Lecture 23 Recursive sorting I

FIT 1008&2085 Introduction to Computer Science



For a list of size N, the worst-case time complexity for insertion sort is...

- A) O(log N)
- B) O(N)
- C) O(N log N)
- D) $O(N^2)$

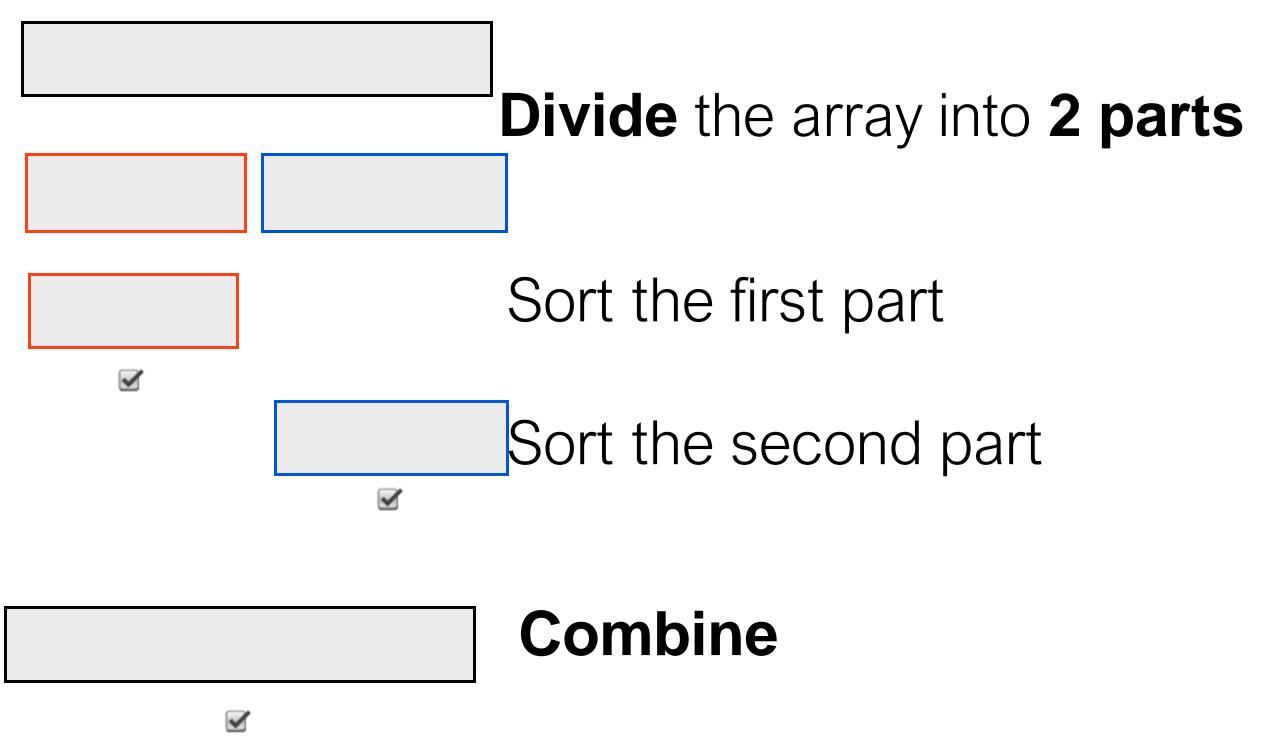
Can we do better?

Overview

- To review what a "divide and conquer" algorithm is
- To review in more depth two different "divide and conquer" sorting algorithms:
 - Merge Sort
 - Quick Sort

 To be able to implement them and compare their efficiency for different classes of inputs

Divide and Conquer: Sorting

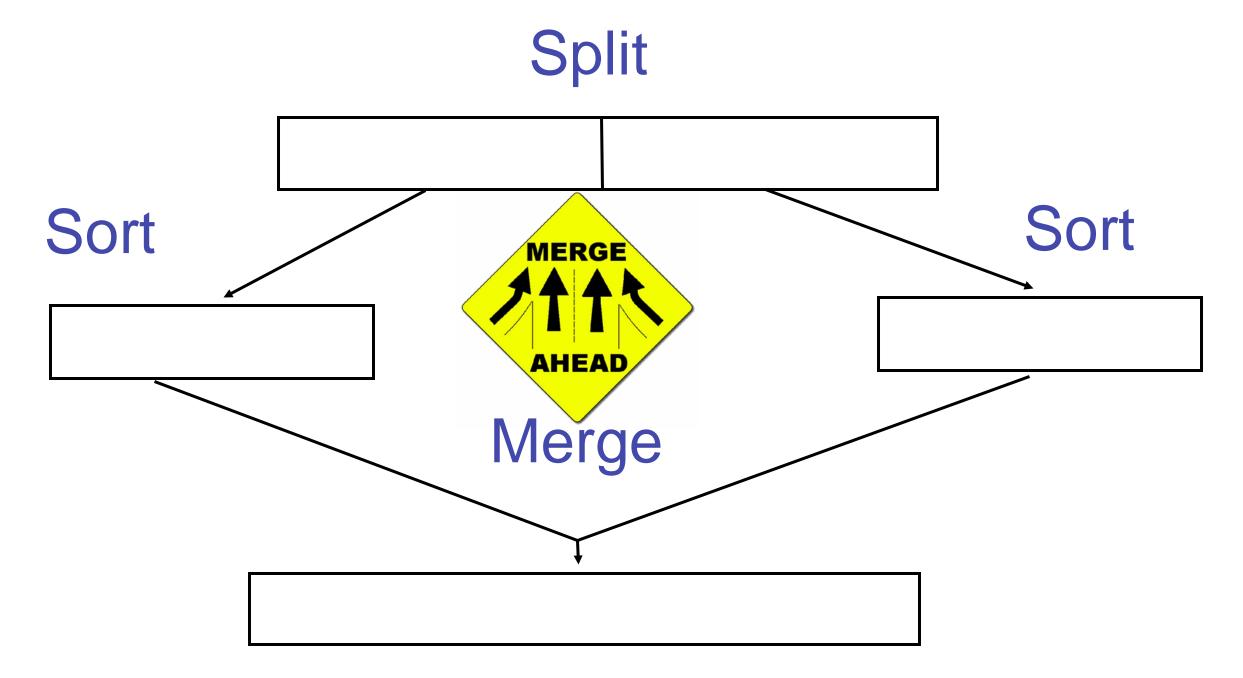


Divide and Conquer: Sorting

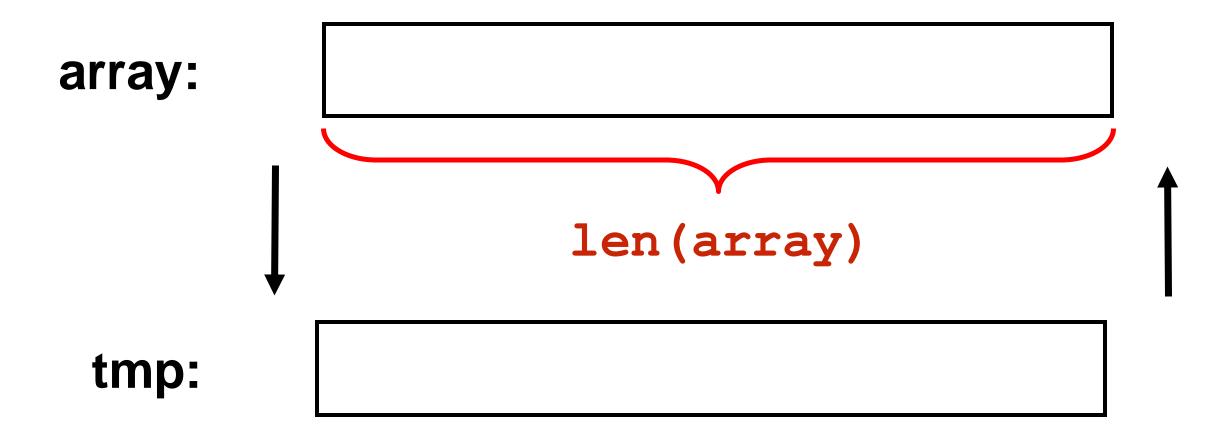
General Idea

```
def sort(array):
    if len(array) > 1:
        split(array, first_part, second_part)
        sort(first_part)
        sort(second_part)
        combine(first_part, second_part)
```

- Merge Sort has a <u>simple split</u> and a <u>elaborate combine</u>
- Quick Sort has a <u>elaborate split</u> and a <u>simple combine</u>



- **Split:** In half.
- Merge: Take two unsorted arrays, produce a sorted one.



Use a temporary array, then copy back to original array.

def merge_sort(array):

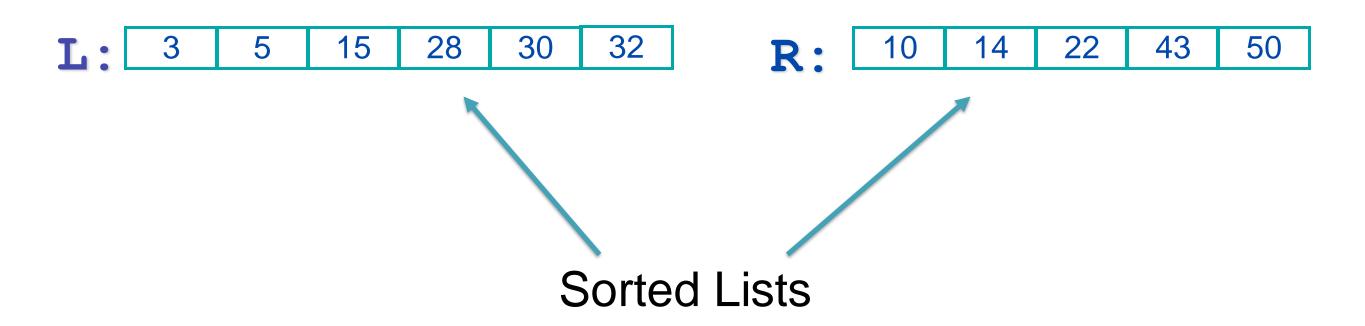
```
def merge_sort(array):
    tmp = [None] * len(array)
    start = 0
    end = len(array)-1
    merge_sort_aux(array, start, end, tmp)

    the array start index end index temporary array

def merge_sort_aux(array, start, end, tmp):
```

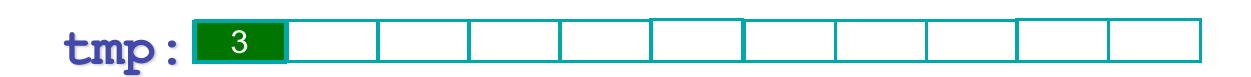
def merge_sort_aux(array, start, end, tmp):

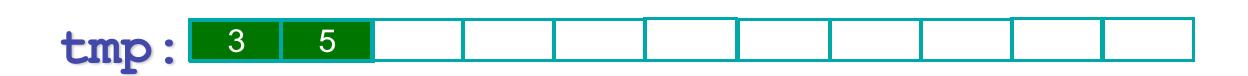
```
def merge_sort_aux(array, start, end, tmp):
    if start < end: # 2 or more still to sort</pre>
        mid = (start + end)//2
        # split into two halves
        merge_sort_aux(array, start, mid, tmp)
        merge_sort_aux(array, mid+1, end, tmp)
        # merge
        merge_arrays(array, start, mid, end, tmp)
        # copy tmp back into the original
        for i in range(start,end+1):
            array[i] = tmp[i]
```



start: 0, mid: 5, end: 10









```
      I:
      3
      5
      15
      28
      30
      32

      R:
      10
      14
      22
      43
      50

      i=2
      j=1
```



```
      I:
      3
      5
      15
      28
      30
      32

      R:
      10
      14
      22
      43
      50

      i=2
      j=2
```



```
      I:
      3
      5
      15
      28
      30
      32

      R:
      10
      14
      22
      43
      50
```



```
L: 3 5 15 28 30 32 R: 10 14 22 43 50
i=3
```



```
      I:
      3
      5
      15
      28
      30
      32

      R:
      10
      14
      22
      43
      50
```



















def merge_arrays(array, start, mid, end, tmp):

```
def merge_arrays(array, start, mid, end, tmp):
    i = start
    j = mid+1
    for k in range(start, end+1):
        if i > mid: # left finished, copy right
           tmp[k] = array[j]
            j += 1
        elif j > end: # right finished, Take remaining elements
            tmp[k] = array[i]
            i += 1
        elif array[i] <= array[j]: # array[i] is the item to copy</pre>
           tmp[k] = array[i]
            i += 1
        else:
            tmp[k] = array[j] # array[j] is the item to copy
            j += 1
Take the smaller from either list and advance
```

Invariant: After the kth iteration, the first k items in tmp are in order Invariant(2): After the kth iteration tmp holds the smallest k items from array 28

If N= end+1-start, what is the time complexity of merge_arrays?

```
def merge_arrays(array, start, mid, end, tmp):
    i = start
    j = mid+1
    for k in range(start, end+1):
        if i > mid: # left finished, copy right
            tmp[k] = array[j]
            j += 1
        elif j > end: # right finished, copy left
            tmp[k] = array[i]
            i += 1
        elif array[i] <= array[j]: # array[i] is the item to copy
        tmp[k] = array[j] # array[j] is the item to copy
        j += 1</pre>
```

- A) O(log N)
- B) O(N)
- C) O(N log N)
- D) $O(N^2)$

Merge Sort Analysis

- Natural: Typically the method that you would use when sorting a pile of books, CDs cards, etc.
- Most of the work is in the merging
- Uses more space than other sorts
- Close to optimal in number of comparisons. Good for languages where comparison is expensive.

If N=end+1-start, what is the (worst-case) time complexity of merge sort?

```
def merge_sort_aux(array, start, end, tmp):
    if start < end: # 2 or more still to sort
        mid = (start + end)//2

    # split into two halves
    merge_sort_aux(array, start, mid, tmp)
    merge_sort_aux(array, mid+1, end, tmp)

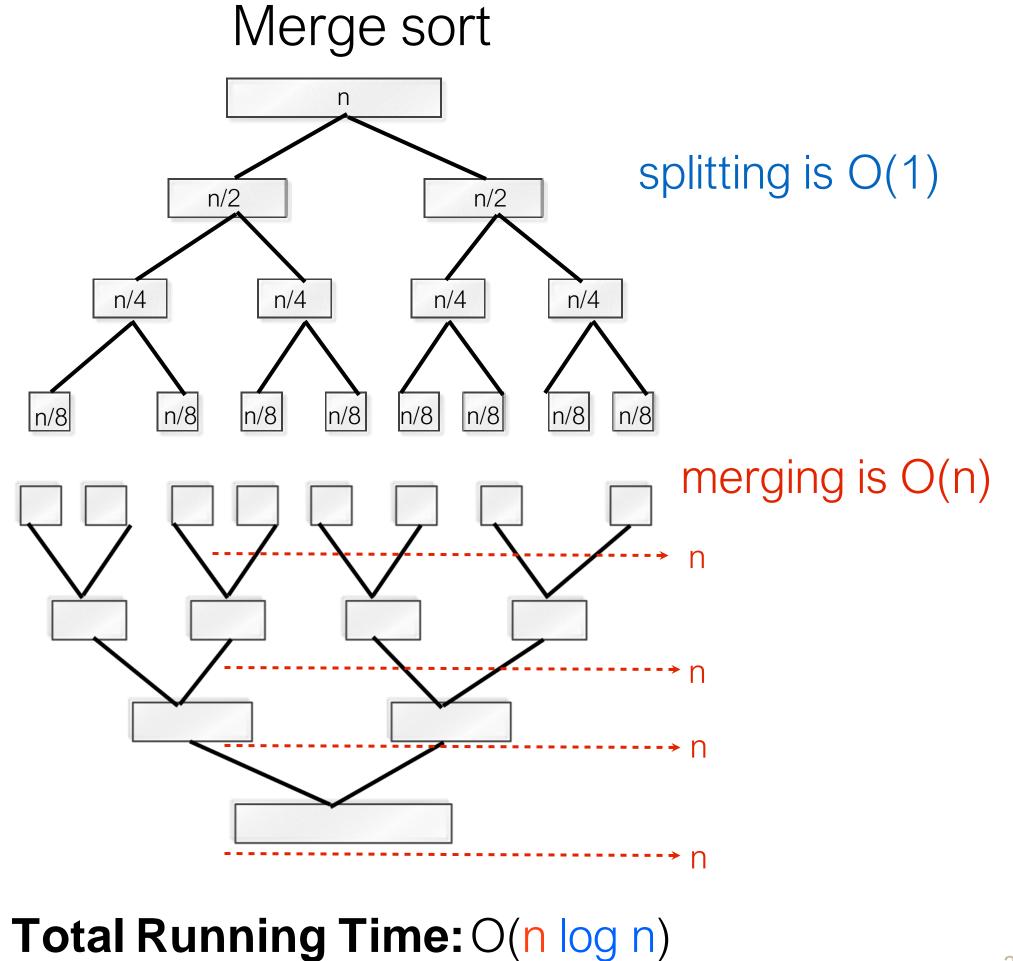
# merge
    merge_arrays(array, start, mid, end, tmp)

# copy tmp back into the original
    for i in range(start,end+1):
        array[i] = tmp[i]</pre>
```

- A) O(log N)
- B) O(N)
- C) O(N log N)
- D) $O(N^2)$

height is O(log n)

height is O(log n)



If N=end+1-start, what is the **best-case** time complexity of merge sort?

- A) O(log N)
- B) O(N)
- C) O(N log N)
- D) $O(N^2)$

Summary

	Best case	Worst case
Quicksort	O(?)	O(?)
Mergesort	O(n log n)	O(n log n)

Summary

Divide and Conquer and Recursive Algorithms (for sorting).

Merge Sort

- Easy: Split
- Elaborate: merge method