Number Systems II:

2’s Complement, Arithmetic, Overflow, & Writing Bit-wise Logical & Shifting Code

PICK UP HANDOUT!

Today's lecture

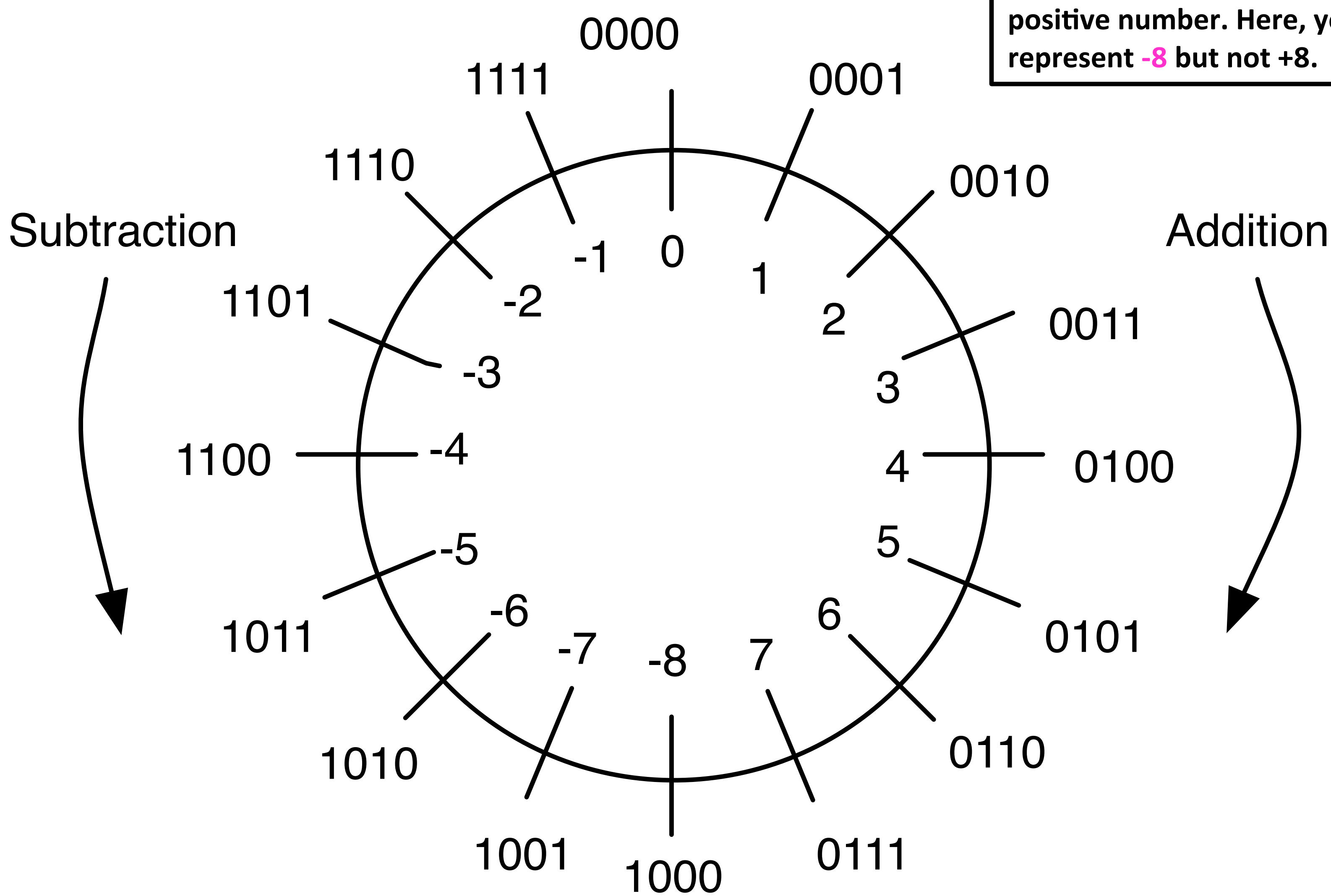
- **Two's complement signed binary representation**
  - Negating numbers in Two's complement
  - Sign extension

- **Bit-wise shift operations**
  - Writing bit-wise logical and shifting code

- **Two's complement arithmetic**
  - Addition
  - Subtraction
  - Overflow

Review: 4-bit 2’s complement

Two’s complement has asymmetric ranges; there is one more negative number than positive number. Here, you can represent -8 but not +8.



Negating Numbers in 2's Complement

- **To negate a number:**
  - Complement each bit and then add 1.

$$\begin{array}{r} 1\,1 \\ 1011 \\ +\;1 \\ \hline 1100 \end{array}$$

- **Example:**

0100	**= +4_{10}**	**(a positive number in 4-bit two's complement)**
1011	**=**	**(invert all the bits)**
1100	**= -4_{10}**	**(and add one)**
0011	**=**	**(invert all the bits)**
0100	**= +4_{10}**	**(and add one)**

$$\begin{array}{r} 1\,1 \\ 0011 \\ +\;1 \\ \hline 0100 \end{array}$$

Sometimes, people talk about "taking the two's complement" of a number. This is a confusing phrase, but it usually means to negate some number that's already in two's complement format.

Negating Numbers in 2's Complement

- **To negate a number:**
  - Complement each bit and then add 1.
- **Example:**

0100	= $+4_{10}$	(a positive number in 4-bit two's complement)
1011	=	(invert all the bits)
1100	= -4_{10}	(and add one)
0011	=	(invert all the bits)
0100	= $+4_{10}$	(and add one)

Converting 2's Complement to Decimal

- **Algorithm 1:**
  - if negative, negate; then do unsigned binary to decimal
- **Algorithm 2:**
  - **Same as with n-bit unsigned binary**
    - **Except, the MSB is worth $-(2^{n-1})$**

$$-b_{n-1}2^{n-1} + \sum_{k=0}^{n-2} b_k 2^k$$

- **Example:**

-8 4 2 1

1100 = -4_{10} (a negative number in 4-bit two's complement)

-8 + 4 = -4

2's Complement Negation

- **If 01101 is the 5-bit representation for 13, what is the 2's complement representation for -13?**
  - A: 10010
  - B: 11010
  - C: 10001
  - D: 10011
  - E: 01101

$$\begin{array}{r} 10010 \\ +1 \\ \hline 10011 \end{array}$$

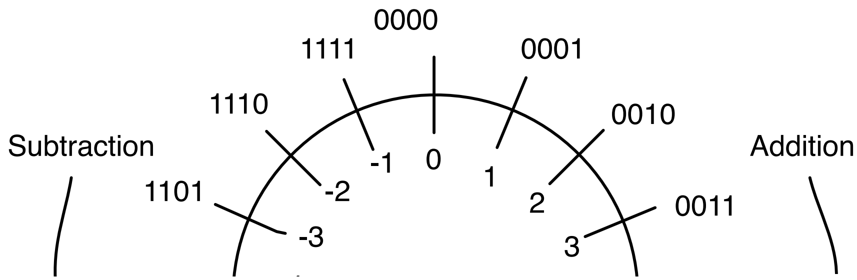
Sign Extension

- In everyday life, decimal numbers are assumed to have an infinite number of 0's in front of them. This helps in "lining up" numbers.
- To subtract 231 and 3, for instance, you can imagine:

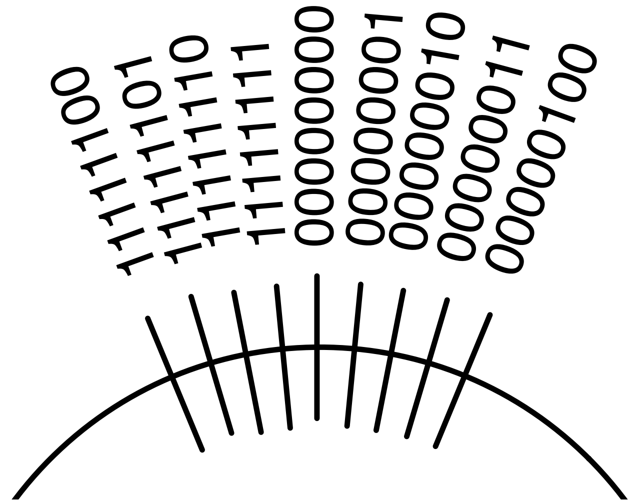
$$\begin{array}{r} 231 \\ -\ 003 \\ \hline 228 \end{array}$$

- This works for ***positive*** 2's complement numbers, but not ***negative*** ones.
- To preserve sign and value for negative numbers, we add more 1's.
- For example, going from 4-bit to 8-bit numbers:
  - 0101 (+5) should become 0000 0101 (+5).
  - But 1100 (-4) should become 1111 1100 (-4).
- The proper way to extend any signed binary number is to replicate the sign bit.

Sign Extension, cont.



4-bit



8-bit

What you need to know for Lab 2.

Review: Bitwise Logical operations

unsigned char a = 0x55; 0 1 0 1 0 1 0 1

unsigned char b = 0x0f; 0 0 0 0 1 1 1 1

- **Last time we introduced bit-wise logical operations:**

unsigned char c = a | b; *(bit-wise OR)*

```
      0 1 0 1  0 1 0 1
OR    0 0 0 0  1 1 1 1
      ----------------
      0 1 0 1  1 1 1 1
```

unsigned char d = a & b; *(bit-wise AND)*

```
      0 1 0 1  0 1 0 1
AND   0 0 0 0  1 1 1 1
      ----------------
      0 0 0 0  0 1 0 1
```

MASKING OPERATION

unsigned char e = a ^ b; *(bit-wise XOR)*

```
      0 1 0 1  0 1 0 1
XOR   0 0 0 0  1 1 1 1
      ----------------
      0 1 0 1  1 0 1 0
```

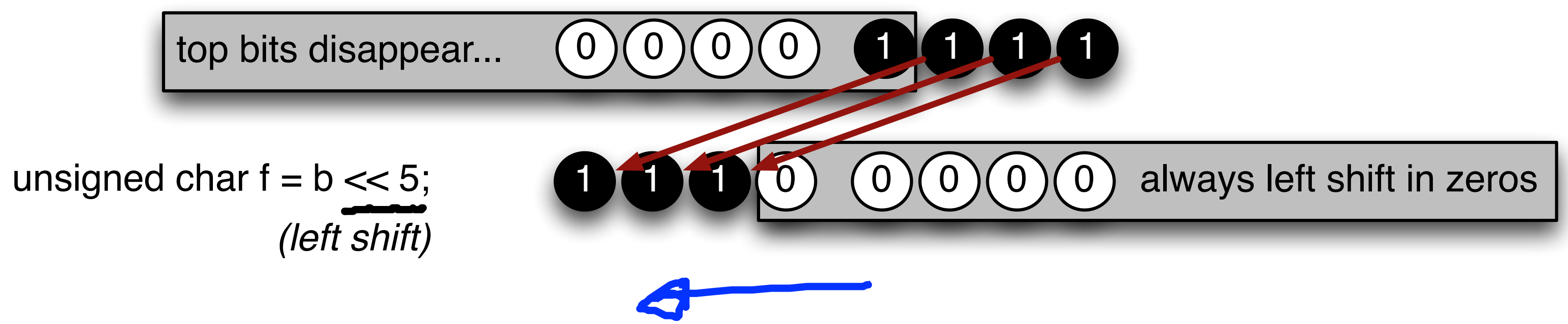
unsigned char n = ~a; *(bit-wise NOT)*

```
NOT   0 1 0 1  0 1 0 1
      ----------------
      1 0 1 0  1 0 1 0
```

Bit-wise shifting

- **When doing bit-wise logical operations, it can be useful to "shift" bits to the left or right within a word.**

- **Left shift:**



We are shifting bits toward the most significant bit (MSB); we call this a left shift because we think of the MSB being on the left.

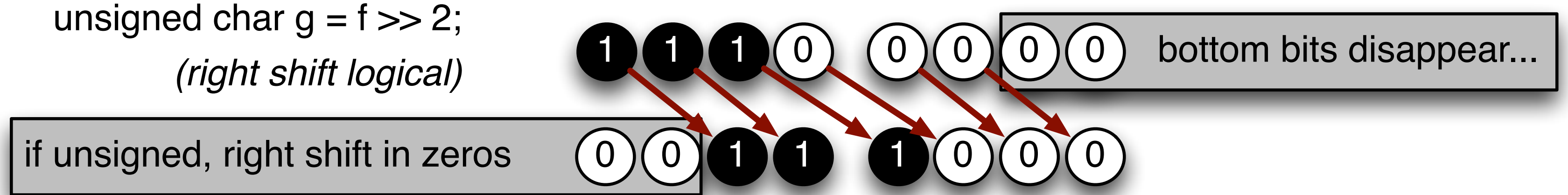
Bit-wise shifting, cont.

- **Two kinds of right shift, depends on type of variable:**

f/4

  - **Unsigned numbers**

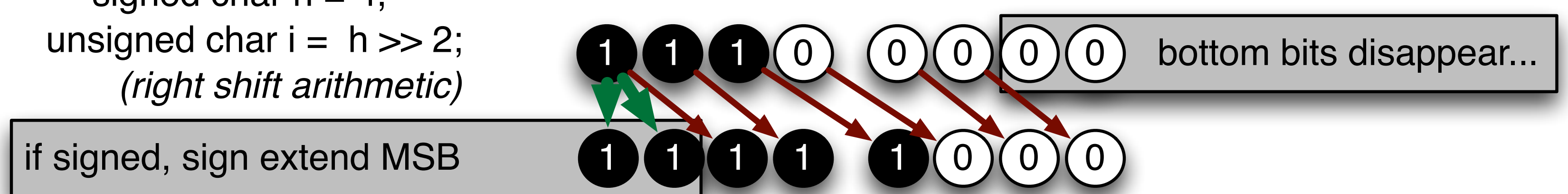
unsigned char g = f >> 2;
(right shift logical)



unsigned char i = (unsigned)(h>>2);

  - **Signed numbers**

signed char h = f;
unsigned char i = h >> 2;
(right shift arithmetic)



Note: x >> 1 not the same as x/2 for negative numbers; compare (-3)>>1 with (-3)/2

Useful for extracting bits

- **We have the unsigned 8-bit word:** x= $b_7b_6b_5b_4b_3b_2b_1b_0$
- **And we want the 8-bit word:** y= $0\,0\,0\,0\,0\,b_5b_4b_3$
  - i.e., we want to extract bits 3-5.

- **We can do this with bit-wise logical & shifting operations**
  - y = (x >> 3) & 0x7;

0x00000007

((x<<2)>>5)

x	$b_7b_6b_5b_4b_3b_2b_1b_0$
x >> 3	$0\,0\,0\,b_7\,b_6\,b_5\,b_4\,b_3$
(x >> 3) & 0x7	$0\,0\,0\,0\,0\,b_5\,b_4\,b_3$

$0xa \rightarrow 1010$

Useful for merging two bit patterns

- We have 2 unsigned 8-bit words: $x = a_7a_6a_5a_4a_3a_2a_1a_0$

  $y = b_7b_6b_5b_4b_3b_2b_1b_0$

- And we want the 8-bit word: $z = a_7b_6a_5b_4a_3b_2a_1b_0$

$(x \,\&\, 0xaa) = a_7 0 a_5 0 a_3 0 a_1 0$

$(y \,\&\, 0x55) = 0 b_6 0 b_4 0 b_2 0 b_0$

$(x \,\&\, 0xaa) \mid (y \,\&\, 0x55)$

Bit-wise Logical & Shifting

iclicker.

- **We have 2 unsigned 8-bit words:** $x = a_7a_6a_5a_4a_3a_2a_1a_0$
  $y = b_7b_6b_5b_4b_3b_2b_1b_0$
- **And we want the 8-bit word:** $z = a_3a_2a_1a_0b_3b_2b_1b_0$

  - A: z = (x >> 4) | (y << 4)
  - B: z = (x & (0x0f << 4)) | (y & 0xf) $a_7a_6a_5a_4b_3b_2b_1b_0$
  - C: z = (x >> 4) | (y & 0xf)
  - D: z = (x & 0xf0) | (y & 0xf)
  - E: z = (x << 4) | (y & 0x0f)

Binary addition with 2's Complement

- **You can add two's complement numbers just as if they are unsigned numbers.**
  - Recall, this was the whole reason for this representation

	1	1					
		0	1	0	1	1	11
+		1	1	1	0	0	+ (-4)
		0	0	1	1	1	7

4 + 2 + 1 = 7

Subtraction

- **We can implement subtraction by negating the 2nd input and then adding:**

	0	1	1	0	1	13
-	0	1	0	1	0	- 10

	1	1				
	0	1	1	0	1	13
+	1	0	1	1	0	+ (-10)
	0	0	0	1	1	3

10101
+1
10110

Why does this work?

- For n-bit numbers, the negation of B in two's complement is 2^n - B (this is alternative way of negating a 2's-complement number).

$$\begin{aligned} A - B &= A + (-B) \\ &= A + (2^n - B) \\ &= (A - B) + 2^n \end{aligned}$$

- If $A \geq B$, then (A - B) is a positive number, and 2^n represents a carry out of 1. Discarding this carry out is equivalent to subtracting 2^n, which leaves us with the desired result (A - B).
- If $A < B$, then (A - B) is a negative number and we have 2^n - (A - B). This corresponds to the desired result, -(A - B), in two's complement form.

2's Complement Subtraction

```
    1  1  0  1
-   1  0  1  0
--------------
```

1
1101
+ 0110
0011

- A: 0111
- B: 0011
- C: 1000
- D: 0101
- E: 1001

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Overflow Review

- **Recall that when we add two numbers the result may be larger than we can represent.**

(in 5b 2's complement we can represent -16 to +15)

	0	1	0	1	1	Augend	(11)
+	0	1	1	1	0	Addend	(14)
	1	1	0	0	1	Sum	(-7)

> 25

- **The same thing can happen when we add negative numbers.**

	1	1	0	0	1	Augend	(-7)
+	1	0	1	0	0	Addend	(-12)
	0	1	1	0	1	Sum	(13)

> -19

How can we know if overflow has occurred?

- **The easiest way to detect signed overflow is to look at all of the sign bits.**

	0100	(+4)		1100	(-4)
+	0101	(+5)	+	1011	(-5)
	1001	(-7)		0111	(+7)

- **Overflow occurs only in the two situations above:**
  - If you add two *positive* numbers and get a *negative* result.
  - If you add two *negative* numbers and get a *positive* result.

- **Overflow cannot occur if you add a positive number to a negative number. Do you see why?**

Overflow

- **In which circumstance can overflow <u>not</u> occur?**
  - A: subtracting a positive number from a negative number
  - B: subtracting a negative number from zero
  - C: adding two negative numbers
  - D: subtracting a negative number from a positive number
  - E: subtracting a negative number from a negative number

Overflow in software (e.g., Java programs)

```
public class overflow {
  public static void main(String[] args) {
      int i = 0;
      while (i >= 0) {
         i++;
      }
      System.out.println("i = " + i);
      i--;
      System.out.println("i = " + i);
      i++;
      System.out.println("i = " + i);
 }}
```

Output:

i = -2147483648 2^{31}

i = 2147483647 $2^{31}-1$

i = -2147483648