

$$1) i. \quad T(n) = \begin{cases} c e_1 & n \leq 1 \\ c e_2 + T(n-1) & n > 1 \end{cases}$$

Suponiamo che  $n > 1$

Paso ①  $c e_2 + T(n-1)$

Paso ②  $c e_2 + [c e_2 + T(n-2)] =$   
 $= 2 c e_2 + T(n-2)$

Paso ③  $2 c e_2 + [c e_2 + T(n-3)] =$   
 $= 3 c e_2 + T(n-3)$

Paso i (paso general)

$$T(n) = i \cdot c e_2 + T(n-i)$$

Caso base

$$\boxed{n-i \leq 1}$$

$$-i \leq 1-n$$

$$i \geq n-1$$

Reemplazo i

$$(n-1) \cdot c e_2 + T(n-(n-1)) =$$

$$= (n-1) c e_2 + c e_1 = T(n)$$

↓  
 $O(n)$

ii.

$$T(n) = \begin{cases} \alpha e_1, & n \leq 1 \\ 2 * T(n-1) + \alpha e_2, & n > 1 \end{cases}$$

Suponiendo que  $n > 1$

Paso ①

$$2 * T(n-1) + \alpha e_2$$

Paso ②

$$2 * [2 * T(n-2) + \alpha e_2] + \alpha e_2 = \\ = 4 T(n-2) + 3 \alpha e_2$$

Paso ③

$$4 [2 * T(n-3) + \alpha e_2] + 3 \alpha e_2 = \\ = 8 T(n-3) + 7 \alpha e_2$$

Paso ④

$$8 [2 * T(n-4) + \alpha e_2] + 7 \alpha e_2 = \\ = 16 T(n-4) + 15 \alpha e_2$$

Paso i (paso general)

$$2^i * T(n-i) + (2^i - 1) * \alpha e_2$$

Caso base

$$n-i \leq 1$$

$$-i \leq 1-n$$

$$i \geq n-1$$

reemplazo i

$$2^{n-1} * T(n-(n-1)) + (2^{n-1} - 1) \alpha e_2 =$$



$$= \left( 2^{n-1} \right) \cdot \sigma e_1 + \left( 2^{n-1} - 1 \right) \sigma e_2 = \tau(n)$$

$$\frac{2^n}{2} \cdot \sigma e_1 + \left( \frac{2^n}{2} - 1 \right) \cdot \sigma e_2 = 2^n \cdot \frac{\sigma e_1}{2} + \left( \frac{2^n}{2} - 1 \right) \cdot \sigma e_2$$

↓

$O(2^n)$  ?