

Lecture 6

Series Problem

Trick ~~(i)~~ Relation $b \mid w$ i (loop) & series

~~(ii)~~ Generate formula.

Category 1 : Basic sum

Q1

$$1 + 2 + 3 + \dots + N = (\text{Sum of natural numbers})$$

$n = 3$

Trick

i	t
1	1
2	= 2
3	= 2
4	4

\rightarrow for ($\text{int } i = 1; i \leq 3; i++$) {

$\sum = \sum + i;$

$$\sum = \sum + 1 + 2 + 3;$$

Dry Run

$\sum = 0$

$i = 1, 2, 3, 4$

$1 + 2 + 3 = 6$

$$i = t =$$

Q2

$$1 * 2 - 1 = t$$

$$1 \times 2 - 1 = 1$$

$$2 \times 2 - 1 = 3$$

5-1 Const. table

i	t
1	$1 * 2 - 1 = 1$
2	$2 * 2 - 1 = 3$
3	$3 * 2 - 1 = 5$
4	$4 * 2 - 1 = 7$

5-2 Formulate

$$2 \times 2 - 1 = 3$$

Shortcut p & q format

int sum = 0;

for (int i = 1; i <= n; i++) {

sum = sum + $i * 2 - 1$;

}

$$t = i * 2 - 1$$

```
int sum = 0;
```

```
for(int i = 1; i <= n; i++) {
```

```
    sum = sum + 2 * 3 - 1
```

}

x

-

Dry Run :

sum = 0

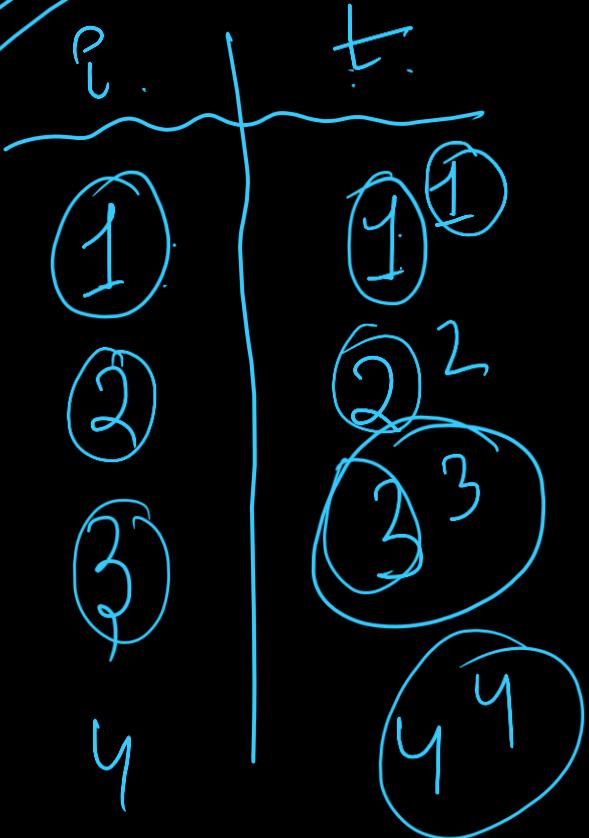
i = 1

x x y

= 9

$$\frac{Q_2}{H(w)} \cdot 1^1 + 2^2 + \underline{3^3} + 4^4 + \dots + n^n$$

~~$$1^2 + 2^2 + \underline{3^2} + 4^2 \dots + n^2 =$$~~



$$\boxed{t = 0}$$

int sum = 0;

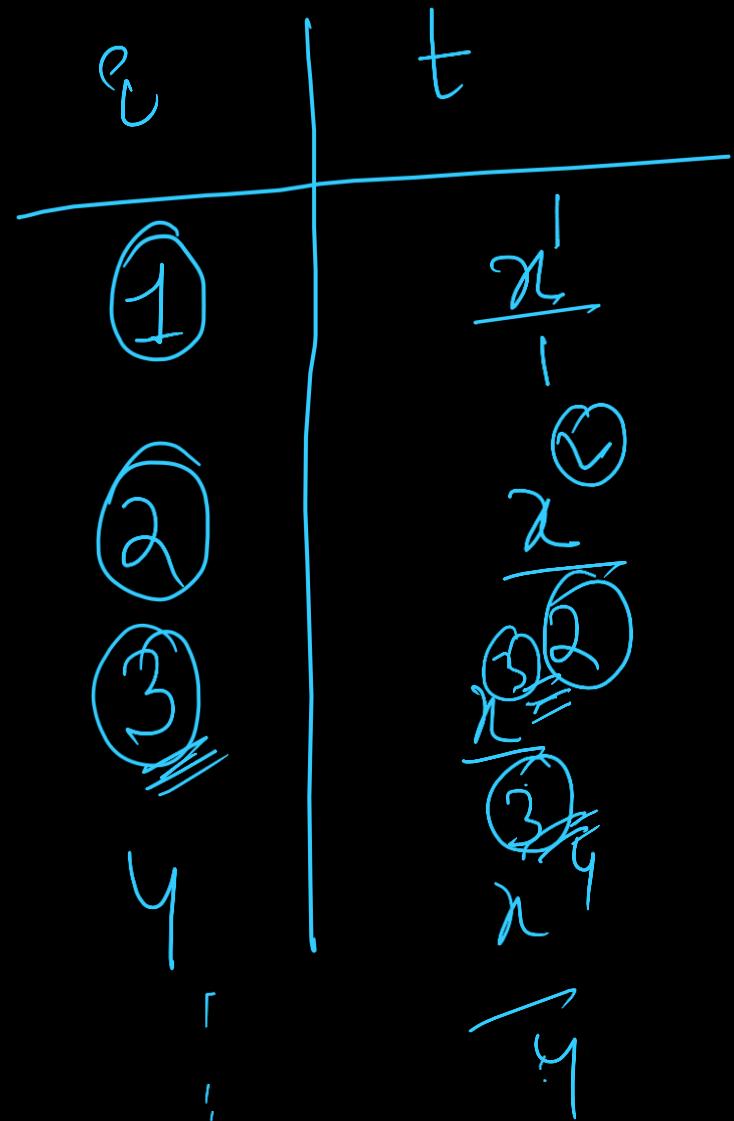
for(int i=1; i<=n; i++){

 sum = sum + pow(l, i)

}

$$1^1 + 2^2 + 3^3$$

$$Q4 \quad 1 + x + \frac{x^2}{2} + \dots + \frac{x^n}{n}$$



```
t = xi int sum = 1;
for(int i = 1; i <= n; i++) {
```

sum = sum + (pow(x, i) / i)

}

Q5

$$1 \times 2 \times 3 \times 4 \cdots n \quad (\text{h!})$$

$$2! = 1 \times 2$$

$t = i$

$$4! = ((\times 2 \times 3 \times 4))$$

int ~~prod = 0;~~ $\frac{1}{=}$
for (int $i = 1; i \leq n; i++$) {

~~prod = prod * i~~

}

#Category 2 : $t = i \times x + y$

Q1 ① + 4 + 7 + 10

$$\begin{array}{c} i \\ \text{---} \\ 1 \end{array}$$

$$1 \times b + a = 1$$

② $2 \times b + a = 4$

③ 7

$$b + a = 1$$

$$2 \times b + a = 4$$

$$- - -$$

$$b = -3$$

$$b = 3$$

$$i \times 3 - 2 = t$$

int sum = 0;

for(int i = 1; i <= n; i++) {

$$\text{sum} = \text{sum} + [i * 3 - 2]$$

Category 3 : Nested Loop.

Trick

① Formula for outer series

② Formula for inner series

③ Relation b/w i & j

Q1

S1 S2
Dry Run
1t

$$1 + (1+3) + (1+3+5) + \dots + n$$

y

int $ns = 0$; // outer sum

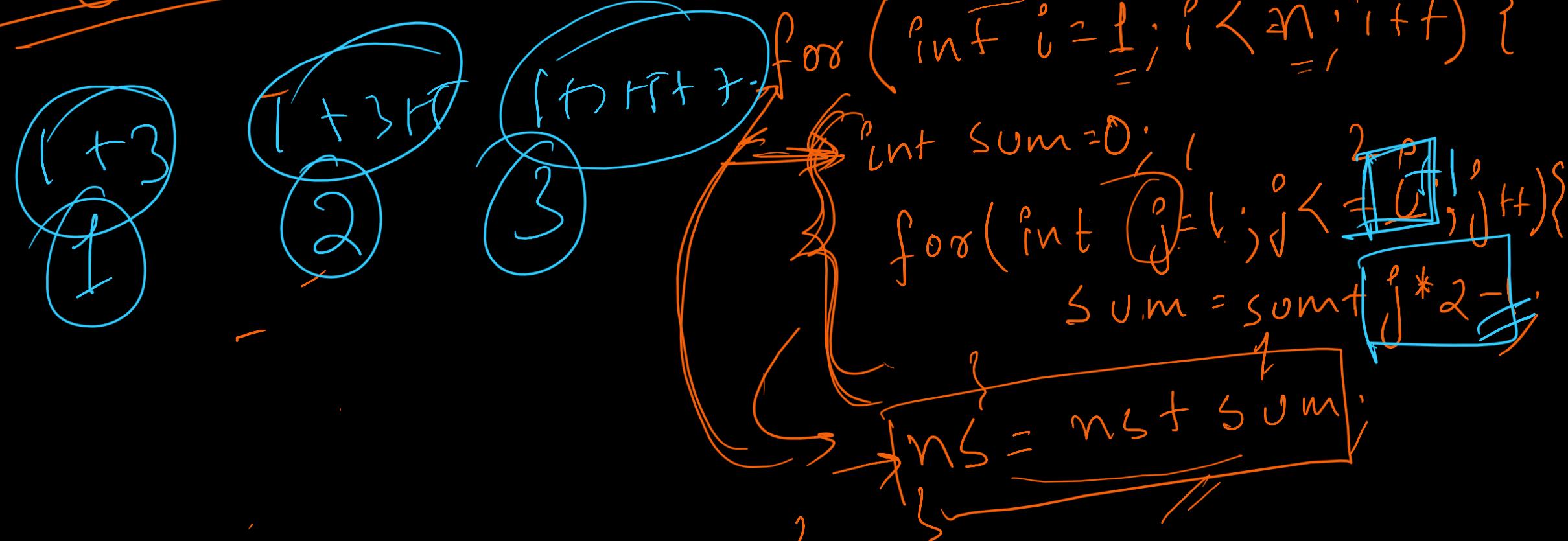
for ($i=1$; $i < n$; $i++$) {

 int $sum = 0$;

 for (int $j=1$; $j < i+1$; $j++$) {

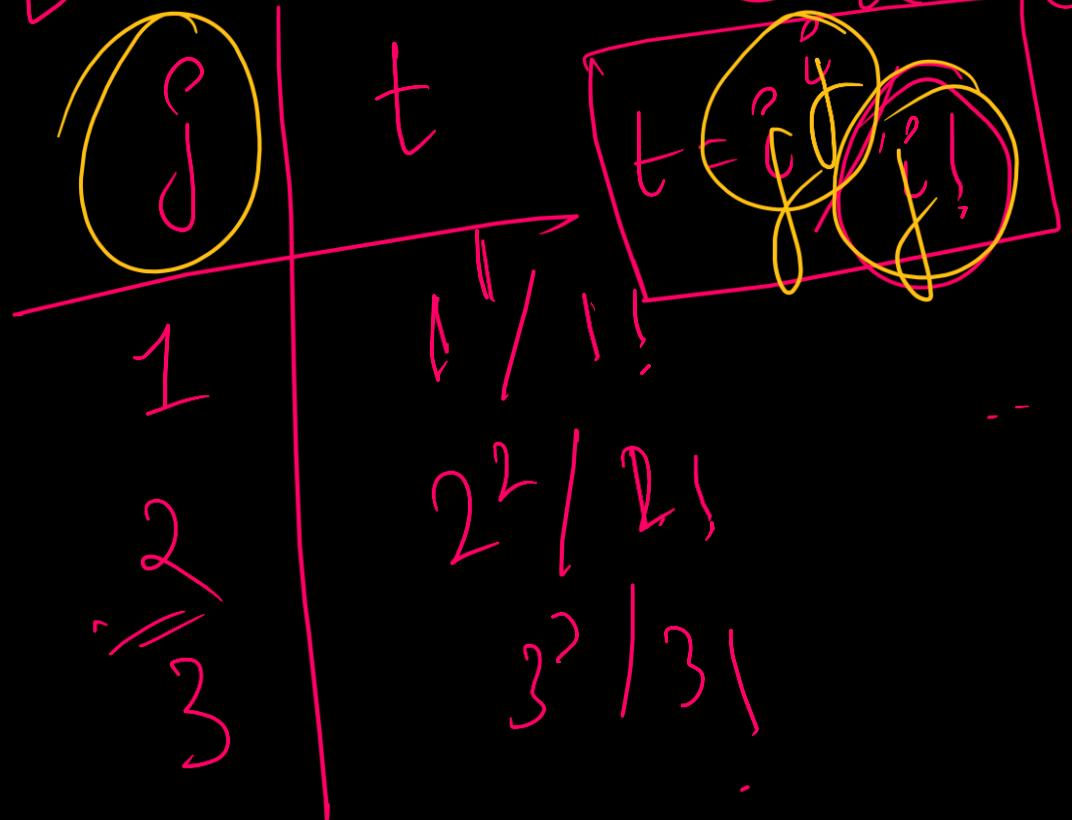
$sum = sum + j * 2 - j$;

$ns = ns + sum$;



$$\frac{1}{1!} + \frac{2^2}{2!} + \frac{3^3}{3!} + \dots + \frac{n^n}{n!}$$

$$t = \sum_{i=1}^n i!$$



int ns = 0;

```
foo(int j){ j < n; j++) {
```

```
{ int prod = 1;
for (int i = 1; i <= j; i++) {
```

```
prod *= i; }
```

$ns = ns + \text{pow}(j, j) / prod$

$$\text{Category 4: } \underline{\underline{2}} * \underline{\underline{P}} = \underline{\underline{10}} + \underline{\underline{n}}$$

Dry Run

$$3^*10+3 + 33^*10+3 + 333 + 3\overline{3}3 =$$

```

int sum = 0;
for (int i = 1; i <= n; i++) {
    sum = sum + i;
    cout << sum;
}

```

The diagram illustrates the execution of the code. A variable *i* is tracked in a blue circle, starting at 1 and increasing by 1 each iteration. The value of *i* is also shown above the assignment statement in each row. The variable *sum* is tracked in a red circle, starting at 0 and increasing by the current value of *i* each iteration. The value of *sum* is also shown above the addition operation in each row.

Iteration	<i>i</i>	<i>sum</i>
1	1	0 + 1 = 1
2	2	1 + 2 = 3
3	3	3 + 3 = 6
4	4	6 + 4 = 10
5	5	10 + 5 = 15
6	6	15 + 6 = 21
7	7	21 + 7 = 28
8	8	28 + 8 = 36
9	9	36 + 9 = 45
10	10	45 + 10 = 55

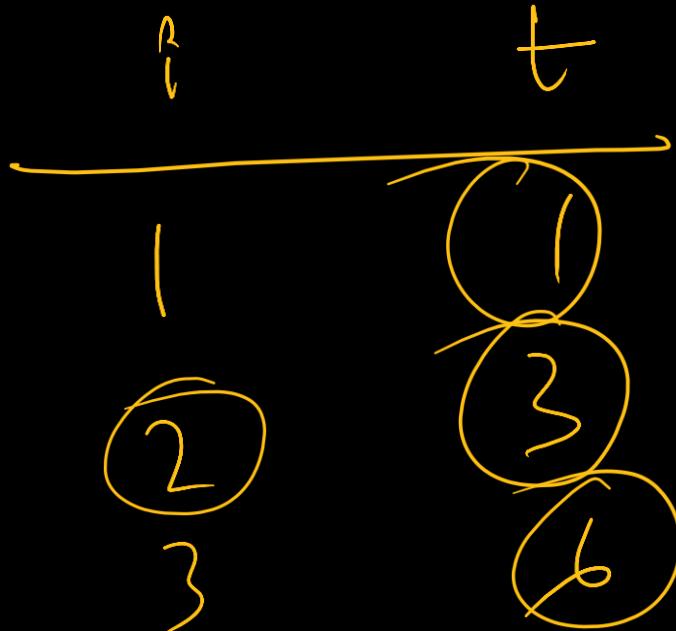
$$\text{Sum} = \underline{\underline{0}} + 3 + 3$$

$$\begin{array}{r} b = 3 \\ \hline c = 1 \end{array}$$

Category 5 : Special Series

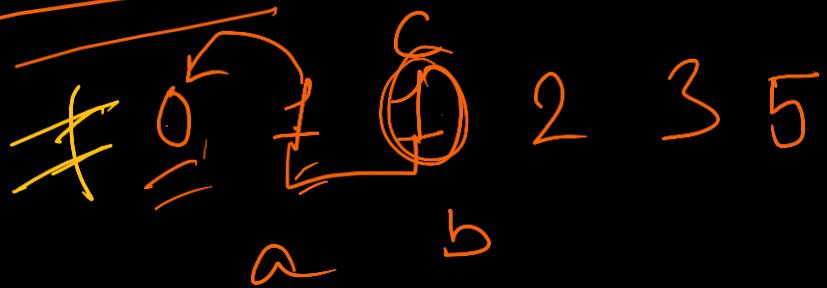
(Triangular Series)

$$1 \quad 3 \quad 6 \quad 10 \quad 15$$
$$\left\{ i^2 * \left(\frac{(i+1)}{2} \right) \right\} / 2 = t$$



```
{ os( int i = 1; i <= n; i++ ) {  
    printf( "%*c(%d + 1) / 2);  
}
```

Fibonacci Series



$$\begin{aligned} a &= 0 \\ b &= 1 \\ c &= 1 \end{aligned}$$

for (int i = 1; i <= n; i++) {
 c = a + b;
 printf("%d", c);
 a = b;

shout
 b = c;



T
1 2 3

c = a + b;
printf(c);
a = b;
b = c

$\Rightarrow -16, -4, 3, -7, +10, -13, 16, -19$
 $\equiv \leftarrow$ $\text{int } p = 1$
 for $\left(\begin{array}{l} \text{int } p \rightarrow (\vdash ? \leq n; ? + f) \\ \text{if } (p \leq 2 = 0) \text{ printf}(p) \\ \text{else } \text{printf}(-p) \end{array} \right)$ {
 }
 $p = p + 3$

Q

10 + 5
~~45~~ 45
 60
~~45~~ 45
 15
~~45~~ 45
 110
~~45~~ 45

int b = 10;
 int a = 5;

for (int i = 1; i <= n; i++) {

if (i % 2 == 0) {

printf("%d",

b = b + 50;

else {

printf("%d",

a = a + 10);

}

6 9 14 21 30

~~3+2 5+2~~ ~~7+2 9~~

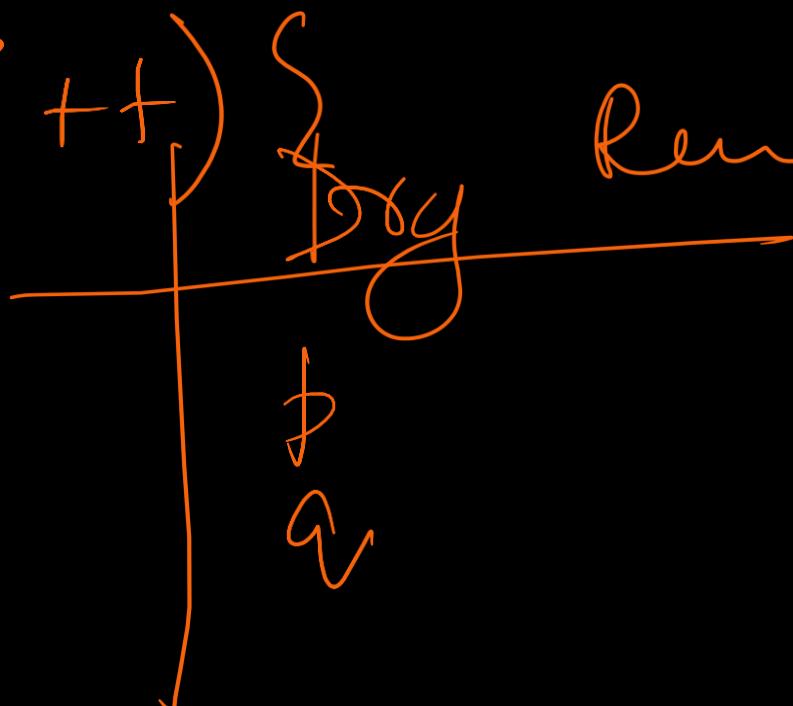
int p=6; int q=3;

for (int i=1; i<=n; i++) {

cout < p;

p = p + @;

q = q + 2;



2

15

41 80
2nd

n_i

p

p = 2

2 * 0 + 13 = 15
2nd 2

15 * 2 + 13 = 41
3rd 2

4 * 3 + 13 = 0
3rd 3

80 * 4 + 0 = ==

b * p * n_i + 13 int p = 2
 $b * p = (int i(1), i < n)$
{+ f}

print(p);

p = b * p + 13
=

}

