

# Today's announcements:

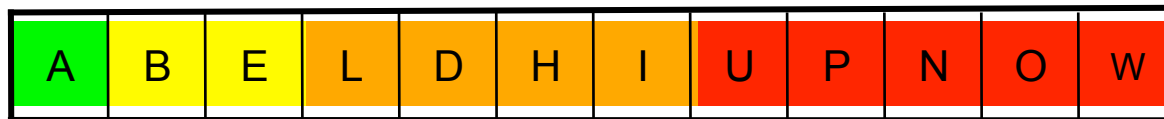
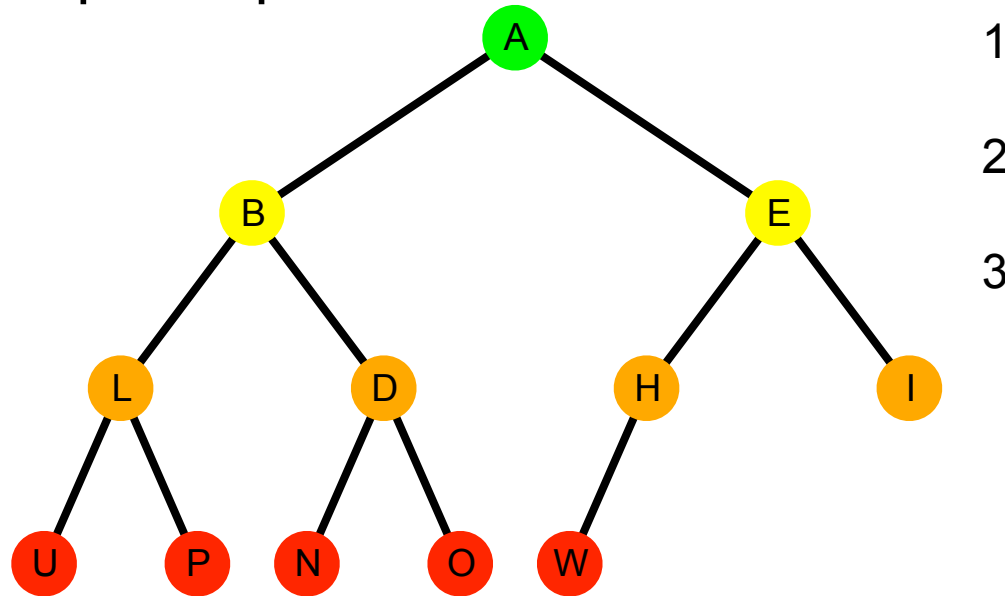
MP7 available, due 4/30, 11:59p. EC due 4/19.

Code Challenge #4, 4/17, 9p, Siebel 0224.



This image reminds us of a \_\_\_\_\_,  
which is one way we can implement ADT \_\_\_\_\_,  
whose functions include \_\_\_\_\_ and \_\_\_\_\_,  
whose running times are \_\_\_\_\_.  
This structure can be built in time \_\_\_\_\_.

(min)Heap: heapSort



Running time?

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Why do we need another  
sorting algorithm?

•

An example:

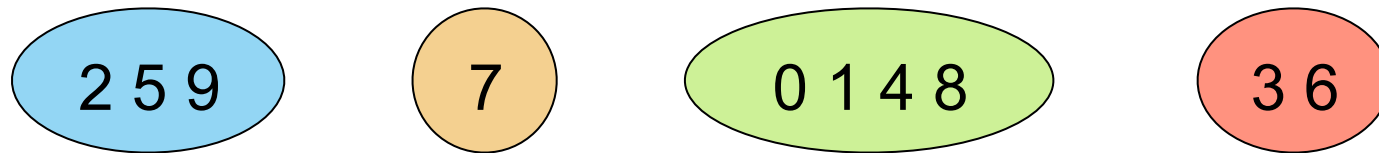
Let  $R$  be an equivalence relation on the set of students in this room, where  $(s, t) \in R$  if  $s$  and  $t$  have the same favorite among  $\{AB, FN, DJ, ZH, \_\_\_\_\_\}$ .

Notation from math:  $[\_\_\_\_]_R = \{x : xR\_\_\_\_\_\}$

One big goal for us: Given  $s$  and  $t$  we want to determine if  $sRt$ .

## A Disjoint Sets example:

Let  $R$  be an equivalence relation on the set of students in this room, where  $(s, t) \in R$  if  $s$  and  $t$  have the same favorite among  $\{AB, FN, DJ, ZH, \_\_\_\}$ .



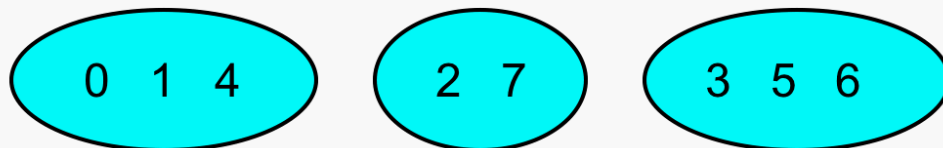
1. Find(4)
2. Find(4)==Find(8)
3. If  $!(\text{Find}(7) == \text{Find}(2))$  then Union(Find(7), Find(2))

## Disjoint Sets ADT

We will implement a data structure in support of “Disjoint Sets”:

- Maintains a collection  $S = \{s_0, s_1, \dots, s_k\}$  of disjoint sets.
- Each set has a representative member.
- Supports functions:  
    `void MakeSet(const T & k);`  
    `void Union(const T & k1, const T & k2);`  
    `T & Find(const T & k);`

A first data structure for Disjoint Sets:



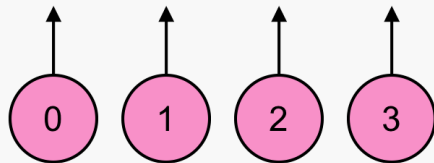
0	1	2	3	4	5	6	7
A	A	B	C	A	C	C	B

Find:

Union:

## A better data structure for Disjoint Sets: UpTrees

- if array value is -1, then we've found a root, o/w value is index of parent.
- x and y are in the same tree iff they are in the same set.



0	1	2	3
-1	-1	-1	-1

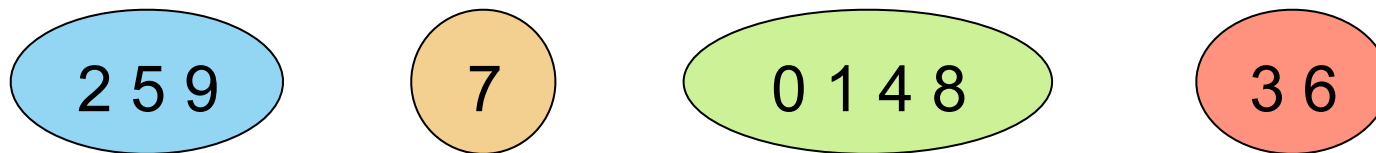
0	1	2	3

0	1	2	3

0	1	2	3

## A Disjoint Sets example:

Let  $R$  be an equivalence relation on the set of students in this room, where  $(s, t) \in R$  if  $s$  and  $t$  have the same favorite among  $\{AB, FN, DJ, ZH, PvZ\}$ .



0	1	2	3	4	5	6	7	8	9
4	8	5	6	-1	-1	-1	-1	4	5

1. Find(4)
2. Find(4)==Find(8)
3. If (!(Find(7)==Find(2))) then Union(Find(7),Find(2))

## A better data structure for Disjoint Sets:

```
int DS::Find(int i) {  
    if (s[i] < 0) return i;  
    else return Find(s[i]);  
}
```

Running time depends on \_\_\_\_\_.

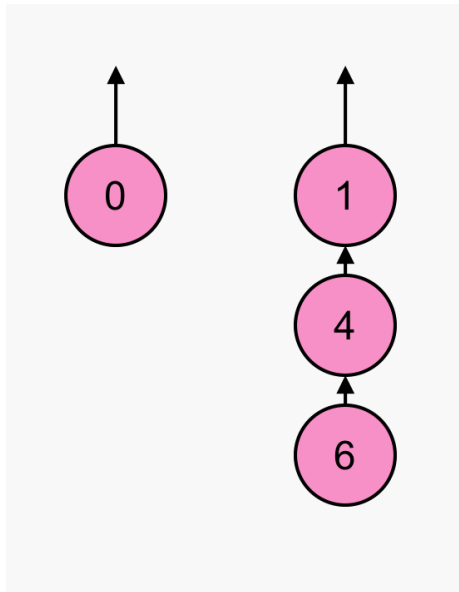
Worst case?

What's an ideal tree?

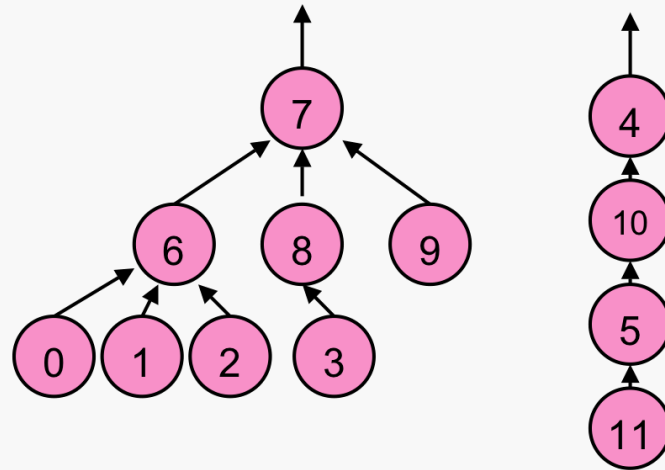
```
void DS::Union(int root1, int root2) {  
    _____;  
}
```



something to consider...



## Smart unions:



Union by height:

0	1	2	3	4	5	6	7	8	9	10	11
6	6	6	8		10	7		7	7	4	5

*Keeps overall height of tree as small as possible.*

Union by size:

0	1	2	3	4	5	6	7	8	9	10	11
6	6	6	8		10	7		7	7	4	5

*Increases distance to root for fewest nodes.*

Both of these schemes for Union guarantee the height of the tree is \_\_\_\_\_.

## Smart unions:

```
int DS::Find(int i) {  
    if (s[i] < 0) return i;  
    else return Find(s[i]);  
}
```

```
void DS::UnionBySize(int root1, int root2) {  
    int newSize = s[root1]+s[root2];  
    if (isBigger(root1,root2)) {  
        s[root2]= root1;  
        s[root1]= newSize;  
    }  
    else {  
        s[root1] = root2;  
        s[root2]= newSize;  
    }  
}
```