

CS411 Database Systems

Kevin C. Chang

Announcements

• If you are waiting to register for this class.

gane override.

Question 1: How to "query" a database?

Goal: specify what we want from our database
 Find all the employees who earn more than \$50,000 and pay taxes in Champaign-Urbana.

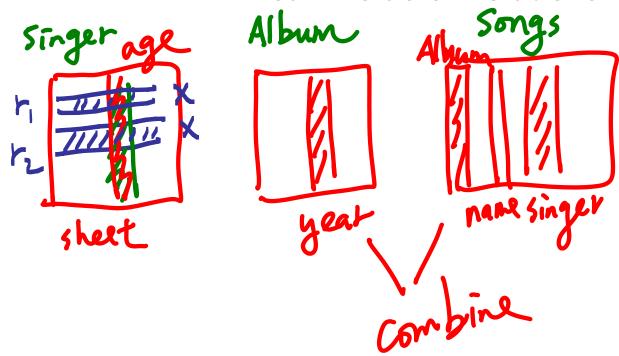
Write in C++/Java?

What I want guages: Not How

- Use high-level query languages:
 - Theoretical: Relational Algebra, Datalog
 - Practical: SQL

• Relational algebra: a basic set of operations on relations that provide the *basic principles*.

Question 2: What kinds of "computation" can we do on relations?



What is an "Algebra"
Relation Mathematical system consisting of: • Operands -- variables or values from which new values can be constructed. **Operators --- symbols denoting procedures that construct new values from given values. Relations (and only

Q: Example algebra?

- Arithmetic algebra. Linear algebra.
- What are operands?
- What are operators?

What is Relational Algebra?

- An algebra whose operands are relations or variables that represent relations.
- Operators are designed to do common things that we need to do with relations in a database.
- The result is an algebra that can be used as a *query language* for relations.

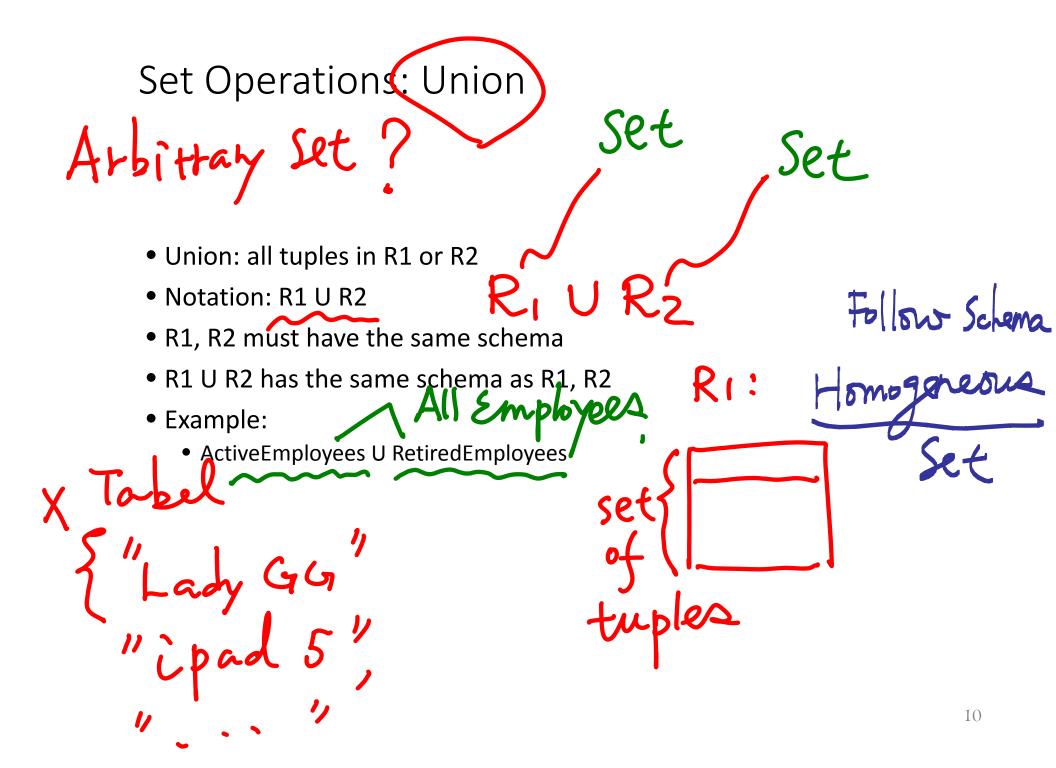
Relational Algebra at a Glance

Operators: relations as input, new relation as output

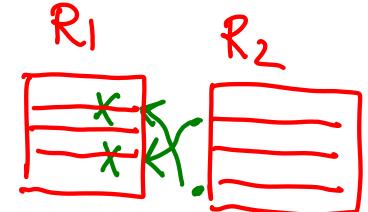
- Basic RA operators:
 - Basic Set Operations
 - union, difference (no intersection, no complement)
 - Selection: σ
 - Projection: π
 - Cartesian Product: X
 - Renaming: ρ
- **Derived** perations:
 - Intersection, complement
 - Joins (natural, equi-join, theta join, semi-join)

Set Operations

- Union, difference
- Binary operations





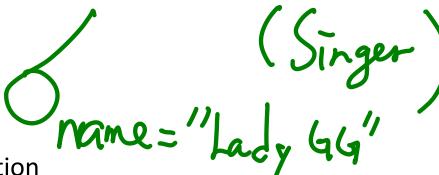


- Difference: all tuples in R1 and not in R2
- Notation: R1 R2
- R1, R2 must have the same schema
- R1 R2 has the same schema as R1, R2
- Example
 AllEmployees RetiredEmployees

Active Employees

Store?





• Returns all tuples which satisfy a condition

• Notation: $\sigma_c(R)$

• *c* is a condition: =, <, >, and, or, not

• Output schema: same as input schema

• Find all employees with salary more than \$40,000:

• $\sigma_{salary>40000}$ (Employee)



C: salary > 40000

Selection Example

SSN	Name	DepartmentID	Salary
92881	John	1	30,000
32877	Tony	1	32,000
44544	Amy	2	45,000

$$\sigma_{salary>40000}(Employee) = ?$$

Projection π

- Unary operation: returns certain columns
- Eliminates duplicate tuples!
- Notation: $\pi_{A_1,...,A_n}(R)$
- Condition:
 - Suppose input schema $R(B_1, ..., B_m)$
 - $\bullet \{A_1, \dots, A_n\} \subseteq \{B_1, \dots, B_m\}$
- Output schema $S(A_1, ..., A_n)$
- Example: $\pi_{SSN,Name}(Employee)$

Projection Example

SSN	Name	DepartmentID	Salary
92881	John	1	30,000
32877	Tony	1	32,000
44544	Amy	2	45,000

$$\pi_{SSN,Name}(Employee) = ?$$



- Think of relation as a table.
- How are they similar?
- How are they different?
- Why do you need both?

Cartesian Product

- Each tuple in R1 with each tuple in R2
- Notation: $R_1 \times R_2$
- Input schemas $R_1(A_1, ..., A_n)$, $R_2(B_1, ..., B_m)$.
- Condition: $\{A_1, \dots, A_n\} \cap \{B_1, \dots, B_m\} = \emptyset$
- Output schema is $S(A_1, ..., A_n, B_1, ..., B_m)$.
- Example: Employee x Dependents
- Very rare in practice (but joins are very common)

Cartesian Product Example

Employee

Name	SSN
John V	99999999
Tony 🗸	7777777

2



EmployeeSSN	Dname
99999999	Emily Q
77777777 🖊	Joe O

x =4

2

Employee x Dependents

Name 4	SSN	EmployeeSSN	Dname
John 🗸	99999999	99999999	Emily •
John V	99999999	77777777	Joe
	77777777		Emily 2
<i>-</i>	77777777		Joe O
			1