

$$① (\sqrt{n} + 2)(n + 5) =$$

$$= \sqrt{n} \cdot n + 5\sqrt{n} + 2n + 10 =$$

$$O(\sqrt{n} \cdot n)$$

$$\boxed{\text{res} \rightarrow 2}$$

$$② \quad n = 5$$

$$\begin{array}{c} 5 \\ 4 \\ 3 \\ 2 \\ 1 \end{array} \quad \begin{array}{c} ① \\ ② \\ ③ \\ ④ \\ ⑤ \end{array}$$

$$\boxed{\sum_{i=1}^n cte}$$

$$\boxed{\text{res} \rightarrow 4}$$

$$③ \quad n \cdot c$$

$$\boxed{\text{res} \rightarrow 4}$$

$$④ \quad T(n) = a + T(n-3)$$

$$\boxed{\text{res} \rightarrow 4}$$

$$⑤ \quad \boxed{2T \left(\underset{\substack{\downarrow \\ n-1-1}}{n-2} \right) + n-1}$$

$$\boxed{\text{res} \rightarrow 5}$$

$$⑥ \quad T(n) \begin{cases} cte & n = 0 \\ 2T(n-1) + cte & n > 0 \end{cases}$$

se pone en 0 que $n > 0$

paso ① $2T(n-1) + cte$

paso ② $2[2T(n-2) + cte] + cte =$
 $= 4T(n-2) + 3cte$

paso ③ $4[2T(n-3) + cte] + 3cte =$
 $= 8T(n-3) + 4cte$

paso ④ $2^i T(n-i) + (2^i - 1) \cdot cte$

caso base $T(n-i) = 0$
 $n = i$

reemplazo i

$$2^n T(n-n) + (2^n - 1) \cdot cte =$$

$$= 2^n \cdot T(0) + (2^n - 1) \cdot cte =$$

$$= 2^n \cdot cte_1 + 2^n \cdot cte_2 - cte_2 = T(n)$$

$O(2^n) \rightarrow \text{res } 1$

④ res $\rightarrow 2$ $O(1), \log(n), n \log(n), n^2$

⑧ $T(n) = \begin{cases} 1 & \text{si } n = 1 \\ 16T(n/2) + n^4 & \text{si } n \geq 2 \end{cases}$

paso ① $16T\left(\frac{n}{2}\right) + n^4$

caso (2)

$$16 T\left(\frac{n}{4}\right) + \left(\frac{n}{2}\right)^4 + n^4 =$$

$$= 256 T\left(\frac{n}{4}\right) + \cancel{16} \frac{n^4}{\cancel{16}} + n^4 =$$

$$= 256 T\left(\frac{n}{4}\right) + 2n^4$$

caso (1)

$$16^i T\left(\frac{n}{2^i}\right) + i n^4$$

caso base

$$\frac{n}{2^i} = 1$$

$$n = 2^i$$

$$\log_2(n) = i$$

$$16^i = (2^4)^i$$

$$= (2^i)^4 =$$

$$= \boxed{n^4}$$

reemplazo

$$n^4 T\left(\frac{n}{2^{\log_2(n)}}\right) + \log_2(n) \cdot n^4 =$$

$$= n^4 \cdot T(1) + \log_2(n) \cdot n^4 =$$

$$= n^4 \cdot 1 + \log_2(n) \cdot n^4$$

$$O(\log_2(n) \cdot n^4) \quad \text{res} \rightarrow c$$

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TIEMPO while

si $n = 8$

1

$$j = 1$$

2

$$j = 2$$

3

$$j = 4$$

i

$$j = 2^{i-1}$$

$$\boxed{4}$$

caso base

$$2^{i-1} > n$$

$$i-1 > \log_2(n)$$

$$\log_2(n)$$

$$\sum_{i=1}$$

FOR

$$n = 2$$

1

$$4 = n * n$$

2

$$3 = n * n - 1$$

3

$$2 = n * n - 2$$

i

$$n * n - (i-1)$$

$$\boxed{4}$$

caso base

$$n * n - i + 1 < 1$$

$$n * n + 1 < 1 + i$$

$$n * n + \cancel{1} / \cancel{1} < i$$

$$n * n$$

$$n * n < i$$

$$\sum_{i=1}$$

$$T(n) = \sum_{i=1}^{\log_2(n)} \left(\sum_{j=1}^{n/2^i} 1 \right) =$$

$$= \sum_{i=1}^{\log_2(n)} n^2 = \boxed{\log_2(n) \cdot n^2}$$

res \rightarrow d

$$(10) \quad T(n) = \sum_{i=1}^z \left(\sum_{j=1}^{\log_2(n)} 1 \right) =$$

$$= \sum_{i=1}^z \log_2(n) = z \log_2(n)$$

$$n = 1024$$

$$T(1024) = 20$$

$$1.000 \text{ ————— } 1 \text{ seg}$$

$$20 \text{ ————— } 0,02 \text{ seg}$$

res \rightarrow c