

## Tutorial 12

Semester 1, 2019

### Objectives of this tutorial

- To understand Binary Trees and Binary Search Trees.

### Exercise 1 \*

A binary expression tree is a binary tree used to represent algebraic expressions composed of unary and binary operators. The leaves of a binary expression tree are operands, such as constants or variable names, and the other nodes contain operations.

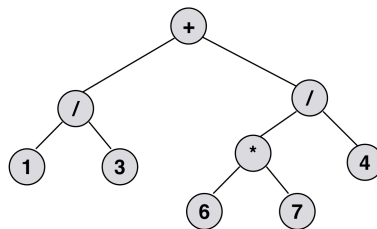


Figure 1: Expression tree

The prefix notation of an algebraic expression results from traversing the corresponding expression tree in pre-order. The infix notation results from traversing the tree in in-order; and the postfix notation (or reverse polish notation) results from traversing the tree in post-order.

Give the unambiguous mathematical expression as well as the prefix, infix and postfix notation of the expression represented by the tree above.

#### Solution

The expression is  $((1/3) + ((6 * 7)/4))$ .

- **Prefix:** + / 1 3 / \* 6 7 4
- **Infix:** 1 / 3 + 6 \* 7 / 4
- **Postfix:** 1 3 / 6 7 \* 4 / +

### Exercise 2 \*

Consider a `BinaryTree` class which defines a binary tree data type implemented using linked nodes, defined as follows:

```
1 class TreeNode:
2     def __init__(self, new_item=None, left=None, right=None):
3         self.item = new_item
4         self.left = left
5         self.right = right
6
7 class BinaryTree:
8     def __init__(self):
9         self.root = None
```

Write down an attribute method for the class that returns the height of the Binary Tree.

#### Solution

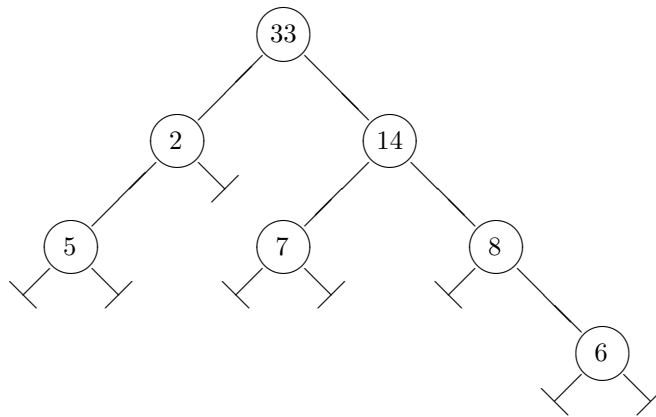
```

1 def __len__(self):
2     return self._aux_len(self.root)
3
4 def _aux_len(self, current):
5     if current is None:
6         return 0
7     else:
8         leftLen = self._aux_len(current.left)
9         rightLen = self._aux_len(current.right)
10        return 1 + max(leftLen, rightLen)

```

### Exercise 3 \*

Add to the class above the method `sum_leaves(self)` which returns 0 if the tree is empty and, otherwise, returns the result of adding the value of every leaf in the tree. For example, for `a.tree` of the form:



the result of `a.tree.sum_leaves()` would be  $5 + 7 + 6 = 18$ .

**Solution**

```

1 def sum_leaves(self):
2     return sum_leaves_aux(self.root);
3
4 def sum_leaves_aux(current):
5     if current is None:
6         return 0
7     elif current.left is None and current.right is None:
8         return current.item
9     else:
10        left = sum_leaves_aux(current.left)
11        right = sum_leaves_aux(current.right)
12        return left+right

```

### Exercise 4 \*

We want to extend the `BinarySearchTree` class defined in the lectures by adding a method `find_min()`. This method returns the minimum key in the tree, or `None` if the tree is empty. In doing so, it does not modify the tree. The following code shows two failed attempts at an implementation of such method:

```

1 def find_min_1(self):
2     return self.find_min_aux_1(self.root)
3
4 def find_min_aux_1(self, current):

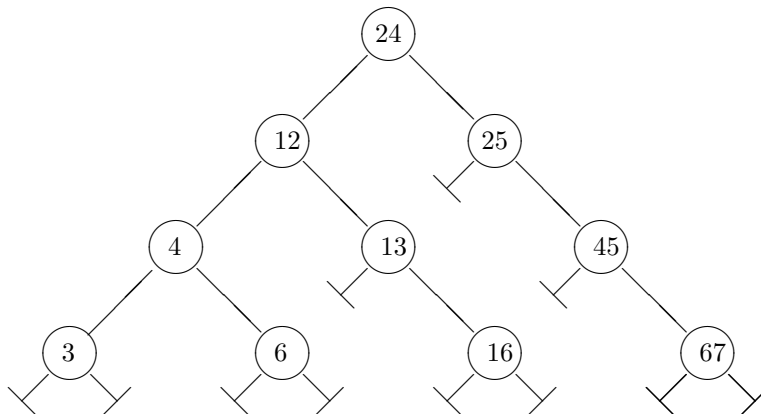
```

```

5     if current is not None:
6         return self.find_min_aux_1(current.left)
7     else:
8         return current
9
10    def find_min_2(self):
11        if self.root is None:
12            return self.root
13        else:
14            return self.find_min_aux_2(root)
15
16    def find_min_aux_2(self, current):
17        if current.left is not None:
18            return current
19        else:
20            return self.find_min_aux_2(current.left)

```

Consider a tree `the_tree` with integer keys with the form:



1. Show the value of result after calling `result = the_tree.find_min()` for each definition above.
2. Provide a correct definition for the above method.

### Solution

The first definition goes all the way down the leftmost branch and then returns `null`. The problem is that the condition for the if-then-else is `current is not None` rather than `current.left is not None`. Additionally, it returns `current`, rather than the key of the node.

The second definition stops right at the root node (since its left child is not empty) and thus returns a reference to the root node. The problem is that the THEN and ELSE branches of the if-then-else are swapped (or the condition negated). It also suffers from the same problem as the one before: it returns `current` rather than `current.key`.

A correct definition is:

```

1    def find_min(self):
2        if self.root is None:
3            return None
4        else:
5            return self.find_min_aux(self.root)
6
7    def find_min_aux(self, current):
8        if current.left is None:
9            return current.key
10       else:

```

```
11         return find_min_aux(current.left)
```

This method is linear (only one recursive call) and direct (calls itself). It is also tail recursive and can therefore be converted to a simple iteration without the need of a stack.

The value returned by this method for the tree given in the figure would be

## Exercise 5

Given a BST with numeric values and two numbers  $a$  and  $b$ , write down a function that returns a list with all items between  $a$  and  $b$ . The idea is to do this without visiting all elements, if possible.

**Solution**

```
1 def tree_range(self, a, b):
2     ans = []
3     self._range(a, b, self.root, ans)
4     return ans
5
6
7 def _range(self, a, b, current, a_list):
8     if current is None:
9         return
10    if a < current.key:
11        self._range(a, b, current.left, a_list)
12    if a <= current.key <= b:
13        a_list.append(current.key)
14    if current.key < b:
15        self._range(a, b, current.right, a_list)
```

## Exercise 6

Given a BST with numeric values and a number  $k > 0$ , write down an algorithm that returns the  $k$ -largest element in the BST. For  $k = 1$  it returns the largest element, for  $k = 2$  the second largest, and so on. The idea is to do this without visiting all elements, if possible.

**Solution**

```
1 def k_largest(self, k):
2     a_list = []
3     self._k_largest(self.root, k, a_list)
4     return a_list[-1]
5
6 def _k_largest(self, current, k, a_list):
7     if current is not None:
8         ans = self._k_largest(current.right, k, a_list)
9         if ans is not None:
10            return ans
11        a_list.append(current.key)
12        if len(a_list) == k:
13            return a_list
14        ans = self._k_largest(current.left, k, a_list)
15        if ans is not None:
16            return ans
```