Tight Bounds on 3-Neighbor Bootstrap Percolation

by

Abel Emanuel Romer B.A.Sc., Quest University Canada, 2017

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We acknowledge with respect the Lekwungen peoples on whose traditional territory the university stands, and the Songhees, Esquimalt, and WSÁNEĆ peoples whose historical relationships with the land continue to this day.

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ABSTRACT

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DEDICATION

Chapter 1 Introduction

Consider the lattice depicted in the leftmost diagram of Figure 1.1. We refer to the elements of this lattice as *cells*. Suppose we have the capacity to infect some cells (colored red) with a disease, and that this disease will, over a period of time, propagate through uninfected cells of the lattice. Let uninfected cells become infected if they are exposed to at least two infected neighboring cells in the vertical and/or horizontal directions. We say that the initial infection is *lethal* if the entire lattice ultimately becomes infected. Here is the puzzle:

Question. What is the fewest number of infected cells necessary to spawn a lethal infection?

Before we present the solution, let us take a moment to examine some properties of infectious sets and attempt to characterize what attributes might correspond to lethality. It should not take too long to observe that if an initial infection is in some way "spread too thin," it will be unable to jump between infected areas, leading to gaps in infection, which we refer to as *immune regions*. The perimeter of the lattice is particularly susceptible to this, as vertices there have fewer neighbors from whom they might be exposed. Heuristically, then, a lethal set must have the ability to effectively span the entire lattice, and must be particularly virulent along the perimeter.

With this criteria in mind, we are able to make some educated guesses regarding the specific structure of sets that are likely to be lethal. In particular, we would like to consider the two starting infections illustrated in Figure 1.2. Notice that while Figure



Figure 1.1: An arbitrary set of initially infected cells in the 10×10 lattice, and the stages of infection.



Figure 1.2: Two lethal sets and their resulting infections after one time-step.

1.2 (b) has far fewer perimeter infections, both (a) and (b) manage to form continuous bands of infected cells that appear to span the entire lattice after one step. Indeed, this holds with our notion of immune regions (or lack thereof), and we see that both infections will continue to propagate outwards from these bands until all cells become infected.

It is clear from Figure 1.2 that we may obtain lethal sets on the $n \times n$ lattice of size n by simply infecting the diagonal. What is less obvious is whether it is possible to improve upon this result. Perhaps the most natural first attempt at this is to remove an infection from one of the cells along the diagonal. However, this seems to form an immune region around the removed cell. After some experimentation, one begins to believe it impossible to simultaneously satisfy the heuristic that a starting infection must span the lattice, while also using fewer than n initial infections. The question therefore becomes: how do we prove it?

Consider the cumulative perimeter of infected regions. For a given infectious set A, let P(A) be the total perimeter of the infected regions of A. Let A_0 be an initial infection, and observe that $P(A_0) \leq 4|A_0|$. (This bound is only tight if no two infected cells are adjacent. Otherwise, the edge between such cells lies within the infected region, and cannot contribute to the infection's perimeter.) Observe that for any uninfected cell to become infected, it must abut at least two infected cells. Upon infection, the edges adjacent to these cells no longer lie on the infection's perimeter; additionally, the remaining edges of this newly infected cell contribute at most 2 to this perimeter. All told, after infection, $P(A_1) \leq P(A_0)$.

If we suppose that A_0 is a lethal set, then at some point in time, the entire grid will become infected. This infection will have a perimeter 4n. Since this perimeter did not increase, A_0 must have originally had a perimeter of at least 4n. Since each cell in A_0 can contribute at most 4 to this perimeter, it must be the case that $|A_0| \ge n$. Our diagonal construction shows that $|A_0| \le n$, and so we are able to conclude that n is best possible.

This proof is an instance of the famous *perimeter argument*, which has belonged to bootstrap percolation folklore since at least the work of Pete [?]. In the following

section, we present generalizations of this argument to higher dimensional rectangular grids.

1.1 Bootstrap Percolation

The study of such cellular infection spread in grids (and more generally in graphs) is known in the literature as bootstrap percolation, and was introduced in the 1970s by Chalupa et al. [2] as a simplified model for the behavior of ferromagnetic fields. In their original 1979 paper, the authors research the stable structure of probabilistically selected initial infections. While this differs from the problem posed in Question 1, the rules for the spread of infection and its broad behavior remain the same. It is worth noting that a large portion of contemporary research on bootstrap problems is focused on questions of probabilistic nature; while these problems are certainly interesting and of merit, they do not fall within the scope of this thesis. Rather, we shall focus on those problems where we have specific control over the structure of the initial infections; in particular, we aim to determine the smallest lethal set on the Cartesian product of paths and cycles.

We define the problem in concrete terms. Let G be a graph and $A_0 \subseteq V(G)$ be a set of initially infected vertices. Iteratively, infect those vertices of G with at least r infected neighbors. For all t > 0, let A_t be the set of infected vertices at time step t. We then have

$$A_t = A_{t-1} \cup \{ v \in V(G) : |N_G(v) \cap A_{t-1}| \ge r \},$$

where $N_G(v)$ is the set of vertices adjacent to v in G. We define the *closure* of A_0 under r-neighbor bootstrap percolation to be $[A_0] = \bigcup_{t=0}^{\infty} A_t$. We say that A_0 percolates or is *lethal* if $[A_0] = V(G)$. We define the smallest percolating set on a graph G under r-neighbor bootstrap percolation by the quantity m(G, r). We note that under these rules, it is not possible for vertices to become uninfected.

While it is possible to study bootstrap percolation on any graph G, most contemporary research focuses on multidimensional grids [?]. We therefore introduce the following notation. For all $n \in \mathbb{N}$, let $[n] = \{1, 2, ..., n\}$. Let the grid graph with vertex set $\prod_i^d [a_i]$ be denoted by $\prod_i^d [a_i]$. Note that $\prod_i^d [a_i] = P_{a_1} \square \cdots \square P_{a_d}$. Furthermore, define:

$$m(a_1,\ldots,a_d,r)=m(\prod_i^d [a_i],r).$$

There are a number of natural generalizations of the problem posed in Question 1. In this thesis, we discuss those obtained by varying the structure of G and the value of r. Below, we outline some of the existing results for graphs that are the Cartesian product of paths and cycles, and $r \in \{0, 1, 2, 3\}$. These results are summarized in Table 1.1.

	Grids								
r					_				
	$[a_1]$	$[a_1] \times [a_2]$	$[n]^2$	$[a_1] \times [a_2] \times [a_3]$	$[n]^3$	• • • •	$\prod_{i=1}^{d} [a_i]$	$[n]^d$	$[2]^d$
r = 0	0	0	0	0	0		0	0	0
r = 1	1	1	1	1	1		1	1	1
r = 2	$\left\lceil \frac{a_1-1}{2} \right\rceil + 1$	$\left\lceil \frac{a_1 + a_2 - 2}{2} \right\rceil + 1$		$\left\lceil \frac{a_1 + a_2 + a_3 - 3}{2} \right\rceil + 1$	$\left\lceil \frac{3(n-1)}{2} \right\rceil + 1$		$\left\lceil \frac{\sum_{i=1}^d (a_i-1)}{2} \right\rceil + 1$	$\left\lceil \frac{d(n-1)}{2} \right\rceil + 1$	$\lceil \frac{d}{2} \rceil + 1$
r = 3	???	???	$\left\lceil \frac{n^2+2n+4}{3} \right\rceil^*$	S.A. bound	n^2		???	???	$\lceil \frac{d(d+3)}{6} \rceil$
÷						4.			
r = d	???	???	???	???	???		S.A. bound	n^{d-1}	???

Table 1.1: A summary of known bootstrap percolation results for grids and the torus, $r \in \{0, 1, 2, 3, d\}$.

1.1.1 Results on grids and tori

In this section, we highlight major bootstrap percolation results on grids and tori. Some of the following bounds are not tight and require supplemental constructions, which are often difficult to obtain. We further sub-divide this discussion into results on the grid (of which there are many), and results on the torus (of which there are few).

Grids

From the puzzle posed in Question 1, we readily obtain variant problems by altering three parameters: the size and shape of the grid G, the grid's dimension d, and the threshold number of neighbors r. We examine each of these problems in turn.

In the prior discussion of the perimeter argument, we showed that for square grids $[n]^2$, $m([n]^2, 2) \ge n$, and verified this to be tight with a diagonal construction. The following result (attributed to Pete) generalizes this result to all rectangular grids $[a_1] \times [a_2]$. A proof is included for completeness.

Theorem 1.1. For $a_1, a_2 \ge 1$,

$$m(a_1, a_2, 2) = \left\lceil \frac{a_1 + a_2 - 2}{2} \right\rceil + 1.$$

Proof. We obtain a lower bound on $m(a_1, a_2, 2)$ by applying the perimeter argument. Note that the perimeter of the $a_1 \times a_2$ grid is $2(a_1+a_2)$, and so the $m(a_1, a_2, 2) \ge \left\lceil \frac{a_1+a_2}{2} \right\rceil$. (We take the ceiling because the size of infected sets must be integral. See Figure 1.4.) For the upper bound, we proceed by induction on $a_1 + a_2$. For $a_1 + a_2 \le 4$, the lethal sets in Figure 1.3 match the lower bounds given by the perimeter argument (1, 2, 2, and 2, respectively). For $a_1 + a_2 > 4$, suppose without loss of generality that $a_1 \le a_2$, and so $a_2 \ge 3$. By hypothesis, $[a_1] \times [a_2 - 2]$ admits a lethal set A_0 at the perimeter bound. We show that A_0 , plus the addition of any infection in the final column of $[a_1] \times [a_2]$, is lethal and matches the perimeter bound.

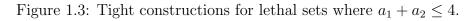




Figure 1.4: Tight constructions for lethal sets on the $[a] \times [b]$ grid.

Observe that A_0 infects all vertices of $[a_1] \times [a_2]$, apart from the final two columns. The additional vertex in the final column is then sufficient to infect all remaining healthy vertices. Finally, by incrementing a_2 by two, the perimeter bound is incremented by exactly one. This completes the proof.

Let us take a moment to examine the issue of integrality in the perimeter bound. Non-integrality occurs either when adjacent vertices are infected in the same generation, or when a vertex is infected by more than r neighbors. Note that in both cases, this decreases the perimeter of infection. One way to think about this is to consider each vertex as having "infectious potential": vertices $v \in A_0$ can infect up to d(v) healthy vertices, whereas vertices $v \in A_i$ for i > 0 can infect at most d(v) - r. An integral perimeter bound mandates that each vertex realize its potential, whereas a non-integral bound leaves a small margin for error. Figure 1.4a illustrates the integral case, where each cell is infected by exactly two neighboring cells; this condition ensures that $P(A_i) = P([A_0])$ for all i. Conversely, in Figure 1.4b, the cell demarcated with an "X" experiences infection on three sides, thereby reducing its infectious potential. The existence of such a cell is guaranteed by the fact the perimeter bound in this case is non-integral.

We can further generalize for the case of r=2. In 2006, Balogh and Bollobas [1] proved the following general form of Theorem 1.1 for all d-dimensional hypercubes $(a_1, \ldots, a_d), a_i \geq 1$:

Theorem 1.2 (Balogh). For $d \ge 1$ and $a_1, \ldots, a_d \ge 1$,

$$m(a_1, \dots, a_d, 2) = \left\lceil \frac{\sum_{i=1}^d (a_i - 1)}{2} \right\rceil + 1.$$

Theorem 1.2 completes the picture for infections with a threshold of two on grids. The next question is whether similar results exist for larger r. Unfortunately, while generalizing to d-dimensional grids yields nice results for r = 2, attempts to obtain

a holistic understanding of $m(a_1, \ldots, a_d, r)$ for arbitrary r have been largely fruitless. Even the case of r = 3 remains stubbornly inaccessible for nearly all large d. However, certain breakthroughs have been made in the following circumstances: d = 2, d = 3, and $G = [2]^d$.

We first consider 3-neighbor percolation on two-dimensional square grids. In 2021, Benevides et al. proved that

$$m(n,n,3) = \left\lceil \frac{n^2 + 2n + 4}{3} \right\rceil$$

for even n, and

$$\left\lceil \frac{n^2 + 2n}{3} \right\rceil \le m(n, n, 3) \le \left\lceil \frac{n^2 + 2n}{3} \right\rceil + 1$$

for odd n. Additionally, they showed that these bounds are tight for odd n: if $n = 5 \pmod{6}$, or $n = 2^k - 1$ for some $k \in \mathbb{N}$, then $m(n, n, 3) = \lceil \frac{n^2 + 2n}{3} \rceil$; and if $n \in \{9, 13\}$, then $m(n, n, 3) = \lceil \frac{n^2 + 2n}{3} \rceil + 1$. Constructions that achieve this bound are illustrated in Chapter 7. We add to this picture with the following theorem, proven in Chapter 6, and corollary:

Theorem 1.3. Suppose that $a, b \ge 1$ such that

$$m(a, b, 3) = \frac{2ab + a + b}{3}.$$

Then there exists $k \ge 1$ such that $a = b = 2^k - 1$.

Corollary 1.4. For all $n \geq 1$,

$$m(n, n, 3) = \begin{cases} \left\lceil \frac{n^2 + 2n + 4}{3} \right\rceil & n \equiv 0 \pmod{2} \\ \frac{n^2 + 2n}{3} & n = 2^k - 1, \ k \in \mathbb{N} \\ \frac{n^2 + 2n + 1}{3} & n \equiv 5 \pmod{6} \\ \frac{n^2 + 2n + 3}{3} & otherwise. \end{cases}$$

Proof. The first three cases follow from Theorem 1 of [?] and the observation that if $n \equiv 5 \pmod{6}$, then $\lceil \frac{n^2+2n}{3} \rceil = \frac{n^2+2n+1}{3}$. In the final case, n is congruent to either 1 or 3 modulo 6. This implies that n^2+2n is divisible by three. From Theorem 1 of [?], we have that $m(n,n,3) \leq \frac{n^2+2n+3}{3}$. Furthermore, since n is not of the form 2^k-1 , it follows from Theorem 1.3 that $m(n,n,3) > \frac{n^2+2n}{3}$. Therefore, $m(n,n,3) = \frac{n^2+2n+3}{3}$.

This result resolves the question of the minimum lethal set for two dimensional square grids. For the more general case of rectangular grids, the problem remains unsolved. However, we are able to achieve an upper bound of $m(a,b,3) \leq \lceil \frac{ab+a+b+6}{3} \rceil$

and a lower bound (described below) of $m(a, b, 3) \ge \lceil \frac{ab+a+b}{3} \rceil$, for all a, b > 1. Further discussion of these results can be found in Chapter 6.

One significant and well-known result for 3-neighbor percolation on two- and three-dimensional grids is the following lower bound, taken as a three-dimensional analogue of the perimeter bound. This result is referenced frequently throughout this document, and referred to interchangeably as the *surface area* or SA bound. We prove the statement in full generality, while noting that we only make use of the case where d=3. We also note that, like the perimeter bound, the following proof belongs to bootstrap percolation folklore, and appears to have been first published in 1997 by Balogh and Pete [?].

Theorem 1.5. For any $d \ge 1$ and $a_1, a_2, ..., a_d \ge 1$,

$$m(a_1, a_2, \dots, a_d, d) \ge \frac{\sum_{j=1}^d \prod_{i \ne j} a_i}{d}.$$

Proof. We apply the same invariant strategy presented in the perimeter argument. For simplicity, consider $\prod_{i=1}^d [a_i]$ to be embedded within the larger graph $\prod_{i=1}^d \{0, \ldots, a_i + 1\}$. Note that in $\prod_{i=1}^d \{0, \ldots, a_i + 1\}$, each vertex $v \in \prod_{i=1}^d [a_i]$ has degree 2d. Let A_0 be a lethal set in $\prod_{i=1}^d [a_i]$ under the d-neighbor bootstrap process. Let A_t be the set of infected vertices in $\prod_{i=1}^d [a_i]$ at generation t. Denote by m_t the number of edges between vertices $u \in A_t$ and $v \in \prod_{i=1}^d [a_i] \setminus A_t$. We show that $m_{t-1} \geq m_t$ for all t > 0.

By definition, each vertex $v \in A_t \setminus A_{t-1}$ has at least d neighbors in A_{t-1} . Therefore, since d(v) = 2d, v has no more than d neighbors outside of A_t . This implies that the number of edges from $A_{t-1} \cup \{v\}$ to $\prod_{i=1}^d [a_i] \setminus A_{t-1} \cup \{v\}$ cannot exceed m_{t-1} . Furthermore, this holds for every vertex $v \in A_{t-1}$, and so $m_{t-1} \ge m_t$.

Since A_0 is lethal, we have that

$$2d|A_0| \ge m_0 \ge m_1 \ge \dots \ge 2\sum_{j=1}^d \prod_{i \ne j} a_i,$$

where the final expression gives the total number of edges between the fully infected grid and the surrounding larger grid. Dividing through by 2d gives the result.

We note that the prior argument is precisely the same as the so-called perimeter argument outlined on Page 2. Here, the quantity m_t is a d-dimensional analogue of the perimeter of infection $P(A_t)$ at time-step t, and the lower bound

$$2\sum_{j=1}^{d} \prod_{i \neq j} a_i$$

is the d-dimensional "perimeter" of the grid. Again, observe that equality can only be

obtained when no vertices of A_0 are adjacent, and all vertices $v \in A_{t>0}$ are infected by exactly d neighbors. Any imprecision causes a reduction in "perimeter" of two units, corresponding to a 1/d increase in the bound.

The primary aim of this thesis is to prove that the surface area bound is tight for sufficiently large grids when r=3. This process employs a number of general constructions (discussed in Chapter 7), as well as a recursive strategy (Chapter 3). In Chapter 5, we prove the following result:

Theorem 1.6. For all $a_1, a_2, a_3 \ge 11$,

$$m(a_1, a_2, a_3, 3) = \left[\frac{a_1a_2 + a_2a_3 + a_1a_3}{3}\right].$$

Unfortunately, the complete resolution of the r=3 case on grids remains elusive. Tight constructions exist for cubes $[n]^3$ and hypercubes $[2]^d$, but in general bounds are difficult to obtain. Worse, for r>3, the only additional result beyond the surface area bound addresses the very specific case of r-dimensional cubes. Open problems abound.

Tori

In addition to varying the parameters r and d, we might also change the very structure of G. It is natural to shift from grids (the Cartesian product of paths) to tori (the Cartesian product of cycles). In fact, it could be argued that bootstrap percolation on the torus is more natural than the grid, since tori are regular and grids are not. This problem has been studied by Benevides et al. In 2021, they obtained the following lower bound for the Cartesian product of two cycles [?]. Their proof is included here for completeness.

Theorem 1.7. For $a, b \geq 1$,

$$m(C_a \square C_b, 3) \ge \left\lceil \frac{ab+1}{3} \right\rceil.$$

Proof. Let $G = C_a \square C_b$, and let I be a lethal set on G. Let $H = V(G) \setminus I$, and note that |H| = ab - |I|. Let m_H be the number of edges in the subgraph of G induced by H, and m_{IH} be the number of edges between vertices in I and vertices in H. Note that m_{IH} is similar to the notion of perimeter on a grid.

Observe that G[H] must be cycle-free: cycles in G[H] constitute immune regions, and contradict the lethality of I. Therefore, G[H] is a forest, and so $m_H = |H| - c$, where c is the number of components in G[H]. Additionally, note that $m_{IH} \leq 4|I|$, since G is 4-regular. Finally, observe that the total degree of G[H] is $2m_H = 4|H| - m_{IH}$.

Chaining together these inequalities, we obtain:

$$4|I| \ge m_{IH} = 4|H| - 2m_H$$

= 4|H| - 2(|H| - c) = 2|H| + 2c
= 2(ab - |I|) + 2c

Combining like terms and simplifying, we have

$$|I| \ge \frac{ab+c}{3} \ge \frac{ab+1}{3}.$$

Observe that the conditions $c \geq 1$ and $m_{IH} \leq 4|I|$ prevent us from obtaining strict equality. Specifically, if I is lethal, G[H] has one component, and no vertices in I are adjacent, then |I| is minimized. Note that these conditions are quite similar to those on grids; the specific difference is that equality in the bound on grids mandates that no vertex be infected by more than r neighbors, whereas equality on three-dimensional tori requires this inefficiency be centered on one particular vertex.

1.2 Other Problems

In this section, we provide a cursory overview of some other related areas of study in bootstrap percolation. We highlight problems regarding the number of iterations necessary to infect all vertices on a graph.

1.3 Structure of this Thesis

As stated by Theorem 1.6, the primary goal of this thesis is to prove a tight bound for 3-neighbor bootstrap percolation on three-dimensional grids of sufficiently large size. This task requires the use of two major lemmas, as well as both original and previously published ideas and constructions. In an effort to present this material in a coherent manner, the paper is structured as follows.

Chapters 2 and 3 are dedicated to building a conceptual and intuitive framework upon which to prove Theorem 1.6. In Chapter 2, we categorize and present percolating sets on both divisible and non-divisible grids, and discuss the differences between these cases. In Chapter 3, we prove two lemmas allowing us to develop large families of lethal sets that match the surface area bound.

Chapters 4 and 5 leverage the results of Chapter 3 to prove Theorem 1.6. In Chapter 4, we show that all divisible grids (a_1, a_2, a_3) , $a_1 \le a_2 \le a_3$ admit lethal sets matching the SA bound, and in Chapter 5, we extend this result to all grids of size 11. Chapter

6 further builds on this, highlighting some new results for 3-neighbor percolation on grids of the form $(a_1, a_2, 1)$.

Chapter 7 proves the constructions presented in Chapter 2.

Finally, Chapters 8, 9 and 10 wrap up the discussion with a summary of the programmatic techniques used to discover lethal sets, helpful software resources for future research, how the results of this thesis may be helpful in pursuit of results on the torus, and recommendations for future research in similar and related problems.

Chapter 2 Conceptual Tools

2.1 Conceptual Tools for Understanding Bootstrap Percolation

As outlined above, it appears that lethal sets in grid graphs adhere to certain fixed rules. We examine these rules here, and explain why they are necessary and helpful for understanding the problem of bootstrap percolation.

2.1.1 Integrality of bounds

Not all grids have integral surface area bounds. We refer to such grids as non-divisibility cases. The non-divisibility cases for three-dimensional grids where r=3 are illustrated in red in Table 2.1.

Percolation in non-divisible grids is inefficient in one of two ways: either each vertex in A_0 contributes, on average, less than 2d units to the final surface area, or at some time-step i, there exists a vertex $v \in V \setminus A_i$ such that $N_{A_i}(v) > r$. In the first case, the inefficiency is a consequence of two adjacent vertices in A_0 ; heuristically, we may view this as poor distribution of infection across the grid. In the second case, the inefficiency results from excessive infection in a particular region; a vertex is over-exposed.

It is insightful to consider this behavior in the context of Theorem 1.7. Instead of thinking of m_{IH} at the number of edges between I and H, we shall consider it to represent the perimeter, or surface area, of the initial infection I.

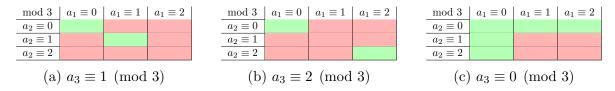


Table 2.1: Integrality of grids by congruence class. Green indicates integral surface area bound.

2.1.2 A different characterization

We shall find it useful to think about the problem of bootstrap percolation in terms of the graph induced by the set of uninfected vertices $V \setminus A_i$ at some time-step t_i .

2.1.3 Attributes of $G[V \setminus A]$

No cycles, no paths between opposite faces. Reference to Benevides.

2.1.4 The lower bound

Let us reconsider the our proof of the lower bound in this new framework.

2.1.5 Surface area, optimality, trees, and components

The surface area bound is not always integral. In cases where it is not, this suboptimality can be traced to a particular time-step where an uninfected cell is infected by more than r neighbors. The proof of the lower bound on the torus considers the number of components in $T[V \setminus A]$, as each such component necessarily collapses in on a vertex where this occurs. In the grid, this circumstance can be avoided by permitting components to abut the perimeter. What is the relationship here?

2.1.6 The bipartition (and when to violate it)

Discuss the heuristic of placing infections on one side of the bipartition and when this technique fails.

2.1.7 A potentially interesting proof framework

A discussion on the broad structure of the proof of the lower bound on the torus as presented in Benevides. Note that this proof relies on a sequence of inequalities, and that two of these inequalities characterize precise conditions that are necessary for a set to be lethal. Discuss how the proof might change as a function of the dimension of the grid.

Chapter 3

Two Recursive Techniques

In this chapter, we shall present two techniques for constructing large *perfect* grids from smaller perfect grids. The first is a recursive technique that builds d-dimensional sets from (d-1)-dimensional components, while the second builds d-dimensional sets from smaller blocks of the same dimension.

3.1 The Three Walls Lemma

Lemma 3.1. Let S be an infected set on $G = \prod_{i=1}^d [a_i]$. Let $\overline{S} = V(G) \setminus S$, and let $H = G[\overline{S}]$ be the subgraph of G induced by \overline{S} . For $1 \le k \le a_j$, let $F_{j,k} = \prod_{i=1}^{j-1} [a_i] \times \{k\} \times \prod_{i=j+1}^d [a_i]$ be the kth face of G in the jth dimension. If H does not contain a path between $F_{j,1}$ and F_{j,a_j} , for all $1 \le j \le d$, then S is lethal on G under d-neighbor percolation.

Proof. We proceed by induction on $|V(H)| = \prod_{i=1}^d a_i - |S|$. If |V(H)| = 0, then all vertices of G are infected and we are done. Suppose |V(H)| > 0, and consider a connected component Y of H. By hypothesis, for all $j \in [d]$, either $V(Y) \cap F_{j,1} = \emptyset$ or $V(Y) \cap F_{j,a_j} = \emptyset$ (or both). Suppose, without loss of generality, that $V(Y) \cap F_{j,a_j} = \emptyset$. For each $j \in [d]$, let x_j be the maximum value such that $V(Y) \cap F_{j,x_j}$ is non-empty. Note that such an x_j must exist since |V(H)| > 0.

Consider the vertex $\vec{x} = (x_1, \dots, x_d) \in V(Y)$, and observe that

$$\{\bigcup_{j\in[d]}F_{j,x_j+1}\}\cap V(Y)=\emptyset.$$

In particular, note that $(x_1 + 1, x_2, ..., x_d), ..., (x_1, ..., x_d + 1) \in N_S(\vec{x})$. Therefore, \vec{x} becomes infected. Furthermore, since $|V(H) \setminus \{\vec{x}\}| < |V(H)|$, the resulting graph percolates by induction. This completes the proof.

Corollary 3.2. Let G be the grid graph $\prod_{i=1}^{d} [a_i]$. For each $j \in [d]$ and some $1 \leq k \leq d$

 a_i , let

$$M = \bigcup_{j,k} F_{j,k}$$

be a subset of the vertices of G formed by the union of mutually orthogonal faces. If a set S is lethal on M, then it is lethal on G.

Proof. Since S is lethal on M, there exists a time t where $M \subseteq A_t$. Therefore, for all $j \in [d]$, the graph $G[\overline{A_t}]$ cannot contain a path between $F_{j,1}$ and F_{j,a_j} . By Lemma 7.2, S is lethal on G.

Corollary 3.2 provides a general characterization of lethal sets on d-dimensional grids in terms of their (d-1)-dimensional faces, provided these faces are mutually orthogonal. Here, we return to the notion first introduced in Chapter 1 of the capacity of a lethal set to "span" a grid. In particular, we see that the set in Figure 1.2a is comprised of lethal sets on the two one-dimensional orthogonal faces $F_{1,1}$ and $F_{2,1}$ of $[10]^2$. In this regard, the problem of obtaining perfect d-neighbor lethal sets on d-dimensional grids is reduced to the problem of determining a "good" union M of mutually orthogonal (d-1)-dimensional faces. In Chapter {CONSTRUCTIONS}, we apply this idea to obtain an infinite family of three-dimensional grids from three orthogonal two-dimensional faces. However, we caution that the challenge of determining a "good" union M is non-trivial in general.

The following corollary (taken as a particular instance of Corollary 3.2) will be useful in our discussion of lethal sets on three-dimensional grids (a_1, a_2, a_3) .

Corollary 3.3. Let G be the grid graph (a_1, a_2, a_3) . If a set S is lethal on $F_{1,1} \cup F_{2,1} \cup F_{3,1}$, then S is lethal on G.

Proof. By hypothesis, S is lethal on $F_{1,1} \cup F_{2,1} \cup F_{3,1}$. Therefore, there exists some time t for which $F_{1,1} \cup F_{2,1} \cup F_{3,1} \subseteq A_t$, and so $G[\overline{A_t}]$ satisfies the conditions of Lemma 7.2. We conclude that S is lethal on G.

3.2 Manifolds

We refer to unions M of perpendicular faces as manifolds. As suggested in the previous section, there are a wide variety of possible manifold structures for a given grid. We highlight two particular structures in Figure 3.1. The following two lemmas show that it is possible to obtain perfect lethal sets on these structures from existing perfect sets on grids of thickness one.

Lemma 3.4. Let G be the grid graph (a_1, a_2, a_3) and

$$M = F_{1,1} \triangle F_{2,1} \triangle F_{3,1}.$$

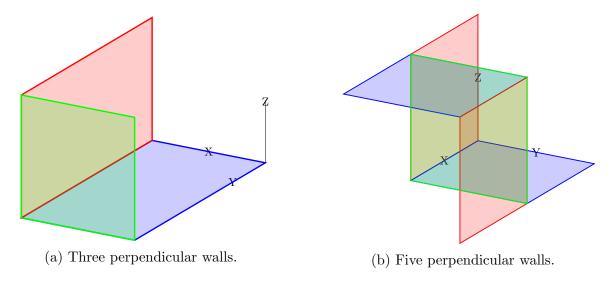


Figure 3.1: Two types of manifold used in our constructions.

Note that G[M] is isomorphic to the disjoint union of grids $(a_1 - 1, a_2 - 1, 1), (a_2 - 1, a_3 - 1, 1), (a_3 - 1, a_1 - 1, 1)$. Let $S_{1,2}, S_{2,3}, S_{3,1}$ be perfect lethal sets under 3-neighbor percolation on each of these grids, respectively. Let S be the set of vertices $S_{1,2} \cup S_{2,3} \cup S_{3,1}$ under isomorphism from $(a_1 - 1, a_2 - 1, 1) \cup (a_2 - 1, a_3 - 1, 1) \cup (a_3 - 1, a_1 - 1, 1)$ to G[H]. Then $S \cup (1, 1, 1)$ is lethal and perfect on G.

Proof. This proof is just algebra. Assume that each of the three faces has a lethal set at the surface area bound, add them all up, and observe that the resulting sum is precisely the surface area bound on the grid, minus one point. \Box

WHOOPS! THE FOLLOWING PROOF DOESN'T QUITE WORK. SHOULD ADD 4/3 OF A POINT, NOT 2 POINTS.

Lemma 3.5. Let G be the grid graph (a_1, a_2, a_3) and let M be the symmetric difference of vertex sets on the red, green and blue faces illustrated in Figure 3.1b. Let S be a perfect lethal set under 3-neighbor percolation on G[M]. Then $S \cup \{$ opposite corners of the green face $\}$ is perfect and lethal on G.

Proof. A bunch of basic algebra shows that the bound match up. The two points on the green face are sufficient to guarantee the lethality of S on the entire manifold. \square

QUESTION: Does this extend to d-neighbor percolation on d-dimensional grids?

QUESTION: The above lemma applies to the circumstance where M is half of the surface area of G. What if M is a different set of mutually orthogonal faces? How do we write this is general, where none of the faces lie on the surface of the grid?

QUESTION: What is the relationship between that case (mentioned above) and the block recursion?

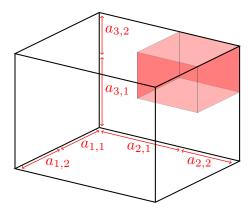


Figure 3.2: A recursively constructed $[b_1] \times [b_2] \times [b_3]$ grid, for n = 2, d = 3.

It bears mention that many lethal sets on manifolds are discovered through trial and error. One strategy is to imagine these manifolds as folded pieces of paper, and to identify lethal sets on their unfolded nets. This technique is discussed in more detail in Chapter {CONSTRUCTIONS}.

3.3 Block Recursion

Note that there are certain broad structures in a cube that, if present, immediately guarantee it become fully infected. Of greatest importance here is the observation that certain configurations of fully infected sub-cubes (which we shall call blocks) will cause the larger brick to become infected.

Furthermore, note that if each of these smaller blocks is infected with a minimum lethal set, the composite larger brick will also be infected with a minimum lethal set (barring some considerations for divisibility).

The proof of this claim makes use of the so-called modified bootstrap process in $[n]^d$, discussed in [?] and [?]. This is a strengthened variation of the problem introduced in the previous chapter, whereby vertices in the $[n]^d$ grid become infected if and only if they are adjacent to infected vertices along edges in each of the d directions. For example, in the $[n]^2$ grid, a vertex that sees infection in one of both the North/South and East/West directions will itself become infected, whereas a vertex with infected neighbors to the East and West (but not North and South) will not.

In particular, the lemma considers composite grids $[n]^d$ where each vertex $x = (x_1, \ldots, x_d) \in [n]^d$ is itself a smaller block G_x . We prove that lethal sets on these grids can be built from the smaller lethal sets on each G_x .

Lemma 3.6. For $n, d \ge 1$, let $A = (a_{i,j})$ be a $d \times n$ matrix of positive integers, and let $b_i = \sum_{j=1}^n a_{i,j}$, for $1 \ge j \ge d$. Let S be a lethal set under the modified process on $[n]^d$, and for each vertex $\vec{x} = (x_1, \dots, x_v) \in S$, let $T_{\vec{x}}$ be a lethal set on $\prod_{i=1}^d [a_{i,x_i}]$ under

d-neighbor percolation. Then

$$m(b_1,\ldots,b_d,d) \leq \sum_{\vec{x}\in S} |T_{\vec{x}}|.$$

Proof. We imagine sub-dividing the $\prod_{i=1}^d [b_i]$ brick into smaller blocks by partitioning each of the d axes into segments $a_{i,1}, a_{i,2}, \ldots, a_{i,n}, 1 \le i \le d$. Each block is given by a unique product of these segments, and represented by a vector $\vec{x} = (x_1, \ldots, x_d) \in [n]^d$. Formally, for each such \vec{x} , let $G_{\vec{x}}$ be the block with vertex set

$$\prod_{i=1}^{d} \{1 + \sum_{j=1}^{x_i-1} a_{i,j}, \dots, \sum_{j=1}^{x_i} a_{i,j}\},\$$

and edges between vertices that differ by one in exactly one coordinate. Figure 3.2 illustrates the block $G_{\vec{x}}$ for $\vec{x} = (1, 2, 2) \in [2]^3$. Observe that $G_{\vec{x}}$ is isomorphic to $\prod_{i=1}^d [a_{i,x_i}]$.

For each $\vec{x} \in S$, let $A_{\vec{x}}$ be the vertices of $G_{\vec{x}}$ corresponding to the vertices of $T_{\vec{x}}$ under isomorphism from $\prod_{i=1}^d [a_{i,x_i}]$ to $G_{\vec{x}}$, and let $A_0 = \bigcup A_{\vec{x}}$. Observe that $|A_0| = \sum_{\vec{x} \in S} |T_{\vec{x}}|$. We show that A_0 is lethal on $\prod_{i=1}^d [b_i]$.

By the definition of $T_{\vec{x}}$, for each $\vec{x} \in S$, $A_{\vec{x}}$ is lethal on $G_{\vec{x}}$. Imagine running the d-neighbor process until all blocks $G_{\vec{x}}$ are fully infected. We claim that this is sufficient to infect all remaining vertices of $\prod_{i=1}^d [b_i]$. Consider the remaining blocks $G_{\vec{x}}$, for $\vec{x} \in [n]^d \setminus S$. Since S is lethal under the modified process, each $G_{\vec{x}}$ is adjacent to fully infected blocks in all d directions. In particular, if we consider expanding out the faces of $G_{\vec{x}}$ towards these infected blocks, the resulting cube has d fully infected faces that share a common corner. By Corollary 7.3, this structure will infect all the vertices of $G_{\vec{x}}$. Repeating this process on each uninfected region of $\prod_{i=1}^d [b_i]$ (as they are exposed under the modified process) ultimately results in all vertices becoming infected. This completes the proof.

We note that although the lemma above is true in full generality, we are only concerned with the particular case where n=2 and d=3. The following corollary proves that the bound in Lemma 3.6 is tight for n=2 and d=3, if lethal sets on at least three of the constituent blocks are perfect.

Corollary 3.7. Let $A = (a_{i,j})$ be a 3×2 matrix of positive integers, and let $b_i = a_{i,1} + a_{i,2}$ for all $1 \le i \le 3$. Then $m(b_1, b_2, b_3, 3)$ is at most

$$m(a_{1,1}, a_{2,1}, a_{3,1}, 3) + m(a_{1,2}, a_{2,2}, a_{3,1}, 3) + m(a_{1,2}, a_{2,1}, a_{3,2}, 3) + m(a_{1,1}, a_{2,2}, a_{3,2}, 3).$$

Furthermore, this bound is tight if at least 3 of the constituent grids are perfect.

Proof. The upper bound on $m(b_1, b_2, b_3, 3)$ is a direct consequence of Lemma 3.6, since (1, 1, 1), (2, 2, 1), (2, 1, 2), (1, 2, 2) is lethal under the modified process on $[2]^3$. If all grids are perfect, then:

$$m(a_{1,1}, a_{2,1}, a_{3,1}, 3) + m(a_{1,2}, a_{2,2}, a_{3,1}, 3) + m(a_{1,2}, a_{2,1}, a_{3,2}, 3) + m(a_{1,1}, a_{2,2}, a_{3,2}, 3)$$

$$= \frac{a_{1,1}a_{2,1} + a_{2,1}a_{3,1} + a_{3,1}a_{1,1}}{3} + \frac{a_{1,2}a_{2,2} + a_{2,2}a_{3,1} + a_{3,1}a_{1,2}}{3} + \frac{a_{1,1}a_{2,2} + a_{2,2}a_{3,2} + a_{3,2}a_{2,1}}{3} + \frac{a_{1,1}a_{2,2} + a_{2,2}a_{3,2} + a_{3,2}a_{1,1}}{3}$$

$$= \frac{(a_{1,1} + a_{1,2})(a_{2,1} + a_{2,2}) + (a_{2,1} + a_{2,2})(a_{3,1} + a_{3,2}) + (a_{3,1} + a_{3,2})(a_{1,1} + a_{1,2})}{3}$$

$$= \frac{b_1b_2 + b_2b_3 + b_3b_1}{3}.$$

Similarly, suppose, without loss of generality, that $(a_{1,1}, a_{2,1}, a_{3,1})$ is optimal and the remaining grids are perfect. Then:

$$m(a_{1,1}, a_{2,1}, a_{3,1}, 3) + m(a_{1,2}, a_{2,2}, a_{3,1}, 3) + m(a_{1,2}, a_{2,1}, a_{3,2}, 3) + m(a_{1,1}, a_{2,2}, a_{3,2}, 3)$$

$$= \left\lceil \frac{a_{1,1}a_{2,1} + a_{2,1}a_{3,1} + a_{3,1}a_{1,1}}{3} \right\rceil + \frac{a_{1,2}a_{2,2} + a_{2,2}a_{3,1} + a_{3,1}a_{1,2}}{3} + \frac{a_{1,1}a_{2,2} + a_{2,2}a_{3,2} + a_{3,2}a_{2,1}}{3} + \frac{a_{1,1}a_{2,2} + a_{2,2}a_{3,2} + a_{3,2}a_{1,1}}{3}$$

$$= \left\lceil \frac{(a_{1,1} + a_{1,2})(a_{2,1} + a_{2,2}) + (a_{2,1} + a_{2,2})(a_{3,1} + a_{3,2}) + (a_{3,1} + a_{3,2})(a_{1,1} + a_{1,2})}{3} \right\rceil$$

$$= \left\lceil \frac{b_1b_2 + b_2b_3 + b_3b_1}{3} \right\rceil.$$

In both cases, we obtain the lower bound $m(b_1, b_2, b_3, 3)$. This completes the proof. \square

3.4 Applying the recursion

Corollary 3.7 provides a proscriptive method for constructing optimal and perfect lethal sets recursively, provided the existence of sufficiently many small building blocks. In the following chapter, we use this technique to obtain perfect lethal sets on all (b_1, b_2, b_3) grids, for $b_1, b_2, b_3 \geq 5$, and optimal lethal sets on all (b_1, b_2, b_3) grids, for $b_1, b_2, b_3 \geq 11$. To facilitate this process, we first present some useful constructions of lethal sets (discussed in greater detail in Chapter [Constructions]), as well as particular applications of Corollary 3.7 that hold for general grids.

Proposition 3.8. For all $k \ge 1$ such that $k \ne 2$, (3,3,k) is perfect.

Proposition 3.9. For all $k \ge 2$, (3, 6, k) is perfect.

Proposition 3.10. For all $k \equiv 3 \pmod{6}$ and $l \equiv 1 \pmod{2}$, (3, k, l) is perfect.

Combining the above propositions with Corollary 3.7, we are able to obtain the following lemmas.

Lemma 3.11. Suppose (b_1, b_2, b_3) is optimal. Then $(b_1 + 3, b_2 + 3, b_3 + 3)$ is optimal.

Proof. By Proposition 3.8, each of $(b_1, 3, 3), (3, b_2, 3), (3, 3, b_3)$ is perfect. Therefore, by Corollary 3.7,

$$m(b_1+3,b_2+3,b_3+3,3) = m(b_1,b_2,b_3,3) + m(b_1,3,3,3) + m(3,b_2,3,3) + m(3,3,b_3,3),$$

and so (b_1+3,b_2+3,b_3+3) is optimal.

3.5 Examples and Notation

3.6 Regional vs. Temporal Infections

Chapter 4

A Tight Bound on Grids of Size ≥ 5

4.1 Introduction and Definitions

Let the ordered tuple (a, b, c) represent the $a \times b \times c$ grid G where $a \geq b \geq c$. We refer to c as the *thickness* of G. For example, the tuple (5, 3, 3) represents a $5 \times 3 \times 3$ grid of thickness 3. We refer to a tuple as *divisible*, or a *divisibility case*, if and only if $ab + bc + ca \equiv 0 \pmod{3}$. If a tuple is divisible and percolates at the lower bound, we refer to it as *perfect*. Observe that the divisibility cases are precisely those grids with integral lower bounds. The divisibility cases of thicknesses belonging to the three residue classes modulo 3 are illustrated in {Figure something}.

In the following lemmas, we use the notation (a, b, c) + (x, y, z) = (a+x, b+y, c+z) to represent respective increases of x, y, and z to the side lengths a, b, and c of G. We note the following:

Remark 4.1. By applying the recursion, (a, b, c) + (x, y, z) percolates at the lower bound when either:

- 1. (a,b,c),(a,y,z),(x,b,z),(x,y,c) all percolate at the lower bound; or
- 2. (x, y, z), (x, b, c), (a, y, c), (a, b, z) all percolate at the lower bound.

We shall call a thickness *complete* if it can be shown that all divisibility cases in that thickness are perfect. In this section, we demonstrate that thickness 5, thickness 6 and thickness 7 are all complete. As these belong to the residue classes 2, 0, and 1 modulo 3, respectively, we then use a recursive construction to show that all larger grids are also complete.

4.2 Completeness of Thickness 5

Leveraging {lemmas from earlier chapters yet to be written}, we show that all divisibility cases in thickness 5 percolate at the lower bound.

NOTE: THE FOLLOWING LEMMAS HOLD ASSUMING WE HAVE A GENERAL CONSTRUCTION FOR (2,3,3k) FOR ALL k.

Lemma 4.2. All divisibility cases for grids of the form (k, 5, 5) are perfect.

Proof. We consider grids obtained from (5,2,2) + (a,3,3), for $a \equiv 0 \pmod{3}$ and a > 3. By remark 4.1, it is sufficient to show that (5,2,2), (5,3,3), (a,2,3), (a,2,3) are all perfect. By {a bunch of constructions}, each of these grids percolates at the lower bound for a > 3. We therefore obtain all grids of the form (k,5,5), for k > 8. The only missing grids are (5,5,5) and (8,5,5), which we have by construction. This completes the proof.

Lemma 4.3. All divisibility cases for grids of the form (k, 6, 5) percolate at the lower bound.

Proof. We consider grids obtained from (6,3,2) + (a,3,3), for $a \equiv 0 \pmod{3}$ and a > 3. By remark 4.1, it is sufficient to show that (6,3,2), (6,3,3), (a,3,3), (a,3,2) are all perfect. By {a bunch of constructions}, each of these grids percolates at the lower bound for a > 3. We therefore obtain all grids of the form (k,6,5), for k > 8. The only missing grid is (6,6,5), which we have by construction. This completes the proof. \square

Lemma 4.4. Thickness 5 is complete.

Proof. Let (a, b, 2) represent an arbitrary (divisible) grid of thickness 2, and let $x = a \pmod{6}$ and $y = b \pmod{6}$. By {some as of yet unwritten construction}, we have that (a, b, 2) percolates at the lower bound for all $x, y \in \{0, 2, 3, 5\}$, where $x \neq y$. We consider two constructions: (a, b, 2) + (6, 3, 3) and (a, b, 2) + (6, 6, 3).

By item (1) of the remark, in order to show that (a, b, 2) + (6, 3, 3) percolates at the lower bound, it is sufficient to show that (a, b, 2), (a, 3, 3), (6, b, 3), (6, 3, 2) all percolate at the bound. By {more unwritten constructions}, this is true for all $x, y \in \{0, 2, 3, 5\}$, where $x \neq y, a, b > 1$, and at least one of $\{a, b\} > 2$. (Note that if a = 2, one of the tuples is (2, 3, 3), which does not percolate at the lower bound; we accommodate for this by re-writing (a, b, 2) + (6, 3, 3) as (a, b, 2) + (3, 6, 3).) The resulting tuple (a', b', 5) is a grid of thickness 5, with a' and b' in the same residue class modulo (a', b', b', 5) and at least one of $(a', b') \geq 9$. From {some figure representing the divisibility cases of thickness 5}, we see that the lower bound on a' and b' omits all grids of the form (5, 5, k) and (5, 6, k), as well as the singular grid (8, 8, 5).

Applying an analogous argument to (a, b, 2) + (6, 6, 3), we must demonstrate that (a, b, 2), (a, 6, 3), (6, b, 3), (6, 6, 2) all percolate at the lower bound. By {some other constructions}, we again find that this holds for all $x, y \in \{0, 2, 3, 5\}$, where $x \neq y$ and a, b > 1. This gives all thickness 5 tuples (a', b', 5) with a' and b' in different residue classes modulo 6, where $a', b' \geq 8$.

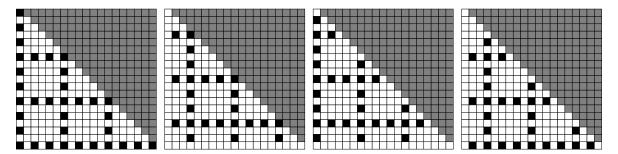


Figure 4.1: Thickness 6 grids with perfect percolating sets as obtained in lemma 4.5 (left), and divisibility cases of thickness 6 (right).

Combining these results, we have completeness for all grids of thickness 5 except those of the form (5,5,k) and (5,6,k), and the singular grid (8,8,5). By lemmas 4.2 and 4.3, and {some construction for (8,8,5)}, these cases are also complete, and so thickness 5 is complete. This completes the proof.

4.3 Completeness of Thickness 6

We shall show that all grids of thickness 6 can be obtained recursively from (3n, m, 3), where $n, m \equiv 1 \pmod{2}$ (this is that general thickness 3 construction), and one of $\{(3,3,3),(6,6,3),(6,3,3),(3,6,3)\}$. We examine each of these cases separately and show that each is complete.

(NOTE (to Peter and Jon): I have struggled a bit with the canonical way to describe grids. I like the tuple representation (a,b,c) where WLOG $a \leq b \leq c$. However, this becomes a bit mucky in the following proofs, because (3n,m,3) potentially violates this rule if n is large and m is small. To accommodate this, I have written "grids of the form (a,b,6), where $a \equiv 0 \pmod 6$ and $b \equiv 0 \pmod 2$, or $b \equiv 0 \pmod 6$ and $a \equiv 0 \pmod 2$," in an attempt to address the circumstance where the ordering of the tuple is flipped because n is large and m is small. However, I think this may just muddy the waters.)

Lemma 4.5. All grids of the form (a, b, 6), where $a \equiv 0 \pmod{6}$ and $b \equiv 0 \pmod{2}$, or $b \equiv 0 \pmod{6}$ and $a \equiv 0 \pmod{2}$, percolate at the lower bound.

Proof. We consider (3n, m, 3) + (3, 3, 3), for $n, m \equiv 1 \pmod{2}$. By remark 4.1, we have that (3n, m, 3) + (3, 3, 3) percolates if (3n, m, 3), (3n, 3, 3), (3, m, 3), (3, 3, 3) all percolate. By construction {yet to be named}, all grids (a, 3, 3) are perfect. Therefore, (3, 3n, m) + (3, 3, 3) is perfect. Note that this grid is of the form (3k, l, 6), where $k, l \equiv 0 \pmod{2}$. This is equivalent to grids of the form (a, b, 6), where $a \equiv 0 \pmod{6}$ and $b \equiv 0 \pmod{2}$, or $b \equiv 0 \pmod{6}$ and $a \equiv 0 \pmod{2}$. This completes the proof. \Box

Lemma 4.6. All grids of the form (a, b, 6), where $a \equiv 3 \pmod{6}$ and $b \equiv 1 \pmod{2}$, or $b \equiv 3 \pmod{6}$ and $a \equiv 1 \pmod{2}$, percolate at the lower bound.

Proof. We apply the same argument as above, this time considering (3n, m, 3) + (6, 6, 3), for $n, m \equiv 1 \pmod{2}$. Again, by remark 4.1, it is sufficient to show that (3n, m, 3), (3n, 6, 3), (6, m, 3), (6, 6, 3) are all perfect. By {more thickness 3 constructions}, each of these grids percolates at the lower bound. The resulting grid, (3n, m, 3) + (6, 6, 3), for $n, m \equiv 1 \pmod{2}$, is of the form (a, b, 6), for $a \equiv 3 \pmod{6}$ and $b \equiv 1 \pmod{2}$, or $b \equiv 3 \pmod{6}$ and $a \equiv 1 \pmod{2}$. This completes the proof.

Lemma 4.7. All grids of the form (a, b, 6), where $a \equiv 3 \pmod{6}$ and $b \equiv 0 \pmod{2}$, or $b \equiv 3 \pmod{6}$ and $a \equiv 0 \pmod{2}$, percolate at the lower bound.

Proof. Similarly to the previous proofs, we consider (3n, m, 3) + (6, 3, 3), for $n, m \equiv 1 \pmod{2}$. We show that (3n, m, 3), (3n, 3, 3), (6, m, 3), (6, 3, 3) are all perfect. By the same thickness 3 constructions, each of these grids percolates at the lower bound. Therefore, by remark 4.1, (3n, m, 3) + (6, 3, 3) is perfect. Furthermore, observe that (3n, m, 3) + (6, 3, 3) is of the form (a, b, 6), where $a \equiv 3 \pmod{6}$ and $b \equiv 0 \pmod{2}$. This completes the proof.

Lemma 4.8. All grids of the form (a, b, 6), where $a \equiv 0 \pmod{6}$ and $b \equiv 1 \pmod{2}$, or $b \equiv 0 \pmod{6}$ and $a \equiv 1 \pmod{2}$, percolate at the lower bound.

Proof. We consider (3n, m, 3)+(3, 6, 3), for $n, m \equiv 1 \pmod{2}$. We show that (3n, m, 3), (3n, 6, 3), (3, m, 3), (3, 6, 3) are all perfect. By {thickness 3 constructions} and remark 4.1, (3n, m, 3) + (3, 6, 3) is perfect. Observe that (3n, m, 3) + (3, 6, 3) is of the form (a, b, 6), where $a \equiv 0 \pmod{6}$ and $b \equiv 1 \pmod{2}$. This completes the proof.

Lemma 4.9. Thickness 6 is complete.

Proof. All divisibility cases for thickness 6 are grids of the form (x, y, 6) such that at least one of $\{x, y\}$ is congruent to 0 modulo 3. Lemmas 4.5, 4.6, and 4.7 cover all such cases. The result follows.

4.4 Completeness of Thickness 7

We show that all divisibility cases for grids of thickness 7 percolate at the lower bound. Observe that divisibility cases for thickness 7 consist of grids of the form (x, y, 7) for x, y in residue classes $\{0, 1, 3, 4\}$ modulo 6. We separate these divisibility cases into the following four categories and show that each category is complete:

- 1. (x, y, 7) for $x, y \in \{1, 4\}$ and $x \equiv y \pmod{6}$;
- 2. (x, y, 7) for $x, y \in \{1, 4\}$ and $x \not\equiv y \pmod{6}$;

- 3. (x, y, 7) for $x, y \in \{0, 3\}$ and $x \equiv y \pmod{6}$;
- 4. (x, y, 7) for $x, y \in \{0, 3\}$ and $x \not\equiv y \pmod{6}$.

Lemma 4.10. All grids of the form (x, y, 7) for $x, y \in \{1, 4\}$ and $x \equiv y \pmod{6}$ are complete.

Proof. Consider the construction (a, b, 2) + (8, 5, 5) for $a, b \in \{2, 5\}$ and $a \not\equiv b \pmod{6}$. Observe that this construction obtains all grids of the form described in (1) above. We show that the grids (a, b, 2), (a, 5, 5), (8, b, 5), (8, 5, 2) are all complete. {The fact that these grids are complete follows from a number of constructions and the observation that thickness 5 is complete.} By remark 4.1, the construction (a, b, 2) + (8, 5, 5) percolates at the lower bound. This completes the proof.

Lemma 4.11. All grids of the form (x, y, 7) for $x, y \in \{1, 4\}$ and $x \not\equiv y \pmod{6}$ are complete.

Proof. Consider the construction (a, b, 2) + (5, 5, 5) for $a, b \in \{2, 5\}$ and $a \not\equiv b \pmod{6}$. Observe that this construction obtains all grids of the form described in (2) above. We consider the grids (a, b, 2), (a, 5, 5), (5, b, 5), (5, 5, 2). By {known constructions and completeness of thickness 5}, each of these grids is perfect, and so by remark 4.1, (a, b, 2) + (5, 5, 5) percolates at the lower bound. This completes the proof.

Lemma 4.12. All grids of the form (x, y, 7) for $x, y \in \{0, 3\}$ and $x \equiv y \pmod{6}$ are complete.

Proof. Consider the construction (a, b, 2) + (6, 9, 5) for $a, b \in \{0, 3\}$ and $a \not\equiv b \pmod{6}$. Observe that this construction contains all grids described in (3) above. We consider the grids (a, b, 2), (a, 9, 5), (6, b, 5), (6, 9, 2). Observe that (3, 9, 5) is perfect by construction, and (a, 9, 5) is perfect in general for a > 3. Similarly, (6, 3, 5) is perfect by construction, and (6, b, 5) is perfect for all other b > 3. Finally, (a, b, 2) and (6, 9, 2) are perfect by construction. Therefore, by remark 4.1, (a, b, 2) + (5, 5, 5) percolates at the lower bound. This completes the proof.

Lemma 4.13. All grids of the form (x, y, 7) for $x, y \in \{0, 3\}$ and $x \not\equiv y \pmod{6}$ are complete.

Proof. Consider the construction (a, b, 2) + (6, 6, 5) for $a, b \in \{0, 3\}$ and $a \not\equiv b \pmod{6}$. Observe that this construction contains all grids described in (4) above. We consider the grids (a, b, 2), (a, 6, 5), (6, b, 5), (6, 6, 2). Observe that (3, 6, 5) is perfect by construction, and (a, 6, 5) is perfect in general for a > 3. Similarly, (6, 3, 5) is perfect by construction, and (6, b, 5) is perfect for all other b > 3. Finally, (a, b, 2) and (6, 6, 2) are perfect by construction. Therefore, by remark 4.1, (a, b, 2) + (5, 5, 5) percolates at the lower bound. This completes the proof.

Lemma 4.14. Thickness 7 is complete.

Proof. By lemmas 4.10, 4.11, 4.12, and 4.13, all divisibility cases for thickness 7 percolate at the lower bound. $\hfill\Box$

4.5 Completeness of Grids of Size ≥ 5

We can get completeness in every residue class modulo 3 by simply considering the grids obtained from (x, y, z) + (3, 3, 3).

Chapter 5 Tight Bound on (a, b, c) Grids for $a \ge b \ge c \ge 11$

Chapter 6 Thickness One

Chapter 7 Constructions

7.1 Introduction

In this chapter, we present diagrammed proofs of lethal sets that percolate at the lower bound. The proofs are organized by the thickness of the grid. Many of the constructions in the following sections belong to infinite families of either optimal or perfect sets. In this case, we shall examine the grids by region, and observe that certain regions can be expanded to arbitrarily large sizes using mathematical induction.

We shall call a thickness *semi-complete* if all divisibility cases are optimal.

7.2 Useful lemmas and observations

We shall see that similar patterns and structures appear with some regularity in optimal sets. These structures always infect entire regions, and it will be helpful to recognize them within larger grids when they appear.

Lemma 7.1. Let $G = P_a \square P_b$ be a graph and $S \subseteq V(G)$ be a subset of the vertices of G. Let $G[\overline{S}]$ be the subgraph of G induced by vertices not in G. Then G is a lethal set if and only if $G[\overline{S}]$ is cycle-free, and has no paths between any two boundary vertices.

Lemma 7.2. Let G be the grid graph (a,b,c) and let A be a set of infected vertices in G. Let $H = G[\overline{A}]$ be the subgraph of G induced by uninfected vertices. Let $\{F,F'\} \subseteq V(G)$ be orthogonal faces of G. If H does not contain a path between vertices $v \in F, v' \in F'$, then G percolates.

Proof. We proceed by induction on |V(H)| = abc - |A|. If |V(H)| = 0, then all vertices of G are infected and we are done. Suppose |V(H)| > 0, and consider a connected component Y of H. By hypothesis, Y does not contain a path between any two orthogonal faces of G. Therefore, without loss of generality, there exists a face $X = \{a\} \times \{1, \ldots, b\} \times \{1, \ldots, c\}$ such that $V(Y) \cap X = \emptyset$. Take the largest i such that

 $X_i = \{a\} \times \{1, \dots b\} \times \{1, \dots, c\}$ contains a vertex of Y. Repeat this process to obtain maximal faces Y_i and Z_k in the b and c directions, respectively.

Observe that this construction gives a vertex $v = (i, j, k) \in V(Y)$, and planes $X_{i+1}, Y_{j+1}, Z_{k+1}$ such that $(X_{i+1} \cup Y_{j+1} \cup Z_{k+1}) \cap V(Y) = \emptyset$. In particular, note that $(i+1, j, k), (i, j+1, k), (i, j, k+1) \in N(v)$. Since these three vertices belong to A and $v \notin A$, v becomes infected. Furthermore, since $|V(H) \setminus \{v\}| < |V(H)|$, the resulting graph percolates by induction. This completes the proof.

Corollary 7.3. Let G be the grid graph (a,b,c). If a set A is lethal on three mutually orthogonal faces of G, then A is lethal on G.

Proof. Let X, Y, Z be three mutually orthogonal faces of G. By hypothesis, $X \cup Y \cup Z \subseteq A_t$ for some time t. Therefore, the graph $H = G[\overline{A_t}]$ cannot contain a path between vertices on orthogonal faces of G. By lemma 7.2, G percolates.

Corollary 7.4. Let G be the grid graph (a, b, c) and let A be a set of infected vertices of G. Let $\{R_1, \ldots, R_k\}$ be a partition of $V(G) \setminus A$ into sub-grids $(a_1, b_1, c_1), \ldots, (a_k, b_k, c_k)$ such that each R_i is bounded by a lethal infection on three mutually orthogonal faces. Then A is lethal on G.

Proof. By hypothesis and corollary 7.3, A is lethal on each R_i . Since $A \cup R_1 \cup \cdots \cup R_k = V(G)$ and each R_i becomes infected, A must be lethal.

Definition 7.5. Let a manifold M of the grid graph G = (a, b, c) be any set of vertices A satisfying the conditions of corollary 7.4.

Definition 7.6. Let a proper unfolding of G = (a, b, c) be a planar representation of the manifold of G. This can be thought of as a special type of folding net of G, such that when assembled the resulting structure satisfies the conditions of corollary 7.4.

Corollary 7.7. Let M be the proper unfolding of a manifold of G = (a, b, c), and let A be a lethal set on M. Then A is lethal on G.

Proof. The proof follows directly from definitions 7.5 and 7.6. \Box

7.3 Thickness 1

There are two general constructions in thickness 1 that percolate at the surface area bound. The first construction is perfect for all $(2^n - 1, 2^n - 1, 1)$ grids, and originates in a 2021 paper by Benevides et al. [?]. The second construction is optimal for all grids (a, b, 1), where $a \equiv 5 \pmod{6}$, $b \equiv 1 \pmod{2}$, and $a, b \geq 5$. As such grids constitute non-divisibility cases, this construction is not perfect.



Figure 7.1: A perfect percolating set for (3,3,1).

7.3.1 Purina

We refer to this construction colloquially as the Purina construction, due to the similarly between its instance on the (3,3,1) grid and the logo of the pet food brand. No funding has been offered, but we are open to the possibility. A more extensive discussion on this pattern can be found in [?].

Construction 7.8. All grids of the form $(2^n - 1, 2^n - 1, 1)$ are perfect.

Proof. This is a recursive construction built from the base component piece shown in figure 7.1. Note that this (3,3,1) construction is lethal under the 3-neighbor bootstrap process, and that it meets the surface area bound:

$$\frac{1}{3} \cdot (ab + bc + ca) = \frac{1}{3} \cdot (9 + 3 + 3) = 5.$$

For larger grids of size $(2^n - 1, 2^n - 1, 1)$, join four copies of $(2^{n-1} - 1, 2^{n-1}, 1)$ about two perpendicular corridors, and infect the vertex at their intersection (figure 7.2). Observe that the resulting set is lethal: each of the four smaller grids is lethal by hypothesis, and the remaining vertices induce a forest with disconnected boundary points, which percolates by lemma 7.1. Furthermore, note that

S.A.
$$(2^{n} - 1, 2^{n} - 1, 1) = \frac{1}{3} \cdot (2^{2n} - 1)$$

= $4 \cdot \frac{1}{3} \cdot (2^{2n-2} - 1) + 1 = 4 \cdot \text{S.A.}(2^{n-1} - 1, 2^{n-1}, 1) + 1$,

and therefore this construction is perfect.

7.3.2 Snakes

As indicated by lemma 7.1, a fundamental characteristic of lethal sets S is the presence of an initially uninfected corridor, bounded by walls of infection. This structure is apparent in the second diagrams of figures 7.2 and 7.5. These corridors correspond to forests in the complement $G[\overline{S}]$ of S. In this subsection, we provide a general method for constructing such corridors in (a, b, 1) grids where $a \equiv 5 \pmod{6}$ and $b \equiv 1 \pmod{2}$.

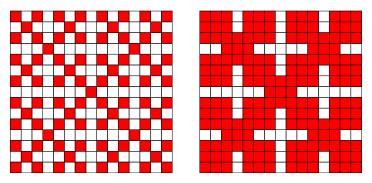


Figure 7.2: A perfect percolating set for (15, 15, 1).

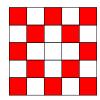


Figure 7.3: An optimal percolating set for (5, 5, 1).

Construction 7.9. All grids of the form (a, b, 1), $a \equiv 5 \pmod{6}$, $b \equiv 1 \pmod{2}$, and $a, b \geq 5$ are optimal.

Proof. For grids of the form (a, b, 1), $a \equiv 5 \pmod{6}$, $b \equiv 1 \pmod{2}$, we construct an optimal infected set and show that it percolates by lemma 7.1. For the base case, consider the (5,5,1) grid G illustrated in figure 7.3. Observe that this construction is optimal. Now consider the grid G' resulting from the insertion of a (5,2k,1) block, as shown in figure 7.4. Note that the subgraph induced by the uninfected vertices of G' satisfies the conditions of lemma 7.1. Furthermore, note that if any (5, n, 1) grid is optimal, the (5, n+2, 1) grid resulting from such a construction has surface area bound S.A.(5, n, 1) + 4, which agrees with the number of infected vertices.

To extend this construction in the vertical direction, we introduce a kink in the snaking infection. This kink requires six rows to produce a repeating pattern. The structure of this design is shown in figure 7.5. For grids of smaller width, the same construction gives optimal percolating sets; however, the snaking pattern is increasingly difficult to recognize in thin grids.

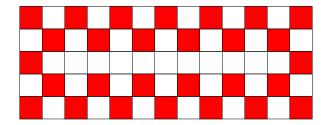


Figure 7.4: An optimal percolating set for (5, 13, 1).

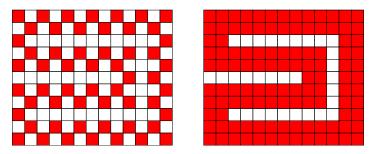


Figure 7.5: An optimal percolating set for (11, 13, 1).

7.4 Thickness 2

In this section we examine four infinite families of perfect grids. We show that each has a manifold that admits a lethal set of perfect size. We note that such lethal sets are likely to exist for nearly all divisibility cases in thickness two; however, constructions are elusive and those presented here are sufficient to prove the main result of this thesis.

Unlike those presented in the previous section, the following proofs all leverage lemma 7.2. As a consequence, their argumentative structure remains broadly the same, even as the constructions themselves appear quite different. For this reason, we shall outline this structure here, before examining the specific proofs.

We begin by demonstrating that a grid G admits a manifold M. To show so, we identify the regions R_1, \ldots, R_k that partition $V(G) \setminus M$. In our diagrams, these regions are represented by the volumes enclosed by three perpendicular blue, green, and red walls. We then identify a proper unfolding H of M and show that H admits a lethal set A, where |A| = SA(G). Finally, we apply corollary 7.7 to prove that G is perfect.

Construction 7.10. All (a, 3, 2) grids with $a \equiv 0 \pmod{6}$ are perfect.

Proof. Let G be an (a,3,2) grid with $a \equiv 0 \pmod{6}$, and let M be a manifold of G and H be its proper unfolding (see figure 7.7). Observe that M is indeed a manifold: it partitions $V(G) \setminus M$ into two sets R_1 and R_2 , both bounded by mutually orthogonal red, green, and blue faces (see figure 7.7). Furthermore, note that H is obtained from M by cutting along seams between red and green faces, and flattening the figure. It follows that H is a proper unfolding of G. We show that H admits a perfect lethal set.

Consider the initial infection A of H illustrated in figure 7.8. Observe that A infects all vertices of H by lemma 7.1, with the exception of two regions on the far left and far right. However, note that upon refolding, the two cells marked with an "X" in H represent the same cell in G. This is sufficient to infect the remaining regions of H, and by corollary 7.7, A is lethal on G. Finally, a simple calculation reveals that |A| matches the surface area bound, and therefore is perfect.

We can extend this construction in the x direction by inserting the repeated structure of six columns; this augmentation percolates by lemma 7.1 and agrees with the

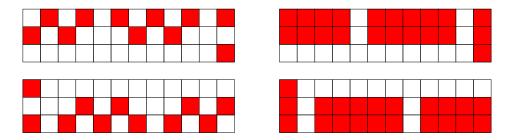


Figure 7.6: A perfect percolating set for (3, 12, 2).

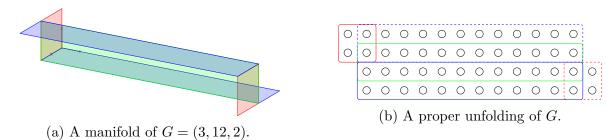


Figure 7.7: A proper unfolding of G = (3, 12, 2). Colored rectangles indicate faces of G. Dashed lines indicate that cells appear on different layers.

surface area bound for all larger grids. This completes the proof. \Box

Construction 7.11. All (a, 3, 2) grids with $a \equiv 3 \pmod{6}$ are perfect.

Proof. (This construction is the same as the previous one, except the the final four columns are augmented slightly to accommodate the $0 \pmod{3}$ requirement. Instead of deriving a proper unfolding, it is probably easier to simply show that this small change is sufficient to guarantee lethality, and agrees with the S.A. bound.)

Construction 7.12. All (a, b, 2) grids with $a, b \in \{2, 5\} \pmod{6}$, $a \neq b \pmod{6}$, and a, b > 2 are perfect.

Proof. Let G be an (a, b, 2) grid with $a, b \in \{2, 5\} \pmod{6}$ and $a \neq b \pmod{6}$, and let M be a manifold of G and H be its proper unfolding (figure 7.10). Note that M partitions the vertices of $V(G) \setminus M$ into two disjoint sets R_1 and R_2 , both bounded by mutually orthogonal red, green, and blue faces. Note, also, that H is obtained from M by cutting along seams between red and green faces, and flattening the figure. Therefore, H is a proper unfolding of G. We show that H admits a perfect lethal set.

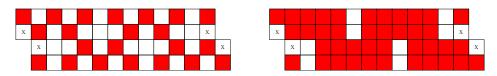


Figure 7.8: A percolating set on the proper unfolding of G = (3, 12, 2).

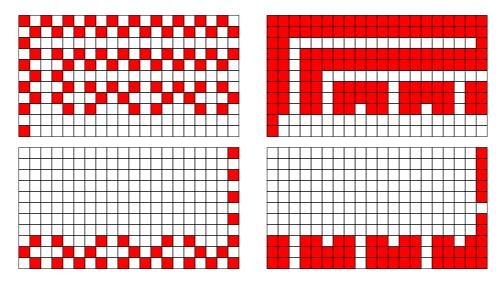


Figure 7.9: A perfect percolating set for (11, 20, 2).

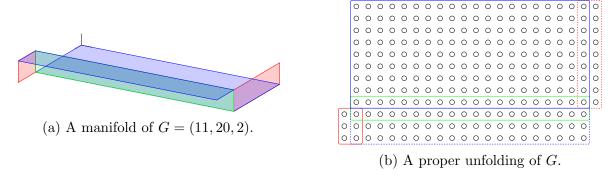


Figure 7.10: A proper unfolding of G = (11, 20, 2). Colored rectangles indicate faces of G. Dashed lines indicate that cells appear on different layers.

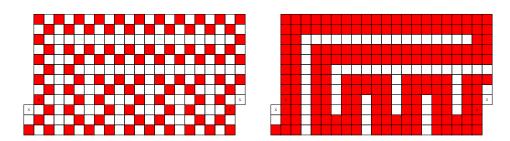


Figure 7.11: A percolating set on the proper unfolding of G = (11, 20, 2).

Consider the initial infection A of H as shown in figure 7.11. By lemma 7.1, A infects all vertices of H, with the exception of two regions on the left- and right-most sides of the grid. However, note that the vertices labeled "X" in figure 7.11 represent that same vertex in G. This permits A to infect the remaining healthy vertices, thereby proving that A is lethal on H. By corollary 7.7, A is lethal on G. Finally, a simple calculation shows that |A| satisfies the surface area bound on G, and so A is perfect.

The above construction holds for $a \geq 5$ and $b \geq 8$. It can be extended in the x direction by inserting a block of width 6, representing the repeating vertical snaking pattern. Similarly, it can be extended in the y direction by inserting a block of height 6, representing the horizontal snaking pattern. Both such augmentations spawn infections that satisfy the conditions of lemma 7.1, and a simple calculation reveals that they agree with the surface area bound. This concludes the proof.

Construction 7.13. All (a, b, 2) grids with $a, b \in \{0, 3\} \pmod{6}$ and $a \neq b \pmod{6}$ are perfect.

Proof. Consider the (21, 12, 2) grid G shown in figure 7.12. Let H be a unfolding of G (figure 7.13). Observe that H is proper: three mutually orthogonal faces of G_1 are shown by blue, green and dashed red regions, and mutually orthogonal faces of G_2 are shown by the red, green and dashed blue regions. We show that H admits a lethal set of size S.A.(12, 21, 2) = 106. Consider such a set, as shown in figure 7.14. (Observe that this is the same set as shown in figure 7.12.) By lemma 7.1, this set percolates with the exception of two G_4 s in the top and bottom of the grid. However, notice that one of these cells is a duplicate of an already infected cell. (This duplication is a consequence of the proper unfolding of G.) Therefore, H admits a lethal set, and by corollary 7.7, G is perfect.

For all larger grids, observe that the snaking corridor in the left side G can be extended by multiples of 6 in both the x and y directions. These resulting grid still percolates under lemma 7.1. A simple calculation verifies that such an alteration produces initial infections at the surface area bound.

7.5 Thickness 3

Construction 7.14. All (a, b, 3) grids G with $a \equiv 3 \pmod{6}$ and $b \equiv 1 \pmod{2}$ are perfect.

Proof. Consider the grid H = (a + 2, b + 2, 1), and observe that such a grid admits an optimal percolating set by construction 7.9. Note that

$$SA(a, b, 3) = [SA(a + 2, b + 2, 1)] - 3.$$

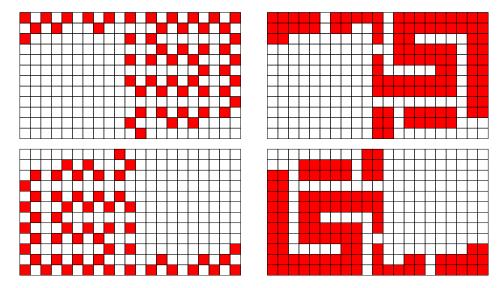


Figure 7.12: A perfect percolating set for (12, 21, 2).

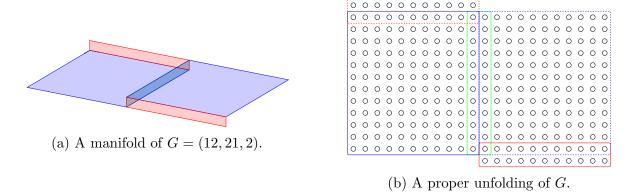


Figure 7.13: A proper unfolding of G = (12, 21, 2). Colored rectangles indicate faces of G. Dashed lines indicate that cells appear on different layers.

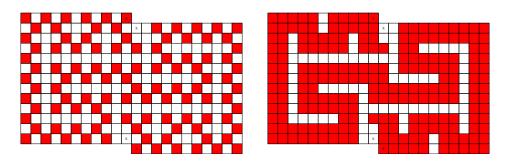


Figure 7.14: A percolating set on the proper unfolding of G = (12, 21, 2).

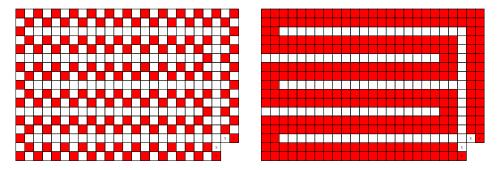


Figure 7.15: A percolating set on the proper unfolding H' of G = (15, 23, 3).

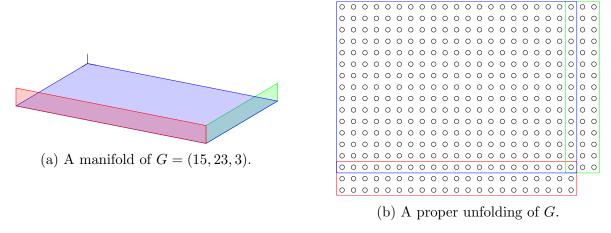


Figure 7.16: A proper unfolding of G = (15, 23, 3). Colored rectangles indicate faces of G.

We show that an unfolding of G can be obtained from a simple augmentation of H. Let H' be the grid obtained by deleting the four vertices in the bottom, right-most corner of H (see figure 7.15). Consider the folding pattern illustrated in figure 7.16, and observe that the pairs of vertices adjacent to the deleted region are duplicates of each other. (In other words, consider folding up the red and green regions in figure 7.16, and notice that this operation causes vertices to overlap.) Taking this into account, the unfolding of G percolates by lemma 7.1. Since H admits an optimal percolating set of size $\lceil SA(a+2,b+2,1) \rceil$, and precisely 3 of the vertices deleted from H to obtain H' were infected, it follows that the unfolding of G percolates at the lower bound. Finally, by lemma 7.7, since the unfolding of G is proper and percolates at the lower bound, G is perfect.

7.6 Individual constructions

In this section, we diagram lethal set constructions for single grids. The initial infection A is colored red, and all other cells are labeled with the time t that they are first

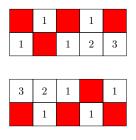


Figure 7.17

infected.

Construction 7.15. The grid (5,2,2) is perfect.

Proof. See figure 7.17.

Construction 7.16. The grid (5,5,5) is perfect.

Proof. See figure 7.18.

Construction 7.17. The grid (8,5,5) is perfect.

Proof. See figure 7.19.

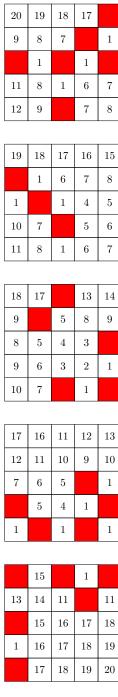


Figure 7.18

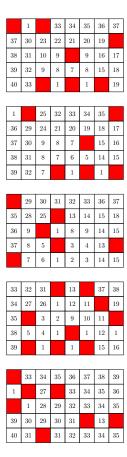


Figure 7.19

Chapter 8 Programmatic Approach

Chapter 9 Torus

9.1 Introduction

I know I shouldn't be working on this now but I have some ideas. We can use Benevides proof of the lower bound for the torus, which states the following:

Lemma 9.1. Let $T = C_n \square C_m$ be the $n \times n$ torus. Let A be a lethal set on T. Then $|A| \ge \frac{nm+c}{3}$, where c is the number of components in the graph $T[\overline{A}]$.

There are some nice observations to be made here. First, in much the same way that loss of surface area contributes to fractional increases in the surface area bound for grid graphs, the number of components can be understood to indicate sub-optimality in the torus. As a particular example, the torus $T = C_1 \otimes C_1 \otimes C_2 \otimes C_1 \otimes C_1 \otimes C_2 \otimes C_2$

There's something interesting happening here. In this particular case, $T[\overline{A}]$ has three components. Assuming the best case scenario (where it has one component), we get the lower bound of 72.33. However, if we instead consider three components, the lower bound sits precisely as 73. As the construction of T simple involves adding an additional vertex to an optimal (not perfect) grid construction, it appears as though the inadequacies of the grid construction (namely, cells that are infected by 4 neighbors) somehow translate directly to inadequacies of the torus construction (multiple components in the complement).

Oh wait, is this just because the surface area argument is fundamentally an argument about the number of components in the complement? I think it is. Every time you lose surface area, you've deleted a tree from the forest (because in order to remove the final vertex in a tree, you need an inefficiency—unless the tree sits on the boundary). Therefore, it makes sense that the best solutions on the torus have exactly one component in the complement.

How does the surface area argument map to an argument about components? Knowing this should allow us to construct a lower bound on the 4D torus.

There are other questions regarding some form of algebraic or numerical equivalence between expressions for the lower bounds of tori and grids. It is clear that

Chapter 10 Conclusion

Bibliography

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