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1 Question-1

1.1 Program-1

Sequences that are accepted by the program:

This program takes input as sequence of numbers and then prints "yes" if the numbers are in arithmetic sequence and prints "no" otherwise.

1.1.1 How I reached this conclusion:

In order to analyze the binary file we first disassemble it by using the Linux command:

objdump -DM intel program1 > output.asm

Now we see many section, following are the important section that tells us what the program actually do :

- .plt
- .main

1.1.2 .plt

```
0000000000401020 <.plt>:
401020: ff 35 e2 2f 00 00    push    QWORD PTR [rip+0x2fe2]    # 404008 <_GLOBAL_OFFSET_TABLE_+0x8>
401026: ff 25 e4 2f 00 00    jmp     QWORD PTR [rip+0x2fe4]    # 404010 <_GLOBAL_OFFSET_TABLE_+0x10>
40102c: 0f 1f 40 00         nop     DWORD PTR [rax+0x0]

0000000000401030 <puts@plt>:
401030: ff 25 e2 2f 00 00    jmp     QWORD PTR [rip+0x2fe2]    # 404018 <puts@GLIBC_2.2.5>
401036: 68 00 00 00 00      push    0x0
40103b: e9 e0 ff ff ff      jmp     401020 <.plt>

0000000000401040 <printf@plt>:
401040: ff 25 da 2f 00 00    jmp     QWORD PTR [rip+0x2fda]    # 404020 <printf@GLIBC_2.2.5>
401046: 68 01 00 00 00      push    0x1
40104b: e9 d0 ff ff ff      jmp     401020 <.plt>

0000000000401050 <_isoc99 scanf@plt>:
401050: ff 25 d2 2f 00 00    jmp     QWORD PTR [rip+0x2fd2]    # 404028 <_isoc99_scanf@GLIBC_2.7>
401056: 68 02 00 00 00      push    0x2
40105b: e9 c0 ff ff ff      jmp     401020 <.plt>
```

The role of this block is that it provides the c function puts, scanf and printf to be used in the code . First we jump to the Global Offset Table (GOT) entry at a particular address like **[rip+0x2fe2]** which stores the required function

and then pushes a value like, **0x0** onto the stack. This is typically used for passing arguments to functions . Then we again jump to .plt.

1.1.3 .main

1.1.4 code before printing

```
401146: 55          push    rbp
401147: 48 89 e5    mov     rbp, rsp
40114a: 48 83 ec 30  sub    rsp, 0x30
40114e: 89 7d dc    mov     DWORD PTR [rbp-0x24], edi
401151: 48 89 75 d0  mov     QWORD PTR [rbp-0x30], rsi
401155: bf 08 20 40 00 mov     edi, 0x402008
40115a: b8 00 00 00 00 mov     eax, 0x0
40115f: e8 dc fe ff ff call    401040 <printf@plt>
```

First rbp is pushed into the stack , then the stack pointer address is stored in it , then stack pointer is moved down to allocate 48 bytes of memory , this is a buffer that stores 3 16 bit integers that are given. Then the value stored in registers edi and rsi are stored as 32 bit value and 64 bit value at positions **[rbp-0x24]** and **[rbp-0x30]** respectively.

Then edi is loaded with an address of formatted string "Enter three or more numbers (Terminate with CTRL+D) : " and eax is loaded with zero so that the string stored in edi is printed when printf is called. Then printf is called and the string is printed.

1.1.5 code after printing

```
401164: c7 45 fc 00 00 00 00 mov     DWORD PTR [rbp-0x4], 0x0
40116b: c7 45 f4 00 00 00 00 mov     DWORD PTR [rbp-0xc], 0x0
401172: c6 45 f3 01    mov     BYTE PTR [rbp-0xd], 0x1
401176: eb 67        jmp     4011df <main+0x99>
```

Then we store three variables and initialise them

- **[rbp-0x4]** initialised to 0 (consider this as some kind of offset o)
- **[rbp-0xc]** initialised to 0 (consider this as the number of inputs given calls)
- **[rbp-0xd]** initialised to 1 (consider this as a flag call f)

The reason for the above variables would be understood later.

1.1.6 Block 4011df

```
4011d5: 01 c2          add    edx,eax
4011d7: 83 e2 01       and    edx,0x1
4011da: 29 c2          sub    edx,eax
4011dc: 89 55 fc       mov    DWORD PTR [rbp-0x4],edx
4011df: 8b 45 fc       mov    eax,DWORD PTR [rbp-0x4]
4011e2: 48 98          cdqe
4011e4: 48 8d 14 85 00 00 00 lea    rdx,[rax*4+0x0]
4011eb: 00
4011ec: 48 8d 45 e4     lea    rax,[rbp-0x1c]
4011f0: 48 01 d0       add    rax,rdx
4011f3: 48 89 c6       mov    rsi,rax
4011f6: bf 40 20 40 00 mov    edi,0x402040
4011fb: b8 00 00 00 00 mov    eax,0x0
401200: e8 4b fe ff ff call   401050 <__isoc99_scanf@plt>
401205: 83 f8 01       cmp    eax,0x1
401208: 0f 84 6a ff ff je     401178 <main+0x32>
```

Then we jump to the address 4011df. There **[rbp-0x4]** i.e offset is loaded in eax ,then it is sign extended so that $rax = eax$, then $rdx = rax*4$ and rax is loaded with a new quantity **[rbp-0x1c]** which holds the address of previous input that was stored in the buffer. Then rax becomes $rax + rdx$ which is equivalent to saying that **$rax = \text{previous address} + \text{offset} * 4$** , we would soon see how offset works.Then the updated address is loaded in rsi and edi is loaded with some address and eax is loaded with zero to make the syscall and then scanf is called which would store the input in the updated address in the buffer. Now if the input given to the scanf was a number then eax becomes 1 and it jumps to then address 401178 which is the main logic of the code , else eax would be zero and we go ahead and check the further code.

1.1.7 Block 401178

```
401178: 83 45 f4 01     add    DWORD PTR [rbp-0xc],0x1
40117c: 83 7d f4 01     cmp    DWORD PTR [rbp-0xc],0x1
401180: 7e 45          jle    4011c7 <main+0x81>
```

This is the code at 401178. This increments the **[rbp-0xc]** i.e our size variable s by 1 as we got a input as a number and then check if the size variable was greater than 1 , if not then it jumps to a very interesting piece of code , the address 4011c7 , else it continues the further code

1.1.8 Block 4011c7 : Increment logic

```
4011c7: 8b 45 fc      mov     eax,DWORD PTR [rbp-0x4]
4011ca: 8d 50 01      lea     edx,[rax+0x1]
4011cd: 89 d0         mov     eax,edx
4011cf: c1 f8 1f      sar     eax,0x1f
4011d2: c1 e8 1f      shr     eax,0x1f
4011d5: 01 c2         add     edx,eax
4011d7: 83 e2 01      and     edx,0x1
4011da: 29 c2         sub     edx,eax
4011dc: 89 55 fc      mov     DWORD PTR [rbp-0x4],edx
```

First the quantity **[rbp-0x4]** i.e offset is loaded in **eax** , then $edx = eax + 1$, then $eax = edx$ then the following two lines basically makes $eax = 1$ if it was odd , else if it was even then eax becomes 0. Then $edx = eax + edx$ i.e $edx = edx + (0 \text{ or } 1)$ then edx and 1 is done which check the parity of edx and then $edx = edx - eax$ and then this edx is loaded in **[rbp-0x4]** as the new offset. If analyzed carefully we can say that if initial offset was zero then after this step it becomes 1 ,and if it was 1 initially then offset becomes zero. I call this block of code as the **increment logic**. After this again we go to block 4011df and take the input the number and if ctrl D is not given then we again go to the block 401178 and it goes on .We can see that this is forming a loop , now when ctrl D is hit we don't go to 401178 , instead we move ahead and do the following:

1.1.9 Exit logic

```
40120e: 83 7d f4 02   cmp     DWORD PTR [rbp-0xc],0x2
401212: 7f 11         jg      401225 <main+0xdf>
401214: bf 48 20 40 00 mov     edi,0x402048
401219: e8 12 fe ff ff call    401030 <puts@plt>
40121e: b8 ff ff ff ff mov     eax,0xffffffff
401223: eb 21         jmp     401246 <main+0x100>
401225: 80 7d f3 00   cmp     BYTE PTR [rbp-0xd],0x0
401229: 74 0c         je      401237 <main+0xf1>
40122b: bf 77 20 40 00 mov     edi,0x402077
401230: e8 fb fd ff ff call    401030 <puts@plt>
401235: eb 0a         jmp     401241 <main+0xfb>
401237: bf 7b 20 40 00 mov     edi,0x40207b
40123c: e8 ef fd ff ff call    401030 <puts@plt>
401241: b8 00 00 00 00 mov     eax,0x0
401246: c9           leave
401247: c3           ret
401248: 0f 1f 84 00 00 00 00 00 nop     DWORD PTR [rax+rax*1+0x0]
40124f: 00
```

Here first **[rbp-0xc]** is checked , if number of entry is greater than two and we already know that ctrl d is pressed then it goes to 401225 where it compares the flag to 0 , if the flag is zero then it prints no and makes `eax = 0` and then leave and return , else if flag is 1 then it prints yes , makes `eax = 0` and then leave and return. If number of entry less than two then print the string that "Enter more than three numbers"

1.1.10 Block after 401178 : main logic

```

401182:      8b 45 f8      mov     eax,DWORD PTR [rbp-0x8]
401185:      89 45 ec      mov     DWORD PTR [rbp-0x14],eax
401188:      8b 45 fc      mov     eax,DWORD PTR [rbp-0x4]
40118b:      48 98          cdqe
40118d:      8b 4c 85 e4    mov     ecx,DWORD PTR [rbp+rax*4-0x1c]
401191:      8b 45 fc      mov     eax,DWORD PTR [rbp-0x4]
401194:      8d 50 01      lea     edx,[rax+0x1]
401197:      89 d0          mov     eax,edx
401199:      c1 f8 1f      sar     eax,0x1f
40119c:      c1 e8 1f      shr     eax,0x1f
40119f:      01 c2          add     edx,eax
4011a1:      83 e2 01      and     edx,0x1
4011a4:      29 c2          sub     edx,eax
4011a6:      89 d0          mov     eax,edx
4011a8:      48 98          cdqe
4011aa:      8b 44 85 e4    mov     eax,DWORD PTR [rbp+rax*4-0x1c]
4011ae:      29 c1          sub     ecx,eax
4011b0:      89 ca          mov     edx,ecx
4011b2:      89 55 f8      mov     DWORD PTR [rbp-0x8],edx
4011b5:      83 7d f4 02    cmp     DWORD PTR [rbp-0xc],0x2
4011b9:      7e 0c          jle     4011c7 <main+0x81>
4011bb:      8b 45 ec      mov     eax,DWORD PTR [rbp-0x14]
4011be:      3b 45 f8      cmp     eax,DWORD PTR [rbp-0x8]
4011c1:      74 04          je      4011c7 <main+0x81>
4011c3:      c6 45 f3 00    mov     BYTE PTR [rbp-0xd],0x0

```

Here first **[rbp-0x8]** which stores the difference of previous two inputs (we would see how) is loaded in `eax` and then `eax` is loaded to **[rbp-0x14]** (consider this as a temp var that stores the previous difference) , now `eax` is loaded with **[rbp-0x4]** the current offset and then sign extension so `rax = eax` , now the line **DWORD PTR [rbp+rax*4-0x1c]** says that `ecx` holds some value stored in array corresponding to current offset , this must be our new number.

Now again `eax` is loaded with current offset and a bunch of same calculation are done which are same as in block 4011c7 , which just inverts the offset , so at the line 4011a6 `eax` holds the inverted offset and then it is sign extended to `rax` then the `eax` is loaded with **DWORD PTR [rbp+rax*4-0x1c]** which would be the number corresponding to the previous offset , so we can say that

this was the previous number stored.

Finally the difference between ecx and eax is made and stored in `[rbp-0x8]` which would be the new difference.

Now if number of elements are greater than 2 then the values of previous difference and current difference are compared, if they are same then jump to increment block else flag becomes 0. If the number of elements were less than or equal to zero then we would have just jumped to the increment block.

1.1.11 Conclusion

In this program as we are inputting the numbers, previous difference was maintained and current difference was calculated, if they were equal then flag remained 1, because it was initialised with value 1, if the difference was not same then flag become 0 for rest of the program and when we hit ctrl-D if the flag was 1 "Yes" was printed else "No" was printed.

1.2 Program-2

1.2.1 What is the function?

The function that was given in program 2 takes a number as an input, if the number is 0 then it returns 1, else if a number n is given then it returns

$$F(n) = \sum_{k=1}^n F(k-1) * F(n-k) \quad (1)$$

1.2.2 How I got to this result ?

Similar to previous question we first disassemble the program and again we look at important parts of code. Again there is a .plt section which is exact same as the previous program, then there are two important block of codes :

- func
- main

We would first analyze the main program and then analyze the func portion.

1.2.3 Analysis of Main

```
00000000004011a7 <main>:
4011a7: 55                push    rbp
4011a8: 48 89 e5          mov     rbp, rsp
4011ab: 48 83 ec 10        sub     rsp, 0x10
4011af: bf 08 20 40 00     mov     edi, 0x402008
4011b4: b8 00 00 00 00     mov     eax, 0x0
4011b9: e8 72 fe ff ff     call    401030 <printf@plt>
4011be: 48 8d 45 f8        lea     rax, [rbp-0x8]
4011c2: 48 89 c6          mov     rsi, rax
4011c5: bf 27 20 40 00     mov     edi, 0x402027
4011ca: b8 00 00 00 00     mov     eax, 0x0
4011cf: e8 6c fe ff ff     call    401040 <__isoc99_scanf@plt>
4011d4: 48 8b 45 f8        mov     rax, QWORD PTR [rbp-0x8]
4011d8: 48 89 c7          mov     rdi, rax
4011db: e8 56 ff ff ff     call    401136 <func>
4011e0: 48 89 c6          mov     rsi, rax
4011e3: bf 2c 20 40 00     mov     edi, 0x40202c
4011e8: b8 00 00 00 00     mov     eax, 0x0
4011ed: e8 3e fe ff ff     call    401030 <printf@plt>
4011f2: b8 00 00 00 00     mov     eax, 0x0
4011f7: c9                leave   %edi
4011f8: c3                ret
4011f9: 0f 1f 80 00 00 00 00 00  nop     DWORD PTR [rax+0x0]
```

Here first the rbp, the base pointer is pushed into the stack and then the rsp, stack pointer is loaded into the rbp then the stack pointer is decremented, creating a 16 bit space where variables could be stored.

Then edi is loaded with an address of a formatted string and eax is loaded with 0 and then printf is called which print "Enter a non-negative integer:" to the terminal. Now the address [rbp-0x8] is stored in rax and then rsi is loaded with rax, then edi and eax are loaded with some constants for the syscall, then the scanf is called which would store the inputted value inside the address that is stored in rsi i.e at the address [rbp-0x8]. Then the value from this address is again moved to the register rdi and the function is called, this effectively means that the inputted number is given as argument to this function.

Now the value that would be returned by the function would be stored at rax. Again move rax into rsi, load the constants inside edi and eax and then printf is called which would print the value inside rsi, then eax would be made zero then return to the main.

1.2.4 Analysis of Func

```

0000000000401136 <func>:
401136: 55                push    rbp
401137: 48 89 e5          mov     rbp, rsp
40113a: 53                push    rbx
40113b: 48 83 ec 28       sub     rsp, 0x28
40113f: 48 89 7d d8       mov     QWORD PTR [rbp-0x28], rdi
401143: 48 83 7d d8 00    cmp     QWORD PTR [rbp-0x28], 0x0
401148: 75 07            jne     401151 <func+0x1b>
40114a: b8 01 00 00 00    mov     eax, 0x1
40114f: eb 50            jmp     4011a1 <func+0x6b>
401151: 48 c7 45 e8 00 00 00 mov     QWORD PTR [rbp-0x18], 0x0
401158: 00
401159: 48 c7 45 e0 01 00 00 mov     QWORD PTR [rbp-0x20], 0x1
401160: 00
401161: eb 30            jmp     401193 <func+0x5d>
401163: 48 8b 45 e0       mov     rax, QWORD PTR [rbp-0x20]
401167: 48 83 e8 01       sub     rax, 0x1
40116b: 48 89 c7          mov     rdi, rax
40116e: e8 c3 ff ff ff   call    401136 <func>
401173: 48 89 c3          mov     rbx, rax
401176: 48 8b 45 d8       mov     rax, QWORD PTR [rbp-0x28]
40117a: 48 2b 45 e0       sub     rax, QWORD PTR [rbp-0x20]
40117e: 48 89 c7          mov     rdi, rax
401181: e8 b0 ff ff ff   call    401136 <func>
401186: 48 0f af c3       imul    rax, rbx
40118a: 48 01 45 e8       add     QWORD PTR [rbp-0x18], rax
40118e: 48 83 45 e0 01    add     QWORD PTR [rbp-0x20], 0x1
401193: 48 8b 45 e0       mov     rax, QWORD PTR [rbp-0x20]
401197: 48 39 45 d8       cmp     QWORD PTR [rbp-0x28], rax
40119b: 73 c6            jae     401163 <func+0x2d>
40119d: 48 8b 45 e8       mov     rax, QWORD PTR [rbp-0x18]
4011a1: 48 8b 5d f8       mov     rbx, QWORD PTR [rbp-0x8]
4011a5: c9                leave   rax
4011a6: c3                ret

```

First the base pointer is loaded in the stack, `rsp` address is loaded into `rbp`, also `rbx` is pushed in the stack, then `rsp` is decremented to store 40 bits of variables, then the input value which is stored in `rdi` is loaded in address `[rbp-0x28]`, then if the input value is 0 then program skips jump to 401151 and makes `eax = 1` then jumps to 4011a1 where `rbx` is loaded with `[rbp-0x8]` and then leave and return, this `eax` value would then be 1 which makes `rax` value to be 1 and then the main prints the value inside `rax`. Else if the input is not 0 then jump to address 401151. Now would analyse block 401151.

1.2.5 Block 401151

```
401151:  48 c7 45 e8 00 00 00    mov     QWORD PTR [rbp-0x18],0x0 (var 1)
401158:  00
401159:  48 c7 45 e0 01 00 00    mov     QWORD PTR [rbp-0x20],0x1 (var 2)
401160:  00
401161:  eb 30                  jmp     401193 <func+0x5d>
```

Following things happens:

- Value 0 is stored in address **[rbp-0x18]** (call this variable ans)
- Value 1 is stored in address **[rbp-0x20]** (call this variable i)
- jump to 401193 address

1.2.6 Block 401193

```
401193:  48 8b 45 e0            mov     rax,QWORD PTR [rbp-0x20]
401197:  48 39 45 d8            cmp     QWORD PTR [rbp-0x28],rax
40119b:  73 c6                  jae     401163 <func+0x2d>
```

- rax is loaded with **[rbp-0x20]**
- If its value is less than or equal to input value then jump to 401163
- Else continue for the termination of the code

1.2.7 Block 401163

```
401163:    48 8b 45 e0    mov     rax,QWORD PTR [rbp-0x20]
401167:    48 83 e8 01    sub     rax,0x1
40116b:    48 89 c7        mov     rdi,rax
40116e:    e8 c3 ff ff ff call     401136 <func>
401173:    48 89 c3        mov     rbx,rax
401176:    48 8b 45 d8    mov     rax,QWORD PTR [rbp-0x28]
40117a:    48 2b 45 e0    sub     rax,QWORD PTR [rbp-0x20]
40117e:    48 89 c7        mov     rdi,rax
401181:    e8 b0 ff ff ff call     401136 <func>
401186:    48 0f af c3    imul    rax,rbx
40118a:    48 01 45 e8    add     QWORD PTR [rbp-0x18],rax
40118e:    48 83 45 e0 01 add     QWORD PTR [rbp-0x20],0x1

401193:    48 8b 45 e0    mov     rax,QWORD PTR [rbp-0x20]
401197:    48 39 45 d8    cmp     QWORD PTR [rbp-0x28],rax
40119b:    73 c6          jae     401163 <func+0x2d>
```

Following things happen:

- First $rax = i$ and then $rax = rax - 1$ and $rdi = rax$
- Then control flow goes to 4001136 to the start of function with argument $rdi = i-1$
- Then the value returned to rax is stored in rbx
- Then rax is loaded with **[rbp-0x28]** i.e $rax = \text{input}$ and then $rax = rax - i$ i.e $rax = \text{input} - i$
- Then again the function is called and its value is stored in rax .
- Then $rax = rax * rbx$ then $ans = ans + rax$ i.e $ans = ans + f(i-1) * f(\text{input} - i)$
- Then again block 401193 is executed if the value of i is less than or equal to input value, if yes then again loop is done else we exit the loop

1.2.8 Conclusion

- If the value given to function is 0 then it returns 1
- If given value n then it loops from 1 to n and in each iteration it adds $f(i-1)*f(n-i)$ to the ans where i is the iterator and ans is the value to be returned
- Hence we got the answer what the function does