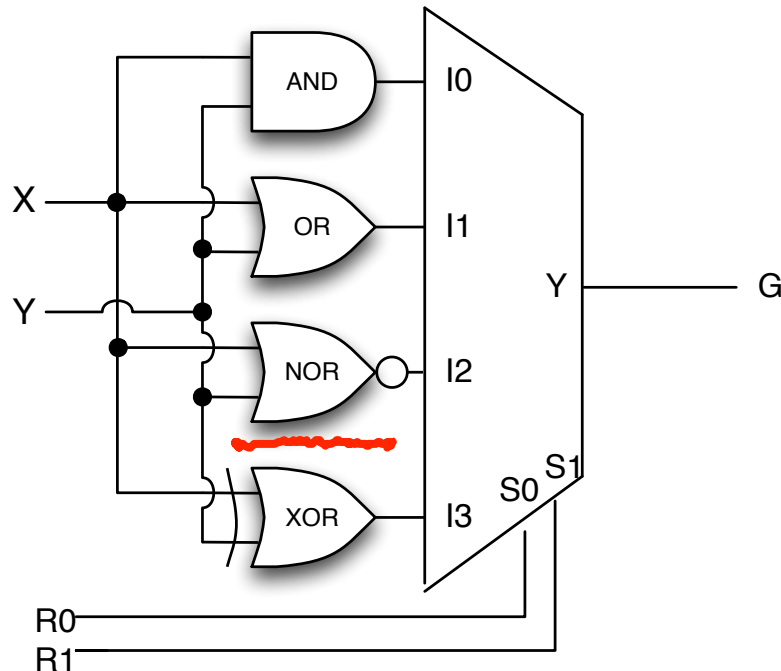


Building an ALU (Part 2):

GRAB A
HANDOUT

Why NOR is more general than NOT:

- We've changed the ALU design to include NOR over NOT



x	NOT(x)
0	1
1	0

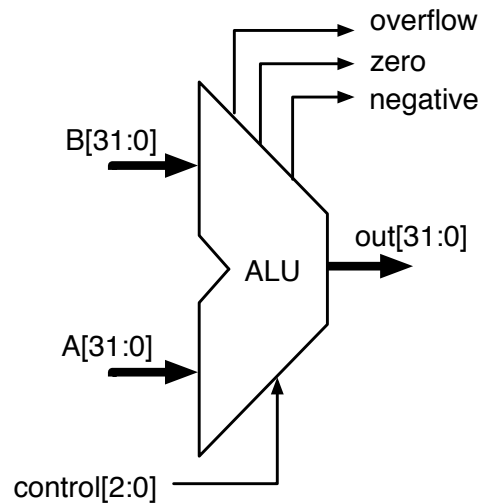
x	y	NOR(x,y)
0	0	1
0	1	0
1	0	0
1	1	0

- We can still implement NOT:
 - NOT(X) is implemented by using NOR(X, Y) & setting Y=0

Today's lecture

- **We'll finish the 32-bit ALU today!**
 - 32-bit ALU specification
- **Complete 1-bit ALU**
- **Assembling them to make 32-bit ALU**
- **Handling flags:**
 - zero, negative, overflow

Building 32-bit ALU



control	out =
0	undefined
1	undefined
<u>2</u>	<u>A + B</u>
3	A - B
4	A AND B
5	A OR B
6	A NOR B
7	A XOR B

Adder
logic unit

```

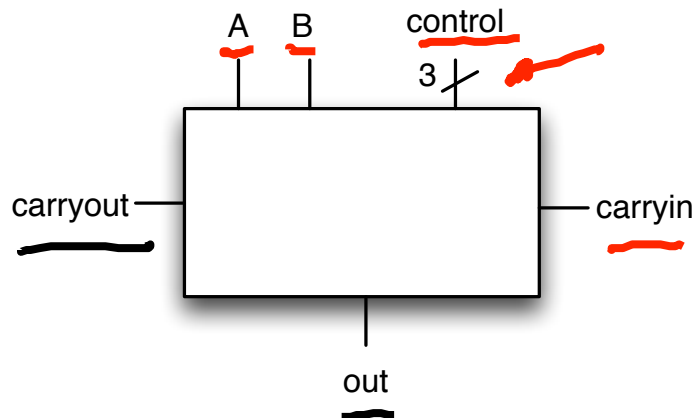
module alu32(out, overflow, zero, negative,
             A, B, control);
    output[31:0] out;
    output      overflow, zero, negative;
    input  [31:0] A, B;
    input   [2:0] control;

```

Did overflow occur?
Is the output equal to zero?
Is the output negative?

We want to create a 1-bit ALU

- Previously we showed 1-bit adder/subtractor, 1-bit logic unit
 - Time to put them together.



control	out _i =
0	undefined
1	undefined
2	$X_i + Y_i$
3	$X_i - Y_i$
4	$X_i \text{ AND } Y_i$
5	$X_i \text{ OR } Y_i$
6	$X_i \text{ NOR } Y_i$
7	$X_i \text{ XOR } Y_i$

↓

010
011
100
101
110
111

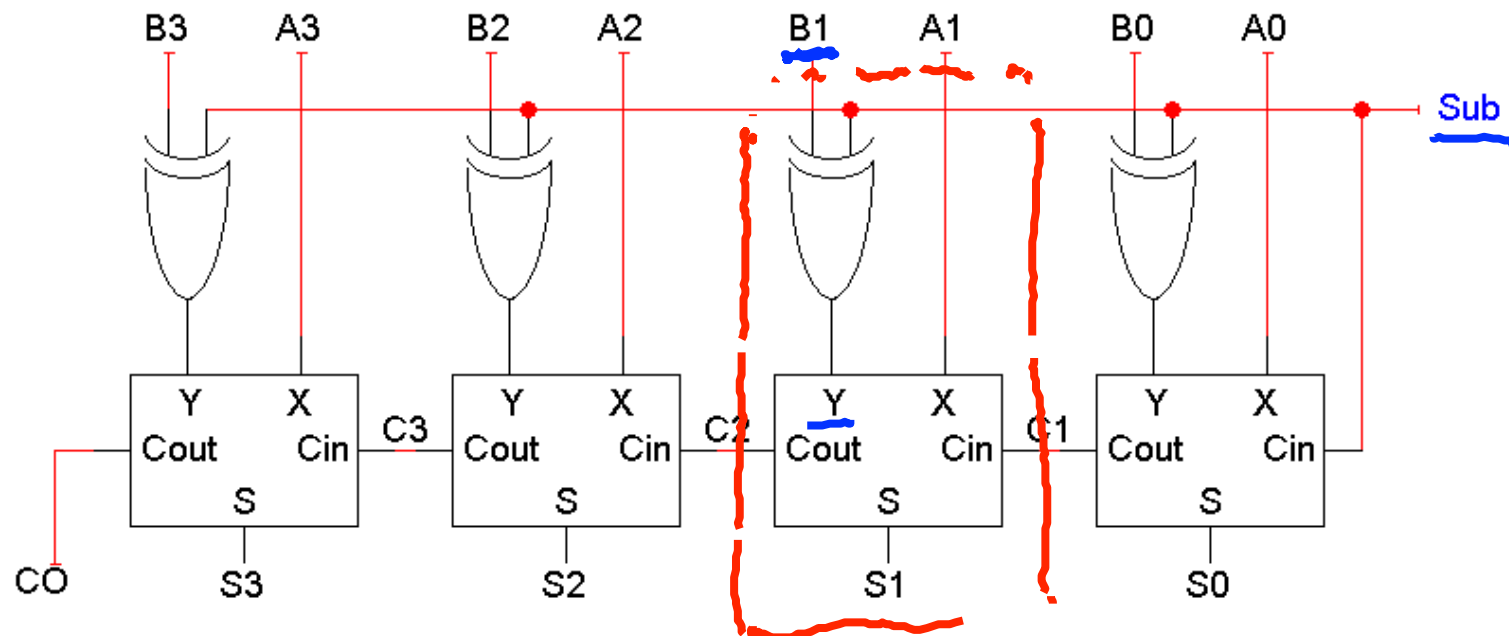
```

module alu1(out, carryout, A, B, carryin, control);
    output out, carryout;
    input A, B, carryin;
    input [2:0] control;

```

Addition + Subtraction in one circuit

- When Sub = 0, Y = B and Cin = 0. Result = $A + B + 0 = A + B$.
- When Sub = 1, Y = ~B and Cin = 1. Result = $A + \sim B + 1 = A - B$.

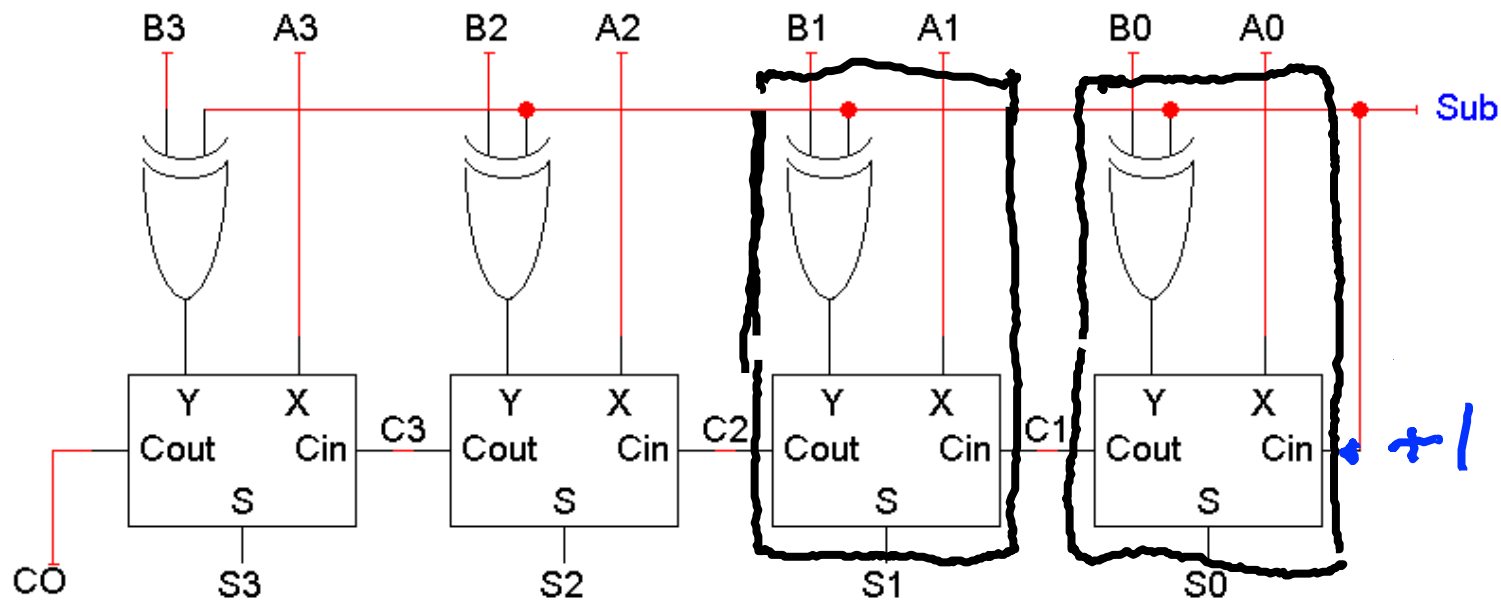


- Which parts belong in inside the 1-bit ALU?
A) the Full Adder, B) the XOR gate, **C) Both**, D) Neither

$$A - B = A + (-B) = A + \sim B + 1$$

Addition + Subtraction in one circuit

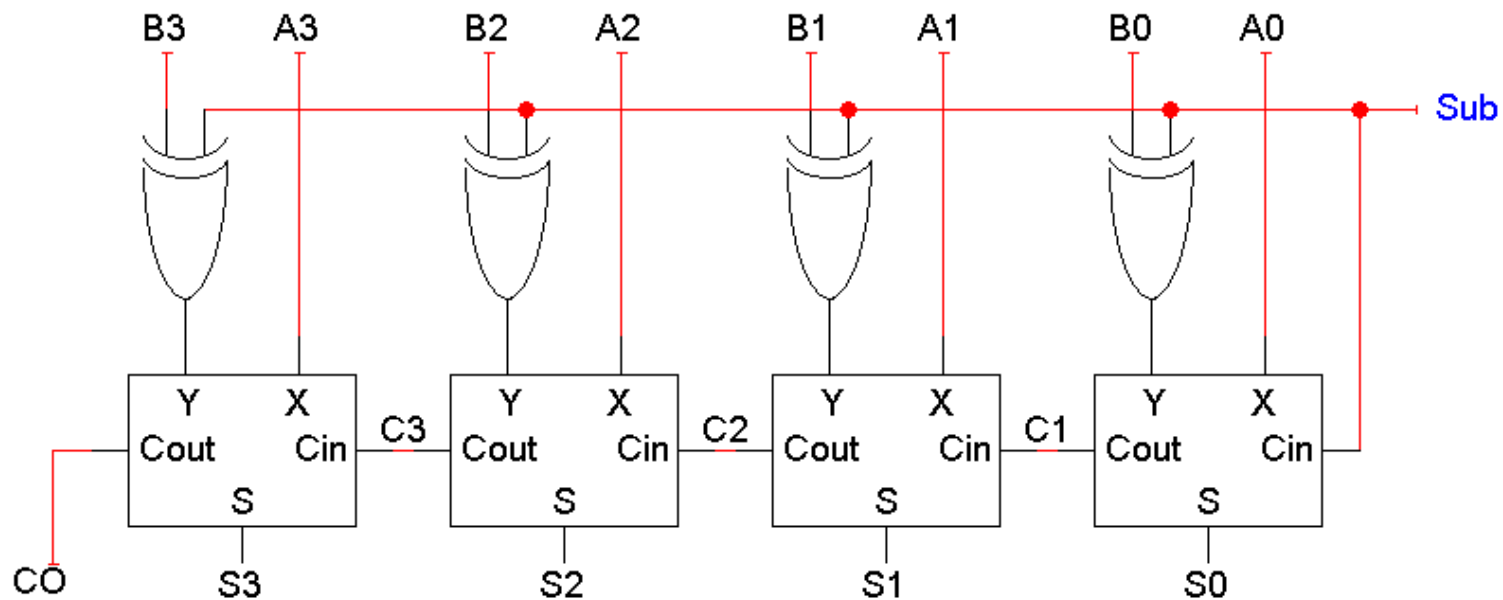
- When $\text{Sub} = 0$, $Y = B$ and $\text{Cin} = 0$. Result = $A + B + 0 = A + B$.
- When $\text{Sub} = 1$, $Y = \sim B$ and $\text{Cin} = 1$. Result = $A + \sim B + 1 = A - B$.



- What should we do with the full adder's Cin input?
 - A) Connect to Sub, B) Connect to 1-bit ALU's carryin

Addition + Subtraction in one circuit

- When $\text{Sub} = 0$, $Y = B$ and $\text{Cin} = 0$. Result = $A + B + 0 = A + B$.
- When $\text{Sub} = 1$, $Y = \sim B$ and $\text{Cin} = 1$. Result = $A + \sim B + 1 = A - B$.



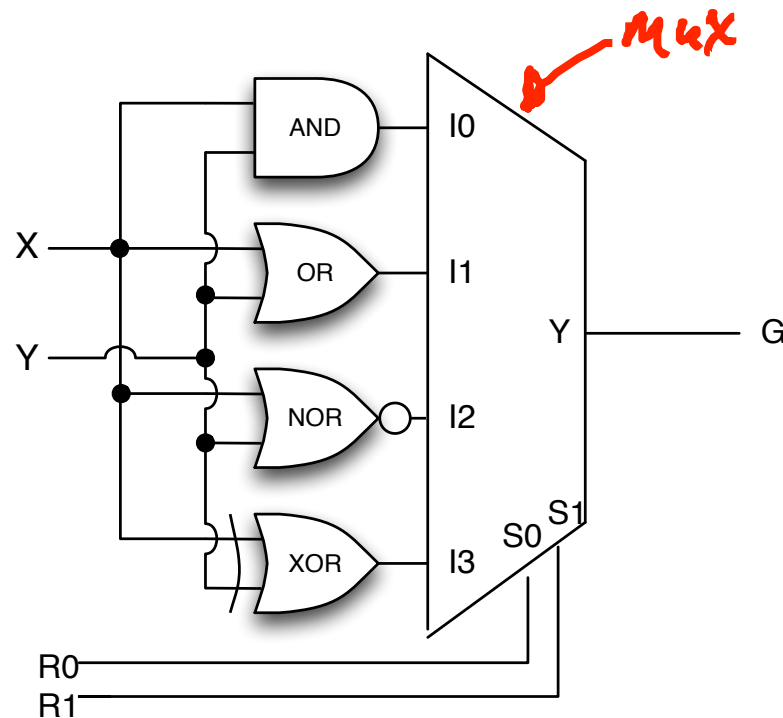
- Where will the “Sub” signal come from?

Control [0]

$$R_1 = \text{Control}[1]$$

$$R_0 = \text{Control}[0]$$

Complete 1-bit Logic Unit



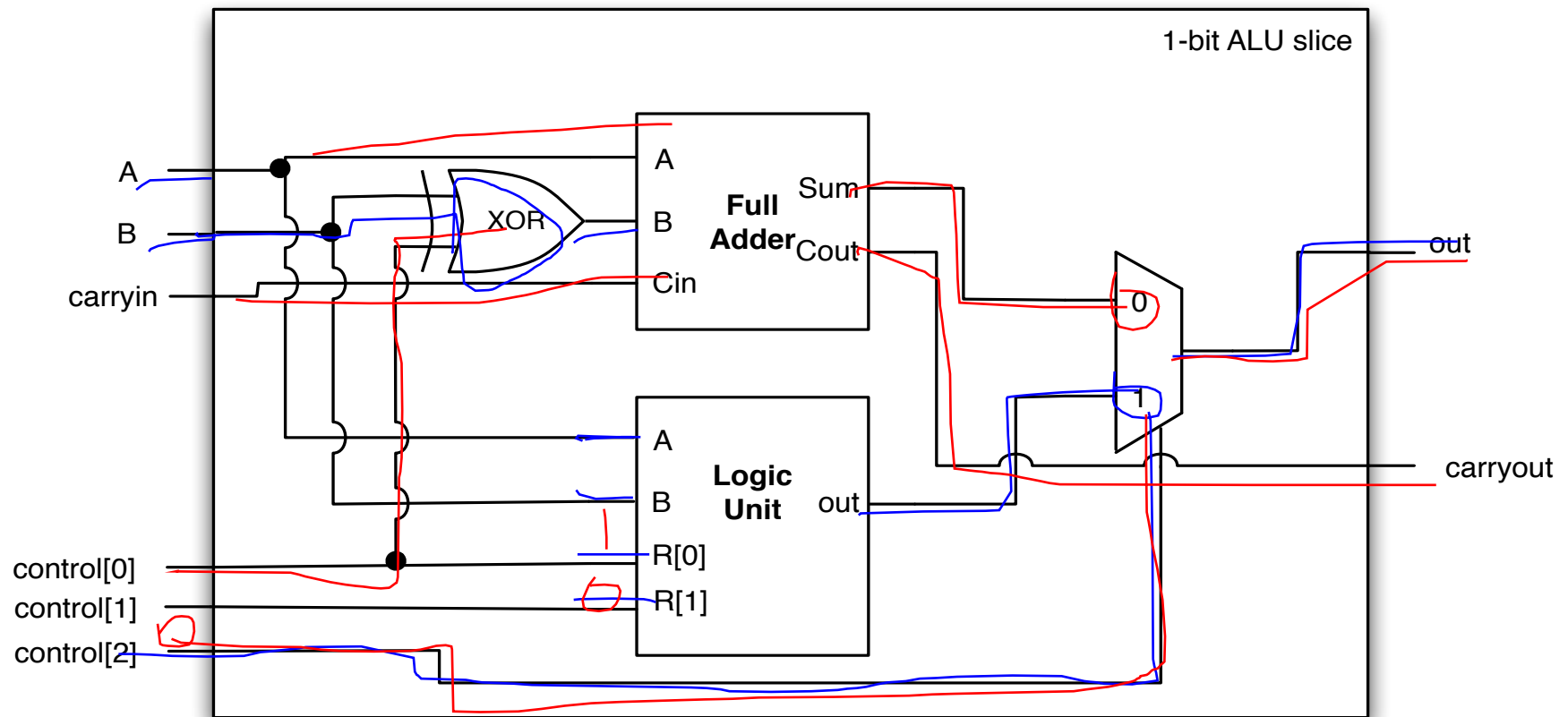
R_1	R_0	Output
0	0	$G_i = X_i Y_i$
0	1	$G_i = X_i + Y_i$
1	0	$G_i = (X_i + Y_i)'$
1	1	$G_i = X_i \oplus Y_i$

- What should the control inputs (R_0 , R_1) connect to?
- How do we select between the adder and the logic unit?
- How do we control the selection?

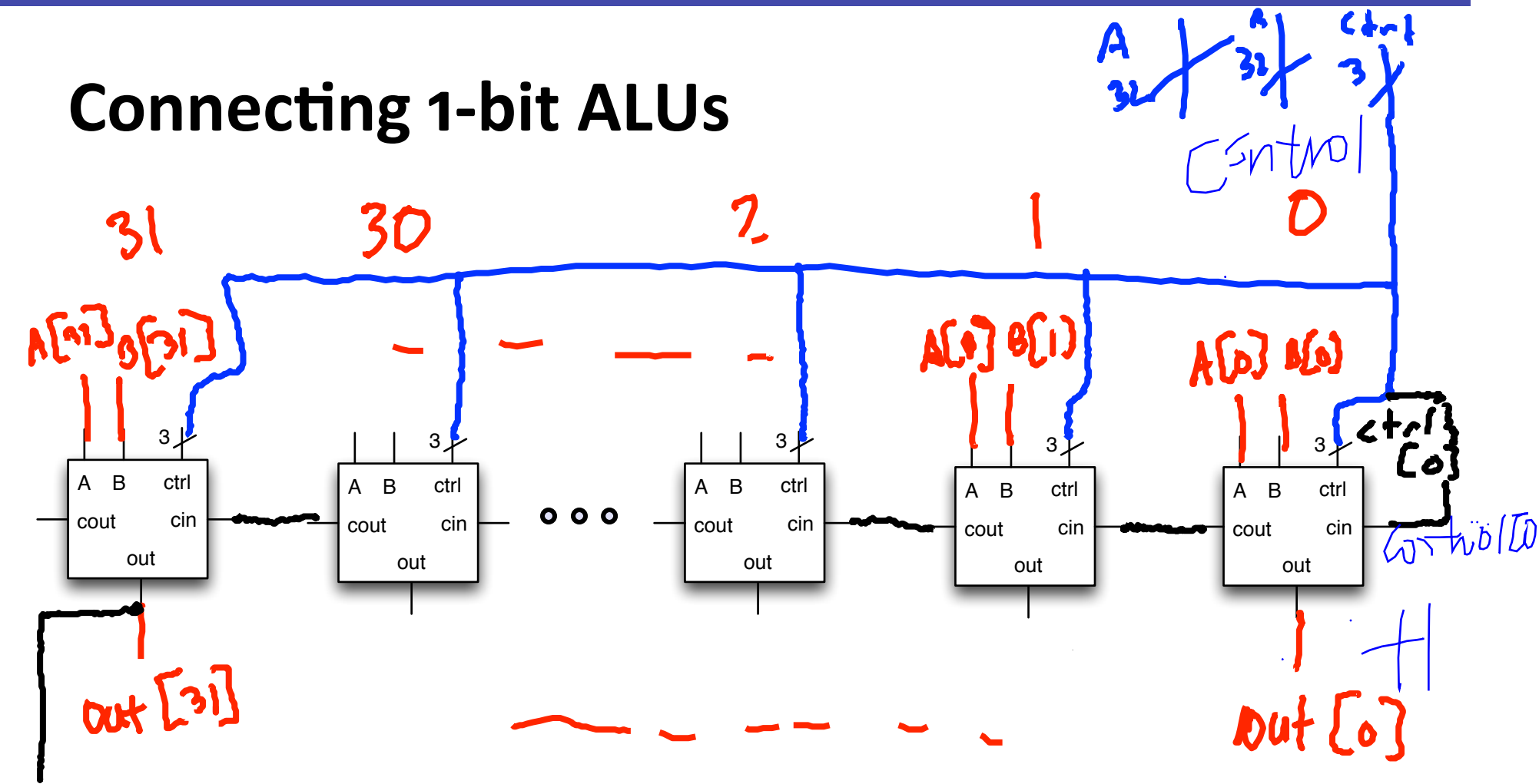
Control [2]

MUX

Complete 1-bit ALU



Connecting 1-bit ALUs



negative

1000101

Δ

$A - B$

$A + (\sim B + 1)$

subtract.

Flags (overflow, zero, **negative**)

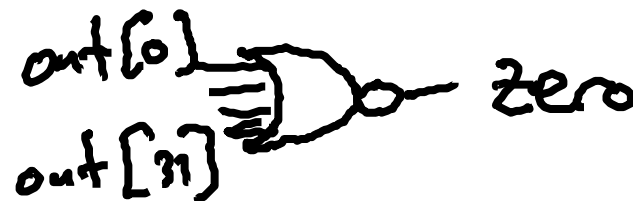
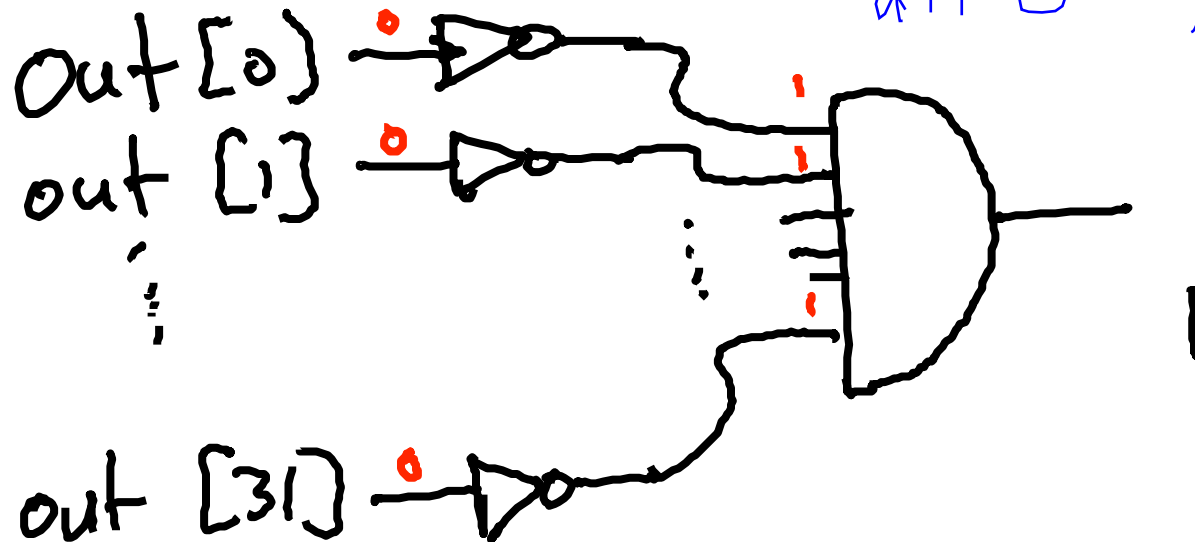
- Let's do negative first; negative evaluates to:
 - 1 when the output is negative, and
 - 0 when the output is positive or zero
- **Negative** =
 - a) carryout[30]
 - b) output[30]
 - c) carryout[31]
 - d) output[31]
 - e) control[0]

Flags (overflow, zero, negative)

- zero evaluates to:

- 1 when the output is equal to zero, else 0

- Zero =



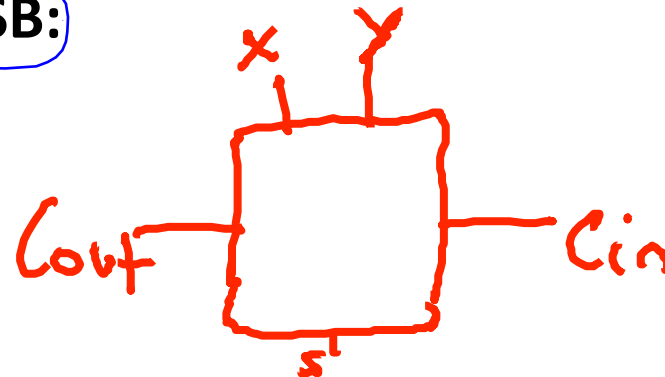
Flags (**overflow**, zero, negative)

- **Overflow evaluates to:**

- 1 when the overflow occurred, else 0
 - adding two positive numbers yields a negative number
 - adding two negative numbers yields a positive number

- **Consider the adder for the MSB:**

X	Y	C _{in}	C _{out}	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1



- a) cin[31] NOR cout[31]
- b) cin[31] AND cout[31]
- c) cin[31] OR cout[31]
- d) cin[31] XOR cout[31]
- e) cin[31] NAND cout[31]

- **Overflow =**