In this project, we were asked to solve the knapsack problem in three ways: once by using an exhaustive search method, once by using dynamic programming, and once by using a method of our choosing.

Firstly, the time and space efficiencies for each implementation will be proven:

*Exhaustive search:*

def **exhaustive**(items, capacity):

possibleSolutions = 2 \*\* len(items)

max\_value = -1

optimal\_items = []

for i in range(possibleSolutions):

decision\_matrix = [int(d) for d in str(bin(i)[2:]).zfill(len(items))]

current\_weight = 0

current\_value = 0

for j in range(len(decision\_matrix)):

if decision\_matrix[j] == 1:

current\_weight += items[j][0]

current\_value += items[j][1]

if current\_weight <= capacity and current\_value > max\_value:

max\_value = current\_value

optimal\_items = decision\_matrix

return optimal\_items

In the exhaustive search, all possible subsets were generated and compared against each other to see which would give an optimal solution.

Because this assignment deals with a 0/1 knapsack, then all of the possible subsets for an item set of *n* items would be the power set of *n*: .

As , that means that the outermost loop is . However, because this particular solution was implemented by getting via a binary representation, the process of using the binary vector to check each solution is O(n), as shown in the inner loop. Therefore, the time efficiency for the exhaustive search is

As for space efficiency, the only storage that changes with *n* is optimal\_items and decision\_matrix. Therefore, the space efficiency is .

*Dynamic Programming:*

def **dynamic**(items, capacity):

memoization\_table =

[[0 for i in range(capacity + 1)] for j in range(len(items) + 1)]

for i in range(1, len(items) + 1):

current\_item = i - 1

for j in range(capacity + 1):

if j - items[current\_item][0] >= 0:

if items[current\_item][1] +

memoization\_table[i - 1][j - items[current\_item][0]] >

memoization\_table[i - 1][j]:

memoization\_table[i][j] =

items[current\_item][1] +

memoization\_table[i - 1][j - items[current\_item][0]]

else:

memoization\_table[i][j] = memoization\_table[i - 1][j]

else:

memoization\_table[i][j] = memoization\_table[i - 1][j]

i = len(items)

j = capacity

optimal\_items = [0 for x in range(len(items))]

while i > 0 and j > 0:

if memoization\_table[i][j] != memoization\_table[i - 1][j]:

optimal\_items[i - 1] = 1

j -= items[i - 1][0]

i -= 1

return optimal\_items

In the dynamic programming approach, a table which represents the subproblems is created. Because the table has dimensions of *n* and capacity *c*, then the time efficiency to create and fill in the table is . This creation and generation of the table is presented by the nested for loops.

The other major aspect of this implementation is the backtracking algorithm. The specific way that this algorithm was implemented was that the program started in the bottom right hand corner of the memoization\_table, go up until the next number is different (thus giving you the smallest weight for the largest value), then move over to the column where the . The process is then repeated, until all the solutions are found.

Even though the size of the matrix is O(nc), there is no possible way that the program will iterate through every cell, simple due to the nature of the algorithm. Therefore, the time efficiency for the backtracking is < O(nc), and therefore can be ignored.

Hence, the time efficiency for the dynamic programming approach is .

As mentioned previously, the memoization\_table, is a table of dimensions *n* x *c*. The only other data storage in this approach is optimal\_items, which is of size *n*. Therefore, the space efficiency for the dynamic programming approach is

*Greedy Algorithm:*

def **greedy**(items, capacity):

complete\_array = []

for i in range(len(items)):

complete\_array.append([i,

items[i][0],

items[i][1],

items[i][1] / items[i][0]])

complete\_array = sorted(complete\_array, key=lambda x: x[3])

remaining\_space = capacity

optimal\_items = [0 for x in range(len(items))]

i = len(complete\_array) - 1

while remaining\_space > 0 and i >= 0:

if remaining\_space - complete\_array[i][1] >= 0:

optimal\_items[complete\_array[i][0]] = 1

remaining\_space -= complete\_array[i][1]

i-= 1

return optimal\_items

The approach that I chose for this program is a greedy algorithm. The basic idea behind the algorithm is to sort the items based on the value to weight ratio. Then, the algorithm will try to put the item with the largest value to weight ratio until no more items can fit.

There are three main aspects of this algorithm: generating the value to weight ratios, sorting the algorithm, and then determining what are the optimal values for the knapsack.

For the generation of the value to weight ratios, the time complexity is fairly straightforward: O(n).

For the sorting of the algorithm, python’s default Timsort was used, which has a time complexity of O(nlogn)

For the determination of what are the optimal values for the knapsack, it is obvious that the worst case is if all the items belong in the knapsack, which would yield a time complexity of O(n).

Therefore, the total time complexity for the greedy algorithm is .

As for the space efficiency, there are only two instances of data storage, complete\_array and optimal\_items. The space efficiency for both is O(n), so the overall space efficiency for the greedy algorithm is O(n).