# Non-interactive zero-knowledge proof of SHE

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### 1 Notations

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G_1 = \langle g_1 \rangle \text{ ; DLP is hard,}
G_2 = \langle g_2 \rangle \text{ ; DLP is hard,}
sk = (s_1, s_2) \text{ ; secret keys,}
pk = (h_1, h_2) \text{ ; public keys where } h_1 = g_1^{s_1}, \ h_2 = g_2^{s_2}.
Enc(m) = (c_1, c_2, c_3, c_4) = (g_1^{\rho}, g_1^{m} h_1^{\rho}, g_2^{\sigma}, g_2^{m} h_2^{\sigma}) \text{ where } \rho, \sigma \leftarrow \mathbb{Z}_p.
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### 2 Equality of DLs

#### 2.1 Prove

$$\begin{split} & r_{\rho}, r_{\sigma}, r_{m} \leftarrow \mathbb{Z}_{p}, \\ & (R_{1}, R_{2}, R_{3}, R_{4}) = ({g_{1}}^{r_{\rho}}, {g_{1}}^{r_{m}}{h_{1}}^{r_{\rho}}, {g_{2}}^{r_{\sigma}}, {g_{2}}^{r_{m}}{h^{r_{\sigma}}}), \\ & c = H(pp, pk, c_{1}, c_{2}, c_{3}, c_{4}, R_{1}, R_{2}, R_{3}, R_{4}), \\ & (s_{\rho}, s_{\sigma}, s_{m}) = (r_{\rho} + c\rho, r_{\sigma} + c\sigma, r_{m} + cm), \\ & \text{output } (c, s_{\rho}, s_{\sigma}, s_{m}). \end{split}$$

### 2.2 Verify

verify 
$$c = H(pp, pk, c_1, c_2, c_3, c_4, R'_1, R'_2, R'_3, R'_4),$$
  
where  $(R'_1, R'_2, R'_3, R'_4) = (g_1^{s_p}/c_1^c, g_1^{s_m} h_1^{s_p}/c_2^c, g_2^{s_\sigma}/c_3^c, g_2^{s_m} h_2^{s_\sigma}/c_4^c).$ 

### 2.3 Correctness

$$\begin{split} R_1' &= g_1{}^{s_\rho-c\rho} = g_1{}^{r_\rho} = R_1, \\ R_2' &= g_1{}^{s_m-cm}h_1{}^{s_\rho-c\rho} = g_1{}^{r_m}h_1{}^{r_\rho} = R_2, \\ R_3' &= g_2{}^{s_\sigma-c\sigma} = g_2{}^{r_\sigma} = R_3, \\ R_4' &= g_2{}^{s_m-cm}h_2{}^{s_\rho-c\rho} = g_2{}^{r_m}h_2{}^{r_\rho} = R_4. \end{split}$$

# 3 m = 0 or 1

### 3.1 Prove

$$\begin{split} &d_{1-m}, s_{\rho,1-m} \leftarrow \mathbb{Z}_p, \\ &R_{1,1-m} = g_1^{s_{\rho,1-m}}/c_1^{d_{1-m}}, \\ &R_{2,1-m} = h_1^{s_{\rho,1-m}}/(c_2/g_1^{1-m})^{d_{1-m}}, \\ &r_{\rho,m}, r_{\rho}, r_{\sigma}, r_m \leftarrow \mathbb{Z}_p, \\ &R_{1,m} = g_1^{r_{\rho,m}}, \\ &R_{2,m} = h_1^{r_{\rho,m}}, \\ &c = H(pp, pk, c_1, c_2, R_{1,0}, R_{2,0}, R_{1,1}, R_{2,1}), \\ &d_m = c - d_{1-m}, \\ &s_{\rho,m} = r_{\rho,m} + d_m \rho, \\ &\text{output } (d_0, d_1, s_{\rho,0}, s_{\rho,1}). \end{split}$$

### 3.2 Verify

$$\begin{split} R'_{1,i} &= g_1{}^{s_{\rho,i}}/c_1{}^{d_i}, \text{ for } i = 0, 1, \\ R'_{2,0} &= h_1{}^{s_{\rho,0}}/c_2{}^{d_0}, \\ R'_{2,1} &= h_1{}^{s_{\rho,1}}/(c_2/g_1)^{d_1}, \\ c &= H(pp, pk, c_1, c_2, R'_{1,0}, R'_{2,0}, R'_{1,1}, R'_{2,1}), \\ \text{verify } c &= d_0 + d_1. \end{split}$$

# 4 m = 0 or 1 and Equality of DLs

### 4.1 Prove

$$\begin{split} &d_{1-m}, s_{\rho,1-m} \leftarrow \mathbb{Z}_p, \\ &R_{1,1-m} = g_1^{s_{\rho,1-m}}/c_1^{d_{1-m}}, \\ &R_{2,1-m} = h_1^{s_{\rho,1-m}}/(c_2/g_1^{1-m})^{d_{1-m}}, \\ &r_{\rho,m}, r_{\rho}, r_{\sigma}, r_m \leftarrow \mathbb{Z}_p, \\ &R_{1,m} = g_1^{r_{\rho,m}}, \\ &R_{2,m} = h_1^{r_{\rho,m}}, \\ &R_3 = g_1^{r_{\rho}}, \\ &R_4 = g_1^{r_m}h_1^{r_{\rho}}, \\ &R_5 = g_2^{r_{\sigma}}, \\ &R_6 = g_2^{r_m}h_2^{r_{\sigma}}, \\ &c = H(pp, pk, c_1, c_2, R_{1,0}, R_{2,0}, R_{1,1}, R_{2,1}, R_3, \dots, R_6), \\ &d_m = c - d_{1-m}, \\ &s_{\rho,m} = r_{\rho,m} + d_m \rho, \\ &s_{\rho} = r_{\rho} + c\rho, \\ &s_{\sigma} = r_{\sigma} + c\sigma, \\ &s_m = r_m + cm, \\ &\text{output } (d_0, d_1, s_{\rho,0}, s_{\rho,1}, s_{\sigma}, s_{\rho}, s_m). \end{split}$$

### 4.2 Verify

$$\begin{split} R'_{1,i} &= g_1{}^{s_{\rho,i}}/c_1{}^{d_i}, \text{ for } i = 0, 1, \\ R'_{2,0} &= h_1{}^{s_{\rho,0}}/c_2{}^{d_0}, \\ R'_{2,1} &= h_1{}^{s_{\rho,1}}/(c_2/g_1)^{d_1}, \\ R'_3 &= g_1{}^{s_\rho}/c_1{}^c, \\ R'_4 &= g_1{}^{s_m}h_1{}^{s_\rho}/c_2{}^c, \\ R'_5 &= g_2{}^{s_\sigma}/c_3{}^c, \\ R'_6 &= g_2{}^{s_m}h_2{}^{s_\sigma}/c_4{}^c, \\ \text{where } c &= d_0 + d_1, \\ \text{verify } c &= H(pp, pk, c_1, c_2, R_{1,0}, R_{2,0}, R_{1,1}, R_{2,1}, R'_3, \dots, R'_6). \end{split}$$

**5** 
$$m \in M := \{ m_1, \dots, m_n \}$$

### 5.1 Notations

$$P$$
 : generator 
$$x: \text{secret key}$$
 
$$Q:=xP: \text{public key}$$
 
$$\operatorname{Enc}(m,r):=(mP+rQ,rP); \text{ ciphertext of } m$$

Properties:

$$\operatorname{Enc}(m_1, r_1) + \operatorname{Enc}(m_2, r_2) = \operatorname{Enc}(m_1 + m_2, r_1 + r_2).$$

### 5.2 Prove

Let  $C := \operatorname{Enc}(m, r)$  and  $m = m_{i'}$ . Select  $\{a_i\}_{i \neq i'}$  and  $\{t_i\}$  randomly.  $R_i := \operatorname{Enc}(a_i(m - m_i), t_i),$ 

$$h := \operatorname{Hash}(P, Q, C, \{ R_i \}),$$

$$a_{i'} := h - \sum_{i \neq i'} a_i,$$

$$b_i := t_i - a_i r.$$

Output a proof  $\pi := \{ a_i, b_i \}$ .

### 5.3 Verify

For given  $P, Q, C, \pi := \{ a_i, b_i \}, M$ ,

$$R_i := a_i(C - \operatorname{Enc}(m_i, 0)) + \operatorname{Enc}(0, b_i),$$
  

$$h := \operatorname{Hash}(P, Q, C, \{R_i\}),$$
  
Verify  $h = a_1 + \dots + a_n$ .

Remark: If the verification is okay,

$$R_i = \operatorname{Enc}(a_i(m-m_i), a_ir + b_i).$$

Let  $t_i := a_i r + b_i$ , then

$$\operatorname{Hash}(P, Q, C, \{ \operatorname{Enc}(a_i(m - m_i), t_i) \}) = a_1 + \dots + a_n.$$

If  $m \notin M$ , it is hard to find  $a_i, t_i$ .