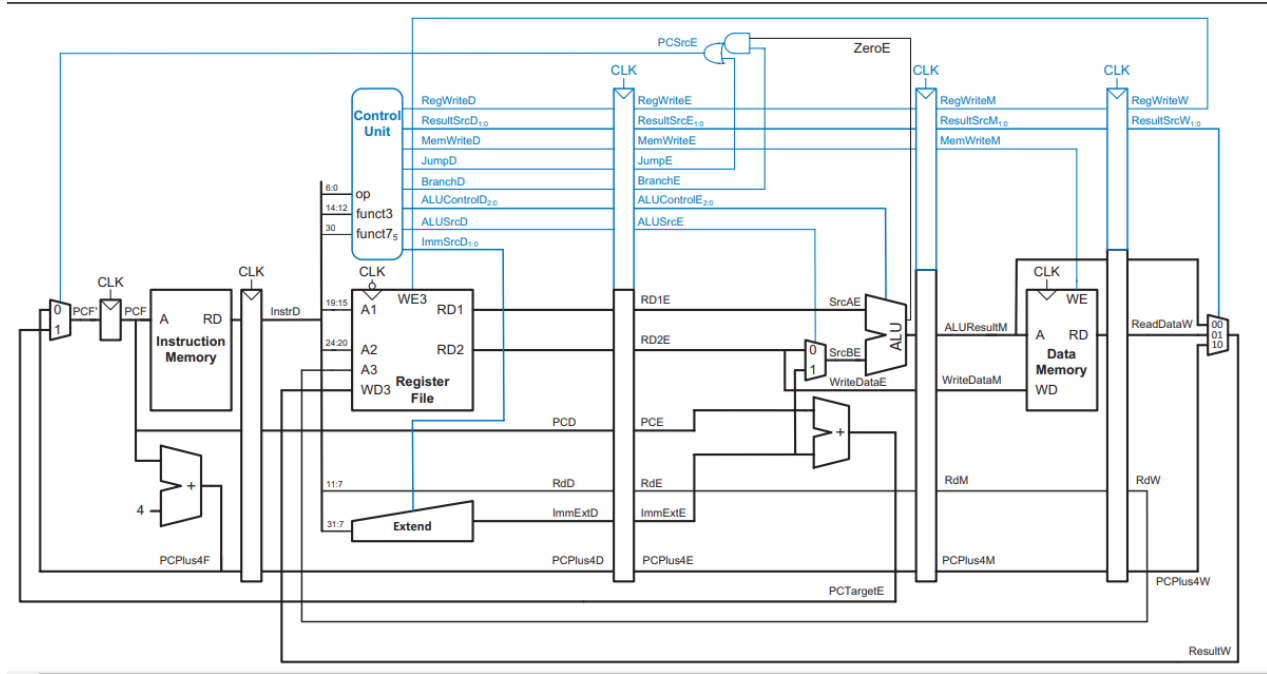


Pipeline Architecture

Datpath:



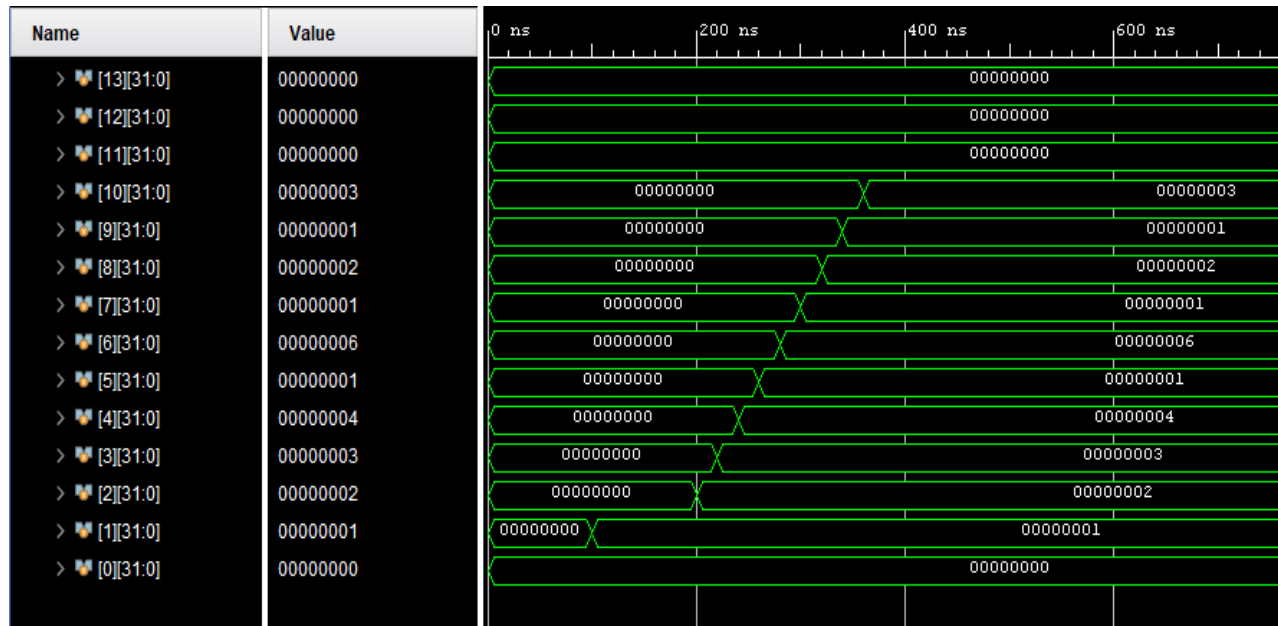
I followed this Datapath and convert my single cycle processor into pipeline architecture by adding four clock driven registers. Then I test it on a instruction memory which contains no hazards.

Testing Instruction Memory

```
32'h00: instruction = 32'h00100093; //addi x1 , x0 , 1
32'h04: instruction = 32'h00200113; //addi x2 , x0 , 2
32'h08: instruction = 32'h00300193; //addi x3 , x0 , 3
32'h0c: instruction = 32'h00400213; //addi x4 , x0 , 4
32'h10: instruction = 32'h401102b3; //sub x5 , x2 , x1
32'h14: instruction = 32'h00508313; //add x6 , x1 , 5
32'h18: instruction = 32'h402183b3; //sub x7 , x3 , x2
32'h1c: instruction = 32'h00138413; //addi x8,x7,1
32'h20: instruction = 32'h407404b3; //sub x9,x8,x7
32'h24: instruction = 32'h00940533; //sub x10,x8,x9

default: instruction = 32'h00000000;
endcase
```

Testbench result:



You can check from the test bench waveform after reset become zero first result is written back after five clock cycles. After that the result is written after every clock cycle.

Pipeline Architecture with Hazard Controller

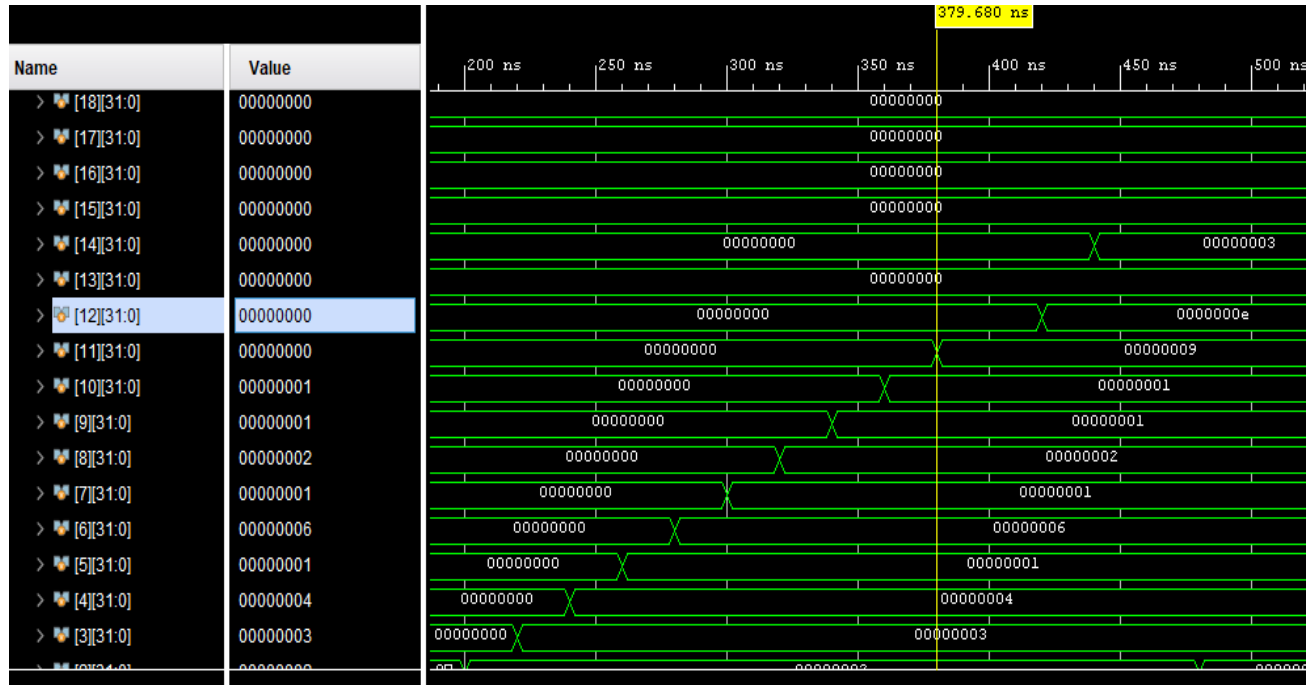
In this part, I added hazard controller unit to control the hazard and test it on a instruction memory that contain hazards.

Testing Instruction Memory:

```
○ default: instruction = 32'h00000000;*/  
○ 32'h00: instruction = 32'h00100093;//addi x1 , x0 , 1  
○ 32'h04: instruction = 32'h00200113;//addi x2 , x0 , 2  
○ 32'h08: instruction = 32'h00300193;//addi x3 , x0 , 3  
○ 32'h0c: instruction = 32'h00400213;//addi x4 , x0 , 4  
○ 32'h10: instruction = 32'h401102b3;//sub x5 , x2 , x1  
○ 32'h14: instruction = 32'h00508313;//add x6 , x1 , 5  
○ 32'h18: instruction = 32'h402183b3;//sub x7, x3 , x2  
○ 32'h1c: instruction = 32'h00138413;//addi x8,x7,1  
○ 32'h20: instruction = 32'h407404b3;//sub x9,x8,x7  
○ 32'h24: instruction = 32'h40940533;//sub x10,x8,x9  
○ 32'h28: instruction = 32'h0030a583;//lw x11,3(x1)  
○ 32'h2c: instruction = 32'h00558613;//addi x12 , x11 , 5  
○ 32'h30: instruction = 32'h00860693;//addi x13 , x12 , 8  
○ 32'h34: instruction = 32'h00368713;//addi x14 , x13 , 3  
○ 32'h38: instruction = 32'h00208663;//beq x1,x2,stop  
○ 32'h3c: instruction = 32'h40110133;//sub x2,x2,x1  
○ 32'h40: instruction = 32'hfc9ff06f;//j label  
○ 32'h44: instruction = 32'h0000006f;//j stop  
○ default: instruction = 32'h00000000;  
    endcase  
  
endmodule
```

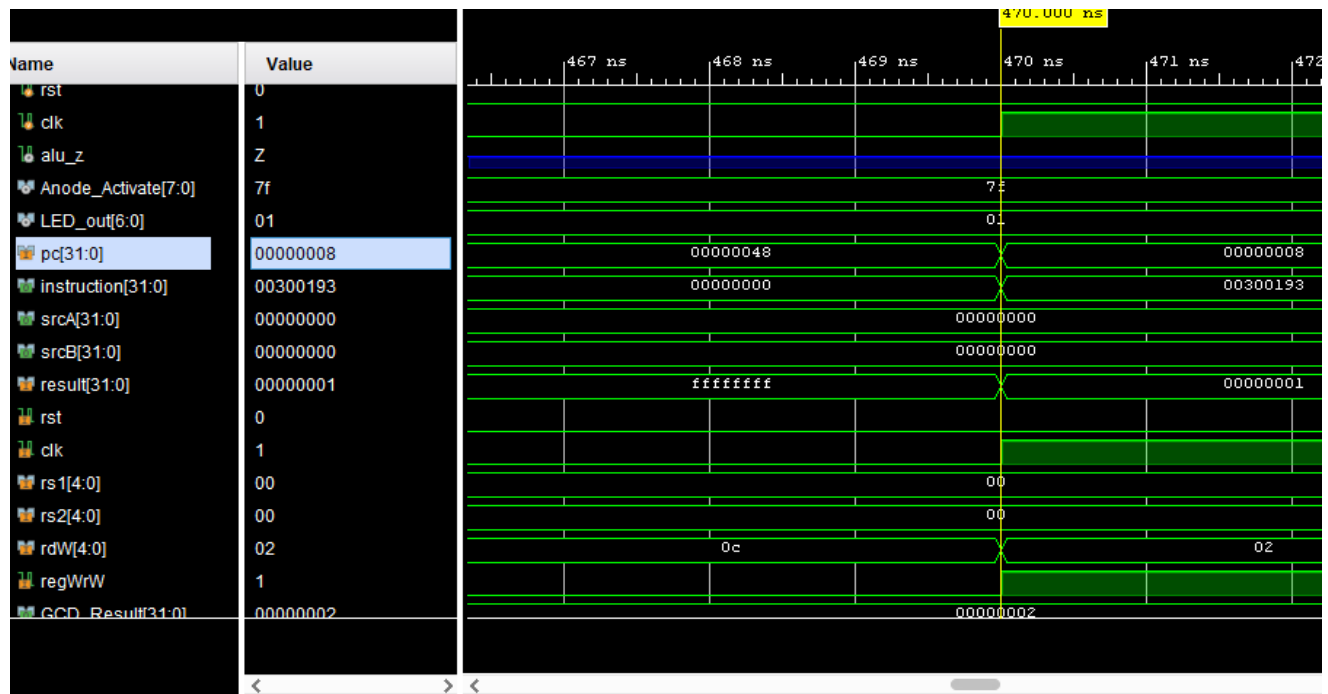
TestBench result:

RAW hazard and stalling and flushing:



Raw hazards are solved by forwarding and there is a one cycle stall when load instruction came. When result is written back to x11 after it there is a one cycle stall due to load instruction.

Controll Hazards:

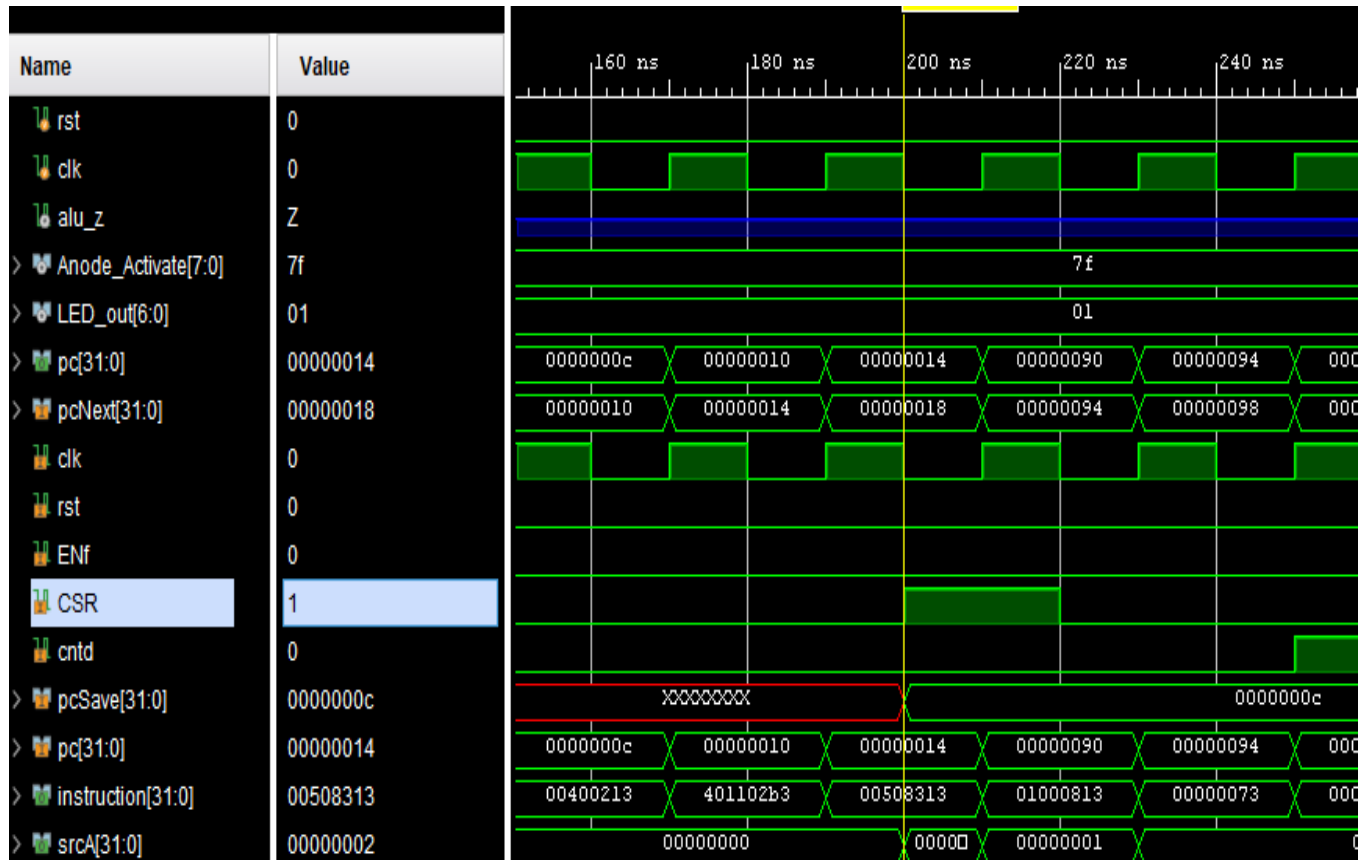


It also supports jump instruction, pc jump from 48 to 08 hen jump instruction came.

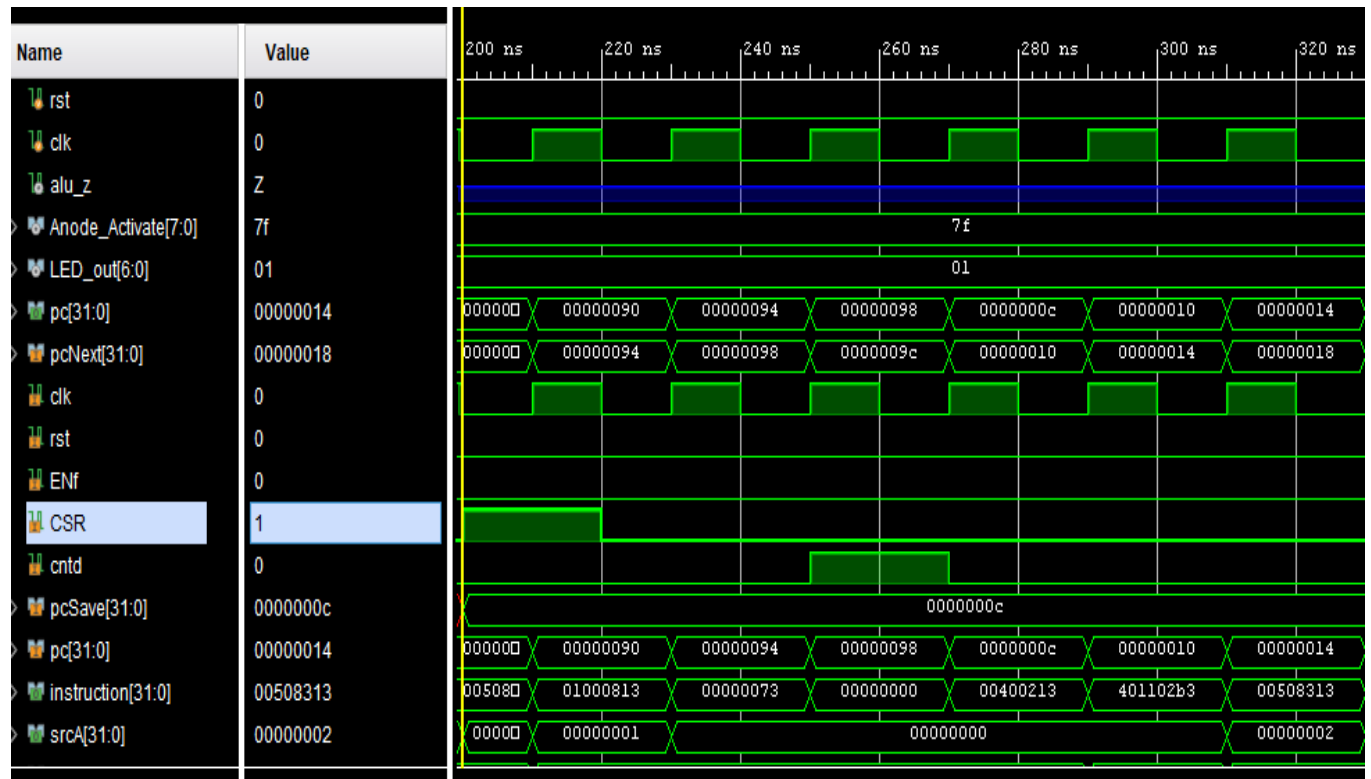


It also supports branch instruction when beq instruction came it stops as shown in instruction memory.

TestBench result:



When CSR flag is turned on pc jumps to a handler hard coded in instruction memory and pc of execution is saved in pcSave.



After the execution of intrupt handler instruction it return back to its normal execution.