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**5GTRANSFORMER**

Orchestrating a brighter world

**NEC**

# Overbooking Network Slices through Yield-driven End-to-End Orchestration

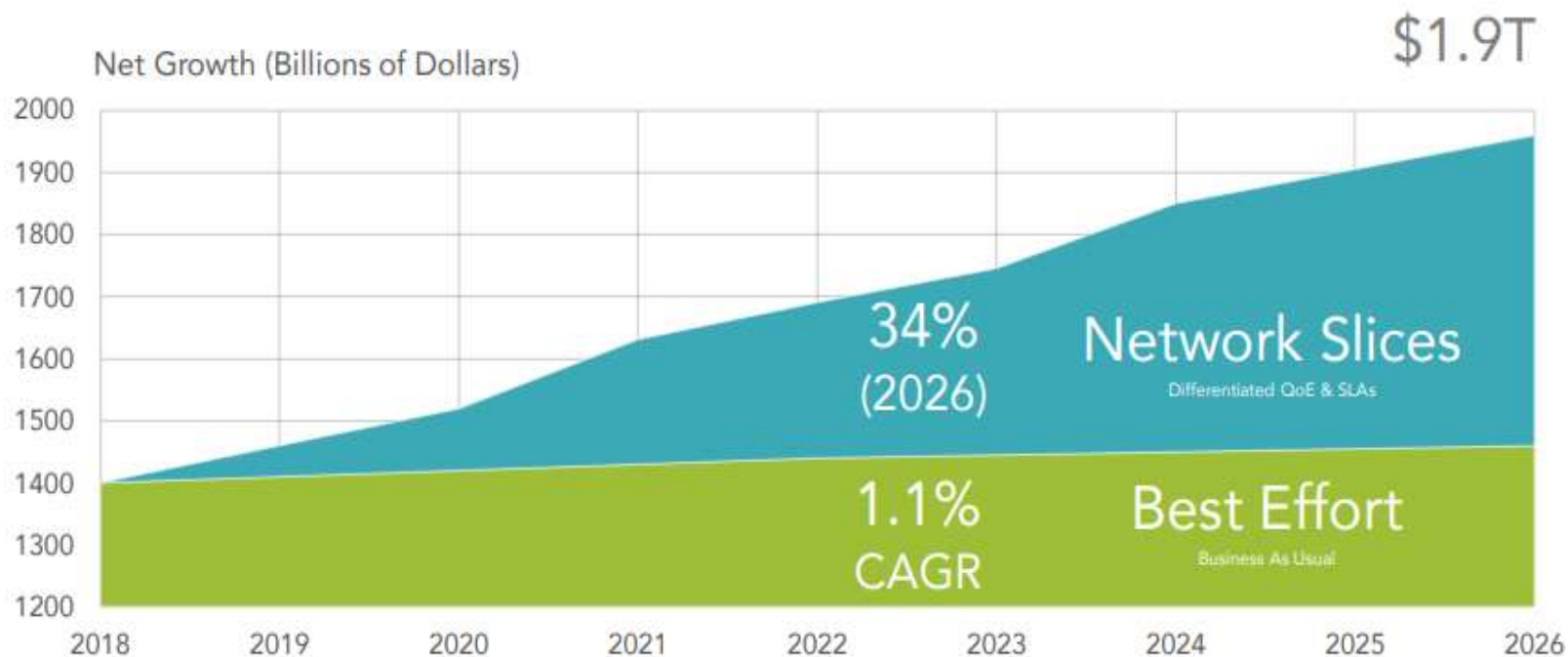
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Vincenzo Sciancalepore  
Xavier Costa-Perez

**NEC Laboratories Europe**

# Introduction: Network Slicing

- (possibly hyper-hyped but...) Telcos and vendors expect that **NS to unlock around \$300bn in business opportunities** with verticals/private enterprises

34% GROWTH TO 2026 FROM MOVING OFF "BEST EFFORT"

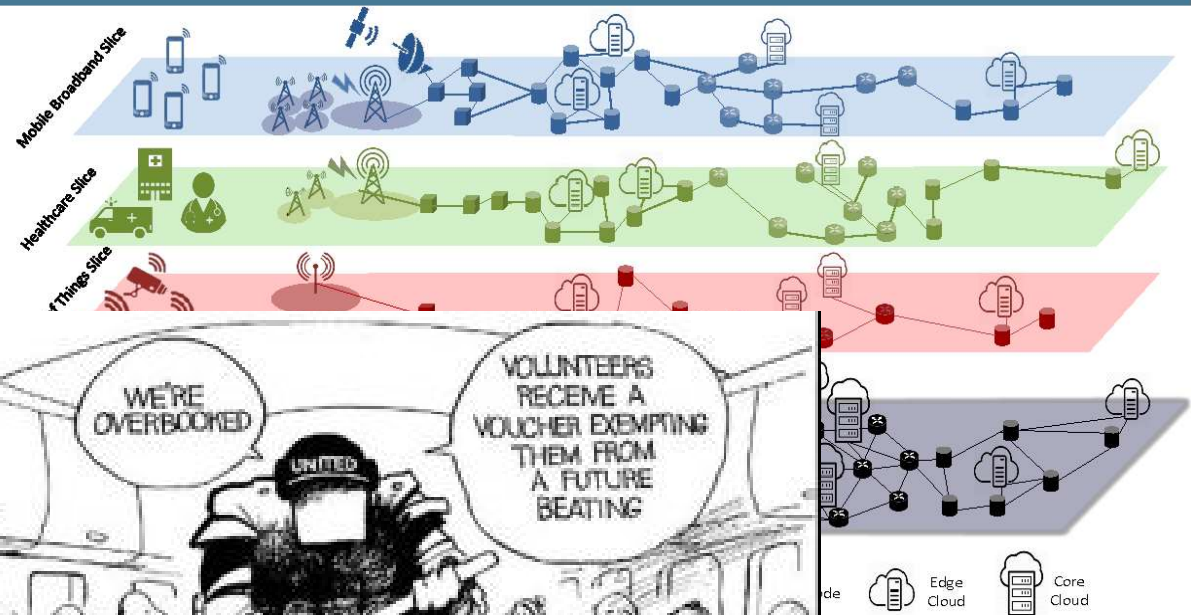


Source: Ericsson, Arthur D. Little

# Challenges and Contributions

## 1. End-to-End Orchestration

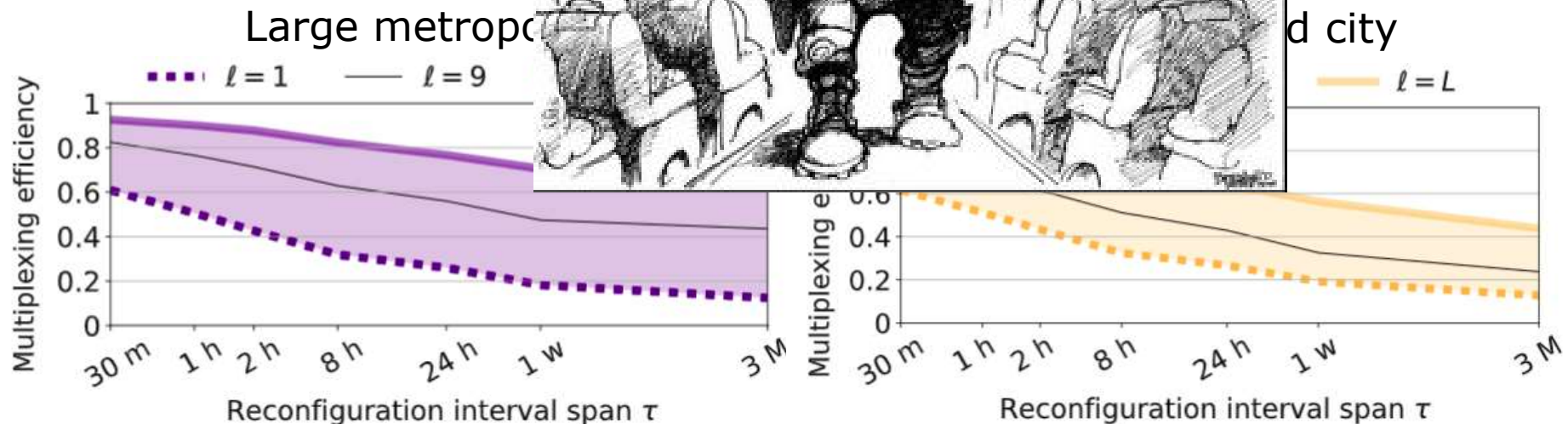
## 2. Network Slice Overbooking (Admission Control and Resource Reservation mechanism)



NFV: Concepts, Architectures, pp. 80-87, May 2017.

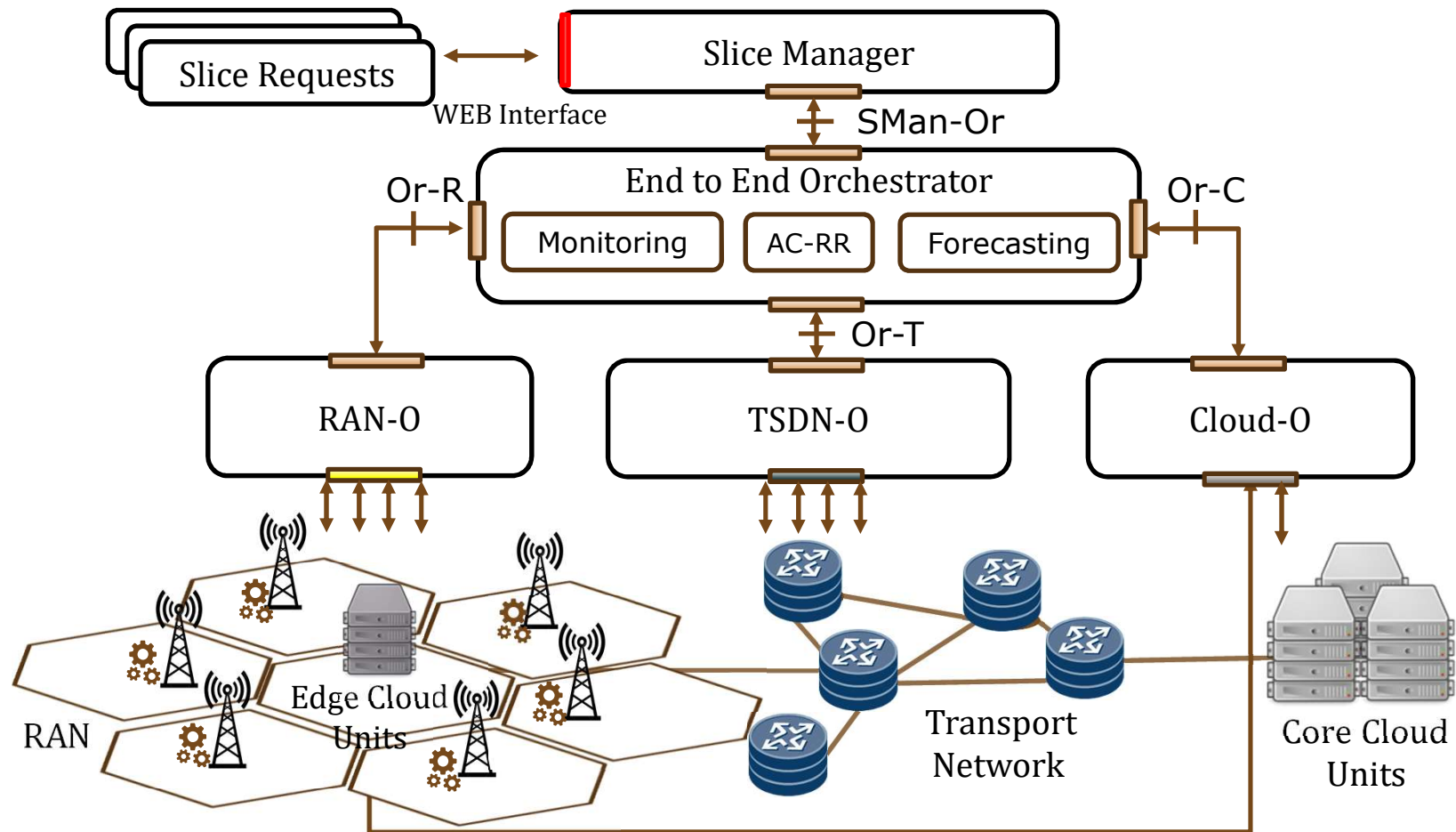
Large metropol...

city

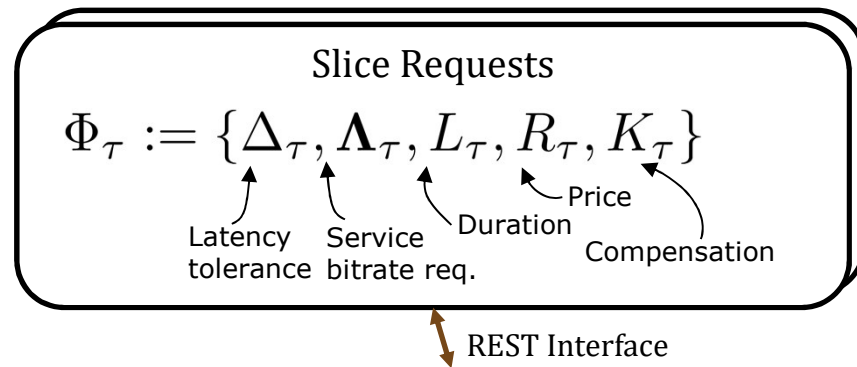


C. Marquez et al. "How Should I Slice My Network?: A Multi-Service Empirical Evaluation of Resource Sharing Efficiency". In ACM MobiCom 2018

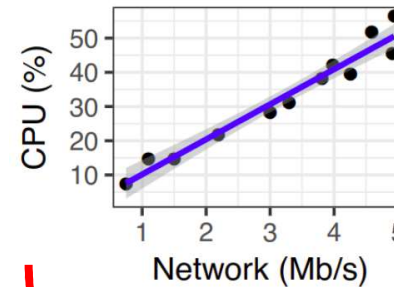
# System Design and Model



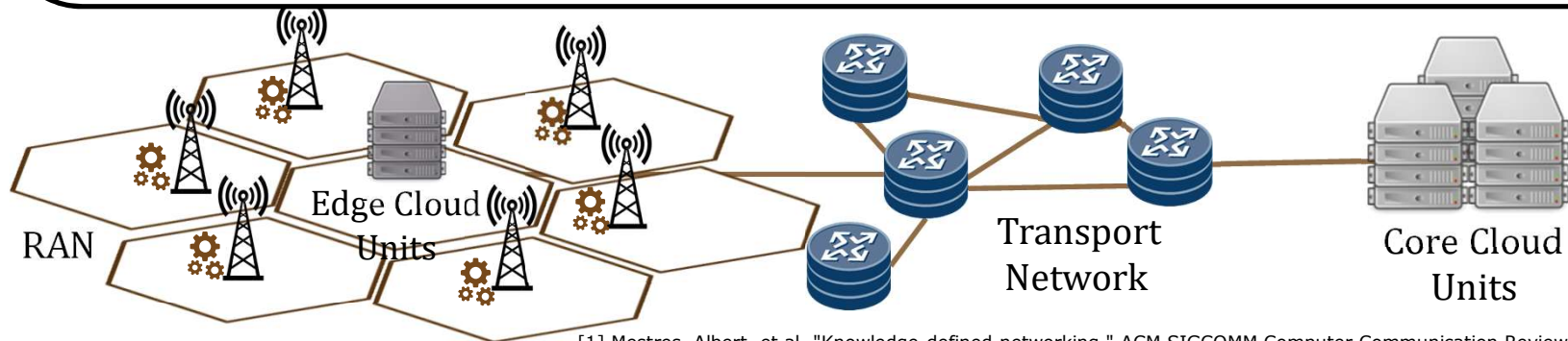
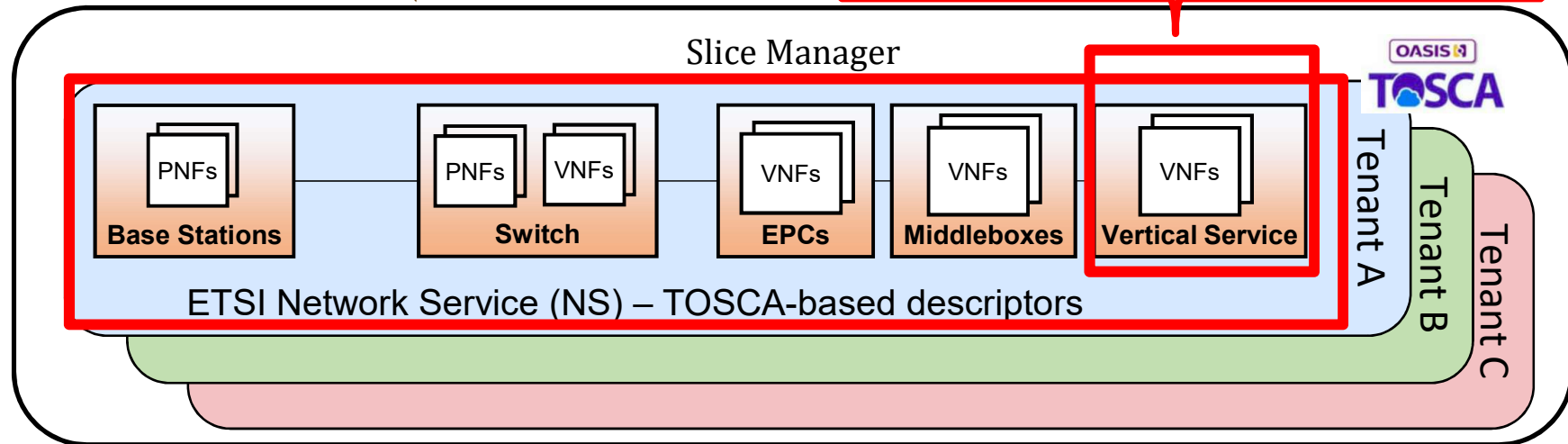
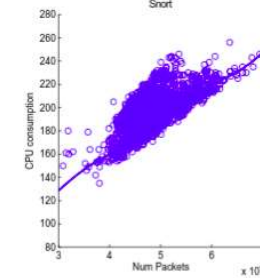
# System Design and Model: NS Descriptors



e.g., OpenFace:



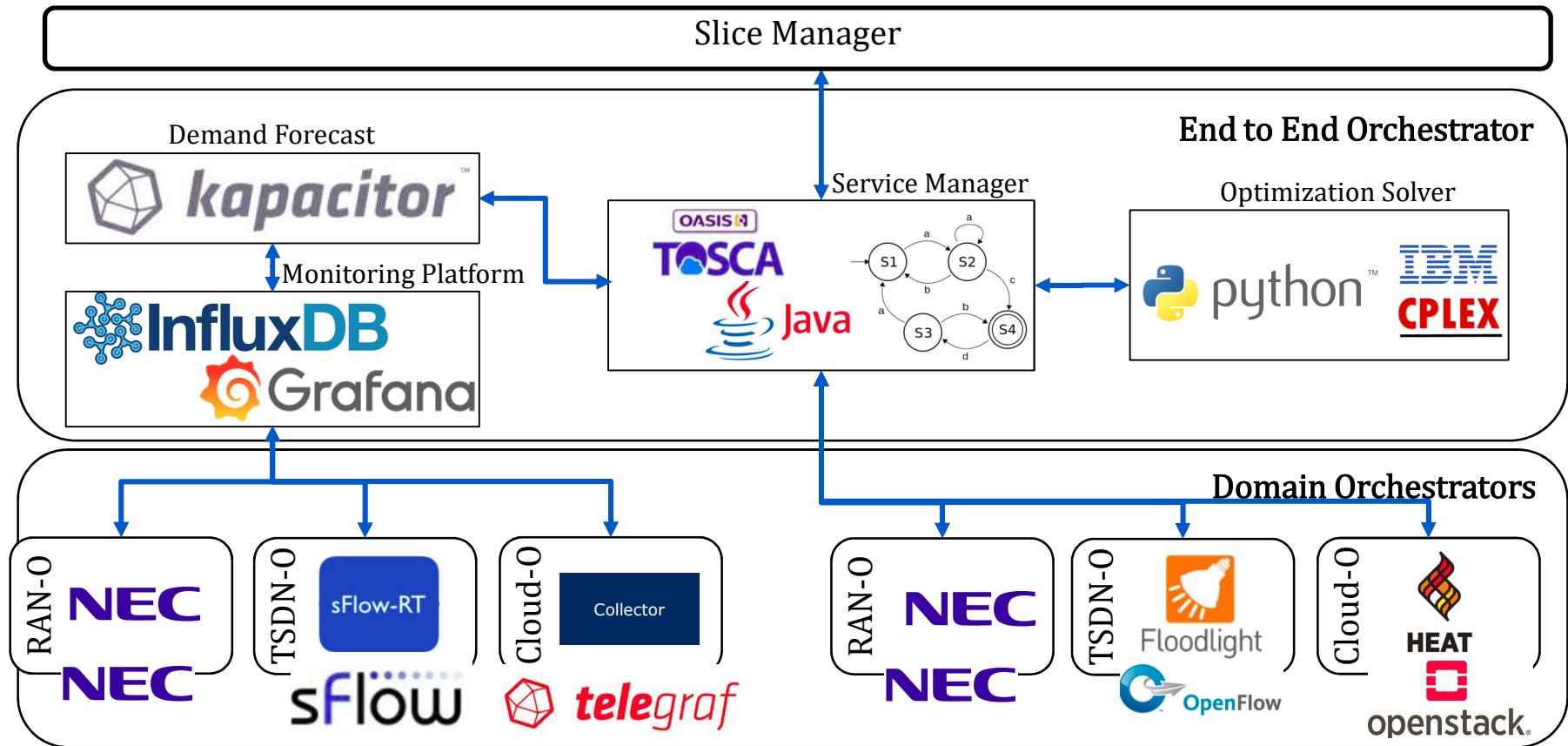
e.g., snort [1]



[1] Mestres, Albert, et al. "Knowledge-defined networking." ACM SIGCOMM Computer Communication Review 47.3 (2017): 2-10.

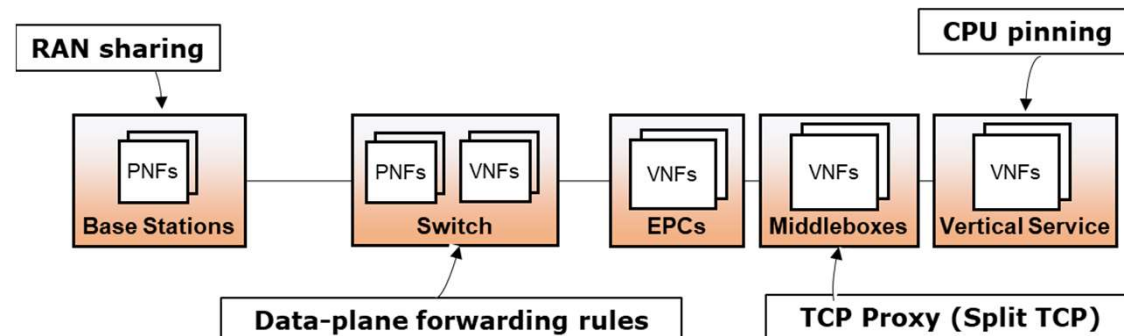


# System Design and Model: E2E Orchestrator



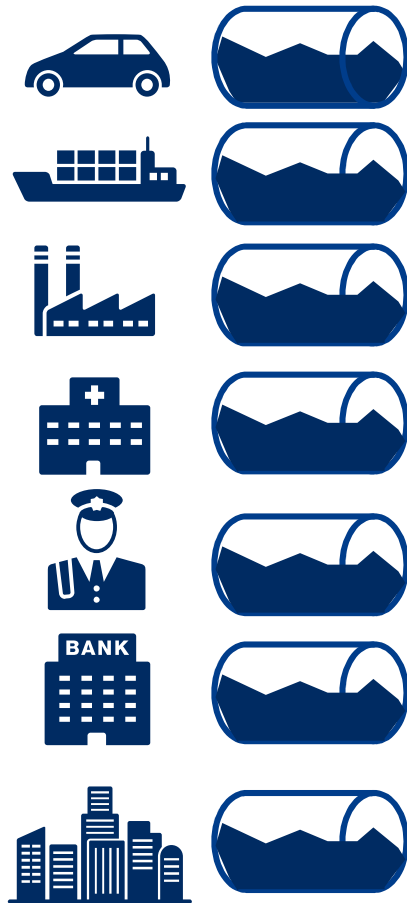
## “Local” orchestrators

1. Create data-plane abstractions
2. Monitoring data-plane usage
3. Enforce orchestration decisions

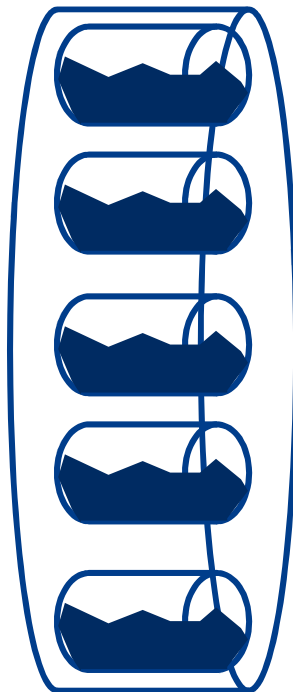


# System Design: Overbooking

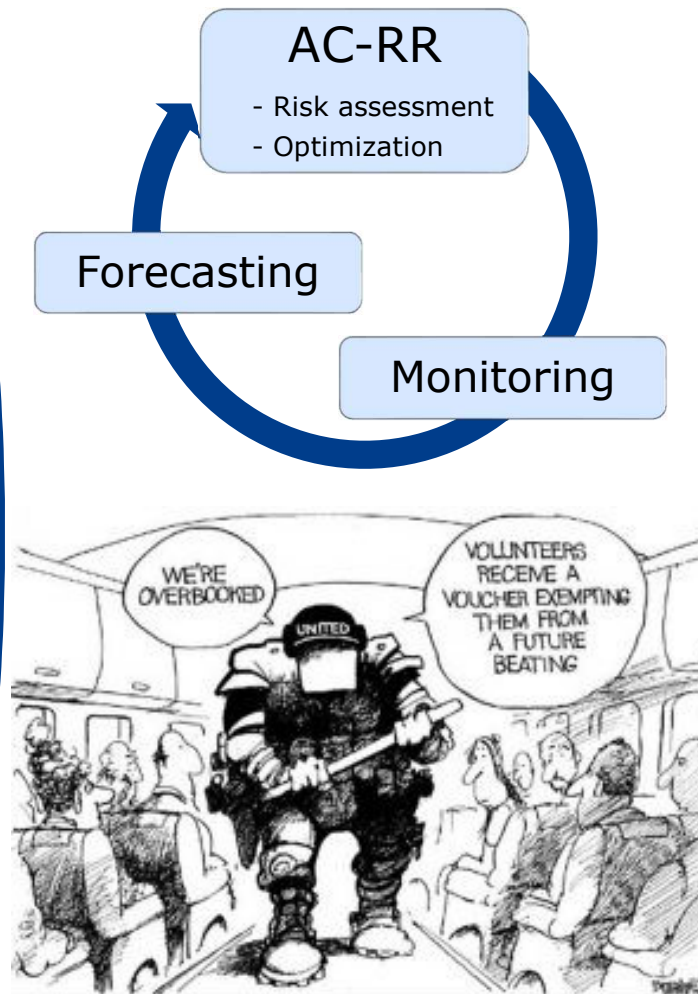
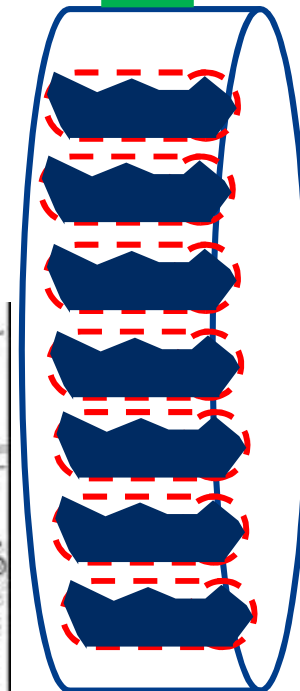
Verticals Reqs



5G System



5G System



Rate Latency CPU Storage

# Problem formulation

**Random processes  
hard to characterize**

Individual penalty of the slice

Individual reward of the slice

$$\min_{x^{(t)} \in \{0,1\}^S, z^{(t)} \in \mathbb{R}_+^S} \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=1}^T \sum_{\tau \in \mathcal{T}^{(t)}} \sum_{\substack{p \in \mathcal{P}_{b,c} \\ \forall b \in \mathcal{B}, c \in \mathcal{C}}} \underbrace{K_\tau x_{\tau,p}^{(t)}}_{\text{Expected penalty}} \underbrace{\Pr \left[ z_{\tau,p}^{(t)} < \lambda_{\tau,p}^{(t)} \right]}_{\text{Probability that the reservation is less than the usage (SLA violation)}} \underbrace{- R_\tau x_{\tau,p}^{(t)}}_{\text{Reward}}$$

Admission Control decisions

Resource Reservation decisions

Set of the slice requests at epoch t

subject to (1) capacity/delay/system constraints (linear and decoupled)

$$(2) z_{\tau,p}^{(t)} \leq x_{\tau,p}^{(t)} \Lambda_\tau, \quad \forall \tau \in \mathcal{T}, \forall p \in \mathcal{P}_{b,c}, \forall b \in \mathcal{B}, \forall c \in \mathcal{C}$$

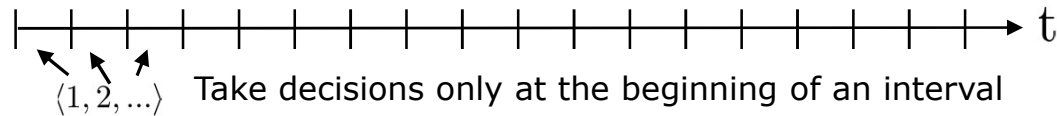
**No more resources than SLA if access is granted..  
... but nothing if access is rejected**



# Problem formulation

## Practical simplifications

- $L_\tau$  is small compared to the system's time horizon.



- We do forecasting. And weight by a deterministic risk-cost function.

### Risk of resource deficit

$$P_{\tau,p} := \frac{\Lambda_{\tau,p} - z_{\tau,p}}{\Lambda_{\tau,p} - \hat{\lambda}_{\tau,p}}, \quad 0 \leq P_{\tau,p} \leq 1$$

### Risk of wrong predictions

(forecast uncertainty amplified by slice duration)

$$\xi_{\tau,p} := \hat{\sigma}_{\tau,p} L_\tau, \quad 0 < \xi_{\tau,p} \leq L_\tau$$

$$\min_{\mathbf{x}^{(t)} \in \{0,1\}^S, \mathbf{z}^{(t)} \in \mathbb{R}_+^S} \quad \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=1}^T \sum_{\tau \in \mathcal{T}^{(t)}} \sum_{\substack{p \in \mathcal{P}_{b,c} \\ \forall b \in \mathcal{B} \ c \in \mathcal{C}}} \underbrace{K_\tau x_{\tau,p}^{(t)} \overbrace{\text{PP}[\hat{\lambda}_{\tau,p}^{(t)}]}^{\text{Expected penalty}}}_{\text{Penalty}} - \underbrace{R_\tau x_{\tau,p}^{(t)}}_{\text{Reward}}$$

subject to (1) Capacity/delay/system constraints (linear and decoupled)

$$(2) \ x_{\tau,p}^{(t)} \hat{\lambda}_{\tau,p}^{(t)} \leq z_{\tau,p}^{(t)} \leq x_{\tau,p}^{(t)} \Lambda_\tau, \quad \forall \tau \in \mathcal{T}, \forall p \in \mathcal{P}_{b,c}, \forall b \in \mathcal{B}, \forall c \in \mathcal{C}$$

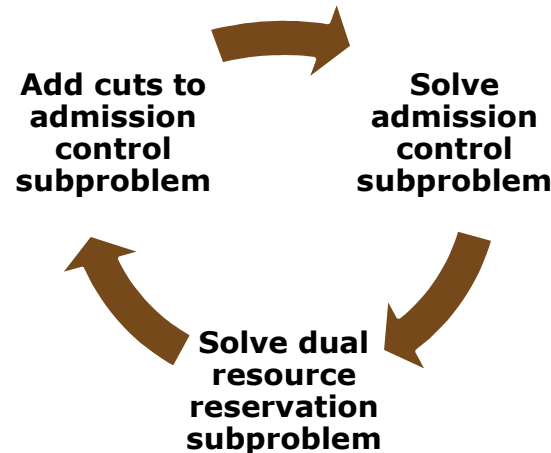
No less resources than forecasted peak demand

No more resources than SLA if access is granted

# Problem formulation

■ We can linearize our problem easily

- ... but still a MILP with coupled constraints (NP-hard)



■ We use two methods to solve this problem

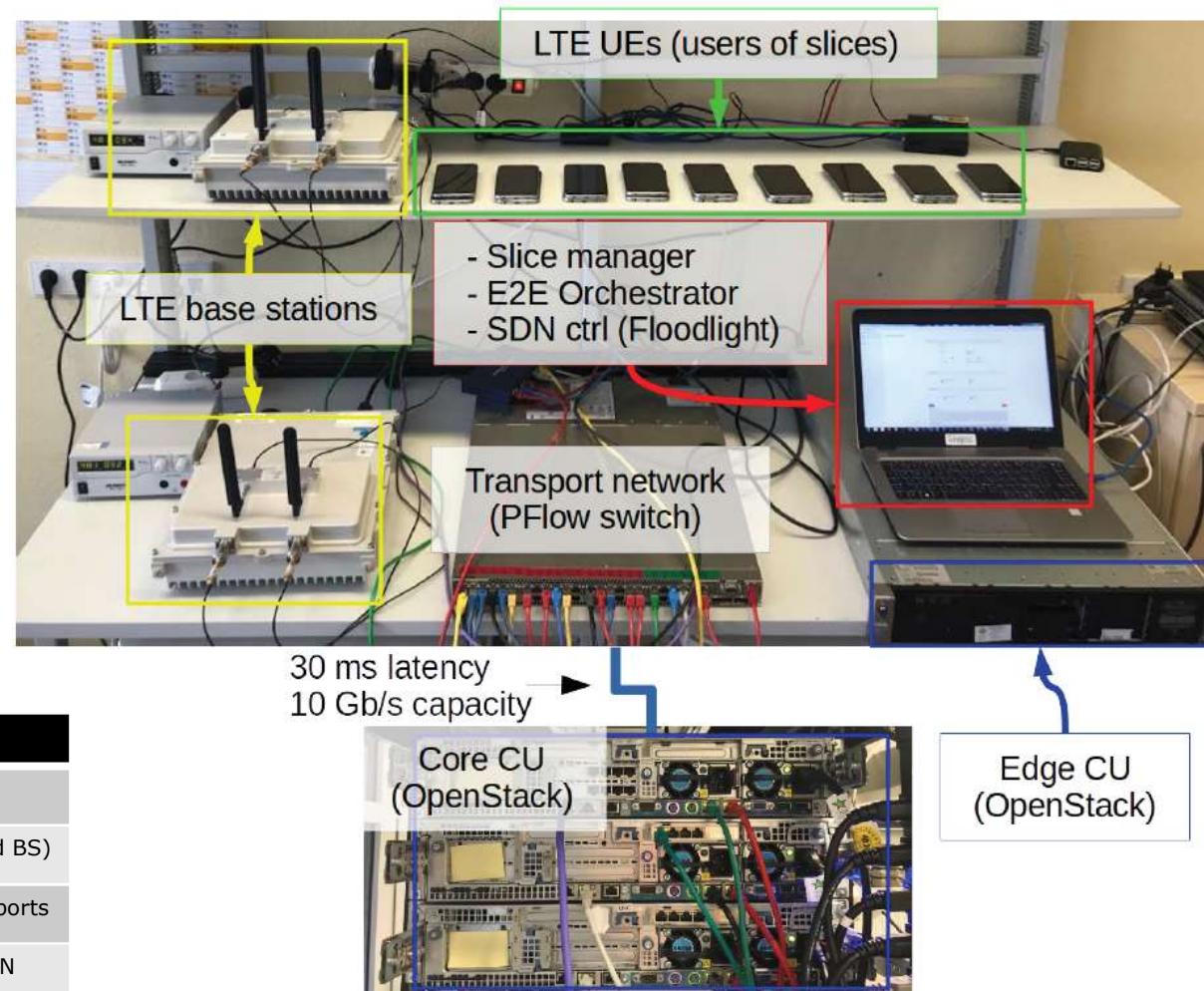
- **Bender's decomposition**

- Developed in 1960s, increasingly popular last years.
- Finds an exact solution, but no guarantees for the time needed

- **Heuristic algorithm**

1. Cast the admission control subproblem into a classical 0-1 Knapsack problem and use an heuristic to solve it.
2. Add only feasibility cuts

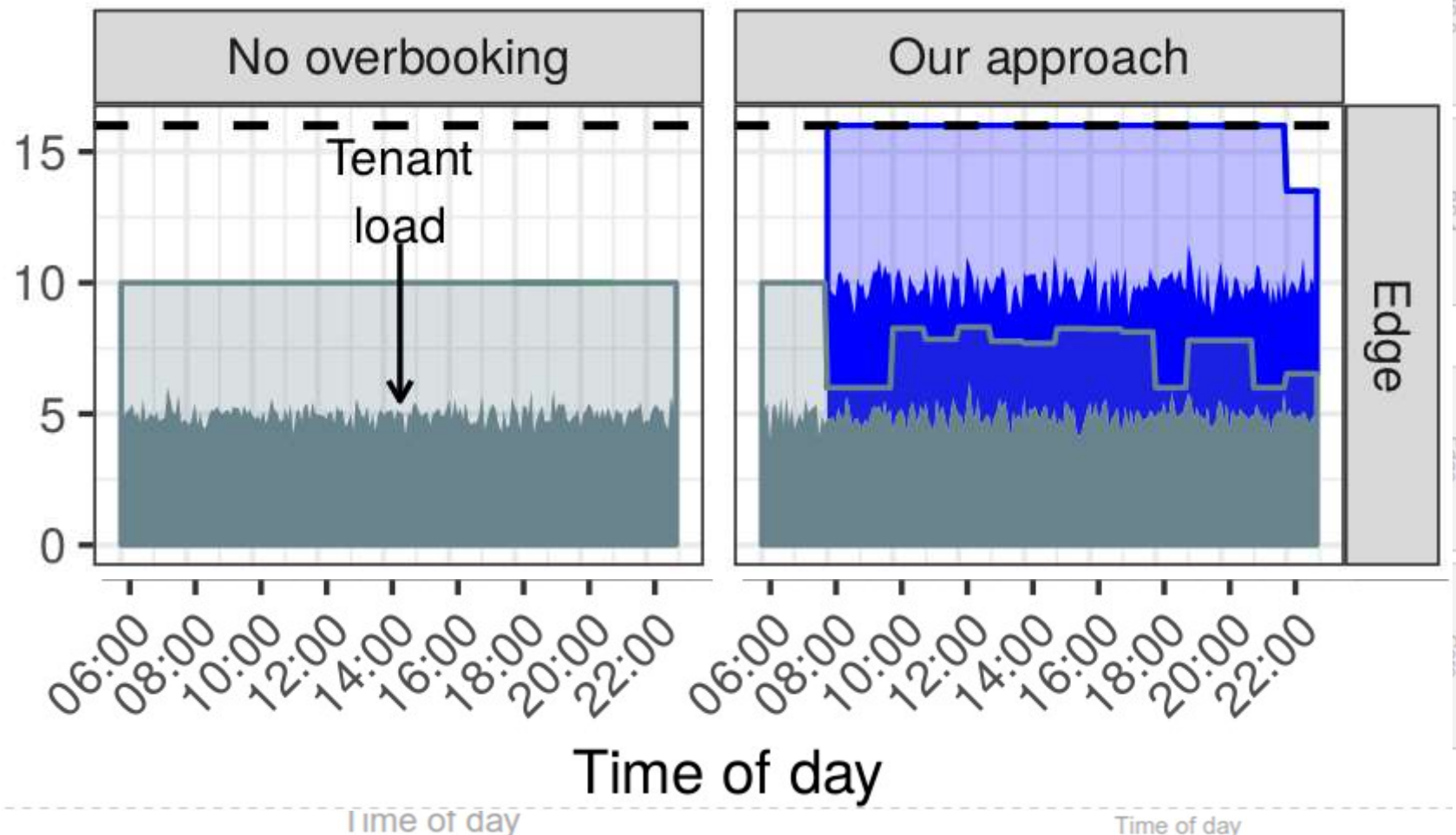
# Evaluation: Experimental proof-of-concept



| Device | Description                                       |
|--------|---|
| vEPCs  | OpenEPC Rel. 7 (1x per slice)                     |
| UEs    | Samsung Galaxy 7 (1x per slice and BS)            |
| BH     | Openflow 1.5 switch w/ 48 1-Gb/s ports            |
| RAN    | 2x 20-MHz NEC small cell w/ RAN sharing           |
| CUs    | OpenStack Queens w/ 16 (Edge) and 64 (Core) cores |

# Evaluation: Experimental proof-of-concept

| Type  | Delay req. (ms) | Service bitrate req. (Mb/s) | CPU req. (number of CPUs) | Price (monetary units) | Penalty (monetary units) |
|-------|-----------------|-----------------------------|---------------------------|------------------------|--------------------------|
| eMBB  | 30              | 50                          | 0                         | 1                      | 1/50                     |
| mMTC  | 30              | 10                          | $2\lambda$                | 3                      | 3/10                     |
| uRLLC | 5               | 25                          | $0.2\lambda$              | 2                      | 2/25                     |

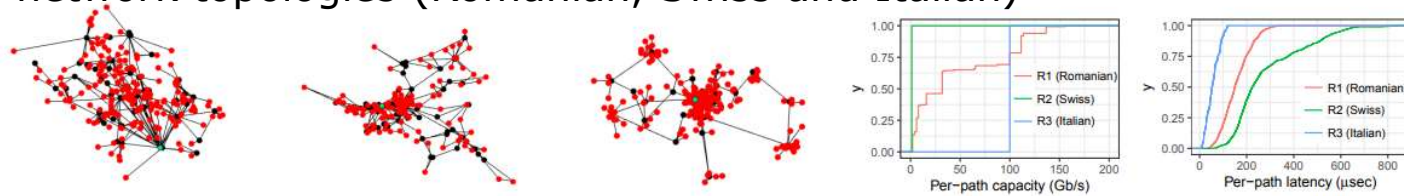




# Evaluation: Larger scale scenarios

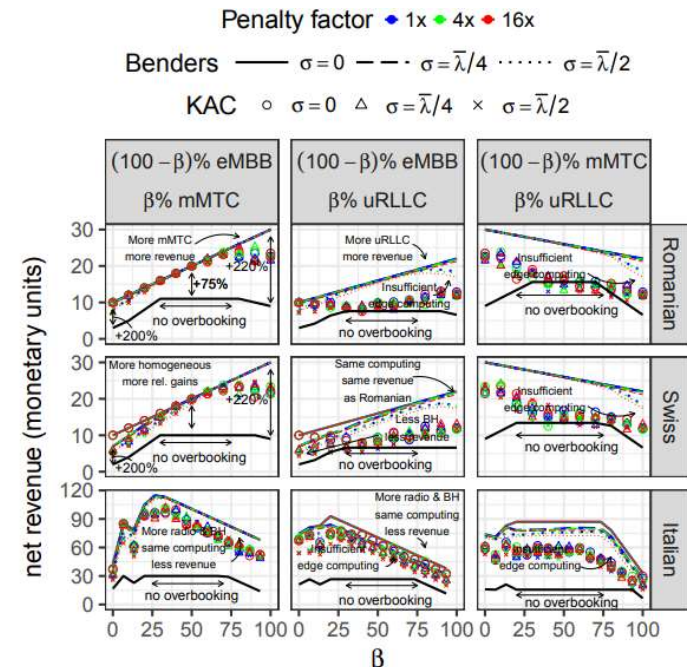
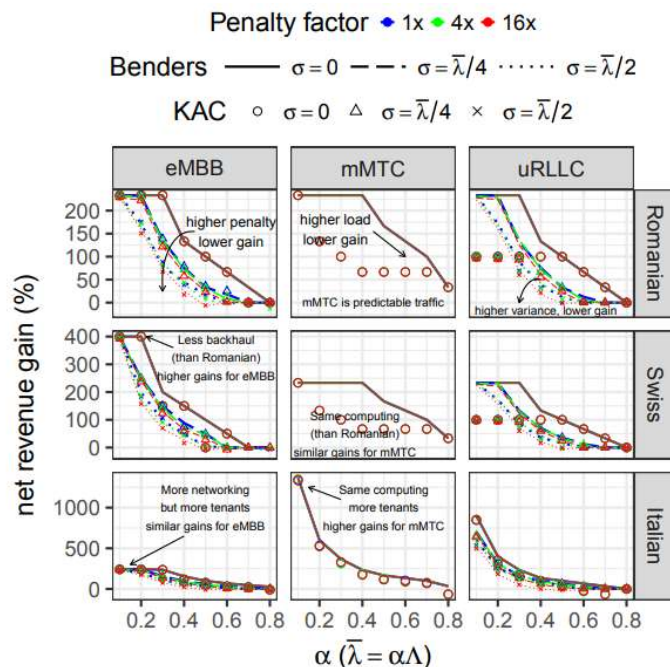
## Overbooking brings significant gains

- 3 real network topologies (Romanian, Swiss and Italian)



(a) Romanian topology (N1). (b) Swiss topology (N2). (c) Italian topology (N3). (d) Path Capacity Distribution (e) Path Delay Distribution

- Evaluation: Wide set of penalty, traffic load and traffic variability
- Up to 200% revenue gains when load is low and is predictable





# Conclusions

## Network Slicing will be a key technology for 5G

- New sources of revenue for mobile operators **and** vendors

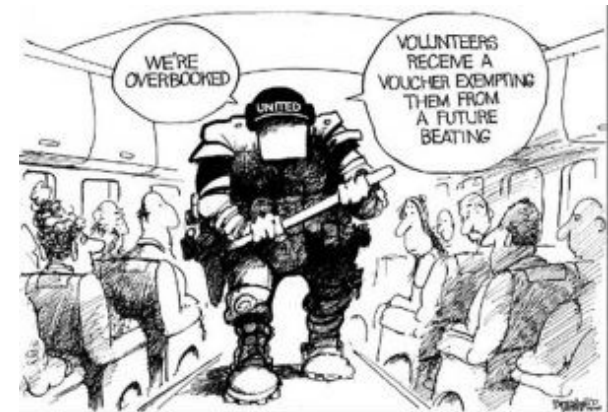
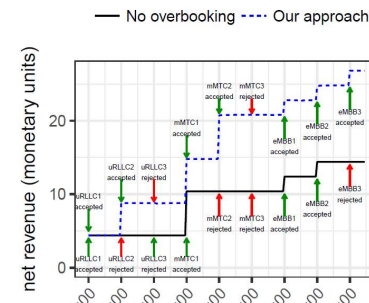
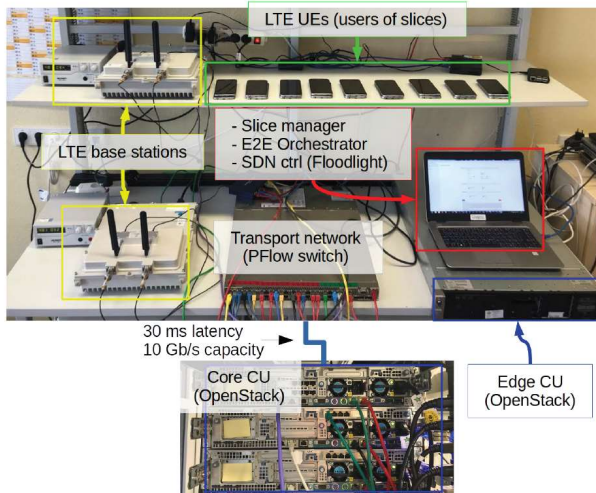
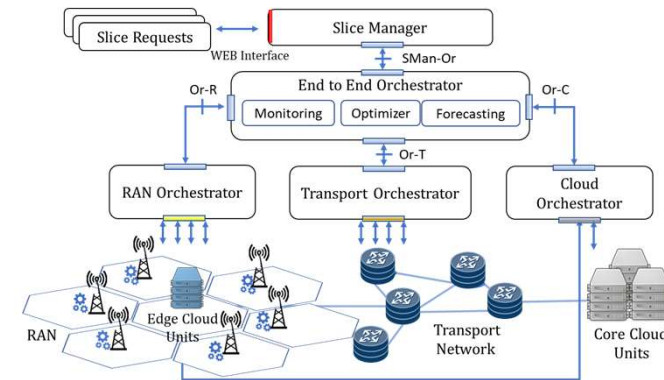
## Two main challenges:

### • Service/Slice **Orchestration Platform**

- Hierarchical orchestration for large-scale
- Feedback loop -> monitoring/control

### • Admission Control and Resource Reservation

- We explore the concept of **slice overbooking**
- 2 algorithms, 1 experimental PoC, 3 large-scale topologies



(a) Romanian topology (N1).

(b) Swiss topology (N2).

(c) Italian topology (N3).

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Heraklion/Crete, Greece

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