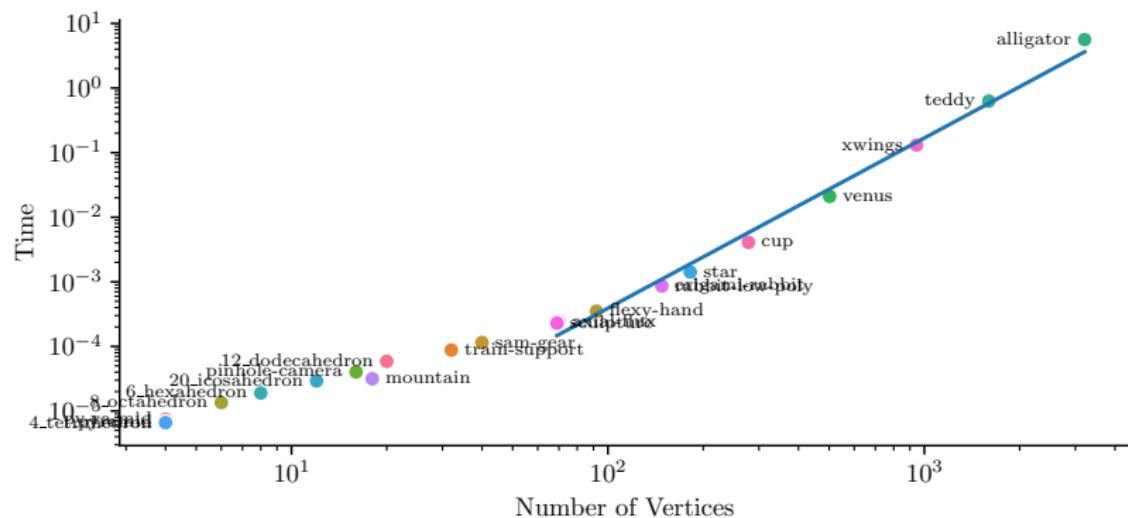
# Window Filtering



We can discard a window  $w$  if:

- $d + \|pb_0\| > d_0 + \|v_0b_1\|$
- $d + \|pb_1\| > d_1 + \|v_1b_1\|$
- $d + \|pb_0\| > d_2 + \|v_2b_0\|$

# Compute Time I



**Figure:** Benchmark of the *Heat method* as a function of the number of vertices. Slope: 2.63,  $r^2$ : 0.99.

# Compute Time II

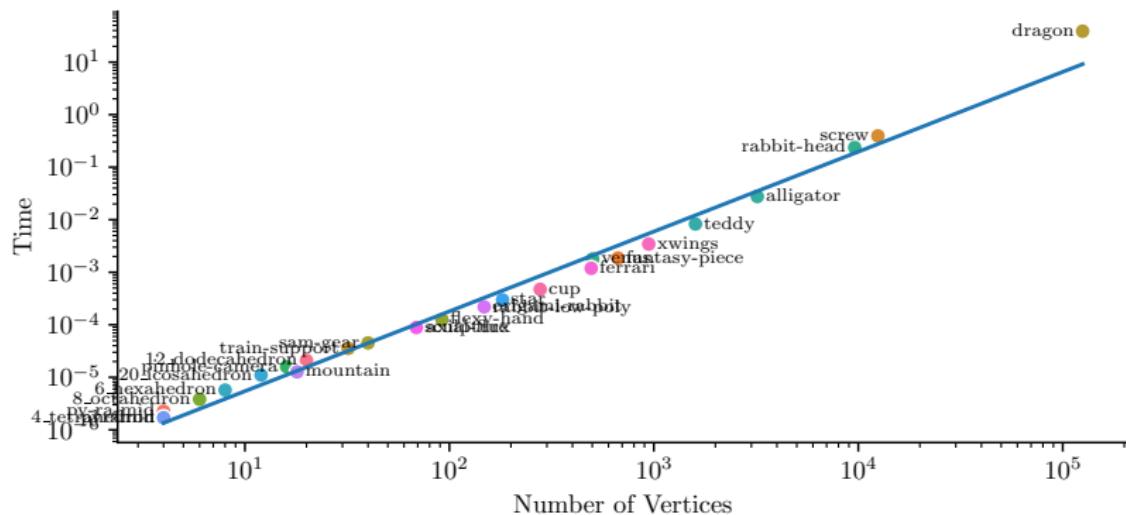


Figure: Benchmark of the *Fast Marching algorithm* as a function of the number of vertices. Slope: 1.51,  $r^2$ : 0.98.

# Compute Time III

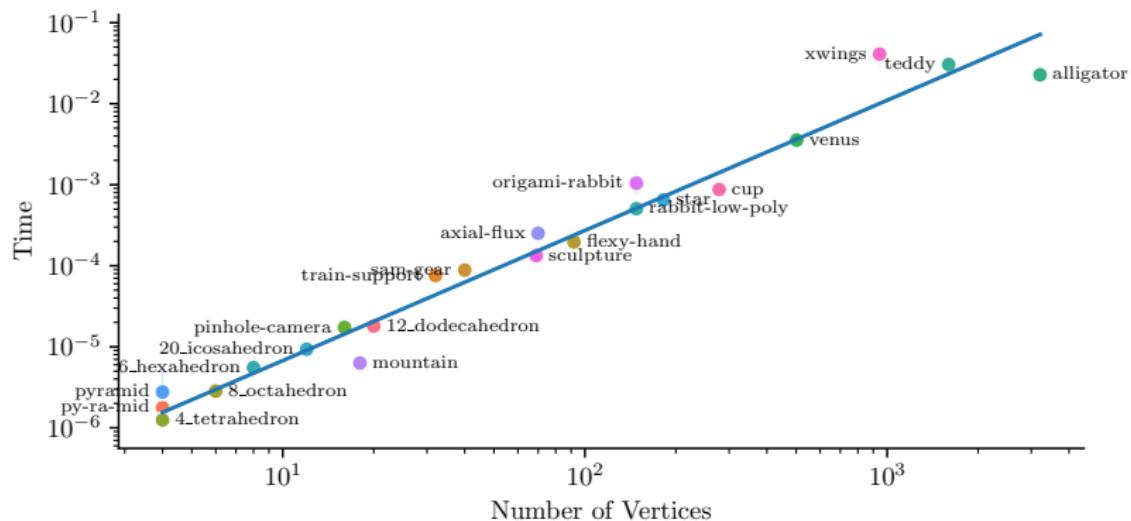


Figure: Benchmark of the *ICH* algorithm as a function of the number of vertices. Slope: 1.60,  $r^2$ : 0.96.

# Comparison of methods

Method	Theoretical Runtime	Actual Runtime	Precision	Limitations
<i>PDE-based methods</i>				
Heat	$\mathcal{O}(n_V^\omega)$	$\mathcal{O}(n_V^\omega)$ $n_V \rightarrow \infty$	Approx.	Stability, runtime, fixed sources, no paths
Poisson	$\mathcal{O}(n_V^\omega)$		Approx.	Stability, not true geodesics, runtime, fixed sources, no paths
Spectral	$\mathcal{O}(n_V^\omega)$		Approx.	Stability, not true geodesics, .

# Bibliography I

---