

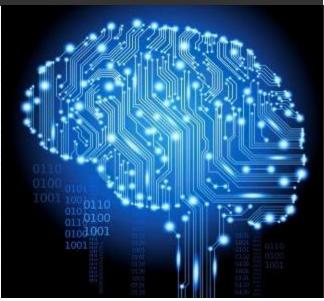
**Information Technology** 

## FIT1008&2085 Lecture 18

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# Queues with Arrays





#### Where were we at?

- Last lecture we saw the Stack abstract data type:
  - Main operations
  - Their complexity
- To be able to
  - Implement them with arrays, use them, modify their operations and reason about their complexity



#### Objectives for the this lecture

- To do the same for the Queue ADT
- To be able to decide when it is appropriate to use a Stack and a Queue



#### **FIFO**

- First In First Out (FIFO) processing:
  - The first element to arrive, is the first to be processed
- In terms of data storage:
  - The first element to be added, is the first to be deleted
- Again: access to any other element is unnecessary (and thus not allowed)



#### The Queue Data Type

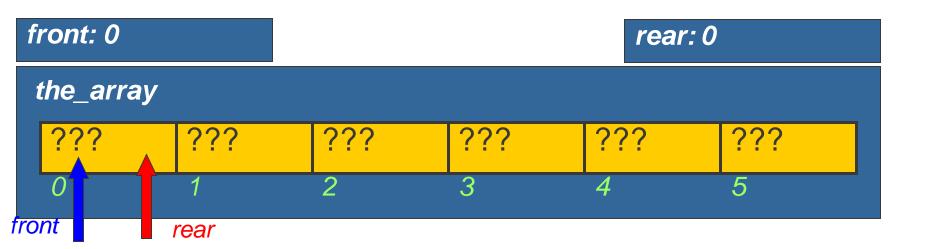
- Follows the FIFO model
- Its operations are defined in its interface:
  - Create the queue (Queue)
  - Add an item to the back (append)
  - Take an item off the front (serve)
  - Is the queue empty?
  - Is the queue full?
  - Empty the queue (reset)
- Remember: you can only access the element at the front of the queue (first item inserted that is still in)



- We need to:
  - Add items at the rear
  - Take items from the front
- A single marker is not going to be enough
- Lets try implementing queues using:
  - An array to store the items in the order they arrive.
  - An integer marking the front of the queue
    - Points to the first element to be served
  - An integer marking the rear of the queue
    - Points to the first empty cell at the rear
- Invariant: valid data appears in front .. rear-1

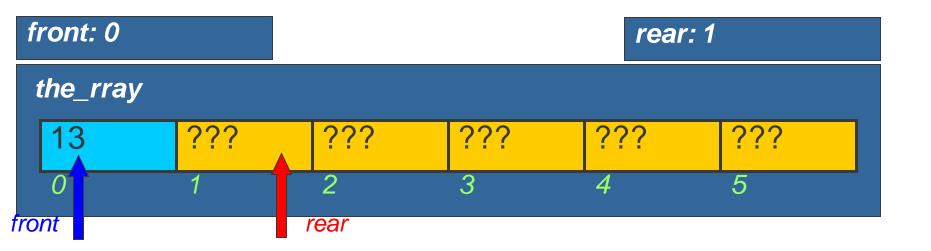


Create a new queue: initially no items



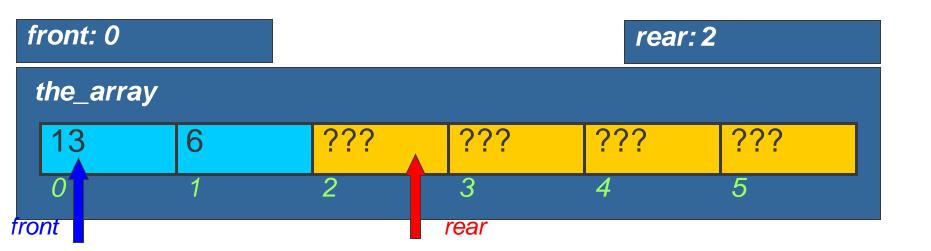


- Create a new queue: initially no items
- Append item 13





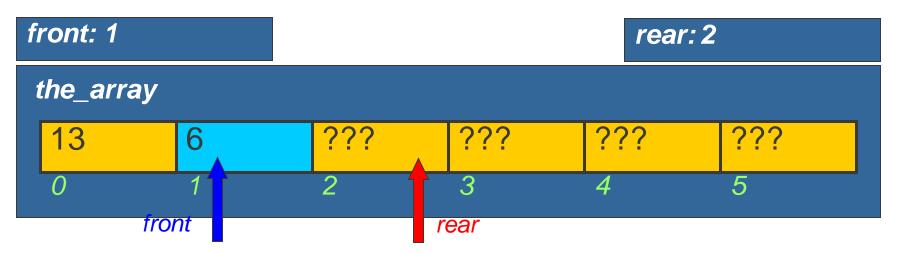
- Create a new queue: initially no items
- Append item 13
- Append item 6





- Create a new queue: initially no items
- Append item 13
- Append item 6
- Serve item

13





#### Implementation for a Linear Queue

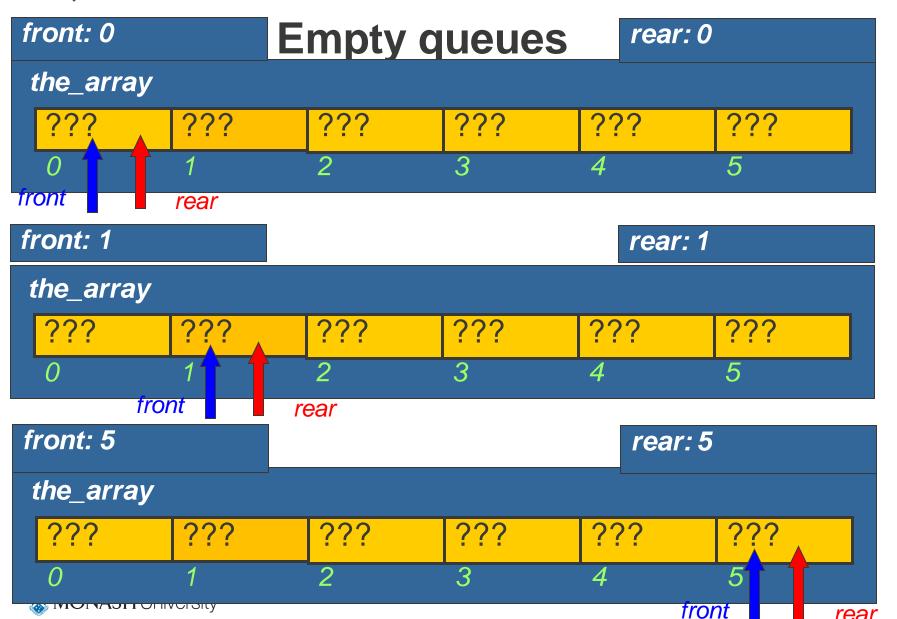
```
def Queue(max size):
    the array = [None] * max_size
    front = 0
    rear = 0
    return [front, rear, the array]
def size(the queue):
    [front, rear, ] = the queue
    return front - rear
def is empty(the queue):
    [front,rear,_] = the_queue
    return rear == front
```

front and rear must be equal

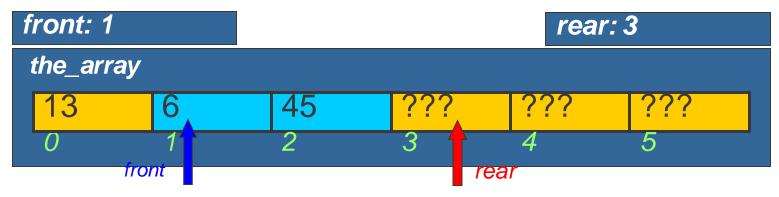
Why? Remember the invariant

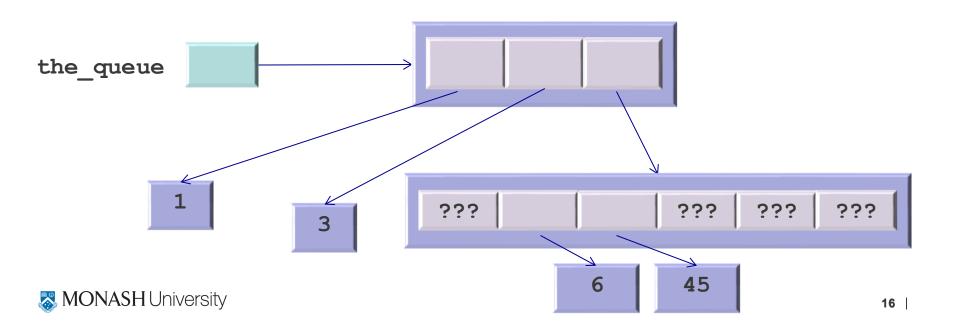
Big O?





## How is the queue represented in memory?





#### Implementation for a linear Queue

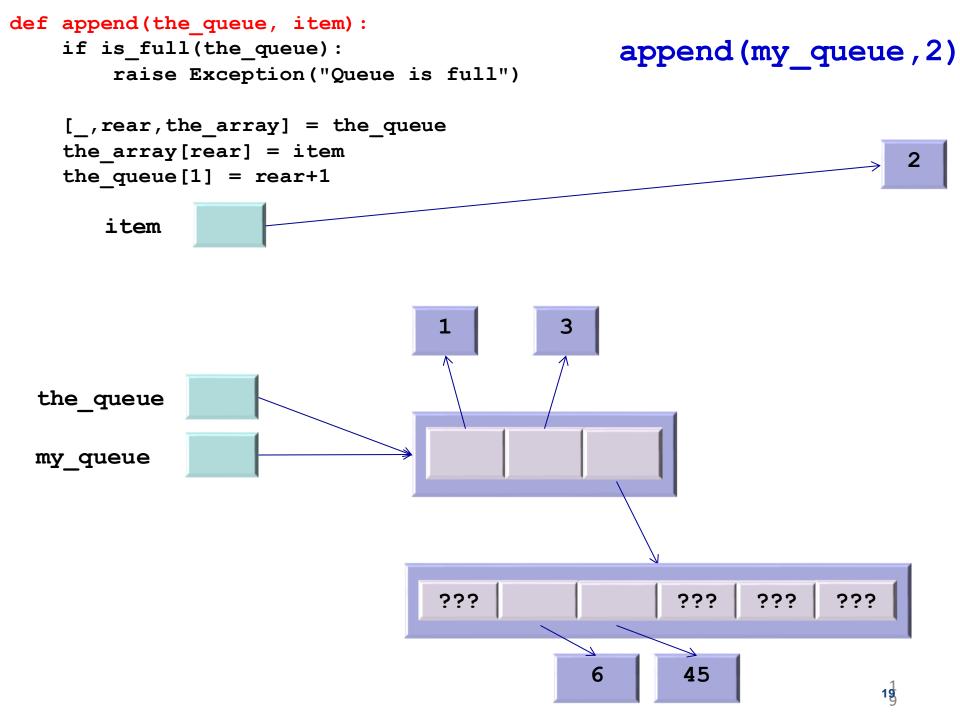
```
def is_full(the_queue):
    [_,rear,the_array] = the_queue
    return rear == len(the_array)

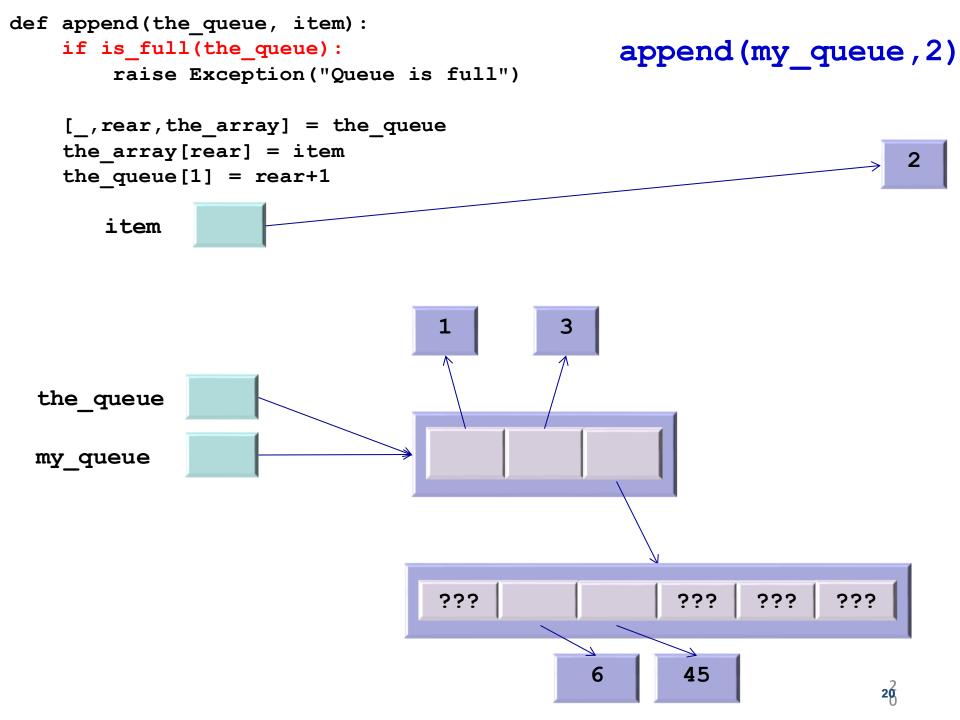
def append(the_queue, item):
    if is_full(the_queue):
        raise Exception("Queue is full")

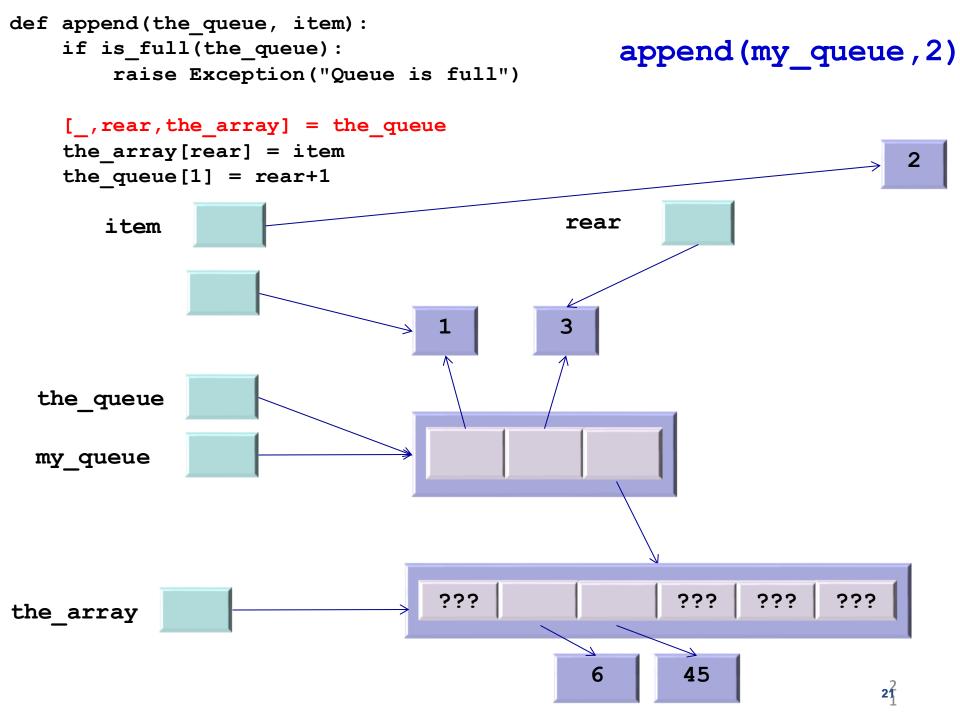
    [_,rear,the_array] = the_queue
    the_array[rear] = item
    the_queue[1] = rear+1
```

Big O?

```
def append(the queue, item):
    if is full(the queue):
                                              append (my_queue, 2)
       raise Exception("Queue is full")
    [_,rear,the_array] = the_queue
    the_array[rear] = item
                                                                     2
    the_queue[1] = rear+1
 my_queue
                                                  33.3 33.3 33.3
                                ???
                                            6
                                                   45
```







```
def append(the queue, item):
    if is full(the queue):
                                              append (my_queue, 2)
       raise Exception("Queue is full")
    [ ,rear,the array] = the queue
   the array[rear] = item
    the queue[1] = rear+1
                                          rear
       item
  the queue
 my queue
                                                       333 333
the array
                                            6
                                                   45
```

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    the queue[1] = rear+1
                                          rear
       item
  the queue
 my queue
                                                       333 333
the array
                                            6
                                                   45
```

#### Implementation for a linear Queue

```
def serve(the_queue):
    if is_empty(the_queue):
        raise Exception("Queue is empty")

    [front,_,the_array] = the_queue
    item = the_array[front]
    the_queue[0] = front+1
    return item

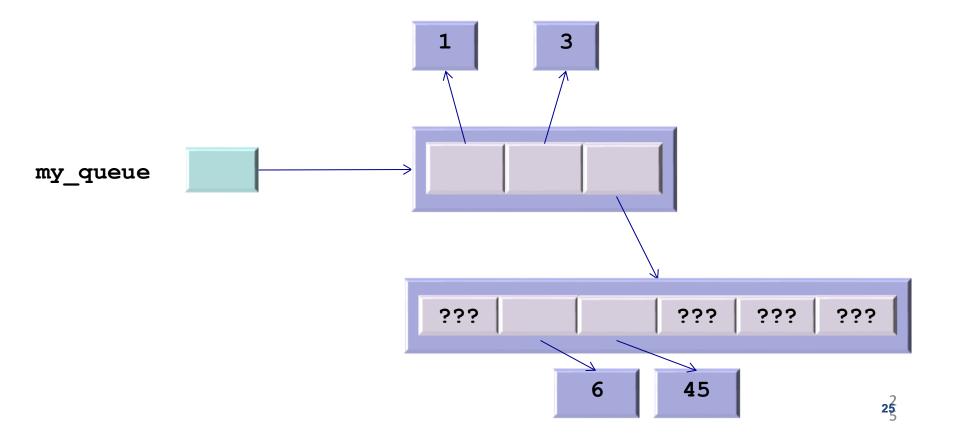
Again, front
    and rear
    must be equal

def reset(the_queue):
    the_queue[0] = 0
    the_queue[1] = 0
```



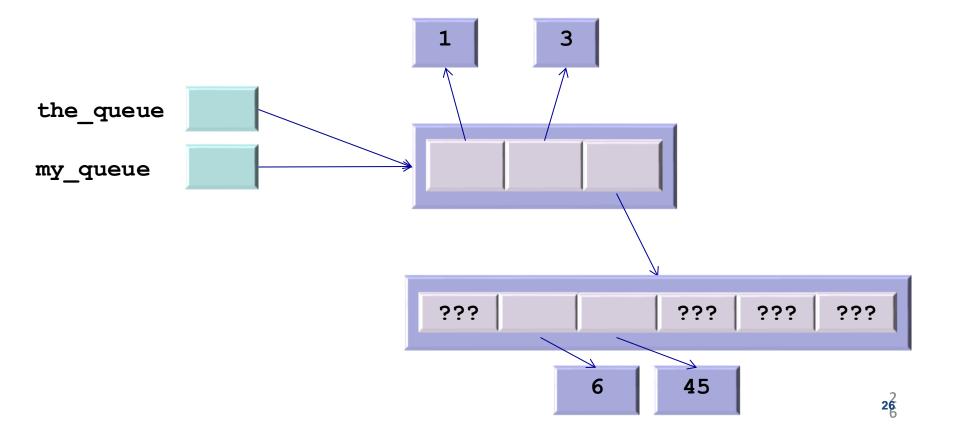
```
def serve(the_queue):
    if is_empty(the_queue):
        raise Exception("Queue is empty")

    [front,_,the_array] = the_queue
    item = the_array[front]
    the_queue[0] = front+1
    return item
```



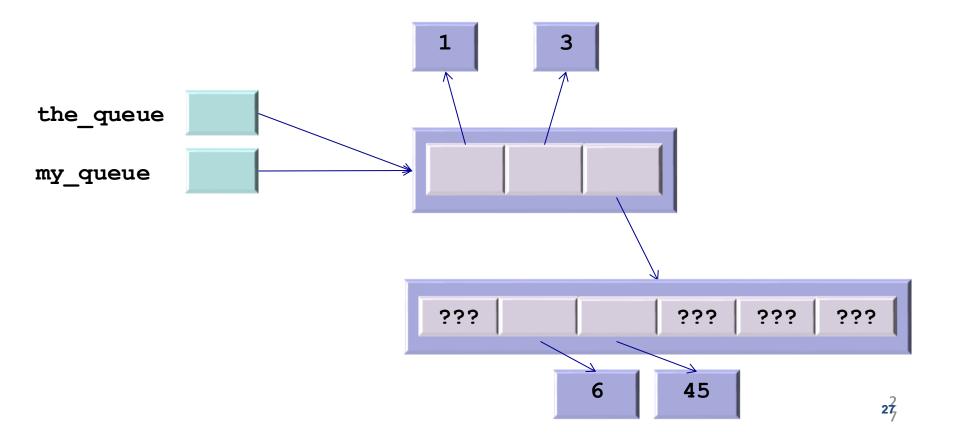
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def serve(the_queue):
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    [front,_,the_array] = the_queue
    item = the_array[front]
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    return item
```

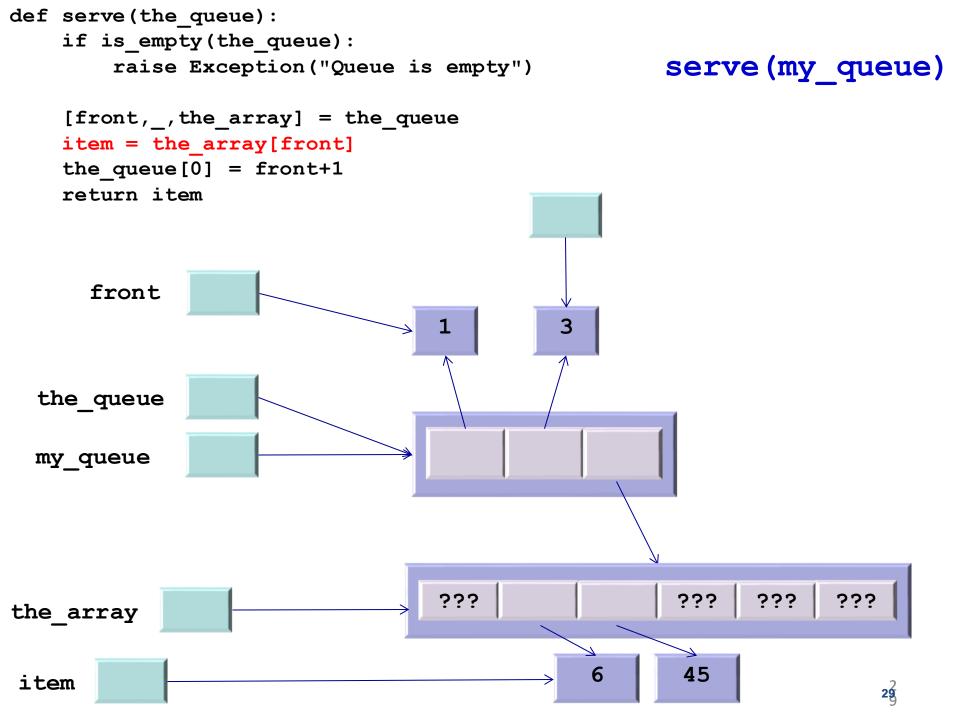


```
def serve(the_queue):
    if is_empty(the_queue):
        raise Exception("Queue is empty")

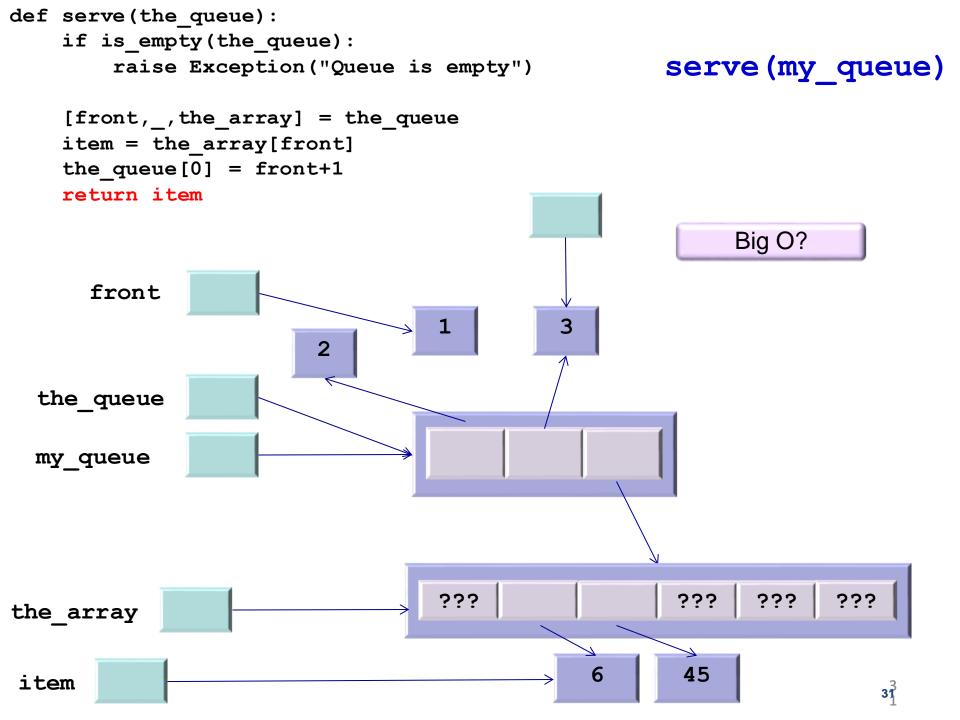
    [front,_,the_array] = the_queue
    item = the_array[front]
    the_queue[0] = front+1
    return item
```



```
def serve(the queue):
   if is_empty(the_queue):
                                              serve(my queue)
       raise Exception("Queue is empty")
   [front,_,the_array] = the_queue
   item = the array[front]
   the queue[0] = front+1
   return item
     front
 the queue
 my queue
                                               333
the_array
                                         6
                                                45
```



```
def serve(the queue):
    if is_empty(the_queue):
                                                  serve (my_queue)
       raise Exception("Queue is empty")
    [front,_,the_array] = the_queue
    item = the array[front]
    the_queue[0] = front+1
    return item
      front
  the queue
 my queue
                                                   33.3 33.3 33.3
                                333
the array
                                            6
                                                   45
item
```



#### Implementation problem

- When rear reaches the end of the array
- No more elements can be added
- Even if there is space left!

front: 3

the\_array

13 6 ???? 24 36 7

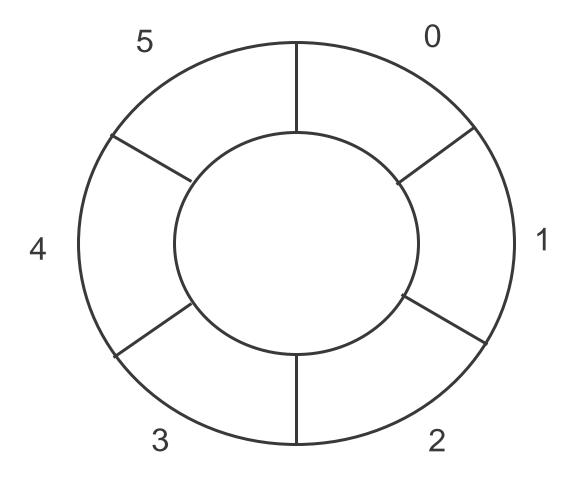
0 1 2 3 4 5

rear

Waste!

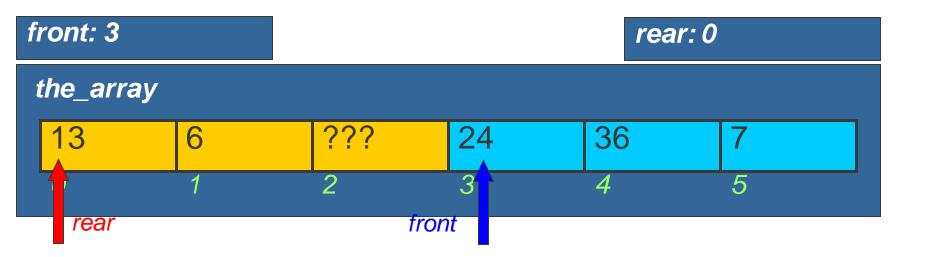


#### **Solution: Circular Queues**



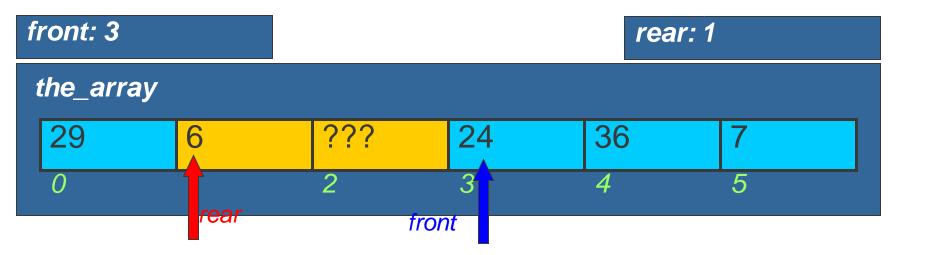
Simulated by allowing rear and front to wrap around each other

after appending 7 the situation would be



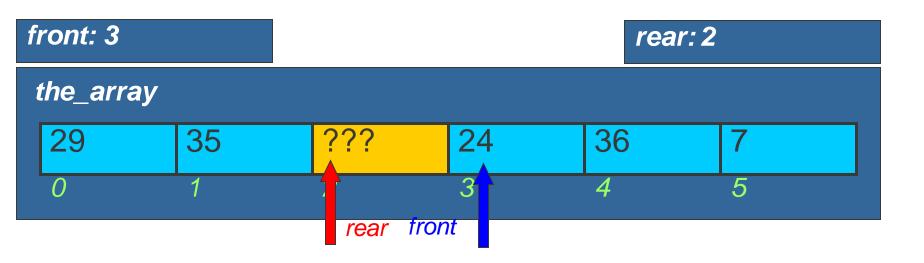


- after appending 7 the situation would be
- append item 29



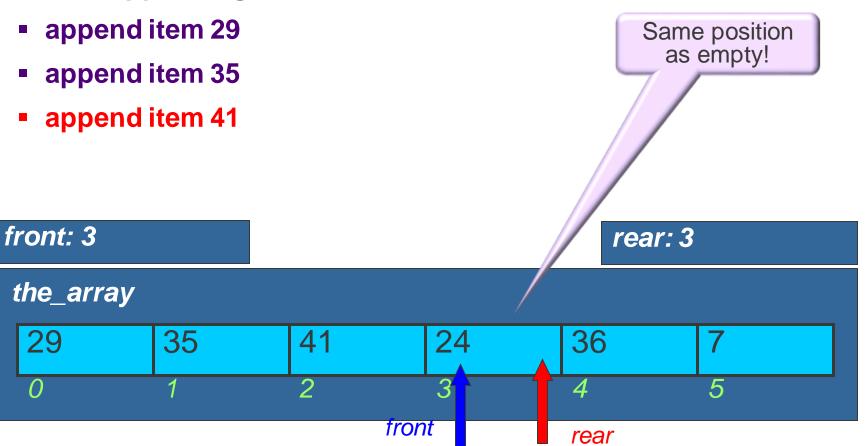


- after appending 7 the situation would be
- append item 29
- append item 35





after appending 7 the situation would be





- When front == rear, we need to know whether the queue is empty or full
- Some solutions:
  - Add a counter for number of items
  - Add one cell which is never used
- The former is simpler (to find the size of the queue...)



#### Implementation for a circular Queue

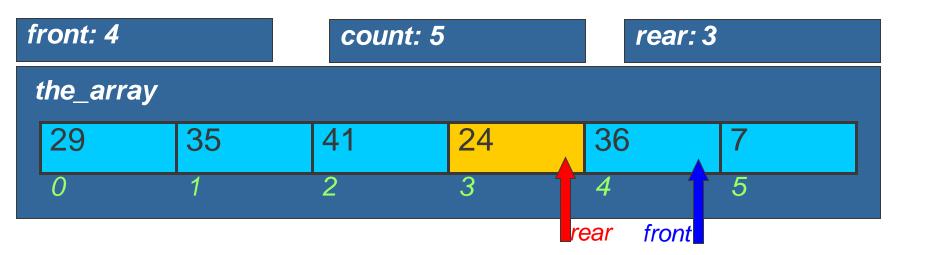
```
def Queue(max size):
    the array = [None] * max_size
    front = 0
    rear = 0
    count = 0
    return [front, rear, the array, count]
                                         Same as returning count
def size(the queue):
    return the queue[3]
                                              Or count == 0, that
def is empty(the queue):
                                                is, a very similar
    return size(the queue) == 0
                                               definition to the one
                                                 used for stacks
def is full(the queue):
    [ , ,the array,count] = the queue
    return count == len(the array)
                                               Big O?
```

#### Implementation for a circular Queue

Big O?

### Circular queue: both front and rear wrap

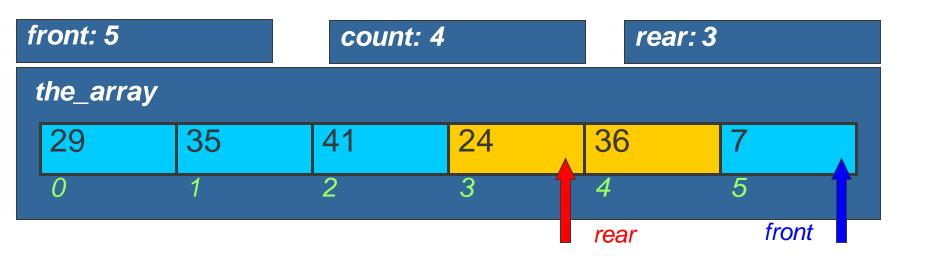
serve item (returns 24)





#### Circular queue: both front and rear wrap

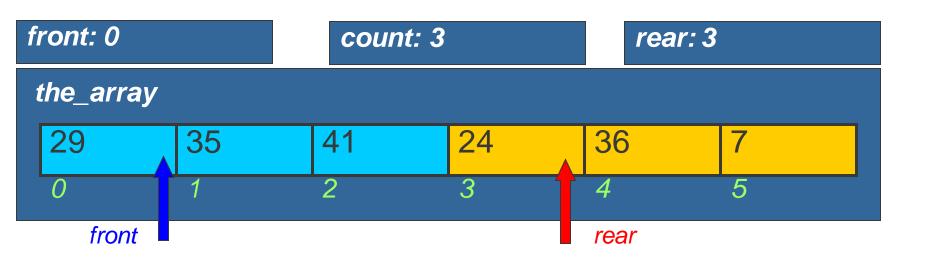
- serve item (returns 24)
- serve item (returns 36)





### Circular queue: both front and rear wrap

- serve item (returns 24)
- serve item (returns 36)
- serve item (returns 7)





#### Implementation for a circular Queue

Big O?

#### Example: print elements from front to rear

- Lets implement it as a function within the Queue ADT
  - So, it has access to the implementation
- Do not modify the queue, just print its elements



### **Complexity of circular Queue operations**

- Single loop
- Loop always executed count times
  - Best = worst
- Inside the loop, the number of operations is fixed except for ...
  - print
  - Their number of operations depends on the size m of the item,
     e.g., the length of a string, the number of integers in a matrix, etc
- count\*(K\*m) ≈ count\*m
- Which gives best = worst = O(count\*m)

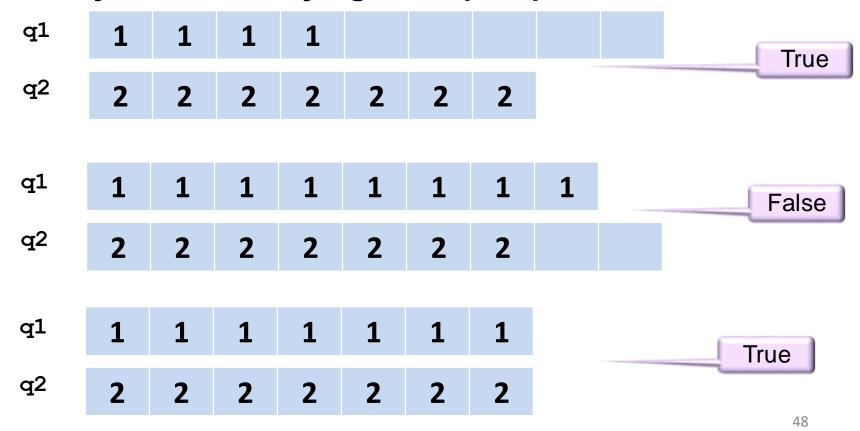
```
def print_items(the_queue):
    [front,_,the_array,count] = the_queue
    index = front
    for _ in range(count):
        print(the_array[index])
        index = (index+1)%len(the array)
```

def greater(q1,q2)

#### Returns true if and only if

- q2 is at least as long as q1
- AND every element in q1 is less or equal than the element in the same position in q2

#### Don't worry about modifying the input queues



Make sure you use it as an ADT (no access to the implementation)

Try first using only is\_empty and serve

```
def greater(q1,q2):
    while not is_empty(q1) and not is_empty(q2):
        if serve(q1) > serve(q2):
            return False
    return is empty(q1)
```

Try now using size

#### **Some Queue Applications**

- Scheduling and buffering
  - Printers
  - Keyboards
  - Executing asynchronous procedure calls



## Summary

#### Queues

- Array implementation
- Basic operations
- Their complexity