How do we evaluate computer architectures?

- Think of 5 characteristics that differentiate computers?
 - Can some processors compute things that others can't? N_0



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Two notions of performance

| Aircraft | DC to Paris | Passengers | |
|----------|-------------|------------|--------------|
| 747 | 6 hours | 500 | - throughput |
| Concorde | 3 hours | 125 | - latercy |

- Which has higher performance?
- From a passenger's viewpoint: latency (time to do the task)
 - hours per flight, execution time, response time
- From an airline's viewpoint: throughput (tasks per unit time)
 - passengers per hour, bandwidth
- Latency and throughput are often in opposition

Some Definitions

Relative performance: "x is N times faster than y"

$$\frac{\text{Performance}(x)}{\text{Performance}(y)} = N$$

If we are primarily concerned with latency,

Performance(x) =
$$\frac{1}{\text{Latency}(x)}$$

If we are primarily concerned with throughput,

Performance(x) = throughput(x)

CPU performance

- The obvious metric: how long does it take to run a test program? This depends upon three factors:
- 1. The number of *dynamic* instructions N in the program
 - Executing more instructions tends to take longer.
- 2. The kind of instructions in the program
 - Some instructions take more CPU cycles than others
 - Let(c)be the average number of cycles per instruction (CPI)
- 3. The time t per CPU clock cycle (clock-cycle time)

CPU time = Instructions executed
$$\times$$
 CPI \times Clock cycle time

The three components of CPU performance

- Instructions executed:
 - the dynamic instruction count (#instructions actually executed)
 - not the (static) number of lines of code
- Average Cycles per instruction:
 - function of the machine and program
 - CPI(floating-point operations) > CPI(integer operations)
 - Improved processor may execute same instructions in fewer cycles
 - Single-cycle machine: each instruction takes 1 cycle (CPI = 1)
 - CPI can be > 1 due to memory stalls and slow instructions
 - CPI can be < 1 on superscalar machines
- Clock cycle time: 1 cycle = minimum time it takes the CPU to do any work
 - clock cycle time = 1/ clock frequency
 - 500MHz processor has a cycle time of 2ns (nanoseconds)
 - 2GHz (2000MHz) CPU has a cycle time of just 0.5ns
 - higher frequency is usually better

Execution time, again

CPU time = Instructions executed × CPI × Clock cycle time

Make things faster by making any component smaller!

| | Program | Compiler | ISA | Organization | Technology |
|----------------------|---------|----------|--------|--------------|------------|
| Instruction Executed | X | X | X | | |
| СРІ | X | × | X | X | ~ |
| Clock Cycle Time | ~ | N | \sim | × | X |

Often easy to reduce one component by increasing another

A: Yes B: No

Example 1: ISA-compatible processors

- Let's compare the performances two x86-based processors.
 - An 800MHz AMD Duron, with a CPI of 1.2 for an MP3 compressor.
 - A 1GHz Pentium III with a CPI of 1.5 for the same program.
- Compatible processors implement identical instruction sets and will use the same executable files, with the same number of instructions.
- But they implement the ISA differently, which leads to different CPIs.

Example 2: Comparing across ISAs

- Intel's Itanium (IA-64) ISA is designed facilitate executing multiple instructions per cycle. If an Itanium processor achieves an average CPI of .3 (3 instructions per cycle), how much faster is it than a Pentium4 (which uses the x86 ISA) with an average CPI of 1?
 - a) Itanium is three times faster
 - b) Itanium is one third as fast
 - c) Not enough information

Clock freq